

# Introduction to Computer Networks

## Rate and Delay Guarantees (§5.4.4)



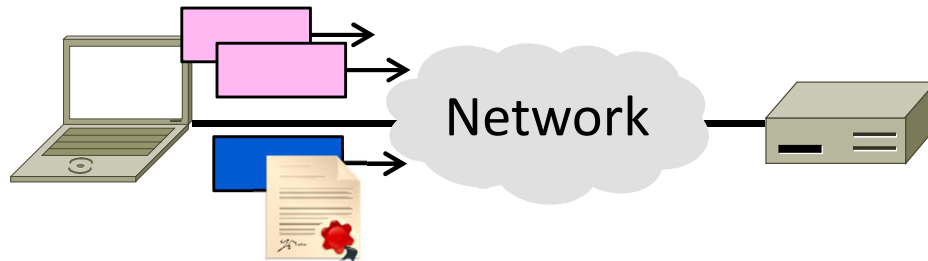
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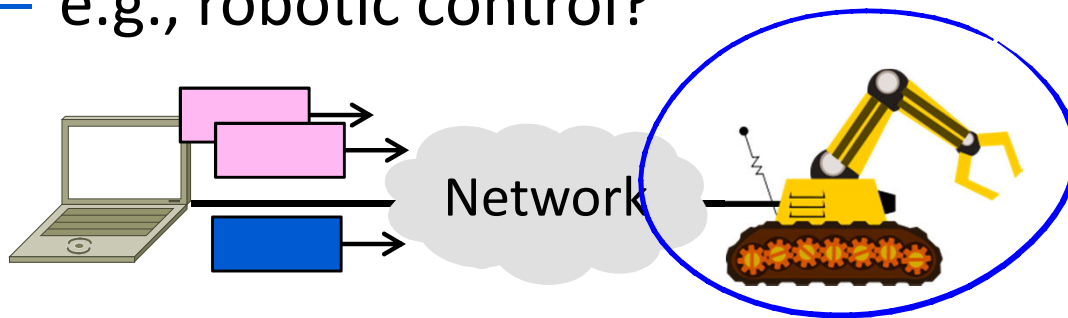
# Topic

- Guaranteeing performance for traffic flows across in the network
  - This is “hard QOS” with a firm guarantee for a traffic flow



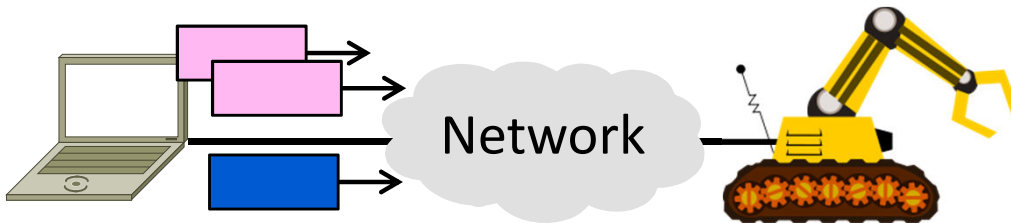
# Motivation

- Sometimes we want guaranteed service – like the telephone network
  - Minimum rate and maximum delay regardless of how other flows behave
  - e.g., robotic control?



## Motivation (2)

- Could provision a dedicated circuit (or build a network), but expensive
- Can we have statistical multiplexing together with hard guarantees?

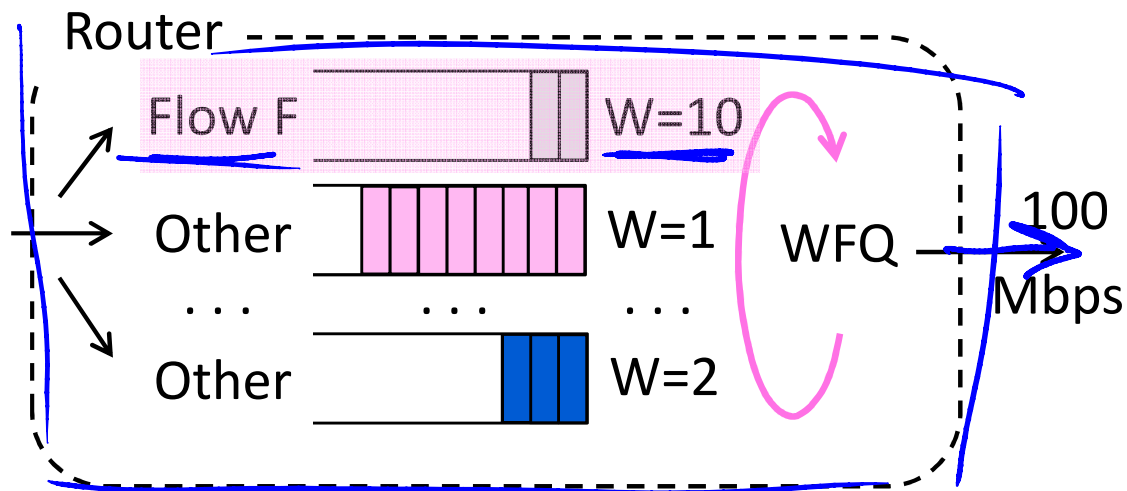


# Admission Control

- Suppose we have a flow F that needs rate  $\geq R$  Mbps and delay  $\leq D$  secs
- We must decide whether to admit or reject it from the network
  - This is admission control
  - Rejecting should be infrequent
- Key point is we need the ability to control load to make guarantees

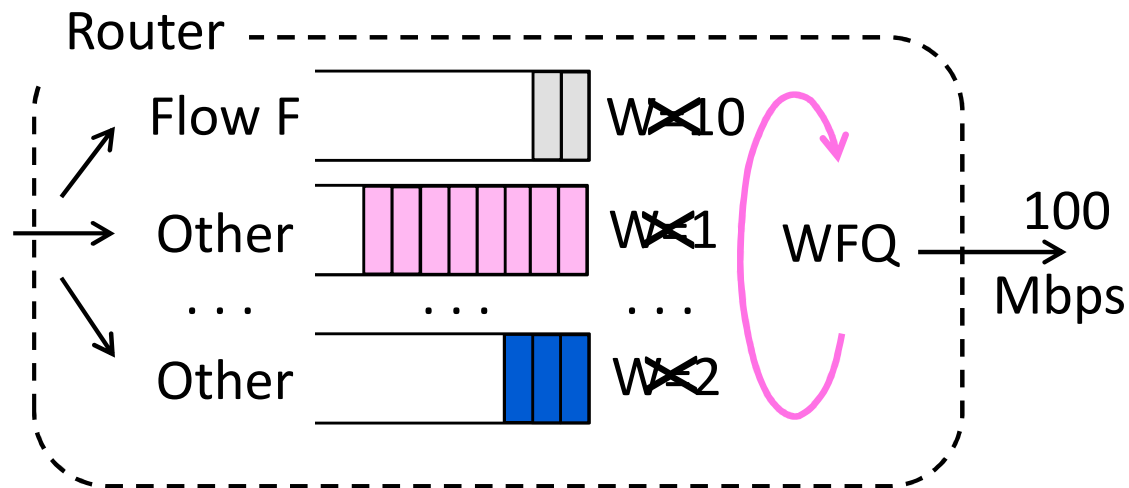
# Router Rate Guarantee

- WFQ can guarantee rate at a router
  - What rate will Flow F get?



# Router Rate Guarantee (2)

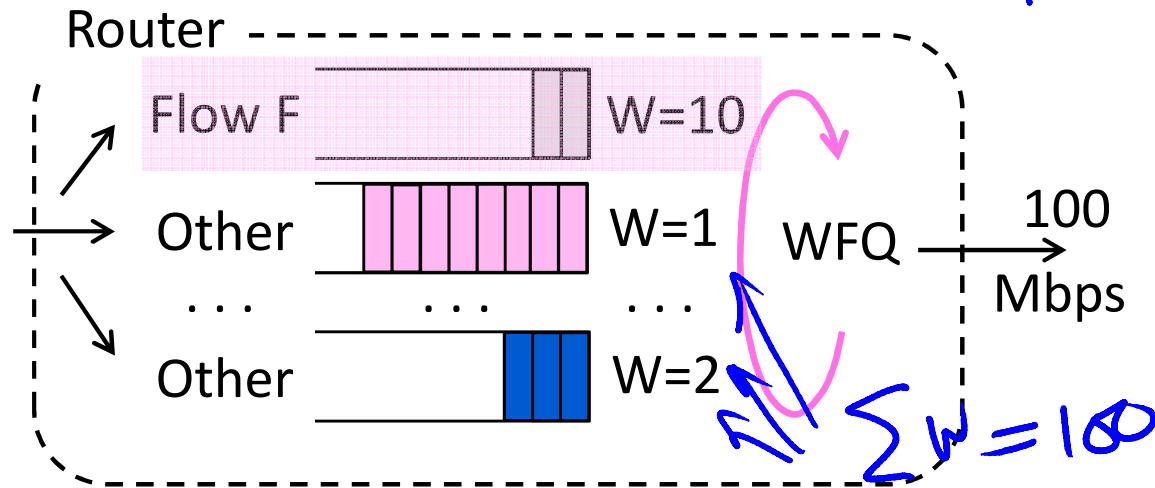
- Consider  $N$  flows with weight 1
  - Each flow gets  $1/N$ th share under load
  - Or at least  $100/N$  Mbps



# Router Rate Guarantee (3)

- Consider flow F with weight 10
  - Suppose weight of all flows is 100

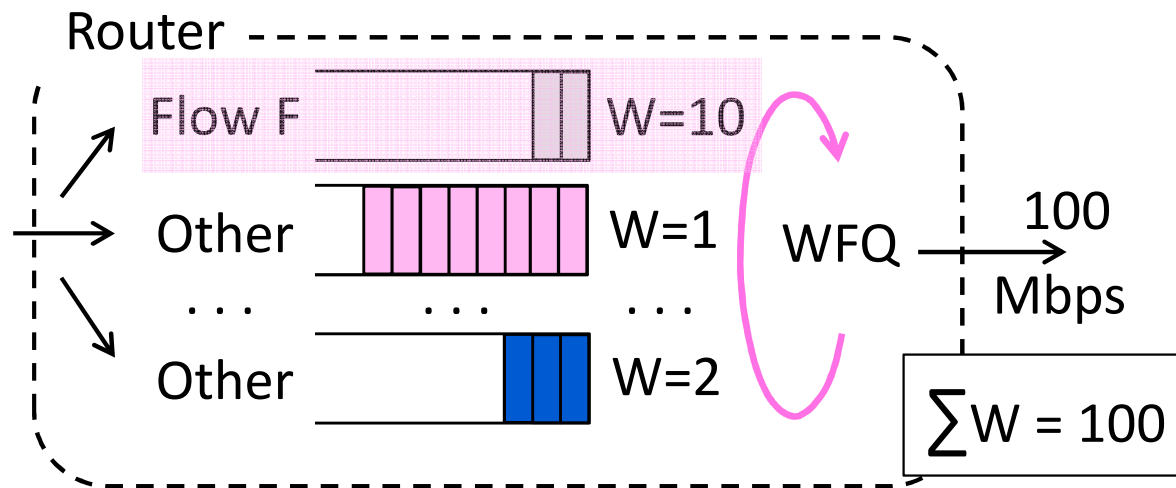
$$\text{Flow F} \rightarrow 10/100 \times 100\text{Mbps} = 10\text{ Mbps}$$





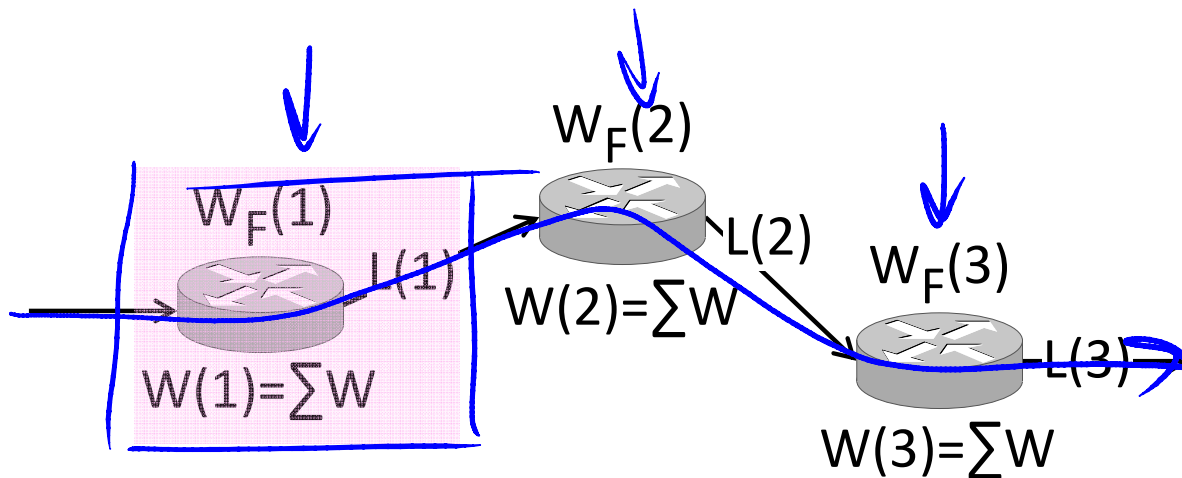
# Router Rate Guarantee (4)

- Consider flow F with weight 10
  - Flow F gets  $\geq (10/100) \cdot 100 = 10$  Mbps



# Network Rate Guarantee

- We can guarantee a minimum rate for a network path by guaranteeing it at each router

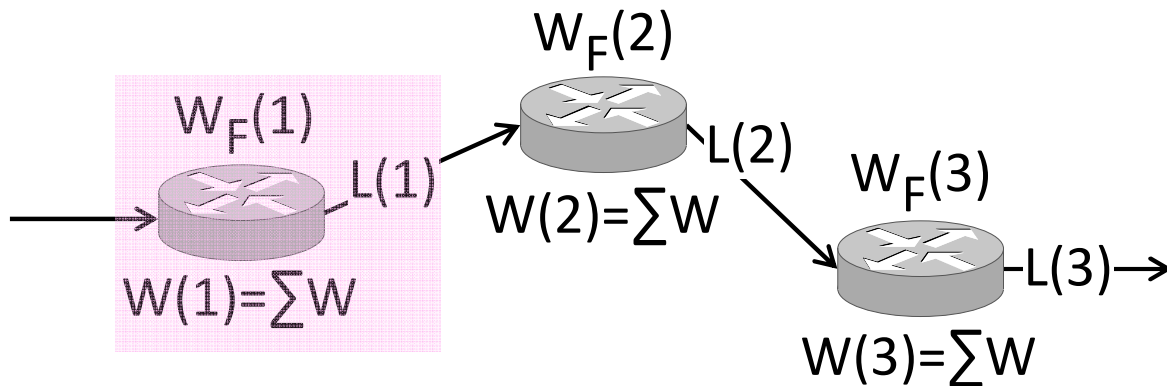


# Network Rate Guarantee (2)

- Condition for each router:

For all routers  $i$ :

$$\underline{W_F(i)} / \underline{W(i)} * \underline{L(i)} \geq R \text{ Mbps}$$

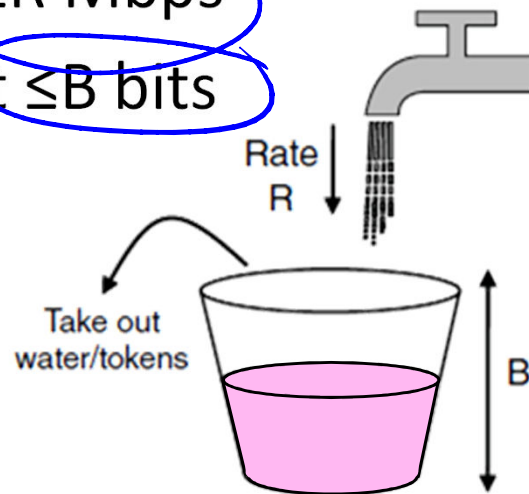


# Delay Guarantee

- What about the queuing delay?
  - How much larger than latency might the delay be, given rate guarantee?
- It depends on the traffic flow
  - If exceeds R Mbps then queues may build and delay will grow ...
- Need to shape traffic for guarantee
  - We'll use token buckets ☺

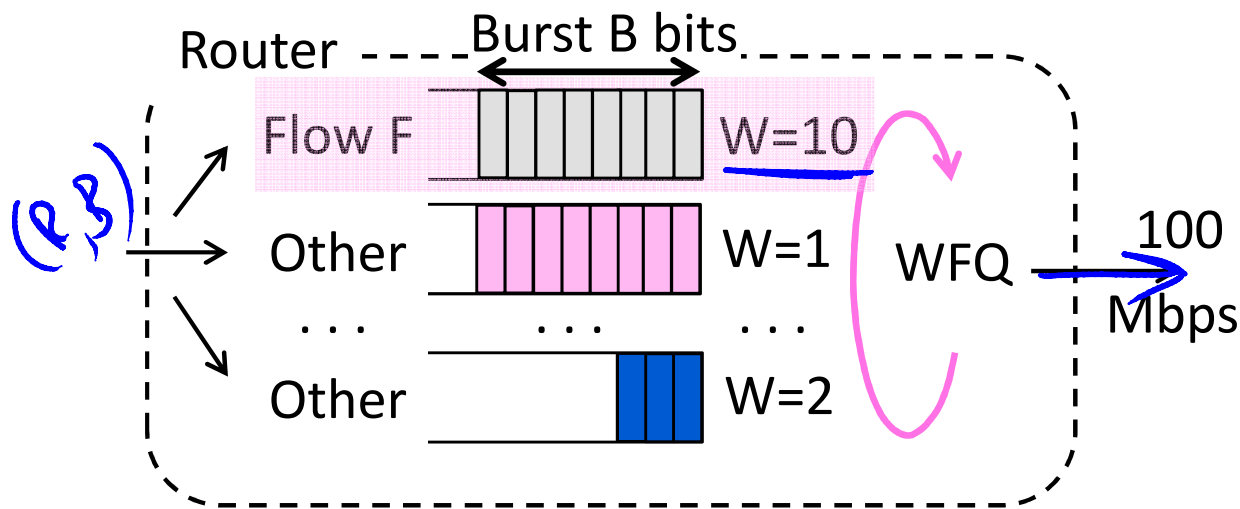
# Router Delay Guarantee

- Assume traffic flow F is shaped by an (R, B) token bucket
  - Long-term rate  $\leq R$  Mbps
  - Short-term burst  $\leq B$  bits



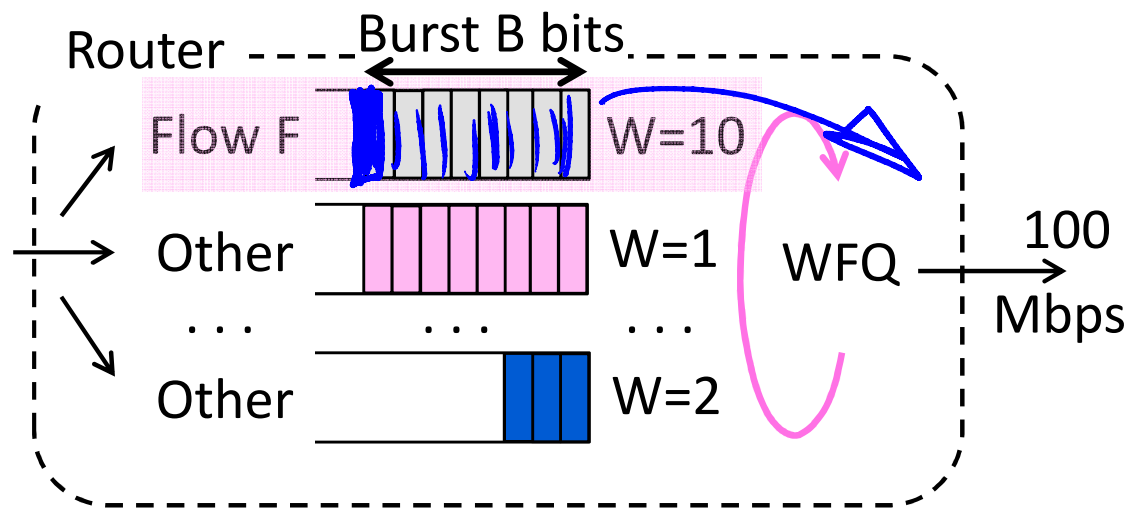
# Router Delay Guarantee (2)

- What is delay of flow F at a router?
  - Traffic shaped by  $(R, B)$  token bucket
  - WFQ with weight set for rate  $\geq R$  Mbps



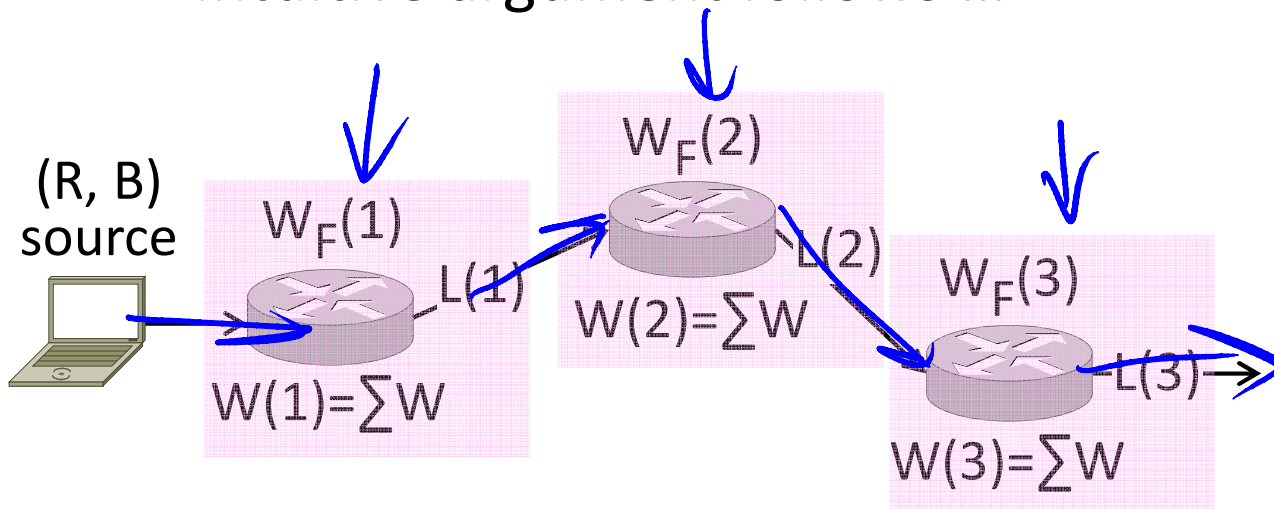
# Router Delay Guarantee (3)

- In worst case B arrives all at once
  - So queuing delay is  $\leq B/R$  seconds



# Network Delay Guarantee

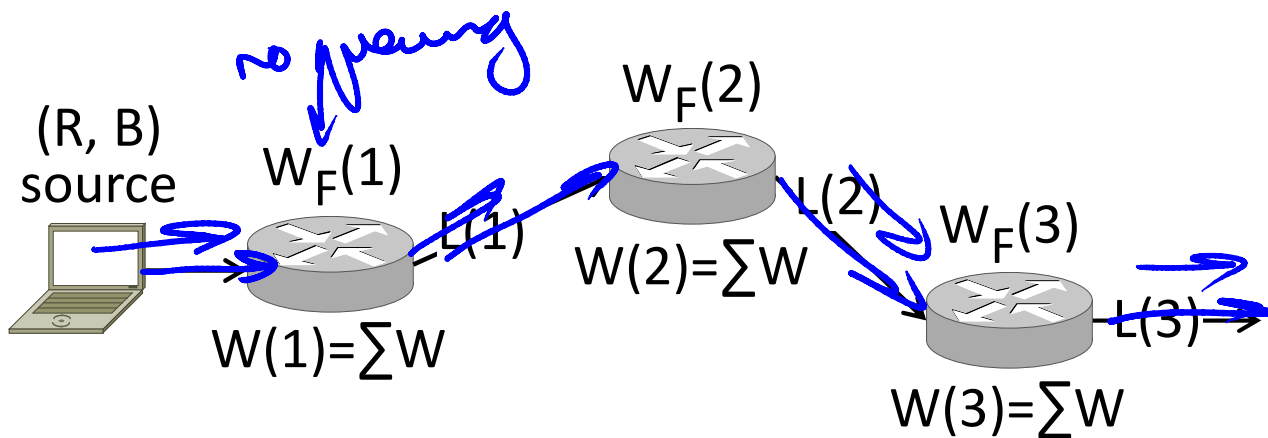
- What is the delay across  $N$  routers?
  - This is tricky! Each router add delays
  - Bound of  $N*B/R$  is too loose
  - Intuitive argument follows ...





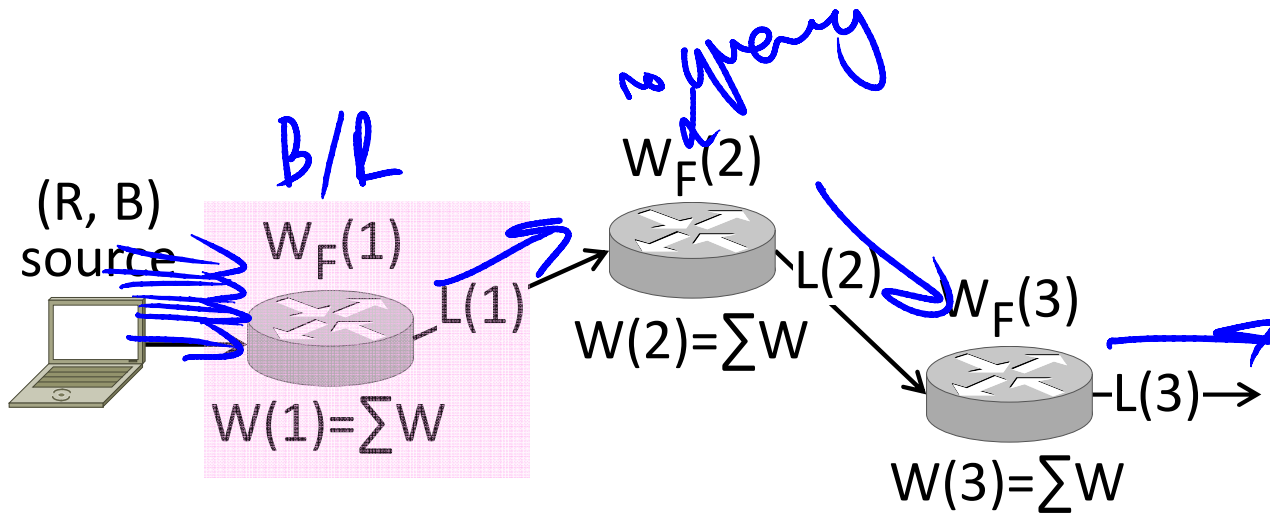
# Network Delay Guarantee (2)

- If traffic is perfectly smooth at rate  $R$  (no bursts) then queuing delay is zero
  - Packet enters router just in time to leave
  - Delay is latency (propagation, transmission)



# Network Delay Guarantee (3)

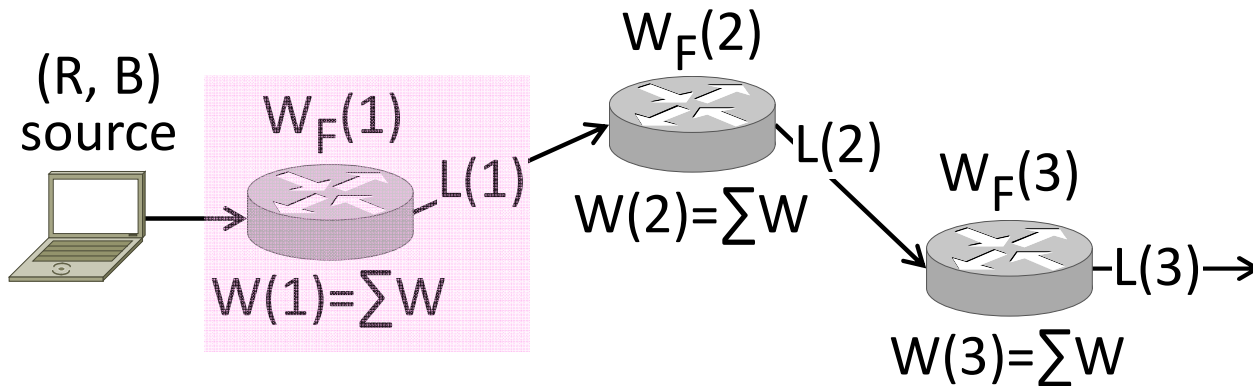
- Observe if traffic pays for burst B at one router, it is smoothed for the next  
➔ Burst delay is only paid once!



# Network Delay Guarantee (4)

- Delay across N routers:

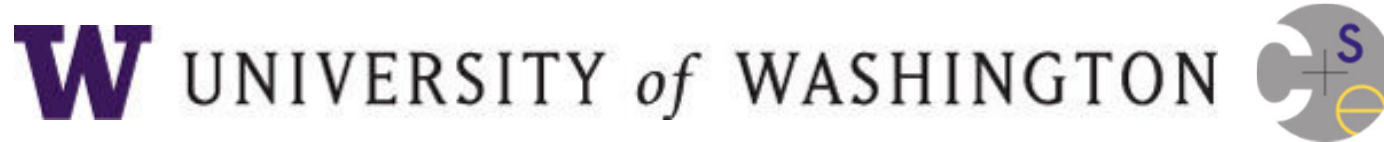
$$\text{Delay} \leq \text{Latency terms} + B/R$$



# Rate/Delay Guarantee

- Given a network with:
  - (R, B) shaped traffic flow
  - WFQ routers with proper weights
  - Sharing via statistical multiplexing
- We can guarantee the flow a minimum rate and maximum delay
  - Rate is  $\geq R$  Mbps
  - Delay is  $\leq \text{latency} + B/R$  secs
- Regardless of how other flows behave

# END



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