

# [CSE3081(2반)] 알고리즘 설계와 분석

2020학년도 2학기

강의자료

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# Mathematical Induction & Proof of Correctness

- Proof by induction

$P(k)$  :  $k$  번째 항은 ... 하다.  $\begin{matrix} T \\ \text{statement} \\ \text{(predicate)} \end{matrix}$

$0 \ 1 \ 2 \ \dots \ k \ k+1 \ \dots$

$\uparrow$  Base case (Basis)

$\uparrow$  Induction step

$k \geq 0$

- Proof of correctness: MSS (1D)

```

/* 1*/   ThisSum = MaxSum = 0;
/* 2*/   for( j = 0; j < N; j++ )
{
/* 3*/       ThisSum += A[ j ];

/* 4*/       if( ThisSum > MaxSum )
/* 5*/           MaxSum = ThisSum;
/* 6*/       else if( ThisSum < 0 )
/* 7*/           ThisSum = 0;
}
/* 8*/   return MaxSum;
    
```

$P(j)$  : for-loop이  $j$  번 수행한 직후에  
 ThisSum 변수는 ( ) 값을,  
 MaxSum 변수는 ( ) 값을  
 가지고 있다.

$k = 0, 1, 2, \dots, j-1, j, \dots, N$

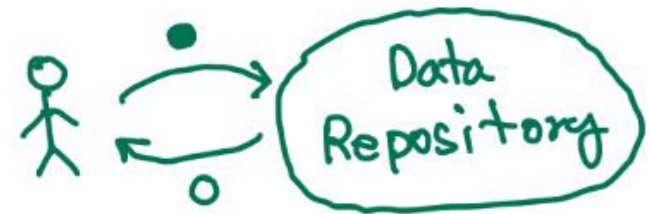
$\uparrow$   $j \geq 1$

## [주제 2]

# Heap-based Priority Queues and Heap Sort (Review)

# A Variety of Priority Queue Implementations

- [Priority Queue 1: Max(Min) Heap]
- [Priority Queue 2: Min-Max Heap]
- [Priority Queue 3: Heap and Hashing]
- [Priority Queue 4: Deap]
- [Priority Queue 5: Leftist Tree]
- [Priority Queue 6: Binomial Heap]
- [Priority Queue 7: Fibonacci Heap]

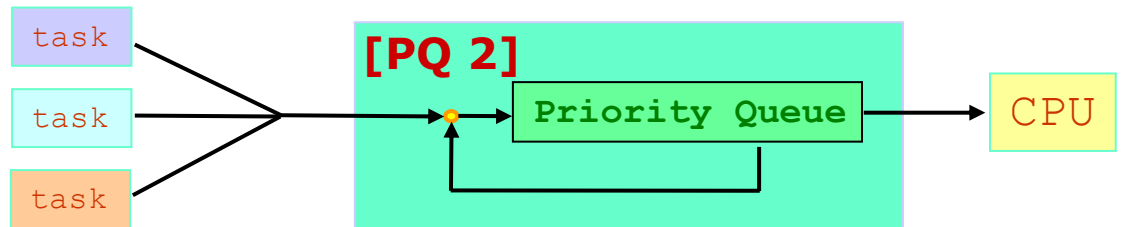
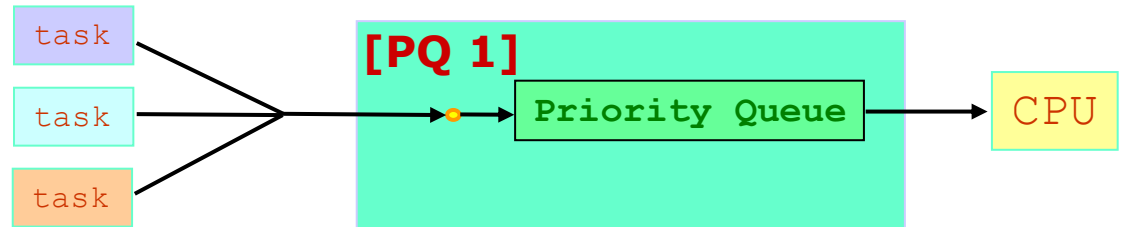
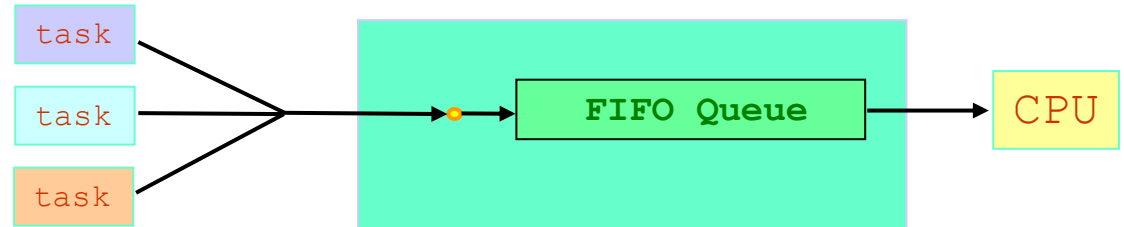


when retrieved,  
FIFO - based,  
LIFO - based,  
priority - based,  
⋮

# [Job Scheduling Example: Priority Queue]

- Consider the following sequence of requests in an operating system:

```
8134 kim 32
8137 kim 15
8138 jpark 24
8142 root 58
8244 ihm 45
8246 jpark 44
8251 root 34
8252 root 58
8255 sylee 21
8256 root 58
8261 jpark 58
8262 kim 32
8264 bhchoi 8
8267 root 58
8268 root 30
8269 root 58
8271 ihm 47
```

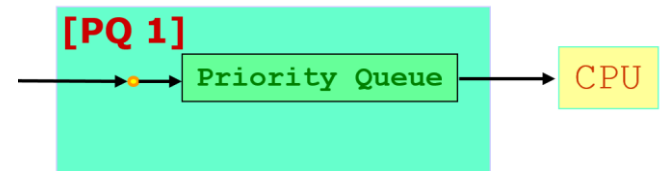


- **Requirement 1**

- CPU executes the process with **the highest priority** first.

☺ Use a heap structure – a simple max heap.

```
typedef struct _process {  
    int proc_id;  
    char *owner;  
    int priority;  
} Process;  
  
static Process *_proc_heap;  
static int _proc_heap_size = 0;  
static int _proc_heap_ptr = 0;  
  
int PH_create(int n);  
int PH_full();  
int PH_empty();  
int PH_insert(Process item);  
int PH_delete(Process *item);
```



- **Requirement 2**

- The priority of processes **can be modified** after they are placed in the priority queue.

```
int H_change_priority(int proc_id, int new_priority);
```

- ☹ This function requires locating a particular process in the heap, but the basic heap operations provide no efficient way to do it.
- ☺ Employ an auxiliary data structure such as a hash table that keeps track of the location of each process in the heap structure.

- Once the two requirements are satisfied, the operating system can process the following basic commands efficiently:

```
INSERT <proc_id> <owner> <priority>
```

```
DELETE
```

```
CHANGEPR <proc_id> <new_priority>
```

```
PRINTHEAP
```

```
END
```



# Priority Queue 1: Max(Min) Heap

[Horowitz 5.6.2]

[Neapolitan 7.6]

- **Problem**

- The following operations must be performed as mixed in data processing:
  - Store a record with a key in an arbitrary order.
  - Fetch the record with the current largest key.
- A solution: **Design a data structure** that offers an efficient implementation of the following operations:
  - **Insert an element with an arbitrary key.**
  - **Delete an element with the largest key.**

# An Array Implementation

[Sedgewick 9.2]

```
void PQinit();
int  PQempty();
void PQinsert(int);
int  PQdelmin();
void PQdec(int)
-----
#include <stdlib.h>
static int *pq;
static int N;
#define MAX_N 10000;
void PQinit() {
    pq = malloc(MAX_N*sizeof(int));
    N = 0;
}

int PQempty() {
    return N == 0;
}

void PQinsert(int v) {
    pq[N++] = v;
}
```

```
int PQdelmin() {
    int j, min = 0;
    for (j = 1; j < N; j++)
        if (less(pq[min], pq[j]))
            min = j;
    exch(pq[min], pq[N-1]);
    return pq[--N];
}

int less(int i, int j) {
    return ... ;
}

void exch(int i, int j) {
    ...
}

void PQdec(int k) {
    ...
}
```

- What will be the worst-case time complexity of each operation?

# Max(Min) Heap: Definitions

[Horowitz 5.6.2]

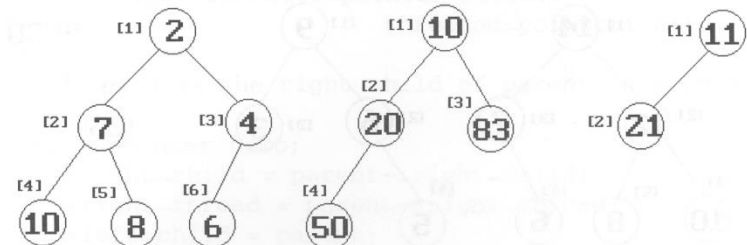
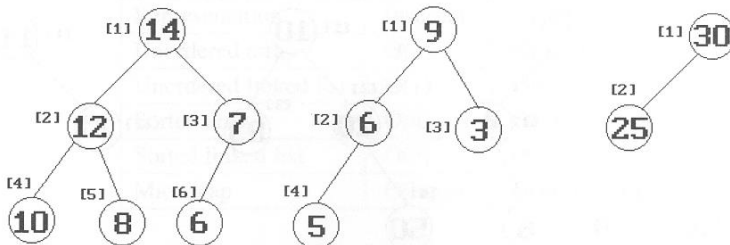
[Neapolitan 7.6]

- **Definition 1**

- A **max(min) heap** is a complete binary tree where the key value in each internal node is no **smaller(larger)** than the key values in its children.

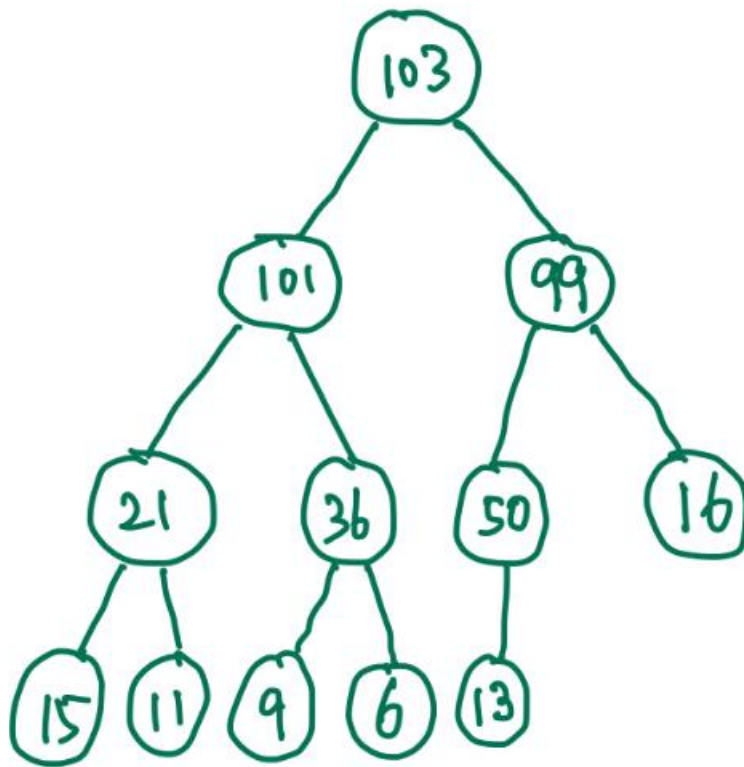
- **Definition 2**

- A binary tree has the **max(min) heap property** if and only if
  - ① The number of nodes of the tree is either 0 or 1, *or*
  - ② For the tree that has at least two nodes, the key in the root is no **smaller(larger)** than that in each child and the subtree rooted at the child has the **max(min) heap property**.
- A **max(min) heap** is a complete binary tree that has the **max(min) heap property**.



# Brainstorming on Max Heap Operations

## Max Heap Example



## Deletion Example 1

