Lecture 6: Smooth unconstrained minimization

- terminology
- general descent method
- gradient & steepest descent methods
- Newton's method
- quasi-Newton methods
- self-concordance & Newton's method

Terminology

unconstrained minimization problem

minimize f(x)

 $f: \mathbf{R}^n \to \mathbf{R}$, convex, differentiable (hence $\operatorname{\mathbf{dom}} f$ is open . . .)

minimizing sequence: $x^{(k)}$, $k \to \infty$

$$f(x^{(k)}) \to f^{\star}$$

optimality condition

$$\nabla f(x^{\star}) = 0$$

set of nonlinear equations, usually no analytical solution; more generally, if $\nabla^2 f(x) \succeq mI$, then

$$f(x) - f^* \le \frac{1}{2m} \|\nabla f(x)\|^2$$

 \dots yields stopping criterion (if you know m)

Examples

unconstrained quadratic minimization

$$\text{minimize } \boldsymbol{x}^T \boldsymbol{P} \boldsymbol{x} + 2\boldsymbol{q}^T \boldsymbol{x} + \boldsymbol{r}$$

$$(P = P^T \succeq 0)$$

unconstrained geometric programming

minimize
$$\log \sum_{i=1}^{m} e^{a_i^T x + b_i}$$

analytic center of linear inequalities

$$\mathsf{minimize} \ - \sum_i \log(b_i - a_i^T x)$$

$$(\mathbf{dom} \ f = \{x | a_i^T x < b_i, \ i = 1, \dots, m\})$$

Descent method

- 1. Compute a search direction v
- 2. Line search. Choose step size t>0
- 3. Update. x := x + tv

until stopping criterion is satisfied

Descent method: $f(x^{(k+1)}) < f(x^{(k)})$

Since f convex, v must be a **descent direction**: $\nabla f(x^{(k)})^T v^{(k)} < 0$

examples

- $\bullet \ v^{(k)} = -\nabla f(x^{(k)})$
- $v^{(k)} = -H^{(k)} \nabla f(x^{(k)}), H^{(k)} = H^{(k)T} > 0$
- $v^{(k)} = -\nabla^2 f(x^{(k)})^{-1} \nabla f(x^{(k)})$

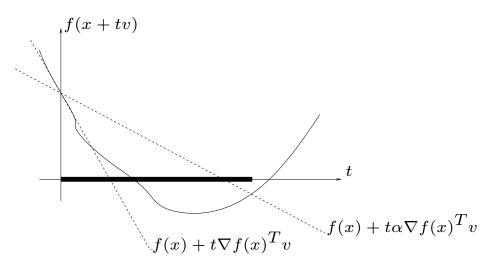
Line search types

two simple & effective line search types:

exact line search: $t = \operatorname{argmin}_{s>0} f(x + sv)$

backtracking line search ($0 < \beta < 1$, $0 < \alpha < 0.5$)

- starting with t = 1, $t := \beta t$
- until $f(x + tv) \le f(x) + t\alpha \nabla f(x)^T v$



Gradient method

- 1. Compute search direction $v = -\nabla f(x)$
- 2. Line search. Choose step size t
- 3. Update. x := x + tv

until stopping criterion is satisfied

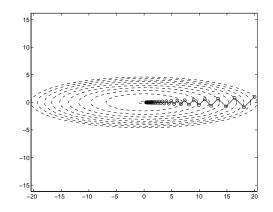
- converges with exact or bactracking line search
- can be very slow
- rarely used in practice

Example

$$\text{minimize } \frac{1}{2}(x_1^2 + Mx_2^2)$$

where M>0; optimal point is $x^{\star}=0$

- use exact line search
- start at $x^{(0)} = (M, 1)$ (to simplify formulas)



iterates are then

$$x^{(k)} = \left(M\left(\frac{M-1}{M+1}\right)^k, \left(-\frac{M-1}{M+1}\right)^k\right)$$

convergence is

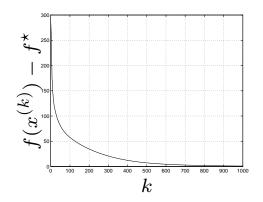
- fast if M close to 1
- ullet slow, zig-zagging if $M\gg 1$ or $M\ll 1$

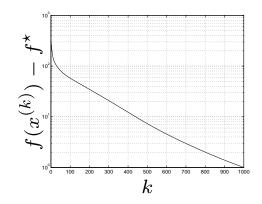
Numerical example: gradient method

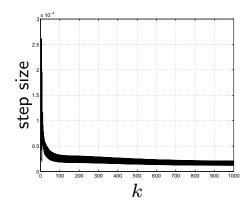
$$\text{minimize } c^T x - \sum_{i=1}^m \log(a_i^T x + b_i)$$

$$m = 100, n = 50$$

gradient method with exact line search







slow convergence; zig-zagging

Steepest descent direction

first-order approximation of f at x:

$$f(x+z) \approx f(x) + \nabla f(x)^T z$$

 $\nabla f(x)^T z$ gives approximate decrease in f for (small) step z

steepest descent direction for general norm $\|\cdot\|$:

$$v_{\mathrm{sd}} = \operatorname{argmin}\{\nabla f(x)^T v \mid ||v|| = 1\}$$

gives greatest (approximate) decrease in f, per length of step (measured by $\|\cdot\|$)

Euclidean norm: $v_{\rm sd} = -\nabla f(x) / \|\nabla f(x)\|$

quadratic norm: $||z||_P = \left(z^T P z\right)^{1/2}$, $P = P^T \succ 0$

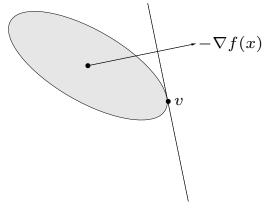
$$v_{\rm sd} = -\left(\nabla f(x)^T P^{-1} \nabla f(x)\right)^{-1/2} P^{-1} \nabla f(x)$$

can express $v_{
m sd}$ as

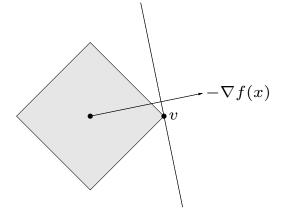
$$v_{\mathrm{sd}} = \operatorname{argmin} \{ \nabla f(x)^T v \mid ||v|| \le 1 \}$$

geometric interpretation: go as far as possible in direction $-\nabla f(x)$, while staying in unit ball

quadratic norm:



 ℓ_1 -norm:



Steepest descent method

- 1. Compute steepest descent direction $v_{\rm sd} = \operatorname{argmin}\{\nabla f(x)^T v \mid ||v|| = 1\}$
- 2. Line search. Choose a step size t
- 3. Update. x := x + tv

until stopping criterion is satisfied

- converges with exact or backtracking line search
- ullet sometimes $v_{
 m sd}$ is scaled between 1 and 2
- can be very slow
- ullet used in special cases where v is cheap to compute

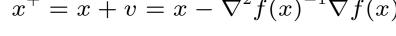
The Newton step

the Newton step (at x) is

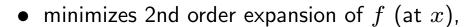
$$v = -\nabla^2 f(x)^{-1} \nabla f(x)$$

the Newton iteration (at x) is

$$x^{+} = x + v = x - \nabla^{2} f(x)^{-1} \nabla f(x)$$



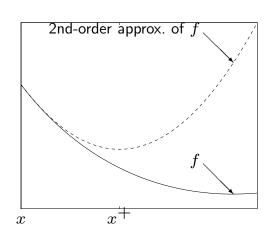


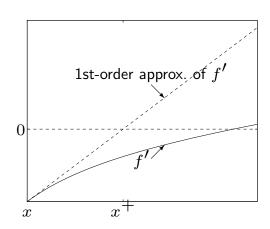


$$f(x) + \nabla f(x)^T (y-x) + \frac{1}{2} (y-x)^T \nabla^2 f(x) (y-x)$$

solves linearized optimality condition:

$$0 = \nabla f(x) + \nabla^2 f(x)(y - x)$$





Local convergence of Newton iteration

assumptions: $\nabla^2 f(x) \succeq mI$ and Hessian satisfies Lipschitz condition:

$$\left\| \nabla^2 f(x) - \nabla^2 f(y) \right\| \le L \|x - y\|$$

(L small means f nearly quadratic)

result

$$\frac{L}{2m^2} \left\| \nabla f(x^+) \right\| \le \left(\frac{L}{2m^2} \left\| \nabla f(x) \right\| \right)^2$$

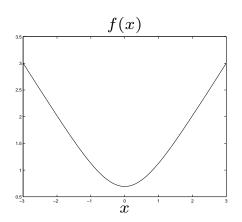
ullet region of **quadratic convergence**: $\|\nabla f(x)\|$ (hence, $f(x)-f^\star$) decreases very rapidly if

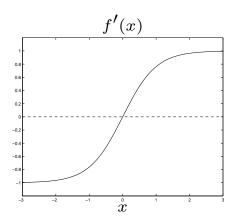
$$\left\| \nabla f(x^{(0)}) \right\| < m^2 / L$$

- bound on #iterations for accuracy $f(x) f^* \le \epsilon$: $\log_2 \log_2(\epsilon_0/\epsilon)$, $\epsilon_0 = m^3/L^2$
- practical rule of thumb: 5-6 iterations

Global behavior of Newton iteration

Newton iteration can diverge **example:** $f(x) = \log(e^x + e^{-x})$, start at $x^{(0)} = 1.1$





k	$x^{(k)}$	$f(x^{(k)}) - f^{\star}$
1	$-1.129 \cdot 10^{0}$	$5.120 \cdot 10^{-1}$
2	$1.234\cdot 10^0$	$5.349 \cdot 10^{-1}$
3	$-1.695\cdot10^{0}$	$6.223 \cdot 10^{-1}$
4	$5.715\cdot 10^0$	$1.035\cdot 10^0$
5	$-2.302\cdot10^4$	$2.302\cdot 10^4$

Newton's method

given starting point $x \in \operatorname{dom} f$ repeat

- 1. Compute Newton direction $v = -\nabla^2 f(x)^{-1} \nabla f(x)$
- 2. Line search. Choose a step size t
- 3. Update. x := x + tv

until stopping criterion is satisfied

(also called damped or guarded Newton method)

- global convergence with backtracking or exact line search
- quadratic local convergence (hence, stopping criterion not an issue)

Affine invariance of Newton method

- use new coords $x = T\bar{x}$, $\det T \neq 0$
- $\bullet \;$ apply Newton to $g(\bar{x}) = f(T\bar{x})$
- then $x^{(k)} = T\bar{x}^{(k)}$

e.g., Newton method not affected by variable scaling (cf. gradient, steepest descent)

Convergence analysis

assumptions: $mI \preceq \nabla^2 f(x) \preceq MI$ and Lipschitz condition

$$\left\| \nabla^2 f(x) - \nabla^2 f(y) \right\| \le L \|x - y\|$$

results: two phases

1. damped Newton phase: $\|\nabla f(x)\| \ge \eta_1$: $f(x^+) \le f(x) - \eta_2$, hence

$$\# \mathsf{iterations} \leq \eta_2^{-1}(f(x^{(0)}) - f^\star)$$

2. quadratically convergent phase: $\|\nabla f(x)\| < \eta_1$

#iterations
$$\leq \log_2 \log_2(\epsilon_0/\epsilon)$$

total #iterations bounded by

$$\eta_2^{-1}(f(x^{(0)}) - f^*) + \log_2 \log_2(\epsilon_0/\epsilon)$$

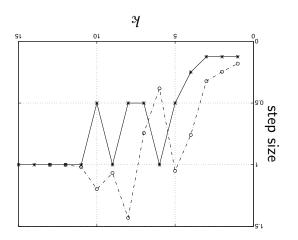
 η_1 , η_2 , ϵ_0 depend on m, M, L (and α , β for backtracking)

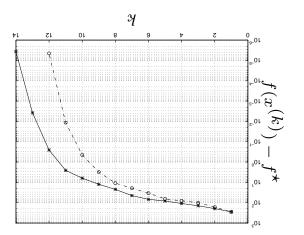
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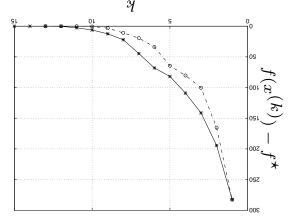
Numerical example: Newton method

minimize
$$c^T x - \sum_{i=i}^m \log(a_i^T x + b_i)$$

 $0\delta = n ,001 = m$







solid line: backtracking $(\beta=0.5,\,\alpha=0.2)$; dashed line: exact line search; (iters are more expensive than gradient method)

Quasi-Newton methods

idea: replace $\nabla^2 f(x)$ by approximation H

given starting point $x \in \operatorname{dom} f$, $H \succ 0$ repeat

1. Compute quasi-Newton direction.

$$v = -H^{-1}\nabla f(x)$$

- $v=-H^{-1}
 abla f(x)$ 2. Line search. Choose a step size t3. Update H.
 4. Update x. x:=x+tv

until stopping criterion is satisfied

many update rules $H \to H^+$, which all satisfy:

- $H = H^T \succ 0$
- secant condition: $\nabla f(x^+) \nabla f(x) = H^+(x^+ x)$
- $H^{-1}\nabla f(x)$ more easily computed (than $\nabla^2 f(x)^{-1}\nabla f(x)$)

advantages (compared to Newton method)

- don't need to evaluate $\nabla^2 f(x)$
- ullet v can be computed in $O(n^2)$ operations

disadvantage (compared to Newton method)

• local convergence fast, but not quadratic

quasi-Newton methods

- \bullet converge in n steps for (cvx) quadratic f on \mathbf{R}^n
- widely used in practice

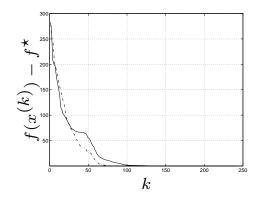
most common is Broyden-Fletcher-Goldfarb-Shanno (BFGS):

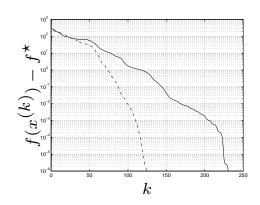
$$y = \nabla f(x^{+}) - \nabla f(x), \quad s = x^{+} - x, \quad H^{+} = H + \frac{yy^{T}}{y^{T}s} - \frac{Hss^{T}H}{s^{T}Hs}$$

Numerical example: BFGS method

$$\text{minimize } c^Tx - \sum_{i=1}^m \log(a_i^Tx + b_i)$$

m = 100, n = 50





solid line: backtracking ($\beta=0.5,~\alpha=0.2$); dashed line: exact line search when comparing with Newton method, remember:

- ullet BFGS is $O(n^2)$ per iter (plus finding ∇f)
- Newton is $O(n^3)$ per iter (plus finding ∇f , $\nabla^2 f$)

Self-concordance: motivation

drawbacks of classical analysis of Newton's method:

- Newton's method is affinely invariant, but convergence analysis is not
- ullet never know m, M, L in practice
- \bullet m, M, L can depend on starting point

Nesterov & Nemirovsky's analysis of Newton's method

- is affinely invariant
- involves no unknown constants
- is valid for many (but not all) functions f

Self-concordance: definition

(Nesterov & Nemirovsky)

ullet function $f: \mathbf{R} \to \mathbf{R}$ is self-concordant if

$$|f'''(t)| \le 2f''(t)^{3/2}$$

 \bullet $f: \mathbf{R}^n \to \mathbf{R}$ is self-concordant if restriction to arbitrary line is self-concordant

SC condition

- limits third derivative in terms of second
- ullet is affinely invariant: f is SC $\Longleftrightarrow f(Tx)$ is SC

examples of self-concordant fcts

• linear, (convex) quadratic functions

- $\bullet \log x \text{ on } \{x | x > 0\}$
- $-\log \det X$ on $\{X|X=X^T\succ 0\}$
- $-\log(t^2 x^T x)$ on $\{(x, t) \mid ||x|| < t\}$

simple properties:

ullet affine transformation of domain: $f \ \mathsf{SC} \Longrightarrow g(z) = f(Az+b) \ \mathsf{SC}$

 \bullet sums: f, \tilde{f} SC $\Longrightarrow f + \tilde{f}$ SC

• scaling: f SC, $\alpha \geq 1 \Longrightarrow \alpha f$ SC

hence, e.g.,

- $\bullet \ -\sum_i \log(b_i a_i^T x)$ is SC
- $-\log \det (F_0 + x_1F_1 + \cdots + x_nF_n)$ is SC

Convergence analysis via SC

Newton method, with backtracking or exact line search:

$$\# \text{iterations} \leq \frac{f(x^{(0)}) - f^\star}{\eta_2} + \log_2 \log_2(2/\epsilon)$$

where η_2 depends only on backtracking parameters:

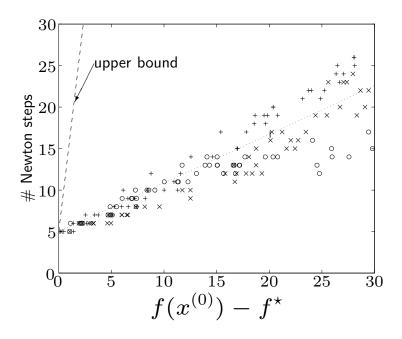
$$\eta_2 = \beta \frac{\alpha (1/2 - \alpha)^2}{5 - 2\alpha}$$

e.g., for $\alpha=0.2$, $\beta=0.7$, we have $1/\eta_2 pprox 365$

(a more refined analysis yields smaller bound)

example:

$$f(x) = \log \det (F_0 + x_1 F_1 + \dots + x_n F_n)^{-1}$$



conclusion:

- ullet $f(x^{(0)}) f^{\star}$ gives upper bound on #iterations
- $f(x^{(0)}) f^*$ is also good measure in practice