MAT 226B Large Scale Matrix Computation Final Project

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Problem 1:

(a) We know from nonsymmetric Lanczos process that

$$MV_k = V_k T_k + \beta_{k+1} [0 \dots 0 v_{k+1}]$$

We can multiply the above by e_1 to extract the first column (v_1) from V_k before multiplying it by M and the result is

$$MV_k e_1 = V_k T_k e_1 + \beta_{k+1} [0 \dots 0 v_{k+1}] e_1$$

 $MV_k e_1 = V_k T_k e_1 + 0$

Note that $MV_ke_1=Mv_1$. Now, we can easily give the proof as

$$M^j r = M^j(\beta_1 v_1) = \beta_1 M^j v_1$$

$$M^{j}r = \beta_{1}V_{k}T_{k}^{j}e_{1}, \quad \forall j = 0, 1, \dots, k-1$$
 (1)

For the second part, we note that $e_k^T T_k^{k-1} e_1 = 0$. Thus, the second sum has no effect. We can let j = k in 1 and we get

$$M^{k}r = \beta_{1}V_{k}T_{k}^{k}e_{1} + \beta_{1}\beta_{k+1}(e_{k}^{T}T_{k}^{k-1}e_{1})v_{k+1}$$

(b) We follow the same steps as in (a). First we have

$$M^{T}W_{k} = W_{k}\hat{T}_{k} + \gamma_{k+1}[0\dots 0w_{k+1}]$$

$$M^{T}W_{k}e_{1} = W_{k}\hat{T}_{k}e_{1} + \gamma_{k+1}[0\dots 0w_{k+1}]e_{1}$$

$$M^{T}w_{1} = W_{k}\hat{T}_{k}e_{1} + 0$$

Taking the transpose of the above, we get

$$(M^T w_1)^T = w_1^T M = e_1^T \hat{T}_k^T W_k^T$$

We also know that $\hat{T_k}^T = D_k T_k D_k^{-1}$. Thus,

$$w_1^T M = e_1^T D_k T_k D_k^{-1} W_k^T$$

Now, we can give the proof as

$$c^{T}M^{j} = \gamma_{1}w_{1}^{T}M^{j}$$

$$c^{T}M^{j} = \gamma_{1}e_{1}^{T}D_{k}T_{k}^{j}D_{k}^{-1}W_{k}^{T}$$

$$c^{T}M^{j} = \gamma_{1}\delta_{1}e_{1}^{T}T_{k}^{j}D_{k}^{-1}W_{k}^{T}$$

(c) We can write the Z(s) as

$$Z(s) = \sum_{j=0}^{\infty} \sigma^j c^T M^j r \tag{2}$$

We can find two values positive j_1 and j_2 such that $j_1 + j_2 = j$. Then, we can write 2 as

$$Z(s) = \sum_{j=0}^{\infty} \sigma^{j} c^{T} M^{j_{1}} M^{j_{2}} r$$

$$Z(s) = \sum_{j=0}^{\infty} \sigma^{j} (\gamma_{1} \delta_{1} e_{1}^{T} T_{k}^{j_{1}} D_{k}^{-1} W_{k}^{T}) (\beta_{1} V_{k} T_{k}^{j_{2}} e_{1})$$

$$Z(s) = \sum_{j=0}^{\infty} \sigma^{j} \gamma_{1} \delta_{1} \beta_{1} e_{1}^{T} T_{k}^{j_{1}} D_{k}^{-1} W_{k}^{T} V_{k} T_{k}^{j_{2}} e_{1}$$
(3)

We know from Lanczos process that $W_k V_k = D_k$. In addition, we have $c^T r = (\gamma_1 w_1)^T (\beta_1 v_1) = \gamma_1 \beta_1 w_1^T v_1 = \gamma_1 \delta_1 \beta_1$. We can plug this relations in 3 to get

$$Z(s) = \sum_{j=0}^{\infty} \sigma^{j}(c^{T}r)e_{1}^{T}T_{k}^{j_{1}}D_{k}^{-1}D_{k}T_{k}^{j_{2}}e_{1}$$

$$Z(s) = \sum_{j=0}^{\infty} \sigma^{j}(c^{T}r)e_{1}^{T}T_{k}^{j_{1}}T_{k}^{j_{2}}e_{1} = \sum_{j=0}^{\infty} \sigma^{j}(c^{T}r)e_{1}^{T}T_{k}^{j}e_{1}$$

Problem 2:

Here we are required to find an efficient way to compute q = Mv and $q = M^Tv$ for $v \in \mathbb{C}^n$ where $M = (A - s_0 E)^{-1} E$. We can compute the matrix-vector multiplication efficiently using LU factorization. We first can write the multiplication as

$$q = (A - s_0 E)^{-1} E v = \underbrace{(A - s_0 E)^{-1}}_{W} \underbrace{E v}_{f}$$

$$q = W^{-1} f \quad \Rightarrow \quad W q = f \quad \Rightarrow \quad \underbrace{P D^{-1} W Q}_{LU} \underbrace{Q^T q}_{d} = P D^{-1} f$$

Thus, we can fist solve $Lc = PD^{-1}f$ for $c \in \mathbb{C}^n$ via forward substitution, then solve Ud = c for $d \in \mathbb{C}^n$ via backward substitution, and finally set q = Qd.

We can use the same LU factorization to compute $q=M^Tv$ efficiently. We first not that transposing the LU factorization for a given matrix W is $U^TL^T=Q^TW^TD^{-T}P^T$ We can write this multiplication as

$$q = ((A - s_0 E)^{-1} E)^T v = E^T \underbrace{(A - s_0 E)^{-T} v}_{g}$$

$$q = W^{-T} v \quad \Rightarrow \quad W^T g = v \quad \Rightarrow \quad \underbrace{Q^T W^T D^{-T} P^T}_{U^T L^T} \underbrace{(D^{-T} P^T)^{-1} g}_{d} = Q^T v$$

Thus, we can first solve $U^Tc=Q^Tv$ for c via forward substitution, then solve $L^Td=c$ for d via backward substitution, and then set $g=D^{-T}P^Td$. Finally, we multiply g from the left by E^T to get q. The functions Mv and transposeMv implements these operations as discussed.

Problem 3:

The leading 2k moments $\mu_j = c^T M^j r$ for $j = 0, 1, \dots, 2k-1$ can be computing as follows. Let $f_j = M^j r$. It is easy to see that $f_j = M f_{j-1}$ from which we can compute the moment at j as $\mu_j = c^T f_j$ and compute f_j recursively. We can use the same LU factorization to compute r and used the function Mv to compute f_j . The function compute Moments compute the moments as discussed here.

We wrote another function textbookAlgo that utilizes computeMoments to implement the textbook algorithm for computing $Z_k(s)$. More precisely, it compute the coefficient of the polynomials $p(\sigma)$ and $q(\sigma)$ such that $Z_k(s) = \frac{p(\sigma)}{q(\sigma)}$ where $p(\sigma) = \alpha_0 + \alpha_1 \sigma + \cdots + \alpha_{k-1} \sigma^{k-1}$, $q(\sigma) = \beta_0 + \beta_1 \sigma + \cdots + \beta_k \sigma^k$, $\alpha_0, \ldots, \alpha_{k-1}, \beta_1 \ldots \beta_k \in \mathbb{C}$, and $\beta_0 = 1$. The output of this function is two vectors α and β containing the coefficients.

Problem 4:

We wrote the function zkViaLanczos which computes Z_k given T_k , s and s_0 . T_k is computed from our previous implementation of the nonsymmetric Lanczos in Homework 3 which feed in with the efficient implementation of the Mv and M^Tv from Problem 2.

System Specs: All our experiments run on Intel(R) Xeon(R) CPU E3-1280 v5 with 3.70 GHz and 32 GB of RAM on 64-bit operating system running Windows 7.

 s_0 with fast convergence: We test our implementation of the textbook and Lanczos-based algorithm for different values of s_0 and found that it runs fairly fast for the small input given in FP_Ex1.mat; it takes less than a second even for large i.e., k < 100.

We followed the recommendation given in the lectures for how to pick s_0 . We choose $s_0 = 1e5 + 2\pi i 5.5e8$.

Comparison: We run both our implementation for different values of k and the above s_0 and compared between both. Table 1 show the average different ($\|\cdot\|^2$) and the maximum (absolute) different between the two vectors containing the output of both algorithms for different values of s. Function Table_1 () in driver.m file generates these data. We can see that when k>13, the two algorithms will give difference numerical results.

Explanation: We believe the reason why the textbook algorithm does not perform well is because it depends on computing $M^j r$ (to compute the moments) for increasing values of j which convergences quickly to the eigenvector of M with largest eigenvalue. Thus, the information it contains comes from a single eigenvector where the information should comes from all eigenvectors of M. In contrast, Lanczos's T_k represents oblique projection of M onto the $K_k(M,r)$ Krylov subspace which contains information about k eigenvectors.

k	Average Difference	Maximum Difference	
2	2.082747e - 28	$2.109424e{-15}$	
3	7.919421e - 27	$6.439294e{-15}$	
4	5.026520e - 25	$4.618528e{-14}$	
5	8.547583e - 26	$2.664535e{-14}$	
6	$3.052289e{-23}$	$4.112266e{-13}$	
7	$1.148403e{-19}$	$4.235057e{-11}$	
8	2.656897e - 17	3.190033e-09	
9	$6.690131e{-15}$	1.812676e - 08	
10	6.292776e - 12	3.008865e - 07	
11	7.619719e - 11	$6.315981e{-06}$	
12	4.360106e - 04	1.650155e - 02	
13	9.107709e - 04	1.494378e - 02	
14	7.052574e + 00	3.945073e - 01	
15	5.904627e+01	1.279984e+00	
16	1.779754e + 02	1.488481e+00	
17	1.105689e + 02	1.625567e+00	
18	1.100795e+02	1.632208e+00	
19	1.219843e+02	1.688382e+00	
20	1.075413e + 02	9.725151e - 01	
21	1.108709e+02	1.217680e+00	
22	4.712087e + 02	1.762950e+00	
23	1.160591e+03	2.593543e+00	
24	3.507356e + 03	4.102138e+00	
25	7.329648e + 03	5.568756e + 00	
26	2.245490e+04	8.889336e+00	
27	2.633420e+04	9.681861e+00	
28	4.881768e+04	1.252447e+01	
29	7.066637e + 04	1.471937e+01	
30	1.151561e + 05	1.801101e+01	

Table 1: Average and maximum (absolute) difference between the results of the textbook algorithm and Lanczos approach for different k values.

s_0	k	Time	Average Difference
 			J

Table 2: Lanczos approach using different k and s_0 values and comparing it with the exact solution.

Lanczos approach with difference s_0 :

Problem 5:

(a)