

Boolean combinations of cellular complexes as chain operations *

Alberto Paoluzzi

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Contents

1	Introduction	2
1.1	Preview of the Boolean algorithm	2
1.2	Remarks	2
2	Step 1: merging discrete spaces	3
2.1	Requirements	3
2.2	Implementation	4
3	Step 2: splitting cells	6
3.1	Requirements	6
3.2	Implementation	7
4	Step 3: cell labeling	12
4.1	Requirements	12
5	Step 4: greedy cell gathering	12
6	Exporting the library	12
7	Tests and examples	13
A	Appendix: utility functions	18
A.1	Numeric utilities	18

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1 Introduction

In this module a novel approach to Boolean operations of cellular complexes is defined and implemented. The novel algorithm may be summarised as follows.

First we compute the CDC (Common Delaunay Complex) of the input LAR complexes A and B , to get a LAR of the *simplicial* CDC.

Then, we split the cells intersecting the boundary faces of the input complexes, getting the final *polytopal* SCDC (Split Common Delaunay Complex), whose cells provide the basis for the linear coordinate representation of both input complexes, upon the same space decomposition.

Afterwards, every Boolean result is computed by bitwise operations, between the coordinate representations of the transformed A and B input.

Finally a greedy assembly of SCDC cells is executed, in order to return a polytopal complex with a reduced number of cells.

1.1 Preview of the Boolean algorithm

The goal is the computation of $A \diamond B$, with $\diamond \in \{\cup, \cap, -\}$, where a LAR representation of both A and B is given. The Boolean algorithm works as follows.

1. Embed both cellular complexes A and B in the same space (say, identify their common vertices) by $V_{ab} = V_a \cup V_b$.
2. Build their CDC (Common Delaunay Complex) as the LAR of *Delaunay triangulation* of the vertex set V_{ab} , and embedded ∂A and ∂B in it.
3. Split the (highest-dimensional) cells of CDC crossed by ∂A or ∂B . Their lower dimensional faces remain partitioned accordingly. We name the resulting complex SCDC (Split Common Delaunay Complex).
4. With respect to the SCDC basis of d -cells C_d , compute two coordinate chains $\alpha, \beta : C_d \rightarrow \{0, 1\}$, such that:

$$\begin{aligned}\alpha(cell) &= 1 & \text{if } |cell| \subset A; & \quad \text{else } \alpha(cell) = 0, \\ \beta(cell) &= 1 & \text{if } |cell| \subset B; & \quad \text{else } \beta(cell) = 0.\end{aligned}$$

5. Extract accordingly the SCDC chain corresponding to $A \diamond B$, with $\diamond \in \{\cup, \cap, -\}$.

1.2 Remarks

You may make an analogy between the SCDC (*Split* CDC) and a CDT (Constrained Delaunay Triangulation). In part they coincide, but in general, the SCDC is a polytopal complex, and is not a simplicial complex as the CDC.

The more complex algorithmic step is the cell splitting. Every time, a single d -cell c is split by a single hyperplane (cutting its interior) giving either two splitted cells c_1 and c_2 , or just one output cell (if the hyperplane is the affine hull of the CDC facet) whatever the input cell dimension d . After every splitting of the cell interior, the row c is substituted (within the CV matrix) by c_1 , and c_2 is added to the end of the CV matrix, as a new row.

The splitting process is started by “splitting seeds” generated by $(d - 1)$ -faces of both operand boundaries. In fact, every such face, say f , has vertices on CDC and *may* split some incident CDC d -cell. In particular, starting from its vertices, f must split the CDC cells in whose interior it passes through.

So, a dynamic data structure is set-up, storing for each boundary face f the list of cells it must cut, and, for every CDC d -cell with interior traversed by some such f , the list of cutting faces. This data structure is continuously updated during the splitting process, using the adjacent cells of the split ones, who are to be split in turn. Every split cell may add some adjacent cell to be split, and after the split, the used pair (**cell,face**) is removed. The splitting process continues until the data structure becomes empty.

Every time a cell is split, it is characterized as either internal (1) or external (0) to the used (oriented) boundary facet f , so that the two resulting subcells c_1 and c_2 receive two opposite characterization (with respect to the considered boundary).

At the very end, every (polytopal) SCDC d -cell has two bits of information (one for argument A and one for argument B), telling whether it is internal (1) or external (0) or unknown (-1) with respect to every Boolean argument.

A final recursive traversal of the SCDC, based on cell adjacencies, transforms every -1 into either 0 or 1, providing the two final chains to be bitwise operated, depending on the Boolean operation to execute.

2 Step 1: merging discrete spaces

2.1 Requirements

The *join* of two sets $P, Q \subset \mathbb{E}^d$ is the set $PQ = \{\alpha\mathbf{x} + \beta\mathbf{y} \mid \mathbf{x} \in P, \mathbf{y} \in Q\}$, where $\alpha, \beta \in \mathbb{R}$, $\alpha, \beta \geq 0$, and $\alpha + \beta = 1$. The join operation is associative and commutative.

Input Two LAR models of two non-empty “solid” d -spaces A and B , denoted as (**V1,CV1**) and (**V2,CV2**).

Output The LAR representation (**V,CV**) of Delaunay triangulation (simplicial d -complex) of the set $\text{conv } AB \subset \mathbb{E}^d$, convex hull of the join of A and B , named Common Delaunay Complex (CDC) in the following.

Auxiliary data structures This software module returns also:

1. a dictionary **vertDict** of V vertices, with *key* the symbolic representation of vertices v returned by expressions $vcode(v)$, $v \in V$, and with values the finite ordinal numbers of the vertices;
2. the numbers $n1$, $n12$, $n2$ of the elements of $V1$, $V1 \cap V2$, and $V2$, respectively. Notice that the following assertions must hold (see Figure 1):

$$n1 - n12 + n2 = n \quad (1)$$

$$0 < n - n2 \leq n1 < n \quad (2)$$

3. the input boundary complex (V, BC) , with $BC = BC1 + BC2$, i.e. the union of the two boundary $(d-1)$ -complexes $(V, BC1)$ and $(V, BC2)$, defined on the common vertices.

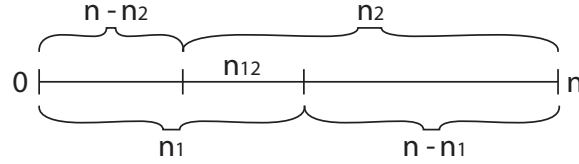


Figure 1: Relationships inside the orderings of CDC vertices

2.2 Implementation

⟨ Compute model boundaries of complex of convex cells 4a ⟩ ≡

```
""" Compute model boundaries of complex of convex cells """
```

```
def larFacetsOfPolytopalComplex(vertDict, cells, facets):
    (V1, CV1), (V2, CV2) = model1, model2
    for cell in CV1:
        Vcell = [V1[v] for v in cell]
```

◇

Macro never referenced.

⟨ Merge two dictionaries with keys the point locations 4b ⟩ ≡

```
""" Merge two dictionaries with keys the point locations """
```

```
def mergeVertices(model1, model2):

    (V1, CV1), (V2, CV2) = model1, model2

    n = len(V1); m = len(V2)
    def shift(CV, n):
```

```

    return [[v+n for v in cell] for cell in CV]
CV2 = shift(CV2,n)

vdict1 = defaultdict(list)
for k,v in enumerate(V1): vdict1[vcode(v)].append(k)
vdict2 = defaultdict(list)
for k,v in enumerate(V2): vdict2[vcode(v)].append(k+n)
vertDict = defaultdict(list)
for point in vdict1.keys(): vertDict[point] += vdict1[point]
for point in vdict2.keys(): vertDict[point] += vdict2[point]

case1, case12, case2 = [],[],[]
for item in vertDict.items():
    key,val = item
    if len(val)==2: case12 += [item]
    elif val[0] < n: case1 += [item]
    else: case2 += [item]
n1 = len(case1); n2 = len(case12); n3 = len(case2)

invertedindex = list(0 for k in range(n+m))
for k,item in enumerate(case1):
    invertedindex[item[1][0]] = k
for k,item in enumerate(case12):
    invertedindex[item[1][0]] = k+n1
    invertedindex[item[1][1]] = k+n1
for k,item in enumerate(case2):
    invertedindex[item[1][0]] = k+n1+n2

V = [eval(p[0]) for p in case1] + [eval(p[0]) for p in case12] + [eval(
    p[0]) for p in case2]
CV1 = [sorted([invertedindex[v] for v in cell]) for cell in CV1]
CV2 = [sorted([invertedindex[v] for v in cell]) for cell in CV2]

return V,CV1,CV2, n1+n2,n2,n2+n3

```

◇

Macro referenced in [12a](#).

⟨ Make Common Delaunay Complex 5 ⟩ ≡

```

""" Make Common Delaunay Complex """
def makeCDC(arg1,arg2):

    (V1,basis1), (V2,basis2) = arg1,arg2
    (facets1,cells1),(facets2,cells2) = basis1[-2:],basis2[-2:]
    model1, model2 = (V1,cells1),(V2,cells2)

    V, _,_, n1,n12,n2 = mergeVertices(model1, model2)

```

```

n = len(V)
assert n == n1 - n12 + n2

CV = sorted(AA(sorted)(Delaunay(array(V)).simplices))
vertDict = defaultdict(list)
for k,v in enumerate(V): vertDict[vcode(v)] += [k]

BC1 = signedCellularBoundaryCells(V1,basis1)
BC2 = signedCellularBoundaryCells(V2,basis2)
BC = [[ vertDict[vcode(V1[v])][0] for v in cell] for cell in BC1] + [
      [ vertDict[vcode(V2[v])][0] for v in cell] for cell in BC2]

return V,CV,vertDict,n1,n12,n2,BC

```

◇

Macro referenced in [12a](#).

3 Step 2: splitting cells

The goal of this section is to transform the CDC simplicial complex, into the polytopal Split Common Delaunay Complex (SCDC), by splitting the d -cells of CDC crossed in their interior by some cell of the input boundary complex.

3.1 Requirements

We call here for a sequential implementation, following every $(d - 1)$ -facet `lambda` in `BC` (for *Boundary Cells*). We start the splitting with `COVECTOR(lambda)` from `cell`, one of the CDC d -cells incident on a vertex of `lambda`, and continue the splitting on the d -cells $(d - 1)$ -adjacent to `cell`, where (a) `COVECTOR(lambda)` either crosses the `cell`'s interior or contains one of `cell`'s $(d - 1)$ -facets and (b) such that the intersection with `lambda` is not empty, until the queue (or stack) of d -cells to intersect with `covector` is not empty.

Best computational strategy First associate to each cutting facet the list of cells it may cut; then execute all the cuts. In this way we can compute the adjacency matrix just one time at the beginning of the procedure, and do not need to update it after every split.

Input The output of previous algorithm stage.

Output The LAR representation (W, PW) of the SCDC,

Auxiliary data structures This software module returns also a dictionary `splitFacets`, with keys the input boundary faces and values the list of pairs `(covector, fragmentedFaces)`.

3.2 Implementation

Computing the adjacent cells of a given cell To perform this task we make only use of the **CV** list. In a more efficient implementation we should make direct use of the sparse adjacency matrix, to be dynamically updated together with the **CV** list. The computation of the adjacent d -cells of a single d -cell is given here by extracting a column of the $\text{CSR}(M_d M_d^t)$. This can be done by multiplying $\text{CSR}(M_d)$ by its transposed row corresponding to the query d -cell.

\langle Computing the adjacent cells of a given cell 6 $\rangle \equiv$

```

""" Computing the adjacent cells of a given cell """
def adjacencyQuery (V,CV):
    dim = len(V[0])
    def adjacencyQuery0 (cell):
        nverts = len(CV[cell])
        csrCV = csrCreate(CV)
        csrAdj = matrixProduct(csrCV,csrTranspose(csrCV))
        cellAdjacencies = csrAdj.indices[csrAdj.indptr[cell]:csrAdj.indptr[cell+1]]
        return [acell for acell in cellAdjacencies if dim <= csrAdj[cell,acell] < nverts]
    return adjacencyQuery0

```

◇

Macro referenced in 12a.

Updating the adjacency matrix At every step of the CDC splitting, generating two output cells **cell11** and **cell12** from the input **cell**, the element of such index in the list **CV** is restored with the **cell11** vertices, and a new (last) element is created in **CV**, to store the **cell12** vertices. Therefore the row of index **cell** of the symmetric adjacency matrix must be recomputed, being the **cell** column updated consequently. Also, a new last row (and column) must be added to the matrix.

\langle Updating the adjacency matrix 7a $\rangle \equiv$

```

""" Updating the adjacency matrix """
pass

```

◇

Macro never referenced.

Relational inversion (characteristic matrix transposition) The operation could be executed by simple matrix transposition of the CSR (Compressed Sparse Row) representation of the sparse characteristic matrix $M_d \equiv \text{CV}$. A simple relational inversion using Python lists is given here. The **invertRelation** function is given here, linear in the size of the **CV** list, where the complexity of each cell is constant and small in most cases.

\langle Characteristic matrix transposition 7b $\rangle \equiv$

```

""" Characteristic matrix transposition """
def invertRelation(CV):
    columnNumber = max(AA(max)(CV))+1
    VC = [[] for k in range(columnNumber)]
    for k,cell in enumerate(CV):
        for v in cell:
            VC[v] += [k]
    return VC

```

◇

Macro referenced in [12a](#).

Computation of splitting tests In order to compute, in the simplest and more general way, whether each of the two split d -cells is internal or external to the splitting boundary $d-1$ -facet, it is necessary to consider the oriented covector ϕ (or one-form) canonically associated to the facet f by the covector representation theorem, i.e. the corresponding oriented hyperplane. In this case, the internal/external attribute of the split cell will be computed by evaluating the pairing $\langle v, \phi \rangle$.

$\langle \text{Splitting tests 7c} \rangle \equiv$

```

""" Splitting tests """
def testingSubspace(V,covector):
    def testingSubspace0(vcell):
        inout = SIGN(sum([INNERPROD([1.]+V[v],covector)] for v in vcell)))
        return inout
    return testingSubspace0

def cuttingTest(covector,polytope,V):
    signs = [INNERPROD([covector, [1.]+V[v]]) for v in polytope]
    signs = eval(vcode(signs))
    return any([value<-0.001 for value in signs]) and any([value>0.001 for value in signs])

def tangentTest(covector,facet,adjCell,V):
    common = list(set(facet).intersection(adjCell))
    signs = [INNERPROD([covector, [1.]+V[v]]) for v in common]
    count = 0
    for value in signs:
        if -0.0001<value<0.0001: count +=1
    if count >= len(V[0]):
        return True
    else:
        return False

```

◇

Macro referenced in [8](#).

Elementary splitting test Let us remember that the adjacency matrix between d -cells is computed via SpMSpM multiplication by the double application

`adjacencyQuery(V,CV)(cell),`

where the first application `adjacencyQuery(V,CV)` returns a partial function with buffering of the adjacency matrix, and the second application to `cell` returns the list of adjacent d -cells sharing with it a $(d - 1)$ -dimensional facet.

\langle Elementary splitting test 8 $\rangle \equiv$

\langle Splitting tests **7c** \rangle

```
""" Elementary splitting test """
def dividenda(V,CV, cell,facet,covector,unchosen):
    out = []
    adjCells = adjacencyQuery(V,CV)(cell)
    for adjCell in set(adjCells).difference(unchosen):
        if (cuttingTest(covector,CV[adjCell],V) and \
            cellFacetIntersecting(facet,adjCell,covector,V,CV)) or \
            tangentTest(covector,facet,CV[adjCell],V): out += [adjCell]
    return out
```

◇

Macro referenced in **12a**.

CDC cell splitting with one or more facets

\langle CDC cell splitting with one or more cutting facets 9a $\rangle \equiv$

```
""" CDC cell splitting with one or more cutting facets """
def fragment(cell,cellCuts,V,CV,BC):
    vcell = CV[cell]
    cellFragments = [[V[v] for v in vcell]]

    for f in cellCuts[cell]:
        facet = BC[f]
        plane = COVECTOR([V[v] for v in facet])
        for k,fragment in enumerate(cellFragments):

            #if not tangentTest(plane,facet,fragment,V):
            [below,equal,above] = SPLITCELL(plane,fragment)
            if below != above:
                cellFragments[k] = below
                cellFragments += [above]

    return cellFragments
```

◇

Macro referenced in **12a**.

SCDC splitting with every boundary facet The function `makeSCDC` is used to compute the LAR model (W, CW) of the SCDC. It takes as input the LAR model (V, CV) of the CDC, and the LAR model (V, BC) of the input Boolean Complex, and returns also the vertex-cell relation VC , i.e. the transposed of CV .

For every $k \in BC$, a list `cellsToSplit`

```

⟨SCDC splitting with every boundary facet 9b⟩ ≡
""" SCDC splitting with every boundary facet """
def makeSCDC(V,CV,BC):
    index = -1
    defaultValue = -1
    VC = invertRelation(CV)
    CW = []
    Wdict = dict()
    BCellcovering = boundaryCover(V,CV,BC,VC)
    cellCuts = invertRelation(BCellcovering)
    for k in range(len(CV) - len(cellCuts)): cellCuts += [[]]

    for k,frags in enumerate(cellCuts):
        if cellCuts[k] == []:
            cell = []
            for v in CV[k]:
                key = vcode(V[v])
                if Wdict.get(key,defaultValue) == defaultValue:
                    index += 1
                    Wdict[key] = index
                    cell += [index]
            else:
                cell += [Wdict[key]]
            CW += [cell]
        else:
            cellFragments = fragment(k,cellCuts,V,CV,BC)
            for cellFragment in cellFragments:
                cellFrag = []
                for v in cellFragment:
                    key = vcode(v)
                    if Wdict.get(key,defaultValue) == defaultValue:
                        index += 1
                        Wdict[key] = index
                        cellFrag += [index]
                    else:
                        cellFrag += [Wdict[key]]
                CW += [cellFrag]
    W = sorted(zip( Wdict.values(), Wdict.keys() ))
    W = AA(eval)(TRANS(W)[1])
    return W,CW,VC,BCellcovering,cellCuts

```

◇

Macro referenced in 12a.

Computation of boundary facets covering with CDC cells

⟨ Computation of boundary facets covering with CDC cells 10 ⟩ ≡

```

""" Computation of boundary facets covering with CDC cells """
def boundaryCover(V,CV,BC,VC):
    cellsToSplit = list()
    boundaryCellCovering = []
    for k,facet in enumerate(BC):
        covector = COVECTOR([V[v] for v in facet])
        seedsOnFacet = VC[facet[0]]
        cellsToSplit = [dividenda(V,CV, cell,facet,covector,[]) for cell in seedsOnFacet ]
        cellsToSplit = set(CAT(cellsToSplit))
        while True:
            newCells = [dividenda(V,CV, cell,facet,covector,cellsToSplit) for cell in cellsToSplit]
            if newCells != []: newCells = CAT(newCells)
            covering = cellsToSplit.union(newCells)
            if covering == cellsToSplit:
                break
            cellsToSplit = covering
        boundaryCellCovering += [list(covering)]
    return boundaryCellCovering

```

◇

Macro referenced in 12a.

Cell-facet intersection test

⟨ Cell-facet intersection test 11 ⟩ ≡

```

""" Cell-facet intersection test """
def cellFacetIntersecting(boundaryFacet,cell,covector,V,CV):
    points = [V[v] for v in CV[cell]]
    vcell1,newFacet,vcell2 = SPLITCELL(covector,points,tolerance=1e-4,ntry=4)
    boundaryFacet = [V[v] for v in boundaryFacet]
    translVector = boundaryFacet[0]

    # translation
    newFacet = [ VECTDIFF([v,translVector]) for v in newFacet ]
    boundaryFacet = [ VECTDIFF([v,translVector]) for v in boundaryFacet ]

    # linear transformation: boundaryFacet -> standard (d-1)-simplex
    d = len(V[0])
    transformMat = mat( boundaryFacet[1:d] + [covector[1:]] ).T.I

```

```

# transformation in the subspace x_d = 0
newFacet = (transformMat * (mat(newFacet).T)).T.tolist()
boundaryFacet = (transformMat * (mat(boundaryFacet).T)).T.tolist()

# projection in E^{d-1} space and Boolean test
newFacet = MKPOL([ AA(lambda v: v[:-1])(newFacet), [range(1,len(newFacet)+1)], None ])
boundaryFacet = MKPOL([ AA(lambda v: v[:-1])(boundaryFacet), [range(1,len(boundaryFacet)+1)] ])
verts,cells,pols = UKPOL(INTERSECTION([newFacet,boundaryFacet]))

if verts == []: return False
else: return True

```

◇

Macro referenced in [12a](#).

4 Step 3: cell labeling

The goal of this stage is to label every cell of the SCDC with two bits, corresponding to the input spaces A and B , and telling whether the cell is either internal (1) or external (0) to either spaces.

4.1 Requirements

Input The output of previous algorithm stage.

Output The array `cellLabels` with *shape* $\text{len}(\text{PW}) \times 2$, and values in $\{0,1\}$.

5 Step 4: greedy cell gathering

The goal of this stage is to make as lower as possible the number of cells in the output LAR of the space AB , partitioned into convex cells.

Input The LAR model (W, PW) of the SCDC and the array `cellLabels`.

Output The LAR representation (W, RW) of the final fragmented and labeled space AB .

6 Exporting the library

```

"lib/py/bool1.py" 12a ≡
    """ Module for Boolean ops with LAR """
    <Initial import of modules 17b>
    DEBUG = False

```


Macro referenced in 13, 14ab, 15ab, 16abc, 17a.

"test/py/bool1/test1.py" 13 ≡

```
import sys
""" import modules from larcc/lib """
sys.path.insert(0, 'lib/py/')
from bool1 import *

""" Definition of Boolean arguments """
V1 = [[3,0],[11,0], [13,10], [10,11], [8,11], [6,11], [4,11], [1,10], [4,3], [6,4],
      [8,4], [10,3]]
FV1 = [[0,1,8,9,10,11],[1,2,11], [3,10,11], [4,5,9,10], [6,8,9], [0,7,8], [2,3,11],
       [3,4,10], [5,6,9], [6,7,8]]
EV1 = [[0,1],[0,7],[0,8],[1,2],[1,11],[2,3],[2,11],[3,4],[3,10],[3,11],[4,5],[4,10],[5,6],[5,9]]
VV1 = AA(LIST)(range(len(V1)))

V2 = [[0,3],[14,2], [14,5], [14,7], [14,11], [0,8], [3,7], [3,5]]
FV2 = [[0,5,6,7], [0,1,7], [4,5,6], [2,3,6,7], [1,2,7], [3,4,6]]
EV2 = [[0,1],[0,5],[0,7],[1,2],[1,7],[2,3],[2,7],[3,4],[3,6],[4,5],[4,6],[5,6],[6,7]]
VV2 = AA(LIST)(range(len(V2)))

arg1 = V1,(VV1,EV1,FV1)
arg2 = V2,(VV2,EV2,FV2)
⟨Debug input and vertex merging 12b⟩
◇
```

"test/py/bool1/test2.py" 14a ≡

```
import sys
""" import modules from larcc/lib """
sys.path.insert(0, 'lib/py/')
from bool1 import *

V1 = [[3,0],[11,0], [13,10], [10,11], [8,11], [6,11], [4,11], [1,10], [4,3], [6,4],
      [8,4], [10,3]]

FV1 = [[0,1,8,9,10,11],[1,2,11], [3,10,11], [4,5,9,10], [6,8,9], [0,7,8]]
EV1 = [[0,1],[0,7],[0,8],[1,2],[1,11],[2,11],[3,10],[3,11],[4,5],[4,10],[5,9],[6,8],[6,9],[7,8]]
VV1 = AA(LIST)(range(len(V1)))

V2 = [[0,3],[14,2], [14,5], [14,7], [14,11], [0,8], [3,7], [3,5]]
FV2 = [[0,5,6,7], [0,1,7], [4,5,6], [2,3,6,7], [1,2,7], [3,4,6]]
EV2 = [[0,1],[0,5],[0,7],[1,2],[1,7],[2,3],[2,7],[3,4],[3,6],[4,5],[4,6],[5,6],[6,7]]
VV2 = AA(LIST)(range(len(V2)))
```

```

    arg1 = V1,(VV1,EV1,FV1)
    arg2 = V2,(VV2,EV2,FV2)
    <Debug input and vertex merging 12b>
    ◇

"test/py/bool1/test3.py" 14b ≡

import sys
""" import modules from larcc/lib """
sys.path.insert(0, 'lib/py/')
from bool1 import *

V1 = [[3,0],[11,0], [13,10], [10,11], [8,11], [6,11], [4,11], [1,10], [4,3], [6,4],
      [8,4], [10,3]]

FV1 = [[0,1,8,9,10,11],[1,2,11], [3,10,11], [4,5,9,10], [6,8,9], [0,7,8]]
EV1 = [[0,1],[0,7],[0,8],[1,2],[1,11],[2,11],[3,10],[3,11],[4,5],[4,10],[5,9],[6,8],[6,9],[7,8]]
VV1 = AA(LIST)(range(len(V1)))

V2 = [[0,3],[14,2], [14,5], [14,7], [14,11], [0,8], [3,7], [3,5]]
FV2 = [[0,5,6,7], [0,1,7], [4,5,6], [2,3,6,7]]
EV2 = [[0,1],[0,5],[0,7],[1,7],[2,3],[2,7],[3,6],[4,5],[4,6],[5,6],[6,7]]
VV2 = AA(LIST)(range(len(V2)))

arg1 = V1,(VV1,EV1,FV1)
arg2 = V2,(VV2,EV2,FV2)
<Debug input and vertex merging 12b>
◇

"test/py/bool1/test4.py" 15a ≡

import sys
""" import modules from larcc/lib """
sys.path.insert(0, 'lib/py/')
from bool1 import *

V1 = [[0,0],[10,0],[10,10],[0,10]]
FV1 = [range(4)]
EV1 = [[0,1],[1,2],[2,3],[0,3]]
VV1 = AA(LIST)(range(len(V1)))

V2 = [[2.5,2.5],[12.5,2.5],[12.5,12.5],[2.5,12.5]]
FV2 = [range(4)]
EV2 = [[0,1],[1,2],[2,3],[0,3]]
VV2 = AA(LIST)(range(len(V2)))

arg1 = V1,(VV1,EV1,FV1)

```

```

    arg2 = V2, (VV2, EV2, FV2)
    <Debug input and vertex merging 12b>
    ◇

"test/py/bool1/test5.py" 15b ≡

import sys
""" import modules from larcc/lib """
sys.path.insert(0, 'lib/py/')
from bool1 import *

V1 = [[0,0],[5,0],[5,5],[0,5]]
FV1 = [range(4)]
EV1 = [[0,1],[1,2],[2,3],[0,3]]
VV1 = AA(LIST)(range(len(V1)))

V2 = [[5,0],[10,0],[10,5],[5,5]]
FV2 = [range(4)]
EV2 = [[0,1],[1,2],[2,3],[0,3]]
VV2 = AA(LIST)(range(len(V2)))

arg1 = V1, (VV1, EV1, FV1)
arg2 = V2, (VV2, EV2, FV2)
<Debug input and vertex merging 12b>
◇

"test/py/bool1/test6.py" 16a ≡

import sys
""" import modules from larcc/lib """
sys.path.insert(0, 'lib/py/')
from bool1 import *

V1 = [[0,0,0],[10,0,0],[10,10,0],[0,10,0],[0,0,10],[10,0,10],[10,10,10],[0,10,10]]
V1, [VV1, EV1, FV1, CV1] = larCuboids((1,1,1), True)
V1 = [SCALARVECTPROD([5,v]) for v in V1]

V2 = [SUM([v, [2.5, 2.5, 2.5]]) for v in V1]
[VV2, EV2, FV2, CV2] = [VV1, EV1, FV1, CV1]

arg1 = V1, (VV1, EV1, FV1, CV1)
arg2 = V2, (VV2, EV2, FV2, CV2)
<Debug input and vertex merging 12b>
◇

"test/py/bool1/test7.py" 16b ≡

```



```

import sys
""" import modules from larcc/lib """
sys.path.insert(0, 'lib/py/')
from bool1 import *

V1 = [[0,0],[10,0],[10,10],[0,10]]
FV1 = [range(4)]
EV1 = [[0,1],[1,2],[2,3],[0,3]]
VV1 = AA(LIST)(range(len(V1)))

V2 = [[2.5,2.5],[7.5,2.5],[7.5,7.5],[2.5,7.5]]
FV2 = [range(4)]
EV2 = [[0,1],[1,2],[2,3],[0,3]]
VV2 = AA(LIST)(range(len(V2)))

arg1 = V1,(VV1,EV1,FV1)
arg2 = V2,(VV2,EV2,FV2)
⟨Debug input and vertex merging 12b⟩
◇

"test/py/bool1/test8.py" 16c ≡

import sys
""" import modules from larcc/lib """
sys.path.insert(0, 'lib/py/')
from bool1 import *

n = 48
V1 = [[5*cos(angle*2*PI/n)+2.5, 5*sin(angle*2*PI/n)+2.5] for angle in range(n)]
FV1 = [range(n)]
EV1 = TRANS([range(n),range(1,n+1)]); EV1[-1] = [0,n-1]
VV1 = AA(LIST)(range(len(V1)))

V2 = [[4*cos(angle*2*PI/n), 4*sin(angle*2*PI/n)] for angle in range(n)]
FV2 = [range(n)]
EV2 = EV1
VV2 = AA(LIST)(range(len(V2)))

arg1 = V1,(VV1,EV1,FV1)
arg2 = V2,(VV2,EV2,FV2)
⟨Debug input and vertex merging 12b⟩
◇

"test/py/bool1/test9.py" 17a ≡

```

```

import sys
""" import modules from larcc/lib """
sys.path.insert(0, 'lib/py/')
from bool1 import *

V1 = [[0,0],[15,0],[15,14],[0,14]]
FV1 = [range(4)]
EV1 = [[0,1],[1,2],[2,3],[0,3]]
VV1 = AA(LIST)(range(len(V1)))

V2 = [[1,1],[7,1],[7,6],[1,6],[8,1],[14,1],[14,7],[8,7],[1,7],[7,7],[7,13],[1,13],[8,8],[14,8],[14,14],[7,14],[1,14]]
FV2 = [range(4),range(4,8),range(8,12),range(12,16)]
EV2 = [[0,1],[1,2],[2,3],[0,3],[4,5],[5,6],[6,7],[4,7],[8,9],[9,10],[10,11],[8,11],[12,13],[13,14],[11,14]]
VV2 = AA(LIST)(range(len(V2)))

arg1 = V1,(VV1,EV1,FV1)
arg2 = V2,(VV2,EV2,FV2)
⟨Debug input and vertex merging 12b⟩
◇

```

A Appendix: utility functions

⟨Initial import of modules 17b⟩ ≡

```

from pyplasm import *
from scipy import *
import sys
""" import modules from larcc/lib """
sys.path.insert(0, 'lib/py/')
from lar2psm import *
from simplexn import *
from larcc import *
from largrid import *
from myfont import *
from mapper import *
◇

```

Macro referenced in 12a.

A.1 Numeric utilities

A small set of utility functions is used to transform a *point* representation, given as array of coordinates, into a string of fixed format to be used as point key into python dictionaries.

⟨Symbolic utility to represent points as strings 18⟩ ≡

```

""" TODO: use package Decimal (http://docs.python.org/2/library/decimal.html) """
global PRECISION
PRECISION = 4.95

def verySmall(number): return abs(number) < 10**-(PRECISION/1.15)

def prepKey (args): return "[" + ".join(args) + "]"

def fixedPrec(value):
    out = round(value*10**(PRECISION*1.1))/10**(PRECISION*1.1)
    if out == -0.0: out = 0.0
    return str(out)

def vcode (vect):
    """
    To generate a string representation of a number array.
    Used to generate the vertex keys in PointSet dictionary, and other similar operations.
    """
    return prepKey(AA(fixedPrec)(vect))

```

◇

Macro never referenced.

References

- [CL13] CVD-Lab, *Linear algebraic representation*, Tech. Report 13-00, Roma Tre University, October 2013.