Accelerated intersection of geometric objects *

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Abstract

This module contains the first experiments of a parallel implementation of the intersection of (multidimensional) geometric objects. The first installment is being oriented to the intersection of line segment in the 2D plane. A generalization of the algorithm, based on the classification of the containment boxes of the geometric values, will follow quickly.

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1 Introduction

An easily parallelizable implementation of the accelerated intersection of geometric objects is given in this module. Our first aim is to implement a specialized version for simplices, that generalizes the nD-trees of points (that are 0-simplices), to (d-1)-dimensional simplices in d-space, starting with the intersection of line segments in the plane. Our plan is to follow with an implementation for intersection of general $non\ convex$ sets.

2 Implementation

The first implementation of this module concerns the computation of the intersection points among a set of line segment in the 2D plane. The containment boxes of the input segments are iteratively classified against the 1-dimensional centroid of smaller and smaller buckets of data.

At the end of the classification, where the same geometric object may be inserted in several different buckets, a *brute-force* intersection is applied to each final subset. Finally, the duplicated intersection points are removed, and a 1-dimensional LAR data structure is generated, with 1-cells given by the split line segments.

A complete LAR of the plane partition generated by the arrangment of lines is then computed by: (a) generating the maximal 2-connected components of such 1-dimensional graph; and (b) by traversing in counter-clockwise order the generated subgraphs to report the 2-dimensional cells of the plane partition.

The splitting algorithm may be easily parallelized, since both during their generation and at the end of this one, the various buckets of data can be dispatched to different processors for independent computation, followed by elimination of duplicates. In particular, a standard *map-reduce* software infrastructure may be used for this parallelization purpose.

2.1 Construction of independent buckets

Containment boxes Given as input a list randomLineArray of pairs of 2D points, the function containment2DBoxes returns, in the same order, the list of containment boxes of the input lines. A containment box of a geometric object of dimension d is defined as the minimal d-cuboid, equioriented with the reference frame, that contains the object. For a 2D line it is given by the tuple (x1, y1, x2, y2), where (x1, y1) is the point of minimal coordinates, and (x2, y2) is the point of maximal coordinates.

 $\langle \text{Containment boxes 2a} \rangle \equiv$

```
""" Containment boxes """
def containment2DBoxes(randomLineArray):
    boxes = [eval(vcode([min(x1,x2),min(y1,y2),max(x1,x2),max(y1,y2)]))
        for ((x1,y1),(x2,y2)) in randomLineArray]
    return boxes
```

Splitting the input above and below a threshold

```
⟨Splitting the input above and below a threshold 2b⟩ ≡
""" Splitting the input above and below a threshold """
def splitOnThreshold(boxes, subset, coord):
    theBoxes = [boxes[k] for k in subset]
    threshold = centroid(theBoxes, coord)
    ncoords = len(boxes[0])/2
    a = coord%ncoords
    b = a+ncoords
    below, above = [],[]
    for k in subset:
        if boxes[k][a] <= threshold: below += [k]
    for k in subset:
        if boxes[k][b] >= threshold: above += [k]
    return below, above
```

Macro referenced in 34b.

Iterative splitting of box buckets

```
for h,inters in enumerate(geomInters):
        geomInters[h] = geomInters[h].difference([h])
   return AA(list)(geomInters)
def boxBuckets(boxes):
   bucket = range(len(boxes))
   splittingStack = [bucket]
   finalBuckets = []
   while splittingStack != []:
        bucket = splittingStack.pop()
        below,above = splitOnThreshold(boxes,bucket,1)
        below1,above1 = splitOnThreshold(boxes,above,2)
        below2,above2 = splitOnThreshold(boxes,below,2)
        splitting(above,below1,above1, finalBuckets,splittingStack)
        splitting(below,below2,above2, finalBuckets,splittingStack)
        finalBuckets = list(set(AA(tuple)(finalBuckets)))
   parts = geomPartitionate(boxes,finalBuckets)
   return AA(sorted)(parts)
   #return finalBuckets
```

2.2 Brute force intersection within the buckets

Intersection of two line segments

```
\langle Intersection of two line segments 3b \rangle \equiv
     """ Intersection of two line segments """
     def segmentIntersect(boxes,lineArray,pointStorage):
         def segmentIntersectO(h):
             p1,p2 = lineArray[h]
             line1 = '['+ vcode(p1) +','+ vcode(p2) +']'
              (x1,y1),(x2,y2) = p1,p2
              B1,B2,B3,B4 = boxes[h]
              def segmentIntersect1(k):
                  p3,p4 = lineArray[k]
                  line2 = '['+ vcode(p3) +','+ vcode(p4) +']'
                  (x3,y3),(x4,y4) = p3,p4
                  b1,b2,b3,b4 = boxes[k]
                  if not (b3<B1 or B3<b1 or b4<B2 or B4<b2):
                  #if True:
                      m23 = mat([p2,p3])
                      m14 = mat([p1,p4])
                      m = m23 - m14
                      v3 = mat([p3])
                      v1 = mat([p1])
```

```
v = v3-v1
                      a=m[0,0]; b=m[0,1]; c=m[1,0]; d=m[1,1];
                      det = a*d-b*c
                      if det != 0:
                          m_{inv} = mat([[d,-b],[-c,a]])*(1./det)
                           alpha, beta = (v*m_inv).tolist()[0]
                          #alpha, beta = (v*m.I).tolist()[0]
                           if -0.0 \le alpha \le 1 and -0.0 \le beta \le 1:
                               pointStorage[line1] += [alpha]
                               pointStorage[line2] += [beta]
                               return list(array(p1)+alpha*(array(p2)-array(p1)))
                  return None
              return segmentIntersect1
         return segmentIntersect0
Macro referenced in 34b.
Brute force bucket intersection
\langle Brute force bucket intersection 4\rangle \equiv
     """ Brute force bucket intersection """
     def lineBucketIntersect(boxes,lineArray, h,bucket, pointStorage):
         intersect0 = segmentIntersect(boxes,lineArray,pointStorage)
         intersectionPoints = []
         intersect1 = intersect0(h)
         for line in bucket:
              point = intersect1(line)
              if point != None:
                  intersectionPoints.append(eval(vcode(point)))
         return intersectionPoints
Macro referenced in 34b.
Accelerate intersection of lines
\langle Accelerate intersection of lines 5\rangle \equiv
     """ Accelerate intersection of lines """
     def lineIntersection(lineArray):
         lineArray = [line for line in lineArray if len(line)>1]
         from collections import defaultdict
         pointStorage = defaultdict(list)
         for line in lineArray:
              print "line =",line
              p1,p2 = line
```

```
key = '['+ vcode(p1) +','+ vcode(p2) +']'
pointStorage[key] = []

boxes = containment2DBoxes(lineArray)
buckets = boxBuckets(boxes)
intersectionPoints = set()
for h,bucket in enumerate(buckets):
    pointBucket = lineBucketIntersect(boxes,lineArray, h,bucket, pointStorage)
    intersectionPoints = intersectionPoints.union(AA(tuple)(pointBucket))

frags = AA(eval)(pointStorage.keys())
params = AA(COMP([sorted,list,set,tuple,eval,vcode]))(pointStorage.values())
return intersectionPoints,params,frags ### GOOD: 1, WRONG: 2 !!!
```

2.3 Generation of LAR representation of split segments

The function lines2lar is used to generate a 1-dimensional LAR complex from an array of lines, i.e. of pairs of 2D points. For every *line* in frags is computed an *ordered* list outline of *symbolic* intersection points, including the first and last vertex of the line, and every interior point generated by the list params[k].

Then, for every symbolic representation key of a point in outline, a dictionary vertex is either created or retrieved, and a corresponding edge is orderly created, using the index of the point. At the same time, the vertices created in this way are accumulated within the V array. Finally, each edge in EV is extended to contain a second vertex index using the subsequent edge.

The third stage finalizes the vertex set of the output LAR, by identifying the closest vertices, i.e. those at distance less or equal to the current resolution, set to 10**(-PRECISION), by searching via the scipy.spatialKDTree the pairs of vertices at less than this distance.

A fourth stage identifies the possibly duplicated edges. Some of these could appear, e.g., when importing a set of adjacent boxes from some drawing program, to generate an array of lines, to be mutually intersected and transformed into a LAR data structure.

Create the LAR of fragmented lines

```
⟨ Create the LAR of fragmented lines 6⟩ ≡
    """ Create the LAR of fragmented lines """
    from scipy import spatial

def lines2lar(lineArray):
    _,params,frags = lineIntersection(lineArray)
    vertDict = dict()
```

```
index, default Value, V, EV = -1, -1, [], []
for k,(p1,p2) in enumerate(frags):
    outline = [vcode(p1)]
    if params[k] != []:
        for alpha in params[k]:
            if alpha != 0.0 and alpha != 1.0:
                p = list(array(p1)+alpha*(array(p2)-array(p1)))
                outline += [vcode(p)]
    outline += [vcode(p2)]
    edge = []
    for key in outline:
        if vertDict.get(key,defaultValue) == defaultValue:
            index += 1
            vertDict[key] = index
            edge += [index]
            V += [eval(key)]
        else:
            edge += [vertDict[key]]
        EV.extend([[edge[k],edge[k+1]] for k,v in enumerate(edge[:-1])])
model = (V, EV)
return larSimplify(model)
```

2.4 Biconnected components of a 1-complex

An implementation of the Hopcroft-Tarjan algorithm [HT73] for computation of the biconnected components of a graph is given here.

Biconnected components

```
⟨Biconnected components 7a⟩ ≡

""" Biconnected components """

⟨Adjacency lists of 1-complex vertices 7b⟩

⟨Main procedure for biconnected components 7c⟩

⟨Hopcroft-Tarjan algorithm 8a⟩

⟨Output of biconnected components 8b⟩

⋄
```

Adjacency lists of 1-complex vertices

```
⟨ Adjacency lists of 1-complex vertices 7b⟩ ≡
    """ Adjacency lists of 1-complex vertices """
    def vertices2vertices(model):
        V,EV = model
        csrEV = csrCreate(EV)
        csrVE = csrTranspose(csrEV)
        covV = matrixProduct(csrVE,csrEV)
        cooVV = csrVV.tocoo()
        data,rows,cols = AA(list)([cooVV.data, cooVV.row, cooVV.col])
        triples = zip(data,rows,cols)
        VV = [[] for k in range(len(V))]
        for datum,row,col in triples:
            if row != col: VV[col] += [row]
        return AA(sorted)(VV)
```

Macro referenced in 7a.

Main procedure for biconnected components

```
\langle Main procedure for biconnected components 7c\rangle \equiv
     """ Main procedure for biconnected components """
     def biconnectedComponent(model):
         W_{,-} = model
         V = range(len(W))
         count = 0
         stack, out = [],[]
         visited = [None for v in V]
         parent = [None for v in V]
         d = [None for v in V]
         low = [None for v in V]
         for u in V: visited[u] = False
         for u in V: parent[u] = []
         VV = vertices2vertices(model)
         for u in V:
              if not visited[u]:
                  DFV_visit( VV,out,count,visited,parent,d,low,stack, u )
         return W, [component for component in out if len(component) > 1]
```

Hopcroft-Tarjan algorithm

Macro referenced in 7a.

```
\langle Hopcroft-Tarjan algorithm 8a\rangle \equiv
```

```
""" Hopcroft-Tarjan algorithm """
def DFV_visit( VV,out,count,visited,parent,d,low,stack,u ):
   visited[u] = True
   count += 1
   d[u] = count
   low[u] = d[u]
   for v in VV[u]:
        if not visited[v]:
            stack += [(u,v)]
            parent[v] = u
            DFV_visit( VV,out,count,visited,parent,d,low,stack, v )
            if low[v] >= d[u]:
                out += [outputComp(stack,u,v)]
            low[u] = min( low[u], low[v] )
        else:
            if not (parent[u] == v) and (d[v] < d[u]):
                stack += [(u,v)]
                low[u] = min(low[u], d[v])
```

Macro referenced in 7a.

Output of biconnected components

```
⟨ Output of biconnected components 8b⟩ ≡
    """ Output of biconnected components """
    def outputComp(stack,u,v):
        out = []
        while True:
            e = stack.pop()
            out += [list(e)]
            if e == (u,v): break
        return list(set(AA(tuple)(AA(sorted)(out))))
```

Macro referenced in 7a.

2.5 2D cells from biconnected components

It is very easy, using the LAR representation of topology, to compute the 2-cells of the plane partitions (see Figures 1b and 1c) induced by the biconnected components extracted from a graph (1-complex).

In particular, let us consider the CSR (Compressed Sparse Row) representation of the characteristic matrix M_1 , here usually denoted as EV, in order to remark that we represent the edges on the rows, and the vertices on the columns of the matrix. As such it is a binary matrix. So, we can readily reconstruct the topology of 2-cells by associating to



Figure 1: Two random line arrangements, and the biconnected components extracted by their LAR 1-complexes.

each non-zero (sparse) matrix element $angle_k EV(h, k)$ the angle in radians that the edge e_h forms with the orizontal line, when it incides on the vertex v_k .

```
Of course, if e_h = (v_{k_1}, v_{k_2}), then it will be
```

```
angle_EV(h, k_2) = angle_EV(h, k_1) + \pi = -angle_EV(h, k_1)
```

Therefore, the columns of angle_EV, i.e. the rows of angle_VE := angle_EV^t, after being sorted on their angles α , and associated with the angle differences $\Delta \alpha$, will provide a basis of elementary 1 - cochains that evaluate to zero for each closed 1-cochain, i.e. for every cycle supported by the linear space of 1-chains on the given line arrangement.

Slope of edges

Circular ordering of edges around vertices

```
\langle Slope of edges 9 \rangle \equiv
     """ Circular ordering of edges around vertices """
     def edgeSlopeOrdering(model):
         V,EV = model
         VE, VE_angle = invertRelation(EV),[]
         for v,ve in enumerate(VE):
              ve_angle = []
             if ve != []:
                  for edge in ve:
                      v0,v1 = EV[edge]
                                      x,y = list(array(V[v1]) - array(V[v0]))
                      if v == v0:
                                      x,y = list(array(V[v0]) - array(V[v1]))
                      elif v == v1:
                      angle = math.atan2(y,x)
                      ve_angle += [180*angle/PI]
              pairs = sorted(zip(ve_angle,ve))
              #VE_angle += [TRANS(pairs)[1]]
              VE_angle += [[pair[1] for pair in pairs]]
         return VE_angle
```

Macro referenced in 34b.

Ordered incidence relationship vertices to edges. As we have seen, the VE_angle list of lists reports, for every vertex in V, the list of incident edges, counterclockwise ordered around the vertex. Therefore the ordered_csrVE function, given below, returns the "compressed sparse row" matrix, row-indexed by vertices and column-indexed by edges, and such that in position (v, e) contains the index ℓ of the next edge (after e, say) in the counterclockwise ordering of edges around v.

 \langle Ordered incidence relationship of vertices and edges 11a \rangle \equiv

```
""" Ordered incidence relationship of vertices and edges """
def ordered_csrVE(VE_angle):
    triples = []
    for v,ve in enumerate(VE_angle):
        n = len(ve)
        for k,edge in enumerate(ve):
            triples += [[v, ve[k], ve[ (k+1)%n ]]]
    csrVE = triples2mat(triples,shape="csr")
    return csrVE
```

Faces from biconnected components Since edges in the plane partition induced by a line arrangement are (d-1)-cells, they are located on the boundary of $two\ d$ -cells (faces) of the partition. Hence, the traversal algorithm of the data structure storing the relevant information may be driven by signing the two extremes (vertices) of each edge as either already visited or not.

```
\langle Faces from biconnected components 11b\rangle \equiv
     """ Faces from biconnected components """
     def firstSearch(visited):
         for edge,vertices in enumerate(visited):
             for v,vertex in enumerate(vertices):
                  if visited[edge,v] == 0.0:
                      visited[edge,v] = 1.0
                      return edge, v
         return -1,-1
     def facesFromComps(model):
         V,EV = model
         # Remove zero edges
         EV = list(set([ tuple(sorted([v1,v2])) for v1,v2 in EV if v1!=v2 ]))
         FV = []
         VE_angle = edgeSlopeOrdering((V,EV))
         csrEV = ordered_csrVE(VE_angle).T
         visited = zeros((len(EV),2))
         edge,v = firstSearch(visited)
         vertex = EV[edge][v]
         fv = []
         while True:
             if (edge, v) == (-1, -1):
                  break #return [face for face in FV if face != None]
             elif (fv == []) or (fv[0] != vertex):
```

```
fv += [vertex]
        nextEdge = csrEV[edge,vertex]
        v0,v1 = EV[nextEdge]
        try:
            vertex, = set([v0,v1]).difference([vertex])
        except ValueError:
            print 'ValueError: too many values to unpack'
            break
        if v0==vertex: pos=0
        elif v1==vertex: pos=1
        if visited[nextEdge, pos] == 0:
            visited[nextEdge, pos] = 1
            edge = nextEdge
    else:
        FV += [fv]
        fv = []
        edge,v = firstSearch(visited)
        vertex = EV[edge][v]
    FV = [face for face in FV if face != None]
return V,FV,EV
```

Txample The ordered csrVE (vertex-edge) matrix generated by the example of file test/py/inters/test07.py is shown in dense format in the example script below. Let us notice the each non-zero element csrVE(k,h) stores the index of the previous edge inciding on the vertex v_k before the edge e_h . The traversal of the data structure is made accordingly, in order to extract the vertices of all the faces (minimal edge cycles) generated by a line arrangement in the plane.

 $\langle \text{Example of VE matrix with nextEdge indices } 13 \rangle \equiv$

```
csr2DenseMatrix(csrVE)
>>> array([
                   Ο,
                                                 Ο,
    [12, 0,
               0,
                       0,
                            0,
                                0,
                                    0,
                                        0,
                                             0,
                                                      0, 11,
                                                              0,
                                                                       0],
                                                                       0],
    [ 1,
         2,
               0,
                   0,
                       0,
                            0,
                                0,
                                    0,
                                        0,
                                             0,
                                                 0,
                                                      0,
                                                          0,
                                                              0,
                                                                   Ο,
                                                     Ο,
    [ 0, 14,
               0,
                   0,
                       Ο,
                           0,
                                0,
                                    0,
                                        0,
                                             0,
                                                 0,
                                                          0,
                                                              0,
                                                                   1,
                                                                       0],
    [0,0,
               6,
                   5,
                       Ο,
                           2,
                                3,
                                    Ο,
                                        Ο,
                                             0,
                                                 0,
                                                     0,
                                                                       0],
                                                          Ο,
                                                              0,
                                        Ο,
                                                     Ο,
    [ 0,
          Ο,
               0, 10,
                       Ο,
                           Ο,
                                Ο,
                                    0,
                                             3,
                                                 9,
                                                          0,
                                                              0,
                                                                   Ο,
                                                                       0],
                                0,
                                        0,
                                                     Ο,
    [ 0,
          Ο,
              0,
                   0, 15,
                           0,
                                    0,
                                             0,
                                                 0,
                                                          Ο,
                                                              0,
                                                                   0,
                                                                       4],
              Ο,
                                    Ο,
                                        Ο,
    [ 0,
                   0, 12,
                           4,
                               Ο,
                                             0,
                                                 0,
                                                          5,
                                                     0,
                                                              0,
                                                                   0,
                                                                       0],
    [0, 0, 0, 0, 0, 0, 7,
                                    8, 6,
                                             0, 0,
                                                         Ο,
```

```
Ο,
            0,
                    7, 0,
                            0,
        Ο,
                Ο,
                    0, 10,
                            Ο,
                                8,
                                     Ο,
    0,
            0,
                                         0,
               Ο,
   0,
        0,
            0,
                    0,
                        0,
                            9,
                                0,
                                     0,
                                         0,
0,
   0,
        0,
            0,
               0,
                    0, 0,
                            0, 0, 13,
                                         0, 14, 11,
        0.
            0.
                0.
                    0,
                        0.
                            0.
                                0.
    0.
                                    0.
                                         0.15.
```

Macro never referenced.

2.6 Containments between non intersecting cycles

In this section we compute the containment relation between non-intersecting cycles generated on 2-faces by the incident faces. This step is preparatory to the representation of fragmented 2-faces — embedded in 2D — as LAR data structures, to be subsequently restored in the ambient 3D and sticked together to generate the the 2-skeleton of the Boolean complex.

For this purpose, the set of non-intersecting cycles, as lists of edges, are returned in the EVs array by the biconnectedComponent function. The pair V, EVs is therefore passed as input to the latticeArray function, that returns a dense matrix with elements in $\{-1,0,1\}$, where the element i,j either contains 0 if anyone (and hence all) of vertices of cycle j-th is external to cycle i-th, or contains 1 if anyone (and hence all) of vertices of cycle j-th is internal to cycle i-th, or finally contains -1 if anyone (and hence all) of vertices of cycle j-th is on boundary of cycle i-th. The last condition may hold only for diagonal elements i,j where i=j.

The returned testArray matrix of dimension $n \times n$, where n is the number of non-intersecting cycles on a fragmented 2D face, is the used by the cellsFromCycles function, that return a list of lists of cycles, each one defining the boundary of a single connected but possibly non path-connected 2-cell in the LAR representation of the fragmented 2-complex.

Classification of non intersecting cycles The function latticeArray takes as input the pair V, EVs, where V is the list of vertices and EVs is the list of cycles of a fragmented 2-cell in 2D. Each EVs element is given as a list of edges, given os pairs of integer vertex indices. The returned testArray matrix — characterizing the incidences between cycles — has dimension $n \times n$.

```
⟨ Classification of non intersecting cycles 14⟩ ≡

""" Classification of non intersecting cycles """

def latticeArray(V,EVs):
    n = len(EVs)
    testArray = []
    for k,ev in enumerate(EVs):
        row = []
        classify = pointInPolygonClassification((V,ev))
        for h in range(0,n):
```

```
i = EVs[h][0][0]
point = V[i]
test = classify(point)
if test=="p_in": row += [1]
elif test=="p_out": row += [0]
elif test=="p_on": row += [-1]
else: print "error: in cycle classification"
testArray += [row]
return testArray
```

Extraction of path-connected boundaries The function cellsFromCycles, given by the script below, takes as input the cycle-incidence matrix testArray, and return the list out of lists of cycles, providing the boundaries of a decomposition into connected (but possibly non path-connected) 2-cells of the fragmented 2-face whose non-intersecting cycles were computed as output of the biconnectedComponent function.

First the sons of each cycle are computed, i.e. the indices of cycles containd in it, as well as the level of every cycle, i.e. their position within the lattice of the containment relation (that is a partial order). The level of a cycle is computed as the sum of elements in its matrix column. The roots of the lattice, i.e. the more external cycles have level -1; the following levels have values 0, 1, 2, ..., respectively.

Then the cycles are ordered in the encreasing value of their level (or rank), and finally the significant subsets of disjoint cycles are extracted within the out list, starticg from the root cycle(s). The important properties exploited for the extraction are the following: (a) the first element of each sublist is the external boundary cycle, whereas the following cycles, if any, are its internal boundaries; (b) the rank difference between each external and internal boundary must be less or equal to one. It will be 0 only whe the sublist contains only one element (the external boundary) wich is being compared with itself.

Finally the sublists are pruned, by eliminating those whose first element has benn previously used within some of the previous ones (of course: was already used as an internal cycle).

```
⟨Extraction of path-connected boundaries 15a⟩ ≡

""" Extraction of path-connected boundaries """

def cellsFromCycles (testArray):
    n = len(testArray)
    sons = [[h]+[k for k in range(n) if row[k]==1] for h,row in enumerate(testArray)]
    level = [sum(col) for col in TRANS(testArray)]

def rank(sons): return [level[x] for x in sons]
    preCells = sorted(sons,key=rank)
```

```
def levelDifference(son,father): return level[son]-level[father]
root = preCells[0] [0]
out = [[son for son in preCells[0] if (levelDifference(son,root)<=1) ]]
for k in range(1,n):
    father = preCells[k] [0]
    inout = [son for son in preCells[k] if levelDifference(son,father)<=1 ]
    if not (inout[0] in CAT(out)):
        out += [inout]
return out</pre>
```

Testing containments between non intersecting cycles The test code for verifying the approch to computation of the containment lattice between non intersecting cycles is given below. Three test files, named respectively lattice, lattice1 and lattice2, with the different situations shown in Figure ??, may be imported from the directory test/svg/inters/. Other tests may be easily generated by inserting in this directory some .svg files generated with a drawing program.

```
"test/py/inters/test15.py" 15b \equiv
     """ Testing containments between non intersecting cycles """
     from larlib import *
     filename = "test/svg/inters/facade.svg"
     lines = svg2lines(filename)
     VIEW(STRUCT(AA(POLYLINE)(lines)))
     V,EV = lines2lar(lines)
     V,EVs = biconnectedComponent((V,EV))
     # candidate face
     FVs = AA(COMP([list,set,CAT]))(EVs)
     testArray = latticeArray(V,EVs)
     for k in range(len(testArray)):
        print k,testArray[k]
     print "\ncells = ", cellsFromCycles(testArray),"\n"
     VV = AA(LIST)(range(len(V)))
     submodel = STRUCT(MKPOLS((V,EV)))
     VIEW(larModelNumbering(1,1,1)(V,[VV,EV,FVs],submodel,0.15))
```

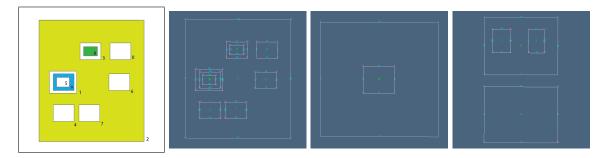


Figure 2: Some examples of nested non intersecting cycles. The corresponding solutions are given in the text.

2.7 Reduction of multiple cycles to a single polyline

The reduction of a lattice of non-intersecting 1-cycles on the boundary of a 2-cell into a single polyline is performed using a scan-line algorithm.

In order to filter the complications induced by edges aligned with the reference axes, first we perform a transformation of vertices from Cartesian to polar coordinates (see Figure 3).

Then a specialized scan-line algorithm is executed, producing a set of *bridge-edges* [YT85] that, added in double instance to the set EV of the LAR of the 2-cell, allow for a triangulation of its interior using the algorithm provided by the poly2tria module.

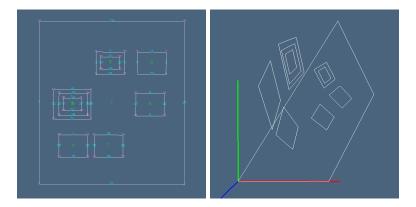


Figure 3: A set of non-intersecting boundary cycles in Cartesian and polar coordinates, using (improperly) an Euclidean metric in the transformed space.

Transforming to polar coordinates The transformation from Cartesian to polar coordinates of the vertices of a two-dimensional LAR model is given in the following script. An image of such transformation is shown in Figure 3. It is only used here to put the vertex coordinates in general position, so simplifying the scan-line algorithm.

```
⟨Transforming to polar coordinates 17a⟩ ≡
    """ Transforming to polar coordinates """
    def cartesian2polar(V):
        Z = [[sqrt(x*x + y*y),math.atan2(y,x)] for x,y in V]
        VIEW(STRUCT(MKPOLS((Z,EV))))
        return Z
```

Macro never referenced.

Scan line algorithm

```
\langle Scan line algorithm 17b \rangle \equiv
     """ Scan line algorithm """
     def scan(V,FVs, group,cycleGroup,cycleVerts):
         bridgeEdges = []
         scannedCycles = []
         for k,(point,cycle,v) in enumerate(cycleGroup[:-2]):
             nextCycle = cycleGroup[k+1][1]
             n = len(FVs[group][cycle])
             if nextCycle != cycle:
                  if not ((nextCycle in scannedCycles) and (cycle in scannedCycles)):
                      print "k =",k
                      scannedCycles += [nextCycle]
                      m = len(FVs[group][nextCycle])
                      v1, v2 = v, cycleGroup[k+1][2]
                      minDist = VECTNORM(VECTDIFF([V[v1],V[v2]]))
                      for i in FVs[group][cycle]:
                          for j in FVs[group][nextCycle]:
                              dist = VECTNORM(VECTDIFF([V[i],V[j]]))
                               if dist < minDist:</pre>
                                  minDist = dist
                                  v1,v2 = i,j
                      bridgeEdges += [(v1,v2)]
         return bridgeEdges[:-1]
```

Macro referenced in 34b.

Scan line algorithm input/output The two-dimensional LAR model \equiv V,EV was previously created by a lines2lar(lines) expression.

```
⟨ Scan line algorithm input/output 18⟩ ≡
    """ Scan line algorithm input/output """
    def connectTheDots(model):
        V,EV = model
```

```
V,EVs = biconnectedComponent((V,EV))
FV = AA(COMP([sorted,set,CAT]))(EVs)
testArray = latticeArray(V,EVs)
cells = cellsFromCycles(testArray)
FVs = [[FV[cycle] for cycle in cell] for cell in cells]

indexedCycles = [zip(FVs[h],range(len(FVs[h]))) for h,cell in enumerate(cells)]
indexedVerts = [CAT(AA(DISTR)(cell)) for cell in indexedCycles]
sortedVerts = [sorted([(V[v],c,v) for v,c in cell]) for cell in indexedVerts]

bridgeEdges = []
for (group,cycleGroup,cycleVerts) in zip(range(len(cells)),sortedVerts,indexedVerts):
    bridgeEdges += [scan(V,FVs, group,cycleGroup,cycleVerts)]
return cells,bridgeEdges
```

Orientation of component cycles of unconnected boundaries This step of the algorithm needs some preliminary computation, starting from the list EV of edges by vertices. Just notice that, at this point of the implementation (as the edges of a set of non intersecting cycles) they all are boundary edges, and hence the edgeBoundary variable can be filled with consecutive integers. The edgeCycles returned by boundaryCycles are very well characterized, according to how discussed in the description of the Edge cycles associated to a closed chain of edges paragraph. The corresponding vertexCycles are first oriented accordingly to boundaryCycles, then rotated (as equivalent permutations) in order to put their element of minimum index in first position (accordingly to the numeration of cycles used in the variable FVs). Then, the CVs will be set to contain the circularly ordered vertices of the various boundary cycles of each LAR 2-cell, in a manner analogous to FVs, where the vertices are not circularly ordered. Finally, the first (i.e. the external) cycle of each cell is counterclockwise oriented, whereas the other cycles (i.e. the internal ones), are oriented clockwise.

```
⟨Orientation of component cycles of unconnected boundaries 19⟩ ≡
    """ Orientation of component cycles of unconnected boundaries """
    def rotatePermutation(inputPermutation,transpositionNumber):
        n = transpositionNumber
        perm = inputPermutation
        permutation = range(n,len(perm))+range(n)
        return [perm[k] for k in permutation]

def canonicalRotation(permutation):
        n = permutation.index(min(permutation))
        return rotatePermutation(permutation,n)
```

```
def setCounterClockwise(h,k,cycle,areas,CVs):
         if areas[cycle] < 0.0:
             chain = copy.copy(CVs[h][k])
             CVs[h][k] = canonicalRotation(REVERSE(chain))
     def setClockwise(h,k,cycle,areas,CVs):
         if areas[cycle] > 0.0:
             chain = copy.copy(CVs[h][k])
             CVs[h][k] = canonicalRotation(REVERSE(chain))
     def orientBoundaryCycles(model,cells):
         V, EV = model
         edgeBoundary = range(len(EV))
         edgeCycles = boundaryCycles(edgeBoundary,EV)
         vertexCycles = [[ EV[e][1] if e>0 else EV[-e][0] for e in cycle ] for cycle in edgeCycles]
         rotations = [cycle.index(min(cycle)) for cycle in vertexCycles]
         theCycles = sorted([rotatePermutation(perm,n) for perm,n in zip(vertexCycles,rotations)])
         CVs = [[theCycles[cycle] for cycle in cell] for cell in cells]
         areas = surfIntegration((V,theCycles,EV),signed=True)
         for h,cell in enumerate(cells):
             for k,cycle in enumerate(cell):
                 if k == 0: setCounterClockwise(h,k,cycle,areas,CVs)
                 else: setClockwise(h,k,cycle,areas,CVs)
         return CVs
Macro referenced in 34b.
```

From nested boundary cycles to triangulation

```
\(\text{From nested boundary cycles to triangulation 20} \) \(\text{=} \)
\(\text{""" From nested boundary cycles to triangulation """ \)
\(\text{def boundaryCycles2triangulation( (V,EV) ): model = V,EV \)
\(\text{cells,bridgeEdges = connectTheDots(model)} \)
\(\text{CVs = orientBoundaryCycles(model,cells)} \)
\(\text{polygons = [[[V[u] for u in cycle] for cycle in cell] for cell in CVs] \)
\(\text{triangleSet = []} \)
\(\text{for polygon in polygons: triangledPolygon = [] \)
\(\text{externalCycle = polygon[0] } \)
\(\text{polyline = [] } \)
\(\text{for p in externalCycle: polyline.append(Point(p[0],p[1]))} \)
\(\text{cells} \)
\(\text{polyline.append(Point(p[0],p[1]))} \)
\(\text{orients} \)
\(\text{cells} \)
\(\text{cells,bridgeEdges = connectTheDots(model) } \)
\(\text{cells} \)
\(\text{cells,bridgeEdges = connectTheDots(model) } \)
\(\text{cells} \)
\(\text{cells} \)
\(\text{cells} \)
\(\text{cells,bridgeEdges = connectTheDots(model) } \)
\(\text{cells} \)
\(\
```

```
cdt = CDT(polyline)

internalCycles = polygon[1:]
for cycle in internalCycles:
    hole = []
    for p in cycle:
        hole.append(Point(p[0],p[1]))
    cdt.add_hole(hole)

triangles = cdt.triangulate()
trias = [ [[t.a.x,t.a.y,0],[t.c.x,t.c.y,0],[t.b.x,t.b.y,0]] for t in triangles ]
triangleSet += [AA(REVERSE)(trias)]
return triangleSet
```

2.8 Monocyclic polygons using bridge-edges

This subsection, even if correct, is not currently used, since the used triangulation algorithm, from the poly2tri package, cannot handle repeated vertices.

Generation of 1-boundaries as vertex permutation In algebraic topology a k-cycle is a k-chain whose boundary is empty. Also, an unconnected k-cycle is the direct sum of two or more k-cycles. A good formal representation of every simplicial k-cycle, where each component k-simplex has k+1 (k-1)-adjacent k-simplices is a (k+1)-array, indexed by k-simplices, i.e. the Winged Representation [PBCF93].

In the case of oriented 1-cycles, a good representation is given by considering the (ordering of) 0-faces (vertices) as a permutation of n integers, i.e. as a bijective function $\pi:[0,n]\to[0,n]$ that can be represented as an array verts of integers indexed on integers, and the 1-faces (edges) as a dictionary (mapping) nextVert $verts \to verts$.

In order to join two component cycles using one of bridgeEdges, say (u, v), computed by the function connectTheDots previously given, we must save $\pi(u)$ and $\pi(v)$, say, within x and y, respectively

```
⟨Generation of 1-boundaries as vertex permutation 21a⟩ ≡

""" Generation of 1-boundaries as vertex permutation """

def boundaryCycles2vertexPermutation( model ):

    V,EV = model
    cells,bridgeEdges = connectTheDots(model)

    CVs = orientBoundaryCycles(model,cells)

    verts = CAT(CAT( CVs ))
    n = len(verts)
    W = copy.copy(V)
    assert len(verts) == sorted(verts)[n-1]-sorted(verts)[0]+1
```

Wire-frame LAR to boundary polygons The 2-dimensional LAR model (W,EW) returned by the function boundaryCycles2vertexPermutation is pretty special, since it is at the same time both a standard LAR model, i.e. a pair (vertices, edges_by_vertices), and a permutation of vertex indices providing implicitly the ordered cycles on the boundary of a 2-complex. In other words, EW is both a (possibly unconnected) 1-cycle and an 1-boundary.

In the following script we extract the list of connected boundaries (including bridge-edges) from the EW permutation of the first m integers.

```
\langle Wire-frame LAR to boundary polygons 21b \rangle \equiv
     """ lar2boundaryPolygons """
     def lar2boundaryPolygons(model):
         W,EW = boundaryCycles2vertexPermutation( model )
         EW = AA(list)(EW)
         polygons = []
         for k,edge in enumerate(EW):
              polygon = []
              if edge[0]>=0:
                  first = edge[0]
                  done = False
              while (not done) and edge[0] >= 0:
                  polygon += [edge[0]]
                  edge[0] = -edge[0]
                  edge = EW[edge[1]]
                  if len(polygon)>1 and polygon[-1] == first:
                      EW[first][0] = -float(first)
                      break
              if polygon != []:
                  if polygon[0] == polygon[-1]: polygon=polygon[:-1]
                  polygons += [polygon]
         return W, polygons
```

 \Diamond

Macro referenced in 34b.

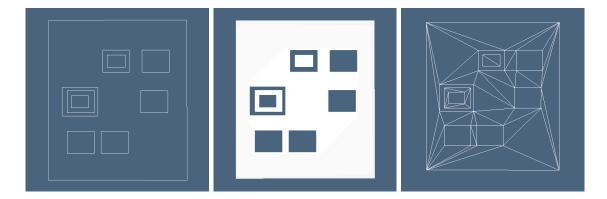


Figure 4: (a) Closed 1-chain c, i.e. such that $\partial c = 0$, that it is both cycle and boundary at the same time; (b) corresponding 2-chain $h = \partial c$; (c) 1-skeleton of a triangulation of h.

2.9 Transformation of an array of lines in a 2D LAR complex

Transformation of an array of lines in a 2D LAR complex The whole transformation of an array of lines into a two-dimensional LAR complex is executed by the function larFromLines. The function returns the model triple V,FV,EV. The last element in FV is the *ordered* boundary chain. just notice that

 \langle Transformation of an array of lines in a 2D LAR complex 23a \rangle \equiv

```
""" Transformation of an array of lines in a 2D LAR complex """

def larFromLines(lines):
    V,EV = lines2lar(lines)
    #VIEW(EXPLODE(1.2,1.2,1.2)(MKPOLS((V,EV))))
    V,EVs = biconnectedComponent((V,EV))
    if EVs != []:
        EV = list(set(AA(tuple)(AA(sorted)(max(EVs, key=len))))) ## NB
        V,EV = larRemoveVertices(V,EV)
        V,FV,EV = facesFromComps((V,EV))
        areas = integr.surfIntegration((V,FV,EV))
        orderedFaces = sorted([[area,FV[f]]] for f,area in enumerate(areas)])
        interiorFaces = [face for area,face in orderedFaces[:-1]]
        return V,interiorFaces,EV
    else: return None
```

2.10 Pruning LAR models from parts out of proper resolution

Pruning of clusters of too close vertices is executed by taking a LAR model as input, executing the following computations, and producing a new simplified LAR model.

Pruning away clusters of close vertices First, reduce the array of vertices pts to its quotient set with respect to the transitive closure of the relation of "nearness". Two vertices are "near" when their (Euclidean) distance is less than a given RADIUS. The subgraphs of the graph of this relation are contracted in a single point, set to the centroid of the vertices of the subgraph. The function W takes as input the array pts of vertex points, and returns: (a) the array newV of new vertices; (b) the list of lists close of sorted indices of pairs of close vertices, removed from duplicates; (c) the list of clusters of pts indices; (d) the integer vmap array, mapping old vertex indices to new vertex indices.

```
\langle Pruning away clusters of close vertices 23b \rangle \equiv
     """ Pruning away clusters of close vertices """
     from scipy.spatial import cKDTree
     def pruneVertices(pts,radius=0.001):
         tree = cKDTree(pts)
         a = cKDTree.sparse_distance_matrix(tree,tree,radius)
         print a.keys()
         close = list(set(AA(tuple)(AA(sorted)(a.keys()))))
         import networkx as nx
         G=nx.Graph()
         G.add_nodes_from(range(len(pts)))
         G.add_edges_from(close)
         clusters, k, h = [], 0, 0
         subgraphs = list(nx.connected_component_subgraphs(G))
         V = [None for subgraph in subgraphs]
         vmap = [None for k in xrange(len(pts))]
         for k,subgraph in enumerate(subgraphs):
             group = subgraph.nodes()
             if len(group)>1:
                  V[k] = CCOMB([pts[v] for v in group])
                  for v in group: vmap[v] = k
                  clusters += [group]
             else:
                  oldNode = group[0]
                  V[k] = pts[oldNode]
                  vmap[oldNode] = k
         return V, close, clusters, vmap
```

Export a simplified LAR model Next, update the arrays of compressed characteristic matrices of a Linear Algebraic Representation. The standard approach is to read row-wise the arrays of matrices of incidence of cells on vertices; translate every index using the vmap array, mapping old vertex indices to new ones; remove repeated indices and substitute them with a single instance; check if the new index list has length greater or equal to the number of vertices of the simplex of the proper dimension. Finally, write an output cell if and only if the previous test is true.

```
\langle \text{Return a simplified LAR model 24} \rangle \equiv
     """ Return a simplified LAR model """
     def larSimplify(model,radius=0.001):
         if len(model) == 2: V, CV = model
         elif len(model)==3: V,CV,FV = model
         else: print "ERROR: model input"
         W,close,clusters,vmap = pruneVertices(V,radius)
         celldim = DIM(MKPOL([V,[[v+1 for v in CV[0]]],None]))
         newCV = [list(set([vmap[v] for v in cell])) for cell in CV]
         CV = list(set([tuple(sorted(cell)) for cell in newCV if len(cell) >= celldim+1]))
         CV = sorted(CV,key=len) # to get the boundary cell as last one (in most cases)
         if len(model)==3:
             celldim = DIM(MKPOL([V,[[v+1 for v in FV[0]]],None]))
             newFV = [list(set([vmap[v] for v in facet])) for facet in FV]
             FV = [facet for facet in newFV if len(facet) >= celldim]
             FV = list(set(AA(tuple)(AA(sorted)(FV))))
             return W,CV,FV
         else: return W,CV
```

Test of pruning clusters of close vertices Here a list of random 2D points is generated. Then the set of vertices is pruned by updating it to its quotient set with respect to the transitive closure of a relation of "nearness" within an Euclidean distance of given RADIUS. The pruning of vertices is performed by the pruneVertices function, with input the array pts of points. The dictionary vmap

```
"test/py/inters/test13.py" 25a =

""" Test of pruning clusters of close vertices """

from larlib import *

from scipy import rand

from scipy.spatial import cKDTree

POINTS = 1000

RADIUS = 0.01
```

```
pts = [rand(2).tolist() for k in range(POINTS)]
VIEW(STRUCT(AA(MK)(pts)))
V,close,clusters,vmap = pruneVertices(pts,RADIUS)
circles = [T([1,2])(pts[h])(CIRCUMFERENCE(RADIUS)(18)) for h,k in close]
convexes = [JOIN(AA(MK)([pts[v] for v in cluster])) for cluster in clusters]
W = COLOR(CYAN)(STRUCT(AA(MK)(V)))
VIEW(STRUCT(AA(MK)(pts)+AA(COLOR(YELLOW))(circles)))
VIEW(STRUCT(AA(COLOR(RED))(convexes)+AA(MK)(pts)+AA(COLOR(YELLOW))(circles)+[W]))
```

Test for exporting a simplified LAR model

```
"test/py/inters/test14.py" 25b \equiv
     """ Test for exporting a simplified LAR model """
     from larlib import *
     filename = "test/svg/inters/closepoints.svg"
     lines = svg2lines(filename)
     VIEW(STRUCT(AA(POLYLINE)(lines)))
     V,FV,EV = larFromLines(lines)
     VIEW(EXPLODE(1.2,1.2,1)(MKPOLS((V,FV[:-1]+EV)) + AA(MK)(V)))
     VV = AA(LIST)(range(len(V)))
     submodel = STRUCT(MKPOLS((V,EV)))
     VIEW(larModelNumbering(1,1,1)(V,[VV,EV,FV[:-1]],submodel,0.5))
     V.EV = lines2lar(lines)
     VIEW(EXPLODE(1.2,1.2,1)(MKPOLS((V,EV))))
     pts = V
     RADIUS = 0.05
     V,close,clusters,vmap = pruneVertices(pts,RADIUS)
     circles = [T([1,2])(pts[h])(CIRCUMFERENCE(RADIUS)(18)) for h,k in close]
     convexes = [JOIN(AA(MK)([pts[v] for v in cluster])) for cluster in clusters]
     W = COLOR(CYAN)(STRUCT(AA(MK)(V)))
     VIEW(STRUCT(AA(COLOR(RED)))(convexes)+AA(MK)(pts)+AA(COLOR(YELLOW))(circles)+[W]))
```

2.11 Point in polygon classification

```
\begin{split} \langle \, \text{Point in polygon testing 26a} \, \rangle &\equiv \\ & \langle \, \text{Half-line crossing test 28} \, \rangle \\ & \langle \, \text{Tile codes computation 26b} \, \rangle \\ & \langle \, \text{Point-in-polygon classification algorithm 26c} \, \rangle \end{split}
```

Tile codes computation

```
⟨Tile codes computation 26b⟩ ≡
""" Tile codes computation """
def setTile(box):
    tiles = [[9,1,5],[8,0,4],[10,2,6]]
    b1,b2,b3,b4 = box
    def tileCode(point):
        x,y = point
        code = 0
        if y>b1: code=code|1
        if y<b2: code=code|2
        if x>b3: code=code|4
        if x<b4: code=code|8
        return code
    return tileCode
</pre>
```

Macro referenced in 26a.

Point in polygon testing

```
\langle Point-in-polygon classification algorithm 26c\rangle \equiv
     """ Point in polygon classification """
     def pointInPolygonClassification(pol):
         V,EV = pol
         # edge orientation
         FV = [sorted(set(CAT(EV)))]
         orientedCycles = boundaryPolylines(Struct([(V,FV,EV)]))
         EV = []
         for cycle in orientedCycles:
              EV += zip(cycle[:-1],cycle[1:])
         def pointInPolygonClassificationO(p):
              x,y = p
              xmin, xmax, ymin, ymax = x, x, y, y
              tilecode = setTile([ymax,ymin,xmax,xmin])
              count, status = 0,0
              for k,edge in enumerate(EV):
                  p1,p2 = edge[0],edge[1]
                  (x1,y1),(x2,y2) = p1,p2
                  c1,c2 = tilecode(p1),tilecode(p2)
                  c_{edge}, c_{un}, c_{int} = c1^c2, c1|c2, c1&c2
                  if c_edge == 0 and c_un == 0: return "p_on"
```

```
elif c_edge == 12 and c_un == c_edge: return "p_on"
        elif c_edge == 3:
            if c_int == 0: return "p_on"
            elif c_int == 4: count += 1
        elif c_edge == 15:
            x_{int} = ((y-y2)*(x1-x2)/(y1-y2))+x2
            if x_{int} > x: count += 1
            elif x_int == x: return "p_on"
        elif c_edge == 13 and ((c1==4) or (c2==4)):
                crossingTest(1,2,status,count)
        elif c_edge == 14 and (c1==4) or (c2==4):
                crossingTest(2,1,status,count)
        elif c_edge == 7: count += 1
        elif c_edge == 11: count = count
        elif c_edge == 1:
            if c_int == 0: return "p_on"
            elif c_int == 4: crossingTest(1,2,status,count)
        elif c_edge == 2:
            if c_int == 0: return "p_on"
            elif c_int == 4: crossingTest(2,1,status,count)
        elif c_edge == 4 and c_un == c_edge: return "p_on"
        elif c_edge == 8 and c_un == c_edge: return "p_on"
        elif c_edge == 5:
            if (c1==0) or (c2==0): return "p_on"
            else: crossingTest(1,2,status,count)
        elif c_edge == 6:
            if (c1==0) or (c2==0): return "p_on"
            else: crossingTest(2,1,status,count)
        elif c_edge == 9 and ((c1==0) or (c2==0)): return "p_on"
        elif c_edge == 10 and ((c1==0) or (c2==0)): return "p_on"
    if ((round(count)%2)==1): return "p_in"
    else: return "p_out"
return pointInPolygonClassificationO
```

Macro referenced in 26a.

Half-line crossing test

```
⟨Half-line crossing test 28⟩ ≡

""" Half-line crossing test """

def crossingTest(new,old,count,status):
    if status == 0:
        status = new
        count += 0.5

else:
    if status == old: count += 0.5
```

```
else: count -= 0.5 status = 0
```

Macro referenced in 26a.

2.12 Drawing multiply-connected boundary polylines

A multiply-connected boundary polyline gives the boundary edges of a 2-cell. In LAR, 2-cells must be connected, but not necessarily simply-connected (or path-connected, or homotopic to a point). The "solid" drawing of such a generally non-convex 2-cells is not easy. In particular, they must be decomposed into a coherently-oriented simplicial 2-complex, i.e. into a set of equioriented triangles. The library function used at this purpose (), contained in the poly2tria module, only accepts a single boundary polyline. Hence the purpose of this section is to transform a list of cycles, corresponding orderly to the exterior and the interior boundaries of a 2-cell, into a single polyline, using the so-called bridge-edges [YT85] mechanism. The transformation from a set of non-intersecting boundary cycles to a single connected polyline is given in Section 2.7.

Generating the LAR of a set of non-intersecting cycle We use the test/svg/lattice.svg file to test the exporting of different boundary chains. The example test15.py aims to prepare the computational environment for writing down the LAR of the 2-complex generated by any set of non-intersecting 1-cles in 2D.

Let us remember that any d-cell in the domain of the LAR scheme must be connected, but not necessarily contractible to a point, i.e. may contain internal holes. The goal af this test and of the previous one, is hence to generalize the creation of 2-complexes (sets of polygons) starting from several non-intersecting boundary cycles, instead than starting from just one closed polyline.

The motivation arises from situations created by the Boolean algorithm, as well as from the imput of a 2-complex from general wire-frame drawings.

```
"test/py/inters/test16.py" 29 \=
    """ Generating the LAR of a set of non-intersecting cycles """
    from larlib import *

    sys.path.insert(0, 'test/py/inters/')
    from test15 import *

    cells = cellsFromCycles(testArray)
    CV = AA(COMP([list,set,CAT]))(EVs)
    EVdict = dict(zip(EV,range(len(EV))))
    FE = [[EVdict[edge] for edge in cycle] for cycle in EVs]
    edges = [CAT([FE[cycle] for cycle in cell]) for cell in cells]
    FVs = [[CV[cycle] for cycle in cell]]
```

LAR of the 2-complex generated by non-intersecting cycles Therefore, the LAR of the 2-complex generated from the input get from test15.py is (V,FV,EV) with the components given by the the following:

```
V = [[0.7808, 0.6751], [0.6044, 0.6751], [0.319, 0.5804], [0.319, 0.3994], [0.8936, 0.0],
[0.0,0.0],[0.6044,0.8123],[0.7808,0.8123],[0.3495,0.6751],[0.3495,0.8123],
[0.1218, 0.3036], [0.2983, 0.3036], [0.0886, 0.3994], [0.0886, 0.5804], [0.2581, 0.4505],
[0.1495, 0.4505], [0.7717, 0.4213], [0.5952, 0.4213], [0.7717, 0.5585], [0.0, 1.0],
[0.3403, 0.1664], [0.3403, 0.3036], [0.5952, 0.5585], [0.5168, 0.1664], [0.5168, 0.3036],
[0.5259, 0.8123], [0.1495, 0.5293], [0.5259, 0.6751], [0.8936, 1.0], [0.2983, 0.1664],
[0.1218, 0.1664], [0.2581, 0.5293], [0.4965, 0.7815], [0.4965, 0.7059], [0.2983,
0.5585], [0.2983, 0.4213], [0.1218, 0.4213], [0.3789, 0.7815], [0.1218, 0.5585],
[0.3789,0.7059]]
FV = [[0,1,2,3,4,5,6,7,8,9,10,11,12,13,16,17,18,19,20,21,22,23,24,25,27,
28,29,30],[32,33,37,39],[34,35,36,38,26,15,14,31]]
EV = [(0,1),(0,7),(1,6),(2,3),(2,13),(3,12),(4,5),(4,28),(5,19),(6,7),(8,9),
(8,27),(9,25),(10,11),(10,30),(11,29),(12,13),(14,15),(14,31),(15,26),
(16,17),(16,18),(17,22),(18,22),(19,28),(20,21),(20,23),(21,24),(23,24),
(25,27),(26,31),(29,30),(32,33),(32,37),(33,39),(34,35),(34,38),(35,36),
(36,38),(37,39)
```

Of course, this LAR representation gives full control of the complex topology, including k-(co)chains and (co)boundary operators. For example, in Figure 5 we show a solid image of the three 2-cells in FV, and, drawn in red the boundary 1-cells in the complex, corresponding to the 2-chains of coordinates [1,0,0], [0,1,0], [0,0,1], and [1,1,0], respectively.

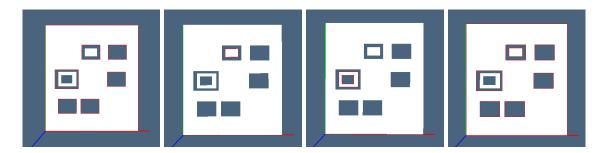


Figure 5: Some examples of boundaries (in red) of 2-chains from nested non intersecting cycles.

2.13 Boundary polylines

Generating the LAR of a set of non-intersecting cycles The example provided by test/py/inters/test17.py is completed here by showing the solid drawing of the generated LAR data structure, and superimposing to it the boundary 1-chains generated by several 2-chains.

```
"test/py/inters/test17.py" 31a \equiv
     """ Generating the LAR of a set of non-intersecting cycles """
     from larlib import *
     sys.path.insert(0, 'test/py/inters/')
     from test16 import *
     lar = (V, FV, EV)
     bcycles = boundaryCycles(range(len(EV)),EV)
     polylines = [[V[EV[e][1]] if e>0 else V[EV[-e][0]] for e in cycle ] for cycle in bcycles]
     polygons = [polyline + [polyline[0]] for polyline in polylines]
     complex = SOLIDIFY(STRUCT(AA(POLYLINE)(polygons)))
     for chain in chains:
         chainBoundary = COLOR(RED)(STRUCT(MKPOLS((V,[EV[e]
                             for e in chain2BoundaryChain(FV,EV)(chain)])))
         VIEW(STRUCT([complex, chainBoundary]))
"test/py/inters/test18.py" 31b \equiv
     """ Orienting a set of non-intersecting cycles """
     from larlib import *
     sys.path.insert(0, 'test/py/inters/')
     from test17 import *
```

```
cells,bridgeEdges = connectTheDots((V,EV))
     CVs = orientBoundaryCycles((V,EV),cells)
     print "\nCVs =",CVs
"test/py/inters/test19.py" 31c \equiv
     """ Generating the LAR of a set of non-intersecting cycles """
     from larlib import *
     sys.path.insert(0, 'test/py/inters/')
     from test17 import *
     W,EW = boundaryCycles2vertexPermutation( (V,EV) )
     VIEW(EXPLODE(1.2,1.2,1.2)(MKPOLS((W,EW))))
"test/py/inters/test20.py" 31d \equiv
     """ Generating the Triangulation of a set of non-intersecting cycles """
     from larlib import *
     sys.path.insert(0, 'test/py/inters/')
     from test17 import *
     triangleSet = boundaryCycles2triangulation( (V,EV) )
     VIEW(STRUCT(AA(JOIN)(AA(AA(MK))(CAT(triangleSet)))))
     VIEW(SKEL_1(STRUCT(AA(JOIN)(AA(AA(MK))(CAT(triangleSet))))))
     11 11 11
     model = V,EV
     W,FW = lar2boundaryPolygons(model)
     polygons = [[W[u] for u in poly] for poly in FW]
     VIEW(STRUCT(AA(POLYLINE)(polygons)))
     triangleSet,triangledFace = [],[]
     for polygon in polygons:
         polyline = [Point(p[0],p[1]) for p in polygon]
         cdt = CDT(polyline)
         triangles = cdt.triangulate()
         trias = [[[t.c.x,t.c.y,0,1],[t.b.x,t.b.y,0,1],[t.a.x,t.a.y,0,1]] for t in triangles ]
         triangledFace += [[v[:-1] for v in triangle] for triangle in trias]
         triangleSet += [triangledFace]
     .....
```

From structures to boundary polylines

From LAR oriented boundary model to polylines

```
\langle From LAR boundary model to polylines 32b\rangle \equiv
     """ From LAR oriented boundary model to polylines """
     def boundaryModel2polylines(model):
         V, EV = model
         polylines = []
         succDict = dict(EV)
         visited = [False for k in range(len(V))]
         nonVisited = [k for k in succDict.keys() if not visited[k]]
         while nonVisited != []:
             first = nonVisited[0]; v = first; polyline = []
             while visited[v] == False:
                  visited[v] = True;
                  polyline += V[v],
                 v = succDict[v]
             polyline += [V[first]]
             polylines += [polyline]
             nonVisited = [k for k in succDict.keys() if not visited[k]]
         return polylines
```

From Struct object to LAR boundary model

```
⟨From Struct object to LAR boundary model 33⟩ ≡
    """ From Struct object to LAR boundary model """
    def structFilter(obj):
        if isinstance(obj,list):
            if (len(obj) > 1):
                return [structFilter(obj[0])] + structFilter(obj[1:])
            return [structFilter(obj[0])]
        if isinstance(obj,Struct):
            if obj.category in ["external_wall", "internal_wall", "corridor_wall"]:
```

```
return Struct(structFilter(obj.body),obj.name,obj.category)
   return obj
def structBoundaryModel(struct):
   filteredStruct = structFilter(struct)
   #import pdb; pdb.set_trace()
   V,FV,EV = struct2lar(filteredStruct)
   edgeBoundary = boundaryCells(FV,EV)
   cycles = boundaryCycles(edgeBoundary,EV)
   edges = [signedEdge for cycle in cycles for signedEdge in cycle]
   orientedBoundary = [ AA(SIGN)(edges), AA(ABS)(edges)]
   cells = [EV[e] if sign==1 else REVERSE(EV[e]) for (sign,e) in zip(*orientedBoundary)]
    if cells[0][0] == cells[1][0]: # bug badly patched! ... TODO better
        temp0 = cells[0][0]
        temp1 = cells[0][1]
        cells[0] = [temp1, temp0]
   return V, cells
```

Edge cycles associated to a closed chain of edges The output cycles are returned as lists of oriented edges, given in consecutive sequence in each list. Each cycle is given in the opposite ordering of its first edge. The first edge of each cycle is the one of minimum index in the cycle. The list of output cycles is returned ordered for increasing (or better, in decreasing order, since negative) order of its first element. The first output cycle starts from index 0, that cannot oriented directly, since -0 is not allowed for integer indices. It must be considered as negative, i.e. as oriented as the opposite of its canonical orientation.

```
e = list(set(verts2edges[v]).difference([e]))[0]
    cbe.remove(e)
    v = list(set(EV[e]).difference([v]))[0]
    if (e,v)==cycle[0]:
        break
    n = len(cycle)
    cycles += [[e if EV[e]==(cycle[(k-1)%n][1],cycle[k%n][1]) else -e
        for k,(e,v) in enumerate(cycle)]]
    return cycles
```

3 Exporting the module

```
"larlib/larlib/inters.py" 34b \equiv
     """ Module for pipelined intersection of geometric objects """
     from larlib import *
     from scipy import mat
     DEBUG = True
      ⟨ Coding utilities 46b⟩
      (Generation of random lines 47a)
      Containment boxes 2a
      (Splitting the input above and below a threshold 2b)
      Box metadata computation?
      (Iterative splitting of box buckets 3a)
      (Intersection of two line segments 3b)
      Brute force bucket intersection 4
      (Accelerate intersection of lines 5)
       Edge cycles associated to a closed chain of edges 34a
      From Struct object to LAR boundary model 33
      From structures to boundary polylines 32a
      From LAR boundary model to polylines 32b
       Point in polygon testing 26a
       Classification of non intersecting cycles 14 \
      Extraction of path-connected boundaries 15a
       Scan line algorithm 17b
       Scan line algorithm input/output 18 >
       Orientation of component cycles of unconnected boundaries 19)
       From nested boundary cycles to triangulation 20 \
       Generation of 1-boundaries as vertex permutation 21a
       Wire-frame LAR to boundary polygons 21b
       Create the LAR of fragmented lines 6
      (Biconnected components 7a)
      (Slope of edges 9)
```

```
\label{eq:condition} $$ \langle$\mbox{ Ordered incidence relationship of vertices and edges $11a$} \rangle$$ $$ \langle$\mbox{ Faces from biconnected components $11b$} \rangle$$ $$ \langle$\mbox{SVG input parsing and transformation $43$} \rangle$$ $$ \langle$\mbox{ Transformation of an array of lines in a 2D LAR complex $23a$} \rangle$$ $$ \langle$\mbox{ Pruning away clusters of close vertices $23b$} \rangle$$ $$ \langle$\mbox{ Return a simplified LAR model $24$} \rangle$$
```

4 Examples

Generation of random line segments and their boxes

```
"test/py/inters/test01.py" 35 ==
    """ Generation of random line segments and their boxes """
    from larlib import *

    randomLineArray = randomLines(200,0.3)
    VIEW(STRUCT(AA(POLYLINE)(randomLineArray)))

    boxes = containment2DBoxes(randomLineArray)
    rects= AA(box2rect)(boxes)
    cyan = COLOR(CYAN)(STRUCT(AA(POLYLINE)(randomLineArray)))
    yellow = COLOR(YELLOW)(STRUCT(AA(POLYLINE)(rects)))
    VIEW(STRUCT([cyan,yellow]))
```

Split segment array in four independent buckets

```
"test/py/inters/test02.py" 36a \(\text{ } \)
    """ Split segment array in four independent buckets """
    from larlib import *

    randomLineArray = randomLines(200,0.3)
    VIEW(STRUCT(AA(POLYLINE)(randomLineArray)))
    boxes = containment2DBoxes(randomLineArray)
    bucket = range(len(boxes))
    below,above = splitOnThreshold(boxes,bucket,1)
    below1,above1 = splitOnThreshold(boxes,above,2)
    below2,above2 = splitOnThreshold(boxes,below,2)

    cyan = COLOR(CYAN)(STRUCT(AA(POLYLINE)(randomLineArray[k] for k in below1)))
    yellow = COLOR(YELLOW)(STRUCT(AA(POLYLINE)(randomLineArray[k] for k in above1)))
    red = COLOR(RED)(STRUCT(AA(POLYLINE)(randomLineArray[k] for k in below2)))
    green = COLOR(GREEN)(STRUCT(AA(POLYLINE)(randomLineArray[k] for k in above2)))

    VIEW(STRUCT([cyan,yellow,red,green]))
```

Generation and random coloring of independent line buckets

```
"test/py/inters/test03.py" 36b \equiv
     """ Generation and random coloring of independent line buckets """
     from larlib import *
     lines = randomLines(200,0.3)
     VIEW(STRUCT(AA(POLYLINE)(lines)))
     boxes = containment2DBoxes(lines)
     buckets = boxBuckets(boxes)
     colors = [CYAN, MAGENTA, WHITE, RED, YELLOW, GRAY, GREEN, ORANGE, BLACK, BLUE, PURPLE, BROWN]
     sets = [COLOR(colors[k%12])(STRUCT(AA(POLYLINE)([lines[h]
                 for h in bucket]))) for k,bucket in enumerate(buckets) if bucket!=[]]
     VIEW(STRUCT(sets))
Construction of LAR = (V,EV) of random line arrangement
"test/py/inters/test04.py" 37a \equiv
     """ LAR of random line arrangement """
     from larlib import *
     lines = randomLines(300,0.2)
     VIEW(STRUCT(AA(POLYLINE)(lines)))
     intersectionPoints,params,frags = lineIntersection(lines)
     marker = CIRCLE(.005)([4,1])
     markers = STRUCT(CONS(AA(T([1,2]))(intersectionPoints))(marker))
     VIEW(STRUCT(AA(POLYLINE)(lines)+[COLOR(RED)(markers)]))
     V,EV = lines2lar(lines)
     marker = CIRCLE(.01)([4,1])
     markers = STRUCT(CONS(AA(T([1,2]))(V))(marker))
     #markers = STRUCT(CONS(AA(T([1,2]))(intersectionPoints))(marker))
     polylines = STRUCT(MKPOLS((V,EV)))
     VIEW(STRUCT([polylines]+[COLOR(MAGENTA)(markers)]))
```

Splitting of othogonal lines

[&]quot;test/py/inters/test05.py" $37b \equiv$

```
""" LAR from splitting of othogonal lines """
from larlib import *
    ⟨Orthogonal example 37c⟩
    ⟨Orthogonal example 37c⟩

⟨Orthogonal example 37c⟩

lines = [[[0,0],[6,0]], [[0,4],[10,4]], [[0,0],[0,4]], [[3,0],[3,4]], [[6,0],[6,8]], [[3,2],[6,2]], [[10,0],[10,8]], [[0,8],[10,8]]]

VIEW(EXPLODE(1.2,1.2,1)(AA(POLYLINE)(lines)))

V,EV = lines2lar(lines)
VIEW(EXPLODE(1.2,1.2,1)(MKPOLS((V,EV))))
    ⟨
```

Macro referenced in 37b, 38b, 40a.

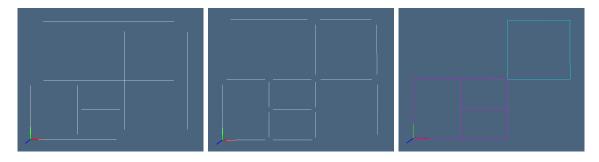


Figure 6: Splitting of orthogonal lines: (a) exploded input; (a) exploded output; (c) biconnected components.

Random coloring of the generated 1-complex LAR

```
"test/py/inters/test06.py" 38a =
    """ Random coloring of the generated 1-complex """
    from larlib import *

lines = randomLines(800,0.2)
    VIEW(STRUCT(AA(POLYLINE)(lines)))

V,EV = lines2lar(lines)
    colors = [CYAN, MAGENTA, WHITE, RED, YELLOW, GRAY, GREEN, ORANGE, BLACK, BLUE, PURPLE, BROWN]
    sets = [COLOR(colors[k%12])(POLYLINE([V[e[0]],V[e[1]]])) for k,e in enumerate(EV)]

VIEW(STRUCT(sets))
```



Figure 7: Splitting of intersecting lines: (a) random input; (a) splitted and colored LAR output.



Figure 8: The intersection of 5000 random lines in the unit interval, with ${\tt scaling}$ parameter equal to ${\tt 0.1}$

Biconnected components from orthogonal LAR model

```
"test/py/inters/test07.py" 38b \equiv
     """ Biconnected components from orthogonal LAR model """
     from larlib import *
     colors = [CYAN, MAGENTA, WHITE, RED, YELLOW, GREEN, ORANGE, BLACK, BLUE, PURPLE]
     \langle \text{Orthogonal example } 37c \rangle
     model = V,EV
     V,EVs = biconnectedComponent(model)
     HPCs = [STRUCT(MKPOLS((V,EV))) for EV in EVs]
     sets = [COLOR(colors[k%10])(hpc) for k,hpc in enumerate(HPCs)]
     VIEW(STRUCT(sets))
     VIEW(STRUCT(MKPOLS((V,CAT(EVs)))))
     #V,EV = larRemoveVertices(V,CAT(EVs))
2-complex from orthogonal line segments
```

```
"test/py/inters/test08.py" 40a \equiv
     """ 2-complex from orthogonal line segments """
     from larlib import *
     colors = [CYAN, MAGENTA, WHITE, RED, YELLOW, GREEN, ORANGE, BLACK, BLUE, PURPLE]
     \langle \text{Orthogonal example } 37c \rangle
     model = V,EV
     V,EVs = biconnectedComponent(model)
     HPCs = [STRUCT(MKPOLS((V,EV))) for EV in EVs]
     sets = [COLOR(colors[k%10])(hpc) for k,hpc in enumerate(HPCs)]
     VIEW(STRUCT(sets))
     EV = sorted(CAT(EVs))
     VIEW(STRUCT(MKPOLS((V,EV))))
     V,FV,EV = facesFromComps((V,EV))
     areas = surfIntegration((V,FV,EV))
     boundaryArea = max(areas)
     FV = [FV[f] for f,area in enumerate(areas) if area!=boundaryArea]
     VIEW(EXPLODE(1.2,1.2,1)(MKPOLS((V,FV+EV)) + AA(MK)(V)))
```

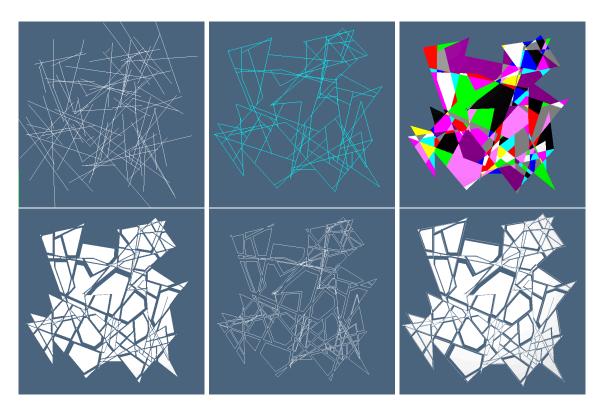


Figure 9: LAR complex generation random lines. (a) the input random lines; (b) maximal biconnected graph extracted from the 1D LAR of intersected lines; (c) 2D cells of such *regularized* 2-complex; (d) 2-cells, drawn exploded; (e) boundaries of 2D cells; (f) regularized cellular 2-complex extracted from lines.

Biconnected components from random LAR model

```
"test/py/inters/test09.py" 40b \equiv
     """ Biconnected components from orthogonal LAR model """
     from larlib import *
     colors = [CYAN, MAGENTA, YELLOW, RED, GREEN, ORANGE, PURPLE, WHITE, BLACK, BLUE]
     lines = randomLines(100,.8)
     V,EV = lines2lar(lines)
     model = V,EV
     VIEW(STRUCT(AA(POLYLINE)(lines)))
     V, EVs = biconnectedComponent(model)
     HPCs = [STRUCT(MKPOLS((V,EV))) for EV in EVs]
     sets = [COLOR(colors[k%10])(hpc) for k,hpc in enumerate(HPCs)]
     VIEW(STRUCT(sets))
     EV = CAT(EVs)
     V,EV = larRemoveVertices(V,EV)
     V,FV,EV = facesFromComps((V,EV))
     areas = surfIntegration((V,FV,EV))
     boundaryArea = max(areas)
     FV = [FV[f] for f, area in enumerate(areas) if area!=boundaryArea]
     polylines = [[V[v] for v in face+[face[0]]] for face in FV]
     VIEW(EXPLODE(1.2,1.2,1)(MKPOLS((V,EV)) + AA(MK)(V) + AA(FAN)(polylines) ))
     colors = [CYAN, MAGENTA, WHITE, RED, YELLOW, GRAY, GREEN, ORANGE, BLACK, BLUE, PURPLE, BROWN]
     sets = [COLOR(colors[k%12])(FAN(pol)) for k,pol in enumerate(polylines)]
     VIEW(STRUCT(sets))
     VIEW(EXPLODE(1.2,1.2,1)((AA(FAN)(polylines))))
     VIEW(EXPLODE(1.2,1.2,1)((AA(POLYLINE)(polylines))))
     VV = AA(LIST)(range(len(V)))
     submodel = STRUCT(MKPOLS((V,EV)))
     VIEW(larModelNumbering(1,1,1)(V,[VV,EV],submodel,0.1))
     \Diamond
```

SVG input parsing and transformation We postulate here that the input file test/py/inters/test.svg should contain only primitives, so we skip any other content. Such primitives are parsed by matching against regular expressions, and their x1,y1,x2,y2 attributes are extracted and stored into the lines variable. An isomorphic window-viewport transformation is then performed, to transform the data within the standard unit 2D square [0, 1]². The input vertices are finally set to a fixed resolution, using the vcode function.

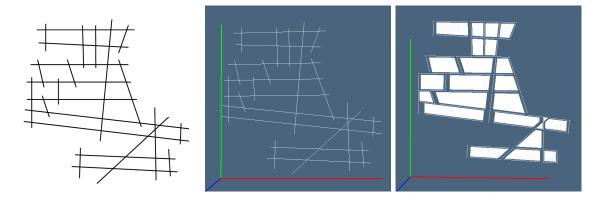


Figure 10: LAR complex generation from SVG file. (a) the input set of lines; (b) imported in pyplasm environment; (c) the extracted regularized 2-complex, drawn exploded.

```
\langle SVG input parsing and transformation 43\rangle \equiv
    """ SVG input parsing and transformation """
    from larlib import *
    import re # regular expression
    def svg2lines(filename,containmentBox=[],rect2lines=True):
       stringLines = [line.strip() for line in open(filename)]
       # SVG <line> primitives
       lines = [string.strip() for string in stringLines if re.match("<line ",string)!=None]</pre>
       outLines = ""
       for line in lines:
           if searchObj:
               outLines += "[["+searchObj.group(4)+","+searchObj.group(6)+"], ["+searchObj.group(
       if lines != []:
           lines = list(eval(outLines))
       # SVG <rect> primitives
       rects = [string.strip() for string in stringLines if re.match("<rect ",string)!=None]</pre>
       outRects,searchObj = "",False
       for rect in rects:
           if searchObj:
               outRects += "[["+searchObj.group(2)+","+searchObj.group(4)+"], ["+searchObj.group(
       if rects != []:
           rects = list(eval(outRects))
           if rect2lines:
```

Macro referenced in 34b.

Macro referenced in 43.

SVG input normalization transformation The normalization transformation maps the input lines to the $[0,1]^2$ viewport, i.e. to the standard unit square.

```
\langle SVG input normalization transformation 44a\rangle \equiv
     """ SVG input normalization transformation """
     # window-viewport transformation
     xs,ys = TRANS(CAT(lines))
     box = [min(xs), min(ys), max(xs), max(ys)]
     # viewport aspect-ratio checking, setting a computed-viewport 'b'
     b = [None for k in range(4)]
     if (box[2]-box[0])/(box[3]-box[1]) > 1:
         b[0]=0; b[2]=1; bm=(box[3]-box[1])/(box[2]-box[0]); b[1]=.5-bm/2; b[3]=.5+bm/2
     else:
         b[1]=0; b[3]=1; bm=(box[2]-box[0])/(box[3]-box[1]); b[0]=.5-bm/2; b[2]=.5+bm/2
     # isomorphic 'box -> b' transform to standard unit square
     lines = [[[
     ((x1-box[0])*(b[2]-b[0]))/(box[2]-box[0]),
     ((y1-box[1])*(b[3]-b[1]))/(box[1]-box[3]) + 1], [
     ((x2-box[0])*(b[2]-b[0]))/(box[2]-box[0]),
     ((y2-box[1])*(b[3]-b[1]))/(box[1]-box[3]) + 1]]
           for [[x1,y1],[x2,y2]] in lines]
     # line vertices set to fixed resolution
     lines = eval("".join(['['+ vcode(p1) +','+ vcode(p2) +'], ' for p1,p2 in lines]))
```

2-complex extraction from svg file The input lines arrangments produces a 1-dimensional complex stored into the LAR model V, EV. Then the *dangling edges* are removed from EV_, and the whole data set is renumbered, in order to remove the unused vertices,

using the larRemoveVertices function. Finally the 2-cells are computed and stored in FV, and the positive areas of every 2cells are computed, so allowing for identify and removal of the exterior face, corresponding to the boundary of the complex. The polygonal boundary of the complex is finally drawn.

```
"test/py/inters/test10.py" 44b =
    """ Biconnected components from orthogonal LAR model """
    from larlib import *

filename = "test/py/inters/plan.svg"
    #filename = "test/py/inters/building.svg"
    #filename = "test/py/inters/complex.svg"
    lines = svg2lines(filename)
    VIEW(STRUCT(AA(POLYLINE)(lines)))

V,FV,EV = larFromLines(lines)
    VIEW(EXPLODE(1.2,1.2,1)(MKPOLS((V,FV[:-1]+EV)) + AA(MK)(V)))

VV = AA(LIST)(range(len(V)))
    submodel = STRUCT(MKPOLS((V,EV)))
    VIEW(larModelNumbering(1,1,1)(V,[VV,EV,FV[:-1]],submodel,0.05))

verts,faces,edges = polyline2lar([[ V[v] for v in FV[-1] ]])
    VIEW(STRUCT(MKPOLS((verts,edges))))
```

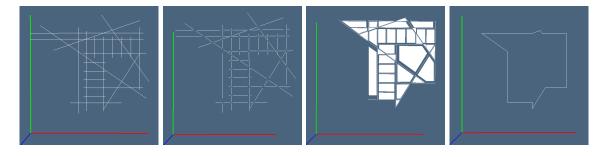


Figure 11: LAR complex generation from SVG file. (a) the input set of lines parsed from an SVG file; (b) the intersection of lines; (c) the extracted *regularized* 2-complex, drawn exploded; (d) the boundary LAR.

```
"test/py/inters/test11.py" 45 \equiv """ Fast Polygon Triangulation based on Seidel's Algorithm """ # data generated by test10.py on file polygon.svg
```

```
from larlib import *
     V, FV, EV = ([[0.222, 0.889],
       [0.722, 1.0],
       [0.519, 0.763],
       [1.0, 0.659],
       [0.859, 0.233],
       [0.382, 0.119],
       [0.519, 0.348],
       [0.296, 0.53],
       [0.0, 0.059]],
      [[0, 1, 2, 3, 4, 5, 6, 7, 8]],
      [[2, 3], [6, 7], [0, 8], [3, 4], [1, 2], [7, 8], [4, 5], [5, 6], [0, 1]])
     VV = AA(LIST)(range(len(V)))
     submodel = STRUCT(MKPOLS((V,EV)))
     VIEW(larModelNumbering(1,1,1)(V,[VV,EV],submodel,0.5))
     xord = TRANS(sorted(zip(V,range(len(V)))))[1]
     trapezoids = zip(xord[:-1],xord[1:])
     vert2forw_trap = dict()
     vert2back_trap = dict()
     for k,(a,b) in enumerate(trapezoids[1:-1]):
         print k, (a,b)
         vert2back_trap[a]=k
         vert2forw_trap[a]=k+1
         vert2back_trap[b]=k+1
         vert2forw_trap[b]=k+2
     vert2forw_trap[trapezoids[0][0]] = 0
     vert2back_trap[trapezoids[-1][1]] = len(trapezoids)-1
"test/py/inters/test12.py" 46a \equiv
     """ Biconnected components from orthogonal LAR model """
     from larlib import *
     V = [[0.395, 0.296], [0.593, 0.0], [0.79, 0.773], [0.671, 0.889], [0.79, 0.0], [0.593, 0.296],
     FV = [[0, 5, 4, 1], [1, 9, 0], [8, 7, 0, 9], [7, 8, 3, 2, 4, 5, 6]]
     EV = [[0, 1], [8, 9], [6, 7], [4, 5], [1, 4], [3, 8], [5, 6], [2, 3], [1, 9], [0, 9], [0, 5],
     polylines = [[V[v] for v in face+[face[0]]] for face in FV]
      VIEW(EXPLODE(1.1,1.1,1)(MKPOLS((V,EV)) + AA(MK)(V) + AA(FAN)(polylines))) 
     VV = AA(LIST)(range(len(V)))
     submodel = STRUCT(MKPOLS((V,EV)))
     VIEW(larModelNumbering(1,1,1)(V,[VV,EV,FV],submodel,.6))
```

```
VIEW(EXPLODE(1.1,1.1,1)(AA(POLYLINE)(polylines)))
```

A Code utilities

Coding utilities Some utility functions used by the module are collected in this appendix. Their macro names can be seen in the below script.

```
⟨Coding utilities 46b⟩ ≡

""" Coding utilities """

⟨Generation of a random point 47b⟩

⟨Generation of a random line segment 47c⟩

⟨Transformation of a 2D box into a closed polyline 48a⟩

⟨Computation of the 1D centroid of a list of 2D boxes 48b⟩

⟨Pyplasm XOR of FAN of ordered points 49⟩

⋄
```

Macro referenced in 34b.

Generation of random lines The function randomLines returns the array randomLineArray with a given number of lines generated within the unit 2D interval. The scaling parameter is used to scale every such line, generated by two randow points, that could be possibly located to far from each other, even at the distance of the diagonal of the unit square.

The arrays xs and ys, that contain the x and y coordinates of line points, are used to compute the minimal translation v needed to transport the entire set of data within the positive quadrant of the 2D plane.

```
⟨Generation of random lines 47a⟩ ≡
""" Generation of random lines """
def randomLines(numberOfLines=200,scaling=0.3):
    randomLineArray = [redge(scaling) for k in range(numberOfLines)]
    [xs,ys] = TRANS(CAT(randomLineArray))[:2]
    xmin, ymin = min(xs), min(ys)
    v = array([-xmin,-ymin])
    randomLineArray = [[list(v1[:2]+v), list(v2[:2]+v)] for v1,v2 in randomLineArray]
    return randomLineArray
```

Macro referenced in 34b.

Generation of a random point A single random point, codified in floating point format, and with a fixed (quite small) number of digits, is returned by the rpoint2d() function, with no input parameters.

Generation of a random line segment A single random segment, scaled about its centroid by the scaling parameter, is returned by the redge() function, as a tuple of two random points in the unit square.

```
⟨Generation of a random line segment 47c⟩ ≡
    """ Generation of a random line segment """
    def redge(scaling):
        v1,v2 = array(rpoint2d()), array(rpoint2d())
        c = (v1+v2)/2
        pos = rpoint2d()
        v1 = (v1-c)*scaling + pos
        v2 = (v2-c)*scaling + pos
        return tuple(eval(vcode(v1))), tuple(eval(vcode(v2)))
        ◊
Macro referenced in 46b.
```

Transformation of a 2D box into a closed polyline The transformation of a 2D box into a closed rectangular polyline, given as an ordered sequence of 2D points, is produced by the function box2rect

```
⟨Transformation of a 2D box into a closed polyline 48a⟩ ≡
    """ Transformation of a 2D box into a closed polyline """
    def box2rect(box):
        x1,y1,x2,y2 = box
        verts = [[x1,y1],[x2,y1],[x2,y2],[x1,y2],[x1,y1]]
        return verts
        ◊
```

Macro referenced in 46b.

Computation of the 1D centroid of a list of 2D boxes The 1D centroid of a list of 2D boxes is computed by the function given below. The direction of computation (either x or y) is chosen depending on the value of the xy parameter.

 \langle Computation of the 1D centroid of a list of 2D boxes 48b $\rangle \equiv$

```
""" Computation of the 1D centroid of a list of 2D boxes """

def centroid(boxes,coord):
    delta,n = 0,len(boxes)
    ncoords = len(boxes[0])/2
    a = coord%ncoords
    b = a+ncoords
    for box in boxes:
        delta += (box[a] + box[b])/2
    return delta/n
```

Macro referenced in 46b.

Pyplasm XOR of FAN of ordered points

```
⟨Pyplasm XOR of FAN of ordered points 49⟩ ≡
    """ XOR of FAN of ordered points """

def FAN(points):
    pairs = zip(points[1:-2],points[2:-1])
    triangles = [MKPOL([[points[0],p1,p2],[[1,2,3]],None]) for p1,p2 in pairs]
    return XOR(triangles)

if __name__ == "__main__":
    pol = [[0.476,0.332],[0.461,0.359],[0.491,0.375],[0.512,0.375],[0.514,0.375],
    [0.527,0.375],[0.543,0.34],[0.551,0.321],[0.605,0.314],[0.602,0.307],[0.589,
    0.279],[0.565,0.244],[0.559,0.235],[0.553,0.227],[0.527,0.239],[0.476,0.332]]

VIEW(EXPLODE(1.2,1.2,1)(FAN(pol)))
```

Macro referenced in 46b.

References

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- [PBCF93] A. Paoluzzi, F. Bernardini, C. Cattani, and V. Ferrucci, *Dimension-independent modeling with simplicial complexes*, ACM Trans. Graph. **12** (1993), no. 1, 56–102.

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