Boolean combinations of cellular complexes as chain operations *

Alberto Paoluzzi

August 18, 2014

Contents

1 Introduction

1.1 Preview of the algorithm

- 1. Embed both cellular complexes A and B in the same space (say, identify their common vertices) by $V_{ab} = V_a \cup V_b$.
- 2. Build their CDC (Common Delaunay Complex) as the LAR of Delaunay triangulation of the vertex set V_{ab} .
- 3. Split the (highest-dimensional) cells of CDC crossed by ∂A or ∂B ;
- 4. With respect to the (split) CDC basis, compute two coordinate chains $\alpha, \beta : CDC \rightarrow \{0,1\}$, such that:

$$\alpha(cell) = 1$$
 if $|cell| \subset A$; else $\alpha(cell) = 0$,

$$\beta(cell) = 1$$
 if $|cell| \subset B$; else $\beta(cell) = 0$.

5. Extract accordingly the CDC chain corresponding to AopB, with $op \in \{\cup, \cap, -\}$.

You may think to the split CDC as a CDT (constrained Delaunay Complex). In part they coincide, but in general, CDC is a polytopal complex (not a simplicial complex).

^{*}This document is part of the *Linear Algebraic Representation with CoChains* (LAR-CC) framework [?]. August 18, 2014

2 Merging arguments

2.1 Reordering of vertex coordinates

A global reordering of vertex coordinates is executed as the first step of the Boolean algorithm, in order to eliminate the duplicate vertices, by substituting duplicate vertex copies (coming from two close points) with a single instance.

Two dictionaries are created, then merged in a single dictionary, and finally split into three subsets of (vertex,index) pairs, with the aim of rebuilding the input representations, by making use of a novel and more useful vertex indexing.

The union set of vertices is finally reordered using the three subsets of vertices belonging (a) only to the first argument, (b) only to the second argument and (c) to both, respectively denoted as V_1, V_2, V_{12} . A top-down description of this initial computational step is provided by the set of macros discussed in this section.

@D Place the vertices of Boolean arguments in a common space @""" First step of Boolean Algorithm """ @; Initial indexing of vertex positions @; @; Merge two dictionaries with keys the point locations @; @; Filter the common dictionary into three subsets @; @; Compute an inverted index to reorder the vertices of Boolean arguments @; @; Return the single reordered pointset and the two d-cell arrays @; @

2.1.1 Re-indexing of vertices

Initial indexing of vertex positions The input LAR models are located in a common space by (implicitly) joining V1 and V2 in a same array, and (explicitly) shifting the vertex indices in CV2 by the length of V1. @D Initial indexing of vertex positions @from collections import defaultdict, OrderedDict

```
""" TODO: change defaultdict to OrderedDefaultdict """ class OrderedDefaultdict (collections.OrderedDict): def _{init}_{(self,*args,**kwargs):ifnotargs:self.default_factory=Noneelse:ifn} def _{missing}_{(self,key):ifself.default_factoryisNone:raiseKeyError(key)self[key]=default=self.default_factory()returndefault} def _{reduce}_{(self):optional,forpicklesupportargs=(self.default_factory,)ifself.default_factoryelsetuple()returnself._{class}_{,args,None,None,self.iters} def vertexSieve(model1, model2): from lar2psm import larModelBreak V1,CV1 = lar-ModelBreak(model1) V2,CV2 = larModelBreak(model2) n = len(V1); m = len(V2) def shift(CV, n): return [[v+n for v in cell] for cell in CV] CV2 = shift(CV2,n) @
```

Merge two dictionaries with point location as keys Since currently CV1 and CV2 point to a set of vertices larger than their initial sets V1 and V2, we index the set V1 \cup V2 using a Python defaultdict dictionary, in order to avoid errors of "missing key". As dictionary keys, we use the string representation of the vertex position vector provided by the vcode function given in the Appendix. @D Merge two dictionaries with keys the point

 $\begin{aligned} & \text{locations} @ \ \text{vdict1} = \text{defaultdict(list)} \ \text{for} \ k, v \ \text{in} \ \text{enumerate(V1):} \ vdict1[vcode(v)].append(k) \\ & \text{vdict2} = \ \text{defaultdict(list)} \ \text{for} \ k, v \ \text{in} \ \text{enumerate(V2):} \ vdict2[vcode(v)].append(k+n) \end{aligned}$

vertdict = defaultdict(list) for point in vdict1.keys(): vertdict[point] += vdict1[point]
for point in vdict2.keys(): vertdict[point] += vdict2[point] @

Example of string coding of a vertex position The position vector of a point of real coordinates is provided by the function **vcode**. An example of coding is given below. The *precision* of the string representation can be tuned at will.

```
>>> vcode([-0.011660381062724849, 0.297350056848685860])
'[-0.0116604, 0.2973501]'
```

Filter the common dictionary into three subsets Vertdict, dictionary of vertices, uses as key the position vectors of vertices coded as string, and as values the list of integer indices of vertices on the given position. If the point position belongs either to the first or to second argument only, it is stored in case1 or case2 lists respectively. If the position (item.key) is shared between two vertices, it is stored in case12. The variables n1, n2, and n12 remember the number of vertices respectively stored in each repository. @D Filter the common dictionary into three subsets @ case1, case12, case2 = [],[],[] for item in vertdict.items(): key,val = item if len(val)==2: case12 += [item] elif val[0]; n: case1 += [item] else: case2 += [item] n1 = len(case1); n2 = len(case12); n3 = len(case2) @

Compute an inverted index to reorder the vertices of Boolean arguments The new indices of vertices are computed according with their position within the storage repositories case1, case2, and case12. Notice that every item[1] stored in case1 or case2 is a list with only one integer member. Two such values are conversely stored in each item[1] within case12. @D Compute an inverted index to reorder the vertices of Boolean arguments @ invertedindex = list(0 for k in range(n+m)) for k,item in enumerate(case1): invertedindex[item[1][0]] = k for k,item in enumerate(case12): invertedindex[item[1][0]] = k+n1 invertedindex[item[1][1]] = k+n1 for k,item in enumerate(case2): invertedindex[item[1][0]] = k+n1+n2 @

2.1.2 Re-indexing of d-cells

Return the single reordered pointset and the two d-cell arrays We are now finally ready to return two reordered LAR models defined over the same set V of vertices, and where (a) the vertex array V can be written as the union of three disjoint sets of points C_1, C_{12}, C_2 ; (b) the d-cell array CV1 is indexed over $C_1 \cup C_{12}$; (b) the d-cell array CV2 is indexed over $C_{12} \cup C_{2}$.

The vertexSieve function will return the new reordered vertex set $V = (V_1 \cup V_2) \setminus (V_1 \cap V_2)$, the two renumbered s-cell sets CV1 and CV2, and the size len(case12) of $V_1 \cap V_2$. @D Return the single reordered pointset and the two d-cell arrays @ V =

[eval(p[0]) for p in case1] + [eval(p[0]) for p in case12] + [eval(p[0]) for p in case2] CV1= [sorted([invertedindex[v] for v in cell]) for cell in CV1] CV2 = [sorted([invertedindex[v] for v in cell]) for cell in CV2] return V, CV1, CV2, len(case12) @

2.1.3 Example of input with some coincident vertices

In this example we give two very simple LAR representations of 2D cell complexes, with some coincident vertices, and go ahead to re-index the vertices, according to the method implemented by the function vertexSieve.

Example discussion The aim of the vertexSieve function is twofold: (a) eliminate vertex duplicates before entering the main part of the Boolean algorithm; (b) reorder the input representations so that it becomes less expensive to check whether a 0-cell can be shared by both the arguments of a Boolean expression, so that its coboundaries must be eventually split. Remind that for any set it is:

$$|A \cup B| = |A| + |B| - |A \cap B|.$$

Let us notice that in the previous example

$$|V| = |V_1 \cup V_2| = 12 \le |V_1| + |V_2| = 6 + 9 = 15,$$

and that

$$|V_1| + |V_2| - |V_1 \cup V_2| = 15 - 12 = 3 = |C_{12}| = |V_1 \cap V_2|,$$

where C_{12} is the subset of vertices with duplicated instances. @D Output from test/py/boolean/test02.py @V = [[3.0,1.0],[2.0,1.0],[3.0,3.0],[1.0,1.0],[1.0,3.0],[2.0,3.0],[3.0,2.0],[2.0,0.0],[2.0,2.0],[0.0,0.0],[3.0,0.0],[0.0,1.0]] CV1 = [[3,5,1,4],[2,0,5,1]] CV2 = [[8,6,7,10],[3,4,5,8,11,9,7]] @ Notice also that V has been reordered in three consecutive subsets C_1, C_{12}, C_2 such that CV1 is indexed within $C_1 \cup C_{12}$, whereas CV2 is indexed within $C_{12} \cup C_2$. In our example we have $C_{12} = \{3,4,5\}$: @D Reordering of vertex indexing of cells @ $\vdots \vdots \vdots$ sorted(CAT(CV1)) [0, 1, 1, 2, 3, 4, 5, 5] $\vdots \vdots \vdots$ sorted(CAT(CV2)) [3, 4, 5, 6, 7, 7, 8, 8, 9, 10, 11] @

Cost analysis Of course, this reordering after elimination of duplicate vertices will allow to perform a cheap O(n) discovering of (Delaunay) cells whose vertices belong both to V1 and to V2. Actually, the same test can be now used both when the vertices of the input arguments are all different, and when they have some coincident vertices. The total cost of such pre-processing, executed using dictionaries, is $O(n \ln n)$.

2.1.4 Example

Building a covering of common convex hull @D Building a covering of common convex hull @def covering(model1,model2,dim=2,emptyCellNumber=1): V, CV1, CV2, n12 = vertexSieve(model1,model2) $_{,}EEV1 = larFacets((V,CV1),dim,emptyCellNumber),EEV2 = larFacets((V,CV2),dim,emptyCellNumber)ifemptyCellNumber! = 0 : CV1 = CV1[: -emptyCellNumber]ifemptyCellNumber! = 0 : CV2 = CV2[: -emptyCellNumber]VV = AA(LIST)(range(len(V)))returnV, [VV, EEV1, EEV2, CV1, CV2], n12@$

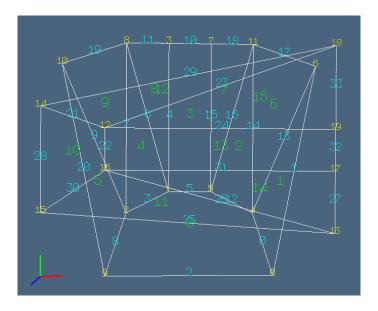


Figure 1: Set covering of the two Boolean arguments.

3 Selecting cells to split

Relational inversion (Characteristic matrix transposition) @D Characteristic matrix transposition @""" Characteristic matrix transposition """ def invertRelation(V,CV): VC = [[] for k in range(len(V))] for k,cell in enumerate(CV): for v in cell: VC[v] += [k] return VC @

@D Look for cells in Delaunay, with vertices in both operands @""" Look for cells in Delaunay, with vertices in both operands """ def mixedCells(CV,CV1,CV2,n12): n0,n1 = 0, $\max(AA(\max)(CV1))$ vertices in CV1 (extremes included) m0,m1 = n1+1-n12, $\max(AA(\max)(CV2))$ vertices in CV2 (extremes included) return [list(cell) for cell in CV if $any([n0]=v_i=n1$ for v in cell]) and $any([m0]=v_i=m1$ for v in cell])] @

@D Look for cells in cells12, with vertices on boundaries @"" Look for cells in cells12, with vertices on boundaries """ def mixedCellsOnBoundaries(cells12,BV): cells12BV =

[cell for cell in cells12 if len(list(set(cell).intersection(BV))) != 0] return cells12BV @

@D Build intersection tasks @""" Build intersection tasks """ def cutting Test (cutting Hyperplane, polytope, signs = [INNERPROD([cutting Hyperplane, V[v]+[1.]]) for v in polytope] signs = eval(vcode(signs)) return any([value;-0.001 for value in signs]) and any([value;0.001 for value in signs])

$$\label{eq:cutVerts} \begin{split} & \text{def splittingTasks}(V, \text{pivots}, BV, BC, VBC, CV, VC): } & \text{tasks} = [] \text{ for pivotCell in pivots: } \\ & \text{cutVerts} = [\text{v for v in pivotCell if v in BV}] \text{ for v in cutVerts: } & \text{cutFacets} = VBC[\text{v}] \text{ cells2cut} \\ & = VC[\text{v}] \text{ for face,cell in CART}([\text{cutFacets,cells2cut}]): } & \text{polytope} = CV[\text{cell}] \text{ points} = [V[\text{w}] \\ & \text{for w in BC[face]}] \text{ dim} = \text{len(points[0]) } & \text{theMat} = \text{Matrix}([(\text{dim}+1)^*[1.]] + [\text{p}+[1.]] \text{ for p} \\ & \text{in points}]) & \text{cuttingHyperplane} = [(-1)^{**}(\text{col})^* & \text{theMat.minor}(0,\text{col}). \\ & \text{determinant}() \text{ for col in} \\ & \text{range}(\text{dim}+1)] & \text{if cuttingTest}(\text{cuttingHyperplane,polytope}, V): \\ & \text{tasks} + = [[\text{face,cell,cuttingHyperplane}]] \\ & \text{tasks} = \text{AA(eval)}(\text{set}(\text{AA(str)}(\text{tasks}))) & \text{tasks} = \text{TrivialIntersection}(\text{tasks}, V, BC, CV) & \text{return} \\ & \text{tasks} & @ \\ \end{aligned}$$

facet-cell trivial intersection filtering A final filtering is applied to the pairs (cutting-Hyperplane, polytope in the tasks array, in order to remove those facets (pairs in 2D) whose intersection reduces to a single point, i.e. to the comman vertex between the boundary (d-1)-face, having cuttingHyperplane as affine hull, and the polytope d-cell.

For this purpose, it is checked that at least one of the facet vertices, transformed into the common-vertex-based coordinate frame, have all positive coordinates. This fact guarantees the existence of a non trivial intersection between the (d-1)-face and the d-cell.

@D Trivial intersection filtering @""" Trivial intersection filtering """ def Trivial-Intersection(tasks,V,EEV,CV): out = [] for face,cell,affineHull in tasks: faceVerts, cellVerts = EEV[face], CV[cell] v0 = list(set(faceVerts).intersection(cellVerts))[0] v0 = common vertex transformMat = mat([VECTDIFF([V[v],V[v0]])) for v in cellVerts if v != v0]).T.I vects = (transformMat * (mat([VECTDIFF([V[v],V[v0]]))) for v in faceVerts if v != v0]).T.tolist() if any([all([x;0 for x in list(vect)])) for vect in vects]): out += [[face,cell,affineHull]] return out @

4 Splitting cells traversing the boundaries

In the previous section we computed a set of "slitting seeds", each made by a boundary facet and by a Delaunay cell to be splitted by the facet's affine hull. Here we show how to partition at each such cells into two cells, according to Figure ??, where the boundary facets of the two boolean arguments are shown in yellow color.

In the example in Figure ??, the set of pairs (facet, cell) to be used as splitting seeds are given below.

[[25, 3], [1, 3], [29, 18], [20, 22], [1, 19], [25, 10], [20, 10], [29, 22]]

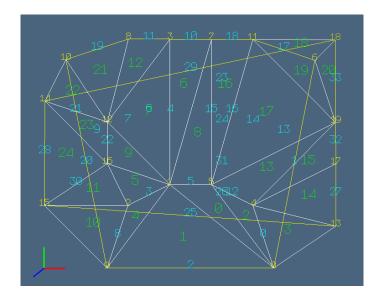


Figure 2: example caption

4.1 Cell splitting

A cell will be split by pyplasm intersection with a suitable rotated and translated instance of a (large) d-cuboid with the superior face embedded in the hyperplane z = 0.

Splitting a cell with an hyperplane The macro below defines a function cellSplitting, with input the index of the face, the index of the cell to be bisected, the covector giving the coefficients of the splitting hyperplane, i.e. the affine hull of the splitting face, and the arrays V, EEV, CV, giving the coordinates of vertices, the (accumulated) facet to vertices relation (on the input models), and the cell to vertices relation (on the Delaunay model), respectively.

The actual subdivision of the input cell onto the two output cells cell1 and cell2 is performed by using the pyplasm Boolean operations of intersection and difference of the input with a solid simulation of the needed hyperspace, provided by the rototranslSubspace variable. Of course, such pyplasm operators return two Hpc values, whose vertices will then extracted using the UKPOL primitive.

```
@D Cell splitting @""" Cell splitting in two cells """ def cellSplitting(face,cell,covector,V,EEV,CV): \dim = \operatorname{len}(V[0]) \text{ subspace} = (\operatorname{T}(\operatorname{range}(1,\dim+1))(\dim^*[-50])(\operatorname{CUBOID}(\dim^*[100]))) \text{ normal} = \operatorname{covector}[:-1] \text{ if len}(\operatorname{normal}) == 2 : \quad 2D \text{ complex rotatedSubspace} = \operatorname{R}([1,2])(\operatorname{ATAN2}(\operatorname{normal}) - \operatorname{PI}/2)(\operatorname{T}(2)(-50)(\operatorname{subspace})) \text{ elif len}(\operatorname{normal}) == 3 : \quad 3D \text{ complex rotatedSubspace} = \operatorname{R}()()(\operatorname{subspace}) \text{ else: print "rotation error" t} = \operatorname{V}[\operatorname{EEV}[\operatorname{face}][0]] \text{ rototranslSubspace} = \operatorname{T}(\operatorname{range}(1,\dim+1))(t)(\operatorname{rotatedSubspace}) \text{ cellHpc} = \operatorname{MKPOL}([V,[[v+1 \text{ for v in } \operatorname{CV}[\operatorname{cell}]]],\operatorname{None}]) 
\operatorname{cell1} = \operatorname{INTERSECTION}([\operatorname{cellHpc},\operatorname{rototranslSubspace}]) \text{ tolerance} = 0.0001 \text{ use}_{o}ctree =
```

 $Falsecell 1 = Plasm.boolop(BOOL_CODE_AND, [cellHpc, rototranslSubspace], tolerance, plasm_config.maxn UKPOL(cell1)cell 1 = AA(vcode)(verts)$

 $cell2 = DIFFERENCE([cellHpc,rototranslSubspace]) cell2 = Plasm.boolop(BOOL_CODE_DIFF, [cellHpc,rototranslSubspace]) cell2 = Plasm.boolop(BOOL_CODE_DIFF, [cellHpc,rototr$

return cell1,cell2 @

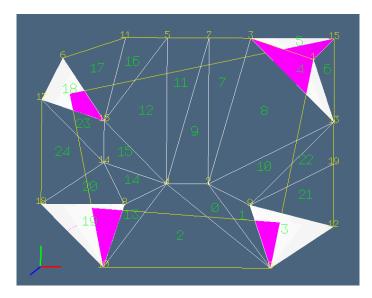


Figure 3: example caption

4.2 Cross-building of two task dictionaries

The correct and efficient splitting of the combined Delaunay complex (CDC) with the (closed and orientable) boundaries of two Boolean arguments, requires the use of two special dictionaries, respectively named dict_fc (for face-cell), and dict_cf (for cell-face).

On one side, for each splitting facet ((d-1)-face), used as key, we store in dict_fc the list of traversed d-cells of CDC, starting in 2D with the two cells containing the two extreme vertices of the cutting edge, and in higher dimensions, with all the d-cells containing one of vertices of the splitting (d-1)-face.

On the other side, for each d-cell to be split, used as key, we store in $dict_cf$ the list of cutting (d-1)-cells, since a single d-cell may be traversed and split by more than one facet.

```
Example of face-cell and cell-face dictionaries @D Example of face-cell and cell-face dictionaries @""" Example of face-cell and cell-face dictionaries """ tasks (face,cell) = [ [0, 4, [-10.0, 2.0, 110.0]], [31, 5, [3.0, -14.0, 112.0]], [17, 18, [10.0, 2.0, -30.0]], [22, 3, [-1.0, -14.0, 42.0]], [17, 19, [10.0, 2.0, -30.0]], [31, 18, [3.0, -14.0, 112.0]], [22, 19, [-1.0, -14.0, 42.0]], [0, 3, [-10.0, 2.0, 110.0]]]
```

```
 \begin{array}{l} {\rm tasks}\;({\rm dict}_f c) = default dict (< type'list'>, 0:[(4,[-10.0,2.0,110.0]),(3,[-10.0,2.0,110.0])],17:[(18,[10.0,2.0,10.0]),(2,[-1.0,-14.0,42.0])],18:[(17,[10.0,2.0,-30.0]),(22,[-1.0,-14.0,42.0])],18:[(17,[10.0,2.0,-30.0]),(22,[-1.0,-14.0,42.0])],18:[(17,[10.0,2.0,-30.0]),(22,[-1.0,-14.0,42.0])],18:[(17,[10.0,2.0,-30.0]),(22,[-1.0,-14.0,42.0])],18:[(17,[10.0,2.0,-30.0]),(22,[-1.0,-14.0,42.0])],18:[(17,[10.0,2.0,-30.0]),(22,[-1.0,-14.0,42.0])],18:[(17,[10.0,2.0,-30.0]),(22,[-1.0,-14.0,42.0])],18:[(17,[10.0,2.0,-30.0]),(22,[-1.0,-14.0,42.0])],18:[(17,[10.0,2.0,-30.0]),(22,[-1.0,-14.0,42.0])],18:[(17,[10.0,2.0,-30.0]),(22,[-1.0,-14.0,42.0])],18:[(17,[10.0,2.0,-30.0]),(22,[-1.0,-14.0,42.0])],18:[(17,[10.0,2.0,-30.0]),(22,[-1.0,-14.0,42.0])],18:[(17,[10.0,2.0,-30.0]),(22,[-1.0,-14.0,42.0])],18:[(17,[10.0,2.0,-30.0]),(22,[-1.0,-30.0]),(22,[-1.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(22,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,[-3.0,-30.0]),(23,
```

4.3 Updating the vertex set and dictionary

In any dimension, the split of a *d*-cell with an hyperplane (crossing its interior) produces two *d*-cells and some new vertices living upon the splitting hyperplane.

When the d-cell c is contained in only one seed of the CDC decomposition, i.e. when $dict_cf[c]$ has cardinality one (in other words: it is crossed only by one boundary facet), the two generated cells vcell1, vcell2 can be safely output, and accommodated in two slots of the CV list.

Conversely, when more than one facet crosses c, much more care must be taken to guarantee the correct fragmentation of this cell.

Managing the splitting dictionaries The function splittingControl takes care of cells that must be split several times, as crossed by several boundary faces.

If the dictionary item dict_cf[cell] has *length* one (i.e. is crossed *only* by one face) the CV list is updated and the function returns, in order to update the dict_fc dictionary.

Otherwise, the function subdivides the facets cutting cell between those to be associated to vcell1 and to vcell2. For each pair aface, covector in dict_cf[cell] and in position following face in the list of pairs, check if either vcell1 or vcell2 or both, have intersection with the subset of vertices shared between cell and aface, and respectively put in alist1, in alist2, or in both. Finally, store vcell1 and vcell2 in CV, and alist1, alist2 in dict_cf.

```
@D Managing the splitting dictionaries @""" Managing the splitting dictionaries """ def splittingControl(face,cell,covector,vcell,vcell1,vcell2,dict_fc, dict_cf, V, BC, CV, VC, CVbits, lenBC1): boundaryFacet = BC[face] translVector = V[boundaryFacet[0]] tcovector = [cv+tv*covector[-
```

1] for (cv,tv) in zip(covector[:-1],translVector)]+[0.0] print "tcovector =",tcovector

 ${\tt c1,c2} = {\tt cell,cell\ if\ not\ haltingSplitTest(cell,vcell1,vcell2,boundaryFacet,translVector,tcovector,V)}$

only one facet covector crossing the cell cell Verts = CV[cell] CV[cell] = vcell1 CV += [vcell2] print "¿¿¿¿ cell,len(CV) =",cell,len(CV) CVbits += [copy(CVbits[cell])] c1,c2 = cell,len(CV)-1

firstCell,secondCell = AA(testingSubspace(V,covector))([vcell1,vcell2]) if face i lenBC1 and firstCell==-1: face in boundary(op1) CVbits[c1][0] = 0 CVbits[c2][0] = 1 elif face

4.4 Updating the split cell and the queues of seeds

When a d-cell of the combined Delaunay complex (CDC) is split into two d-cells, the first task to perform is to update its representation as vertex list, and to update the list of d-cells. In particular, as cell, and cell1, cell2 are the input d-cell and the two output d-cells, respectively, we go to substitute cell with cell1, and to add the cell2 as a new row of the $CSR(M_d)$ matrix, i.e. as the new terminal element of the CV array. Of course, the reverse relation VC must be updated too.

Updating the split cell First of all notice that, whereas cell is given as an integer index to a CV row, cell1, cell2 are returned by the cellSplitting function as lists of lists of coordinates (of vertices). Therefore such vectors must be suitably transformed into dictionary keys, in order to return the corresponding vertex indices. When transformed into two lists of vector indices, cell1, cell2 will be in the form needed to update the CV and VC relations.

Updating the vertex set of split cells The code in the macro below provides the splitting of the CDC along the boundaries of the two Boolean arguments. This function, and the ones called by its, provide the dynamic update of the two main data structures, i.e. of the LAR model (V,CV).

@D Updating the vertex set of split cells @""" Updating the vertex set of split cells """ def testingSubspace(V,covector): def testingSubspace0(vcell): inout = SIGN(sum([INNERPROD([V[v]+[1.],cofor v in vcell])) return inout return testingSubspace0

```
def splitCellsCreateVertices(vertdict,dict<sub>f</sub>c, dict_cf, V, BC, CV, VC, lenBC1): DEBUG = False; CVbits = [[-1, -1]forkinrange(len(CV))]nverts = len(V); cellPairs = []; twoCellIndices = []
```

```
[]; while any([tasks! = []forface, tasks indict_fc.items()]): forface, tasks indict_fc.items(): fortask intasks: cell, covector = task vcell = <math>CV[cell]
```

if cuttingTest(covector,vcell,V): cell1,cell2 = cellSplitting(face,cell,covector,V,BC,CV) print ",face,cell1,cell2",cell,face,cell1,cell2 if cell1 == [] or cell2 == []: print "cell1,cell2 =",cell1,cell2 else: adjCells = adjacencyQuery(V,CV)(cell)

print "cell,adjCells =",cell,adjCells,cell1,cell2

vcell1 = [] for k in cell1: if vertdict[k] = = []: vertdict[k] + = [nverts] V + = [eval(k)] nverts + = 1 vcell1 + = [vertdict[k]]

 $vcell1 = CAT(vcell1) \ vcell2 = CAT([vertdict[k] \ for \ k \ in \ cell2])$

V,CV,CV bits, $dict_cf$, $dict_fc$, $twoCells = splittingControl(face, cell, covector, vcell, vcell1, vcell2, dict_fc, divoCells[1]:$

for adjCell in adjCells: $dict_f c[face] + = [(adjCell, covector)]dict_c f[adjCell] + = [(face, covector)]cellPairs [[vcell1, vcell2]]twoCellIndices + = [[twoCells]]$

 $if \ DEBUG: showSplitting(V,cellPairs,BC,CV) \ else: \ dict_fc[face]. remove((cell,covector)) remove the splitted properties of the context of the context$

Test for split halting along a boundary facet The cell splitting is operated by the facet's hyperplane H(f), that we call *covector*, and the splitting with it may continues outside $f ext{ ... }!!$

This fact may induce some local errors in the decision procedure (attributing either 0 or 1 to each split cell pair). So, when splitting a pair (cell,face) — better: (cell,covector) — already stored in the data structure, and then computing its adjacent pairs, we should check if the common facet f_{12} between c_1 and c_2 is (or is not) at least partially internal to f.

If this fact is not true, and hence f_{12} is out(f) in the induced topology of the H(f) hyperplane, the split process on that pair must be halted: c_1 and c_2 are not stored, and their adjacent cells not split.

@D Test for split halting along a boundary facet @""" Test for split halting along a boundary facet """ def haltingSplitTest(cell,vcell1,vcell1,vcell2,boundaryFacet,translVector,tcovector,V): newFacet = list(set(vcell1).intersection(vcell2)) print ",vcell1,vcell2 = ",vcell,vcell1,vcell2 print "newFacet,boundaryFacet = ",newFacet,boundaryFacet

$$\label{eq:translation} \begin{split} & translation \ newFacet = [\ eval(vcode(VECTDIFF([V[v],translVector]))) \ for \ v \ in \ newFacet] \ boundaryFacet = [\ eval(vcode(VECTDIFF([V[v],translVector]))) \ for \ v \ in \ boundaryFacet] \ print "newFacet,boundaryFacet = ",newFacet,boundaryFacet] \end{split}$$

linear transformation: new Facet -; standard (d-1)-simplex transform Mat = mat(boundary Facet[1:] + [tcovector[:-1]]). T.I print "transform Mat = ",transform Mat

 $\begin{array}{l} \operatorname{transformation\ in\ the\ subspace}\ \mathbf{x}_d = 0 new Facet = AA(COMP([eval,vcode]))((transformMat*(newFacet).T)).T.tolist())boundary Facet = AA(COMP([eval,vcode]))((transformMat*(mat(boundaryFacet).T)).T.tolist())print" newFacet, boundary Facet = ", newFacet, boundary Facet | ", newFacet, boundary$

projection in E^{d-1} space and Boolean test new Facet = MKPOL([AA(lambdav:v[:-1])(new Facet), [range])

1)], None]) boundary Facet = MKPOL([AA(lambdav : v[: -1])(boundary Facet), [range(1, len(boundary Facet), [vertex, cells, range = VKPOL([NTEPSECTION([new Facet, boundary Facet]))]) <math>boundary Facet

1)], None]) verts, cells, pols = UKPOL(INTERSECTION([newFacet, boundaryFacet])) print" verts, cells, policy and print the print of the print that the print the prin

```
".verts, cells, pols, ""ifverts == []:print(')print"*****cell = ", cellsys.exit()returnTrueelse:returnFalse
```

$$\label{eq:cell1} \begin{split} \text{cell1} &= \text{INTERSECTION}([\text{cellHpc}, \text{rototranslSubspace}]) \ \ \text{tolerance} = 0.0001 \ \ \text{use}_{o}ctree = \\ Falsecell1 &= Plasm.boolop(BOOL_CODE_AND, [cellHpc, rototranslSubspace], tolerance, plasm_config.maxnUKPOL(cell1)cell1 &= AA(vcode)(verts)if@ \end{split}$$

Test about halting the split along a facet @D Test about halting the split along a facet @""" Test about halting the split along a facet """ def haltingSplitTest(face,vcell,adjCell,CV,V,printout=facet = list(set(vcell).intersection(CV[adjCell]))

```
\begin{aligned} & points = [V[v] \ for \ v \ in \ facet] \ dim = len(points[0]) \ theMat = Matrix(\ [(dim+1)^*[1.]] \ + \\ & [p+[1.] \ for \ p \ in \ points] \ ) \ covector = [(-1)^{**}(col)^*theMat.minor(0,col).determinant() \ for \ col \ in \ range(dim+1)] \ if \ printout: \ print \ "face,vcell,CV[adjCell],facet,covector = ",face,vcell,CV[adjCell],facet,covector \ haltingTest = (not \ cuttingTest(covector,CV[face],V)) \ if \ printout: \ print \ "haltingTest = ",haltingTest \ print \ "(set(facet).intersection(CV[adjCell])!=set()) = ",(set(facet).intersection(CV[adjCell])!=set()) \ print \ "(not \ cuttingTest(covector,CV[face],V)) = ",(not \ cuttingTest(covector,CV[face],V)) \ return \ haltingTest \ @ \end{aligned}
```

4.5 Updating the cells adjacent to the split cell

Once the list of d-cells has been updated with respect to the results of a split operation, it is necessary to consider the possible update of all the cells that are adjacent to the split one. It particular we need to update their lists of vertices, by introducing the new vertices produced by the split, and by updating the dictionaries of tasks, by introducing the new (adjacent) splitting seeds.

Computing the adjacent cells of a given cell To perform this task we make only use of the CV list. In a more efficient implementation we should make direct use of the sparse adjacency matrix, to be dynamically updated together with the CV list. The computation of the adjacent d-cells of a single d-cell is given here by extracting a column of the $CSR(M_d M_d^t)$. This can be done by multiplying $CSR(M_d)$ by its transposed row corresponding to the query d-cell.

@D Computing the adjacent cells of a given cell @""" Computing the adjacent cells of a given cell """ def adjacencyQuery (V,CV): dim = len(V[0]) def adjacencyQuery0 (cell): nverts = len(CV[cell]) csrCV = csrCreate(CV) csrAdj = matrixProduct(csrCV,csrTranspose(csrCV)) cellAdjacencies = csrAdj.indices[csrAdj.indptr[cell]:csrAdj.indptr[cell+1]] return [acell for acell in cellAdjacencies if dim ;= csrAdj[cell,acell] ; nverts] return adjacencyQuery0 @

Updating the adjacency matrix At every step of the CDC splitting, generating two output cells cell1 and cell2 from the input cell, the element of such index in the list CV is restored with the cell1 vertices, and a new (last) element is created in CV, to store the cell2 vertices. Therefore the row of index cell of the symmetric adjacency matrix must

be recomputed, being the cell column updated consequently. Also, a new last row (and column) must be added to the matrix.

@D Updating the adjacency matrix @""" Updating the adjacency matrix """ pass @

5 Reconstruction of results

5.1 The Boolean algorithm flow

Show the process of CDC splitting @D Show the process of CDC splitting @""" Show the process of CDC splitting """ def showSplitting(V,cellPairs,BC,CV): VV = AA(LIST)(range(len(V))) boundaries = $COLOR(RED)(SKEL_1(STRUCT(MKPOLS((V,BC)))))$ submodel = COLOR(CYAN)(STRUCT(MKPOLS((V,BC))))) submodel = COLOR(CYAN)(STRUCT(SIS)) = COLOR(SIS) = COLO

Computation of bits of split cells In order to compute, in the simplest and more general way, whether each of the two split d-cells is internal or external to the splitting boundary d-1-facet, it is necessary to consider the oriented covector ϕ (or one-form) canonically associated to the facet f by the covector representation theorem, i.e. the corresponding oriented hyperplane. In this case, the internal/external attribute of the split cell will be computed by evaluating the pairing $\langle v, \phi \rangle$.

5.2 Final traversal of the CDC

Several cells of the split CDC are characterised as either internal or external to the Boolean arguments A and B according to the splitting process. Such characterisation is stored within the CVbits array of pairs of values in $\{-1,0,1\}$, where CVbits[k][h], with $k \in \text{range}(\text{len}(C_d))$ and $h \in \text{range}(2)$, has the following meanings:

```
\begin{array}{ll} -1, & \quad \text{if position of } c_k \in C_d \text{ is } unknown \text{ w.r.t. complex } K_h \\ \text{CVbits[k][h]} = & 0, & \quad \text{if cell } c_k \in C_d \text{ is } external \text{ w.r.t. complex } K_h \\ & 1, & \quad \text{if cell } c_k \in C_d \text{ is } internal \text{ w.r.t. complex } K_h \end{array}
```

Therefore, a double d-cell visit of CDC must be executed, starting from some d-cell interior to either A or B, and traversing from a cell to its untraversed adjacent cells, but without crossing the complex boundary, until all cells have been visited.

The initial computation of chains of Boolean arguments The initial setting of CVbits[k][h] values is done within the splitting process by the splitCellsCreateVertices function, and mainly by the splittingControl function.

The traversal of Boolean arguments Let us remember that the adjacency matrix between d-cells is computed via SpMSpM multiplication by the double application

```
adjacencyQuery(V,CV)(cell),
```

where the first application adjacencyQuery(V,CV) returns a partial function with bufferization of the adjacentcy matrix, and the second application to cell returns the list of adjacent d-cells sharing with it a (d-1)-dimensional facet.

Traversing a Boolean argument within the CDC A recursive function booleanChainTraverse is given in the script below, where

@D Traversing a Boolean argument within the CDC @"" Traversing a Boolean argument within the CDC """ def booleanChainTraverse(h,cell,V,CV,CVbits,value): adj-Cells = adjacencyQuery(V,CV)(cell) for adjCell in adjCells: if CVbits[adjCell][h] == -1: CVbits[adjCell][h] = value CVbits = booleanChainTraverse(h,adjCell,V,CV,CVbits,value) return CVbits @

Boolean fragmentation and classification of CDC @D Boolean fragmentation and classification of CDC @""" Boolean fragmentation and classification of CDC """ def booleanChains(arg1,arg2): (V1,basis1), (V2,basis2) = arg1,arg2 model1, model2 = (V1,basis1[-1]), (V2,basis2[-1]) V,[VV,,,CV1,CV2], n12 = covering(model1, model2, 2, 0)CV = sorted(AA(sorted)(Delaudefaultdict(list)fork, vinenumerate(V) : vertdict[vcode(v)] + = [k]

BC1 = signedCellularBoundaryCells(V1,basis1) BC2 = signedCellularBoundaryCells(V2,basis2)

 $submodel1 = mkSignedEdges((V1,BC1)) \ submodel2 = mkSignedEdges((V2,BC2)) \ VIEW(STRUCT([submodel2] \ submodel2) \ submodel2) \ submodel2 = mkSignedEdges((V2,BC2)) \ VIEW(STRUCT([submodel2] \ submodel2) \ submodel2) \ submodel2 = mkSignedEdges((V2,BC2)) \ VIEW(STRUCT([submodel2] \ submodel2) \ submodel2) \ submodel2 = mkSignedEdges((V2,BC2)) \ VIEW(STRUCT([submodel2] \ submodel2) \ submodel2) \ submodel2 = mkSignedEdges((V2,BC2)) \ VIEW(STRUCT([submodel2] \ submodel2) \ submodel2) \ submodel2 = mkSignedEdges((V2,BC2)) \ VIEW(STRUCT([submodel2] \ submodel2) \ submodel2) \ submodel2 = mkSignedEdges((V2,BC2)) \ VIEW(STRUCT([submodel2] \ submodel2) \ submodel2) \ submodel2 = mkSignedEdges((V2,BC2)) \ VIEW(STRUCT([submodel2] \ submodel2) \ submodel2) \ submodel2 \ su$

BC = sorted([[vertdict[vcode(V1[v])][0] for v in cell] for cell in BC1] + [[vertdict[vcode(V2[v])][0] for v in cell] for cell in BC2]) BV = list(set(CAT([v for v in BC]))) VV

 $= AA(LIST)(range(len(V))) \text{ submodel} = SKEL_1(STRUCT(MKPOLS((V,CV))))VIEW(larModelNumber STRUCT([SKEL_1(STRUCT(MKPOLS((V,CV)))), COLOR(RED)(STRUCT(MKPOLS((V,BC))))])VIEW(larModelNumber STRUCT([SKEL_1(STRUCT(MKPOLS((V,CV)))), COLOR(RED)(STRUCT(MKPOLS((V,BC))))])VIEW(larModelNumber STRUCT([SKEL_1(STRUCT(MKPOLS((V,CV))), COLOR(RED)(STRUCT(MKPOLS((V,BC))))])VIEW(larModelNumber STRUCT([SKEL_1(STRUCT(MKPOLS((V,CV))), COLOR(RED)(STRUCT(MKPOLS((V,BC))))])VIEW(larModelNumber STRUCT([SKEL_1(STRUCT(MKPOLS((V,CV))), COLOR(RED)(STRUCT(MKPOLS((V,BC))))])VIEW(larModelNumber STRUCT([SKEL_1(STRUCT(MKPOLS((V,CV))), COLOR(RED)(STRUCT(MKPOLS((V,BC))))])VIEW(larModelNumber STRUCT([SKEL_1(STRUCT(MKPOLS((V,BC)), COLOR(RED)(STRUCT(MKPOLS((V,BC))))])VIEW(larModelNumber STRUCT([SKEL_1(STRUCT(MKPOLS((V,BC)), COLOR(RED)(STRUCT(MKPOLS((V,BC))))])VIEW(larModelNumber STRUCT([SKEL_1(STRUCT(MKPOLS((V,BC)), COLOR(RED)(STRUCT(MKPOLS((V,BC))))]))VIEW(larModelNumber STRUCT([SKEL_1(STRUCT(MKPOLS((V,BC)), COLOR(RED)(STRUCT(MKPOLS((V,BC))))]))VIEW(larModelNumber STRUCT([SKEL_1(STRUCT(MKPOLS((V,BC)), COLOR(RED)(STRUCT(MKPOLS((V,BC))))]))VIEW(larModelNumber STRUCT(STRUCT(MKPOLS((V,BC)))))))VIEW(larModelNumber STRUCT(S$

 $cells 12 = mixed Cells (CV, CV1, CV2, n12) \ pivots = mixed Cells On Boundaries (cells 12, BV)$

 $\label{eq:VBC} VBC = \text{invertRelation}(V,BC) \ VC = \text{invertRelation}(V,CV) \ tasks = \text{splittingTasks}(V,\text{pivots},BV,BC,VBC,CV,V) \ dict_f c, \\ dict_c f = initTasks(tasks)$

CVbits,cellPairs,twoCellIndices = splitCellsCreateVertices(vertdict,dict $_fc$, $dict_cf$, V, BC, CV, VC, len(BC) for k in range(len(CV)): print ",CVbits[k],CV[k] = ",k,CVbits[k],CV[k]

for cell in range(len(CV)): if CVbits[cell][0] == 1: CVbits = booleanChainTraverse(0,cell,V,CV,CVbits,1) if CVbits[cell][0] == 0: CVbits = booleanChainTraverse(0,cell,V,CV,CVbits,0) if CVbits[cell][1]

== 1: CVbits = booleanChainTraverse(1,cell,V,CV,CVbits,1) if CVbits[cell][1] == 0:

CVbits = booleanChainTraverse(1,cell,V,CV,CVbits,0)

for k in range (len(CV)): print ",CVbits[k],CV[k] = ",k,CVbits[k],CV[k]

 ${\it chain 1, chain 2} = {\it TRANS}({\it CVbits}) \; {\it print} \; "_c f", \\ {\it dict_c f print} "_f c", \\ {\it dict_f c}, ""return V, CV, \\ {\it chain 1, chain 2, CVbits} \; ({\it CVbits}) \; {\it print} \; "_c f", \\ {\it dict_f c, ""return V, CV, chain 1, chain 2, CVbits} \; ({\it CVbits}) \; {\it print} \; "_c f", \\ {\it dict_f c, ""return V, CV, chain 1, chain 2, CVbits} \; ({\it CVbits}) \; {\it print} \; "_c f", \\ {\it dict_f c, ""return V, CV, chain 1, chain 2, CVbits} \; ({\it CVbits}) \; {\it print} \; "_c f", \\ {\it dict_f c, ""return V, CV, chain 1, chain 2, CVbits} \; ({\it CVbits}) \; {\it print} \; "_c f", \\ {\it dict_f c, ""return V, CV, chain 1, chain 2, CVbits} \; ({\it CVbits}) \; {\it CVbits} \; ({\it CVbits}) \; ({\it CVbi$

6 Exporting the library

@O lib/py/bool.py @""" Module for Boolean ops with LAR """ from matrix import * @; Initial import of modules @¿ @; Symbolic utility to represent points as strings @¿ @; Place the vertices of Boolean arguments in a common space @¿ @; Building a covering of common convex hull @¿ @; Building a partition of common convex hull of vertices @¿ @; Characteristic matrix transposition @¿ @; Look for cells in Delaunay, with vertices in both operands @¿ @; Look for cells in cells12, with vertices on boundaries @¿ @; Build intersection tasks @¿ @; Trivial intersection filtering @¿ @; Cell splitting @¿ @; Init face-cell and cell-face dictionaries @¿ @; Updating the split cell @¿ @; Updating the vertex set of split cells @¿ @; Managing the splitting dictionaries @¿ @; Test for split halting along a boundary facet @¿ @; Computing the adjacent cells of a given cell @¿ @; Show the process of CDC splitting @¿ @; Traversing a Boolean argument within the CDC @¿ @; Boolean fragmentation and classification of CDC @¿ @

7 Tests

7.1 2D examples

7.1.1 First examples

Three sets of input 2D data are prepared here, ranging from very simple to a small instance of the hardest kind of dataset, known to produce an output of size $O(n^2)$.

```
@D First set of 2D data: Fork-0 input @""" Definition of Boolean arguments """ V1 = [[3,0],[11,0], [13,10], [10,11], [8,11], [6,11], [4,11], [1,10], [4,3], [6,4], [8,4], [10,3]] FV1 = [[0,1,8,9,10,11],[1,2,11], [3,10,11], [4,5,9,10], [6,8,9], [0,7,8], [2,3,11], [3,4,10], [5,6,9], [6,7,8]] EV1 = [[0,1],[0,7],[0,8],[1,2],[1,11],[2,3],[2,11],[3,4],[3,10],[3,11],[4,5],[4,10],[5,6],[5,9],[6,7],[6,8],[6,9],[7,8],[8,9],[9,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10], [1,10
```

 $[2,3,6,7], [1,2,7], [3,4,6]] \ EV2 = [[0,1],[0,5],[0,7],[1,2],[1,7],[2,3],[2,7],[3,4],[3,6],[4,5],[4,6],[5,6],[6,7]]$ VV2 = AA(LIST)(range(len(V2))) @ $@D \ Eigst \ set \ of \ 2D \ data; \ Early 1 \ input \ @""" \ Definition \ of \ Packet parameter \ """ \ V1$

@D First set of 2D data: Fork-1 input @""" Definition of Boolean arguments """ V1 = [[3,0],[11,0], [13,10], [10,11], [8,11], [6,11], [4,11], [1,10], [4,3], [6,4], [8,4], [10,3]]

 $\begin{aligned} & \text{FV1} = [[0,1,8,9,10,11],[1,2,11],\ [3,10,11],\ [4,5,9,10],\ [6,8,9],\ [0,7,8]] \ \text{EV1} = [[0,1],[0,7],[0,8],[1,2],[1,11],[2,11],[3,10,11],\\ & \text{VV1} = \text{AA(LIST)}(\text{range}(\text{len}(\text{V1}))) \end{aligned}$

 $\begin{array}{l} V2 = [[0,3],[14,2],\ [14,5],\ [14,7],\ [14,11],\ [0,8],\ [3,7],\ [3,5]] \ FV2 = [[0,5,6,7],\ [0,1,7],\ [4,5,6], \\ [2,3,6,7],\ [1,2,7],\ [3,4,6]] \ EV2 = [[0,1],[0,5],[0,7],[1,2],[1,7],[2,3],[2,7],[3,4],[3,6],[4,5],[4,6],[5,6],[6,7]] \\ VV2 = AA(LIST)(range(len(V2))) @ \\ \end{array}$

Input and visualisation of Boolean arguments @D Computation of lower-dimensional cells @""" Computation of edges an input visualisation """ model1 = V1,FV1 model2 =

```
 \begin{aligned} &V2,FV2\;basis1 = [VV1,EV1,FV1]\;basis2 = [VV2,EV2,FV2]\;submodel12 = STRUCT(MKPOLS((V1,EV1)) + MVIEW(larModelNumbering(V1,basis1,submodel12,4))\;VIEW(larModelNumbering(V2,basis2,submodel12,4))\\ &@ \end{aligned}
```

Exporting test file @D Bulk of Boolean task computation @""" Bulk of Boolean task computation """ @; Computation of lower-dimensional cells @;

 $V,CV,chain1,chain2,CVbits = booleanChains((V1,basis1), (V2,basis2)) \ for \ k \ in \ range(len(CV)): print \ ",CVbits[k],CV[k] = ",k,CVbits[k],CV[k]$

VIEW(EXPLODE(1.2,1.2,1)(MKPOLS((V,[cell for cell,c in zip(CV,chain1) if c==1)

-)))) VIEW(EXPLODE(1.2,1.2,1)(MKPOLS((V,[cell for cell,c in zip(CV,chain2) if c==1]
-)))) VIEW(EXPLODE(1.2,1.2,1)(MKPOLS((V,[cell for cell,c1,c2 in zip(CV,chain1,chain2)
- if c1+c2==1])))) VIEW(EXPLODE(1.2,1.2,1)(MKPOLS((V,[cell for cell,c1,c2 in zip(CV,chain1,chain2) if c1+c2;=1])))) @

@O test/py/bool/test01.py @import sys """ import modules from larcc/lib """ sys.path.insert(0, 'lib/py/') from bool import * @i First set of 2D data: Fork-0 input @i @i Bulk of Boolean task computation @i @

7.1.2 Two squares

```
@o test/py/bool/test03.py @""" import modules from larcc/lib """ import sys sys.path.insert(0, 'lib/py/') from bool import *
```

$$V1 = [[0,0],[10,0],[10,10],[0,10]] \text{ FV1} = [\text{range}(4)] \text{ EV1} = [[0,1],[1,2],[2,3],[3,0]] \text{ VV1} = \text{AA(LIST)}(\text{range}(\text{len}(\text{V1})))$$

$$V2 = [[2.5, 2.5], [12.5, 2.5], [12.5, 12.5], [2.5, 12.5]] \; FV2 = [range(4)] \; EV2 = [[0, 1], [1, 2], [2, 3], [3, 0]]$$

$$VV2 = AA(LIST)(range(len(V2)))$$
 @; Bulk of Boolean task computation @; @

@o test/py/bool/test04.py @""" import modules from larcc/lib""" import sys sys.path.insert(0, 'lib/py/') from bool import *

$$V1 = [[0,0],[10,0],[10,10],[0,10]] \text{ FV1} = [\text{range}(4)] \text{ EV1} = [[0,1],[1,2],[2,3],[3,0]] \text{ VV1} = \text{AA}(\text{LIST})(\text{range}(\text{len}(\text{V1})))$$

$$V2 = [[2.5, 2.5], [7.5, 2.5], [7.5, 7.5], [2.5, 7.5]] \text{ FV2} = [\text{range}(4)] \text{ EV2} = [[0,1], [1,2], [2,3], [3,0]]$$

$$VV2 = AA(LIST)(range(len(V2)))$$
 @; Bulk of Boolean task computation @; @

@o test/py/bool/test05.py @""" import modules from larcc/lib""" import sys sys.path.insert(0, 'lib/py/') from bool import *

$$V1 = [[2.5, 2.5], [7.5, 2.5], [7.5, 7.5], [2.5, 7.5]] \text{ FV1} = [\text{range}(4)] \text{ EV1} = [[0, 1], [1, 2], [2, 3], [3, 0]]$$

 $VV1 = AA(LIST)(\text{range}(\text{len}(V1)))$

$$V2 = [[2.5, 2.5], [7.5, 2.5], [7.5, 7.5], [2.5, 7.5]] \text{ FV2} = [\text{range}(4)] \text{ EV2} = [[0,1], [1,2], [2,3], [3,0]]$$

$$VV2 = AA(LIST)(range(len(V2)))$$
 @; Bulk of Boolean task computation @; @

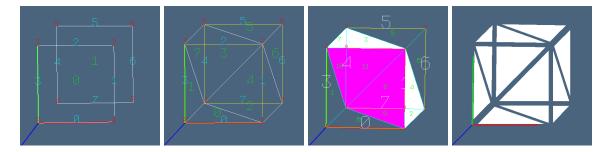


Figure 4: Partitioning of the CDC (Common Delaunay Complex): (a) the two Boolean arguments merged in a single covering; (b) the CDC together with the two (yellow) boundaries; (c) the split CDC cells; (d) the exploded CDC partition.

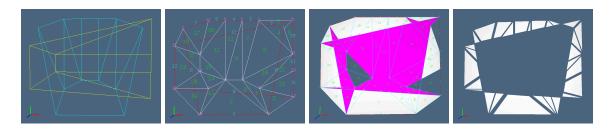


Figure 5: Partitioning of the CDC (Common Delaunay Complex): (a) the two Boolean arguments merged in a single covering; (b) the CDC together with the two (yellow) boundaries; (c) the split CDC cells; (d) the XOR of Boolean arguments.

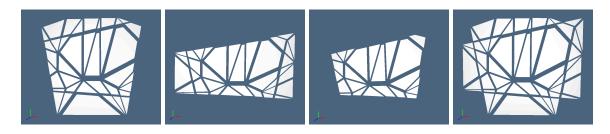


Figure 6: Some chains defined on the CDC (Common Delaunay Complex): (a) the first Boolean argument; (b) the second Boolean argument; (c) the intersection chain; (d) the union chain.

A Appendix: utility functions

@D Initial import of modules @from pyplasm import * from scipy import * import sys """ import modules from larcc/lib """ sys.path.insert(0, 'lib/py/') from lar2psm import * from simplexn import * from larcc import * from largrid import * from myfont import * from mapper import * @

A.1 Numeric utilities

A small set of utility functions is used to transform a point representation as array of coordinates into a string of fixed format to be used as point key into python dictionaries.

@D Symbolic utility to represent points as strings @""" TODO: use package Decimal (http://docs.python.org/2/library/decimal.html) """ global PRECISION PRECISION = 4

```
def prepKey (args): return "["+", ".join(args)+"]" def fixedPrec(value): out = round(value*10**PRECISION)/10**PRECISION if out == -0.0: out = 0.0 return str(out)
```

def vcode (vect): """ To generate a string representation of a number array. Used to generate the vertex keys in PointSet dictionary, and other similar operations. """ return prepKey(AA(fixedPrec)(vect)) @