Boolean combinations of cellular complexes as chain operations *

Alberto Paoluzzi

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1 Introduction

In this module a novel approach to Boolean operations of cellular complexes is defined and implemented. The novel algorithm may be summarised as follows.

First we compute the CDC (Common Delaunay Complex) of the input LAR complexes A and B, to get a LAR of the *simplicial* CDC.

Then, we split the cells intersecting the boundary faces of the input complexes, getting the final *polytopal* SCDC (Split Common Delaunay Complex), whose cells provide the basis for the linear coordinate representation of both input complexes, upon the same space decomposition.

Afterwards, every Boolean result is computed by bitwise operations, between the coordinate representations of the transformed A and B input.

Finally a greedy assembly of SCDC cells is executed, in order to return a polytopal complex with a reduced number of cells.

1.1 Preview of the Boolean algorithm

The goal is the computation of $A \diamond B$, with $\diamond \in \{\cup, \cap, -\}$, where a LAR representation of both A and B is given. The Boolean algorithm works as follows.

- 1. Embed both cellular complexes A and B in the same space (say, identify their common vertices) by $V_{ab} = V_a \cup V_b$.
- 2. Build their CDC (Common Delaunay Complex) as the LAR of *Delaunay triangulation* of the vertex set V_{ab} , and embedded ∂A and ∂B in it.
- 3. Split the (highest-dimensional) cells of CDC crossed by ∂A or ∂B . Their lower dimensional faces remain partitioned accordingly. We name the resulting complex SCDC (Split Common Delaunay Complex).
- 4. With respect to the SCDC basis of d-cells C_d , compute two coordinate chains α, β : $C_d \to \{0, 1\}$, such that:

```
\alpha(cell) = 1 if |cell| \subset A; else \alpha(cell) = 0, \beta(cell) = 1 if |cell| \subset B; else \beta(cell) = 0.
```

5. Extract accordingly the SCDC chain corresponding to $A \diamond B$, with $\diamond \in \{\cup, \cap, -\}$.

1.2 Remarks

You may make an analogy between the SCDC (*Split* CDC) and a CDT (Constrained Delaunay Triangulation). In part they coincide, but in general, the SCDC is a polytopal complex, and is not a simplicial complex as the CDC.

The more complex algorithmic step is the cell splitting. Every time, a single d-cell c is split by a single hyperplane (cutting its interior) giving either two splitted cells c_1 and c_2 , or just one output cell (if the hyperplane is the affine hull of the CDC facet) whatever the input cell dimension d. After every splitting of the cell interior, the row c is substituted (within the CV matrix) by c_1 , and c_2 is added to the end of the CV matrix, as a new row.

The splitting process is started by "splitting seeds" generated by (d-1)-faces of both operand boundaries. In fact, every such face, say f, has vertices on CDC and may split some incident CDC d-cell. In particular, starting from its vertices, f must split the CDC cells in whose interior it passes though.

So, a dynamic data structure is set-up, storing for each boundary face f the list of cells it must cut, and, for every CDC d-cell with interior traversed by some such f, the list of cutting faces. This data structure is continuously updated during the splitting process, using the adjacent cells of the split ones, who are to be split in turn. Every split cell may add some adjacent cell to be split, and after the split, the used pair (cell,face) is removed. The splitting process continues until the data structure becomes empty.

Every time a cell is split, it is characterized as either internal (1) or external (0) to the used (oriented) boundary facet f, so that the two resulting subcells c_1 and c_2 receive two opposite characterization (with respect to the considered boundary).

At the very end, every (polytopal) SCDC d-cell has two bits of information (one for argument A and one for argument B), telling whether it is internal (1) or external (0) or unknown (-1) with respect to every Boolean argument.

A final recursive traversal of the SCDC, based on cell adjacencies, transforms every -1 into either 0 or 1, providing the two final chains to be bitwise operated, depending on the Boolean operation to execute.

2 Step 1: merging discrete spaces

2.1 Requirements

The *join* of two sets $P, Q \subset \mathbb{E}^d$ is the set $PQ = \{\alpha \mathbf{x} + \beta \mathbf{y} | \mathbf{x} \in P, \ \mathbf{y} \in Q\}$, where $\alpha, \beta \in \mathbb{R}$, $\alpha, \beta \geq 0$, and $\alpha + \beta = 1$. The join operation is associative and commutative.

Input Two LAR models of two non-empty "solid" d-spaces A and B, denoted as (V1,CV1) and (V2,CV2).

Output The LAR representation (V,CV) of Delaunay triangulation (simplicial d-complex) of the set conv $AB \subset \mathbb{E}^d$, convex hull of the join of A and B, named Common Delaunay Complex (CDC) in the following.

Auxiliary data structures This software module returns also:

- a dictionary vertDict of V vertices, with key the symbolic representation of vertices v returned by expressions vcode(v), v ∈ V, and with values the finite ordinal numbers of the vertices;
- 2. the numbers n1, n12, n2 of the elements of V1, $V1 \cap V2$, and V2, respectively. Notice that the following assertions must hold (see Figure 1):

$$n1 - n12 + n2 = n \tag{1}$$

$$0 < \mathbf{n} - \mathbf{n} 2 \le \mathbf{n} 1 < \mathbf{n} \tag{2}$$

3. the input boundary complex (V,BC), with BC = BC1+BC2, i.e. the union of the two boundary (d-1)-complexes (V,BC1) and (V,BC2), defined on the common vertices.

2.2 Implementation

2.2.1 Summary

 $\langle \text{First Boolean step 4a} \rangle \equiv$

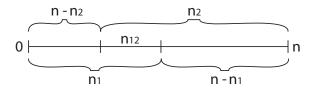


Figure 1: Relationships inside the orderings of CDC vertices

```
""" First Boolean step """
def larBool1():
    V, CV1,CV2, n1,n12,n2 = mergeVertices(model1,model2)
    VV = AA(LIST)(range(len(V)))
    V,CV,vertDict,n1,n12,n2,BC,nbc1,nbc2 = makeCDC(arg1,arg2, brep)
    W,CW,VC,BCellCovering,cellCuts,boundary1,boundary2,BCW = makeSCDC(V,CV,BC,nbc1,nbc2)
    assert len(VC) == len(V)
    assert len(BCellCovering) == len(BC)
    return W,CW,VC,BCellCovering,cellCuts,boundary1,boundary2,BCW
```

Macro referenced in 25.

2.2.2 Detail functions

```
⟨Compute model boundaries of complex of convex cells 4b⟩ ≡
    """ Compute model boundaries of complex of convex cells """

def larFacetsOfPolytopalComplex(vertDict,cells,facets):
    (V1,CV1),(V2,CV2) = model1,model2
    for cell in CV1:
        Vcell = [V1[v] for v in cell]

◇

Macro never referenced.

⟨Merge two dictionaries with keys the point locations 4c⟩ ≡
    """ Merge two dictionaries with keys the point locations """
    def mergeVertices(model1, model2):
        (V1,CV1),(V2,CV2) = model1, model2

        n = len(V1); m = len(V2)
        def shift(CV, n):
            return [[v+n for v in cell] for cell in CV]
        CV2 = shift(CV2,n)
```

```
vdict1 = defaultdict(list)
        for k,v in enumerate(V1): vdict1[vcode(v)].append(k)
        vdict2 = defaultdict(list)
        for k,v in enumerate(V2): vdict2[vcode(v)].append(k+n)
        vertDict = defaultdict(list)
        for point in vdict1.keys(): vertDict[point] += vdict1[point]
        for point in vdict2.keys(): vertDict[point] += vdict2[point]
        case1, case12, case2 = [],[],[]
        for item in vertDict.items():
           key, val = item
           if len(val)==2: case12 += [item]
           elif val[0] < n: case1 += [item]</pre>
           else: case2 += [item]
        n1 = len(case1); n2 = len(case12); n3 = len(case2)
        invertedindex = list(0 for k in range(n+m))
        for k,item in enumerate(case1):
           invertedindex[item[1][0]] = k
        for k,item in enumerate(case12):
           invertedindex[item[1][0]] = k+n1
           invertedindex[item[1][1]] = k+n1
        for k, item in enumerate(case2):
           invertedindex[item[1][0]] = k+n1+n2
        V = [eval(p[0]) \text{ for p in case1}] + [eval(p[0]) \text{ for p in case12}] + [eval(p[0]) \text{ for p in case12}]
                  p[0]) for p in case2]
        CV1 = [sorted([invertedindex[v] for v in cell]) for cell in CV1]
        CV2 = [sorted([invertedindex[v] for v in cell]) for cell in CV2]
        return V,CV1,CV2, n1+n2,n2,n2+n3
Macro referenced in 27.
\langle Make Common Delaunay Complex 5\rangle \equiv
     """ Make Common Delaunay Complex """
     def makeCDC(arg1,arg2, brep):
        (V1,basis1), (V2,basis2) = arg1,arg2
        (facets1,cells1),(facets2,cells2) = basis1[-2:],basis2[-2:]
        model1, model2 = (V1,cells1),(V2,cells2)
        V, _,_, n1,n12,n2 = mergeVertices(model1, model2)
        n = len(V)
        assert n == n1 - n12 + n2
```

```
CV = sorted(AA(sorted)(Delaunay(array(V)).simplices))

vertDict = defaultdict(list)
for k,v in enumerate(V): vertDict[vcode(v)] += [k]

if brep == False:
    signs1,BC1 = signedCellularBoundaryCells(V1,basis1)

BC1pairs = zip(*signedCellularBoundaryCells(V1,basis1))
BC1 = [basis1[-2][face] if sign>0 else swap(basis1[-2][face]) for (sign,face) in BC1pair

BC2pairs = zip(*signedCellularBoundaryCells(V2,basis2))
BC2 = [basis2[-2][face] if sign>0 else swap(basis2[-2][face]) for (sign,face) in BC2pair

else:
    BC1,BC2 = basis1[-1],basis2[-1]

BC = [[ vertDict[vcode(V1[v])][0] for v in cell] for cell in BC1] + [
        [ vertDict[vcode(V2[v])][0] for v in cell] for cell in BC2] #+ qhullBoundary(V)

return V,CV,vertDict,n1,n12,n2,BC,len(BC1),len(BC2)
```

Macro referenced in 27.

 \Diamond

3 Step 2: splitting cells

The goal of this section is to transform the CDC simplicial complex, into the polytopal Split Common Delaunay Complex (SCDC), by splitting the d-cells of CDC crossed in their interior by some cell of the input boundary complex.

3.1 Requirements

We call here for a sequential implementation, following every (d-1)-facet lambda in BC (for Boundary Cells). We start the splitting with COVECTOR(lambda) from cell, one of the CDC d-cells incident on a vertex of lambda, and continue the splitting on the d-cells (d-1)-adjacent to cell, where (a) COVECTOR(lambda) either crosses the cell's interior or contains one of cell's (d-1)-facets and (b) such that the intersection with lambda is not empty, until the queue (or stack) of d-cells to intersect with covector is not empty.

Best computational strategy First associate to each cutting facet the list of cells it may cut; then execute all the cuts. In this way we can compute the adjacency matrix just one time at the beginning of the procedure, and do not need to update it after every split.

Input The output of previous algorithm stage.

Output The LAR representation (W,PW) of the SCDC,

Auxiliary data structures This software module returns also a dictionary splitFacets, with keys the input boundary faces and values the list of pairs (covector, fragmentedFaces).

3.2 Implementation

3.2.1 Summary

```
\(\text{Second Boolean step 7a} \) \equiv \(
\text{""" Second Boolean step """} \\
\text{def larBool2(boundary1,boundary2):} \\
\text{dim = len(W[0])} \\
\text{WW = AA(LIST)(range(len(W)))} \\
\text{FW = convexFacets (W,CW)} \\
\text{_,EW = larFacets((W,FW), dim=2)} \\
\text{boundary1,boundary2,FWdict = makeFacetDicts(FW,boundary1,boundary2)} \\
\text{if dim == 3:} \\
\text{_,EW = larFacets((W,FW), dim=2)} \\
\text{bases = [WW,EW,FW,CW]} \\
\text{elif dim == 2: bases = [WW,FW,CW]} \\
\text{elif dim == 2: bases,boundary1,boundary2,FW,BCW}
\(\phi\)
\
```

Macro referenced in 25.

3.2.2 Detail functions

Computing the adjacent cells of a given cell To perform this task we make only use of the CV list. In a more efficient implementation we should make direct use of the sparse adjacency matrix, to be dynamically updated together with the CV list. The computation of the adjacent d-cells of a single d-cell is given here by extracting a column of the $CSR(M_d M_d^t)$. This can be done by multiplying $CSR(M_d)$ by its transposed row corresponding to the query d-cell.

```
⟨ Computing the adjacent cells of a given cell 7b⟩ ≡

""" Computing the adjacent cells of a given cell """

def adjacencyQuery (V,CV):
    dim = len(V[0])
    csrCV = csrCreate(CV)
    csrAdj = matrixProduct(csrCV,csrTranspose(csrCV))
    def adjacencyQueryO (cell):
```

```
nverts = len(CV[cell])
  cellAdjacencies = csrAdj.indices[csrAdj.indptr[cell]:csrAdj.indptr[cell+1]]
  return [acell for acell in cellAdjacencies if dim <= csrAdj[cell,acell] < nverts]
  return adjacencyQueryO
</pre>
```

Macro referenced in 27.

Relational inversion (characteristic matrix transposition) The operation could be executed by simple matrix transposition of the CSR (Compressed Sparse Row) representation of the sparse characteristic matrix $M_d \equiv \text{CV}$. A simple relational inversion using Python lists is given here. The invertRelation function is given here, linear in the size of the CV list, where the complexity of each cell is constant and small in most cases.

Computation of splitting tests In order to compute, in the simplest and more general way, whether each of the two split d-cells is internal or external to the splitting boundary d-1-facet, it is necessary to consider the oriented covector ϕ (or one-form) canonically associated to the facet f by the covector representation theorem, i.e. the corresponding oriented hyperplane. In this case, the internal/external attribute of the split cell will be computed by evaluating the pairing $\langle v, \phi \rangle$.

```
\( \text{Splitting tests 8b} \) \( \)
\( \text{""" Splitting tests """} \)
\( \text{def testingSubspace(V,covector):} \)
\( \text{def testingSubspace0(vcell):} \)
\( \text{inout = SIGN(sum([INNERPROD([[1.]+V[v],covector]) for v in vcell]))} \)
\( \text{return inout} \)
\( \text{return testingSubspace0} \)
\( \)
```

Macro referenced in 9.

Elementary splitting test Let us remember that the adjacency matrix between d-cells is computed via SpMSpM multiplication by the double application

```
adjacencyQuery(V,CV)(cell),
```

where the first application adjacencyQuery(V,CV) returns a partial function with bufferisation of the adjacency matrix, and the second application to cell returns the list of adjacent d-cells sharing with it a (d-1)-dimensional facet.

```
⟨Elementary splitting test 9⟩ ≡

⟨Splitting tests 8b⟩

""" Elementary splitting test """

def dividenda(V,CV, cell,facet,covector,unchosen):
    out = []
    adjCells = adjacencyQuery(V,CV)(cell)
    for adjCell in set(adjCells).difference(unchosen):
        if (cuttingTest(covector,CV[adjCell],V) and \
            cellFacetIntersecting(facet,adjCell,covector,V,CV)) or \
            tangentTest(covector,facet,CV[adjCell],V): out += [adjCell]
        return out
```

Macro referenced in 27.

CDC cell splitting with one or more facets When splitting a d-cell with some hyperplanes, we need to return not only either the two cut parts or the cell itself when the hyperplane is tangent to a (d-1)-face, but also the facet lying on the hyperplane. In the first cade it is directly computed by the SPLITCELL function, and returned as the equal set of points. In the second case, the cell is transformed by the map that sends the hyperplane in the $x_d = 0$ subspace (z = 0 in 3D), and the searched facet is returned as the (back-transformed) set of cell vertices on this subspace.

Actually, the process is strongly complicated by the fact that the input cell (and its facets) may be cut by several hyperplanes. By now, we resort to the simplest computation, even if more time-expensive: to compare each vertex of each cell fragment, against every hyperplanes. This approach will adapt well to the writing of a computational kernel on the GPU.

```
\langle \text{CDC cell splitting with one or more cutting facets } 10a \rangle \equiv
     """ CDC cell splitting with one or more cutting facets """
     def fragment(cell,cellCuts,V,CV,BC):
        vcell = CV[cell]
        cellFragments = [[V[v] for v in vcell]]
        for f in cellCuts[cell]:
            facet = BC[f]
            plane = COVECTOR([V[v] for v in facet])
            for k,fragment in enumerate(cellFragments):
               #if not tangentTest(plane,facet,fragment,V):
               [below,equal,above] = SPLITCELL(plane,fragment,tolerance=1e-4,ntry=4)
               if below != above:
                  cellFragments[k] = below
                  cellFragments += [above]
            facets = facetsOnCuts(cellFragments,cellCuts,V,BC)
        return cellFragments
Macro referenced in 27.
```

SCDC splitting with every boundary facet The function makeSCDC is used to compute the LAR model (W,CW) of the SCDC. It takes as input the LAR model (V,CV) of the CDC, and the LAR model (V,BC) of the input Boolean Complex, and returns both a new LAR model (W,CW) and the vertex-cell relation VC, i.e. the transposed of CV.

```
For every k \in BC, a list cellsToSplit
```

```
⟨ SCDC splitting with every boundary facet 10b ⟩ ≡
   """ SCDC splitting with every boundary facet """
   def makeSCDC(V,CV,BC,nbc1,nbc2):
```

```
index, defaultValue = -1, -1
VC = invertRelation(CV)
CW,BCfrags = [],[]
Wdict = dict()
BCellcovering = boundaryCover(V,CV,BC,VC)
cellCuts = invertRelation(BCellcovering)
for k in range(len(CV) - len(cellCuts)): cellCuts += [[]]
def verySmall(number): return abs(number) < 10**-5.5</pre>
for k,frags in enumerate(cellCuts):
   if cellCuts[k] == []:
      cell = []
      for v in CV[k]:
         key = vcode(V[v])
         if Wdict.get(key,defaultValue) == defaultValue:
            index += 1
            Wdict[key] = index
            cell += [index]
         else:
            cell += [Wdict[key]]
      CW += [cell]
   else:
      cellFragments = fragment(k,cellCuts,V,CV,BC)
      for cellFragment in cellFragments:
         cellFrag = []
         for v in cellFragment:
            key = vcode(v)
            if Wdict.get(key,defaultValue) == defaultValue:
               index += 1
               Wdict[key] = index
               cellFrag += [index]
               cellFrag += [Wdict[key]]
         CW += [cellFrag]
         BCfrags += [ (h, [Wdict[vcode(w)] for w in cellFragment if verySmall(
                     PROD([ COVECTOR( [V[v] for v in BC[h]] ), [1.]+w ])) ] )
                   for h in cellCuts[k]]
BCW = [ [ Wdict[vcode(V[v])] for v in cell ] for cell in BC]
W = sorted(zip( Wdict.values(), Wdict.keys() ))
W = AA(eval)(TRANS(W)[1])
dim = len(W[0])
boundary1,boundary2 = boundaryEmbedding(BCfrags,nbc1,dim)
```

```
Macro referenced in 27.
\langle Boolean argument boundaries embedding in SCDC 11 \rangle \equiv
     """ Boolean argument boundaries embedding in SCDC """
     def boundaryEmbedding(BCfrags,nbc1,dim):
        boundary1,boundary2 = defaultdict(list),defaultdict(list)
        for h, frags in BCfrags:
            if h < nbc1: boundary1[h] += [frags]</pre>
            else: boundary2[h] += [frags]
        boundarylist1,boundarylist2 = [],[]
        for h,facets in boundary1.items():
            boundarylist1 += [(h, AA(eval)(set([str(sorted(f))
                            for f in facets if len(set(f)) >= dim])) )]
        for h, facets in boundary2.items():
            boundarylist2 += [(h, AA(eval)(set([str(sorted(f))
                            for f in facets if len(set(f)) >= dim])) )]
        boundary1,boundary2 = dict(boundarylist1),dict(boundarylist2)
        return boundary1, boundary2
Macro referenced in 27.
\langle Make facets dictionaries 12a\rangle \equiv
     """ Make facets dictionaries """
     def makeFacetDicts(FW,boundary1,boundary2):
        FWdict = dict()
        for k,facet in enumerate (FW): FWdict[str(facet)] = k
        for key,value in boundary1.items():
            value = [FWdict[str(facet)] for facet in value]
           boundary1[key] = value
        for key,value in boundary2.items():
            value = [FWdict[str(facet)] for facet in value]
           boundary2[key] = value
        return boundary1, boundary2, FWdict
Macro referenced in 27.
Computation of boundary facets covering with CDC cells
\langle Computation of boundary facets covering with CDC cells 12b\rangle \equiv
     """ Computation of boundary facets covering with CDC cells """
     def boundaryCover(V,CV,BC,VC):
        cellsToSplit = list()
        boundaryCellCovering = []
```

return W,CW,VC,BCellcovering,cellCuts,boundary1,boundary2,BCW

```
for k,facet in enumerate(BC):
           covector = COVECTOR([V[v] for v in facet])
           seedsOnFacet = VC[facet[0]]
           cellsToSplit = [dividenda(V,CV, cell,facet,covector,[])
                          for cell in seedsOnFacet ]
           cellsToSplit = set(CAT(cellsToSplit))
           if cellsToSplit == set(): cellsToSplit=set(seedsOnFacet) ## NB !!!
           while True:
              newCells = [dividenda(V,CV, cell,facet,covector,cellsToSplit)
                          for cell in cellsToSplit ]
              if newCells != []: newCells = CAT(newCells)
              covering = cellsToSplit.union(newCells)
              if covering == cellsToSplit:
                 break
              cellsToSplit = covering
           boundaryCellCovering += [list(covering)]
        return boundaryCellCovering
Macro referenced in 27.
Cell-facet intersection test
\langle Cell-facet intersection test 13\rangle \equiv
     """ Cell-facet intersection test """
     def cellFacetIntersecting(boundaryFacet,cell,covector,V,CV):
        points = [V[v] for v in CV[cell]]
        vcell1,newFacet,vcell2 = SPLITCELL(covector,points,tolerance=1e-4,ntry=4)
        boundaryFacet = [V[v] for v in boundaryFacet]
        translVector = boundaryFacet[0]
        # translation
        newFacet = [ VECTDIFF([v,translVector]) for v in newFacet ]
        boundaryFacet = [ VECTDIFF([v,translVector]) for v in boundaryFacet ]
        # linear transformation: boundaryFacet -> standard (d-1)-simplex
        d = len(V[0])
        transformMat = mat( boundaryFacet[1:d] + [covector[1:]] ).T.I
        # transformation in the subspace x_d = 0
        newFacet = (transformMat * (mat(newFacet).T)).T.tolist()
        boundaryFacet = (transformMat * (mat(boundaryFacet).T)).T.tolist()
        # projection in E^{d-1} space and Boolean test
        newFacet = MKPOL([ AA(lambda v: v[:-1])(newFacet),
                           [range(1,len(newFacet)+1)], None ])
        boundaryFacet = MKPOL([ AA(lambda v: v[:-1])(boundaryFacet),
```

```
[range(1,len(boundaryFacet)+1)], None ])
verts,cells,pols = UKPOL(INTERSECTION([newFacet,boundaryFacet]))
if verts == []: return False
else: return True
```

4 Step 3: cell labeling

The goal of this stage is to label every cell of the SCDC with two bits, corresponding to the input spaces A and B, and telling whether the cell is either internal (1) or external (0) to either spaces.

4.1 Requirements

Macro referenced in 27.

Input The output of previous algorithmic stage.

Output The array cellLabels with *shape* len(PW) \times 2, and values in $\{0,1\}$.

4.2 Implementation

The labelling of LAR of the SCDC may be decomposed in five consecutive steps. The first step was actually executed during the splitting stage, by accumulating a single facet of every split cells embedded on the affine hull (the covector hyperplane) of the splitting boundary facet. The second step provides the computation of the sparse matrix of the linear coboundary operator $\delta_{d-1}: C_{d-1} \to C_d$. The third step operates upon the previous two pieces of information, in order to compute the coboundary chain of the boundary chain of both input Boolean arguments. The fourth step attaches a IN/OUT label to each d-cell of the previously computed d-chain. Finally, the fifth step spreads around the labels to cover all the d-cells of SCDC. This knowledge allows for the computation of every interesting Boolean expressions between the input complexes.

4.2.1 Summary

```
⟨Third Boolean step 14⟩ ≡
    """ Third Boolean step """
    def larBool3():
        coBoundaryMat = signedCellularBoundary(W,bases).T
        boundaryMat = coBoundaryMat.T
        CWbits = [[-1,-1] for k in range(len(CW))]
        CWbits = cellTagging(boundary1,boundaryMat,CW,FW,W,BCW,CWbits,0)
        CWbits = cellTagging(boundary2,boundaryMat,CW,FW,W,BCW,CWbits,1)
```

```
for cell in range(len(CW)):
    if CWbits[cell][0] == 1:
        CWbits = booleanChainTraverse(0,cell,W,CW,CWbits,1)
    if CWbits[cell][0] == 0:
        CWbits = booleanChainTraverse(0,cell,W,CW,CWbits,0)
    if CWbits[cell][1] == 1:
        CWbits = booleanChainTraverse(1,cell,W,CW,CWbits,1)
    if CWbits[cell][1] == 0:
        CWbits = booleanChainTraverse(1,cell,W,CW,CWbits,0)
    chain1,chain2 = TRANS(CWbits)

chain = [k for k,cell in enumerate(chain1) if cell==1]
    _,bound1 = chain2complex(W,CW,chain,boundaryMat)

chain = [k for k,cell in enumerate(chain2) if cell==1]
    _,bound2 = chain2complex(W,CW,chain,boundaryMat)

return W,CW,FW,boundaryMat,bound1,bound2,chain1,chain2,CWbits
```

Macro referenced in 25.

Macro referenced in 27.

4.2.2 Detail functions

Computation of boundary cells embedded in SCDC

Coboundary operator on SCDC space decomposition In this section we develop a stronger characterisation of the boundaries, by fully tagging in SCDC the internal coboundary of boundaries of A and B Boolean arguments. This novel strategy should allow the recursive tagging extension to work correctly in all cases.

As we know, the coboundary operators $\delta_{k-1}: C_{k-1} \to C_k$ are the transpose of the boundary operators $\partial_k: C_k \to C_{k-1}$ $(1 \le k \le d)$. We therefore proceed to the construction of the operator δ_{d-1} , according to the procedure illustrated in []. For this purpose we need to use both the C_d and the C_{d-1} bases of SCDC. The first basis is generated as CV array

during the splitting. The second basis will be built from C_d using the proper d-adjacency algorithm from [].

Let us remember that a (co)boundary operator may be applied to any chain from the linear space of chains defined upon a cellular complex. In our case we have already generated the (d-1)-chains ∂A and ∂B while building the SCDC, by accumulating, in the course of the splitting phase, the (d-1)-facets discovered while tracking the boundaries of A and B. We just need now to tag (a subset of) $\delta_{d-1}\partial_d A$ and $\delta_{d-1}\partial_d B$.

Computation of facets of a connected polytopal LAR model The (d-1)-facets of a d-dimensional polytopal and convex LAR model are computed by the below convexFacets function by using both the algorithm codified by the function larFacets for the facets internal to the complex, and the convexBoundary algorithm for the facets on the convex boundary of the complex.

Conversely, the larConvexFacets function can be used to compute the (d-1)-facets of a possibly non-connected and/or non-convex polytopal complex.

```
\langle Coboundary operator on the convex decomposition of common space 16\rangle
     """ Coboundary operator on the convex decomposition of common space """
     from scipy.spatial import ConvexHull
     def qhullBoundary(V):
        points = array(V)
        hull = ConvexHull(points)
        out = hull.simplices.tolist()
        return sorted(out)
     def facetDimensionTest(V,facet,covector):
        covector = eval(covector)
        return all([-0.01 < INNERPROD([[1.]+W[v], covector]) < 0.01 for v in facet ])
     def convexFacets (V,CV,dim=2):
        dim = len(V[0])
        model = V,CV
        V,FV = larFacets(model,dim)
        FV = AA(eval)(list(set(AA(str)(AA(sorted)(FV + convexBoundary(V,CV)))))))
        return FV
     def larConvexFacets (V,CV):
        vdict = defaultdict(list)
        for k,v in enumerate(V): vdict[vcode(v)] += [k]
        FV = []
        for cell in CV:
           fv = convexFacets([V[v] for v in cell],[range(len(cell))])
           FV += [tuple([cell[v] for v in facet]) for facet in fv]
```

```
return sorted(AA(list)(set(FV)))

if __name__ == "__main__":
    V = [[0,0],[1,0],[1,1],[0.5,1],[0,1]]
    CV = [[0,1,2,3,4]]
    FV = convexFacets(V,CV)

if __name__ == "__main__":
    V,CV = larCuboids((10,10,10))
    FV = convexFacets(V,CV,2)
    #EV = convexFacets(V,FV,1)
```

Macro referenced in 27.

Computation of boundary operator The computation of the boundary operator ∂_d on the SCDC d-basis (W,CW) requires the knowledge of the (d-1)-basis (W,FW). The goal of this section is hence the—partially incremental—computation of FW. This set can be partitioned into *internal* cells, that have 2 cofaces, and *boundary* cells, that have only 1 coface. The first subset is easily computed by the larFacets function; the computation of the second subset requires some more work, specified in the following.

First, we compute the 0-chain of boundary vertices of the SCDC, using qHull, and take advantage of the CV matrix to extract the chain of d-cells sharing with the boundary a (d-1)-facet. Second, using the partial boundary operator generated by using only the interior (d-1)-facets, and the associated (d-2)-boundary operator, we select the sub-chain made by the non-closed d-cells of this subset. Third, the boundary facet of each of them is finally selected, added to the (d-1)-basis of SCDC, and the corresponding row is added at the bottom line of the matrix of ∂_{d-1} .

```
if all([-0.01 < INNERPROD([[1.]+V[v], covector]) < 0.01 for v in facet]):
                 boundaryFacets += [ facet ]
                 splitFacets = [[] for k in range(len(boundaryEquations))]
                 for v in facet:
                     for k,equation in enumerate(boundaryEquations):
                        if -0.01 < INNERPROD([V[v]+[1.], equation]) < 0.01:
                           splitFacets[k] += [v]
              boundaryFacets += [f for f in splitFacets if f != [] and len(f)>=dim ]
        return boundaryFacets
Macro referenced in 27.
Coboundary of boundary chains
\langle Coboundary of boundary chain 18a\rangle \equiv
     """ Coboundary of boundary chain """
Macro never referenced.
Labeling seeds
\langle Writing labelling seeds on SCDC 18b\rangle \equiv
     """ Writing labelling seeds on SCDC """
     def cellTagging(boundaryDict,boundaryMat,CW,FW,W,BC,CWbits,arg):
        dim = len(W[0])
        for face in boundaryDict:
           for facet in boundaryDict[face]:
              cofaces = list(boundaryMat[facet].tocoo().col)
              if len(cofaces) == 1:
                 CWbits[cofaces[0]][arg] = 1
              elif len(cofaces) == 2:
                 v0 = list(set(CW[cofaces[0]]).difference(FW[facet]))[0]
                  v1 = list(set(CW[cofaces[1]]).difference(FW[facet]))[0]
                 # take d affinely independent vertices in face (TODO: use pivotSimplices()
                  simplex0 = BC[face][:dim] + [v0]
                  simplex1 = BC[face][:dim] + [v1]
                  sign0 = sign(det([W[v]+[1] for v in simplex0]))
                 sign1 = sign(det([W[v]+[1] for v in simplex1]))
                 if sign0 == 1: CWbits[cofaces[0]][arg] = 1
                  elif sign0 == -1: CWbits[cofaces[0]][arg] = 0
                  if sign1 == 1: CWbits[cofaces[1]][arg] = 1
                  elif sign1 == -1: CWbits[cofaces[1]][arg] = 0
```

else:

Recursive diffusion of labels A recursive function booleanChainTraverse is given in the script below, where

Macro referenced in 27.

5 Step 4: greedy cell gathering

The goal of this stage is to make as lower as possible the number of cells in the output LAR of the space AB, partitioned into convex cells.

Input The LAR model (W,PW) of the SCDC and the array cellLabels.

Output The LAR representation (W,RW) of the final fragmented and labeled space AB.

5.1 Requirements

The algorithm proposed here for d-cell gathering into bigger polytopes is local and greedy. Starting from an initial random d-cell, a (d-1)-connected d-chain is built, by attaching, one at a time, single cells to the boundary of the chain, after (local) verification that the support of the new chain will remain a convex set.

In case of failure of the test, the facets of the current chain boundary are checked for the gluing of their adjacent and external d-coface, until either a new convex is built, or no single cell can be attached convexly, so that the attachment process relative to that chain stops, and its boundary vertices are written in the LAR of a new complex, to gather a single new polytope generated by them.

Actually, during the stage of boundary checking for finding a new cocell to glue, only a subchain is checked, obtained by subtraction from the boundary of the cutting facets, where attachments are not possible, without violating the topology of Boolean results.

Two main algorithm components are needed here. The first one concerns the extraction of the current d-chain boundary, the subtraction from it of the splitting facets, and the selection of the facet where to glue another d-cell; the second one deals with the convexity test of the candidate (chain + boundary cocell) pair.

The local convexity test will extract, using the (co)boundary matrix of the current chain, the coboundary of the boundary of the candidate facet, and, selected the matrix of hyperplanes associated to it, will compute the centroid of the facet and the vector of signs exposed by the point transformed by right product with this matrix. The local test of convexity is satisfied if and only if all new vertices expose the same signs (or zero), when transformed by this matrix. In other words, the test is satisfied if all new vertices remain internal (or non external) to the cone generated by such set of boundary hyperplanes.

Every time that a new cell has been selected to join the current chain, the cell is also signed as already used, and hence as no more available for other choices. Of course, the algorithm terminates when all the input d-cells have been selected and signed.

5.2 Implementation

A synthetic view of the simplification process is given by the script below. The first tool provides a mapping from (d-1)-facets of SCDC to their embedding hyperplanes, i.e. to their affine hulls of codimension 1. The second one compute the boundary (d-1)-complex of the SCDC d-chain currently transformed into a single convex cell. The algorithmic bulk of the simplification process is contained in the script entitled Sticking cells together. The last script provides the high-level interface to transform the generated SCDC into a strongly simplified polytopal complex.

High-level description

```
⟨Simplification of the output polytopal complex 20a⟩ ≡

⟨Mapping from facets to hyperplanes 21a⟩

⟨Building the boundary complex of the current chain 21b⟩

⟨Sticking cells together 21c⟩

⟨Gathering and writing a polytopal complex 23⟩

⋄

Macro referenced in 27.
```

5.2.1 Summary

```
\langle Fourth Boolean step 20b \rangle \equiv
```

```
""" Fourth Boolean step """
def larBool4(W,CW,FW,boundaryMat,boundary1,boundary2,CWbits):
    X,CX,CXbits = gatherPolytopes(W,CW,FW,boundaryMat,boundary1,boundary2,CWbits)
    FX = larConvexFacets (X,CX)
    return X,CX,FX,CXbits
```

Macro referenced in 25.

5.2.2 Detail functions

Mapping from facets to hyperplanes The function facet2covectors return the list of hyperplane covectors, with first term homogeneous, i.e. the row vector (c, a, b) for the line equation ax + by + c = 0, or the row vector (d, a, b, c) for the plane equation ax + by + cz + d = 0.

```
⟨Mapping from facets to hyperplanes 21a⟩ ≡
    """ Mapping from hyperplanes to lists of facets """
    def facet2covectors(W,FW):
        return [COVECTOR([W[v] for v in facet]) for facet in FW]

def boundaries(boundary1,boundary2):
    #return set(CAT(boundary1.values() + boundary2.values()))
    return boundary1.union(boundary2)
    ◊

Macro referenced in 20a, 27.
```

Building the complex of the current chain The function chain2complex returns the boundary complex of the current chain, minus the facet in constraints, where non d-cell may be attached to the current chain. It is computed via multiplication between the matrix of boundary operator and the coordinate representation chainCoords of chain. The constraint set is finally subtracted to the result.

```
⟨Building the boundary complex of the current chain 21b⟩ ≡
   from scipy.sparse import csc_matrix
""" Building the boundary complex of the current chain """
   def chain2complex(W,CW,chain,boundaryMat,constraints=[]):
        chainCoords = csc_matrix((len(CW), 1))
        for cell in chain: chainCoords[cell] = 1
        boundaryCells = set((boundaryMat * chainCoords).tocoo().row)
        envelope = boundaryCells.difference(constraints)
        return envelope,boundaryCells
```

Macro referenced in 20a.

Sticking cells together

```
⟨Sticking cells together 21c⟩ ≡

""" Sticking cells together """

⟨Testing the convexity of a single added vertex 22a⟩

⟨Testing the convexity when attaching a cell to a chain 22b⟩

⟨Elongate a chain while supports a convex set 22c⟩

◇

Macro referenced in 20a.
```

Testing the convexity of a single added vertex A single cell is possibly attached to the boundary envelope of the current chain. In case of success the function protrudeChain returns True; otherwise returns False.

Testing the convexity of current chain

Macro referenced in 21c.

Chain elongation while is convex

```
\langle Elongate a chain while supports a convex set 22c\rangle \equiv
     """ Elongate a chain while supports a convex set """
     def protrudeChain (W,CW,FW,chain,boundaryMat,covectors,usedCells,constraints):
        verts = []
        while True:
           changed = False
           envelope,boundaryFacets = chain2complex(W,CW,chain,boundaryMat,constraints)
           for facet in envelope:
              success = False
              chainCoords = csr_matrix((1,len(FW)))
              chainCoords[0,facet] = 1
              cocells = list((chainCoords * boundaryMat).tocoo().col)
              if len(cocells)==2:
                 if cocells[0] in chain: cell = cocells[1]
                 elif cocells[1] in chain: cell = cocells[0]
                 if not usedCells[cell]:
                     success = testAttachment(cell,usedCells,facet,chain, \
                              W,CW,FW,boundaryMat,boundaryFacets,covectors)
                 if success:
                     changed = True
                    usedCells[cell] = True
                    chain += [cell]
              else: print "error: in protrudeChain (len(cocells) not equal to 2)"
              chainCoords = csc_matrix((len(CW),1))
              for cell in chain:
                 chainCoords[cell,0] = 1
                 usedCells[cell] = True
              boundaryFacets = list((boundaryMat*chainCoords).tocoo().row)
           if not changed: break
        verts = [FW[facet] for facet in boundaryFacets]
        verts = sorted(list(set(CAT(verts))))
        return verts, usedCells
```

Macro referenced in 21c.

Gathering and writing a polytopal complex The task of the gatherPolytopes function, given below, is to return the LAR (X,CX) of the SCDC (W,CW) generated by the previous phases of the Boolean algorithm, after reducing its representation to a much smaller size, (a) by gathering subsets of cells into single bigger polytopal cells within the characteristic matrix CX, and (b) by assembling their boundary vertices into the (reduced) vertex set X. Of course, while reducing the number of polytopal cells, the procedure should

not change the Boolean structure of the input complex, i.e. the support spaces $|C_A|, |C_B|$ of proper chains C_A and C_B and of their Boolean combinations.

```
\langle Gathering and writing a polytopal complex 23\rangle \equiv
     """ Gathering and writing a polytopal complex """
     def gatherPolytopes(W,CW,FW,boundaryMat,bounds1,bounds2,CWbits):
        usedCells = [False for cell in CW]
        covectors = facet2covectors(W,FW)
        constraints = boundaries(bounds1,bounds2)
        Xdict,index,CX,defaultValue,CXbits = dict(),0,[],-1,[]
        while not all(usedCells):
           for k,cell in enumerate(CW):
               if not usedCells[k]:
                  chain = [k]
                  usedCells[k] = True
                  verts,usedCells = protrudeChain(W,CW,FW,chain,boundaryMat,
                                  covectors, usedCells, constraints)
                  CX += [ verts ]
                  CXbits += [ CWbits[k] ]
        X,CX = larRemoveVertices(W,CX)
        return X,CX,CXbits
        #return W,CX,CXbits
```

Macro referenced in 20a.

5.2.3 Final removal of redundant vertices

After the simplification step, that replaces two or more convex d-cells with a single one cell, and the subsequent computation of the facets of the new cells, some vertices may become redundant, since are not intersection of at least d affine hulls supporting the (d-1)-facets.

Removal of redundant vertices from simplified LAR model The input to the function larVertexRemoval is the triple X, CX, FX of vertices, cells-by-vertices, and facets-by-vertices, where some of vertices may be redundant. Therefore, for each vertex, we compute the subset of incident facets, and then the subset of supporting covectors, i.e. the affine functions defining their affine hulls. If their number is greater or equal to the dimension of the embedding space, i.e. to the number of coordinates of vertices, then the vertex is non-redundant, and cannot be eliminated from the LAR model. If the vertex k is redundant, the corresponding value X[k] is substituted by the empty list. At the very end a rewriting vertex dictionary is generated and used to produce the novel output V, CV, and FV. Let us notice that the array affineHullNumber represents, for each vertex, the number of incident affine hulls. Hence, when affineHullNumber[k] is less that dim, the vertex X[k] is redundant, and can be eliminated.

```
\langle Removal of redundant vertices from simplified LAR model 24\rangle \equiv
     """ Removal of redundant vertices from simplified LAR model """
     def facetCovectors(X,FX):
         covectors = defaultdict(list)
         for k, facet in enumerate(FX):
            covect = list(COVECTOR([X[v] for v in facet]))
            normalizedCovect = UNITVECT([ h*SIGN(covect[0]) for h in covect])
            for h,comp in enumerate(normalizedCovect):
               if not isclose(0.0, comp):
                  theSign = SIGN(comp)
            normalizedCovect = [x*theSign if x!=abs(0.0) else x for x in normalizedCovect]
            covectors[vcode(normalizedCovect)] += [k]
         return covectors
     def larVertexRemoval(X,CX,FX):
         dim = len(X[0])
         covectors = facetCovectors(X,FX)
         CovectF = covectors.values()
         FCovect = invertRelation(CovectF)
         XF = invertRelation(FX)
         affineHullNumber = [len([FCovect[face] for face in vertFaces]) for vertFaces in XF]
         Y = [X[k] if val>=dim else [] for k,val in enumerate(affineHullNumber)]
        newIndex, Z = 0, dict()
         for oldIndex, vertex in enumerate(Y):
            if vertex != []:
               Z[oldIndex] = newIndex # (old,new) vertex indices
               newIndex += 1
         V = [None for k in range(len(Z))]
         for old, new in Z.items():
            V[new] = X[old]
         FV = [[Z[v] \text{ for } v \text{ in facet if } v \text{ in } Z] \text{ for facet in } FX]
         CV = [[Z[v] \text{ for } v \text{ in cell if } v \text{ in } Z] \text{ for cell in } CX]
         return V,CV,FV
```

Macro referenced in 27.

6 The main Boolean procedure

6.1 Goal: generating the Boolean complex

6.2 Implementation

```
\langle\, \mbox{Boolean Algorithm 25}\,\rangle \equiv """ Boolean Algorithm """
```

```
def larBool(arg1,arg2, brep=False):
   V1,basis1 = arg1
  V2,basis2 = arg2
  cells1 = basis1[-1]
   cells2 = basis2[-1]
   model1, model2 = (V1, cells1), (V2, cells2)
   (First Boolean step 4a)
   W,CW,VC,BCellCovering,cellCuts,boundary1,boundary2,BCW = larBool1()
   (Second Boolean step 7a)
   W,CW,dim,bases,boundary1,boundary2,FW,BCW = larBool2(boundary1,boundary2)
   ⟨Third Boolean step 14⟩
   V,CV,FV,boundaryMat,boundary1,boundary2,chain1,chain2,CWbits = larBool3()
   submodel = SKEL_1(STRUCT(MKPOLS((V,CV))))
   VV = AA(LIST)(range(len(V)))
   if DEBUG:
      VIEW(larModelNumbering(1,1,1)(V,[VV,FV,CV],submodel,1))
      VIEW(EXPLODE(1.2,1.2,1)(MKPOLS((V,[cell for k,cell in enumerate(CV) if sum(CWbits[k])==2
   ⟨ Fourth Boolean step 20b⟩
   W,CX,FX,CXbits = larBool4(V,CV,FV,boundaryMat,boundary1,boundary2,CWbits)
   W,CX,FX = larVertexRemoval(W,CX,FX)
   chain1,chain2 = TRANS(CXbits)
   boundaryMat = boundary(CX,FX)
   def theBoundary(boundaryMat,CX,coords):
      chainCoords = csc_matrix((len(CX), 1))
      for cell in coords: chainCoords[cell,0] = 1
      boundaryCells = list((boundaryMat * chainCoords).tocoo().row)
      orientations = list((boundaryMat * chainCoords).tocoo().data)
      orientedBoundary = [ FX[face] for (sign, face) in zip(orientations, boundaryCells) if sig
      return orientedBoundary
   def larBool0(op):
      if op == "union":
         ucoords,uchain = TRANS([(k,cell) for k,(cell,c1,c2) in enumerate(zip(CX,chain1,chain2
         return W, CW, uchain, CX, FX, the Boundary (boundary Mat, CX, ucoords)
      elif op == "intersection":
```

data = TRANS([(k,cell) for k,(cell,c1,c2) in enumerate(zip(CX,chain1,chain2)) if c1*c

```
if data != []:
                  icoords,ichain = data
                  return W, CW, ichain, CX, FX, the Boundary (boundary Mat, CX, icoords)
                  icoords,ichain = [],[]
                  return W,CW,[],[],[],[]
            elif op == "xor":
               xcoords,xchain = TRANS([(k,cell) for k,(cell,c1,c2) in enumerate(zip(CX,chain1,chain2
               return W, CW, xchain, CX, FX, the Boundary (boundary Mat, CX, xcoords)
            elif op == "difference":
               data = TRANS([(k,cell) for k,(cell,c1,c2) in enumerate(zip(CX,chain1,chain2)) if c1==
               if data != []:
                  icoords,ichain = data
                  return W, CW, ichain, CX, FX, the Boundary (boundary Mat, CX, icoords)
                  icoords,ichain = [],[]
                  return W,CW,[],[],[],[]
           else: print "Error: non implemented op"
        return larBool0
Macro referenced in 27.
```

7 LAR simplification

Occasionally, we may need to simplify

8 Exporting the library

```
"lib/py/bool1.py" 27 \(\equiv \) Module for Boolean ops with LAR """ \(\langle \text{Initial import of modules 42a} \rangle \) from splitcell import *

DEBUG = False
\(\langle \text{Symbolic utility to represent points as strings 42b} \rangle \) Merge two dictionaries with keys the point locations 4c \(\rangle \text{Make Common Delaunay Complex 5} \rangle \) Cell-facet intersection test 13 \(\rangle \text{Cell-facet intersection test 13} \rangle \) Computing the adjacent cells of a given cell 7b \(\rangle \text{Computation of boundary facets covering with CDC cells 12b} \rangle \) CDC cell splitting with one or more cutting facets 10a \(\rangle \text{Boolean argument boundaries embedding in SCDC 11} \rangle \) Make facets dictionaries 12a \(\rangle \)
```

9 Tests and examples

```
\langle \text{ Debug via visualization 28a} \rangle \equiv
     """ Debug via visualization """
     boolean = larBool(arg1,arg2)
     W,CW,chain,CX,FX,orientedBoundary = boolean("xor")
     glass = MATERIAL([1,0,0,0.6, 0,1,0,0.6, 0,0,1,0.6, 0,0,0,0.6, 100])
     VIEW(glass(EXPLODE(1.2,1.2,1.2)(MKPOLS((W,chain)))))
     if DEBUG:
         VIEW(SKEL_1(EXPLODE(1.1,1.1,1.1)(MKPOLS((W,orientedBoundary)))))
         W,CW,chain,CX,FX,orientedBoundary = boolean("union")
         VIEW(EXPLODE(1.1,1.1,1)(MKPOLS((W,chain))))
         VIEW(SKEL_1(EXPLODE(1.1,1.1,1.1)(MKPOLS((W,orientedBoundary)))))
         W,CW,chain,CX,FX,orientedBoundary = boolean("intersection")
         if chain != []:
           VIEW(EXPLODE(1.1,1.1,1)(MKPOLS((W,chain))))
           VIEW(SKEL_1(EXPLODE(1.1,1.1,1.1)(MKPOLS((W,orientedBoundary)))))
         W,CW,chain,CX,FX,orientedBoundary = boolean("difference")
         if chain != []:
           VIEW(EXPLODE(1.1,1.1,1)(MKPOLS((W,chain))))
           VIEW(SKEL_1(EXPLODE(1.1,1.1,1.1)(MKPOLS((W,orientedBoundary)))))
           VIEW(EXPLODE(1.1,1.1,1.1)(MKPOLS((W,CX))))
     submodel = SKEL_1(STRUCT(MKPOLS((W,FX))))
     VV = AA(LIST)(range(len(W)))
     VIEW(larModelNumbering(1,1,1)(W,[VV,FX,CX],submodel,1))
```

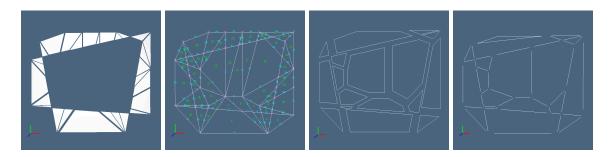


Figure 2: 2D example of file test/py/bool1/test1.py. (a) The cell numbering of SCDC; (b) the XOR of Boolean arguments; (c) the boundaries of exploded 2-cells of reduced SCDC; (d) exploded 1-cells of reduced SCDC.

```
"test/py/bool1/test0.py" 28b \equiv
     import sys
     """ import modules from larcc/lib """
     sys.path.insert(0, 'lib/py/')
     from bool1 import *
     """ Definition of Boolean arguments """
     mod_1 = AA(LIST)(range(n)), [[2*k,2*k+1] for k in range(n/2)]
     squares1 = larModelProduct([mod_1,mod_1])
     mod_2 = AA(LIST)([0.5+k*2 \text{ for } k \text{ in } range(n/2)]), [[2*k,2*k+1] \text{ for } k \text{ in } range(n/4)]
     squares2 = larModelProduct([mod_2,mod_2])
     V1 = squares1[0]
     V2 = squares2[0]
     VV1 = AA(LIST)(range(len(V1)))
     VV2 = AA(LIST)(range(len(V2)))
     EV1 = larConvexFacets (*squares1)
     EV2 = larConvexFacets (*squares2)
     FV1 = squares1[1]
     FV2 = squares2[1]
     arg1 = V1, (VV1, EV1, FV1)
     arg2 = V2, (VV2, EV2, FV2)
     (Debug via visualization 28a)
```

```
"test/py/bool1/test0b.py" 29 \equiv
     import sys
     """ import modules from larcc/lib """
     sys.path.insert(0, 'lib/py/')
     from bool1 import *
     """ Definition of Boolean arguments """
     mod_1 = AA(LIST)(range(n)), [[2*k,2*k+1] for k in range(n/2)]
     squares1 = INSR(larModelProduct)([mod_1,mod_1,mod_1])
     mod_2 = AA(LIST)([0.5+k*2 \text{ for } k \text{ in } range(n/2)]), [[2*k,2*k+1] \text{ for } k \text{ in } range(n/4)]
     squares2 = INSR(larModelProduct)([mod_2,mod_2,mod_2])
     V1 = squares1[0]
     V2 = squares2[0]
     VV1 = AA(LIST)(range(len(V1)))
     VV2 = AA(LIST)(range(len(V2)))
     FV1 = larConvexFacets (*squares1)
     FV2 = larConvexFacets (*squares2)
     CV1 = squares1[1]
     CV2 = squares2[1]
     arg1 = V1, (VV1, FV1, CV1)
     arg2 = V2, (VV2, FV2, CV2)
     (Debug via visualization 28a)
"test/py/bool1/test1.py" 30 \equiv
     import sys
     """ import modules from larcc/lib """
     sys.path.insert(0, 'lib/py/')
     from bool1 import *
     """ Definition of Boolean arguments """
     V1 = [[3,0],[11,0],[13,10],[10,11],[8,11],[6,11],[4,11],[1,10],[4,3],[6,4],
            [8,4],[10,3]]
     FV1 = [[0,1,8,9,10,11],[1,2,11],[3,10,11],[4,5,9,10],[6,8,9],[0,7,8],[2,3,
            11],[3,4,10],[5,6,9],[6,7,8]]
     EV1 = [[0,1],[0,7],[0,8],[1,2],[1,11],[2,3],[2,11],[3,4],[3,10],[3,11],[4,
           5], [4,10], [5,6], [5,9], [6,7], [6,8], [6,9], [7,8], [8,9], [9,10], [10,11]]
     VV1 = AA(LIST)(range(len(V1)))
     V2 = [[0,3],[14,2],[14,5],[14,7],[14,11],[0,8],[3,7],[3,5]]
```

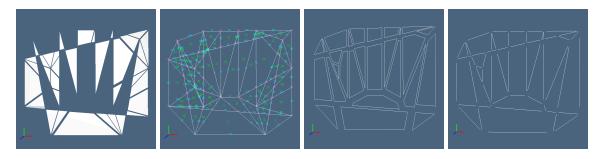


Figure 3: 2D example of file test/py/bool1/test2.py. (a) The cell numbering of SCDC; (b) the XOR of Boolean arguments; (c) the boundaries of exploded 2-cells of reduced SCDC; (d) exploded 1-cells of reduced SCDC.

"test/py/bool1/test2.py" $31 \equiv$

[5,6],[6,7]] VV2 = AA(LIST)(range(len(V2)))

import sys

EV2 = [[0,1],[0,5],[0,7],[1,2],[1,7],[2,3],[2,7],[3,4],[3,6],[4,5],[4,6],

V2 = [[0,3],[14,2],[14,5],[14,7],[14,11],[0,8],[3,7],[3,5]]FV2 = [[0,5,6,7],[0,1,7],[4,5,6],[2,3,6,7],[1,2,7],[3,4,6]]

```
arg1 = V1, (VV1,EV1,FV1)
arg2 = V2, (VV2,EV2,FV2)
⟨Debug via visualization 28a⟩
⋄
```

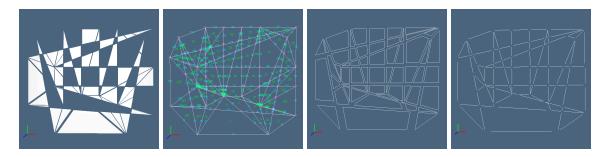


Figure 4: 2D example of file test/py/bool1/test3.py. (a) The cell numbering of SCDC; (b) the XOR of Boolean arguments; (c) the boundaries of exploded 2-cells of reduced SCDC; (d) exploded 1-cells of reduced SCDC.

```
"test/py/bool1/test3.py" 32a \equiv
```

```
import sys
""" import modules from larcc/lib """
sys.path.insert(0, 'lib/py/')
from bool1 import *
V1 = [[3,0],[11,0],[13,10],[10,11],[8,11],[6,11],[4,11],[1,10],[4,3],[6,4],
      [8,4],[10,3]]
FV1 = [[0,1,8,9,10,11],[1,2,11],[3,10,11],[4,5,9,10],[6,8,9],[0,7,8]]
EV1 = [[0,1],[0,7],[0,8],[1,2],[1,11],[2,11],[3,10],[3,11],[4,5],[4,10],[5,
      9],[6,8],[6,9],[7,8],[8,9],[9,10],[10,11]]
VV1 = AA(LIST)(range(len(V1)))
V2 = [[0,3],[14,2],[14,5],[14,7],[14,11],[0,8],[3,7],[3,5]]
FV2 = [[0,5,6,7],[0,1,7],[4,5,6],[2,3,6,7]]
EV2 = [[0,1],[0,5],[0,7],[1,7],[2,3],[2,7],[3,6],[4,5],[4,6],[5,6],[6,7]]
VV2 = AA(LIST)(range(len(V2)))
arg1 = V1, (VV1, EV1, FV1)
arg2 = V2, (VV2, EV2, FV2)
⟨ Debug via visualization 28a⟩
```

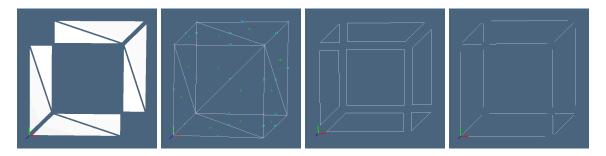


Figure 5: 2D example of file test/py/bool1/test4.py. (a) The cell numbering of SCDC; (b) the XOR of Boolean arguments; (c) the boundaries of exploded 2-cells of reduced SCDC; (d) exploded 1-cells of reduced SCDC.

```
"test/py/bool1/test4.py" 32b \equiv
     import sys
     """ import modules from larcc/lib """
     sys.path.insert(0, 'lib/py/')
     from bool1 import *
     V1 = [[0,0],[10,0],[10,10],[0,10]]
     FV1 = [range(4)]
     EV1 = [[0,1],[1,2],[2,3],[0,3]]
     VV1 = AA(LIST)(range(len(V1)))
     V2 = [[2.5,2.5],[12.5,2.5],[12.5,12.5],[2.5,12.5]]
     FV2 = [range(4)]
     EV2 = [[0,1],[1,2],[2,3],[0,3]]
     VV2 = AA(LIST)(range(len(V2)))
     arg1 = V1, (VV1, EV1, FV1)
     arg2 = V2, (VV2, EV2, FV2)
     ⟨ Debug via visualization 28a⟩
```

ERROR Problems remain with facet extraction from a (too) small convex complex, both if made of simplices (the automatically generated FW is wrong) and if made of cuboids (the automatically generated FX is wrong too). Errors to solve in the implementation of automatic extraction of facets.

[&]quot;test/py/bool1/test5a.py" $34a \equiv$

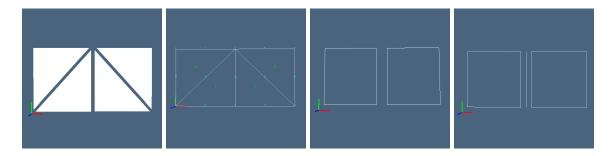


Figure 6: 2D example of file test/py/bool1/test5.py. (a) The cell numbering of SCDC; (b) the XOR of Boolean arguments; (c) the boundaries of exploded 2-cells of reduced SCDC; (d) exploded 1-cells of reduced SCDC. (ERRORS in the images)

```
import sys
     """ import modules from larcc/lib """
     sys.path.insert(0, 'lib/py/')
     from bool1 import *
     V1 = [[0,0],[5,0],[5,5],[0,5]]
     FV1 = [range(4)]
     EV1 = [[0,1],[1,2],[2,3],[0,3]]
     VV1 = AA(LIST)(range(len(V1)))
     V2 = [[5,0],[10,0],[10,5],[5,5]]
     FV2 = [range(4)]
     EV2 = [[0,1],[1,2],[2,3],[0,3]]
     VV2 = AA(LIST)(range(len(V2)))
     arg1 = V1, (VV1, EV1, FV1)
     arg2 = V2, (VV2, EV2, FV2)
     ⟨ Debug via visualization 28a⟩
"test/py/bool1/test5b.py" 34b \equiv
     import sys
     """ import modules from larcc/lib """
     sys.path.insert(0, 'lib/py/')
     from bool1 import *
     V1 = [[0,0],[5,0],[5,5],[0,5]]
     FV1 = [range(4)]
     EV1 = [[0,1],[1,2],[2,3],[0,3]]
```

```
VV1 = AA(LIST)(range(len(V1)))
     V2 = [[5,2],[10,2],[10,7],[5,7]]
     FV2 = [range(4)]
     EV2 = [[0,1],[1,2],[2,3],[0,3]]
     VV2 = AA(LIST)(range(len(V2)))
     arg1 = V1, (VV1, EV1, FV1)
     arg2 = V2, (VV2, EV2, FV2)
     ⟨ Debug via visualization 28a⟩
"test/py/bool1/test5c.py" 35a \equiv
     import sys
     """ import modules from larcc/lib """
     sys.path.insert(0, 'lib/py/')
     from bool1 import *
     V1 = [[0,0],[5,0],[5,5],[0,5]]
     FV1 = [range(4)]
     EV1 = [[0,1],[1,2],[2,3],[0,3]]
     VV1 = AA(LIST)(range(len(V1)))
     V2 = [[5,5],[10,5],[10,10],[5,10]]
     FV2 = [range(4)]
     EV2 = [[0,1],[1,2],[2,3],[0,3]]
     VV2 = AA(LIST)(range(len(V2)))
     arg1 = V1, (VV1, EV1, FV1)
     arg2 = V2, (VV2, EV2, FV2)
     ⟨ Debug via visualization 28a⟩
"test/py/bool1/test5d.py" 35b \equiv
     import sys
     """ import modules from larcc/lib """
     sys.path.insert(0, 'lib/py/')
     from bool1 import *
     V1 = [[0,0],[5,0],[5,5],[0,5]]
     FV1 = [range(4)]
     EV1 = [[0,1],[1,2],[2,3],[0,3]]
     VV1 = AA(LIST)(range(len(V1)))
```

```
V2 = [[5,6],[10,6],[10,11],[5,11]]

FV2 = [range(4)]

EV2 = [[0,1],[1,2],[2,3],[0,3]]

VV2 = AA(LIST)(range(len(V2)))

arg1 = V1,(VV1,EV1,FV1)

arg2 = V2,(VV2,EV2,FV2)

\( \text{Debug via visualization 28a} \)
```

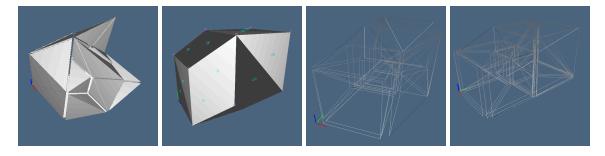


Figure 7: 2D example of file test/py/bool1/test6.py. (a) The cell numbering of SCDC; (b) the XOR of Boolean arguments; (c) the boundaries of exploded 2-cells of reduced SCDC; (d) exploded 1-cells of reduced SCDC. (ERRORS in the images)

ERROR Problems remain with tagging of 3D cell as internal/external to Boolean boundaries. Errors to solve in the implementation of general (no simplicial) signed boundary operator matrix.

"test/py/bool1/test6.py" $36a \equiv$

```
import sys
""" import modules from larcc/lib """
sys.path.insert(0, 'lib/py/')
from bool1 import *

V1 = [[0,0,0],[10,0,0],[10,10,0],[0,10,0],[0,0,10],[10,0,10],[10,10,10]]
V1,[VV1,EV1,FV1,CV1] = larCuboids((1,1,1),True)
V1 = [SCALARVECTPROD([5,v]) for v in V1]

V2 = [SUM([v,[2.5,2.5,2.5]]) for v in V1]
[VV2,EV2,FV2,CV2] = [VV1,EV1,FV1,CV1]
```

```
arg1 = V1,(VV1,EV1,FV1,CV1)
arg2 = V2,(VV2,EV2,FV2,CV2)
⟨Debug via visualization 28a⟩
⋄
```

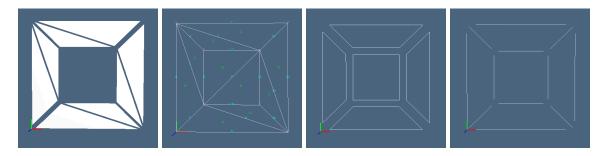


Figure 8: 2D example of file test/py/bool1/test7.py. (a) The cell numbering of SCDC; (b) the XOR of Boolean arguments; (c) the boundaries of exploded 2-cells of reduced SCDC; (d) exploded 1-cells of reduced SCDC.

```
"test/py/bool1/test7.py" 36b \equiv
     import sys
     """ import modules from larcc/lib """
     sys.path.insert(0, 'lib/py/')
     from bool1 import *
     V1 = [[0,0],[10,0],[10,10],[0,10]]
     FV1 = [range(4)]
     EV1 = [[0,1],[1,2],[2,3],[0,3]]
     VV1 = AA(LIST)(range(len(V1)))
     V2 = [[2.5,2.5],[7.5,2.5],[7.5,7.5],[2.5,7.5]]
     FV2 = [range(4)]
     EV2 = [[0,1],[1,2],[2,3],[0,3]]
     VV2 = AA(LIST)(range(len(V2)))
     arg1 = V1, (VV1, EV1, FV1)
     arg2 = V2, (VV2, EV2, FV2)
     \langle\, \text{Debug via visualization 28a}\,\rangle
```

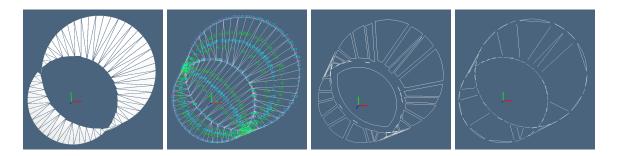


Figure 9: 2D example of file test/py/bool1/test8.py. (a) The cell numbering of SCDC; (b) the xor of Boolean arguments; (c) the boundaries of exploded 2-cells of reduced SCDC; (d) exploded 1-cells of reduced SCDC. (ERRORS: Numeric (?) errors in the splitting procedure?)

```
"test/py/bool1/test8.py" 37 \equiv
     import sys
     """ import modules from larcc/lib """
     sys.path.insert(0, 'lib/py/')
     from bool1 import *
     n = 48
     V1 = [[5*\cos(\text{angle}*2*PI/n)+2.5, 5*\sin(\text{angle}*2*PI/n)+2.5] \text{ for angle in range(n)}]
     FV1 = [range(n)]
     EV1 = TRANS([range(n), range(1, n+1)]); EV1[-1] = [0, n-1]
     VV1 = AA(LIST)(range(len(V1)))
     V2 = [[4*cos(angle*2*PI/n), 4*sin(angle*2*PI/n)]  for angle in range(n)]
     FV2 = [range(n)]
     EV2 = EV1
     VV2 = AA(LIST)(range(len(V2)))
     arg1 = V1, (VV1, EV1, FV1)
     arg2 = V2, (VV2, EV2, FV2)
     (Debug via visualization 28a)
"test/py/bool1/test9.py" 38 \equiv
     import sys
     """ import modules from larcc/lib """
     sys.path.insert(0, 'lib/py/')
```

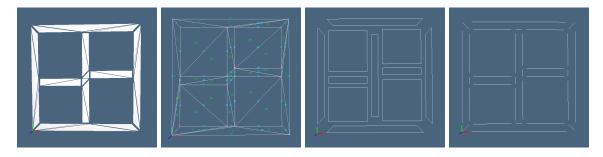


Figure 10: 2D example of file test/py/bool1/test10.py. (a) The cell numbering of SCDC; (b) the XOR of Boolean arguments; (c) the boundaries of exploded 2-cells of reduced SCDC; (d) exploded 1-cells of reduced SCDC.

```
"test/py/bool1/test10.py" 39 \(\text{import sys}\)
    """ import modules from larcc/lib """
    sys.path.insert(0, 'lib/py/')
    from bool1 import *

V1 = [[0,0],[15,0],[15,14],[0,14]]
    FV1 = [range(4)]
    EV1 = [[0,1],[1,2],[2,3],[0,3]]
    VV1 = AA(LIST)(range(len(V1)))
```

9.1 Random data input

Macro never referenced.

Random points in unit disk First we generate a set of n random points in the unit D^d disk centred on the origin, to be subsequently used to generate a random Delaunay complex of variable granularity.

```
⟨Generation of n random points in the unit d-disk 40b⟩ ≡

def randomPointsInUnitCircle(n=200,d=2, r=1):
    points = random.random((n,d)) * ([2*math.pi]+[1]*(d-1))
    return [[SQRT(p[1])*COS(p[0]),SQRT(p[1])*SIN(p[0])] for p in points]
    ## TODO: correct for $d$-sphere

if __name__=="__main__":
    VIEW(STRUCT(AA(MK)(randomPointsInUnitCircle())))

◇
```

Macro referenced in 40a.

Random points in the standard d-cuboid A set of n random d-points is then generated within the standard d-cuboid, i.e. withing the d-dimensional interval with a vertex on the origin.

 \langle Generation of *n* random points in the standard *d*-cuboid 40c \rangle \equiv

```
def randomPointsInUnitCuboid(n=200,d=2):
    return random.random((n,d)).tolist()

if __name__=="__main__":
    VIEW(STRUCT(AA(MK)(randomPointsInUnitCuboid())))
```

Macro referenced in 40a.

Triangulation of random points The Delaunay triangulation of randomPointsInUnitCircle is generated by the following macro.

Macro referenced in 40a.

```
"test/py/bool1/test11.py" 41b \equiv
     """ Union of 2D non-structured grids """
     import sys
     """ import modules from larcc/lib """
     sys.path.insert(0, 'lib/py/')
     from bool1 import *
     model1 = randomTriangulation(100,2,'disk')
     V1,CV1 = model1
     VIEW(EXPLODE(1.5,1.5,1)(MKPOLS(model1)+cellNames(model1,CV1,MAGENTA)))
     FV1 = convexFacets (V1,CV1)
     VV1 = AA(LIST)(range(len(V1)))
     model2 = randomTriangulation(100,2,'cuboid')
     V2,CV2 = model2
     V2 = larScale([2,2])(V2)
     model2 = V2,CV2
     VIEW(EXPLODE(1.5,1.5,1)(MKPOLS(model2)+cellNames(model2,CV2,RED)))
     FV2 = convexFacets (V2,CV2)
     VV2 = AA(LIST)(range(len(V2)))
     arg1 = V1, (VV1, FV1, CV1)
     arg2 = V2, (VV2, FV2, CV2)
     (Debug via visualization 28a)
```

A Appendix: utility functions

```
⟨Initial import of modules 42a⟩ ≡
    from pyplasm import *
    from scipy import *
    import sys
    """ import modules from larcc/lib """
    sys.path.insert(0, 'lib/py/')
    from lar2psm import *
    from simplexn import *
    from larcc import *
    from largrid import *
    from myfont import *
    from mapper import *
```

A.1 Numeric utilities

A small set of utility functions is used to transform a *point* representation, given as array of coordinates, into a string of fixed format to be used as point key into python dictionaries.

```
(Symbolic utility to represent points as strings 42b) =
    """ TODO: use package Decimal (http://docs.python.org/2/library/decimal.html) """
    global PRECISION
PRECISION = 4.

def verySmall(number): return abs(number) < 10**-(PRECISION)

def prepKey (args): return "["+", ".join(args)+"]"

def fixedPrec(value):
    out = round(value*10**(PRECISION))/10**(PRECISION)
    if out == -0.0: out = 0.0
        return str(out)

def vcode (vect):
        """
        To generate a string representation of a number array.
        Used to generate the vertex keys in PointSet dictionary, and other similar operations.
        """
        return prepKey(AA(fixedPrec)(vect))</pre>
```

References

Macro referenced in 27.

[CL13] CVD-Lab, *Linear algebraic representation*, Tech. Report 13-00, Roma Tre University, October 2013.