Time complexity of recursive function inside for loop

Asked 4 years, 6 months ago Modified 1 year, 4 months ago Viewed 6k times



If we have a function:-

```
2    int x=0;
    int fun(int n)
    {
        if(n==0)
            return 1;
        for(int i=0; i<n;i++)
            x += fun(i);
    }
```

According to me, time complexity can be calculated as:-

```
T(n) = T(n-1) + T(n-2) + ... + T(0).

T(n) = nT(n-1).

T(n) = O(n!).
```

Am I correct?

time-complexity big-o

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edited Aug 19, 2017 at 9:14

Tony Tannous

12.9k 8 44 79

asked Aug 19, 2017 at 5:47



1,401

01 3 17



If you're measuring the number of function calls (or additions -- it turns out the same), the correct recurrence relations are:

4

$$T(0) = 0$$

 $T(n) = T(0) + T(1) + T(2) + ... + T(n-1) + n$



You can compute the first few values:



$$T(0) = 0$$

$$T(1) = 1$$

$$T(2) = 3$$

$$T(3) = 7$$

T(4) = 15

You can guess from this that $T(n) = 2^n - 1$, and that's easy to check by a proof by induction.

In some sense you are right that T(n) = O(n!) since $n! > 2^n$ for n > 3, but $T(n) = O(2^n)$ is a tighter bound.

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answered Aug 19, 2017 at 9:18



Paul Hankin

48.8k 11 84 106

1 I think if we count the initial function call also then the recurrence would be :- T(n) = T(n-1) + T(n-2) + ... + 1 and not n at the end. For n=3, T(3) = T(0) + T(1) + T(2) + 1 = 1 + 2 + 4 + 1 = 8. – Zephyr Aug 19, 2017 at 10:22



1.
$$T(n) = T(n-1) + T(n-2) + T(n-3) + + T(0)$$

4

// Replace n with n-1
2.
$$T(n-1) = T(n-2) + T(n-3) + + T(0)$$



Replace $T(n-2) + T(n-3) + \dots + T(0)$ with T(n-1) in 1st Equation

```
T(n) = T(n-1) + T(n-1)
= 2 * T(n-1)
= 2 * 2 * T(n-2) // Using T(n-1) = 2 * T(n-2)
= 2^n * T(n-n)
= 2^n * T(0) // Consider T(0) = 1
= 2^n
```

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edited Sep 27, 2019 at 11:43

answered Sep 23, 2019 at 13:54

