



Construct a balanced BST from the given keys

Given an unsorted integer array that represents binary search tree (BST) keys, construct a height-balanced BST from it. For each node of a height-balanced tree, the difference between its left and right subtree height is at most 1.

For example,

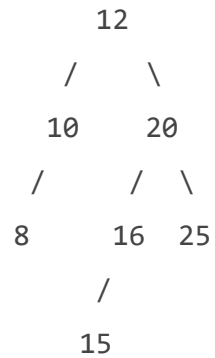
Input: keys = [15, 10, 20, 8, 12, 16, 25]

Output:

```
      15
     /  \
    /    \
   10     20
```

8 12 16 25

OR



OR

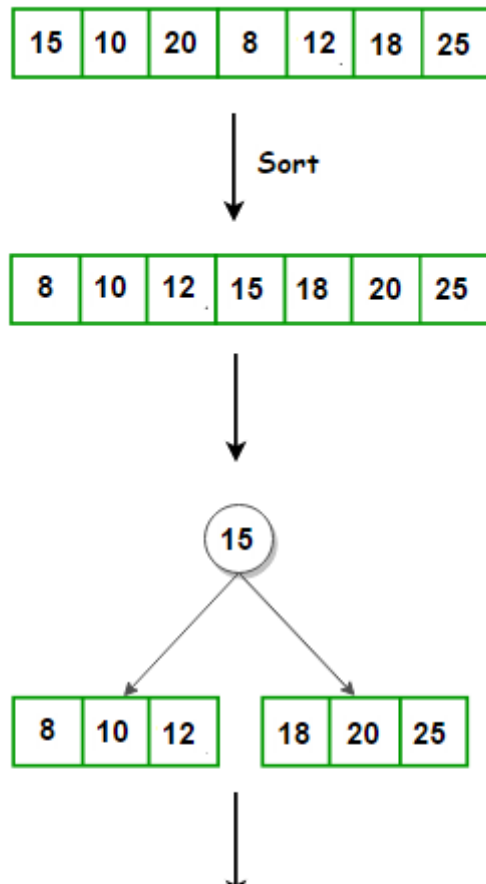
Any other possible representation.

Practice this problem

We have already discussed how to [insert a key into a BST](#). The height of such BST in the worst-case can be as much as the total number of keys in the BST. The worst case happens when given keys are sorted in ascending or descending order, and we get a skewed tree where all the nodes except the leaf have one and only one child.

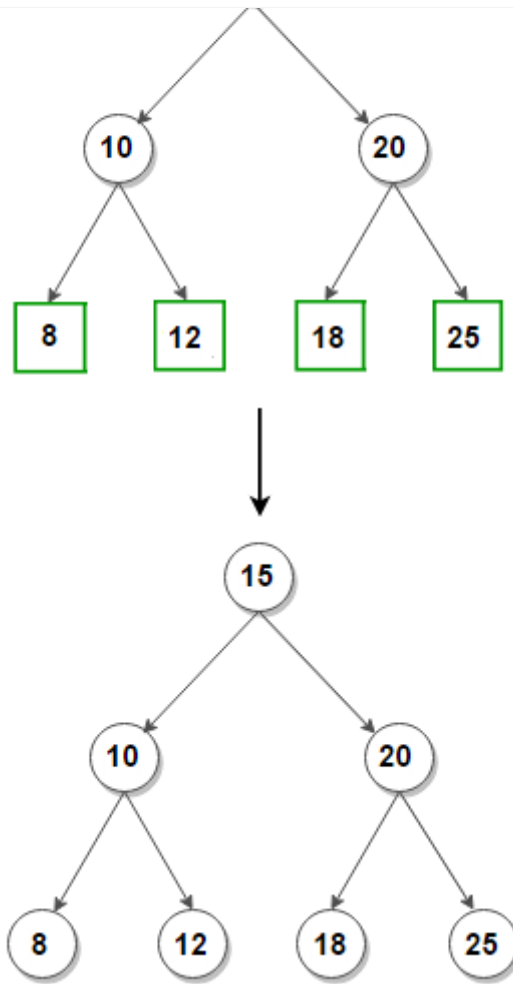
better than the linear time required to find items by key in an (unsorted) array or unbalanced trees.

We can easily modify the solution to get height-balanced BSTs if all keys are known in advance. The idea is to [sort the given keys](#) first. Then the root will be the middle element of the sorted array, and we recursively construct the left subtree of the root by keys less than the middle element and the right subtree of the root by keys more than the middle element. For example,



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Following is the C++, Java, and Python implementation of the idea:

C++

Java

Python

```
3  #include <algorithm>
4  using namespace std;
5
6  // Data structure to store a BST node
7  struct Node
8  {
9      int data;
10     Node* left = nullptr, *right = nullptr;
11
12     Node() {}
13     Node(int data): data(data) {}
14 };
15
16 // Function to perform inorder traversal on the tree
17 void inorder(Node* root)
18 {
19     if (root == nullptr) {
20         return;
21     }
22
23     inorder(root->left);
24     cout << root->data << " ";
25     inorder(root->right);
26 }
27
28 // Recursive function to insert a key into a BST
29 Node* insert(Node* root, int key)
30 {
31     // if the root is null, create a new node and return it
32     if (root == nullptr) {
33         return new Node(key);
34     }
35
36     // if the given key is less than the root node, recur for the left subtree
37     if (key < root->data) {
38         root->left = insert(root->left, key);
39     }
40
41     // if the given key is more than the root node, recur for the right subtree
```

```
45     return root;
46 }
47
48 // Function to construct balanced BST from the given sorted array.
49 // Note that the root of the tree is passed by reference here
50 void convert(vector<int> const &keys, int low, int high, Node* &root)
51 {
52     // base case
53     if (low > high) {
54         return;
55     }
56
57     // find the middle element of the current range
58     int mid = (low + high) / 2;
59
60     // construct a new node from the middle element and assign it to the root
61     root = new Node(keys[mid]);
62
63     // left subtree of the root will be formed by keys less than middle element
64     convert(keys, low, mid - 1, root->left);
65
66     // right subtree of the root will be formed by keys more than the middle element
67     convert(keys, mid + 1, high, root->right);
68 }
69
70 // Function to construct balanced BST from the given unsorted array
71 Node* convert(vector<int> keys)
72 {
73     // sort the keys first
74     sort(keys.begin(), keys.end());
75
76     // construct a balanced BST
77     Node* root = nullptr;
78     convert(keys, 0, keys.size() - 1, root);
79
80     // return root node of the tree
81     return root;
82 }
83 }
```

```
87 // input keys
88 vector<int> keys = { 15, 10, 20, 8, 12, 16, 25 };
89
90 // construct a balanced binary search tree
91 Node* root = convert(keys);
92
93 // print the keys in an inorder fashion
94 inorder(root);
95
96 return 0;
97 }
```

[Download](#) [Run Code](#)**Output:**

8 10 12 15 16 20 25

The time complexity of the above solution is $O(n \cdot \log(n))$, where n is the size of the BST, and requires space proportional to the tree's height for the call stack.

📁 [BST, Sorting](#)

🔧 [Amazon](#), [Easy](#), [Recursive](#)