

# $O(N+M)$ time complexity

Asked 7 years, 6 months ago   Modified 2 years ago   Viewed 34k times



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I'm working through some practice problems where I'm given a target time complexity and space complexity. One of them gives a target time complexity of  $O(N+M)$ . I'm having some trouble with the intuition of what an  $O(N+M)$  algorithm would look like. Does anyone have an example of an algorithm like this or can explain it clearly? Every example I try to think of seems like  $O(N*M)$  to me.

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edited Sep 11, 2014 at 20:27

asked Sep 11, 2014 at 20:18



[user137717](#)

**1,665**   4   20   43

- 1 Can you post an example of a case with  $O(N+M)$ ? – [Javi](#) Sep 11, 2014 at 20:19
- 1 Yes, so we can figure out what  $N$  and  $M$  are in a specific case. All I can think about right now is something like: find the minimum of 2 vector, once of size  $N$  and the other one of size  $M$ , so that it would be  $O(M+N)$  – [Javi](#) Sep 11, 2014 at 20:21
- 1 I'd like to avoid it if I can. The answer below is a good example and I don't want to post the question for fear of someone giving too much of the answer away. Thanks for your help! – [user137717](#) Sep 11, 2014 at 20:31
- 1 In case you want the formal definition: [en.wikipedia.org/wiki/Big\\_O\\_notation#Multiple\\_variables](http://en.wikipedia.org/wiki/Big_O_notation#Multiple_variables) – [Tavian Barnes](#) Sep 11, 2014 at 20:50

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A simple example of algorithm that is  $O(m+n)$  :

32



```
int sum(int[] nArr, int[] mArr) {
    int sum = 0;
    for(int i : nArr) {
        sum += i;
    }
}
```



```

    }
    for(int i : mArr) {
        sum += i;
    }
    return sum;
}

```

To compute the sum, you need to go through all elements in `nArr` (size `n`) and all elements in `mArr` (size `m`), so the overall complexity is  $O(m+n)$

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answered Sep 11, 2014 at 20:24



[Jean Logeart](#)

**50.3k** 11 80 116

- 1 Complexity-wise, that is impossible to do better :) – [Jean Logeart](#) Sep 11, 2014 at 20:27
- 2 Let's say array `nArr` with size `N` is larger. Would it still be correct to say that this algorithm is  $O(N)$ ? – [user137717](#) Sep 11, 2014 at 20:30
- 4 No this is not correct, unless `m` is constant or unless `m` depends on `n` such that  $m = O(n)$  – [Jean Logeart](#) Sep 11, 2014 at 20:31
- 2 Is it correct to think about it as an  $O(n)$  algorithm followed by an  $O(m)$  algorithm and together they are  $O(n+m)$  and abstract this to an arbitrary number of arrays? – [user137717](#) Sep 11, 2014 at 20:41
- 9 Complexity annotation (such as  $O()$ ) describes the way an algorithm's performance (e.g. time or space) is dependent on the size of the input. It is "blind" to a specific input. When you write  $O(m+n)$  it means that the algorithm will take  $O(m)$  time when  $m > n$  and  $O(n)$  when  $m < n$ . In any way,  $O(m+n) = O(\max(m, n))$  . – [Shaked](#) Sep 11, 2014 at 20:49



Quick and simple example of an  $O(n + m)$  algorithm:

18



```

for (i = 0; i < n; i++)
{
    // do something but don't loop or invoke recursive functions
    // only constant  $O(c)$  complexity is allowed: a simple series of commands
}

for (i = 0; i < m; i++)
{

```

```
// idem
}
```

Complexity is commutative when added ( $O(n + m) == O(m + n)$ ) this means you may invert the two `for()` without affecting complexity. Obviously, on an algorithmic level the inverted one *MAY* not be equivalent to the straight one.

As an additional help, here is an example of  $O(n * m)$  algorithm:

```
for (i = 0; i < n; i++)
{
    for (j = 0; j < m; j++)
    {
        // do something but don't loop or invoke recursive functions
        // only constant  $O(c)$  complexity is allowed: a simple series of commands
    }
}
```

Again, you may invert inside with outside loop without affecting complexity ( $O(n * m) == O(m * n)$ ). The same obvious considerations apply.

The limitation on what you may put into the `for()` bodies is because the big-o notation constraints the upper bound. If it were a lower bound (little-o notation) you may have put more complex stuff in, but it could never get less than that.

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edited Sep 11, 2014 at 21:19

answered Sep 11, 2014 at 21:12



pid

11.2k

5

35

58



So, in order to expand the other replies, I will try to add a example of such problems to help you understand:

2

- Find a min/max in a  $N$  sized array, and then look for this value in an  $M$  sized array. Since you need to perform first min/max search, you cannot do it at once.



For instance, summing up the elements of 2 vectors can be done in  $O(M+N)$ , but it can be thought as  $O(N)$  (assuming  $N > M$ ) or  $O(M)$  (if  $M > N$ ).

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edited Sep 11, 2014 at 21:01

answered Sep 11, 2014 at 20:56



Javi

3,114

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2 Assuming  $N < M$ , it is  $O(M)$ , so  $O(M+N)$  is simply  $\max(O(M), O(N))$ , which is still linear. – Mephy Sep 11, 2014 at 21:00

1 Well I think that is kind of obvious. However, I have updated the answer.  $O(M+N)$  is linear, no matter what the problem is. – Javi Sep 11, 2014 at 21:02



2



```
for(int i=0; i < n;i++){
    for(int j=0; j < m;j++){
        //some_code
    }
}
```

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edited Feb 22, 2020 at 16:15

answered Feb 22, 2020 at 15:19



Mort

1,321

3

16

47



m.mashaly

81

6

1 Doesn't [this answer](#) already cover this? – kaya3 Feb 22, 2020 at 15:26

Oops didn't see it – m.mashaly Feb 23, 2020 at 16:56



1



The intuition of this problem is that you have two unique variables  $n$  and  $m$ . Now imagine these two unique variables independently increasing, approaching infinity.

If this were an  $O(n)$  problem (i.e. BIG-O), the upward boundary of the complexity of this problem would be linear at least. You could say that  $O(n) = n^2$ . But an  $O(n)$  problem would never even get close to that  $n^2$  limit as  $n$  (the input) approaches infinite.

Likewise, the behavior for  $m$  would be the same.  $O(m)$  can be  $m^2$ . But it is more accurate to say that  $O(m) = m$ . The complexities of these two problems are **linear**.

Now, if you just do  $O(n+m)$ , is that really  $n^2$ ? It shouldn't be. Even if  $n=m$ , the sum would be  $2n$  or  $2m$ . The complexity of this problem is still **linear** because the size of the output is still **proportional** to the inputs  $n$  and  $m$ . Therefore, *the most precise answer to this problem* would be  $O(n+m) = n+m$ .

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answered Sep 11, 2014 at 20:48



Caleb Faruki

2,374 3 27 50



1



One instructive example which does something non-trivial is to take two sorted arrays of size  $M$  and  $N$  and output a new sorted array with all of these elements. This is the basis of merge-sort and will take  $O(M+N)$  comparisons.

You can find an example anywhere or do it yourself.

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answered Sep 12, 2014 at 0:50



user1952500

6,265 3 22 36