Relational Database Design

Module 8: Higher Normal Forms

Hugo Kornelis hugo@perFact.info



Outline

Basic normal forms

- □ First Normal Form (1NF)
- Second Normal Form (2NF)
- □ Third Normal Form (3NF)

Higher normal forms

EKNF, BCNF, 4NF, 5NF, DK/NF, ONF, 6NF

Denormalization

Standard patterns for

- Referencing a range
- Storing historic data

Third Normal Form (3NF)

- Every non-key attribute depends on each candidate key
 - (nothing more, nothing less)
- But what about key attributes?
 - Every single attribute key always depends on each candidate key
 - But a part of a composite key may depend on other attributes as well!

Boyce-Codd Normal Form (BCNF)

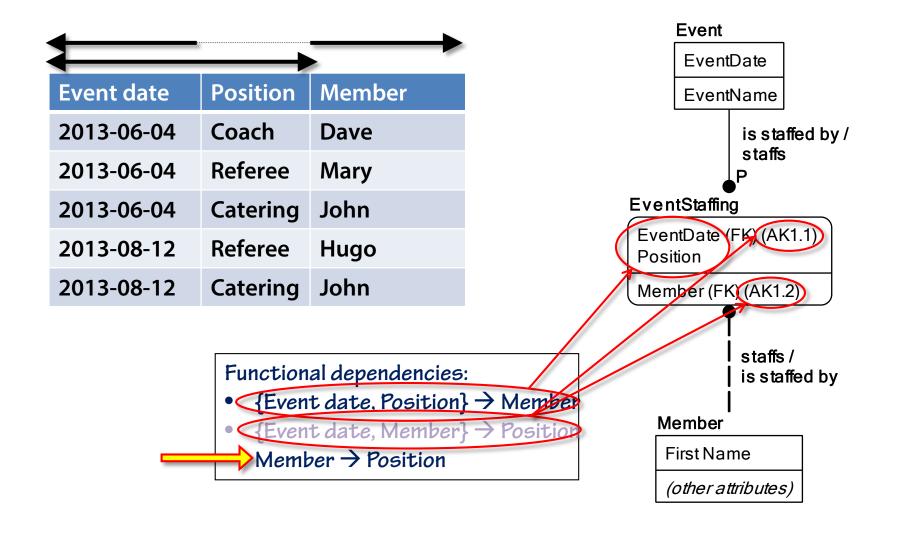
- Every attribute (key or non-key!) depends on each candidate key
 - (nothing more, nothing less)

Example





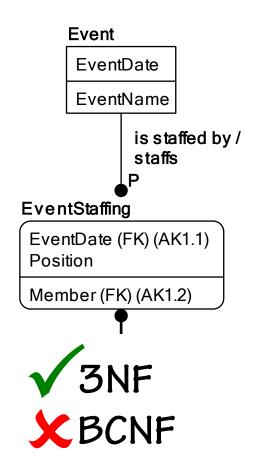




=		
Event date	Position	Member
2013-06-04	Coach	Dave
2013-06-04	Referee	Mary
2013-06-04	Catering	John
2013-08-12	Referee	Hugo
2013-08-12	Catering	John
2013-08-12	Coach	Mary

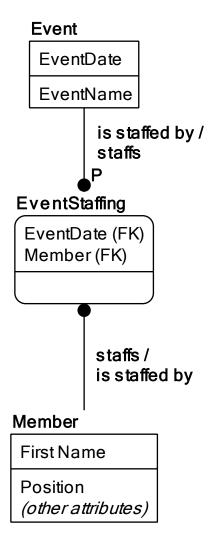
Functional dependencies:

- {Event date, Position} → Member
- Member → Position



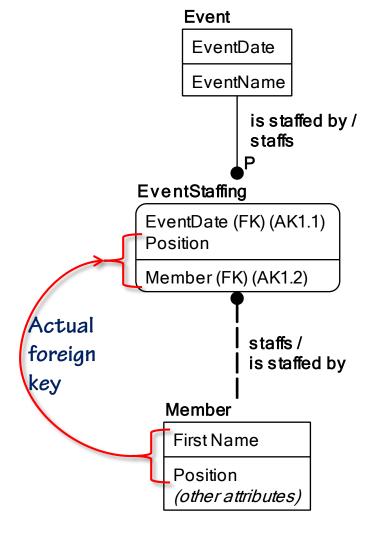
•	—	
	Event date	Member
	2013-06-04	Hugo
	2013-06-04	Mary
	2013-06-04	John
	2013-08-12	Hugo
	2013-08-12	John

Member	Position
Dave	Coach
Mary	Referee
John	Catering
Hugo	Referee



=		\rightarrow
Event date	Position	Member
2013-06-04	Coach	Dave
2013-06-04	Referee	Mary
2013-06-04	Catering	John
2013-08-12	Referee	Hugo
2013-08-12	Catering	John

	-
Member	Position
Dave	Coach
Mary	Referee
John	Catering
Hugo	Referee



- Boyce-Codd Normal Form (BCNF)
 - □ Every attribute (key **or** non-key!) depends on each candidate key
 - (nothing more, nothing less)
 - Can't always be achieved
 - □ (But usually can!)

Elementary Key Normal Form

Boyce-Codd Normal Form (BCNF)

- Every attribute (key or non-key!) depends on each candidate key
 - (nothing more, nothing less)
- Can't always be achieved
 - □ (But usually can!)

Elementary Key Normal Form (EKNF)

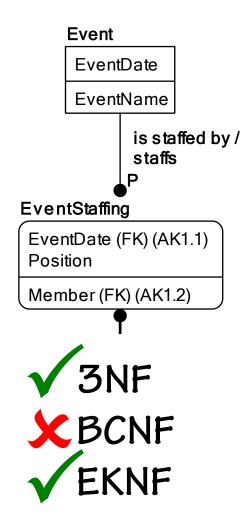
- Between 3NF and BCNF
 - Every non-key attribute depends on each candidate key
 - Every *elementary* key attribute depends on each candidate key
- Does not solve all problems

Elementary Key Normal Form

=		
Event date	Position	Member
2013-06-04	Coach	Dave
2013-06-04	Referee	Mary
2013-06-04	Catering	John
2013-08-12	Referee	Hugo
2013-08-12	Catering	John
2013-08-12	Coach	Mary

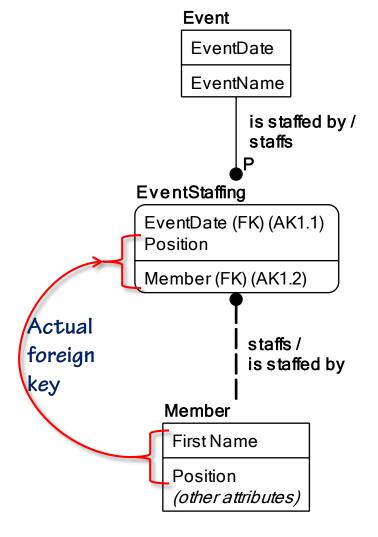
Functional dependencies:

- {Event date, Position} → Member
- Member → Position



=		\rightarrow
Event date	Position	Member
2013-06-04	Coach	Dave
2013-06-04	Referee	Mary
2013-06-04	Catering	John
2013-08-12	Referee	Hugo
2013-08-12	Catering	John

+	
Member	Position
Dave	Coach
Mary	Referee
John	Catering
Hugo	Referee



Elementary Key Normal Form

Boyce-Codd Normal Form (BCNF)

- Every attribute (key or non-key!) depends on each candidate key
 - (nothing more, nothing less)
- Can't always be achieved
 - (But usually can!)

Elementary Key Normal Form (EKNF)

- Between 3NF and BCNF
 - Every non-key attribute depends on each candidate key
 - Every *elementary* key attribute depends on each candidate key
- Does not solve all problems
- Can always be achieved
 - Using Bernstein's algorithm for synthesis of a Third Normal Form schema
 - (which actually generates a schema in Elementary Key Normal Form)

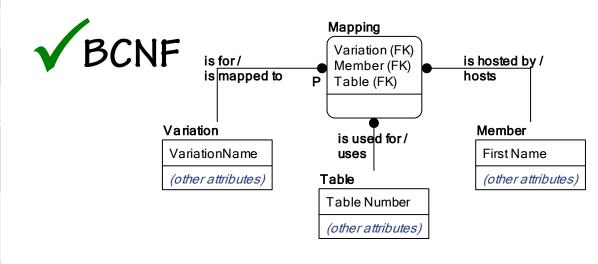


	←		\rightarrow
	Variation	Host	Table
	Pooker	Mary	2
	Pooker	Dave	2
	Pooker	Wary	1
	Pooker	Dave	1
	Sneaker	John	1
(Sneaker	John	2
•	Sneaker	John	3
	Sneaker	Mary	1
	Sneaker	Mary	2
	Sneaker	Mary	3
	Quick Snook	Mary	3
	Quick Snook	Dave	3
	Quick	Mary	4

Variation Pooker is hosted by Mary on table

Variation Pooker is hosted by Mary.

Variation Pooker can be played on table 2.



—		\rightarrow
Variation	Host	Table
Pooker	Mary	2
Pooker	Dave	2
Pooker	Mary	1
Pooker	Dave	1
Sneaker	John	1
Sneaker	John	2
Sneaker	John	3
Sneaker	Mary	1
Sneaker	Mary	2
Sneaker	Mary	3
Quick Snook	Mary	3
Quick Snook	Dave	3

	•		\rightarrow	
	Variation	Host	Table	
	Pooker	Mary	2	
	Pooker	Dave	1	
	Sneaker	John	1	
	Sneaker	Mary	2	
	Sneaker	?	3	
	Quick Snook	Mary	3	
BCI	Quick Spook	Dave	?	is hosted by /
X 4N	is mapped to		lle (FK)	hosts
	Variation VariationName (other attributes)	is u use Table	sed for/	Member First Name (other attributes)
		Table	e Number r attributes)	

Multivalued dependency

- Possible values for dependent columns depend on determinant only
 - No influence from other columns
 - Special case: always one possible value this is a functional dependency
- Functional dependency:
 - □ $\{a, b\} \rightarrow \{c, d\}$ equivalent to $\{a, b\} \rightarrow c$ and $\{a, b\} \rightarrow d$
- Multivalued dependency:
 - $\neg \{a, b\} \rightarrow \{c, d\} \text{ NOT }$ equivalent to $\{a, b\} \rightarrow c \text{ and } \{a, b\} \rightarrow d$

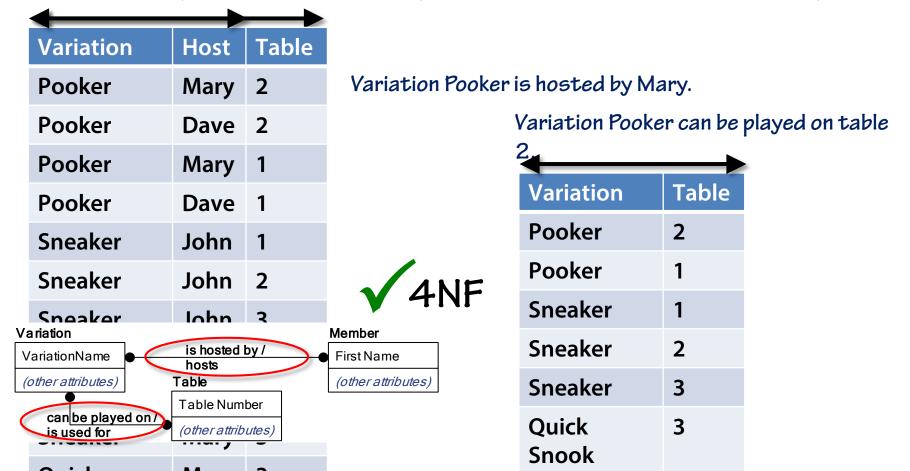
Not relevant for 4NF

- □ Trivial multivalued dependency: <
 - □ Every set of columns → the combination of <u>ALL</u> other columns

- Requirements for Fourth Normal Form (4NF):
 - Table must be in Boyce-Codd Normal Form
 - □ For every non-trivial dependency, the determinant is (a superset of) a key

•		\rightarrow
Variation	Host	Table
Pooker	Mary	2
Pooker	Dave	2
Pooker	Mary	1
Pooker	Dave	1
Sneaker	John	1
Sneaker	John	2
Sneaker	John	3
Sneaker	Mary	1
Sneaker	Mary	2
Sneaker	Mary	3

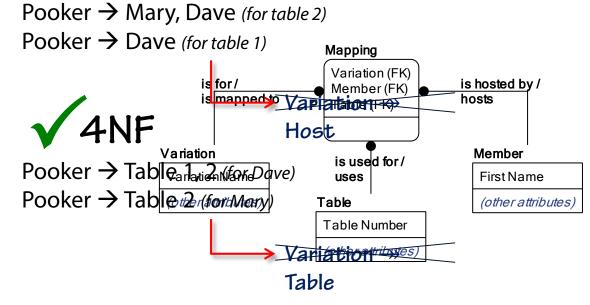
- Requirements for Fourth Normal Form (4NF):
 - Table must be in Boyce-Codd Normal Form
 - For every non-trivial dependency, the determinant is (a superset of) a key



- Requirements for Fourth Normal Form (4NF):
 - Table must be in Boyce-Codd Normal Form
 - For every non-trivial dependency, the determinant is (a superset of) a key

		\rightarrow
Variation	Host	Table
Pooker	Mary	2
Pooker	Dave	2
Pooker	Mary	1
Pooker	Dave	1
Sneaker	John	1
Sneaker	John	2
Sneaker	John	3
Sneaker	Mary	1
Sneaker	Mary	2
Sneaker	Mary	3

Variation Pooker is hosted by Mary on table 2.



		\rightarrow
Variation	Host	Table
Pooker	Mary	2
Pooker	Dave	2
Pooker	Mary	1
Pooker	Dave	1
Sneaker	John	1
Sneaker	John	2
Sneaker	John	3
Sneaker	Mary	1
Sneaker	Mary	2
Sneaker	Mary	3
Quick Snook	Mary	3
Quick Snook	Dave	3

Variation Pooker is hosted by Mary.

Variation Pooker can be played on table2.



4		\rightarrow
Variation	Host	Table
Pooker	Mary	2
Pooker	Dave	2
Pooker	Mary	1
Pooker	Dave	1
Sneaker	John	1
Sneaker	John	2
Sneaker	John	3
Sneaker	Mary	1
Sneaker	Mary	2
Sneaker	Mary	3
Quick Snook	Mary	3
Quick Snook	Dave	3

Variation Pooker is hosted by Mary.

Variation Pooker can be played on table2.

Variation Pooker is hosted by Mary on table 2.



Variation	Host	Table
Pooker	Mary	2
Pooker	Dave	2
Pooker	Mary	1
Pooker	Dave	1
Sneaker	John	1
Sneaker	John	2
Sneaker	John	3
Sneaker	Mary	1
Sneaker	Mary	2
Sneaker	Mary	3
Quick Snook	Mary	3
Quick Snook	Dave	3

Variation Pooker is hosted by Mary.

Variation Pooker can be played on table 2.

Variation Pooker is hosted by Mary on table 2.

Mary can host games on table 2. Mary can host games on table 3.



Variation →
Host
Variation ->
Table
Host-» Variation
Host » Table
Table ->>
Variation

4		\rightarrow
Variation	Host	Table
Pooker	Mary	2
Pooker	Dave	2
Pooker	Mary	1
Pooker	Dave	1
Sneaker	John	1
Sneaker	John	2
Sneaker	John	3
Sneaker	Mary	1
Sneaker	Mary	2
Sneaker	Mary	3
Quick Snook	Mary	3
Quick Snook	Dave	3

Variation Pooker can be played on table 2. ><variation> can be played on .

←		\rightarrow
Variation	Host	Table
Pooker	Mary	2
Pooker	Dave	2
Pooker	Mary	1
Pooker	Dave	1
Sneaker	John	1
Sneaker	John	2
Sneaker	John	3
Sneaker	Mary	1
Sneaker	Mary	2
Sneaker	Mary	3
Quick Snook	Mary	3
Quick Snook	Dave	3

Variation Pooker is hosted by Mary on table 2. ><variation> is hosted by <host> on .

-		\rightarrow
Variation	Host	Table
Pooker	Mary	2
Pooker	Dave	2
Pooker	Mary	1
Pooker	Dave	1
Sneaker	John	1
Sneaker	John	2
Sneaker	John	3
Sneaker	Wary	1
Sneaker	Mary	2
Sneaker	Mary	3
Quick Snook	Mary	3
Quick Snook	Dave	3

0.14100	. I Г						
ormal Form			Pooker		M	lary	
				Pooker		D	ave
Variatio	n Pooker is l	nosted	by l	Sneaker		Jo	hn
	variation		tec	Sneaker		M	lary
	Variation	Table		Quick		M	lary
Variatio	Pooker	2	yed	Snook			
2	Pooker	1	; pla	Quick Snook		D	ave
	Sneaker	1		SHOOK	/_		
Variation	Sneaker	2	by	Mary on	tabl	e	
2	Sneaker	3	ted	by khos	Hos	1	Tabl
Mary ca	Quick Snook	3	ıble	2.			e
. 1				00 00 44	Mar	У	2
	> <host> car</host>	1 11051	gam	es on <1	Dav	e	2
					Dav	e	1
	JOII				Johi	n	1
	JUII				Johi	n	2
					Johi	n	3
					Mar	у	3

Variation

Host

Dave 3

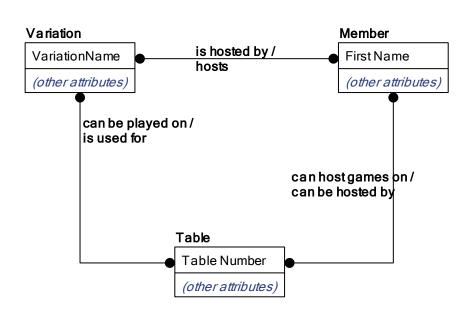
Fifth N

-		\rightarrow
Variation	Host	Table
Pooker	Mary	2
Pooker	Dave	2
Pooker	Mary	1
Pooker	Dave	1
Sneaker	John	1
Sneaker	John	2
Sneaker	John	3
Sneaker	Mary	1
Sneaker	Mary	2
Sneaker	Mary	3
Quick Snook	Mary	3
Quick Snook	Dave	3

la una s	A Eaum			Variation	า	Н	ost	
Iormal Form			Pooker		M	ary		
				Pooker		D	ave	
Variatio	n Pooker is l	10sted	by l	Sneaker		Jo	hn	
	variation		tec	Sneaker		M	ary	
	Variation	Table		Quick		M	ary	
Variatio	Pooker	2	yed	Snook				
<u>4.</u>	Pooker	1	; pla	Quick Snook		D	ave	
	Sneaker	1		SHOOK	/_			
Variatio	Sneaker	2	by	Mary on	tabl	le		
2	Sneaker	3	ted	bykhos	Hos	t	Tab	
Manyos	Quick	3	ıble		1103		e	
Mary ca					Mar	У	2	
	> <host> car</host>	1 host	gam	<i>es o</i> n <1	Dav	e	2	
					Dav	e	1	
	TOH	X			Johi	n	1	
	JOHN				Johi	n	2	
					Johi	n	3	

Mary 3

Dave 3



4	
Variation	Table
Pooker	2
Pooker	1
Sneaker	1
Sneaker	2
Sneaker	3
Quick Snook	3

4	\rightarrow
Variation	Host
Pooker	Mary
Pooker	Dave
Sneaker	John
Sneaker	Mary
Quick Snook	Mary
Quick Snook	Dave

4-	\longrightarrow
Host	Tabl e
Mary	2
Dave	2
Dave	1
John	1
John	2
John	3
Mary	3
Dave	3

Join dependency

- Possible to split table in multiple new tables, ...
- ... such that joining them reconstructs original table
- Attributes represent independent facts
- Trivial join dependency:
 - One of the new tables is equal to original table
- Join dependency implied by candidate key:
 - Candidate key of original table included in every new table
- Join dependency with two new tables:
 - Equivalent to multivalued dependency
 - □ Shared columns → distributed columns

- Requirements for Fifth Normal Form (5NF):
 - Table must be in Fourth Normal Form
 - □ For every non-trivial dependency, the determinant is (a superset of) a key
- Fifth Normal Form is also known as PJ/NF
 - Project-Join Normal Form
 - Projection-Join Normal Form

- Requirements for Domain-Key Normal Form (DK/NF):
 - NOT based on dependencies
 - Based on: Valid values for an attribute (column)
 Domains Candidate keys
 Constraints
 - Rules for valid data
 - Every constraint must be implied by the keys and domains
 - (Implies Fifth –and lower– Normal Form)

Relevance of Domain-Key Normal Form

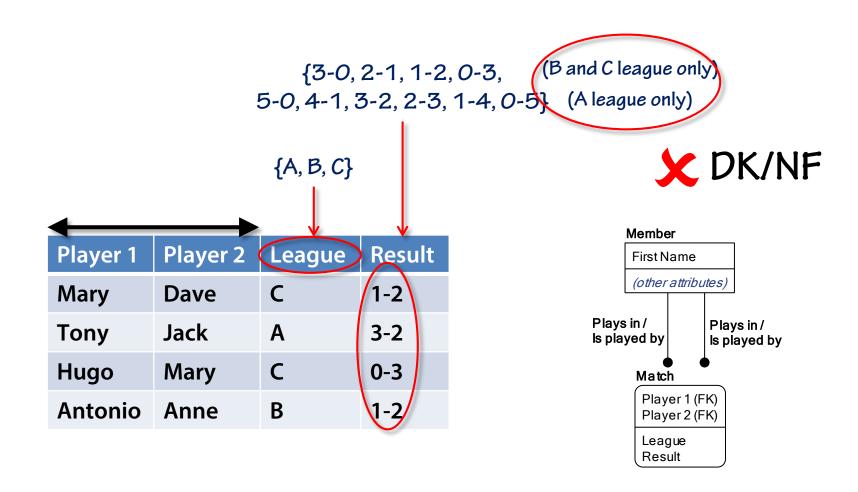
- □ Domains → declared, enforced (no code needed)
- □ Keys → declared, enforced (no code needed)
- Other constraints -> code needed to enforce
 - □ Code = cost factor:
 - Time to write
 - Time to test and debug
 - □ Future maintenance

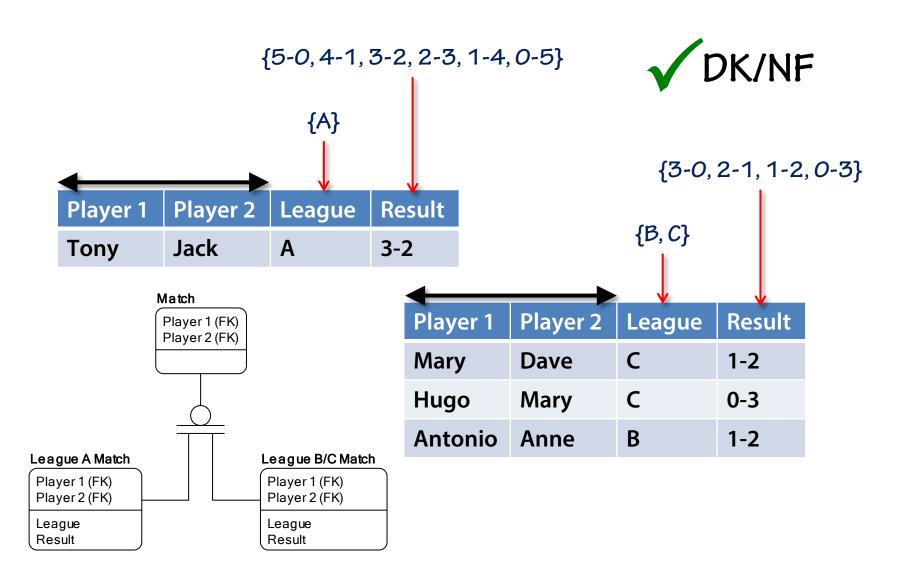
More code for constraints, less code for queries

Achievability of Domain-Key Normal Form

- Sometimes impossible
- Otherwise often requires extra tables (subtypes)
 - Introduces need for more (and more complex) code

Less code for constraints, more code for querying





Sixth Normal Form

Sixth Normal Form (6NF)

Has historically sometimes been used for DK/NF

Fifth Normal Form:

Non-trivial join dependency not allowed, unless implied by a candidate key

Sixth Normal Form:

- Non-trivial join dependency not allowed, unless implied by a candidate key
- Any 6NF table has candidate key + at most one other column

Sixth Normal Form

Relevance of Sixth Normal Form

No NULLs needed in storage

Fifth Normal Form

Member	Birthdate	League	Board position
Mary (NULL	C	Secretary
Jack	15 Dec 1960	Α (NULL
Anthony	8 Apr 1975	В	NULL

Sixth Normal Form

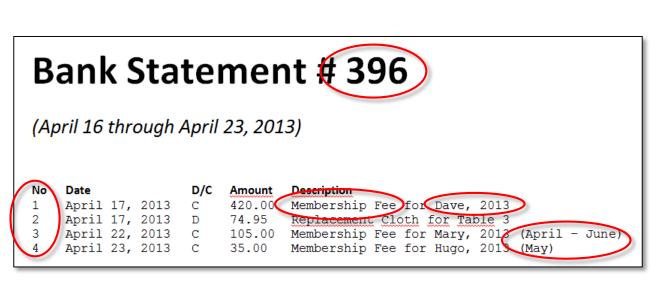
$ \longleftarrow $		
Member	League	
Mary	C	
Jack	Α	
Anthony	В	

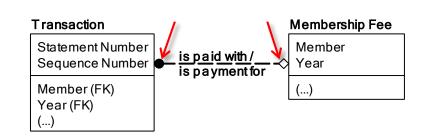
\longleftrightarrow		
Member	Birthdate	
Jack	15 Dec	
	1960	
Anthony	8 Apr 1975	

\longleftrightarrow	
Member	Board position
Mary	Secretary

Sixth Normal Form

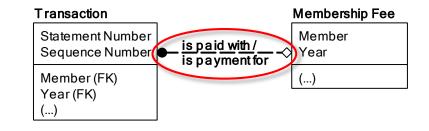
- Relevance of Sixth Normal Form
 - No NULLs needed in storage
- Consequence of Sixth Normal Form
 - Huge number of tables
 - Many joins required for most queries
 - Simple constraints can be difficult to enforce





Sixth Normal Form

- Relevance of Sixth Normal Form
 - No NULLs needed in storage
- Consequence of Sixth Normal Form
 - Huge number of tables
 - Many joins required for most queries
 - Simple constraints can be difficult to enforce



				$\qquad \qquad \longrightarrow$	FOREIGN KEY (Member, Year)		
	Fifth Normal Form				REFERENCES MembershipFees (Member, Year		
	StmtNbr	SeqNbr	Member	Year	()		
	396	1	Dave	2013	()		
	396	2	NULL	NULL	()		
	396	3	Mary	2013	()		
	396	4	Hugo	2013	()		

CONCUDATIVE ToDasmon+For

Sixth Normal Form



No NULLs needed in storage

Consequence of Sixth Normal Form

- Huge number of tables
- Many joins required for most queries
- Simple constraints can be difficult to enforce

Sixth Normal Form

StmtNbr	SeqNbr	Member
396	1	Dave
396	3	Mary
396	4	Hugo

	V	
StmtNbr	SeqNbr	Year
396	1	2013
396	3	2013
396	4	2013

Transaction

Member (FK)

Year (FK)

Statement Number

Sequence Number



Membership Fee

Member

(X) Year

(...)

is paid with

is payment for

Optimal Normal Form

Optimal Normal Form (ONF)

- Based on fact-based modeling methods (e.g. ORM, NIAM)
- Every fact type becomes a table

		
StmtNbr	SeqNbr	Amount
396	1	420.00
396	2	74.95
396	3	105.00
396	4	35.00

\longrightarrow			
StmtNbr	SeqNbr	TranDate	
396	1	2013-04- 17	
396	2	2013-04- 17	
396	3	2013-04- 22	
396 CONSTE	4 RAINT IsPa	2013-04-	

$\overline{}$	\rightarrow			
StmtNbr	SeqNbr	Member	Year	3
396	1	Dave	201	
396	3	Mary	2013	3
396	4	Hugo	2013	3

FOREIGN KEY (Member, Year)

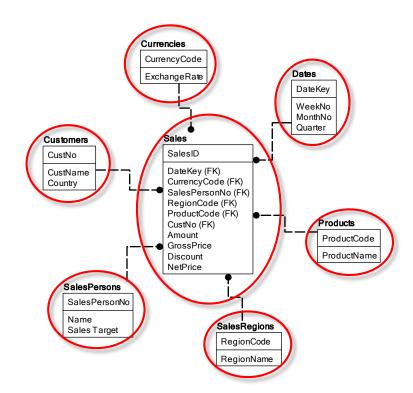
REFERENCES MembershipFees (Member, Year)

- OLTP (Online Transaction Processing)
 - Normalization is always good
- OLAP (Online Analytical Processing) and reporting
 - Normalization may harm performance
 - Solution: denormalize
 - NOTE: Unnormalized is not denormalized!
 - Denormalization is done after normalization

Storing the result of a query

- "Do it yourself"
 - Store query results in table
 - Add application logic to update stored results when underlying data changes
- Feature in some RDBMS's
 - Materialized view
 - Indexed view
- Typical use cases
 - Data stored in separate tables but often queries together
 - Pre-aggregating data

- Star schema
 - Fact table
 - Dimension tables
- Snowflake schema



OLAP cubes

- Specifically designed to support Online Analytical Processing
- Pre-aggregates data on several levels
- Usually built on top of a star or snowflake schema

- Non-First Normal Form (N1NF or NF²)
 - Attribute can be a single value (1NF)
 - OR attribute can be a complete table (violates 1NF)
 - Results in nested tables
 - Rigid design still applies
 - Multi-valued databases comes close

Variation	Host	ing
Pooker	Host	Table
	Mary	2
	Dave	1
	Dave	2
Sneaker	Host	Table
	John	1
	John	2
	Mary	2
	Mary	3
Quick	Host	Table
Snook	Mary	3
	Dave	3

Referencing a range



+ 37 ranking points

Total ranking points: 1,496

Rankingpoints	League
0 - 500	С
500 - 1500	В
1500+	Α

League: B

Functional dependency:

Ranking Points → League

Referencing a range

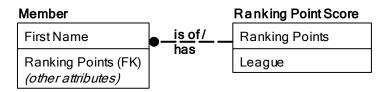
Member			
First Name			
League Ranking Points			
(other attributes)			

\longrightarrow			
FirstName	League	RankingPoints	()
Dave	C	478	()
Hugo	C	271	()
Jack	Α	2080	()
Mary	В	772	()

Functional dependency:

• Ranking Points → League

Referencing a range



		
FirstName	RankingPoints	()
Dave	478	()
Hugo	271	()
Jack	2080	()
Mary	772	()

$\qquad \qquad \longrightarrow$	
RankingPoints	League
0	C
1	C
•••	•••
500	C
501	В
•••	•••
1500	В
1501	Α
•••	•••

Functional dependency:

Ranking Points → League

Referencing a range

The ranking point range that runs from O to 500 corresponds to the C league. The ranking point range that runs from 500 to 1,500 corresponds to the B league.

The ranking point range that runs from 1,500 corresponds to the C league.

Rankingpoints	League
0 – 500	С
500 – 1500	В
1500+	Α

```
(not an entity type) (no duplicates allowed)

(not an entity type) (no duplicates allowed)
```

The ranking point range that runs from <Score> ends at <Score>.

```
(entity type) (no duplicates allowed)

(not an entity type) (no duplicates allowed)
```

The ranking point range that runs from <Score> corresponds to <League>.

Referencing a range

Member
First Name
Ranking Points
(other attributes)

Start score End score (AK1.1) League (AK2.1)

\longleftrightarrow		
FirstName	RankingPoints	()
Dave	478	()
Hugo	271	()
Jack	2080	()
Mary	772	()

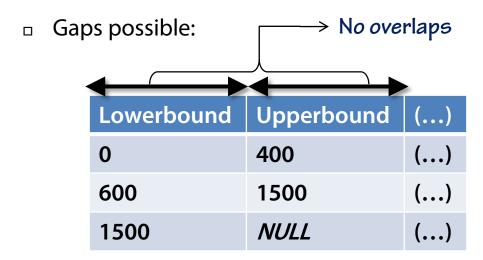
\longleftrightarrow	\longleftrightarrow	\longleftrightarrow
Start score	End score	League
0	500	C
500	1500	В
1500	NULL	A

Referencing a range

□ No gaps:

\longleftarrow	
Lowerbound	()
0	()
500	()
1500	()

No overlaps No gaps		
Lowerbound	Upperbound	()
0	500	()
500	1500	()
1500	NULL	()



Referencing a range

\longleftrightarrow		
FirstName	RankingPts	()
Dave	478	()
Hugo (-271	()
Jack	2080	()
Mary	772	()

Member

First Name

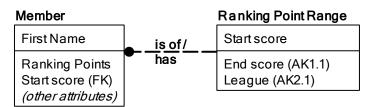
Ranking Points
(other attributes)

Ranking Point Range

Start score
End score (AK1.1)
League (AK2.1)

\longleftrightarrow	\longleftrightarrow	\longleftrightarrow
StartScore	EndScore	League
0	500	C
500	1500	В
1500	NULL	Α

Referencing a range



\longleftrightarrow			
FirstName	RankingPts	StartScore	()
Dave	478	0	()
Hugo	271	0	()
Jack	2080	500	()
Mary	772	500	()

\leftarrow	\longleftrightarrow	\longleftrightarrow
StartScore	EndScore	League
0	500	C
500	1500	В
1500	NULL	Α

Referencing a range

772

Mary



	StartScore	Lilascoic	League
FirstName RankingPts StartScore EndS	0	500	C
Dave 478 0 500	500	1500	В
	1500	NULL	Α
Hugo 271 0 500 Jack 2080 1500 <i>NUL</i>	(,		

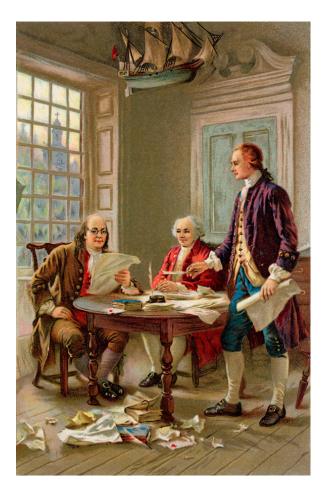
1500

FOREIGN KEY (StartScore, EndScore)
REFERENCES RankingPointRange (StartScore, EndScore)

500

CHECK (RankingPts BETWEEN StartScore AND EndScore)

Storing historic data



- Storing historic data
 - OLTP:
 - Mostly current data
 - □ OLAP:
 - Lots of historic data
 - Standard design pattern ("Slowly changing dimensions")

- Storing historic data
 - OLTP:
 - Mostly current data
 - Some historic data needed as well

$\qquad \qquad \longrightarrow$			
ProductID	Product Name	Price	
123	Snooker cue with extension	208.5	K
456	Standard snooker ball set	65.00	67.50
789	Replacement felt	195.0 0	

- Storing historic data
 - OLTP:
 - Mostly current data
 - Some historic data needed as well

ProductID	Product Name	Price	OldPrice	DatePriceChange
123	Snooker cue with extension	208.5	NULL	NULL
456	Standard snooker ball set	65.00	67.50	Oct 5, 2013
789	Replacement felt	195.0 0	NULL	NULL

- Storing historic data
 - OLTP:
 - Mostly current data
 - Some historic data needed as well

—			\rightarrow
ProductID	Product Name	Price	ValidTo
123	Snooker cue with extension	208.5	NULL
456	Standard snooker ball set	65.00	Oct 5, 2013
456	Standard snooker ball set	67.50	NULL
789	Replacement felt	195.0 0	NULL

- Storing historic data
 - □ OLTP:
 - Mostly current data
 - Some historic data needed as well

Sold "Standard snooker ball set" on August 12, 2013

ProductID	Product Name	Price	ValidTo
123	Snooker cue with extension	208.5	Dec 31, 9999
456	Standard snooker ball set	65.00	Oct 5, 2013
456	Standard snooker ball set	67.50	Dec 31, 9999
789	Replacement felt	195.0 0	Dec 31, 9999

Storing historic data

OLTP:

Mostly current data

Some historic data needed

ProductID	Product Name	Price
123	Snooker cue with extension	208.5
456	Standard snooker ball set	67.50
789	Replacement felt	195.0 0

ProductID	Product Name	Price	ValidTo
123	Snooker cue with extension	208.5	Dec 31, 9999
456	Standard snooker ball set	65.00	Oct 5, 2013
456	Standard snooker ball set	67.50	Dec 31, 9999
789	Replacement felt	195.0 0	Dec 31, 9999

Storing historic data

OLTP:

Mostly current data

Some historic data needed

ProductID	Product Name	Price
123	Snooker cue with extension	208.5
456	Standard snooker ball set	67.50
789	Replacement felt	195.0 0

ProductID	Product Name	Pri	ce	ValidTo
456	Standard snooker ba	all set 65	5.00	Oct 5, 2013

ProductID	Product Name	Price	ValidTo
123	Snooker cue with extension	208.5	NULL
456	Standard snooker ball set	65.00	Oct 5, 2013
456	Standard snooker ball set	67.50	NULL
789	Replacement felt	195.0	NULL

Storing historic data

□ OLTP:

Mostly current data

Some historic data needed

ProductID	Product Name	Price
123	Snooker cue with extension	208.5
456	Standard snooker ball set	67.50
789	Replacement felt	190.0
		0

ProductID	Product Name	Price	ValidTo	
456	Standard snooker ball set	65.00	Oct 5, 2013	
789	Replacement felt	195.0	Oct 27, 2013	То
		0		,

123	extension	0	IVOLL
456	Standard snooker ball set	65.00	Oct 5, 2013
456	Standard snooker ball set	67.50	NULL
789	Replacement felt	195.0	Oct 27, 2013

- Storing historic data
 - OLTP:
 - Mostly current data
 - Some historic data needed as well

	ProductID	ValidTo	Price
	123	Dec 31, 9999	208.5
	456	Oct 5, 2013	65.00
	456	Dec 31, 9999	67.50
	789	Oct 27, 2013	195.0 0
it	789	Dec 31, 9999	190.0 0

ProductID	ValidTo	Product Name	
123	Dec 31, 9999	Snooker cue wit 789 extension	
456	Dec 31, 9999	Standard snooker ball set	
789	Dec 31, 9999	Replacement felt	

Summary

Higher normal forms

- Elementary Key Normal Form (EKNF)
- Boyce-Codd Normal Form (BCNF)
- Fourth Normal Form (4NF)
- Fifth Normal Form (5NF)
- Domain/Key Normal Form (DK/NF)
- Sixth Normal Form (6NF)
 - Optimal Normal Form (ONF)

Denormalization

Non-standard problems

- Referencing a range
- Storing historic data

References

Further reading:

- Database normalization on Wikipedia:
 http://en.wikipedia.org/wiki/Database_normalization
 or
 http://tinyurl.com/Norm-DB
- Pluralsight Course Library
 - Business Intelligence section