

Discussion #6

Covers: Chapter 6

Questions:

1. The following pseudocode illustrates the basic push() and pop() operations of an array-based stack. Assuming that this algorithm could be used in a concurrent environment, answer the following questions:
 - a. What data have a race condition?
 - b. How could the race condition be fixed?

```
push(item) {
    if (top < SIZE) {
        stack[top] = item;
        top++;
    }
    else
        ERROR
}

pop() {
    if (!is empty()) {
        top--;
        return stack[top];
    }
    else
        ERROR
}

is empty() {
    if (top == 0)
        return true;
    else
        return false;
}
```

2. Consider how to implement a mutex lock using the compare and swap() instruction. Assume that the following structure defining the mutex lock is available:

```
typedef struct {
    int available;
} lock;
```

The value (available == 0) indicates that the lock is available, and a value of 1 indicates that the lock is unavailable. Using this struct, illustrate how the following functions can be implemented using the compare and swap() instruction:

- void acquire(lock *mutex)
- void release(lock *mutex)

Be sure to include any initialization that may be necessary.

3. The first known correct software solution to the critical-section problem for two processes was developed by Dekker. The two processes, P0 and P1, share the following variables:

```
boolean flag[2];
int turn;
```

The structure of process P_i ($i == 0$ or 1) is shown below. The other process is P_j ($j == 1$ or 0). Prove that the algorithm satisfies all three requirements for the critical-section problem.

```
do
{
    flag[i] = true;
    while (flag[j])
    {
        if (turn == j)
        {
            flag[i] = false;
            while (turn == j)
                ; /* do nothing */
            flag[i] = true;
        }
    }
    /* critical section */
    turn = j;
    flag[i] = false;
    /* remainder section */
} while (true);
```

4. The following algorithm shows how to use the test_and_set instruction to implement the critical section problem for n processes (consider that the code is for process P_i). The common data structures used by the algorithm are:

```
boolean waiting[n];
boolean lock;
```

The elements in the waiting array are initialized to false, and lock is initialized to false. Show that the solution satisfies the bounded waiting requirement.

```
do {
    waiting[i] = true;
    key = true;
    while (waiting[i] && key)
        key = test_and_set(&lock);
    waiting[i] = false;
    /* critical section */
    j = (i + 1) % n;
    while ((j != i) && !waiting[j])
        j = (j + 1) % n;
    if (j == i)
        lock = false;
    else
        waiting[j] = false;
    /* remainder section */
} while (true);
```