

Q.3)

Lagrange:

Assumption:

We are NOT creating the polynomial; we are interpolating on a point or more.

Data:

Input: $n+1$

Output: N

Analysis:

There are two main computations involved PER Point:

(1) The Coefficients, $L_{n,k}(x)$:

- Requires n -additions, n -multiplications in the numerator, and the same in the denominator.

$$\Rightarrow O(4*(n+1)) = O(n) \text{ PER Coefficient.} \quad (1.1)$$

- There are n -Coefficients to be computed,

$$\Rightarrow O(n+1) \quad (1.2)$$

(2) The polynomial, $\sum L_{n,k}(x) * y_k$:

- Requires n -additions and n -multiplications

$$\Rightarrow O(2*(n+1)) = O(n+1) \quad (2.1)$$

To compute 1-Output:

$$\Rightarrow ((1.1) * (1.2)) + (2.1)$$

$$= O(4*(n+1)) * (n+1) + O(2*(n+1))$$

$$= O(4n^2) + O(2n)$$

$$= O(4n^2 + 2n)$$

$$= O(n^2 + n)$$

$$= \mathbf{O(n^2)} \quad \textbf{(I)}$$

To compute N-Outputs:

$$\Rightarrow \textbf{(I)} * O(N)$$

$$= O(n^2) * O(N)$$

$$= O(n^2 * N)$$

$$= \mathbf{O(N * n^2)} \quad \textbf{(II)}$$

Notice: To construct the polynomial, the required complexity is $O(n^2 * \log^2(n))$

Barycentric:

Assumption:

1. We are NOT creating the polynomial; we are interpolating on a point or more.
2. We are saving the weights and coefficient (pre-computing) .

Data:

Input: $n+1$
Output: N

Analysis:

There are three main computations involved:

(1) The Coefficients, $\ell(x)$:

- Requires n -additions and n -multiplications
 $\Rightarrow O(2*(n+1)) = O(n)$ (1.1)

(2) The weights, w_j :

- Requires n -additions and n -multiplications
 $\Rightarrow O(2*(n+2)) = O(n)$ PER weight (2.1)

- There are n -weights to compute:
 $\Rightarrow O(2*(n+2)) = O(n)$ PER point (2.2)

(3) The polynomial, $\sum w_j/x-x_j * y_k$:

- Requires $2n$ -additions and n -multiplications
 $\Rightarrow O(3*(n+1)) = O(n)$ (3.1)

To compute 1-Output:

$$\begin{aligned} &\Rightarrow (1.1) + ((2.1) * (2.2)) + (3.1) \\ &= O(2*(n+1)) + O(2*(n+2)) * (2*(n+2)) + O(3*(n+1)) \\ &= O(2n) + O(4n^2) + O(3n) \\ &= O(4n^2 + 5n) \\ &= O(n^2 + n) \\ &= \mathbf{O(n^2)} \end{aligned} \quad \text{(I)}$$

To compute N-Outputs:

$$\begin{aligned} &\Rightarrow (3.1) * O(N) \\ &= O(n) * O(N) \\ &= O(n * N) \\ &= \mathbf{O(N * n)} \end{aligned} \quad \text{(II)}$$

Q.4)

- There are two sides to this question.
 - Why? Because convergence depends heavily on
 - (1) Point Nodes
 - (2) The Function
- The first function:
 - is absolutely continuous on the interval $[-1,1]$ and, thus,
 - has the property that the interpolating polynomial constructed on Chebyshev nodes will converge uniformly to it as $n \rightarrow \infty$.
- The second function:
 - is absolutely continuous on the interval $[-1,1]$ and, thus,
 - has the property that the interpolating polynomial constructed on Chebyshev nodes will converge to it as $n \rightarrow \infty$.
- The polynomial, using the linspace nodes, will converge to Neither of the two functions. [Runge's Phenomena]
 - linspace nodes are equidistant
 - Convergence of polynomials using equidistant nodes is not guaranteed, even, for infinitely differentiable functions.
 - The error term could potentially go to infinity as n goes to infinity.
 - The first function has an absolute value term, which is very hard to capture.
 - The second function has ill-behaved higher derivatives.
 - In-fact, I can show that the error tends to infinity as $n \rightarrow \infty$.
- Finally, notice that, the Banach–Steinhaus theorem tells us that for there to exist a table of nodes that allow the interpolating polynomial to converge to the function as $n \rightarrow \infty$ for all the functions in C^1 you need the operator norm of the projection operator on the vector space of polynomials of degree less than or equal to n is to be bounded.
- However, Lebesgue constant tells us that the projection operator norm is bigger than or equal to $\frac{2}{\pi} \log(n + 1) + \text{const.}$
- Thus, it is impossible to find a table of nodes that allows the

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lagrange polynomial to converge to all functions.