

Hw 12 Scalability

Case 1:

Computation Cost:

→ 1 iteration, work per processor:

$$T_{\text{comp}} = \gamma \cdot \frac{n^2}{p}$$

→ $n = 10^6$, $p = 1000$, $\gamma = 10^{-9}$

$$T_{\text{comp}} = 10^{-9} \times \frac{(10^6)^2}{1000} = 10^{-9} \times \frac{10^{12}}{10^3} = 10^{-9} \times 10^9 = 1 \text{ s}$$

Communication Cost:

→ each process must exchange boundary data with its top and bottom neighbors. Size of each boundary is n data points.

→ Single message:

$$T_{\text{msg}} = a + \beta n = 10^{-6} + 10^{-9} \times 10^6 = 10^{-6} + 10^{-3} = 0.001001 \text{ s}$$

→ both top and bottom:

$$T_{\text{comm}} \approx 2 \times T_{\text{msg}} \approx 0.002 \text{ s}$$

Weak Scaling Consideration

→ $N_{\text{local}} = \frac{n^2}{p} = \text{constant}$, therefore $n \propto \sqrt{p}$, thus,

$$T_{\text{comm}} \approx a + \beta n \propto a + \sqrt{p}$$

∴ this grows as \sqrt{p} , and in turn does not scale weakly

Case 2

→ each processor Subdomain Size:

$$\frac{n}{\sqrt{p}} \times \frac{n}{\sqrt{p}}, \quad N_{\text{local}} = \frac{n^2}{p}$$

Computation Cost:

$$T_{\text{comp}} = \gamma \times \frac{n^2}{p} = 1 \text{ s} \quad (\text{same as case 1})$$

Communication Cost:

→ each process exchanges boundary data across 4 edges: Each edge length is $\frac{n}{\sqrt{p}}$, thus:

$$T_{\text{edge}} = \alpha + \beta \frac{n}{\sqrt{p}}$$

$$\rightarrow n = 10^6, \sqrt{p} = \sqrt{1000} \approx 31.62$$

$$T_{\text{edge}} = 10^{-6} + 10^{-9} \times \frac{10^6}{31.62} = 3.26 \times 10^{-5} \text{ s}$$

→ assuming 4 neighbors:

$$T_{\text{comm}} = 4 \times T_{\text{edge}} = 1.3 \times 10^{-4} \text{ s}$$

Weak Scaling Consideration:

→ $N_{\text{local}} = \frac{n^2}{p} = \text{constant}$, therefore $n \propto \sqrt{p}$ so that $\frac{n}{\sqrt{p}} = \text{constant}$

→ Thus communication cost remains constant as p increases and does scale weakly