

Assiut University

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1	Data structures	1	8.5 Derivatives/Integrals	else
	1.1 Ordered Set	1	8.6 Sums	res = merg
	1.2 Sparse Table	1	8.7 Series	first = fals 1 += (1 << i
	1.3 Fenwick Tree	1	8.8 Probability theory	}
	1.4 Segment Tree	2	8.8.1 Discrete distributions	return res;
	1.5 Treap	3	8.8.2 Continuous distributions	} };
	1.6 Line Container	5	0.0.2 Continuous distributions 20	, , , , , , , , , , , , , , , , , , ,
	1.7 MO	5	Data structures (1)	1.3 Fenwick
	1.1 1/10	0	$\underline{\text{Data structures}}$ (1)	FenwickTree.h
2	Graph	5	1.1 Ordered Set	template <typename t=""></typename>
	2.1 Graph Data structures	5	OrderedSet h	struct fenwick_tree
	2.2 Flow	8	52819a, 10 lines	<pre>int n; vector<t> BIT;</t></pre>
	2.2.1 Bipartite Matching	8	<pre>#include<ext assoc_container.hpp="" pb_ds=""> // keep-include #include<ext pb_ds="" tree_policy.hpp=""> // keep-include</ext></ext></pre>	fenwick_tree(int
	2.2.2 Max Flow	8	using namespacegnu_pbds;	T getAccum(int i
	2.2.3 Min Cost Max Flow	9	template <typename key=""></typename>	T sum = 0; while (idx)
	2.3 Shortest Path	9	<pre>using ordered_set = tree<key, less<key="" null_type,="">, rb_tree_tag , tree_order_statistics_node_update>;</key,></pre>	sum += B
	2.4 Tarjan	10	// find_by_order(k) :	idx -= (
	·		// It returns to an iterator to the k-th element (counting	} return sum;
3	Strings	11	from zero) in the set in O(logn) time. // To find the first element k must be zero.	}
	3.1 Rolling Hash	11	$//$ order_of_key(k) :	void add(int idx
	3.2 String matching	11	// It returns to the number of items that are strictly smaller	assert(idx ! while (idx <
	3.3 Suffix structures	12	than our item k in $O(logn)$ time.	BIT[idx]
	3.4 Manacher Hash	14	1.2 Sparse Table	idx += (
			SparseTable.h	}
4	Math	14	Sparse Lable. II fb52af, 41 lines	T getValue(int i
-	4.1 Prime	14	template <typename t=""></typename>	return getAc
	4.2 ModInverse	15	<pre>struct sparse_table { vector<vector<t>> sparseTable;</vector<t></pre>	// ordered stati
	4.3 Number Theory	15	<pre>using F = function<t(t,t)>;</t(t,t)></pre>	// get index the
	4.4 Combinatorics	16	F merge;	// values must b
	4.5 Matrix	16	<pre>static int LOG2(int x) { //floor(log2(x)) return 31builtin_clz(x);</pre>	<pre>int getIdx(T acc int start =</pre>
	4.6 FFT	17	}	while (start
	1.0 111	11	<pre>sparse_table(vector<t> &v, F merge) : merge(merge) {</t></pre>	<pre>int mid T val =</pre>
5	Geometry	18	<pre>int n = v.size();</pre>	if (val
•	5.1 Basic Geometry	18	<pre>int logN = LOG2(n);</pre>	rt =
	5.2 Circles	19	<pre>sparseTable = vector < vector < T >> (logN + 1); sparseTable[0] = v;</pre>	else star
	5.3 Polygons	20	for (int k = 1, len = 1; k <= logN; k++, len <<= 1) {	}
	0.0 Torygons	20	sparseTable[k].resize(n);	return rt;
6	Misc	21	<pre>for (int i = 0; i + len < n; i++) sparseTable[k][i] = merge(sparseTable[k - 1][i],</pre>	//need review (f
U	WHISE	21	sparseTable[k - 1][i + len]);	//last index les
7	Mathematics Notes	22	}	int find(T accum
•	7.1 Sums	22	T query(int 1, int r) {	<pre>int i = 1, i while ((1 <<</pre>
	7.2 Congruence	22	int $k = LOG2(r - 1 + 1)$; // max $k \Longrightarrow 2^k \le length$ of	i <<= 1;
	7.3 $E(x^2)$	$\frac{22}{22}$	range	for (; i > 0
	1.0 D (a)	44	//check first 2^k from left and last 2^k from right // overlap	int tidx if (tidx
Q	KACTL Maths Notes	22	<pre>return merge(sparseTable[k][1], sparseTable[k][r - (1 << k)</pre>	cont
0		22 22	+ 1]);	if (accu
	8.1 Equations		T query_shifting(int 1, int r) {	idx accu
	8.2 Recurrences	22	T res;	}
	8.3 Trigonometry	22	<pre>bool first = true; for (int i = (int) sparseTable.size() - 1; i >= 0; i)</pre>	} if (!accum
	8.4 Geometry	23	if (1 + (1 << i) - 1 <= r) {	return idx +
	8.4.1 Triangles	23	<pre>if (first)</pre>	}
	8.4.2 Spherical coordinates	23	<pre>res = sparseTable[i][1];</pre>	};

```
ge(res, sparseTable[i][l]);
se;
i);
```

Tree

a2e116, 57 lines

```
n) : n(n), BIT(n + 1) {}
idx) {
BIT[idx];
(idx & -idx);
x, T val) {
!= 0);
<= n) {
+= val;
(idx & -idx);
idx) {
ccum(idx) - getAccum(idx - 1);
istics tree
at has value >= accum
be positive
cum) {
1, end = n, rt = -1;
<= end) {
= start + end >> 1;
getAccum(mid);
>= accum)
= mid, end = mid - 1;
rt = mid + 1;
from topcoder)
ss than or equal accum O(\log n)
m) {
idx = 0;
< i) <= n)
); i >>= 1) {
x = idx + i;
( > n)
tinue;
um >= BIT[tidx]) {
= tidx;
um -= BIT[tidx];
|| idx + 1 > n)return -1;
1;
```

```
FenwickTreeUpdateRange.h
                                                     e3d8f2, 47 lines
// x[i] = a[i] - a[i-1] //a is original array
// y[i] = x[i]*(i-1)
// sum(1,3) = a[1] + a[2] + a[3]
// = (x[1]) + (x[2] + x[1]) + (x[3] + x[2] + x[1])
// = 3*(x[1] + x[2] + x[3]) - 0*x[1] - 1*x[2] - 2*x[3] \Rightarrow same
     equation but more complex
// = sum X(1,3) * 3 - sum Y(1,3)
// so sum(1,n) = sumX(1,n)*n - sumY(1,n)
// x[l] += val, x[r+1] = val
// y[l] += val *(l-1), y[r+1] -= r*val
template<typename T>
class fenwick_tree {
  int n;
  vector<T> x, y;
  T getAccum(vector<T> &BIT, int idx) {
   T sum = 0;
    while (idx) {
     sum += BIT[idx];
     idx = (idx & -idx);
    return sum;
  void add(vector<T> &BIT, int idx, T val) {
    assert(idx != 0);
    while (idx <= n) {
     BIT[idx] += val;
     idx += (idx & -idx);
  T prefix_sum(int idx) {
   return getAccum(x, idx) * idx - getAccum(y, idx);
public:
  fenwick_tree(int n) :
     n(n), x(n + 1), y(n + 1) {
  void update_range(int 1, int r, T val) {
    add(x, l, val);
    add(x, r + 1, -val);
    add(y, 1, val * (1 - 1));
    add(y, r + 1, -val * r);
  T range_sum(int 1, int r) {
    return prefix_sum(r) - prefix_sum(l - 1);
};
```

1.4 Segment Tree

SegmentTree.h

c0adcc, 85 lines

```
// for efficient memory (2*n)
// #define LEFT (idx+1)
// #define MID ((start+end)>>1)
// #define RIGHT (idx+((MID-start+1)<<1))
template<typename T>
class segment_tree {//1-based
#define LEFT (idx<<1)
#define RIGHT (idx<<1|1)
#define MID ((start+end)>>1)
  vector<T> tree, lazy;
  T merge(const T& left, const T& right) {}
  inline void pushdown (int idx, int start, int end) {
   if (lazy[idx] == 0) return;
   //update tree[idx] with lazy[idx]
```

```
tree[idx] += lazy[idx];
    if (start != end) {
     lazv[LEFT] += lazv[idx];
     lazy[RIGHT] += lazy[idx];
    //clear lazy
    lazy[idx] = 0;
  inline void pushup(int idx) {
    tree[idx] = merge(tree[LEFT], tree[RIGHT]);
  void build(int idx, int start, int end) {
    if (start == end) return;
   build(LEFT, start, MID);
   build(RIGHT, MID + 1, end);
    pushup(idx);
  void build(int idx, int start, int end, const vector<T>& arr)
    if (start == end) {
     tree[idx] = arr[start];
      return;
    build(LEFT, start, MID, arr);
   build(RIGHT, MID + 1, end, arr);
    pushup(idx);
 T query(int idx, int start, int end, int from, int to) {
    pushdown(idx, start, end);
    if (from <= start && end <= to)</pre>
      return tree[idx];
    if (to <= MID)
      return query (LEFT, start, MID, from, to);
    if (MID < from)</pre>
      return query (RIGHT, MID + 1, end, from, to);
    return merge (query (LEFT, start, MID, from, to),
      query(RIGHT, MID + 1, end, from, to));
  void update (int idx, int start, int end, int lq, int rq,
    const T & val) {
    pushdown(idx, start, end);
    if (rg < start || end < lg)</pre>
    if (lq <= start && end <= rq) {
     lazy[idx] += val; //update lazy
     pushdown(idx, start, end);
      return;
    update(LEFT, start, MID, lq, rq, val);
    update(RIGHT, MID + 1, end, lq, rq, val);
    pushup(idx);
public:
  segment\_tree(int n) : n(n), tree(n << 2), lazy(n << 2) {
  segment_tree(const vector<T>& v) {
   n = v.size() - 1;
   tree = vector<T>(n << 2);
   lazy = vector < T > (n << 2);
   build(1, 1, n, v);
 T query(int 1, int r) {
    return query(1, 1, n, 1, r);
 void update(int 1, int r, const T& val) {
    update(1, 1, n, 1, r, val);
#undef LEFT
#undef RIGHT
```

```
#undef MID
};
ExtendedSegmentTree.h
                                                     16b868, 64 lines
struct segtree {
  segtree *left = nullptr, *right = nullptr;
  int mx = 0;
  segtree(int val = 0) :
      mx(val) {
  void extend() {
    if (left == nullptr) {
     left = new segtree();
      right = new segtree();
  void pushup() {
    mx = max(left->mx, right->mx);
  ~segtree() {
    if (left == nullptr)return;
    delete left;
    delete right;
};
class extened segment tree {
#define MID ((start+end)>>1)
  void update(segtree *root, int start, int end, int pos, int
    if (pos < start || end < pos)</pre>
      return;
    if (start == end) {
      root->mx = max(root->mx, val);
      return:
    root->extend();
    update(root->left, start, MID, pos, val);
    update(root->right, MID + 1, end, pos, val);
    root->pushup();
  int query(segtree *root, int start, int end, int 1, int r) {
    if (root == nullptr || r < start || end < 1)</pre>
      return 0;
    if (1 <= start && end <= r)
      return root->mx;
    return max(query(root->left, start, MID, 1, r),
        query(root->right, MID + 1, end, 1, r));
public:
  int start, end;
  seatree *root;
  extened_segment_tree() {
  ~extened_segment_tree() {
    delete root:
  extened segment tree (int start, int end) :
      start(start), end(end) {
    root = new segtree();
  void update(int pos, int val) {
    update(root, start, end, pos, val);
  int query(int 1, int r) {
    return query (root, start, end, 1, r);
#undef MID
```

PersistentSegmentTree.h

```
52a616, 86 lines
```

```
struct segtree {
  static segtree *sentinel;
  segtree *left, *right;
  bool dirty = false;
  11 \text{ sum} = 0;
  11 lazy = 0;
  segtree(11 val = 0):
     sum(val) {
    left = right = this;
  segtree (segtree *left, segtree *right) :
     left(left), right(right) {
    sum = left->sum + right->sum;
};
segtree *segtree::sentinel = new segtree();
class persistent segment tree {
#define MID ((start+end)>>1)
  segtree* apply(segtree *root, int start, int end, 11 val) {
    segtree *rt = new segtree(*root);
    rt->dirty = true;
    rt->sum += (end - start + 1) \star val;
    rt->lazv += val;
    return rt:
  void pushdown(segtree *root, int start, int end) {
    if (root->dirty == false || start == end)
    root->left = apply(root->left, start, MID, root->lazy);
    root->right = apply(root->right, MID + 1, end, root->lazy);
    root -> lazv = 0;
    root->dirty = 0;
  segtree* build(int start, int end, const vector<int> &v) {
    if (start == end)
      return new segtree(v[start]);
    return new segtree (build (start, MID, v), build (MID + 1, end
  segtree* Set(segtree *root, int start, int end, int pos, 11
      new_val) {
    pushdown (root, start, end);
    if (pos < start || end < pos)</pre>
      return root;
    if (pos <= start && end <= pos)</pre>
      return new segtree (new_val);
    return new segtree (Set (root->left, start, MID, pos, new_val
        Set(root->right, MID + 1, end, pos, new_val));
  segtree* update(segtree *root, int start, int end, int 1, int
       r, 11 val) {
    pushdown (root, start, end);
    if (r < start || end < 1)
      return root;
    if (1 <= start && end <= r)</pre>
      return apply (root, start, end, val);
    return new segtree (update (root->left, start, MID, 1, r, val
        update(root->right, MID + 1, end, 1, r, val));
  11 query(segtree *root, int start, int end, int 1, int r) {
    pushdown (root, start, end);
```

```
if (r < start || end < 1)</pre>
      return 0:
    if (1 <= start && end <= r)
      return root->sum;
    return query (root->left, start, MID, 1, r)
        + query(root->right, MID + 1, end, 1, r);
public:
 int start, end;
 vector<segtree*> versions;
 persistent_segment_tree(int start, int end) :
      start(start), end(end) {
    versions.push_back(segtree::sentinel);
 persistent_segment_tree(const vector<int> &v) :
      start(0), end(v.size() - 1) {
    versions.push_back(build(start, end, v));
  void update(int 1, int r, 11 val) {
    versions.push_back(update(versions.back(), start, end, 1, r
         , val));
  11 query(int time, int 1, int r) {
    return query(versions[time], start, end, 1, r);
#undef MID
};
IterativeSegmentTree.h
                                                     3f1e4d, 67 lines
struct node {
  node() { //set Default value
  node (const node &a, const node &b) {
 void apply(int val) {
};
struct segment_tree {
  int n; //0 to n-1
  vector<node> tree;
  segment_tree(int n) {
    resize(n);
    build();
  template<typename T>
  segment_tree(const vector<T> &arr) {
    resize(arr.size());
    for (int i = 0; i < arr.size(); i++)</pre>
     tree[n + i] = arr[i];
    build();
  void resize(int n) {
    int p = 1;
    while (p < n)</pre>
     p <<= 1;
    this->n = p;
    tree = vector < node > (p << 1);
  void build() {
    for (int i = n - 1; i > 0; i--)
     tree[i] = node(tree[i << 1], tree[i << 1 | 1]);
  template<typename T>
  void update(int p, const T &value) {
```

```
tree[p += n].apply(value);
    for (int i = p / 2; i > 0; i >>= 1)
      tree[i] = node(tree[i << 1], tree[i << 1 | 1]);</pre>
 node query(int 1, int r) { //[l, r]
    node resl, resr; //set default value in node
    for (1 += n, r += n + 1; 1 < r; 1 >>= 1, r >>= 1) {
     if (1 & 1) {
        resl = node(resl, tree[1]);
     if (r & 1) {
       r--;
        resr = node(tree[r], resr);
    return node (resl, resr);
 int kth_one(int k, int i = 1) {
   if (k > tree[i])
      return -1;
    if (i >= n)
      return i - n;
    if (tree[i << 1] >= k)
      return kth_one(k, i << 1);</pre>
    return kth_one(k - tree[i << 1], i << 1 | 1);</pre>
};
```

Treap 1.5

ImplicitTreap.h

```
97bfaf, 167 lines
enum DIR {
    L, R
};
template<typename T>
struct cartesian tree {
    static cartesian_tree<T> *sentinel;
    int priority = 0, size = 0;
    bool reverse = false;
    cartesian_tree *child[2];
    cartesian_tree *parent;
    cartesian_tree() {
        key = T();
        priority = 0;
        child[L] = child[R] = parent = this;
    cartesian_tree(const T &x, int y) :
            key(x), priority(y) {
        size = 1;
        child[L] = child[R] = sentinel;
        parent = sentinel;
    void push_down() {
        if (!reverse)
            return;
        reverse = 0;
        child[L]->doReverse();
        child[R]->doReverse();
    void doReverse() {
        reverse ^= 1;
        swap(child[L], child[R]);
    void push_up() {
        if (child[L] != sentinel)child[L]->parent = this;
        if (child[R] != sentinel)child[R]->parent = this;
```

```
size = child[L]->size + child[R]->size + 1;
};
template<typename T>
cartesian tree<T> *cartesian tree<T>::sentinel = new
    cartesian tree<T>();
template<typename T, template<typename > class cartesian_tree>
class implicit_treap { //1 based
    typedef cartesian_tree<T> node;
    typedef cartesian_tree<T> *nodeptr;
#define emptyNode cartesian_tree<T>::sentinel
    nodeptr root;
    void split (nodeptr root, nodeptr &1, nodeptr &r, int
        firstXElment) {
        if (root == emptyNode) {
            1 = r = emptyNode;
            return:
        root->push down();
        if (firstXElment <= root->child[L]->size) {
            split(root->child[L], 1, root->child[L],
                 firstXElment);
            r = root;
       } else {
            split(root->child[R], root->child[R], r,
                  firstXElment - root->child[L]->size - 1);
            1 = root;
        root->push_up();
    nodeptr merge(nodeptr 1, nodeptr r) {
        1->push_down();
        r->push_down();
        if (1 == emptyNode | | r == emptyNode)
            return (1 == emptyNode ? r : 1);
        if (l->priority > r->priority) {
            1->child[R] = merge(1->child[R], r);
            1->push_up();
            return 1;
        r->child[L] = merge(l, r->child[L]);
        r->push up();
        return r;
    vector<nodeptr> split_range(int s, int e) { // [x < s, s < = x < = e]
         , e < x /
        nodeptr 1, m, r, tmp;
        split(root, 1, tmp, s - 1);
        split(tmp, m, r, e - s + 1);
        return {1, m, r};
   map <T, nodeptr> mp;
public:
    implicit_treap() :
            root(emptyNode) {
    int size() {
        return root->size;
    void insert(int pos, const T &key) {
       nodeptr tmp = new node(key, rand());
        nodeptr 1, r;
        split(root, 1, r, pos - 1);
        root = merge(merge(1, tmp), r);
```

```
void push_back(const T &value) {
       nodeptr tmp = new node(value, rand());
       mp[value] = tmp;
       root = merge(root, tmp);
   T getByIndex(int pos) {
       vector<nodeptr> tmp = split_range(pos, pos);
       nodeptr l = tmp[0], m = tmp[1], r = tmp[2];
       T rt = m->key;
        root = merge(merge(1, m), r);
       return rt;
    void erase(int pos) {
       vector<nodeptr> tmp = split_range(pos, pos);
       nodeptr 1 = tmp[0], m = tmp[1], r = tmp[2];
       delete m;
       root = merge(1, r);
    void cyclic_shift(int s, int e) { //to the right
       vector<nodeptr> tmp = split_range(s, e);
       nodeptr l = tmp[0], m = tmp[1], r = tmp[2];
       nodeptr first, second;
        split(m, first, second, e - s);
       root = merge(merge(merge(1, second), first), r);
   void reverse_range(int s, int e) {
       vector<nodeptr> tmp = split_range(s, e);
       nodeptr l = tmp[0], m = tmp[1], r = tmp[2];
       m->reverse ^= 1;
       root = merge(merge(1, m), r);
   node range_query(int s, int e) {
       vector<nodeptr> tmp = split_range(s, e);
       nodeptr 1 = tmp[0], m = tmp[1], r = tmp[2];
       node rt = *m;
        root = merge(merge(1, m), r);
    int getIndexByValue(const T &value) {
        nodeptr cur = mp[value];
        vector<nodeptr> path;
       path.push_back(cur);
        while (cur != root) {
            cur = cur->parent;
           path.push_back(cur);
        for (int i = path.size() - 1; i >= 0; i--)
           path[i]->push down();
       for (auto &it: path)
           it->push_up();
       cur = mp[value];
       int cnt = cur->child[L]->size + 1;
       while (cur != root) {
           nodeptr par = cur->parent;
           if (par->child[R] == cur)cnt += par->child[L]->size
                 + 1;
           cur = par;
       return cnt;
};
implicit_treap<long long, cartesian_tree> tp;
```

```
TreapOrderedMultiset.h
```

619894, 157 lines

```
enum DIR {
 L, R
};
template<typename T>
struct cartesian tree {
  static cartesian_tree<T> *sentinel;
  int priority = 0, frequency = 0, size = 0;
  cartesian tree *child[2];
  cartesian_tree() {
    key = T();
    priority = 0;
    child[L] = child[R] = this;
  cartesian_tree(const T &x, int y) :
      key(x), priority(y) {
    size = frequency = 1;
    child[L] = child[R] = sentinel;
 void push down() {
  void push up() {
    size = child[L]->size + child[R]->size + frequency;
};
template<typename T>
cartesian_tree<T> *cartesian_tree<T>::sentinel = new
    cartesian tree<T>();
template<typename T>
void split(cartesian_tree<T> *root, T key, cartesian_tree<T> *&
    cartesian_tree<T> *&r) {
  if (root == cartesian tree<T>::sentinel) {
    l = r = cartesian tree<T>::sentinel;
    return;
  root->push_down();
  if (root->key <= key) {</pre>
    split(root->child[R], key, root->child[R], r);
    1 = root;
    split(root->child[L], key, l, root->child[L]);
    r = root;
  root->push_up();
template<typename T>
cartesian_tree<T>* merge(cartesian_tree<T> *1, cartesian_tree<T</pre>
    > *r) {
 1->push_down();
  r->push down();
  if (1 == cartesian_tree<T>::sentinel || r == cartesian_tree<T</pre>
      >::sentinel)
    return (1 == cartesian_tree<T>::sentinel ? r : 1);
  if (l->priority > r->priority) {
    1->child[R] = merge(1->child[R], r);
    1->push_up();
    return 1:
  r->child[L] = merge(1, r->child[L]);
  r->push_up();
  return r;
```

```
template<typename T, template<typename > class cartesian tree>
class treap {
  typedef cartesian_tree<T> node;
  typedef node *nodeptr;
#define emptyNode node::sentinel
  nodeptr root;
  void insert(nodeptr &root, nodeptr it) {
    if (root == emptyNode) {
      root = it;
    } else if (it->priority > root->priority) {
      split(root, it->key, it->child[L], it->child[R]);
      root = it:
      insert(root->child[root->key < it->key], it);
    root->push_up();
  bool increment (nodeptr root, const T &key) {
    if (root == emptyNode)
     return 0;
    if (root->key == key) {
     root->frequency++;
     root->push_up();
     return root;
   bool rt = increment(root->child[root->key < key], key);</pre>
    root->push up();
    return rt;
  nodeptr find(nodeptr root, const T &key) {
    if (root == emptyNode || root->key == key)
     return root;
    return find(root->child[root->key < key], key);</pre>
  void erase (nodeptr &root, const T &key) {
    if (root == emptyNode)
     return;
    if (root->kev == kev) {
      if (--(root->frequency) == 0)
        root = merge(root->child[L], root->child[R]);
     erase(root->child[root->key < key], key);
    root->push_up();
  T kth(nodeptr root, int k) {
    if (root->child[L]->size >= k)
     return kth(root->child[L], k);
    k -= root->child[L]->size;
    if (k <= root->frequency)
     return root->key;
    return kth(root->child[R], k - root->frequency);
  int order_of_key(nodeptr root, const T &key) {
    if (root == emptyNode)
     return 0:
    if (key < root->key)
     return order_of_key(root->child[L], key);
    if (key == root->key)
     return root->child[L]->size;
    return root->child[L]->size + root->frequency
        + order_of_key(root->child[R], key);
public:
      root(emptyNode) {
  void insert(const T &x) {
```

```
if (increment(root, x)) //change it to find(x) to make it
         as \ a \ set
      return;
    insert(root, new node(x, rand()));
  void erase(const T &x) {
    erase(root, x);
  bool find(const T &x) {
    return (find(root, x) != emptyNode);
  int get_kth_number(int k) {
    assert(1 <= k && k <= size());
    return kth(root, k);
  int order_of_key(const T &x) {
    return order_of_key(root, x);
 int size() {
    return root->size;
};
int main() {
 treap<int, cartesian_tree> tp;
```

1.6 Line Container

LineContainer.h

a90417, 44 lines

```
struct Line {
 11 m, c;
 mutable 11 p; //p is intersection between cur and next
 bool operator<(const Line &o) const {</pre>
    //change to (m > 0.m, c < 0.c) to get min
    if (m != o.m)
     return m < o.m;
    return c > o.c;
 bool operator<(11 x) const {</pre>
    return p < x;
struct LineContainer: multiset<Line, less<>>> {
 // (for doubles, use inf = INFINITY, div(a,b) = a/b)
 static const ll inf = LLONG_MAX;
 ll div(ll a, ll b) { // floored division
    return a / b - ((a ^ b) < 0 && a % b);
 bool isect(iterator x, iterator y) {
   if (v == end())
     return x->p = inf, 0;
    if (x->m == y->m)
     x->p = inf;
      x->p = div(y->c - x->c, x->m - y->m);
    return x->p >= y->p;
 void add(ll m, ll c) {
   auto z = insert( \{ m, c, 0 \}), y = z++, x = y;
   while (isect(y, z))
     z = erase(z);
    if (x != begin() && isect(--x, y))
     isect(x, y = erase(y));
    while ((y = x) != begin() \&\& (--x)->p >= y->p)
      isect(x, erase(y));
```

```
11 query(11 x) {
    assert(!empty());
    auto 1 = *lower_bound(x);
    return 1.m * x + 1.c;
    }
};
```

1.7 MO

```
MO.h
```

d279dd, 40 lines

```
int sqrtN; //use a constent value
struct query {
  int 1, r, qIdx, block;
  query(int 1, int r, int qIdx) :
      l(l), r(r), qIdx(qIdx), block(l / sqrtN) {
  bool operator <(const query &o) const {</pre>
    if (block != o.block)
      return block < o.block;
    return (block % 2 == 0 ? r < o.r : r > o.r);
};
int curL, curR, ans;
vector<query> q;
void add(int index);
void remove(int index);
void solve(int 1, int r) {
  while (curL > 1)
    add(--curL);
  while (curR < r)
    add(++curR);
  while (curL < 1)
    remove(curL++):
  while (curR > r)
    remove (curR--);
vector<int> MO() {
  vector<int> rt(q.size());
  ans = curL = curR = 0;
  add(0); //v[0]
  sort(q.begin(), q.end());
  for (auto it : q) {
    solve(it.1, it.r);
    rt[it.qIdx] = ans;
  return rt;
```

$\overline{\text{Graph}}$ (2)

2.1 Graph Data structures

DSU.h

5b2539, 32 lines

```
struct DSU {
  vector<int> rank, parent, size;
  int forsets;
  DSU(int n) {
    size = rank = parent = vector<int>(n + 1, 1);
    forsets = n;
  for (int i = 0; i <= n; i++) {
      parent[i] = i;
    }
  }
  int find_set(int v) {</pre>
```

void link(int par, int node) {

return parent[v] = find set(parent[v]);

if (v == parent[v])

return v;

DSURollback LCA HLD

// this function save the current trees (merged) and

```
parent[node] = par;
    size[par] += size[node];
   if (rank[par] == rank[node])
     rank[par]++;
    forsets--:
  bool union_sets(int v, int u) {
   v = find_set(v), u = find_set(u);
    if (v != u) {
     if (rank[v] < rank[u])</pre>
       swap(v, u);
     link(v, u);
    return v != u;
};
DSURollback.h
                                                      ef3ba1, 62 lines
struct dsu_save {
    int v, rnkv, u, rnku;
    dsu save() {}
    dsu_save(int _v, int _rnkv, int _u, int _rnku)
        : v(_v), rnkv(_rnkv), u(_u), rnku(_rnku) {}
struct dsu with rollbacks {
    vector<int> p, rnk;
    int comps:
    stack<dsu_save> op;
    dsu with rollbacks() {}
    dsu_with_rollbacks(int n) {
       p.resize(n + 1);
       rnk.resize(n + 1);
        for (int i = 1; i <= n; i++) {</pre>
           p[i] = i;
            rnk[i] = 0;
        comps = n;
    int find_set(int v) {
        return (v == p[v]) ? v : find_set(p[v]);
   bool same_group(int v, int u) {
       v = find_set(v);
        u = find_set(u);
        if (v == u)
            return false;
        comps--;
        if (rnk[v] > rnk[u])
            swap(v, u);
        op.push(dsu_save(v, rnk[v], u, rnk[u]));
        p[v] = u;
        if (rnk[u] == rnk[v])
            rnk[u]++;
        return true;
    void snapshot() {
```

```
don't rollback them any more
        while (!op.empty())
            op.pop();
    void rollback() {
        // you can erase the while loop if you want to rollback
              just the last merge
        while (!op.empty()) {
            dsu_save x = op.top();
            op.pop();
            comps++;
            p[x.v] = x.v;
            rnk[x.v] = x.rnkv;
            p[x.u] = x.u;
            rnk[x.u] = x.rnku;
    }
};
LCA.h
                                                     15e<u>721, 61 lines</u>
class LCA {
  int n, logN, root = 1;
  vector<int> depth;
  vector<vector<int>> adj, lca;
  void dfs(int node, int parent) {
    lca[node][0] = parent;
    depth[node] = (~parent ? depth[parent] + 1 : 0);
    for (int k = 1; k <= logN; k++) {</pre>
      int up_parent = lca[node][k - 1];
      if (~up_parent)
        lca[node][k] = lca[up_parent][k - 1];
    for (int child : adi[node])
      if (child != parent)
        dfs(child, node);
public:
  LCA(const vector<vector<int>> &_adj, int root = 1) :
      root(root), adj(_adj) {
    adj = _adj;
    n = adj.size() - 1;
    logN = log2(n);
    lca = vector<vector<int>>(n + 1, vector<int>(logN + 1, -1))
    depth = vector<int>(n + 1);
    dfs(root, -1);
  int get_LCA(int x, int y) {
    if (depth[x] < depth[y])</pre>
      swap(x, y);
    for (int k = logN; k >= 0; k--)
      if (depth[x] - (1 \ll k) >= depth[y])
        x = lca[x][k];
    if (x == y)
      return x;
    for (int k = logN; k >= 0; k--) {
      if (lca[x][k] != lca[y][k]) {
        x = lca[x][k];
        y = lca[y][k];
    return lca[x][0];
  int get_distance(int u, int v) {
    return depth[u] + depth[v] - 2 * depth[get_LCA(u, v)];
  int kth_ancestor(int node, int dist) {
```

```
if (dist & (1 << i))
        node = lca[node][i];
    return node;
  edge get_path(int u, int LCA) {
    edge rt;
    for (int k = logN; k \ge 0; k--)
      if (depth[u] - (1 << k) >= depth[LCA]) {
        rt = merge(rt, lca[u][k]);
        u = lca[u][k].to;
    return rt;
};
HLD.h
// 1-based, if value in node, just update it after build chains
// don't forget to call build_chains after add edges.
class heavy_light_decomposition {
  int n, is value in edge;
  vector<int> parent, depth, heavy, root, pos_in_array,
       pos_to_node, size;
  const static int merge(int a, int b); //implement it
  struct array_ds { //implement it
    array_ds(int n) :
        n(n) {
    void update(int pos, int value);
    int query(int 1, int r);
  } seq;
  struct TREE {
    int cnt_edges = 1;
    vector<vector<int>> adj;
    //need for value in edges
    vector<vector<int>> edge idx;
    //edge_to need for undirected tree //end of edge in
         directed tree
    vector<int> edge_to, edge_cost;
    TREE(int n) :
        adj(n + 1), edge_idx(n + 1), edge_to(n + 1), edge_cost(
            n + 1) {
    void add_edge(int u, int v, int c) {
      adj[u].push_back(v);
      adj[v].push_back(u);
      edge_idx[u].push_back(cnt_edges);
      edge_idx[v].push_back(cnt_edges);
      edge_cost[cnt_edges] = c;
      cnt_edges++;
  } tree;
  int dfs_hld(int node) {
    int size = 1, max_sub_tree = 0;
    for (int i = 0; i < (int) tree.adj[node].size(); i++) {</pre>
      int ch = tree.adj[node][i], edge_idx = tree.edge_idx[node
      if (ch != parent[node]) {
        tree.edge_to[edge_idx] = ch;
        parent[ch] = node;
        depth[ch] = depth[node] + 1;
        int child_size = dfs_hld(ch);
        if (child_size > max_sub_tree)
          heavy[node] = ch, max_sub_tree = child_size;
        size += child_size;
```

for (int i = logN; i >= 0 && ~node; i--)

VirtualTree CentroidDecomposition

```
return size:
vector<tuple<int, int, bool>> get_path(int u, int v) { //l,r,
    must\_reverse?
 vector<pair<int, int>> tmp[2];
 bool idx = 1;
  while (root[u] != root[v]) {
   if (depth[root[u]] > depth[root[v]]) {
     swap(u, v);
      idx = !idx;
    //if value in edges ,you need value of root[v] also (
        connecter edge)
   tmp[idx].push_back( { pos_in_array[root[v]], pos_in_array
        [V] });
   v = parent[root[v]];
  if (depth[u] > depth[v]) {
   swap(u, v);
   idx = !idx;
  if (!is_value_in_edge || u != v)
   tmp[idx].push_back( { pos_in_array[u] + is_value_in_edge,
        pos_in_array[v] });
  reverse(all(tmp[1]));
  vector<tuple<int, int, bool>> rt;
  for (int i = 0; i < 2; i++)
   for (auto &it : tmp[i])
      rt.emplace_back(it.first, it.second, i == 0);
  return rt; //u is LCA
heavy_light_decomposition(int n, bool is_value_in_edge) :
   n(n), is_value_in_edge(is_value_in_edge), seg(n + 1),
        tree(n + 1) {
  heavy = vector<int>(n + 1, -1);
  parent = depth = root = pos_in_array = pos_to_node = size =
       vector<int>(
     n + 1);
void add edge(int u, int v, int c = 0) {
 tree.add_edge(u, v, c);
void build_chains(int src = 1) {
 parent[src] = -1;
  dfs hld(src);
  for (int chain_root = 1, pos = 1; chain_root <= n;</pre>
      chain root++) {
    if (parent[chain root] == -1
        || heavy[parent[chain_root]] != chain root)
      for (int j = chain_root; j != -1; j = heavy[j]) {
        root[j] = chain_root;
        pos_in_array[j] = pos++;
        pos_to_node[pos_in_array[j]] = j;
  if (is_value_in_edge)
    for (int i = 1; i < n; i++)</pre>
     update_edge(i, tree.edge_cost[i]);
void update_node(int node, int value) {
  seg.update(pos_in_array[node], value);
void update_edge(int edge_idx, int value) {
 update_node(tree.edge_to[edge_idx], value);
void update_path(int u, int v, ll c) {
 vector<tuple<int, int, bool>> intervals = get_path(u, v);
```

```
for (auto &it : intervals)
     seg.update(get<0>(it), get<1>(it), c);
 node query_in_path(int u, int v) {
    vector<tuple<int, int, bool>> intervals = get_path(u, v);
    //initial\ value, check\ if\ handling\ u == v
    node query res = 0;
    for (auto &it : intervals) {
     int 1, r;
     bool rev:
     tie(1, r, rev) = it;
      node cur = seg.query(1, r);
      if (rev)
       cur.reverse();
      query_res = node(query_res, cur);
    return query_res;
};
VirtualTree.h
                                                     702679, 62 lines
int subsize[N], depth[N], dfs_num[N], id[N], have[N], timer, n;
vector<vector<int>> adj;
void calc(int node) {
    dfs_num[node] = ++timer;
    subsize[node] = 1;
    for (auto ch: adj[node]) {
        adi[ch].erase(find(all(adj[ch]), node));
        depth[ch] = depth[node] + 1;
        subsize[node] += subsize[ch];
bool parent (int node, int par) {
    return dfs_num[par] <= dfs_num[node] && dfs_num[node] <</pre>
         dfs num[par] + subsize[par];
struct virtual tree {
    vector<int> nodes:
    vector<vector<int>> adj;
    virtual_tree(const vector<int> &v) : nodes(v) {
        for (auto &it: nodes) have[it] = true;
       sort(all(nodes), [&](int a, int b) {
            return dfs_num[a] < dfs_num[b];</pre>
       });
       int tmp = nodes.size();
       for (int j = 0; j + 1 < tmp; j++) {
            int lca = tree.get_LCA(nodes[j], nodes[j + 1]);
            nodes.push_back(lca);
       nodes.push_back(1);
       sort (all (nodes));
       nodes.erase(unique(all(nodes)), nodes.end());
       sort(all(nodes), [&](int a, int b) {
            return dfs_num[a] < dfs_num[b];</pre>
       });
        int cnt = 0:
        for (auto &it: nodes)id[it] = cnt++;
        stack<int> stk;
```

```
CentroidDecomposition.h
                                                    689adb, 84 lines
class centroid_decomposition {
 vector<bool> centroidMarked;
  vector<int> size;
 void dfsSize(int node, int par) {
    size[node] = 1;
    for (int ch : adj[node])
     if (ch != par && !centroidMarked[ch]) {
       dfsSize(ch, node);
        size[node] += size[ch];
 int getCenter(int node, int par, int size_of_tree) {
    for (int ch : adj[node]) {
     if (ch == par || centroidMarked[ch])
        continue;
      if (size[ch] * 2 > size of tree)
       return getCenter(ch, node, size_of_tree);
    return node:
 int getCentroid(int src) {
    dfsSize(src, -1);
    int centroid = getCenter(src, -1, size[src]);
    centroidMarked[centroid] = true;
   return centroid;
 int decomposeTree(int root) {
    root = getCentroid(root);
    solve(root);
    for (int ch : adj[root]) {
     if (centroidMarked[ch])
       continue;
      int centroid_of_subtree = decomposeTree(ch);
      //note: root and centroid_of_subtree probably not have a
           direct edge in adj
      centroidTree[root].push_back(centroid_of_subtree);
      centroidParent[centroid_of_subtree] = root;
    return root;
 void calc(int node, int par) {
    //TO-DO
    for (int ch : adj[node])
     if (ch != par && !centroidMarked[ch])
       calc(ch, node);
 void add(int node, int par) {
    //TO-DO
    for (int ch : adj[node])
      if (ch != par && !centroidMarked[ch])
```

int test id;

bool canMatch(int i) {

if (vis[i] == test id)

```
add(ch, node);
  void remove(int node, int par) {
    //TO-DO
    for (int ch : adj[node])
     if (ch != par && !centroidMarked[ch])
        remove(ch, node);
  void solve(int root) {
    //add root
    for (int ch : adj[root])
     if (!centroidMarked[ch]) {
        calc(ch, root);
        add(ch, root);
    //TO-DO
    //remove root
    for (int ch : adj[root])
     if (!centroidMarked[ch])
        remove(ch, root);
public:
  int n, root;
  vector<vector<int>> adj, centroidTree;
  vector<int> centroidParent;
  centroid_decomposition(vector<vector<int>> &adj) :
     adj(adj) {
    n = (int) adj.size() - 1;
    size = vector<int>(n + 1);
    centroidTree = vector<vector<int>> (n + 1);
    centroidParent = vector<int>(n + 1, -1);
    centroidMarked = vector<bool>(n + 1);
    root = decomposeTree(1);
};
Sack.h
                                                     ebd0f0, 20 lines
void add(int node, int par, int big);
void remove(int node, int par, int big);
void dfs(int node, int par, bool keep) {
    int mx = -1, big = -1;
    for (int ch: adj[node]) {
        if (ch != par && subsize[ch] > mx)
            mx = subsize[ch], big = ch;
    for (int ch: adj[node]) {
        if (ch != par && ch != big)
            dfs(ch, node, false);
    if (big != -1) dfs(big, node, true);
    add(node, par, big);
    // SOLVE
    if (!keep)
        remove (node, par, -1);
      Flow
```

2.2.1 Bipartite Matching

MaximumBipartiteMatching.h

```
// O(VE) in worst case, but in practice it runs in O(E) or O(
     sqrt(V)E)
vector<vector<int>> adj;
vector<int> rowAssign, colAssign, vis; //make vis array instance
```

```
return false;
  vis[i] = test_id;
 for (int j : adj[i]) {
    if (colAssign[j] == -1) {
      colAssign[j] = i;
      rowAssign[i] = j;
      return true;
 for (int j : adj[i]) {
   if (canMatch(colAssign[j])) {
     colAssign[j] = i;
      rowAssign[i] = j;
      return true;
 return false;
// O(rows * edges) //number of operation could by strictly less
int maximum_bipartite_matching(int rows, int cols) {
 int maxFlow = 0;
 rowAssign = vector<int>(rows, -1);
  colAssign = vector<int>(cols, -1);
  vis = vector<int>(rows);
  for (int i = 0; i < rows; i++) {</pre>
   test id++;
    if (canMatch(i))
     maxFlow++;
  vector<pair<int, int>> matches;
  for (int j = 0; j < cols; j++)
    if (~colAssign[j])
      matches.push_back( { colAssign[j], j });
  return maxFlow;
HopcroftKarp.h
//Hopcroft-Karp algorithm for maximum bipartite matching
//O(sqrt(V) * E)
struct Hopcroft_Karp {//1-based
#define NIL 0
#define INF INT MAX
 int n, m;
 vector<vector<int>> adj;
 vector<int> rowAssign, colAssign, dist;
 bool bfs() {
    queue<int> q;
    dist = vector<int>(adj.size(), INF);
    for (int i = 1; i <= n; i++)</pre>
      if (rowAssign[i] == NIL) {
        dist[i] = 0;
        q.push(i);
    while (!q.empty()) {
      int cur = q.front();
      q.pop();
      if (dist[cur] >= dist[NIL])break;
      for (auto& nxt : adj[cur]) {
        if (dist[colAssign[nxt]] == INF) {
          dist[colAssign[nxt]] = dist[cur] + 1;
          q.push(colAssign[nxt]);
```

```
return dist[NIL] != INF;
 bool dfs(int i) {
    if (i == NIL)
      return true;
    for (int j : adj[i]) {
      if (dist[colAssign[j]] == dist[i] + 1 && dfs(colAssign[j])
        colAssign[j] = i;
        rowAssign[i] = j;
        return true;
    dist[i] = INF;
    return false;
  Hopcroft_Karp(int n, int m)
    :n(n), m(m), adj(n + 1), rowAssign(n + 1), colAssign(m + 1)
  void addEdge(int u, int v) {
    adj[u].push_back(v);
  int maximum_bipartite_matching() {
    int rt = 0;
    while (bfs()) {
      for (int i = 1; i <= n; i++)</pre>
        if (rowAssign[i] == NIL && dfs(i))
    return rt:
};
```

2.2.2 Max Flow

Dinic.h

```
4fa46d, 70 lines
struct Dinic {
  struct flowEdge {
   int from, to;
    11 \text{ cap, flow} = 0;
    flowEdge(int from, int to, ll cap) :
        from (from), to (to), cap(cap) {
 vector<flowEdge> edges;
 int n, m = 0, source, sink;
 vector<vector<int>> adj;
 vector<int> level, ptr;
 Dinic(int n, int source, int sink) :
      n(n), source(source), sink(sink), adj(n), level(n), ptr(n
 void addEdge(int u, int v, ll cap) {
    edges.emplace_back(u, v, cap);
   edges.emplace_back(v, u, 0);
   adj[u].push_back(m);
   adj[v].push_back(m + 1);
   m += 2;
 bool bfs() {
    queue<int> q;
    level = vector<int>(n, -1);
    level[source] = 0;
    q.push (source);
    while (!q.empty()) {
     int cur = q.front();
      q.pop();
```

EdmondsKarp MCMF Bellmanford

```
for (auto &id : adj[cur]) {
        if (edges[id].cap - edges[id].flow <= 0)</pre>
          continue;
        int nxt = edges[id].to;
        if (level[nxt] != -1)
          continue;
        level[nxt] = level[cur] + 1;
        q.push(nxt);
    return level[sink] != -1;
  11 dfs(int node, ll cur_flow) {
    if (cur flow == 0 || node == sink)
      return cur_flow;
    for (int &cid = ptr[node]; cid < adj[node].size(); cid++) {</pre>
      int id = adj[node][cid];
      int nxt = edges[id].to;
      if (level[node] + 1 != level[nxt]
          || edges[id].cap - edges[id].flow <= 0)</pre>
      11 tmp = dfs(nxt, min(cur_flow, edges[id].cap - edges[id
           1.flow));
      if (tmp == 0)
        continue;
      edges[id].flow += tmp;
      edges[id ^ 1].flow -= tmp;
      return tmp;
    return 0;
  ll flow() {
    11 \text{ max flow} = 0;
    while (bfs()) {
      fill(ptr.begin(), ptr.end(), 0);
      while (11 pushed = dfs(source, INF))
        max_flow += pushed;
    return max flow;
};
```

EdmondsKarp.h

```
fd0a7c, 39 lines
//O(V*E*E)
int n;
int capacity[101][101];
int getPath(int src, int dest, vector<int> &parent) {
  parent = vector<int>(n + 1, -1);
  queue<pair<int, int>> q;
  q.push( { src, INF });
  while (q.size()) {
   int cur = q.front().first, flow = q.front().second;
    q.pop();
    if (cur == dest)
     return flow;
    for (int i = 1; i <= n; i++)</pre>
     if (parent[i] == -1 && capacity[cur][i]) {
       parent[i] = cur;
        q.push( { i, min(flow, capacity[cur][i]) });
       if (i == dest)
          return q.back().second;
  return 0;
int Edmonds_Karp(int source, int sink) {
  int max flow = 0;
```

```
int new flow = 0;
vector<int> parent(n + 1, -1);
while (new flow = getPath(source, sink, parent)) {
  max_flow += new_flow;
  int cur = sink;
  while (cur != source) {
    int prev = parent[cur];
    capacity[prev][cur] -= new_flow;
    capacity[cur][prev] += new_flow;
    cur = prev;
  };
return max_flow;
```

2.2.3 Min Cost Max Flow

MCMF.h

da97ce, 89 lines

```
struct MCMF { //0-based
 struct edge {
    int from, to, cost, cap, flow, backEdge;
     from = to = cost = cap = flow = backEdge = 0;
    edge (int from, int to, int cost, int cap, int flow, int
        from (from), to (to), cost (cost), cap (cap), flow (flow),
            backEdge (
            backEdge) {
   bool operator <(const edge &other) const {</pre>
     return cost < other.cost;</pre>
 };
 int n, src, dest;
 vector<vector<edge>> adj;
 const int 00 = 1e9;
 MCMF(int n, int src, int dest) :
     n(n), src(src), dest(dest), adj(n) {
 void addEdge(int u, int v, int cost, int cap) {
   edge e1 = edge(u, v, cost, cap, 0, adj[v].size());
   edge e2 = edge(v, u, -cost, 0, 0, adj[u].size());
   adj[u].push_back(e1);
    adj[v].push_back(e2);
 pair<int, int> minCostMaxFlow() {
   int maxFlow = 0, cost = 0;
   while (true) {
      vector<pair<int, int>> path = spfa();
     if (path.empty())
       break:
      int new flow = 00;
      for (auto &it : path) {
       edge &e = adj[it.first][it.second];
       new_flow = min(new_flow, e.cap - e.flow);
      for (auto &it : path) {
       edge &e = adj[it.first][it.second];
       e.flow += new_flow;
       cost += new_flow * e.cost;
       adj[e.to][e.backEdge].flow -= new_flow;
     maxFlow += new_flow;
    return {maxFlow,cost};
 enum visit {
   finished, in_queue, not_visited
```

```
vector<pair<int, int>> spfa() {
    vector<int> dis(n, 00), prev(n, -1), from_edge(n), state(n,
        not visited);
    deque<int> q;
    dis[src] = 0;
    q.push_back(src);
    while (!q.empty()) {
      int u = q.front();
      q.pop_front();
      state[u] = finished;
      for (int i = 0; i < adj[u].size(); i++) {</pre>
        edge e = adj[u][i];
        if (e.flow >= e.cap || dis[e.to] <= dis[u] + e.cost)</pre>
          continue;
        dis[e.to] = dis[u] + e.cost;
        prev[e.to] = u;
        from_edge[e.to] = i;
        if (state[e.to] == in_queue)
          continue;
        if (state[e.to] == finished
            || (!q.empty() && dis[q.front()] > dis[e.to]))
          g.push front(e.to);
        el se
          q.push_back(e.to);
        state[e.to] = in_queue;
    if (dis[dest] == 00)
      return {};
    vector<pair<int, int>> path;
    int cur = dest;
    while (cur != src) {
      path.push_back( { prev[cur], from_edge[cur] });
      cur = prev[cur];
    reverse(path.begin(), path.end());
    return path;
};
```

2.3 Shortest Path

Bellmanford.h

b759d9, 70 lines

```
#define on 0x3f3f3f3f3fJJ
struct edge {
 int from, to, weight;
 edge() {
    from = to = weight = 0;
  edge (int from, int to, int weight) :
      from(from), to(to), weight(weight) {
 bool operator <(const edge &other) const {</pre>
    return weight > other.weight;
};
vector<edge> edgeList;
//O(V*E)
void bellmanford(int n, int src, int dest = -1) {
 vector<int> dis(n + 1, oo), prev(n + 1, -1);
 dis[src] = 0;
 bool negativeCycle = false;
 int last = -1, tmp = n;
  while (tmp--) {
   last = -1;
    for (edge e : edgeList)
      if (dis[e.to] > dis[e.from] + e.weight) {
```

SPFA EdgeClassification SCC ArticulationPointsAndBridges

```
dis[e.to] = dis[e.from] + e.weight;
        prev[e.to] = e.from;
        last = e.to;
    if (last == -1)
     break;
    if (tmp == 0)
     negativeCycle = true;
  if (last != −1) {
    for (int i = 0; i < n; i++)</pre>
     last = prev[last];
    vector<int> cycle;
    for (int cur = last; cur != last || cycle.size() > 1; cur =
         prev[cur])
     cycle.push_back(cur);
    reverse(cycle.begin(), cycle.end());
  vector<int> path;
  while (dest !=-1) {
   path.push_back(dest);
   dest = prev[dest];
  reverse(path.begin(), path.end());
void difference_constraints() {
 int m;
  cin >> m;
  int cnt = 1;
  while (m--) {
    string x1, x2;
   int w; // x1 - x2 <= w
    cin >> x1 >> x2 >> w;
    map<string, int> id;
    if (id.find(x1) == id.end())
     id[x1] = cnt++;
    if (id.find(x2) == id.end())
     id[x2] = cnt++;
    edgeList.emplace_back(id[x2], id[x1], w);
  for (int i = 1; i < cnt; i++)</pre>
    edgeList.emplace_back(cnt, i, 0);
  bellmanford(cnt, cnt);
SPFA.h
                                                     b6cec5, 43 lines
struct edge {
  int from, to;
 11 weight:
  edge() {
   from = to = weight = 0;
  edge (int from, int to, ll weight) :
      from(from), to(to), weight(weight) {
  bool operator <(const edge &other) const {</pre>
    return weight < other.weight;
};
vector<vector<edge>> adj;
void spfa(int src) {
  enum visit {
   finished, in_queue, not_visited
  int n = adj.size();
```

```
vector<int> dis(n, INF), prev(n, -1), state(n, not_visited);
  dis[src] = 0:
 deque<int> q;
 q.push_back(src);
 while (!q.empty()) {
   int u = q.front();
   q.pop_front();
   state[u] = finished;
    for (auto &e : adj[u]) {
     if (dis[e.to] > dis[e.from] + e.cost) {
        dis[e.to] = dis[e.from] + e.cost;
       prev[e.to] = e.from;
       if (state[e.to] == not_visited) {
          q.push_back(e.to);
        } else if (state[e.to] == finished) {
         q.push_front(e.to);
        state[e.to] = in_queue;
 }
2.4 Tarjan
```

EdgeClassification.h

543a6a, 20 lines

int v = -1;

```
vector<vector<int>> adj;
vector<int> start, finish;
int timer;
void dfsEdgeClassification(int node) {
 start[node] = timer++;
 for (int child : adj[node]) {
   if (start[child] == -1)
     dfsEdgeClassification(child);
     if (finish[child] == -1)
      ; // Back Edge
     else if (start[node] < start[child])</pre>
      ; // Forward Edge
     else
       ; // Cross Edge
 finish[node] = timer++;
```

```
SCC.h
                                                     26b1e6, 48 lines
vector<vector<int>> adj, scc;
vector<set<int>> dag;
vector<int> dfs_num, dfs_low, compId;
vector<bool> inStack;
stack<int> stk:
int timer;
void dfs(int node) {
 dfs_num[node] = dfs_low[node] = ++timer;
 stk.push(node);
 inStack[node] = 1;
  for (int child : adj[node])
   if (!dfs_num[child]) {
      dfs(child);
      dfs_low[node] = min(dfs_low[node], dfs_low[child]);
    } else if (inStack[child])
      dfs_low[node] = min(dfs_low[node], dfs_num[child]);
  //can\ be\ dfs\_low[node] = min(dfs\_low[node],\ dfs\_low[child]);
  if (dfs_low[node] == dfs_num[node]) {
    scc.push_back(vector<int>());
```

```
while (v != node) {
     v = stk.top();
      stk.pop();
      inStack[v] = 0;
      scc.back().push_back(v);
      compId[v] = scc.size() - 1;
void SCC() {
 timer = 0:
  dfs_num = dfs_low = compId = vector<int>(adj.size());
 inStack = vector<bool>(adj.size());
  scc = vector<vector<int>>();
  for (int i = 1; i < adj.size(); i++)</pre>
   if (!dfs_num[i])
      dfs(i);
void DAG() {
 dag = vector<set<int>>(scc.size());
 for (int i = 1; i < adj.size(); i++)</pre>
    for (int j : adj[i])
      if (compId[i] != compId[j])
        dag[compId[i]].insert(compId[j]);
ArticulationPointsAndBridges.h
                                                     1e3317, 50 lines
vector<vector<int>> adj;
vector<int> dfs num, dfs low;
vector<bool> articulation_point;
vector<pair<int, int>> bridge;
stack<pair<int, int>> edges;
vector<vector<pair<int, int>>> BCC; //biconnected components
int timer, cntChild;
// O(n + m)
void dfs(int node, int par) {
 dfs num[node] = dfs low[node] = ++timer;
 for (int child : adj[node]) {
    if (par != child && dfs_num[child] < dfs_num[node])</pre>
      edges.push( { node, child });
    if (!dfs_num[child]) {
      if (par == -1)
        cntChild++;
      dfs(child, node);
      if (dfs_low[child] >= dfs_num[node]) {
        articulation_point[node] = 1;
        //get biconnected component
        BCC.push_back(vector<pair<int, int>>());
        pair<int, int> edge;
          edge = edges.top();
          BCC.back().push_back(edge);
          edges.pop();
        } while (edge.first != node || edge.second != child);
      if (dfs_low[child] > dfs_num[node]) //can be (dfs_low[
           child = dfs_num(child)
        bridge.push_back( { node, child });
```

dfs_low[node] = min(dfs_low[node], dfs_low[child]);

dfs_low[node] = min(dfs_low[node], dfs_num[child]);

else if (child != par)

```
Ala7lam_team (Assiut University)
void articulation_points_and_bridges() {
  timer = 0;
  dfs_num = dfs_low = vector<int>(adj.size());
  articulation_point = vector < bool > (adj.size());
  bridge = vector<pair<int, int>>();
  for (int i = 1; i < adj.size(); i++)</pre>
    if (!dfs_num[i]) {
      cntChild = 0;
      dfs(i, -1);
      articulation_point[i] = cntChild > 1;
2SAT.h
"SCC.h"
                                                       fc4f77, 30 lines
int n:
int Not(int x) {
  return (x > n ? x - n : x + n);
void addEdge(int a, int b) {
  adi[Not(a)].push back(b);
  adj[Not(b)].push_back(a);
void add xor edge(int a, int b) {
  addEdge(Not(a), Not(b));
  addEdge(a, b);
bool 2SAT(vector<int> &value) {
  SCC();
  for (int i = 1; i <= n; i++)</pre>
    if (compId[i] == compId[Not(i)])
```

ShortestCvcle.h

return true;

return false;

if (assign[i] == -1) {

for (int i = 1; i <= n; i++)</pre>

value[i] = assign[compId[i]];

assign[i] = true;

vector<int> assign(scc.size(), -1);

for (int i = 0; i < scc.size(); i++)</pre>

assign[compId[Not(scc[i].back())]] = false;

3a7390, 30 lines

```
//Shortest cycle starting from node O(n^2)
int get_shortest_cycle(int start) {
    vector<int> dist(n + 1, INF), parent(n + 1), group(n + 1);
    vector<bool> vis(n + 1);
    dist[start] = 0;
    group[start] = start;
    while (true) {
        int cur = 0;
        for (int i = 1; i <= n; i++)</pre>
            if (!vis[i] && dist[i] < dist[cur])cur = i;</pre>
       if (dist[cur] == (int) INF)break;
       vis[cur] = true;
        for (int ch = 1; ch <= n; ch++) {
            if (dist[ch] > dist[cur] + adj[cur][ch]) {
                dist[ch] = dist[cur] + adj[cur][ch];
                parent[ch] = cur;
                group[ch] = (cur == start ? ch : group[cur]);
```

```
int mn = INF;
for (int i = 1; i <= n; i++)</pre>
    for (int j = i + 1; j <= n; j++) {</pre>
        if (parent[i] == j || parent[j] == i)continue;
        if (group[i] == group[j]) continue;
        if (dist[i] == (int) INF || dist[j] == (int) INF)
        mn = min(mn, dist[i] + dist[j] + adj[i][j]);
return mn;
```

Strings (3)

3.1 Rolling Hash

```
Hashing.h
                                                      4dc92d, 23 lines
struct hashing {
   int MOD, BASE;
    vector<int> Hash, modInv;
   hashing(string s, int MOD, int BASE, char first_char = 'a')
            MOD(MOD), BASE(BASE), Hash(sz(s) + 1), modInv(sz(s)
                  + 1) {
        11 BASE_INV = power(BASE, MOD - 2, MOD);
       modInv[0] = 1;
       11 \text{ base} = 1;
        for (int i = 1; i <= sz(s); i++) {</pre>
            Hash[i] = (Hash[i-1] + (s[i-1] - first\_char +
                 1) * base) % MOD;
            modInv[i] = (modInv[i - 1] * BASE_INV) % MOD;
            base = (base * BASE) % MOD;
    int getHash(int 1, int r) { //1-based
        return (1LL * (Hash[r] - Hash[l - 1] + MOD) % MOD *
             modInv[1 - 1]) % MOD;
};
/MOD = 1e9 + 9 , BASE = 31
/MOD = 2000000011 ,BASE = 53 \rightarrow careful of overflow
/MOD = 998634293,BASE = 953
/MOD = 986464091, BASE = 1013
```

String matching

KMP.h

```
98770d, 92 lines
struct KMP {
 string pattern;
 vector<int> longestPrefix;
 KMP(const string& str) :pattern(str) {
   failure_function();
 int fail(int k, char nxt) {
   while (k > 0 \&\& pattern[k] != nxt)
     k = longestPrefix[k - 1];
   if (nxt == pattern[k])
     k++;
   return k;
 void failure function() {
   int n = pattern.size();
   longestPrefix = vector<int>(n);
    for (int i = 1, k = 0; i < n; i++)
     longestPrefix[i] = k = fail(k, pattern[i]);
```

```
void match(const string& str) {
    int n = str.size();
    int m = pattern.size();
    for (int i = 0, k = 0; i < n; i++) {
      k = fail(k, str[i]);
      if (k == m) {
        cout << i - m + 1 << endl; //0-based
        k = longestPrefix[k - 1]; // if you want next match
};
vector<br/>vector<br/>suffix_pal(string s) { //[i..n-1] pal?
 string r = s;
  reverse(all(r));
  vector<bool> v(s.size());
  v[0] = (s == r);
  string pattern = r + "#" + s;
  int n = pattern.size();
  vector<int> longestPrefix(n);
  int k = 0;
  for (int i = 1; i < n; i++) {</pre>
    while (k > 0 && pattern[k] != pattern[i])
      k = longestPrefix[k - 1];
    if (pattern[i] == pattern[k])
    longestPrefix[i] = k;
  while (k > 0) {
    v[s.size() - k] = true;
    k = longestPrefix[k - 1];
  return v;
vector<br/>bool> prefix_pal(string s) { // [0..i] pal?
  string r = s;
  reverse(all(r));
  vector<bool> v(s.size());
 v.back() = (s == r);
  string pattern = s + "#" + r;
  int n = pattern.size();
  vector<int> longestPrefix(n);
  int k = 0:
  for (int i = 1; i < n; i++) {</pre>
    while (k > 0 && pattern[k] != pattern[i])
      k = longestPrefix[k - 1];
    if (pattern[i] == pattern[k])
      k++;
    longestPrefix[i] = k;
  while (k > 0) {
   v[k-1] = true;
    k = longestPrefix[k - 1];
 return v;
//frq[i] = number of occur s[0..i] in s
vector<int> build_fre_prefix(const string& s) {
 KMP kmp(s);
 kmp.failure_function();
 vector<int> f = kmp.longestPrefix;
 int n = sz(s);
 vector<int> frq(n);
  for (int i = n - 1; i >= 0; i--)
    if (f[i])
      frq[f[i] - 1] += frq[i] + 1;
```

for (auto& it : frq)it++;

return frq;

AhoCorasick ZAlgo SuffixArray

```
AhoCorasick.h
                                                     d54f05, 82 lines
#define all(v) v.begin(), v.end()
struct aho corasick {
  struct trie node {
   vector<int> pIdxs; //probably take memory limit
    map<char, int> next;
   int fail;
   trie node() :
        fail(0) {
   bool have_next(char ch) {
     return next.find(ch) != next.end();
    int& operator[](char ch) {
     return next[ch];
  };
  vector<trie_node> t;
  vector<string> patterns;
  vector<int> end of pattern;
  vector<vector<int>> adj;
  int insert(const string &s, int patternIdx) {
    int root = 0;
    for (const char &ch : s) {
     if (!t[root].have next(ch)) {
       t.push_back(trie_node());
       t[rootl[ch] = t.size() - 1;
     root = t[root][ch];
    t[root].pIdxs.push_back(patternIdx);
    return root:
  int next_state(int cur, char ch) {
    while (cur > 0 && !t[cur].have_next(ch))
     cur = t[cur].fail;
    if (t[cur].have_next(ch))
     return t[cur][ch];
   return 0;
  void buildAhoTree() {
    queue<int> q;
    for (auto &child : t[0].next)
     q.push (child.second);
    while (!q.empty()) {
     int cur = q.front();
      q.pop();
     for (auto &child : t[cur].next) {
       int k = next_state(t[cur].fail, child.first);
       t[child.second].fail = k;
        vector<int> &idxs = t[child.second].pIdxs;
        //dp[child.second] = max(dp[child.second], dp[k]);
       idxs.insert(idxs.end(), all(t[k].pIdxs));
        q.push (child.second);
  void buildFailureTree() {
    adj = vector<vector<int>>(t.size());
   for (int i = 1; i < t.size(); i++)</pre>
      adj[t[i].fail].push_back(i);
  aho_corasick(const vector<string> &_patterns) {
```

```
t.push_back(trie_node());
    patterns = _patterns;
    end of pattern = vector<int>(patterns.size());
    for (int i = 0; i < patterns.size(); i++)</pre>
      end_of_pattern[i] = insert(patterns[i], i);
    buildAhoTree();
    buildFailureTree();
  vector<vector<int>> match(const string &str) {
    int k = 0:
    vector<vector<int>>> rt(patterns.size());
    for (int i = 0; i < str.size(); i++) {</pre>
     k = next_state(k, str[i]);
      for (auto &it : t[k].pIdxs)
        rt[it].push_back(i);
    return rt;
};
ZAlgo.h
                                                     90c964, 18 lines
// z[i] equal the length of the longest substring starting from
      s[i] which is also a prefix of s
vector<int> z_algo(string s) {
  int n = s.size();
  vector<int> z(n);
  z[0] = n;
  for (int i = 1, L = 1, R = 1; i < n; i++) {
    int k = i - L;
    if (z[k] + i >= R) {
     L = i;
      R = max(R, i);
      while (R < n \&\& s[R - L] == s[R])
       R++;
      z[i] = R - L;
    } else
      z[i] = z[k];
  return z:
3.3 Suffix structures
SuffixArrav.h
                                                     76f545, 220 lines
#define all(v) v.begin(), v.end()
class suffix array {
    const static int alpha = 128;
    int getOrder(int a) const {
        return (a < (int) order.size() ? order[a] : 0);</pre>
    void radix_sort(int k) {
        vector<int> frq(n), tmp(n);
        for (auto &it : suf)
            frq[getOrder(it + k)]++;
        for (int i = 1; i < n; i++)</pre>
            frq[i] += frq[i - 1];
        for (int i = n - 1; i >= 0; i--)
            tmp[--frq[getOrder(suf[i] + k)]] = suf[i];
        suf = tmp;
public:
    int n;
    vector<int> suf, lcp, order; // order store position of
         suffix i in suf array
    suffix_array(const string &s) :
            n(s.size() + 1), s(s) {
        suf = order = lcp = vector<int>(n);
```

```
vector<int> bucket_idx(n), newOrder(n), newsuff(n);
    vector<int> prev(n), head(alpha, -1);
    for (int i = 0; i < n; i++) {</pre>
        prev[i] = head[s[i]];
        head[s[i]] = i;
   int buc = -1, idx = 0;
    for (int i = 0; i < alpha; i++) {</pre>
        if (head[i] == -1)
            continue;
        bucket_idx[++buc] = idx;
        for (int j = head[i]; ~j; j = prev[j])
            suf[idx++] = j, order[j] = buc;
   int len = 1;
   do {
        auto cmp = [&](int a, int b) {
            if (order[a] != order[b])
                return order[a] < order[b];</pre>
            return getOrder(a + len) < getOrder(b + len);</pre>
        for (int i = 0; i < n; i++) {</pre>
            int j = suf[i] - len;
            if (j < 0)
                continue;
            newsuff[bucket_idx[order[j]]++] = j;
        for (int i = 1; i < n; i++) {</pre>
            suf[i] = newsuff[i];
            bool cmpres = cmp(suf[i - 1], suf[i]);
            newOrder[suf[i]] = newOrder[suf[i - 1]] +
                 cmpres;
            if (cmpres)
                bucket_idx[newOrder[suf[i]]] = i;
        order = newOrder;
        len <<= 1;
    } while (order[suf[n - 1]] != n - 1);
   buildLCP();
 * longest common prefix
 * lcp[i] = lcp(suf[i], suf[i-1])
void buildLCP() {
   lcp = vector<int>(n);
    int k = 0:
    for (int i = 0; i < n - 1; i++) {
        int pos = order[i];
        int j = suf[pos - 1];
        while (s[i + k] == s[j + k])
           k++;
        lcp[pos] = k;
        if (k)
            k--;
int LCP_by_order(int a, int b) {
   if (a > b)
        swap(a, b);
   int mn = n - suf[a] - 1;
    for (int k = a + 1; k \le b; k++)
        mn = min(mn, lcp[k]);
   return mn;
//LCP(i,j): longest common prefix between suffix i and
     suffix i
int LCP(int i, int j) {
```

```
//return LCP_by_order(order[i],order[j]);
        if (order[j] < order[i])</pre>
            swap(i, j);
        int mn = n - i - 1;
        for (int k = order[i] + 1; k <= order[j]; k++)</pre>
            mn = min(mn, lcp[k]);
        return mn;
    /\!/compare \ s[a.first..a.second] \ with \ s[b.first..b.second]
    //-1:a < b , 0:a = b , 1:a > b
    int compare_substrings(pair<int, int> a, pair<int, int> b)
        int lcp = min(
                 { LCP(a.first, b.first), a.second - a.first +
                      1, b.second
        a.first += lcp;
        b.first += lcp;
        if (a.first <= a.second) {</pre>
            if (b.first <= b.second) {</pre>
                 if (s[a.first] == s[b.first])
                     return 0;
                 return (s[a.first] < s[b.first] ? -1 : 1);</pre>
             return 1;
        return (b.first <= b.second ? -1 : 0);
    pair<int, int> find string(const string &x) {
        int st = 0, ed = n;
        for (int i = 0; i < sz(x) && st < ed; i++) {
            auto cmp = [&](int a, int b) {
                 if (a == -1)
                     return x[i] < s[b + i];</pre>
                 return s[a + i] < x[i];</pre>
            st = lower_bound(suf.begin() + st, suf.begin() + ed
                 , -1, cmp)
                  - suf.begin();
             ed = upper_bound(suf.begin() + st, suf.begin() + ed
                  , -1, cmp)
                  - suf.begin();
        return {st,ed-1};
11 number_of_different_substrings(string s) {
    int n = s.size();
    suffix array sa(s);
    11 \text{ cnt} = 0;
    for (int i = 0; i <= n; i++)</pre>
        cnt += n - sa.suf[i] - sa.lcp[i];
    return cnt;
```

```
string longest common substring (const string &s1, const string
    &s2) {
    suffix_array sa(s1 + "#" + s2);
    vector<int> suf = sa.suf, lcp = sa.lcp;
    auto type = [&](int idx) {
        return idx <= s1.size();</pre>
    int mx = 0, idx = 0;
    int len = s1.size() + 1 + s2.size();
    for (int i = 1; i <= len; i++)</pre>
        if (type(suf[i - 1]) != type(suf[i]) && lcp[i] > mx) {
            mx = lcp[i];
            idx = min(suf[i - 1], suf[i]);
    return sl.substr(idx, mx);
int bongest_common_substring(const vector<string> &v) {
    int n = v.size();
    fritsten = n - 1;
    for (auto &it : v)
    + len += it.size();
    string s(len, '.');
    \hat{\mathbf{v}}ector<\hat{\mathbf{int}}> type(len + 1, n), frq(n + 1);
    for (int i = 0, j = 0; i < v.size(); i++) {</pre>
    } if (i)
            s[j] = 'z' + i;
        for (char ch : v[i]) {
            s[j] = ch;
            type[j] = i;
            j++;
    suffix_array sa(s);
    vector<int> suf = sa.suf, lcp = sa.lcp;
    monoqueue q;
    int st = 0, ed = 0, cnt = 0, mx = 0;
    while (st <= s.size()) {</pre>
        while (ed <= s.size() && cnt < v.size()) {
            q.push(lcp[ed], ed);
            if (++frq[type[suf[ed]]] == 1)
                cnt++;
            ed++;
        q.pop(st);
        if (cnt == v.size())
            mx = max(mx, q.getMin()); //st+1,ed
        if (--frq[type[suf[st]]] == 0)
            cnt--;
        st++;
    return mx;
string kth_substring(string s, int k) { //1-based, repated
    int n = s.size();
    suffix_array sa(s);
    vector<int> suf = sa.suf, lcp = sa.lcp;
    for (int i = 1; i <= n; i++) {</pre>
        int len = n - suf[i];
        int cnt = 0:
        for (int 1 = 1; 1 <= len; 1++) {</pre>
            cnt++:
            int st = i + 1, ed = n, ans = i;
            while (st <= ed) {
                int md = st + ed >> 1;
                if (sa.LCP_by_order(i, md) >= 1)
                     st = md + 1, ans = md;
                 else
```

```
ed = md - 1;
         cnt += ans - i;
         if (cnt \geq= k)
             return s.substr(suf[i], 1);
    k \rightarrow len;
assert(0);
```

9d6a18, 116 lines

13

```
SuffixAutomaton.h
struct suffix automaton {
 struct state {
   int len, link = 0, cnt = 0;
   bool terminal = false, is_clone = false;
    map<char, int> next;
    state(int len = 0) :
       len(len) {
   bool have_next(char ch) {
      return next.find(ch) != next.end();
   void clone(const state &other, int nlen) {
     len = nlen:
     next = other.next;
     link = other.link;
     is clone = true;
 };
 vector<state> st;
 int last = 0;
 suffix automaton() {
   st.push back(state());
   st[0].link = -1;
 suffix_automaton(const string &s) :
      suffix automaton() {
    for (char ch : s)
     extend(ch);
   mark terminals();
 void extend(char c) {
   int cur = st.size();
   st.push_back(state(st[last].len + 1));
   st[cur].cnt = 1;
   int p = last;
   last = cur:
    while (p != -1 && !st[p].have_next(c)) {
      st[p].next[c] = cur;
     p = st[p].link;
   if (p == -1)
      return;
    int q = st[p].next[c];
    if (st[p].len + 1 == st[q].len) {
     st[cur].link = q;
      return;
    int clone = st.size();
    st.push_back(state());
    st[clone].clone(st[q], st[p].len + 1);
    while (p != -1 \&\& st[p].next[c] == q) {
     st[p].next[c] = clone;
      p = st[p].link;
    st[q].link = st[cur].link = clone;
```

Manacher Prime MillerRabinPrimalityTest

```
void mark terminals() {
 for (int cur = last; cur > 0; cur = st[cur].link)
   st[cur].terminal = true;
void calc_number_of_occurrences() {
 vector<vector<int>> lvl(st[last].len + 1);
  for (int i = 1; i < st.size(); i++)</pre>
   lvl[st[i].len].push_back(i);
  for (int i = st[last].len; i >= 0; i--)
   for (auto cur : lvl[i])
     st[st[cur].link].cnt += st[cur].cnt;
vector<ll> dp;
11 Count (int cur) { //count number of paths
 11 &rt = dp[cur];
 if (rt)
   return rt;
 rt = 1;
 for (auto ch : st[cur].next)
  rt += Count(ch.second);
  return rt;
//1-based, different substring, 0 = ""
string kth_substring(ll k) {
 assert(k <= Count(0));
 string rt;
 int cur = 0;
  while (k > 0) {
   for (auto ch : st[cur].next) {
     if (Count(ch.second) < k)</pre>
       k -= Count (ch.second);
      else {
       rt += ch.first;
       cur = ch.second;
       k--;
       break;
 return rt;
string longest common substring(const string &t) {
 int cur = 0, 1 = 0, mx = 0, idx = 0;
  for (int i = 0; i < t.size(); i++) {</pre>
   while (cur > 0 && !st[cur].have_next(t[i])) {
     cur = st[cur].link;
     1 = st[cur].len;
   if (st[cur].have next(t[i])) {
     cur = st[cur].next[t[i]];
     1++;
   if (1 > mx) {
     mx = 1:
      idx = i;
  return t.substr(idx - mx + 1, mx);
```

3.4 Manacher Hash

Manacher.h

b65384, 32 lines

```
void Manacher(string& s) {
  int n = sz(s);
  vector<int> d1(n);
  //d1[i] and d2[i], denoting the number of
  //palindromes accordingly with odd and even lengths with
```

```
//centers in the position i
// with d2 the center of aa is pos 1
// 0-based
for (int i = 0, l = 0, r = -1; i < n; i++) {
  int k = (i > r) ? 1 : min(d1[1 + r - i], r - i + 1);
  while (0 \le i - k \&\& i + k \le n \&\& s[i - k] == s[i + k])  {
  d1[i] = k--;
  if (i + k > r) {
   1 = i - k;
   r = i + k;
vector<int> d2(n);
for (int i = 0, l = 0, r = -1; i < n; i++) {
  int k = (i > r) ? 0 : min(d2[1 + r - i + 1], r - i + 1);
  while (0 \le i - k - 1 \&\& i + k \le n \&\& s[i - k - 1] == s[i + k]
        k]) {
   k++;
  d2[i] = k--;
  if (i + k > r) {
  1 = i - k - 1;
   r = i + k;
```

Math (4)

4.1 Prime

Prime.h

20dd4a, 69 lines

```
//linear sieve
const int N = 1e7:
int lpf[N + 1];
vector<int> prime;
void sieve() {
  for (int i = 2; i <= N; i++) {</pre>
    if (lpf[i] == 0) {
      lpf[i] = i;
      prime.push_back(i);
    for (int j : prime) {
     if (j > lpf[i] || 1LL * i * j > N)break;
      lpf[i * j] = j;
 }
// return number of Divisors(n) using prime factorization
11 numOfDivisors(primeFactors mp) {
 11 \text{ cnt} = 1;
  for (auto it : mp) cnt *= (it.second + 1);
  return cnt;
// return sum of Divisors(n) using prime factorization
11 sumOfDivisors(primeFactors mp) {
  for (auto it : mp) sum *= sumPower(it.first, it.second);
  return sum;
ll phi_function(ll n) {
  11 result = n;
  primeFactors pf = prime_factors(n);
```

```
for (auto &it : pf) {
    11 p = it.first;
    result -= (result / p);
  return result;
void phi_1_to_n(int n) {
  for (int i = 0; i <= n; i++)</pre>
    phi[i] = i;
  for (int i = 2; i <= n; i++)</pre>
    if (phi[i] == i)
      for (int j = i; j <= n; j += i)
        phi[j] -= phi[j] / i;
char mob[N];
bool prime[N];
void moebius() {
  memset (mob, 1, sizeof mob);
  memset (prime + 2, 1, sizeof (prime) - 2);
  mob[0] = 0;
  mob[2] = -1;
  for (int i = 4; i < N; i += 2) {
    mob[i] *= (i & 3) ? -1 : 0;
    prime[i] = 0;
  for (int i = 3; i < N; i += 2)
    if (prime[i]) {
      mob[i] = -1;
      for (int j = 2 * i; j < N; j += i) {
       mob[j] *= j % (1LL * i * i) ? -1 : 0;
        prime[j] = 0;
```

if (!(n & 1)) return false;

```
MillerRabinPrimalityTest.h
                                                      5697aa, 47 lines
const int ITER = 4;
mt19937 rng(chrono::steady_clock::now().time_since_epoch().
     count());
__int128 power(__int128 x, __int128 y, 11 mod) {
    if (v == 0)
        return 1;
    if (y & 1)return (x * power(x, y - 1, mod)) % mod;
    _{int128} r = power(x, y >> 1, mod);
    return (r * r) % mod;
bool millerTest(ll n, ll d) {
    __int128 a = uniform_int_distribution<11>(2, n - 2)(rng);
    a = power(a, d, n);
    if (a == 1 || a == n - 1)
        return true;
    d <<= 1:
    while (d != n - 1) {
        a = a * a % n;
        if (a == 1) return false;
        if (a == n - 1) return true;
        d <<= 1;
    return false;
bool is_prime(ll n) {
    if (n <= 1) return false;</pre>
    if (n <= 3) return true;</pre>
```

```
11 d = n - 1;
    while (!(d & 1))
        d >>= 1;
    for (int i = 0; i < ITER; i++)</pre>
        if (!millerTest(n, d))
             return false:
    return true;
bool isPrimeSquare(ll n) {
    11 \text{ sq} = \text{sqrt}(n);
    if (sq * sq < n) {
         if (sq * sq != n) return false;
    return is_prime(sq);
```

4.2 ModInverse

ModInverse.h

043bf3, 48 lines

```
#define 11 long long
ll power(ll x, ll y, int mod) {
 if (y == 0)
   return 1;
  if (y == 1)
    return x % mod;
  11 r = power(x, y >> 1, mod);
  return (((r * r) % mod) * power(x, y & 1, mod)) % mod;
// return a \land 1 + a \land 2 + a \land 3 + \ldots a \land k
11 sumPower(ll a, ll k, int mod) {
  if (k == 1) return a % mod;
  ll half = sumPower(a, k / 2, mod);
  ll p = half * power(a, k / 2, mod) % mod;
  p = (p + half) % mod;
  if (k & 1) p = (p + power(a, k, mod)) % mod;
  return p;
ll modInverse(ll b, ll mod) { // if mod is Prime
 return power(b, mod - 2, mod);
ll modInverse(ll b, ll mod) { // if mod is not Prime, gcd(a,b)
     must be equal 1
  return power(b, phi_function(mod) - 1, mod);
// (a^n)%p=result, return n
int getPower(int a, int result, int mod) {
  int sq = sqrt(mod);
  map<int, int> mp;
  11 r = 1;
  for (int i = 0; i < sq; i++) {
   if (mp.find(r) == mp.end())
     mp[r] = i;
    r = (r * a) % mod;
  11 tmp = modInverse(r, mod);
  11 cur = result;
  for (int i = 0; i <= mod; i += sq) {</pre>
   if (mp.find(cur) != mp.end())
     return i + mp[cur];
    cur = (cur * tmp) % mod; //val/(a^sq)
  return INF;
```

4.3 Number Theory

ExtendedEuclidean.h

e4e457, 84 lines

```
ll egcd(ll a, ll b, ll& x, ll& y) {
    if (a < 0) {
       auto g = egcd(-a, b, x, y);
       x \star = -1;
        return q;
    if (b < 0) {
        auto g = egcd(a, -b, x, y);
       v *= -1:
        return g;
    if (!b) {
       x = 1:
       v = 0;
        return a;
    11 x1, v1;
    11 q = eqcd(b, a % b, x1, y1);
    x = y1, y = x1 - y1 * (a / b);
    return q;
//O(n * log(m)) Memory & Time; coefficients.size() <= n,
     coefficients[i] \le m
// 0-based implementation
template<typename T>
T extended_euclidean(const deque<T>& cof, deque<T>& var) {
    int n = cof.size();
    if (!cof.back()) {
        int cnt = 0, id = 0;
        for (int i = 0; i < n; i++)</pre>
            if (!cof[i]) {
                cnt++;
                var[i] = 0;
            } else id = i;
        if (cnt >= n - 1) {
            var[id] = 1;
            return cof[id];
        deque<T> new_cof, new_var;
        for (int i = 0; i < n; i++)</pre>
            if (cof[i]) {
                new_cof.push_back(cof[i]);
                new_var.push_back(var[i]);
        T q = extended_euclidean(new_cof, new_var);
        for (int i = 0; !new_var.empty(); i++)
            if (cof[i]) {
                var[i] = new_var.front();
                new_var.pop_front();
        return g;
    deque<T> new_cof = cof;
    for (int i = 0; i < n - 1; i++)</pre>
       new_cof[i] %= new_cof.back();
    new_cof.push_front(new_cof.back());
   new cof.pop back();
    var.push_front(var.back());
    var.pop_back();
    T g = extended_euclidean(new_cof, var);
    var.push_back(var.front());
    var.pop front();
```

```
for (int i = 0; i < n - 1; i++)
        var.back() -= cof[i] / cof.back() * var[i];
    return q;
template<typename T>
vector<T> find any solution(const vector<T>& cof, T rhs) {
    int n = cof.size();
    if (!n)
        return vector<T>();
    deque<T> deque_cof(cof.begin(), cof.end()), deque_var(n);
    T q = extended_euclidean(deque_cof, deque_var);
    if (q && rhs % q)
        return vector<T>();
    vector<T> var(deque_var.begin(), deque_var.end());
    if (g) {
        rhs /= g;
        for (auto& it : var)
            it *= rhs:
    return var;
Diophantine.h
                                                     41fbe7, 96 lines
// return false if there is no solution
// return true if there exist a solution
// x, y are the solutions and q is the gcd between a and b
bool Diophantine Solution(11 a, 11 b, 11 c, 11& x, 11& y, 11& q
    if (!a && !b) {
        if (c) return false;
        x = y = q = 0;
        return true;
    g = egcd(a, b, x, v);
    if (c % q) return false;
    x *= c / q;
    v *= c / q;
    return true;
void shift_solution(ll& x, ll& y, ll a, ll b, ll cnt) {
    x += b * cnt;
    y -= a * cnt;
```

// find all number of solutions of ax + by = c

if (minx > maxx || miny > maxy) return 0;

if (!Diophantine_Solution(a, b, c, x, y, g))

11 Diophantine_Solutions(11 a, 11 b, 11 c, 11 minx, 11 maxx, 11

return (maxx - minx + 1) * (maxy - miny + 1);

return (num >= miny && num <= maxy) * (maxx - minx + 1)</pre>

return (num >= minx && num <= maxx) * (maxy - miny + 1)</pre>

// x in range {minx, maxx}

// y in range {miny, maxy}

miny, 11 maxy) {

if (!a && !b && !c)

if (!a) {

if (!b) {

11 x, y, g;

return 0:

if (!a && !b) return 0;

11 num = c / b;

11 num = c / a;

if (c % b) **return** 0;

if (c % a) **return** 0;

```
11 1x1, 1x2, rx1, rx2;
// a * x + b * y = c
//(a/g) * x + (b/g) * y = c/g
a /= g, b /= g, c /= g;
q = 1;
int sign_a = (a > 0 ? 1 : -1);
int sign_b = (b > 0 ? 1 : -1);
// x + k * b >= minx
// k * b >= minx - x
// k >= (minx - x) / b
// k >= ceil((minx - x) / b)
shift_solution(x, y, a, b, (minx - x) / b);
// if x is less than minx so we need to increase it by one
     step only
// from the upove equation x + k * b >= minx
// if b is positive so choose k equal to 1 to increase x
// if b is negative so choose k equal to -1 because -1 *
    -1 = 1 ans also increase x
if (x < minx)
    shift_solution(x, y, a, b, sign_b);
if (x > maxx) return 0;
1x1 = x;
// x + k * b \le maxx
// k * b \le maxx - x
// k \le (maxx - x) / b
shift_solution(x, y, a, b, (maxx - x) / b);
if (x > maxx)
    shift_solution(x, y, a, b, -sign_b);
// y - k * a >= miny
// y - miny >= k * a
// k * a \le y - miny
// k \le (y - miny) / a
shift_solution(x, y, a, b, (y - miny) / a);
if (y < miny)</pre>
   shift_solution(x, y, a, b, -sign_a);
if (v > maxy) return 0;
1x2 = x;
// y - k * a \le maxy
// y - maxy \le k * a
// k * a >= y - maxy
// k >= (y - maxy) / a
shift_solution(x, y, a, b, (y - maxy) / a);
if (y > maxy)
   shift_solution(x, y, a, b, sign_a);
if (1x2 > rx2) swap(1x2, rx2);
// because we calculate the equations lx2, rx2 from
     shifting y
// not from shifting x directly
11 1x = max(1x1, 1x2);
11 \text{ rx} = \min(\text{rx1, rx2});
if (lx > rx) return 0;
return (rx - lx) / abs(b) + 1;
```

CRT.h

c447d1, 15 lines

```
11 CRT(vector<11>& a, vector<11>& m) {
 11 lcm = m[0], rem = a[0];
  int n = a.size();
  for(int i = 1; i < n; i++) {</pre>
   11 gcd = extended_euclidean(lcm, m[i], x, y);
   if((a[i] - rem) % gcd) return -1;
   ll tmp = m[i] / qcd, f = (a[i] - rem) / <math>qcd;
          x = ((x % tmp) * (f % tmp)) % tmp;
```

```
rem += lcm * x;
    lcm = lcm * m[i] / gcd;
   rem = (rem % lcm + lcm) % lcm;
 return rem;
4.4 Combinatorics
Combinatorics.h
                                                      534525, 77 lines
typedef unsigned long long ull;
// nCr = n!/((n-r)! * r!)
// nCr(n,r) = nCr(n,n-r)
// nPr = n!/(n-r)!
// nPr\_circle = nPr/r
// nCr(n,r) = pascal[n]/r
// \operatorname{catalan}[n] = n\operatorname{Cr}(2n,n)/(n+1)
ull nCr(int n, int r) {
 if (r > n)
    return 0;
  r = max(r, n - r);
  ull ans = 1, div = 1, i = r + 1;
  while (i <= n) {
    ans *= i++:
    ans /= div++;
  return ans;
ull nPr(int n, int r) {
 if (r > n)
   return 0;
  ull p = 1, i = n - r + 1;
  while (i <= n)</pre>
   p *= i++;
  return p;
vector<vector<ull>> pascalTriangle(int n) {
  vector<vector<ull>>> pascal(n + 1, vector<ull>(n + 1));
  for (int i = 0; i <= n; i++) {</pre>
    pascal[i][i] = pascal[i][0] = 1;
    for (int j = 1; j < n; j++)
      pascal[i][j] = pascal[i - 1][j] + pascal[i - 1][j - 1];
 return pascal;
// return catalan number n-th using dp O(n^2)//max = 35 then
     overflow
vector<ull> catalanNumber(int n) {
 vector<ull> catalan(n + 1);
 catalan[0] = catalan[1] = 1;
  for (int i = 2; i <= n; i++) {</pre>
   ull &rt = catalan[i];
    for (int j = 0; j < n; j++)
      rt += catalan[j] * catalan[n - j - 1];
  return catalan;
// count number of paths in matrix n*m
// go to right or down only
ull countNumberOfPaths(int n, int m) {
  return nCr (n + m - 2, n - 1);
```

```
const int N = 1e5 + 100;
const int MOD = 1e9 + 7;
11 fact[N];
11 inv[N]; //mod inverse for i
11 invfact[N]; //mod inverse for i!
void factInverse() {
 fact[0] = inv[1] = fact[1] = invfact[0] = invfact[1] = 1;
  for (long long i = 2; i < N; i++) {</pre>
    fact[i] = (fact[i - 1] * i) % MOD;
    inv[i] = MOD - (inv[MOD % i] * (MOD / i) % MOD);
    invfact[i] = (inv[i] * invfact[i - 1]) % MOD;
11 nCr(int n, int r) {
 if (r > n)
    return 0;
  return (((fact[n] * invfact[r]) % MOD) * invfact[n - r]) %
```

```
4.5 Matrix
Matrix.h
                                                   bb4c31, 104 lines
#define 11 long long
#define sz(s) (int)s.size()
#define REP(i,n) for(int i = 0;i<n;i++)
struct matrix {
 using T = int;
 using row = vector<T>;
 vector<vector<T>> v;
 matrix() {
 matrix(int n, int m, T val = 0) :
     v(n, row(m, val)) {
 int size() const {
   return v.size();
 int cols() const {
   return v[0].size();
 matrix operator*(T a) const {
   matrix rt = *this;
   REP(i,rt.size())
     REP(j,rt.cols())
       rt.v[i][j] *= a;
   return rt;
 friend matrix operator*(T a, const matrix &b) {
    return (b * a);
 friend matrix operator+(const matrix &a, const matrix &b) {
   assert(a.size() == b.size() && a.cols() == b.cols());
   matrix rt(a.size(), a.cols());
   REP(i,rt.size())
     REP(j,rt.cols())
       rt.v[i][j] = a.v[i][j] + b.v[i][j];
   return rt:
 friend matrix operator*(const matrix &a, const matrix &b) {
   assert(a.cols() == b.size());
   matrix rt(a.size(), b.cols());
   REP(i,rt.size())
     REP(k,a.cols())
       if (a.v[i][k] == 0)
         continue;
```

```
REP(j,rt.cols())
          rt.v[i][j] += a.v[i][k] * b.v[k][j];
    return rt;
};
matrix identity(int n) {
 matrix r(n, n);
  for (int i = 0; i < n; i++)</pre>
    r.v[i][i] = 1;
  return r;
matrix addIdentity(const matrix &a) {
  matrix rt = a;
  REP(i,a.size())
   rt.v[i][i]++;
  return rt;
matrix power (matrix a, long long y) {
  assert(v >= 0 \&\& a.size() == a.cols());
  matrix rt = identity(a.size());
  while (y > 0) {
   if (y & 1)
     rt = rt * a;
    a = a * a;
   y >>= 1;
  return rt:
matrix sumPower(const matrix &a, 11 k) {
  if (k == 0)
    return matrix(sz(a), sz(a));
  if (k & 1)
    return a * addIdentity(sumPower(a, k - 1));
  return (sumPower(a, k >> 1) * addIdentity(power(a, k >> 1)));
/* return matrix contains
 a^k
a^1+a^2...a^k I
matrix sumPowerV2(const matrix &a, 11 k) {
  int n = sz(a);
  matrix rt(2 * n, 2 * n);
  REP(i,n)
   REP(j,n)
     rt.v[i][j] = a.v[i][j];
     rt.v[i + n][j] = a.v[i][j];
  for (int i = n; i < 2 * n; i++)
    rt.v[i][i] = 1;
  return power(rt, k);
```

4.6 FFT

FFT.h

41ce38, 96 lines

```
typedef valarray<complex<double>> polynomial;
const int LGN = 20;
vector<complex<double>> CM1[3][LGN + 1];
const double PI = acos(-1);
void prepare() {
   for (int sign = -1; sign <= 1; sign += 2) {</pre>
```

```
for (int i = 0; i <= LGN; i++) {</pre>
            int N = 1 << i;
            double theta = sign * 2 * PI / N;
            complex<double> cm1 = 1;
            complex<double> cm2(cos(theta), sin(theta));
            for (int j = 0; j < N / 2; j++) {
                CM1[sign + 1][i].push_back(cm1);
                cm1 *= cm2;
    }
void fft(polynomial &a, int sign = -1) {
    int N = a.size();
    int lqn = log2(N);
    for (int m = N; m >= 2; m >>= 1, lqn--) {
        int mh = m >> 1;
        for (int i = 0; i < mh; i++) {</pre>
            const complex<double> &w = CM1[sign + 1][lgn][i];
            for (int j = i; j < N; j += m) {</pre>
                int k = j + mh;
                complex<double> x = a[j] - a[k];
                a[i] += a[k];
                a[k] = w * x;
        }
    int i = 0;
    for (int j = 1; j < N - 1; j++) {
        for (int k = N >> 1; k > (i ^= k); k >>= 1);
        if (j < i)
            swap(a[i], a[j]);
valarray<ll> inv_fft(polynomial&& a) {
    complex<double> N = a.size();
    fft(a, 1);
    a /= N;
    valarrav<ll> rt(a.size());
    for (int i = 0; i < a.size(); i++)</pre>
        rt[i] = round(a[i].real());
    return rt;
valarray<int> mul(const valarray<int> &a, const valarray<int> &
  int adeg = (int) a.size() - 1, bdeg = (int) b.size() - 1;
  int N = 1;
  while (N <= adeg + bdeg)
   N <<= 1;
  polynomial A(N), B(N);
  for (int i = 0; i < a.size(); i++)</pre>
   A[i] = a[i];
  for (int i = 0; i < b.size(); i++)</pre>
   B[i] = b[i];
  fft(A);
  fft(B);
  polynomial m = A * B;
  inv fft(m);
  return rt;
valarray<int> mul_with_mod(const vector<int> & a, const vector<</pre>
     int>\& b, int MOD = 1e9 + 7) {
    int adeg = (int) a.size() - 1, bdeg = (int) b.size() - 1;
    int N = 1:
    while (N <= adeg + bdeg)
        N <<= 1;
    int C = sqrt(MOD);
    polynomial a1(N), a2(N);
    polynomial b1(N), b2(N);
```

```
for (int i = 0; i < a.size(); ++i) {</pre>
        al[i] = a[i] % C;
        a2[i] = a[i] / C;
    for (int i = 0; i < b.size(); ++i) {</pre>
        b1[i] = b[i] % C;
        b2[i] = b[i] / C;
    fft(a1), fft(a2);
    fft(b1), fft(b2);
    valarray<11> m11 = inv_fft(a1 * b1) % MOD;
    valarray<11> m12 = inv_fft(a1 * b2) % MOD;
    valarray<11> m21 = inv_fft(a2 * b1) % MOD;
    valarray<11> m22 = inv_fft(a2 * b2) % MOD;
    valarray<11> res = m11 % MOD;
    res += \tilde{C} * (m12 + m21) % MOD;
    res += C * (C * m22 % MOD) % MOD;
    res %= MOD;
    valarray<int> rt(res.size());
    for (int i = 0; i < res.size(); i++)</pre>
        rt[i] = res[i];
    return rt;
NTT.h
                                                     0493f2, 178 lines
const int LGN = 20;
struct modint {
#define CUR (*this)
 int val;
  modint(const long long& a = 0) {
    val = a % MOD;
    if (val < 0)
      val += MOD;
 modint& operator+=(const modint& a) {
    if ((val += a.val) >= MOD)
     val -= MOD;
    return CUR;
 modint operator+(const modint& a) const {
    modint c = CUR;
    c += a;
    return c:
  modint& operator-=(const modint& a) {
    if ((val -= a.val) < 0)</pre>
     val += MOD;
    return CUR;
 modint operator-(const modint& a) const {
    modint c = CUR;
    c -= a;
    return c:
 modint operator*(const modint& a) const {
    return modint((1LL * this->val * a.val) % MOD);
  modint& operator *= (const modint& a) {
    CUR = CUR * a;
    return CUR;
  modint operator/(const modint& a) {
    return CUR * power(a, MOD - 2);
 modint& operator/=(const modint& a) {
    CUR = CUR / a;
    return CUR;
```

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```
static modint power(modint x, long long y) {
    modint res = 1;
    while (y > 0) {
     if (y & 1)
       res *= x;
      x *= x;
      y >>= 1;
    return res;
  friend ostream& operator<<(ostream& out, const modint& a) {</pre>
   out << a.val;
    return out;
#undef CUR
};
typedef valarray<modint> polynomial;
vector<modint> CM1[2][LGN + 1];
bool validRoot(modint root) {
  modint rootinv = modint(1) / root;
  for (int invert = 0; invert <= 1; invert++) {</pre>
    for (int i = 1; i <= LGN; i++) {</pre>
     int N = 1 << i;</pre>
      assert ((MOD -1) % N == 0);
      int C = (MOD - 1) / N;
      modint cm2 = modint::power(invert ? root : rootinv, C);
      if (cm2.val <= 1)
        return false;
  return true;
void prepare() {
  modint root = 2;
  while (!validRoot(root))
   root += 1:
  modint rootinv = modint(1) / root;
  for (int invert = 0; invert <= 1; invert++) {</pre>
    for (int i = 0; i <= LGN; i++) {</pre>
      int N = 1 << i;</pre>
      int C = (MOD - 1) / N;
      modint cm2 = modint::power(invert ? root : rootinv, C);
      modint cm1 = 1;
      set<int> st;
      for (int j = 0; j < N / 2; j++) {
        CM1[invert][i].push_back(cm1);
        cm1 \star = cm2;
void fft(polynomial& a, bool invert = 0) {
  int N = a.size();
  int lgn = log2(N);
  for (int m = N; m >= 2; m >>= 1, lqn--) {
    int mh = m >> 1;
    for (int i = 0; i < mh; i++) {</pre>
      const modint& w = CM1[invert][lqn][i];
      for (int j = i; j < N; j += m) {</pre>
       int k = j + mh;
        modint x = a[j] - a[k];
       a[j] += a[k];
        a[k] = w * x;
```

```
int i = 0:
 for (int j = 1; j < N - 1; j++) {
   for (int k = N >> 1; k > (i ^= k); k >>= 1)
   if ( † < i)
     swap(a[i], a[j]);
void inv_fft(polynomial& a) {
 int N = a.size();
 fft(a, 1);
 a /= N;
valarray<modint> mul(const polynomial& a, const polynomial& b)
 int adeg = (int)a.size() - 1, bdeg = (int)b.size() - 1;
 int N = 1;
 while (N <= adeg + bdeg)</pre>
  N <<= 1;
 polynomial A(N), B(N);
 for (int i = 0; i < a.size(); i++)</pre>
   A[i] = a[i];
 for (int i = 0; i < b.size(); i++)</pre>
   B[i] = b[i];
 fft(A);
 fft(B);
 polynomial rt = A * B;
 inv fft(rt);
 return rt;
int main() {
 run();
 prepare();
 int t:
 cin >> t;
 while (t--) {
   int n;
    cin >> n:
    polynomial f1(n + 1), f2(n + 1);
    for (int i = 0; i < n; i++) {</pre>
     int x, f;
     cin >> x >> f;
     f1[x] += f;
     f2[n - x] += f;
   polynomial res = mul(f1, f2);
   vector<modint> out(res.size());
   out[0] = res[0];
    for (int i = 1; i < sz(res); i++)
     out[i] = res[i] + out[i - 1];
   modint inv2 = modint(1) / 2;
   int a:
   cin >> q;
    while (q--) {
     int 1, r; cin >> 1 >> r;
     modint res = out[n + r] - out[n + 1 - 1];
     res += out[n - 1];
     if (n - r - 1 \ge 0) res -= out [n - r - 1];
     cout << res * inv2 << endl;
 }
```

Geometry (5)

5.1 Basic Geometry

Point.h

```
<br/>
<br/>
dits/stdc++.h>
                                                     54fd2e, 46 lines
using namespace std;
#define 11 long long
typedef complex<double> point; // it can be long long not
     double
template<class T>
istream& operator>>(istream& is, complex<T>& p) {
 T value:
 is >> value;
  p.real(value);
  is >> value;
  p.imag(value);
  return is:
#define PI acos(-1.0)
#define EPS 1e-8
#define X real()
#define Y imag()
#define angle(a) (atan2((a).imag(), (a).real())) // angle with
      orignial
#define length(a) (hypot((a).imag(), (a).real()))
#define vec(a,b) ((b)-(a)) // return diff x1-x2, y1-y2
#define dp(a,b) ( (conj(a)*(b)).real() )
// a*b cos(T), if zero \Rightarrow prep dot product A.B
#define cp(a,b) ( (conj(a)*(b)).imag() )
// a*b sin(T), if zero \rightarrow parllel\ cross\ product = area\ of
     parllelogram
#define normalize(a) (a)/length(a)
// norm(a) // return x^2 + y^2 //a is point //can use dp(a,a)
bool same (point p1, point p2) {// check to points same or not
 return dp(vec(p1, p2), vec(p1, p2)) < EPS;
point rotate(point p, double angle, point around = point(0, 0))
  p -= around;
  return (p * exp(point(0, angle))) + around;
// Refelect v around m
point reflectO(point v, point m) {
  return conj(v / m) * m;
// Refelect point p around l1-l2
point reflect (point p, point 11, point 12) {
  point z = p - 11, w = 12 - 11;
  return conj(z / w) * w + 11;
Triangles.h
                                                      c25cbe, 71 lines
```

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```
"point.h"
// \sin(A)/a = \sin(B)/b = \sin(C)/c
// a^2 = b^2 + c^2 - 2b*c*cos(A)
// \sin(A+B) = \sin(A) * \cos(B) + \sin(B) * \cos(A)
// \sin(A-B) = \sin(A) * \cos(B) - \sin(B) * \cos(A)
//\cos(A+B) = \cos(A) * \cos(B) - \sin(A) * \sin(B)
//\cos(A-B) = \cos(A) * \cos(B) + \sin(A) * \sin(B)
\frac{1}{1/2} \tan(A+B) = \frac{(\tan(A) + \tan(B))}{(1 - \tan(A) \cdot \tan(B))}
// \tan(A-B) = (\tan(A) - \tan(B))/(1 - \tan(A) * \tan(B))
double fixAngle(double A) {
```

```
return A > 1 ? 1 : (A < -1 ? -1 : A);
// return min angle: aOb / bOa
// dp(v1, v2) = |v1| * |v2| * cos(theta)
double angleO(point a, point O, point b) {
  point v1(a - 0), v2(b - 0);
  return acos(fixAngle(dp(v1, v2) / dist(v1) / dist(v2)));
double getSide a bAB(double b, double A, double B) {
  return (sin(A) * b) / sin(B);
double getAngle A abB(double a, double b, double B) {
 return asin(fixAngle((a * sin(B)) / b));
// wr answer in team formation :D
double getAngle_A_abc(double a, double b, double c) {
 return acos(fixAngle((b * b + c * c - a * a) / (2 * b * c));
double triangleArea(double a, double b, double c) {
  double s = (a + b + c) / 2.0;
  return sqrt((s - a) * (s - b) * (s - c) * s);
double triangleArea(point p0, point p1, point p2) {
  double a = length(vec(p1, p0)), b = length(vec(p2, p0)),
   c = length(vec(p2, p1));
  return triangleArea(a, b, c);
double triangleArea3points(const point &a, const point &b,
    const point &c) {
  return fabs(cross(a,b) + cross(b,c) + cross(c,a)) / 2;
bool pointInTriangle(point a, point b, point c, point pt) {
 11 s1 = fabs(cp(vec(a,b), vec(a,c)));
  11 	ext{ s2} = fabs(cp(vec(pt,a), vec(pt,b)))
    + fabs(cp(vec(pt, b), vec(pt, c)))
    + fabs(cp(vec(pt, a), vec(pt, c)));
  return s1 == s2;
// largest circle inside a triangle
//A triangle with area A and semi-perimeter s has an inscribed
    circle (incircle) with
//radius r = A/s
bool circleInTriangle(point a, point b, point c, point& ctr,
    double& r) {
  double ab = length(a - b), bc = length(b - c),
   ca = length(c - a);
  double s = 0.5 * (ab + bc + ca);
  r = triangleArea(ab, bc, ca) / s;
  if (fabs(r) < EPS) return 0; // no inCircle center</pre>
  double ratio = length(a - b) / length(a - c);
  point p1 = b + (vec(b, c) * (ratio / (1 + ratio)));
  ratio = length(b - a) / length(b - c);
  point p2 = a + (vec(a, c) * (ratio / (1 + ratio)));
  return intersectSegments(a, p1, b, p2, ctr); // get their
       intersection point
```

Ala7lam_team (Assiut University)

```
Lines.h
"point.h"
                                                     a29fdd, 106 lines
//o = anlge(a) - angle(b)
 //if \ o = 0 \ | \ | \ o = 180 \ is Collinear
  //else not
bool isCollinear(point a, point b, point c) {
  return fabs(cp(vec(a, b), vec(a, c))) < EPS;
//o = anlge(a) - angle(b)
//if \ o = 0 \ isPointOnRay \ a \rightarrow b
//else_not
bool isPointOnRay(point a, point b, point c) {
  if (!isCollinear(a, b, c))
    return false:
  return dcmp(dp(vec(a, b), vec(a, c)), 0) == 1;
bool isPointOnRay(point a, point b, point c) {
  if (length(vec(a, c)) < EPS) return true;</pre>
  return same (normalize (vec(a, b)), normalize (vec(a, c)));
bool isPointOnSegment(point a, point b, point c) {
  return isPointOnRay(a, b, c) && isPointOnRay(b, a, c);
bool isPointOnSegment(point a, point b, point c) {
  double acb = length(vec(b, a)), ac = length(vec(c, a)), cb =
      length(vec(c, b));
  return demp(acb - (ac + cb), 0) == 0;
// dist point p2 to line p0-p1
double distToLine(point p0, point p1, point p2) {
  return fabs(cp(vec(p0, p1), vec(p0, p2)) / length(vec(p1, p0)
      )); // area = 0.5*b*h
//minimum distance from point p2 to segment p0-p1
double distToSegment(point p0, point p1, point p2) {
  double d1, d2;
  point v1 = p1 - p0, v2 = p2 - p0;
  if ((d1 = dp(v1, v2)) \le 0) return length(vec(p0, p2));
  if ((d2 = dp(v1, v1)) \le d1) return length(vec(p1, p2));
  double t = d1 / d2;
  return length(vec((p0 + v1 * t), p2));
// minimum point in segment po-p1 to point p2
point pointToSegment(point p0, point p1, point p2) {
  double d1, d2;
  point v1 = p1 - p0, v2 = p2 - p0;
  if ((d1 = dp(v1, v2)) <= 0) return p0;
  if ((d2 = dp(v1, v1)) <= d1) return p1;</pre>
  double t = d1 / d2;
  return (p0 + v1 * t);
// return point intersect in line a-b with c-d using parametric
      equations
bool intersectSegments(point a, point b, point c, point d,
     point& intersect) {
  double d1 = cp(vec(b, a), vec(c, d)),
    d2 = cp(vec(c, a), vec(c, d)),
    d3 = cp(vec(b, a), vec(c, a));
  if (fabs(d1) < EPS)</pre>
    return false; // Parllel || identical
  double t1 = d2 / d1, t2 = d3 / d1;
```

```
intersect = a + (b - a) * t1;
  if (t1 < -EPS || t2 < -EPS || t2 > 1 + EPS)
    return false; //e.g ab is ray, cd is segment ... change to
          whatever
  return true:
// return 1 if point c is counter-clockwise about segment a-b
// -1 if point c is clockwise about segment a-b
// 0 if c is is Collinear about a-b
int ccw(point a, point b, point c) {
 point v1(b - a), v2(c - a);
  double t = cp(v1, v2);
 if (t > EPS)
   return 1;
  if (t < -EPS)
   return -1:
 if (v1.X * v2.X < -EPS | | v1.Y * v2.Y < -EPS)
   return -1;
  if (norm(v1) < norm(v2) - EPS)
    return 1;
 return 0;
bool intersect (point p1, point p2, point p3, point p4) {
  // special case handling if a segment is just a point
  bool x = (p1 == p2), y = (p3 == p4);
 if (x && y) return p1 == p3;
 if (x) return ccw(p3, p4, p1) == 0;
 if (y) return ccw(p1, p2, p3) == 0;
  return ccw(p1, p2, p3) * ccw(p1, p2, p4) <= 0 &&
    ccw(p3, p4, p1) * ccw(p3, p4, p2) <= 0;
bool lineInsideRectangle (double x1, double x2, double y1,
    double y2, point st, point ed) {
 if (x2 < x1) swap(x1, x2);
 if (y2 < y1) swap(y1, y2);
  double mnX = min(st.X, ed.X), mxX = max(st.X, ed.X),
   mnY = min(st.Y, ed.Y), mxY = (st.Y, ed.Y);
  return dcmp(x1, mnX) \le 0 \&\& dcmp(x2, mxX) >= 0
 && dcmp(v1, mnY) \le 0 && dcmp(v2, mxY) >= 0;
```

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5.2 Circles

Circles.h

```
"point.h", "triangles.h" //getAngle_A_abc, "lines.h" //intersectSegments bb9f4a. 72 lines
// 2 points has infinite circles
// Find circle passes with 3 points, some times, there is no
     circle! (in case colinear)
// Draw two perpendicular lines and intersect them
pair<double, point> findCircle(point a, point b, point c) {
  //create median, vector, its prependicular
  point m1 = (b + a) * 0.5, v1 = b - a, pv1 = point(v1.Y, -v1.X)
  point m2 = (b + c) * 0.5, v2 = b - c, pv2 = point(v2.Y, -v2.X)
  point end1 = m1 + pv1, end2 = m2 + pv2, center;
  intersectSegments(m1, end1, m2, end2, center);
  return make_pair(length(vec(center, a)), center);
// If line intersect circle at point p, and p = p0 + t(p1-p0)
// Then (p-c)(p-c) = r^2 substitute p and rearrange
//(p1-p0)(p1-p0)t^2 + 2(p1-p0)(p0-C)t + (p0-C)(p0-C) = r*r; \rightarrow
      Quadratic
```

MinimumEnclosingCircle Polygon

```
vector<point> intersectLineCircle(point p0, point p1, point C,
    double r) {
  double a = dp(vec(p0, p1), vec(p0, p1)),
   b = 2 * dp(vec(p0, p1), vec(C, p0)),
   c = dp(vec(C, p0), vec(C, p0)) - r * r;
  double f = b * b - 4 * a * c;
  vector<point> v;
  if (demp(f, 0) >= 0) {
   if (demp(f, 0) == 0) f = 0;
   double t1 = (-b + sqrt(f)) / (2 * a);
   double t2 = (-b - sqrt(f)) / (2 * a);
   v.push_back(p0 + t1 * (p1 - p0));
   if (dcmp(f, 0) != 0)
     v.push_back(p0 + t2 * (p1 - p0));
  return v;
vector<point> intersectCircleCircle(point c1, double r1, point
    c2, double r2) {
  // Handle infinity case first: same center/radius and r > 0
  if (same(c1, c2) && dcmp(r1, r2) == 0 && dcmp(r1, 0) > 0)
    return vector<point>(3, c1); // infinity 2 same circles
        (not points)
    // Compute 2 intersection case and handle 0, 1, 2 cases
  double ang1 = angle(vec(c1, c2)),
    ang2 = getAngle_A_abc(r2, r1, length(vec(c1, c2)));
  if (::isnan(ang2)) // if r1 or d = 0 \Rightarrow nan in qetAngle\_A\_abc
    ang2 = 0; // fix corruption
  vector<point> v(1, polar(r1, ang1 + ang2) + c1);
  // if point NOT on the 2 circles = no intersection
  if (dcmp(dp(vec(c1, v[0]), vec(c1, v[0])), r1 * r1) != 0 ||
   dcmp(dp(vec(c2, v[0]), vec(c2, v[0])), r2 * r2) != 0)
    return vector<point>();
  v.push_back(polar(r1, ang1 - ang2) + c1);
  if (same(v[0], v[1])) // if same, then 1 intersection only
    v.pop_back();
  return v;
ld circleCircleIntersectionArea(point cen1, ld r1, point cen2,
  ld dis = hypot(cen1.X - cen2.X, cen1.Y - cen2.Y);
  if (dis > r1 + r2) return 0;
  if (dis <= fabs(r2 - r1) && r1 >= r2)
   return PI * r2 * r2;
  if (dis <= fabs(r2 - r1) && r1 < r2)</pre>
   return PI * r1 * r1;
  1d \ a = r1 * r1, \ b = r2 * r2;
  ld ang1 = acos((a + dis * dis - b) / (2 * r1 * dis)) * 2;
  1d \ ang2 = acos((b + dis * dis - a) / (2 * r2 * dis)) * 2;
 1d ret1 = .5 * b * (ang2 - sin(ang2));
 1d ret2 = .5 * a * (ang1 - sin(ang1));
  return ret1 + ret2;
```

MinimumEnclosingCircle.h

7fd612, 71 lines

```
//init p array with the points and ps with the number of points
//cen and rad are result circle
//you must call random_shuffle(p,p+ps); before you call mec
typedef complex<double> point;
#define perp(a) (point(-(a).Y, (a).X))
```

```
#define vec(a,b) ((b) - (a))
#define mid(a,b) (((a) + (b)) / point(2, 0))
enum STATE {
    IN, OUT, BOUNDRY
STATE circlePoint(const point &cen, const double &r, const
    double lensgr = lengthSqr(vec(cen,p));
    if (fabs(lensqr - r * r) < EPS)
        return BOUNDRY;
    if (lensqr < r * r)
        return IN;
    return OUT;
void circle2(const point &p1, const point &p2, point &cen,
     double &r) {
    cen = mid(p1, p2);
    r = length(vec(p1, p2)) / 2;
bool circle3 (const point &p1, const point &p2, const point &p3,
             point& cen, double& r) {
    point m1 = mid(p1, p2);
    point m2 = mid(p2, p3);
    point perp1 = perp(vec(p1, p2));
    point perp2 = perp(vec(p2, p3));
    bool res = intersect(m1, m1 + perp1, m2, m2 + perp2, cen);
    r = length(vec(cen, p1));
    return res;
#define MAXPOINTS 100000
point p[MAXPOINTS], r[3], cen;
int ps, rs;
double rad;
//init p array with the points and ps with the number of points
//cen and rad are result circle
//you must call random_shuffle(p,p+ps); before you call mec
void mec() {
    if (rs == 3) {
        circle3(r[0], r[1], r[2], cen, rad);
        return;
    if (rs == 2 && ps == 0) {
        circle2(r[0], r[1], cen, rad);
        return;
    if (!ps) {
        cen = r[0];
        rad = 0;
        return;
    ps--;
    if (circlePoint(cen, rad, p[ps]) == OUT) {
        r[rs++] = p[ps];
        mec();
        rs--:
    ps++;
```

| 5.3 Polygons

Polygon.h

```
"Lines.h" //isPointOnSegment, CCW
                                                     49191e, 118 lines
struct cmp {
 point about;
  cmp(point c) {
    about = c;
 bool operator()(const point &p, const point &q) const {
    double cr = cp (vec(about, p), vec(about, q));
    if (fabs(cr) < EPS)
      return make_pair(p.Y, p.X) < make_pair(q.Y, q.X);</pre>
    return cr < 0;
};
void sortAntiClockWise(vector<point> &pnts) {
 point mn = pnts[0];
  for (int i = 0; i < sz(pnts); i++)</pre>
    if (make_pair(pnts[i].Y, pnts[i].X) < make_pair(mn.Y, mn.X)</pre>
      mn = pnts[i];
  sort(all(pnts), cmp(mn));
// CCW function must return 0 if the 3 points are collinear
bool isConvex(vector<point>& v) {
 int n = v.size(), m = n, sum = 0;
 v.push_back(v[0]);
 v.push_back(v[1]);
  char tmp;
  for (int i = 0; i < n; i++) {</pre>
    tmp = ccw(v[i], v[i + 1], v[i + 2]);
    if (tmp) sum += tmp;
    else m--:
 v.pop_back();
 v.pop_back();
  return abs(sum) == m;
//Area(p) = interal\_points + (boundry\_point/2) - 1
//2*interal\_points = 2*Area(p) - 2*(boundry\_point/2) + 2
ll picksTheorm(vector<point> &p) { //point = complex<int>
 11 \text{ area} = 0;
 11 \text{ bound} = 0;
 for (int i = 0; i < sz(p); i++) {
   int j = (i + 1 < sz(p) ? i + 1 : 0);
   area += cp(p[i], p[j]);
    point v = vec(p[i], p[j]);
    bound += abs(\_gcd(v.X, v.Y));
 return (abs(area) - 2 * (bound / 2) + 2) / 2;
bool pointInPolygon(const vector<point> &p, point p0) {
 int wn = 0; // the winding number counter
 for (int i = 0; i < sz(p); i++) {</pre>
    point cur = p[i], nxt = p[(i + 1) % sz(p)];
    if (isPointOnSegment(cur, nxt, p0))
      return true;
    if (cur.Y <= p0.Y) { // Upward edge
      if (nxt.Y > p0.Y && cp(nxt - cur, p0 - cur) > EPS)
                             // Downward edge
    } else {
      if (nxt.Y <= p0.Y && cp(nxt - cur, p0 - cur) < -EPS)</pre>
```

```
return wn != 0;
//to check if the points are sorted anti-clockwise or clockwise
//remove the fabs at the end and it will return -ve value if
    clockwise
double polygonArea(const vector<point> &p) {
 double res = 0;
  for (int i = 0; i < sz(p); i++) {</pre>
     int j = (i + 1) % sz(p);
     res += cp(p[i],p[j]);
  return fabs(res) / 2;
// return the centroid point of the polygon
// The centroid is also known as the "centre of gravity" or the
      "center of mass". The position of the centroid
// assuming the polygon to be made of a material of uniform
    density.
point polygonCentroid(vector<point> &polygon) {
   point res(0, 0);
    double a = 0;
    for (int i = 0; i < (int) polygon.size(); i++) {</pre>
       int j = (i + 1) % polygon.size();
        res = res + (polygon[i] + polygon[j]) * cp(polygon[i],
            polygon[j]);
       a += cp(polygon[i], polygon[j]);
    return res / 3.0 / a;
// P need to be counterclockwise convex polygon
pair<vector<point>, vector<point>> polygonCut (vector<point> &p,
    point A, point B) {
    vector<point> left, right;
   point intersect;
    for (int i = 0; i < sz(p); ++i) {
       point cur = p[i], nxt = p[(i + 1) % sz(p)];
   bool in1 = cp(B-A, cur-A) >= 0;
   bool in2 = cp(B-A, nxt-A) >= 0;
       if (in1) right.push_back(cur);
        //NOTE adust intersectSegments should handled AB as
        if (intersectSegments(A, B, cur, nxt, intersect)) {
            right.push back(intersect);
            left.push back(intersect);
       if (in2)left.push_back(cur);
    return make_pair(left, right);
ConvexHull.h
```

e39775, 48 lines

```
bool cmp (point a, point b) {
  return a.X < b.X || (a.X == b.X && a.Y < b.Y);</pre>
11 cross(point a, point b, point c) {
  return cp (vec(a,b), vec(a,c));
```

```
bool cw(point a, point b, point c) {
  return cp(vec(a,b), vec(b,c)) < 0;
bool ccw(point a, point b, point c) {
  return cp(vec(a,b), vec(b,c)) > 0;
//with collinear points, to remove collinears check if cross ==
      0 when pop
vector<point> convex_hull(vector<point> &p) {
 if (p.size() == 1)
    return p;
  sort(p.begin(), p.end(), &cmp);
  point p1 = p[0], p2 = p.back();
  vector<point> up, down;
  up.push_back(p1);
  down.push_back(p1);
  for (int i = 1; i < (int) p.size(); i++) {</pre>
    if (i == p.size() - 1 || cw(p1, p[i], p2)) {
      while (up.size() >= 2
          && !cw(up[up.size() - 2], up[up.size() - 1], p[i]))
        up.pop_back();
      up.push_back(p[i]);
    if (i == p.size() - 1 || ccw(p1, p[i], p2)) {
      while (down.size() >= 2
          && !ccw(down[down.size() - 2], down[down.size() - 1],
               p[i]))
        down.pop_back();
      down.push_back(p[i]);
  vector<point> convex;
  for (int i = 0; i < (int) down.size(); i++)</pre>
    convex.push back(down[i]);
  for (int i = up.size() - 2; i > 0; i--)
    convex.push_back(up[i]);
  return convex;
PointInPolygon.h
                                                     100ac7, 38 lines
11 cross(point a, point b, point c) {
  return cp(vec(a,b), vec(a,c));
void prepare (vector<point> &polygon) {
  int pos = 0;
  for (int i = 0; i < sz(polygon); i++) {</pre>
    if (make_pair(polygon[i].X, polygon[i].Y)
        < make_pair(polygon[pos].X, polygon[pos].Y))
  rotate(polygon.begin(), polygon.begin() + pos, polygon.end())
bool isPointOnSegment(point a, point b, point c) {
  double acb = length(a - b), ac = length(a - c), cb = length(b)
  return demp(acb - (ac + cb), 0) == 0;
bool pointInConvexPolygon(const vector<point> &polygon, point
```

```
if (isPointOnSegment(polygon[0], polygon.back(), pt))
  return true;
if (cross(polygon[0], polygon.back(), pt) > 0)
  return false;
if (cross(polygon[0], polygon[1], pt) < 0)</pre>
  return false;
if (polygon.size() == 2)
  return false;
int st = 2, ed = sz(polygon) - 2, ans = 1;
while (st <= ed) {</pre>
  int md = st + ed >> 1;
  if (cross(polygon[0], polygon[md], pt) >= 0)
    st = md + 1, ans = md;
  else
    ed = md - 1;
return cross(polygon[ans], polygon[ans + 1], pt) >= 0;
```

Misc (6)

LIS.h

8ca62c, 47 lines

```
//without build
//make upper_bound if can take equal elements
int LIS(const vector<int> &v) {
 vector<int> lis(v.size());//put value less than zero if
       needed
  int 1 = 0;
  for (int i = 0; i < sz(v); i++) {
    int idx = lower_bound(lis.begin(), lis.begin() + 1, v[i]) -
         lis.begin();
    if (idx == 1) 1++;
    lis[idx] = v[i];
 return 1;
void LIS_binarySearch(vector<int> v) {
 int n = v.size();
 vector<int> last(n), prev(n, -1);
 int length = 0;
 auto BS = [&](int val) {
    int st = 1, ed = length, md, rt = length;
    while (st <= ed) {</pre>
      md = st + ed >> 1;
      if (v[last[md]] >= val)
        ed = md - 1, rt = md;
      else
        st = md + 1;
    return rt;
  for (int i = 1; i < n; i++) {</pre>
    if (v[i] < v[last[0]])
     last[0] = i;
    else if (v[i] > v[last[length]]) {
      prev[i] = last[length];
      last[++length] = i;
    } else {
      int index = BS(v[i]);
      prev[i] = last[index - 1];
      last[index] = i;
  cout << length + 1 << "\n";
  vector<int> out;
  for (int i = last[length]; i >= 0; i = prev[i])
```

for (auto it : out)cout << it << endl;</pre>

static uint64_t splitmix64(uint64_t x) {

 $x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;$

 $x = (x ^ (x >> 27)) * 0x94d049bb133111eb;$

now().time_since_epoch().count();

now().time_since_epoch().count(); return splitmix64(x + FIXED RANDOM);

second + FIXED RANDOM) >> 1);

x += 0x9e3779b97f4a7c15;

return x ^ (x >> 31);

out.push_back(v[i]); reverse(out.begin(), out.end());

Random CustomHash

$S(0,n) + 2^{n+1} = 2^0 + (2^1 + 2^2 + \dots + 2^{n+1})$ $S(0,n) + 2^{n+1} = 1 + 2 \times S(0,n)$

$S(0,n) = 2^{n+1} - 1$

ax + by = e cx + dy = f $x = \frac{ed - bf}{ad - bc}$ $y = \frac{af - ec}{ad - bc}$ $\binom{n}{0}^2 + \binom{n}{1}^2 + \binom{n}{2}^2 + \ldots + \binom{n}{n}^2 = \binom{2n}{n}$

In general, given an equation Ax = b, the solution to a variable x_i is given by

$$x_i = \frac{\det A_i'}{\det A}$$

where A'_i is A with the i'th column replaced by b.

8.2 Recurrences

The extremum is given by x = -b/2a.

If $a_n = c_1 a_{n-1} + \cdots + c_k a_{n-k}$, and r_1, \ldots, r_k are distinct roots of $x^k - c_1 x^{k-1} - \cdots - c_k$, there are d_1, \ldots, d_k s.t.

$$a_n = d_1 r_1^n + \dots + d_k r_k^n.$$

Non-distinct roots r become polynomial factors, e.g. $a_n = (d_1 n + d_2)r^n.$

Trigonometry

$$\sin(v+w) = \sin v \cos w + \cos v \sin w$$
$$\cos(v+w) = \cos v \cos w - \sin v \sin w$$

$$\tan(v+w) = \frac{\tan v + \tan w}{1 - \tan v \tan w}$$
$$\sin v + \sin w = 2\sin\frac{v+w}{2}\cos\frac{v-w}{2}$$
$$\cos v + \cos w = 2\cos\frac{v+w}{2}\cos\frac{v-w}{2}$$

$$(V+W)\tan(v-w)/2 = (V-W)\tan(v+w)/2$$

where V, W are lengths of sides opposite angles v, w.

$$a\cos x + b\sin x = r\cos(x - \phi)$$
$$a\sin x + b\cos x = r\sin(x + \phi)$$

where $r = \sqrt{a^2 + b^2}$, $\phi = \operatorname{atan2}(b, a)$.

CustomHash.h

struct custom hash {

Random.h f7aba2, 9 lines

size_t operator()(pair<uint64_t, uint64_t> x) const { // for

static const uint64_t FIXED_RANDOM = chrono::steady_clock::

return splitmix64(x.first + FIXED RANDOM) ^ (splitmix64(x.

size_t operator()(uint64_t x) const { // for single element

static const uint64 t FIXED RANDOM = chrono::steady clock::

Congruence

f21ae4, 17 lines

- IF $ax = by \mod n$, THEN $x = y \mod n$ IFF gcd(a, n) = 1(Division Rule)
- IF $a = b \mod m$, THEN $a^n = b^n \mod m$ (Powers Rule)
- $x^n = x^{n\%phi(m)} \mod m \text{ IFF } n \ge \log_2 m... \text{ otherwise}$ calculate without a modulo (Euler Totient)
- ALL nC_r for $n = \text{prime}, r \neq 0, r \neq n$ is divisible by n (nC_r Rule)
- Distribution is ok IFF p is prime: $(x+y)^p = x^p + y^p$ $\mod p$ (Powers Distribution Rule)
- IF $a = b \mod m$, THEN $a^{m^k} = b^{m^k}$ $\mod m^{k+1}; k = 0, 1, 2, 3, \dots$

KACTL Maths Notes (8)

$$\phi(n!) = (n - isPrime(n)) \times \phi((n-1)!)$$

$$\sum_{d|n} \phi(d) = n$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} gcd(i,j) = \sum_{d=1}^{n} \sum_{i=1}^{n} \sum_{j=1}^{n} \phi(d) \ [d|i] \ [d|j]$$

$$= \sum_{d=1}^n \phi(d) \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{n}{d} \rfloor} 1 = \sum_{d=1}^n \phi(d) \lfloor \frac{n}{d} \rfloor^2$$

7.3 $E(x^2)$

8.1 Equations

- To get the expected value of x^2 , you need to represent the random variable as a **length** variable \rightarrow get the product of probabilities of every pair \rightarrow Convert every loop to an equation.
- Some problems require going through all pairs using DP (if possible) \rightarrow vou can compute the result for every individual variable separately like the problem of bears and humans.

 $ax^2 + bx + c = 0 \Rightarrow x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$

<u>Mathematics Notes</u> (7)

7.1 Sums

};

- To get the summation of a^n sequence from x to y, transform a^n form to $a^{\{n\}}$ form (factorial polynomial $= a \times (a-1) \times (a-2) \times \ldots \times (a-n+1)$, then integrate the equation with limits from x to (y+1)
- To get the summation law from known values substitute and compute the equation:

$$S(0,n) + x[n+1] = x[0] + S(1,n+1)$$

, then take a factor from S(1, n+1) to transform it to S(0,n), so you can substitute and compute S(0,n) easily. You can use replacement in the original formula and the bounds of the summation (like replacing every (k-i) by j). You can also multiply the summand by k - it might help. For example:

$$S(0,n) = 2^0 + 2^1 + \ldots + 2^n$$

8.4 Geometry

8.4.1 Triangles

Side lengths: a, b, c

Semiperimeter: $p = \frac{a+b+c}{2}$

Area: $A = \sqrt{p(p-a)(p-b)(p-c)}$

Circumradius: $R = \frac{abc}{4A}$

Inradius: $r = \frac{A}{1}$

Length of median (divides triangle into two equal-area triangles): $m_a = \frac{1}{2}\sqrt{2b^2 + 2c^2 - a^2}$

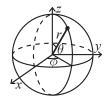
Length of bisector (divides angles in two):

$$s_a = \sqrt{bc \left[1 - \left(\frac{a}{b+c}\right)^2\right]}$$

Law of sines: $\frac{\sin \alpha}{a} = \frac{\sin \beta}{b} = \frac{\sin \gamma}{c} = \frac{1}{2R}$ Law of cosines: $a^2 = b^2 + c^2 - 2bc \cos \alpha$

Law of tangents: $\frac{a+b}{a-b} = \frac{\tan \frac{\alpha+\beta}{2}}{\tan \frac{\alpha-\beta}{2}}$

8.4.2 Spherical coordinates



$$\begin{aligned} x &= r \sin \theta \cos \phi & r &= \sqrt{x^2 + y^2 + z^2} \\ y &= r \sin \theta \sin \phi & \theta &= \arccos(z/\sqrt{x^2 + y^2 + z^2}) \\ z &= r \cos \theta & \phi &= \operatorname{atan2}(y, x) \end{aligned}$$

Derivatives/Integrals

$$\frac{d}{dx}\arcsin x = \frac{1}{\sqrt{1-x^2}} \qquad \frac{d}{dx}\arccos x = -\frac{1}{\sqrt{1-x^2}}$$

$$\frac{d}{dx}\tan x = 1 + \tan^2 x \qquad \frac{d}{dx}\arctan x = \frac{1}{1+x^2}$$

$$\int \tan ax = -\frac{\ln|\cos ax|}{a} \qquad \int x\sin ax = \frac{\sin ax - ax\cos ax}{a^2}$$

$$\int e^{-x^2} = \frac{\sqrt{\pi}}{2}\operatorname{erf}(x) \qquad \int xe^{ax}dx = \frac{e^{ax}}{a^2}(ax-1)$$

Integration by parts:

$$\int_{a}^{b} f(x)g(x)dx = [F(x)g(x)]_{a}^{b} - \int_{a}^{b} F(x)g'(x)dx$$

8.6 Sums

$$c^{a} + c^{a+1} + \dots + c^{b} = \frac{c^{b+1} - c^{a}}{c - 1}, c \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(2n+1)(n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \frac{n^{2}(n+1)^{2}}{4}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

8.7Series

$$e^{x} = 1 + x + \frac{x^{2}}{2!} + \frac{x^{3}}{3!} + \dots, (-\infty < x < \infty)$$

$$\ln(1+x) = x - \frac{x^{2}}{2} + \frac{x^{3}}{3} - \frac{x^{4}}{4} + \dots, (-1 < x \le 1)$$

$$\sqrt{1+x} = 1 + \frac{x}{2} - \frac{x^{2}}{8} + \frac{2x^{3}}{32} - \frac{5x^{4}}{128} + \dots, (-1 \le x \le 1)$$

$$\sin x = x - \frac{x^{3}}{3!} + \frac{x^{5}}{5!} - \frac{x^{7}}{7!} + \dots, (-\infty < x < \infty)$$

$$\cos x = 1 - \frac{x^{2}}{2!} + \frac{x^{4}}{4!} - \frac{x^{6}}{6!} + \dots, (-\infty < x < \infty)$$

Probability theory 8.8

Let X be a discrete random variable with probability $p_X(x)$ of assuming the value x. It will then have an expected value (mean) $\mu = \mathbb{E}(X) = \sum_{x} x p_X(x)$ and variance $\sigma^2 = V(X) = \mathbb{E}(X^2) - (\mathbb{E}(X))^2 = \sum_x (x - \mathbb{E}(X))^2 p_X(x)$ where σ is the standard deviation. If X is instead continuous it will have a probability density function $f_X(x)$ and the sums above will instead be integrals with $p_X(x)$ replaced by $f_X(x)$.

Expectation is linear:

$$\mathbb{E}(aX + bY) = a\mathbb{E}(X) + b\mathbb{E}(Y)$$

For independent X and Y,

$$V(aX + bY) = a^2V(X) + b^2V(Y).$$

8.8.1 Discrete distributions Binomial distribution

The number of successes in n independent yes/no experiments, each which yields success with probability p is $Bin(n, p), n = 1, 2, ..., 0 \le p \le 1.$

$$p(k) = \binom{n}{k} p^k (1-p)^{n-k}$$

$$\mu = np, \, \sigma^2 = np(1-p)$$

Bin(n, p) is approximately Po(np) for small p.

First success distribution

The number of trials needed to get the first success in independent yes/no experiments, each which yields success with probability p is Fs(p), $0 \le p \le 1$.

$$p(k) = p(1-p)^{k-1}, k = 1, 2, \dots$$

$$\mu = \frac{1}{p}, \sigma^2 = \frac{1-p}{p^2}$$

Poisson distribution

The number of events occurring in a fixed period of time t if these events occur with a known average rate κ and independently of the time since the last event is $Po(\lambda)$, $\lambda = t\kappa$.

$$p(k) = e^{-\lambda} \frac{\lambda^k}{k!}, k = 0, 1, 2, \dots$$
$$\mu = \lambda, \sigma^2 = \lambda$$

8.8.2 Continuous distributions Uniform distribution

If the probability density function is constant between a and b and 0 elsewhere it is U(a, b), a < b.

$$f(x) = \begin{cases} \frac{1}{b-a} & a < x < b \\ 0 & \text{otherwise} \end{cases}$$

$$\mu = \frac{a+b}{2}, \, \sigma^2 = \frac{(b-a)^2}{12}$$

Exponential distribution

The time between events in a Poisson process is $\operatorname{Exp}(\lambda), \lambda > 0.$

$$f(x) = \begin{cases} \lambda e^{-\lambda x} & x \ge 0\\ 0 & x < 0 \end{cases}$$
$$\mu = \frac{1}{\lambda}, \, \sigma^2 = \frac{1}{\lambda^2}$$

Normal distribution

Most real random values with mean μ and variance σ^2 are well described by $\mathcal{N}(\mu, \sigma^2)$, $\sigma > 0$.

$$f(x) = \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{(x-\mu)^2}{2\sigma^2}}$$

If $X_1 \sim \mathcal{N}(\mu_1, \sigma_1^2)$ and $X_2 \sim \mathcal{N}(\mu_2, \sigma_2^2)$ then

$$aX_1 + bX_2 + c \sim \mathcal{N}(\mu_1 + \mu_2 + c, a^2\sigma_1^2 + b^2\sigma_2^2)$$

thinking techniques

Appendix (A)

IO re-representation

Think in bipartite graph

thinking.txt

174 lines

```
Reading problem statement
1. Read carefully, make sure no statements conflicts.
2. Rewrite important details /mark it.
3. If text is not small, Re-read the problem statement again.
    Make sure you have the full picture.
4. Extract constraints info.
   Never ignore any constraints, especially unusual one (e.g.
        2*(a+b) < c).
    Sometimes constraints are not direct. Find triangle angle
        with 2 precision -> 360 * 10^2 //brute force
5. Trace Samples as long as they are traceable.
6. Think in Missed cases, smallest boundaries & largest
    boundaries & Especial cases.
7. Write it on paper (to test it in your idea).
8. don't make assumptions.
9. Revise carefully output section and output formats.
Investigate
Analysis
 Problem Constraints
 Problem domain(s)
  Search Space Size (size of unique solutions)
 Nature of target Function
  Output Bounding
Problem Abstraction
    never to drop the original problem, sometimes your
        abstraction drop some important domain consideration
Problem Simplification
    Problem to Sub-Problems: you are JUST solving a sub-problem
    Incrementally: think in a special problem/case, and then
         try to update the solution for general problem/case.
    Simplification by Assumptions
    think in general case
Problem Reverse
    f(x) = y , search (x)
    atmost vs exect
    Property Reverse
       probability(x) = 1 - probability(!x)
        Subset with x = total - (subset without x)
Problem Domain re-interpretation
Thinking
    Think on papers not in pc
    The Brute Force solution
        Think in BF ITERATIVE and RECURSIVE.
        Think about Search space & Search State
   Problem Domain re-interpretation
    Concretely, Symbolically, Pictorially
   Forward and backward
   Brainstorm - Rank - Approach
   Divide & Conquer problem
Observations Discovery
    used pc if better (SIMPLE code only, don't take a lot of
        time)
Some popular properties:
   number of states
    Symmetry
    Inference
    Redundancy
    Independency
```

```
DAG: topological sort
        Tree is a bipartite graph
    Canonical Form
    Cycle tricks
    Input Function Nature
    Tricks
        Inference value
        convert to table:
        Applied data structures
        Adhoc trick
    Dp optimizations (D&C, Knuth, Convex Hull)
    Patterns
    Cyclic Function (repeated after X step)
1. Try to re-state the problem definition in 2 or 3 different
     ways and see if this helps.
2. Still Stuck? Do BF & Observe. Write the impractical BF
     solution, and try to find a pattern for answer / useful
3. Still Stuck? Iterate on your algorithms list, see if it
     could be the solution.
4. Be careful in analysis for solutions. You may discard a
     correct solution
    E.g., Calculated O(n^3) although it is O(n^2logn)
    E.g., Calculated recursion depth wrongly.
    Check Exact # of operation
    Check for Reduced variables & Constrained input
         combinations
    Check duplicates and unique values
    SQRT tricks
    Preprocessing
        Think in all kinds of precomputation and try to utilize
              anv of them.
        Next array
        Calc on machine a small temp array that help you in run
    In guery problems, solve them offline if that help
    Order of loops and data such that loops break so early as
         much as possible.
    Reference of locality tricks
        Order loops such that: loops don't pass over arrays is
            in depth. Watch out from Col-order accessing
        Maybe duplicate some memory to switch some col order to
Solution Verification
But before you verify, you should remember (keep it simple)!
    Could we find something simpler!
verifying the solution requires:
1. Test cases Verification (Ones in statement, boundaries,
     vours)
2. Solution Order: Time & Memory
3. Full logic & intuitive Revision
4. Is valid (BS/TS)?
5. Correctness, at least good intuitive - Assumption's
    validations
6. In case recursive code, does depth fit?
7. In case (*+^) operations, Any overflow (intermediate &
     output)
8. In case Double /, Is precision fine? Using long double
    instead of double?
9. In case Double /, can we not use it??
After Implementation
1. Revise code order & logic. Make sure it matches what you
     intended
```

```
2. Challenge every block of code. Never to read in a way that
    drop even ONE character
3. Think how this block of code maybe fail.
4. Revise data types
5. Double comparison, precision of +- numbers [(int)(a +/- EPS)
6. Return statement in functions
Test special cases. Testing the boundaries + revise SPECIAL
Failed? Check Error Inspection List
Error Inspection
General
 Do you read all input file?
  Initialization (between test cases)
  TYPO, variable names
  Conditions \function base case\return statement in function
  Overflow
  avoid double operations if possible
  Corner cases
 Arrays boundary
Wrong Answers
 Review constrains
  Review code again
  Re-read the problem statement
   Tricky text description
  Geometry
    Double precision
    Are there duplicate points? Does it matter? Co-linearity?
    Polygon: convex? concave?
    Connected or disconnected?
    Directed or Undirected?
    Self Loops?
    Multiple edges
Precision
    Watchout -0.0
    int x = (int)(a +/- EPS) depends on a > 0 \mid a < 0.
Time Limit Exceed
    May be bug and just infinite loop
    Can we precompute the results?
    Function calls may need reference variables.
    % is used extensively?
        If mod is 2^p-1, use bitwise
    What is blocks of code that represent order? Do we just
         need to optimize it?
    Big Input file
        Need scanf & printf?
        Optimize code operations
        Switch to arrays and char[]
        Do you really need to clear each time?
        The base case order is not O(1)
        Use effective base conditions E.g. If you are sure Dp
             (0, M) is X, do not wait until Dp(0,0)
    Cyclic recurrence?
    Backtracking If you have different ways to do it, try to do
          what minimize stack depth
Run time error
    Make sure to have correct array size.
    Make sure no wrong indexing < 0 \mid \mid x >= n
    In DP, check you access dimensions correctly
    Stack overflow
    Using incorrect compare function (e.g. return that return (
         A, B) same answer as (B, A))
```

techniques.txt

09 lines

Data structures STL bitmask bitset monoqueue BIT update range 2d Persistent Segment tree lazy dynamic merge sort persistent wavelet tree sparse table SORT Decomposition ordered set DSU Treaps Strings KMP trie aho-corasick suffix array rolling hash suffix automaton Recursion Divide and conquer Greedy algorithm Graph theory DFS BFS bipartite graph Topological sorting Strongly connected components Aritculation Points and Bridges biconnected components 2-SAT MST Prim Kruskal bfs dijkstra Bellman-Ford's floyd Flow Augmenting paths Edmonds-Karp Bipartite matching Dinic Min-cost max flow Hopcroft-Karp Euler cycles Trees bipartite graph Diameter DP on tress SACK Data structues LCA HLD Centroid Decomposition Virtual trees Dynamic programming Knapsack

Knapsack SORT trick Coin change Longest common subsequence Longest increasing subsequence Number of paths in a dag Shortest path in a dag Max 1D, 2D Range Sum DP consecutive Ranges DP nested Ranges Push loop to paramter Push parameter to loop DP General Ranges Counting Building output DP Bitmask (TSP) Bitonic cycle DP on trees DP SOS (sum of submasks) Speed-up: Convex Hull Optimization Speed-up: Divide and Conquer Optimization Speed-up: Knuth Optimization Techniques Binary search Ternary search Sliding Window State-Space Search Meet-in-the-middle Backtracking Binary Search ternary search divide and conquer Matrix power Game theory Nim Grundy numbers Mini-max Alpha-beta pruning cycles and patterns mirror strategy Combinatorics Inclusion/exclusion Catalan number