Lecture 2

Introduction to distributed databases system

Generic DBMS Architecture

- The functions performed by a DBMS can be layered.
- Taking a top-down approach, the layers are the **interface**,

control, compilation, execution, data access, and consistency management.

1. The interface layer

manages the interface to the applications.

E.g. SQL statement in RDBMS.

2. The control layer

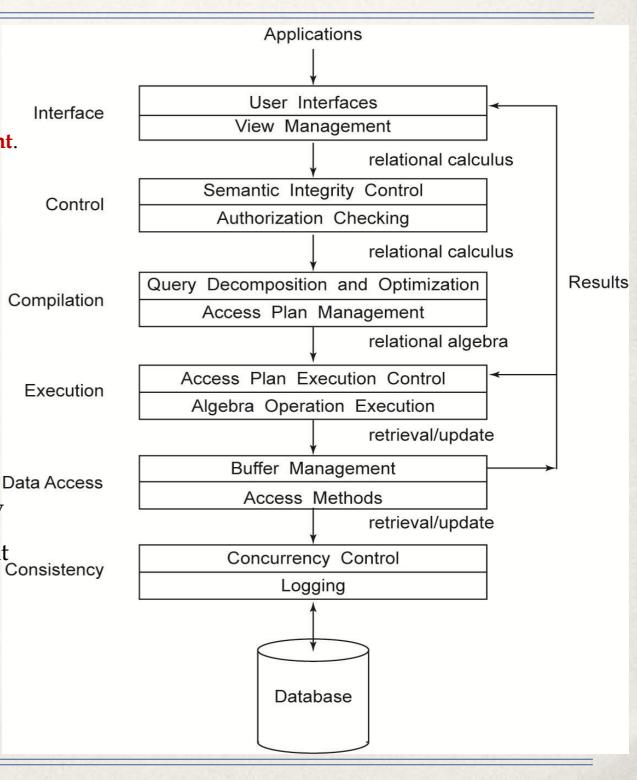
semantic integrity control:

checks integrity constraints as unique key constraints, and Foreign key constraint.

 Semantic integrity control ensures database consistency by rejecting update transactions that lead to inconsistent Cons database states

Authorization Checking

Check whether the user can execute the operation on the object.



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Generic DBMS Architecture (cont)

3. The query processing (or compilation) layer

- maps the query into an optimized sequence of lower-level operations.
- > The output of this layer is a query expressed in lower-level code (algebra operations).
- > This layer is concerned with performance.

The execution layer

- directs the execution of the access plans, including transaction management (commit, restart) and synchronization of algebra operations.
- It interprets the relational operations by calling the data access layer through the retrieval and update requests.

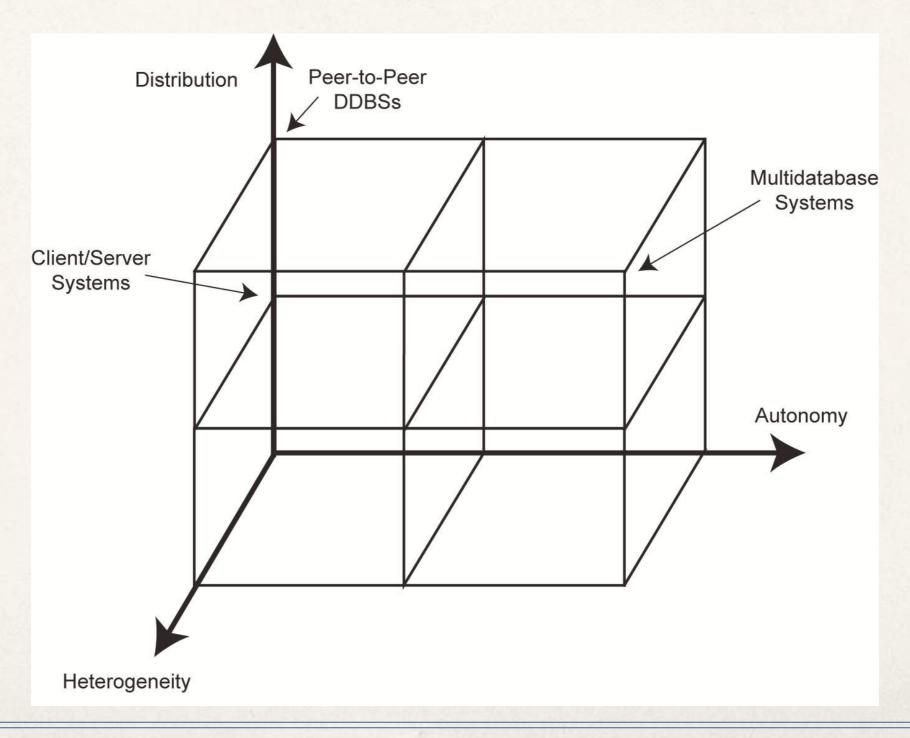
• The data access layer:

- manages the data structures that implement the files, indices, etc.
- → It also manages the buffers by caching the most frequently accessed data.
- Careful use of this layer minimizes the access to disks to get or write data.

Finally, the consistency layer:

- manages concurrency control and logging for update requests.
- → This layer allows transaction, system, and media recovery after failure.

DBMS Implementation Alternatives



Dimensions of the Problem

- Autonomy: Requirements of an autonomous system have been specified in D-DBMS as follows:
 - → The local operations of the individual DBMSs are not affected by their participation.
 - → The manner in which the individual DBMSs process queries and optimize them should not be affected by the execution of global queries in the distributed system.
 - System consistency or operation should not be compromised when individual DBMSs join or leave the distributed system.
 - **→** Various versions of autonomy
 - ◆ Design autonomy: Ability of a component DBMS to decide on issues related to its own design.
 - ◆ Communication autonomy: Ability of a component DBMS to decide whether and how to communicate with other DBMSs.
 - ◆ Execution autonomy: Ability of a component DBMS to execute local operations in any manner it wants to.

Dimensions of the Problem

Distribution

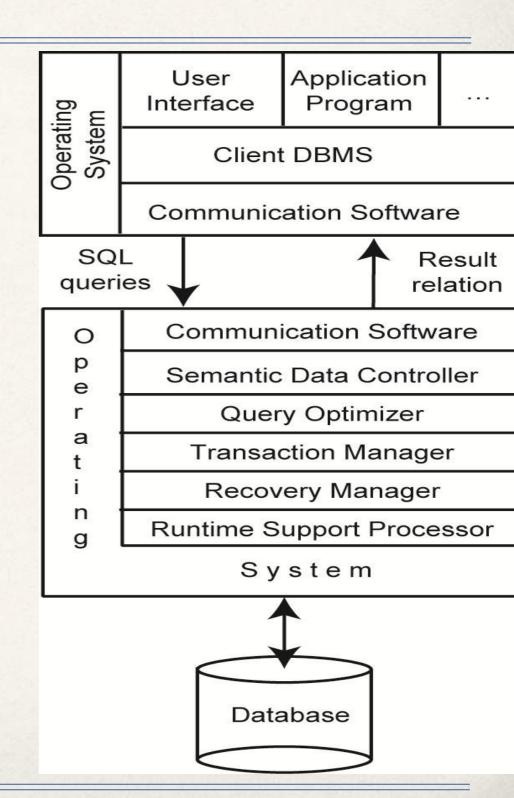
- → Whether the components of the system are located on the same machine or not and the physical distribution of data over multiple sites
 - ◆ The client/server distribution concentrates data management duties at servers while the clients focus on providing the application environment including the user interface.
 - ❖ In peer-to-peer systems, there is no distinction of client machines versus servers. Each machine has full DBMS functionality and can communicate with other machines to execute queries and transactions.

Heterogeneity

- → Various levels (hardware, Networking protocols, operating system)
- DBMS important one
 - data model, query language, transaction management algorithms

Client/Server Architecture

- the client passes SQL queries to the server without trying to understand or optimize them. The server does most of the work and returns the result relation to the client.
- DBMS client module is responsible for
 - managing the data that is cached to the client
 - managing the transaction locks



Multiple client/single server approach

Client servers

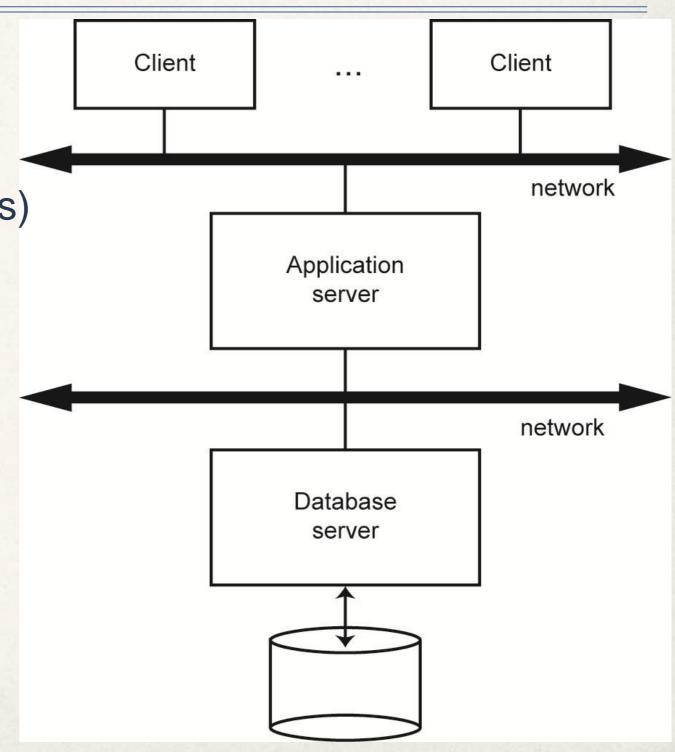
run the user interface(e.g. web servers)

Application servers

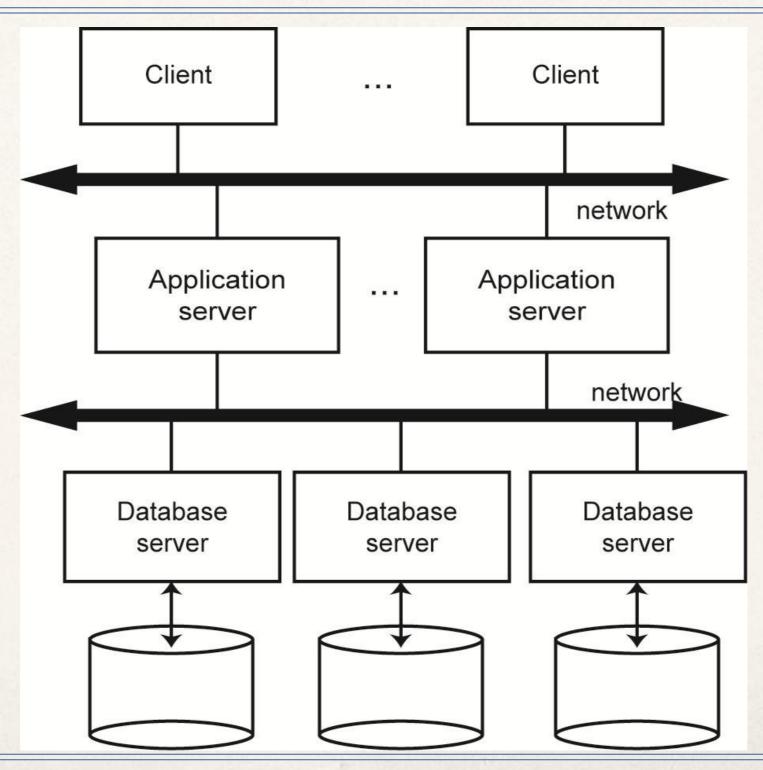
run application programs.

Database servers

run database management functions.

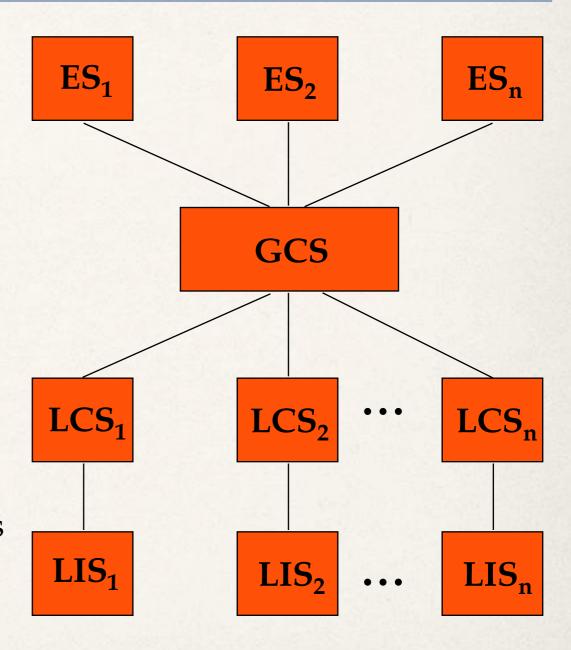


Multiple client/multiple server approach



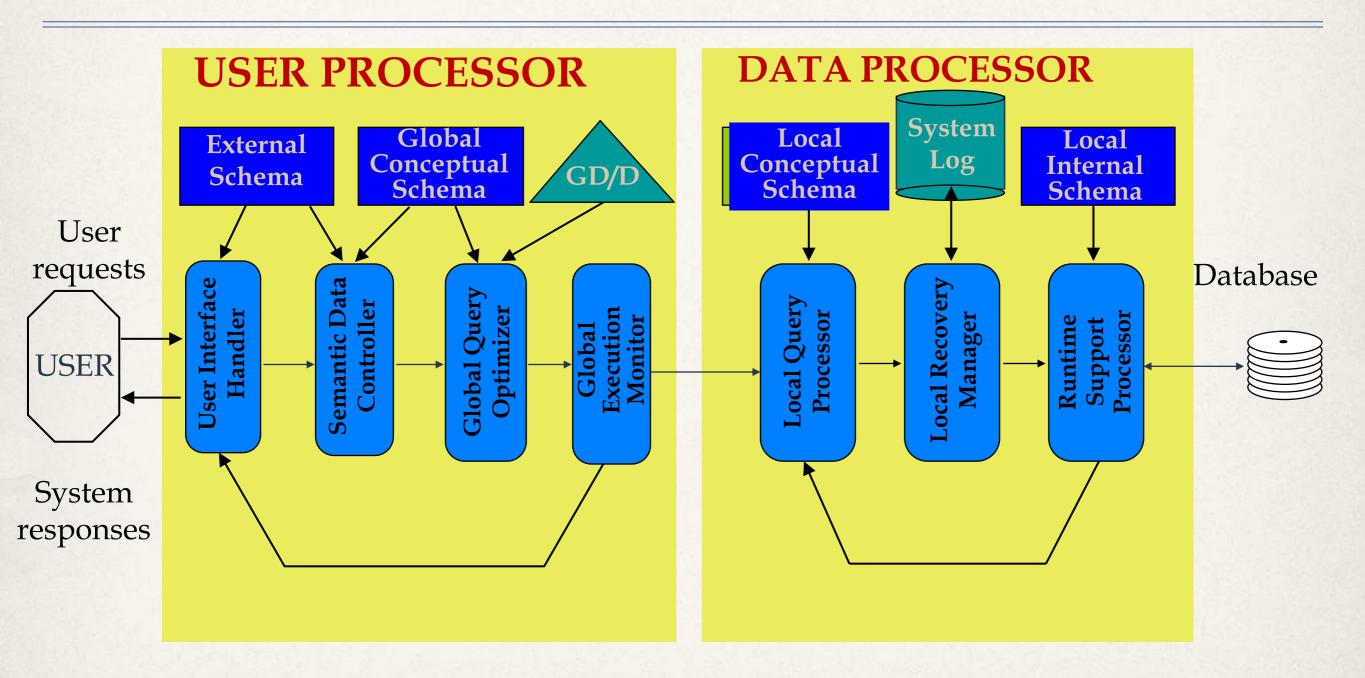
Distributed DBMS based peer- to peer Architecture

- local internal schema (LIS).
 - → Internal schema definition at each site
- conceptual schema (GCS).
 - → The enterprise view of the data.
- local conceptual schema (LCS).
 - → handle data fragmentation and replication at each site
- external schemas (ESs).
 - Supports user applications and user access to the database



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Peer-to-Peer Component Architecture



user processor consists of four elements

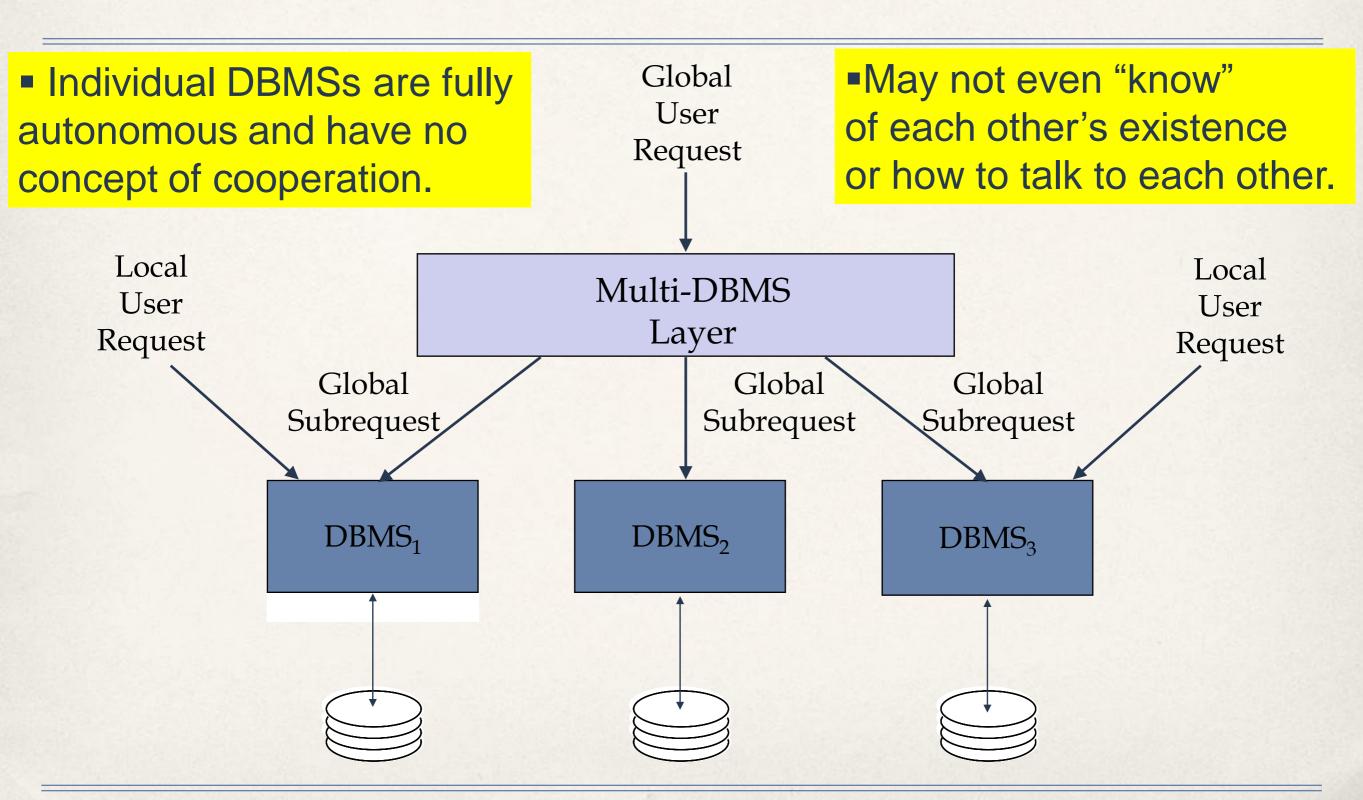
- 1.The user interface handler is responsible for interpreting user commands as they come in, and formatting the result data as it is sent to the user.
- 2. The semantic data controller uses the integrity constraints and authorizations that are defined as part of the global conceptual schema to check if the user query can be processed.
- 3. The global query optimizer and decomposer determine the best strategy to execute distributed join operations to minimize a cost function.
- 4. The distributed execution monitor coordinates the distributed execution of the user request. The execution monitor is also called the distributed transaction manager.

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data processor and consists of three elements:

- 1. The local query optimizer, acts as the access path selector, is responsible
 for choosing the best access path to access any data item
- 2. The local recovery manager is responsible for making sure that the local database remains consistent even when failures occur.
- 3. The run-time support processor
 - physically accesses the database according to the physical commands in the schedule generated by the query optimizer.
 - → is the interface to the operating system and contains the database buffer (or cache) manager, which is responsible for maintaining the main memory buffers and managing the data accesses.

MDBS Components & Execution



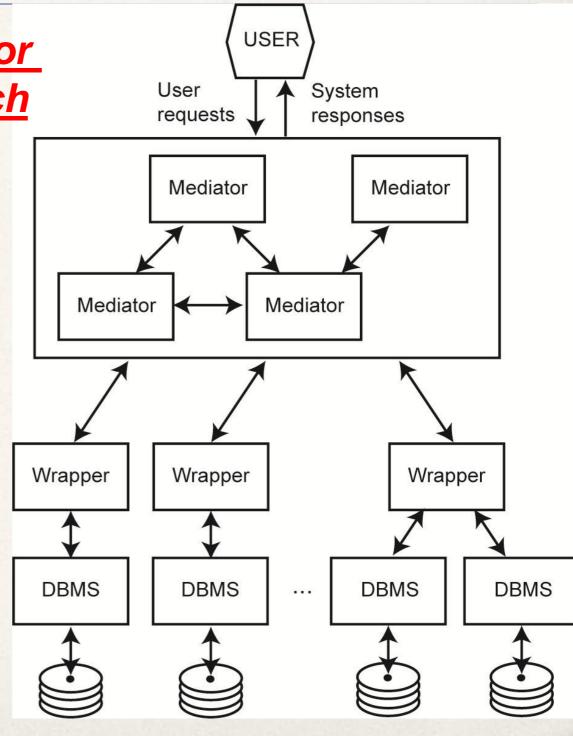
Mediator/Wrapper Architecture

A popular implementation architecture for MDBSs is the mediator/wrapper approach

A mediator "is a software module that exploits encoded knowledge about certain Sets data to create information for a higher layer of applications.

wrappers provide a mapping between a source DBMSs view and the mediators' For example,

if the source DBMS is a relational one, but the mediator implementations are object-oriented, the required mappings are established by the wrappers.



Questions

