

*Pharos University in Alexandria*  
*Faculty of Computer Science & Artificial Intelligence*  
*Course Title: Theory of Computation*  
*Code: CS 307*



# Theory of Computation

**Lecturer: Sherine Shawky**

Text Books

1. Introduction to formal languages and automata, Peter Linz, 6th edition, 2017.

Week 2

# Finite Automata and Regular Languages

# Finite Automata

- Finite automata are good models for computers with an extremely limited amount of memory.
- A *finite automaton* is a 5-tuple  $(Q, \Sigma, \delta, q_0, F)$ ,

where

$Q$  is a finite set called the *states*,

$\Sigma$  is a finite set called the *alphabet*,

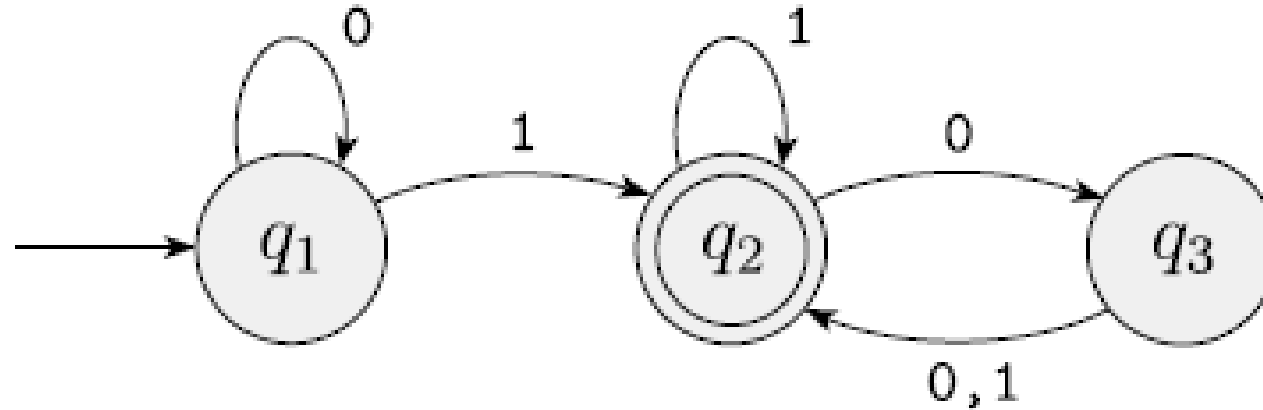
$\delta : Q \times \Sigma \rightarrow Q$  is the *transition function*

$q_0 \in Q$  is the *start state*, and

$F \subseteq Q$  is the *set of accept states*.

# Finite Automata

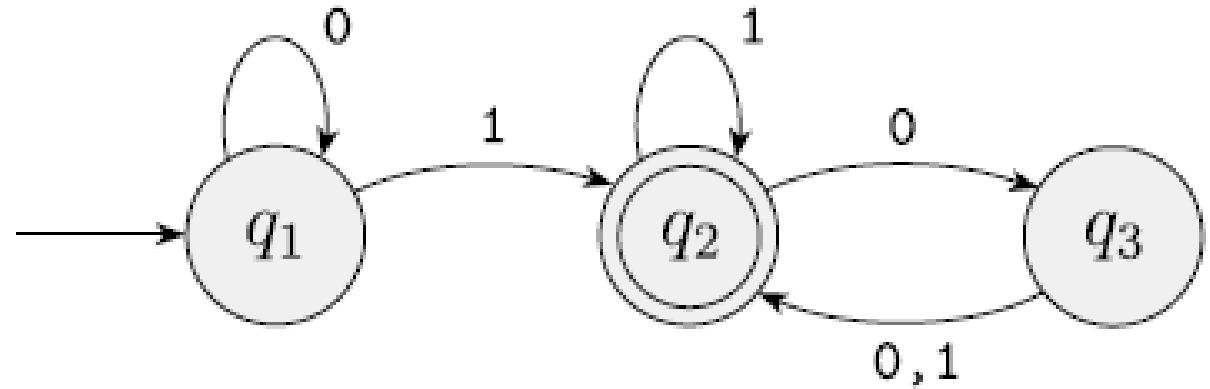
- Shown figure called the *state diagram* of M1.
- Example, with input string 1101 into M1, find finite automata parameters.



# Finite Automata

- $M1 = (Q, \Sigma, \delta, q_1, F)$ , where  
     $Q = \{q_1, q_2, q_3\}$ ,  
     $\Sigma = \{0, 1\}$ ,  
     $\delta$  is described as

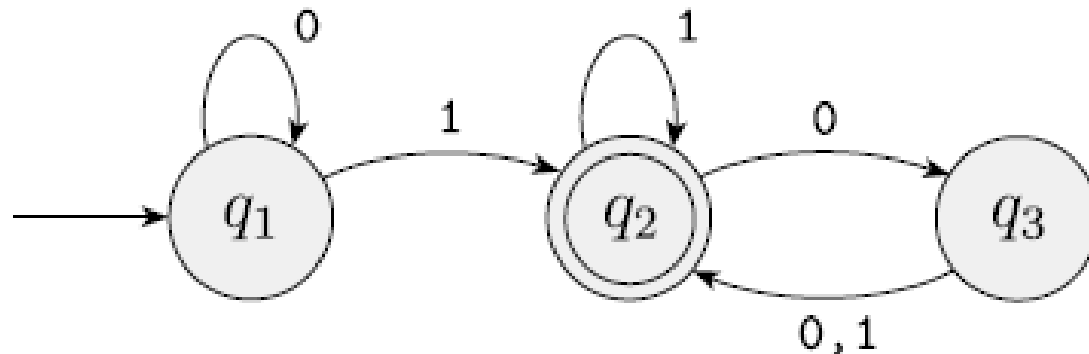
	0	1
$q_1$	$q_1$	$q_2$
$q_2$	$q_3$	$q_2$
$q_3$	$q_2$	$q_2$



# Finite Automata

$q_1$  is the start state, and  
 $F = \{q_2\}$ .

- In our example, let  $A = \{w \mid w \text{ contains at least one } 1 \text{ and an even number of } 0\text{s follow the last } 1\}$ .
- Then  $L(M_1) = A$ , or equivalently,  $M_1$  recognizes  $A$ .

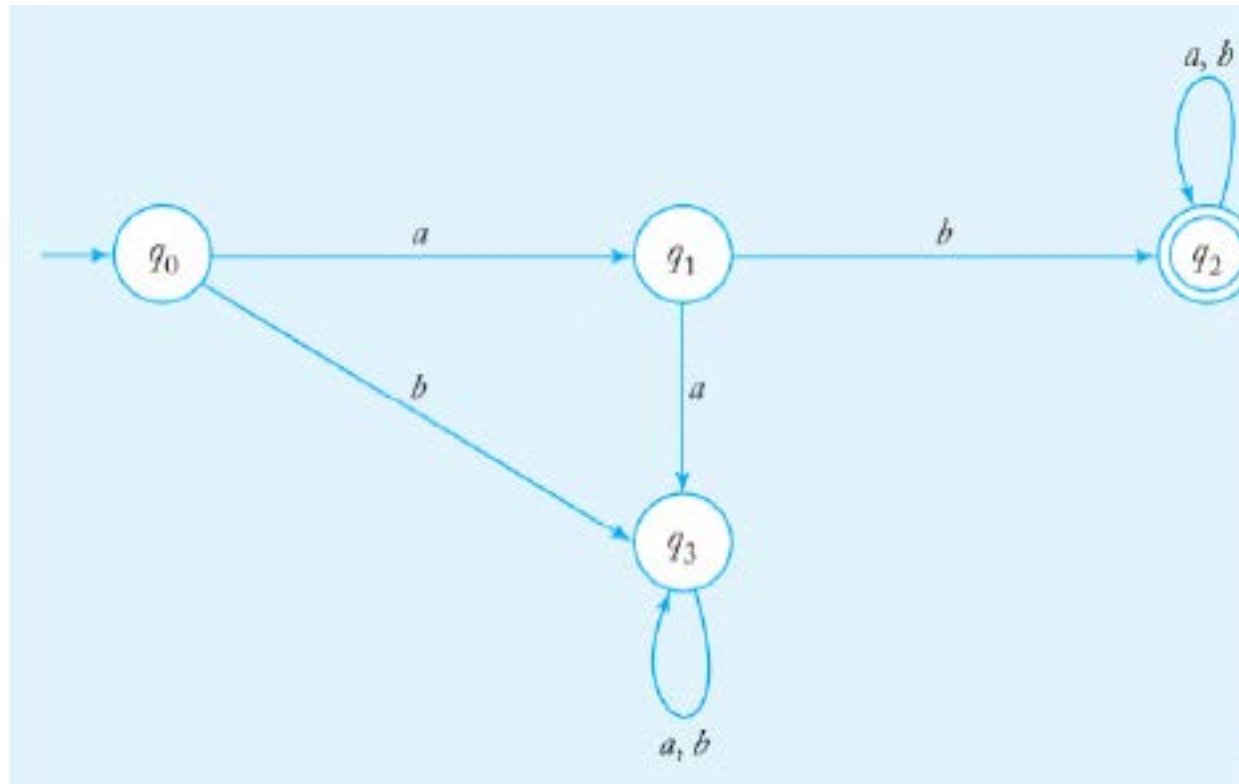


# Finite Automata

- Definition 2: the language is the set of all the strings accepted by the automaton.
- The language accepted by DFA  $M = (Q, \Sigma, \delta, q_0, F)$  is the set of all strings on  $\Sigma$  accepted by  $M$ .  
$$L(M) = \{w \in \Sigma^* : \delta^*(q_0, w) \in F\}.$$
- $\delta^*$ , be total functions. At each step, a unique move is defined, so that we are justified in calling such an automaton deterministic.
- DFA will process every string in  $\Sigma^*$  and either accept it or not accept it. Non-acceptance means that the DFA stops in a non-final state, so that  
$$L(M) = \{w \in \Sigma^* : \delta^*(q_0, w) \notin F\}.$$

# DFA Problem

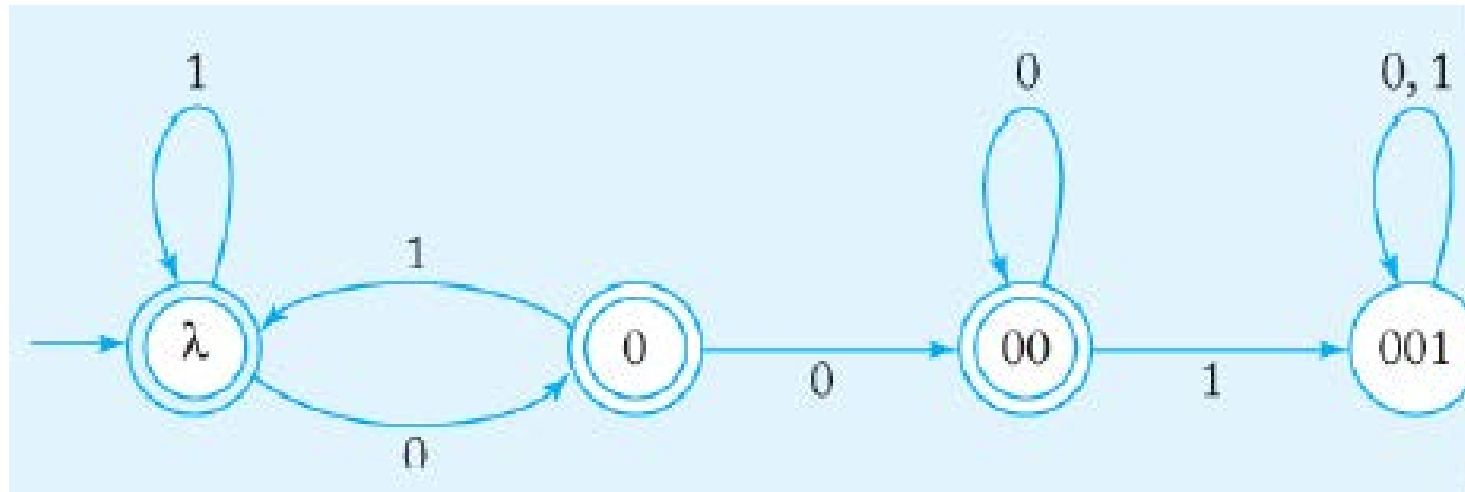
- Find a deterministic finite accepter that recognizes the set of all strings on  $\Sigma = \{a, b\}$  starting with the prefix  $ab$ .





# DFA Problem

- Find a DFA that accepts all the string on  $\{0,1\}$ , except those containing the substring 001.

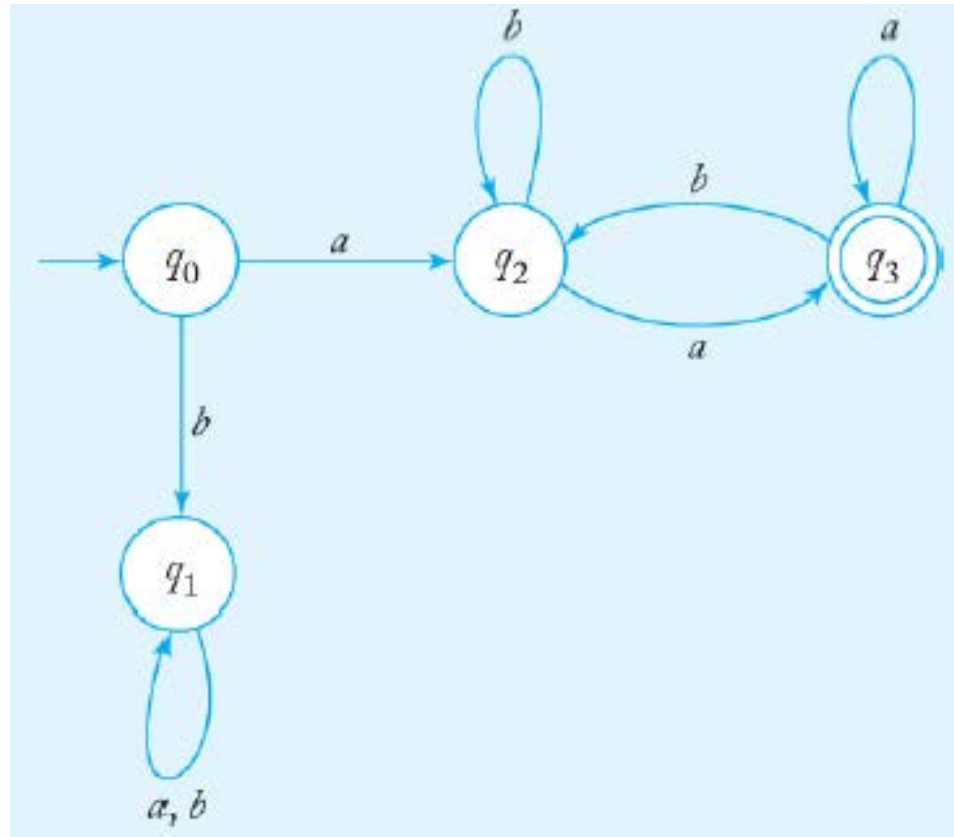


# Regular Language

- A language is called a *regular language* if some finite automaton recognizes it.
- A language  $L$  is called regular if and only if there exists some deterministic finite acceptor  $M$  such that  $L = L(M)$ .

# RL Problem

- Show that the language  $L = \{awa : w \in \{a, b\}^*\}$  is regular.



# RL Problem

- $L2 = \{aw1aaw2a : w1, w2 \in \{a, b\}^*\}$ .

