

*Pharos University in Alexandria*  
*Faculty of Computer Science & Artificial Intelligence*  
*Course Title: Theory of Computation*  
*Code: CS 307*



# Theory of Computation

**Lecturer: Sherine Shawky**

Text Books

1. Introduction to formal languages and automata, Peter Linz, 6th edition, 2017.

Week 1

# Introduction to Theory of Computation

# Agenda

- Hi ☺
- Marks
- Class Rules
- Importance
- Definition
- Mathematical Notation
- Languages
- Grammar
- Automaton

# Importance

- It is the most important course in CS.
- It helps us to understand how people thought of CS as a science in last 50 years.
- It deals with kind of things that can be computed mechanically showing how fast and required space that it needed.
- It provides concepts and principles that help us understand how to design machine.
- What are the mathematical properties of computer hardware and software?
- What are the limitations of computers? Can everything be computed?

# Definition

- **Theory of computation** is the branch that deals with whether and how efficiently problems can be solved on a model of computation, using an algorithm.
- It is divided into three major branches:
  - automata theory
  - computability theory
  - complexity theory
- There are several models in use, but the most commonly examined is the Turing machine.

# Mathematical Notation

- Sets
- Functions
- Graphs and Trees
- Proof Techniques

# Mathematical Notation

- Sets:
  - A set is a collection of elements, without any structure other than membership.
  - $S = \{0, 1, 2\}$

# Mathematical Notation

- Union
  - Union ( $\cup$ ): the combination of the distinct elements from both sets.
- Intersection:
  - Intersection ( $\cap$ ): the elements common to both sets.
- Difference:
  - A difference of two sets is the elements in one set that are NOT in the other.
- Complement :
  - A complement of a set is all the elements that are NOT in the set.
  - Complement symbol is a bar above the set name.



# Mathematical Notation

- Power set:

- The power set of a set  $S$  (written as  $2^S$ ) is the set of all the subsets of the set  $S$ .

If  $S$  is the set  $\{a, b, c\}$ , then its powerset is

$$2^S = \{\emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}\}$$

- Cartesian product:

- The Cartesian product of two sets  $S$  and  $T$ , is the set of all the ordered pairs created by choosing one element of  $S$  and one element of  $T$ .

Let  $S_1 = \{2, 4\}$  and  $S_2 = \{2, 3, 5, 6\}$ . Then

$$S_1 \times S_2 = \{(2, 2), (2, 3), (2, 5), (2, 6), (4, 2), (4, 3), (4, 5), (4, 6)\}$$

# Mathematical Notation

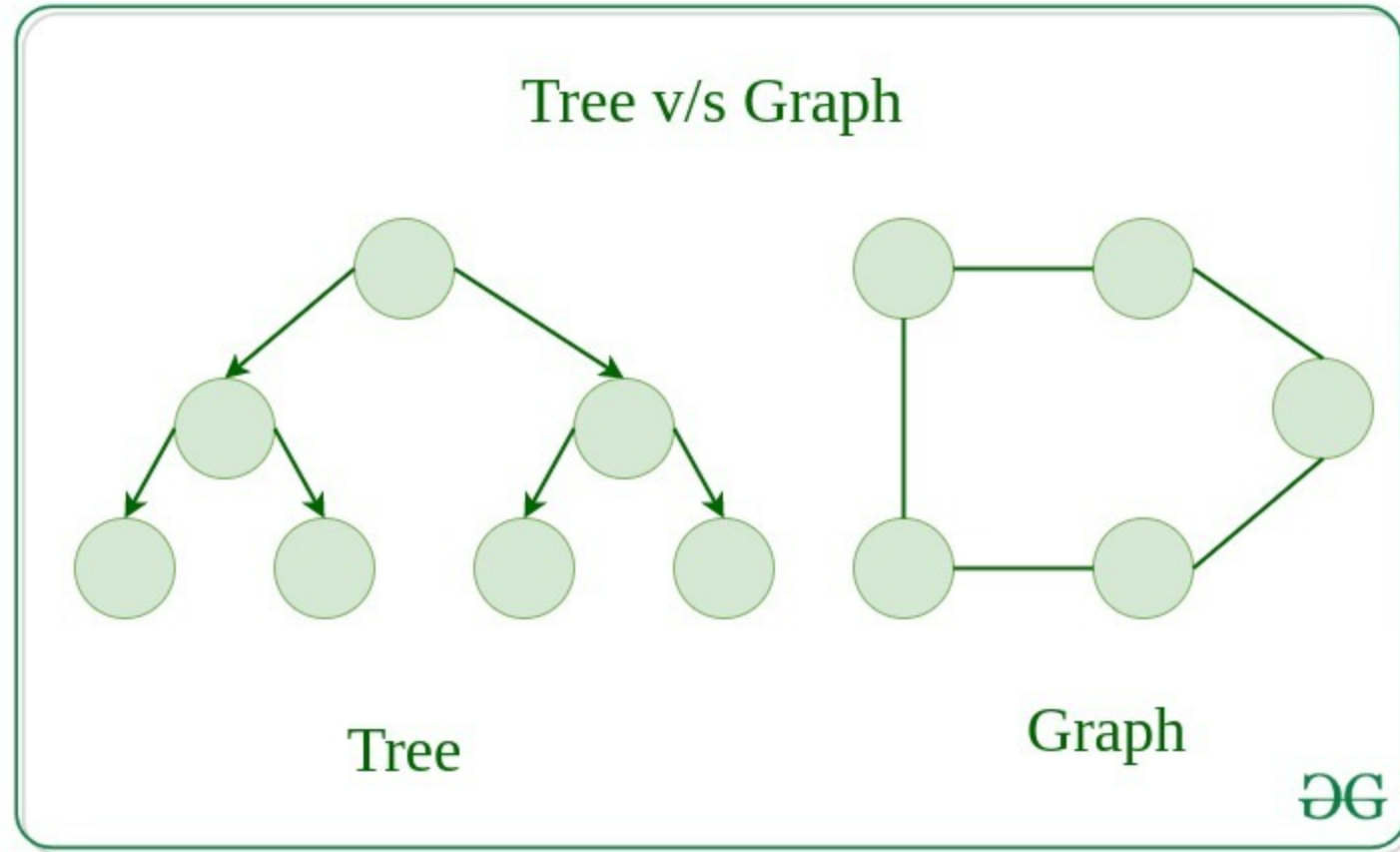
- Functions:
  - A function  $f$ : is the mapping of elements of  $S$  to unique elements of  $T$ .

A **function** is a rule that assigns to elements of one set a unique element of another set. If  $f$  denotes a function, then the first set is called the **domain** of  $f$ , and the second set is its **range**. We write

$$f : S_1 \rightarrow S_2$$

to indicate that the domain of  $f$  is a subset of  $S_1$  and that the range of  $f$  is a subset of  $S_2$ .

# Mathematical Notation



<https://www.geeksforgeeks.org/difference-between-graph-and-tree/>

# Mathematical Notation

- Proof Techniques:
  - Direct Proof
  - Proof by contrapositive
  - Proof by contradiction.
  - Proof by induction.

# Languages

- Language: a set of strings based on a certain alphabet.
  - Alphabet (usually denoted by  $S$ )
    - Finite
    - Non-Empty
    - Consists of a Set of Symbols
    - Usually show with lowercase letters like a, b, c, ...
    - Example:  $S = \{e, u, t, r\}$ ,  $L = \{eut, tr, rrr, \dots\}$
  - Symbols :
    - Symbols are indivisible objects or entity that cannot be defined.
    - That is, symbols are the atoms of the world of languages.
    - A symbol is any single object such as ,  $a$ , 0, 1, #, begin, or do.

# Languages

- **Concatenation**

- Example:  $w = a_1 a_2 \dots a_n$

$$u = b_1 b_2 \dots b_m$$

- $w u = a_1 a_2 \dots a_n b_1 b_2 \dots b_m$

- $\lambda w = w \lambda = w$

- $|uv| = |u| + |v|$  “Length”

# Languages

- **Star Closure (Universe)**
  - Set of all strings over S (alphabet)

Let  $\Sigma = \{a, b\}$ . Then

$$\Sigma^* = \{\lambda, a, b, aa, ab, ba, bb, aaa, aab, \dots\}.$$

The set

$$\{a, aa, aab\}$$

is a language on  $\Sigma$ . Because it has a finite number of sentences, we call it a finite language. The set

$$L = \{a^n b^n : n \geq 0\}$$

# Grammars

- **Grammars:**

A grammar  $G$  is defined as a quadruple

$$G = (V, T, S, P),$$

- **Definition:**

- **V** = Set of Variables (non-empty)
- **T** = Set of Terminal Symbols
- **S** = Start Symbol
- **P** = Set of Productions



# Grammars

- Example 1:

Consider the grammar:  $G = (\{S\}, \{a, b\}, S, P)$ ,

with  $P$  given by

$$S \rightarrow aSb,$$

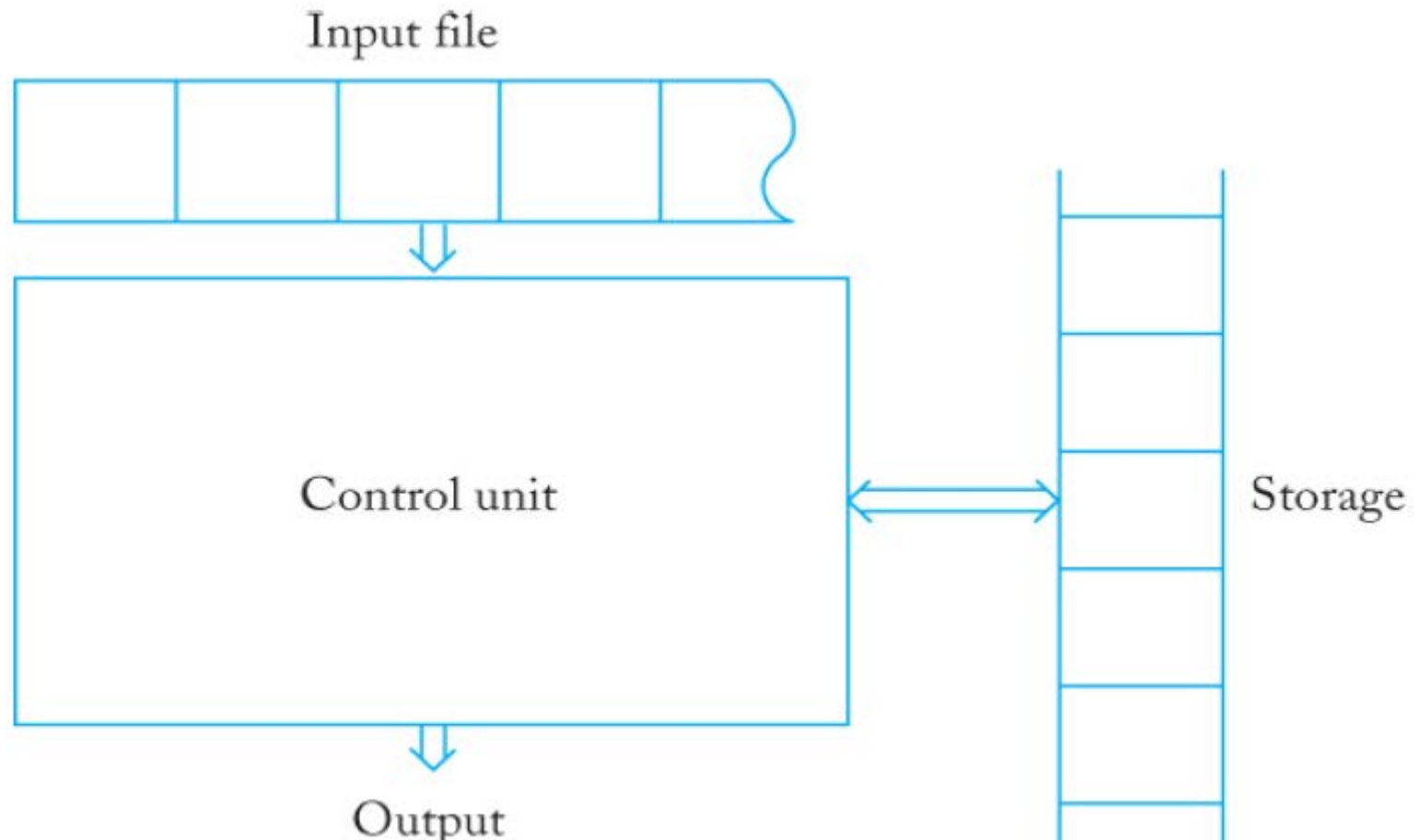
$$S \rightarrow \lambda.$$

*It generates the language:*

$$L(G) = \{\lambda, ab, aabb, aaabbb, \dots\}$$

$$L(G) = \{a^n b^n : n \geq 0\},$$

# Automaton



# Automaton

