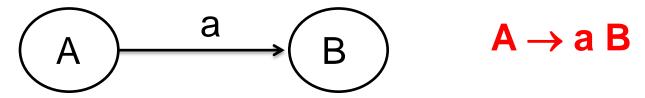
Automata and Language Theory Chapter 6 (Finite Automata)

Part 1

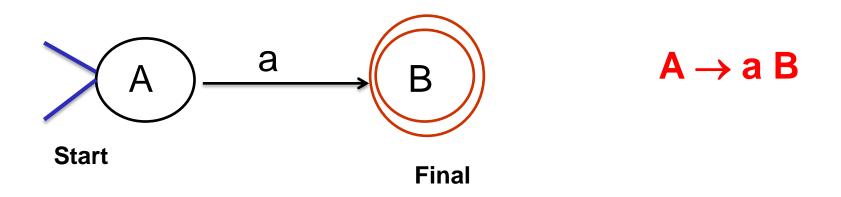
Dr. Doaa Shebl
Faculty of Computers and Artificial Intelligence
Beni-Suef University

Deterministic Finite Automaton (DFA)

- An informal definition
 - A diagram with a finite number of states represented by circles
 - An arrow points to one of the states, the unique start state
 - Double circles mark any number of the states as accepting states



Deterministic Finite Automaton (DFA)



- For every state, for every symbol in Σ , there is exactly one arrow labeled with that symbol going to another state (or back to the same state).
- Given any string over Σ , a DFA can read the string and follow its state-to-state transitions.
- At the end of the string, if it is in an accepting state, we say it accepts the string. Otherwise it rejects.
- The language defined by a DFA is the set of strings in Σ^* that it accepts.

Definition

The state diagram of a DFA $M = (Q, \Sigma, \delta, q_0, F)$ is a labeled graph G defined by the following conditions:

- i) The nodes of G are the elements of Q.
- ii) The labels on the arcs of G are elements of Σ .
- iii) q_0 is the start node, depicted \times .
- iv) F is the set of accepting nodes; each accepting node is depicted \bigcirc .
- v) There is an arc from node q_i to q_j labeled a if $\delta(q_i, a) = q_j$.
- vi) For every node q_i and symbol $a \in \Sigma$, there is exactly one arc labeled a leaving q_i .

δ is the transition function

A function $\delta(q,a)$ that takes the current state q and next input symbol a, and returns the next state

Example 1:

- This DFA defines $\{xa \mid x \in \{a,b\}^*\}$
- Formally, $M = (Q, \Sigma, \delta, q_0, F)$, where

$$- Q = \{q_0, q_1\}$$

$$-\Sigma = \{a,b\}$$

$$- F = \{q_1\}$$

- Start sate q_0





$\begin{array}{c} & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\$			
	a	b	
G o	q_1	q_0	

 \mathbf{q}_1

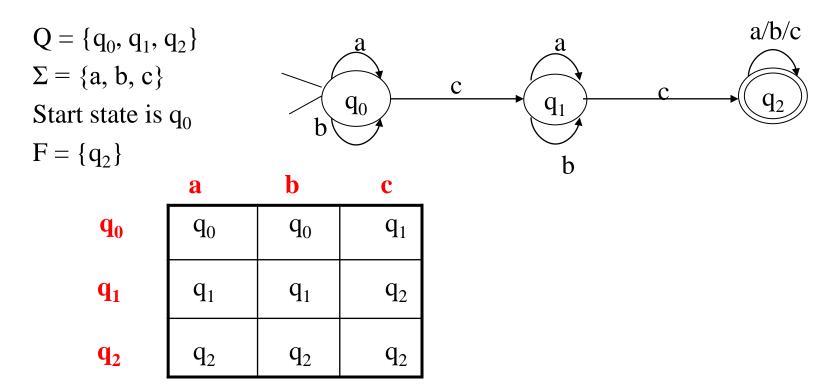
 $\mathbf{q_0}$

 \mathbf{q}_1

$$-\delta(q_0,a) = q_1, \, \delta(q_0,b) = q_0, \, \delta(q_1,a) = q_1, \, \delta(q_1,b) = q_0$$

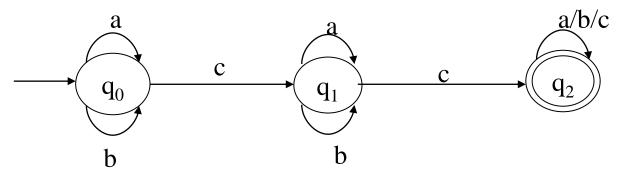
• We could just say $M = (\{q_0, q_1\}, \{a, b\}, \delta, q_0, \{q_1\})$

Example 2:



• Since δ is a function, at each step M has exactly one option.

Example 3:



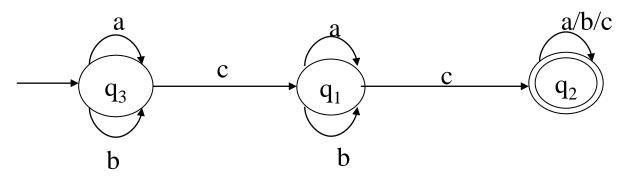
Which of the string acccb and aac accepted by Machine M???

Acccb

- $\vdash [q_0, acccb]$
- $\vdash [q_0, cccb]$
- \vdash [q₁, ccb]
- $\vdash [q_2, cb]$
- $\vdash [q_2, b]$
- $\vdash [q_2, \lambda]$

accepted

Example 3:



aac

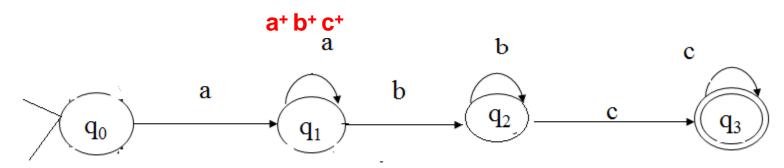
- $\vdash [q_0, aac]$
- $\vdash [q_0, ac]$
- $\vdash [q_0, c]$
- $\vdash [q_1, \lambda]$

rejected

Accepts those strings that contain at least two c's

Example 4:

Give finite automata for the set of strings over {a, b, c } which all the a's precede the b's, which in turn precede the c's. without the null string.



G:
$$q_0 \rightarrow aq_1$$

 $q_1 \rightarrow a q_1 / b q_2$
 $q_2 \rightarrow b q_2 / c q_3$
 $q_3 \rightarrow c q_3 / \lambda$

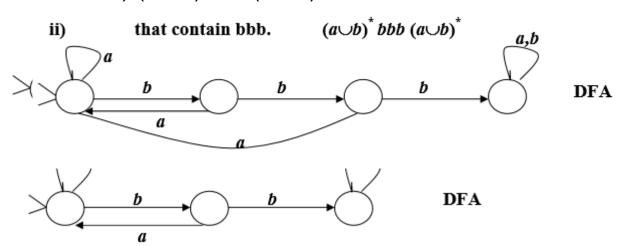
Example 5:

Give a finite automata for the set of strings over {a, b}:

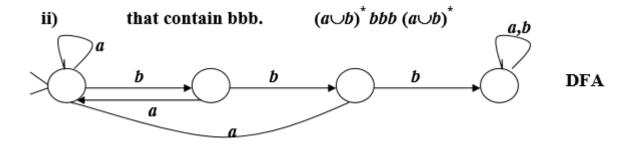
i) that contain bb. ii) that contain bbb.

Solution

i) (a∪ b)* bb (a∪ b)*

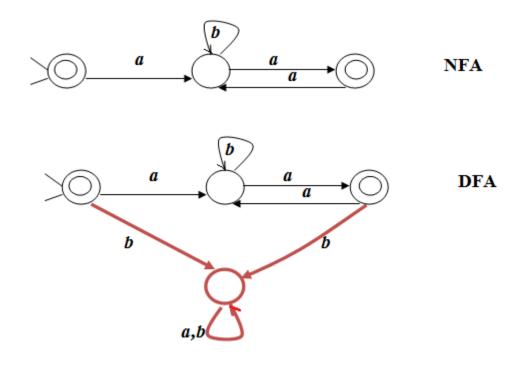


Example 5:



Example 6:

Give a state diagram of NFA and DFA for the following: $(ab^*a)^*$.



DFA Applications

We have seen how DFAs can be used to define formal languages. In addition to this formal use, DFAs have practical applications. DFAbased pieces of code lie at the heart of many commonly used computer programs.

DFA Applications

- Programming language processing
 - Scanning phase: dividing source file into "tokens" (keywords, identifiers, constants, etc.), skipping whitespace and comments
- Command language processing
 - Typed command languages often require the same kind of treatment
- Text pattern matching
 - Unix tools like awk, egrep, and sed, mail systems like ProcMail, database systems like MySQL, and many others

More DFA Applications

- Signal processing
 - Speech processing and other signal processing systems
 use finite state models to transform the incoming signal
- Controllers for finite-state systems
 - Hardware and software
 - A wide range of applications, from industrial processes to video games