# Automata and Language Theory Chapter 3(Context Free Grammar)

Dr. Doaa Shebl
Faculty of Computers and Artificial Intelligence
Beni-Suef University

## **Definition 3.1.1**

A context-free grammar is a quadruple  $(V, \Sigma, P, S)$  where V is a finite set of variables,  $\Sigma$  (the alphabet) is a finite set of terminal symbols, P is a finite set of rules, and S is a distinguished element of V called the start symbol. The sets V and  $\Sigma$  are assumed to be disjoint.

A rule, often called a production,

G: 
$$S \rightarrow a S b /ab / A$$
  
 $A \rightarrow a A / a$   
 $\Sigma = \{ a, b \}$   
 $N = V = \{ S, A \}$   
 $P = \{ S \rightarrow a S b, S \rightarrow a b, S \rightarrow A, A \rightarrow a A, A \rightarrow a \}$   
 $S = S$ : Start Symbol

# **Derivation**

 $G: S \rightarrow a S b / SS / ab$ 

$$S \Rightarrow ab$$

$$S \Rightarrow aSb$$
$$\Rightarrow aabb$$

$$S \Rightarrow aSb$$

$$\Rightarrow aaSbb$$

$$\Rightarrow \cdots aa...bb$$

Let G be the grammar :  $S \rightarrow AA$ 

 $A \rightarrow AAA/bA/Ab/a$ 

Construct the left most and right most derivation of the string <u>baba</u>

#### **Left most derivation**

 $S \Rightarrow AA$ 

 $\Rightarrow$  bAA

 $\Rightarrow$  b a A

⇒b a bA

⇒ba ba

#### Right most derivation

 $S \Rightarrow AA$ 

 $\Rightarrow$  A a

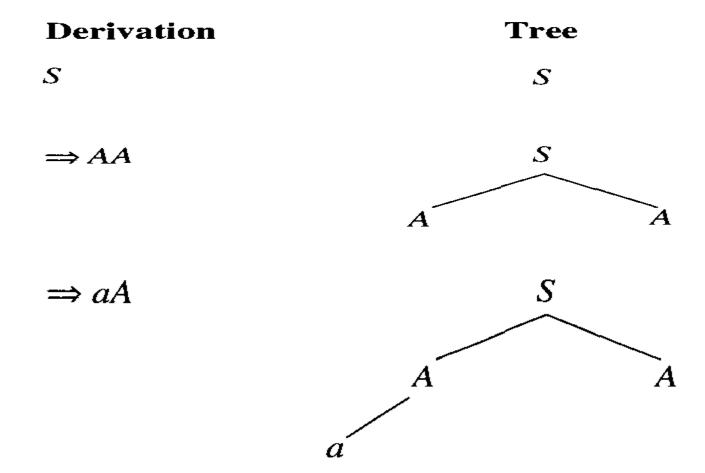
 $\Rightarrow$  Ab a

 $\Rightarrow$ b A ba

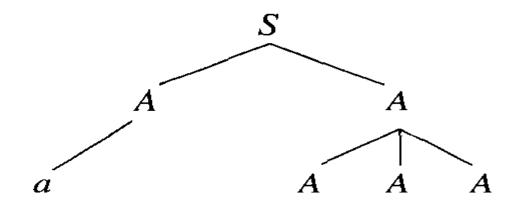
⇒ba ba

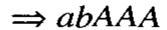
G = (V, 
$$\Sigma$$
, P, S)  
V = {S, A}  
 $\Sigma = \{a, b\}$   
P:  $S \rightarrow AA$   
 $A \rightarrow AAA \mid bA \mid Ab \mid a$ 

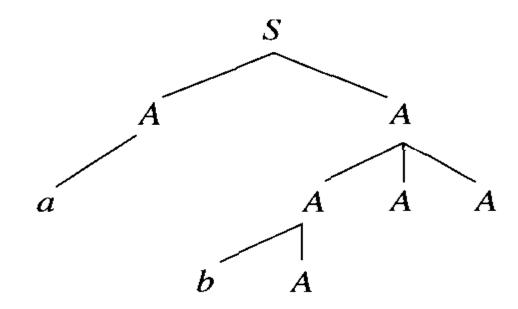
$$S \Rightarrow AA$$
  $S \Rightarrow AA$   $S \Rightarrow AA$   $\Rightarrow AA$   $\Rightarrow$ 



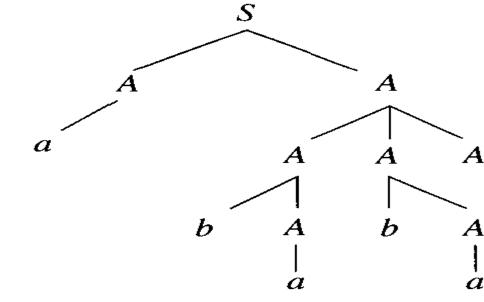




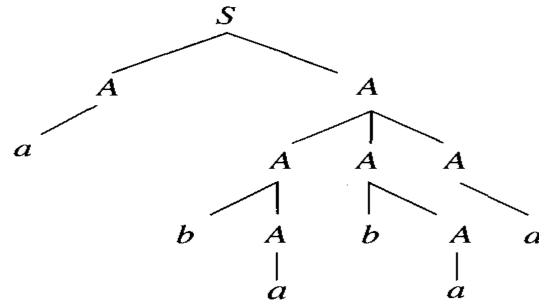




 $\Rightarrow ababaA$ 



⇒ ababaa

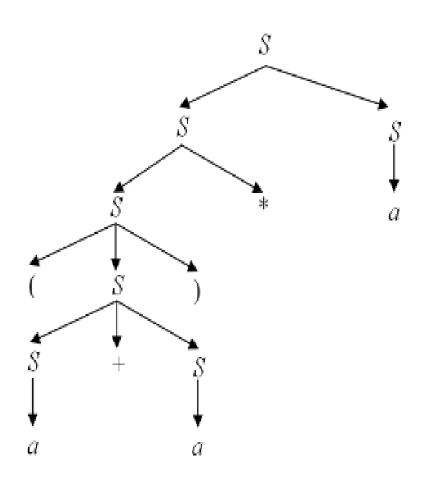


For the  $S \rightarrow S + S \mid SS \mid (S) \mid S^* \mid a$  with string  $(a + a)^* a$ :

- i) Give a leftmost derivation for the string.
- ii) Give a rightmost derivation for the string.
- iii) Give a parse tree for the string.

i) 
$$S => S S => S * S => (S) * S => (S+S) * S$$
  
 $=> (a+S) * S => (a+a) * S => (a+a) * a$ .  
ii)  $S => S S => S a => S * a => (S) * a => (S+S) * a$   
 $=> (S+a) * a => (a+a) * a$ .

iii)



Parse tree