

TSKS14

Multiple Antenna Communications

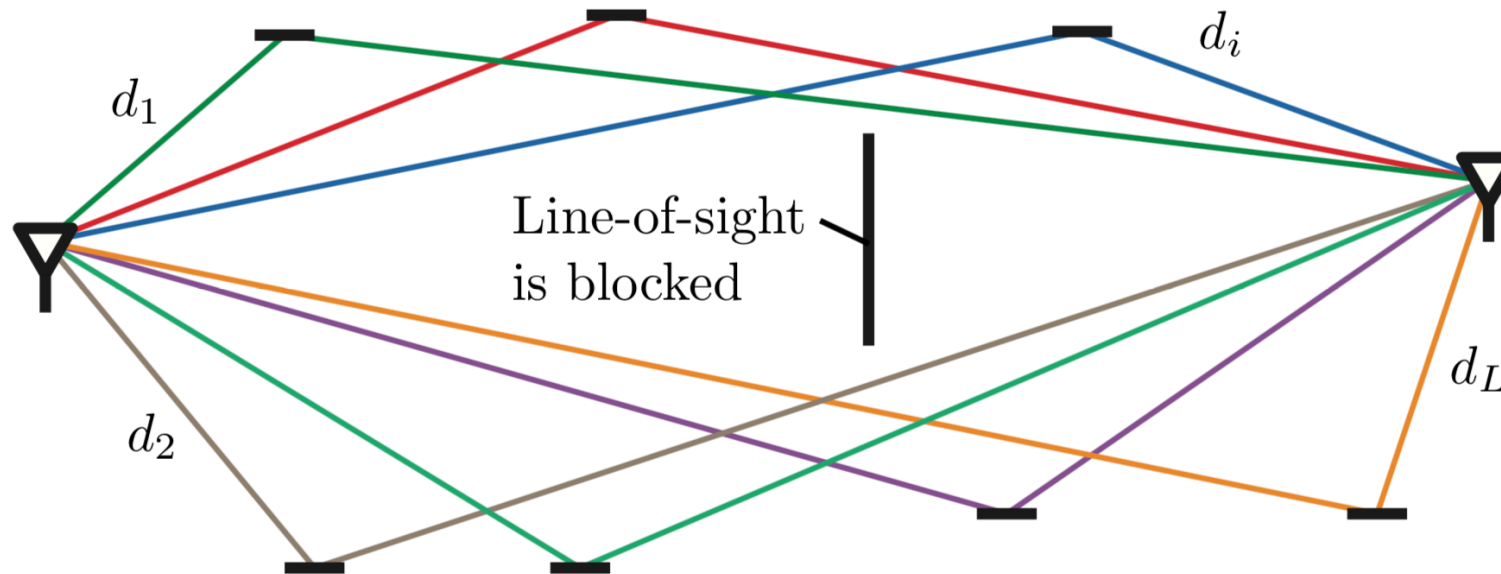
Lecture 3, 2020

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Outline of this lecture

- Multipath propagation and Rayleigh fading
- Slow fading
 - Outage probability
 - Outage capacity
 - Spatial diversity
- Fast fading
 - Ergodic capacity
 - Channel hardening

Multipath propagation



Non-line-of-sight channel:

- Scattering

- Channel with L propagation paths:

$$g = \sum_{i=1}^L \sqrt{\alpha_i} e^{-j2\pi \frac{d_i - d}{\lambda}}$$

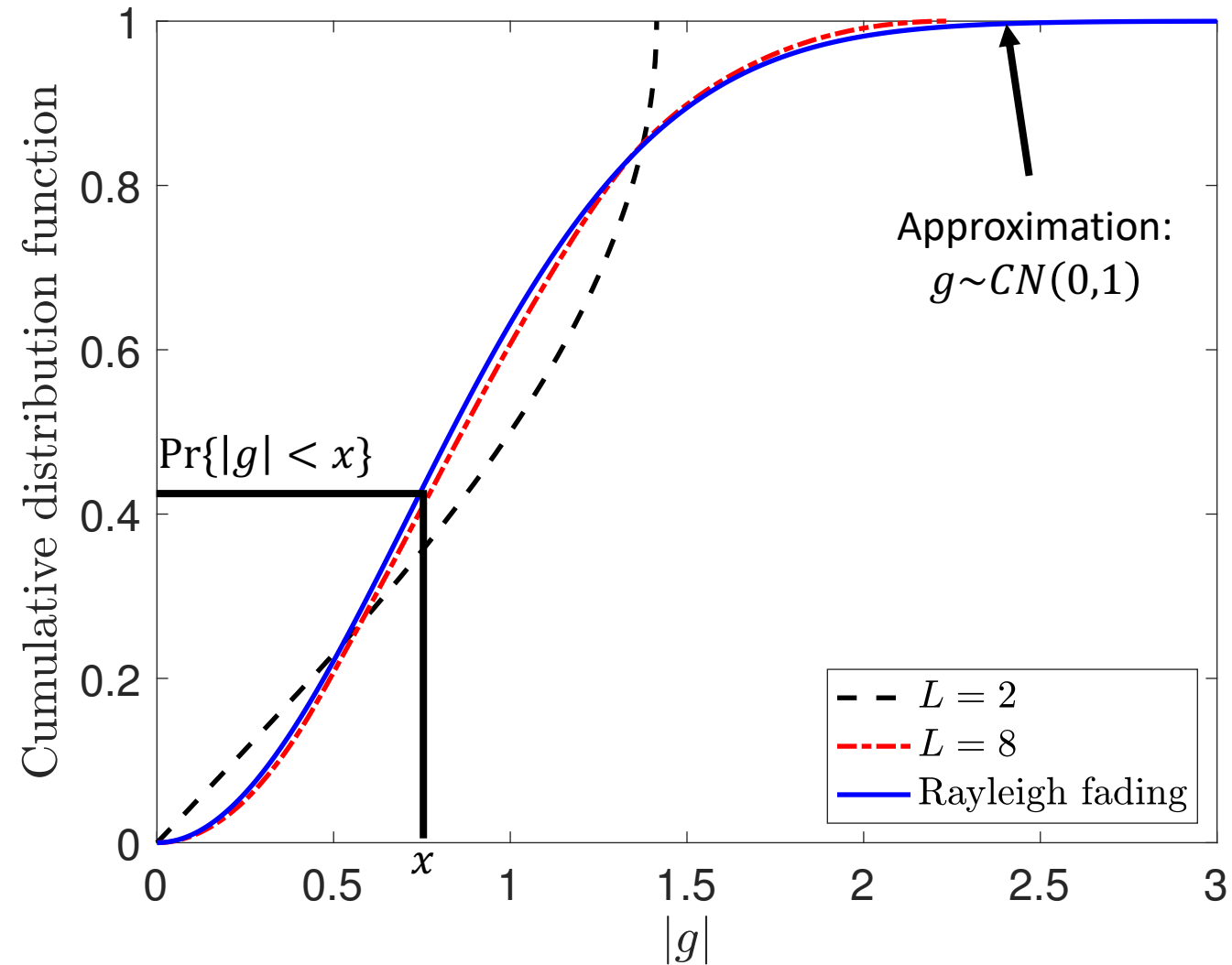
Channel gain
Reference distance

Wavelength

Multipath fading

- Example:
 - $\alpha_i = \frac{1}{L}$
 - $\theta_i = 2\pi \frac{d_i - d}{\lambda} \sim U(0, 2\pi)$
- Channel magnitude:

$$|g| = \left| \sum_{i=1}^L \sqrt{\frac{1}{L}} e^{-j\theta_i} \right|$$



Rich scattering: Rayleigh fading

Central limit theorem

Let X_1, \dots, X_L be a sequence of L real-valued independent and identically distributed random variables with zero mean and variance σ^2 . As $L \rightarrow \infty$,

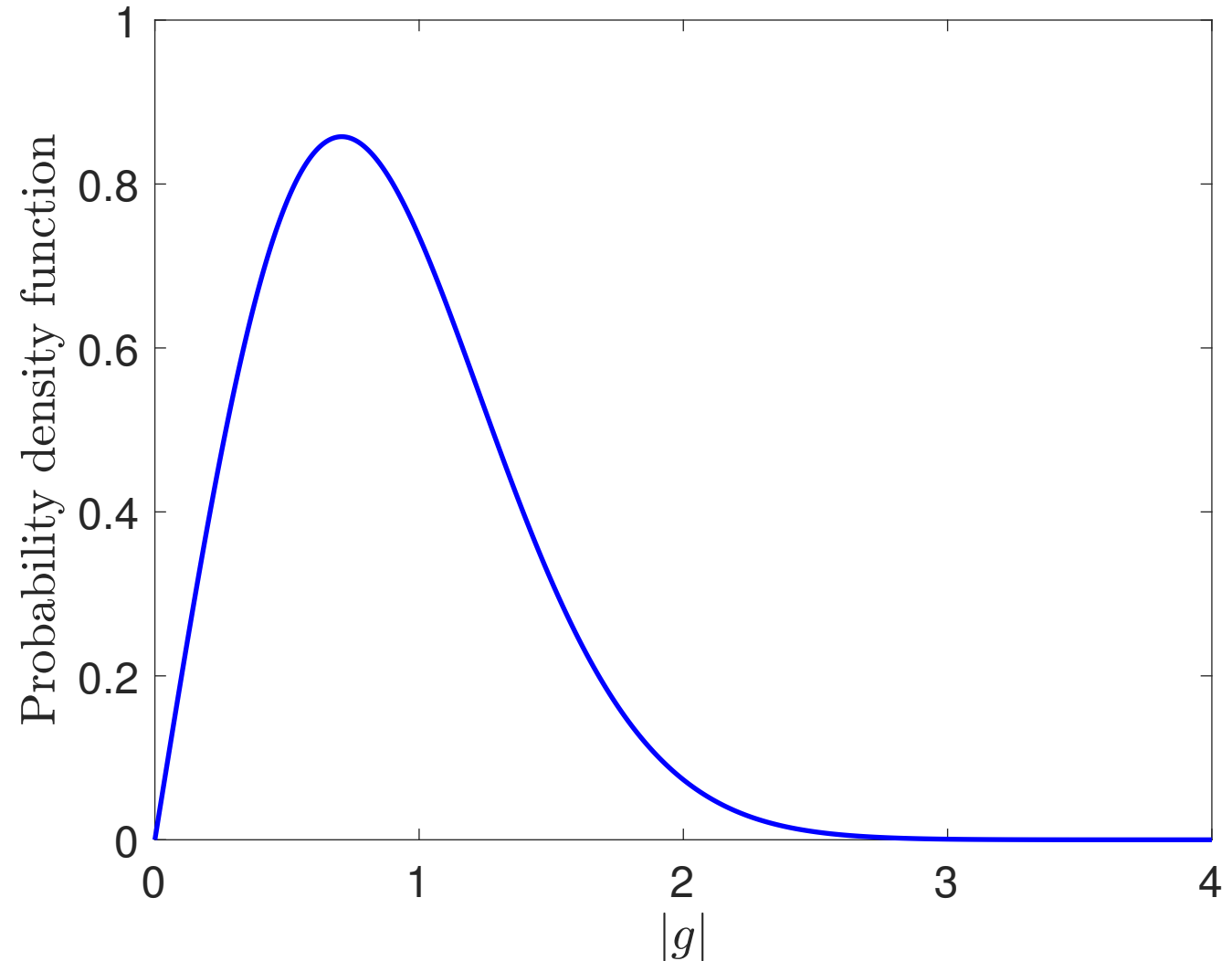
$$\frac{1}{\sqrt{L\sigma^2}} \sum_{i=1}^L X_i$$

converges to a standard Gaussian distribution $N(0,1)$.

- Rich multipath propagation
 - Very large number paths: Gaussian distribution
 - Channel gain: $g \sim CN(0, \beta)$
 - Called *Rayleigh fading* since $|g| \sim \text{Rayleigh}(\sqrt{\beta/2})$

Rayleigh fading, $|g| \sim \text{Rayleigh}(1/\sqrt{2})$

- Channel gain changes over time
- In this case: $g \sim \text{CN}(0,1)$

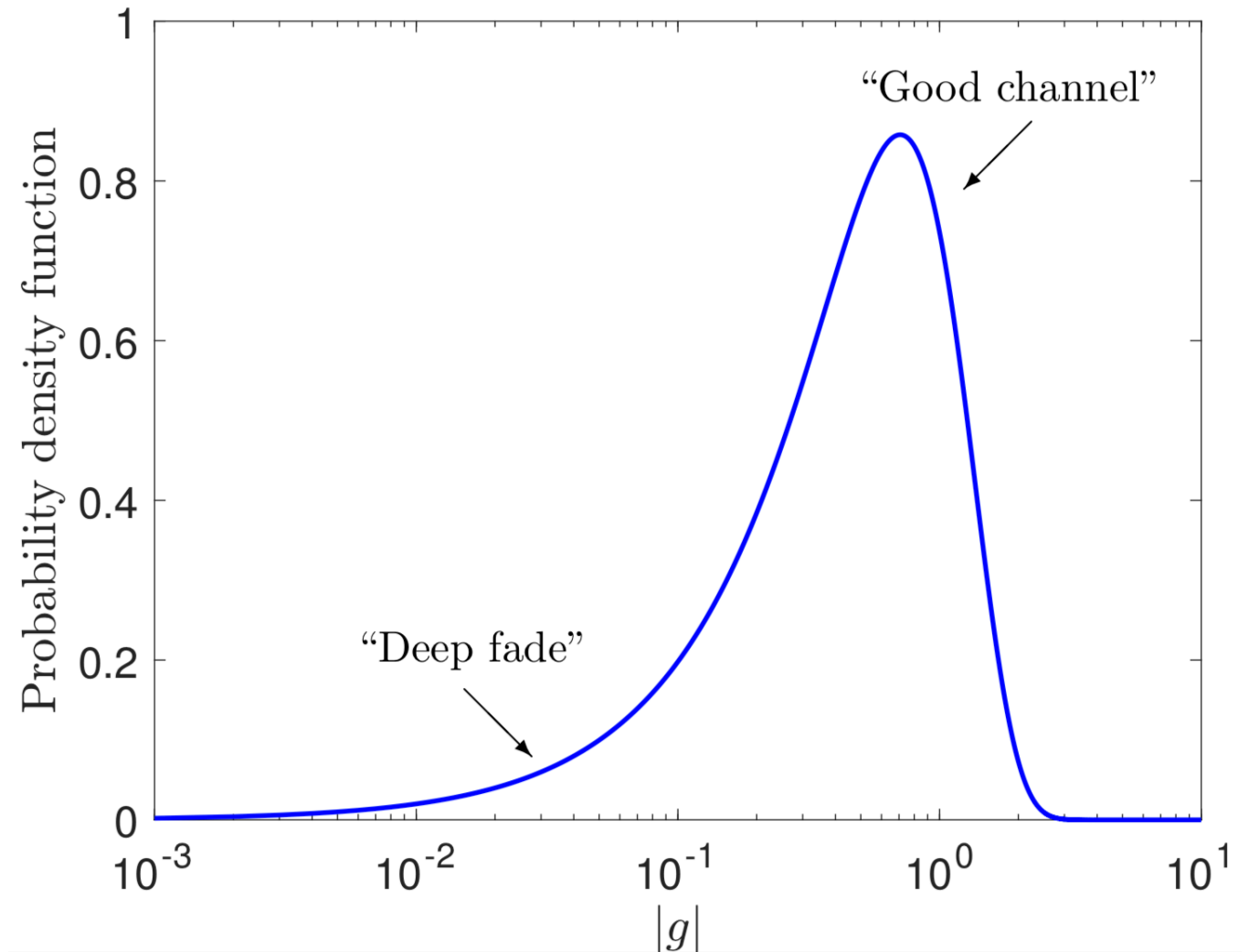


Rayleigh fading, zooming in on tail

- Risk of very small channel gain

Two issues:

- Variations in channel quality
- Unpredictable



Capacity of fading channel

- AWGN channel with a random channel response $g[l]$:

$$y[l] = g[l] \cdot x[l] + n[l]$$

- $x[l] \sim CN(0, q)$, energy per sample: $q = P/B$
 - $n[l] \sim CN(0, N_0)$
- Two categories:
 - Slow fading: $g[l]$ takes one realization during communication
 - Fast fading: $g[l]$ takes “all” realizations during communication

Reality might be somewhere in between

Slow fading


- Received signal

$$y[l] = g \cdot x[l] + n[l]$$

- Fixed channel $g[l] = g$ for the entire transmission
- Assumption: Receiver knows g , but not the transmitter

- Capacity for a realization g :

$$C_g = \log_2(1 + |g|^2 \text{SNR})$$

q/N_0 

Transmitter does not know C_g
Cannot encode the data to achieve it!

Opportunistic transmission

- Suppose transmitter encode using the rate R bit/s/Hz
- Two possible events:
 - If $R \leq C_g$: Successful transmission
 - If $R > C_g$: Large error probability

System is outage if $R > C_g$

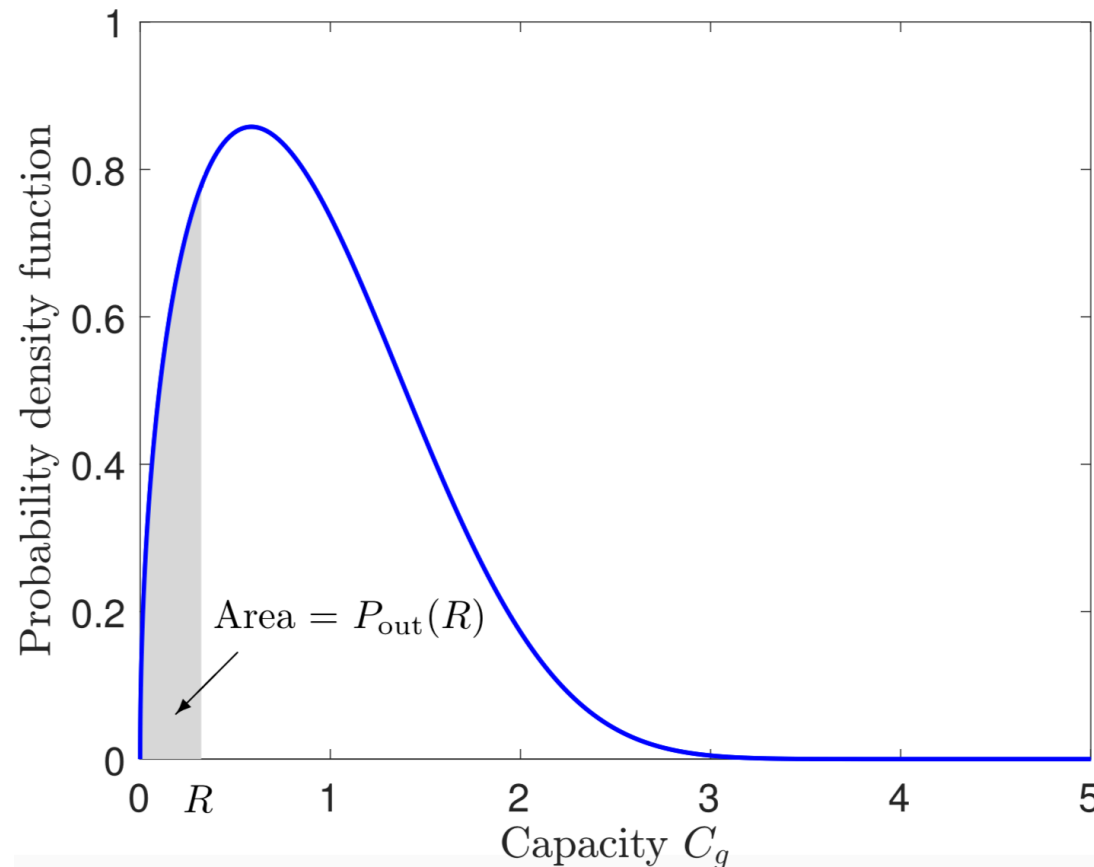
- Outage probability for rate R

$$p_{out}(R) = \Pr\{C_g < R\} = \Pr\{\log_2(1 + |g|^2 \text{SNR}) < R\}$$

Outage probability with $g \sim \mathcal{CN}(0,1)$

- Outage probability for rate R :

$$p_{out}(R) = \Pr\{C_g < R\} = 1 - e^{-\frac{2^R - 1}{\text{SNR}}} \stackrel{\text{High SNR}}{\approx} \frac{2^R - 1}{\text{SNR}}$$



Outage probability
decays with
 $\text{SNR} = q/N_0$
as SNR^{-1}

Outage capacity

- Difference from AWGN channel
 - Only $R = 0$ can guarantee zero error probability
 - Capacity is zero
- ϵ -Outage capacity C_ϵ :
 - Largest rate R such that $p_{out}(R) \leq \epsilon$

Interpretation:

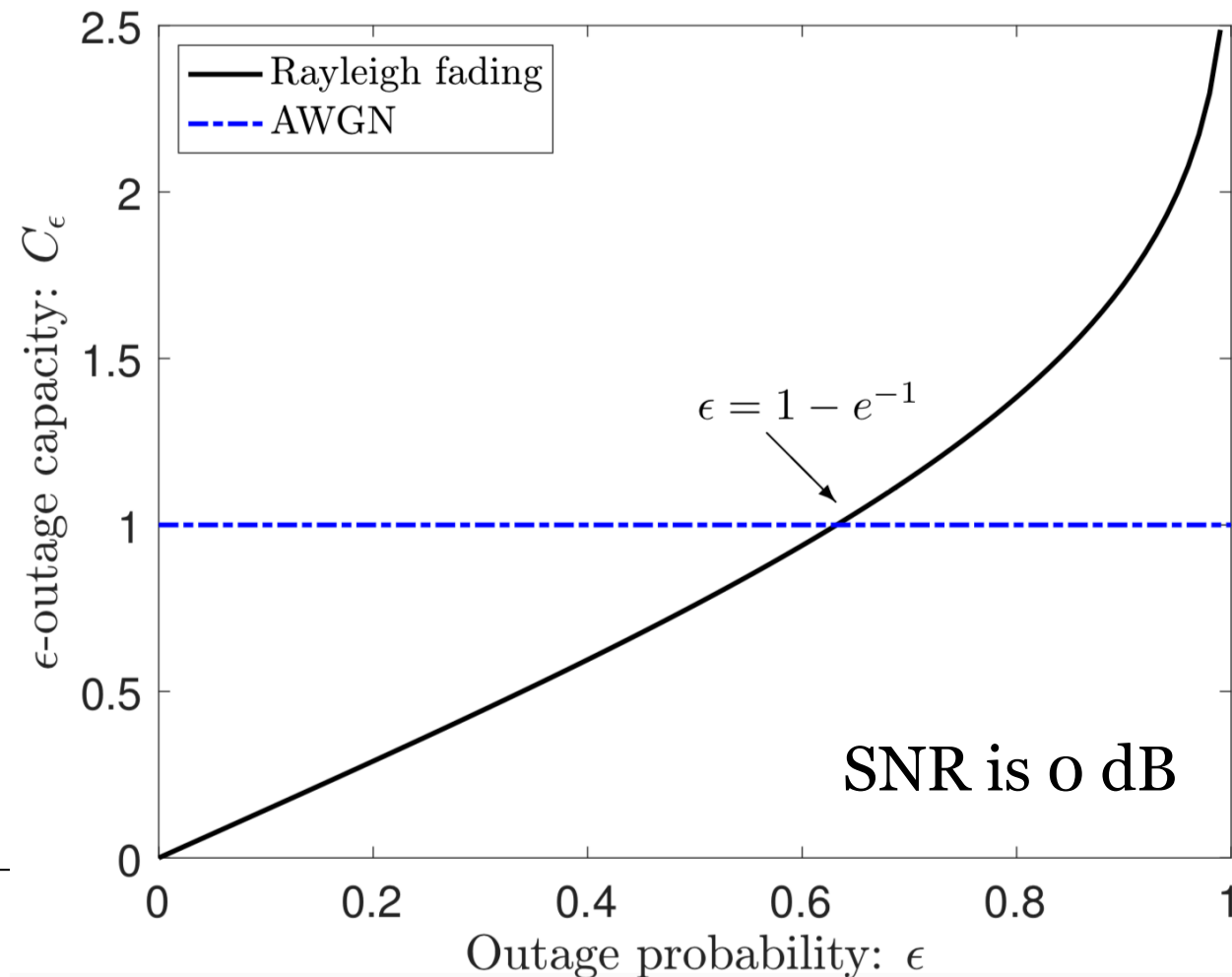
With probability $1 - \epsilon$, we can communicate at C_ϵ with is zero error probability

Outage capacity with $g \sim \mathcal{CN}(0,1)$

$$C_\epsilon = \log_2 \left(1 + \underbrace{\text{SNR} \ln((1 - \epsilon)^{-1})}_{\text{Difference from AWGN channel}} \right)$$

Difference from
AWGN channel

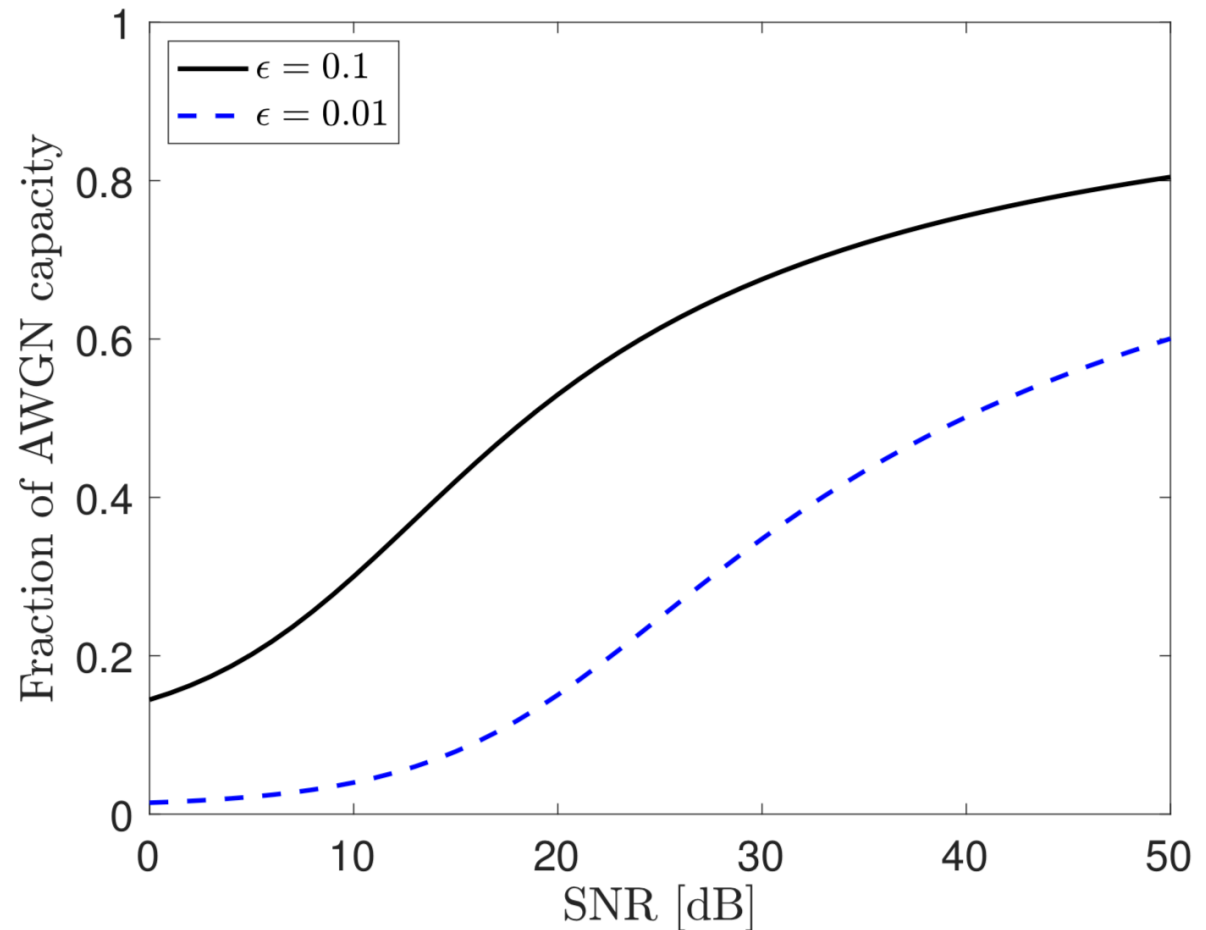
- Low ϵ : Better with AWGN channel
- High ϵ : Better with fading channel



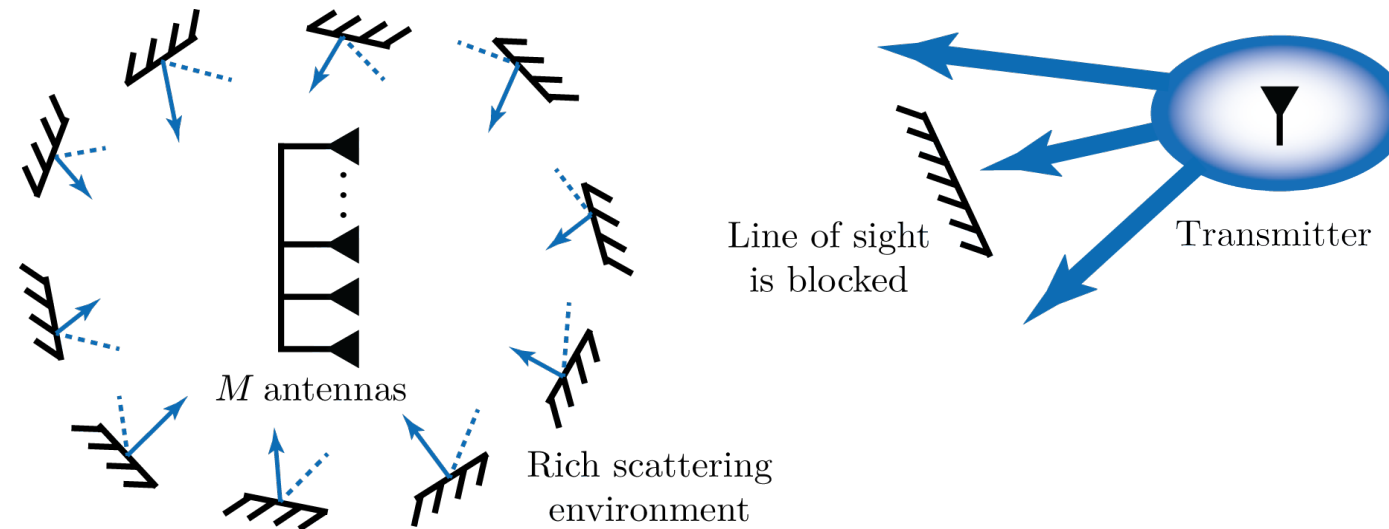
Outage capacity with small outage probability

- Fraction of AWGN capacity:
$$\frac{\log_2(1 + \text{SNR} \ln((1 - \epsilon)^{-1}))}{\log_2(1 + \text{SNR})}$$
- Much lower capacity than with AWGN channel

Can we improve the situation?



Fading multiple antenna channels



- Independent and identically distributed Rayleigh fading
 - Channel gain: $\mathbf{g} \sim \mathcal{CN}(\mathbf{0}, \beta \mathbf{I}_M)$

- Distribution of $\|\mathbf{g}\|^2$:
$$f_{\|\mathbf{g}\|^2}(x) = \frac{x^{M-1} e^{-\frac{x}{\beta}}}{(M-1)! \beta^M}$$

Independent:
Uniform linear array
with $\Delta = \lambda/2$

M receive antennas and i.i.d. Rayleigh fading

- Outage probability

$$p_{out}(R) = \Pr\{\log_2(1 + \|\mathbf{g}\|^2 \text{SNR}) < R\} = \int_0^{\frac{2^R - 1}{\text{SNR}}} f_{\|\mathbf{g}\|^2}(x) dx \stackrel{\text{High SNR}}{\approx} \left(\frac{2^R - 1}{\text{SNR}}\right)^M \frac{1}{M!}$$

Spatial diversity gain

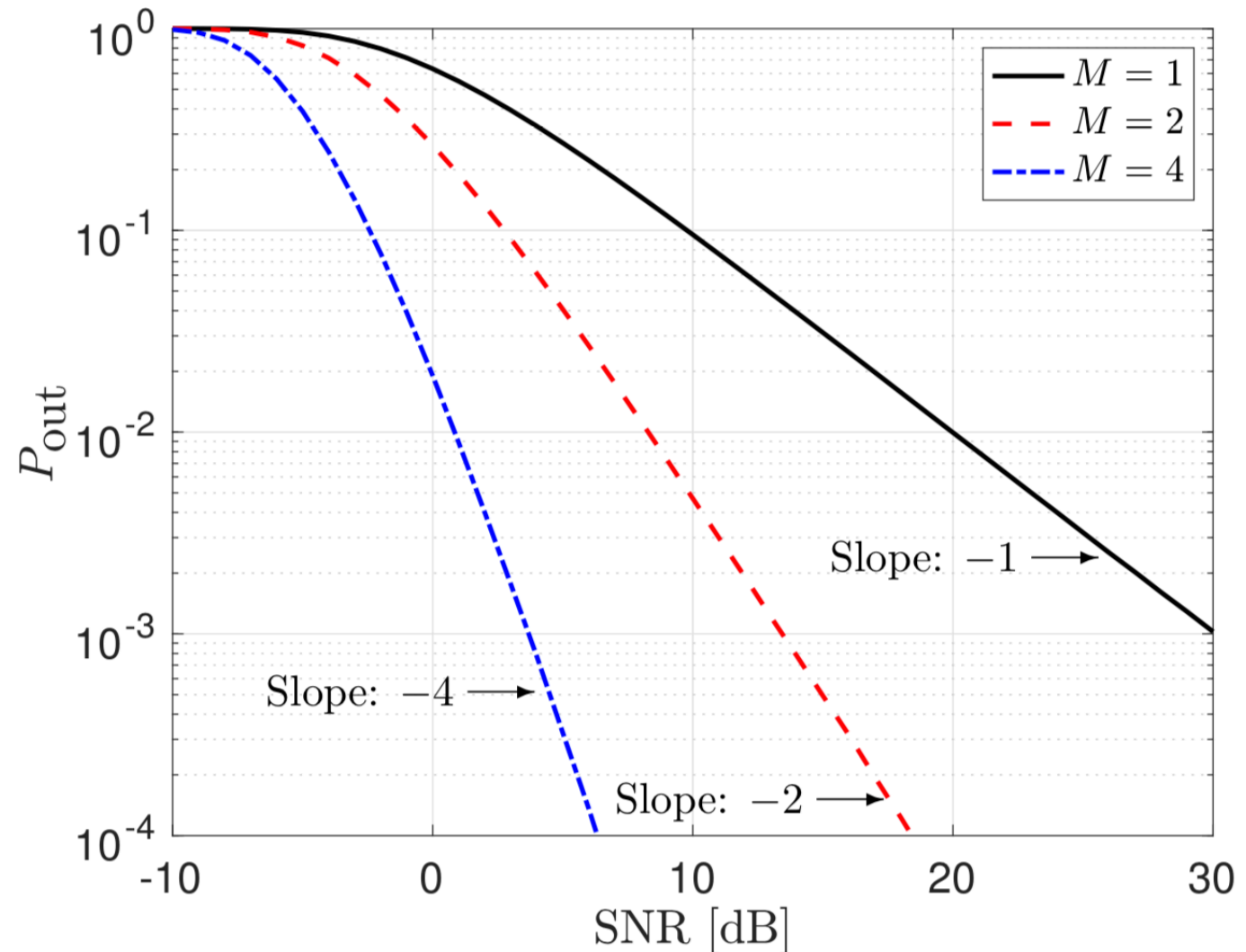
$p_{out}(R)$ proportional to SNR^{-M}
 M is the diversity order

Outage probability with M receive antennas

- Outage probability decays as SNR^{-M}
- Makes a huge difference!

Multiple receive antennas gives:

- Beamforming gain
- Diversity gain



Fast fading

- Received signal

$$y[l] = g[l] \cdot x[l] + n[l]$$

- Block fading
 - One realization of channel $g[l]$ per l (or a finite-sized block of symbols)
 - New independent realization every time (ergodic process)

Opportunistic transmission

- Suppose transmitter encode using the rate R bit/s/Hz
 - There are L fading realization: $g[1], \dots, g[L]$

- Reliable communication if

$$\sum_{l=1}^L \log_2(1 + \text{SNR} |g[l]|^2) > LR$$

- Many fading realizations:

$$R < \frac{1}{L} \sum_{l=1}^L \log_2(1 + \text{SNR} |g[l]|^2) \rightarrow \mathbb{E}\{\log_2(1 + |g|^2 \text{SNR})\}$$

As $L \rightarrow \infty$

Mean value with respect
to channel fading

Ergodic capacity

- This is called ergodic capacity:

$$\mathbb{E}\{\log_2(1 + |g|^2 \text{SNR})\}$$

- Deterministic: Transmitter knows it even if g is unknown
- There are no outage issues!

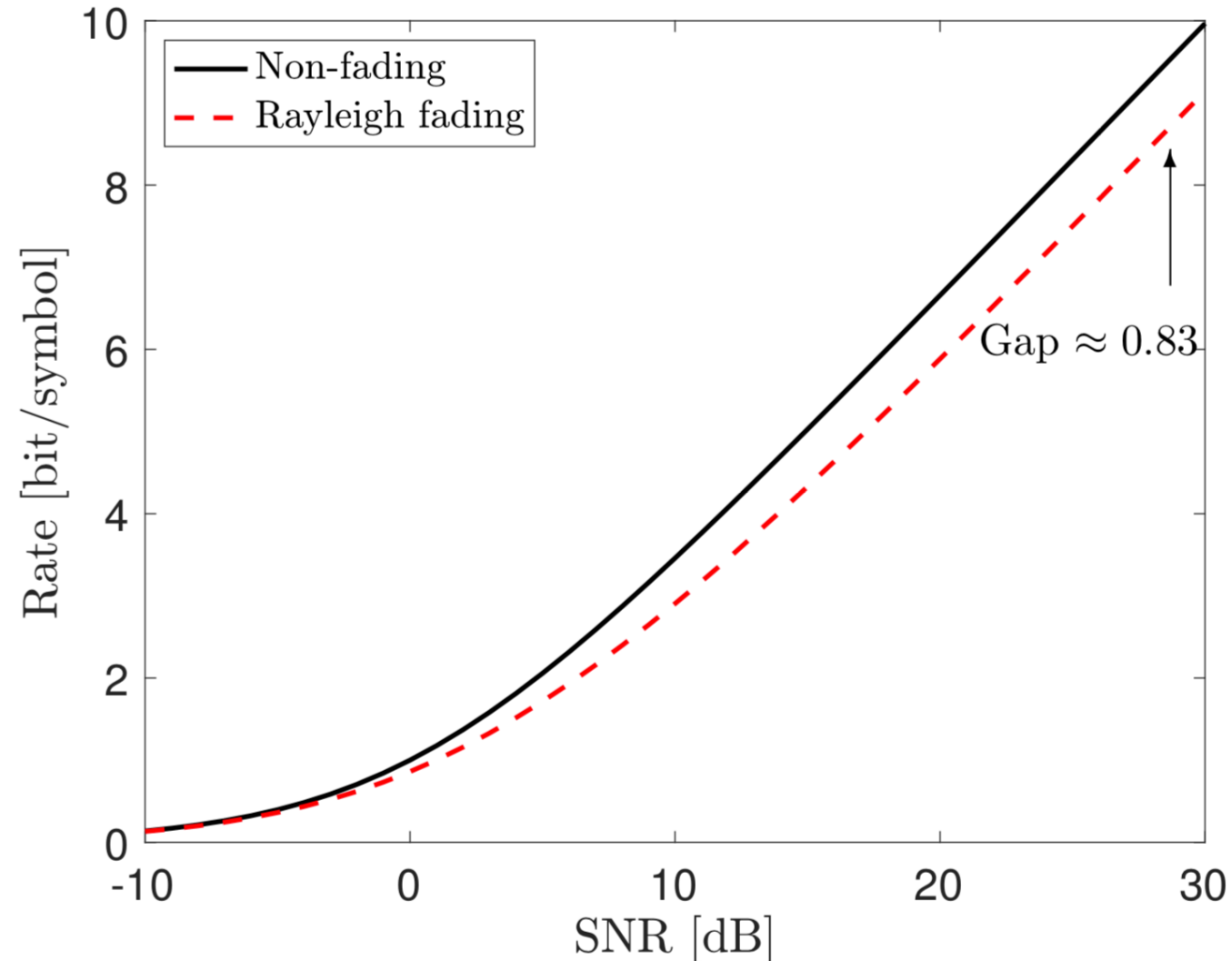
- Extension to SIMO case with channel:

$$\mathbb{E}\{\log_2(1 + \|\mathbf{g}\|^2 \text{SNR})\}$$

Comparison with AWGN channel

- AWGN channel:
 $\log_2(1 + \text{SNR})$
- Rayleigh fading:
 $\mathbb{E}\{\log_2(1 + |g|^2 \text{SNR})\}$

Low SNR: Little difference
High SNR: Ergodic capacity is lower

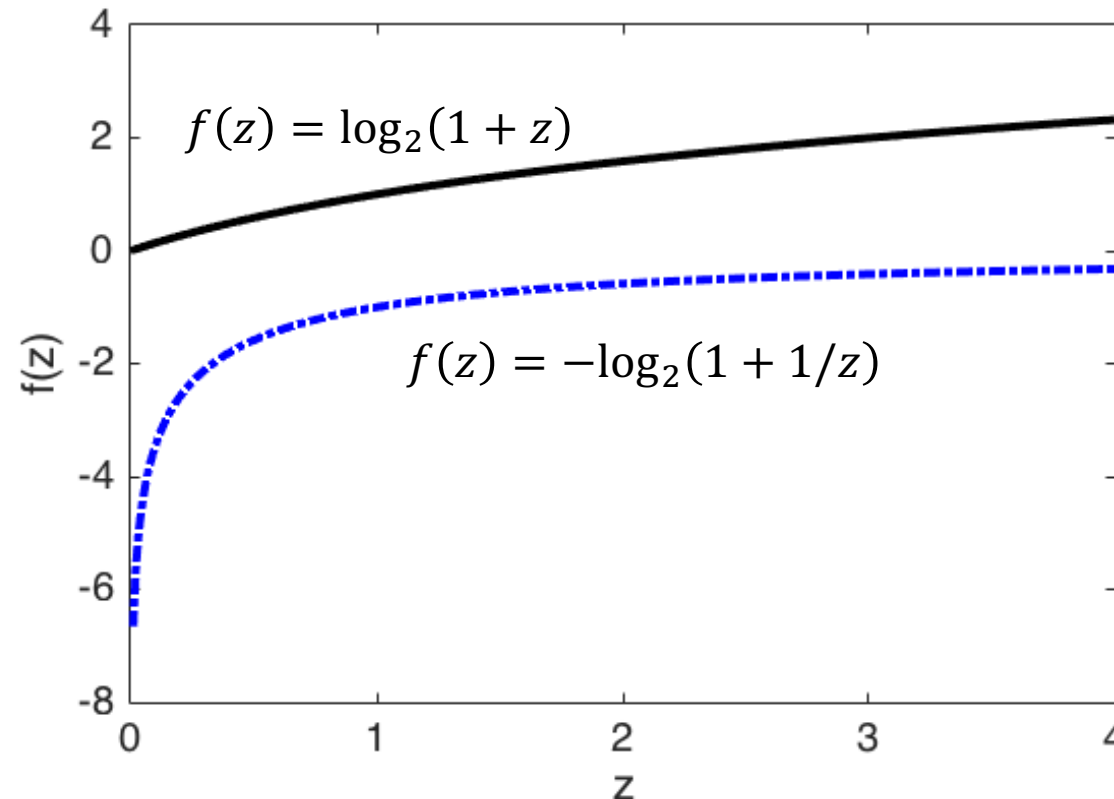


Jensen's inequality and concave functions

- For any random variable z and concave function $f(\cdot)$,
$$\mathbb{E}\{f(z)\} \leq f(\mathbb{E}\{z\})$$

A function is concave if

- Any line between two points on the curve is below the curve
- Second derivative is negative




Ergodic capacity with SIMO channel


- Can be used to prove

$$\log_2 \left(1 + \frac{\text{SNR}}{\mathbb{E}\{\|\mathbf{g}\|^{-2}\}} \right) \leq \mathbb{E}\{\log_2(1 + \|\mathbf{g}\|^2 \text{SNR})\} \leq \log_2(1 + \mathbb{E}\{\|\mathbf{g}\|^2\} \text{SNR})$$

Jensen's inequality with
 $f(z) = -\log_2(1 + 1/z)$



Jensen's inequality with
 $f(z) = \log_2(1 + z)$



- Rayleigh fading with \mathbf{g} having i.i.d. $CN(0,1)$ elements:

$$\log_2(1 + (M - 1)\text{SNR}) \leq \mathbb{E}\{\log_2(1 + \|\mathbf{g}\|^2 \text{SNR})\} \leq \log_2(1 + M \text{SNR})$$

Non-fading channel with $\|\mathbf{g}\|^2 = M - 1$



Non-fading channel with $\|\mathbf{g}\|^2 = M$

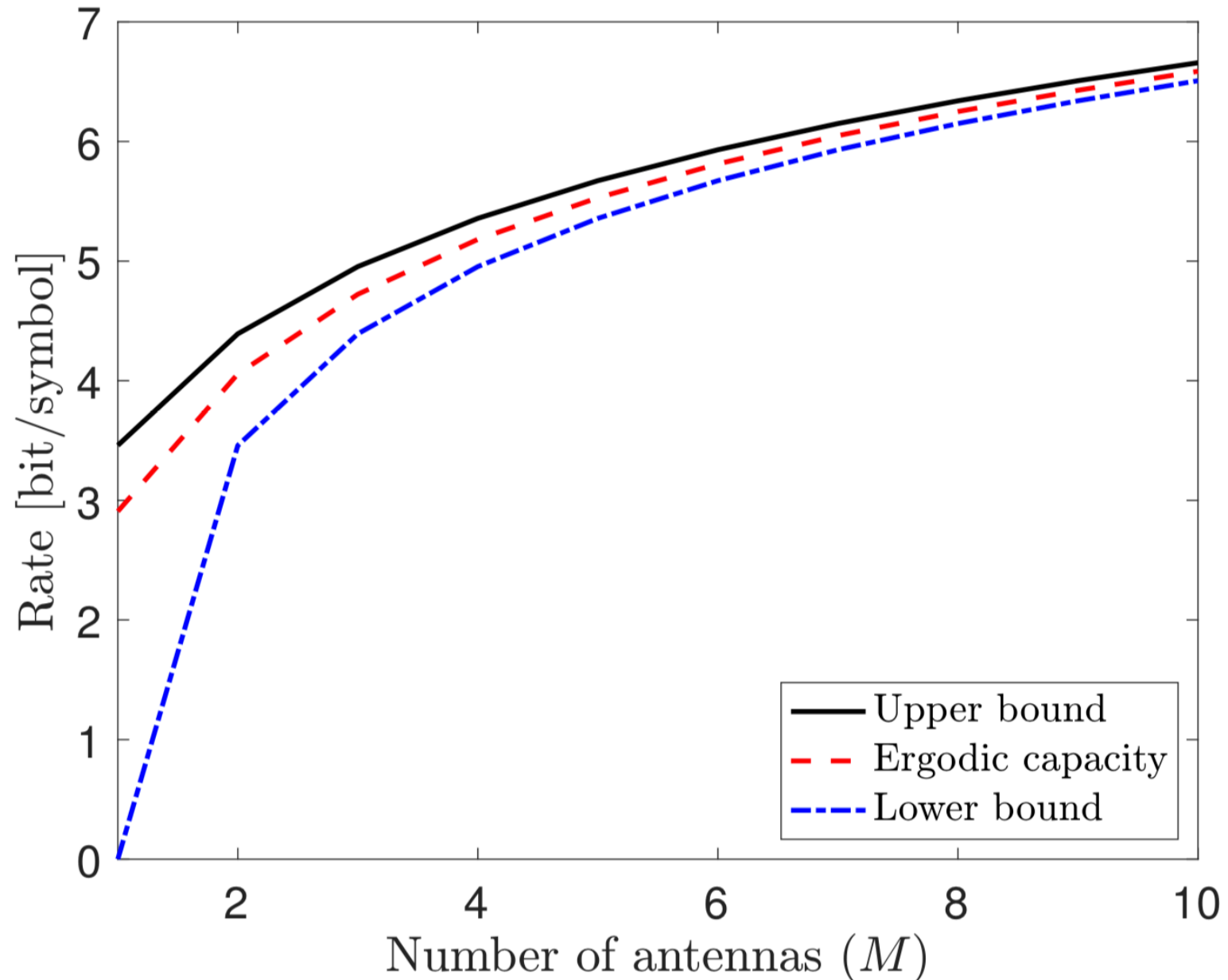


Comparison

- Small M
 - Large loss from channel fading
- Larger M
 - Small loss

Channel hardening:

When M is large, no penalty from channel fading



Summary

- **Slow fading:** One realization during transmission
 - Outage probability, outage capacity
 - Reliable communication → Large performance loss
 - Multiple antennas give more reliability
- **Fast fading:** Many realizations during transmission
 - Ergodic capacity with averaging over fading
 - No reliability issue, but performance loss
 - Multiple antennas give similar capacity as with non-fading channels

End of Lecture 3

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