

TSKS14

Multiple Antenna Communications

Lecture 8, 2020

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Outline of this lecture

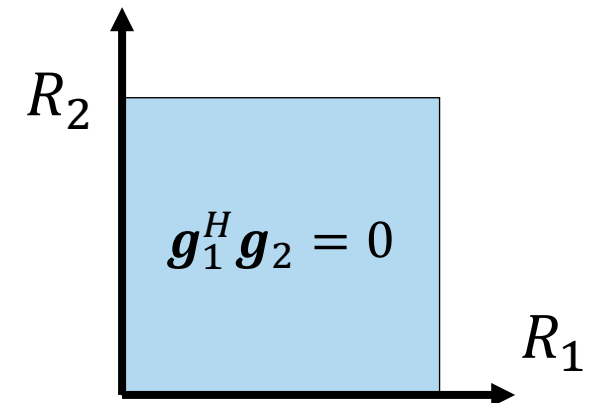
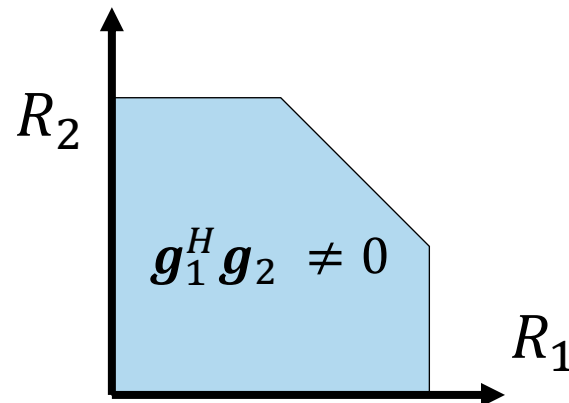
- Simple signal processing
 - Maximum ratio processing
 - Favorable propagation
- Simple capacity lower bound
 - Use-and-then-forget technique
 - Expression with maximum ratio processing

Recall: Sum Capacity with $K = 2$

- Recall: Sum Capacity with $K = 2$ and $\mathbf{G} = [\mathbf{g}_1 \ \mathbf{g}_2]$:

$$\begin{aligned}
 R_1 + R_2 &= \log_2(\det(\mathbf{I}_2 + \rho_{ul} \mathbf{G}^H \mathbf{G})) = \log_2 \left(\det \left(\mathbf{I}_2 + \rho_{ul} \begin{bmatrix} \|\mathbf{g}_1\|^2 & \mathbf{g}_1^H \mathbf{g}_2 \\ \mathbf{g}_2^H \mathbf{g}_1 & \|\mathbf{g}_2\|^2 \end{bmatrix} \right) \right) \\
 &= \log_2 \left((1 + \rho_{ul} \|\mathbf{g}_1\|^2)(1 + \rho_{ul} \|\mathbf{g}_2\|^2) - \rho_{ul}^2 |\mathbf{g}_1^H \mathbf{g}_2|^2 \right) \\
 &\leq \log_2(1 + \rho_{ul} \|\mathbf{g}_1\|^2) + \log_2(1 + \rho_{ul} \|\mathbf{g}_2\|^2)
 \end{aligned}$$

Equality if and only if $\mathbf{g}_1^H \mathbf{g}_2 = 0$



Favorable propagation

- A collection of channel vectors $\{\mathbf{g}_k\}$ are said to offer *favorable propagation* if
$$\mathbf{g}_k^H \mathbf{g}_i = 0, \quad k, i = 1, \dots, K, \quad k \neq i$$

- Never satisfied exactly in practice

- *Asymptotically favorable propagation* if

$$\frac{1}{M} \mathbf{g}_k^H \mathbf{g}_i \rightarrow 0, \quad M \rightarrow \infty, \quad k, i = 1, \dots, K, \quad k \neq i$$

- One cannot physically let $M \rightarrow \infty$ in practice

Law of large numbers

- Consider a sequence X_1, X_2, \dots of independent and identically distributed random variables
 - Assume $E\{X_i\} = \mu$ for $i = 1, 2, \dots$
 - Assume $Var\{X_i\} = \sigma^2 < \infty$ for $i = 1, 2, \dots$

Law of large numbers

The sample average

$$\bar{X}_n = \frac{X_1 + \dots + X_n}{n}$$

converges to the expected value

$$\bar{X}_n \rightarrow \mu \text{ as } n \rightarrow \infty$$

$$Var\{\bar{X}_n\} = \frac{Var\{X_1\} + \dots + Var\{X_n\}}{n^2} = \frac{\sigma^2}{n}$$

Properties of Rayleigh fading channels

- Channels independently distributed as $\mathbf{g}_k \sim \mathcal{CN}(\mathbf{0}, \beta_k \mathbf{I}_M)$

- Offer channel hardening:

$$\frac{1}{M} \|\mathbf{g}_k\|^2 \rightarrow \beta_k, \quad M \rightarrow \infty, \quad k = 1, \dots, K$$

- Offer favorable propagation:

$$\frac{1}{M} \mathbf{g}_k^H \mathbf{g}_i \rightarrow 0, \quad M \rightarrow \infty, \quad k, i = 1, \dots, K, \quad k \neq i$$

Approximations when M is large:

$$\frac{1}{M} \|\mathbf{g}_k\|^2 \approx \beta_k \text{ and } \frac{1}{M} \mathbf{g}_k^H \mathbf{g}_i \approx 0$$

Recall: Estimates of channels

- MMSE estimate of g_k^m from user k to antenna m

- Estimate: $\hat{g}_k^m = E\{g_k^m | \mathbf{Y}'_p\} = \frac{\sqrt{\tau_p \rho_{ul} \beta_k}}{1 + \tau_p \rho_{ul} \beta_k} [\mathbf{Y}'_p]_{mk} \sim CN(0, \gamma_k)$

- Estimation error: $\tilde{g}_k^m = \hat{g}_k^m - g_k^m \sim CN(0, \beta_k - \gamma_k)$
where

$$\gamma_k = \frac{\tau_p \rho_{ul} \beta_k^2}{1 + \tau_p \rho_{ul} \beta_k}$$

- Vector notation: $\hat{\mathbf{g}}_k = \begin{bmatrix} \hat{g}_k^1 \\ \vdots \\ \hat{g}_k^M \end{bmatrix} \sim CN(\mathbf{0}, \gamma_k \mathbf{I}_M), \quad \tilde{\mathbf{g}}_k = \begin{bmatrix} \tilde{g}_k^1 \\ \vdots \\ \tilde{g}_k^M \end{bmatrix} \sim CN(\mathbf{0}, (\beta_k - \gamma_k) \mathbf{I}_M)$

Properties of estimated Rayleigh fading channels

- Estimated channels independently distributed as $\hat{\mathbf{g}}_k \sim CN(\mathbf{0}, \gamma_k \mathbf{I}_M)$

- Offer channel hardening:

$$\frac{1}{M} \|\hat{\mathbf{g}}_k\|^2 \rightarrow \gamma_k, \quad M \rightarrow \infty, \quad k = 1, \dots, K$$

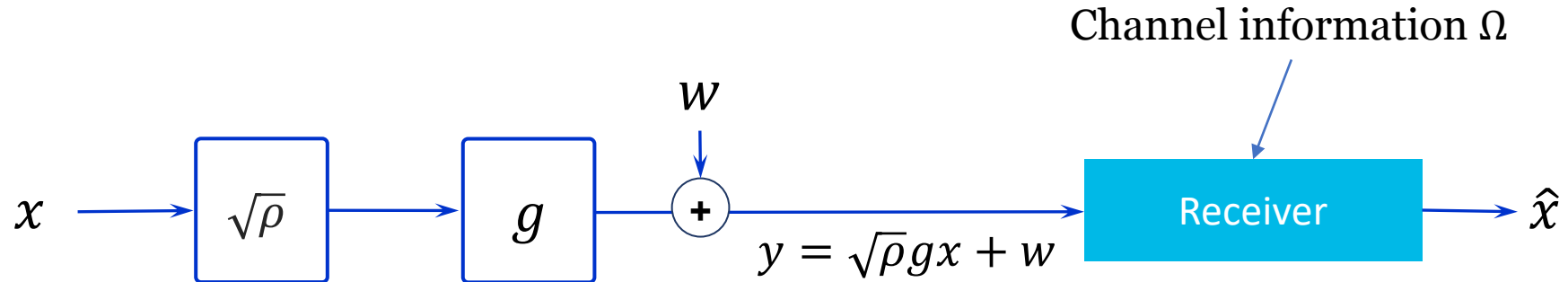
- Offer favorable propagation:

$$\frac{1}{M} \hat{\mathbf{g}}_k^H \hat{\mathbf{g}}_i \rightarrow 0, \quad M \rightarrow \infty, \quad k, i = 1, \dots, K, \quad k \neq i$$

Approximations when M is large:

$$\frac{1}{M} \|\hat{\mathbf{g}}_k\|^2 \approx \gamma_k \text{ and } \frac{1}{M} \hat{\mathbf{g}}_k^H \hat{\mathbf{g}}_i \approx 0$$

Recall: Capacity lower bound

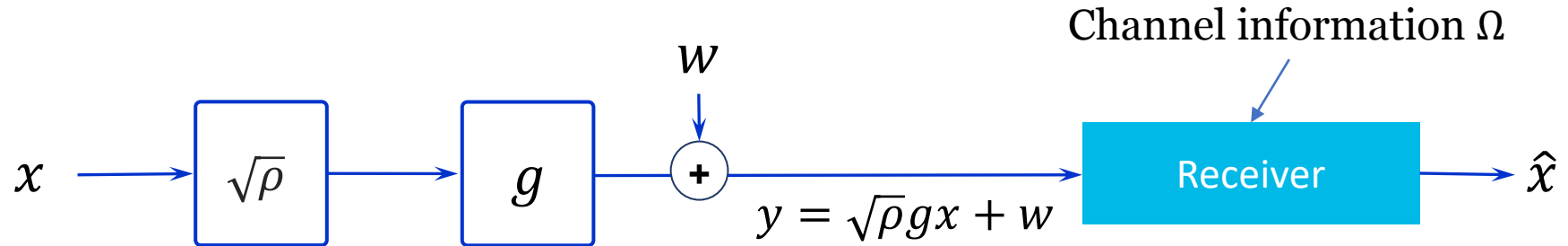


- Desired signal x , power ρ , and g and w uncorrelated
- Channel coefficient g , known channel information Ω

Capacity lower bound:

$$C \geq E \left\{ \log_2 \left(1 + \frac{\rho |E\{g|\Omega\}|^2}{\rho \text{Var}\{g|\Omega\} + \text{Var}\{w|\Omega\}} \right) \right\}$$

Capacity lower bound with deterministic channel



- Desired signal x with unit power, transmit power ρ
- Deterministic and known channel coefficient g ($\Omega = \{g\}$)

Capacity lower bound:

$$C \geq E \left\{ \log_2 \left(1 + \frac{\rho |E\{g|\Omega\}|^2}{\rho \text{Var}\{g|\Omega\} + \text{Var}\{w|\Omega\}} \right) \right\} = \log_2 \left(1 + \frac{\rho |g|^2}{\text{Var}\{w\}} \right)$$

Revisiting the received uplink signal

- Received signal:

$$\mathbf{y} = \sum_{k=1}^K \mathbf{g}_k \sqrt{\rho_{ul}\eta_k} q_k + \mathbf{w} = \underbrace{\sum_{k=1}^K \hat{\mathbf{g}}_k \sqrt{\rho_{ul}\eta_k} q_k}_{\text{Useful part}} - \underbrace{\sum_{k=1}^K \tilde{\mathbf{g}}_k \sqrt{\rho_{ul}\eta_k} q_k}_{\mathbf{w}' : \text{Unusable part}} + \mathbf{w}$$

- Assign receiver filter \mathbf{a}_i for user i :

$$\mathbf{a}_i^H \mathbf{y} = \underbrace{\mathbf{a}_i^H \hat{\mathbf{g}}_i \sqrt{\rho_{ul}\eta_i} q_i}_{\text{Desired part}} + \underbrace{\sum_{k=1, k \neq i}^K \mathbf{a}_i^H \hat{\mathbf{g}}_k \sqrt{\rho_{ul}\eta_k} q_k}_{\text{Interference}} + \mathbf{a}_i^H \mathbf{w}'$$

Focusing on the desired part

- Consider $\hat{\mathbf{g}}_i \sim \mathcal{CN}(\mathbf{0}, \gamma_i \mathbf{I}_M)$
- Which value of \mathbf{a}_i maximizes the ratio $\frac{|\mathbf{a}_i^H \hat{\mathbf{g}}_i|}{\|\mathbf{a}_i\|}$?

Cauchy-Schwartz inequality

$$\frac{|\mathbf{a}_i^H \hat{\mathbf{g}}_i|}{\|\mathbf{a}_i\|} \leq \frac{\|\mathbf{a}_i\| \|\hat{\mathbf{g}}_i\|}{\|\mathbf{a}_i\|} = \|\hat{\mathbf{g}}_i\|$$

with equality if $\mathbf{a}_i = c \hat{\mathbf{g}}_i$ for some constant $c \neq 0$

- We now call $\mathbf{a}_i = c \hat{\mathbf{g}}_i$ maximum ratio (MR) processing
 - Same thing as MRC for $c = 1/\|\hat{\mathbf{g}}_i\|$

Received signal when using MR processing

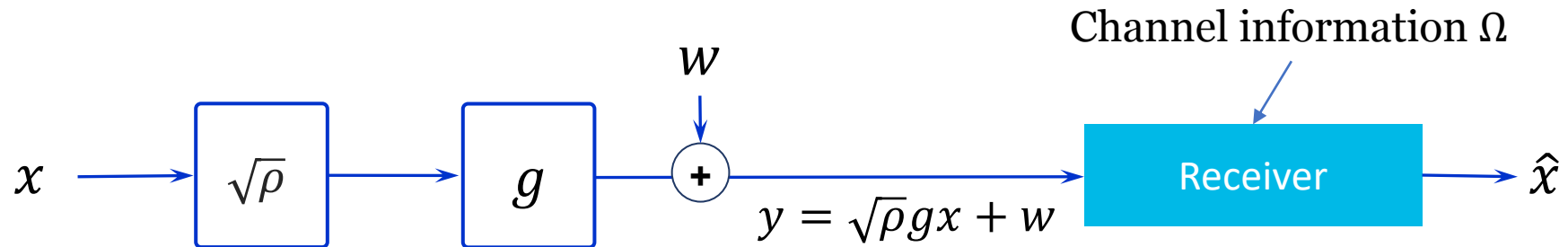
- Set $\mathbf{a}_i = \frac{1}{M} \hat{\mathbf{g}}_i$:

$$\mathbf{a}_i^H \mathbf{y} = \underbrace{\frac{\hat{\mathbf{g}}_i^H \hat{\mathbf{g}}_i}{M}}_{\approx \gamma_i} \sqrt{\rho_{ul} \eta_i} q_i + \sum_{k=1, k \neq i}^K \underbrace{\frac{\hat{\mathbf{g}}_i^H \hat{\mathbf{g}}_k}{M}}_{\approx 0} \sqrt{\rho_{ul} \eta_k} q_k + \underbrace{\frac{\hat{\mathbf{g}}_i^H \mathbf{w}'}{M}}_{\approx 0}$$

- Use-and-then-forget technique:

$$\mathbf{a}_i^H \mathbf{y} = \underbrace{\gamma_i \sqrt{\rho_{ul} \eta_i} q_i}_{\text{Desired part with deterministic channel!}} + \underbrace{\left(\frac{\hat{\mathbf{g}}_i^H \hat{\mathbf{g}}_i}{M} - \gamma_i \right) \sqrt{\rho_{ul} \eta_i} q_i + \sum_{k=1, k \neq i}^K \frac{\hat{\mathbf{g}}_i^H \hat{\mathbf{g}}_k}{M} \sqrt{\rho_{ul} \eta_k} q_k + \frac{\hat{\mathbf{g}}_i^H \mathbf{w}'}{M}}_{w: \text{Uncorrelated interference and noise}}$$

Using the capacity bound with deterministic channel



- Desired signal $x = q_i$, transmit power $\rho = \rho_{ul}\eta_i$
- Deterministic and known channel coefficient $g = \gamma_i$

Capacity lower bound:

$$C \geq \log_2 \left(1 + \frac{\rho |g|^2}{\text{Var}\{w\}} \right)$$

$$\rho |g|^2 = \rho_{ul} \eta_i \gamma_i^2$$

$$\text{Var}\{w\} = \dots = \frac{\gamma_i}{M} \left(\sum_{k=1}^K \rho_{ul} \eta_k \beta_k + 1 \right)$$

Capacity bound with MR and use-and-then-forget technique

$$C \geq \log_2 \left(1 + \frac{M \rho_{ul} \eta_i \gamma_i}{\sum_{k=1}^K \rho_{ul} \eta_k \beta_k + 1} \right)$$

- Interpretation
 - Small-scale fading is not visible in this bound
 - **Numerator:**
Coherent beamforming gain, grows with antennas M , power $\rho_{ul} \eta_i$ and estimation quality γ_i
 - **Denominator:**
Sum of non-coherent interference from all users plus noise variance

Summary

- When having many antennas:
 - Channel hardening: $\frac{1}{M} \|\hat{\mathbf{g}}_k\|^2 \approx \gamma_k$
 - Favorable propagation: $\frac{1}{M} \hat{\mathbf{g}}_k^H \hat{\mathbf{g}}_i \approx 0$
 - Motivation for maximum ratio processing
- Use-and-then-forget technique
 - Pretend as if channel is deterministic
 - Compute a closed-form capacity lower bound

End of Lecture 8

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