

9/3/23

(1) Primal Simplex

The problem starts at a basic feasible solⁿ. Successive iterations continue to be feasible until optimality is reached in the last iteration.

(2) Dual Simplex

L.P starts at a better than optimal solⁿ which is infeasible. Successive iterations continue to be optimal and feasibility is restored at the last iteration.

Dual Simplex Algorithm

(i) Feasibility:

The leaving var. x_r is the basic var. having the most -ve value (ties are broken arbitrarily). If all the basic var. are non-ve, algorithm ends.

(ii) Optimality condⁿ:

Given that x_r is the leaving var., let \bar{c}_j be the reduced cost of non-basic var. x_j and α_{rj} is the constraint coeff. in x_r row and x_j column. The entering var. is the non-basic var. with $\alpha_{rj} < 0$.

$$\min_{x_j \text{ non-basic}} \left\{ \begin{array}{l} \bar{c}_j \\ \alpha_{rj} \end{array} ; \alpha_{rj} < 0 \right\}$$

If $\alpha_{ij} \geq 0 \quad \forall$ non-basic x_j , the problem has infeasible solⁿ.

The L.P model starts with 2 requirements:
(i) obj. f^n must satisfy the optimality condⁿ of regular simplex method.

(ii) All constraints should be \leq . If the constraints are of \geq type then we multiply both sides by -1 to convert to equality.
If constraint is an equality then it is converted to this as eg: $x_1 + x_2 = 1$

$$\hookrightarrow \begin{cases} x_1 + x_2 \leq 1 \\ x_1 + x_2 \geq 1 \rightarrow -x_1 - x_2 \leq -1 \end{cases}$$

Dual Simplex

(Q.) Min $z = 3x_1 + 2x_2 + x_3$

subject to $3x_1 + x_2 + x_3 \geq 3$
 $-3x_1 + 3x_2 + x_3 \geq 6$
 $x_1 + x_2 + x_3 \leq 3$
 $x_1, x_2, x_3 \geq 0$

Ans.)

$-3x_1 - x_2 - x_3 \leq -3$
 $3x_1 - 3x_2 - x_3 \leq -6$
 $x_1 + x_2 + x_3 \leq 3$

$x_1, x_2, x_3 \geq 0$

* All consters. are \leq , \therefore starting basic var. will be slack var.

Basic	x_1	$x_2 \downarrow$	x_3	s_1	s_2	s_3	Sol ⁿ
Z	-3	-2	-1	0	0	0	0
s_1	-3	-1	-1	1	0	0	-3
$\leftarrow s_2$	3	-3	-1	0	1	0	-6
s_3	1	1	1	0	0	1	3
Z	-5	0	-1/3	0	-2/3	0	4
$\leftarrow s_1$	-4	0	-2/3	1	-1/3	0	-1
x_2	-1	1	1/3	0	-1/3	0	2
s_3	2	0	2/3	0	1/3	1	1

* Its a minimiz. prob. and obj. fⁿ coeff. are -ve which implies optimality.



$$\min_{\text{non-basic } x_j} \left\{ \frac{\bar{E}_j}{\alpha_{rj}} ; \alpha_{rj} < 0 \right\}$$

Non-basic var.'s (1st iteration)

$$x_1: \left| \frac{-3}{3} \right| = X \quad (\text{as } \alpha_{rj} \text{ not } < 0)$$

$$x_2: \left| \frac{-2}{-3} \right| = \frac{2}{3} \rightarrow \min. \therefore \text{it enters } (x_2)$$

$$x_3: \left| \frac{-1}{-1} \right| = 1$$

Non-basic var.'s (2nd iteration)

$$x_1: \left| \frac{-5}{-4} \right| = \frac{5}{4} //$$

$$x_3: \left| \frac{-1/3}{-2/3} \right| = \frac{1}{2} //$$

$$s_2 = \left| \frac{-2/3}{-1/3} \right| = 2 //$$

	x_1	x_2	x_3	s_1	s_2	s_3	Sol ⁿ
z	-3	0	0	-1/2	-1/2	0	9/2
x_3	6	0	1	-3/2	1/2	0	3/2
x_2	-3	1	0	1/2	-1/2	0	5/2
s_3	-2	0	0	1	0	1	0

\therefore optimal & feasible.

13/3/23

(Q7) Min $z = 5x_1 + 6x_2$

s.t $x_1 + x_2 \geq 2$

$4x_1 + x_2 \geq 4$

Use dual simplex.

Ans.)

$-x_1 - x_2 \leq -2$

$-4x_1 - x_2 \leq -4$

Basic	$x_1 \downarrow$	x_2	s_1	s_2	Soln
z	-5	-6	0	0	0
s_1	-1	-1	1	0	-2
$\leftarrow s_2$	<u>-4</u>	-1	0	1	<u>-4</u>
z	0	-19/4	0	-5/4 \downarrow	5
$\leftarrow s_1$	0	-3/4	<u>1</u>	<u>-1/4</u>	<u>-1</u>
x_1	1	1/4	0	-1/4	1
z	0	-1	-5	0	10
s_2	0	3	-4	1	4
x_1	1	1	-1	0	2

$\therefore z = 10$

$x_1 = 2, x_2 = 0, s_1 = 0, s_2 = 4$

\therefore optimal & feasible.

Iteration - 1

$$x_1 = \left| \frac{-5}{-4} \right| = \left(\frac{5}{4} \right) \rightarrow \text{entering var}, \quad x_2 = \left| \frac{-6}{-1} \right| = \underline{\underline{6}}$$

Iteration - 2

$$x_2 = \left| \frac{-19/4}{-3/4} \right| = \frac{19}{3} //, \quad s_2 = \left| \frac{-5/4}{-1/4} \right| = \left(\underline{\underline{5}} \right) \checkmark$$

(Q.) Use dual simplex.

$$\text{Min } z = 4x_1 + 2x_2$$

$$\text{s.t.} \quad x_1 + x_2 = 1$$

$$3x_1 - x_2 \geq 2$$

$$\boxed{x_1 + x_2 \leq 1}$$

$$+x_1 + x_2 \geq 1 \Rightarrow \boxed{-x_1 - x_2 \leq -1}$$

$$\boxed{-3x_1 + x_2 \leq -2}$$

Iteration - 1 $\rightarrow x_1 = \left| \frac{-4}{-3} \right| = \left(\frac{4}{3} \right) \checkmark, \quad x_2 = \left| \frac{-2}{1} \right| = \left(\frac{2}{1} \right) \times > 0$

Iteration - 2 $\rightarrow s_3 = \left| \frac{-4/3}{-1/3} \right| = \underline{\underline{4}}, \quad x_2 = \left| \frac{-10/3}{-4/3} \right| = \left(\underline{\underline{\frac{5}{2}}} \right) \checkmark$

Basic	$x_1 \downarrow$	x_2	s_1	s_2	s_3	Sol ⁿ
Z	-4	-2	0	0	0	0
s_1	1	1	1	0	0	1
s_2	-1	-1	0	1	0	-1
$\leftarrow s_3$	-3	1	0	0	1	-2
Z	0	$-10/3 \downarrow$	0	0	$-4/3$	$8/3$
s_1	0	$4/3$	1	0	$+1/3$	$1/3$
$\leftarrow s_2$	0	$-4/3$	0	1	$-1/3$	$-1/3$
x_1	1	$-1/3$	0	0	$-1/3$	$2/3$
Z	0	0	0	$-5/2$	$-1/2$	$3\frac{1}{2}$ $7/2$
s_1	0	0	1	1	0	0
x_2	0	1	0	$-3/4$	$1/4$	$1/4$
x_1	1	0	0	$-1/4$	$-1/4$	$3/4$

\therefore optimal & feasible

$$\therefore Z = 7/2$$

$$\underline{s_2 = s_3 = s_1 = 0}, \quad \underline{x_2 = 1/4}, \quad \underline{x_1 = 3/4}$$