

Data Structures

Heap Insertion

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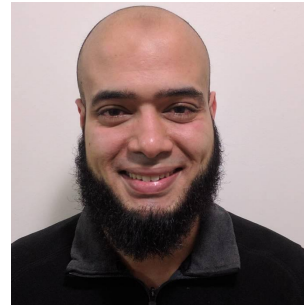
Teaching, Training and Coaching since more than a decade!

Artificial Intelligence & Computer Vision Researcher

PhD from Simon Fraser University - Canada

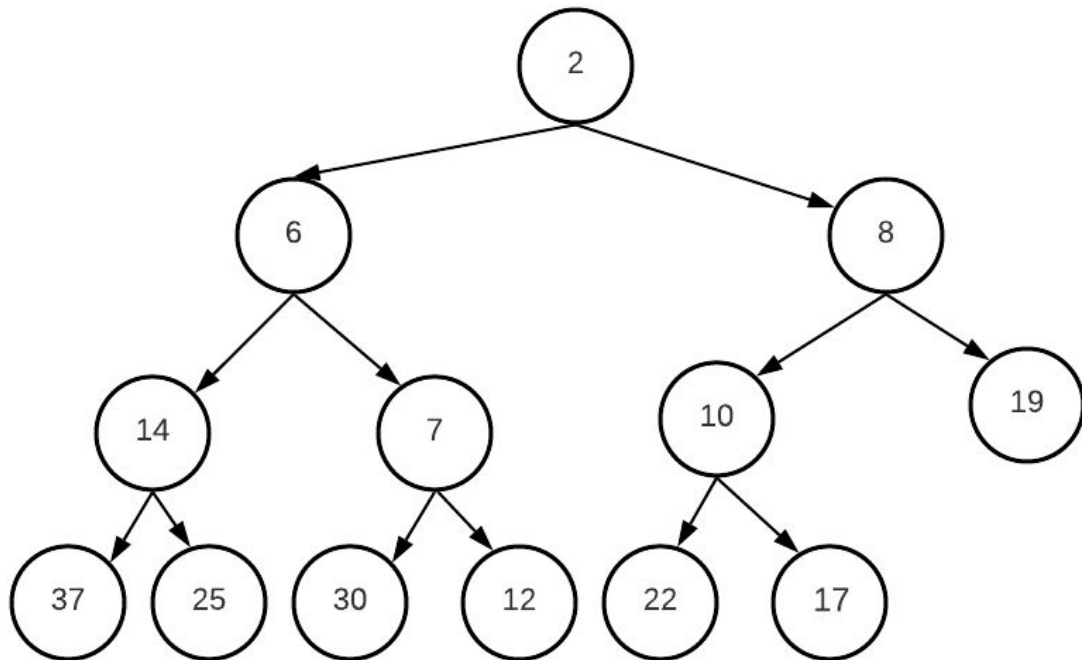
Bachelor / Msc from Cairo University - Egypt

Ex-(Software Engineer / ICPC World Finalist)



Let's insert 5

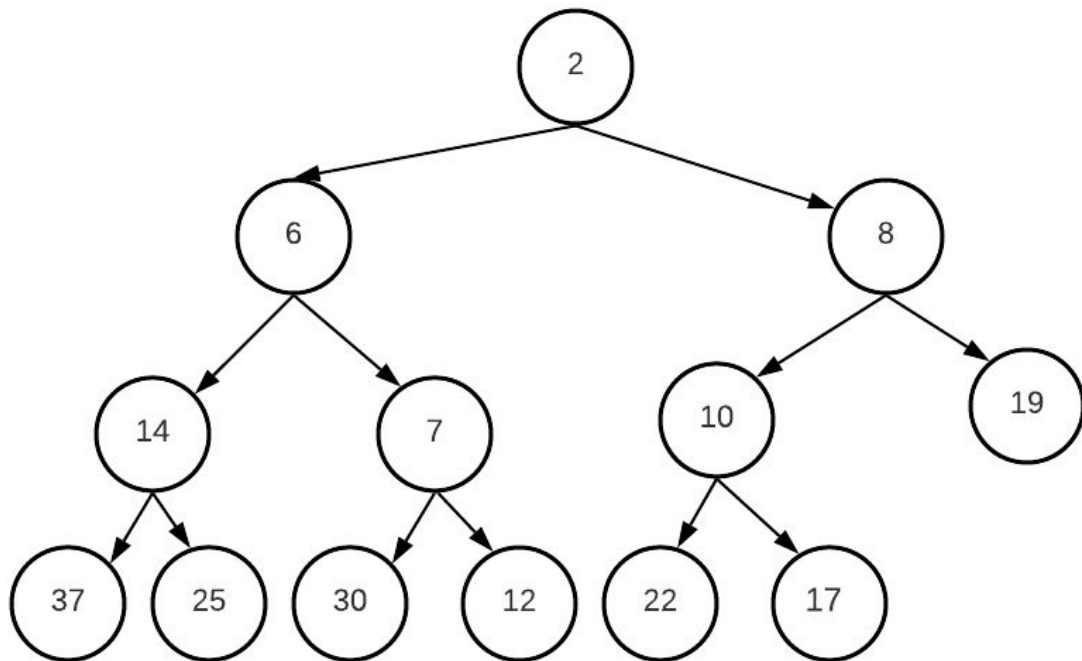
- To insert a value in heap, we depend on a trick a nice idea
- We **add** to the tree the item, and then **fix** the corruption
 - A smart approach, but sadly not widely applicable



0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
2	6	8	14	7	10	19	37	25	30	12	22	17		

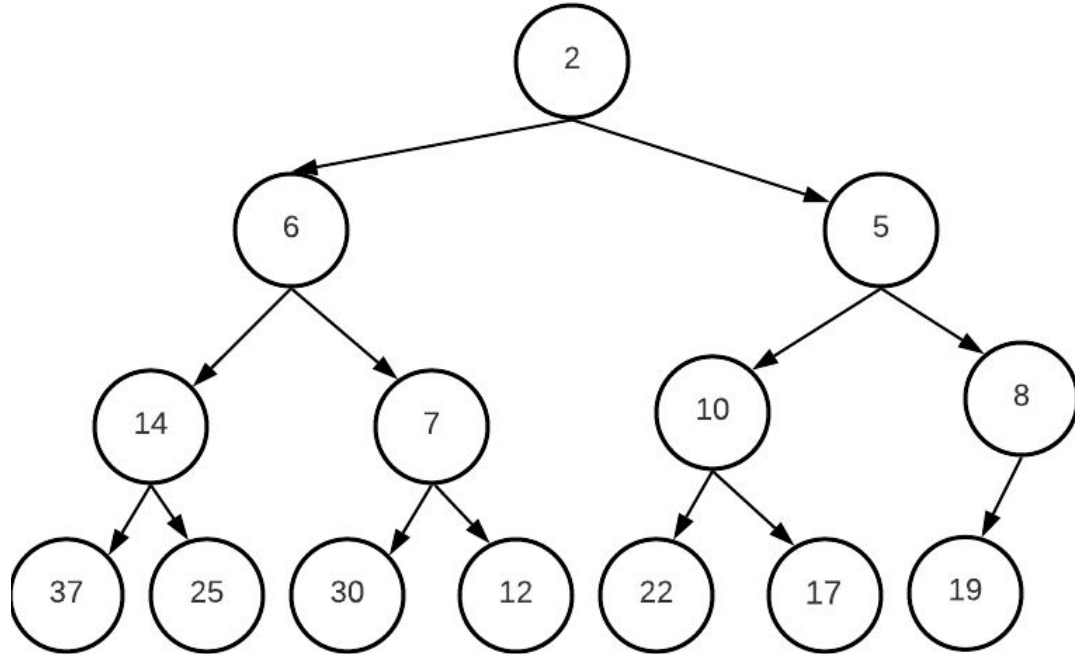
Let's insert 5

- What is the first available node in this tree?
 - Left of 19
- Add 5
- What is the chain of its parents
 - [5, 19, 8, 2]
- Shift up 5 to be in its **right location** (decreasing seq)
 - [19, 8, 5, 2]



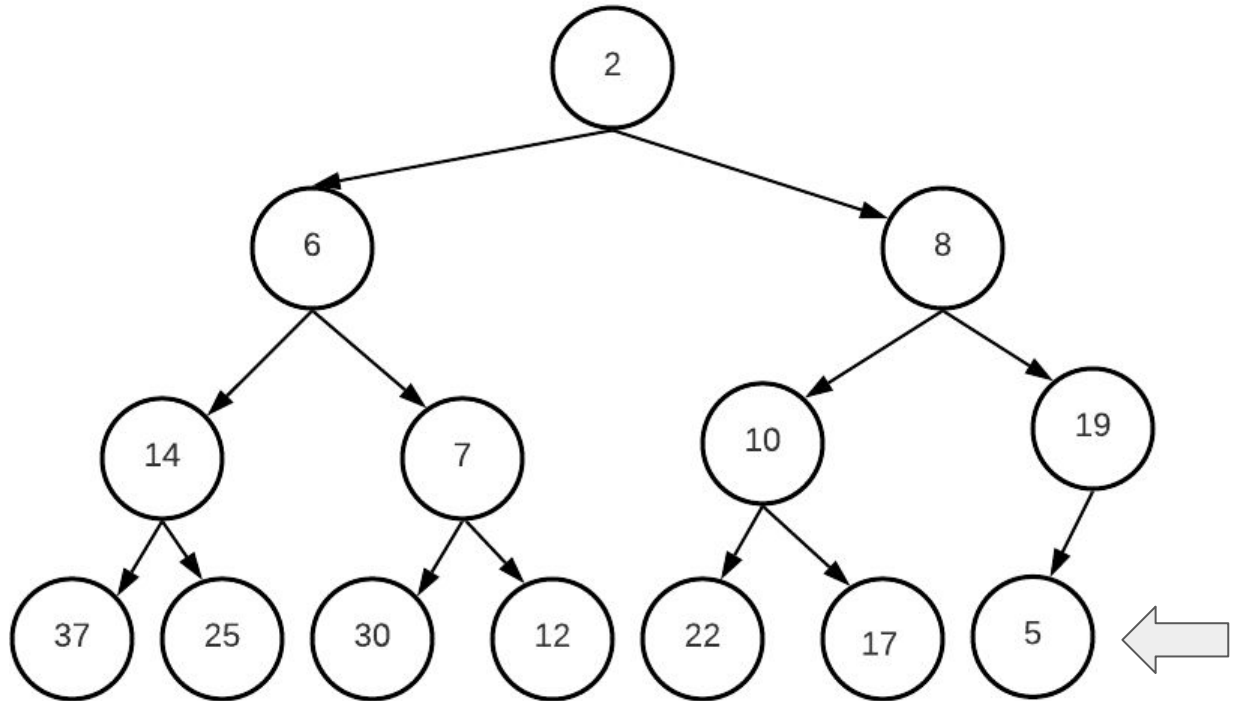
Heapify Up

- How to simply implement that?
- Start from the new added node index
- As long as its parent is greater, shift the parent to the current child node and move up
 - Observe: 19 and 8 are moved down



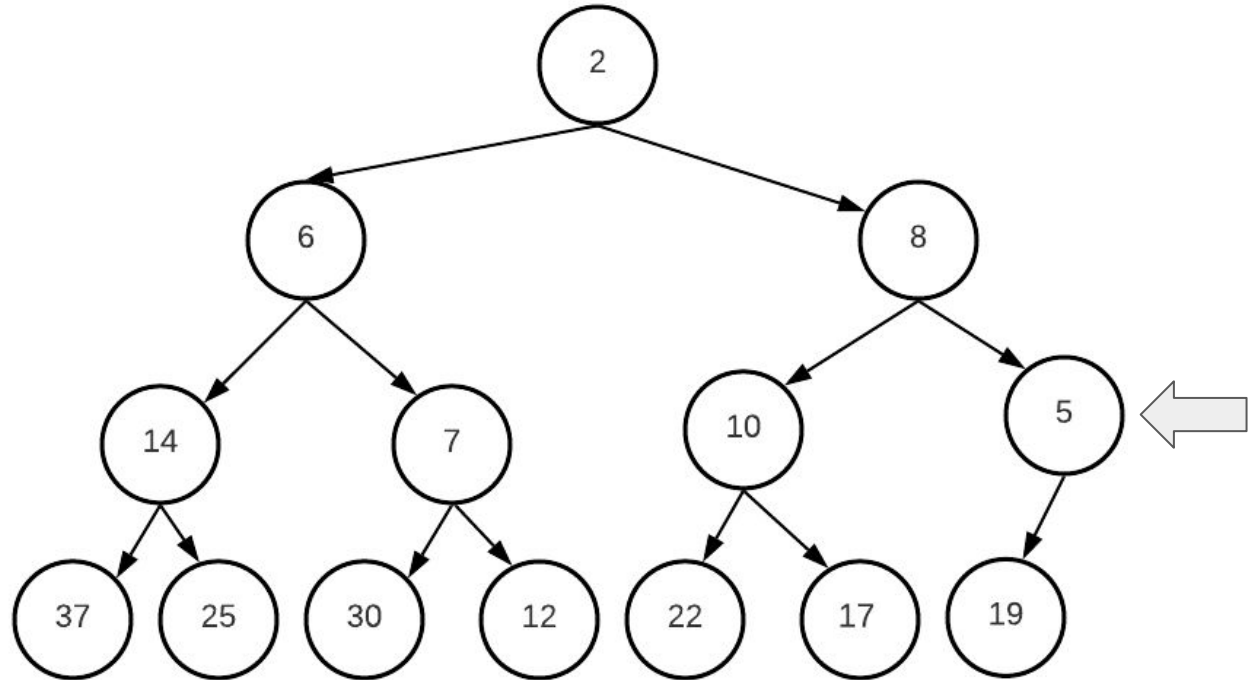
Simulation: Add 5 at available position

- 5 index = 13
 - Parent idx = 6
 - Parent value 19
- $5 < 19$
 - Push 19 down



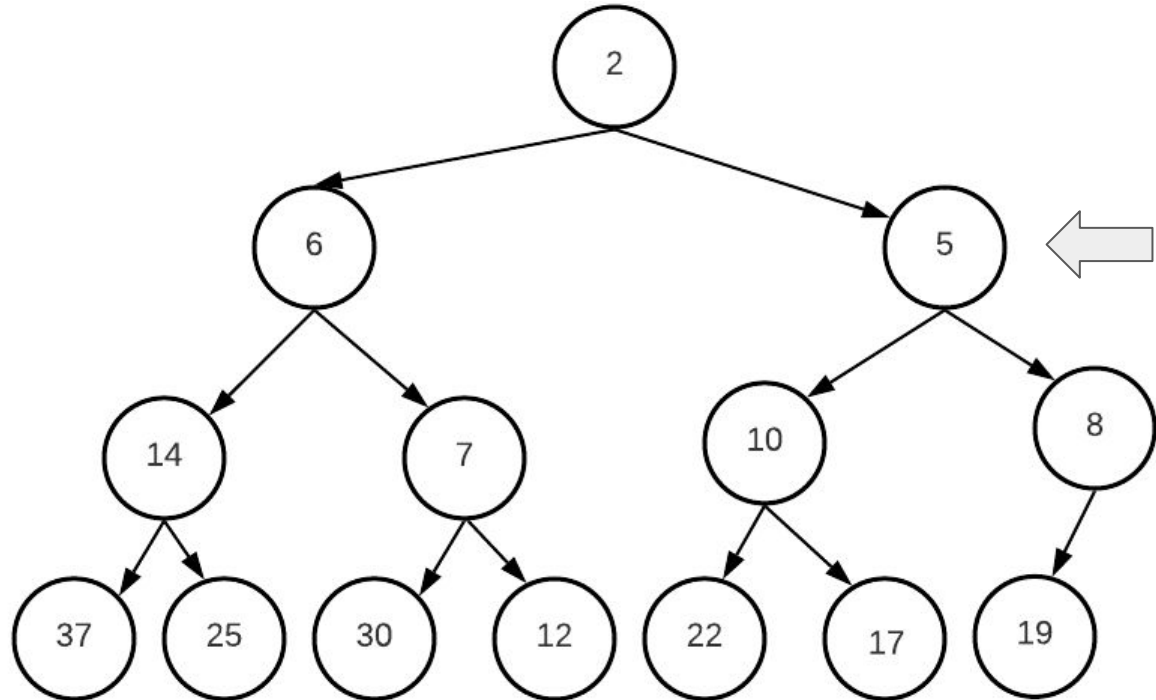
Simulation: Compare with parent

- 5 index = 6
 - Parent idx = 2
 - Parent value 8
- $5 < 8$
 - Push 8 down



Simulation: Compare with parent

- 5 index = 2
 - Parent idx = 0
 - Parent value 2
- $5 > 2$
 - Perfect heap
 - Stop
- Take 10 min to code it



Min Heap Class

- Let's create a class
- Assume for simplicity some internal capacity
- left, right, parent functions added from last time

```
6 class MinHeap {  
7     private:  
8         int *array { };  
9         int size { };  
10        int capacity { 1000 };  
11  
12    public:  
13        MinHeap() {  
14            array = new int[capacity] { };  
15            size = 0;  
16        }  
17  
18        ~MinHeap() {  
19            delete[] array;  
20            array = nullptr;  
21        }
```


Insertion Implementation

- Add element in the end of array, then refix this last position
- As tree is complete, its height is $O(\log(n))$
- Given n elements to insert in a heap, we need $O(n\log(n))$

```
void heapify_up(int child_pos) {
    // stop when parent is smaller (or no parent)
    int par_pos = parent(child_pos);
    if (child_pos == 0 || array[par_pos] < array[child_pos])
        return;

    swap(array[child_pos], array[par_pos]);
    heapify_up(par_pos);
}

void push(int key) {
    assert(size + 1 <= capacity);
    array[size++] = key;
    heapify_up(size - 1);
}

int top() {
    assert(!isempty());
    return array[0];
}
```

You turn

- What if we want to remove the smallest element?
- It is the root
- How can we fix the tree? Think in a similar top-down procedure

“Acquire knowledge and impart it to the people.”

“Seek knowledge from the Cradle to the Grave.”