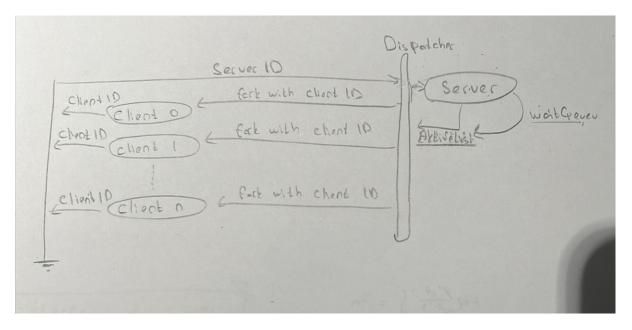
CSE 344 MIDTERM

Project description

This project aims to develop server-client relationships with using fork to serve a lot of clients and sheared memory to communicate clients.

Project Design



Firstly I design server side and client side with connect each other via Server PID. Wen I create a server in a terminal I take server PID and create a client with that. Server clients write own client PID to Fifo that named server PID. After that if server accept the client, server create a fork. This fork serve to spesific client through client Fifo named client PID.

Therefore server add user activeUsers list. Of course this list has a limit that giving by user when server created. If activeUsers list is full, server add the client to waitQueue lists. (I use list structure instead of queue structure but its work same as queue.) I use list because client who stayed in the middle of queue can be kill itself. If it is does, we cannot erase from queue. So logic is queue but structure is list.

My implementation

There are some screen shot for server and clients for their communication.

Client.h; this header file show my design basically. So I share this screenshot.

activeUser and waitQueue structure; User can be inside of the activeUser list or waitQueue. If conditions are available to adding new client to server, it add activeUser list. Otherwise client should be wait in waitQueue. There are headers of the functions to demonstrate struct of the design.

```
void initClientLists(); // Initialize the active clients and waiting queue
void addClientToList(Client **list, int *count, int *capacity, Client newClient); // Add a client to the list
void addClientToWaiting(Client newClient); // Add a client to the active clients list
void addClientToWaiting(Client newClient); // Add a client to the waiting queue
int removeClientFromList(Client *list, int *count, int pid); // Remove a client from the list
void activateClientFromQueue(); // Activate a client from the waiting queue

void manageClient(char* clientReturnFifo); // Manage the client
char* getLine(FILE *file, int lineNumber); // Get a specific line from a file
```

Here are a lot user connect same server simultaneously;

This server side main loop; read server fifo and if detected clients try to add active users otherwise add wait queue.

```
mile[1];
// open fifo for reading connection
int serverResponsed = open(mainFifo, 0_ROOMLY);
// open fifo for reading connection
int serverResponsed = open(mainFifo, 0_ROOMLY);
// Road client PID
int clientPid;
// read = consection =
```

This client main loop; after connecting server it wait accepted response. When client take this response it has 2 operation these sending request and taking response.

Client doesn't know anything about management of server side. It is just use server properties and server's allows.

```
char request[1024], response[1024];
while(1){

// Read request from user
memset(request, 0, sizeof(request); // clear the tempMssg
printf(">> Enter comment: ");
fgets(request, sizeof(request), stdin);

requestFd = open(clientFifoName, 0_WRONLY); // open client fifo for request from server
if(requestFd < 0){
perror("client open fifo failed");
return -1;
}

// Send request to server
if (write(requestFd, &request, sizeof(request)) == -1){
perror("server requesting error.");
return -1;
}

close(requestFd);

// Read response from server
responseFd = open(clientFifoName, 0_RDONLY); // open client fifo for reading response from server
if(responseFd < 0){
perror("client open fifo failed");
return -1;
}

// Read response from server for ensure connection is okay
while (read(responseFd, &response, sizeof(response)) > 0){
printf("%s\n", client.response.content); // clear the tempMssg
}

close(repsonseFd);
}

sumlink(clientFifoName).
```

Semaphores and shared memory; We use a lot of clients in the server. Therefore they are use server properties commonly. In addition for their management server fork should know their states(active or not). Therefore we use same memory places for some variables. (all process has own memory address spaces)

Of course, this situation cause a race condition. For solving them we combine shared memory with semaphores. I used activeUserCounter in server side. When user in there I increase this value, otherwise decrease. For doing that I create semaphore and below some screenshot to show that.

```
//// shared memory Design with semaphoers///////
int shm fd = shm_open(counterSharedMemory, 0_CREAT | 0_RDWR, 0666); // define shared memory file descriptor

if (shm_fd == -1) {
    perror("shm_open");
    return EXIT_FAILURE;
}

if (ftruncate(shm_fd, sizeof(int)) == -1) { // set size of shared memory
    perror("ftruncate");
    return EXIT_FAILURE;
}

sharedSequentialClientCounter = mmap(NULL, sizeof(int), PROT_READ | PROT_WRITE, MAP_SHARED, shm_fd, 0); // map shared memory

if (sharedSequentialClientCounter == MAP_FAILED) {
    perror("mmap");
    return EXIT_FAILURE;
}

**sharedSequentialClientCounter = 0; // initialize shared memory

// create semaphore

sem = sem_open(semaphoreName, 0_CREAT, 0666, 1);

if (sem == SEM_FAILED) {
    perror("sem_open");
    return EXIT_FAILURE;
}

//// shared memory Design with semaphoers////////
```

```
/////// shared memory Design with semaphoers Constants////////
int *sharedSequentialClientCounter; // shared memory for sequentialClientCounter

const char *semaphoreName = "mysemaphore"; // semaphore name

sem_t *sem; // define semaphore

const char *counterSharedMemory = "/clientCounter"; // shared memory name

/////// shared memory Design with semaphoers Constants////////

void handleInterruptSignal(int sig) {
    printf("SIGTSTP received, unlinking FIFO...\n");

// Unlink the FIFO

if (unlink(mainFifo) == -1) { // if already unlinked, it will return -1
    perror("Failed to unlink FIFO");
}

// kill clients and child process
for (int i = 0; i < activeCount; i++) {
    kill(activeClients[i].clientPid, SIGKILL);
}

// release shared memory
munmap(sharedSequentialClientCounter, sizeof(int));
shm_unlink(counterSharedMemory);

// release semaphores
sem_close(sem);
sem_unlink(semaphoreName);
kill(getpid(), SIGKILL); // kill the server process

kill(getpid(), SIGKILL); // kill the server process
}
```

Here are definition of semaphores **with signal interrupt**; Signal interrupts have critical role because when client forcely quit terminal server should know that.

When server send a kill request to **kill** server, server should kill whole child process (clients), otherwise they are hang. Also this is cause optimization problems. Client signal shown below, too.

```
void handleInterruptSignal(int sig) {
    printf("SIGTSTP received, unlinking FIFO...\n");

// send quit signal to server for killing the client serverside
int requestFd = open(clientFifoName, 0_WRONLY); // open client fifo for request from server
if(requestFd < 0){
    perror("In signal client open fifo failed");
    return;
}

// Send request to server
if (write(requestFd, "quit", sizeof("quit")) == -1){
    perror("Server requesting error.");
    return;
}

close(requestFd);

// Unlink the FIFO
if (unlink(clientFifoName) == -1) { // if already unlinked, it will return -1
    perror("Failed to unlink FIFO");
}

printf("killed cleint");
kill(getpid(), SIGKILL); // kill the client process
exit(EXIT_SUCCESS);
}
</pre>
```

Some utilities

Server have a lot of utilities like list directories, help about properties, upload/download, quit and kill server etc.

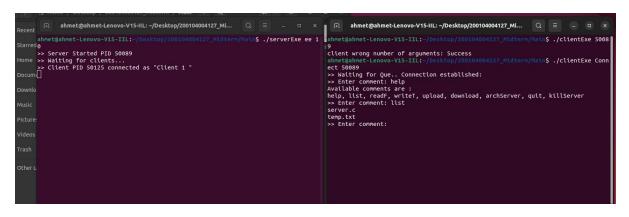
Here are example of help and list request code block and their terminal results;

Help Codeblock;

```
// - help: display the list of possible client request
if (strrmp(nendeq[0], "help") = 0) {
// Help request
if (newReq[1] == NULL){
snpintf(lempAssg, sizeof(tempAssg), "Available comments are :\nhelp, list, readF, writeT, upload, download, archServer, quit, killServer");
printf("burasi temp ness] %s\n", tempAssg);
} clies iffstrcap(newGeq[0], "readF") == 0){
printf("burasi temp ness] %s\n", tempAssg);
} clies iffstrcap(newGeq[1], "readF") == 0){
printf("annas");
snpintf(tempAssg, sizeof(tempAssg), "readF -file> + Nr requests to display the # line of the <file>, if no line number is given \n the whole contents of the file is requested (and displayed on the client side)");
} clies iffstrcap(newGeq[1], "witeT") == 0){
strcpy(tempAssg, "writeT <file> + Strcpy(tempAssg, "world of the wuploads the file from the current working directory of client to the Servers directory \n(beware of the cases no file in clients current working directory and file with the same\n name on Servers side()");
} clies iffstrcap(newGeq[1], "wolldow) == 0){
strcpy(tempAssg, "download <file> \nrequest to receive <file> from Servers directory to client side");
} clies iff(newGeq[1] = "killServer") == 0}{
strcpy(tempAssg, "archServer <fleName>.tar \n Using fork, exce and tar utilities create a child process that will collect all the files currently \navailable on the the Server side and store them in the <filename>.tar archive");
} clies iff(newGeq[1] == "killServer"){
strcpy(tempAssg, "killServer \nSends a kill request to the Server');
} clies iff(newGeq[1] == "quit"){
strcpy(tempAssg, "quit \nSend write request to Server side log file and quits");
} if (clientTd < 0) {
perror('server open fifo failed');
exit(ExIT_FALUBE);
} if (write(clientFd, StempAssg) sizeof(tempAssg)) < 0) {
perror('server write failed');
e
```

List Codeblock;

Their Terminal Results;



Release, free and kill; These process save us to leak memory

For managing whole operation in a same root directory I create makefile.

```
Main > M makefile

1 All: compile run clean

2 compile: server/server.c client/client.c

4 @gcc server/server.c -o serverExe

5 @gcc client/client.c -o clientExe

6 run:

8 ./serverExe neHosServer 1 & ./clientExe Connect 10

9 10 clean:

11 @rm -f *.txt exe
```