

Computer Architecture Lab Exam

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Implementation details:

I followed the same pseudocode given in the question. I maintained 3 loops, and 4 counters

- Row_index_A
- Column_index_A
- Row_index_B
- Column_index_B

Now, to find the address of the required $A[i][k]$ and $B[k][j]$ as per the pseudocode, I did the following:

- We know that the values are given in row major format, which means that we could convert this 2d matrix to 1d array.
- To do so, we would need the indexing in 1D : $\text{array_A}[\text{current}] = i * \text{colA} + k$, where colA refers to the number of columns in A, and i, k are the current indexes.
- Similarly I have done for matrix B as well
- Now to access the value of $A[i][k]$, I have basically added the base address of A with $8 * \text{array_A}[\text{current}]$
- Similarly for B

To handle error cases:

- If colA not equal rowB, then error
- If any of rows or columns are 0, then error

For testing I followed the following tests:

1. .dword 3, 4, 4, 2, 3, 1, 2, 4, 1, 1, 1, 2, 2, 3, 1, 1, 1, 2, 2, 2, 3, 4, 4, 1 -> Given in the question
2. .dword 2, 2, 2, 2, 1, 1, 1, 1, 1, 1, 1, 1
3. .dword 3, 3, 3, 3, 1, 0, 0, 0, 1, 0, 0, 0, 1, 1, 0, 0, 0, 1, 0, 0, 0, 1 -> Unit matrix
4. .dword 3, 3, 4, 3, 1, 0, 0, 0, 1, 0, 0, 0, 1, 1, 0, 0, 0, 1, 0, 0, 0, 1 -> Error case
5. .dword 0, 3, 3, 3, 1, 0, 0, 0, 1, 0, 0, 0, 1, 1, 0, 0, 0, 1, 0, 0, 0, 1 -> Error case

Part-2

Below attached photo shows the configuration of the cache as specified by the question.

