

# Linear Block Code

Course Title: Computer Networks



**Dr. Nazib Abdun Nasir**  
**Assistant Professor**  
**CS, AIUB**  
**[nazib.nasir@aiub.edu](mailto:nazib.nasir@aiub.edu)**

# Lecture Outline



1. Linear block code



# Linear Block Code

## Generator Matrix

Linear Block Code: A code in which addition of any two codewords gives another codeword [2].

Message,  $M$ :  $k$  bits long  
 Redundant bits,  $Q$ :  $q$  bits long  
 Codeword length,  $N$ :  $k+q$  bits long

Generator matrix,  $G = [P_{k \times q} I_k]$

For  $k = 3$  and  $q = 3$ ,

$$G = \begin{bmatrix} 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 0 & 1 \end{bmatrix}$$

$$P_{3 \times 3} = \begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \end{bmatrix}$$

$$I_3 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

Then, it is a  $(n, k) = (6, 3)$  block code



# Linear Block Code....

## Codeword calculation

The codeword for the message [ 0 1 1 ] is

$$C = M \times G$$

$$C = [0 \quad 1 \quad 1] \times \begin{bmatrix} 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 0 & 1 \end{bmatrix}$$

$$C = [\underbrace{1 \quad 1 \quad 0}_Q \quad \underbrace{0 \quad 1 \quad 1}_M]$$

## Modulo-2 summation

$$0 \times 1 \oplus 0 \times 1 \oplus 1 \times 1 = 1$$

$$0 \times 1 \oplus 1 \times 1 \oplus 1 \times 0 = 1$$

$$0 \times 0 \oplus 1 \times 1 \oplus 1 \times 1 = 0$$

$$0 \times 1 \oplus 1 \times 0 \oplus 1 \times 0 = 0$$

$$0 \times 0 \oplus 1 \times 1 \oplus 1 \times 0 = 1$$

$$0 \times 0 \oplus 1 \times 0 \oplus 1 \times 1 = 1$$



# Linear Block Code....

Error-detection

Receiving end

Parity check matrix,

$$H = [I_q \ P_{k \times q}^T]$$

$$H = [I_3 \ P_{3 \times 3}^T]$$

$$H = \begin{bmatrix} 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$$

$$P_{3 \times 3} = \begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \end{bmatrix}$$

$$P_{3 \times 3}^T = \begin{bmatrix} 1 & 0 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \end{bmatrix}$$

$P_{3 \times 3}^T$  is the transpose of  $P_{3 \times 3}$



# Linear Block Code....

Error-detection....

Suppose that there is no error in the received sequence.

Hence the received sequence,  $\mathbf{r}$ , is the same as the transmit sequence,  $\mathbf{C}$ .

$$\mathbf{r} = \mathbf{C}$$

$$\mathbf{r} = [1 \quad 1 \quad 0 \quad 0 \quad 1 \quad 1]$$

$$\text{Syndrome, } \mathbf{s} = \mathbf{rH}^T$$

$$\mathbf{s} = [1 \quad 1 \quad 0 \quad 0 \quad 1 \quad 1] \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \end{bmatrix}$$

$$\mathbf{s} = [0 \quad 0 \quad 0]$$

$$\mathbf{H} = \begin{bmatrix} 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$$

$$\mathbf{H}^T = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \end{bmatrix}$$

The all-zero syndrome indicates a correct reception !



# Linear Block Code....

Error-detection....

Suppose that there is an error in the received sequence.  
The second bit (from left side) has altered from 1 to 0

$$r = [1 \quad 0 \quad 0 \quad 0 \quad 1 \quad 1]$$

Syndrome,  $s = rH^T$

$$s = [1 \quad 0 \quad 0 \quad 0 \quad 1 \quad 1] \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \end{bmatrix}$$

$$s = [0 \quad 1 \quad 0]$$

The non-zero syndrome indicates an erroneous reception !



# Linear Block Code....

## Error-correction

### How to correct the error?

$$H^T = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \end{bmatrix}$$

1. Syndrome ,  $s = [0 \ 1 \ 0]$
2. Locate the syndrome in  $H^T$
3. It is in second row
4. So, the second element in the received sequence,  $r = [1 \ 0 \ 0 \ 0 \ 1 \ 1]$  is erroneous.
4. Alter the second bit from 0 to 1.
5. So the correct received sequence is  $[1 \ 1 \ 0 \ 0 \ 1 \ 1]$ .

Note: The given generator matrix enables correction of at most 1 bits.

It is possible to correct more bits , but it requires quite a lot work! No Free Lunch!



# Homework



❖ Consider a  $(7, 4)$  code whose generator matrix is given by

$$G = \begin{bmatrix} 1 & 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 1 \end{bmatrix}$$

- (a) Find all the codewords of the code
- (b) Find the parity-check matrix
- (c) Find the syndrome for the received vector  $[1 \ 1 \ 0 \ 1 \ 0 \ 1 \ 0]$ . Is it a valid codeword?



# References

- [1] W. Stallings, *Data and Computer Communication*, 10<sup>th</sup> ed., Pearson Education, Inc., 2014, USA, pp. 194 - 196.
- [2] B. Sklar, *Digital Communications*, 2<sup>nd</sup> ed., Prentice Hall. 2017, USA, pp. 328 - 345.



# Recommended Books

1. **Data Communications and Networking**, *B. A. Forouzan*, McGraw-Hill, Inc., Fourth Edition, 2007, USA.
2. **Computer Networking: A Top-Down Approach**, *J. F. Kurose, K. W. Ross*, Pearson Education, Inc., Sixth Edition, USA.
3. **Official Cert Guide CCNA 200-301 , vol. 1**, *W. Odom*, Cisco Press, First Edition, 2019, USA.
4. **CCNA Routing and Switching**, *T. Lammle*, John Wiley & Sons, Second Edition, 2016, USA.
5. **TCP/IP Protocol Suite**, *B. A. Forouzan*, McGraw-Hill, Inc., Fourth Edition, 2009, USA.
6. **Data and Computer Communication**, *W. Stallings*, Pearson Education, Inc., Tenth Edition, 2013, USA.