

# Cyclic Redundancy Check

Course Title: Computer Networks



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# Lecture Outline



1. Cyclic redundancy check



# Cyclic Redundancy Check

## Introduction

- ❖ What if the transmitted bits get altered on the way?
  - Is there any technique to detect the error?

Yes, using Cyclic Redundancy Check (CRC)

### ❑ CRC

- In CRC, some redundant bits are sent in addition to the message bits.
- The purpose of the redundant bits is to facilitate detecting error.
- *The redundant bits are called frame check sequence (FCS)*

How is FCS generated?



# Cyclic Redundancy Check

## Introduction....

- Strength of the CRC depends on the number of redundant bits (that is, FCS length)
- Longer FCS length results in better accuracy in detecting error

### □ Required two sequence

- *Message sequence,  $M$* 
  - The desired data to be sent
  - Can be of any length
- *Pattern sequence,  $P$* 
  - Known to both sender and receiver
  - If we want to use  $K$  bits FCS, we need a pattern bit sequence,  $P$ , of length  $K+1$  bits.

# Cyclic Redundancy Check

## Generation of FCS



1. Decide how many FCS bits,  $K$ , you are going to use.
2. Append  $K$  zeros at the end of the message bits to generate  $M+K$  bits long sequence  $S$ .
3. Select a  $K+1$  bits long pattern sequence,  $P$ .
4. Divide the sequence  $S$  by the pattern sequence  $P$  to find the  $K$  bits of the remainder,  $R$ .
5. Remove the appended zeros from  $S$  and append the calculated remainder  $R$ .  
Thus, the  $N$  bits message bits and  $K$  bits remainder constitutes the transmitting sequence,  $T$ .

# Cyclic Redundancy Check

Error detection at the receiver



1. At the destination, the received sequence,  $T'$ , is divided by the same pattern sequence,  $P$ .
2. If at this step there is no remainder, the data unit is assumed to be correct and is therefore accepted.
3. A remainder indicates that the data unit has been damaged on the way and therefore must be rejected.



# Cyclic Redundancy Check

## Example 1

□ Generate FCS if the message polynomial and generator polynomial are  $X^3 + X^2 + 1$  And  $X^3 + X + 1$ , respectively.

Let  $M(x)$  be the **message polynomial**

Let  $P(x)$  be the **generator polynomial/Pattern sequence**

Let  $P(x) = X^3 + X + 1 \rightarrow 1011$

Let  $M(x) = X^3 + X^2 + 1 \rightarrow 1101$

1. Consider the case where  $M=1101$  and  $P=1011$ .
2. Since  $P$  consists of 4 bits, append  $K=3$  bits zeros (000) at the end of  $M$ ,  $S=1101000$
3. Divide  $S$  by  $P$  to get 3 bits remainder.



# Cyclic Redundancy Check

## Example 1

At sender

Transmit sequence  
or codeword  
 $T = 1101001$

$$\begin{array}{r}
 \begin{array}{c} p \rightarrow 1011 \end{array} \left| \begin{array}{r}
 1111 \\
 1101000 \leftarrow S \\
 1011 \\
 \hline
 1100 \\
 1011 \\
 \hline
 1110 \\
 1011 \\
 \hline
 1010 \\
 1011 \\
 \hline
 001 \leftarrow R
 \end{array}
 \end{array}$$



# Cyclic Redundancy Check

## Example 1

$$\begin{array}{r}
 1011 \overline{) 1101001} \\
 \underline{1011} \phantom{00} \\
 1100 \\
 \underline{1011} \\
 1110 \\
 \underline{1011} \\
 1011 \\
 \underline{1011} \\
 0000
 \end{array}$$

At Receiver

Since the remainder is zero, there is no error in the received sequence



# Cyclic Redundancy Check

## Example 1

### What if any bit gets altered in the channel?

Suppose that the second bit (red) has altered from 1 to 0.

$$\begin{array}{r}
 1011 \overline{) 1 \textcolor{red}{0} 0 1 0 0 1} \\
 \underline{1 \ 0 \ 1 \ 1} \phantom{0} \\
 1 \ 0 \ 0 \ 0 \\
 \underline{1 \ 0 \ 1 \ 1} \\
 1 \ 1 \ 1
 \end{array}$$

*The nonzero remainder indicates an erroneous reception.*

The frame will not be acknowledged.

The sender will resend the frame.

# Cyclic Redundancy Check

## Example 2



- Message  $M = 1010001101$
- Pattern  $P = 110101$
- Length of  $P=6$
- Append  $K=6-1=5$  zeros at the end of  $M$
- $S=101000110100000$
- Now divide  $S$  by  $P$  to find 5 bits remainder [1].

- Transmitted sequence,  $T=101000110101110$
- At the receiving end,  $T$  is divided by  $P$  to see if the remainder is zero. The zero remainder indicates error free reception.

## Example 1

Because there is no remainder, it is assumed that there have been no errors.

# Homework



1. Detect whether the received sequence 101110101 is error free if the pattern sequence is 1010.



# References

- [1] W. Stallings, *Data and Computer Communication*, 10<sup>th</sup> ed., Pearson Education, Inc., 2014, USA, pp. 194 - 196.
- [2] B. Sklar, *Digital Communications*, 2<sup>nd</sup> ed., Prentice Hall. 2017, USA, pp. 328 - 345.



# Recommended Books

1. **Data Communications and Networking**, *B. A. Forouzan*, McGraw-Hill, Inc., Fourth Edition, 2007, USA.
2. **Computer Networking: A Top-Down Approach**, *J. F. Kurose, K. W. Ross*, Pearson Education, Inc., Sixth Edition, USA.
3. **Official Cert Guide CCNA 200-301 , vol. 1**, *W. Odom*, Cisco Press, First Edition, 2019, USA.
4. **CCNA Routing and Switching**, *T. Lammle*, John Wiley & Sons, Second Edition, 2016, USA.
5. **TCP/IP Protocol Suite**, *B. A. Forouzan*, McGraw-Hill, Inc., Fourth Edition, 2009, USA.
6. **Data and Computer Communication**, *W. Stallings*, Pearson Education, Inc., Tenth Edition, 2013, USA.