CoU_Skadooosh Comilla University

Ahsanul Anam Saboj Ahnaf Hasan Shifat Ahmed Jaki

(Allah is the best planner)

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Important formulas

1.1 Prime

- The number of prime numbers less than or equal to n is approximately $\frac{n}{\ln n}$
- The k-th prime number approximately equals $k \ln k$

Properties of phi:

- 1. If gcd(i,n) = d; where $1 \le i \le n-1$ then, there are $phi\left(\frac{n}{d}\right)$ possible values of i.
- 2. $phi(p^k) = p^k p^{k-1}$
- 3. $phi(ab) = phi(a)phi(b)\frac{d}{phi(d)}$; where $d = \gcd(a, b)$
- 4. $x^n = x^{\varphi(m) + [n \mod \varphi(m)]} \mod m$; where n>= log₂m
- 5. Sum of integers that are coprime to n equals to $(\phi(n)\times n)/2$
- 6. For a given integer N, the sum of Euler Phi of each of the divisors of N equals to N
- 7. Given a number N, let d be a divisor of N. Then the number of pairs \mathbf{a} , N, where $1 \le a \le N$ and $\gcd(a,N)=d$, is $\phi(N/d)$

1.3 Properties of mod:

- 1. $ac = bc \pmod{m}$, then $a = b \mod \frac{m}{a(c,m)}$
- 2. $mn \pmod{n} = 0$ then the smallest number m is equal to lcm(n, d)
- 3. if p is prime, then $(x + y)^p = x^p + y^p \pmod{p}$.
- 4. $ab \pmod{ac} = a \pmod{c}$

Properties of Digitsum:

- a. DigitSum(x + y) = DigitSum(DigitSum(x) + DigitSum(y))
- b. DigitSum(x y) = DigitSum(DigitSum(x) DigitSum(y))
- c. DigitSum(x * y) = DigitSum(DigitSum(x) * y)
- d. DigitSum(x * y) = DigitSum(x * DigitSum(y))
- e. DigitSum(x * y) = DigitSum(DigitSum(x) * DigitSum(y))
- f. $DigitSum(x^y) = DigitSum(DigitSum(x)^y)$
- g. $DigitSum(x^y) = DigitSum(x^{y\%\varepsilon})$ where $DigitSum(x^{\varepsilon})=1$

Properties of FLOOR CEIL:

- 1. $\left[\frac{n}{m}\right] = \left[\frac{n+m-1}{m}\right] = \left[\frac{n-1}{m}\right] + 1$ 2. $\left[\frac{n}{m}\right] = \left[\frac{n+m-1}{m}\right] = \left[\frac{n-1}{m}\right] + 1$ 3. $\sum_{k=1}^{n-1} \left[\frac{kn}{m}\right] = \frac{(m-1)(n-1) + \gcd(m,n) 1}{2}$

1.6 Bit Manipulation Macros

```
#define least_one_pos(x) __builtin_ffs(x)
#define leading_zeros(x) __builtin_clz(x)
#define tailing_zeros(x) __builtin_ctz(x)
#define num_of_one(x) __builtin_popcount(x)
#define msb(x) 32-leading zeros(x)
```

1.7 Important Series and properties

1.
$$\sum_{n\geq 0} a^n z^n = \frac{1}{1-az}$$

2. $\sum_{n\geq 0} {m \choose n} z^n = (1+z)^m \quad (m \in \mathbb{Z})$

3.
$$\sum_{n\geq 0} {m+n-1 \choose n} z^n = \frac{1}{(1-z)^m} \quad (m \in \mathbb{Z})$$

3.
$$\sum_{n\geq 0} {m+n-1 \choose n} z^n = \frac{1}{(1-z)^m} \quad (m \in \mathbb{Z})$$

4. $\sum_{n\geq 0} {m+n \choose n} z^n = \frac{1}{(1-z)^{m+1}} \quad (m \in \mathbb{Z})$

5.
$$\sum_{n\geq 0} {n \choose m} z^n = \frac{z^m}{(1-z)^{m+1}} \quad (m \in \mathbb{N})$$

6. $\sum_{n\geq 0} {p+n \choose m} z^n = \frac{z^{m-p}}{(1-z)^{m+1}}$

6.
$$\sum_{n\geq 0} {p+n \choose m} z^n = \frac{z^{m-p}}{(1-z)^{m+1}}$$

$$7. \quad \sum_{n\geq 0} \frac{z^n}{n!} = e^z$$

8.
$$\sum_{n\geq 0} -1^{n-1} \frac{z^n}{n} = \log(1+z)$$

9. Vandermonde
$$\binom{x+y}{n} = \sum_{k=0}^{n} \binom{x}{k} \binom{y}{n-k}$$

10.
$$\sum_{m=0}^{n} {m \choose k} = \sum_{k=0}^{n} {i \choose k} = {n+1 \choose k+1}$$

11.
$$\sum_{k=0}^{n} \binom{n}{k} = 2^n$$

12.
$$\sum_{k=0}^{m} {n+k \choose k} = {m+n+1 \choose m}$$

13.
$$\sum_{k=0}^{n} {n \choose k}^2 = {2n \choose n}$$

14. $\sum_{k=1}^{n} k {n \choose k} = n2^{n-1}$

14.
$$\sum_{k=1}^{n} k \binom{n}{k} = n2^{n-1}$$

Binomial Coefficient (NCR)

```
long long ncr (long long n, long long r)
{
  if(n<r)
    return 0;
  II ans=1;
  ans *= fact[n];
  II d=fact[r];
  d*=fact[n-r];
  d%=MOD;
  an*=binpow(d,MOD-2,MOD);
  an%=MOD;
  return an;
}
```

Modular Inverse

```
// A and M need to be coprime
int x = bigmod(a, m - 2, m); //(ax)%m = 1
// When M is not prime
int modInv ( int a, int m ) {
int x, y;
ext_gcd( a, m, &x, &y );
// Process x so that it is between 0 and m-1
x \% = m;
if (x < 0) x += m;
return x;
```

2 Number Theory

2.1 Sieve

```
vector<long long> prime;
bitset<100000>mark;
inline void sieve( long long n)
    mark[0] = mark[1] = 1;
    long long i,j,limit=sqrt(n*1.0)+2;
    prime.emplace back(2);
    for (i=4; i<=n; i+=2)</pre>
         mark[i]=1;
    for(i=3; i<=n; i+=2)
         if(!mark[i])
             prime.emplace back(i);
             if (i<=limit)</pre>
                  for (j=i*i; j<=n;</pre>
j+=i*2)
                      mark[j]=1;
             }
    }
}
```

2.2 Sum of divisor

```
Lucas Theorem
void init(int n, int mod)
//first calucalte i^(-1)
  inv_f[1] = 1;
  for(int i = 2; i <= n; i++)
    inv_f[i] = mod - 1LL * (mod / i) * inv_f[mod % i] % mod;
//Calculate Inverse factorial and factorial
  inv_f[0] = f[0] = f[1] = 1;
  for(int i = 2; i <= n; i++)
    f[i] = (1LL * f[i - 1] * i) % mod;
    inv f[i] = (1LL * inv f[i] * inv f[i-1]) % mod;
int nCr(int n, int r, int mod)
  if(r > n) return 0;
  return (((1LL * f[n] * inv_f[n - r]) % mod) * inv_f[r]) % mod;
//Convert n to some base
vector<int> toBase(II n, int base)
  vector<int> digits;
  while(n)
     digits.push_back(n % base);
     n /= base;
  return digits;
int lucas(II n, II r, int mod)
  if(r > n) return 0;
// convert n and r to base mod
  vector<int> N = toBase(n, mod);
  vector<int> R = toBase(r, mod);
//make lengths equal by filling leading digits of R with zeros
  while(R.size() < N.size())
     R.push_back(0);
//Calculate answer
  int ans = 1:
  for(int i = 0; i < N.size(); i++)
     ans = (1LL * ans * nCr(N[i], R[i], mod)) % mod;
  return ans;
```

2.3 Miller Robin

```
///Miller - Rabin primality test starts frome
here
const int N=1e6;
using u64 = uint64_t;
using u128 = __uint128_t;
u64 power(u64 base, u64 e, u64 md) {
  u64 result = 1;
  base %= md;
  while (e) {
    if (e & 1)
       result = (u128)result * base % md;
    base = (u128)base * base % md;
    e >>= 1;
  return result;
bool MillerRabin(u64 n, u64 a, u64 d, int s) {
  u64 x = power(a, d, n);
  if (x == 1 | | x == n - 1)
     return false;
  for (int r = 1; r < s; r++) {
    x = (u128)x * x % n;
    if (x == n - 1)
       return false;
  }
  return true;
};
bool isPrime(u64 n, int iter=5)
 // returns true if n is probably prime, else returns false.
  if (n < 4)
     return n == 2 || n == 3;
  int s = 0;
  u64 d = n - 1;
  while ((d \& 1) == 0) \{
    d >>= 1;
    S++;
  for (int i = 0; i < iter; i++) {
    int a = 2 + rand() \% (n - 3);
    if (MillerRabin(n, a, d, s))
       return false;
  return true;
///Miller - Rabin primality test ends here
```

Dynamic Programming DP (Prefix Sum) void solve(int tc) cin>>n; a.resize(n+1); b.resize(n+1); for(int i=1; i<=n; i++) cin>>a[i]; for(int i=1; i<=n; i++) cin>>b[i]; // memset(dp, -1, sizeof(dp)); for(int i=a[n]; i<=b[n]; i++) { dp[n][i] = 1;csum[n][i] += csum[n][i-1] + 1;for(int i=n-1; i>=1; i--) for(int j=a[i]; j<=b[i]; j++) int I = max(j, a[i+1]);int r = b[i+1]; $if(l \le r)$ dp[i][j] += csum[i+1][r] - csum[i+1][l-1];dp[i][j] %= mod; if(dp[i][j] < 0) dp[i][j] += mod;csum[i][j] = csum[i][j-1] + dp[i][j];csum[i][j] %= mod; } } II ans = 0; for(int i=a[1]; i<=b[1]; i++) { ans += dp[1][i];ans %= mod; cout<<ans<<endl;

2.4 Number of divisor

```
long long NumberOfDivisor(long long n)
{
  long long ans=1;
  for(long long i=0;prime[i]*prime[i]<=n;i++)
    {
     long long counter=0;
     while(n%prime[i]==0)
     {
          n/=prime[i];
          counter++;
     }
     ans*=(counter+1);
     }
     if(n>1)ans*=2;
     return ans;
}
```

2.5 Euler Phi

2.6 Euler phi from 1 to MAX

```
#define MAX 100000
long long phi[MAX + 7];
void generatePhi()
{
  phi[1] = 0;
  for (long long i = 2; i <= MAX; i++)
{
    if(!phi[i])
    {
      phi[i] = i-1;
      for(long long j=(i<<1);j<=MAX;j+=i)
    {
      if(!phi[j])
      phi[j] = j;
      phi[j] = phi[j] * (i-1) / i;
      }
    }
}</pre>
```

0-1 Knapsack (log2 trick)

```
vol.pb(0); profit.pb(0);
for(int i=0; i<N; i++)
  Il n, v, p;
  cin>>n>>v>>p;
  II now = 1, cntSum = 0;
  while(cntSum + now <= n)
// if(v * now > 5000) break;
    vol.pb(v * now);
    profit.pb(p * now);
    cntSum += now;
    now *= 2;
  if(n - cntSum >= 0)
    vol.pb(v * (n - cntSum));
    profit.pb(p * (n - cntSum));
  }
}
```

LIS (nlogn)

```
void solve(int tc)
  int n;
  cin>>n;
  vii v(n+1);
  for(int i=1; i<=n; i++) cin>>v[i];
  int inf = 1e9;
  vii ar(n+2, inf);
  v[0] *= -1;
  for(int i=1; i<=n; i++)
    int I = upper_bound( all(ar), v[i]) - ar.begin();
    if(ar[I-1] < v[i] \&\& v[i] < ar[I])
       ar[l] = v[i];
  int ans = 0;
  for(int i=0; i<=n; i++)
    if(ar[i] < inf) ans = i;
  cout<<ans+1<<endl;
```

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3 Data structure

3.1 Segment tree

```
//p=1 begin=0 end=n-1
vector<long long> tree, arr, lazy;
void build(int p, int begin, int end) {
    if (begin == end) {
        tree[p] = arr[begin];
        return;
    int left = p << 1;</pre>
    int right = (p << 1) + 1;</pre>
    int mid = (begin + end) >> 1;
    build(left, begin, mid);
    build(right, mid + 1, end);
    tree[p] = min(tree[left], tree[right]);
void update_lazy(int p, int begin, int end) {
    tree[p] += lazy[p];
    if (begin != end) {
        int left = p << 1;</pre>
        int right = (p << 1) + 1;
        lazy[left] += lazy[p];
        lazy[right] += lazy[p];
    lazy[p] = 0;
void update(int p, int begin, int end, int l, int r, long long value) {
    if (lazy[p] != 0)
        update lazy(p, begin, end);
    if (1 > end || r < begin)
        return;
    if (begin >= 1 && end <= r) {</pre>
        lazy[p] += value;
        update lazy(p, begin, end);
        return;
    int left = p << 1;</pre>
    int right = (p << 1) + 1;</pre>
    int mid = (begin + end) >> 1;
    update(left, begin, mid, l, r, value);
    update(right, mid + 1, end, 1, r, value);
    tree[p] = min(tree[left], tree[right]);
long long query(int p, int begin, int end, int 1, int r) {
    if (lazy[p] != 0)
        update lazy(p, begin, end);
    if (1 > end || r < begin)
        return LLONG MAX;
    if (begin >= 1 && end <= r)
        return tree[p];
    int left = p << 1;</pre>
    int right = (p << 1) + 1;</pre>
    int mid = (begin + end) >> 1;
    long long a = query(left, begin, mid, l, r);
    long long b = query(right, mid + 1, end, l, r);
    return min(a, b);
void segment tree(vector<long long> temp) {
    arr = temp;
    tree.resize(4 * arr.size());
    build(1, 0, arr.size() - 1);
    lazy.assign(4 * arr.size(), OLL);
}
```

3.2 Sqrt Decomposition

```
vector<long long> vcr;
vector<vector<long long >> blocks;
long long N, block size;
void initialize() {
    block size = sqrt(N);
    long long block no = -1;
    for (int i = 0; i < N; ++i) {</pre>
        if (i % block size == 0) {
            block no++;
            vector<long long> s;
            blocks.emplace back(s);
        blocks[block no].push back(vcr[i]);
    for (auto &i: blocks) {
        sort(i.begin(), i.end());
    }
void query(int 1, int r, long long v, int p, long long u) {
    long long k = 0;
    int 11 = 1;
    while (1 % block size && 1 <= r) {
        if (vcr[1] < v)
            k++;
        1++;
    while (l + block size <= r) {</pre>
        int sz=1 / block size;
        k += lower bound(blocks[sz].begin(), blocks[sz].end(), v) -
blocks[sz].begin();
        1 += block size;
    while (1 <= r) {
        if (vcr[1] < v)
            k++;
        1++;
    int sz=p / block size;
    int x = lower bound(blocks[sz].begin(), blocks[sz].end(), vcr[p])-
blocks[sz].begin();
    blocks[sz][x] = (u * k) / (r - 11 + 1);
    vcr[p] = (u * k) / (r - l1 + 1);
    sort(blocks[sz].begin(), blocks[sz].end());
void print array() {
    for (auto &i: vcr) {
        cout << i << endl;
    }
void sqrt Decomposition(vector<long long> &vc) {
   N = vc.size();
   vcr = vc;
    initialize();
}
```

3.3 DSU

```
vector<long long> parent, siz;
void disjointSet(long long n) {
    parent.resize(n), siz.resize(n, 1);
    iota(parent.begin(), parent.end(), 0);
long long find root(long long i) {
    while (parent[i] != i) {
        parent[i] = parent[parent[i]];
        i = parent[i];
    return i;
}
void weighted union(long long a, long long b) {
    long long root a = find root(a);
    long long root b = find root(b);
    if (root a == root b)
        return;
    if (siz[root a] >= siz[root b])
        swap(root a, root b);
    parent[root a] = parent[root b], siz[root b] += siz[root a];
}
bool is connected(long long a, long long b) {
    return find root(a) == find root(b);
}
```

3.4 Trie

```
vector<vector<int>> trie tree;
int min val = '0', total nodes = 0;
vector<int> newnode;
void Trie(int keys) {
    newnode.resize(keys, -1);
    trie tree.emplace back(newnode);
void push(string &s) {
    int level = 0;
    for (int i = 0; i < s.size(); ++i) {</pre>
        if (trie tree[level][s[i] - min val] == -1) {
            trie tree[level][s[i] - min val] = ++total nodes;
            trie_tree.emplace_back(newnode);
        level = trie_tree[level][s[i] - min_val];
    }
long long search(string &s) {
    long long level = 0, value = 0, j = s.size() - 1;
    for (int i = 0; i < s.size(); ++i, --j) {</pre>
        if (trie_tree[level][(s[i] - min val) ^ 1] == -1)
            level = trie_tree[level][(s[i] - min val)];
        else {
            value = value | (1LL << j);</pre>
            level = trie tree[level][(s[i] - min_val) ^ 1];
    }
    return value;
}
```

3.5 Ordered Set

```
#include<ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
template<typename T> using ordered_set =
tree<T,null_type,less<T>,rb_tree_tag,tree_order_statistics_node_update>;
gp_hash_table<int, int> table;
```

3.6 Sparse Table

```
vector<long long >ara;
vector<vector<long long >>BiT;
long long lim, N;
void compute ST()
    for (int i=0; i<N; i++) BiT[0][i]=i;</pre>
    for (long long k=1; (1<<k) <N; k++) {</pre>
         for (long long i=0;i+(1<<k)<=N;i++) {</pre>
             long long x=BiT[k-1][i];
             long long y=BiT[k-1][i+(1<< k-1)];
             BiT[k][i]=ara[x] \le ara[y] ? x : y;
         }
void Sparse table(long long N, vector<long long >&ara)
    ara=ara;
    N=N;
    \lim_{M\to\infty} 64- builtin clz(N);
    BiT.resize(lim, vector<long long >(N));
    compute ST();
long long query(long long i,long long j)
    long long k=log2(j-i);
    long long x=BiT[k][i];
    long long y=BiT[k][j-(1<< k)+1];
    return ara[x] <= ara[y] ? x: y;</pre>
}
```

4 Geometry

4.1 Triangle

6'	aha
Circumcircle	$r = \frac{abc}{}$
	$\sqrt{(a+b+c)(b+c-a)(c+a-b)(a+b-c)}$
	$r = \frac{abc}{}$
	$r = \frac{1}{4 \times AreaOfTriangle}$
Incircle Radius	1 , , , , , , , , , , , ,
	$\frac{1}{2} \times r(a + b + c) = AreaOfTriangle$
Excircle Radius (If the circle is	a+b+c
tangent to side a of the	$r = IncircleRadius \times \frac{a+b+c}{(b+c-a)}$
	(b + c - a)
triangle)	r – 2 × AreaUfTriangle
	$r = 2 \times \frac{AreaOfTriangle}{b + c - a}$
Heron's Formula	$\sqrt{s(s-a)(s-b)(s-c)}$
	$\sqrt{s(s-u)(s-v)(s-v)}$
Sine & Cosine rule	a b c
	$\frac{a}{\sin A} = \frac{b}{\sin B} = \frac{c}{\sin C} = 2R$
	$a^2 = b^2 + c^2 - 2bcCosA$

4.2 Circle

Arc Length	$s = r\theta$ (angle in radian)
Sector Area	area $=\frac{\theta}{2} \times r^2$ (angle in radian)
Chord length	$d = 2 \times r \times \sin\left(\frac{\theta}{2}\right) \text{ (angle in radian)}$ $d = 2 \times \sqrt{r^2 - x^2} (x$ = Perpendicular Distance from the Centre to Chord)
Outside one another	C1C2 > r1 + r2
Touching externally	C1C2 = r1 + r2
Intersecting at 2 points	r1 + r2 < C1C2 < r1 + r2
Touching internally	C1C2 = r1 - r2
One inside the other	C1C2 < r1 - r2

4.3 Others

Cube	$area = 6a^2$ $volume = a^3$
Cylinder	$area = 2\pi rh + 2\pi r^2$ $volume = \pi r^2 h$
Cone	$area = \pi r l$ $volume = \frac{1}{3} \pi r^2 h$
sphere	$area = 4\pi r 2$ $volume = 4 3 \pi r 3$

Operation to Make Palindrome

```
int func(int i, int j)
{
    if(i <= 0 || j <= 0) return 0;
    if(dp[i][j] != -1) return dp[i][j];
    int sum = 0;
    if(s[i-1] == ss[j-1])
    {
        sum = max(sum, func(i-1, j-1) + 1);
    }
    else
    {
        sum = max(sum, max(func(i-1, j), func(i, j-1)));
    }
    return dp[i][j] = sum;
}</pre>
```

Longest Common Subsequence

```
int func (int i, int j)
{
    if(i <= 0 | | j <= 0) return 0;
    if(dp[i][j] != -1) return dp[i][j];
    int sum = 0;
    if(s[i-1] == ss[j-1])
    {
        sum = max( sum, func( i-1, j-1) + 1);
    }
    else
    {
        sum = max( sum, max( func(i-1, j), func(i, j-1)));
    }
    return dp[i][j] = sum;
}</pre>
```

5 Graph

5.1 Articulation point

```
bitset<10017> is visited;
vector<long long> low, dtime;
set<long long>artipoint;
vector<vector<long long>> adjlist;
int minutes;
void articulationpoints (long long u, long long p = -1) {
    ++minutes;
    is visited[u] = true;
    low[u] = dtime[u] = minutes;
    int child = 0;
    for (auto i:adjlist[u]) {
        if (i == p)
            continue;
        if (is visited[i]) {
            low[u] = min(low[u], dtime[i]);
        } else {
            articulationpoints(i, u);
            low[u] = min(low[u], low[i]);
            if (dtime[u] <= low[i] && p != -1)</pre>
                 artipoint.insert(u);
            child++;
    if (p == -1 \&\& child > 1)
        artipoint.insert(u);
}
```

5.2 Dijkstra

```
vector<long long> dis;
vector<int> parent;
vector<vector<pair<int, int>>> adjlist;
void Dijkstra(int node, int source = 0) {
    dis.assign(node, LLONG MAX);
    parent.assign(node, -1);
    dis[source] = 0;
    priority queue<pair<long long, int>> pq;
    pq.push({0, source});
    bitset<100007> processed;
    while (!pq.empty()) {
        int cur node = pq.top().second;
        pq.pop();
        if (processed[cur node])
            continue;
        processed[cur node] = 1;
        for (auto &i : adjlist[cur node]) {
            int x = i.first;
            long long w = i.second;
            if (dis[cur node] + w < dis[x]) {</pre>
                dis[x] = dis[cur node] + w;
                parent[x] = cur node;
                pq.push({-dis[x], x});
            }
        }
    }
}
```

5.3 Bellmenford

```
vector<long long>Node[100005], cost[100005];
long long n, m, i, j, cc=0, k;
long long dis[100005], parent[100005];
long long inf=10e9;
void bellmenford(long long s,long long f)
    for(i=1;i<=n;i++){
         if (i==s) dis[i]=0;else dis[i]=inf;
         parent[i]=-1;
    for (i=1; i<n; i++) {</pre>
        bool done=true;
         for (j=1; j<=n; j++) {
             for (k=0; k<Node[j].size(); k++) {</pre>
                  long long u=j,v=Node[j][k],uv=cost[j][k];
                  if(dis[u]+uv<dis[v]) {
                      dis[v]=dis[u]+uv;
                      parent[v]=u;
                      done=false;
                  }
             }
         if(done)break;/// there was nothing to update ;
    /// Looking for Cycle ;
    bool found=true;
    for(i=1;i<=n;i++){
         for (j=0; j<Node[i].size(); j++) {</pre>
             long long u=i, v=Node[i][j], uv=cost[i][j];
             if(dis[u]+uv<dis[v]){
                  cout<<"Found Negative Cycle"<<endl;</pre>
                  found=false;
                 return;
         if(!found)break;
    for (i=1; i<=n; i++)</pre>
         cout<<"NODE : "<<i<" distance : "<<dis[i]<<endl;</pre>
}
```

5.4 Floyed Warshal

5.5 Ford Fulkerson

```
const int maX=1e5+5;
typedef vector<vector<long long>>v1;
v1 Graph;
long long capacity[1000][1000];
long long n,m;
void init(int N)
    Graph=v1(N+1);
long long bfs(long long s,long long t,vector<long long>&parent)
    fill(parent.begin(),parent.end(),-1);
    parent[s]=-2;
    queue<pair<long long,long long>>q;
    q.push({s, INT MAX});
    while(!q.empty()){
        long long u=q.front().first;
        long long flow=q.front().second;
        q.pop();
        for(long long i=0;i<Graph[u].size();i++){</pre>
            long long v=Graph[u][i];
            if (parent[v] == -1 && capacity[u][v]) {
                 parent[v]=u;
                 long long new flow=min(flow,capacity[u][v]);
                 cout<<v<" ";
                 if(v==t)return new flow;
                 q.push({v,new flow});
            }
        }
    return 0;
long long max flow(long long s,long long t)
    vector<long long>parent(n+1);
    long long flow=0;
    long long new flow;
    while( new flow=bfs(s,t,parent)) {
        cout << endl;
        cout<<new flow<<endl;</pre>
        flow+=new flow;
        long long u=t;
        while(s != u) {
            long long prev=parent[u];
            capacity[prev][u] -= new flow;
            capacity[u][prev]+=new flow;
            u=prev;
        }
    return flow;
}
```

5.6 Prim

```
const int maX=1e5+5;
long long nodes, edges;
bool visit[maX];
vector<pair<long long,long long>>adj[maX];
long long prim(long long x)
{ long long i, j, minimumcost=0, cost;
    priority queue<pair<long long,long long>, vector<pair<long</pre>
    long,long long>>, greater<pair<long long,long long>>> Q;
    pair<long long, long long> p;
    Q.push(\{0, x\});
    while(! Q.empty()){
        p=Q.top();
        Q.pop();
        x=p.second;
        if (visit[x] == true) continue;
        visit[x]=true;
        minimumcost+=p.first;
        for (i=0; i < adj [x].size(); i++) {</pre>
             long long y=adj[x][i].second;
             if (visit[y] == false) Q.push(adj[x][i]);
    return minimumcost;
}
```

5.7 Kruskal

```
const int maX=1e5+5;
long long id[maX], nodes, edges;
pair<long long,pair<long long,long long>>p[maX];
void initialize()
    for (int i=1; i < maX; i++) id[i] = i;</pre>
long long root(long long x)
{
    while(x != id[x])id[x]=id[id[x]], x=id[x];
    return x;
void union1(long long x,long long y)
    long long p=root(x);long long q=root(y);id[p]=id[q];
long long kruskal(pair<long long,pair<long long,long long>>p[])
    long long x,y,cost,minimumcost=0,i;
    for(i=0;i<edges;i++){
        x=p[i].second.first; y=p[i].second.second;cost=p[i].first;
        if(root(x) != root(y)) {
            minimumcost+=cost;
            union1(x, y);
    return minimumcost;
}
```

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// do something

is removed[centroid] = true;

for (int child: G[centroid])

if (is_removed[child])

build_centroid_decomp(child, centroid);

continue;

} }

5.8 LCA

```
const int LOG = 20;
vii G[N], depth(N, 0);
int up[N][LOG];
void dfs (int a)
{
  for( auto it: G[a])
     depth[it] = 1 + depth[a];
     // Binary Lifting
     up[it][0] = a;
     for (int j=1; j<LOG; j++)</pre>
       up[it][j] = up[ up[it][j-1] ][j-1];
     }
     dfs(it);
  }
}
int get lca( int a, int b)
{
  if(depth[a] < depth[b]) swap(a, b);</pre>
  int k = depth[a] - depth[b];
  for ( int j=LOG-1; j>=0; j--)
     If (checkbit(k, j)) a = up[a][j];
  if(a == b) return a;
  for( int j=LOG-1; j>=0; j--)
     If ( up[a][j] != up[b][j] )
       a = up[a][j];
       b = up[b][j];
     }
  }
  return up[a][0];
}
```

```
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Centroid Decomposition
vector<vector<int>> G;
vector<bool> is_removed;
vector<int> sub, par;
int get_subtree_size(int node, int parent = -1)
  sub[node] = 1;
  for (int child: G[node])
    if (child == parent | | is removed[child])
      continue;
    sub[node] += get_subtree_size(child, node);
  }
  return sub[node];
}
int get_centroid(int node, int tree_size, int parent = -1)
  for (int child: G[node])
    if (child == parent || is_removed[child])
    {
      continue;
    if (sub[child] * 2 > tree_size)
      return get_centroid(child, tree_size, node);
    }
  return node;
}
void build_centroid_decomp( int node = 0, int parent=-1)
{
  int centroid = get_centroid(node, get_subtree_size(node));
  par[centroid] = parent;
```

5.9 SCC

```
const int maX=1e5+5;
vector<long long>Graph[maX], Re Graph[maX], check[maX];
long long visit[maX];
stack<long long>ans;
void dfs(long long u)
{
    visit[u]=1;
    for(long long i=0;i<Graph[u].size();i++){</pre>
        long long v=Graph[u][i];
        if(!visit[v]){
            dfs(v);
    ans.push(u);
void dfs2(long long u,long long mark)
    check[mark].emplace back(u);
    visit[u]=1;
    for(long long i=0;i<Re Graph[u].size();i++) {</pre>
        long long v=Re Graph[u][i];
        if(visit[v]==0) {
            dfs2(v,mark);
    }
}
```

5.10 Topsort

```
const int maX=1e5+5;
vector<long long>ara[maX], cost[maX];
bitset<maX> visit;
long long start[maX], finish[maX];
long long cnt=0;
vector<long long>ans;
void dfs(long long source)
{
    cnt++;
    visit[source]=1;
    start[source]=cnt;
    for(int i=0;i<ara[source].size();i++){</pre>
        long long y=ara[source][i];
        if(visit[y] == 0) {
            dfs(y);
    }
    cnt++;
    ans.emplace back(source);
    finish[source]=cnt;
    return;
}
```

5.

6 String

6.1 KMP

```
template<typename T>
class kmp {
    vector<int> indx;
public:
    void lps(T &patt) {
        indx.resize(patt.size(), 0);
        int i = 0, j = 1;
        while (j < patt.size()) {</pre>
             if (patt[i] == patt[j])
                 indx[j] = ++i, j++;
             else {
                 if (i != 0)
                      i = indx[i - 1];
                 else
                      indx[j] = 0, j++;
             }
         }
    }
    bool match(T &text, T &patt) {
         int i = 0, j = 0;
        while (j < text.size()) {</pre>
             if (patt[i] == text[j])
                 i++, j++;
             else {
                 if (i != 0)
                     i = indx[i - 1];
                 else
                      j++;
             if (i == patt.size())
                 return true;
        return false;
    }
    int frequency(T &text, T &patt) {
         int i = 0, j = 0, cnt = 0;
        while (j < text.size()) {</pre>
             if (patt[i] == text[j])
                 i++, j++;
             else {
                 if (i != 0)
                      i = indx[i - 1];
                 else
                     j++;
             if (i == patt.size())
                 cnt++, i = indx[i -
1];
        return cnt;
    }
};
```

Some Properties/Techniques of Number Theory

```
1. How many numbers are coprime, from 1 to N*M, where N and
M are coprime ,GCD(N,M)=1
    1. With N but not with M
    2. With M but not with N
    3. With both M and N
    Ans1 = (phi(N)*M)-phi(N*M)
    Ans2 = (phi(M)*N)-phi(N*M)
    Ans3 = phi(M*N)
2. Divisors sum of every number from 1 to 2*10^9?
Ans = sod(1) + sod(2) + sod(3) + sod(4) + ..... + sod(n);
long long sum_all_divisors(long long num)
  long long sum = 0;
  for (long long i = 1; i \le sqrt(num); i++) {
    long long t1 = i * (num / i - i + 1);
    long long t2 = (((num / i) * (num / i + 1)) / 2) - ((i * (i + 1)) / 2);
    sum += t1 + t2;
  }
  return sum;
3. If gcd(x,n) = 1 then gcd(n-x,n) = 1;
4.\log_k(number) = \frac{\log_{10}(number)}{\log_{10}k} // This is used for base
conversion from decimal to k base.
```

```
int power(long long n, long long k, const int mod) {///Shifater HAshing
 int ans = 1 % mod;
 n %= mod;
 if (n < 0) n += mod;
 while (k) {
 if (k & 1) ans = (long long) ans * n % mod;
  n = (long long) n * n % mod;
  k >>= 1:
 return ans;
const int MOD1 = 127657753, MOD2 = 987654319;
const int p1 = 137, p2 = 277;
int ip1, ip2;
pair<int, int> pw[N], ipw[N];
void prec() {
 pw[0] = \{1, 1\};
 for (int i = 1; i < N; i++) {
  pw[i].first = 1LL * pw[i - 1].first * p1 % MOD1;
  pw[i].second = 1LL * pw[i - 1].second * p2 % MOD2;
 ip1 = power(p1, MOD1 - 2, MOD1);
 ip2 = power(p2, MOD2 - 2, MOD2);
 ipw[0] = \{1, 1\};
 for (int i = 1; i < N; i++) {
  ipw[i].first = 1LL * ipw[i - 1].first * ip1 % MOD1;
  ipw[i].second = 1LL * ipw[i - 1].second * ip2 % MOD2;
struct Hashing {
 int n;
 string s; // 0 - indexed
 vector<pair<int, int>> hs; // 1 - indexed
 Hashing() {}
 Hashing(string _s) {
  n = _s.size();
  s = _s;
  hs.emplace_back(0, 0);
  for (int i = 0; i < n; i++) {
   pair<int, int> p;
   p.first = (hs[i].first + 1LL * pw[i].first * s[i] % MOD1) % MOD1;
   p.second = (hs[i].second + 1LL * pw[i].second * s[i] % MOD2) %
MOD2;
   hs.push_back(p);
  }
 }
 pair<int, int> get_hash(int I, int r) { // 1 - indexed
  assert(1 \le 1 \&\& 1 \le r \&\& r \le n);
  pair<int, int> ans;
  ans.first = (hs[r].first - hs[l - 1].first + MOD1) * 1LL * ipw[l - 1].first %
  ans.second = (hs[r].second - hs[l - 1].second + MOD2) * 1LL * ipw[l -
1].second % MOD2;
  return ans;
 pair<int, int> get_hash() {
  return get_hash(1, n);
 }
};
```

```
#include <bits/stdc++.h> ////Jakir Hashing
#define ff first
#define ss second
#define mp make_pair
#define II long long
using namespace std;
typedef long long LL;
typedef pair<LL, LL> PLL;
///============///
const PLL M = mp(1088888881, 1481481481); ///Should be large primes
const LL base = 347;
                          ///Should be a prime larger than highest value
const int N = 1e6 + 7;
                          ///Highest length of string
ostream& operator<<(ostream& os, PLL hash)
  return os << "(" << hash.ff << ", " << hash.ss << ")";
PLL operator+ (PLL a, LL x)
  return mp(a.ff + x, a.ss + x);
PLL operator- (PLL a, LL x)
  return mp(a.ff - x, a.ss - x);
PLL operator* (PLL a, LL x)
  return mp(a.ff * x, a.ss * x);
PLL operator+ (PLL a, PLL x)
  return mp(a.ff + x.ff, a.ss + x.ss);
PLL operator- (PLL a, PLL x)
  return mp(a.ff - x.ff, a.ss - x.ss);
PLL operator* (PLL a, PLL x)
  return mp(a.ff * x.ff, a.ss * x.ss);
PLL operator% (PLL a, PLL m)
  return mp(a.ff % m.ff, a.ss % m.ss);
PLL power (PLL a, LL p)
  if (p == 0) return mp(1, 1);
  PLL ans = power(a, p / 2);
  ans = (ans * ans) % M;
  if (p % 2) ans = (ans * a) % M;
  return ans;
PLL inverse(PLL a)///Magic!!!!!!!
  return power(a, (M.ff - 1) * (M.ss - 1) - 1);
PLL pb[N]; ///powers of base mod M
PLL invb;
void hashPre()///Call pre before everything
  pb[0] = mp(1, 1);
  for (int i = 1; i < N; i++)
    pb[i] = (pb[i - 1] * base) % M;
  invb = inverse(pb[1]);
PLL Hash (string s)///Calculates Hash of a string
  PLL ans = mp(0, 0);
  for (int i = 0; i < s.size(); i++)
    ans = (ans * base + s[i]) % M;
  return ans;
PLL append(PLL cur, char c)///appends c to string
  return (cur * base + c) % M;
```

6.4 Manachar Algorithm to Find longest Palindrome

```
#define olta(a)
reverse(a.begin(),a.end())
#define mem(a,b)
                         memset(a,b,sizeof(a))
#define rsrt(v) sort(v.rbegin(), v.rend());
#definegsrt(a) sort(a.begin(), a.end(),
greater<11>())
#define vp
                      vector<pair<11, 11> >
#define v_min(a)
*min_element(a.begin(),a.end())
#define v_max(a)
                    *max element(a.begin(),a.end())
#define v_mini(v)
                    min element(v.begin(), v.end())
- v.begin();
#define v_maxi(v) max_element(v.begin(),v.end()) -
v.begin();
#define v_sum(a)
                   accumulate(a.begin(),a.end(),0)
#define un(a)
a.erase(unique(a.begin(),a.end()),a.end())
#define delete(a)
a.erase(a.begin(),a.end())
#define Sort(a)
sort(a.begin(),a.end())
#define is(a)
is sorted(a.begin(),a.end())
#define Saboj4632
ios base::sync with stdio(0);cin.tie(0);cout.tie(0)
#define lcm(a,b)
                                ((a) * (b)) /gcd (a, b)
#define pi
                                3.141592653589793
```

```
PLL prepend(PLL cur, int k, char c)///prepends c to string with size k
  return (pb[k] * c + cur) % M;
PLL replace(PLL cur, int i, char a, char b)///replaces the i-th (0-indexed) character
from right from a to b;
  cur = (cur + pb[i] * (b - a)) % M;
  return (cur + M) % M;
PLL pop_back(PLL hash, char c)///Erases c from the back of the string
  return (((hash - c) * invb) % M + M) % M;
PLL pop_front(PLL hash, int len, char c)///Erases c from front of the string with size
  return ((hash - pb[len - 1] * c) % M + M) % M;
PLL concat(PLL left, PLL right, int k)///concatenates two strings where length of the
right is k
  return (left * pb[k] + right) % M;
///Calculates hash of string with size len repeated cnt times
///This is O(log n). For O(1), pre-calculate inverses
PLL repeat(PLL hash, int len, LL cnt)
  PLL mul = (pb[len * cnt] - 1) * inverse(pb[len] - 1);
  mul = (mul \% M + M) \% M;
  PLL ans = (hash * mul) % M;
  if (pb[len].ff == 1) ans.ff = hash.ff * cnt;
  if (pb[len].ss == 1) ans.ss = hash.ss * cnt;
  return ans;
///Calculates hashes of all prefixes of s including empty prefix
vector<PLL> hashList(string &s)
  int n = s.size();
  vector<PLL> ans(n + 1);
  ans[0] = mp(0, 0);
  for (int i = 1; i <= n; i++) ans[i] = (ans[i - 1] * base + s[i - 1]) % M;
///Calculates hash of substring s[l..r] (1 indexed)
PLL substringHash(const vector<PLL> &hashlist, int I, int r)
  int len = (r - l + 1);
  return ((hashlist[r] - hashlist[l - 1] * pb[len]) % M + M) % M;
void solve()
  II n, q;
  cin>>n; string s; cin>>s;
  vector<PLL>Hash_A = hashList(s);
  /// NOW HEAT THE PROBLEM...
int main()
  hashPre();
  solve();
  return 0;
```

```
#Dijkstra
```

```
int main()
    ios base::sync with stdio(0);
    cin.tie(0);
    cout.tie(0);
    11 tst = 0;
    11 n, k, m, i, j, x, y, w;
    cin >> n >> m;
    11 dis[n+1];
    11 vis[n+1];
    map<11,11>path;
    for(i=0; i<=n; i++)</pre>
         dis[i] = INF;
         vis[i] = 0;
    vp v[n+1];
    for(i=0; i<m; i++)</pre>
         cin>>x>>y>>w;
         v[x].pb(mp(y,w));
         v[y].pb(mp(x,w));
    }
         cout << "Saboj" << endl;
priority queue<pair<11,11>,vp,greater<pair<1</pre>
1,11>>>s;
    s.push (\mathbf{mp}(0,1));
    dis[1] = 0;
    while(!s.empty())
         11 vv,d;
         vv = s.top().S;
         d = s.top().F;
         s.pop();
             // cout<<vv<<" "<<d<endl;
         for(auto child:v[vv])
             x = child.F;
             y = child.s;
             //cout<<x<<" "<<y<<endl;
             if((dis[vv]+y)<dis[x])</pre>
             {
                  path[x] = vv;
                  dis[x] = dis[vv] + y;
                  s.push(mp(dis[x],x));
             }
         }
    vecl ans:
    if (dis[n] == INF) cout << -1 << endl;</pre>
    else {
         x = n;
         ans.pb(x);
         while (x>1) {
             x = path[x];
             ans.pb(x);
         for (i=ans.size()-1;i>=0;i--
)cout<<ans[i]<<" ";
         cout << endl;
    return 0;
```

Merge Sort Tree (Jaki's)

```
const II N = 3e4+3;
vector<II> Tree[4*N], v(N);
void mergee(II node,II I,II mid,II r)
  II n1 = mid-I+1;
  II n2 = r-mid;
  Il ar1[n1+1], ar2[n2+1];
  for(|| i=1; i<=n1; i++) ar1[i]=v[l+i-1];
  for(|| i=1; i<=n2; i++) ar2[i]=v[mid+i];
  II i=1, j=1, cnt=1;
  while(i<=n1 && j<=n2)
     if(ar1[i]<ar2[j])
       v[l+cnt-1] = ar1[i];
       i++, cnt++;
    }
    else
       v[l+cnt-1]=ar2[j];
       j++, cnt++;
while(i<=n1){ v[l+cnt-1]=ar1[i]; cnt++; i++;}
while(j<=n2){ v[l+cnt-1]=ar2[j]; cnt++; j++; }
for(II i=I; i<=r;i++) Tree[node].push back(v[i]);
void mergeSort(II node,II I,II r)
if(l==r) {Tree[node].push back(v[l]); return;}
     II mid = (I+r)/2;
     II left = 2*node;
     II right = left+1;
    mergeSort(left, I, mid);
     mergeSort(right, mid+1, r);
     mergee(node, l, mid, r);
  }
II querry(II node, II I, II r, II i, II j, II x)
  if(i>r || j<l) return 0;
  if(i <= 1 && j>=r)
  Il up = upper_bound(Tree[node].begin(),
  Tree[node].end(), x)-Tree[node].begin();
  II nn = r-l+1;
  return nn-up;
  II left = node*2;
  II right = left+1;
  II mid = (l+r)/2;
  Il a = querry(left, I, mid, i, j, x);
  Il b = querry(right, mid+1, r, i, j, x);
  return a+b;
```

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```
//KMP
char
txt[1000009],pat[1000009];
void lps ar(char *pat,int
M, int *lps)
{
 int len=0; lps[0]=0; int i=1;
    while(i<M)
   if (pat[i] == pat[len])
        len++, lps[i] = len, i++;
        else
        {
        if(len!=0)
len=lps[len-1];
         else lps[i]=0,i++;
    }
void KMPsearch(char *txt,char
*pat)
    int N=strlen(txt);
    int M=strlen(pat);
    int lps[M];
    lps ar(pat,M,lps);
    int i=0, j=0; total=0;
    while(i<N)
         //cout<<"kmp"<<endl;</pre>
 if (pat[j] == txt[i]) i++, j++;
        if(j==M)
           //to print how many
times match
        total++;
//cout << "found pattern at
index: "<<i-j<<endl;</pre>
             j=lps[j-1];
else if(i<N &&
pat[j]!=txt[i])
     if(j!=0) j=lps[j-1];
       else i++;
```

//STRING MULTIPLACATION

```
string multyply(string a,int b)

{
    int carry = 0;
    ans = "";
for(int i=0;i<a.size();i++)
{
    carry=((a[i]-'0')*b+carry);
    ans += carry % 10 + '0';
        carry /= 10;
    }
    while(carry != 0) {
    ans += carry % 10 + '0';
        carry /= 10;
    }
    return ans;
}</pre>
```

```
//find nCr
ll nCr(ll n,ll r)
    ll p=1, q=1;
    r=min(r,n-r);
    if(r!=0)
         while(r)
             p*=n;q*=r;
             ll x=\underline{gcd(p,q)};
             p/=x;q/=x;
             n--;r--;
    else p=1;
    return p;
///print power
ll power(ll x,ll n)
    11 \text{ res}=1;
    while(n)
         if(n&1) res*=x;
        x*=x;
        n >> = 1;
    return res;
///print power mod
ll power mod(ll a,ll b)
    11 \text{ res}=1;
    while(b)
         if(b&1)
res=(res*1LL*a)%MOD;
         a = (a*1LL*a) %MOD;
        b >> 1;
    return res;
Farmat Little Theorem
Power mod(n,m-2)
Here we should pass the mod as
power.
//GCD
ll gcd(ll a,ll b) {
    if(b==0 || a==0) return 0;
    if (b%a==0) return a;
    else return gcd(b%a,a);
```

```
//o 1 bfs
void zeronebfs(ll x,ll
y, 11 r, 11 c)
    for(i=0; i<=r; i++)</pre>
for (j=0; j<=c; j++)</pre>
dis[i][j]=INT MAX;
    dis[0][0]=0;
    deque<pair<ll, ll>>q;
    q.push back(\{x,y\});
    while(!q.empty())
         auto
it=q.front();
         q.pop front();
         ll a=\overline{i}t.F;
         11 b=it.S;
         for(i=0; i<4;
i++)
     ll e=a+fx[i];
      ll d=b+fy[i];
if(e)=0 \&\& e< r \&\& d>=0 \&\&
d<c)
  {
           11 z=0;
if(ar[a][b]!=ar[e][d])
z=1;
if (dis[a][b]+z<dis[e][d])</pre>
                  {
dis[e][d]=dis[a][b]+z;
if(z==0)
q.push front({e,d});
     else
q.push back({e,d});
                  }
         cout << dis[r-1][c-
1]<<endl;
```

```
//EULET TOTIENT ( 1 to N) (nlog(n))
void EulerTotient()
    phi[1] = 1;
    for (int i=2; i<MAX; i++)</pre>
        if (!phi[i])
             phi[i] = i-1;
 for (int j = i*2; j < MAX; j+=i)
      if (!phi[ j]) phi[j] = j;
   phi[j] = (phi[j]/i)*(i-1);
            } } } }
// EULET TOTIENT (sqrt(n) * log(n))
int phi(int n) {
    int result = n;
    for(int i=2;i*i<= n; i++) {</pre>
        if (n % i == 0) {
      while (n \% i == 0) n /= i;
        result -= result / i;
        } }
    if (n > 1)
      result -= result / n;
    return result;
```

Fibonacci of nth Using Matrix multiplication

```
11 fib(11 n) {
    if(n==1)return 0;
    if (n==2) return 1;
    11 b = n-2;
    11 x, y, z, w;
    ll f[2][2] = \{\{1,1\},\{1,0\}\};
    ll r[2][2] = \{\{1,0\},\{0,1\}\};
    if(b<0){
        return 0;
    while (b>0) {
        if(b&1){
            x = ((r[0][0]*f[0][0])*MAX + (r[0][1]*f[1][0])*MAX)*MAX;
            y = ((r[0][0]*f[0][1])%MAX + (r[0][1]*f[1][1])%MAX)%MAX;
            w = ((r[1][0]*f[0][0])%MAX + (r[1][1]*f[1][0])%MAX)%MAX;
            z = ((r[1][0]*f[0][1])%MAX + (r[1][1]*f[1][1])%MAX)%MAX;
            r[0][0] = x;
            r[0][1] = y;
            r[1][0] = w;
            r[1][1] = z;
            //cout<<r[0][0]<<" r"<<endl;
        // cout<<" b "<<b<<endl;
        x = ((f[0][0]*f[0][0]) %MAX + (f[0][1]*f[1][0]) %MAX) %MAX;
        y = ((f[0][0]*f[0][1])*MAX + (f[0][1]*f[1][1])*MAX)*MAX;
        W = ((f[1][0]*f[0][0])%MAX + (f[1][1]*f[1][0])%MAX)%MAX;
        z = ((f[1][0]*f[0][1])%MAX + (f[1][1]*f[1][1])%MAX)%MAX;
        // cout<<"X "<<x<" y "<<y<<" w "<<w<<" "<<z<<endl;
        f[0][0] = x;
        f[0][1] = y;
        f[1][0] = w;
        f[1][1] = z;
        // cout<<"f[0][0] "<<f[0][0]<< " "<<f[0][1]<<endl;
        b>>=1;
    }
    return r[0][0];
```

```
const int fx[]=\{+1,-1,+0,+0\};//graph\ move const int fy[]=\{+0,+0,+1,-1\};//graph\ move const int fx[]=\{+0,+0,+1,-1,-1,+1,-1,+1\};//kings\ move const int fy[]=\{-1,+1,+0,+0,+1,+1,-1,-1\};//kings\ move const int fx[]=\{-2,-2,-1,-1,+1,+1,+2,+2\};//knight's\ move const int fx[]=\{-1,+1,-2,+2,-2,+2,-1,+1\};//knight's\ move
```