Sort type	Pseudo Code	Java Code
Bubble Sort  Time Complexity: O(n^2) Space Complexity: O(1)	<pre>bubbleSort(int arr[]):     for i = 0 to (arr.length-1)         do: for j = 0 to (arr.length-i-1)         do: if arr[j] &gt; arr[j+1]</pre>	<pre>void bubbleSort(int arr[]) {    int i, j, n = arr.length;    for (i = 0; i &lt; n-1; i++)    // Last i elements are already in place       for (j = 0; j &lt; n-i-1; j++)         if (arr[j] &gt; arr[j+1]) {         int temp = arr[j];         arr[j] = arr[j+1];         arr[j+1] = temp;     } }</pre>
Selection Sort Time Complexity: 0(n^2) Space Complexity: 0(1) Find minimum and put in the start.	<pre>selectionSort(int arr[]):     for i = 0 to (arr.length-1)         do: min_id &lt; i         for j = (i+1) to (arr.length)             do: if (arr[j] &lt; arr[min_id])             then: min_id &lt; j         temp &lt; arr[i]         arr[i] &lt; arr[min_id]         arr[min_id] &lt; temp</pre>	<pre>void selectionSort(int arr[]){    int i, j, min_id, n = arr.length;    // One by one move boundary of    // unsorted subarray    for (i = 0; i &lt; n-1; i++){         // Find the minimum element         // in unsorted array         min_id = i;         for (j = i+1; j &lt; n; j++)             if (arr[j] &lt; arr[min_id])             min_id = j;         // Swap the found minimum element         // with the first element         int temp = arr[i];         arr[i] = arr[min_id];         arr[min_id] = temp;    } }</pre>

```
void insertionSort(int arr[]){
                                                                                int i, key, j, n = arr.length;
                                                                                for (i = 1; i < n; i++){
      Insertion Sort
                                                                                    key = arr[i];
                            insertionSort(int arr[]):
                                                                                    i = i - 1:
                                for i = 1 to (arr.length)
                                                                                    // Move elements of arr[0..i-1], that are
  Time Complexity: O(n^2)
                                     do: key <-- arr[i]</pre>
                                                                                    // greater than key, to one position ahead
  Space Complexity: 0(1)
                                     j <-- (i - 1)
                                                                                    // of their current position
  while(i \ge 0 and arr[i] > key)
                                                                                    while (j \ge 0 \&\& arr[j] > key){
   Pick values from the
                                        do: arr[i+1] <-- arr[i]</pre>
                                                                                        arr[j + 1] = arr[j];
unsorted part and place at
                                        j <-- j - 1
                                                                                        j = j - 1;
the correct position in the
                                    arr[j+1] <-- key
       sorted part.
                                                                                    arr[j + 1] = key;
                            merge(A[],B[]):
                                                                            int[] merge(int[] a, int[] b){
                                n1 = A.length, n2 = B.length,
                                                                                int p = a.length, q = b.length;
                                Create C[1 ... (n1+n2)]
                                                                                int c[] = new int[p + q];
                                i = 0, i = 0, k = 0
                                                                                int i = 0, j = 0, k = 0;
                                while i < n1 and j < n2
                                                                                while(i<p && j<q){
                                     do: if A[i] < B [i]</pre>
                                                                                    if(a[i] < b[j]) c[k++] = a[i++];
                                        then: C[k] = A[i]
                                                                                    else c[k++] = b[j++];
                                        i = i+1
        Merge Sort
                                     else
                                                                                while(i < p) c[k++] = a[i++];
                                        do: C[k] = B[j]
                                                                                while(j < q) c[k++] = b[j++];
                                        j = j+1
                                                                                return c;
Time Complexity: O(n lg(n))
                                    k = k+1
  Space Complexity: O(n)
                                while i < n1
                                                                            int[] sort(int[] arr){
                                    do: C[k] = A[i]
                                                                                int l = 0, r = arr.length-1;
  k = k+1
                                                                                if(1 < r){
    Divide and Conquer
                                    i = i+1
                                                                                    int m = (1 + r) / 2; // middle point
Divide input array in two
halves, call sort function
                                while j < n2
                                                                                    int n1 = m - 1 + 1, n2 = r - m;
for the two halves and then
                                    do. C[k] = R[i]
                                                                                    // Crosto tomo arrave
```

```
tor the two naives and then
                                     uo: Ե[K] = Ե[]]
                                                                                     // Create Lemp arrays
                                                                                     int[] L = new int[n1], R = new int[n2];
merge the two sorted halves.
                                     k = k+1
Division runs until each sub
                                     j = j+1
                                                                                    // Copy data to temp arrays
array is of length 1, then
                                 return C
                                                                                    for (int i=0; i<n1; ++i) L[i] = arr[l+i];
                             sort(Arr[]):
                                                                                    for (int j=0; j<n2; ++j) R[j] = arr[m+1+<math>j];
     merge takes place
                                 if Arr.length = 1
                                                                                    // Sort first and second halves
                                     then return
                                                                                    L = sort(L):
                                                                                    R = sort(R);
                                     else
                                     Create L[] with A[0 to (n/2)]
                                                                                    // Merge the sorted halves
                                     Create R[] with A[(n/2+1) to n]
                                                                                    arr = merge(L,R);
                                     L = sort(L)
                                                                                } // l==r
                                     R = sort(R)
                                                                                return arr;
                                     A = merge(L,R)
                             build_max_heap(A)
                                 heap_size[A] = length[A]
                                 for i = (length[A]/2) downto 1
                                     do max_heapify(A, i)
          Heap Sort
                             max heapifv(A. i)
                                 l = left(i)
                                 r = right(i)
                                 if l \le heap\_size[A] and A[l] > A[i]
Time Complexity: O(n lg(n))
                                                                            This link has full heap class and implementation:
                                     then: max = 1
   Space Complexity: 0(1)
                                                                            https://drive.google.
  else: max = i
                                                                            com/file/d/1hAN94NQ7zmZmfqK7BEwZvYf7oWkdjQQX/view?
 Based on Binary Heap data
                                 if r \le heap\_size[A] and A[r] > A[max]
                                                                            usp=sharing
         structure.
                                     then: max = r
  Find the maximum element
                                 if max != i
                                                                            It can simply be used by keeping in the same package
   and place the maximum
                                     then: exchange A[i] with A[max]
                                                                            as driver class.
                                     max_heapify(A, max)
    element at the end.
 Very Similar to Selection
                             heapsort(A[])
            Sort.
                                 build_max_heap(A)
                                 for i = (A.length) downto 2
                                     do: exchange A[1] with A[i]
                                     heap_size[A] = heap_size[A] - 1
                                     max_heapify(A, 1)
```

Quick Sort

Time Complexity:
Average: O(n lg(n))
Worst: O(n^2)
Space Complexity: ~~~

Divide and Conquer
Pick an element as pivot and partition the array around the pivot with lesser elements on one side and greater elements on the other

The key process is partition()
In Randomized\_QuickSort, pivots are just picked randomly

```
int partition(int[] A, int p, int q) {
   int x = A[p]; // pivot
   int i = p; // first index
   for (int j = p+1; j <= q; j++) {
        if(A[j]<=x){
           i = i+1;
           int t = A[i];
           A[i] = A[j];
           A[j] = t;
        }
   // swap first to put in between
   // lesser and greater parts
   int t = A[i];
   A[i] = A[p];
   A[p] = t;
   return i;
int[] quick_sort(int[] A, int p, int r) {
   if(p<r){
        int q = partition(A,p,r);
        quick_sort(A,p,q-1);
        quick_sort(A,q+1, r);
   return A;
```