

<p style="text-align: center;">Insertion Sort</p> <p>..... Time Complexity: $O(n^2)$ Space Complexity: $O(1)$ Pick values from the unsorted part and place at the correct position in the sorted part.</p>	<pre> insertionSort(int arr[]): for i = 1 to (arr.length) do: key <-- arr[i] j <-- (i - 1) while(j >= 0 and arr[j] > key) do: arr[j+1] <-- arr[j] j <-- j - 1 arr[j+1] <-- key </pre>	<pre> void insertionSort(int arr[]){ int i, key, j, n = arr.length; for (i = 1; i < n; i++){ key = arr[i]; j = i - 1; // Move elements of arr[0..i-1], that are // greater than key, to one position ahead // of their current position while (j >= 0 && arr[j] > key){ arr[j + 1] = arr[j]; j = j - 1; } arr[j + 1] = key; } } </pre>
<p style="text-align: center;">Merge Sort</p> <p>..... Time Complexity: $O(n \lg(n))$ Space Complexity: $O(n)$ Divide and Conquer Divide input array in two halves, call sort function for the two halves and then</p>	<pre> merge(A[],B[]): n1 = A.length, n2 = B.length, Create C[1 ... (n1+n2)] i = 0, j = 0, k = 0 while i < n1 and j < n2 do: if A[i] < B[j] then: C[k] = A[i] i = i+1 else do: C[k] = B[j] j = j+1 k = k+1 while i < n1 do: C[k] = A[i] k = k+1 i = i+1 while j < n2 do: C[k] = B[j] k = k+1 j = j+1 </pre>	<pre> int[] merge(int[] a, int[] b){ int p = a.length, q = b.length; int c[] = new int[p + q]; int i = 0, j = 0, k = 0; while(i<p && j<q){ if(a[i]<b[j]) c[k++] = a[i++]; else c[k++] = b[j++]; } while(i<p) c[k++] = a[i++]; while(j<q) c[k++] = b[j++]; return c; } int[] sort(int[] arr){ int l = 0, r = arr.length-1; if(l < r){ int m = (l + r) / 2; // middle point int n1 = m - l + 1, n2 = r - m; // Create temp arrays </pre>

<p>For the two halves and then merge the two sorted halves. Division runs until each sub array is of length 1, then merge takes place</p>	<pre> do: C[k] = B[j] k = k+1 j = j+1 return C sort(Arr[]): if Arr.length = 1 then return else Create L[] with A[0 to (n/2)] Create R[] with A[(n/2+1) to n] L = sort(L) R = sort(R) A = merge(L,R) </pre>	<pre> // Create temp arrays int[] L = new int[n1], R = new int[n2]; // Copy data to temp arrays for (int i=0; i<n1; ++i) L[i] = arr[l+i]; for (int j=0; j<n2; ++j) R[j] = arr[m+1+j]; // Sort first and second halves L = sort(L); R = sort(R); // Merge the sorted halves arr = merge(L,R); } // l==r return arr; } </pre>
<p>Heap Sort</p> <p>.....</p> <p>Time Complexity: $O(n \lg(n))$</p> <p>Space Complexity: $O(1)$</p> <p>.....</p> <p>Based on Binary Heap data structure.</p> <p>Find the maximum element and place the maximum element at the end.</p> <p>Very Similar to Selection Sort.</p>	<pre> build_max_heap(A) heap_size[A] = length[A] for i = (length[A]/2) downto 1 do max_heapify(A, i) max_heapify(A, i) l = left(i) r = right(i) if l <= heap_size[A] and A[l] > A[i] then: max = l else: max = i if r <= heap_size[A] and A[r] > A[max] then: max = r if max != i then: exchange A[i] with A[max] max_heapify(A, max) heapsort(A[]) build_max_heap(A) for i = (A.length) downto 2 do: exchange A[1] with A[i] heap_size[A] = heap_size[A] - 1 max_heapify(A, 1) </pre>	<p>This link has full heap class and implementation: https://drive.google.com/file/d/1hAN94NQ7zmZmfqK7BEwZvYf7oWkdjQQX/view?usp=sharing</p> <p>It can simply be used by keeping in the same package as driver class.</p>

