

School of Informatics, Computing, and Cyber Systems

ALGORITHMIC STRATEGIES

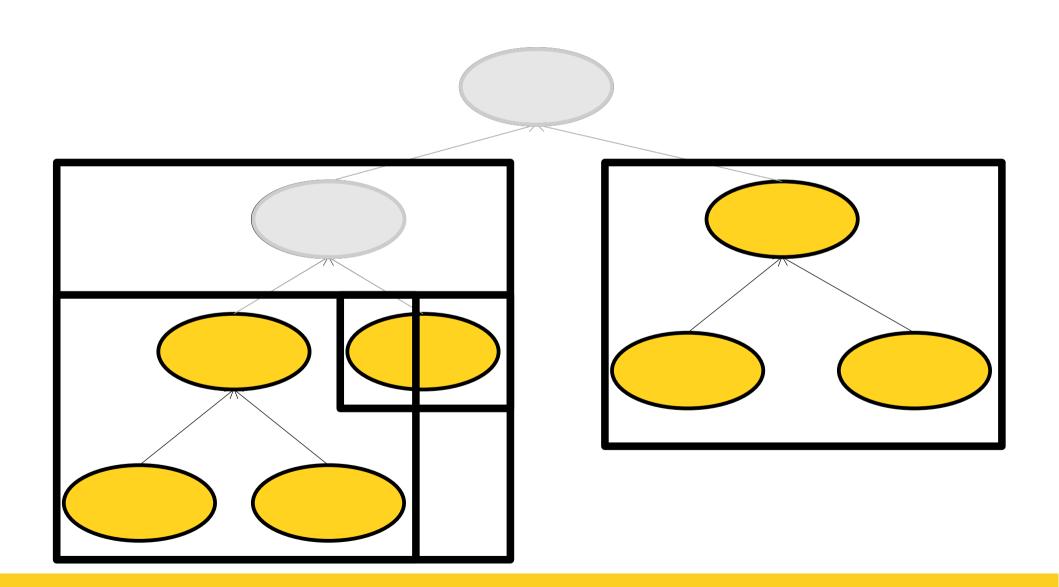
Igor Steinmacher
INF 502

ALGORITHMIC STRATEGIES

- Known techniques:
 - Recursion
 - Divide and Conquer
 - Try and Error
 - Dynamic Programming
 - Greedy algorithms

"To understand recursion, you need to understand recursion"

- Recursive algorithm
 - Algorithm that calls itself direct or indirectly
 - Useful when the problem is naturally recursive or uses recursive data structures (like trees)



- Write an algorithm that calculates the factorial (n!) of a number n provided
 - Non-recursive algorithm:

```
• factorial n = (n)*(n-1)*(n-2)...*2
```

```
n = int(input())
fact = 1
if (n > 1):
    for i in range (2, n+1):
        fact = fact * i
```

Testing

```
n = int(input())
fact = 1
if (n > 1):
    for i in range (2, n+1):
        fact = fact * i
```

n	fact	i	
5	1	5	
	5	4	
	20	3	
	60	2	
	120	1	

Output: 120

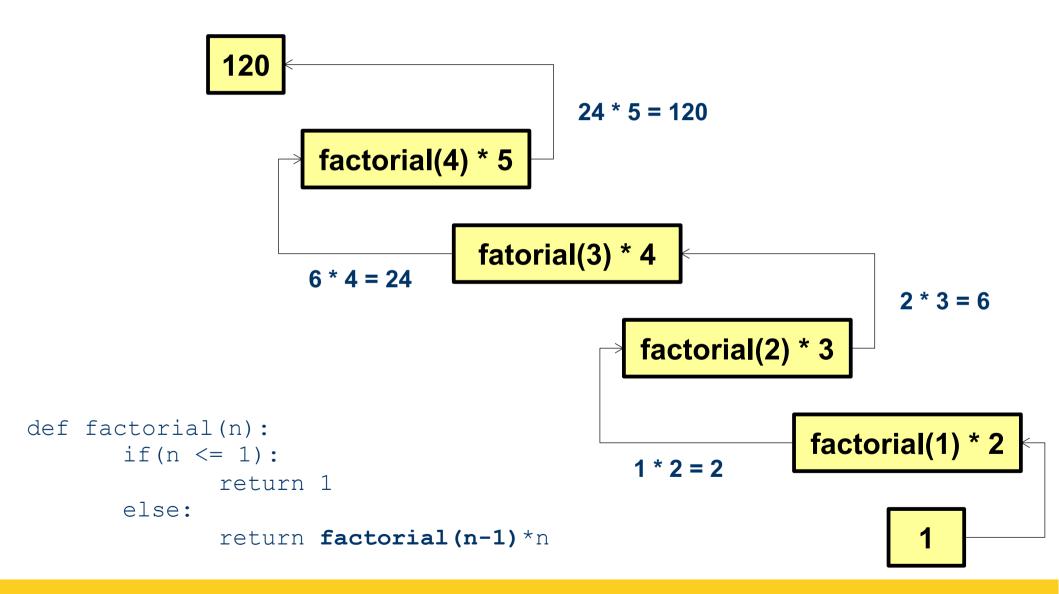
- Recursive solution
 - factorial (n) = factorial(n-1)*n

```
def factorial(n):
    if(n <= 1):
        return 1
    else:
        return factorial(n-1)*n</pre>
```

- Algorithm execution
 - The function calls itself recursively
 - The parameter is smaller in each iteration (n-1)
 - The function multiplies the result of its call by n untill reaching the base case
- Base case
 - Stop condition
 - In the exemple, base case is n ≤ 1
 - We need to guarantee that the input will reach the base case

- For n starting in 5:
 - factorial(4) * 5
 - factorial(3) * 4
 - factorial(2) * 3
 - factorial(1) * 2
 - 1

```
def factorial(n):
    if(n <= 1):
        return 1
    else:
        return factorial(n-1)*n</pre>
```



- How to build a recursive algorithm?
 - Find the base case
 - Change the input until you reach the base case
 - Solve the base case
 - Return the result until you come back to the first call



- Basic Steps
 - Divide
 - Divide the problem into subproblems (smaller)
 - Conquer
 - Calculate the result of the smaller problem
 - Combine
 - Combine the results to get the global solution

Generic algorithm

- Did our factorial algorithm use this technique??
- Zero-cost combination
 - The result of a subproblem is the solution
- High-cost combination
 - To combine it is necessary to analyze the previous results and (usually) loop through them

- Exponential:
 - First solution using weak induction:
 - Base case: n = 0, $a^0 = 1$
 - Induction hypothesis:
 - For any integer n > 0, I know how to calculate a^{n-1}
 - Induction step:
 - Prove that it is possible to calculatee aⁿ, for n > 0.
 - Using our induction hypothesis, I know how to calculate aⁿ, aⁿ⁻¹ by a

$$a^n = a^{n-1} * a$$

- For a = 5 and n = 4, we want to calculate $5^4 = 5^{3*}5$
 - Induction steps:
 - Do I know how to calculate $5^4 = 5^{3*}5$ (final solution)?
 - No, but, inductively i know how to calculate $5^3 = 5^{3-1} * 5 \rightarrow 5^2 * 5$
 - Do I know how to calculate 5³ = 5²*5?
 - No, but, inductively i know how to calculate $5^2 = 5^{2-1} * 5 \rightarrow 5^1 * 5$
 - Do I know how to calculate 5² = 5¹*5?
 - No, but, inductively i know how to calculate $5^1 = 5^{1-1} * 5 \rightarrow 5^0 * 5$
 - Do I know how to calculate $5^1 = 5^{0*}5$?
 - YES! $5^0 = 1$ (base casee) .: 1*5 = 5
 - Now, we can combine our solutions

- For a = 5 and n = 4, we want to calculate $5^4 = 5^{3*}5$
 - Induction steps
 - Do I know how to calculate $5^4 = 5^{3*}5$ (final solution)?
 - No, but, inductively i know how to calculate $5^3 = 5^{3-1} * 5 \rightarrow 5^2 * 5$
 - Do I know how to calculate $5^3 = 5^{2*}5$?
 - No, but, inductively i know how to calculate $5^2 = 5^{2-1} * 5 \rightarrow 5^1 * 5$
 - Do I know how to calculate 5² = 5¹*5?
 - YES, $5^1 * 5 = 5 * 5 = 25$
 - Do I know how to calculate $5^1 = 5^{0*}5$?
 - YES. $5^0 = 1$ (base case) .: 1*5 = 5

- For a = 5 and n = 4, we want to calculate $5^4 = 5^{3*}5$
 - Induction steps
 - Do I know how to calculate $5^4 = 5^{3*}5$ (final solution)?
 - No, but, inductively i know how to calculate $5^3 = 5^{3-1} * 5 \rightarrow 5^2 * 5$
 - Do I know how to calculate 5³ = 5²*5?
 - YES, $5^3 = 5^2 * 5 = 25 * 5 = 125$
 - Do I know how to calculate 5² = 5¹*5?
 - YES, $5^1 * 5 = 5 * 5 = 25$
 - Do I know how to calculate 5¹ = 50*5?
 - YES. $5^0 = 1$ (base case) .: 1*5 = 5

- For a = 5 and n = 4, we want to calculate $5^4 = 5^{3*}5$
 - Induction steps
 - Do I know how to calculate $5^4 = 5^{3*}5$ (final solution)?

$$-$$
 YES, $5^4 = 5^3 * 5 => 5^3 * 5 = 625$

• Do I know how to calculate $5^3 = 5^{2*}5$?

$$-$$
 YES, $5^3 = 5^2 * 5 = 25 * 5 = 125$

- Do I know how to calculate 5² = 5¹*5?
 - YES, $5^1 * 5 = 5 * 5 = 25$
- Do I know how to calculate 5¹ = 50*5?
 - YES. $5^0 = 1$ (base case) .: 1*5 = 5

- For a = 5 and n = 4, we want to calculate $5^4 = 5^{3*}5$
 - Induction steps
 - Do I know how to calculate $54 = 5^{3*}5$ (final solution)?

$$-$$
 YES, $54 = 53 * 5 => 5^3 * 5 = 625$

• Do I know how to calculate $5^3 = 5^{2*}5$?

$$-$$
 YES, $5^3 = 5^2 * 5 = 25 * 5 = 125$

- Do I know how to calculate 5² = 5¹*5?
 - Final Solution: 625
- Do I know how to calculate 5¹ = 50*5?
 - YE\$. $5^0 = 1$ (base case) .: 1*5 = 5

Let's write a divide-and-conquer algorithm that, given a list S of n
 >2 numbers, identifies the smallest element of S

We are building it, not using the min()

```
def calc minimum(lst, start, end):
   min1 = 0
  min2 = 0
   if (end - start <= 1):
      if (lst[start] < lst[end]):</pre>
      # list with 2 values or less
         return lst[start]
      else:
         return lst[end]
   else:
      middle = int((start + end) / 2)
      min1 = calc minimum(lst, middle, end)
      min2 = calc minimum(lst, start, middle-1)
      if (min1 <= min2):
         return min1
      else:
         return min2
numbers = [2, 9, 1, 7, 8, 3]
print(calculate minimum(numbers, 0, len(numbers)-1))
```

2 9 1 7 8 3

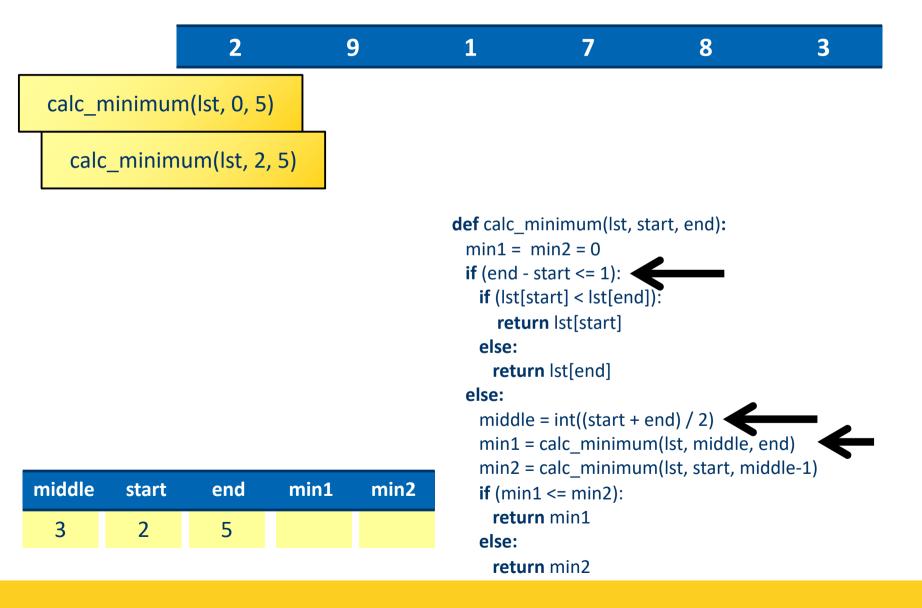
return min2

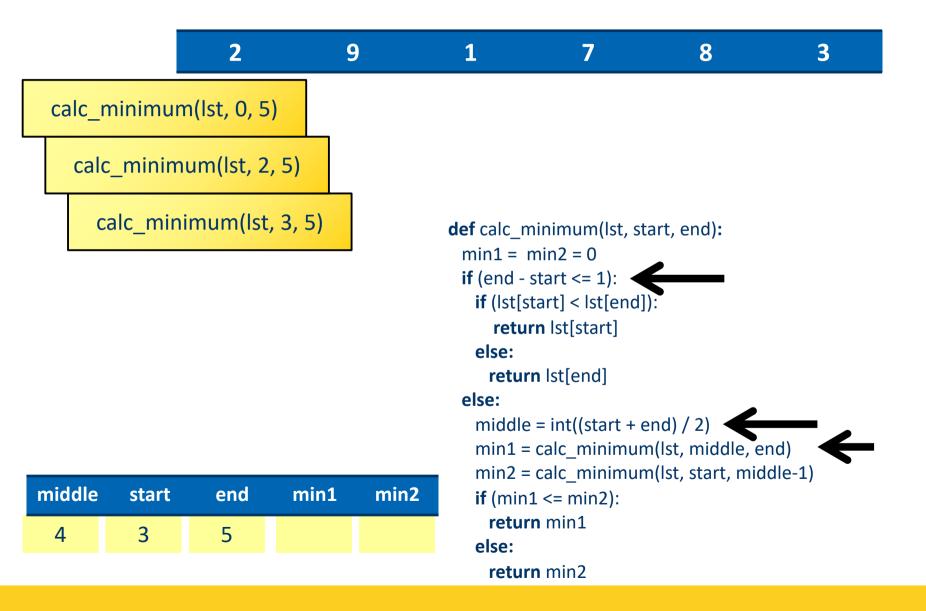
def calc minimum(lst, start, end):

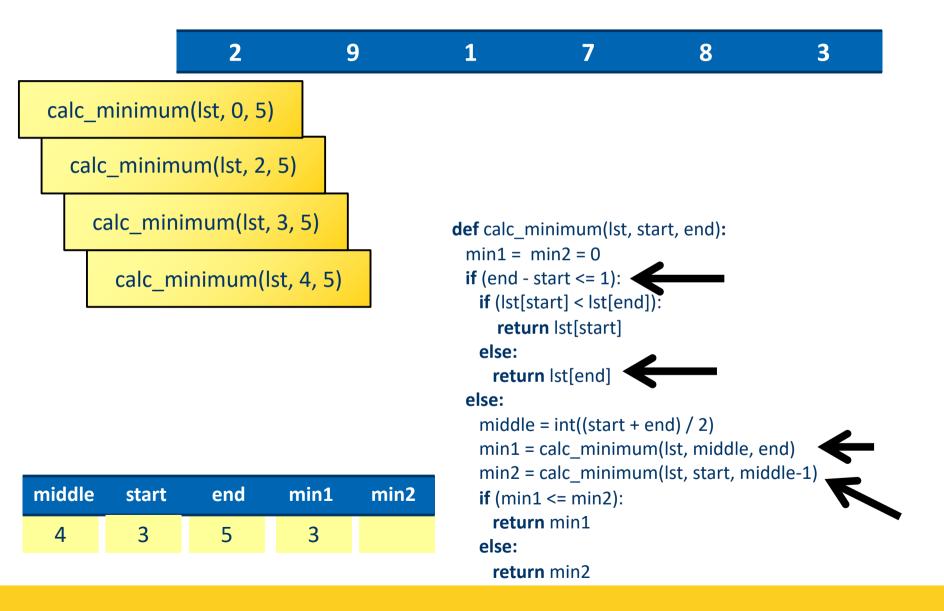
calc_minimum(lst, 0, 5)

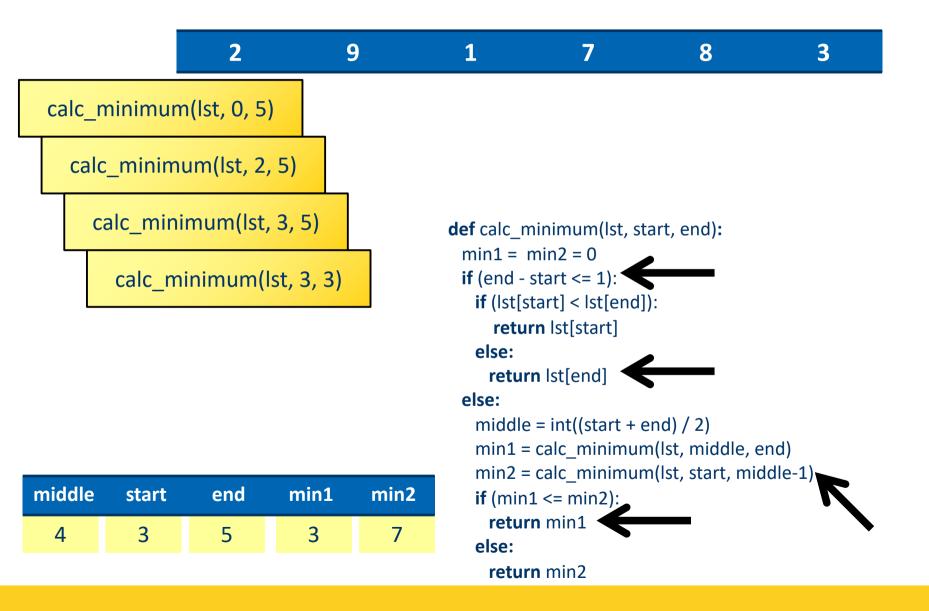
```
min1 = min2 = 0
if (end - start <= 1):
    if (lst[start] < lst[end]):
        return lst[start]
    else:
        return lst[end]
else:
    middle = int((start + end) / 2)
    min1 = calc_minimum(lst, middle, end)
    min2 = calc_minimum(lst, start, middle-1)
    if (min1 <= min2):
        return min1
    else:</pre>
```

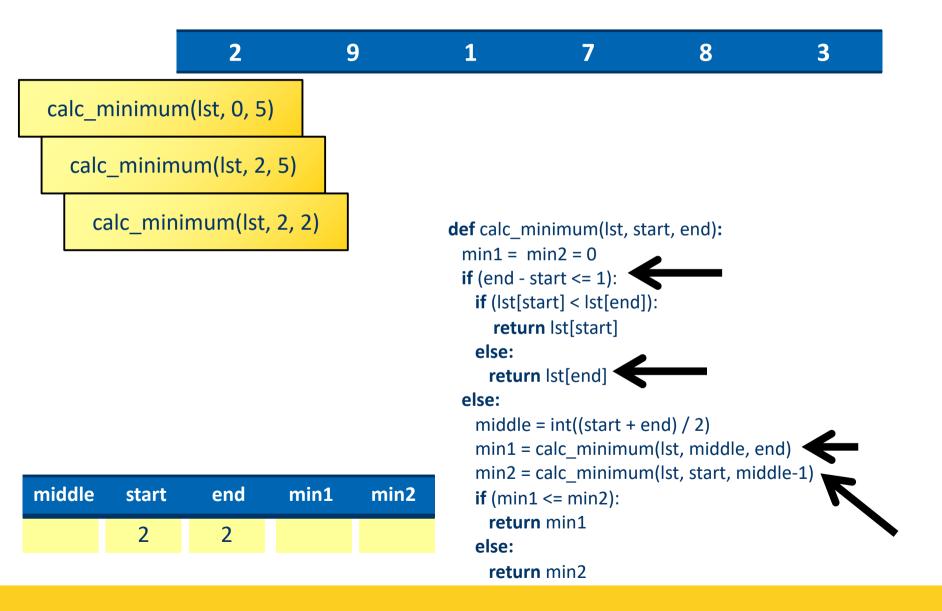
```
middlestartendmin1min2205
```

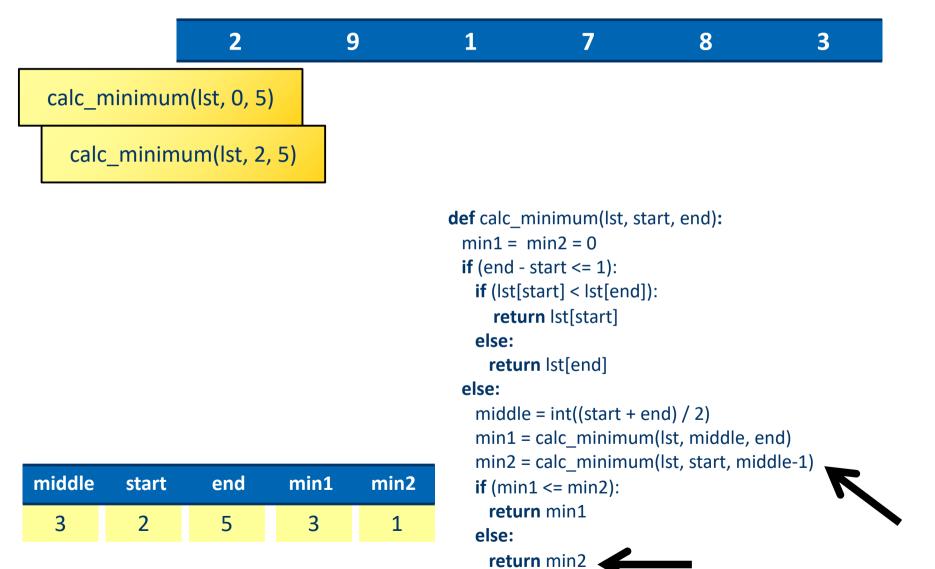












2 9 1 7 8 3

min1 = min2 = 0

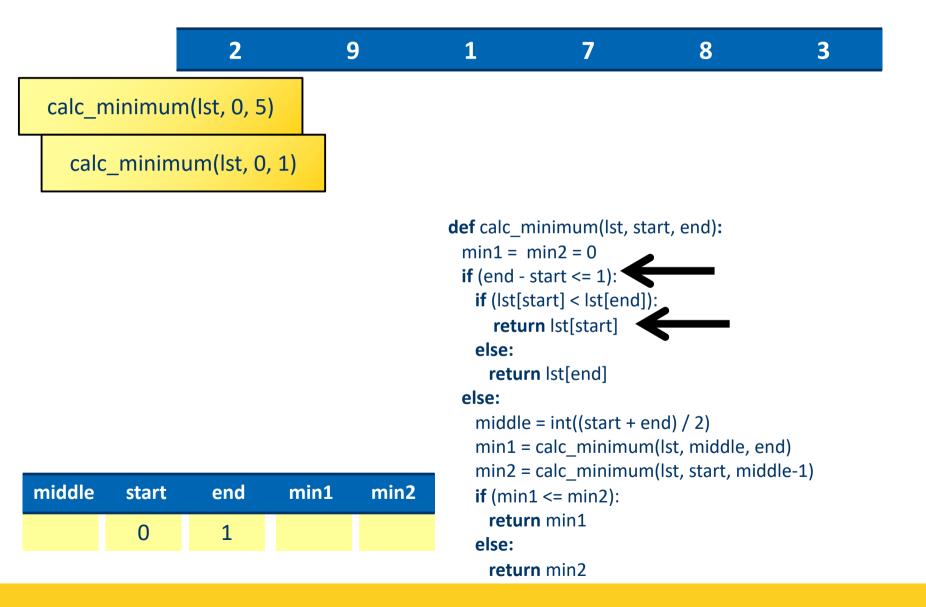
return min2

def calc minimum(lst, start, end):

calc_minimum(lst, 0, 5)

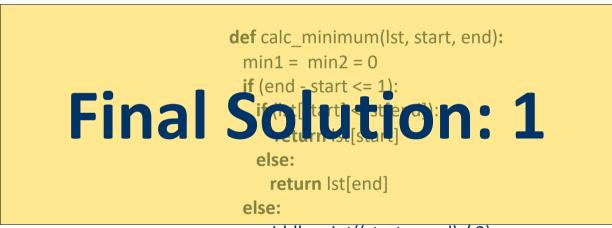
```
if (end - start <= 1):
    if (lst[start] < lst[end]):
        return lst[start]
    else:
        return lst[end]
else:
    middle = int((start + end) / 2)
    min1 = calc_minimum(lst, middle, end)
    min2 = calc_minimum(lst, start, middle-1)
    if (min1 <= min2):
        return min1
    else:</pre>
```

```
middlestartendmin1min22051
```



2 9 1 7 8 3

calc_minimum(lst, 0, 5)



return min2

middle	start	end	min1	min2
2	0	5	1	2

middle = int((start + end) / 2)
min1 = calc_minimum(lst, middle, end)
min2 = calc_minimum(lst, start, middle-1)
if (min1 <= min2):
 return min1
else:





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TRY AND ERROR

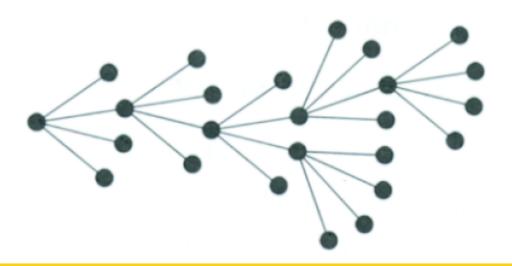
TRY AND ERROR

"Try everything at least once"

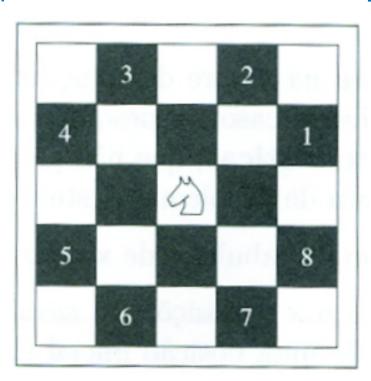
- Uses recursion to solve the problems for which the solution is trying differente alternatives
- Consists of decomposing the process in a finite number of partial subtasks

TRY AND ERROR

- How it Works:
 - Steps towards a solution are attempted and recorded
 - If the steps do not take to the solution, they can be rolled back
 - The search in the solution tree can grow very quickly (exponentially)



- Knight's tour problem
 - Given a n x n board, the Knight makes moves according to the chess rules
 - Given an initial position (x_0, y_0) , the problem is to make the Knight visit every square on the chess board exactly once

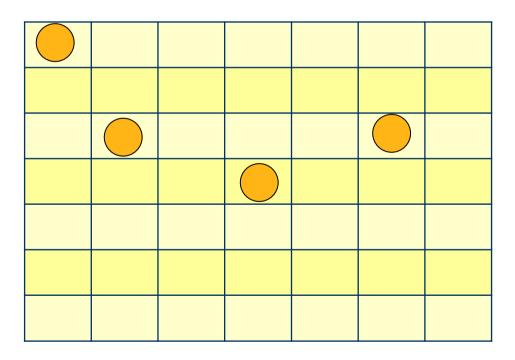


- Computational solution
 - BOARD → nxn matrix
- Each square situation:
 - t[x,y] = 0, <x,y> not visited
 - t[x,y] = i, $\langle x,y \rangle$ visited in the "ith" movement
 - $-1 <= i <= n^2$

Defining the possible moves

$$dr = \{2, 1, -1, -2, -2, -1, 1, 2\}$$

$$dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$$



$$k = 0$$

y + 2 / x + 1

$$k = 1$$

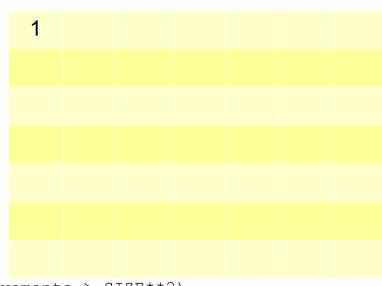
y + 1 / x + 2

$$k = 2$$

y - 1 / x + 2

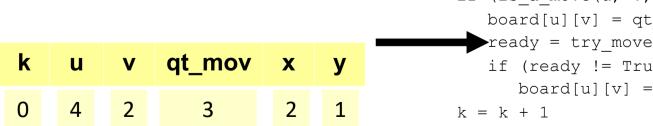
try_move(2,0,0)

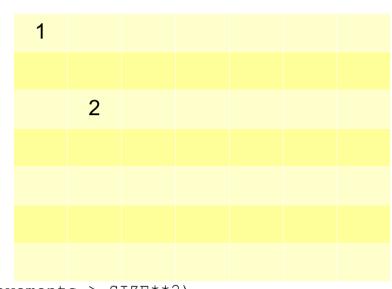
k	u	V	qt_mov	X	у
0	2	1	2	0	0



try_move(2,0,0)

try_move(3,2,1)





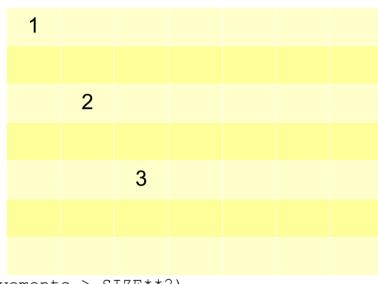
int[]
$$dl = \{2, 1, -1, -2, -2, -1, 1, 2\}$$

int[] $dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$

try_move(2,0,0)

try_move(3,2,1)

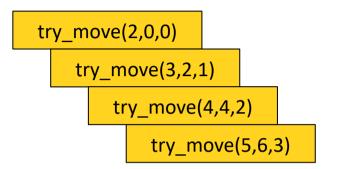
try_move(4,4,2)



 k
 u
 v
 qt_mov
 x
 y

 0
 6
 3
 4
 4
 2

int[] dl = {2, 1, -1, -2, -2, -1, 1, 2} int[] dc = {1, 2, 2, 1, -1, -2, -2, -1}

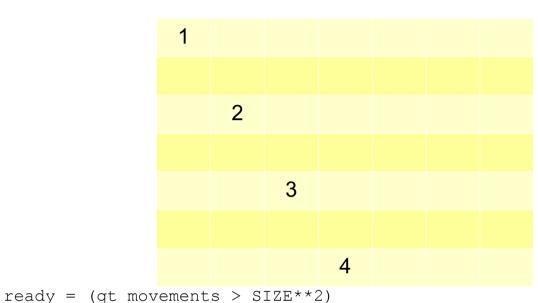


v qt_mov

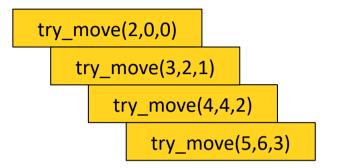
6

k u

8

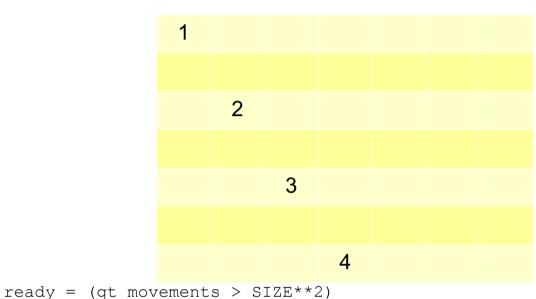


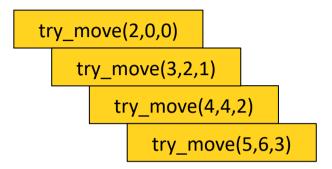
int[] dl = {2, 1, -1, -2, -2, -1, 1, 2} int[] dc = {1, 2, 2, 1, -1, -2, -2, -1}

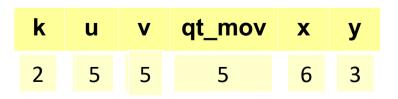


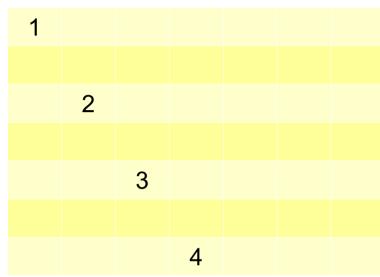
u v qt_mov

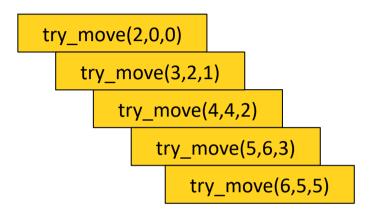
6









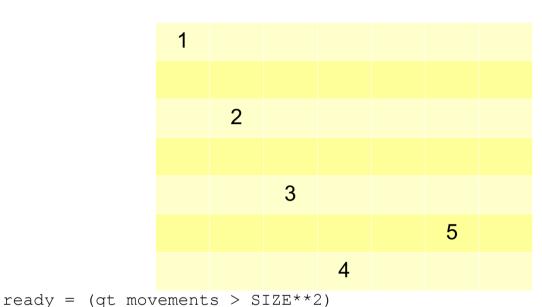


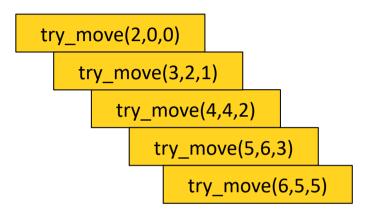
v qt_mov

5

5

k



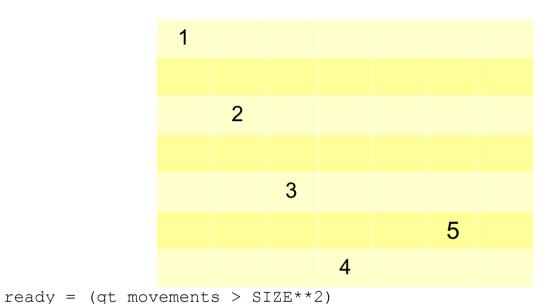


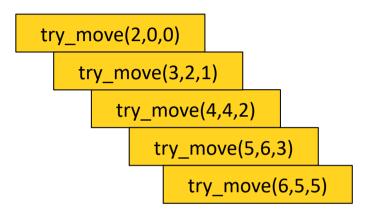
v qt_mov

5

k

6

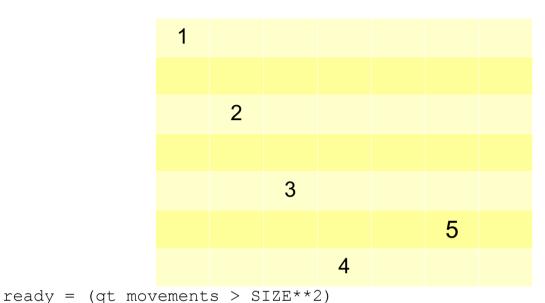


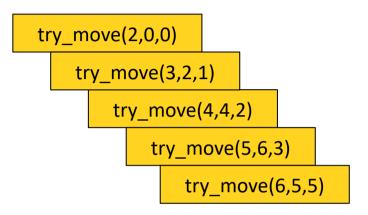


v qt_mov

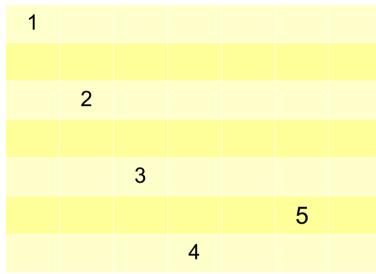
5

k



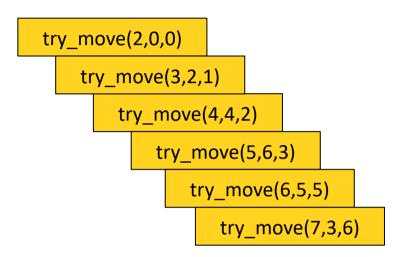


k	u	V	qt_mov	X	у
3	3	6	6	5	5



int[]
$$dl = \{2, 1, -1, -2, -2, -1, 1, 2\}$$

int[] $dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$



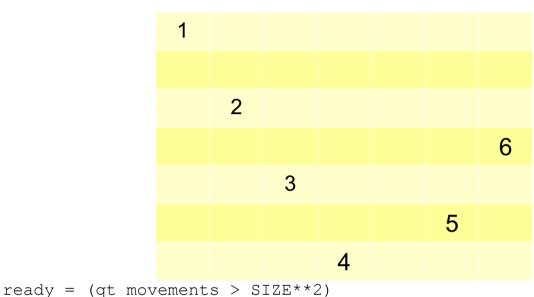
v qt_mov

3

6

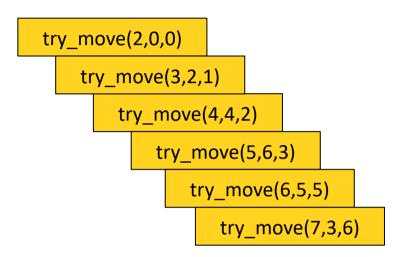
k

0



int[]
$$dl = \{2, | 1, -1, | -2, | -2, -1, | 1, | 2\}$$

int[] $dc = \{1, | 2, | 2, | 1, | -1, | -2, | -2, | -1\}$



qt_mov

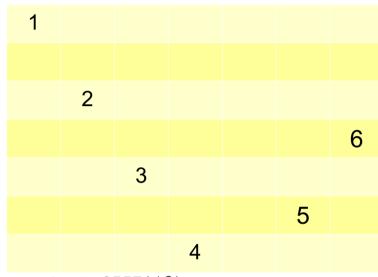
3

6

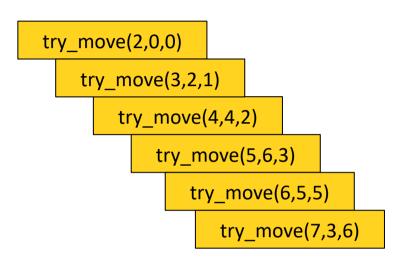
8

4

k



```
int[] dl = {2, 1, -1, -2, -2, -1, 1, 2}
int[] dc = {1, 2, 2, 1, -1, -2, -2, -1}
```

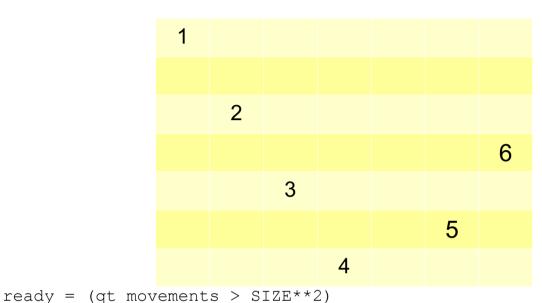


v qt_mov

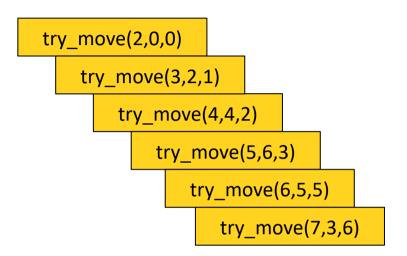
3

8

k

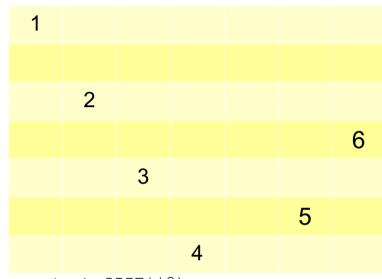


```
int[] dl = {2, 1, -1, -2, -2, -1, 1, 2}
int[] dc = {1, 2, 2, 1, -1, -2, -2, -1}
```

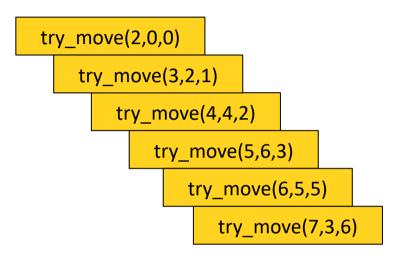


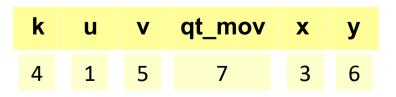
v qt_mov

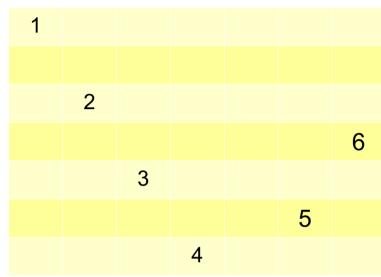
k



7 3 6
$$k = k + 1$$
 return ready

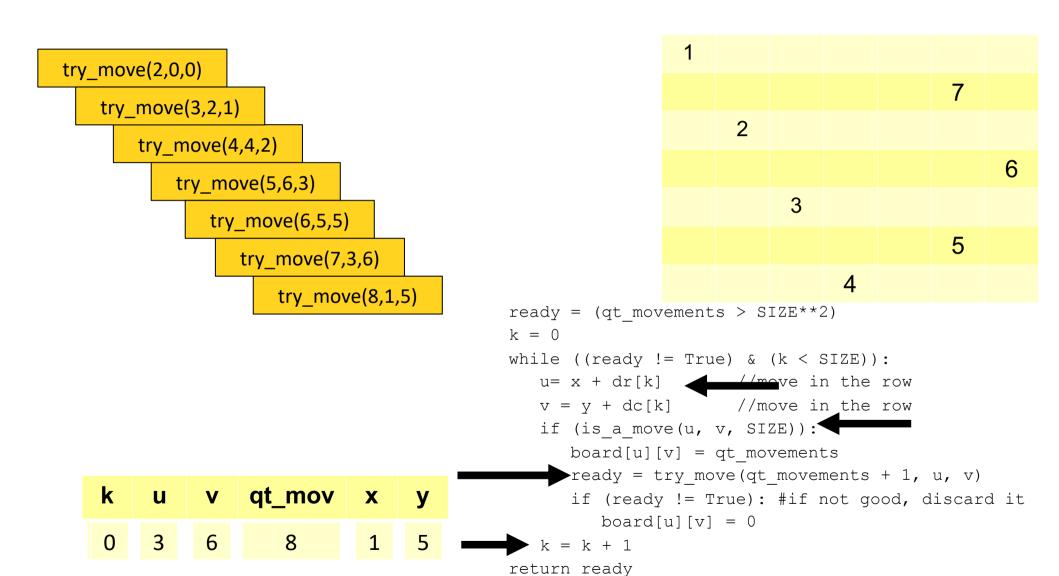






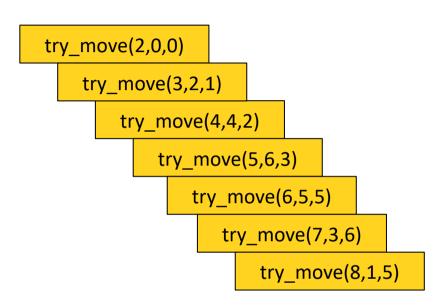
int[]
$$dl = \{2, 1, -1, -2, -2, -1, 1, 2\}$$

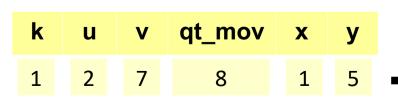
int[] $dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$

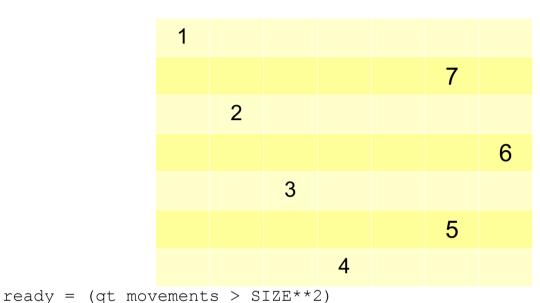


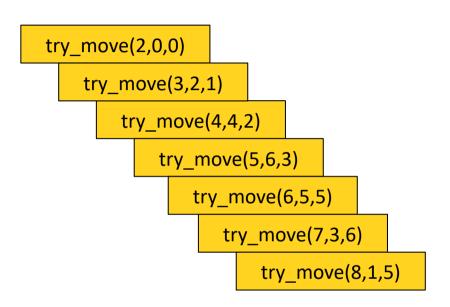
 $int[] dl = \{2, |1, -1, -2, |-2, |-1, 1, 2\}$

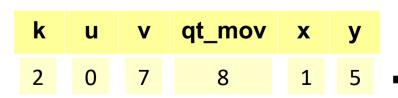
 $int[] dc = |\{1, | 2, 2, 1, |-1, | -2, -2, -1\}|$

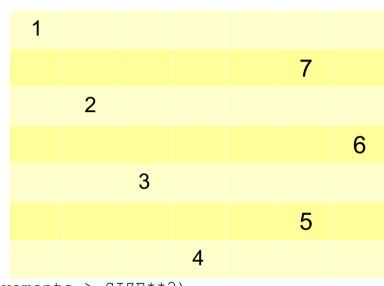


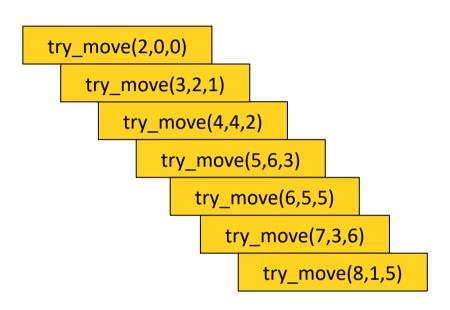


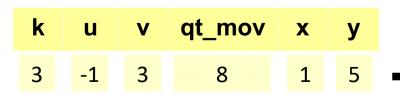


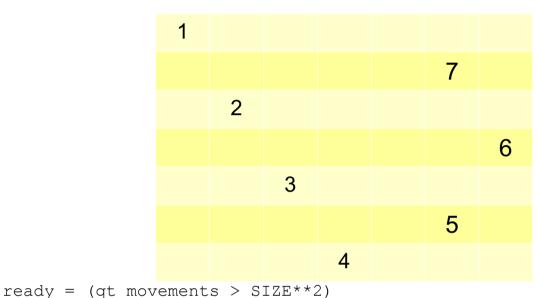






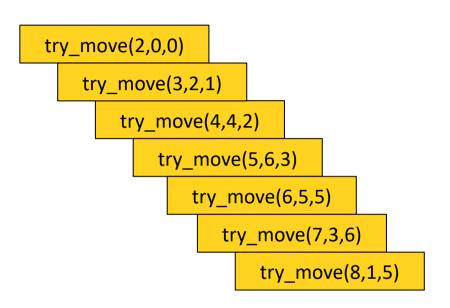


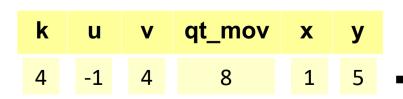


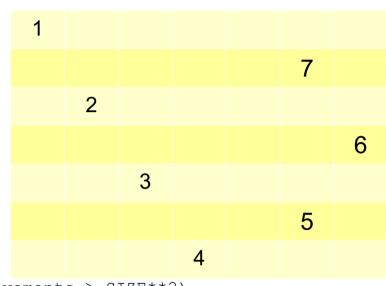


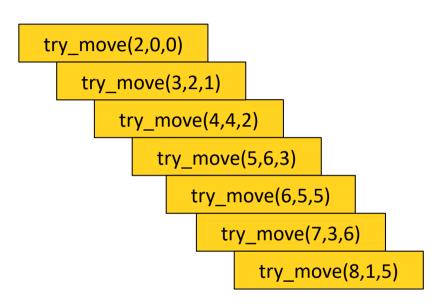
int[]
$$dl = \{2, 1, -1, -2, -2, -1, 1, 2\}$$

int[] $dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$

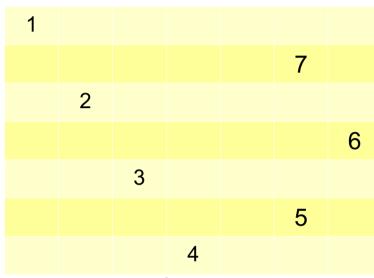


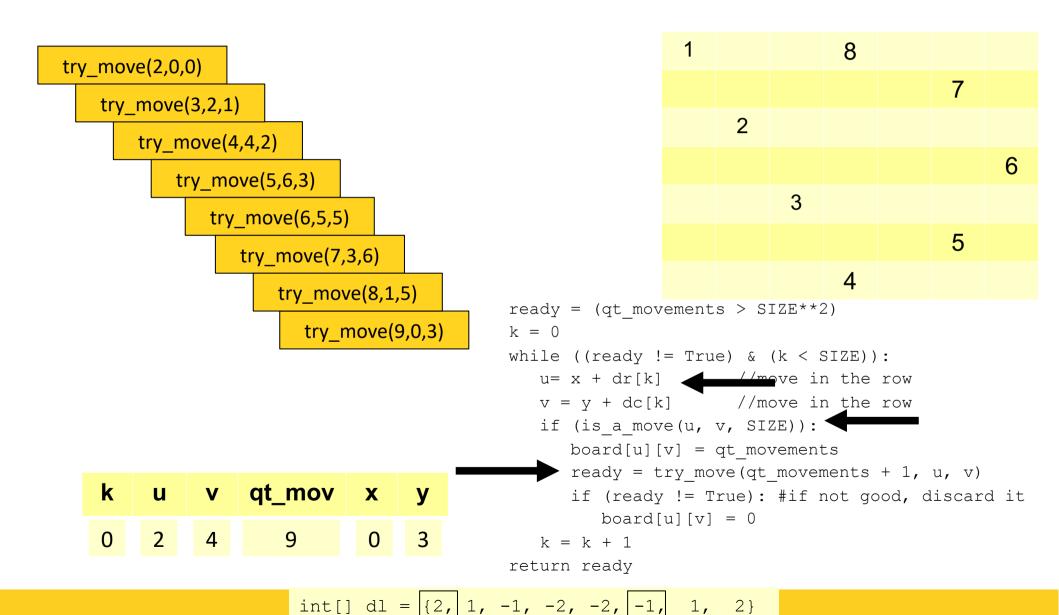




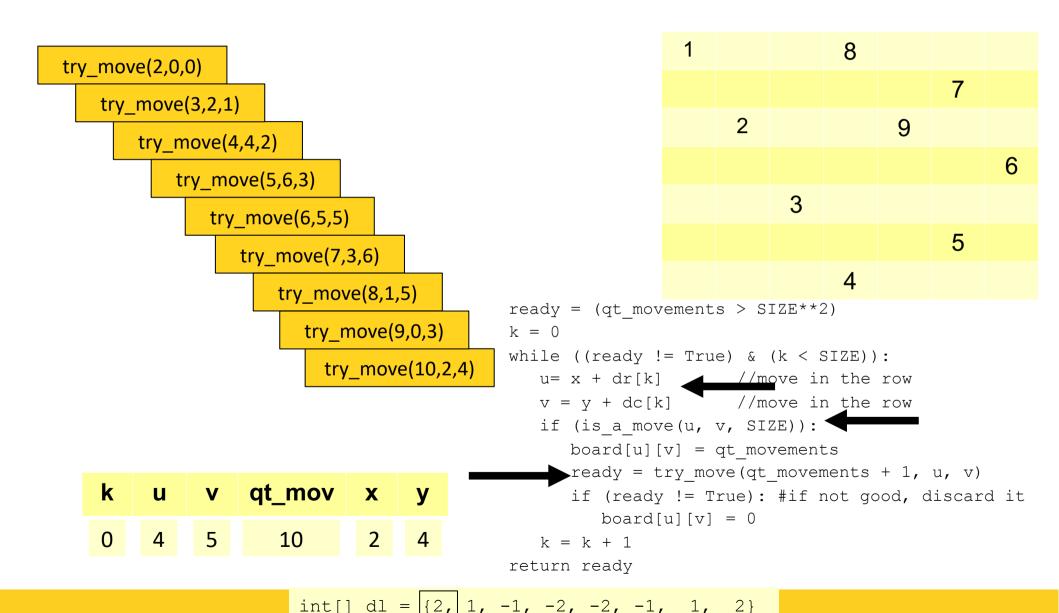


k	u	V	qt_mov	X	у
5	0	3	8	1	5

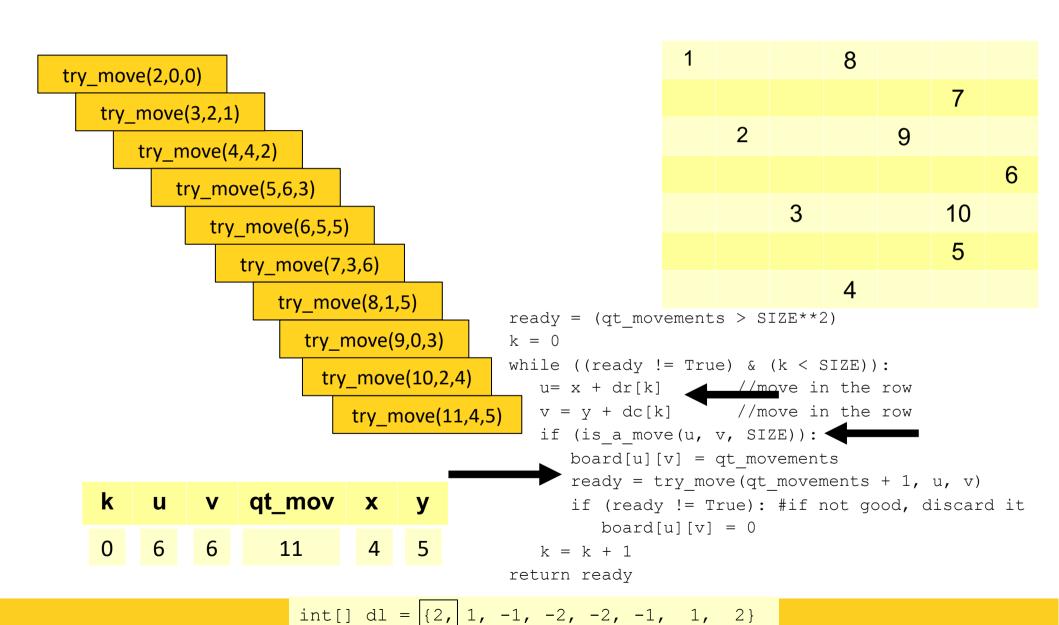




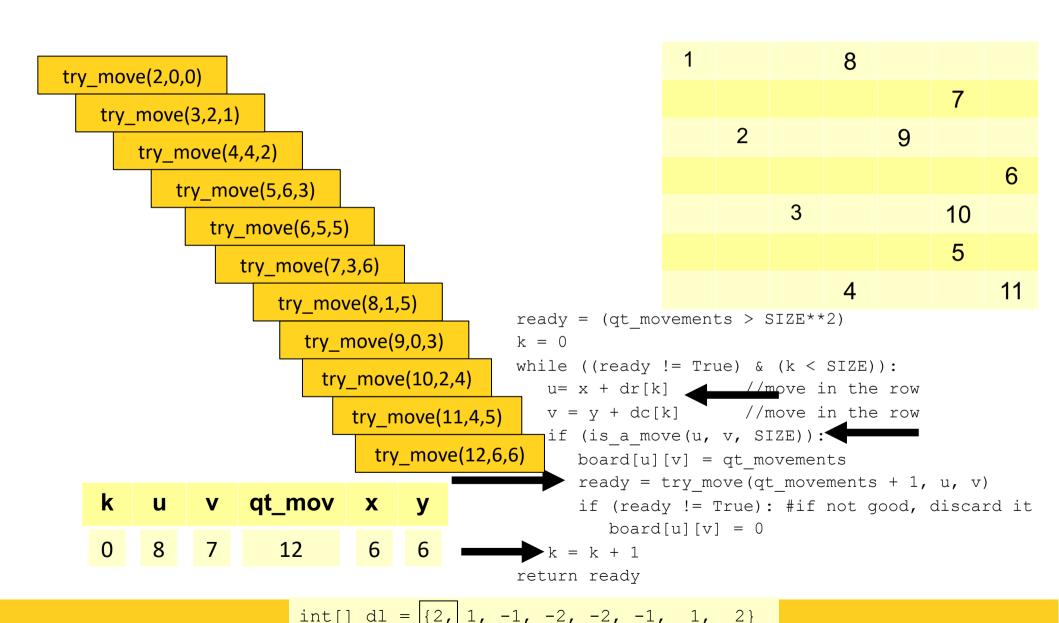
int[] $dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$



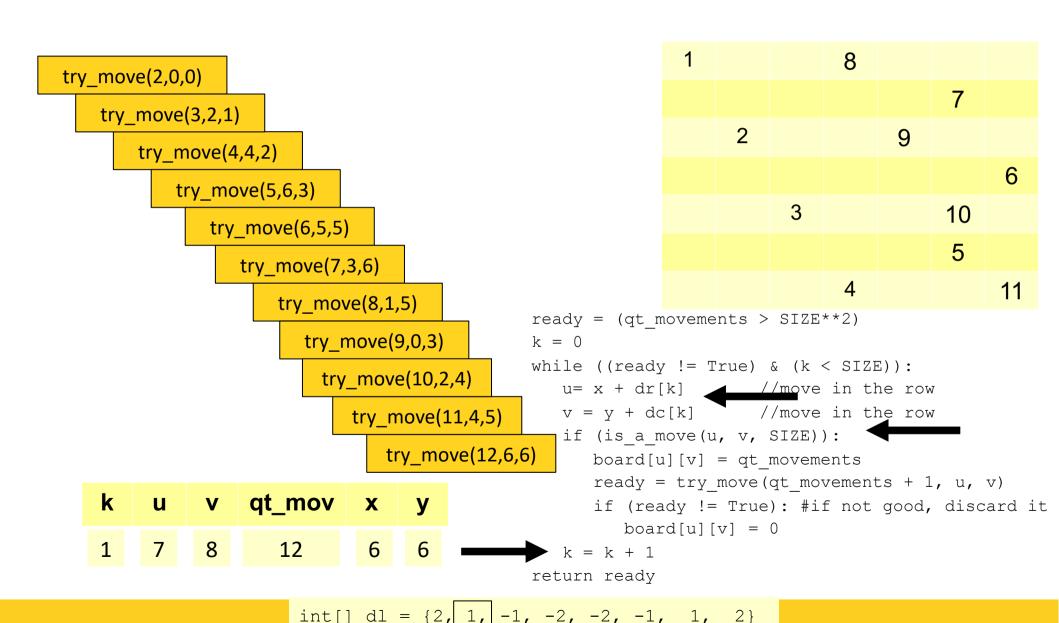
 $int[] dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$



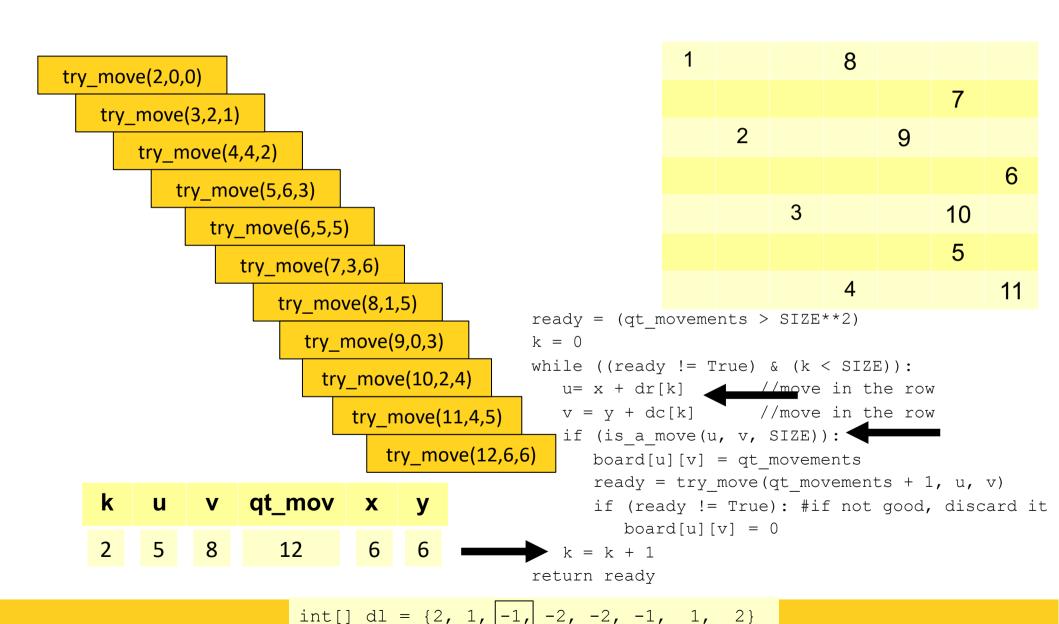
 $int[] dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$



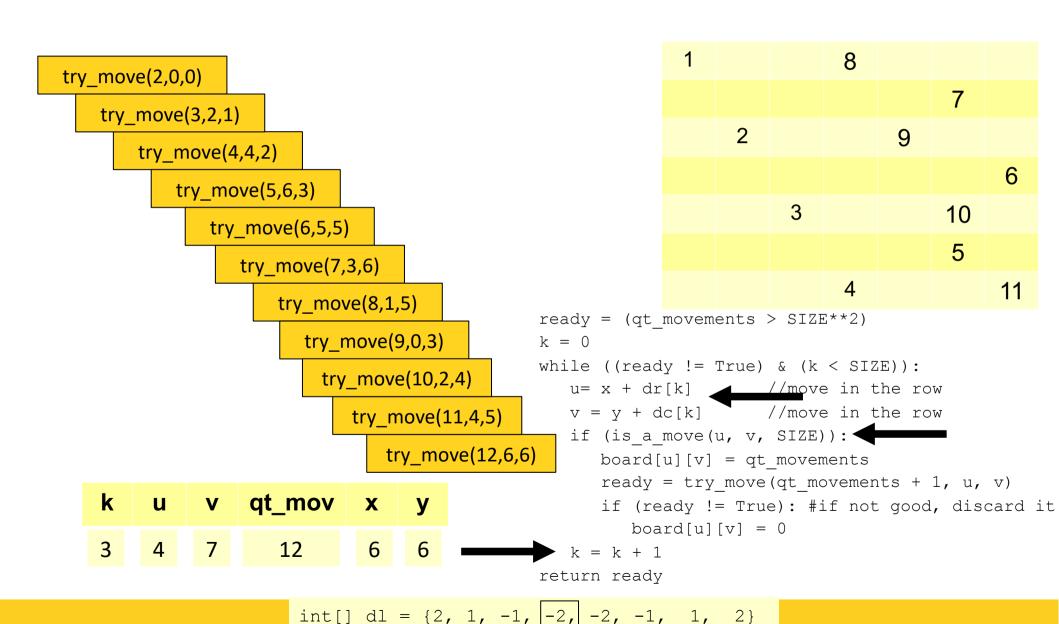
 $int[] dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$



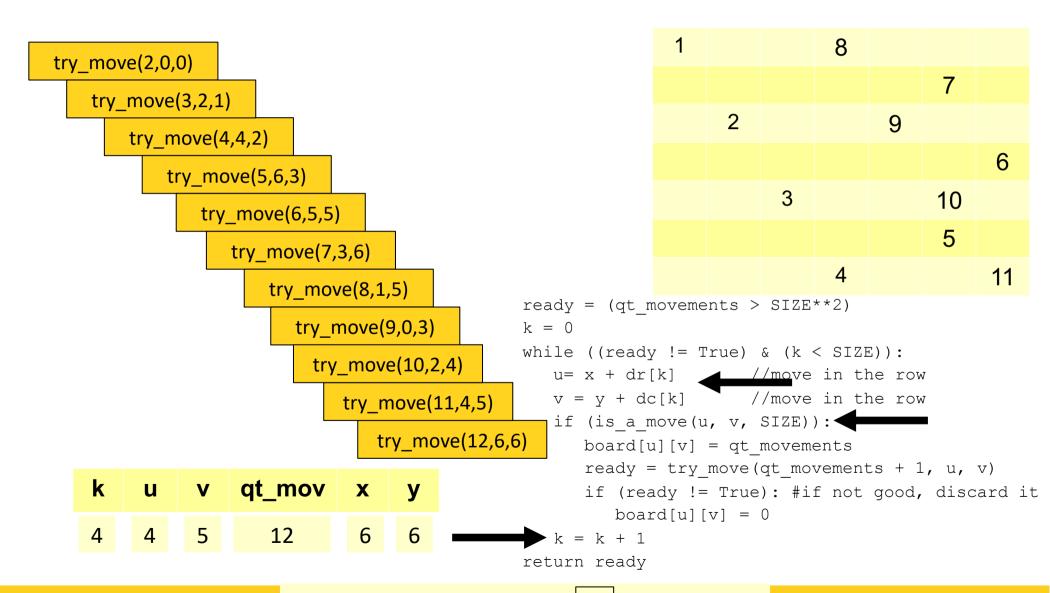
 $int[] dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$



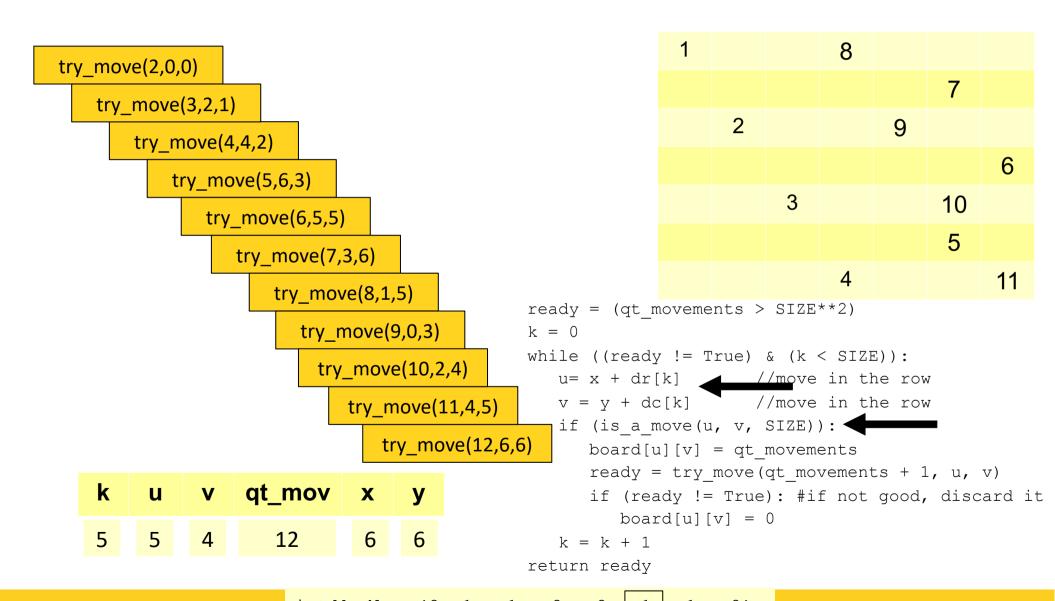
 $int[] dc = \{1, 2, | 2, | 1, -1, -2, -2, -1\}$



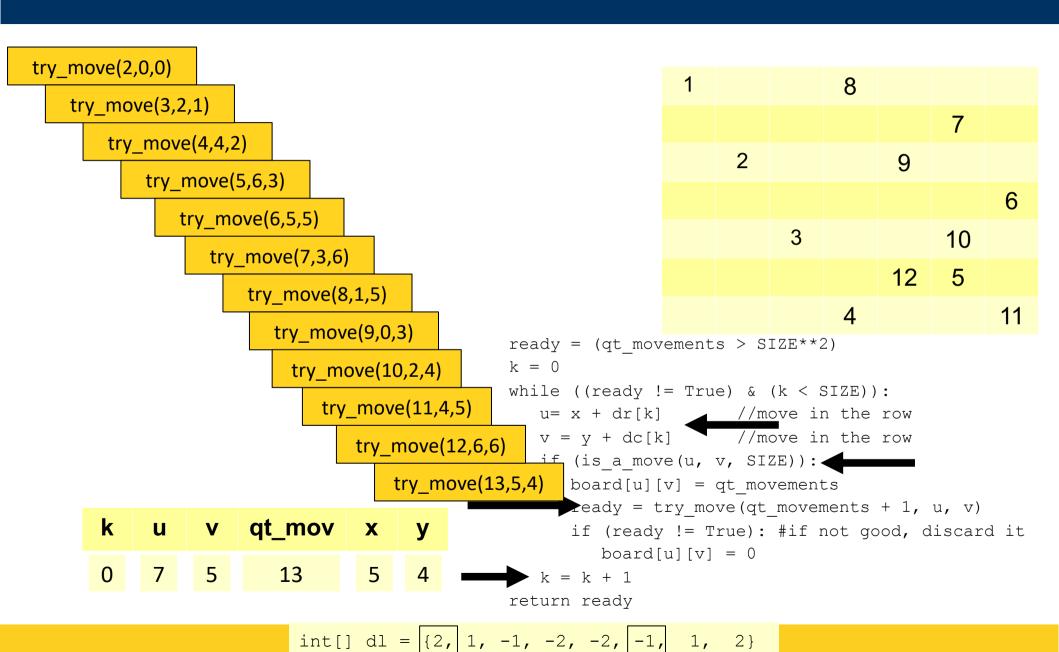
 $int[] dc = \{1, 2, 2, |1, -1, -2, -2, -1\}$



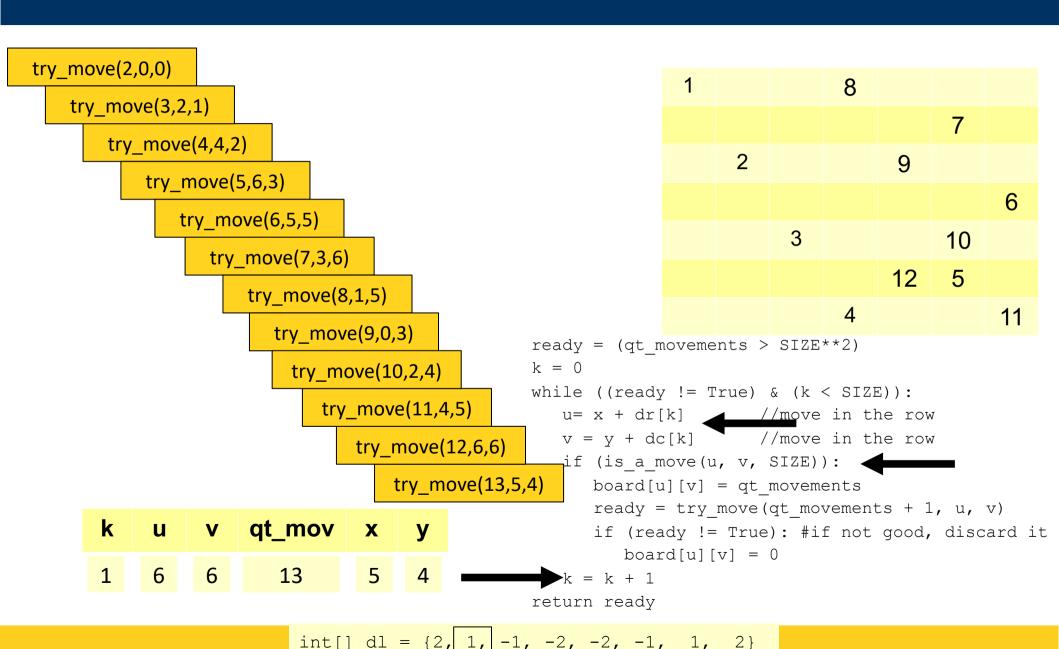
int[] $dl = \{2, 1, -1, -2, -2, -1, 1, 2\}$ int[] $dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$



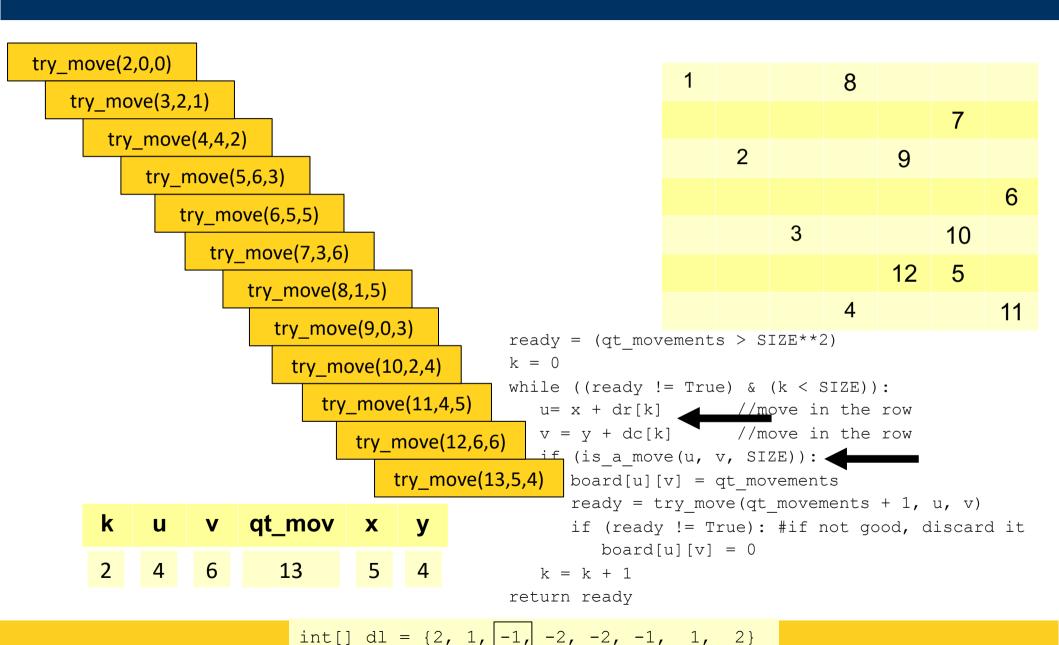
int[] $dl = \{2, 1, -1, -2, -2, -1, 1, 2\}$ int[] $dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$



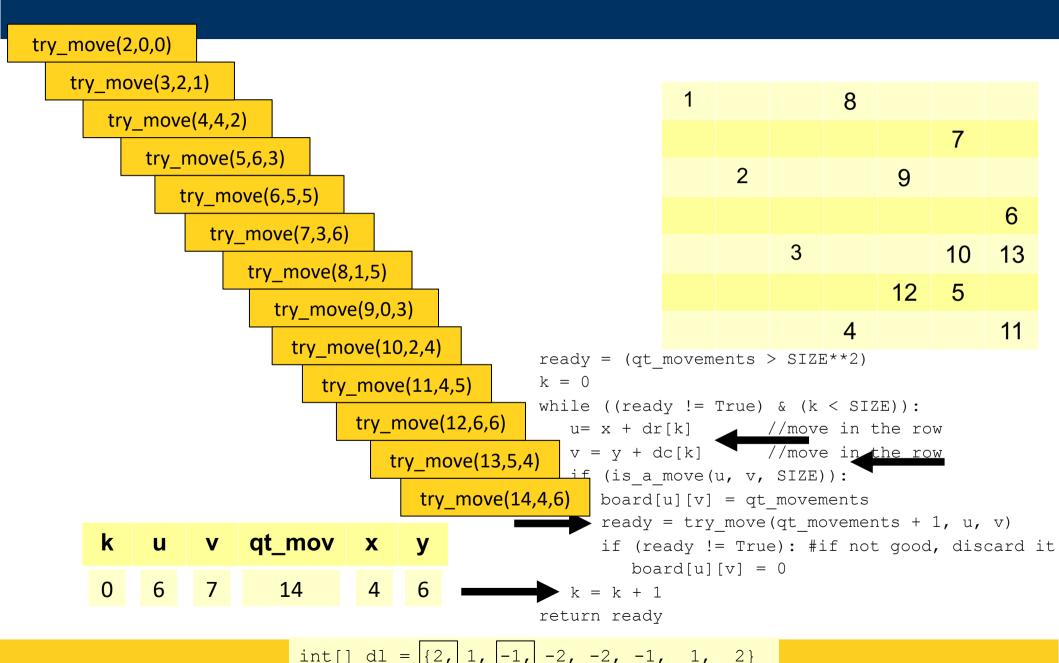
 $int[] dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$

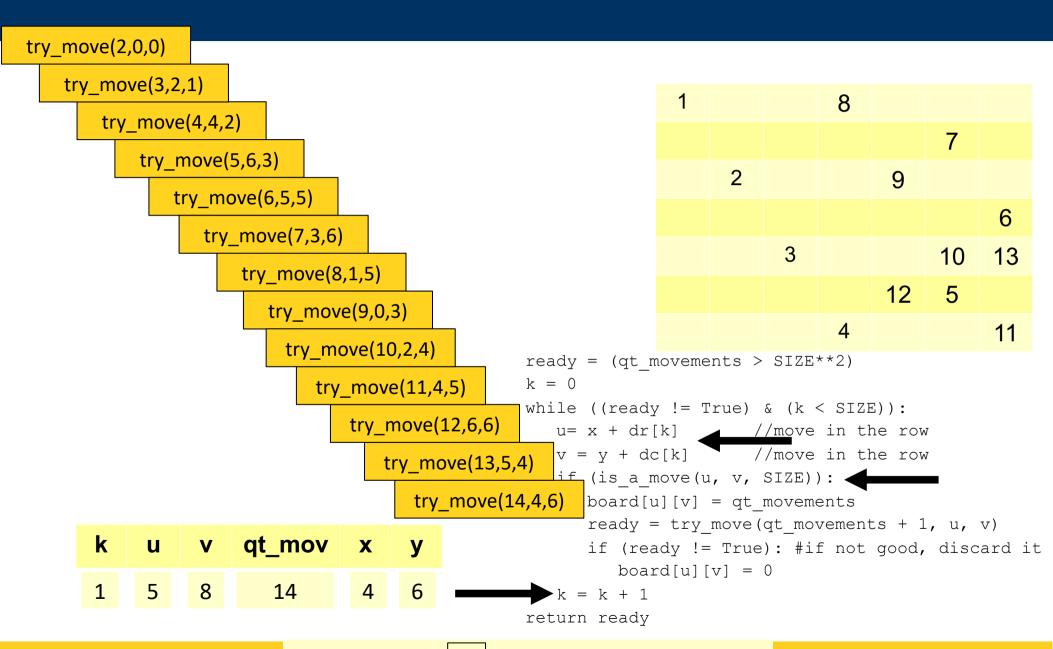


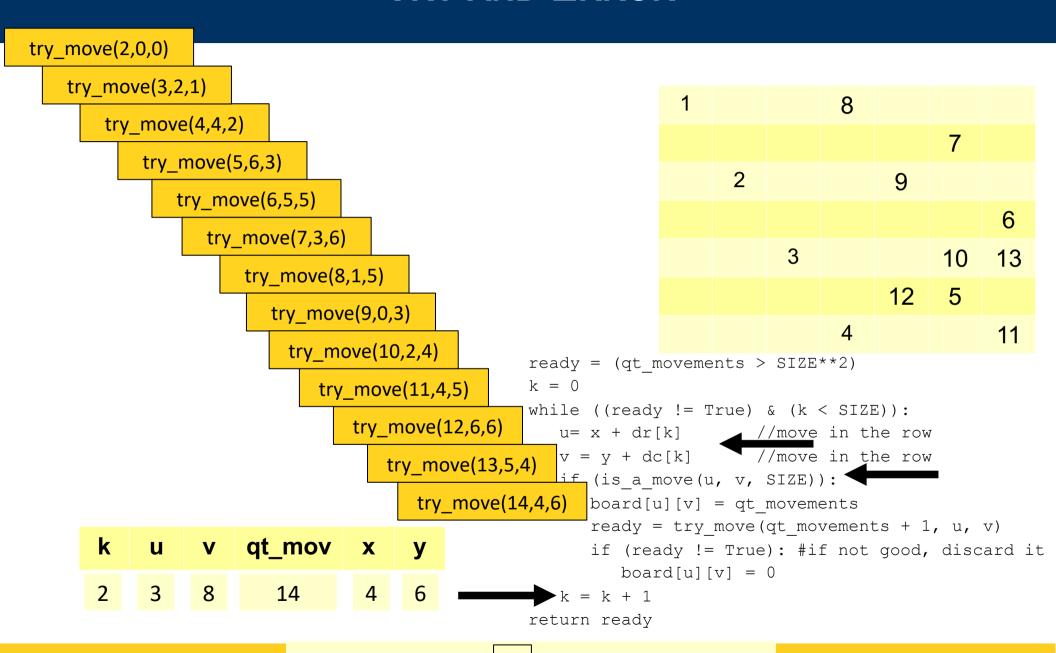
 $int[] dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$

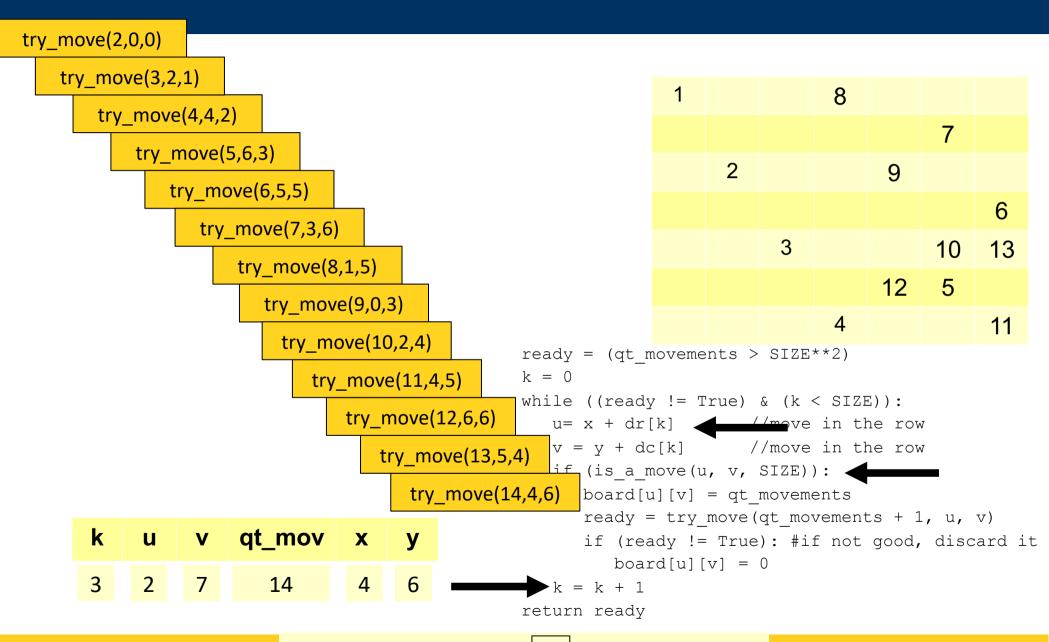


 $int[] dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$

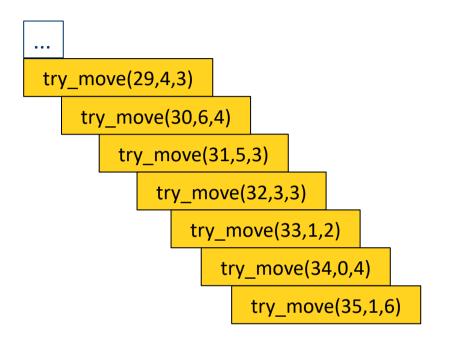








Fast-forwarding...

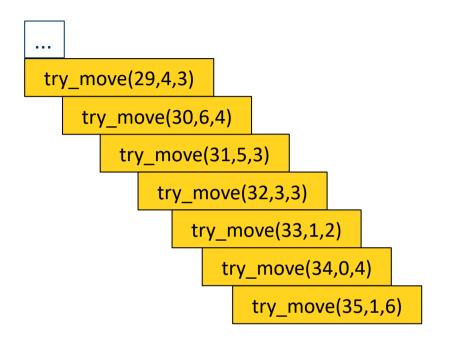


k	u	V	qt_mov	X	у	
0	3	7	35	1	6	

26		8	33	24	15
	32	25	26	7	34
2	27		9	14	23
		31	22	17	6
	3	28	19	10	13
	30	21	12	5	18
		4	29	20	11
		32 2 27 3 30	32 25 2 27 31 3 28 30 21	32 25 26 2 27 9 31 22 3 28 19 30 21 12 4 29	32 25 26 7 2 27 9 14 31 22 17 3 28 19 10 30 21 12 5 4 29 20

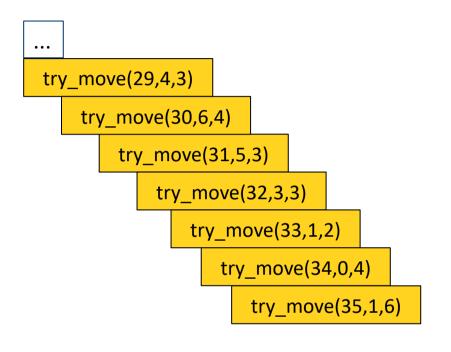
int[]
$$dl = \{2, | 1, -1, -2, -2, -1, 1, 2\}$$

int[] $dc = \{1, | 2, 2, 1, -1, -2, -2, -1\}$



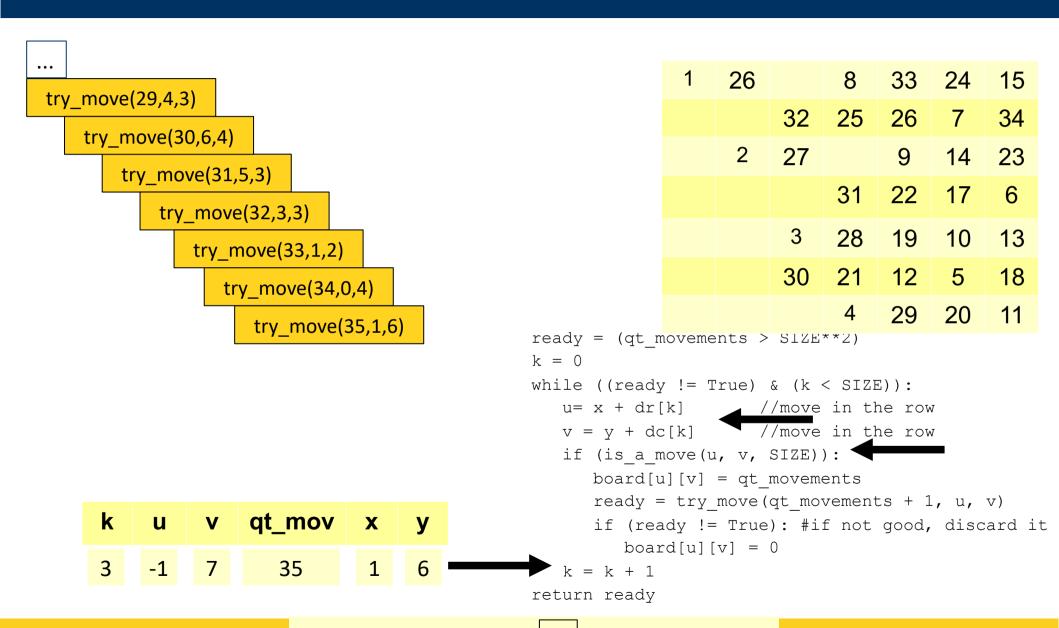
k	u	V	qt_mov	X	у	
1	2	8	35	1	6	

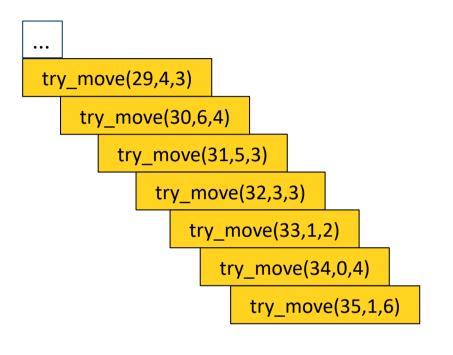
1	26		8	33	24	15
		32	25	26	7	34
	2	27		9	14	23
			31	22	17	6
		3	28	19	10	13
		30	21	12	5	18
			4	29	20	11



k	u	V	qt_mov	X	у	
2	0	8	35	1	6	

1	26		8	33	24	15
		32	25	26	7	34
	2	27		9	14	23
			31	22	17	6
		3	28	19	10	13
		30	21	12	5	18
			4	29	20	11



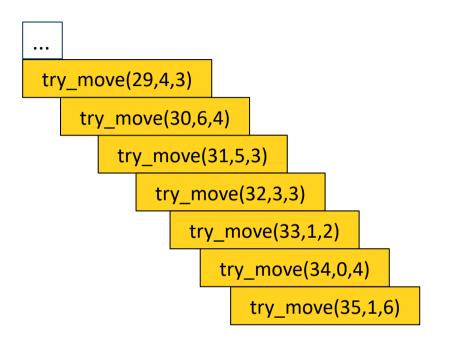


k	u	V	qt_mov	X	у
4	-1	5	35	1	6

1	26		8	33	24	15
		32	25	26	7	34
	2	27		9	14	23
			31	22	17	6
		3	28	19	10	13
		30	21	12	5	18
			4	29	20	11

int[]
$$dl = \{2, 1, -1, -2, -2, -1, 1, 2\}$$

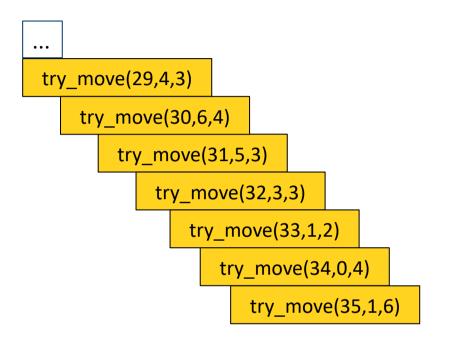
int[] $dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$



k	u	V	qt_mov	X	У
5	0	4	35	1	6

```
1
    26
                  33
                       24
                          15
              8
        32
             25
                  26
                           34
        27
                       14
                           23
                  9
             31
                  22
                       17
         3
             28
                  19
                       10
                            13
        30
                            18
             21
                  12
              4
                  29
                       20
                            11
```

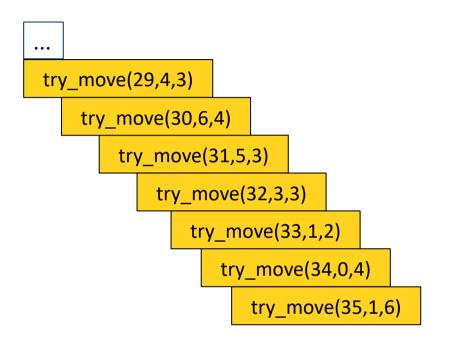
int[] dl = {2, 1, -1, -2, -2, -1, 1, 2} int[] dc = {1, 2, 2, 1, -1, -2, -2, -1}



k	u	V	qt_mov	X	у
6	2	4	35	1	6

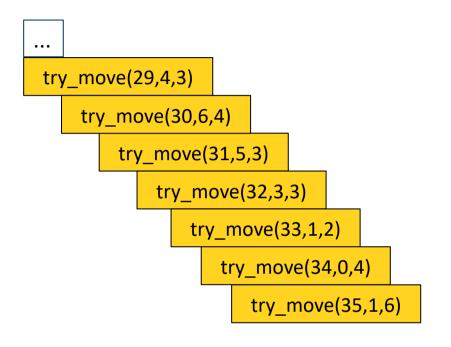
1	26		8	33	24	15
		32	25	26	7	34
	2	27		9	14	23
			31	22	17	6
		3	28	19	10	13
		30	21	12	5	18
			4	29	20	11

int[] dl = {2, 1, -1, -2, -2, -1, 1, 2} int[] dc = {1, 2, 2, 1, -1, -2, -2, -1}



k	u	V	qt_mov	X	У
7	3	5	35	1	6

1	26		8	33	24	15
		32	25	26	7	34
	2	27		9	14	23
			31	22	17	6
		3	28	19	10	13
		30	21	12	5	18
			4	29	20	11

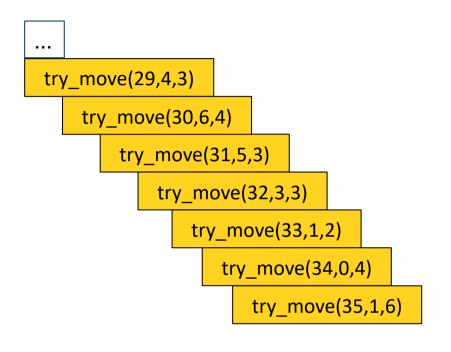


k	u	V	qt_mov	X	у	
8	3	5	35	1	6	

1	26		8	33	24	15
		32	25	26	7	34
	2	27		9	14	23
			31	22	17	6
		3	28	19	10	13
		30	21	12	5	18
			4	29	20	11

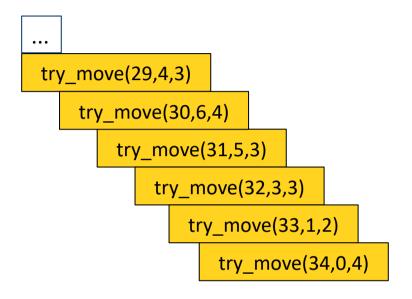
int[]
$$dl = \{2, 1, -1, -2, -2, -1, 1, 2\}$$

int[] $dc = \{1, 2, 2, 1, -1, -2, -2, -1\}$



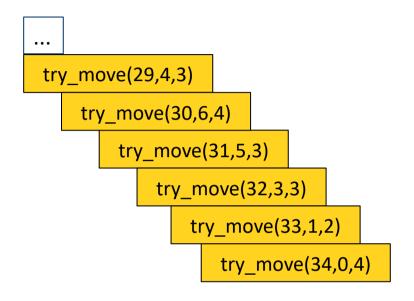
k	u	V	qt_mov	X	У	
1	1	6	34	0	4	

	1	26		8	33	24	15
			32	25	26	7	34
		2	27		9	14	23
				31	22	17	6
			3	28	19	10	13
			30	21	12	5	18
					29	20	11
$ready = (qt_n$	noveme	ents >	SIZE	**2)			



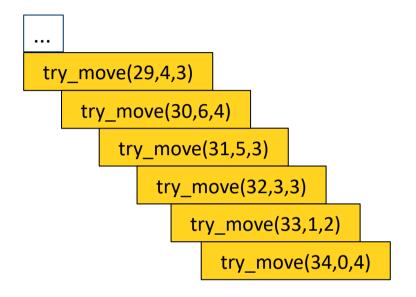
k	u	V	qt_mov	X	У
1	1	6	34	0	4

1	26		8	33	24	15
		32	25	26	7	0
	2	27		9	14	23
			31	22	17	6
		3	28	19	10	13
		30	21	12	5	18
			4	29	20	11



k	u	u v qt_mov		X	У
1	1	6	34	0	4

1	26		8	33	24	15
		32	25	26	7	
	2	27		9	14	23
			31	22	17	6
		3	28	19	10	13
		30	21	12	5	18
			4	29	20	11



			3	28	19	10	13	
			30	21	12	5	18	
				4	29	20	11	
ready = (qt movements > SIZE**2)								
$\zeta = 0$								
nile ((ready	/ != T	rue)	& (k	< SIZI	Ξ)):			
u = x + dr	[k]	/	/move	in th	ne row	J		
v = y + dc								
if (is a r	nove (u	l, V,	SIZE)): -	lacktriangle			
board[u][v] = qt movements								
ready = try move(qt movements + 1, u, v)								
if (rea	_	_	_					t

board[u][v] = 0

```
        k
        u
        v
        qt_mov
        x
        y

        2
        -1
        6
        34
        0
        4
```

int[] dl = {2, 1, -1, -2, -2, -1, 1, 2} int[] dc = {1, 2, 2, 1, -1, -2, -2, -1}

k = k + 1 return ready

24 15

It is possible to make the 34th move with k=7

1	26		8	33	24	15
		32	25	26	7	
	2	27	34	9	14	23
			31	22	17	6
		3	28	19	10	13
		30	21	12	5	18
			4	29	20	11

• Then, add 35 to (0,2)

1	26	35	8	33	24	15
		32	25	26	7	
	2	27	34	9	14	23
			31	22	17	6
		3	28	19	10	13
		30	21	12	5	18
			4	29	20	11

• Then, add 37 to (1, 0)

1	26	35	8	33	24	15
37		32	25	26	7	
	2	27	34	9	14	23
			31	22	17	6
		3	28	19	10	13
		30	21	12	5	18
			4	29	20	11

• Then, add 38 to (3, 1)

1	26	35	8	33	24	15
37		32	25	26	7	
	2	27	34	9	14	23
	38		31	22	17	6
		3	28	19	10	13
		30	21	12	5	18
			4	29	20	11

• And 39 to (5, 0)

1	26	35	8	33	24	15
37		32	25	26	7	
	2	27	34	9	14	23
	38		31	22	17	6
		3	28	19	10	13
39		30	21	12	5	18
			4	29	20	11

• $40 \rightarrow (6, 2)$

1	26	35	8	33	24	15
37		32	25	26	7	
	2	27	34	9	14	23
	38		31	22	17	6
		3	28	19	10	13
39		30	21	12	5	18
		40	4	29	20	11

• $42 \rightarrow (2, 0)$

1	26	35	8	33	24	15
37		32	25	26	7	
42	2	27	34	9	14	23
	38		31	22	17	6
	41	3	28	19	10	13
39		30	21	12	5	18
		40	4	29	20	11

• $41 \rightarrow (4, 1)$

1	26	35	8	33	24	15
37		32	25	26	7	
	2	27	34	9	14	23
	38		31	22	17	6
	41	3	28	19	10	13
39		30	21	12	5	18
		40	4	29	20	11

• $43 \rightarrow (3, 2)$

1	26	35	8	33	24	15
37		32	25	26	7	
42	2	27	34	9	14	23
	38	43	31	22	17	6
	41	3	28	19	10	13
39		30	21	12	5	18
		40	4	29	20	11

• $44 \rightarrow (1, 1)$

1	26	35	8	33	24	15
37	44	32	25	26	7	
42	2	27	34	9	14	23
	38	43	31	22	17	6
	41	3	28	19	10	13
39		30	21	12	5	18
		40	4	29	20	11

• 45 to (3, 0)

1	26	35	8	33	24	15
37	44	32	25	26	7	
42	2	27	34	9	14	23
45	38	43	31	22	17	6
	41	3	28	19	10	13
39		30	21	12	5	18
		40	4	29	20	11

• $46 \rightarrow (5, 1)$

1	26	35	8	33	24	15
37	44	32	25	26	7	
42	2	27	34	9	14	23
45	38	43	31	22	17	6
	41	3	28	19	10	13
39	46	30	21	12	5	18
		40	4	29	20	11

- Impossible to make 47!!!!
 - Rollback 46 to 0 and try another move from 45

1	26	35	8	33	24	15
37	44	32	25	26	7	
42	2	27	34	9	14	23
45	38	43	31	22	17	6
	41	3	28	19	10	13
39	0	30	21	12	5	18
		40	4	29	20	11

- Not possible to make another 46
 - Rollback 45, and go to 44 to try another move

1	26	35	8	33	24	15
37	44	32	25	26	7	
42	2	27	34	9	14	23
0	38	43	31	22	17	6
	41	3	28	19	10	13
39		30	21	12	5	18
		40	4	29	20	11

- Oh no!!!! Impossible to make 45
 - Rollbak 44 and go back to 43

1	26	35	8	33	24	15
37	0	32	25	26	7	
42	2	27	34	9	14	23
	38	43	31	22	17	6
	41	3	28	19	10	13
39		30	21	12	5	18
		40	4	29	20	11

- Yahoo!!! Made it
 - And tries to make 45
 - It will come to na end one day...

1	26	35	8	33	24	15
37		32	25	26	7	
42	2	27	34	9	14	23
	38	43	31	22	17	6
44	41	3	28	19	10	13
39		30	21	12	5	18
		40	4	29	20	11

Solution

1	38	31	8	19	36	15
32	29	20	37	16	7	18
39	2	33	30	9	14	35
28	25	40	21	34	17	6
41	22	3	26	45	10	13
24	27	48	43	12	5	46
49	42	23	4	47	44	11



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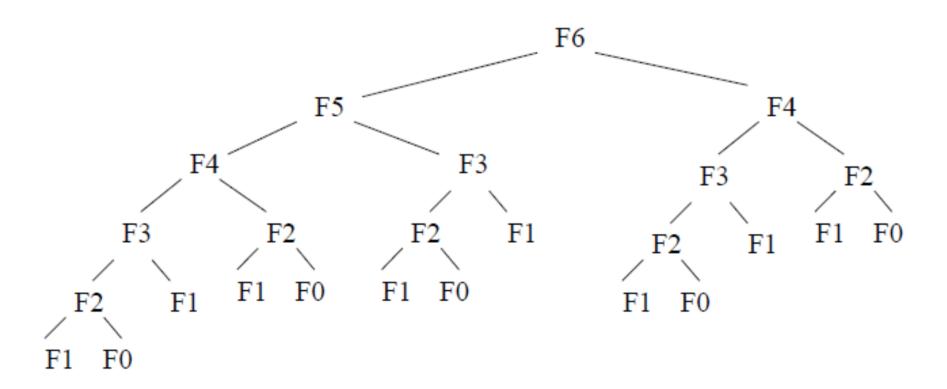
"Life can only be understood **backwards**; but it must be lived **forwards**." (Kierkegaard)

- Most used method in optimization problems
- Applicable to problems in which the optimal solution can be computed from the combination of optimal solutions previously calculated and memorized

- Characteristics of algorithms
 - Optimal substructure
 - A global optimal solution is composed of optimal solution to the subproblems
 - Overlapping subproblems
 - The global solution uses the solution of same subproblems multiple times

- Principle of optimality
 - The complete problem can be solved if the values of each subproblems' best solutions had been previously determined
 - Example
 - If the shortest path between Flagstaff and Tucson goes through Phoenix, so the path between Phoenix and Tucson is the shortest
 - The path between Flagstaff and Phoenix too!

- Fibonacci Algorithm
 - To calculate the Fibonacci of a number, several values are calculated more than once.



- Strategy
 - Store the values we've already calculated
 - When an already-calculated value is requested, no operation is necessary
 - Only new values are calculated

