

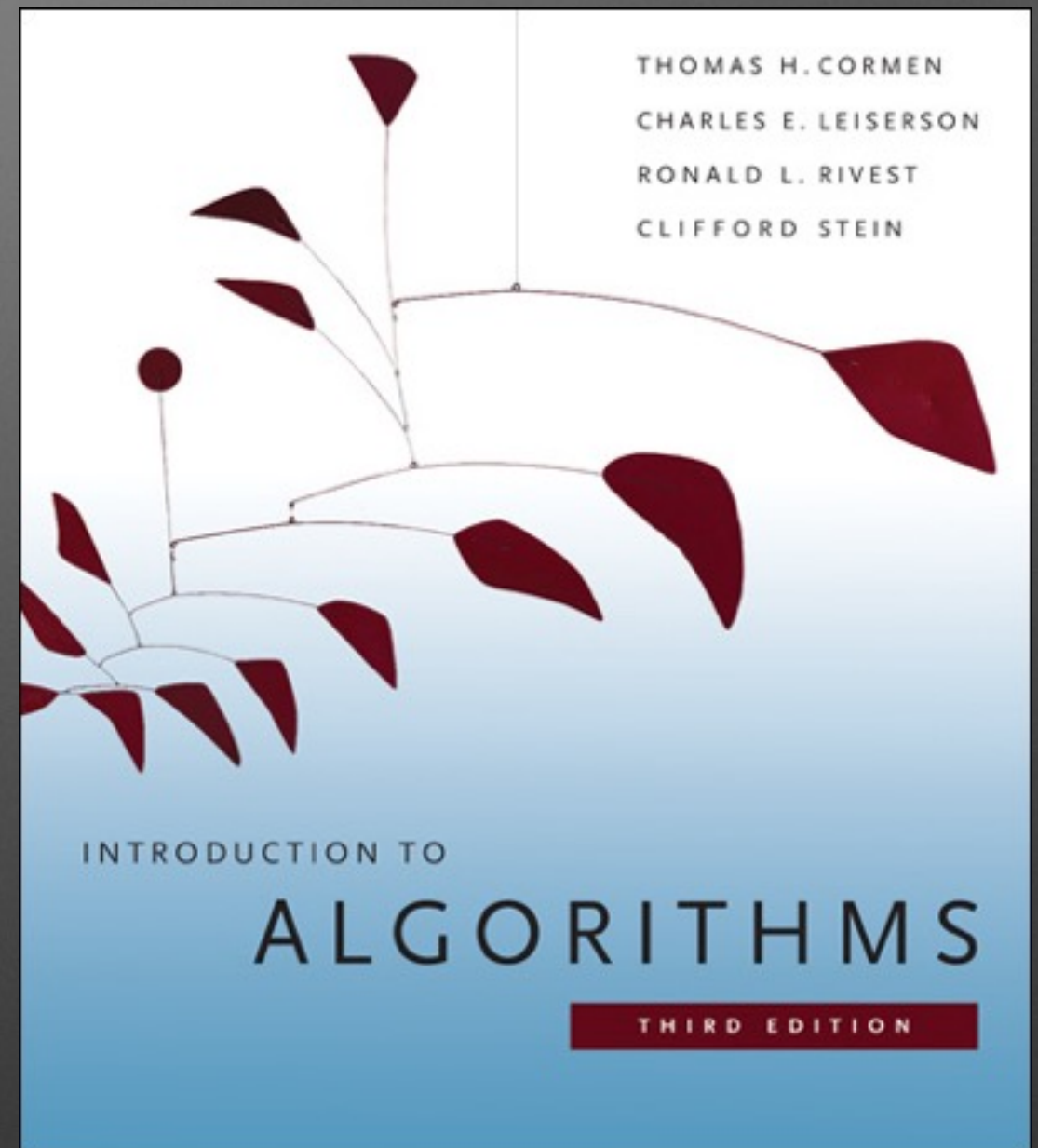
Algorithms

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Chapter 0

Reference

聖經



Language

- C/C++
- Pseudo code



Plan

- Artificial Intelligence
- Machine Learning



Chapter 1

The Role of Algorithms in Computing

What are algorithms?

- Computational problem
- Input/Output
- A procedure is designed for achieving that input/output relation

**What kinds of problems are
solved by algorithms?**

Search engine

- Input : keyword
- Output : Website list

Traveling salesman problem

- Input : road map, target city, cost
- Output : a path cost is minimum

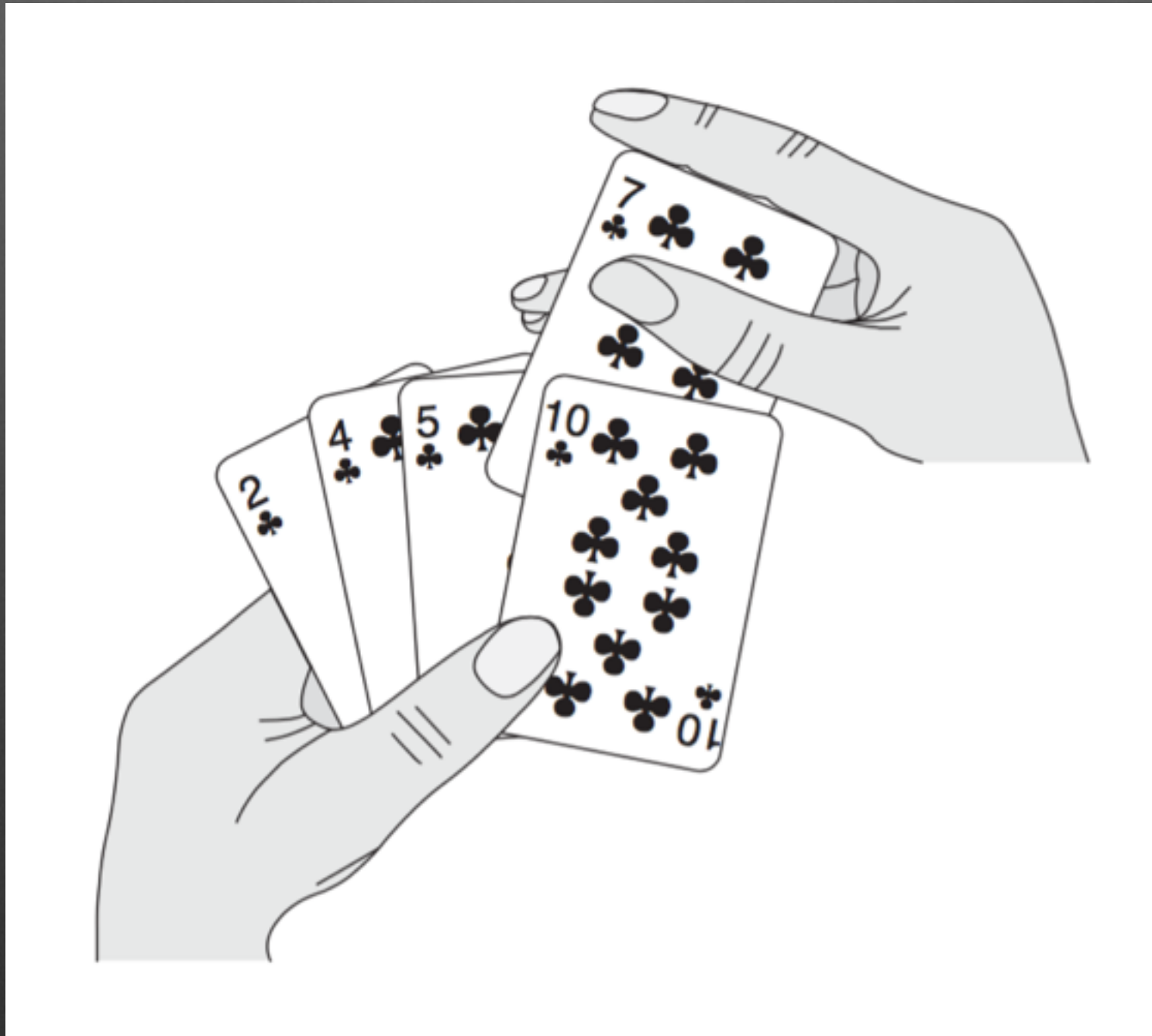
Sorting problem

- Input : A sequence of n numbers $a[1 \dots n]$
- Output : A permutation $a[1 \dots n]$ of the input sequence such $a[1] \leq a[2] \leq \dots \leq a[n]$

Chapter 2

getting started

Insertion sort



Insertion sort

- Insertion-Sort(A)
 - for $j = 2$ to $A.length$
 - key = $A[j]$
 - $i = j - 1$
 - while $i > 0$ and $A[i] > key$
 - $A[i + 1] = A[i]$
 - $i = i - 1$
 - $A[i + 1] = key$

Loop invariants

- Initialization: It is true prior to the first iteration of the loop
- Maintenance: If it is true before an iteration of the loop, it remains true before the next iteration
- Termination: When the loop terminates, the invariant gives us a useful property that helps show that the algorithm is correct

Loop invariants

- Initialization：描述 Invariant Condition 在迴圈執行第一個 iteration前，就成立
- Maintenance：描述 Invariant Condition 在迴圈的任一 iteration執行前跟執行後都維持成立
- Termination：迴圈執行結束後，Invariant Condition 能夠展示整體演算法的正確性

Loop invariants

- subarray $A[1..j-1]$ is sorted
- Initialization : start at $j = 2$, $A[1]$ is sorted
- Maintenance : every loop works by moving $A[j-1]$... until it finds position for $A[j]$
- Termination : terminal at $j = A.length + 1 \Rightarrow A[1..A.length]$ is sorted

Analysis of insertion sort

- Insertion-Sort(A)
 - for $j = 2$ to $A.length$
 - $key = A[j]$
 - $i = j - 1$
 - while $i > 0$ and $A[i] > key$
 - $A[i + 1] = A[i]$
 - $i = i - 1$
 - $A[i + 1] = key$

Chapter 3

Growth of Function

大學以前有教

Chapter 4

Divide-and-Conquer

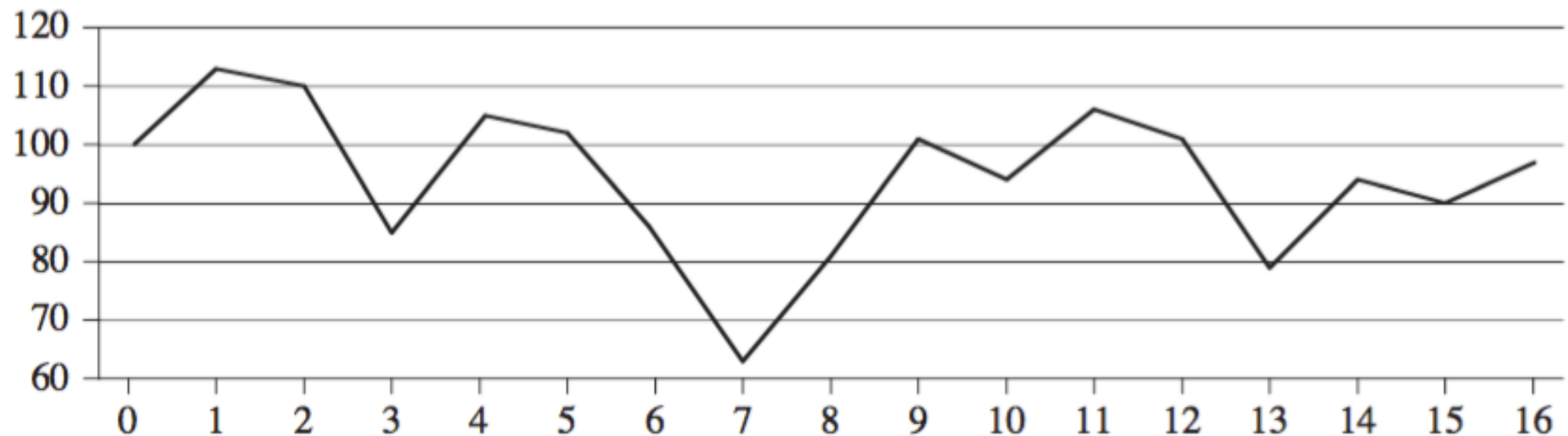
Divide-and-Conquer

- Divide the problem into a number of subproblems that are smaller instances of the same problem.
- Conquer the subproblems by solving them recursively. If the subproblem sizes are small enough, however, just solve the subproblems in a straightforward manner.
- Combine the solutions to the subproblems into the solution for the original problem.

Divide-and-Conquer

- Divide 問題使原本的問題變成多個數量級小一點的相同問題
- Conquer 使用遞迴拆解問題直到問題夠小能直接獲得答案
- Combine 把所有子問題的答案組合起來成為原本問題的答案

The maximum-subarray problem



Day	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
Price	100	113	110	85	105	102	86	63	81	101	94	106	101	79	94	90	97
Change		13	-3	-25	20	-3	-16	-23	18	20	-7	12	-5	-22	15	-4	7

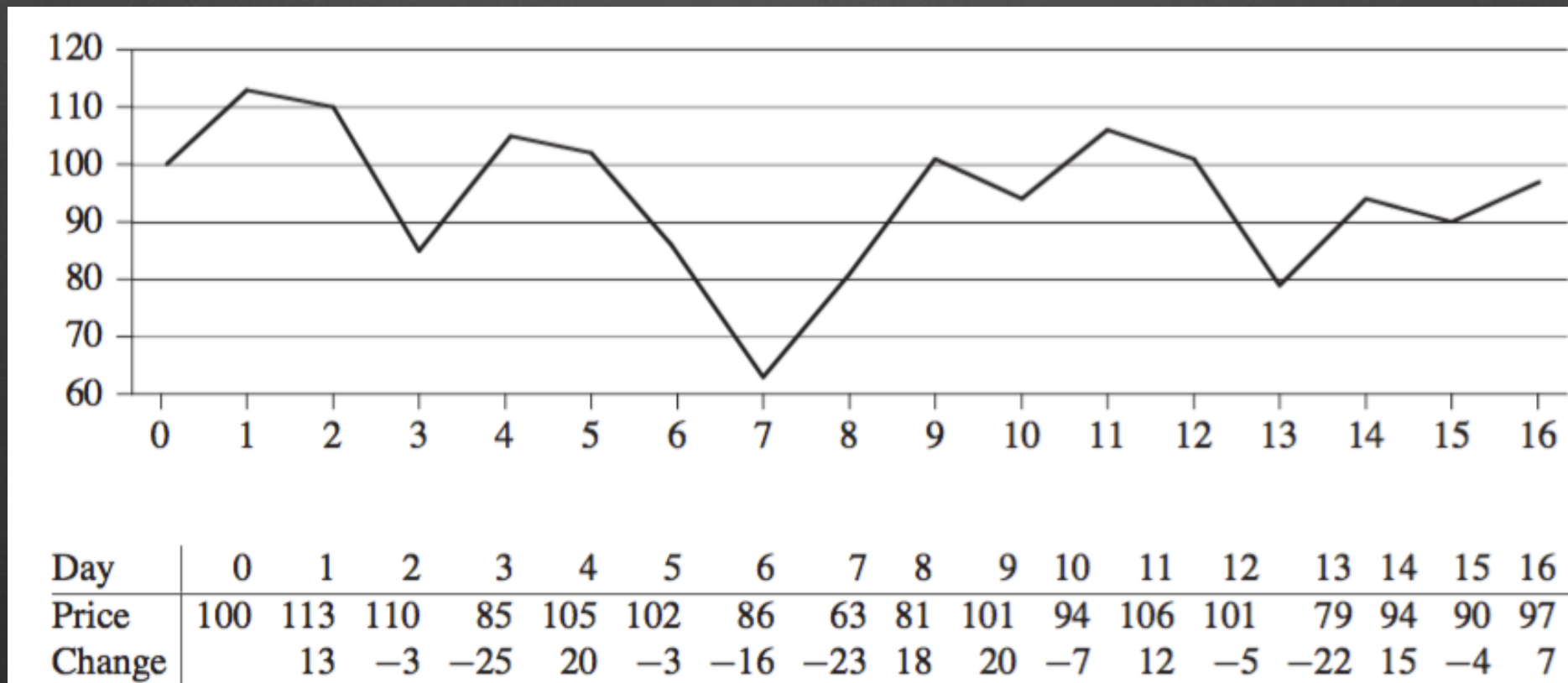
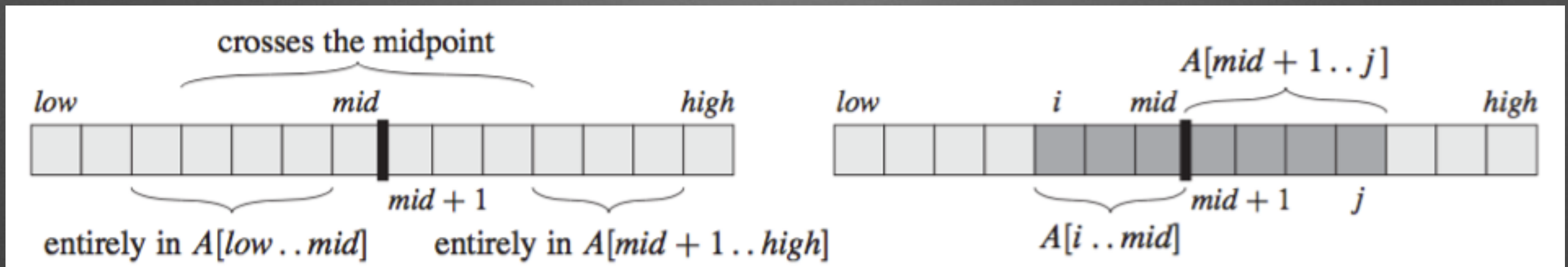
The maximum-subarray problem

```
1  int DAYS = 17;
2  int price[DAYS] = {100,113,110,85,105,102,86,63,81,101,94,106,101,79,94,90};
3  int change[DAYS] = {0};
4  int max =0;
5  int buyDay=0;
6  int sellDay=0;
7  class Point{
8      int max = 0;
9      int buyDay = 0;
10     int sellDay = 0;
11     Point(int max,int buyDay,int sellDay){
12         this.max = max;
13         this.buyDay = buyDay;
14         this.sellDay = sellDay;
15     }
16
17 }
```

The maximum-subarray problem (Brute-force)

```
18 Point brute-force-solution(){
19     Point ans(0,0,0);
20     for(int i = 0 ; i < DAYS-1 ; i++){
21         for(int j =i+1 ; j < DAYS ; j++){
22             int temp = price[j]-price[i];
23             if(ans.max < temp){
24                 ans.max = temp;
25                 ans.buyDay = i;
26                 ans.sellDay = j;
27             }
28         }
29     }
30     return ans;
31 }
```


The maximum-subarray problem (divide-and-conquer)



The maximum-subarray problem (divide-and-conquer)

```
32 void init(){
33     for(int i=1;i<DAYS;i++){
34         change[i] = price[i]-price[i-1];
35     }
36 }
37 Point divide(){
38     init();
39     return findMaxSubArray(0,DAYS-1);
40 }
```

The maximum-subarray problem (divide-and-conquer)

```
41 Point findMaxCrossingSubArray(int low, int mid, int high){
42     int leftSum = INT_MIN, maxLeft = mid;
43     int sum = 0;
44     Point a;
45     for(int i = mid ; i >= low ; i--){
46         sum += change[i];
47         if(sum > leftSum){
48             leftSum = sum;
49             a.buyDay = i;
50         }
51     }
52     int rightSum = INT_MIN, maxRight = mid;
53     sum = 0;
54     for(int i = mid+1 ; i < high ; i++){
55         sum += change[i];
56         if(sum > rightSum){
57             rightSum = sum;
58             a.sellDay = i;
59         }
60     }
61     a.max = rightSum + leftSum;
62     return a;
63 }
```


The maximum-subarray problem (divide-and-conquer)

```
Point findMaxSubArray(int low, int high){
    if(low == high) return new Point(change[low], low, high);
    int mid = (low + high)/2;
    Point leftP = findMaxSubArray(low, mid);
    Point rightP = findMaxSubArray(mid+1, high);
    Point midP = findMaxCrossingSubArray(low, mid, high);
    return leftP.max >= midP.max ? (leftP.max >= rightP.max ? leftP : rightP) : (rightP.max >= midP.max ? rightP : midP);
}
```

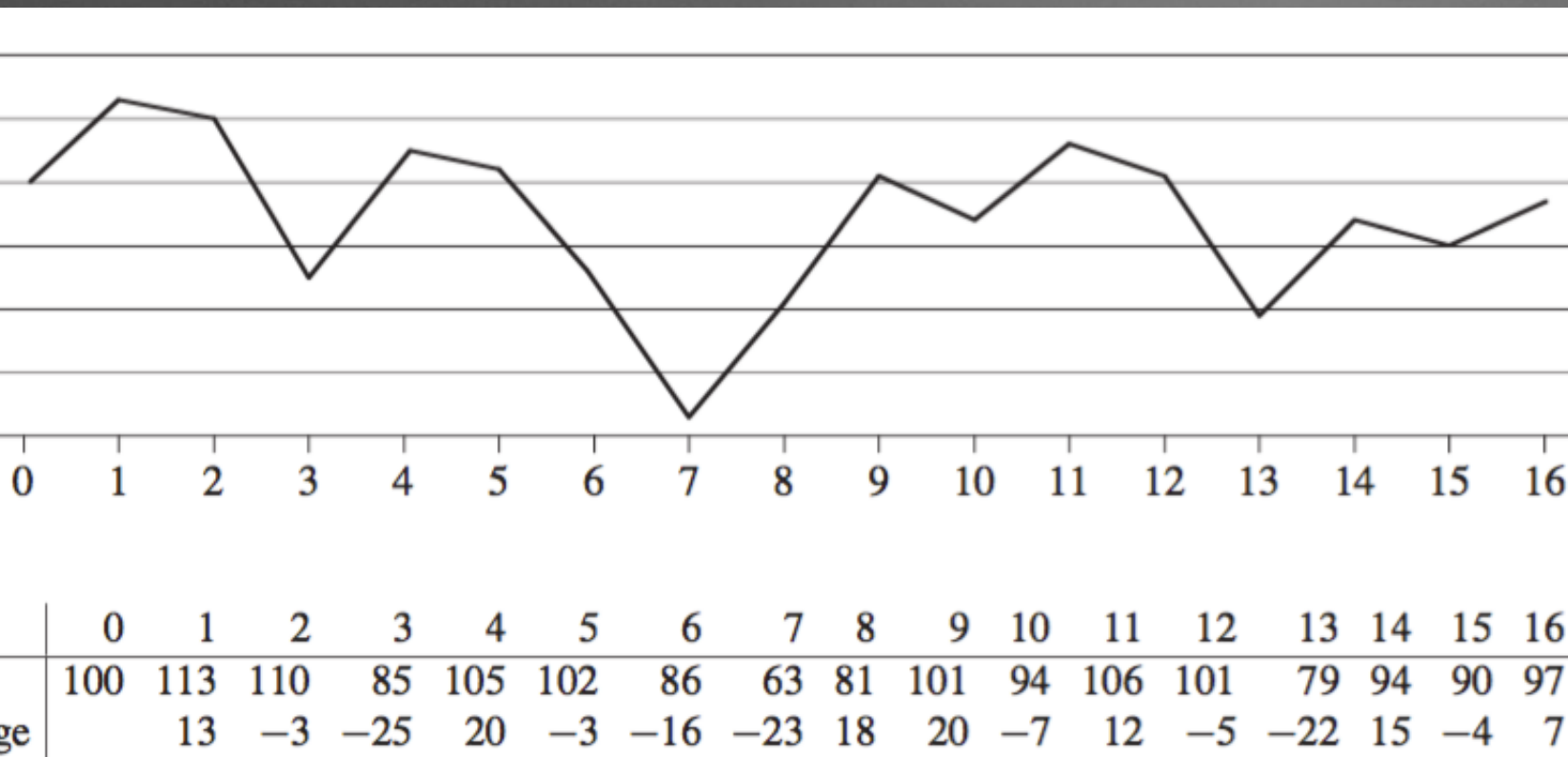

The maximum-subarray problem (divide-and-conquer)

```
32 void init(){
33     for(int i=1;i<DAYS;i++){
34         change[i] = price[i]-price[i-1];
35     }
36 }
37 Point divide(){
38     init();
39     return findMaxSubArray(0,DAYS-1);
40 }
```

Analyzing

- $T(1) = O(1)$
- $T(n) = 2T(n/2) + O(n) = O(n \lg n)$

The maximum-subarray problem (dynamic programming)



• 0,13,10,0,20,17,1,0,18,38,31,43,38,16,31,27,34

```
Point findMaxPointUsingDP(){
    int temp[DAYS] = {0};
    Point ans(0,0,0);
    for(int i=1; i<DAYS; i++){
        temp[i] += change[i];
        if(temp[i] == 0) temp[i] = 0;
        if(temp[i] > max){
            a.max = temp[i];
            a.sell = i;
        }
    }
    for(int i = sell; i >= 0; i--){
        if(temp[i] == 0){
            a.buy = i;
            break;
        }
    }
    return ans;
}
```

Analyzing

- Time $O(n)$
- Space $O(n)$

312. Burst Balloons

- Input : nums[]
- output : maxCoins
- $\text{nums} = [3, 1, 5, 8] \rightarrow [3, 5, 8] \rightarrow [3, 8] \rightarrow [8] \rightarrow []$
- $\text{coins} = 3 * 1 * 5 + 3 * 5 * 8 + 1 * 3 * 8 + 1 * 8 * 1 = 167$

312. Burst Balloons (divide-and-conquer)

```
int maxCoins(int left,int right,vector<int>& nums){  
    if(right - left == 1){  
        return 0;  
    }  
    int max = 0;  
    for(int mid = left+1 ;mid < right; mid++){  
        int temp = nums[mid]*nums[left] * nums[right] +  
                    maxCoins(left,mid,nums) +  
                    maxCoins(mid,right,nums);  
        if(max < temp ){  
            max = temp;  
        }  
    }  
    return max;  
}
```

Analyzing

```
int maxCoins(int left,int right,vector<int>& nums){
    if(right - left == 1){
        return 0;
    }
    int max =0;
    for(int mid = left+1 ;mid < right; mid++){
        int temp = nums[mid]*nums[left] * nums[right] +
                    maxCoins(left,mid,nums) +
                    maxCoins(mid,right,nums);
        if(max < temp ){
            max = temp;
        }
    }
    return max;
}
```

- Time $T(n) = n * 2 * (T(1) + T(2) + \dots T(n-1)) + O(1)$
 $T(n-1) = (n-1) * 2 * (T(1) + \dots + T(n-2)) + O(1)$
 $T(n) - T(n-1) \rightarrow T(n) = S + (n-1) * (T(n-1))$
 $n * T(n) = n * S + n * (n-1) * (T(n-1)) , T(n) = n * S$
 $T(n) = n * (n-1) * T(n-1) / (n-1) = n * T(n-1)$
 $T(n) = n!$

312. Burst Balloons

(dynamic programming)

```
int maxCoins(int left,int right,vector<int>& nums,int* dp,int size){
    if(dp[left + right * size]>0)return dp[left + right * size];
    if(right - left == 1){
        return 0;
    }
    int max =0;
    for(int mid = left+1 ;mid < right; mid++){
        int temp = nums[mid]*nums[left] * nums[right] +
                    maxCoins(left,mid,nums,dp,size) +
                    maxCoins(mid,right,nums,dp,size);
        if(max <temp ){
            max = temp;
        }
    }
    dp[left + right * size] = max;
    return max;
}
```


Analyzing

```
int maxCoins(int left,int right,vector<int>& nums,int* dp,int size){
    if(dp[left + right * size]>0)return dp[left + right * size];
    if(right - left == 1){
        return 0;
    }
    int max =0;
    for(int mid = left+1 ;mid < right; mid++){
        int temp = nums[mid]*nums[left] * nums[right] +
                    maxCoins(left,mid,nums,dp,size) +
                    maxCoins(mid,right,nums,dp,size);
        if(max <temp ){
            max = temp;
        }
    }
    dp[left + right * size] = max;
    return max;
}
```

- Time $T(n) = O(n*n) + O(n*n) + \dots + O(n*n)$
 $= O(n^3)$
- Space $O(n^2)$

- Chapter 1 : What algorithm is
- Chapter 2 : Sample
- Chapter 3 : 大學以前有教
- Chapter 4 : Divide-and-Conquer

Q & A