
FOUR PROOFS OF EUCLID'S THEOREM

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ABSTRACT

Many scientists are familiar with Euclid's theorem of the infinitude of primes. But, some proofs are more insightful than others. In this article, I shall expand upon four different proofs which emphasise different aspects of the infinitude of primes.

1 Euler's proof of the infinitude of primes

A proof by Euler relies on the fundamental theorem of arithmetic, that each integer has a unique prime factorisation. His proof is as follows:

$$\prod_{p \in \mathbb{P}} \frac{1}{1 - \frac{1}{p}} = \prod_{p \in \mathbb{P}} \sum_{k \geq 0} \frac{1}{p^k} = \sum_n \frac{1}{n} \quad (1)$$

Now, since the Harmonic series diverges and $\frac{p}{p-1} \leq 2$ we must conclude that there are infinitely many primes.

2 Erdős' proof

Erdős' insight is that if there were finitely many primes there would not be enough to generate infinitely many integers. To be precise, let's suppose the set of primes \mathbb{P} is finite so we have $\mathbb{P} = p_{i=1}^k$ and each integer N may be expressed as:

$$N = \prod_{i=1}^K p_i^{\alpha_i} \quad (2)$$

and as α_i is either odd or even, we may express N as:

$$N = \left(\prod_{i=1}^K p_i^{e_i} \right) \cdot s^2 \quad (3)$$

where $e_i \in \{0, 1\}$ so $\prod_{i=1}^K p_i^{e_i}$ is square-free.

Now, since $s \leq \sqrt{N}$ and there are only 2^k possible products $\prod_{i=1}^K p_i^{e_i}$ we may conclude that:

$$\forall n \leq N, n \in [1, 2^k \cdot \sqrt{N}] \implies N \leq 2^k \cdot \sqrt{N} \quad (4)$$

which brings us to a contradiction as there are infinitely many integers. Although this concludes the proof, it is worth adding that $\prod_{i=1}^K p_i^{e_i}$ may express at most k bits of information which suggests the possibility of an information-theoretic proof.

3 Chaitin's proof of the infinitude of primes

3.1 Lemma: Almost all integers are incompressible

Given that \mathbb{N} is countable and the space of binary strings of finite length $\{0, 1\}^*$ is also countable, we may construct a bijection from \mathbb{N} to $\{0, 1\}^*$. It follows that every positive integer has a unique binary encoding.

Moreover, almost all positive integers are incompressible since:

$$\forall n \in \mathbb{N}^* \forall k < n, |x \in \{0, 1\}^* : K(x) \leq n - k| \geq 2^n(1 - 2^{-k}) \quad (5)$$

where $|x| = n$, the binary length of x , which may be understood as the machine-code representation of an integer.

Proof:

Let's suppose an integer with binary encoding x has an algorithmic complexity $K(x) \leq n - k$. Given that the number of binary strings of binary length less than $n - k$ is $2^{n-k} - 1 < 2^{n-k}$ we have:

$$2^n - 2^{n-k} = 2^n(1 - 2^{-k}) \quad (6)$$

integers with an algorithmic complexity greater than or equal to $n - k$. As an immediate consequence, for $n > k \geq 10$, more than 99.9% of integers have an algorithmic complexity greater than $n - k$ so less than 1% of integers are compressible.

3.2 Chaitin's proof of Euclid's theorem

Each integer $n \in \mathbb{N}$ may be described by its prime factorisation:

$$n = \prod_{i=1}^k p_i^{\alpha_i} \quad (7)$$

so $\pi(n) = k$ assuming that some exponents α_i equal zero.

Given that $p_i \geq 2$,

$$\forall i \in [1, k], \alpha_i \leq \log_2 n \quad (8)$$

and so any exponent may be described using $\log_2 \log_2 n$ bits.

Now, assuming that the value of $\log_2 \log_2 n$ is known the integer n may be described using:

$$k \cdot \log_2 \log_2 n \quad (9)$$

bits.

However, given that most integers are incompressible there are integers of binary length $l = \log_2 n$ which can't be described in fewer than l bits. So we may deduce that:

$$\pi(n) \cdot \log_2 \log_2 n \geq \log_2 n \quad (10)$$

for almost all positive integers $n \in \mathbb{N}$.

This allows us to deduce a useful lower bound on the prime counting function for almost all n :

$$\pi(n) \geq \frac{\log_2 n}{\log_2 \log_2 n} \quad (11)$$

which implies that there are infinitely many prime numbers.

4 Euclid's theorem via the irrationality of Archimedes' constant

If we let $\alpha = -x^2$ where $|\alpha| < 1$,

$$\sum_{n=0}^{\infty} \alpha^n = \frac{1}{1-\alpha} \quad (12)$$

Given that $\arctan(x) = \int_0^x \frac{dx}{1+x^2}$, if we combine the integral form of $\arctan(x)$ with (10) we have:

$$\arctan(x) = \int_0^x \frac{dx}{1+x^2} = \sum_{k=0}^{\infty} (-1)^k \int_0^x x^{2k} dx = \sum_{k=0}^{\infty} \frac{(-1)^k \cdot x^{2k+1}}{2k+1} \quad (13)$$

so for $x = 1$,

$$\frac{\pi}{4} = \arctan(1) = 1 - \frac{1}{3} + \frac{1}{5} - \frac{1}{7} + \dots \quad (14)$$

Upon closer inspection, (11) simplifies to:

$$\frac{\pi}{4} = \sum_{n=0}^{\infty} \frac{\chi(n)}{n} \quad (15)$$

where $\chi(n) = 0$ if $n \bmod 2 = 0$, $\chi(n) = 1$ if $n \bmod 4 = 1$ and $\chi(n) = -1$ if $n \bmod 4 = 3$.

Now, we may observe:

$$\sum_{n=0}^{\infty} -\frac{1}{p^n} = \frac{1}{1 - (-\frac{1}{p})} = \frac{1}{1 + \frac{1}{p}} = \frac{p}{p+1} \quad (16)$$

Finally, since $\chi(p=2) = 0$ and for $p > 2$, $p \bmod 4$ is either 1 or 3 we may define:

$$P = \{p \in \mathbb{P} : p \equiv 1 \pmod{4}\} \quad (17)$$

$$P' = \{p \in \mathbb{P} : p \equiv 3 \pmod{4}\} \quad (18)$$

so (13) may be expressed as:

$$\sum_{n=0}^{\infty} \frac{\chi(n)}{n} = \left(\prod_{p \in P} \frac{p}{p-1} \right) \cdot \left(\prod_{p \in P'} \frac{p}{p+1} \right) \quad (19)$$

and therefore:

$$\frac{\pi}{4} = \frac{3}{4} \cdot \frac{5}{4} \cdot \frac{7}{8} \cdot \frac{11}{12} \dots \quad (20)$$

where each numerator is an odd prime and each denominator is the nearest multiple of 4 with respect to the numerator. It follows that if there were only finitely many primes, π would be a rational number. QED.

References

References:

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