

Fenwick Tree

(binary indexed tree)

The Problem

- There are several boxes
 - Labeled from 1 to N
- We can
 - Add N marble(s) into i^{th} box
 - We say box $\#i$ has frequency N
- We want to know
 - Total number of marbles in box $\#1$ to $\#j$

Fenwick Tree

- Operation

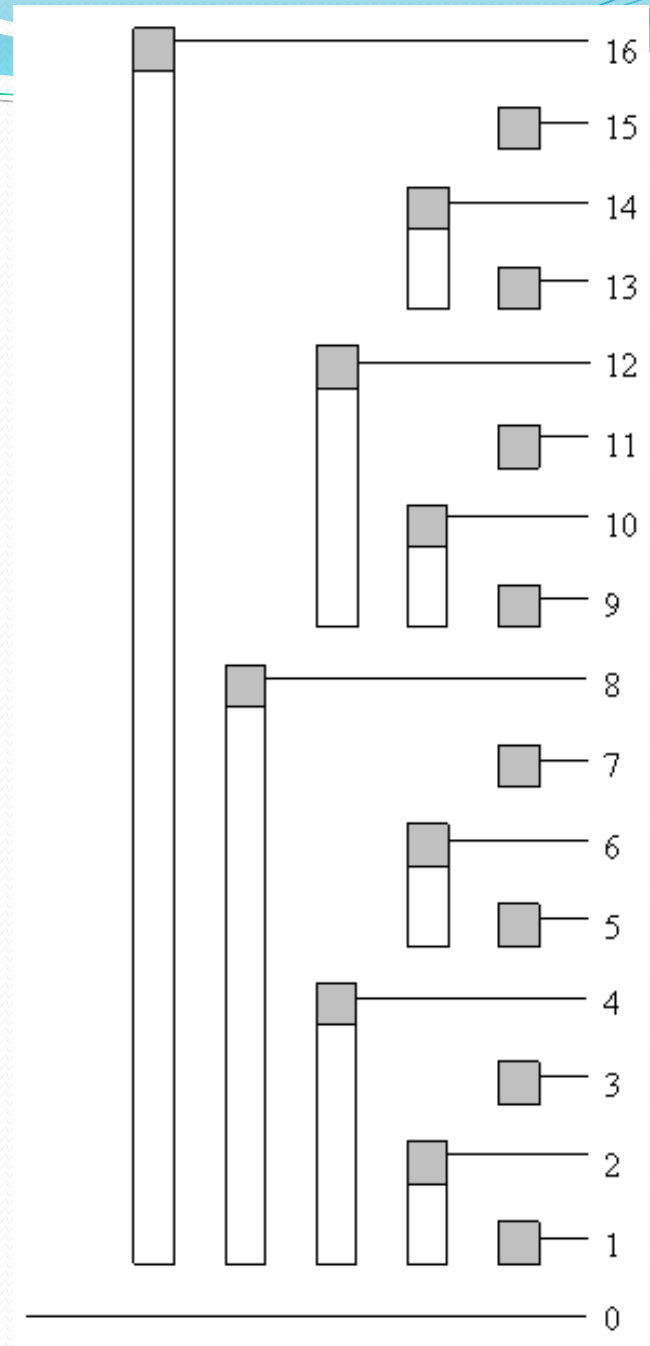
- `void create(int n);` $O(N)$
- `void update(int idx, int val);` $O(\log N)$
- `int freqTo(int idx);` $O(\log N)$
- `int freqAt(int idx);` $O(\log N)$

Storage

- Data
 - An int array of size N

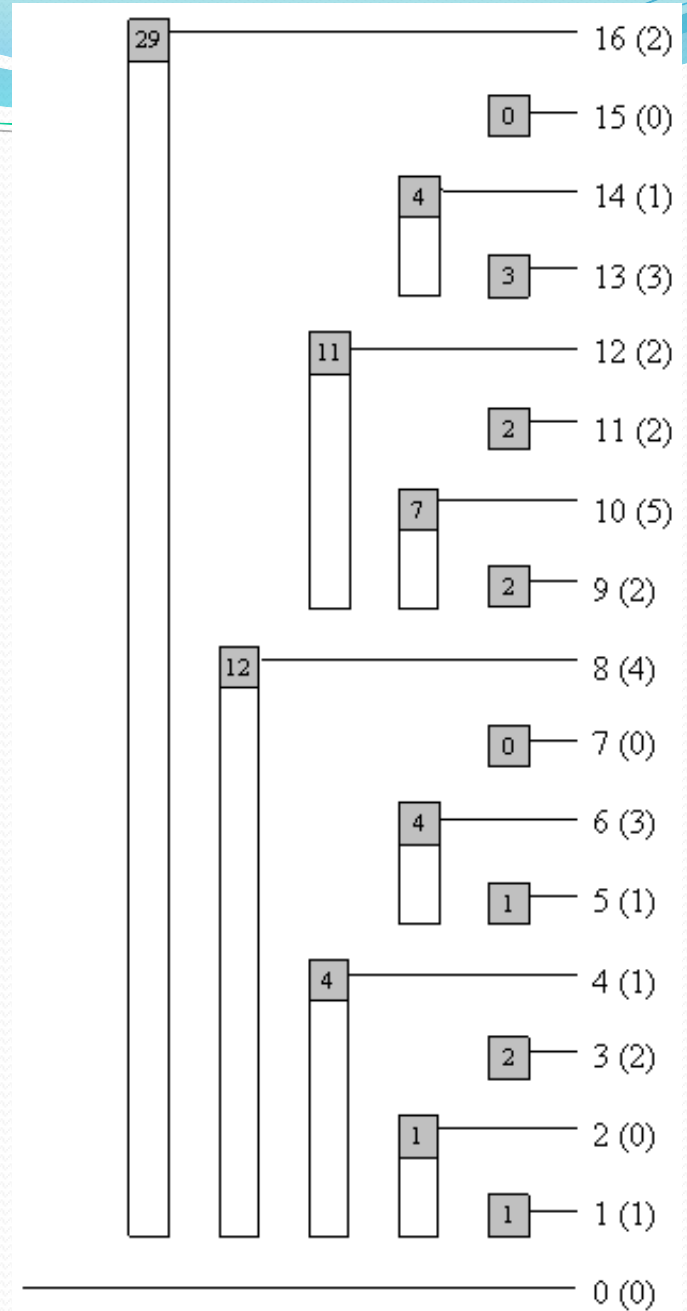
Fenwick Tree

- How it works?
 - Each element in the array stores cumulative frequency of consecutive list of boxes
 - Range of boxes that is stored is related to “binary value” of the index



Define

- $f(x)$ = number of marble in box x
- $c(x)$ = summation of number of marble in box #1 to box # x
- $\text{tree}[x]$ = element x in the array



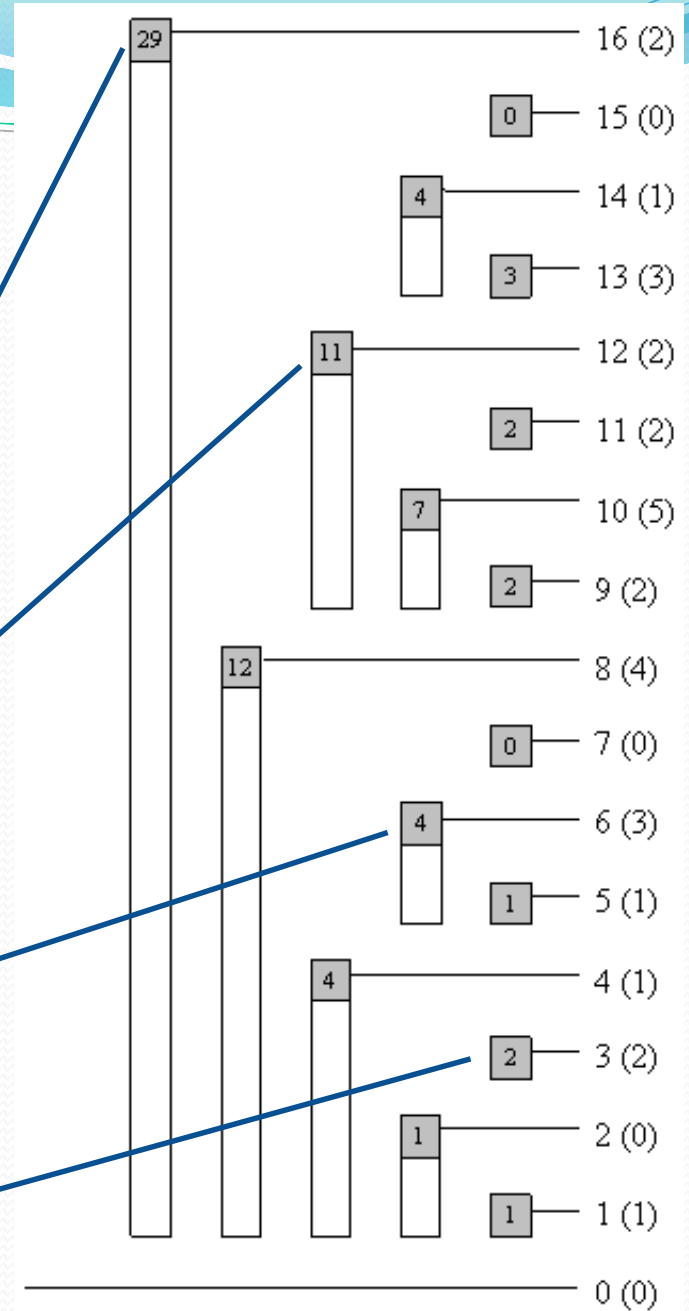
Storage Solution

$$\text{Tree}[16] = f(1) + f(2) + \dots + f(16)$$

$$\text{Tree}[12] = f(9) + f(10) + \dots + f(12)$$

$$\text{Tree}[6] = f(5) + f(6)$$

$$\text{Tree}[3] = f(3)$$



$$f(16) = 2$$

Cumulative Freq

Actual
frequency

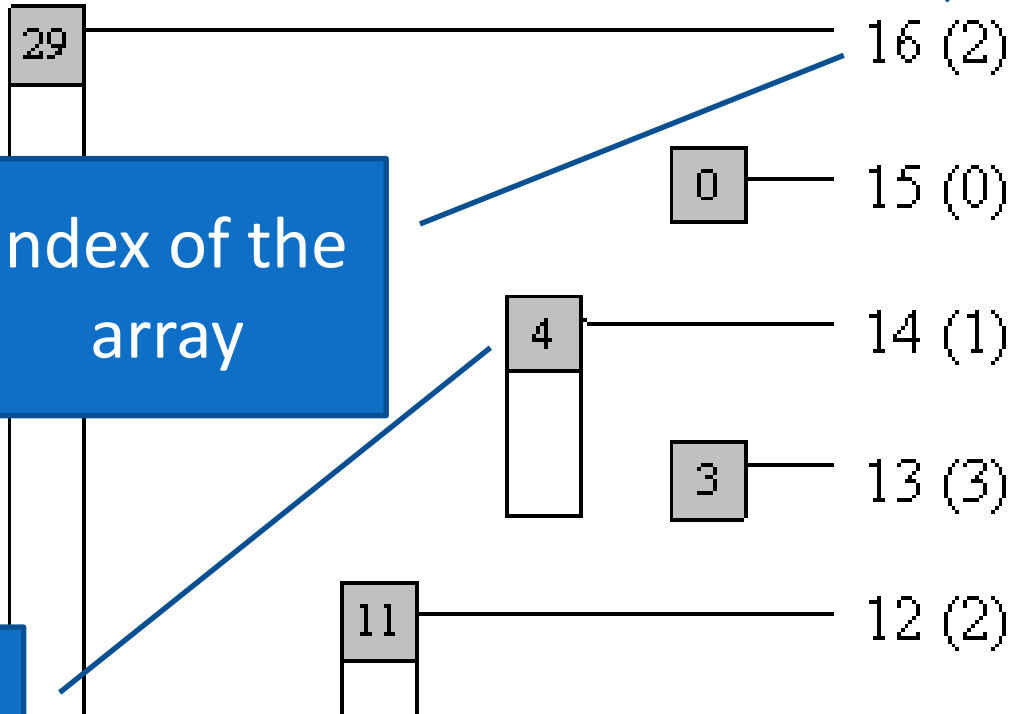
$\text{tree}[16] = 29$

Cumulative
frequency
From 1 to 16

$\text{tree}[14]$

Cumulative
frequency
From 13 to 14

Index of the
array



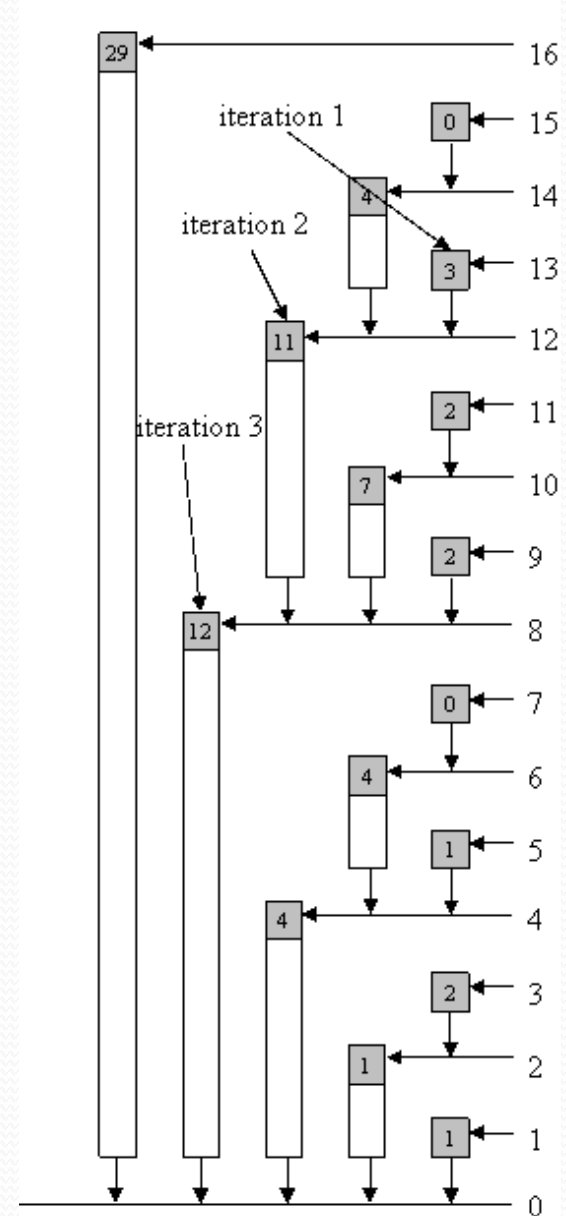
The last 1

- A node at the index X will store freq of boxes in the range
 - $X - 2^r + 1$ to X
 - Where r is the position of the last digit of 1
- Ex
 - $X = 12$ $(1100)_2$
 - Node will store freq from 9 to 12
 - The last 1 of 12 is at position 2 (0-indexed)
 - $12 - 2^2 + 1 = 9 = (1001)_2$

Read Cumulative Freq

$c(13) =$
 $tree[13] +$
 $tree[12] +$
 $tree[8]$

In base-2
 $c(1101_2) =$
 $tree[1101_2] +$
 $tree[1100_2] +$
 $tree[1000_2]$



Update Freq

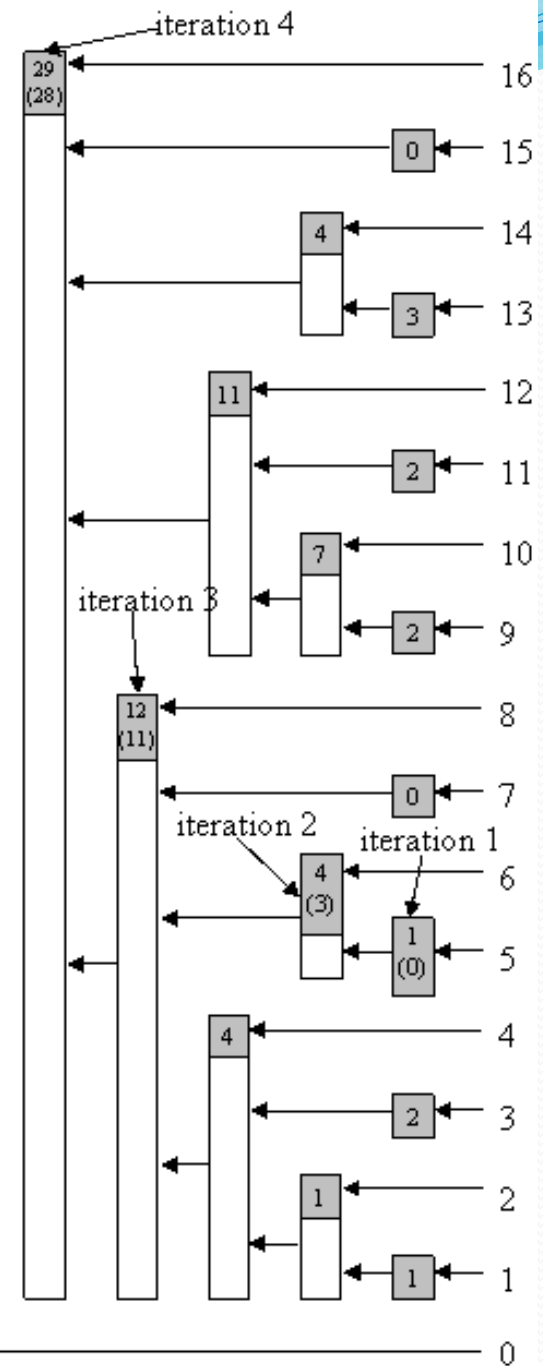
Update $f(5)$ by -1
involve

Tree[16] (10000_2)

Tree[8] (01000_2)

Tree[6] (00110_2)

Tree[5] (00101_2)



Read actual Freq

What is $f(12)$?

Easy, it's $c(12) - c(11)$

easier

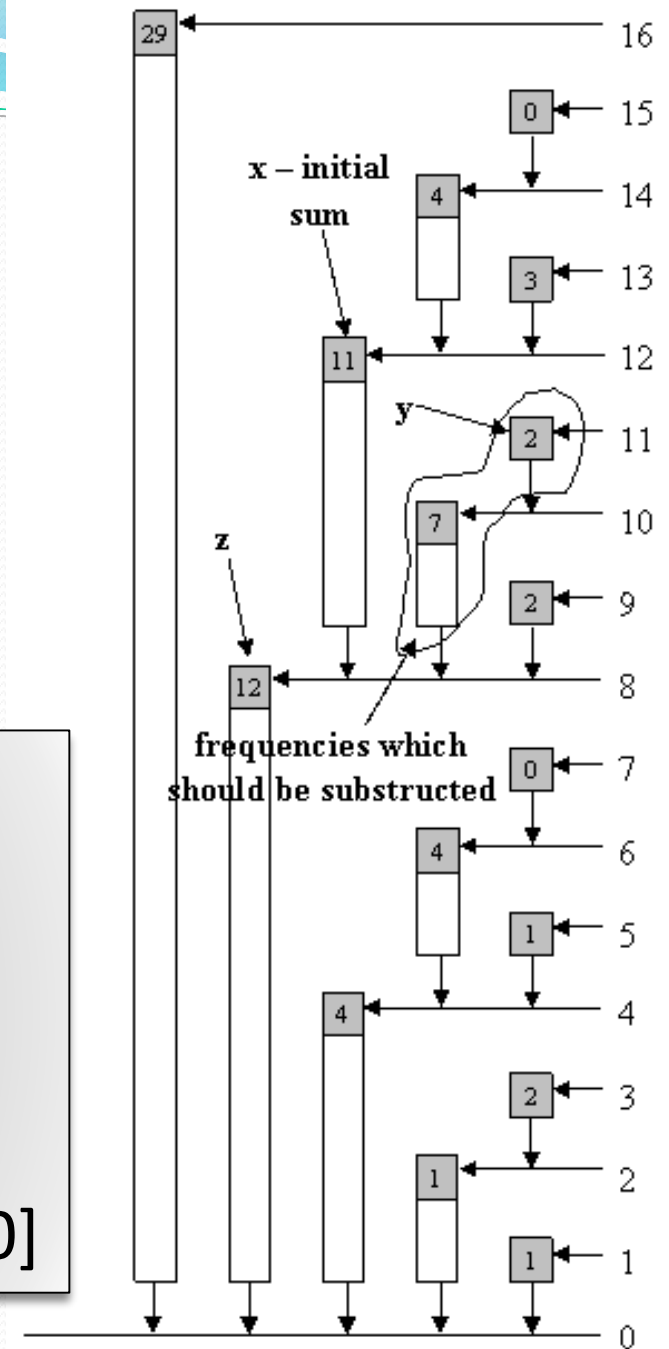
$$\text{Tree}[12] = f(9) + f(10) + f(11) + f(12)$$

$$\text{Tree}[11] = f(11)$$

$$\text{Tree}[10] = f(9) + f(10)$$

Hence,

$$f(12) = \text{Tree}[12] - \text{Tree}[11] - \text{Tree}[10]$$



Two's complement

- A method to represent negative
 - A two's complement of X is
 - (compliment of x) + 1
 - Ex.. 2's Complement of 7 is
 - 0111 → 1000 → 1001
- Finding the last 1
- $x = a1b$
 - b = consecutive of 0
- Ex... $X = 4 = 0100$
 - a = 0 b = 00

0111	7
0110	6
0101	5
0100	4
0011	3
0010	2
0001	1
0000	0
1111	-1
1110	-2
1101	-3
1100	-4
1011	-5
1010	-6
1001	-7
1000	-8

Two's complement

- Now, let's see two's complement more closely
- $-x$
 - $= (a1b)^- + 1$
 - $= a^-0b^- + 1$
 - $= a^-0(0...0)^- + 1$
 - $= a^-0(1...1) + 1$
 - $= a^-1(0...0)$
 - $= a^-1b.$
- So, if we "&" $-x$ and x
 - a^-1b & $a1b.$
 - We got the last 1

0111	7
0110	6
0101	5
0100	4
0011	3
0010	2
0001	1
0000	0
1111	-1
1110	-2
1101	-3
1100	-4
1011	-5
1010	-6
1001	-7
1000	-8

Code

```
int freqTo(int idx) {  
    int sum = 0;  
    while (idx > 0){  
        sum += tree[idx];  
        idx -= (idx & -idx);  
    }  
    return sum;  
}
```

```
void update(int idx ,int val) {  
    while (idx <= MaxVal){  
        tree[idx] += val;  
        idx += (idx & -idx);  
    }  
}
```

Code

```
int freqAt(int idx){
    int sum = tree[idx];
    if (idx > 0) {
        int z = idx - (idx & -idx);
        y = idx - 1;
        while (y != z){
            sum -= tree[y];
            y -= (y & -y);
        }
    }
    return sum;
}
```


2D BIT

- Box is arranged at x-y coordinate
- Operation
 - $\text{Update}(x, y, \text{val})$ (add “val” marble in position (x, y))
 - How many points in the range (x_1, y_1) to (x_2, y_2)

2D BIT

