



Worksheet 5 (Solved)

HoTTEST Summer School 2022

The HoTTEST TAs

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1 (★)

Consider a function $f : A \rightarrow B$. Recall that a *retraction* of f is a function $g : B \rightarrow A$ such that $g \circ f \sim \text{id}_A$. Construct a function

$$\text{retr}(f) \rightarrow \left(\prod_{a, a' : A} f(a) = f(a') \rightarrow a = a' \right).$$

This means that if f has a retraction, then it is an injection.

Let $\gamma : \text{retr}(f)$. By Σ -induction, we may take $\gamma \doteq \langle g, r \rangle$, where $g : B \rightarrow A$ and $r : g \circ f \sim \text{id}_A$. Let $a, a' : A$ and $p : f(a) = f(a')$. We must find a term of type $a = a'$. By the action of g on p , we have that $g(f(a)) = g(f(a'))$. We also have that

$$\begin{array}{ll} a = g(f(a)) & (r(a)^{-1}) \\ g(f(a')) = a'. & (r(a')) \end{array}$$

Finally, by transitivity of equality, we have an identity

$$r(a)^{-1} \cdot \text{ap}_g(p) \cdot r(a') : a = a'.$$

2 **(★★)**

Consider a commuting triangle

$$\begin{array}{ccc} A & \xrightarrow{h} & B \\ & \searrow f & \swarrow g \\ & X & \end{array} .$$

1. Suppose that h has a section $s : B \rightarrow A$. Prove that $f \circ s \sim g$ and that $\mathbf{sec}(f) \leftrightarrow \mathbf{sec}(g)$.
2. Suppose that g has a retraction $r : X \rightarrow B$. Prove that $r \circ f \sim h$ and that $\mathbf{retr}(f) \leftrightarrow \mathbf{retr}(h)$.
3. Prove that if any two of f , g , and h are equivalences, then so is the third.
4. Prove that any retraction or section of an equivalence is itself an equivalence.

1. First, note that

$$f(s(b)) = g(h(s(b))) = g(b)$$

for all $b : B$. This means that $f \circ s \sim g$.

Next, suppose that f has a section $s_f : X \rightarrow A$. Then

$$g(h(s_f(x))) = f(s_f(x)) = x$$

for all $x : X$. Hence $h \circ s_f$ is a section of g .

Conversely, suppose that g has a section $s_g : X \rightarrow B$. Then

$$f(s(s_g(x))) = g(s_g(x)) = x$$

for all $x : X$. Hence $s \circ s_g$ is a section of f .

2. First, note that

$$r(f(a)) = r(g(h(a))) = h(a)$$

for all $a : A$. This means that $r \circ f \sim h$.

Next, suppose that f has a retraction $r_f : X \rightarrow A$. Then

$$r_f(g(h(a))) = r_f(f(a)) = a$$

for all $a : A$. Hence $r_f \circ g$ is a retraction of h .

Conversely, suppose that h has a retraction $r_h : B \rightarrow A$. Then

$$r_h(r(f(a))) = r_h(h(a)) = a$$

for all $a : A$. Hence $r_h \circ r$ is a retraction of f .

3. We have three cases to consider. For any equivalence e , let r_e and s_e denote its retraction and section, respectively.

- Suppose that h and g are equivalences. Then f has a section by part (1) and a retraction by part (2). Thus, it's an equivalence.
- Suppose that f and h are equivalences. Part (1) implies that g has a section. Moreover, since $f \circ s_h \sim g$ by part (1),

$$h(r_f(g(b))) = h(r_f(f(s_h(b)))) = h(s_h(b)) = b$$

for all $b : B$. Thus, $h \circ r_f$ is a retraction of g , so that g is an equivalence.

- Suppose that f and g are equivalences. Part (2) implies that h has a retraction. Moreover, since $r_g \circ f \sim h$ by part (2),

$$h(s_f(g(b))) = r_g(f(s_f(g(b)))) = r_g(g(b)) = b$$

for all $b : B$. Thus, $s_f \circ g$ is a section of h , so that h is an equivalence.

4. Notice a retraction r_e and a section s_e of an equivalence $e : N \rightarrow M$ are commuting triangles $\text{id}_N \sim r_e \circ e$ and $\text{id}_M \sim e \circ s_e$, respectively.

3 (★★)

Consider the type `Bool`, generated by

`true` : `Bool`
`false` : `Bool`.

Define the type family `Eq-bool` : `Bool` \rightarrow `Bool` \rightarrow \mathcal{U}_0 by

`Eq-bool`(`true`, `true`) := $\mathbb{1}$
`Eq-bool`(`true`, `false`) := \emptyset
`Eq-bool`(`false`, `false`) := $\mathbb{1}$
`Eq-bool`(`false`, `true`) := \emptyset .

For every $b, b' : \text{Bool}$, define $\varphi_{b,b'} : (b = b') \rightarrow \text{Eq-bool}(b, b')$ by path induction. Prove that $\varphi_{b,b'}$ is an equivalence.

Let's construct an inverse of $\varphi_{b,b'}$ by double induction on `Bool`:

$\psi_{\text{true}, \text{true}} : \mathbb{1} \rightarrow (\text{true} = \text{true})$
 $\psi(*) := \text{refl}_{\text{true}}$

$\psi_{\text{true}, \text{false}} : \emptyset \rightarrow (\text{true} = \text{false})$
 $\psi(x) := \emptyset\text{-ind}(x)$

$\psi_{\text{false}, \text{false}} : \mathbb{1} \rightarrow (\text{false} = \text{false})$
 $\psi(*) := \text{refl}_{\text{false}}$

$\psi_{\text{false}, \text{true}} : \emptyset \rightarrow (\text{false} = \text{true})$
 $\psi(x) := \emptyset\text{-ind}(x)$.

By double induction on `Bool`, it is straightforward to check that $\psi_{b,b'}$ is an inverse of $\varphi_{b,b'}$.

It is easy to show that $\neg(\mathbb{1} = \emptyset)$. As a consequence, we can prove that $\neg(b = \text{neg-bool}(b))$ for every $b : \text{Bool}$.

4 (★★)

Prove that for all $b : \text{Bool}$,

$$\neg \text{is-equiv}(\text{const}_b).$$

Also, prove that

$$\text{Bool} \not\cong \mathbb{1}.$$

To prove that $\neg \text{is-equiv}(\text{const}_b)$, suppose that const_b is an equivalence. In particular, it has a retraction, hence is injective by Problem 1. Since

$$\text{const}_b(\text{true}) \doteq b \doteq \text{const}_b(\text{false}),$$

it follows that $\text{true} = \text{false}$. But we know that

$$\text{true} \neq \text{neg-bool}(\text{true}),$$

where $\text{neg-bool}(\text{true}) \doteq \text{false}$. This gives us an element of \emptyset , as required.

Likewise, to prove that $\text{Bool} \not\cong \mathbb{1}$, suppose that we have an equivalence $e : \text{Bool} \rightarrow \mathbb{1}$ with retraction $r_e : \mathbb{1} \rightarrow \text{Bool}$. By induction on $\mathbb{1}$, it's easy to check that $r_e(x) = r_e(*)$ for all $x : \mathbb{1}$. Then

$$\text{true} = r_e(e(\text{true})) = r_e(*) = r_e(e(\text{false})) = \text{false}.$$

Again, this gives us an element of \emptyset .

5 **(★)**

Let A be a type and B be a type family over A . For each $x, y : A$, construct an inverse of the function

$$\text{inv}_{x,y} : (x = y) \rightarrow (y = x).$$

Further, for each $p : x = y$, construct an inverse of the function

$$\text{tr}_B(p) : B(x) \rightarrow B(y).$$

We may define these inverses by path induction on $x = y$. Specifically, define $\text{inv}_{x,y}^{-1} : (y = x) \rightarrow (x = y)$ by

$$\text{inv}_{x,x}^{-1}(\text{refl}_x) := \text{refl}_x$$

and define $\text{tr}_B(p)^{-1} : B(y) \rightarrow B(x)$ by

$$\text{tr}_B(\text{refl}_x)^{-1} := \text{id}_{B(x)}.$$

Our definition of $\text{tr}_B(p)^{-1}$ gives us a homotopy

$$\text{tr}_B(p)^{-1} \sim \text{tr}_B(p^{-1})$$

of functions $B(y) \rightarrow B(x)$.

6 (★)

Let $f, g : A \rightarrow B$ and $H : f \sim g$. Prove that $\text{is-equiv}(f) \leftrightarrow \text{is-equiv}(g)$.

To define a function $\text{is-equiv}(f) \rightarrow \text{is-equiv}(g)$, let $\langle h, t_h, k, t_k \rangle : \underbrace{\text{sec}(f) \times \text{retr}(f)}_{\text{is-equiv}(f)}$.

Note that for all $b : B$ and $a : A$,

$$\begin{aligned} g(h(b)) &= f(h(b)) & (H(h(b))^{-1}) \\ &= b & (t_h(b)) \end{aligned}$$

$$\begin{aligned} k(g(a)) &= k(f(a)) & (\text{ap}_k(H(a)^{-1})) \\ &= a. & (t_k(a)) \end{aligned}$$

Thus, h is a section of g , and k is a retraction of g . Now we may take

$$\langle h, \lambda b. H(h(b))^{-1} \cdot t_h(b), k, \lambda a. \text{ap}_k(H(a)^{-1}) \cdot t_k(a) \rangle : \underbrace{\text{sec}(g) \times \text{retr}(g)}_{\text{is-equiv}(g)}$$

as our element of $\text{is-equiv}(g)$. The function in the other direction is similar.

7 (★★)

Suppose that $e, e' : A \rightarrow B$ are equivalences and that $H : e \sim e'$. Let s and s' denote the sections of e and e' , respectively. Prove that s and s' are homotopic.

Recall that any section of an equivalence is also a retraction of it. Therefore, we see that

$$s(b) = s'(e'(s(b))) = s'(e(s(b))) = s'(b)$$

for all $b : B$.