

Algorithm Analysis

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Code Complexity

Primitive Operations:

- Addition etc.
 - 1 operation
- Calling a method or returning from a method
 - 1 operation
- Index in an array
 - 1 operation
- Comparison
 - 1 operation

Code Complexity

Loops:

- FOR
 - $(\# \text{ of iterations}) * (\text{running time of statements inside the loop})$
- Nested loops
 - product of loop complexities
- Consecutive Statements
 - The sum of running time of each segment
- if/Else
 - The testing time + Larger of the running time of the cases

Q1 (Big-Oh)

```
bool checkForOddOrEven(int x{
    if (x%2)
        print("Odd");
    else
        print("Even");
}
```

Q2 (Big-Oh)

```
int findMax(int[] x){  
    int max=0;  
    for(int i=0;i<x.length;i++){  
        if(x[i] > max)  
            max = x[i];  
    }  
    return max;  
}
```

Q3 (Big-Oh)

```
bool checkForDuplicates(int[] x){  
    for(int i=0;i<x.length;i++){  
        for(int j=0;j<x.length;j++){  
            if (i == j) continue;  
            if(x[i] == x[j])  
                return true;  
        }  
    }  
    return false;  
}
```

Q4 (Big-Oh)

```
int calculateStrangeSum(int[] x) {  
    int h=x.length;  
    int sum = 0;  
    while(h>0) {  
        sum += x[h];  
        h = h/2; }  
    return sum;  
}
```

Q5 (Big-Oh)

```
int mixed(int[] x){  
    print(x.length);  
    for(int i=0;i<x.length;i++){  
        for(int j=0;j<x.length;j++){  
            if (i == j) continue;  
            if(x[i] == x[j])  
                print("Duplicate");  
        }  
    }  
}
```


Q6

Order the following functions by growth rate:

$$N, \sqrt{N}, N^{1.5}, N^2, N \log N$$

$$N \log \log N, N \log^2 N, N \log(N^2), 2/N, 2^N, 2^{N/2}, 37, N^2 \log N, N^3.$$

Indicate which functions grow at the same rate.

Q7

Suppose $T_1(N) = O(f(N))$ and $T_2(N) = O(f(N))$. Which of the following are true?

- a. $T_1(N) + T_2(N) = O(f(N))$
- b. $T_1(N) - T_2(N) = o(f(N))$
- c. $T_1(N) / T_2(N) = O(1)$
- d. $T_1(N) = O(T_2(N))$

Q8

Give an analysis of the running time (Big-Oh will do).

```
(1) sum = 0;
    for( i = 0; i < n; i++ )
        sum++;
```

```
(2) sum = 0;
    for( i = 0; i < n; i++ )
        for( j = 0; j < n; j++ )
            sum++;
```

```
(3) sum = 0;
    for( i = 0; i < n; i++ )
        for( j = 0; j < n * n; j++ )
            sum++;
```

```
(4) sum = 0;
    for( i = 0; i < n; i++ )
        for( j = 0; j < i; j++ )
            sum++;
```

```
(5) sum = 0;
    for( i = 0; i < n; i++ )
        for( j = 0; j < i * i; j++ )
            for( k = 0; k < j; k++ )
                sum++;
```

```
(6) sum = 0;
    for( i = 1; i < n; i++ )
        for( j = 1; j < i * i; j++ )
            if( j % i == 0 )
                for( k = 0; k < j; k++ )
                    sum++;
```

Q9

Implement the code in Java, and give the running time for several values of N.

```
(1) sum = 0;
    for( i = 0; i < n; i++ )
        sum++;
```

```
(2) sum = 0;
    for( i = 0; i < n; i++ )
        for( j = 0; j < n; j++ )
            sum++;
```

```
(3) sum = 0;
    for( i = 0; i < n; i++ )
        for( j = 0; j < n * n; j++ )
            sum++;
```

```
(4) sum = 0;
    for( i = 0; i < n; i++ )
        for( j = 0; j < i; j++ )
            sum++;
```

```
(5) sum = 0;
    for( i = 0; i < n; i++ )
        for( j = 0; j < i * i; j++ )
            for( k = 0; k < j; k++ )
                sum++;
```

```
(6) sum = 0;
    for( i = 1; i < n; i++ )
        for( j = 1; j < i * i; j++ )
            if( j % i == 0 )
                for( k = 0; k < j; k++ )
                    sum++;
```

Q10

Compare your analysis with the actual running times.

```
(1) sum = 0;
    for( i = 0; i < n; i++ )
        sum++;
```

```
(2) sum = 0;
    for( i = 0; i < n; i++ )
        for( j = 0; j < n; j++ )
            sum++;
```

```
(3) sum = 0;
    for( i = 0; i < n; i++ )
        for( j = 0; j < n * n; j++ )
            sum++;
```

```
(4) sum = 0;
    for( i = 0; i < n; i++ )
        for( j = 0; j < i; j++ )
            sum++;
```

```
(5) sum = 0;
    for( i = 0; i < n; i++ )
        for( j = 0; j < i * i; j++ )
            for( k = 0; k < j; k++ )
                sum++;
```

```
(6) sum = 0;
    for( i = 1; i < n; i++ )
        for( j = 1; j < i * i; j++ )
            if( j % i == 0 )
                for( k = 0; k < j; k++ )
                    sum++;
```

Q11

Programs A and B are analyzed and found to have worst-case running times no greater than $150N \log_2 N$ and N^2 , respectively. Answer the following questions, if possible:

- a. Which program has the better guarantee on the running time
 - for large values of N ($N > 10,000$)?
 - for small values of N ($N < 100$)?
- b. Which program will run faster on average for $N = 1,000$?

Q12 Maximum Subsequence Sum Problem

Given (possibly negative) integers A_1, A_2, \dots, A_N , find the maximum value of $\sum_{k=i}^j A_k$. (For convenience, the maximum subsequence sum is 0 if all the integers are negative.) [1]

Example:

For input $-2, 11, -4, 13, -5, -2$?

Maximum Subsequence Sum Problem

Input Size	Algorithm Time			
	1 $O(N^3)$	2 $O(N^2)$	3 $O(N \log N)$	4 $O(N)$
$N = 100$	0.000159	0.000006	0.000005	0.000002
$N = 1,000$	0.095857	0.000371	0.000060	0.000022
$N = 10,000$	86.67	0.033322	0.000619	0.000222
$N = 100,000$	NA	3.33	0.006700	0.002205
$N = 1,000,000$	NA	NA	0.074870	0.022711

Figure 2.2 Running times of several algorithms for maximum subsequence sum (in seconds)

Maximum Subsequence Sum Problem

Cubic Maximum Contiguous Subsequence Sum

```
public static int maxSubSum1( int [ ] a )  
  
{  
  
    int maxSum = 0;  
  
    /**  
  
    *TODO  
  
    */  
  
    return maxSum;  
  
}
```

Maximum Subsequence Sum Problem

Quadratic Maximum Contiguous Subsequence Sum

```
public static int maxSubSum2( int [ ] a )  
  
{  
  
    int maxSum = 0;  
  
    /**  
  
    *TODO  
  
    */  
  
    return maxSum;  
  
}
```

Maximum Subsequence Sum Problem

Recursive Maximum Contiguous Subsequence Sum

```
private static int maxSumRec( int [ ] a, int left, int right ){  
  
    /**  
  
    *TODO  
  
    */  
  
    return max3( maxLeftSum, maxRightSum,maxLeftBorderSum + maxRightBorderSum );  
  
}  
  
public static int maxSubSum3( int [ ] a ){  
  
    return maxSumRec( a, 0, a.length - 1 );  
  
}
```

Maximum Subsequence Sum Problem

Linear Time Maximum Contiguous Subsequence Sum

```
public static int maxSubSum4( int [ ] a )  
  
{  
  
    int maxSum = 0; thisSum=0;  
  
    /**  
  
    *TODO  
  
    */  
  
    return maxSum;  
  
}
```

References

[1]Mark Allen Weiss. 2012. Data Structures and Algorithm Analysis in Java. Mark Allen Weiss. Pearson Education.

[2]Ozlem Simsek, CMPE250 Algorithm Analysis Slides

[2]CMPE250 Algorithm Analysis Lecture Slides