

COMP90042 LECTURE 8

CONTEXT-FREE GRAMMARS

SYNTACTIC CONSTITUENTS

- Sequential models like HMMs assume entirely flat structure
- But language clearly isn't like that

[A man] [saw [a dog] [in [the park]]]

- Words group together to form syntactic constituents
 - Can be replaced, or moved around as a unit
- Grammars allow us to formalize these intuitions
 - Symbols correspond to syntactic constituents

TESTING FOR CONSTITUENCY

- Various tests for constituency, based on linguistic intuition
 - Only constituents can answer a question

<u>Trevor gave a lecture on grammar Who gave the lecture on grammar? Trevor gave a lecture on grammar Trevor did what with the lecture on grammar?*gave (fails)</u>
Trevor gave a lecture <u>on grammar</u> What topic was Trevor's lecture on? <u>on grammar</u>

Only constituents can be coordinated with others (of same type)

Trevor gave a lecture <u>on grammar</u> and <u>on parsing</u>
Trevor gave a lecture on <u>grammar</u> and <u>parsing</u>
Trevor gave <u>a lecture on grammar</u> and <u>a treatise on parsing</u>
Trevor <u>gave a lecture on grammar</u> and <u>ate a tasty pie</u>
#Trevor gave <u>a lecture on</u> and <u>a treatise about</u> grammar

▶ More tests, e.g., *topicalisation*, *clefting* and *coordination*.

OUTLINE

- ► The context-free grammar formalism
- Parsing with CFGs
- Representing English with CFGs

BASICS OF CONTEXT-FREE GRAMMARS

- Symbols
 - ► Terminal: word such as *book*
 - Non-terminal: syntactic label such as NP or NN
 - Convention to use upper and lower-case to distinguish, or else "quotes" for terminals
- Productions (rules)

$$W \rightarrow X Y Z$$

- Exactly one non-terminal on left-hand side (LHS)
- An ordered list of symbols on right-hand side (RHS)
 - can be Terminals or Non-terminals

REGULAR EXPRESSIONS AS CFGS

- Regular expressions match simple patterns
 - E.g., [A-Z][a-z]* words starting with a capital
- Can rewrite as a grammar
 - ightharpoonup S
 ightharpoonup C S
 ightharpoonup C RS
 - $ightharpoonup C
 ightharpoonup "A" \qquad C
 ightharpoonup "B" \qquad \dots \qquad C
 ightharpoonup "Z"$
 - ightharpoonup RS
 ightharpoonup R RS ightharpoonup R RS
 - $ightharpoonup R
 ightharpoonup "a" \qquad R
 ightharpoonup "b" \qquad \dots \qquad R
 ightharpoonup "z"$
- In fact, a regex is a way of specifying a regular language (language = set of strings matching pattern)
 - The class of regular languages is a subset of the **context-free languages**, which are specified using a CFG

CFGS VS REGULAR GRAMMARS

- ► CFGs (and regexs) used to describe a set of strings, aka a "language"
- Regular grammars
 - describe a smaller class of languages, e.g., a*b*c*
 - can be parsed using finite state machine
 - ► HMMs are a model for a (weighted) regular grammar (exercise: show the grammar for a POS tagging HMM)
- CFGs
 - can describe hierarchical groupings e.g., matching brackets (aⁿbⁿ) & many recursive tree structures found in language
 - requires push-down automata to parse
- Context sensitive grammars are even more expressive (and intractable)

A SIMPLE GRAMMAR

Terminal symbols: rat, the, ate, cheese

Non-terminal symbols: S, NP, VP, DT, VBD, NN

Productions:

 $S \rightarrow NP VP$

 $NP \rightarrow DT NN$

 $VP \rightarrow VBD NP$

 $DT \rightarrow the$

 $NN \rightarrow rat$

 $NN \rightarrow cheese$

 $VBD \rightarrow ate$

GENERATING SENTENCES WITH CFGS

Always start with S (the sentence/start symbol)

S

Apply rule with S on LHS (S \rightarrow NP VP), i.e substitute RHS

NP VP

Apply rule with NP on LHS (NP \rightarrow DT NN)

DT NN VP

Apply rule with DT on LHS (DT $\rightarrow the$)

the NN VP

Apply rule with NN on LHS (NN $\rightarrow rat$)

the rat VP

GENERATING SENTENCES WITH CFGS

Apply rule with VP on LHS (VP \rightarrow VBD NP)

the rat VBD NP

Apply rule with VBD on LHS (VBD $\rightarrow ate$)

the rat ate NP

Apply rule with NP on LHS (NP \rightarrow DT NN)

the rat ate DT NN

Apply rule with DT on LHS (DT $\rightarrow the$)

the rat ate the NN

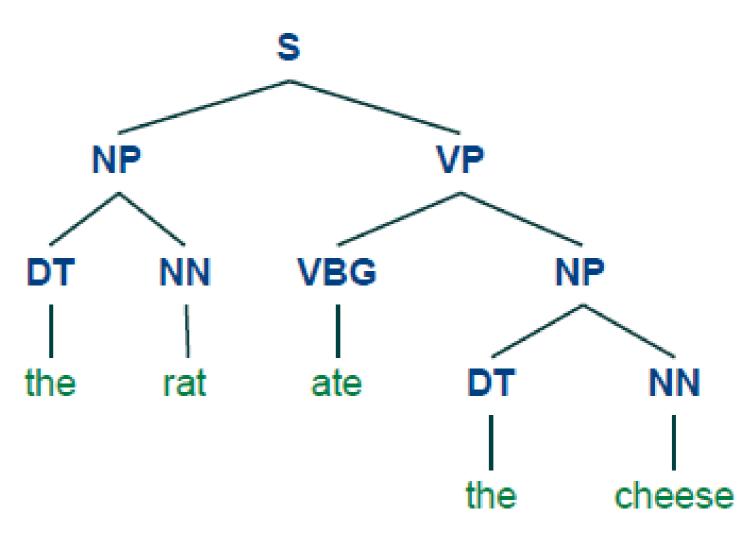
Apply rule with NN on LHS (NN \rightarrow *cheese*)

the rat ate the cheese

CFG TREES

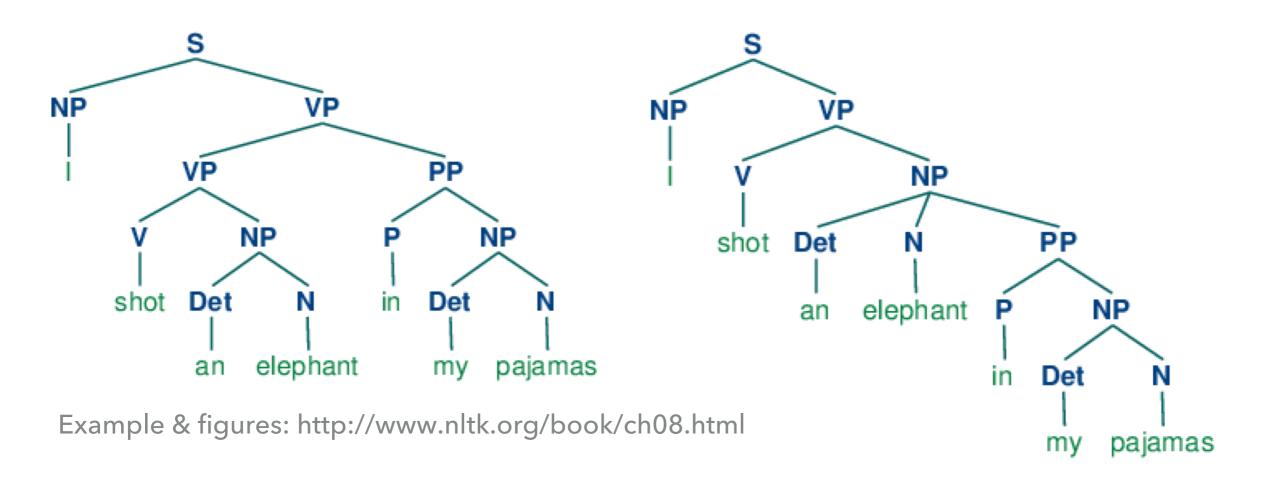
- Generation corresponds to a syntactic tree
- Non-terminals are internal nodes
- Terminals are leaves

Parsing is the reverse process



PARSE AMBIGUITY

- Often more than one tree can describe a string
- While hunting in Africa, I shot an elephant in my pajamas. How he got into my pajamas, I don't know." Animal Crackers (1930)



PARSING CFGS

- Parsing: given string, identify possible structures
- Brute force search is intractable for non-trivial grammars
 - Good solutions use dynamic programming
- Two general strategies
 - Bottom-up
 - Start with words, work up towards S
 - CYK parsing
 - Top-down
 - Start with S, work down towards words
 - Earley parsing (not covered)

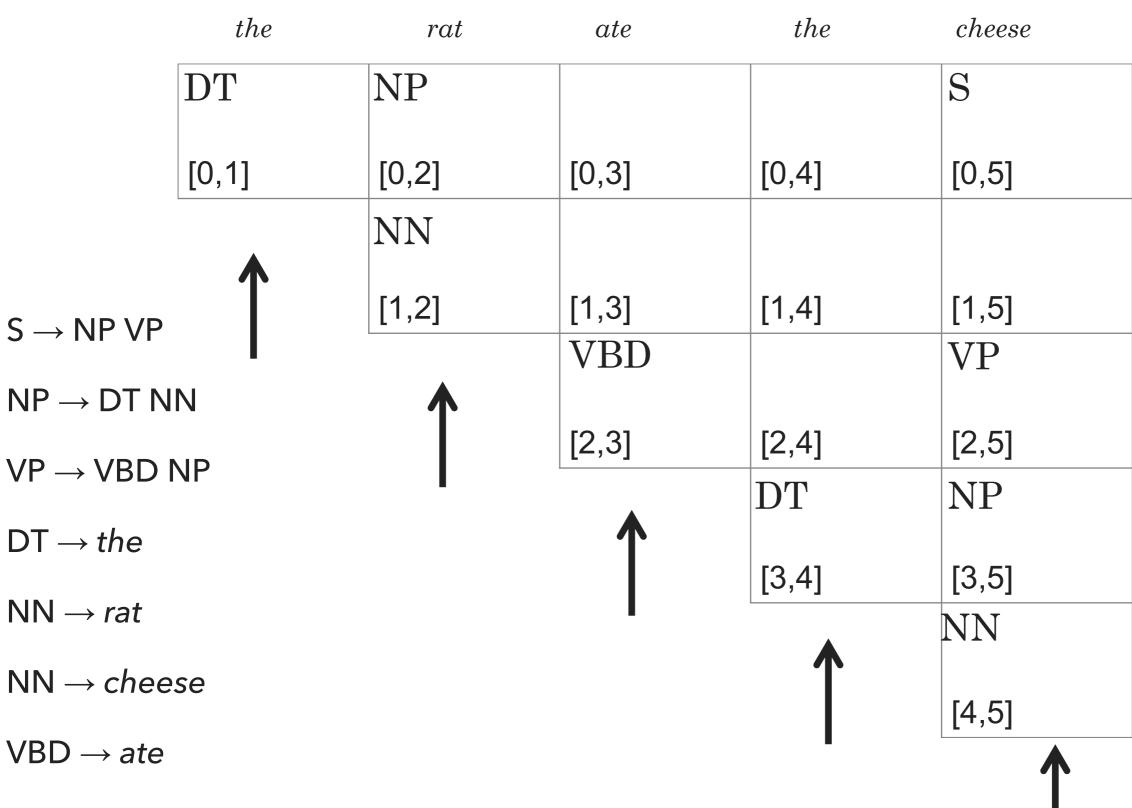
THE CYK PARSING ALGORITHM

- Convert grammar to Chomsky Normal Form (CNF)
- Fill in a parse table
- Use table to derive parse
- Convert result back to original grammar

CONVERT TO CNF

- Change grammar so all rules of form $A \rightarrow B C \text{ or } A \rightarrow a$
- Step 1: Convert rules of form $A \to B c$ into pair of rules $A \to B X$, $X \to c$
 - Not usually necessary in POS-based grammars
- ▶ Step 2: Convert rules $A \rightarrow B C D$ into $A \rightarrow B Y, Y \rightarrow C D$
 - Usually necessary, but not for our toy grammar
- \triangleright X, Y are new symbols we have introduced
- ▶ Unary rules $A \rightarrow B$ can be permitted, but have to factor out cycles)

PARSE TABLE



CYK ALGORITHM

```
function CKY-Parse(words, grammar) returns table

for j \leftarrow from 1 to Length(words) do

for all \{A \mid A \rightarrow words[j] \in grammar\}

table[j-1,j] \leftarrow table[j-1,j] \cup A

for i \leftarrow from j-2 downto 0 do

for k \leftarrow i+1 to j-1 do

for all \{A \mid A \rightarrow BC \in grammar \text{ and } B \in table[i,k] \text{ and } C \in table[k,j]\}

table[i,j] \leftarrow table[i,j] \cup A
```

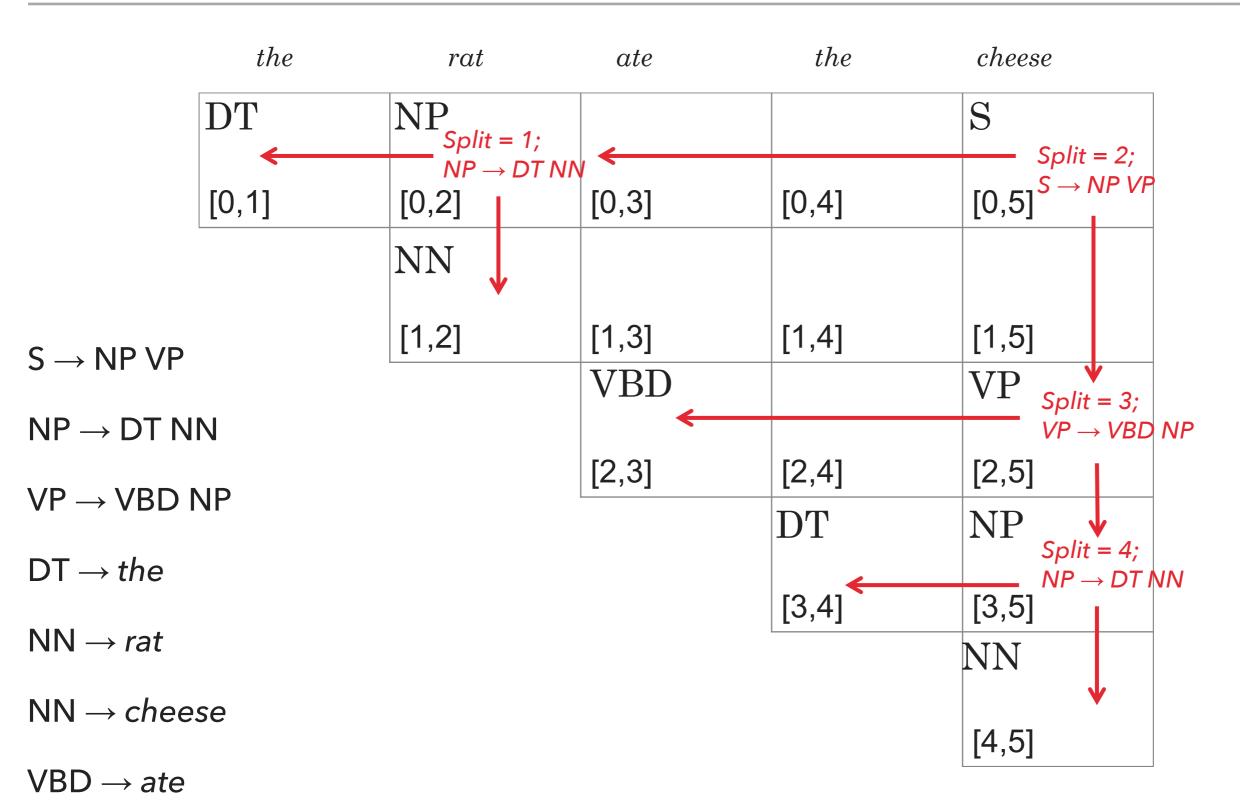
Figure 12.5 The CKY algorithm.

▶ JM3, Ch 12

CYK: RETRIEVING THE PARSES

- S in the top-left corner of parse table indicates success
- ► To get parse(s), follow pointers back for each match
- Convert back from CNF by removing new non-terminals

PARSE TABLE WITH BACKPOINTERS



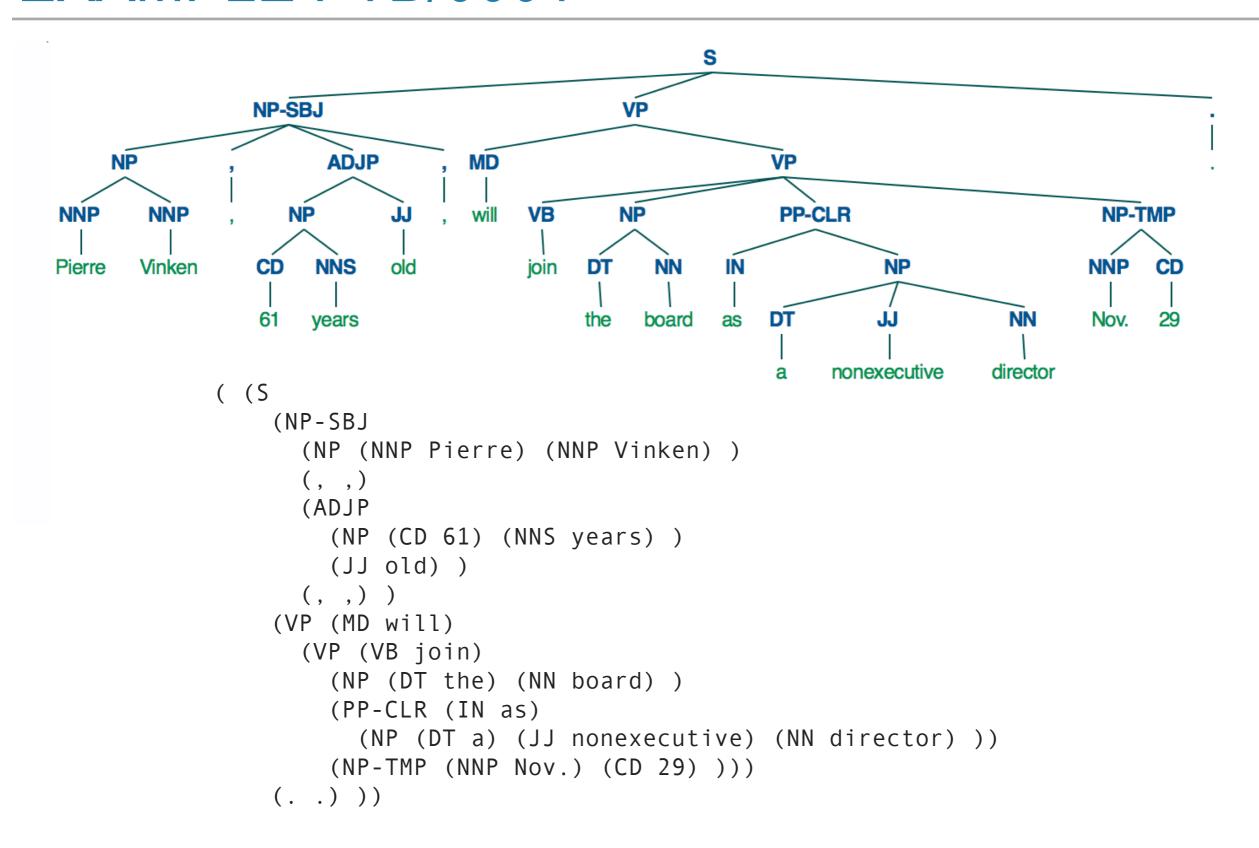
FROM TOY GRAMMARS TO REAL GRAMMARS

- Toy grammars with handful of productions good for demonstration or extremely limited domains
- For real texts, we need real grammars
- Hundreds or thousands of production rules

KEY CONSTITUENTS IN PENN TREEBANK

- Sentence (S)
- Noun phrase (NP)
- Verb phrase (VP)
- Prepositional phrase (PP)
- Adjective phrase (AdjP)
- Adverbial phrase (AdvP)
- Subordinate clause (SBAR)

EXAMPLE PTB/0001



BASIC ENGLISH SENTENCE STRUCTURES

- ▶ Declarative sentences $(S \rightarrow NP VP)$
 - E.g. The rat ate the cheese
- ▶ Imperative sentences $(S \rightarrow VP)$
 - E.g. Eat the cheese!
- \blacktriangleright Yes/no questions (S \rightarrow VB NP VP)
 - E.g. did the rat eat the cheese?
- \blacktriangleright Wh-subject-questions (S \rightarrow WH VP)
 - ▶ Who ate the cheese?
- ▶ Wh-object-questions (S \rightarrow WH VB NP VP)
 - What did the rat eat?

ENGLISH NOUN PHRASES

- Pre-modifiers
 - ▶ DT, CD, ADJP, NNP, NN
 - E.g. the two very best Philly cheese steaks
- Post-modifiers
 - ▶ PP, VP, SBAR
 - ► A delivery from Bob coming today that I don't want to miss

 $NP \rightarrow DT?$ CD? ADJP? (NN | NNP)+ PP* VP? SBAR?

 $NP \rightarrow PRP$

VERB PHRASES

- Auxiliaries
 - ► MD, AdvP, VB, TO
 - E.g should really have tried to wait

 $VP \rightarrow (MD | VB | TO) AdvP? VP$

- Arguments and adjuncts
 - ► NP, PP, SBAR, VP, AdvP
 - E.g told him yesterday that I was ready
 - E.g. gave John a gift for his birthday to make amends

 $VP \rightarrow VB NP? NP? PP* AdvP* VP? SBAR?$

OTHER CONSTITUENTS

- Prepositional phrase
 - ▶ $PP \rightarrow IN NP$ (in the house)
- Adjective phrase
 - ightharpoonup AdjP
 ightharpoonup (AdvP) JJ (really nice)
- Adverb phrase
 - AdvP
 ightharpoonup (AdvP) RB (not too well)
- Subordinate clause
 - ▶ SBAR \rightarrow (IN) S (since I came here)
- Coordination
 - ▶ NP \rightarrow NP CC NP; VP \rightarrow VP CC VP; etc. (Jack and Jill)
- Complex sentences
 - ightharpoonup S
 igh

A FINAL WORD

- Context-free grammars can represent linguistic structure
- ► There are relatively fast dynamic programming algorithms to retrieve this structure
- But what about ambiguity?
 - Extreme ambiguity will slow down parsing
 - ► If multiple possible parses, which is best?

REQUIRED READING

- ▶ J&M3 Ch. 11.1-11.5, Ch. 12.1-12.2
- Constituency tests http://people.umass.edu/nconstan/201/Constituency%20Te sts.pdf