# Key-Value Store with Akka actor model

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## 1. Description of the system

### 1.1. Processes

The aim of this project is to create a key-value store. To do this, we have implemented a multi-writer multi-reader atomic registers system. The principle is therefore to create a system containing several registers. Each register can perform two theoretical actions: read its local value corresponding to a given key, and write a local value for a given key (hence the name 'key-value').

#### 1.1.1. Attributes

This system is composed of `N` processes. Each process runs using nine main attributes, as shown below.

- `processes`: all other processes references
- `mailbox`: a mailbox for storing the incoming messages
- `values`: a local key-value hashmap
- `timestamps`: a local key-timestamp hashmap
- `state`: the current process state
- `proposal`: proposed value in `PUT` operations
- `seqNumber`: the sequence number corresponding to the current operation
- `ackNumber`: the number of received acknowledgments

#### 1.1.2. States

The process can be into five different states, described below.

- `faulty`: This state simulates a process that failed and cannot respond.
- `get`: The process is in this state all along a `GET` request.
- `put`: The process is in this state when it begins a `PUT` request.
- `wait_write`: At the end of a `PUT` request, the process passes to this state until it receives all write responses.
- `none`: By default, and when no operation is running, the process is in this state.

#### 1.1.3. Messages

The processes can receive eight different types of messages. At each message processing, the process executes the following operations:

- `Members`: set local processes references
- `Fail`: pass to state `faulty`
- Operations
  - `Get`: launch a `GET` request
  - `Put`: launch a `PUT` request
- Requests and responses
  - `ReadRequest`: launch a read request to all processes
  - `ReadResponse`: process an incoming read response from a previous read request
  - `WriteRequest`: launch a write request to all processes
  - `WriteResponse`: process an incoming write response from a previous write request

## 1.2. Messages processing

The main method of a process is `onReceive()`. It is the method called each time a process receives a message. According to the type of message, the corresponding private method will be called to process it.

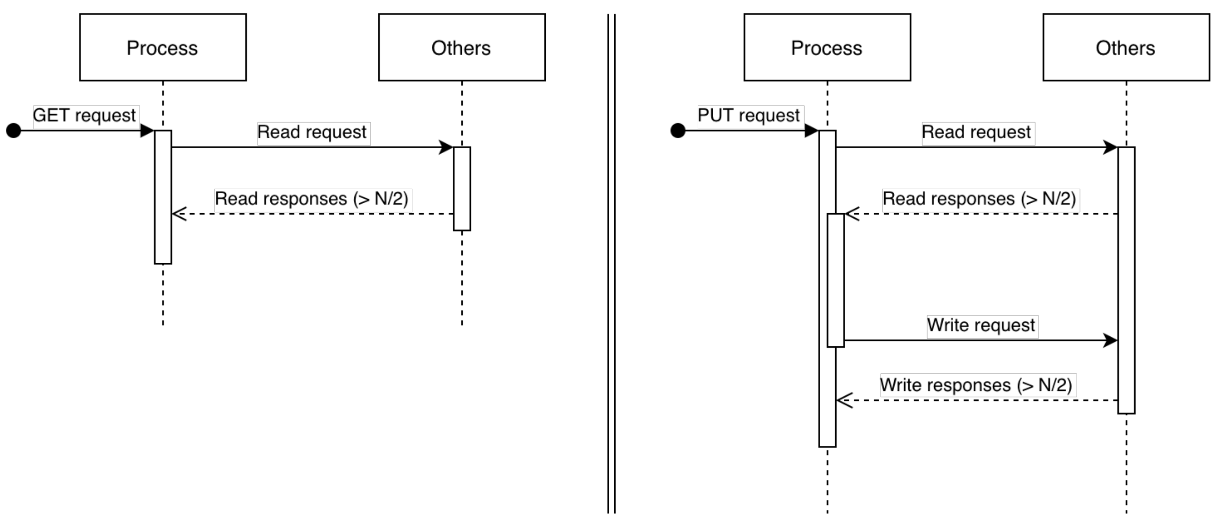
Before discussing how each message is processed, let's see how message arrivals are handled. Each process has a mailbox, as seen before, which is a queue containing all waiting operations, i.e. `GET` and `PUT` messages. As soon as the process terminates processing an operation, it picks the next one in the mailbox.

```
this.state = State.NONE;
nextOperation();
```

Concerning requests and responses, the process processes them immediatly as they arrive, without stocking them. Let's now see it more in details.

### 1.2.1. Operations

The two possible operations are `GET` and `PUT`, and run as follows:




**1.2.1.1. Get** A `Get` message only contains the requested key.

When launching a `GET` request, the process passes to `GET` state and just sends `Read` requests to all other processes. A read request is sent with the key as well as a sequence number, in order to recognise corresponding responses.

**1.2.1.2 Put** A `Put` message contains the requested key and a proposal value to write.

When launching a `PUT` request, the process passes to `PUT` state. Then, it first sends `Read` requests with the key to all other processes.

### 1.2.2. Requests and responses

**1.2.2.1 Read request** A `Read` request contains the requested key and the sequence number of the initial request.

When receiving a read request, the process reads in its values and timestamps maps to find the ones corresponding to the key received. After that, the process sends back a read response with the corresponding sequence number, key, value and timestamp. If no value was found, the process returns a null value and timestamp.

```
sendRequests(Request.READ, msg.key);
```

**1.2.2.2. Read response** A `Read` response contains the initially sent key, the sequence number, as well as the found value and timestamp.

This is the most important method, divided in two cases as explained below.

Before verifying these cases, the first thing to do is to increment `ackNumber`, the number of acknowledgements for the read request.

Then, if the received value is not null, the process verifies if the timestamp is greater than his own, in which case it will update his own value and timestamp. If the timestamps are equal, the process will keep the highest value.
This procedure permits to have consistent results, i.e. that all processes agree on a single value-timestamp couple.
We chose to directly write the values as read requests are received, considering a `Read` request as an update throughout the system for a given key.

At this point, if enough responses have been received, i.e. a majority ($\geqslant \frac{N}{2}$), there are two cases:

- **`GET` state**

If the process is in `GET` state, the operation is finished. In fact, either its local map contains a value for the corresponding key and we can say the process has got a value for the given key, or it does not contain a value, in which case the process failed to get a value for this key. This last case can only happen when no process received a `PUT` request for the key.
We choose not to send additionnal write requests, as the process already received a majority of correct responses. Ask to write to other processes does not seem to be in the scope of a `GET` request. Moreover, and we will discuss that later, our testing script approves that the algorithm we use works as is.

- **`PUT` state**

If the process is in `PUT` state, the operation is obviously not finished. As shown on the previous diagram, the process still has to send write requests to alla other processes, with the given proposal value and the highest got timestamp incremented by one.
If all the read responses returned null, the default timestamp to put is 1.

```
putValue(msg.key, this.proposal, putTimestamp + 1);
```

#### 1.2.2.3. Write request

A `Write` request contains the key, the sequence number, as well as the proposal value and the new timestamp.

The principle here is that the process updates its local value if the same conditions as previously are met: the new timestamp is greater to local one, or if it is equal the proposal is greater than the local value.
Of course, if the value is not found in the local map, the processes writes it.

```
msg.timestamp > localTimestamp ||
(msg.timestamp == localTimestamp && msg.proposal > localValue)
```

Finally, the process sends a confirmation `Write` response, only containing the sequence number and the wrote key.

#### 1.2.2.4. Write response

As explained, a `Write` response contains the sequence number of the initial request and the key.

When receiving a `Write` response, the process only has to check the number of acknowledgements. If the majority of the processes in the system responded ($\geqslant \frac{N}{2}$), the `PUT` operation is done and the process can pass to next one in the mailbox.

# 2. Proof of correctness

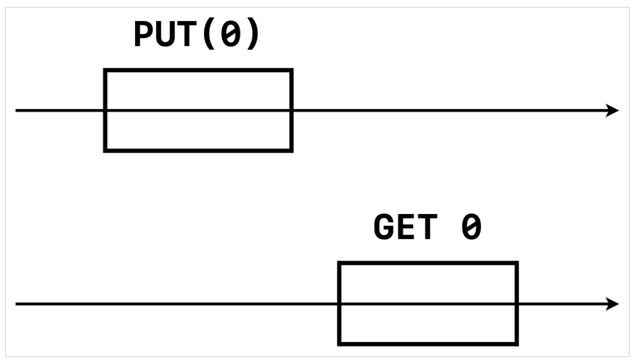
## 2.1. Correctness

Before making a proof of correctness, we need to define what is correctness. We can say that an execution is correct when it satisfies 2 properties: liveness and safety. Let's see the signification of these properties.

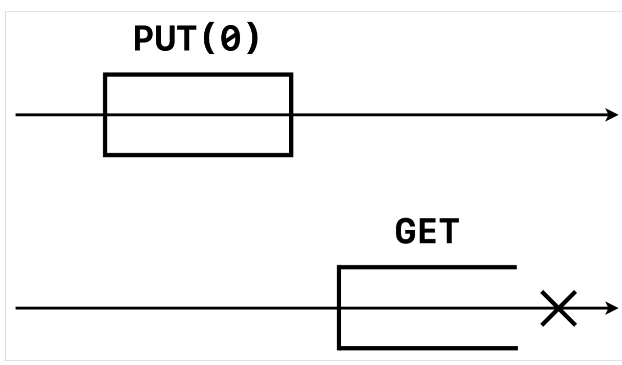
### 2.1.1. Liveness

The first one, liveness, insures that every operation invoked eventually returns. So every started operation must finish some time after.

Here is an exemple of an execution that respects liveness:



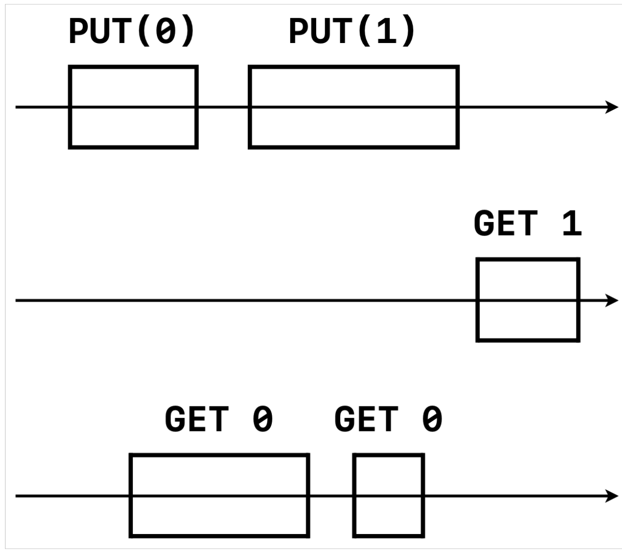

On the contrary, the get operation never ends in the below diagram. This execution does not respect liveness:



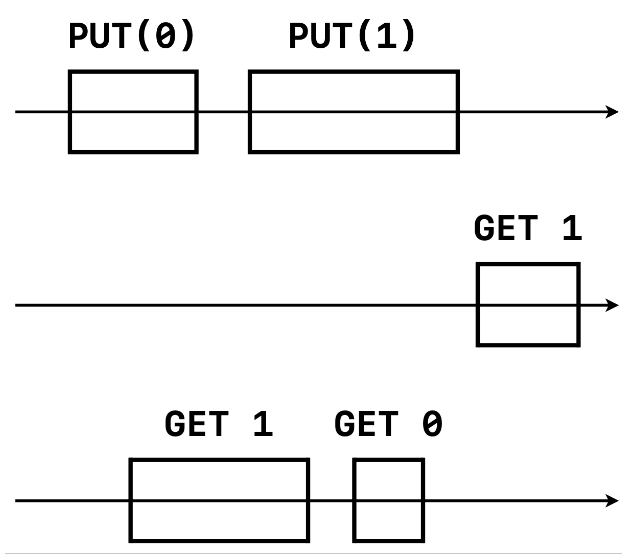

### 2.1.2. Safety

Then, safety insures that the history of the execution is linearizable. That is to say that operations can be totally ordered, preserving legality and precedence. For example, if read1 returns v and read2 returns v', and read1 precedes read2, then write(v') cannot precede write(v).

Here is the history of a safe execution:




On the contrary, the next history does not represent a safe execution due to new-old inversion:




## 2.2. How we check correctness

To check the correctness of our implementation, we execute our program nine times with the combination of these parameters: N = 3, 10, 100 and M = 3, 10, 100. For each instance, we print the logs to a text file and analyze them with a Python program we made.

All the steps are automated:

1. The program launches one by one the nine instances with different N and M parameter values. To perform this, it executes the Java program with the corresponding command line arguments (N and

M values). The output is redirected to a text file.
2. Then, it reads the text file created (which contains the logs) during the execution.
3. From there, the history of the execution is created by parsing the logs and storing the results in objects (operation and history classes). Liveness is verified during this phase. If each operation launched terminates, the execution is lively.
4. After that, we verify if the execution is safety. To perform this, we verify for each get that its return value is either the one of the last put or the one of a concurrent put and that a new-old inversion didn't happen.
5. If liveness and safety are respected, the execution is correct and we display the performance data. We will study that in the next part.

### 2.3. Output

When launching our Python program, we get an output of this form:

```
$ python correctness.py
Testing with N = 3 and M = 3
Lively!
Safe!
Total computation time: 0.029564 sec
Put median duration: 1260.5 us
Get median duration: 482.0 us

Testing with N = 3 and M = 10
Lively!
Safe!
Total computation time: 0.032373 sec
Put median duration: 640.0 us
Get median duration: 233.5 us

Testing with N = 3 and M = 100
Lively!
Safe!
Total computation time: 0.132201 sec
Put median duration: 193.0 us
Get median duration: 106.0 us

Testing with N = 10 and M = 3
Lively!
Safe!
Total computation time: 0.030162 sec
Put median duration: 2372.0 us
Get median duration: 1485.5 us

Testing with N = 10 and M = 10
Lively!
Safe!
Total computation time: 0.06442 sec
Put median duration: 2303.5 us
Get median duration: 1096.0 us

Testing with N = 10 and M = 100
Lively!
```

```
Safe!
Total computation time: 0.209297 sec
Put median duration: 867.0 us
Get median duration: 442.0 us

Testing with N = 100 and M = 3
Lively!
Safe!
Total computation time: 0.190535 sec
Put median duration: 21983 us
Get median duration: 11605 us

Testing with N = 100 and M = 10
Lively!
Safe!
Total computation time: 0.325907 sec
Put median duration: 17072.5 us
Get median duration: 7837.5 us

Testing with N = 100 and M = 100
Lively!
Safe!
Total computation time: 1.095113 sec
Put median duration: 5313.5 us
Get median duration: 2617.5 us
```

We can see that our implementation is correct for every N and M combination.

# 3. Performance analysis

## 3.1. How we measure durations

`//TODO`

## 3.2. Results and analysis

- Total computation time (in sec)

| N\M | 3 | 10 | 100 |
|---|---|---|---|
| 3 | 0.029564 | 0.032373 | 0.132201 |
| 10 | 0.030162 | 0.064420 | 0.209297 |
| 100 | 0.190535 | 0.325907 | 1.095113 |

As we can see, the total computation time increases as N increases and also as M increases. This seems reasonable. Moreover, our program has a very low latency: the highest computation time is only about a second.

- Put median duration (in $\mu$s)

| N\M | 3 | 10 | 100 |
|---|---|---|---|
| 3 | 1260.5 | 640.0 | 193.0 |
| 10 | 2372.0 | 2303.5 | 867.0 |
| 100 | 21983.0 | 17072.5 | 5313.5 |

For the put median duration, we remark that it increases when N increases but it decreases as M increases. This can be explained by the fact that operations are performed faster and faster as the execution is going on.

- Get median duration (in $\mu$s)

| N\M | 3 | 10 | 100 |
|---|---|---|---|
| 3 | 482.0 | 233.5 | 106.0 |
| 10 | 1485.5 | 1096.0 | 442.0 |
| 100 | 11605.0 | 7837.5 | 2617.5 |

For the get median duration, we can make the same observations as for the put median duration. The difference is that the median duration is smaller for get operations. Indeed, it only reads through other processes to get the value. A put operation reads through other processes to get the maximum timestamp and then write the new value to all. So it is reasonable to see this difference.