

LAB - 1

- 1) Implement tic-tac-toe game, 2 players

Algorithm:

- 1) Initialize a 3×3 game board using lists in Python and initialize all cells as empty.
- 2) Write a function to display the board by printing the current state of the board.
- 3) Write a function to check for win, check all rows, columns and diagonals.

row

```
for row in board
    if all(cell == player for cell in row)
        return True
```

col

```
for col in range(3):
    if all(board[row][col] == player for row in range(3)):
        return True
```

diagonal

```
if all(board[i][i] == player for i in range(3)) or
all(board[i][2-i] == player for i in range(3)):
    return True
```

- A) Check the current-player. If current player is human, call get-player-move (board, current-player) else call ai-move (board)
- 5) After every move check if the player wins. If there is no win-condition, change the current player.
- 6) For every move, we check if the board is full. If it is full, display the game is a tie.

2) Vacuum cleaner agent

def vacuum-cleaner-agent(location, status):

$x, y = \text{location}$

if status[x][y] == 'Dirty':

'return f"Vacuum cleaner is at ({x}, {y}) and
it is dirty. Cleaning"

else:

'return f"The vacuum cleaner is at ({x}, {y}) and
it is clean. Moving"

status

status = [[{'Diry': 1, 'Clean': 0}, {'Diry': 0, 'Clean': 1}], [{}]]

location = (0, 0)

while True:

action = vacuum-cleaner-agent(location, status)

print(action)

$x, y = \text{location}$

if status[x][y] == 'Diry':

status[x][y] = 'Clean'.

If status[0][0] == 'Clean' and status[0][1] == 'Clean'

and status[1][0] == 'Clean' and status[1][1] == 'Clean':

print ("All clean").

break

if $y < 1$:

location = (x, y + 1)

elif $x < 1$

location = (x + 1, 0)

Output

VC is at (0,0); dirty, clean

Vr is at (0,1), clean, Many

VC is at (1,0), dirty, clean

VC is at (1,1), dirty, clean

All clean

Diagnosis [P] = 0.177 - 0.001

(0,0) = normal

1st stage

(0,1) = normal

(1,0) = normal

1st stage = 0.001

2nd stage = 0.001

3rd stage = 0.001

4th stage = 0.001

5th stage = 0.001

6th stage = 0.001

7th stage = 0.001

8th stage = 0.001

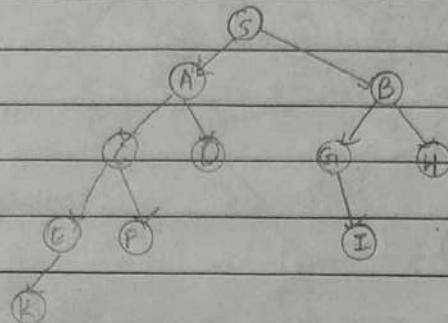
9th stage = 0.001

10th stage = 0.001

LAB-2

Solutions for the given search problems and pseudocode

1. Perform BFS

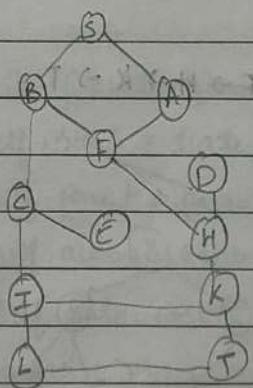


Frontier: S A B C D G H E F I K

Explored: (S) A B C D G H E I ↴ goal reached

Backtracking: S → A → C → E ↴ A ← C ← E ← S ← goal reached

2. Performing DFS both tree search and graph search



Frontier:

S

B A

C F A

S E F A

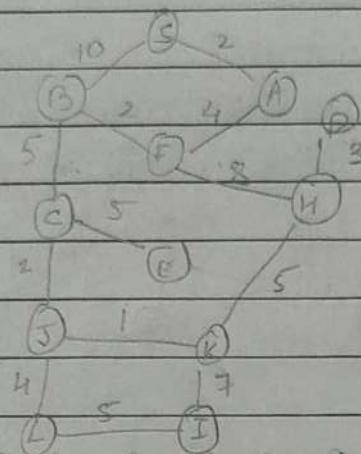
✓ K E F A

T K E F A

↓ goal reached

explored: S B C S L

3. Perform Uniform Cost Search



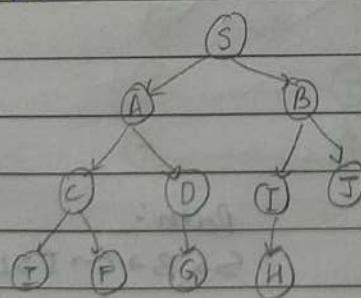
Frontier: $S(0)$ $A(2)$ $B(10)$ $F(2)$ $H(8)$ $C(5)$ $K(15)$
 $D(17)$ $J(17)$ $E(20)$ $L(21)$ $T(26)$

Explored: $S(0)$ $A(2)$ $F(2)$ $B(10)$ $H(8)$ $C(5)$ $J(17)$
 $K(15)$ $T(26)$

Total cost = 26

Path: $S \rightarrow A \rightarrow F \rightarrow H \rightarrow K \rightarrow T$

4. Performing DLS



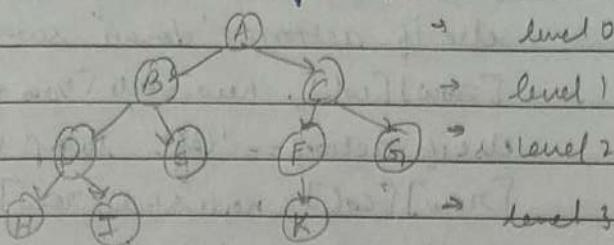
Frontier:

<input checked="" type="checkbox"/>	<input type="checkbox"/> A B	<input type="checkbox"/> D B	<input type="checkbox"/> D B	<input type="checkbox"/> I I	<input type="checkbox"/> I I	<input type="checkbox"/> I I
-------------------------------------	---------------------------------	---------------------------------	---------------------------------	---------------------------------	---------------------------------	---------------------------------

Explored: S A C D B I J

Path: $S \rightarrow B \rightarrow J$

S. Performing Iterative deepening search



Search process:

- i) A
- ii) $A \rightarrow B \rightarrow C$
- iii) $A \rightarrow B \rightarrow D \rightarrow E \rightarrow C$
- iv) $A \rightarrow B \rightarrow D \rightarrow E \rightarrow C \rightarrow F \rightarrow G$

Path: $A \rightarrow C \rightarrow G$

- i) Solve 8-puzzle problem.

Pseudocode:

Class Node:

```

function init_(state, parent, action, path-cost = 0):
    set self.state = state
    set self.parent = parent
    set self.action = action
    set self.path-cost = path-cost
  
```

function expand():

Create children

Set row, col = find-blank()

Create possible-actions

If row > 0 then add 'up' to possible-actions

If row < 2 then add 'down' to possible-actions

If col > 0 then add 'left' to possible-actions

If col < 2 then add 'right' to possible-actions

For action in possible-actions:

Create new-state as a copy } Set state
If action == 'up' Then swap new-state[0][0] }
[row][col]

with new-state [row-1][col]
 else if action == "down" swap new-state
 [row][col], new-state [row+1][col]
 elseif action == "left" swap new-state
 [row][col], new-state [row][col-1]
 else if action == "right" swap new-state
 [row][col], new-state [row][col+1]
 append new Node (new-state, self.parent, action,
 self.path_cost + 1) to children
 return children
 function find-blank():
 for row from 0 to 2:
 for col from 0 to 2:
 if self.state [row][col] == 0 then
 return row, col
 function dept-first-search (initial-state; goal-state):
 set frontier = [Node (initial-state)]
 set explored = empty set
 while frontier is not empty:
 set node = frontier.pop()
 if node.state == goal-state then
 return node
 add tuple of node state to explored
 for child in node.expand():
 if tuple of child state not in explored
 then append child to frontier
 return none
 function print-solution (node):
 create path
 while node is not none:
 append node.action, node.state to path
 set node = node.parent
 reverse path.

for (action, state) in paths:

If action is not none then print "action:"
 action

print state

Print "

Set initial-state = [[1, 2, 3], [0, 4, 6], [7, 5, 8]]

Set goal-state = [[1, 2, 3], [4, 5, 6], [7, 8, 0]]

Set solution = depth-first-search(initial-state,
 goal-state)

If solution is not none then

print "Sol found".

call print-solution(solution)

else

print ("Sol not found")

(2)

Implement iterative deepening search algorithm.

function iterative-deepening-search(initial-state, goal-state,
 max-depth):

for depth from 0 to max-depth:

set results = depth-limited-search(initial-state,
 goal-state, depth)

If result is not none then

return result

return none

function depth-limited-search(node, goal-state, limit):

If node.state == goal-state then

return node

If node.depth >= limit then

return none

for each child in expand(node):

set result = depth-limited-search(child,
 goal-state, limit)

if result is not none then
 return result
return none
set initial state, goal state, max-depth
set solution = iterative-deepening-search(initial
state, goal state, max-depth)
if solution is not none then print solution
else print "no solution found"

LAB. 04

i) 8 puzzle problem using heuristic approach (A*)

Initial State:

2	8	3
1	6	4
7	5	

$$g=0$$

$$h=5$$

Goal State:

1	2	3
8		4
7	6	5

2	8	3
1	6	4
7	5	

R

h=4	2	8	3
g=1	1	6	4
f=5	7	5	

2	8	3
6	4	
1	7	5

$$h=5$$

$$g=1$$

$$f=6$$

L

R

U

h=5	2	8	3	2	8	3	2	8	3
g=2	1	6	4	1	6	4	1	8	4
f=7	7	5		7	5		7	6	5

$$h=5, g=2, f=7$$

X

R

U

D

h=3	2	8	3	2	8	3	2	8	3
g=3	1	4		1	4		1	8	4
f=5	7	6	5	7	6	5	7	6	5

$$h=4$$

$$h=4, g=3, f=7$$

$$h=3, g=3, f=7$$

R U D

L

D

2	8	3	2	3	2	8	3
1	4		2	1	4	7	1
7	6	5	7	6	5	6	5

$$h=3, g=4, f=7$$

$$h=3, g=4, f=8$$

$$h=3, g=3, f=6$$

$$h=6, g=4, f=8$$

$$h=3, g=4, f=7$$

R

D

$$h=3$$

$$h=3$$

$$h=1$$

$$g=5$$

$$g=5$$

$$g=5$$

$$f=8$$

$$f=8$$

$$f=6$$

R

U

P

1	2	3
8	4	
7	6	5

2	3
1	8
7	6

1	2	3
7	4	
6	5	

goal reached.

(ii) 8 puzzle problem using Manhattan Distance approach

Initial state:

2	3	8
1	6	4
7	5	

Goal state:

1	2	3
8		4
7	6	5

2	8	3
1	6	4
7	5	

$h=6$
 $g=0$
 $f=h+g = 6$

1 2 6 8

$1+1+1+2 = 5 = h$

$g=1$

R U

2 8 3 5

1 6 4 4

7 5 1 7 5

1 2 6 7 8

$2+1+1+1+2 = 7 = h$

$g=1$

2	8	3
1	6	4
7	5	

2	8	3
1		4
7	6	5

$g=2$

1 2 6 7 8

$1+1+1+1+2 = 6 = h$

1 2 5 6 7 8

$1+1+1+1+2 = 6 = g$

$= h = g$

1 2 8

$1+1+2 = 4 = h$

2	8	3
1	4	
7	6	5

2	8	3
1		4
7	6	5

2	8	3
1	6	4
7	5	

2	3
1	8
7	6

$h=5$

$n=5$

$h=5$

$n=5$

$g=3$

2	8	3
1		4
7	6	5

2	3
1	8
7	6

2	3
1	8
7	6

$h=4$

$h=2$

$n=4$

$g=4$

1	2	3
	8	4
7	6	5

2	3
1	8
7	6

$h=3$

$g=5$

2	3
8	4
7	6

2	3
1	8
7	6

1	2	3
	8	4
7	6	5

$h=0$

goal state reached

$g=6$

i) Pseudocode for 8-puzzle using Heuristic approach.

function solve_puzzle(initial-state, goal-state):

priority-queue = [heuristic(initial-state, goal-state, 0, initial-state, [])]

visited = empty set

while priority-queue is not empty:

(f-cost, g-cost, current-state, parent-path) =

pop element with smallest f-cost from priority-queue.

[CCT, 2018-19]

if current-state is equal to goal-state:

return current-path + [current-state]

if current-state is in visited:

continue

add current-state to visited

(with break)

for next-state, action in get-possible-moves(current-state):

new-g-cost = g-cost + 1

new-f-cost = new-g-cost + heuristic(next-state, goal-state).

push(new-f-cost, new-g-cost, next-state,

current-path + [current-state, action]) into

priority-queue

return [newer, visited]

function heuristic(state, goal-state):

calculate the heuristic value by summing up

no. of moves to be taken by misplaced tiles to

come to correct position.

by checking the number of misplace tiles.

+ 1 for blank - goal - state] + no. of misplaced tiles

blank - goal

no. of misplaced tiles

function get-possible-moves(state):

 find the position of 0 tile and

 generate a list of possible moves (Up, Down, Left,
 right) by swapping.

 blank the tile with its adjacent

 Return list of (new-state, action) pairs.

ii) pseudocode for 8 puzzle using Manhattan approach

function solve-8-puzzle(initial-state, goal-state):

 priority-queue = [(manhattan-distance(initial-state, goal-state),
 0, initial-state, [])]

 visited = empty set

 while priority-queue is not empty:

 (f-cost, g-cost, current-state, current-path) = pop. Smallest
 f-cost from queue

 if current-state == goal-state : return current-path +
 [current-state]

 if current-state in visited : continue

 visited.add(current-state)

 for next-state, action in get-moves(current-state):

 new-g-cost = g-cost + 1

 new-f-cost = new-g-cost + manhattan dist

 (next-state, goal-state)

 add (new-f-cost, new-g-cost, next-state, current-path
 + [current-state, action]) to queue

 return None

function manhattan-dist(state, goal-state):

 distance = 0

 for each tile in state:

 if tile != 0 :

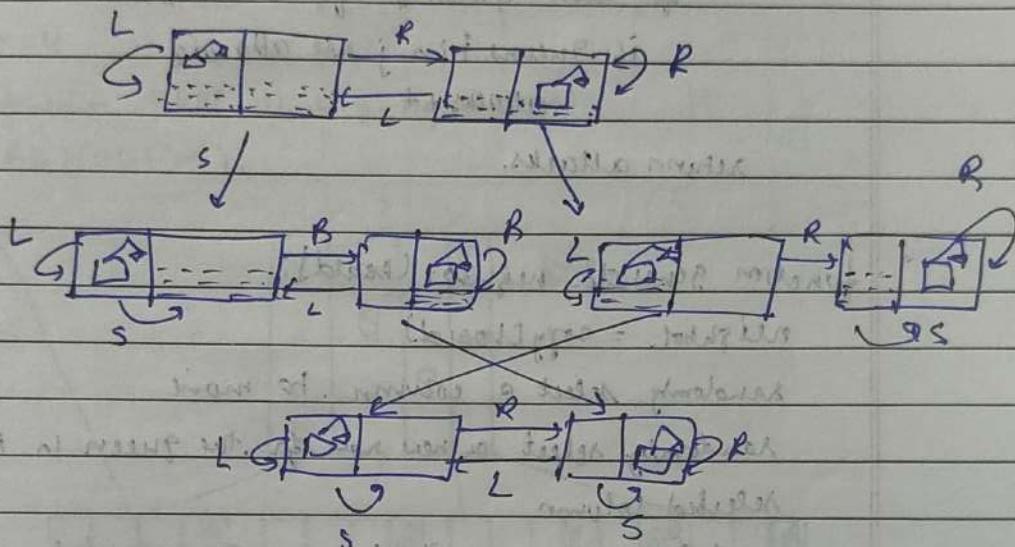
 distance += |goal-row - current-row| +

 |goal-col - current-col|

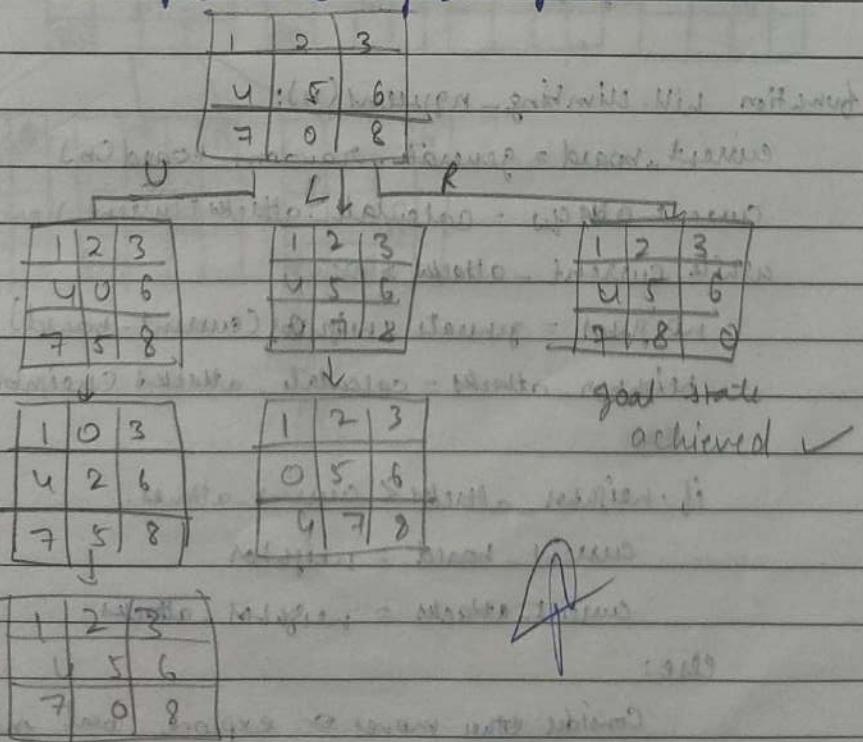
 return distance.

function get-moves(state):
 find blank tile position
 generate possible moves (up, down, left, right) by
 swapping blank tile with adjacent tiles
 return list of (new-state, action) pairs.

2) VACUUM-CLEANER AGENT: (Space tree)



3) State space tree - 8-puzzle problem (DFS):



LAB 5:

Implement Hill Climbing search algorithm to solve N-queens problem:

Function to calculate the number of attacking pairs.

function calculate_attacks (board):

 attacks = 0

 for each queen i in board:

 for each queen j after i in board:

 if queens i in j are attacking:

 attacks++

 return attacks.

function generate_neighor (board):

 neighbor = copy (board)

 randomly select a column to move

 randomly select a new row for the queen in the selected column

 update neighbor with the new queen position

 return neighbor

function hill_climbing_nqueens (n):

 current_board = generate_random_board (n)

 current_attacks = calculate_attacks (current_board)

 while current_attacks > 0:

 neighbor = generate_neighor (current_board)

 neighbor_attacks = calculate_attacks (neighbor)

 if neighbor_attacks < current_attacks:

 current_board = neighbor

 current_attacks = neighbor_attacks

 else:

 Consider other moves or explore local minima

return current_board

function generate_random_board(n):

board = new array of size n

for each columns in board:

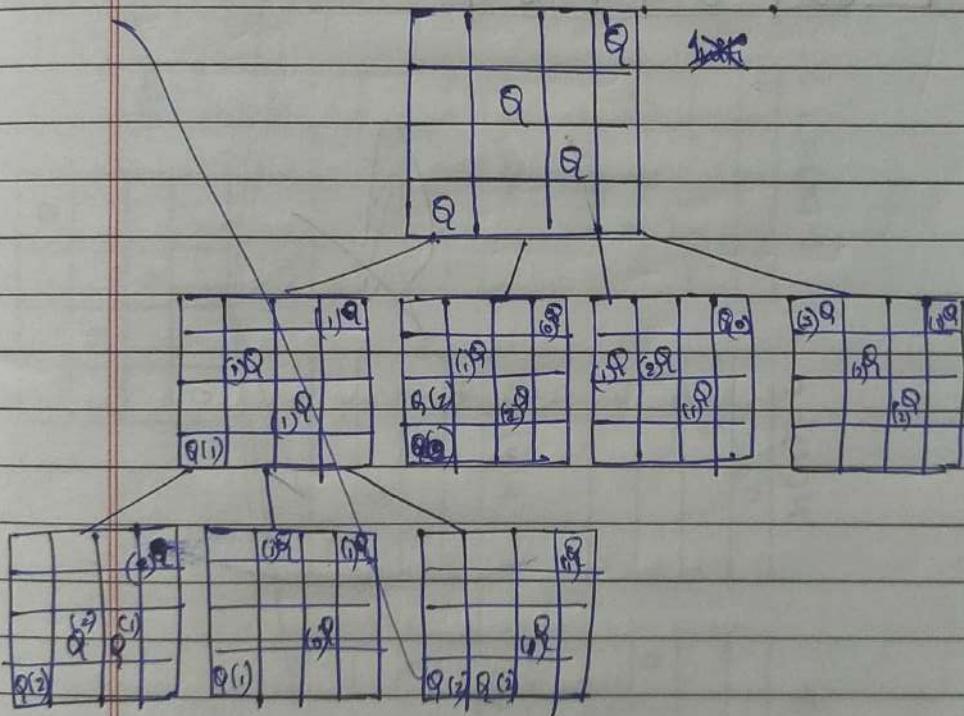
Place a queen in a random row in the column

return board

$n=4$

solution = hill-climbing-nqueens(n)

print(solution).



Initial board

..... Q
.. Q ..
.. Q ..
Q .. .

Final board

..... Q
.. Q ..
.. Q ..
Q .. .

Conflicts = 2

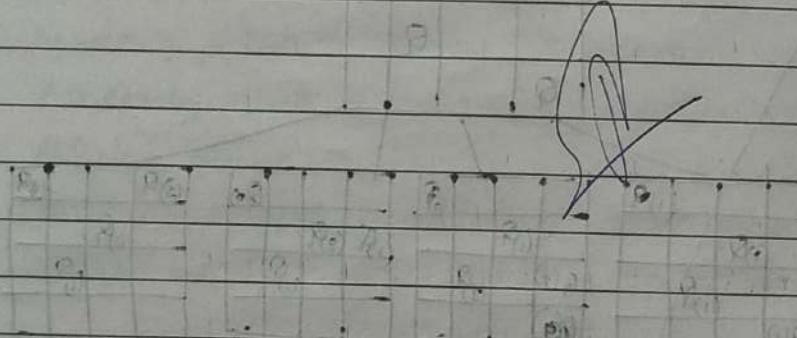
Conflicts = 2

solution: [3, 1, 2, 0]

Conflict = 2

Current heuristic = (1+1+1)+1 = 4

8	4	5	(Q)
4	(Q)	7	6
6	8	(Q)	4
(Q)	6	4	8



$n(m)=2$ Overall cost

0	0	0	0
0	0	0	0
0	0	0	0
0	0	0	0

0	0	0
0	0	0
0	0	0

$$h(m) = 1$$

0	0	0
0	0	0
0	0	0

$$h(m) = 0.$$

0	0	0
0	0	0
0	0	0

$$h(m) = 1$$

$$h(m) = 4$$

$\frac{2}{2} \times \frac{1}{2} \times \frac{1}{2}$

0	0	0	0
0	0	0	0
0	0	0	0
0	0	0	0

$$h(m) = 6$$

$$h(m) = 1$$

0	0	0	0
0	0	0	0
0	0	0	0

$$h(m) = 1$$

$$h(m) = 2$$

0	0	0	0
0	0	0	0
0	0	0	0

$$h(m) = 1$$

$$h(m) = 2$$

LAB-6

1. Analyse and Implement N Queen Problem wiy simulated annealing technique

Pseudocode:

function simulated annealing (n , initial-temp, coolig_rat, iterations) :

 current_board = create_board(n)

 current_conflicts = calculate_conflicts (current_board)

 best_board = current_board

 best_conflict = current_conflicts

 temperature = initial_temperature

 for i in range (iterations):

 position \rightarrow one queen

 neighbour_board \rightarrow generate_neighbor (current_board)

 neighbour_conflicts = calculate_conflicts

 delta_e = $(\text{neighbour board}) - (\text{current board})$

 neighbour_conflicts - current_conflicts

 state if its better or with a probability

 that decreases with temperature if the

 difference to conflicts

If $\Delta E < 0$ or random number ($0, 1$) $< \exp(-\Delta E / \text{temperature})$:

 current_board = neighbour_board

 current_conflicts = neighbour_conflicts

To current_conflicts \leftarrow best_conflicts?

 best_board = current_board

 best_conflicts = current_conflicts

 temperature = $\log(\text{cooling rate})$

return best-board, best-conflict

function create-board(n):

with n queens

return a list of n random integers between
0 and $n-1$ (inclusive)

function calculate-conflicts(board):

$n = \text{length of board}$

conflicts = 0

for each pair of queens (i, j), where $i \neq j$:

if queens i and j are in the same row
or diagonal:

conflicts = conflicts + 1.

return conflicts.

2. Propositional logic - Truth Table Enumeration.

Program:

```
import itertools.
```

```
def evaluate_formula(formula, valuation):
```

```
    formula = formula.replace('P', str(valuation['P']))
```

```
    formula = formula.replace('Q', str(valuation['Q']))
```

```
    return eval(formula)
```

```
def extract_variables(formula):
```

```
    variables = set()
```

```
    for char in formula:
```

```
        if ((char, isalpha())):

```

```
            variables.add(char)
```

```
    return list(variables)
```

```
def generate_truth_table(KB, query):
```

```
    variables = extract_variables(KB) + extract_variables(query)
```

```
    variables = list(set(variables))
```

```
print("Truth Table")
```

```
print("1", join(variables + [KB, query]))
```

```
print("-" * (len(variables) * 4 + 12))
```

entails - query - True

```
for assignment in itertools.product([False, True], repeat = len(variables)):
```

```
    valuation = dict(zip(variables, assignment))
```

KB-truth = evaluate-formula (Query, valuation)

row = [str ('T' if valuation [var] else 'F') for var in variables]

now.append(str ('T' if query-truth else 'F'))

print

if KB-truth and not Query-truth:

entails - query - False

KB = input ("Enter KB ")

query = input ("Enter query ")

generateTruthTable (KB, query)

OUTPUT:

Enter KB (eg : 'P and (q1 = 2)') : 'P & 2

Enter the query : ?

Truth Table:

P | q1 | KB | Query

F | T | False | False, because it just "P"

F | T | T | (False, T) is True

T | F | T | F | because ? (p, q)

T | T | T | T | and right is always true

KB entails query : False

LAB-7

First Order Logic:

1. "John is a human"

FOL:

$$H(\text{John})$$

Explanation:

$H(n)$: n is human

"John" is a constant representing ~~an individual~~ an individual named john

2. "Every human is mortal"

FOL:

$$\forall n (H(n) \rightarrow M(n))$$

$H(n)$: n is a human

$M(n)$: n is mortal

$\forall n$: For all individual n , if n is a human then n is mortal

3. "John loves Mary"

$$L(\text{John}, \text{Mary})$$

$L(n, y)$: n loves y

John and Mary are individuals.

4. "There is someone who loves Mary"

$$\exists n (L(n, \text{Mary}))$$

$L(n, y)$: n loves y

$\exists n$: There exists an individual n such that n loves Mary

5. All dogs are animals.

FOL

$$\forall n (D(n) \rightarrow A(n))$$

$D(n)$: n is a dog

$A(n)$: n is an animal

$\forall n$: For all individuals n , if n is dog $\rightarrow n$ is animal

6. "Some dogs are brown".
 $(\exists x)(D(x) \wedge B(x))$

$D(x)$: x is a dog.

$B(x)$: x is brown

$\exists x$: There exists an individual x such that x is a dog and x is brown.

1)

Let:

S: John is stupid

L: John is lazy

~S: John is not stupid

~L: John is not lazy

Either John isn't stupid and he is lazy, or he is stupid

$(\neg S \wedge L) \vee S$

John is stupid.

Therefore, John is not lazy. $\therefore L$

S	L	$\neg S \wedge L$	$\neg S \wedge \neg L$	$(\neg S \wedge L) \vee S$	S	$\neg L$
T	F	F	F	T	T	F
T	F	F	F	T	+ T	T
F	T	T	T	T	F	F
F	F	T	F	F	F	T

$(P \rightarrow Q) \vee (\neg P \wedge R)$

P	Q	R	$P \rightarrow Q$	$\neg P$	$\neg P \wedge R$	$(P \rightarrow Q) \vee (\neg P \wedge R)$
T	T	F	F	F	F	T
T	F	T	F	F	F	T
T	F	F	F	T	F	T
T	F	F	F	F	F	F
F	T	T	T	T	F	T
F	T	F	T	F	F	T
F	F	T	T	T	T	T
F	F	F	F	F	F	F

function translate-to-fol(sentence):

sentence = lowercase_and_trim(sentence)

if sentence contains "is-a-human":

return translate-to-a-human(sentence)

else if sentence contains "is-mortal":

return translate-mortal(sentence)

else if sentence contains "loves":

return translate-loves(sentence)

else if sentence contains "envy":

return trans-envy(sentence)

else if sentence contains "fele-exists" or "there-is":

return trans-exist(sentence)

else if sentence contains "not-all":

return trans-not-all(sentence)

else if sentence contains "if" question:

return trans-if-question(sentence)

else if sentence contains "and":

return trans-conjunction(sentence)

else

return "invalid"

function translate-to-human(sentence):

if sentence-matches-the-pattern "subject is-a-human"

sub = Extracted subject

return "H(subject)"

else

return "invalid"

function main():

while user input is not "exit":

get sentence from user

fol-translation = translate-to-fol(sentence)

main()

Lab-08

Implement unification in first order logic.

FUNCTION unify (expr1, expr2, substitutions):

if expr1 == expr2:

return substitutions

if is-variable(expr1):

return unify-variable(expr1, expr2, substitutions)

if is-variable(expr2):

return unify-variable(expr1, expr2, substitutions)

if is-compound(expr1) AND is-compound(expr2):

if expr1[0] != expr2[0] OR len(expr1[:]) != len(expr2[:]):

raise "UnificationError"

for each pair (arg1, arg2) in zip(expr1[1:], expr2[1:]):

substitutions = unify(arg1, arg2, substitutions)

return substitutions

raise "UnificationError"

function unify-variable(var, expr, substitutions):

if var == expr:

return true

if is-compound(expr):

for sub in expr:

if occurs-check(var, sub, substitutions):

return false

if expr[0] in substitutions:

return occurs-check(var, substitutions[expr], substitutions)

substitution)

1. Unify the following $P(x, a, b)$ and $P(y, z, b)$

Substitution required: $y \rightarrow z, z \rightarrow a$

Unified predicate is $P(z, a, b)$

2. $\forall x(\alpha, f(x); P(z, f(a)))$

Substitution required: $z \rightarrow a, y \rightarrow z$

Unified predicate is $P(y, f(a))$

3. To prove

i. Ravi enjoys a wide variety of foods

FOL: $\forall n \text{ Food}(n) \rightarrow \text{Enjoys}(\text{Ravi}, n)$

CNF: $\neg \text{Food}(n) \vee \text{Enjoys}(\text{Ravi}, n)$

ii. Bananas are food

FOL: $\text{Food}(\text{Banana})$

CNF: $\text{Food}(\text{Banana})$

iii. Pizza is food

$\text{Food}(\text{Pizza}) \rightarrow \text{FOL}(\text{CNF})$

iv. A food is anything that anyone consumes and isn't harmed by

FOL $\forall x (\exists y (\text{consumes}(y, x) \wedge \neg \text{HarmsBy}(y, x)))$

$\rightarrow \text{Food}(x)$

CNF $\exists y (\text{consumes}(y, x) \wedge \neg \text{HarmsBy}(y, x)) \vee \text{Food}(x)$

$\neg \text{consumes}(\text{Lem}, n) \vee \text{HarmfulBy}(\text{con}, n) \vee \text{Food}(n)$

5. Sam eats Idli and is still alive

FOL (NF): $\text{consumes}(\text{Sam}, \text{Idli}) \wedge \neg \text{HarmfulBy}(\text{con}, \text{Idli})$

6. Bill eats everything Sam eats

FOL: $\forall x \text{consumes}(\text{Sam}, x) \rightarrow \text{consumes}(\text{Bill}, x)$

CNF: $\neg \text{consumes}(\text{Sam}, n) \vee \text{consumes}(\text{Bill}, n)$

Goal: Ravi likes Idli: $\text{enjoys}(\text{Ravi}, \text{Idli})$

Proof

$\text{enjoys}(\text{Ravi}, \text{Idli})$

$\text{Food}(\text{Idli})$

$\text{consumes}(\text{Ravi}, \text{Idli}) \wedge \neg \text{HarmfulBy}(\text{Ravi}, \text{Idli})$

(P) Junktive

(S) Junktive

(Konditionale)

not food

spurkelt

(front) - will A

durch

Implementation of Forward Chaining

Representation in FOL and generation of proof tree by forward chaining:

Facts:

American (Robert) (or not) wanted to sell

Enemy (A, America)

Owns (A, T1) wants to sell weapons to them

Missile (T1) hostile with most others

Sells Weapon (Robert, T1)

Rules:

Missile (M) \rightarrow weapon (n)

Enemy (A, America) \rightarrow Hostile (x)

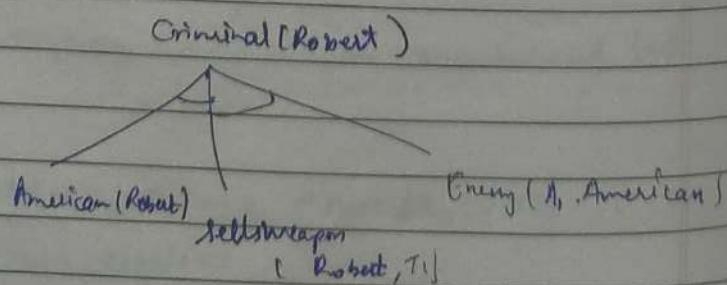
America (P) & Sells Weapon (P, Q) & Enemy (Q, America) \rightarrow
Criminal (P)

Goal: Criminal (Robert)

FOL:

$\forall P \forall Q (\text{American}(P) \wedge \text{SellsWeapon}(P, Q) \wedge \text{Enemy}(Q, \text{America})) \rightarrow$
Criminal (P)

\rightarrow Proof Tree



Thus: Criminal(Robert) \Rightarrow True.

Yes Robert is a Criminal.

(Implementation of Resolution.

\rightarrow Representation in FOL and generation of proof tree by resolution.

1. It's a crime for an American to sell weapons to hostile nations

$\forall p \forall q \forall r (\text{American}(p) \wedge \text{Weapon}(q) \wedge \text{Sells}(p, q, r) \wedge \neg \text{Hostile}(q) \rightarrow \text{Criminal}(p))$

- American(p) $\vee \neg \text{Weapon}(q) \vee \neg \text{Sells}(p, q, r) \vee \neg \text{Hostile}(q) \vee \neg \text{Criminal}(p)$

2. Country A has some missiles

$\exists n (\text{Country}(A, n) \wedge \text{Missile}(n))$

$\text{Owns}(n, T_1) \wedge \text{Missile}(T_1)$

3. All missiles were sold to country A by Robert

$\forall n \text{Missile}(n) \wedge \text{Owns}(A, n) \rightarrow \text{Sells}(\text{Robert}, n, A)$

$\neg \text{Missile}(n) \vee \neg \text{Owns}(A, n) \vee \neg \text{Sells}(\text{Robert}, n, A)$

4. Missiles are weapons

$\forall n (\text{Missile}(n) \Rightarrow \text{Weapon}(n))$

$\neg \text{Missile}(n) \vee \neg \text{Weapon}(n)$

5. Enemy of America is North Korea

$\forall n (\text{Enemy}(n, \text{America}) \Rightarrow \text{Hostile}(n))$

$\neg \text{Enemy}(n, \text{America}) \vee \neg \text{Hostile}(n)$

6. Robert is an American:
American (Robert)

7. Country A is an enemy of America:
Enemy (A, America)

Goal: Criminal (Robert)

Resolution: ~Criminal (Robert)

Proof Tree:

~American (P) V ~Weapon (Q) V ~Sells (P, Q, R) ?Criminal (Robert)

V ~Hostile (R) V ~Criminal (P) ?

American (P) V ~Weapon (Q) V ~Sells (P, Q, R) ?

(V, ~Hostile (R)) E

?American (P) V ~Weapon (Q) V ~Hostile (R) ?

?Missile (T) V ?Down (A, T) ?

- American (P) V ~Weapon (Q) V

? Hostile

?Weapon (Q) V ~Sells (P, Q, R) V ~Hostile (R) ?

missile (T) V

? Sells (P, Q, R) V ~Hostile (R) V ? Missile (T) ?

(V, ~Hostile (R)) E

(V, ~Hostile (R)) E

Analyze alpha-beta pruning method.

Function : alpha-beta (node, depth, alpha, beta, maximizing-player path):

IF depth is 0 OR node is a terminal node THEN
RETURN node's value, path

IF maximizing-player THEN

max-eval = negative infinity

optimal-path = null

for each child of node DO

child-value, child-path = alpha-beta prunij -

(child, depth-1, alpha, beta; FALSE, path +
child's name)

IF child-value > max-eval THEN

max-eval = child-value

optimal-path = child-path

alpha = maximum(alpha, max-eval)

IF beta <= alpha THEN,

break

return

else

min-eval = positive infinity

optimal-path = null

for each of node DO

child-value, child-path = alpha-beta prunij -

(child, depth-1, alpha, beta; TRUE,
path + child's name)

IF child-value < min-eval THEN

min-eval = child-value

optimal-path = child-path

beta = minimum(beta, min-eval)

IF beta <= alpha THEN

BREAK

RETURN min-level, optimal-path

maximizing-player = TRUE

initial-alpha = negative infinity

initial-beta = positive infinity

depth = 3

optimal value, optimal-path = alpha-beta(root node, depth, initial beta, maximizing-player)

PRINT "The optimal value is" optimal value

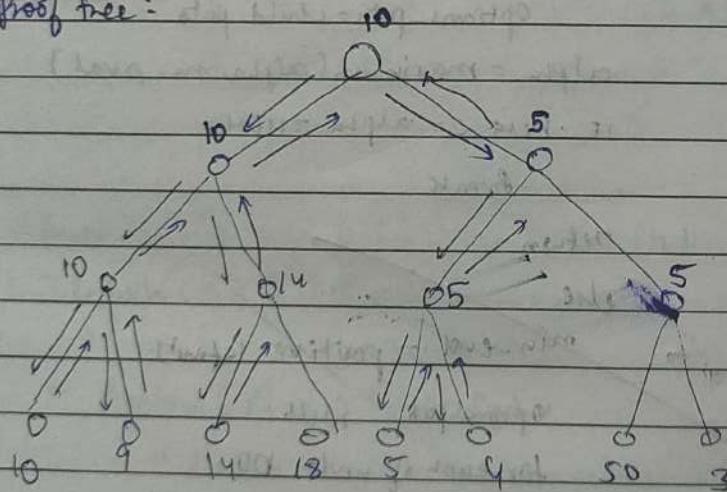
PRINT "The optimal path is" optimal path.

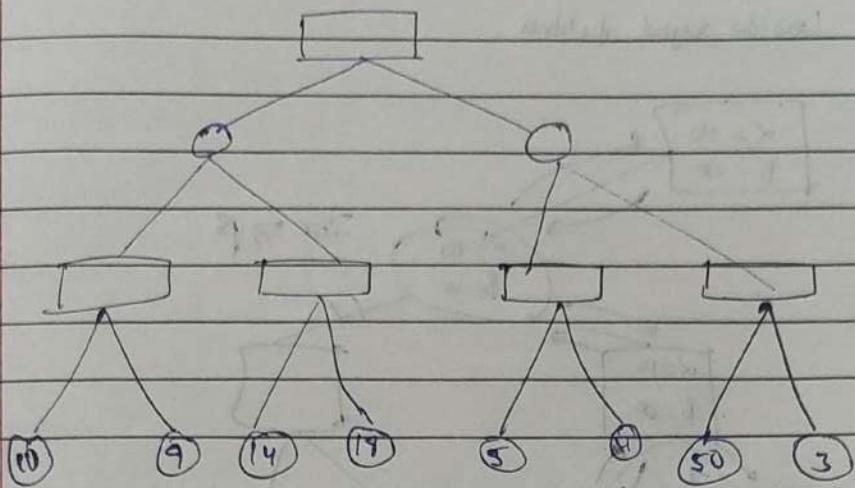
OUTPUT:

The optimal value is: 10

The optimal path is: A → B → D → H

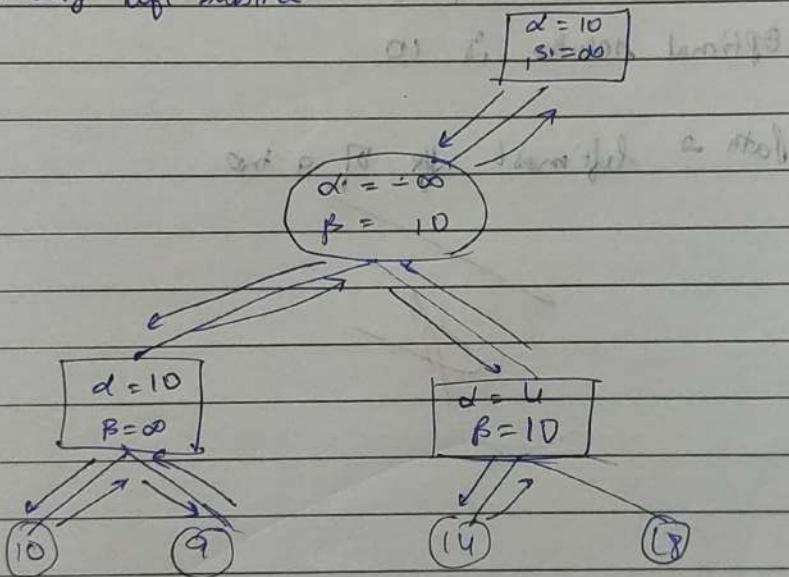
Proof tree:





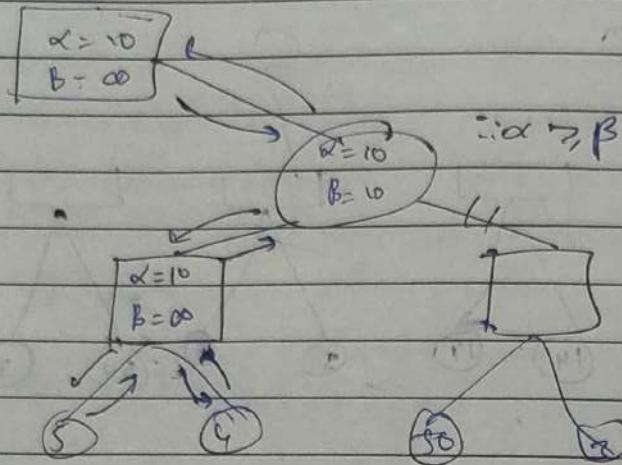
Applying DFS to left sub tree And initializing $\alpha = \infty$, $\beta = \infty$

considering left subtree



Gain
17/12/24

Lensitivity report analysis



Here $\alpha = 10, \beta = 0$

Optimal solution \rightarrow 10

Path \rightarrow left most off δ of a tree.