

# Algorithm Assignment 6

10185101210 陈俊潼

## 22.1-6

Observe a adjacency matrix with a universal sink:

	1	2	3	4
1	0	1	0	0
2	0	0	0	0
3	0	1	0	0
4	0	1	0	0

We found that when a vertex is an universal sink,

Let  $M$  be the adjacency matrix of size  $[V, V]$ . Start from point  $M[i, j]$ .

- If  $M[i, j] == 0$ , then check  $M[i, j+1]$
- If  $M[i, j] == 1$ , then check  $M[i + 1, j]$

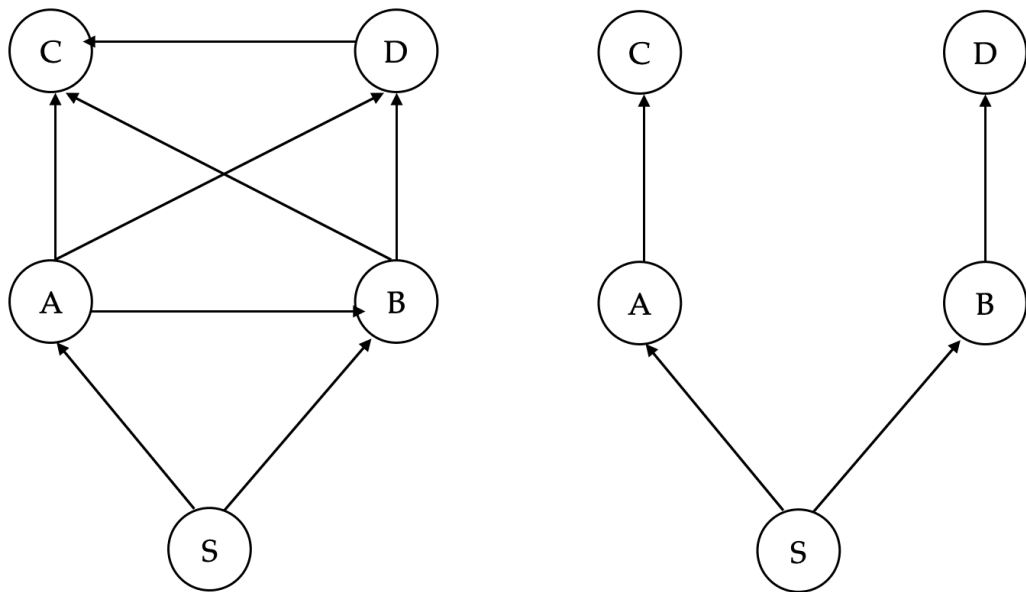
Repeat the process above, until when we found  $i == V$  or  $j == V$ . This might a universal sink.

Then check if the row is all 0. If true, then this is a universal sink. If such point can't be found when  $i$  or  $j$  out of boundary, there doesn't exist a universal sink.

The complexity of this algorithm is  $O(V) + 2O(V) = O(V)$ .

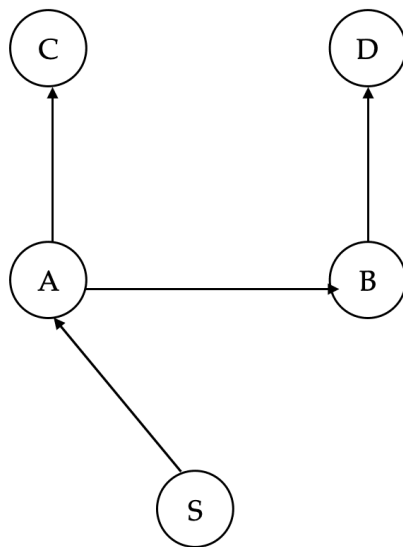
## 22.2-6

This example is shown as follows:

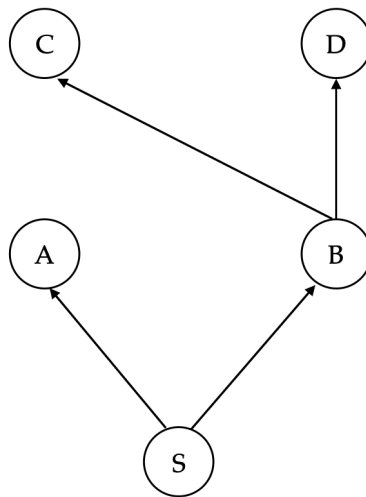


No matter how the vertex is ordered in the adjacency matrix, it's not possible to generate a tree structure as in the right image.

If A is first, then it will be:

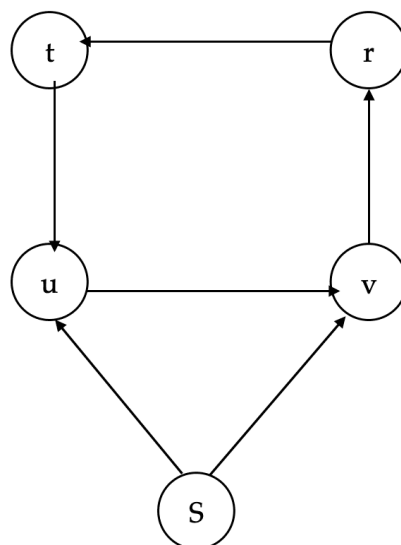


If B is first, then it will be:



## 22.3-9

An example is as follows:



When DFS visit the above graph, if v is visited first, then the steps are:

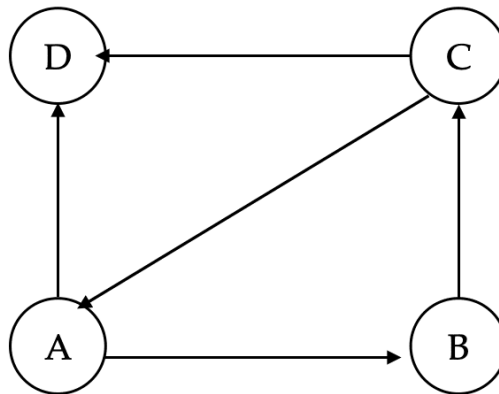
- v.d set
- r.d set
- t.d set
- u.d set
- u.f set
- t.f set
- r.f set

- $v.f$  set

In this case,  $u.f < v.d$ .

## 22.4-4

This is not true. Consider the following case:



Use `TOPOLOGICAL-SORT(A)` on this graph, assume D precedes B.

Now that we have the order BCAD, there are two bad edges: AB, and CD. However, the best case could be ABCD. where there's only one bad edge which is CA.

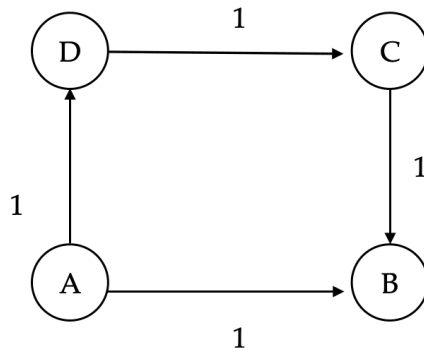
## 23.1-3

Because the edge  $(u, v)$  is already contained in some MST, so we can remove the edge  $(u, v)$  and then get two different MSTs. Where  $V_1$  is the vertices of the first MST and  $V_2$  the vertices of the second MST. In this case, we can cut the tree into two separate the tree into  $T_1$  and  $T_2$ . The edge  $(u, v)$  is then the light edge across  $T_1$  and  $T_2$ .

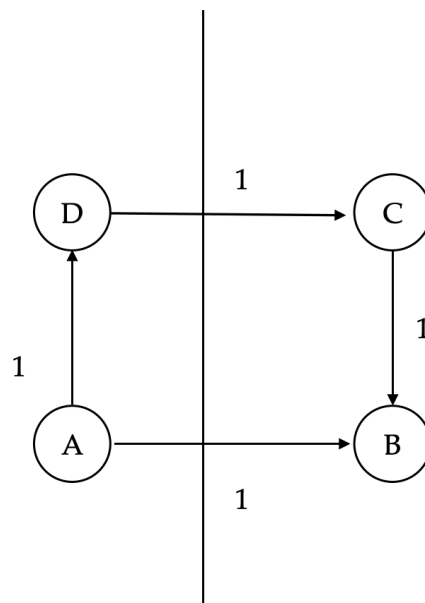
To further verify this, we assume edge  $(u, v)$  is not the light edge. So there would be another edge  $E$ , which is the light edge. But in this case, the former MST would be  $V_1$  and  $V_2$  and an edge  $E$  connecting them, contradicting to our given MST. So  $(u, v)$  must be the light edge.

## 23.1-4

The example is shown as follows:



Here for all the edges (AB, BC, CD, DA), there exists a cut (S, V-S) that (u, v) is a light edge. For example, for edge A->B:



The cut is {A, D}, here AB and DC are both light edges. The situation is the same for edge AD, BC, and CD. Hence the edge set will be all four edges, which of course, is not a minimum spanning tree of the original graph.

## 23.2-4

If the weights of the graph is in the range from 1 to  $|V|$ , then first we can use counting sort to sort all the edges, gives the complexity of  $O(V + E)$ . Then using the Kruskal's algorithm, it will cost  $O(V + E + V \lg V)$ .

If the weights of the graph is the range from 1 to  $W$ , then the different part is the time it costs in counting sort, which gives the complexity of  $O(W + E)$ . Then using the Kruskal's algorithm, it will cost  $O(W + E + V \lg V)$  time.

