

Algorithm Assignment 3

10185101210 陈俊潼

9.3-5

The black box median program is denoted as `MEDIAN(A, p, r)`. The pseudocode is shown as follows.

```
1  LINEAR_SELECTION(A, p, r, i)
2      if p == r
3          return A[p]
4      // Get the median number position
5      mp = MEDIAN(A, p, r)
6      // Partition with mp as pivot, which is a linear time solution
7      q = PARTITION(A, p, r, mp)
8      k = q - p + 1
9      if k == i:
10         return A[k]
11     else if i < k:
12         LINEAR_SELECTION(A, p, q - i, i)
13     else
14         LINEAR_SELECTION(A, q + 1, r, i - k)
```

9.3-7

Using the linear complexity algorithm `RANDOMIZED_SELECT` to find the median p of the array first, Then create a new array A' , all of which the value is $|a_n - p|$. Then find the k^{th} smallest element K , also using `RANDOMIZED_SELECT`. Finally, iterate the new array A' and select elements whose value is no greater than K .

The total complexity is $O(n) + O(n) + O(n) = O(n)$.

15.1-3

```

1 CUT_ROD_WITH_CUTTING_PRICE(p, n, c)
2   let r[0..n] be a new array
3   r[0] = 0
4   for j = 1 to n
5       q = -inf
6       for i = 1 to j - 1
7           q = max(q, p[i] + r[j-1] - c)
8       r[j] = q
9   else
10

```

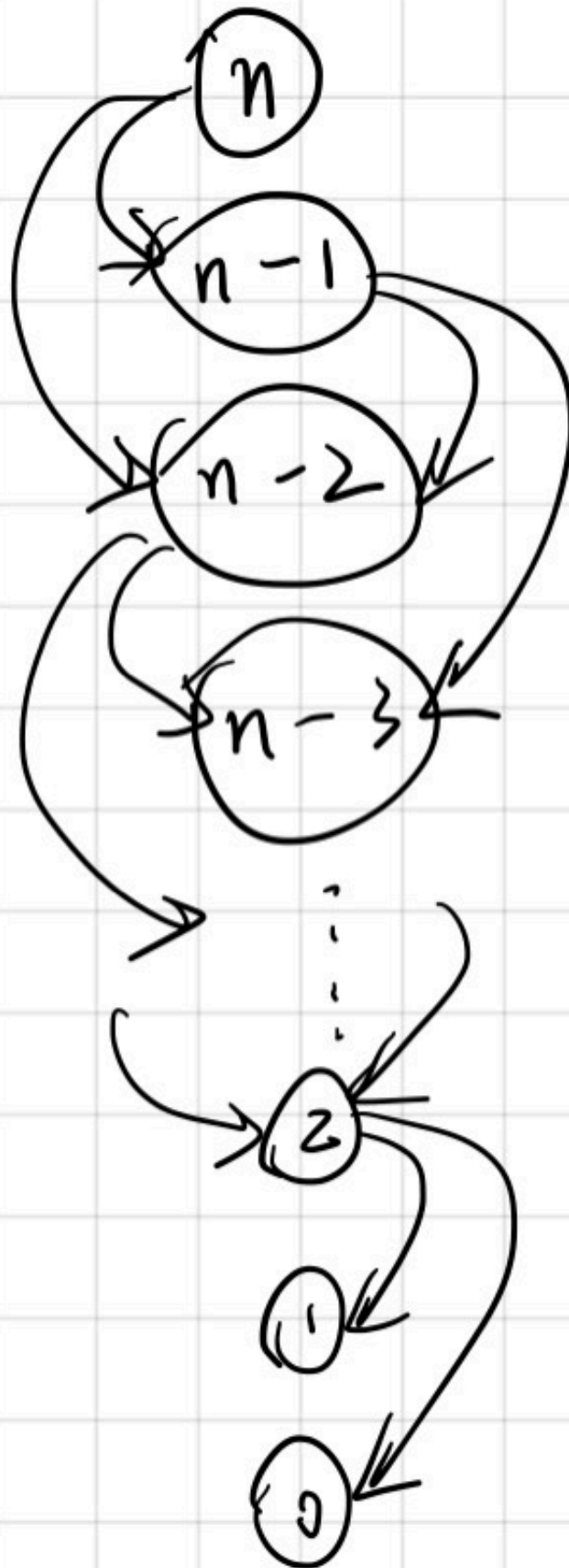
15.1-5

```

1 FIBONACCI(n)
2   let n[0..n] be an array
3   n[0] = 0
4   n[1] = 1
5   if n <= 1
6       return 1
7   else
8       for i = 2 to n
9           n[i] = n[i - 1] + n[i - 2]
10      return n[i]

```

The subproblem graph is denoted as follows:



In which there are $n + 1$ vertices and $2 * (n-1)$ edges.

