748. Shortest Completing Word String Hash Table Array **Leetcode Link** Easy

Problem Description In this LeetCode problem, we are given a string licensePlate and an array of strings words. Our task is to find the shortest word in

words that contains all the letters in licensePlate, disregarding any numbers and spaces, and treating letters as case-insensitive. Notably, if a letter repeats in the licensePlate, it must appear at least as many times in the word. The essence of the problem is to identify the most concise word from an array that encapsulates all alphabetic characters of the given license plate string.

To arrive at the solution for this problem, we breakdown the requirements into smaller steps and handle each part. First, we create a function called count that converts a given word into a frequency array that represents how many times each letter appears in the word.

Intuition

licensePlate.

Once we have count, we need a way to determine if a word in words can be considered a "completing word". To do this, we craft a check function that compares two frequency arrays: one from the cleaned licensePlate and one from a candidate word. The check function goes over each letter's frequency and ensures the candidate word has equal or more occurrences of each letter found in

Then we iterate over each word in words and apply the count function to it. If it passes the check with the licensePlate's frequency array, it means we've found a potential completing word. We maintain a variable for the shortest word found (ans) and its length (n). If a new completing word is shorter than the current best, we update ans with this new word. This way, by prioritizing words that meet the completing criteria and are shorter than any others we've seen, we can solve the problem effectively.

The solution approach for finding the shortest completing word involves the use of frequency arrays and efficient checks for each word in the words array against the licensePlate. Here's a step-by-step explanation of the implementation:

We define a count function that takes a word and transforms it into a frequency array, which is an array of 26 integers

and so on up to 'z'.

The count function

Solution Approach

ord function is used to get the ASCII value of a character, and we subtract the ASCII value of 'a' to map 'a' to index 0, 'b' to index 1, 1 def count(word):

(corresponding to the 26 letters of the English alphabet), each element representing the count of a specific letter from 'a' to 'z'. The

The check function

counter[ord(c) - ord('a')] += 1 return counter

The check function compares two frequency arrays: one from the license plate (counter1) and the other from the current word

The main part of the algorithm is a loop through each word in words. It initializes two variables: ans, to store the currently found

(counter2). If for any letter the count in the license plate is greater than the count in the word, check returns False indicating that the

word is not a completing word. Otherwise, if all letter counts are equal or higher in the word than in the license plate, it returns True. 1 def check(counter1, counter2): for i in range(26): if counter1[i] > counter2[i]: return False return True

if check(counter, t): n = len(word) ans = word

The main loop

2 ans, n = None, 16

for word in words:

if n <= len(word):</pre>

continue

t = count(word)

Example Walkthrough

there's one 'a', one 'b', and one 'c'.

but it's longer than "stripe".

and 'c', so it is a potential completing word.

counter = [0] * 26

for char in word:

for i in range(26):

return counter

return True

for word in words:

continue

Iterating Over Each Word

'a', 'b', and 'c'.

problem.

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C++ Solution

1 #include <string>

2 #include <vector>

5 class Solution {

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return true;

// Helper function to check if a character is an alphabetic letter

words list and the average length of each word. Here is the breakdown:

return (code >= 65 && code <= 90) || (code >= 97 && code <= 122); // A-Z or a-z

function isAlphabetic(char: string): boolean {

const code: number = char.charCodeAt(0);

Python Solution

shortest completing word, and n, to store its length.

The process for each word in the loop is as follows:

and optimizing the process of determining completing words.

1 counter = count(c.lower() for c in licensePlate if c.isalpha())

2. We generate a frequency array of the word (t) using count. 3. We use check to see if t meets the requirements compared to counter, which is the frequency array for licensePlate. 4. If it does, this is the new shortest completing word, and we update ans with this word and n with its length.

The loop proceeds until all words have been checked, and the shortest completing word (ans) is returned as the answer.

1. We skip the word if its length is not less than the current shortest completing word's length (n) to optimize performance.

To illustrate the solution approach, let's take a small example. Suppose licensePlate is "aBc 1123" and words is ["step", "steps", "stripe", "stepple"]. The first step would be to create a frequency array for the cleaned licensePlate, which disregards numbers and spaces and is caseinsensitive. This means "aBc" becomes "abc" and our count function generates an array [1, 1, 1, 0, ..., 0], indicating that

This algorithm effectively leverages data structures (arrays) and functions to compare character frequencies, significantly simplifying

 For "step", we get the frequency array [0, 0, 0, 0, 1, 0, ..., 1], which has an 'e' and a 'p', but lacks 'a' and 'b'. • For "steps", the frequency array is [0, 0, 0, 0, 1, 0, ..., 1, 1]. This word also covers 's', but still misses 'a' and 'b'. • "stripe" has a frequency array [0, 0, 0, 0, 1, 0, ..., 1, 1, 1], containing 'e', 'p', 'r', 's', 't', and 'i', and it does have at least one

We then iterate over each word in words and check if it's a completing word for licensePlate.

Initialize a counter for 26 letters of the alphabet.

Increment the letter's corresponding counter.

Helper function to check if the word contains all the

Iterate through each letter in the alphabet.

if license_counter[i] > word_counter[i]:

then it's not a completing word.

def contains_needed_letters(license_counter, word_counter):

If the license has more letters than the word,

Initialize the variables for tracking the shortest completing word.

min_length = 16 # Given constraint - words length doesn't exceed 15.

Count the frequency of letters in the current word.

with at least the frequencies in the license plate.

Check if the current word has all the required letters

if contains_needed_letters(license_counter, word_counter):

Iterate through the words list to find the shortest completing word.

Skip the iteration if the current word is not shorter than min_length.

counter[ord(char) - ord('a')] += 1

needed letters with required frequency.

return False

shortest_completing_word = None

if min_length <= len(word):</pre>

word_counter = count_letters(word)

// Return the shortest completing word found

// Counts the frequency of each letter in a given word

// Iterate through each character in the word

// Increment the counter for the letter if it's a letter

// Checks if wordCounter covers all the letters in licensePlateCounter

if (licensePlateCounter[i] > wordCounter[i]) {

// If all letter counts are covered, return true

// the word does not cover the license plate, return false

private boolean doesWordCoverLicensePlate(int[] licensePlateCounter, int[] wordCounter) {

// Finds the shortest word in 'words' that contains all the letters in 'licensePlate'

string shortestCompletingWord(string licensePlate, vector<string>& words) {

// If current word is longer than the shortest found, skip it

// Check if the word contains all letters in the license plate

// Update minimum length and answer if this word is shorter

// Checks if 'counterSource' contains at least as many of each letter as 'counterTarget'

bool containsAllLetters(vector<int>& counterSource, vector<int>& counterTarget) {

// If the target has more of a letter than the source, return false

// Count the frequency of each letter in the license plate

vector<int> licenseCounter = countLetters(licensePlate);

// Count the frequency of each letter in the word

if (containsAllLetters(licenseCounter, wordCounter)) {

vector<int> wordCounter = countLetters(word);

// If the current letter's count in licensePlateCounter exceeds that in wordCounter,

return shortestCompletingWord;

private int[] countLetters(String word) {

int[] letterCounter = new int[26];

for (char c : word.toCharArray()) {

if (Character.isLetter(c)) {

// Iterate over each letter count

for (int i = 0; i < 26; ++i) {

return false;

#include <cctype> // For isalpha and tolower

int minLength = INT_MAX;

for (auto& word : words) {

// Iterate over each word in the list

minLength = word.size();

// Counts the frequency of each letter in a word

++counter[tolower(c) - 'a'];

if (counterSource[i] > counterTarget[i]) {

vector<int> countLetters(string& word) {

for (int i = 0; i < 26; ++i) {

return false;

shortestWord = word;

if (minLength <= word.size()) continue;</pre>

string shortestWord;

return shortestWord;

return counter;

return letterCounter;

return true;

letterCounter[c - 'a']++;

min_length = len(word)

Finding the Shortest Completing Word Since "stripe" is shorter than "stepple" and it contains all the required letters (ignoring the unnecessary 'i', 'r', and 't'), "stripe" becomes our ans and its length, 6, becomes n.

As no other word in words is both a completing word for "aBc 1123" and shorter than "stripe", "stripe" is the answer we return. It

successfully encapsulates all alphabetic characters of the license plate "aBc 1123" in the most concise form found in the list words.

This example demonstrates how the count and check functions, together with an iterating loop, can be used to efficiently solve the

• Finally, "stepple" has the array [0, 0, 0, 0, 1, 0, ..., 2, 2, 1], where there is more than one occurrence of some letters,

Through our check function, we see that only "stripe" has equal or more occurrences of the letters in licensePlate, which are 'a', 'b',

class Solution: def shortestCompletingWord(self, licensePlate: str, words: List[str]) -> str: # Helper function to count the frequency of each letter in a word. def count_letters(word):

23 24 # Transforms the license plate into lowercase and filters out non-alphabetic characters. 25 # Then, counts the frequency of each letter in the license plate. 26 license_counter = count_letters(char.lower() for char in licensePlate if char.isalpha())

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                     shortest_completing_word = word
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             return shortest_completing_word
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Java Solution
  1 class Solution {
         // Finds the shortest completing word in words that covers all letters in licensePlate
         public String shortestCompletingWord(String licensePlate, String[] words) {
             // Count the frequency of letters in the license plate, ignoring case and non-letters
             int[] licensePlateCounter = countLetters(licensePlate.toLowerCase());
  6
             // Initialize the answer variable to hold the shortest completing word
             String shortestCompletingWord = null;
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  9
             // Initialize the shortest length found so far to an upper bound
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             int minLength = 16; // As per the problem, words are at most 15 letters long
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             // Iterate through each word in the array
             for (String word : words) {
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                 // Skip the word if its length is not less than the shortest length found so far
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                 if (minLength <= word.length()) {</pre>
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                     continue;
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 20
                 // Count the frequency of letters in the current word
                 int[] wordCounter = countLetters(word);
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                 // Check if the current word covers all letters in the license plate
 24
                 if (doesWordCoverLicensePlate(licensePlateCounter, wordCounter)) {
 25
                     // Update the shortest completing word and the minimum length found so far
                     minLength = word.length();
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 27
                     shortestCompletingWord = word;
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```

Update min_length and shortest_completing_word with the current word's length and the word itself.

35 vector<int> counter(26, 0); // Initiate a counter for 26 letters 36 37 // Increment the count for each letter in the word 38 for (char& c : word) { 39 if (isalpha(c)) { 40

private:

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             // If the source contains at least as many of each letter, return true
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             return true;
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 57 };
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Typescript Solution
    function shortestCompletingWord(licensePlate: string, words: string[]): string {
      // Count the frequency of each letter in the license plate
       const licenseCounter: number[] = countLetters(licensePlate);
       let minLength: number = Number.MAX_SAFE_INTEGER;
       let shortestWord: string = '';
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       // Iterate over each word in the list
       for (let word of words) {
  8
        // If the current word is longer than the shortest so far, skip it
  9
 10
         if (minLength <= word.length) continue;</pre>
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 12
         // Count the frequency of each letter in the current word
 13
         const wordCounter: number[] = countLetters(word);
 14
         // Check if the word contains all the letters in the license plate
 15
 16
         if (containsAllLetters(licenseCounter, wordCounter)) {
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           // Update minimum length and shortest word if this word is shorter
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           minLength = word.length;
           shortestWord = word;
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       return shortestWord;
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    // Helper function to count the frequency of each letter in a word
     function countLetters(word: string): number[] {
       const counter: number[] = new Array(26).fill(0); // Initialize a counter for 26 letters
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      // Increment the count for each alphabet letter in the word
       for (let char of word) {
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        if (isAlphabetic(char)) {
           counter[char.toLowerCase().charCodeAt(0) - 'a'.charCodeAt(0)]++;
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       return counter;
 36
 37
    // Helper function to verify if sourceCounter contains at least
 39 // as many of each letter as targetCounter
    function containsAllLetters(sourceCounter: number[], targetCounter: number[]): boolean {
       for (let i = 0; i < 26; ++i) {
 42
        // If the target has more of any particular letter than the source, return false
        if (sourceCounter[i] < targetCounter[i]) {</pre>
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           return false;
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      // If source contains at least as many of each letter, return true
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• The count function has a time complexity of O(k) where k is the length of the word. It iterates once through all the characters of the word. • This count function is called for every character in the normalized licensePlate once, resulting in a time complexity of O(m) where m is the number of alphabetic characters in licensePlate.

constant size.

In summary:

Time and Space Complexity

and k is the average length of the words. The function check has a time complexity of 0(1) because it checks a fixed size array (26 letters of English alphabet). • Since check is called for every word, we multiply it by n, but it doesn't affect the overall time complexity significantly due to the

• The counter arrays used to store character frequencies for the licensePlate and each word. Since these are of fixed size (26),

• For each word in words, the count function is called again, giving a time complexity of 0(k * n) where n is the number of words

The time complexity of the shortestCompletingWord function primarily depends on two factors: the number of words in the input

- Combining the above steps, the total time complexity for the function is 0(m + k * n). The space complexity is determined by:
- The space used to store the answer and the length of the shortest word found so far (n and ans) is 0(1). Thus, the overall space complexity is 0(1) because it does not scale with the size of the input.

they occupy 0(1) space.

 Time Complexity: 0(m + k * n) • Space Complexity: 0(1)