1480. Running Sum of 1d Array

Easy | Array | Prefix Sum | Leetcode Link

Problem Description

In this problem, you're given an array `nums` which contains a list of integers. The task is to calculate a *running sum* of this array. The running sum is a new array where each element at index `i` is the sum of all the numbers in the `nums` array up to and including the number at index `i`.

More formally, the running sum of `nums` is an array where the value at the `i-th` index is equal to `sum(nums[0]…nums[i])`. So for `runningSum[i]`, we add up all the elements from the start of the array to the current element at `i`.

Your goal is to return the running sum of the given array `nums`.

Intuition

To solve this problem, we can use a common algorithm called *Prefix Sum*. The idea behind Prefix Sum is simple – as we move through the array element by element, we keep updating the current element to be the sum of itself and all previous elements.

So how do we arrive at this solution? Let's think about the process step by step. If we were doing this by hand, we'd start at the first element – nothing to add here since it's the start. We note this number down. Then, we'd move to the next number, add it to the first, and note this new sum down. Then, for the third number, we'd add it to the new sum we just computed, and so on.

This is exactly what the Prefix Sum algorithm does. We start with the first element, then for each subsequent element at index `i`, we add the current element `nums[i]` to the sum we have obtained so far, which is the previous element `nums[i-1]`. This updated sum becomes the new value for `nums[i]`.

By following this approach, we only need to traverse the array once, making our algorithm efficient and straightforward.

Solution Approach

The solution provided uses a built-in Python function `accumulate` from the `itertools` module, which essentially applies the Prefix Sum algorithm automatically for us.

Here's the step by step approach using the Prefix Sum concept:

1. Initialize a variable, let's call it `prefixSum`, to store the sum of elements we've encountered so far. Initially `prefixSum` is set to 0.
2. Traverse the array `nums` from the first to the last element.
3. For each element at index `i`, update `prefixSum` by adding the current element `nums[i]` to it. This will hold the sum of all elements up to the current index.
4. Assign the value of `prefixSum` back to `nums[i]` to reflect the running sum till the current element.
5. Repeat steps 3-4 for the entire length of the array.
6. Once the loop ends, we'd have transformed the original `nums` array into the running sum array.

Now, in the provided solution, the step-by-step process is neatly wrapped up by Python's `accumulate()` function. This function takes an iterable (in our case, the array `nums`) and returns an iterator that yields accumulated sums (or other binary functions if specified).

The solution could also be implemented more explicitly, without the use of `accumulate()`, which would look something like this:

```
class Solution:
    def runningSum(self, nums: List[int]) -> List[int]:
        for i in range(1, len(nums)):
            nums[i] += nums[i-1]
        return nums
```

In the above explicit implementation:

- We use a `for` loop to traverse the array starting from the second element (since the first element's running sum is the element itself).
- We then continuously update each element with the sum of itself and all previous elements. This is done by adding the current element `nums[i]` with the previous element `nums[i-1]` and storing the result back in `nums[i]`.
- After traversing through the array, we return the `nums` array which now contains the running sum.

Both approaches accomplish the task using the powerful concept of Prefix Sum, but one is more succinct by leveraging Python's built-in functionality.

Example Walkthrough

Let's walk through a small example to make the solution approach clear. Suppose the `nums` array is `[3, 1, 2, 10]`.

Here are the steps we would take to get the running sum:

1. Start with the first element. Since there are no elements before it, the running sum is just the element itself. So `runningSum[0]` would be `3`.
2. Move to the second element, `1`. We add this to the previous running sum (`3`), giving us `4`. So `runningSum[1]` is `3 + 1` which equals `4`.
3. Next, the third element is `2`. Adding this to the running sum we have so far (`4`), we get `6`. Now, `runningSum[2]` is `4 + 2` which equals `6`.
4. Finally, take the fourth element, `10`. We add this to the last running sum (`6`), which gives us `16`. Therefore, `runningSum[3]` is `6 + 10` which equals `16`.

The final output, our running sum array, is `[3, 4, 6, 16]`.

Implementing this in Python code without the use of `accumulate()` from the `itertools` module would look like this:

```
class Solution:
    def runningSum(self, nums):
        for i in range(1, len(nums)):
            nums[i] += nums[i-1]
        return nums
```

If we use this code with our example `nums` array `[3, 1, 2, 10]`, here's what happens step by step in the loop:

- `i = 1`: `nums[1]` (which is `1`) is updated to `nums[1] + nums[0]` (which is `1 + 3`), so `nums[1]` becomes `4`.
- `i = 2`: `nums[2]` (which is `2`) is updated to `nums[2] + nums[1]` (which is `2 + 4`), so `nums[2]` becomes `6`.
- `i = 3`: `nums[3]` (which is `10`) is updated to `nums[3] + nums[2]` (which is `10 + 6`), so `nums[3]` becomes `16`.

And we get the final updated `nums` array `[3, 4, 6, 16]`, which is the same result we calculated manually.

Python Solution

```
from itertools import accumulate  # Import accumulate function from itertools module

class Solution:
    def runningSum(self, nums: List[int]) -> List[int]:
        # Calculate the running sum of the list of numbers using accumulate
        running_sum = list(accumulate(nums))
        # Return the running sum as a list
        return running_sum

```

Java Solution

```
class Solution {

    // Method to calculate the running sum of the given array
    public int[] runningSum(int[] nums) {
        // Loop through the array starting from the second element
        for (int index = 1; index < nums.length; index++) {
            // Update the current element by adding the previous element's sum to it
            nums[index] += nums[index - 1];
        }
        // Return the modified array containing the running sum
        return nums;
    }
}

```

C++ Solution

```
#include <vector> // Include vector library for using the vector class

class Solution {
public:
    // Function to calculate the running sum of a 1D vector
    vector<int> runningSum(vector<int>& nums) {
        // Iterate over the vector starting from the second element
        for (size_t i = 1; i < nums.size(); ++i) {
            // Add the previous element's value to the current element
            nums[i] += nums[i - 1];
        }
        // Return the modified vector with the running sum
        return nums;
    }
};

```

Typescript Solution

```
/**
 * Computes the running sum of an array of numbers.
 * For every index in the array, it modifies the value at that index
 * to be the sum of all elements up to and including that index.
 *
 * @param {number[]} numbers - The array of numbers to calculate the running sum for.
 * @returns {number[]} - The array of numbers with the running sum computed.
 */
function runningSum(numbers: number[]): number[] {
    // Start iterating from the first index because the running sum at index 0 is just the element itself
    for (let index = 1; index < numbers.length; ++index) {
        // Update the element at the current index to be the sum of the element at the current index
        // and the element at the previous index which now contains the running sum up to the previous index
        numbers[index] += numbers[index - 1];
    }
    // Return the modified array with the running sums
    return numbers;
}

```

Time and Space Complexity

The time complexity of the code is `O(n)`, where `n` is the length of the input array. This is because `accumulate` function goes through each element of the array only once to calculate the running sum.

The space complexity in the given code should be `O(n)`, not `O(1)` as the reference answer suggests. This difference arises because although `accumulate` itself may operate with `O(1)` additional space, calling `list(accumulate(nums))` stores the results of `accumulate(nums)` in a new list, which takes up `O(n)` space based on the number of elements in `nums`.