



Problem Description

In this problem, we are given an array nums which contains distinct integers. It is 0-indexed, which means the indexing of elements in the array starts with 0. Among all the elements in the nums array, there is one element with the lowest value, termed as the minimum, and one with the highest value, known as the maximum. Our objective is to remove these two elements from the array.

A deletion involves removing an element either from the front (beginning) of the array or from the back (end) of the array. The task is to determine the minimum number of such deletions required to remove both the minimum and maximum elements from the array.

To succeed, we need to find a strategy that minimizes the total number of deletions by intelligently deciding whether to start from the front or the back of the array, as well as considering the possibility of removing elements from both ends.

Intuition

idea hinges upon the realization that the number of deletions needed is influenced by the relative locations of these two elements. Firstly, we identify the indexes of both the minimum and maximum elements by scanning through the array.

When approaching the solution, we focus on the positions (indexes) of the minimum and maximum elements in the array. The main

Then, our strategy involves considering three scenarios:

Remove elements from the front until both the minimum and maximum elements are deleted.

- Remove elements from the back until both elements are deleted. • Remove some elements from the front to delete one of the elements and some from the back to delete the other one.
- To achieve the minimum number of deletions, we calculate the number of deletions needed for each of the scenarios and return the smallest one.

The insight here is to use the indexes of minimum and maximum elements efficiently to minimize the number of operations, as direct removals would be inefficient for large arrays.

Solution Approach

The solution approach for this problem leverages a straightforward scan of the array to determine the positions (indexes) of the minimum and maximum elements.

found.

The following steps are taken in the implementation:

1. Initialize two variables, mi and mx, to store the indexes of the minimum and maximum elements, respectively. We start by

assuming the first element (at index 0) is both the minimum and maximum, hence mi = mx = 0.

2. Loop through every element num in the nums list, keeping track of the current index i. During each iteration:

- If the current element is less than nums [mi], update mi with the current index i because a new minimum has been found. Similarly, if the current element is greater than nums [mx], update mx with the current index i as a new maximum has been
- 3. After identifying the indexes of the minimum and maximum elements, ensure that mi is less than or equal to mx by swapping them if that is not the case. This simplifies the calculation that follows, because we will be dealing with three scenarios
- represented by contiguous sections of the array to be deleted. 4. Calculate the total number of deletions required for each of the three above-mentioned deletion strategies:

Strategy 1: Remove elements from the beginning (front) up to and including the maximum index (mx). The number of

- deletions would be mx + 1, since array indexing starts at 0. Strategy 2: Remove elements from the end (back) up to and including the minimum index (mi). The number of deletions here
 - would be len(nums) mi. Strategy 3: Remove elements from the front up to and including the minimum index (mi + 1) and from the back to the
- maximum index (len(nums) mx), combining these two gives mi + 1 + len(nums) mx. 5. Return the minimum value out of the three calculated strategies to achieve the minimum number of deletions needed to remove
- mentioned three results as arguments. The Python code is elegant and concise due to its use of list enumeration for index tracking, conditional statements for dynamic

both the minimum and maximum elements from the array. This is done using the built-in min() function with the above

Example Walkthrough Let's illustrate the solution approach with a small example:

1. We initialize mi = mx = 0. Currently, nums[mi] = nums[mx] = 3.

Python Solution

class Solution:

10

11

12

13

14

15

16

17

15

16

17

18

19

20

21

22

23

24

25

26

27

28

29

30

31

33

34

35

36

37 }

2. We start looping through the elements:

Assume we have the array nums = [3, 2, 5, 1, 4].

updates, and the min() function for direct comparison.

o num = 2, i = 1: 2 is not less than 3, so mi remains unchanged. 2 is not greater than 3, so mx also remains unchanged. \circ num = 5, i = 2:5 is greater than 3, so we update mx to 2.5 is not less than 3, so mi remains unchanged.

```
o num = 4, i = 4: 4 is not less than 1, and it's not greater than 5, so both mi and mx remain unchanged.
After the loop, mi = 3 (minimum element 1 at index 3) and mx = 2 (maximum element 5 at index 2).
 3. Since mi is not less than mx, we do not swap them.
```

∘ Strategy 2: Removing from the back to the minimum index mi is len(nums) - mi, which is 5 - 3, resulting in 2 deletions. ○ Strategy 3: Removing from the front up to mi and from the back beyond mx is mi + 1 + len(nums) - mx, which is 3 + 1 + 5

 \circ num = 1, i = 3: 1 is less than 3, so we update mi to 3. 1 is not greater than 5, so mx remains 2.

Strategy 1: Removing from the front to the maximum index mx is mx + 1, which equals 3 deletions.

2, totaling 7 deletions. 5. Finally, return the minimum value out of the calculated strategies: min(3, 2, 7) which equals 2.

4. Calculate the number of deletions for each strategy:

Therefore, the minimum number of deletions required to remove both the minimum and maximum elements from the array [3, 2, 5,

Initialize indices for the minimum and maximum value elements

if num > nums[max_index]: # A new maximum element found

Ensure that min_index is always less than or equal to max_index

Calculate the minimum number of deletions using three strategies:

// Update the index of the maximum element if a larger element is found

// Make sure the minIndex is always less than maxIndex for ease of calculation

// Calculate the minimum deletions needed, considering three different scenarios:

// 1. Removing elements from the beginning up to the maxIndex

minIndex + 1 + n - maxIndex);

// 3. Removing elements from both ends surrounding the min and max

// 2. Removing elements from the minIndex to the end

return Math.min(Math.min(maxIndex + 1, n - minIndex),

// Return the smallest of these three options

1, 4] is 2. This occurs by removing the last two elements (strategy 2).

Iterate over the list to find indices of the minimum and maximum elements for i, num in enumerate(nums): if num < nums[min_index]: # A new minimum element found</pre> min index = i9

min_index, max_index = max_index, min_index

if (nums[i] > nums[maxIndex]) {

// Swap minIndex with maxIndex

maxIndex = i;

if (minIndex > maxIndex) {

int temp = maxIndex;

maxIndex = minIndex;

minIndex = temp;

def minimumDeletions(self, nums: List[int]) -> int:

min_index = max_index = 0

 $max_index = i$

if min_index > max_index:

```
# 1. Delete from the front to the max element
19
           # 2. Delete from the min element to the end of the list
           # 3. Delete from both ends to the min and max elements
20
           # Then return the strategy that results in the fewest deletions
21
22
           return min(
23
               max_index + 1, # Deleting from the front to the max element
                len(nums) - min_index, # Deleting from the min element to the end
24
25
               min_index + 1 + len(nums) - max_index # Deleting from both ends
26
27
Java Solution
 1 class Solution {
       public int minimumDeletions(int[] nums) {
           // Initialize the variables to store the index of minimum and maximum element in array
           int minIndex = 0;
           int maxIndex = 0;
           // Length of the array
           int n = nums.length;
           // Iterate through the array to find the indices of the minimum and maximum elements
 9
10
           for (int i = 0; i < n; ++i) {
               // Update the index of the minimum element if a smaller element is found
11
               if (nums[i] < nums[minIndex]) {</pre>
12
13
                   minIndex = i;
14
```

```
38
C++ Solution
 1 class Solution {
2 public:
        int minimumDeletions(vector<int>& nums) {
            int minIndex = 0; // Index of the smallest number
            int maxIndex = 0; // Index of the largest number
            int n = nums.size(); // The total number of elements in nums
           // Loop to find indices of the smallest and largest numbers
           for (int i = 0; i < n; ++i) {
                if (nums[i] < nums[minIndex]) { // Check for new minimum</pre>
                    minIndex = i;
11
12
13
               if (nums[i] > nums[maxIndex]) { // Check for new maximum
14
                   maxIndex = i;
15
16
17
18
           // Ensure minIndex is smaller than maxIndex for easier calculation
19
           if (minIndex > maxIndex) {
20
               int temp = minIndex;
               minIndex = maxIndex;
21
22
               maxIndex = temp;
23
24
25
           // Three possible ways to remove the min and max, take the minimum deletions needed
26
            return min(
               min(
27
                    maxIndex + 1, // Delete all from start to maxIndex, inclusive
                    n - minIndex // Delete all from minIndex to end
29
30
31
               minIndex + 1 + n - maxIndex // Delete from start to minIndex and from maxIndex to end
32
           );
33
34 };
35
```

13 14 15

9

10

Typescript Solution

function minimumDeletions(nums: number[]): number {

// Find the number of elements in the array

let minIndex = nums.indexOf(Math.min(...nums));

let maxIndex = nums.indexOf(Math.max(...nums));

// If the array contains only one element, one deletion is required

// Find the indices of the minimum and maximum elements in the array

const numElements = nums.length;

if (numElements === 1) return 1;

```
11
       // Calculate the leftmost and rightmost positions among the min and max indices
12
       let leftIndex = Math.min(minIndex, maxIndex);
       let rightIndex = Math.max(minIndex, maxIndex);
       // Calculate the possible numbers of deletions
16
       // if removing from the left towards the right
17
       let deleteLeftRight = leftIndex + 1 + numElements - rightIndex;
18
       // if removing from the left only
19
       let deleteLeftOnly = rightIndex + 1;
20
       // if removing from the right only
       let deleteRightOnly = numElements - leftIndex;
22
23
24
       // Return the minimum number of deletions amongst the calculated options
       return Math.min(deleteLeftRight, deleteLeftOnly, deleteRightOnly);
25
26 }
27
Time and Space Complexity
The time complexity of the given code is O(n), where n is the length of the input list nums. This is because the code consists of a
```

single loop that traverses the entire list to find the minimum and maximum elements' indices, which takes linear time in the size of the input.

The space complexity of the code is 0(1). The extra space used by the code includes only a fixed number of integer variables mi

(minimum index) and mx (maximum index), which do not depend on the input size, therefore the space used is constant.