400. Nth Digit Medium Math Binary Search

Problem Description

The problem asks us to return the n^th digit of the concatenation of all positive integers in order. The sequence starts with 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, ... and goes on infinitely. Each number in the sequence is laid out consecutively to form an infinitely long number string. For example, if n equals 3, the function should return 3, as 3 is the third digit of the sequence. If n equals 11, we should return 0, which is the 11th digit in the sequence.

Intuition

- Going digit by digit would be quite inefficient for large n, as the numbers grow larger and the sequence quickly becomes very long. A better method is to calculate the position of the target digit. Let's break the problem into manageable parts: Determine the length of the number k where the n^th digit lies. Numbers with the same number of digits are in the same
- block (e.g., the number 12 belongs to the block where each number has two digits 10 to 99). Calculate the actual number where our digit is present. Once we know the block where the digit is, we can compute the
- number itself. Determine the exact digit within that number which is our answer.
- Here is the step-by-step approach based on the intuition mentioned:

We first find out the length k of the digits where n falls. Numbers with 1 digit are from 1 to 9, with 2 digits are from 10 to 99,

and so on. We can figure out the count of all such numbers with k digits as $9 * 10^{(k-1)}$. We iterate from length 1 upwards, subtracting the count of all the k-digit numbers from n until n is less than the count of •

the next set of numbers with k+1 digits. Each time we subtract, we're effectively skipping a whole block of digits.

- After finding the length k of numbers where our digit is, we identify the number itself. We do this by adding 10^(k-1) to the result of (n-1)//k -- 10^(k-1) gets the first number of the block and (n-1)//k finds how far from the first number the target
- Finally, we find the n^th digit within the number we've identified. The index of our target digit in this number is (n-1)%k. We then convert the number to a string and retrieve our target digit using this index.
- The given solution code follows these steps to find the answer in a time-efficient manner.

The implementation in the solution code follows the steps outlined in the intuition to find the n^th digit efficiently.

• The count cnt is multiplied by 10 to reflect the count for the next size of numbers (k+1 digits).

• We calculate the index within the number where the n^{th} digit lies using idx = (n - 1) % k.

of digit lengths and being clever with arithmetic to pinpoint the exact location of the desired digit.

We then increase k by 1 to reflect the next block's digit length.

scope to the exact location of the desired digit.

Variables Initialization: k is initialized to 1, representing the length of numbers we're currently dealing with, and cnt is set

Solution Approach

one is.

to 9 because there are 9 numbers of length 1.

While Loop to Determine Block: • The condition k * cnt < n is used to continue searching for the right block. If n is greater than the total number of digits in the current block (k * cnt), we advance to the next block of numbers with one more digit.

 \circ We subtract the number of digits in the current block from n with n = k * cnt. This accounts for the digits we have skipped.

- Finding the Number: Once we know the correct block:
- We calculate the actual number using num = 10 ** (k 1) + (n 1) // k. 10 ** (k 1) gives us the starting number of the block. Then (n - 1) // k determines how many numbers into the block our target digit is.

Determining the Exact Digit:

• We convert the number num to a string and retrieve the target digit using return int(str(num)[idx]). The conversion to a string allows us to easily access any digit by its index.

Algorithm Usage: The implementation mainly involves arithmetic operations and string manipulation, which is an illustrative

example of mathematical computing and iterative search techniques. It avoids brute force iteration by narrowing down the

Data Structures Used: No complex data structures are needed for this solution; basic integer and string operations suffice. Patterns Used: The code employs a direct mathematical approach to determine the position rather than using search patterns or sorting methods.

The whole solution uses an efficient method to bypass a significant amount of unnecessary computation which would result from

- iterating through all the numbers in the sequence one by one. The core part of the solution hinges on understanding the pattern
- **Example Walkthrough**

Let's illustrate the solution approach with an example where n = 15. Our task is to determine the 15th digit in the concatenated sequence of positive integers.

We initialize k to 1, which is the length of numbers we are currently considering, starting with numbers of length 1 (1 to 9). There

• Since n > k * cnt (15 > 1 * 9), it implies that the 15th digit is not in the first block of single-digit numbers. • Thus, we subtract the count of single-digit numbers from $n: n \rightarrow 1 * 9$, resulting in n = 6.

the two-digit numbers.

Step 1: Determining the Length of Numbers

are 9 such numbers, so cnt = 9.

Step 2: While Loop to Determine Block

• Now k is 2 and cnt is 90. The condition 2 * 90 < 6 is not true, so we know that the 15th digit is within the block of two-digit numbers.

• We increase k to 2 for the next block of numbers, which are double-digit numbers (10 to 99), and update cnt to 9 * 10 to reflect the count for

- **Step 3: Finding the Number**
 - We calculate the number containing the 15th digit: num = 10 ** (2 1) + (6 1) // 2.
- This gives us num = 10 + 2, because 10 is the first two-digit number and 2 is how many numbers into the two-digit block our target digit is. So, num = 12.

• We then calculate the index within the number where the 15th digit lies: idx = (6 - 1) % 2. • The result is idx = 1, so we are looking for the second digit of the number num.

Step 4: Determining the Exact Digit

 We convert num into a string and retrieve the target digit: int(str(num)[idx]). • This results in int(str(12)[1]), which evaluates to 2.

So, the algorithm tells us that the 15th digit in the concatenated sequence of positive integers is 2.

Python

Solution Implementation

digit_count = 9

class Solution: def find nth digit(self, n: int) -> int:

loop to find the correct digit length for the given 'n'

get the digit at the calculated index of the number and return it

// Initialize count `digitCount` for the count of numbers with `k` digits

digitLength++; // Move to next digit length

// Calculate the index within the number where the nth digit is located

// k represents the number of digits in the numbers we're currently looking at

// digitCount represents the total number of digits for the current k digits wide numbers

// Once the correct range is found, calculate the actual number where the nth digit is from

let digitLength = 1; // The current number of digits we are getting through (1 for 0-9, 2 for 10-99, etc.)

n -= digitLength * numberCount; // Decrease n by the number of digits covered in this step

let numberCount = 9; // The count of numbers that have digitLength digits (9 for one-digit numbers, 90 for two-digits, etc.)

'digit length' represents the current digit length we are calculating (e.g., 1 for 0-9, 2 for 10-99, etc.)

subtract the total length covered so far

while digit length * digit count < n:</pre>

initialize variables # 'digit length' represents the current digit length we are calculating (e.g., 1 for 0-9, 2 for 10-99, etc.) digit_length = 1

'digit count' represents the count of numbers that can be formed with the current 'digit_length'

print(result) # Output will be 2, which is the 15th digit in the sequence of the number "123456789101112131415..."

n -= digitLength * digitCount; // Reduce n by the number of positions we've covered

digitCount *= 10; // Increase the count for the next range of numbers

n -= digit_length * digit_count # increment the digit length since we move on to numbers with more digits digit_length += 1

increase digit count by a factor of 10 as we move to the next set of numbers

digit_count *= 10 # find the actual number where the result digit is located number = $10 ** (digit_length - 1) + (n - 1) // digit_length$

result = sol.find nth digit(15)

int digitLength = 1;

int digitCount = 9;

Example usage:

class Solution {

int findNthDigit(int n) {

int numDigits = 1;

long digitCount = 9;

// Define variables:

public:

sol = Solution()

Java class Solution { public int findNthDigit(int n) {

// Determine the range where the nth digit lies

while ((long) digitLength * digitCount < n) {</pre>

// Initialize digit length `k` for numbers of k digits

find the index of the digit within 'number'

index_within_number = (n - 1) % digit_length

return int(str(number)[index_within_number])

- // Calculate the actual number where the nth digit is from int number = (int) Math.pow(10, digitLength -1) + (n -1) / digitLength;
- int digitIndex = (n 1) % digitLength; // Extract and return the nth digit from number return String.valueOf(number).charAt(digitIndex) - '0'; C++

#include <cmath> // Include cmath library to use pow function

// Loop to find the range in which n falls

int indexInNumber = (n - 1) % numDigits;

string numberStr = to string(number);

* @return {number} - The nth digit in the sequence.

return parseInt(number.toString()[digitIndex]);

def find nth digit(self, n: int) -> int:

initialize variables

digit_length += 1

digit_count *= 10

result = sol.find nth digit(15)

digit_length = 1

class Solution:

function findNthDigit(n: number): number {

while (digitLength * numberCount < n) {</pre>

#include <string> // Include string library to convert number to string

- // 1 * 9 digits for numbers with 1 digit // 2 * 90 digits for numbers with 2 digits // 3 * 900 digits for numbers with 3 digits and so on. while (1ll * numDigits * digitCount < n) {</pre> n -= numDigits * digitCount; ++numDigits: digitCount *= 10;
- // Return the required digit converting it back to int return numberStr[indexInNumber] - '0'; **}**; **TypeScript** /** * Find the nth digit of the infinite integer sequence.

int number = pow(10, numDigits - 1) + (n - 1) / numDigits;

// Convert the number to a string to easily access any digit

* ϕ

// Loop to find the digitLength in which the nth digit is located

digitLength += 1; // Increase the number length we are looking for

numberCount *= 10; // Increase the count to match the next number length

// Find the index within the number where the nth digit is located

// Calculate the actual number where the nth digit is found const startOfRange = Math.pow(10, digitLength -1); // The first number with digitLength digits const number = startOfRange + Math.floor((n - 1) / digitLength); // Identify the exact number // Find the index within 'number' where the nth digit is located const digitIndex = (n - 1) % digitLength;

// Convert the number to a string and get the digit at digitIndex

'digit count' represents the count of numbers that can be formed with the current 'digit_length' digit_count = 9 # loop to find the correct digit length for the given 'n' while digit length * digit count < n: # subtract the total length covered so far n -= digit_length * digit_count # increment the digit length since we move on to numbers with more digits

find the actual number where the result digit is located

number = 10 ** (digit_length - 1) + (n - 1) // digit_length

get the digit at the calculated index of the number and return it

current k value (number of digits), which is a small constant for practical purposes.

increase digit count by a factor of 10 as we move to the next set of numbers

return int(str(number)[index_within_number]) # Example usage: # sol = Solution()

find the index of the digit within 'number'

index_within_number = (n - 1) % digit_length

Time and Space Complexity

The time complexity of the given code is $O(\log n)$ because the while loop runs proportional to the number of digits in n. Each

print(result) # Output will be 2, which is the 15th digit in the sequence of the number "123456789101112131415..."

increase in the number of digits results in a ten-fold increase in the number range, thus the loop iterates through each digit length once, which is related logarithmically to n. The space complexity is 0(1) because there are a fixed number of integer variables (k, cnt, num, idx) used that do not grow

with the input size n. The conversion of num to a string does not significantly affect the space complexity as it is related to the