```
1948. Delete Duplicate Folders in System
Problem Description
 an array representing an absolute path to the ith folder in the file system.
```

In this problem, we have a file system with many duplicate folders due to a bug. We are given a 2D array paths, where paths [i] is

Two folders (not necessarily on the same level) are identical if they contain the same non-empty set of identical subfolders and underlying subfolder structure. The folders do not need to be at the root level to be identical. If two or more folders are identical, then we need to mark the folders as well as all their subfolders. The file system will delete all the marked folders and their subfolders once and then return the 2D array ans containing the

remaining paths after the deletion. The paths can be returned in any order. **Example:** Let's walk through an example to better understand the problem. Given the input paths:

["a"],

The file system looks like:

As we can see, folders "a" and "c" have the same structure and same subfolders. Hence, we need to mark and delete them along with their subfolders.

After deleting, the remaining file system looks like:

- b

So, the output ans should be: [["b"]].

Solution Explanation

To solve this problem, we can utilize a Trie data structure to represent the folder structure. Each Trie node will have a children map and a deleted flag. The key in the children map will be the folder name (a string) and the value will be the child Trie node representing the subfolder. We can follow these steps to remove the duplicate folders:

1. Create and populate the Trie based on the given input paths. 2. Traverse the Trie recursively and build a unique representation of the subtree rooted at each Trie node. We can use the subtree string representation to maintain a map pointing to the Trie nodes with the same subtree strings.

3. Check the map created in step 2. If any subtree string has multiple Trie nodes, those nodes (and their subfolders) are duplicates, so we mark them as deleted. 4. Starting from the root, traverse the Trie again and construct the remaining paths by ignoring the marked Trie nodes.

Let's discuss the solution with an example. paths = [

["a"], ["c"], ["c", "x"],

["b"], ["c", "x", "y"], ["a", "x", "y"]

First, we populate the Trie based on the paths (step 1): - root

The traversal of the Trie to build subtree string representations (step 2) results in:

root => "((a(x(y)))(b)(c(x(y))))" => "(x(y))"

=> "()" => "(x(y))" The subtree string to Trie nodes map (subtreeToNodes in the code) will have:

"()": [y, b], "(y)": [x], "(x(y))": [a, c], "((a(x(y)))(b)(c(x(y))))": [root]

As we can see, the subtree string "(x(y))" has two Trie nodes, so they are duplicates (step 3). Mark the nodes "a" and "c" as deleted:

root (not deleted) - a (deleted) - x (deleted)

- y (deleted) - b (not deleted) - c (deleted) - x (deleted) - y (deleted)

Now, traverse the Trie again to construct the remaining folder paths (step 4):

["b"] Here is the final C++ implementation of the solution provided:

cpp unordered\_map<string, shared\_ptr<TrieNode>> children; bool deleted = false;

struct TrieNode { class Solution { public: vector<vector<string>> deleteDuplicateFolder(vector<vector<string>>& paths) { vector<vector<string>> ans; vector<string> path; unordered\_map<string, vector<shared\_ptr<TrieNode>>> subtreeToNodes;

This solution's time complexity will be O(N \* L) where N is the number of paths and L is the length of the strings involved. The

sort(begin(paths), end(paths));

for (const vector<string>& path : paths) {

buildSubtreeToRoots(root, subtreeToNodes);

for (const auto& [\_, nodes] : subtreeToNodes)

for (shared\_ptr<TrieNode> node : nodes)

shared\_ptr<TrieNode> root = make\_shared<TrieNode>();

for (const auto& [s, child] : node->children)

subtreeToNodes[subtree].push\_back(node);

for (const auto& [s, child] : node->children)

constructPath(child, path, ans);

Now, let's implement the same solution in Python:

self.children = defaultdict(TrieNode)

subtree\_to\_nodes = defaultdict(list)

node = node.children[s]

for nodes in subtree\_to\_nodes.values():

node.deleted = True

for node in nodes:

self.construct\_path(root, path, ans)

for s, child in node.children.items():

for s, child in node.children.items():

if not child.deleted:

path.append(s)

path.pop()

ans.append(path[:])

Finally, let's implement the solution in JavaScript:

this.children = new Map();

this.deleted = false;

deleteDuplicateFolder(paths) {

const subtreeToNodes = new Map();

const root = new TrieNode();

for (const s of p) {

if (nodes.length > 1) {

this.constructPath(root, path, ans);

buildSubtreeToNodes(node, subtreeToNodes) {

if (!subtreeToNodes.has(subtree))

subtreeToNodes.set(subtree, []);

for (const [s, child] of node.children.entries()) {

this.constructPath(child, path, ans);

subtreeToNodes.get(subtree).push(node);

if (!node.children.has(s))

node = node.children.get(s);

this.buildSubtreeToNodes(root, subtreeToNodes);

for (const nodes of subtreeToNodes.values()) {

for (const node of nodes) {

node.deleted = true;

for (const [s, child] of node.children.entries()) {

subtree += s + this.buildSubtreeToNodes(child, subtreeToNodes);

node.children.set(s, new TrieNode());

for (const p of paths) {

let node = root;

const ans = [];

paths.sort();

return ans;

let subtree = "(";

subtree += ")";

return subtree;

if (subtree !== "()") {

constructPath(node, path, ans) {

if (!child.deleted) {

path.push(s);

ans.push(Array.from(path));

path.pop();

if (path.length > 0)

const path = [];

subtree\_to\_nodes[subtree].append(node)

self.construct\_path(child, path, ans)

self.build\_subtree\_to\_nodes(root, subtree\_to\_nodes)

unordered\_map<string, vector<shared\_ptr<TrieNode>>>& subtreeToNodes) {

space complexity will be O(N \* L) as well due to the Trie data structure.## Python Solution

def deleteDuplicateFolder(self, paths: List[List[str]]) -> List[List[str]]:

def build\_subtree\_to\_nodes(self, node: TrieNode, subtree\_to\_nodes: dict) -> str:

def construct\_path(self, node: TrieNode, path: List[str], ans: List[List[str]]):

This Python solution has the same time complexity O(N \* L) and space complexity O(N \* L).

subtree += s + self.build\_subtree\_to\_nodes(child, subtree\_to\_nodes)

subtree += s + buildSubtreeToRoots(child, subtreeToNodes);

void constructPath(shared\_ptr<TrieNode> node, vector<string>& path,

vector<vector<string>>& ans) {

node->children[s] = make\_shared<TrieNode>();

shared\_ptr<TrieNode> node = root;

if (!node->children.count(s))

for (const string& s : path) {

node = node->children[s];

node->deleted = true;

constructPath(root, path, ans);

shared\_ptr<TrieNode> node,

if (nodes.size() > 1)

string buildSubtreeToRoots(

string subtree = "(";

**if** (subtree != "()")

if (!child->deleted) {

path.push\_back(s);

path.pop\_back();

ans.push\_back(path);

from collections import defaultdict

self.deleted = False

def \_\_init\_\_(self):

ans = []

path = []

paths.sort()

return ans

subtree = "("

subtree += ")"

return subtree

if path:

JavaScript Solution

constructor() {

javascript

class TrieNode {

class Solution {

if subtree != "()":

root = TrieNode()

node = root

for s in p:

if len(nodes) > 1:

for p in paths:

if (!path.empty())

python

class TrieNode:

class Solution:

subtree += ")";

return subtree;

return ans;

private: