Given a binary tree, where each node contains an integer, we are asked to find the largest absolute difference in value between two

Problem Description

descendant of a, what is the maximum absolute difference between a.val and b.val (|a.val - b.val|) that we can find in the tree? An important detail to note is that a node can be considered an ancestor of itself, leading to a minimum absolute difference of 0 in

nodes where one node is an ancestor of the other. In other words, if we pick any node a as an ancestor, and any node b as a

such a scenario. The problem is focusing on finding the maximum difference, hence we need to look for pairs of ancestor and descendant nodes where this difference is the largest. Intuition

The intuition behind the solution is that we can find the maximum difference by thoroughly searching through the tree. We do this using a depth-first search (DFS) algorithm, which will allow us to explore each branch of the tree to its fullest extent before moving

on to the next branch.

During the traversal, we keep track of the minimum (mi) and maximum (mx) values encountered along the path from the root to the current node. At each node, we calculate the absolute difference between root val and both the mi and mx values, updating the

global maximum ans if we find a larger difference. The core idea is to track the range of values (minimum and maximum) on the path from the root to the current node because this range will allow us to compute the required maximum absolute difference at each step. By the time we complete our traversal, we will have examined all possible pairs of ancestor and descendant nodes and thus found the maximum difference.

To implement this, we use a recursive helper function dfs(root, mi, mx) that performs a depth-first search on the binary tree. The mi and mx parameters keep track of the minimum and maximum values respectively, seen from the root to the current node. The function also updates a nonlocal variable ans, which keeps track of the maximum difference found so far.

Finally, we initiate our DFS with the root node and its value as both the initial minimum and maximum, and after completing the traversal, we return the value stored in ans, which will be the maximum ancestor-difference that we were tasked to find. **Solution Approach**

The solution to this problem involves a recursive depth-first search (DFS) algorithm to traverse the binary tree. The critical aspect of

the approach is to maintain two variables, mi and mx, to record the minimum and maximum values found along the path from the root

2. If the current root is None, which means we've reached a leaf node's child, we return, as there are no more nodes to process in

3. The helper function is designed to continuously update a nonlocal variable ans, which holds the maximum absolute difference

Here is a step-by-step breakdown of the implementation details: 1. Define a recursive helper function dfs(root, mi, mx) that will be used for DFS traversal of the tree.

found. 4. At each node, we compare and update ans with the absolute difference of the current node's value root.val with both the

Example Walkthrough

and largest values seen along the path.

result—the maximum difference. The function finally returns ans.

minimum and maximum of the path consisting of just itself.

2. Explore left child (3). Call dfs(3, min(8,3), max(8,3)):

 $- \max(5, abs(1 - 8), abs(8 - 1))$

4. The other child of 3 is 6. Call dfs(6, min(3,6), max(8,6)):

5. Node 6 has a left child 4. Call dfs(4, min(3,4), max(8,4)):

Node 1 is a leaf; the traversal will go back up.

 \circ Here, mi = 1, mx = 8.

Update answer:

 \blacksquare ans = 7

 \circ mi = 3, mx = 8.

 \circ mi = 3, mx = 8.

ans remains 7.

this path.

node to the current node.

minimum (mi) and maximum (mx) values seen so far along the path from the root. 5. We perform this comparison using max(ans, abs(mi - root.val), abs(mx - root.val)). 6. After updating ans, we also update mi and mx for the recursive calls on the children nodes, setting mi to min(mi, root.val) and mx to max(mx, root.val). This ensures that as we go deeper into the tree, our range [mi, mx] remains updated with the smallest

7. Recursive calls are then made to continue the DFS traversal on the left child dfs(root.left, mi, mx) and the right child dfs(root.right, mi, mx) of the current node.

mechanism in Python, where every child node receives the current path's minimum and maximum values to keep the comparison going. After the completion of the DFS traversal, the ans variable, which was kept up-to-date during the traversal, will contain the final

root's value, since initially, the root is the only node in the path. The implementation leverages the default argument-passing

The main function initializes the variable ans to 0 and then calls dfs(root, root.val, root.val). We start with both mi and mx as the

The primary data structure used in this implementation is the binary tree itself. No additional data structures are needed because the recursion stack implicitly manages the traversal, and the updating of minimum and maximum values is done using integer variables. This efficient use of space and recursive traversal makes it a neat and effective solution.

Let's consider a small binary tree to illustrate the solution approach. Our binary tree is as follows:

1. The dfs function is first called with root.val = 8, mi = 8, mx = 8. We are at the root.

3. Go down to the left child of 3, node 1. Call dfs(1, min(3,1), max(8,1)):

 \circ Now mi = 3, mx = 8. Update potential answer compare with previous ans: $- \max(0, abs(3 - 8), abs(8 - 3))$ \blacksquare ans = 5

We want to find the maximum absolute difference between the values of any two nodes where one is an ancestor of the other.

We begin by calling the recursive function dfs on the root node with value 8. We start with mi = mx = 8 since the root is both the

- Update answer: $- \max(7, abs(3-6), abs(8-6))$
- Update answer: $- \max(7, abs(3-4), abs(8-4))$ ans remains 7.

Since 4 is a leaf node, we go back up.

ans remains 7.

 \circ mi = 3, mx = 8. Update answer: $- \max(7, abs(3-7), abs(8-7))$

7. Now explore right child (10) of the root. Call dfs(10, min(8,10), max(8,10)):

 \blacksquare ans becomes 14 - 8 = 6 but since our current ans = 7, there is no update.

■ ans remains 7 because the differences 5 and 1 are smaller than the current ans.

After the recursive depth-first search completes, we find that the maximum absolute difference is 7, which comes from the

Node 7 is also a leaf; traverse back up to 3, then to 8.

 $- \max(7, abs(8 - 10), abs(10 - 8))$

 $- \max(7, abs(8 - 14), abs(14 - 8))$

6. Node 6 also has right child 7. Call dfs(7, min(3,7), max(8,7)):

ans remains 7. 8. Node 10 has right child 14. Call dfs(14, min(8,14), max(10,14)): \circ mi = 8, mx = 14.

Update answer:

 \circ mi = 8, mx = 10.

Update answer:

 \circ mi = 8, mx = 14. Update answer:

 $- \max(7, abs(8 - 13), abs(14 - 13))$

difference between nodes 8 (ancestor) and 1 (descendant).

self.right = right # Right child

def dfs(node, min_value, max_value):

new_min = min(min_value, node.val)

new_max = max(max_value, node.val)

dfs(node.left, new_min, new_max)

dfs(node.right, new_min, new_max)

nonlocal max_difference

9. Node 14 has a left child 13. Call dfs(13, min(8,13), max(14,13)):

Left child

def maxAncestorDiff(self, root: Optional[TreeNode]) -> int:

Helper function to perform Depth-First Search (DFS)

Base case: if the current node is None, return

max_difference = max(max_difference, current_diff)

Calculate the maximum difference for the current node

Recursively call dfs for the left and right subtrees

Start DFS from root with its value as both initial min and max

Initialize max_difference which will hold the result

current_diff = max(abs(min_value - node.val), abs(max_value - node.val))

Update the minimum and maximum values with the value of the current node

class TreeNode: # A class for a binary tree node def __init__(self, val=0, left=None, right=None): self.val = val # Node value

self.left = left

if node is None:

return

max_difference = 0

return max_difference

* Definition for a binary tree node.

TreeNode(int val) {

this.val = val;

this.val = val;

class Solution {

/**

/**

this.left = left;

this.right = right;

private int maxDifference;

if (root == null) {

return maxDifference;

if (node == null) {

return;

return 0;

dfs(root, root.val, root.val)

Return the maximum difference found

TreeNode(int val, TreeNode left, TreeNode right) {

* @param root The root of the binary tree.

public int maxAncestorDiff(TreeNode root) {

* @return The maximum difference calculated.

depthFirstSearch(root, root.val, root.val);

* @param node The current node being visited.

maxVal = Math.max(maxVal, node.val);

Python Solution

class Solution:

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Java Solution

/**

*/

class TreeNode { int val; TreeNode left; TreeNode right; TreeNode() {} 10

* Calculates the maximum difference between values of any two connected nodes in the binary tree.

// Start DFS with the initial value of the root for both minimum and maximum.

* A recursive DFS function that traverses the tree to find the maximum difference.

private void depthFirstSearch(TreeNode node, int minVal, int maxVal) {

* @param minVal The minimum value seen so far in the path from root to the current node.

* @param maxVal The maximum value seen so far in the path from root to the current node.

// Calculate the potential differences between the current node value and the observed min and max values.

```
int currentMaxDifference = Math.max(Math.abs(minVal - node.val), Math.abs(maxVal - node.val));
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           // Update the maxDifference if the current one is greater.
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           maxDifference = Math.max(maxDifference, currentMaxDifference);
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           // Update the min and max values to carry them forward in the DFS.
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           minVal = Math.min(minVal, node.val);
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// Recur for both the left and right subtrees.
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           depthFirstSearch(node.left, minVal, maxVal);
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           depthFirstSearch(node.right, minVal, maxVal);
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64 }
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C++ Solution
    /**
     * Definition for a binary tree node.
     */
    struct TreeNode {
         int val;
         TreeNode *left;
         TreeNode *right;
         TreeNode(): val(0), left(nullptr), right(nullptr) {}
  8
         TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
  9
         TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left), right(right) {}
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 11 };
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 13 class Solution {
    public:
 14
         /* Function to calculate max difference between any ancestor and node value. */
 15
         int maxAncestorDiff(TreeNode* root) {
 16
 17
             int maxDifference = 0; // to store the maximum difference
 18
 19
             // Lambda function for depth-first search starting from 'node'
 20
             // It carries the current minimum and maximum values as 'currentMin' and 'currentMax'
 21
             function<void(TreeNode*, int, int)> dfs = [&](TreeNode* node, int currentMin, int currentMax) {
 22
                 if (node == nullptr) {
 23
                     // If the node is null, return as there is nothing to process
 24
                     return;
 25
 26
                 // Update maxDifference with the maximum of the current difference
 27
                 // and the differences with the current node's value
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                 maxDifference = max({
                     maxDifference,
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                     abs(currentMin - node->val),
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                     abs(currentMax - node->val)
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                 });
 33
 34
                 // Update currentMin and currentMax with respective minimum and maximum values
 35
                 currentMin = min(currentMin, node->val);
                 currentMax = max(currentMax, node->val);
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                 // Continue depth-first search on left and right subtrees
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                 dfs(node->left, currentMin, currentMax);
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                 dfs(node->right, currentMin, currentMax);
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             // Initialize DFS with the value of the root for both min and max
 44
             dfs(root, root->val, root->val);
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 46
             // Return the maximum difference found
 47
             return maxDifference;
```

25 // Primary function to initiate the maxAncestorDiff calculation, given the root of a binary tree function maxAncestorDiff(root: TreeNode | null): number { if (root === null) { // If the tree is empty, the maximum difference is 0 by definition return 0;

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};

Typescript Solution

if (!node) {

return;

let maxDifference: number = 0;

// Global variable to track the maximum difference between an ancestor and a node value

// If node is null, return because we've reached a leaf node's child.

// Update the global maxDifference if the new potential difference is greater

function dfs(node: TreeNode | null, minVal: number, maxVal: number): void {

// Calculate the potential new max differences with the current node

maxDifference = Math.max(maxDifference, potentialMaxDiff);

// Continue the DFS traversal for left and right children

const newMinVal = Math.min(minVal, node.val);

const newMaxVal = Math.max(maxVal, node.val);

dfs(node.left, newMinVal, newMaxVal);

dfs(root, root.val, root.val);

Time and Space Complexity

return maxDifference;

dfs(node.right, newMinVal, newMaxVal);

// Recursive function traverses the tree to find the maximum difference between an ancestor and a node value

const potentialMaxDiff = Math.max(Math.abs(node.val - minVal), Math.abs(node.val - maxVal));

// Update the min and max values seen so far after considering the current node's value

// Since we start at the root, the starting min and max values are the root's value

// After traversing the tree, return the global maxDifference found

The given Python function maxAncestorDiff computes the maximum difference between the values of any two nodes with an ancestor/descendant relationship in a binary tree.

The time complexity of the function is O(N), where N is the number of nodes in the binary tree. This is because the function performs

Time Complexity:

a depth-first search (DFS), visiting each node exactly once. During each visit, it performs a constant amount of work by updating the minimum and maximum values encountered so far and comparing them to the current node's value. **Space Complexity:** The space complexity of the function is O(H), where H is the height of the binary tree. This complexity comes from the call stack used

for recursion during DFS. In the worst case of a skewed tree, where the tree takes the form of a linked list (either every node has only a right child or only a left child), the height H would be equal to N, leading to a worst-case space complexity of O(N). For a balanced tree, the space complexity would be O(log N), as the height H would be logarithmic in the number of nodes N.