846. Hand of Straights

Greedy Array

Problem Description

Medium

Alice has a collection of cards, each with a number written on it. To win the game, Alice needs to organize these cards into several groups where each group must have the following properties:

1. Each group contains exactly groupSize cards. 2. The groupSize cards in each group must be consecutive cards by their numbers.

Hash Table

Sorting

The challenge is to determine if it is possible for Alice to rearrange her cards to meet the criteria above. We are given an integer

groups satisfying Alice's conditions, and we return true.

array hand, which represents the cards Alice has (where each element is the number on the card), and an integer groupSize. The goal is to return true if Alice can rearrange the cards into the desired groups, or false if she cannot. Intuition

The intuition behind the solution is to count each card and then try to form groups starting with the lowest card value. Here's the

thinking process:

- Count the occurrence of each card value using a Counter data structure. This lets us know how many times each card appears in the hand.
- Sort the hand array so that we can process the cards from the smallest to the largest. Sorting is crucial because we want to form groups starting with the smallest cards available.
- Iterate through the sorted cards. For each card value v that still has occurrences left in the counter (i.e., cnt[v] is not zero):
- a. Try to make a group of groupSize starting from v. Check each card value x from v to v + groupSize 1 to ensure they exist (i.e., their count is greater than zero in the counter).

b. If any card value x in this range is not found (i.e., cnt[x] is zero), it means we can't form a group starting from v, and we

- return false.
- c. If the card is found, decrement the count for that card in the counter since it's now part of a group. After attempting to form groups for all card values, if no problems arise, it means it is possible to rearrange the cards into
- groupSize consecutive cards.

By following this process, we can efficiently determine whether or not it is possible to organize Alice's hand into groups of

The implementation of the solution leverages a few key ideas in order to check if the cards in Alice's hand can be rearranged into groups of consecutive numbers. Here is a step-by-step breakdown of the solution with reference to the algorithms, data

Use of Counter from the collections library: The solution begins by counting the frequency of each card using the Counter

Sorting the cards: Before we can form our groups, we sort the array hand. This is essential because we need to look at the

data structure. This allows us to keep track of how many copies of each card we have and efficiently update the counts as we form groups.

Solution Approach

structures, or patterns used:

group and return False.

if cnt[x] == 0:

if cnt[x] == 0:

Example Walkthrough

cnt.pop(x)

cnt[x] -= 1

return False

Iterating through sorted cards: We begin a for loop through each card value v in the sorted hand. for v in sorted(hand):

smallest card first and attempt to group it with the next groupSize - 1 consecutive card numbers.

- Forming groups by checking card availability: For each card v that is present in the counter (i.e., cnt[v] > 0), we look for the next groupSize consecutive numbers. We iterate from v up to v + groupSize, checking if each consecutive card x is available.
- for x in range(v, v + groupSize):

Checking and decrementing count: Each time we find the card x is available (i.e., cnt[x] > 0), we decrement its count in

our counter to indicate that it is now part of a group. If x is not available (i.e., cnt[x] = 0), we cannot form the required

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Optimizing by removing zero counts: Once a card's count reaches zero, we remove it from the counter. This is an
optimization step that reduces the size of the counter object as we continue through the cards. It's not strictly necessary for
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the correctness of the algorithm but can improve performance by reducing the lookup time for the remaining items.

is, we never return False), it means that it is possible to rearrange the cards as desired. In that case, the function returns True. By employing this solution approach, we are able to effectively determine the possibility of arranging Alice's cards into the specified groups, exploiting the capabilities of the Counter data structure for efficient counting and updating, as well as the

Success condition: If we successfully iterate through all the card values without finding a group that cannot be formed (that

Sort the hand array: The sorted hand is [1, 2, 2, 3, 3, 4, 6]. Iterate through the sorted cards:

Let's say Alice has the following cards in her hand: [1, 2, 3, 6, 2, 3, 4], and the group size she wants to form is 3. We want

to determine if she can organize her cards into groups where each group contains exactly three consecutive cards.

• Attempt to find 1, 2, 3 for the first group. All are present, so decrement their counts: {2: 1, 3: 1, 4: 1, 6: 1}.

• Attempt to find 2, 3, 4 for the next group. All are present, so decrement their counts: {3: 0, 4: 0, 6: 1}.

• Counter now contains only \{6: 1\}, which is not enough to form a group of three consecutive numbers.

counts, and concluded that the arrangement is not possible. We would return False in this case.

from collections import Counter

return True

return true;

card_count = Counter(hand)

for card value in sorted(hand):

if card count[card value]:

Python

Java

class Solution {

class Solution:

Encountering an issue forming groups: • We are left with the card 6 which cannot form a group with two other consecutive numbers. • Therefore, it is *not* possible for Alice to organize her cards into the required groups.

Following the steps outlined in the solution approach, we checked each card and its availability to form groups, updated the

Solution Implementation

Count the frequency of each card in the hand

Sort the hand to form groups in ascending order

If the current card can start a group

def is n straight hand(self, hand: List[int], group_size: int) -> bool:

Attempt to form a group of the specified size

card count[next card value] -= 1

If all cards can be successfully grouped, return True

public boolean isNStraightHand(int[] hand, int groupSize) {

return false;

bool isNStraightHand(vector<int>& hand, int groupSize)

unordered map<int, int> cardCounts;

sort(hand.begin(), hand.end());

// Traverse through the sorted hand.

if (cardCounts.count(card)) {

return false;

// If all cards can be grouped, return true.

for (int card : hand) {

for (int card : hand) {

// Create a map to count the frequency of each card.

// Sort the hand to arrange the cards in ascending order.

if (!cardCounts.count(nextCard)) {

if (--cardCounts[nextCard] == 0) {

++cardCounts[card]; // Increment the count for each card.

if (cardCounts.get(currentCard) == 0) {

cardCounts.remove(currentCard);

if card count[next card value] == 0:

card count.pop(next card value)

// Creating a map to count the frequency of each card value in the hand

ordered processing of the cards based on their values.

◦ The count will be {1: 1, 2: 2, 3: 2, 4: 1, 6: 1}.

The next smallest is 2 (since 1 is no longer there).

Start with the smallest value: 1.

Count the occurrences of each card value using Counter:

If the next card isn't available, the group can't be formed if card count[next_card_value] == 0: return False # Decrement the count of the current card forming the group

for next card value in range(card value, card value + group size):

If all instances of this card have been used, remove it from the counter

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Map<Integer, Integer> cardCounts = new HashMap<>();
for (int card : hand) {
    cardCounts.put(card, cardCounts.getOrDefault(card, 0) + 1);
// Sorting the hand array to ensure we create sequential groups starting from the lowest card
Arrays.sort(hand);
// Iterating through sorted hand
for (int card : hand) {
    // If the card is still in cardCounts, it means it hasn't been grouped yet
    if (cardCounts.containsKey(card)) {
        // Creating a group starting with the current card
        for (int currentCard = card; currentCard < card + groupSize; ++currentCard) {</pre>
            // If the current card does not exist in cardCounts, we can't form a group
            if (!cardCounts.containsKey(currentCard)) {
```

// Decrement the count of the current card, as it has been used in the group

// If the count goes to zero, remove the card from the map as it is all used up

cardCounts.put(currentCard, cardCounts.get(currentCard) - 1);

// Function to check if the hand can be arranged in groups of consecutive cards of groupSize.

// If the current card is still in count map (i.e., needed to form a group).

for (int nextCard = card; nextCard < card + groupSize; ++nextCard) {</pre>

// Decrement count for the current card in the sequence.

// If the next card in the sequence is missing, can't form a group.

// Attempt to create a group starting with the current card.

// If we've formed groups with all cards without returning false, return true

C++

public:

#include <vector>

class Solution {

#include <algorithm>

#include <unordered map>

```
return true;
};
TypeScript
// Import necessary functionalities from the 'lodash' library.
import _ from 'lodash';
// Define a function to check if the hand can be arranged in groups of consecutive cards of groupSize.
// The function accepts a hand (array of numbers) and a groupSize (number).
function isNStraightHand(hand: number[], groupSize: number): boolean {
    // Create a map (an object) to count the frequency of each card.
    const cardCounts: { [key: number]: number } = {};
    hand.forEach(card => {
        cardCounts[card] = (cardCounts[card] || 0) + 1; // Increment the count for each card.
    });
    // Create a sorted copy of the hand to arrange the cards in ascending order.
    const sortedHand = _.sortBy(hand);
    // Traverse through the sorted hand.
    for (const card of sortedHand) {
        // If the count for the current card is more than 0 (i.e., needed to form a group).
        if (cardCounts[card] > 0) {
            // Attempt to create a group starting with the current card.
            for (let nextCard = card; nextCard < card + groupSize; ++nextCard) {</pre>
                // If the next card in the sequence is missing or count is 0, can't form a group.
                if (!cardCounts[nextCard]) {
                    return false;
                // Decrement count for the current card in the sequence.
                --cardCounts[nextCard];
    // If all cards can be grouped, return true.
    return true;
from collections import Counter
class Solution:
    def is n straight hand(self, hand: List[int], group_size: int) -> bool:
```

cardCounts.erase(nextCard); // Remove the card from count map if count reaches zero.

Time Complexity The time complexity of the function can be broken down as follows:

groupSize where each group consists of consecutive integers.

Count the frequency of each card in the hand

Sort the hand to form groups in ascending order

Attempt to form a group of the specified size

if card count[next_card_value] == 0:

if card count[next card value] == 0:

card count.pop(next card value)

card count[next card value] -= 1

If all cards can be successfully grouped, return True

for next card value in range(card value, card value + group size):

Decrement the count of the current card forming the group

If all instances of this card have been used, remove it from the counter

If the next card isn't available, the group can't be formed

If the current card can start a group

return False

card_count = Counter(hand)

return True

Time and Space Complexity

for card value in sorted(hand):

if card count[card value]:

Counting Elements (Counter): Constructing a counter for the hand array takes O(N) time, with N being the length of the hand array.

Therefore, the overall worst-case time complexity of the function is $0(N \log N + N * groupSize)$. Since groupSize is a constant,

The given Python function isNStraightHand checks whether an array of integers (hand) can be partitioned into groups of

- **Iteration Over Sorted Hand:** The outer loop runs at most N times. However, within this loop, there is an inner loop that runs up to groupSize times. This results in a total of O(N * groupSize) operations in the worst case.
- it can be simplified to O(N log N). **Space Complexity**

Sorting: The sorted(hand) function has a time complexity of O(N log N) since it sorts the array.

The space complexity pertains to the additional memory used by the algorithm as the input size grows. For this function: Counter Space Usage: The Counter used to count instances of elements in hand will use O(N) space.

Sorted Array Space: The **sorted()** function creates a new list and thus, requires **O(N)** space as well.

Since both the Counter and the sorted list exist simultaneously, the overall space complexity would also be O(N).