

997. Find the Town Judge

EasyGraphArrayHash Table

Leetcode Link

Problem Description

In this problem, we are given a scenario where there is a small town with n people indexed from 1 to n . Among them, there is a rumor of one person being the town judge. The town judge is defined by two distinct characteristics:

- The town judge trusts no one, meaning they do not express trust towards any other person in the town.
- Every other person in the town trusts the judge, that means all individuals aside from the judge themselves trust this judge.
- There is only one individual who satisfies both of the above conditions, making them the town judge.

The problem provides us with trust data in the form of an array, where each subarray `trust[i] = [a_i, b_i]` represents that person `a_i` trusts person `b_i`. This array represents all the trust relationships that exist in the town. If no trust relationship exists between two individuals, it will not be represented in this array.

The goal is to find out who the town judge is, if one exists, based on these trust relationships. The output should be the label of the town judge. If no town judge can be identified, based on the rules described, the output should be `-1`.

Intuition

The solution approach follows from the given properties of the town judge.

- Since every person except for the town judge trusts the judge, the town judge will be the one who is trusted by $n-1$ others (everybody else).
- Since the town judge trusts no one, the judge will not show up as the first element in any trust pair `[a_i, b_i]`.

To find the town judge, we can keep two counts for each person:

- `cnt1[i]`: Count of people that person `i` trusts. This will be increased when person `i` is the `a_i` in a trust pair.
- `cnt2[i]`: Count of people who trust person `i`. This will be increased when person `i` is the `b_i` in a trust pair.

As we iterate through the array, we populate these two count arrays. Then we look for a person `i` for whom `cnt1[i]` is `0` (trusts no one) and `cnt2[i]` is $n-1$ (is trusted by everyone else). This person is the town judge.

If there isn't anyone who meets these conditions, we return `-1` to indicate that there's no identifiable town judge based on the given trust relationships.

Solution Approach

The implementation of the solution uses a simple counting algorithm, which is based on the properties that define who the town judge could be:

- The judge trusts no one, so they will not appear as the first person in any of the trust pairs.
- Everyone trusts the judge, so the judge will be the second person in $n-1$ trust pairs if there are n people in the town.

The solution uses two list data structures, `cnt1` and `cnt2`. Both lists, indexed from 1 to n , are initialized to `0`. The lists are used in the following way:

- `cnt1[i]` represents the number of people that person `i` trusts.
- `cnt2[i]` represents the number of people who trust person `i`.

As we iterate through the `trust` list, we update these counts:

- For each trust pair `[a, b]` found in the `trust` list, we increase `cnt1[a]` by `1` because person `a` is shown to trust another person.
- Simultaneously, we increase `cnt2[b]` by `1` because it shows that person `b` is trusted by someone (person `a` in this case).

When the iteration of `trust` pairs is complete, we need to find the person who meets the criteria for being the town judge:

- Person `i` satisfies the condition `cnt1[i] == 0`, indicating they do not trust anyone else, and
- Person `i` satisfies the condition `cnt2[i] == n - 1`, indicating that they are trusted by everyone else in town.

If such a person exists, their index `i` is returned as the town judge. If no one person satisfies both conditions, the function returns `-1`, signaling there is no identifiable town judge.

The algorithm's time complexity is $O(T + n)$, where T is the number of trust relationships (the length of the trust list) and n is the number of people in the town. The space complexity is $O(n)$ for the two count lists used to store the trust counts for each person.

Example Walkthrough

Let's use a small example to illustrate the solution approach. Suppose the town has $n = 3$ people, and the trust relationships are given by the array `trust = [[1, 3], [2, 3]]`. Here's how the algorithm finds the town judge:

- We initialize two count arrays, `cnt1` and `cnt2`, with size equal to the number of people in the town (3 in this case), filled with zeroes. `cnt1` tracks the number of people a person trusts, while `cnt2` tracks the number of people that trust a person.

At the start: `cnt1 = [0, 0, 0]` `cnt2 = [0, 0, 0]`

- We iterate over each trust relationship. First, we look at `[1, 3]`:
 - Increment `cnt1[1]` by 1 since person 1 trusts person 3.
 - Increment `cnt2[3]` by 1 since person 3 is trusted by person 1.After this step: `cnt1 = [1, 0, 0]` `cnt2 = [0, 0, 1]`

- Next, we process `[2, 3]`:
 - Increment `cnt1[2]` by 1 since person 2 trusts person 3.
 - Increment `cnt2[3]` by 1 since person 3 is trusted by person 2.After this step: `cnt1 = [1, 1, 0]` `cnt2 = [0, 0, 2]`

- Now, with the counts updated, we look for the person who is trusted by $n-1$ others and trusts no one:
 - Person 3 satisfies `cnt1[3] == 0` (trusts nobody) and `cnt2[3] == 2` (is trusted by others).
- Finally, we conclude that person 3 is the town judge, as they meet both conditions for being the judge.

Using the template described:

```
1
2 Imagine a town with 3 people. We are given the trust array `trust = [[1, 3], [2, 3]]`, which tells us person 1 trusts person 3, and p
3
4 1. We start by initializing two arrays to track the trust counts, `cnt1` and `cnt2`. Initially, no one trusts anyone, so both arrays
5
6 2. For the first pair `[1, 3]`, person 1 trusts person 3. We increment `cnt1[1]` because person 1 is showing trust, and `cnt2[3]` bec
7
8 3. Next, we process the second pair `[2, 3]`. Similarly, we increment `cnt1[2]` and `cnt2[3]`, reflecting that person 2 trusts persor
9
10 4. After analyzing the trust pairs, we see that person 3 is trusted by both other people (`cnt2[3] = 2`), and person 3 trusts nobody
11
12 Therefore, the algorithm identifies person 3 as the town judge in this example.
```

Python Solution

```
1 class Solution:
2     def findJudge(self, n: int, trust: List[List[int]]) -> int:
3         # Initialize two lists to count the trust votes.
4         # trust_received[i] will hold the number of people who trust person i.
5         # trust_given[i] will hold the number of people person i trusts.
6         trust_received = [0] * (n + 1)
7         trust_given = [0] * (n + 1)
8
9         # Iterate through each trust relationship in the trust list.
10        for giver, receiver in trust:
11            # Increment the trust count: giver trusts another person,
12            # and the receiver is trusted by another person.
13            trust_given[giver] += 1
14            trust_received[receiver] += 1
15
16        # Iterate over each person to find the potential town judge.
17        for person in range(1, n + 1):
18            # The town judge should not trust anyone (trust_given[person] == 0)
19            # and should be trusted by everyone else (trust_received[person] == n - 1).
20            if trust_given[person] == 0 and trust_received[person] == n - 1:
21                return person # Return the town judge.
22
23        # If no town judge is found, return -1.
24        return -1
```

Note: `List` needs to be imported from `typing` if not done already in the code. To adhere to Python coding standards, the import statement would look like this:

```
1 from typing import List
2
```

Java Solution

```
1 class Solution {
2
3     public int findJudge(int n, int[][] trust) {
4         // Array to keep track of the number of people each person trusts
5         int[] trustCount = new int[n + 1];
6
7         // Array to keep track of the number of people who trust each person
8         int[] trustedByCount = new int[n + 1];
9
10        // Iterate over the trust relationships
11        for (int[] relation : trust) {
12            int personTrusts = relation[0]; // Person that trusts someone
13            int personTrusted = relation[1]; // Person that is trusted
14
15            // Increment the trust count for the person who trusts someone
16            trustCount[personTrusts]++;
17
18            // Increment the count of being trusted for the person who is trusted
19            trustedByCount[personTrusted]++;
20        }
21
22        // Iterate over all the people to find the judge
23        for (int i = 1; i <= n; i++) {
24            // The judge is the person who trusts nobody (trustCount is 0)
25            // and the person who is trusted by everyone else (trustedByCount is n - 1)
26            if (trustCount[i] == 0 && trustedByCount[i] == n - 1) {
27                return i; // The index represents the person, so return it as the judge's identity
28            }
29        }
30
31        // If no judge is found, return -1
32        return -1;
33    }
34 }
35
```

C++ Solution

```
1 class Solution {
2 public:
3     int findJudge(int N, vector<vector<int>>& trust) {
4         // Create two vectors to store trust counts
5         vector<int> trustsOthers(N + 1, 0); // Counts the number of people who trust 'i'
6         vector<int> trustByOthers(N + 1, 0); // Counts the number of people 'i' trusts
7
8         // Calculate trust counts
9         for (const auto& relation : trust) {
10            int truster = relation[0]; // Person who trusts
11            int trustee = relation[1]; // Person being trusted
12            trustsOthers[truster]++; // Increment trust count for the truster
13            trustByOthers[trustee]++; // Increment count of being trusted for the trustee
14        }
15
16        // Find the town judge
17        for (int i = 1; i <= N; ++i) {
18            // The judge is the person who trusts no one else (trustsOthers[i] == 0)
19            // and is trusted by everyone else (trustByOthers[i] == N - 1)
20            if (trustsOthers[i] == 0 && trustByOthers[i] == N - 1) {
21                return i; // Judge found
22            }
23        }
24
25        // If no judge found, return -1
26        return -1;
27    };
28 };
29
```

Typescript Solution

```
1 // Function to find the judge in a town.
2 // In a town of n people, a judge is known by everyone but knows no one.
3 // Parameters:
4 // - n: number - The number of people in the town.
5 // - trust: number[][] - Array of pairwise trusts; trust[i] = [a, b] represents person a trusts person b.
6 // Returns: number - The label of the judge; returns -1 if no judge is found.
7 function findJudge(n: number, trust: number[][]): number {
8     // Array to track the number of people that trust each person.
9     const trustCounts: number[] = new Array(n + 1).fill(0);
10
11    // Array to track the number of people each person trusts.
12    const trustedByCount: number[] = new Array(n + 1).fill(0);
13
14    // Process trust relationships.
15    for (const [truster, trusted] of trust) {
16        trustCounts[truster]++;
17        trustedByCount[trusted]++;
18    }
19
20    // Find the judge who is trusted by everyone but trusts nobody.
21    for (let i = 1; i <= n; i++) {
22        if (trustedByCount[i] === n - 1 && trustCounts[i] === 0) {
23            return i; // Judge found
24        }
25    }
26
27    return -1; // No judge found
28 }
29
```

Time and Space Complexity

Time Complexity

The time complexity of the code can be analyzed as follows:

- There is one loop executed over the `trust` array, which in the worst-case scenario will go through m relationships, where m is the length of the `trust` array. This loop has a time complexity of $O(m)$.
- A second loop is run to go through each person from 1 to n , which does constant work for each person. Hence, the time complexity is $O(n)$ for this loop.

Since these are two separate loops, we add the time complexities resulting in a total time complexity of $O(m + n)$.

Space Complexity

The extra space used by the code includes two arrays `cnt1` and `cnt2`, each of size $n + 1$ to keep counts of trust received and given. Therefore, the space complexity of the code is $O(n + n)$, which simplifies to $O(n)$, as generally the constants are not considered in the Big O notation, and we take the highest order term.

In summary:

- Time Complexity: $O(m + n)$
- Space Complexity: $O(n)$