## 2138. Divide a String Into Groups of Size k

Easy String Simulation

### Problem Description

This problem is about partitioning a string s into several groups of characters. Each group must contain exactly k characters. The partitioning follows a set procedure:

• If the remaining characters at the end of the string are fewer than k, we fill the last group with a fill character fill to make sure it also has

• We start from the beginning of the string and create groups by sequentially taking k characters for each until we reach the end of the string.

- exactly k characters.
- The constraint is that after removing any fill characters added to the last group and then concatenating all groups, we should get the original string s.

The goal is to return an array of strings, where each element is one of the groups created by this procedure.

Intuition

### The solution relies on a straightforward approach:

the result list.

• We use a Python list comprehension to perform this operation concisely. The list comprehension iterates over the range of indices starting from 0 up to the length of the string s, skipping k indices at a time.

• For each iteration, we take a slice of the string s from the current index i up to i + k. If i + k exceeds the length of the string s, Python's

• Iteratively process the string in increments of k characters. For every k characters, we take a substring from the original string s and add it to

- slice operation will just return the substring up to the end of the string, which would be less than k characters.

   To ensure that each substring in the result list has exactly k characters, we use the <code>.ljust()</code> string method. This method pads the string on
- the right with the fill character until the string reaches a specified length (k in our case). If the substring already has k characters, .ljust() leaves it unchanged.
- This solution is efficient and only goes through the string once. It also takes care of the edge case where the last group is shorter than k characters without needing additional conditional checks.
- Solution Approach

  In the given problem, we have a simple yet efficient solution that efficiently partitions the string into the required groups. The

## implementation: • List Comprehension: This is a concise way to create lists in Python. It allows for the construction of a list by iterating over

some sequence and applying an expression to each item in the sequence. In our case, we generate a list that contains each group of the partitioned string s.

• String Slicing: A vital feature of Python that allows us to extract a specific portion of a string using the syntax

solution uses Python list comprehension, string slicing, and the str.ljust method. Let's walk through the key parts of the

exclusive. If end extends beyond the length of the string, Python simply returns the substring up to the end without raising an error.

str.ljust(k, fill) Method: This string method left-justifies the string, using fill (a single character) as the fill character,

string[start:end], where start is the index of the first character inclusive and end is the index of the last character

until the whole string is of length k. If the string is already k characters or longer, no filling occurs. This method is perfect for

ensuring that each group is exactly k characters long, even the last group if it's originally shorter than k characters.

The implementation is straightforward:

We iterate over the indices of the string s with the range() function, starting from 0 and ending at len(s), taking steps of k at a time. This ensures that each iteration corresponds to a group of k characters.

For each index i, we slice the string s from i to i+k. This gives us a substring of s which is at most k characters long.

3. We apply the ljust() method to each substring with the specified fill character to ensure it is k characters long.

with fill character if needed.

The code for creating each group looks like this:

features to achieve the desired result with minimal code.

We can illustrate the solution approach with the following steps:

Following the previous steps, we would get "abc" and "def".

the fill character "x" to get "gxx".

So, the partitioned groups for "abcdefg" will be:

fill character fill = "x" to use if necessary.

- After iterating over all possible starting indices that are k characters apart, we collect all the modified substrings into a list. This list represents the groups after the partitioning process and is the final output of the function.
- [s[i : i + k].ljust(k, fill) for i in range(0, len(s), k)]

Example Walkthrough

Let's assume our input string s is "abcdefghi" and we are to partition this string into groups of k = 3 characters. We also have a

As a result, we get a list of strings, each representing a group of s that is exactly k character long, where the last group is filled

This approach is both simple to understand and effective at solving the problem, utilizing core Python string manipulation

According to the problem description, we would start at the beginning of the string and take groups of k characters in sequence.

Our groups would be "abc," "def," and "ghi." Since "ghi" has exactly 3 characters, there is no need to use the fill character in

#### 1. Start at the beginning of the string, with i = 0.

this example.

Take the substring from i to i + k, which translates to s[0:3]. This gives us "abc".
 Move to the next k characters, increment i by k to i = 3, and take the substring s[3:6], which gives us "def".

Continue the process by setting i to 6 and taking the substring s[6:9], yielding "ghi".

The list comprehension puts this all together elegantly, completing the task in a single line of code.

# It divides the string `s` into chunks of size `k` and if the last chunk

int totalPartitions = (inputLength + partitionSize - 1) / partitionSize;

# Initialize an empty list to store the chunks of the string.

# is smaller than `k`, it fills it with the `fill` character until it is of size `k`.

# If the length of the chunk is less than `k`, fill it with the fill character.

However, if we had a different string, let's say s equals "abcdefg", with the same value for k, we'd have an additional step for the last group:

The next i is 6. The substring s[6:9] would result in "g". Since this group is less than k characters, we need to pad it with

We now have the strings "abc," "def," and "ghi". All of these are of length k, so the ljust method is not necessary in this case.

always has k characters. If the last group is originally shorter than k characters, it gets padded with fill until it's k characters long.

Thus, the use of s[i : i + k].ljust(k, fill) in our list comprehension is to handle cases like this, ensuring the last group

# Solution class containing the method to divide and fill a string.
class Solution:
 def divideString(self, s: str, k: int, fill: str) -> list[str]:
 # This method takes in a string `s`, a chunk size `k`, and a fill character `fill`.

# The `liust` method is used to fill the chunk with the given character until it reaches the desired length.

# Loop through the string in steps of `k` to get each chunk.
for i in range(0. len(s). k):
 # Extract the chunk from the string starting at index `i` up to `i + k`.
 chunk = s[i : i + k]

result.append(filled\_chunk)

int inputLength = input.length();

// Calculate the required number of partitions.

// Divide the string into substrings of length k

substrings.push\_back(s.substr(i \* k, k));

// Function to divide a string into substrings of equal length `k`

function divideString(s: string, k: number, fill: string): string[] {

# Solution class containing the method to divide and fill a string.

def divideString(self, s: str, k: int, fill: str) -> list[str]:

// If string length is not a multiple of `k`, append `fill` characters

// If the length of string `s` is not a multiple of `k`, fill the last substring with the `fill` character

# This method takes in a string `s`, a chunk size `k`, and a fill character `fill`.

return substrings; // Return the vector of substrings

for (int i = 0;  $i < s.size() / k; ++i) {$ 

// Declare an array to store the resulting substrings

// Divide the string into substrings of length `k`

substrings.push(s.substring(i, i + k));

let stringLength: number = s.length;

for (let i = 0; i < s.length; i += k) {</pre>

while (stringLength % k !== 0) {

// Initialize the answer array with the calculated size.

# Return the list of chunks.

# The method returns a list of these chunks.

filled\_chunk = chunk.ljust(k, fill)

# Append the filled chunk to the `result` list.

```
Java
class Solution {
   public String[] divideString(String input, int partitionSize, char fillCharacter) {
      // Determine the length of the input string.
```

• "abc"

"def"

"gxx"

Solution Implementation

result = []

return result

```
String[] partitions = new String[totalPartitions];
        // If the input string is not a multiple of partition size, append fill characters to make it so.
        if (inputLength % partitionSize != 0) {
            input += String.valueOf(fillCharacter).repeat(partitionSize - inputLength % partitionSize);
        // Loop through each partition, filling the partitions array with substrings of the correct size.
        for (int i = 0; i < partitions.length; ++i) {</pre>
            // Calculate the start and end indices for the substring.
            int start = i * partitionSize;
            int end = (i + 1) * partitionSize;
            // Extract the substring for the current partition and assign it to the partitions array.
            partitions[i] = input.substring(start, end);
        // Return the final array of partitions.
        return partitions;
C++
class Solution {
public:
    // Function to divide a string into substrings of equal length k.
    // If the length of string s is not a multiple of k, fill the last substring with the 'fill' char.
    vector<string> divideString(string s, int k, char fill) {
        int stringLength = s.size();
        // If string length is not a multiple of k, append 'fill' characters
        if (stringLength % k) {
            for (int i = 0; i < k - stringLength % k; ++i) {</pre>
                s.push_back(fill);
        vector<string> substrings:
```

```
// Return the array of substrings return substrings;
```

class Solution:

**}**;

**TypeScript** 

const substrings: string[] = [];

s += fill;

stringLength++;

```
# It divides the string `s` into chunks of size `k` and if the last chunk
       # is smaller than `k`, it fills it with the `fill` character until it is of size `k`.
       # The method returns a list of these chunks.
       # Initialize an empty list to store the chunks of the string.
        result = []
       # Loop through the string in steps of `k` to get each chunk.
        for i in range(0, len(s), k):
           # Extract the chunk from the string starting at index `i` up to `i + k`.
            chunk = s[i : i + k]
           # If the length of the chunk is less than `k`, fill it with the fill character.
           # The `ljust` method is used to fill the chunk with the given character until it reaches the desired length.
            filled chunk = chunk.ljust(k, fill)
           # Append the filled chunk to the `result` list.
            result.append(filled chunk)
       # Return the list of chunks.
        return result
Time and Space Complexity
  Time Complexity:
  The time complexity of the provided code is mainly determined by the list comprehension. It consists of a for loop that runs
  every k characters across the string s and an ljust operation within each iteration. Here, n is the length of the input string s.
```

# The for loop iterates n/k times because it jumps k characters in each iteration. The ljust operation potentially runs in 0(k) time in the worst case, when the last segment of the string is less than k characters and needs to

be filled with the fill character.

So, the overall time complexity of the code is O(n).

Space Complexity:

• Since there are n/k elements in the list, the total space for the list is 0(n/k \* k) = 0(n).

The space complexity is determined by the space required to store the output list.
Each element of the list is a string of k characters (or fewer if fill characters are added), so the space taken by each element is 0(k).

Therefore, we consider all iterations to get the time complexity: n/k iterations \* 0(k) per ljust operation = 0(n/k \* k) = 0(n).

Additionally, the space taken by the input string s does not count towards the space complexity, as it is considered given and

not part of the space used by the algorithm. Thus, the overall space complexity of the code is O(n), where n is the length of the input string s.