

# 1175. Prime Arrangements

Easy Math

[Leetcode Link](#)

## Problem Description

The problem is to find out how many different permutations of the numbers 1 to  $n$  can be constructed in such a way that all prime numbers are at prime indices, with indices considered to be 1-indexed (meaning they start from 1, not 0). Note that prime numbers are integers greater than 1 that have no divisors other than 1 and themselves.

It is important to understand that indices in a list and the actual numbers placed at those indices can be prime. Therefore, we need to consider two separate conditions for the permutation to meet the prime placement criteria: (1) The numbers themselves must be prime, and (2) Their positions in the permutation must also be prime index locations.

Given a number  $n$ , there will be a specific count of prime numbers within the range from 1 to  $n$ , and those prime numbers must be arranged in the prime indices. All the non-prime numbers, therefore, will go into the remaining positions.

Since the answer could potentially be a very large number, the problem asks us to return the result modulo  $10^9 + 7$ , a common technique in algorithms to avoid integer overflow and to keep the numbers within a manageable range.

## Intuition

The solution is based on combinatorics. Specifically, we can break down the problem into two separate tasks:

- Counting the prime numbers from 1 to  $n$ . This tells us how many numbers need to be placed in prime index positions.
- Calculating permutations of the prime numbers in the prime indices and the permutations of the non-prime numbers in the non-prime indices.

The Sieve of Eratosthenes algorithm, which is an efficient way to find all primes smaller than a given number, is a good approach to get the count of prime numbers. With this algorithm, we iterate over each number from 2 to  $n$  and mark multiples of each number (which cannot be prime) as non-prime. The numbers that remain unmarked at the end of this process are the prime numbers.

Once we count the number of primes within the range of 1 to  $n$  (let's say there are `cnt` primes), we then must calculate the total number of permutations of these `cnt` prime numbers, which is simply the factorial of `cnt` (`factorial(cnt)`).

Simultaneously, we must also calculate the permutations for the remaining  $n - \text{cnt}$  non-prime numbers, which can be arranged in any order in the remaining positions. The total permutations for these non-prime numbers are `factorial(n - cnt)`.

The total number of prime arrangements is then the product of these two permutations: `factorial(cnt)` for the prime numbers and `factorial(n - cnt)` for the non-prime numbers.

The final step is to return this product modulo  $10^9 + 7$  to ensure the output fits within the expected range of values as required by the problem.

The solution code encloses the prime-counting logic within a function `count`, then calculates the factorials, multiplies them, and applies the modulo operation to return the answer.

## Solution Approach

The implementation of the solution can be broken down into two main parts: counting prime numbers and computing factorials for the permutations.

### 1. Counting Prime Numbers using Sieve of Eratosthenes:

- The `count` function initiates a list `primes` of boolean values, initially set to `True`, representing the numbers from 0 to  $n$ . The boolean values will represent if a number is prime (`True`) or not (`False`).
- It then iterates over the list starting from the first prime number 2. For each number  $i$ , if  $i$  is marked as `True` (indicating it is still considered prime), we:
  - Increment the `cnt` which counts the number of primes.
  - Proceed to mark all multiples of  $i$  as `False` since they are not primes.
- After the loop completes, `cnt` holds the number of prime numbers between 1 and  $n$ .

### 2. Calculating Factorials:

- The number of permissible arrangements of prime numbers is calculated by taking the factorial of the count of prime numbers (`factorial(cnt)`).
- Similarly, for non-prime numbers, we use the factorial of the difference between the total count  $n$  and the count of prime numbers (`factorial(n - cnt)`).

### 3. Calculating the Answer:

- The answer to how many unique permutations of numbers are available such that primes are at prime indices is given by the product of these two factorials: `factorial(cnt) * factorial(n - cnt)`.
- Since the product of these factorials can be very large, the final answer is taken modulo  $10^9 + 7$  to fit within the bounds specified by the problem and to avoid integer overflow.

Throughout the implementation, we see the use of basic data structures like lists (`primes` list for the sieve algorithm) and the use of the `range` function to traverse numbers and multiples. Due to Python's built-in modulo operation and factorial computation, no additional custom data structures or algorithms are required for computing factorials and their modulo.

In programming terms, this implementation leverages memoization implicitly by pre-computing the factorials of the numbers from 1 to  $n$  only once, thereby avoiding redundant calculations. This approach, combined with the efficiency of the Sieve of Eratosthenes, ensures that the algorithm runs in a relatively fast time frame appropriate for the size of input  $n$ .

## Example Walkthrough

Let us illustrate the solution approach by walking through a small example. Consider the number  $n = 6$ . We want to find out how many permutations of the numbers 1 to 6 can be formed where prime numbers occupy prime indices (1-indexed).

Firstly, we need to identify the prime numbers from 1 to 6 and also the prime indices. The prime numbers will be 2, 3, and 5, and the prime indices are 2, 3, and 5 since 1 is not considered prime.

Now, apply **Step 1: Counting Prime Numbers using the Sieve of Eratosthenes:**

- Initialize the `primes` list to track prime numbers up to  $n$  (ignoring 0 and 1). This looks like `[True, True, True, True, True, True]`.
- Starting from 2, mark non-prime multiples as `False`. The updated `primes` list becomes `[True, True, True, True, False, True]`.
- Count the number of `True` values excluding the first position which corresponds to 1. We have three primes 2, 3, 5 (which is `cnt = 3`).

Next, for **Step 2: Calculating Factorials:**

- Calculate the factorial of the count of prime numbers (which is 3): `factorial(3) = 3! = 6` to account for permutations among prime numbers.
- Also calculate the factorial of non-prime numbers ( $n - \text{cnt}$ ): `factorial(6 - 3) = factorial(3) = 3! = 6`.

Finally, in **Step 3: Calculating the Answer:**

- Multiply the results of these factorials to find the total permutations: `6 * 6 = 36`.
- Take the result modulo  $10^9 + 7$  gives `36 mod (10^9 + 7) = 36`.

So, for  $n = 6$ , there are 36 permutations where prime numbers occupy the prime indices.

## Python Solution

```
1 class Solution:
2     def numPrimeArrangements(self, n: int) -> int:
3         from math import factorial
4
5         def count_primes(n: int) -> int:
6             """
7             Count the number of prime numbers less than or equal to n using the Sieve of Eratosthenes algorithm.
8             """
9             count = 0
10            is_prime = [True] * (n + 1) # Initialize a list to track prime numbers
11
12            for i in range(2, n + 1):
13                if is_prime[i]: # If i is a prime number
14                    count += 1
15                    for j in range(i * 2, n + 1, i):
16                        is_prime[j] = False # Mark multiples of i as not prime
17            return count
18
19            prime_count = count_primes(n)
20
21            # Calculate the number of arrangements as the factorial of the prime count,
22            # multiplied by the factorial of the count of non-prime numbers
23            arrangements = factorial(prime_count) * factorial(n - prime_count)
24
25            # Return the number of arrangements modulo (10**9 + 7)
26            return arrangements % (10**9 + 7)
27
```

## Java Solution

```
1 class Solution {
2     private static final int MOD = (int) 1e9 + 7; // Constant for the modulo operation
3
4     // Counts the number of prime arrangements possible up to n
5     public int numPrimeArrangements(int n) {
6         int primeCount = countPrimes(n); // Counts the number of prime numbers up to n
7         long arrangements = factorial(primeCount) * factorial(n - primeCount);
8         // Computes the arrangements as prime! * (n - prime)!
9         return (int) (arrangements % MOD); // Returns the result modulo MOD
10    }
11
12    // Calculates the factorial of a number using the modulo operation
13    private long factorial(int n) {
14        long result = 1;
15        for (int i = 2; i <= n; ++i) {
16            result = (result * i) % MOD; // Calculates factorial with modulo at each step
17        }
18        return result;
19    }
20
21    // Counts the number of prime numbers up to n
22    private int countPrimes(int n) {
23        int count = 0; // Initialize count of primes to 0
24        boolean[] isPrime = new boolean[n + 1]; // Create an array to mark non-prime numbers
25        Arrays.fill(isPrime, true); // Assume all numbers are prime initially
26        for (int i = 2; i <= n; ++i) {
27            if (isPrime[i]) { // Check if the number is marked as a prime
28                ++count; // Increment the count of prime numbers
29                // Mark all multiples of i as non-prime
30                for (int j = i * 2; j <= n; j += i) {
31                    isPrime[j] = false;
32                }
33            }
34        }
35        return count; // Return the count of prime numbers
36    }
37 }
38
```

## C++ Solution

```
1 using ll = long long; // Alias for long long type
2 const int MOD = 1e9 + 7; // Constants should be in uppercase
3
4 // Solution class containing methods for prime arrangements calculation
5 class Solution {
6 public:
7     // Calculates the number of prime arrangements for a given number n
8     int numPrimeArrangements(int n) {
9         int primeCount = countPrimes(n); // Count the number of primes up to n
10        // Calculate the factorial of prime count and non-prime count respectively,
11        // then multiply them together modulo MOD
12        ll arrangements = factorial(primeCount) * factorial(n - primeCount) % MOD;
13        return static_cast<int>(arrangements);
14    }
15
16    // Function to calculate factorial of a number n modulo MOD
17    ll factorial(int n) {
18        ll result = 1;
19        for (int i = 2; i <= n; ++i) {
20            result = (result * i) % MOD;
21        }
22        return result;
23    }
24
25    // Function to count the number of prime numbers up to n
26    int countPrimes(int n) {
27        vector<bool> isPrime(n + 1, true); // Create a sieve initialized to true
28        int primeCount = 0;
29        for (int i = 2; i <= n; ++i) {
30            if (isPrime[i]) {
31                ++primeCount; // Increment count if i is a prime
32                // Mark all multiples of i as non-prime
33                for (int j = i * 2; j <= n; j += i) {
34                    isPrime[j] = false;
35                }
36            }
37        }
38        return primeCount;
39    }
40 };
41
```

## Typescript Solution

```
1 // Define type alias for bigint
2 type BigIntAlias = bigint;
3
4 // Constant for modular arithmetic
5 const MOD: BigIntAlias = BigInt(1e9 + 7);
6
7 /**
8  * Calculates the factorial of a number 'n' modulo 'MOD'.
9  * @param n The number to calculate the factorial of.
10  * @returns The factorial of 'n' modulo 'MOD'.
11  */
12 function factorial(n: number): BigIntAlias {
13     let result: BigIntAlias = BigInt(1);
14     for (let i = 2; i <= n; ++i) {
15         result = (result * BigInt(i)) % MOD;
16     }
17     return result;
18 }
19
20 /**
21  * Counts the number of prime numbers up to 'n'.
22  * @param n The number up to which to count primes.
23  * @returns The count of prime numbers up to 'n'.
24  */
25 function countPrimes(n: number): number {
26     const isPrime: boolean[] = new Array<boolean>(n + 1).fill(true);
27     let primeCount: number = 0;
28     for (let i = 2; i <= n; ++i) {
29         if (isPrime[i]) {
30             primeCount++;
31             for (let j = i * 2; j <= n; j += i) {
32                 isPrime[j] = false;
33             }
34         }
35     }
36     return primeCount;
37 }
38
39 /**
40  * Calculates the number of prime arrangements for a given number 'n'.
41  * @param n The number to calculate prime arrangements for.
42  * @returns The number of prime arrangements for 'n'.
43  */
44 function numPrimeArrangements(n: number): number {
45     const primeCount: number = countPrimes(n);
46     const arrangements: BigIntAlias = (factorial(primeCount) * factorial(n - primeCount)) % MOD;
47     return Number(arrangements);
48 }
49
```

## Time and Space Complexity

### Time Complexity

The time complexity of the `count` function is  $O(n * \log(\log(n)))$ . This is because the sieve of Eratosthenes, which is used to find all prime numbers up to  $n$ , has a time complexity of  $O(n * \log(\log(n)))$ .

The `factorial` function (which is not shown here, but its complexity can be derived from typical implementations) typically has a time complexity of  $O(n)$ . Given that the maximum value for factorial calculation is  $n$ , two such calculations are performed - one for the prime count `cnt` and one for the non-prime count  $n - \text{cnt}$ .

Therefore, the overall time complexity of the code is dominated by the sieve in the `count` function, which gives us the total time complexity:  $O(n * \log(\log(n)) + n + n)$ , simplifying to  $O(n * \log(\log(n)))$ .

### Space Complexity

For space complexity, the `count` function allocates an array `primes` of size  $n + 1$ , which leads to  $O(n)$  space complexity. The space requirement for the calculation of the factorial depends on the implementation, but typically, it can be calculated in  $O(1)$  space.

Therefore, the overall space complexity of the code is  $O(n)$  for the sieve of Eratosthenes array storage.