Problem Description

In this problem, we are presented with a country that has n cities arranged in a tree structure, meaning all cities are connected by roads in such a way that there is exactly one unique path between any two cities, and there are no cycles. The cities are numbered from 0 to n - 1, where city 0 is the capital. The connections between cities are described by the roads array, where each element represents a bidirectional road between two cities.

one car with a certain number of seats which can be used to travel to the meeting. When traveling from one city to another, it costs one liter of fuel. If necessary, representatives can switch cars by riding with another representative from a different city. The goal is to calculate the minimum amount of fuel needed for all representatives to reach the capital city for the meeting.

A meeting is scheduled to take place in the capital city, and every city needs to send a representative to the meeting. Each city has

Intuition

To approach this problem, we have to recognize that because the cities are connected in a tree structure, there is only one path for

optimal arrangement of representatives and cars to minimize fuel consumption. Since this is a tree structure, we can use the Depth First Search (DFS) algorithm. We start from the capital city and explore each branch, that is, each connected city recursively. At each city, we calculate the number of representatives that need to be transferred from that city to the capital city, including the representatives from its connected cities.

each representative to take to the capital city, making the task of finding the path straightforward. The challenge is finding out the

The key insight to minimize fuel is to fill each car to its maximum capacity before sending it to the capital. This ensures the least number of trips is made, and subsequently, the minimum amount of fuel is used. The intuition behind the solution code is that we perform a DFS starting from the capital. During the DFS, we keep track of the total

number of representatives that need to be moved from each branch (i.e., from each city and its descendants). As representatives are

transported to the capital city, we calculate the number of cars needed to move them based on the capacity of each car (seats). Each car that needs to move from a city to the capital city uses one liter of fuel, hence we increase the fuel count by the number of

The dfs function updates the total fuel cost (ans) as it calculates the number of representatives that need to be moved from each city to its parent. The calculation for the number of cars needed is done by (t + seats - 1) // seats, where t is the number of representatives, accounting for the last car that may not be full. Solution Approach

The implementation of the solution primarily relies on a recursive Depth-First Search (DFS) algorithm. The algorithm traverses the input tree, starting from the capital city (node 0), going through each city and figuring out the number of seats required to transport the representatives from the subtree rooted at that city to the capital.

A defaultdict from the collections module to represent the graph which describes the tree structure of the country. Keys are

The dfs function works as follows:

child calls.

is used), and finally returns ans.

The code uses the following elements in its approach:

travel from city b through city a toward the capital.

The tree structure based on the roads array is as follows:

cars needed for that branch.

 A nonlocal variable ans, which keeps track of the total fuel cost calculated during the DFS traversal. The variable is defined as nonlocal because it needs to be accessed and modified within the nested dfs function.

A recursive dfs function, which serves as the DFS traversal method, taking two parameters: the current city a, and its parent city

cities, and values are lists of neighboring cities effectively forming the adjacency list for each node.

fa. The parent city helps in avoiding traversing back on the path we came from.

- 1. For each city a, it initializes a local variable size to 1, which implicitly includes the city's own representative. 2. It iterates over all connected cities b in g[a]. If b is not the city we just visited (fa), it recursively calls dfs(b, a). 3. Within each recursive call, it computes the total number of representatives t (including those from the subtrees) that need to
- seats 1) // seats calculates the smallest number of cars needed to transport t representatives given each car has seats number of seats. It uses integer division rounding up, which is achieved by adding seats - 1 before dividing. 5. The function then returns the size of the subtree rooted at the current city a to its parent call which aggregates sizes from all

4. The fuel cost for the subtree rooted at b is then updated by adding (t + seats - 1) // seats to ans. The expression (t +

This algorithm ensures that each representative is counted exactly once and the minimum number of cars needed to transport representatives from each subtree to the root (capital city) is calculated, thus leading to the calculation of the minimum fuel cost.

After defining the dfs and g, the function calls dfs(0, -1) to start the traversal from the capital city (which has no parent, hence -1

Example Walkthrough Let's illustrate the solution approach with a small example. Suppose we have a country with 5 cities (n = 5) arranged in a tree structure with the following connections and seats in cars:

 Roads: [(0, 1), (0, 2), (1, 3), (1, 4)] Seats: 4 for all cars (for simplicity)

Each city except for the capital has one representative who needs to get to the capital. Now, let's walk through the DFS approach:

the representative from city 3).

these representatives to the capital.

def minimumFuelCost(self, roads, seats):

def dfs(current, parent):

subtree_size = 1

return subtree_size

graph[start].append(end)

graph[end].append(start)

self.total_trips = 0

0 (capital)

2. We explore city 1 (which has not been visited). City 1 has 2 neighbors: city 3 and city 4. 3. We go to city 3 (not visited), and it has no neighbors. We determine one representative needs to travel from city 3 to city 1. Since each car has 4 seats, only one car (and thus one liter of fuel) is used. The DFS function returns a subtree size for city 3 as 1 (just

4. Back to city 1, we next explore city 4 with a similar process as city 3. One liter of fuel is used, and the subtree size is 1.

The answer for our example is 4 liters as the minimum amount of fuel needed for all representatives to reach the capital.

Initialize the size of the current subtree, starting with 1 for the current node.

5. Now back at city 1, we have collected the subtree sizes from city 3 and city 4. The total number of representatives at city 1 is 1

(its own) + 1 (from city 3) + 1 (from city 4) = 3. Only one car with 4 seats is needed, so one more liter of fuel is used to move

6. We go back to the capital and move to city 2. It has no neighbors to explore, so we use one car and one liter of fuel to transport

7. Finally, we collect the total fuel cost as the sum of the fuel needed for the subtrees: 1 (from city 3) + 1 (from city 4) + 1 (from city 1 to capital) + 1 (from city 2 to capital) = 4 liters.

Recursive function to perform depth-first search (DFS) on the graph.

Recursively call dfs for the adjacent node.

self.total_trips += (child_size + seats - 1) // seats

child_size = dfs(next_node, current)

The variable to store the total number of trips required.

Start DFS from node 0 assuming 0 as the root with no parent (-1).

subtree_size += child_size

Return the size of the current subtree.

the representative to the capital. The total size of city 2's subtree is 1.

1. We start the DFS from the capital city (0). City 0 has 2 neighbors: city 1 and city 2.

- **Python Solution** from collections import defaultdict class Solution:
- # Iterate through adjacent nodes of the current node. for next_node in graph[current]: 11 # Avoid traversing the edge leading back to the parent node. if next_node != parent:

Update the total number of trips required to move people from the current subtree,

Update the size of the current subtree by adding the size of the child subtree.

ensuring each trip has no more people than the number of available seats.

26 27 # Initialize graph as an adjacency list representation. 28 graph = defaultdict(list) 29 # Build the graph using roads information. 30 for start, end in roads:

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           dfs(0, -1)
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           # Return the total number of trips calculated.
           return self.total_trips
41
42
43 # Example usage:
44 # sol = Solution()
  # print(sol.minimumFuelCost([[0, 1], [1, 2]], 2)) # Example call to the method.
46
Java Solution
   import java.util.ArrayList;
  2 import java.util.Arrays;
    import java.util.List;
    class Solution {
         private List<Integer>[] graph; // "g" is renamed to "graph" for clarity
         private long totalFuelCost; // "ans" is renamed to "totalFuelCost" for clarity
         private int maxSeats; // "seats" is renamed to "maxSeats" for clarity
  8
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         // Function to calculate the minimum fuel cost
 11
         public long minimumFuelCost(int[][] roads, int seats) {
 12
             int nodeCount = roads.length + 1; // Calculate the number of nodes
 13
             graph = new List[nodeCount]; // Initialize the graph
 14
             Arrays.setAll(graph, k -> new ArrayList<>()); // Create a list for each node
 15
             maxSeats = seats; // Set the maximum number of seats
 16
             for (int[] edge : roads) { // Iterate over all roads
 17
                 int from = edge[0], to = edge[1]; // Get the from and to nodes for each road
 18
                 graph[from].add(to); // Add the to node into the from node's list
 19
                 graph[to].add(from); // Add the from node into the to node's list
 20
 21
             dfs(0, -1); // Perform a DFS traversal starting from node 0 with no parent (-1)
 22
             return totalFuelCost; // Return the total fuel cost
 23
 24
 25
         // Depth-First Search function to traverse the graph and calculate fuel cost
 26
         private int dfs(int node, int parent) {
 27
             int subtreeSize = 1; // Initialize the subtree size (1 for the current node)
 28
             for (int adjacent : graph[node]) { // Iterate over the adjacent nodes
 29
                 if (adjacent != parent) { // Check if it's not coming back to the parent
 30
                     int childTreeSize = dfs(adjacent, node); // Recursively call dfs for the child node
 31
                     // Calculate and add fuel cost for reaching the child node's subtree
 32
                     totalFuelCost += (long) Math.ceil((double) childTreeSize / maxSeats);
 33
                     subtreeSize += childTreeSize; // Update the size of the current subtree
 34
 35
 36
             return subtreeSize; // Return the size of the subtree rooted at the current node
 37
 38 }
```

21 graph[to].emplace_back(from); 22 23 24 // Variable to store the total fuel cost 25 long long totalFuelCost = 0;

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C++ Solution

1 #include <vector>

class Solution {

public:

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2 #include <functional>

using namespace std;

// Function to calculate the minimum fuel cost for a road network

long long minimumFuelCost(vector<vector<int>>& roads, int seats) {

// Number of stops in the network is one more than number of roads

// 'seats' represents the number of seats in the vehicle

// Graph 'g' to represent the network connections

// Building the adjacency list for the graph

int from = edge[0], to = edge[1];

// Traverse all the connected stops

graph[from].emplace_back(to);

int numStops = roads.size() + 1;

for (auto& edge : roads) {

vector<vector<int>> graph(numStops);

// 'roads' represent connections where roads[i] = [a, b] is a road between a and b

// Depth-first search function to calculate the size of each subtree and add fuel costs

function<int(int, int)> dfs = [&](int currentNode, int parentNode) -> int {

int subtreeSize = 1; // Each stop contributes size 1 to the subtree

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32
                 for (int nextNode : graph[currentNode]) {
 33
                     if (nextNode != parentNode) {
 34
                         // Calculate the size of child subtree
 35
                         int childSize = dfs(nextNode, currentNode);
 36
                         // Add fuel cost for trips from this child
 37
                         totalFuelCost += (childSize + seats - 1) / seats;
 38
                         // Increment subtree size with the size of the child
 39
                         subtreeSize += childSize;
 41
 42
                 return subtreeSize;
 43
             };
 44
 45
             // Start DFS from stop 0, considering it as the root and no parent (-1)
 46
             dfs(0, -1);
 47
 48
             // Return the total fuel cost after performing DFS
 49
             return totalFuelCost;
 50
 51 };
 52
Typescript Solution
  1 // Importing necessary functionalities from libraries
  2 import { Vector } from "prelude-ts";
    // Function to calculate the size of each subtree and add fuel costs using DFS
     function depthFirstSearch(currentNode: number, parentNode: number, graph: Vector<Vector<number>>, seats: number): number {
         let subtreeSize: number = 1; // Each stop contributes size 1 to the subtree
  8
         // Get all the connected stops for the current node
         const connectedStops: Vector<number> = graph.get(currentNode).getOrElse(Vector.empty());
  9
 10
 11
         // Traverse all the connected stops and avoid visiting the parent node again
         connectedStops.forEach((nextNode: number) => {
 12
 13
             if (nextNode !== parentNode) {
 14
                 // Calculate the size of child subtree
 15
                 const childSize: number = depthFirstSearch(nextNode, currentNode, graph, seats);
 16
                 // Add fuel cost for trips from this child, round up trips given the number of seats
                 totalFuelCost += Math.ceil(childSize / seats);
 17
                 // Increment subtree size with the size of the child
 18
 19
                 subtreeSize += childSize;
 20
         });
 21
 23
         return subtreeSize;
 24 }
 25
    // Function to calculate the minimum fuel cost for a road network
     function minimumFuelCost(roads: Vector<Vector<number>>, seats: number): number {
        // Number of stops in the network is one more than the number of roads
         const numStops: number = roads.length() + 1;
         // Initialise a graph to represent the network connections as an adjacency list
         let graph: Vector<Vector<number>> = Vector.fill(numStops, () => Vector.empty<number>());
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32 33 34 // Building the adjacency list for the graph from the road connections provided roads.forEach((edge: Vector<number>) => { 35 36 const from: number = edge.get(0).getOrElse(0); 37 const to: number = edge.get(1).getOrElse(0); 38 graph = graph.updateAt(from, (edges: Vector<number>) => edges.append(to)); 39 graph = graph.updateAt(to, (edges: Vector<number>) => edges.append(from)); }); 40 41 42 // Variable to store the total fuel cost, initialised to 0 43 totalFuelCost = 0; 44 45 // Start DFS from stop 0, considering it as the root with no parent (-1)

stops needed for each subtree.

46 depthFirstSearch(0, -1, graph, seats); 47

48 // Return the total fuel cost after performing DFS 49 return totalFuelCost; 50 } 51 52 // Global variable to store the total fuel cost let totalFuelCost: number = 0; 54 Time and Space Complexity

The time complexity of the given Python code is O(N), where N is the number of nodes (intersections) in the road network

represented as a tree. This is because each node is visited exactly once thanks to the Depth-First Search (DFS) approach used in

the function dfs. In each call to dfs, it goes through all the edges connected to a node exactly once to calculate the number of fuel

The space complexity of the code is O(N) as well. The additional space used is for the recursive call stack of the dfs function and the adjacency list g that stores the tree. Since the tree is not heavily skewed, the height of the tree (which impacts the call stack size) will not be more than N in the worst case (a straight line tree). Thus, the recursive call stack size is at most N. The adjacency list also stores up to 2(N - 1) edges (since every edge is stored twice, once for each node it connects), which is within a constant factor of the number of nodes.

So, the overall space complexity remains linear with the number of nodes O(N).