678. Valid Parenthesis String String Medium Greedy **Dynamic Programming** 

## **Problem Description**

characters: '(', ')' and ''. The characters must form a valid sequence based on the following criteria: every '(' character must have a corresponding ')' character, every ')' must have a corresponding '(', the '(' must occur before its corresponding ')', and the ' character can be interpreted as either a '(', a ')', or an empty string. We need to assess whether we can rearrange the '\*' characters in such a way that all the parentheses are correctly matched, and if so, we will return true; otherwise, we return false.

The problem is to determine if a given string s is a valid string based on certain rules. The string s consists of only three types of

Intuition

Approaching the solution involves considering the flexibility that the '' character affords. The 'can potentially balance out any

unmet parentheses if placed correctly, which is crucial to forming a valid string. To solve the problem, we have to deal with the

The algorithm employs two passes to ensure that each parenthesis is paired up. In the first pass, we treat every '' as '(', increasing our "balance" or counter whenever we encounter a '(' or ', and decreasing it when we encounter a ')'. If our balance drops to zero, it means we have matched all parentheses thus far. However, if it drops below zero, there are too many ')' characters without a matching '(' or '\*', so the string can't be valid.

ambiguity of the '\*', determining when it should act as a '(', a ')', or be ignored as an empty string ' "" '.

On the second pass, we reverse the string and now treat every '' as ')'. We then perform a similar count, increasing the balance when we see a ')' or ', and decreasing it for '('. If the balance drops below zero this time, it means that there are too many '(' characters without a matching ')' or '\*', and the string is again invalid. If we complete both passes without the balance going negative, it means that for every '(' there is a corresponding ')', and the

string is valid. We have effectively managed the ambiguity of '\*' by checking that they can serve as a viable substitute for whichever parenthesis is needed to complete a valid set.

Our method ensures that all '(' have a corresponding ')', and vice versa, by treating '\*' as placeholder for the correct parenthesis

The reference solution approach suggests using <u>dynamic programming</u> to solve the problem. Dynamic programming is a method for solving complex problems by breaking them down into simpler subproblems. It is typically used when a problem has overlapping subproblems and a recursive structure that allows solutions to these subproblems to be reused.

However, the provided solution code takes a different, more efficient approach by making two single pass scans of the string: one

## In the first pass (left to right), we utilize a counter x to track the balance between the '(' and ')' characters. We increment x when

elif x:

for c in s[::-1]:

if c in '\*)':

**Example Walkthrough** 

x = 0

x = 1

**Solution Approach** 

needed, thus validating the input string.

from left to right, and another from right to left.

the string twice, without requiring any additional data structures.

1. s[0] = '(' : Since it is a '(', we increment x to 1.)

previous '(' or '\*' treated as '('.

Second Pass (Right to Left):

with a subsequent ')' or '\*' treated as ')'.

def checkValidString(self, s: str) -> bool:

open balance += 1

open balance -= 1

public boolean checkValidString(String s) {

// First pass goes from left to right

for (int i = 0; i < n; ++i) {

--balance;

return false;

// Reset balance for the second pass

// Second pass goes from right to left

char currentChar = s.charAt(i);

// Increment balance for ')' or '\*'

for (int i = n - 1; i >= 0; ---i) {

if (currentChar != '(') {

++balance:

int n = s.length(); // Length of the input string

// A closing ')' without a matching '('

# an open parenthesis.

elif open balance:

return False

Let's take the string s = "(\*)" as an example to illustrate the solution approach.

2. s[1] = '\*': Since it is a '\*', it could be '(', ')', or empty. We treat it as '(', so we increment x to 2.

4. s[3] = ')': Another')'. It can pair with the '(' or '\*' we assumed in step 2. We decrement x to 0.

rearranged into a valid sequence. Therefore, our function would return true for this input.

# If the character is ')', decrease the balance counter as a ')' is closing

# If there's no open parenthesis to balance the closing one, return False.

# Backward iteration over string to check if it can be valid from right to left.

Now, we will reset x to ∅ and traverse from right to left, treating '\*' as ')'.

# Initialization of the balance counter for open parentheses.

# Reinitialization of the balance counter for closed parentheses.

// Method to check if a string with parentheses and asterisks (\*) is valid

int balance = 0; // This will keep track of the balance of open parentheses

// Decrement balance if there is an unmatched '(' before

// If balance is zero, too many opening parentheses are encountered

// If the string passes both forward and backward checks, it's valid

// Function to check if the string with parentheses and asterisks is valid

// Forward pass to ensure there aren't too many closing parentheses

// Increment balance for an opening parenthesis or an asterisk

// If balance is zero, it means we have exact matches, so return true

// Increment balance for closing parenthesis or an asterisk

// If the string passes both forward and backward checks, it's valid

# Initialization of the balance counter for open parentheses.

# Forward iteration over string to check if it can be valid from left to right.

# If all parentheses and asterisks can be balanced in both directions,

# If character is '(' or '\*', it could count as a valid open parenthesis,

// Decrement balance for a closing parenthesis if balance is positive

// If balance is zero, too many closing parentheses are encountered

3. s[2] = ')': We have a ')', so this could potentially pair with the previous '(' or '\*'. We decrement x to 1.

we encounter a '(' or a '', and decrement it when we encounter a ')'. If we encounter a ')' but x is already 0, it means there's no '(' or 'preceding it that can be used to make a valid pair, thus the string is invalid and we return false. x = 0

for c in s: if c in '(\*': x += 1

else: return False In the second pass (right to left), we reset the counter x and treat '' as ')! We increment x for ')' or ', and decrement it for '('. Similar to the first pass, if we encounter a '(' but x is 0, it means there's no ')' or '\*' to pair with it, so the string can't be valid.

```
x += 1
  elif x:
      x -= 1
  else:
      return False
After both scans, if our balance has never gone negative, we can conclude that for every '(' there is a matching ')' (after
considering that '*' could count as either), so the string is valid.
This solution is more efficient than the dynamic programming approach mentioned in the reference solution approach. While
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dynamic programming would have a time complexity of O(n^3) and a space complexity of O(n^2), the provided solution has a

time complexity of 0(n) and a space complexity of 0(1), since it simply uses a single integer for tracking and iterates through

First Pass (Left to Right): We will initialize our balance counter x to 0. We iterate through the string one character at a time.

At the end of this pass, x is not negative, which means that we have not encountered a ')' that could not be matched with a

1. s[3] = ')': We increment x to 1 since it could pair with an '(', or a '\*', treated as '('. 2. s[2] = ')': We increment x to 2—similar reasoning as step 1. 3. s[1] = '\*': Treating it as ')', we increment x to 3. 4. s[0] = '(': We have an '(', so we pair it with one of the ')' or '\*' we treated as ')'. We decrement x to 2.

After the second pass, x is not negative, confirming that we have never encountered a '(' character that could not be matched

Since we finished both passes without the balance x going negative, we can conclude that the string s = "(\*)" can be

# Solution Implementation

open\_balance = 0

else:

closed\_balance = 0

for char in s[::-1]:

if char in '(\*':

**Python** 

class Solution:

# Forward iteration over string to check if it can be valid from left to right. for char in s: # If character is '(' or '\*', it could count as a valid open parenthesis, # so increase the balance counter.

# If character is ')' or '\*', it could count as a valid closed parenthesis, # so increase the closed\_balance counter. if char in '\*)': closed balance += 1 # If the character is '(', decrease the closed\_balance counter as '(' is # potentially closing a prior ')'. elif closed balance: closed balance -= 1 # If there's no closing parenthesis to balance the opening one, return False. else: return False # If all parentheses and asterisks can be balanced in both directions, # the string is considered valid.

### char currentChar = s.charAt(i); if (currentChar != ')') { // Increment balance for '(' or '\*' ++balance; } else if (balance > 0) {

} else {

balance = 0;

return True

Java

class Solution {

```
} else if (balance > 0) {
                // Decrement balance if there is an unmatched ')' after
                --balance;
            } else {
                // An opening '(' without a matching ')'
                return false;
       // If we did not return false so far, the string is valid
        return true;
C++
class Solution {
public:
    // Function to check if the string with parentheses and asterisks is valid
    bool checkValidString(string s) -
        int balance = 0: // Track the balance of the parentheses
        int n = s.size(); // Store the size of the string
        // Forward pass to ensure there aren't too many closing parentheses
        for (int i = 0; i < n; ++i) {
            // Increment balance for an opening parenthesis or an asterisk
            if (s[i] != ')') {
                ++balance;
            // Decrement balance for a closing parenthesis if balance is positive
            else if (balance > 0) {
                --balance;
            // If balance is zero, too many closing parentheses are encountered
            else {
                return false;
       // If only counting opening parentheses and asterisks, balance might be positive
        // So we check in the reverse order for the opposite scenario
        balance = 0; // Reset balance for the backward pass
        for (int i = n - 1; i >= 0; --i) {
            // Increment balance for closing parenthesis or an asterisk
            if (s[i] != '(') {
                ++balance;
            // Decrement balance for an opening parenthesis if balance is positive
            else if (balance > 0) {
                --balance;
```

### // Decrement balance for an opening parenthesis if balance is positive balance--; } else { // If balance is zero, too many opening parentheses are encountered return false;

} else if (balance > 0) {

// Reset balance for the backward pass

def checkValidString(self, s: str) -> bool:

# so increase the balance counter.

for (let i = n - 1; i >= 0; --i) {

if (s[i] !== '(') {

balance++:

open\_balance = 0

for char in s:

else {

return true;

**}**;

**TypeScript** 

let n: number;

let balance: number;

return false;

// Global variable to track the balance of parentheses

// Global variable to store the size of the string

n = s.length; // Get the length of the string

function checkValidString(s: string): boolean {

balance = 0; // Initialize balance

for (let i = 0; i < n; ++i) {

} else if (balance > 0) {

if (s[i] !== ')') {

balance++;

balance--;

return false;

} else {

**if** (balance === 0) {

return true;

balance = 0;

return true;

class Solution:

if char in '(\*': open balance += 1 # If the character is ')', decrease the balance counter as a ')' is closing # an open parenthesis. elif open balance: open balance -= 1 # If there's no open parenthesis to balance the closing one, return False. else: return False # Reinitialization of the balance counter for closed parentheses. closed\_balance = 0 # Backward iteration over string to check if it can be valid from right to left. for char in s[::-1]: # If character is ')' or '\*', it could count as a valid closed parenthesis, # so increase the closed\_balance counter. if char in '\*)': closed balance += 1 # If the character is '(', decrease the closed\_balance counter as '(' is # potentially closing a prior ')'. elif closed balance: closed balance -= 1 # If there's no closing parenthesis to balance the opening one, return False.

## The provided algorithm consists of two for-loops that scan through the string s. Each loop runs independently from the beginning and from the end of the string, checking conditions and updating the variable x accordingly.

Time and Space Complexity

return True

return False

# the string is considered valid.

The time complexity of each pass is O(n), where n is the length of string s, since each character is examined exactly once in each pass.

**Time Complexity** 

Since there are two passes through the string, the total time complexity of the algorithm is 0(n) + 0(n) which simplifies to 0(n).

**Space Complexity** The space complexity of the algorithm depends on the additional space used by the algorithm, excluding the input itself.

In this case, only a single integer x is used to keep track of the balance of parentheses and asterisks, which occupies constant

space. Hence, the space complexity of the algorithm is 0(1), as it does not depend on the size of the input string s and only uses a fixed amount of additional space.