186. Reverse Words in a String II

String

# **Problem Description**

Medium

**Two Pointers** 

a contiguous sequence of characters without spaces, and each word is separated by a single space. The key constraint is that the reversal of the word order must be done in-place, which means the solution cannot use additional memory to store a new array or other data structure. Intuition

The problem provided is one of reversing the order of words within a character array s. Each word within the array is defined as

To solve this problem, the intuition is to think about the desired end state: we want the words in reverse order, but each word itself should remain in its original order. An efficient way to achieve this is by applying a two-step process:

2. Reverse the entire character array. Step 1 ensures that when we perform Step 2, each word will appear in the correct order since they were individually reversed in

1. Reverse each individual word.

the first step. Implementing Step 1 involves iterating over the character array and reversing characters within each word boundary, defined by spaces. When a space is encountered, it signals the end of the current word, prompting a reversal of all characters from the start

of that word to the character immediately before the space. For the last word, since there's no trailing space, a check is needed to trigger the reversal when reaching the end of the array.

After all words are individually reversed, Step 2 simply reverses the entire array from start to end. As each word is already correctly ordered internally, this step places the words in the reverse order.

The solution employs helper method reverse that takes a subsection of the array s between indices i and j, swapping the elements at i and j, then moving i forward and j backward, repeating this process until the subsection is completely reversed.

The solution approach follows a systematic process: Initialize Pointers: We start by initializing two pointers, i and j, along with n that holds the length of the character array s. Here, i will serve as the start index of a word, and j will iterate through the array to find the end of a word.

Iterate Through s: We begin iterating over the array with j. The loop continues until j has reached the end of the array.

Reverse Individual Words: As j encounters a space, we realize that it marks the end of a word. At this point, we call the

### reverse helper function with the current start index i and the end index j - 1. After the reversal of the current word, we move i past the space to the beginning of the next word (i = j + 1).

**Example Walkthrough** 

Solution Approach

calling the reverse function with the indices 0 and n-1.

We initialize a pointer i at 0 and n as the length of the array s, which is 15.

i is then updated to j + 1, which is 4, the start of the next word "sky".

We start iterating forward through the array with a pointer j. Initially, j is also at 0.

['e', 'h', 't', ' ', 's', 'k', 'y', ' ', 'i', 's', ' ', 'b', 'l', 'u', 'e']

• This process is repeated for the remaining words. After reversing "sky" and "is", the array looks like:

['e', 'h', 't', ' ', 'y', 'k', 's', ' ', 's', 'i', ' ', 'b', 'l', 'u', 'e']

- The reverse helper function takes two indices and reverses the subsection of s between these indices. It does this by swapping the elements at i and j, then moves i forward and j backward, repeating the process until i meets or passes j. Edge Case for Last Word: Since there's no trailing space after the last word, the condition j == n - 1 checks if j has
- reached the end of the array, signaling the end of the last word, and a reversal is performed for this last word segment. Reverse Entire Array: Once all words are reversed, the final step is to reverse the entire character array. This is again done by
- reverse each word, followed by the reversal of the entire array. The reversing of elements within the array is based on the classic pattern of using a temporary variable to swap two elements, a fundamental operation in many sorting and reversing algorithms.

This algorithm does not use any extra space and operates entirely in-place, modifying the original array to achieve the desired

result. The data structure used is simply the input array, and the algorithm leverages the two-pointer technique to identify and

Let's use a simple example to illustrate the solution approach. Suppose we have an array s representing the sentence: "the sky is blue" The initial state of the array s:

['t', 'h', 'e', ' ', 's', 'k', 'y', ' ', 'i', 's', ' ', 'b', 'l', 'u', 'e'] Now, let's walk through the process:

#### $\circ$ As j moves forward, it encounters the first space at j = 3. This means the first word "the" (from i = 0 to j - 1 = 2) needs to be reversed. Using the reverse helper function, the array now looks like this:

**Initialize Pointers:** 

**Iterate Through s:** 

Reverse Individual Words:

**Reverse Entire Array:** 

Solution Implementation

Parameters:

def reverseWords(self, s: List[str]) -> None:

This method reverses the words in the input list in-place.

separated by single spaces.

start (int): The starting index of the substring

end (int): The ending index of the substring

# Initialize the starting index of the current word

reverse partial(s, word start, idx - 1)

# If a space is found, reverse the word before the space

# Update the starting index for the next word

# Reverse the whole modified list to get words in correct order

// Iterate through the array to find words and reverse them

// Update 'start' to the next word's starting index

// Helper function to reverse characters in the array between indices i and j.

// Swap characters from both ends moving towards the center

// Function to reverse a portion of the array between indices start and end.

let reverseEnd = (s[end] === ' ') ? end - 1 : end;

// Update 'start' to the next word's starting index.

// Check if the current character is a space or if we're at the end of the array.

// Calculate the position to end the reversal (consider the last word case).

function reverse(arr: string[], start: number, end: number): void {

// Swap elements at indices start and end.

// Function to reverse the words in a given character array.

if (s[end] === ' ' || end === n - 1) {

// Reverse the current word.

start = end + 1;

reverse(s, start, reverseEnd);

// After reversing all the words individually, reverse the entire array to get the final result

for (int start = 0, end = 0; end < n; ++end) {</pre>

// A word is found, reverse it

// Last word is found, reverse it

reverse(s, start, end - 1);

**if** (s[end] == ' ') {

reverse(s, 0, n - 1);

while (i < i) {

++i;

--j;

while (start < end) {</pre>

arr[end] = temp;

start++;

end--;

let temp = arr[start];

arr[start] = arr[end];

**TypeScript** 

start = end + 1;

std::swap(s[i], s[j]);

 $}$  else if (end == n - 1) {

reverse(s, start, end);

void reverse(std::vector<char>& s, int i, int i) {

s (List[str]): A list of characters representing a string with words

section (List[str]): The list of characters which substring to be reversed

section[start], section[end] = section[end], section[start]

data structure.

**Python** 

class Solution:

**Continue Iterating and Reversing Words:** 

 $\circ$  When j reaches the end of the array after "blue", j = n - 1, it triggers the last reversal from i = 11 to j = 14, resulting in: ['e', 'h', 't', ' ', 'y', 'k', 's', ' ', 's', 'i', ' ', 'e', 'u', 'l', 'b']

Now, the character array s correctly represents the sentence "blue is sky the", with the words in reverse order, and each

word internally in the correct order. All this has been achieved in-place, without using any additional memory for a new array or

- $\circ$  Finally, we reverse the entire array from 0 to n-1. After this, the array s looks like: ['b', 'l', 'u', 'e', ' ', 'i', 's', ' ', 's', 'k', 'y', ' ', 't', 'h', 'e']
- 1111111 def reverse\_partial(section: List[str], start: int, end: int) -> None: Helper function that reverses a substring of the list in-place. Parameters:

#### word start = idx + 1# If end of the list is reached, reverse the last word elif idx == length - 1: reverse\_partial(s, word\_start, idx)

length = len(s)

word start = 0

while start < end:</pre>

start += 1

# Traverse the list of characters

reverse\_partial(s, 0, length - 1)

end -= 1

for idx in range(length):

if s[idx] == ' ':

```
Java
class Solution {
    // Function to reverse words in place within a character array
    public void reverseWords(char[] str) {
        int n = str.length;
        // First, reverse each word in the array
        for (int start = 0, end = 0; end < n; ++end) {</pre>
            if (str[end] == ' ') {
                // When we find a space, reverse the previous word
                reverse(str, start, end - 1);
                // Move to the start of the next word
                start = end + 1;
            } else if (end == n - 1) {
                // If this is the end of the last word, reverse it
                reverse(str, start, end);
        // After all words are reversed, reverse the entire array to put words into the correct order
        reverse(str, 0, n - 1);
    // Helper function to reverse a section of the character array from index i to j
    private void reverse(char[] str, int i, int j) {
        // Swap characters from the starting and ending indices until they meet in the middle
        while (i < i) {
            char temp = str[i];
            str[i] = str[i];
            str[j] = temp;
            i++;
#include <vector>
#include <algorithm> // for std::swap
class Solution {
public:
    // Function to reverse words in a given character array.
    void reverseWords(std::vector<char>& s) {
        int n = s.size(); // Get the size of the character array
```

#### function reverseWords(s: string[]): void { const n: number = s.length; // Get the size of the character array // Iterate through the array to find words and reverse them. let start: number = 0; for (let end = 0; end < n; end++) {

```
// After reversing all the words individually, reverse the entire array to get the final result.
   reverse(s, 0, n - 1);
class Solution:
   def reverseWords(self, s: List[str]) -> None:
       This method reverses the words in the input list in-place.
       Parameters:
       s (List[str]): A list of characters representing a string with words
                        separated by single spaces.
       111111
       def reverse_partial(section: List[str], start: int, end: int) -> None:
           Helper function that reverses a substring of the list in-place.
            Parameters:
            section (List[str]): The list of characters which substring to be reversed
            start (int): The starting index of the substring
            end (int): The ending index of the substring
           while start < end:</pre>
                section[start], section[end] = section[end], section[start]
                start += 1
                end -= 1
        lenath = len(s)
       # Initialize the starting index of the current word
       word start = 0
       # Traverse the list of characters
       for idx in range(length):
           # If a space is found, reverse the word before the space
           if s[idx] == ' ':
                reverse partial(s, word start, idx - 1)
                # Update the starting index for the next word
               word start = idx + 1
           # If end of the list is reached, reverse the last word
           elif idx == length - 1:
                reverse_partial(s, word_start, idx)
       # Reverse the whole modified list to get words in correct order
       reverse_partial(s, 0, length - 1)
```

The given Python code is designed to reverse the words in a string in-place. Here's the analysis of its complexity:

The reverse function is a part of the total operation, which is called for each word and once for the entire string. It takes

O(k) time for each word, where k is the length of that word, and O(n) for the entire string, where n is the total length of the

## **Time Complexity:** To analyze the time complexity, we observe each part of the code:

Time and Space Complexity

For each word of length k: 0(k)

**Space Complexity:** 

input string. The while loop runs through each character in the string, which takes 0(n) time.

- Combining these, each character is involved in two operations: once in the main loop, and once when its word is reversed. Including the final reversal of the entire string, we have:
- For the entire string: O(n) Since the sum of the lengths k of all words is n, we can order these operations as O(n) + O(n) = O(2n). However, in Big O

notation, we would drop the constant to simplify the expression to O(n).

Hence, the space complexity is 0(1), which indicates constant space usage.

No additional space is needed except for a few variables (i, j, n), since the reversal is done in place. • The reverse function does not use any additional space and is performed in place.

In conclusion, the code has a time complexity of O(n) and a space complexity of O(1).