2583. Kth Largest Sum in a Binary Tree Medium Tree **Breadth-First Search Binary Search**

Leetcode Link

This LeetCode problem requires finding the kth largest level sum in a binary tree. A level sum is defined as the total value of all nodes

Problem Description

at the same depth, where depth is measured by the distance from the root node. We are given the root of the binary tree and a positive integer k, and we need to return the kth largest value among all level sums. If the tree doesn't have k levels, we must return -1. This problem emphasizes that the notion of level is based on a node's distance

from the root, meaning that all nodes at the same distance (or depth) from the root are on the same level. Intuition

To solve this problem, we will perform a level-order traversal (also known as breadth-first search) to calculate the sum of each level in the tree. Since we traverse the tree level by level, we can easily compute the sum of the nodes at each depth. The level-order

binary tree.

Here's the step-by-step approach to arriving at the solution: 1. Initialize an empty list arr to store the sum of the values at each level and a queue q initialized with the root node of the tree.

 Initialize a temporary sum t to 0 to accumulate the sum of the values for the current level. For all the nodes in the queue (which are all on the same level):

traversal uses a queue to keep track of the nodes as we progress.

2. While the queue is not empty, we proceed to the following:

 Dequeue the node from the front of the queue. Add the node's value to t.

Otherwise, we return the last element from the list obtained in the previous step, which is the kth largest level sum.

Following this intuitive approach provides us with a simple and effective solution to identify the kth largest level sum in the given

- If the node has a left child, enqueue it to the queue. If the node has a right child, enqueue it to the queue.

 - After processing all the nodes at the current level, append the sum t to the list arr.
- 3. After the traversal is complete, we use the nlargest function from Python's heapq module to fetch the k highest values from the
- array. The last element in the list returned by nlargest will be the kth largest value. 4. Finally, we check if the array contains fewer than k elements. If so, return -1 as there are not enough levels in the tree.

the entire list, and nlargest is specifically optimized to find the 'n' largest elements in a dataset.

Dequeue (popleft()) the current node from q and add its value to the temporary sum t.

Check if the current node has a left child (root.left). If so, enqueue this child node to q.

Similarly, check for a right child (root.right) and enqueue it if present.

The solution follows a breadth-first search (BFS) approach, which is perfectly suited for level-by-level traversal in a binary tree. The key data structures and algorithms used are:

• A deque from the collections module: It functions as a double-ended queue allowing for efficient insertion and removal from both ends. In our BFS, it's used to enqueue and dequeue nodes. A list arr to keep the sum of values for each level.

• The nlargest function from Python's heapq module: Heaps are used to obtain the largest (or smallest) elements without sorting

Here's a detailed breakdown of the implementation:

of node values at the current level.

length of arr with k.

Consider the following binary tree:

Level 0 sum: Queue = [5]

Level 1 sum: Queue = [3, 8]

Level 2 sum: Queue = [2, 4, 6, 9]

■ Dequeue 2, add to t = 2.

■ Dequeue 4, add to t = 6.

Level sum = 21, arr = [5, 11, 21].

■ Level sum = 5, so arr = [5].

And let k = 2, meaning we want to find the 2nd largest level sum.

1. Initialize a list arr = []. Initialize a queue q and put the root node (value 5) in it.

Dequeue 3, add to t = 3. Enqueue its left (2) and right (4) children.

■ Dequeue 5, add to temporary sum t = 5. Enqueue its left (3) and right (8) children.

visited.

Solution Approach

The deque structure allows easy access for enqueueing and dequeueing nodes as we traverse the tree level by level. 2. Initiate a while loop that continues as long as there are nodes in the queue q. This loop continues until every node has been

1. Initialize a list arr which will hold the sums of node values at each level and a deque called q with the root as the only element.

3. Inside the loop, set a temporary cumulative sum t to zero before considering each level. This is crucial for accumulating the sum

- 4. Execute another loop over q for the number of nodes at the current level (determined by len(q)). Here are the key steps within this loop:
- 5. After processing all nodes at the current level, append the level sum t to the list arr. 6. Once the while loop is complete (no more nodes to process in the queue), check if we have at least k levels by comparing the
- nlargest(k, arr)[-1]. The nlargest function constructs a min-heap of the k largest elements from arr and returns them as a list. We are only interested in the kth element, which is the last one in the list returned by nlargest. This implementation ensures efficient computation of the solution, leveraging BFS for traversal and a min-heap to effectively get the

7. If there aren't enough levels, we return -1 as specified. Otherwise, we find the kth largest value in the sums collected by calling

Example Walkthrough Let's illustrate the solution approach using a simple example:

kth largest value without the need to sort the whole arr, which could be more expensive in terms of time complexity.

2. Begin the level-order traversal:

Dequeue 8, add to t = 11. Enqueue its left (6) and right (9) children. ■ Level sum = 11, arr = [5, 11].

Python Solution

import heapq

class Solution:

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Java Solution

class TreeNode {

int val;

class Solution {

TreeNode left;

TreeNode() {}

this.val = val;

this.left = left;

this.right = right;

TreeNode right;

from collections import deque

self.left = left

level_sums = []

queue = deque([root])

Queue for Breadth First Search (BFS).

node = queue.popleft()

if node.left:

if node.right:

if len(level_sums) < k:</pre>

1 // Class definition for a binary tree node.

// Constructors to initialize tree nodes.

TreeNode(int val, TreeNode left, TreeNode right) {

// Finds the kth largest level sum in a binary tree.

List<Long> levelSums = new ArrayList<>();

Deque<TreeNode> queue = new ArrayDeque<>();

levelSum += currentNode.val;

if (root->right) {

if (levelSums.size() < k) {</pre>

return levelSums[k - 1];

return -1;

queue.push(root->right);

levelSums.push_back(sumOfCurrentLevel);

std::sort(levelSums.rbegin(), levelSums.rend());

// Return the kth largest level sum.

this.val = val === undefined ? 0 : val;

// Store the sum of the current level in the vector.

1 // Represents a node in a binary tree with a value, and left and right children

constructor(val?: number, left?: TreeNode | null, right?: TreeNode | null) {

// Sort the vector in descending order to find the kth largest level sum.

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};

Typescript Solution

val: number;

left: TreeNode | null;

right: TreeNode | null;

2 class TreeNode {

public long kthLargestLevelSum(TreeNode root, int k) {

// List to record the sum of values at each level.

// Queue to perform level order traversal of the tree.

// Perform level order traversal to calculate level sums.

// Iterate over all nodes at the current level.

TreeNode currentNode = queue.pollFirst();

for (int count = queue.size(); count > 0; --count) {

// Accumulate the values of the nodes at this level.

TreeNode(int val) { this.val = val; }

// Start with the root node.

while (!queue.isEmpty()) {

long levelSum = 0;

queue.offer(root);

return -1

current_level_sum += node.val

queue.append(node.left)

queue.append(node.right)

level_sums.append(current_level_sum)

Perform BFS to traverse the tree level by level.

self.right = right

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■ Dequeue 6, add to t = 12.
■ Dequeue 9, add to t = 21.
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4. We use nlargest to fetch the 2 largest values from arr: nlargest(2, arr) gives [21, 11].

6. Finally, we return the last element in the list returned by nlargest which is the 2nd largest value, 11.

By following these steps, we efficiently solve the problem and find the 2nd largest level sum in our example binary tree.

3. After the traversal, arr = [5, 11, 21] now contains the level sums.

5. Since arr contains more than k elements, we do not return -1.

Definition for a binary tree node. class TreeNode: def __init__(self, val=0, left=None, right=None): self.val = val

def kthLargestLevelSum(self, root: Optional[TreeNode], k: int) -> int:

This list will contain the sum of node values at each level.

Add left child to the queue if it exists.

Add right child to the queue if it exists.

Append the sum of the current level to the list.

19 while queue: 20 # Temporary variable to keep the sum of the current level's nodes. 21 current_level_sum = 0 22 # Iterate over all nodes at the current level. 23 for _ in range(len(queue)):

Use nlargest from heapq to find the kth largest element in the list of level sums. # nlargest(k, level_sums) will return the k largest elements in the list, # and [-1] will get the last element among them, which is the kth largest. 40 return heapq.nlargest(k, level_sums)[-1] 41 42

If there are fewer levels than k, return -1 as it's not possible to find the kth largest sum.

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// Add child nodes for the next level.
40
                   if (currentNode.left != null) {
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42
                        queue.offer(currentNode.left);
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                   if (currentNode.right != null) {
44
                        queue.offer(currentNode.right);
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               // Store the level sum.
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                levelSums.add(levelSum);
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           // If there are fewer levels than k, return -1.
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           if (levelSums.size() < k) {</pre>
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               return -1;
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           // Sort the sums in descending order to find the kth largest sum.
59
           Collections.sort(levelSums, Collections.reverseOrder());
60
           // Return the kth element in the sorted list, adjusting the index as list is 0-based.
61
           return levelSums.get(k - 1);
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64 }
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C++ Solution
  1 #include <algorithm>
  2 #include <queue>
     #include <vector>
    // Definition for a binary tree node.
  6 struct TreeNode {
         int val;
         TreeNode *left;
  8
         TreeNode *right;
  9
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         TreeNode() : val(0), left(nullptr), right(nullptr) {}
         TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
 11
         TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left), right(right) {}
 12
 13 };
 14
 15 class Solution {
    public:
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         // Function to find the kth largest sum of a level in a binary tree.
 18
         long long kthLargestLevelSum(TreeNode* root, int k) {
             // Initialize a vector to store the sum of each level.
 19
 20
             std::vector<long long> levelSums;
 21
 22
             // Queue for level order traversal.
 23
             std::queue<TreeNode*> queue;
 24
             queue.push(root);
 25
 26
             // Perform level order traversal to calculate the sum of each level.
             while (!queue.empty()) {
 27
 28
                 // Temporary variable to hold the sum of the current level.
                 long long sumOfCurrentLevel = 0;
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                 // Go through all nodes at the current level.
 32
                 for (int levelSize = queue.size(); levelSize > 0; --levelSize) {
 33
                     // Get the next node in the queue.
 34
                     root = queue.front();
 35
                     queue.pop();
 36
                     // Add the node's value to the sum of the current level.
 37
 38
                     sumOfCurrentLevel += root->val;
 39
                     // Add left and right children if they exist.
 40
                     if (root->left) {
 41
 42
                         queue.push(root->left);
 43
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// If there are fewer levels than k, return -1 as it's not possible to find the kth largest level sum.

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skewed (i.e., each level contains only one node).

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this.left = left === undefined ? null : left;
             this.right = right === undefined ? null : right;
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 11 }
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    /**
      * Finds the kth largest sum of values at each level of a binary tree.
     * @param {TreeNode | null} root - Root node of a binary tree.
     * @param {number} k - The kth largest element to find in the level sums.
     * @return {number} - The kth largest level sum, or -1 if there are less than k levels.
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     */
     function kthLargestLevelSum(root: TreeNode | null, k: number): number {
         // Array to store the sum of values at each level of the tree
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 21
         const levelSums: number[] = [];
 22
 23
         // Queue for breadth-first traversal, starting with the root node
 24
         const queue: Array<TreeNode | null> = [root];
         // Traverse the tree by levels
         while (queue.length) {
             let levelSum = 0; // Sum of values at the current level
             const levelSize = queue.length;
 30
 31
             // Process each node at the current level
             for (let i = 0; i < levelSize; ++i) {</pre>
 32
                 const currentNode = queue.shift();
 33
 34
                 if (currentNode) {
 35
                     // Add the value of the current node to the level sum
 36
                     levelSum += currentNode.val;
 37
                     // Add the left child to the queue if it exists
 38
                     if (currentNode.left) {
 39
                         queue.push(currentNode.left);
 40
 41
                     // Add the right child to the queue if it exists
 42
                     if (currentNode.right) {
 43
                         queue.push(currentNode.right);
 44
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             // Store the sum of the current level
             levelSums.push(levelSum);
 49
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 52
         // If there are fewer levels than k, return -1
 53
         if (levelSums.length < k) {</pre>
 54
             return -1;
 55
 56
 57
        // Sort the sums in descending order to find the kth largest
 58
         levelSums.sort((a, b) => b - a);
 59
 60
         // Return the kth largest level sum
         return levelSums[k - 1];
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Time and Space Complexity
The time complexity of the given code is O(N + k \log N), where N is the number of nodes in the binary tree. The O(N) term comes from
the BFS traversal of the tree, which processes each node once. The O(klogN) term comes from finding the k-th largest sum using the
nlargest function from the heap module, which is typically implemented as a heap-based selection algorithm with that complexity.
The space complexity is 0(N), as the queue q can contain at most all the nodes of the last level in the worst case (considering a
```

complete binary tree), and the array arr will also contain a sum for each level of the tree, which can be at most N if the tree is