1370. Increasing Decreasing String **String** Hash Table Counting Easy

Problem Description

picking the smallest and largest characters from the string according to given rules and appending them to form a new result string. The rules for picking characters are as follows: 1. Select the smallest character from s and append it to the result. 2. Choose the next-smallest character that is greater than the last appended one and append it.

The given problem requires us to reorder a given string s following a specific algorithm. The algorithm involves repeatedly

- - 3. Continue with step 2 until no more characters can be picked. 4. Select the largest character from s and append it to the result.
 - 5. Choose the next-largest character that is smaller than the last appended one and append it. 6. Continue with step 5 until no more characters can be picked.
 - The algorithm allows for any occurrence of the smallest or largest character to be chosen if there are multiple.

7. Repeat steps 1 to 6 until all characters from s have been picked and appended.

Intuition

3. Use a while loop that continues until the length of ans matches the length of the input string s.

Here's a detailed walk-through of the implementation using Python:

counter[ord(c) - ord('a')] += 1

The solution to this problem uses a frequency counter to keep track of how many times each character appears in the string s. We use an array counter of length 26 to represent the frequency of each lowercase English letter.

Here's how the solution is constructed:

1. Count the frequency of each character in the string s and store it in counter. This allows us to know how many times we need to pick each character during the reordering process. 2. Initialize an empty list ans to build up the result string.

4. Inside the loop, iterate over the counter from start to end to append the smallest character (if available) to ans and decrease its count.

5. Then iterate over the counter from end to start to append the largest character (if available) to ans and decrease its count. 6. The process repeats, alternating between picking the smallest and largest character until all characters are used. 7. Convert the list ans to a string and return it as the final reordered string.

Character Frequency Counting: We start by creating a list of zeroes called counter to maintain a count of each letter in the

string. The length of the list is 26, one for each letter of the English alphabet. We then iterate over each character of the

string s, and for each character, we find its corresponding index (0 for 'a', 1 for 'b', etc.) by subtracting the ASCII value of 'a'

Result String Assembly: We define a list ans to accumulate the characters in the order we choose them based on the

- **Solution Approach**
- The solution implementation utilizes a straightforward approach that involves counting, iterating, and string building.

from the ASCII value of the character. We increment the count at this index in the counter list.

while len(ans) < len(s):</pre>

for i in range(26):

if counter[i]:

counter[i] -= 1

counter[i] -= 1

and the counter becomes [1, 2, 1].

'c', 'a', 'b'] and counter is [0, 1, 0].

ans to ['a', 'c', 'a', 'b', 'b'] and counter to [0, 0, 0].

list updates to [1, 2, 0].

ans.append(chr(i + ord('a')))

ans.append(chr(i + ord('a')))

...

counter = [0] * 26

for c in s:

algorithm rules—the resulting string after the reordering will be formed by concatenating the characters in this list. Main Loop - Building the Result: We use a while loop to repeat the process of picking characters from the string s according to the described algorithm. The loop will continue until the length of ans becomes equal to the length of the original string s, signaling that all characters have been chosen and appended.

```
Picking the Smallest Character: Inside the loop, we iterate over the counter from the start (0) to the end (25) which
corresponds to characters 'a' to 'z'. If the current character's counter is not zero, indicating that it is available to be picked, we
append the corresponding character to ans and decrement its count in counter.
```

counter from the end (25) to the start (0). Again, if the current character's counter is not zero, we append the corresponding character to ans and decrement its count. for i in range(25, -1, -1): if counter[i]:

Picking the Largest Character: We do the same for picking the largest character, but in reverse order, iterating over the

Returning the Result: Finally, after the while loop concludes, we join the list of characters in ans using ''.join(ans) to

return the final string, which is the original string s reordered according to the algorithm described.

Character Frequency Counting: We count the frequency of every character in the string.

position represents 'a', 2 at the second position is for 'b', and 1 at the third position is for 'c'.

Main Loop - Building the Result: We start the while loop since len(ans) < len(s), which is 5 in this case.

```
loop are simple and fast) and clean, leading to a solution that straightforwardly follows the rules laid out in the problem
  statement.
Example Walkthrough
  Let's walk through a small example to illustrate the solution approach. Suppose our input string <code>s</code> is "bacab".
```

For the string "bacab", the frequency counter counter would look like this after counting: [2, 2, 1], where 2 at the first

Picking the Smallest Character: On the first iteration, we look for the smallest character, which is 'a'. We add 'a' to ans,

Picking the Largest Character: We now pick the largest character, which is 'c'. After appending 'c' to ans, the counter

Picking the Smallest Character: We pick 'a' again as it is the next available smallest character. The ans list becomes ['a',

Picking the Largest Character: As per our steps, we continue to pick the largest character left, which is still 'b', updating

Result String Assembly: We initialize an empty list ans to store the characters as we pick them following the algorithm.

The above steps translate the problem's algorithm into Python code in a way that is both efficient (since the actions within the

'c', 'a'] and counter is [0, 2, 0]. Picking the Largest Character: We need to pick the largest character now, which is 'b'. After doing so, the ans list is ['a',

Returning the Result: The while loop exits since len(ans) is now equal to len(s). We join the elements of lans to form the final string. Therefore, the final reordered string is "acabb".

count ensures that the rules of the problem statement are adhered to at every step of the process.

Initialize a list to keep track of the count of each character in the string

Continue until the sorted string's length equals the input string's length

Count the occurrences of each character in the string

Join the list of characters into a string and return it

// Counter array to hold frequency of each character 'a'-'z'

// Fill the frequency array with count of each character

char_count[ord(char) - ord('a')] += 1

Initialize a list to build the sorted string

Solution Implementation **Python**

By following the steps laid out in the solution approach, we successfully applied the algorithm to the example input and achieved

the expected outcome. The method of counting characters, appending the smallest and largest in order, and decrementing their

sorted chars.append(chr(i + ord('a'))) char_count[i] -= 1 # Decrement the count of the added character # Traverse the `char count` list from end to start and add each character once if it's present for i in range(25, -1, -1): if char count[i] > 0: sorted chars.append(chr(i + ord('a')))

Traverse the `char_count` list from start to end and add each character once if it's present

char_count[i] -= 1 # Decrement the count of the added character

```
// StringBuilder to hold the result
StringBuilder sortedString = new StringBuilder();
// Loop until the sortedString's length is less than the original string length
```

class Solution:

Java

class Solution {

def sort string(self, s: str) -> str:

while len(sorted chars) < len(s):</pre>

if char count[i] > 0:

// Method to sort the string in a custom order

while (sortedString.length() < s.length()) {</pre>

// Check if the character is present

for (int i = 0; i < 26; ++i) {

if (frequency[i] > 0) {

for i in range(26):

return ''.join(sorted_chars)

public String sortString(String s) {

int[] frequency = new int[26];

frequency[ch - 'a']++;

// Loop from 'a' to 'z'

for (char ch : s.toCharArray()) {

char count = [0] * 26

for char in s:

sorted_chars = []

```
// Append the character to sortedString
                    sortedString.append((char) ('a' + i));
                    // Decrement the frequency of appended character
                    frequency[i]--;
            // Loop from 'z' to 'a'
            for (int i = 25; i >= 0; --i) {
                // Check if the character is present
                if (frequency[i] > 0) {
                    // Append the character to sortedString
                    sortedString.append((char) ('a' + i));
                    // Decrement the frequency of appended character
                    frequency[i]--;
        // Return the resultant sorted string
        return sortedString.toString();
C++
class Solution {
public:
    // Method to sort the string in a specific pattern
    string sortString(string s) {
        // Create a frequency counter for each letter in the alphabet
        vector<int> frequency(26, 0);
        for (char c : s) {
            ++frequency[c - 'a']; // Increment the count of the current letter
        // Initialize the answer string
        string result = "";
        // Keep building the result until its size matches the original string size
        while (result.size() < s.size()) {</pre>
            // Append characters from 'a' to 'z' to the result string if they are present
            for (int i = 0; i < 26; ++i) {
                if (frequency[i] > 0) { // Check if the character is present
                    result += (i + 'a'); // Convert index to char and append
                    --frequency[i]; // Decrement the frequency of the used character
```

// Append characters from 'z' to 'a' to the result string if they are present

* Sort the string based on the custom order: ascending characters followed by descending characters

--frequency[i]; // Decrement the frequency of the used character

if (frequency[i] > 0) { // Check if the character is present

result += (i + 'a'); // Convert index to char and append

for (int i = 25; i >= 0; --i) { // Start from 'z'

// Return the sorted string

* @return {string} - The sorted string.

function sortString(str: string): string {

let resultString: string = '':

for (let char of str) {

* @param {string} str - The original string to be sorted.

// Count the occurrences of each character in the string

charMap.set(char, (charMap.get(char) || 0) + 1);

const charMap: Map<string, number> = new Map();

return result;

};

/**

TypeScript

```
const keys: string[] = Array.from(charMap.keys());
   keys.sort(); // Sort the keys (characters) in ascending order once
   // Keep constructing the string until the resultString's length equals the input string's length
   while (resultString.length < str.length) {</pre>
       // Append characters in ascending order to the result string
        for (let key of keys) {
           if (charMap.get(key)! > 0) { // Ensure the character count is not zero
                resultString += key;
               charMap.set(key, charMap.get(key)! - 1); // Decrement the count in the map
       // Append characters in descending order to the result string
        for (let i = keys.length - 1; i >= 0; i--) {
           if (charMap.get(keys[i])! > 0) {
               resultString += kevs[i]:
               charMap.set(keys[i], charMap.get(keys[i])! - 1); // Decrement the count in the map
   return resultString;
class Solution:
   def sort string(self, s: str) -> str:
       # Initialize a list to keep track of the count of each character in the string
       char_count = [0] * 26
       # Count the occurrences of each character in the string
       for char in s:
           char_count[ord(char) - ord('a')] += 1
       # Initialize a list to build the sorted string
       sorted_chars = []
       # Continue until the sorted string's length equals the input string's length
       while len(sorted chars) < len(s):</pre>
           # Traverse the `char_count` list from start to end and add each character once if it's present
           for i in range(26):
               if char count[i] > 0:
                   sorted chars.append(chr(i + ord('a')))
                   char_count[i] -= 1 # Decrement the count of the added character
           # Traverse the `char count` list from end to start and add each character once if it's present
           for i in range(25, -1, -1):
               if char count[i] > 0:
                   sorted chars.append(chr(i + ord('a')))
                   char_count[i] -= 1 # Decrement the count of the added character
       # Join the list of characters into a string and return it
       return ''.join(sorted_chars)
```

Time Complexity The provided Python function sortString starts with an initial counting pass over the input string s, incrementing values in

(from 'a' to 'z').

Time and Space Complexity

After that, it enters a loop that continues until the length of ans matches that of s. Within this loop, there are two for-loops: the first iterates in ascending order, the second in descending order. Each of these for-loops iterates over the 26 possible characters

For each character, if that character count is non-zero, it is appended to ans and the count decremented. Since each character

in s is processed exactly once (each is appended and then decremented), and there are two passes for each character (one in

ascending and one in descending order), the total count of operations inside the while-loop is 2n, leading to an additional O(n)

time complexity. Thus, the time complexity of the entire function is O(n).

Space Complexity

counter which takes O(n) time, where n is the length of s.

1. The counter array which is always 26 elements long, thus it is a constant space 0(1).

The space complexity includes:

2. The ans list that will eventually grow to be the same size as s to accommodate all characters, which is O(n). Therefore, the total space complexity is O(n), where n is the length of the input string s.