720. Longest Word in Dictionary

Leetcode Link

Given an array of strings words representing an English Dictionary, return the longest word in words that can be built one character at a time by other words in words.

If there is more than one possible answer, return the longest word with the smallest lexicographical order. If there is no answer, return the empty string.

Note that the word should be built from left to right with each additional character being added to the end of a previous word.

Example 1:

```
input: words = ["w","wo","wor","worl","world"]
Output: "world"
```

Explanation: The word "world" can be built one character at a time by "w", "wo", "wor", and "worl".

Example 2:

```
Input: words = ["a","banana","app","appl","appl","apply","apple"]
```

Output: "apple" Explanation: Both "apply" and "apple" can be built from other words in the dictionary. However, "apple" is lexicographically smaller than "apply".

Constraints:

- $1 \leq words$.length ≤ 1000
- $1 \le words[i]$.length ≤ 30
- words[i] consists of lowercase English letters.

Solution

Brute Force

For this problem, we're asked to find the longest word with smallest lexicographical order such that all its non-empty prefixes exist in words. Let's call a word good if all its prefixes exist in words. We can verify if this is the case by iterating through all prefixes to check if they all exist. Let's denote L as the length of the longest **good** word. Out of all **good** words with length L, we'll return the one with lexicographically least length.

Full Solution

Let's denote x_i as the length of words [i].

We can observe that words [i] is **good** if $x_i=1$ or the prefix of words [i] with length x_i-1 is **good**.

Let's try to find a way to use this idea to process all words efficiently. Since processing a word with length x_i requires a word with length x_i-1 to be processed, we should process words by ${f non-decreasing}$ length. This can be done by sorting and simply iterating through the sorted list. In addition, we'll use a hashmap to act as a lookup table for good words.

In our algorithm, we'll iterate through words by non-decreasing length. For each word, we'll check if it's good with the method mentioned above. If the word is **good**, we'll update it in our <u>hashmap</u>.

Time Complexity

Let's denote N as the length of words and S as the sum of lengths of all words in words.

Since sorting takes $\mathcal{O}(N\log N)$ and our main algorithm takes $\mathcal{O}(S)$ from comparing keys, our final time complexity is $\mathcal{O}(N\log N + S)$.

Time Complexity: $\mathcal{O}(N \log N + S)$

Space Complexity

Since our hashmap has $\mathcal{O}(N)$ memory, our space complexity is $\mathcal{O}(N)$.

Space Complexity: $\mathcal{O}(N)$

C++ Solution

```
class Solution {
       static bool comp(string s, string t) { // sorting comparator
           return s.size() < t.size();</pre>
       string longestWord(vector<string>& words) {
6
            sort(words.begin(), words.end(), comp); // sort words by non-decreasing length
            unordered_map<string, bool> goodWords;
                                                          // lookup for good words
8
            int maxLength = 0;
9
           string ans = "";
10
           for (string word : words) {
                if (word.size() == 1) {
                    goodWords[word] = true;
13
                } else if (goodWords[word.substr(0, word.size() - 1)]) { // word with length - 1 prefix is good
14
                    goodWords[word] = true;
15
16
               if (goodWords[word]) {
17
                    if (maxLength < word.size()) { // find longer word</pre>
                        maxLength = word.size();
19
                        ans = word;
20
21
                    } else if (maxLength == word.size()) { // find lexicographically smaller word
                        ans = min(ans, word);
23
24
25
26
           return ans;
27
28 };
```

Java Solution

```
class Solution {
       public String longestWord(String[] words) {
           Arrays.sort(words, (a, b) -> a.length() - b.length()); // sort words by non-decreasing length
           HashMap<String, Boolean> goodWords = new HashMap(); // lookup for good words
           int maxLength = 0;
           String ans = "";
           for (String word : words) {
                if (word.length() == 1) {
                    goodWords.put(word, true);
               } else if (goodWords.containsKey(word.substring(0, word.length() - 1))) {
                    // word with length - 1 prefix is good
11
                    goodWords.put(word, true);
13
               if (goodWords.containsKey(word)) {
14
                    if (maxLength < word.length()) { // find longer word</pre>
15
                        maxLength = word.length();
16
                        ans = word;
                    } else if (maxLength == word.length()
                        && ans.compareTo(word) > 0) { // find lexicographically smaller word
20
                        ans = word;
23
24
           return ans;
25
26 }
```

Python Solution

class Solution:

Note: A set can be used in python which acts as a hashset and serves the same purpose as a hashmap in this solution.

```
def longestWord(self, words: List[str]) -> str:
           words.sort(key=len) # sort words by non-decreasing length
            goodWords = set() # lookup for good words
           maxLength = 0
           ans = ""
            for word in words:
                if len(word) == 1:
 8
                    goodWords.add(word)
 9
                elif word[:-1] in goodWords: # word with length - 1 prefix is good
10
                    goodWords.add(word)
11
                if word in goodWords:
12
                    if maxLength < len(word): # find longer word</pre>
13
14
                        maxLength = len(word)
                        ans = word
                    elif maxLength == len(word): # find lexicographically smaller word
16
                        ans = min(ans, word)
18
            return ans
19
```