1029. Two City Scheduling

Greedy Array

Problem Description

Medium

2n people to interview and need to send exactly n people to each city. The cost of flying each person to each city is different and is provided in a list known as costs. The costs list is structured such that costs[i] = [aCost_i, bCost_i], with aCost_i representing the cost to fly the ith person to city a and bCost_i representing the cost to fly the ith person to city b. The goal is to determine the minimum possible total cost to fly exactly n people to city a and n people to city b.

A company is conducting interviews and needs to fly interviewees to two different cities, city a and city b. They have a total of

Intuition

The intuition behind the solution is to use a greedy strategy. A greedy algorithm makes the locally optimal choice at each step

with the hope of finding the global optimum. In the context of this problem, we want to minimize the total cost of flying people to

By sorting the costs array based on the difference aCost_i - bCost_i in ascending order, we can easily find the n people to fly to city a (those at the start of the sorted list) and the n people to fly to city b (those at the end of the list). This sorting allows us to consider the relative costs of flying each person to each city and ensure that, overall, we are making the most cost-effective decisions.

in favor of city a, and fly them there. This will ensure that we are saving as much money as possible on the first n flights.

the two cities. To do this, we need to find the n people for whom the difference between flying to city a and city b is the largest

and the last n people to city b. **Solution Approach**

The solution handles the sorting and then computes the total cost by summing up the costs for flying the first n people to city a

The solution approach for the given problem uses a greedy algorithm, which is a common pattern for optimization problems. The

idea here is to make a sequence of choices that are locally optimal in order to reach a solution that is close to or equal to the

Sorting Costs: First, we sort the costs array by the difference in cost between flying to city a and city b (aCost_i -

bCost_i). By doing this, we cluster people who are cheaper to fly to city a at the beginning of the array and those cheaper to

Calculating Half Length: Since we need to send n people to each city and we have 2n people in total, we find n by dividing

Steps Involved:

optimal global solution. Here's a breakdown of how the approach is implemented:

the length of the costs array by 2. Note that we're using a bit shift operation to divide by 2, which is equivalent but slightly faster in most programming languages.

Data Structures Used:

city b towards the end.

n = len(costs) >> 1

costs.sort(key=lambda x: x[0] - x[1])

Determining Minimum Cost: Finally, we take the sum of the costs for flying the first n people to city a and the last n people to city b, which gives us the minimum cost. sum(costs[i][0] + costs[i + n][1] for i in range(n))

Algorithms & Patterns Used:

based on the local optimality of the cost difference.

individual costs to minimize the total expenditure.

person to cities a and b are provided as follows:

costs = [[400, 200], [300, 600], [100, 500], [200, 300]]

costs = [[100, 500], [300, 600], [200, 300], [400, 200]]

For city a: costs[0][0] + costs[1][0] = 100 + 300 = 400

def two city sched cost(self, costs: List[List[int]]) -> int:

Compute the number of people that should go to each city.

of the people to city A and the remaining half to city B.

import java.util.Arrays; // Import the Arrays class to use the sort method

Arrays.sort(costs, $(a, b) \rightarrow (a[0] - a[1]) - (b[0] - b[1]));$

int totalCost = 0: // Initialize total cost to 0

return totalCost; // Return the accumulated total cost

// Function to calculate the minimum cost to send n people to two cities.

// Sort the array of costs based on the difference in cost between sending

// Calculate half the length of the costs array since we're splitting the group in two

// Iterate over the first half of the sorted array for city A costs and the second

// Add the cost of sending the i-th person to city A from the first half

// Add the cost of sending the i-th person to city B from the second half

Sort the costs list based on the cost difference between the two cities (A and B).

Calculate the total minimum cost by adding the cost of sending the first half

This works because the costs array is now sorted in a way that the first half

sum(costs[i + num_people][1] for i in range(num_people))

This helps in figuring out the cheaper city for each person in a way that can minimize the total cost.

costs.sort((firstPair, secondPair) => (firstPair[0] - firstPair[1]) - (secondPair[0] - secondPair[1]));

Since it's a two-city schedule, half the people will go to each city.

Calculate the total minimum cost by adding the cost of sending the first half

This works because the costs array is now sorted in a way that the first half

// Sort the array based on the cost difference between the two cities (a and b)

// for each person. The comparator subtracts the cost of going to city B from the

// cost of going to city A for each person, and sorts based on these differences.

costs.sort(key=lambda cost: cost[0] - cost[1])

Sort the costs list based on the cost difference between the two cities (A and B).

num_people = len(costs) // 2 # Using // for floor division which is standard in Python 3.

• A list of lists is used to store the flying costs to both cities for each person.

• List Comprehension: Python list comprehension is used to sum the costs which result in a concise and readable implementation.

Complexity Analysis:

Example Walkthrough

• The space complexity is 0(1) or constant space, aside from the input, since no additional data structures are used that grow with the size of the input. The algorithm implemented here effectively balances the costs for both cities by understanding and utilizing the differences in

• Greedy Algorithm: The algorithm sorts the array and then selects the first n elements for one city and the last n elements for the other city,

• Sorting: A key step in the greedy approach where the array is sorted based on a specific condition (cost difference here).

• The time complexity of the solution is dominated by the sorting step, which is 0(n log n) where n is the number of people.

Let's consider an example where a company has 2n = 4 people (n = 2) to send to cities a and b. The costs for flying each

After sorting these based on the difference in ascending order, our costs array looks like this:

Now, let's walk through the solution approach step by step using this example: Sorting Costs First, we sort the costs array based on the difference between flying to city a and city b, so we calculate aCost_i - bCost_i for each person:

Person 3: 100 - 500 = -400Person 4: 200 - 300 = -100

Person 1: 400 - 200 = 200

Person 2: 300 - 600 = -300

len(costs) >> 1.

Solution Implementation

cost.

class Solution:

Example usage:

class Solution {

Java

C++

public:

#include <vector>

class Solution {

TypeScript

solution = Solution()

costs = [[10,20],[30,200],[400,50],[30,20]]

public int twoCitySchedCost(int[][] costs) {

#include <algorithm> // Include necessary headers

function twoCitvSchedCost(costs: number[][]): number {

// Initialize an accumulator for the total cost

totalCost += costs[i + halfLength][1];

// half for city B costs to compute the minimum total cost

def two city sched cost(self, costs: List[List[int]]) -> int:

of the people to city A and the remaining half to city B.

total cost = sum(costs[i][0] for i in range(num people)) + \

to city B, and vice versa for the second half.

are the people who have a cheaper cost going to city A compared

// a person to city A and city B.

const halfLength = costs.length / 2;

for (let i = 0; i < halfLength; ++i) {</pre>

totalCost += costs[i][0];

// Return the computed total cost

let totalCost = 0;

return totalCost;

from typing import List

return total_cost

cities (N to each city).

factors are dominant for large n.

costs = [[10.201, [30.200], [400, 50], [30, 20]]]

Now, the first two people on the list are the ones we will fly to city a, and the last two to city b. Calculating Half Length Here, n is already given as 2 (since 2n = 4). In a more general sense, we would compute n as

Determining Minimum Cost Lastly, we will add the cost for flying the first n people to city a and the last n people to city b:

```
For city b: costs[2][1] + costs[3][1] = 300 + 200 = 500
 The total minimum cost to fly all these people is the sum of the costs for city a and city b: 400 + 500 = 900.
```

Through this example, we can see how the company efficiently reduces the total cost for flying out interviewees to two different

cities by implementing a greedy algorithm that focuses on sending each person to the city that results in the lowest additional

Python from typing import List

This helps in figuring out the cheaper city for each person in a way that can minimize the total cost.

are the people who have a cheaper cost going to city A compared # to city B, and vice versa for the second half. total cost = sum(costs[i][0] for i in range(num people)) + \ sum(costs[i + num_people][1] for i in range(num_people)) return total_cost

print(solution.two_city_sched_cost(costs)) # Should print the result of the cost calculation based on the provided costs array.

int n = costs.length / 2; // Calculate n, where n is half the number of people (half will go to city A, half to city B)

```
// Accumulate the minimum cost by sending the first n people to city A
// and the remaining n people to city B based on the sorted order
for (int i = 0; i < n; ++i) {
    totalCost += costs[i][0]; // Add cost for city A
    totalCost += costs[i + n][1]; // Add cost for city B
```

```
int twoCitySchedCost(vector<vector<int>>& costs) {
       // Sort the input vector based on the cost difference for attending city A and city B
        sort(costs.begin(), costs.end(), [](const vector<int>& a, const vector<int>& b) {
            return a[0] - a[1] < b[0] - b[1];
       });
       // The total number of people is twice the number of people we need to send to one city
       int totalPeople = costs.size();
       // Calculate the number of people to send to each city
       int halfPeople = totalPeople / 2;
       // Initialize answer to store the total minimized cost
       int totalCost = 0;
       // Calculate the minimum total cost
        for (int i = 0; i < halfPeople; ++i) {</pre>
           // Add up the cost of sending the first half of people to city A
            totalCost += costs[i][0];
           // Add up the cost of sending the second half of people to city B
            totalCost += costs[i + halfPeople][1];
       // Return the computed total cost
       return totalCost;
};
```

costs.sort(key=lambda cost: cost[0] - cost[1]) # Compute the number of people that should go to each city. # Since it's a two-city schedule, half the people will go to each city. $num_people = len(costs) // 2 # Using // for floor division which is standard in Python 3.$

Example usage:

solution = Solution()

class Solution:

- Time and Space Complexity **Time Complexity** The time complexity of the provided code consists of the following parts: Sorting the costs list: Since the sort() function in Python uses the Timsort algorithm, which has a time complexity of O(n
- iterations (since n = len(costs) >> 1 computes to N). The summing operation inside the list comprehension is O(1) for each element, so the overall time for this part is O(N). Combining these two parts, the total time complexity is 0(n log n) + 0(N), which simplifies to 0(n log n) since logarithmic

log n), where n is the total number of elements in the list. Here n is 2N since we're supposed to send 2N people to two

Summing up the costs: The list comprehension iterates over the first half and the second half of the sorted list, resulting in N

print(solution.two_city_sched_cost(costs)) # Should print the result of the cost calculation based on the provided costs array.

Space Complexity

Hence, the time complexity of the code is $O(n \log n)$.

The space complexity of the code analysis involves:

with the size of the input.

The list comprehension for summing the costs and the range used within: both utilize only a small constant amount of additional space.

The sorted list costs: Sorting is done in-place, so it does not consume additional space proportional to the input size.

- Therefore, the space complexity of the entire operation is 0(1), as no significant additional space is used that would increase
- Overall, the space complexity of the code is 0(1).