Medium **Binary Indexed Tree** Simulation Array **Problem Description** In this problem, we are given two pieces of information: 1. An array of queries containing positive integers between 1 and m. 2. We initially have a permutation P consisting of all integers from 1 to m in ascending order.

We are asked to return an array that contains the result for each queries[i] after processing all the queries.

To solve this problem, our approach focuses on simulating the process as described in the problem statement. Here is the

The objective is to process each query by following these steps: • Identify the position of the current query element (let's call it queries[i]) in the permutation P. This position is to be considered using 0-based indexing.

Record the position of queries[i] as this is our result for the current query.

• Move the element queries[i] from its current position to the beginning of the permutation P.

1409. Queries on a Permutation With Key

thinking process for arriving at the solution: **Initial Setup**: We start by generating the initial permutation P which is simply a list of integers from 1 to m.

Processing Queries: For each value v in queries, we need to perform the following operations:

• Find Position: Locate the index j of v within list P which represents the initial position of v in the permutation. Record Result: The index j is the answer for this query, so it needs to be added to an answer array ans.

 Update Permutation: We then remove v from its current position in P and insert it at the beginning of P. Maintain Permutation State: Each iteration modifies the permutation P according to the specified rules. Hence, the state of P is always maintained after processing each query.

Solution Approach

The implementation of the solution can be walked through as follows using Python's list data structure and some of its built-in methods:

generates an iterable from 1 to m inclusive. The list function is then called to convert this iterable into a list.

List Creation: We begin by creating the list P that represents the permutation. This is done using range(1, m + 1) which

p = list(range(1, m + 1))

Iterating Through Queries: We iterate through each element v in the list of queries with a for loop. for v in queries:

Finding Query Position: The index() method of lists is used to find the position j of the value v within the permutation P.

As the list is 0-indexed, this will give us the index starting from 0.

list contains the original positions of each query value before we moved them.

def processQueries(self, queries: List[int], m: int) -> List[int]:

• We move 3 to the front of P resulting in the new permutation [3, 1, 2, 4, 5].

Position of 1 is now 1 (it moved because of the previous query).

o 1 moves to the front, resulting in [1, 3, 2, 4, 5].

j = p.index(v) **Recording the Position**: We append the found position j to our list ans which will eventually be returned as our result array.

Updating the Permutation: To move the query value to the beginning of P, we first remove v from its current position using pop(j) where j is the index found earlier.

p.pop(j) Then, we insert the value v at the beginning of P by using the insert() function with 0 as the index to insert at.

p.insert(0, v)

ans.append(j)

This solution is efficient because accessing an element's index, removing an element, and inserting an element at the beginning of a list all have a time complexity of O(n), making the overall complexity of the algorithm O(n * q) where n is the number of elements in the permutation and q is the number of queries, since each query is processed independently.

The space complexity of the solution is O(m) where m is the size of the permutation, since it's the space required to store the

Returning the Result: After the loop has finished processing all the elements in queries, our answer list ans is returned. This

permutation.

class Solution:

ans = []

p = list(range(1, m + 1))

j = p.index(v)

ans.append(j)

p.insert(0, v)

for v in queries:

p.pop(j)

return ans

• We record this position.

The next query is 1.

• We record this position.

The third query is 2.

The last query is 1.

We record the position.

• queries: [3, 1, 2, 1]

Example Walkthrough Let's say we're given the following inputs:

• m:5 The initial permutation P from 1 to m (5 in this case) is [1, 2, 3, 4, 5]. Now we process each query one by one: The first query is 3. • We find the position of 3 in P, which is 2.

Position of 2 is 2. • We record the position. o 2 moves to the front, permutation is [2, 1, 3, 4, 5].

Python

class Solution:

results = []

from typing import List

class Solution {

for value in queries:

results.append(index)

p sequence.pop(index)

Solution Implementation

p sequence = list(range(1, m + 1))

Each position we recorded forms our result: [2, 1, 2, 1].

Position of 1 is 1 because it was moved to the front earlier.

1 is already at the front, so the permutation remains [2, 1, 3, 4, 5].

def processQueries(self, queries: List[int], m: int) -> List[int];

Find the index of the current value in the P sequence

The code assumes the existence of `List` type hint imported from `typing` module

// Find the index of the queried number in the permutation

// Return the final array of indices representing the answer

vector<int> processQueries(vector<int>& queries, int m) {

// Loop over each value in the queries.

for (int i = 0; i < m; ++i) {

answer.push_back(foundIndex);

// Return the results of the queries.

// Loop over each value in the queries array.

permutation.splice(foundIndex, 1);

return answer;

let answer: number[] = [];

gueries.forEach(value => {

results = []

return results

from typing import List

for value in queries:

results.append(index)

p sequence.pop(index)

p sequence.insert(0, value)

answer.push(foundIndex);

permutation.unshift(value);

// Return the results of the queries.

if (permutation[i] == value) {

// Add the found index to the answer vector.

// Erase the value from its current position.

permutation.insert(permutation.begin(), value);

function processQueries(queries: number[], m: number): number[] {

// Find the index of the 'value' in 'permutation'.

let foundIndex = permutation.indexOf(value);

// Add the found index to the 'answer' array.

// Remove the value from its current position.

Initialize the answer list to store the results

Process each query from the queries list

Append the index to the results list

Return the results list containing the indices

Remove the current value from its index in P

index = p sequence.index(value)

// Initialize the answer array to store the results of the queries.

permutation.erase(permutation.begin() + foundIndex);

std::iota(permutation.begin(), permutation.end(), 1);

// Search for the value in the permutation array.

// Function to process the gueries and return the indices of each query in the permutation

// Initialize P as a LinkedList to easily support element removal and insertion at the front

If not present in the original code file, it should be added at the top as:

Initialize the P sequence with integers from 1 to m

Remove the current value from its index in P

Initialize the answer list to store the results

Process each query from the queries list

index = p sequence.index(value)

Thus, the final answer after processing all queries is the list of recorded positions [2, 1, 2, 1].

Insert the current value at the beginning of the P sequence p sequence.insert(0, value) # Return the results list containing the indices return results

public int[] processQueries(int[] queries, int m) {

// Fill permutation with elements 1 to m

// Initialize index for placing answers

// Process each query in the array

for (int query : queries) {

for (int num = 1; num <= m; num++) {</pre>

permutation.add(num);

int ansIndex = 0;

return indices;

#include <numeric> // For std::iota

vector<int> answer;

vector<int> permutation(m);

for (int value : gueries) {

int foundIndex = 0;

List<Integer> permutation = new LinkedList<>();

Append the index to the results list

```
// Array to store the answer (indices of each queried element)
int[] indices = new int[queries.length];
```

C++

public:

#include <vector>

class Solution {

Java

int queryIndex = permutation.indexOf(query); // Store the index in the result answer array indices[ansIndex++] = queryIndex; // Remove the queried number from its current position permutation.remove(queryIndex); // Add the gueried number to the front of the permutation permutation.add(0, query);

// This function processes the queries on the permutation array and returns the result.

foundIndex = i; // Store index where value is found.

break; // Exit the loop since we found the value.

// Insert the value at the beginning of the permutation array.

// Initialize the permutation array 'permutation' with elements from 1 to 'm'.

let permutation: number[] = Array.from({ length: m }, (_, index) => index + 1);

// Initialize an index 'foundIndex' to store the position of the value in 'permutation'.

// Initialize the permutation array 'P' with elements from 1 to 'm'.

// Initialize the answer vector to store the results of the queries.

};

TypeScript

});

return answer; class Solution: def processQueries(self, queries: List[int], m: int) -> List[int]: # Initialize the P sequence with integers from 1 to m p sequence = list(range(1, m + 1))

Find the index of the current value in the P sequence

Insert the current value at the beginning of the P sequence

The code assumes the existence of `List` type hint imported from `typina` module

If not present in the original code file, it should be added at the top as:

// Insert the value at the beginning of the permutation array.

Time and Space Complexity **Time Complexity** The time complexity of the provided code primarily depends on two factors:

2. The cost of popping and inserting elements to/from the list which can take up to O(m) time.

Since for every value in queries we perform both indexing and pop-insert operations, we multiply this cost by the number of queries n. Therefore, the time complexity is O(n*m), where n is the length of queries.

1. The cost of searching for an index of a value in the list p which is done in O(m) time where m is the length of p.

Space Complexity The space complexity of the algorithm is to consider the additional space used by the algorithm excluding the input and output.

Here, aside from the space used by the input queries and the output list ans, the code maintains a list p of size m, but since this does not grow with the size of the input queries, the additional space remains constant. Thus, the space complexity is 0(1), which means it is constant space complexity as it doesn't depend on the size of the input queries.