1409. Queries on a Permutation With Key Medium **Binary Indexed Tree** Simulation <u>Array</u> **Problem Description**

In this problem, we are given two pieces of information:

2. We initially have a permutation P consisting of all integers from 1 to m in ascending order. The objective is to process each query by following these steps:

1. An array of queries containing positive integers between 1 and m.

• Identify the position of the current query element (let's call it queries[i]) in the permutation P. This position is to be considered using 0-based indexing.

Record the position of queries[i] as this is our result for the current query.

• Move the element queries [i] from its current position to the beginning of the permutation P. We are asked to return an array that contains the result for each queries [i] after processing all the queries.

To solve this problem, our approach focuses on simulating the process as described in the problem statement. Here is the thinking process for arriving at the solution:

Initial Setup: We start by generating the initial permutation P which is simply a list of integers from 1 to m. **Processing Queries:** For each value v in queries, we need to perform the following operations:

• Find Position: Locate the index j of v within list P which represents the initial position of v in the permutation. Record Result: The index j is the answer for this query, so it needs to be added to an answer array ans.

- Update Permutation: We then remove v from its current position in P and insert it at the beginning of P. Maintain Permutation State: Each iteration modifies the permutation P according to the specified rules. Hence, the state of P
- is always maintained after processing each query.

Iterating Through Queries: We iterate through each element v in the list of queries with a for loop.

- Solution Approach
- The implementation of the solution can be walked through as follows using Python's list data structure and some of its built-in methods:
- List Creation: We begin by creating the list P that represents the permutation. This is done using range(1, m + 1) which generates an iterable from 1 to m inclusive. The list function is then called to convert this iterable into a list.

for v in queries:

j = p.index(v)

p = list(range(1, m + 1))

pop(j) where j is the index found earlier.

p = list(range(1, m + 1))

for v in queries:

Finding Query Position: The index() method of lists is used to find the position j of the value v within the permutation P. As the list is 0-indexed, this will give us the index starting from 0.

Recording the Position: We append the found position j to our list ans which will eventually be returned as our result array. ans.append(j)

Updating the Permutation: To move the query value to the beginning of P, we first remove v from its current position using

p.pop(j)

p.insert(0, v)

Returning the Result: After the loop has finished processing all the elements in queries, our answer list ans is returned. This list contains the original positions of each query value before we moved them.

The space complexity of the solution is O(m) where m is the size of the permutation, since it's the space required to store the

Then, we insert the value v at the beginning of P by using the insert() function with 0 as the index to insert at.

elements in the permutation and q is the number of queries, since each query is processed independently.

def processQueries(self, queries: List[int], m: int) -> List[int]:

• We move 3 to the front of P resulting in the new permutation [3, 1, 2, 4, 5].

Position of 1 is now 1 (it moved because of the previous query).

1 moves to the front, resulting in [1, 3, 2, 4, 5].

Position of 1 is 1 because it was moved to the front earlier.

Each position we recorded forms our result: [2, 1, 2, 1].

This solution is efficient because accessing an element's index, removing an element, and inserting an element at the beginning of a list all have a time complexity of O(n), making the overall complexity of the algorithm O(n * q) where n is the number of

permutation.

class Solution:

ans = []

return ans

• queries: [3, 1, 2, 1]

j = p.index(v) ans.append(j) p.pop(j) p.insert(0, v)

Example Walkthrough Let's say we're given the following inputs:

```
• m: 5
The initial permutation P from 1 to m (5 in this case) is [1, 2, 3, 4, 5].
Now we process each query one by one:
    The first query is 3.
    • We find the position of 3 in P, which is 2.
    • We record this position.
```

The third query is 2.

The next query is 1.

We record this position.

• We record the position.

Solution Implementation

results = []

return results

from typing import List

Java

for value in queries:

results.append(index)

class Solution:

 2 moves to the front, permutation is [2, 1, 3, 4, 5]. The last query is 1.

Position of 2 is 2.

- We record the position. 1 is already at the front, so the permutation remains [2, 1, 3, 4, 5].
- **Python**

Initialize the P sequence with integers from 1 to m

Initialize the answer list to store the results

Process each query from the queries list

Append the index to the results list

Return the results list containing the indices

index = p_sequence.index(value)

p_sequence = list(range(1, m + 1))

def processQueries(self, queries: List[int], m: int) -> List[int]:

Find the index of the current value in the P sequence

The code assumes the existence of `List` type hint imported from `typing` module

If not present in the original code file, it should be added at the top as:

Remove the current value from its index in P p_sequence.pop(index) # Insert the current value at the beginning of the P sequence p_sequence.insert(0, value)

// This function processes the queries on the permutation array and returns the result.

foundIndex = i; // Store index where value is found.

break; // Exit the loop since we found the value.

// Insert the value at the beginning of the permutation array.

// Initialize an index 'foundIndex' to store the position of the value in 'permutation'.

// Initialize the permutation array 'P' with elements from 1 to 'm'.

// Initialize the answer vector to store the results of the queries.

vector<int> processQueries(vector<int>& queries, int m) {

// Loop over each value in the queries.

for (int i = 0; i < m; ++i) {

answer.push_back(foundIndex);

// Loop over each value in the queries array.

permutation.splice(foundIndex, 1);

if (permutation[i] == value) {

// Add the found index to the answer vector.

// Erase the value from its current position.

permutation.insert(permutation.begin(), value);

permutation.erase(permutation.begin() + foundIndex);

// Initialize the answer array to store the results of the queries.

// Insert the value at the beginning of the permutation array.

// Find the index of the 'value' in 'permutation'.

let foundIndex = permutation.indexOf(value);

// Add the found index to the 'answer' array.

// Remove the value from its current position.

std::iota(permutation.begin(), permutation.end(), 1);

// Search for the value in the permutation array.

Thus, the final answer after processing all queries is the list of recorded positions [2, 1, 2, 1].

// Function to process the queries and return the indices of each query in the permutation public int[] processQueries(int[] queries, int m) { // Initialize P as a LinkedList to easily support element removal and insertion at the front List<Integer> permutation = new LinkedList<>();

class Solution {

#include <vector>

class Solution {

public:

#include <numeric> // For std::iota

vector<int> answer;

vector<int> permutation(m);

for (int value : queries) {

int foundIndex = 0;

```
// Fill permutation with elements 1 to m
        for (int num = 1; num <= m; num++) {</pre>
           permutation.add(num);
       // Array to store the answer (indices of each queried element)
        int[] indices = new int[queries.length];
       // Initialize index for placing answers
       int ansIndex = 0;
       // Process each query in the array
        for (int query : queries) {
           // Find the index of the queried number in the permutation
           int queryIndex = permutation.indexOf(query);
           // Store the index in the result answer array
            indices[ansIndex++] = queryIndex;
            // Remove the queried number from its current position
            permutation.remove(queryIndex);
            // Add the queried number to the front of the permutation
           permutation.add(0, query);
       // Return the final array of indices representing the answer
       return indices;
C++
```

```
// Return the results of the queries.
        return answer;
};
```

TypeScript

});

class Solution:

return answer;

results = []

from typing import List

Time Complexity

for value in queries:

results.append(index)

p_sequence.pop(index)

Time and Space Complexity

```
function processQueries(queries: number[], m: number): number[] {
   // Initialize the permutation array 'permutation' with elements from 1 to 'm'.
    let permutation: number[] = Array.from({ length: m }, (_, index) => index + 1);
```

let answer: number[] = [];

queries.forEach(value => {

answer.push(foundIndex);

permutation.unshift(value);

// Return the results of the queries.

p_sequence = list(range(1, m + 1))

```
# Insert the current value at the beginning of the P sequence
            p_sequence.insert(0, value)
       # Return the results list containing the indices
        return results
# The code assumes the existence of `List` type hint imported from `typing` module
```

If not present in the original code file, it should be added at the top as:

Find the index of the current value in the P sequence

def processQueries(self, queries: List[int], m: int) -> List[int]:

Initialize the P sequence with integers from 1 to m

Remove the current value from its index in P

Initialize the answer list to store the results

Process each query from the queries list

Append the index to the results list

index = p_sequence.index(value)

The time complexity of the provided code primarily depends on two factors: 1. The cost of searching for an index of a value in the list p which is done in O(m) time where m is the length of p. 2. The cost of popping and inserting elements to/from the list which can take up to O(m) time.

```
queries n. Therefore, the time complexity is O(n*m), where n is the length of queries.
Space Complexity
```

Since for every value in queries we perform both indexing and pop-insert operations, we multiply this cost by the number of

The space complexity of the algorithm is to consider the additional space used by the algorithm excluding the input and output. Here, aside from the space used by the input queries and the output list ans, the code maintains a list p of size m, but since this does not grow with the size of the input queries, the additional space remains constant. Thus, the space complexity is 0(1), which means it is constant space complexity as it doesn't depend on the size of the input queries.