Given a binary tree, where each node contains an integer, we are asked to find the largest absolute difference in value between two nodes where one node is an ancestor of the other. In other words, if we pick any node a as an ancestor, and any node b as a

Problem Description

An important detail to note is that a node can be considered an ancestor of itself, leading to a minimum absolute difference of 0 in such a scenario. The problem is focusing on finding the maximum difference, hence we need to look for pairs of ancestor and descendant nodes where this difference is the largest.

descendant of a, what is the maximum absolute difference between a, val and b, val ([a, val - b, val]) that we can find in the tree?

Intuition

global maximum ans if we find a larger difference.

using a depth-first search (DFS) algorithm, which will allow us to explore each branch of the tree to its fullest extent before moving on to the next branch. During the traversal, we keep track of the minimum (mi) and maximum (mx) values encountered along the path from the root to the

current node. At each node, we calculate the absolute difference between root val and both the mi and mx values, updating the

The intuition behind the solution is that we can find the maximum difference by thoroughly searching through the tree. We do this

The core idea is to track the range of values (minimum and maximum) on the path from the root to the current node because this range will allow us to compute the required maximum absolute difference at each step. By the time we complete our traversal, we will have examined all possible pairs of ancestor and descendant nodes and thus found the maximum difference.

To implement this, we use a recursive helper function dfs(root, mi, mx) that performs a depth-first search on the binary tree. The mi and mx parameters keep track of the minimum and maximum values respectively, seen from the root to the current node. The function also updates a nonlocal variable ans, which keeps track of the maximum difference found so far.

Finally, we initiate our DFS with the root node and its value as both the initial minimum and maximum, and after completing the traversal, we return the value stored in ans, which will be the maximum ancestor-difference that we were tasked to find. Solution Approach

The solution to this problem involves a recursive depth-first search (DFS) algorithm to traverse the binary tree. The critical aspect of

the approach is to maintain two variables, mi and mx, to record the minimum and maximum values found along the path from the root

2. If the current root is None, which means we've reached a leaf node's child, we return, as there are no more nodes to process in

3. The helper function is designed to continuously update a nonlocal variable ans, which holds the maximum absolute difference

node to the current node. Here is a step-by-step breakdown of the implementation details: 1. Define a recursive helper function dfs(root, mi, mx) that will be used for DFS traversal of the tree.

4. At each node, we compare and update ans with the absolute difference of the current node's value root.val with both the minimum (mi) and maximum (mx) values seen so far along the path from the root.

this path.

found.

5. We perform this comparison using max(ans, abs(mi - root.val), abs(mx - root.val)). 6. After updating ans, we also update mi and mx for the recursive calls on the children nodes, setting mi to min(mi, root.val) and mx to max(mx, root.val). This ensures that as we go deeper into the tree, our range [mi, mx] remains updated with the smallest and largest values seen along the path.

7. Recursive calls are then made to continue the DFS traversal on the left child dfs(root.left, mi, mx) and the right child dfs(root.right, mi, mx) of the current node.

Let's consider a small binary tree to illustrate the solution approach. Our binary tree is as follows:

1. The dfs function is first called with root.val = 8, mi = 8, mx = 8. We are at the root.

result—the maximum difference. The function finally returns ans.

minimum and maximum of the path consisting of just itself.

2. Explore left child (3). Call dfs(3, min(8,3), max(8,3)):

 $- \max(5, abs(1 - 8), abs(8 - 1))$

4. The other child of 3 is 6. Call dfs(6, min(3,6), max(8,6)):

Node 1 is a leaf; the traversal will go back up.

= max(7, abs(3 - 6), abs(8 - 6))

= max(7, abs(3 - 4), abs(8 - 4))

 \circ Now mi = 3, mx = 8.

 \circ Here, mi = 1, mx = 8.

Update answer:

 \bullet ans = 7

 \circ mi = 3, mx = 8.

Update answer:

 \circ mi = 3, mx = 8.

Update answer:

 \circ mi = 3, mx = 8.

Update answer:

 \circ mi = 8, mx = 10.

Update answer:

ans remains 7.

The main function initializes the variable ans to 0 and then calls dfs(root, root.val, root.val). We start with both mi and mx as the root's value, since initially, the root is the only node in the path. The implementation leverages the default argument-passing mechanism in Python, where every child node receives the current path's minimum and maximum values to keep the comparison going.

After the completion of the DFS traversal, the ans variable, which was kept up-to-date during the traversal, will contain the final

This efficient use of space and recursive traversal makes it a neat and effective solution. Example Walkthrough

We want to find the maximum absolute difference between the values of any two nodes where one is an ancestor of the other.

We begin by calling the recursive function dfs on the root node with value 8. We start with mi = mx = 8 since the root is both the

The primary data structure used in this implementation is the binary tree itself. No additional data structures are needed because the

recursion stack implicitly manages the traversal, and the updating of minimum and maximum values is done using integer variables.

 Update potential answer compare with previous ans: $- \max(0, abs(3 - 8), abs(8 - 3))$ ans = 5 3. Go down to the left child of 3, node 1. Call dfs(1, min(3,1), max(8,1)):

- ans remains 7. 5. Node 6 has a left child 4. Call dfs(4, min(3,4), max(8,4)):
- ans remains 7. Since 4 is a leaf node, we go back up.

7. Now explore right child (10) of the root. Call dfs(10, min(8,10), max(8,10)):

■ ans becomes 14 - 8 = 6 but since our current ans = 7, there is no update.

ans remains 7 because the differences 5 and 1 are smaller than the current ans.

6. Node 6 also has right child 7. Call dfs(7, min(3,7), max(8,7)):

= max(7, abs(3 - 7), abs(8 - 7))ans remains 7. Node 7 is also a leaf; traverse back up to 3, then to 8.

 $- \max(7, abs(8 - 10), abs(10 - 8))$

 $- \max(7, abs(8 - 13), abs(14 - 13))$

def __init__(self, val=0, left=None, right=None):

new_min = min(min_value, node.val)

new_max = max(max_value, node.val)

dfs(node.left, new_min, new_max)

dfs(node.right, new_min, new_max)

Recursively call dfs for the left and right subtrees

Start DFS from root with its value as both initial min and max

Initialize max difference which will hold the result

Node value

8. Node 10 has right child 14. Call dfs(14, min(8,14), max(10,14)):

9. Node 14 has a left child 13. Call dfs(13, min(8,13), max(14,13)):

- \circ mi = 8, mx = 14. Update answer: $- \max(7, abs(8 - 14), abs(14 - 8))$
- After the recursive depth-first search completes, we find that the maximum absolute difference is 7, which comes from the difference between nodes 8 (ancestor) and 1 (descendant).

A class for a binary tree node

self.val = val

self.left = left

max difference = 0

return max_difference

* Definition for a binary tree node.

dfs(root, root.val, root.val)

Return the maximum difference found

TreeNode(int val, TreeNode left, TreeNode right) {

Python Solution

class TreeNode:

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C++ Solution

struct TreeNode {

int val;

13 class Solution {

TreeNode *left;

TreeNode *right;

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 \circ mi = 8, mx = 14.

Update answer:

TreeNode() {} 10 TreeNode(int val) { 11 this.val = val; 13

class TreeNode {

int val;

TreeNode left;

TreeNode right;

this.val = val;

this.left = left;

this.right = right;

return maxDifference;

if (node == null) {

return;

* @param node The current node being visited.

minVal = Math.min(minVal, node.val);

maxVal = Math.max(maxVal, node.val);

* Definition for a binary tree node.

int maxAncestorDiff(TreeNode* root) {

return;

});

maxDifference = max({

maxDifference,

dfs(root, root->val, root->val);

// Return the maximum difference found

abs(currentMin - node->val),

currentMin = min(currentMin, node->val);

currentMax = max(currentMax, node->val);

dfs(node->left, currentMin, currentMax);

dfs(node->right, currentMin, currentMax);

abs(currentMax - node->val)

// Recur for both the left and right subtrees.

depthFirstSearch(node.left, minVal, maxVal);

depthFirstSearch(node.right, minVal, maxVal);

TreeNode(): val(0), left(nullptr), right(nullptr) {}

TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}

TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left), right(right) {}

// If the node is null, return as there is nothing to process

// Update currentMin and currentMax with respective minimum and maximum values

// Update maxDifference with the maximum of the current difference

// and the differences with the current node's value

// Continue depth-first search on left and right subtrees

// Initialize DFS with the value of the root for both min and max

/* Function to calculate max difference between any ancestor and node value. */

Java Solution

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*/

Left child self.right = right # Right child class Solution: def maxAncestorDiff(self, root: Optional[TreeNode]) -> int: 9 # Helper function to perform Depth-First Search (DFS) 10 11 def dfs(node, min_value, max_value): # Base case: if the current node is None, return 12 13 if node is None: 14 return 15 16 # Calculate the maximum difference for the current node nonlocal max_difference 17 current_diff = max(abs(min_value - node.val), abs(max_value - node.val)) 18 max_difference = max(max_difference, current_diff) 19 20 21 # Update the minimum and maximum values with the value of the current node 22

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class Solution {
       private int maxDifference;
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        * Calculates the maximum difference between values of any two connected nodes in the binary tree.
        * @param root The root of the binary tree.
        * @return The maximum difference calculated.
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        public int maxAncestorDiff(TreeNode root) {
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            if (root == null) {
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                return 0;
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           // Start DFS with the initial value of the root for both minimum and maximum.
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           depthFirstSearch(root, root.val, root.val);
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// Calculate the potential differences between the current node value and the observed min and max values.

int currentMaxDifference = Math.max(Math.abs(minVal - node.val), Math.abs(maxVal - node.val));

* A recursive DFS function that traverses the tree to find the maximum difference.

private void depthFirstSearch(TreeNode node, int minVal, int maxVal) {

// Update the maxDifference if the current one is greater.

maxDifference = Math.max(maxDifference, currentMaxDifference);

// Update the min and max values to carry them forward in the DFS.

* @param minVal The minimum value seen so far in the path from root to the current node.

* @param maxVal The maximum value seen so far in the path from root to the current node.

17 int maxDifference = 0; // to store the maximum difference 18 19 // Lambda function for depth-first search starting from 'node' 20 // It carries the current minimum and maximum values as 'currentMin' and 'currentMax' 21 function<void(TreeNode*, int, int)> dfs = [&](TreeNode* node, int currentMin, int currentMax) { 22 if (node == nullptr) {

public:

47 return maxDifference; 48 49 }; 50 Typescript Solution // Global variable to track the maximum difference between an ancestor and a node value let maxDifference: number = 0; // Recursive function traverses the tree to find the maximum difference between an ancestor and a node value function dfs(node: TreeNode | null, minVal: number, maxVal: number): void { if (!node) { // If node is null, return because we've reached a leaf node's child. return; 9 10 // Calculate the potential new max differences with the current node 11 12 const potentialMaxDiff = Math.max(Math.abs(node.val - minVal), Math.abs(node.val - maxVal)); 13 // Update the global maxDifference if the new potential difference is greater 14 maxDifference = Math.max(maxDifference, potentialMaxDiff); 16 17 // Update the min and max values seen so far after considering the current node's value const newMinVal = Math.min(minVal, node.val); 18 const newMaxVal = Math.max(maxVal, node.val); 19 20 // Continue the DFS traversal for left and right children 21 22 dfs(node.left, newMinVal, newMaxVal); 23 dfs(node.right, newMinVal, newMaxVal); 24 } 25 // Primary function to initiate the maxAncestorDiff calculation, given the root of a binary tree function maxAncestorDiff(root: TreeNode | null): number { if (root === null) { // If the tree is empty, the maximum difference is 0 by definition return 0; 31 33 // Since we start at the root, the starting min and max values are the root's value dfs(root, root.val, root.val); 34 35 // After traversing the tree, return the global maxDifference found 36 37 return maxDifference;

ancestor/descendant relationship in a binary tree. Time Complexity:

Time and Space Complexity

The time complexity of the function is O(N), where N is the number of nodes in the binary tree. This is because the function performs a depth-first search (DFS), visiting each node exactly once. During each visit, it performs a constant amount of work by updating the minimum and maximum values encountered so far and comparing them to the current node's value.

The given Python function maxAncestorDiff computes the maximum difference between the values of any two nodes with an

Space Complexity:

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The space complexity of the function is O(H), where H is the height of the binary tree. This complexity comes from the call stack used for recursion during DFS. In the worst case of a skewed tree, where the tree takes the form of a linked list (either every node has only a right child or only a left child), the height H would be equal to N, leading to a worst-case space complexity of O(N). For a balanced tree, the space complexity would be O(log N), as the height H would be logarithmic in the number of nodes N.