their appropriate positions based on the k value. Here's how we can think of the approach:

Binary Indexed Tree

# Problem Description

Greedy

represents the height of the person and k represents the number of people who are in front of this person in a queue and are taller or of the same height. The array people does not necessarily represent the order of people in a queue. Our task is to reconstruct the queue in such a way that the attributes of people in the reconstructed queue align with the original array people. The reconstructed queue should also be represented as an array where each element follows the same pair representation [h, k], with each pair denoting the attributes of the people in the queue.

Array

In this problem, we are given an array of people's attributes, where each person is represented by a pair of values [h, k]. The h

Sorting

exact count of people ahead of this person with heights greater than or equal to h\_i.

To be more concrete, if  $people[i] = [h_i, k_i]$ , we need to place the ith person in such a position in the queue that  $k_i$  is the

## The intuition behind the solution lies in the fact that it's easier to place the tallest people first and then insert the shorter people at

Intuition

1. Sort people by their height in descending order; if two people have the same height, sort them by their k value in ascending order. This way, we start with the tallest people, and for those with the same height, we place those with fewer people in front

first. 2. Initialize an empty list ans that will represent our queue.

3. Iterate over the sorted list, inserting each person into the ans list at the index equal to their k value. It works because by

- inserting the tallest people first, their k value corresponds exactly to the position they should be in the queue. As we insert the
- shorter people, we displace the taller ones only if needed, which maintains the integrity of the k values. By following this approach, we ensure that each person is placed in accordance with the number of taller or equally tall people that are supposed to be in front of them according to the original array people.

Solution Approach

### walkthrough of the implementation based on the solution code provided: 1. Sort the people array in a way such that we first sort the people by height in descending order and then, if the heights are equal,

1 ans = []

1 people.sort(key=lambda x: (-x[0], x[1]))

sort by the k value in ascending order. The lambda function in the sort method accomplishes this by using -x [0] for descending height and x[1] for ascending k values. Sorting is done with this line:

The solution uses a greedy algorithm approach along with list operations to reconstruct the queue. Here is a step-by-step

By sorting in this manner, we prioritize taller individuals and, among those of identical height, prioritize those with a smaller k value, which is the count of people in front of them.

```
3. Iterate over the sorted people array. For each person p, we insert them into the ans list at the index specified by their k value.
```

list, since those have been inserted earlier due to the sorting strategy.

2. Initialize an empty list ans which will be used to construct the queue.

1 for p in people: ans.insert(p[1], p) The insert operation places the person at the correct position in the queue, shifting elements that are already there.

This is possible because at the time of their insertion, there will be p[1] (which is k) taller or equal height people already in the

```
1 return ans
```

Using this greedy strategy, we are able to re-order the array so that each person is positioned in accordance with their height and the count of people in front of them with equal or greater height. The use of sorting and list insertions are key aspects of the

algorithm, ensuring that the resulting ans list satisfies the requirements that are stated in the problem.

Example Walkthrough

Let's consider a small example to illustrate the solution approach.

Suppose we are given the following array people:

4. Return the ans list which now represents the reconstructed queue.

#### Following the steps outlined in the solution approach:

sorting, our people array looks like this:

1 people = [[7, 0], [7, 1], [6, 1], [5, 0], [5, 2], [4, 4]]

1 people = [[7, 0], [4, 4], [7, 1], [5, 0], [6, 1], [5, 2]]

```
1 ans = []
```

3. We then iterate over the sorted people array and insert each person into the ans list at the index specified by their k value:

1. We first sort the array people by height in descending order and k value in ascending order when the heights are equal. After

```
    We insert [7, 1] at index 1: ans = [[7, 0], [7, 1]]
```

1 ans = [[5, 0], [7, 0], [5, 2], [6, 1], [4, 4], [7, 1]]

We insert [6, 1] at index 1: ans = [[7, 0], [6, 1], [7, 1]]

def reconstructQueue(self, people: List[List[int]]) -> List[List[int]]:

# Sort the 'people' array in descending order of height (h),

people.sort(key=lambda person: (-person[0], person[1]))

queue.insert(person[1], person)

We insert [7, 0] at index 0: ans = [[7, 0]]

```
    We insert [5, 0] at index 0: ans = [[5, 0], [7, 0], [6, 1], [7, 1]]

    We insert [5, 2] at index 2: ans = [[5, 0], [7, 0], [5, 2], [6, 1], [7, 1]]
```

2. We initialize an empty list ans:

 Finally, we insert [4, 4] at index 4: ans = [[5, 0], [7, 0], [5, 2], [6, 1], [4, 4], [7, 1]] 4. The final ans list is our reconstructed queue which aligns with the original people array attributes:

```
Using the greedy solution approach of sorting the individuals by height and then inserting them into the queue based on their k
value, we have successfully reconstructed a queue that satisfies the requirements of the problem. In this queue, for every person,
the k value is the exact count of people ahead of them with heights greater than or equal to their own.
Python Solution
   from typing import List
```

# and in the case of a tie, in ascending order of the number of people in front (k).

# which represents the number of people in front of them with equal or greater height.

# Initialize an empty list to hold the final queue reconstruction 10 queue = [] 11 # Iterate over the sorted 'people' list 12 for person in people: 13 # Insert each person into the queue. The index for insertion is the k-value, 14

```
# Return the reconstructed queue
18
           return queue
19
```

Java Solution

1 import java.util.Arrays;

15

16

17

29

9

10

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12

13

14

15

21

22

23

24

25

27

26 };

class Solution:

```
2 import java.util.ArrayList;
   import java.util.List;
   class Solution {
 6
       public int[][] reconstructQueue(int[][] people) {
           // Sort the array. First, sort by height in descending order.
           // If heights are equal, sort by the number of people in front (k-value) in ascending order.
           Arrays.sort(people, (person1, person2) ->
10
               person1[0] == person2[0] ? person1[1] - person2[1] : person2[0] - person1[0]);
11
12
13
           // Initialize a list to hold the final queue reconstruction,
14
           // which allows us to insert people at specific indices.
           List<int[]> reconstructedQueue = new ArrayList<>(people.length);
15
16
17
           // Iterate over the sorted array, and insert each person into the list
           // at the index specified by their k-value.
18
           for (int[] person : people) {
19
               // The second value of each person array (person[1])
20
21
               // specifies the index at which this person should be added in the queue
22
               reconstructedQueue.add(person[1], person);
23
24
25
           // Convert the list back to an array before returning it.
26
           return reconstructedQueue.toArray(new int[reconstructedQueue.size()][]);
27
28 }
```

std::vector<std::vector<int>> reconstructQueue(std::vector<std::vector<int>>& people) {

// Initialize an empty vector for the final arrangement.

std::vector<std::vector<int>> arrangedQueue;

// Return the final arrangement of the queue.

return arrangedQueue;

// Sort the people to arrange them according to their height in descending order.

// If two people are of the same height, order them by their "k-value" in ascending order.

return personA[0] > personB[0] || (personA[0] == personB[0] && personA[1] < personB[1]);</pre>

std::sort(people.begin(), people.end(), [](const std::vector<int>& personA, const std::vector<int>& personB) {

#### // Insert each person into the arrangedQueue. 16 // The index to insert the person is determined by the person's "k-value", 17 18 // which indicates the number of people in front of this person who have a height greater than or equal to this person's heig for (const std::vector<int>& person : people) { 19 arrangedQueue.insert(arrangedQueue.begin() + person[1], person); 20

C++ Solution

1 #include <vector>

2 #include <algorithm>

class Solution {

});

public:

```
Typescript Solution
   function reconstructQueue(people: number[][]): number[][] {
       // Sort the `people` array to arrange them by their height in descending order.
       // If two people have the same height, order them by their k-value in ascending order.
       people.sort((a, b) => {
           return b[0] - a[0] || a[1] - b[1];
 6
       });
       // Initialize an empty array for the final arrangement.
       let arrangedQueue: number[][] = [];
10
       // Insert each person into the `arrangedQueue`.
       // The index at which to insert is determined by the person's k-value,
       // which indicates the number of people in front of them with a height greater than or equal.
13
       for (let person of people) {
14
           arrangedQueue.splice(person[1], 0, person);
15
16
17
18
       // Return the final arrangement of the queue.
       return arrangedQueue;
19
20 }
21
22 // Example of usage:
  // let peopleInput: number[][] = [[7, 0], [4, 4], [7, 1], [5, 0], [6, 1], [5, 2]];
   // let queue = reconstructQueue(peopleInput);
   // console.log(queue); // Outputs arranged queue based on the given conditions.
26
```

# Time Complexity

1. Sorting the people array:

Time and Space Complexity

### This operation uses a sorting algorithm, typically Timsort in Python, which has a time complexity of O(n log n) for an array of n elements.

The insert operation within a list has a time complexity of O(n) in the worst case because it may require shifting the elements of the list.

The time complexity of the code can be broken down into two major operations:

case complexity of  $O(n^2)$  as it performs n insertions into a structure that can be up to n in size.

2. Reconstructing the queue by inserting elements:

By combining these, the overall time complexity of the code is  $0(n \log n) + 0(n^2)$ , which simplifies to  $0(n^2)$  since  $0(n^2)$  is the dominating term when n is large.

The space complexity of the code is the amount of extra space used by the algorithm, not counting the space occupied by the input

Since insert is called for each person in the sorted array and the list can grow up to n elements, this operation will have a worst-

1. Sorted people array:

Space Complexity

itself. In this case:

Since the sorting operation is done in-place, it does not contribute to additional space complexity beyond what is taken by the

input people array. 2. The ans list:

This list is created to store the output and grows to contain n elements, which are the same elements from the input people

array. Thus, the space required is O(n). Therefore, the overall space complexity of the code is O(n).