2046. Sort Linked List Already Sorted Using Absolute Values

Problem Description

Linked List

Two Pointers

Sorting

The given problem asks us to take the head of a singly-linked list, where the list is sorted in non-decreasing order based on the absolute values of its nodes. The goal is to re-sort this list, but this time according to the actual values of its nodes, still maintaining a non-decreasing order. This challenge essentially involves sorting negative numbers (where applicable) before the non-negative ones while adhering to the non-decreasing sequence.

Intuition

The intuition behind the solution is to leverage the fact that the list is already sorted by absolute values. This characteristic

Medium

implies that all negative values, if any, will be at the beginning of the list - but they will be in decreasing order because their absolute values are sorted in non-decreasing order. Conversely, the non-negative values will follow in non-decreasing order. Therefore, the solution approach involves iterating through the list and moving any encountered negative values to the front.

Since we know that the list is effectively partitioned into a non-increasing sequence of negative values followed by a non-

decreasing sequence of non-negative values, moving each negative value we encounter to the front of the list will naturally sort it. By iterating just once through the linked list, whenever a negative value is found, it's plucked out of its position and inserted

before the current head. This insert-at-head operation ensures that the ultimately larger (less negative) values encountered later are placed closer to the head, thereby achieving a non-decreasing order of actual values. The key points of the approach are:

 Iterating just once throughout the list is sufficient. • A simple insert-at-head operation for negative values during iteration results in the desired sorted order.

The solution preserves the initial list's order in terms of absolute values while reordering only where necessary to get the actual

• No need for a complex sorting algorithm since the list is already sorted by absolute values.

values sorted, hence achieving the re-sorting in an efficient manner.

Negative values must precede non-negative ones for actual value sorting.

additional data structures. The algorithm follows these steps: 1. Initialize two pointers prev and curr. prev starts at the head of the list, and curr starts at the second element (head.next), since there is no need

The provided solution uses a simple while loop and direct manipulation of the linked list nodes to achieve the task without

2. Iterate through the linked list with a while loop, which continues until curr is None, meaning we've reached the end of the list. 3. For each node, we check if curr. val is negative.

As for complexity:

Example Walkthrough

 $2 \rightarrow -3 \rightarrow -4 \rightarrow 5 \rightarrow 6$

to move the head itself.

Solution Approach

o If it is negative, we conduct the following re-linking process: a. Detach the curr node by setting previnent to currinext. This removes the current node from its position in the list. b. Move the curr node to the front by setting curr next to head. This places the current node at the

list by setting it to t (the node following the moved node). o If currival is not negative, we simply move forward in the list by setting prev to curr and curr to currinext, as no changes are needed for

non-negative or zero values. 4. Repeat this process until the curr pointer reaches the end of the list.

solution very efficient. There are no complicated data structures or algorithms needed; the solution makes direct alterations to

beginning of the list. c. Update head to be curr since curr is now the new head of the list. d. Move the curr pointer to the next node in the

The algorithm uses the two-pointer technique to traverse and manipulate the linked list in place without requiring extra space for sorting. The key insight here is that the insertion of negative nodes at the head can be done in constant time, which makes the

the list's nodes, which takes advantage of the properties of linked lists for efficient insertion operations.

• The time complexity is O(n), where n is the number of nodes in the list, since every node is visited once. • The space complexity is O(1), as no additional space proportional to the input size is used.

This approach is efficient and takes advantage of the already partially sorted nature of the list based on absolute values to

According to the problem statement, our goal is to reorder this list based on the actual values, preserving non-decreasing order. Here's a step-by-step walkthrough using the solution approach provided:

3. Since the value of curr (-3) is negative, we insert it before the head: a. Detach curr: we remove -3 from the list. The list becomes: 2 -> -4 -> 5

-> 6 b. Move curr to the front: we insert -3 before the current head. The list becomes: -3 -> 2 -> -4 -> 5 -> 6 c. Update head to curr: -3 is

7. We've reached the end of the list. The list is now correctly sorted by actual values in non-decreasing order: -4 -> -3 -> 2 -> 5 -> 6

4. Repeat the process for the next curr, which again has a negative value (-4): a. Detach curr: we remove -4 from the list. The list becomes: -3 -> 2 -> 5 -> 6 b. Move curr to the front: we insert -4 before the head. The list becomes: -4 -> -3 -> 2 -> 5 -> 6 c. Update head to curr: -4 is

Solution Implementation

Definition for a singly-linked list node.

if current.val < 0:</pre>

if previous is None:

continue

current = current.next

Move on to the next node

current = previous.next

// Continue traversal from the next node

// If the current value is not negative, move both pointers forward

// Return the updated list with all negative values at the front

current = temp;

previous = current;

current = current.next;

* Definition for singly-linked list.

ListNode() : val(0), next(nullptr) {}

ListNode* sortLinkedList(ListNode* head) {

ListNode* current = head->next;

// Iterate through the list

} else {

while (current != nullptr) {

if (current->val < 0) {</pre>

head = current;

previous = current;

current = current->next;

return head; // Return the sorted list

current = temp;

current->next = head;

// Initialize pointers for traversal

ListNode(int x) : val(x), next(nullptr) {}

ListNode(int x, ListNode *next) : val(x), next(next) {}

ListNode* previous = head; // Pointer to track the node just before the current one

// When the current node has a negative value, reposition it to the start of the list

// If the current node is non-negative, move both pointers forward

ListNode* temp = current->next; // Hold the next node to continue traversal later

previous->next = temp; // Remove the current node from its original position

// Pointer to track the current node during traversal

// Update the head to the new first node

// Continue traversal from the next node

// Insert the current node at the beginning of the list

} else {

return head;

* struct ListNode {

class Solution {

int val;

ListNode *next;

C++

/**

*

*/

* };

public:

previous = current

def __init__(self, val=0, next=None):

self.val = val # Value of the node

Python

class ListNode:

5. Now curr points to a positive value (5), so we just move the prev to curr and curr to its next node, which is 6. 6. As curr points to another positive value (6), we again just advance prev and curr accordingly.

the new head of the list. d. Move curr to next node: curr now points to 5.

the new head of the list. d. Move curr to next node: curr now points to -4.

achieve the fully sorted list by actual values with minimal operations.

Let's consider the following singly-linked list sorted by absolute values:

```
While iterating through the list, we only moved the negative values, which were already sorted in non-increasing order, to the
front. We didn't need to touch the positive values because they were already in non-decreasing order. This process respects the
```

previously sorted absolute values while ordering the elements by their actual values.

self.next = next # Reference to the next node in the list

Check if the current node's value is negative.

Traverse the linked list starting from the head until there are no more nodes.

If this is the first node (head of the list), no need to move it.

1. We initialize prev as the head (node with value 2) and curr as the second node (node with value -3).

2. We start the while loop and iterate over the list. In our example, the first curr is negative.

class Solution: def sortLinkedList(self, head: Optional[ListNode]) -> Optional[ListNode]: # Initialize pointers. # `previous` will track the node before `current`. previous = None current = head

If current node is not negative, just move `previous` and `current` pointers one step forward.

If current node is negative, shift it to the beginning of the list. previous.next = current.next # Link the previous node to the next node, effectively removing `current` from its current.next = head # Make `current` node point to the current head head = current # Update `head` to be the `current` node

else:

while current:

```
current = current.next
       return head # Return the modified list with all negative values moved to the beginning
Java
/**
* Definition for singly-linked list.
class ListNode {
 int val;
 ListNode next;
 ListNode() {}
 ListNode(int val) { this.val = val; }
 ListNode(int val, ListNode next) { this.val = val; this.next = next; }
class Solution {
 /**
  * This method sorts a linked list such that all nodes with negative values
  * are moved to the front.
  * @param head The head of the singly-linked list.
  * @return The head of the sorted singly-linked list.
 public ListNode sortLinkedList(ListNode head) {
   // Previous and current pointers for traversal
   ListNode previous = head;
   ListNode current = head.next;
   // Iterate through the list
   while (current != null) {
     // When a negative value is encountered
     if (current.val < 0) {</pre>
       // Detach the current node with the negative value
       ListNode temp = current.next;
       previous.next = temp;
       // Move the current node to the front of the list
       current.next = head;
       head = current;
```

```
};
```

TypeScript

```
/**
   * Definition for singly-linked list.
  interface ListNode {
    val: number;
    next: ListNode | null;
  /**
   * Sorts a singly linked list such that all negative values are moved to the front.
   * @param {ListNode | null} head - The head of the linked list to be sorted.
   * @returns {ListNode | null} The head of the sorted linked list.
   */
  function sortLinkedList(head: ListNode | null): ListNode | null {
    // There is nothing to sort if the list is empty or has only one node.
    if (!head || !head.next) {
      return head;
    // Initialize pointers for traversal.
    let previous: ListNode | null = head; // Pointer to track the node just before the current one.
    let current: ListNode | null = head.next; // Pointer to track the current node during traversal.
    // Iterate through the list.
    while (current !== null) {
      // When the current node has a negative value, reposition it to the start of the list.
      if (current.val < 0) {</pre>
        let temp: ListNode | null = current.next; // Hold the next node to continue traversal later.
        previous.next = temp; // Remove the current node from its original position.
        current.next = head; // Insert the current node at the beginning of the list.
        head = current; // Update the head to the new first node.
        current = temp; // Continue traversal from the next node.
      } else {
        // If the current node is non-negative, move both pointers forward.
        previous = current;
        current = current.next;
    // Return the head of the sorted list.
    return head;
# Definition for a singly-linked list node.
class ListNode:
   def __init__(self, val=0, next=None):
        self.val = val # Value of the node
        self.next = next # Reference to the next node in the list
class Solution:
   def sortLinkedList(self, head: Optional[ListNode]) -> Optional[ListNode]:
        # Initialize pointers.
        # `previous` will track the node before `current`.
        previous = None
        current = head
       # Traverse the linked list starting from the head until there are no more nodes.
        while current:
            # Check if the current node's value is negative.
            if current.val < 0:</pre>
                # If this is the first node (head of the list), no need to move it.
                if previous is None:
                    current = current.next
                    continue
```

previous.next = current.next # Link the previous node to the next node, effectively removing `current` from its curre

Time and Space Complexity

Time Complexity

Space Complexity

else:

adjustments which take constant time. Hence, no matter the input size of the singly-linked list, the algorithm guarantees that each element is checked exactly once. The time complexity of the code is O(n), where n is the number of nodes in the linked list.

If current node is negative, shift it to the beginning of the list.

current.next = head # Make `current` node point to the current head

return head # Return the modified list with all negative values moved to the beginning

head = current # Update `head` to be the `current` node

Move on to the next node

current = previous.next

current = current.next

previous = current

The space complexity of the code is related to the extra space used by the algorithm, excluding the space for the input. The provided code uses a constant amount of space: variables prev, curr, and t. No additional data structures are used that

If current node is not negative, just move `previous` and `current` pointers one step forward.

The given code is intended to sort a singly-linked list by moving all negative value nodes to the beginning of the list in one pass.

It iterates through the list only once, with a fixed set of actions for each node. The main operations within the loop involve pointer

Thus, the space complexity of the code is 0(1) since it uses constant extra space.

grow with the input size. All changes are made in-place by adjusting the existing nodes' next pointers.