

Problem Description

defined as a sequence of the same character that appears at least three times consecutively. The goal is to find the starting index and the ending index for each of these large groups.

The given problem revolves around identifying the 'large' groups of consecutive identical characters in a string s. A 'large' group is

Consider this example: In the string s = "abbxxxxzyy", the following groups are formed: "a", "bb", "xxxx", "z", and "yy". Out of these, only "xxxx" qualifies as a large group because it contains the same character ('x') four times in a row. The interval for this

group is [3, 6], where 3 is the starting index, and 6 is the ending index of the group within the string.

The expected output is a list of intervals with each interval representing a large group and the list is sorted by the starting index of each interval.

Intuition

we encounter a different character or reach the end of the string, we check if the current character sequence qualifies as a large group by ensuring the sequence length is at least three characters long. If it does, we record the start and the end indices of this group. Next, we update our tracker to the current character's index and continue scanning until we have examined all characters in the string. Iterating only once through the string results in an efficient solution with a linear time complexity (O(n)), where n is the length of the

To solve this problem, the approach is to iterate over the string while keeping track of the start of a potential large group. Whenever

suitable for large input strings as well. The key takeaway is that we maintain the current character sequence length and start index while traversing the string and only record intervals that satisfy the 'large group' criteria.

string. We can assert that the solution is efficient because we're only scanning the string once without any nested loops, making it

Solution Approach

The solution's implementation utilizes a two-pointer pattern, which is often used in problems involving arrays or strings where we need to track subarrays or substrings based on a certain condition. Here, the two pointers are denoted as i (the start pointer) and j

(the end pointer). The approach begins with initializing the starting pointer i to 0 and an empty list ans to collect the intervals representing the large groups.

The main algorithm can be described as follows: 1. Begin iterating over the string with the starting pointer i, aiming to pinpoint the start of a character sequence.

2. Create an inner loop that increments the end pointer j for as long as the character at position j is the same as the character at

- the start pointer i. 3. Upon exit from the inner loop, check if the difference between j and i is greater than or equal to 3, indicating a large group.
- 4. If it is indeed a large group, append the interval [i, j-1] to the list ans, since j-1 is the end index of the large group. 5. Finally, update the start pointer i to j to commence scanning for the next potential group.
- The code maintains a while loop that will run until i reaches the end of the string (i < n). Each time a large group is found, the
- interval is appended to the list before resetting i to continue searching from the end of the last found group. The use of the continuous while loop paired with the two-pointer technique ensures all large groups are identified without

Once the while loop is completed, the function returns ans, which now contains all intervals of the identified large groups, arranged by their start indices as per the requirement.

1 while i < n: # Main loop iterating through the string j = i # Initializing end pointer while j < n and s[j] == s[i]: # Inner loop to increment 'j' as long as characters match

Move start pointer 'i' to the end of the last found group i = j

defined as a sequence with at least three occurrences of the same character.

Here is an excerpt showcasing the two-pointer technique in the given code:

repetitively checking characters, which efficaciously reduces the time complexity.

```
This solution is efficient as it involves a single pass of the input string, making the time complexity O(n), where n is the size of the
string.
Example Walkthrough
```

if j - i >= 3: # If the length of the group is at least 3, it is a large group

ans.append([i, j-1]) # Append interval to answer list as a large group is found

Let's walk through an example to illustrate the solution approach:

Consider the string s = "aabbbccdeee". We need to find the large groups of consecutive identical characters, where a large group is

1. Initialize the start pointer i to 0 and create an empty list ans for storing our intervals.

j += 1

3. The end pointer j is also initialized to the same value as i.

5. No interval is added to ans after the first loop since j - i < 3.

2. Initiate the main while loop, where i starts at index 0.

4. Start the inner while loop that compares characters at positions j and i and increments j as long as they are identical. For the first loop, we compare characters 'a' at index 0 and 1, and since they are not identical, j stops at 1.

6. Now, i is set to j, and our pointers are i = 1, j = 1. The comparison starts with 'a' and 'b'. Since they are different characters, we immediately look for the next sequence.

index because the string is traversed from left to right, naturally maintaining the order.

Initialize start index `start` and length of the string `length`

def largeGroupPositions(self, s: str) -> List[List[int]]:

while end < length and s[end] == s[start]:</pre>

Return the final list of positions of large groups

if (endIndex - startIndex >= 3) {

largeGroups.add(Arrays.asList(startIndex, endIndex - 1));

return largeGroups; // Return the list of all large groups found.

Initialize an empty list to store the answer

9. Update i to the current value of j, which is 4.

start, length = 0, len(s)

end = start

return answer

end += 1

answer = []

- 7. Next iteration, i is still at 1, and 'b' at position 1 matches 'b' at positions 2 and 3. The inner loop increments j to 4. 8. Now j - i >= 3 - it's a large group, so we add the interval [1, 3] to ans.
- 10. Continue this process for the rest of the string. \circ i = 4, j = 4: Match 'c' with subsequent 'c' until j = 6. j - i = 2, not a large group.
 - o i = 6, j = 6: Skip 'd' as it's a single character. \circ i = 7, j = 7: Characters 'e' at index 7, 8, and 9 match. j moves to 10. j - i >= 3, we add the interval [7, 9] to ans.
- 11. Finally, our list ans contains [[1, 3], [7, 9]], representing the large groups 'bbb' and 'eee' with their respective starting and ending indices.

The implementation efficiently identifies the large groups in a single traversal of the string s. The output list is sorted by starting

Python Solution class Solution:

Iterate over the string characters by index while start < length:</pre> 9 # Initialize the end index `end` at the same position as start 10

Move `end` forward as long as the subsequent characters are the same as s[start]

```
# If the length of the group is 3 or more, add it to the answer list
15
               if end - start >= 3:
16
                    answer.append([start, end - 1])
17
               # Update start to the next character group
18
19
               start = end
```

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Java Solution
1 import java.util.ArrayList;
2 import java.util.Arrays;
   import java.util.List;
   class Solution {
       // Method to find the starting and ending indices of all large groups.
       public List<List<Integer>> largeGroupPositions(String s) {
            int strLength = s.length(); // Length of the input string.
9
           int startIndex = 0; // Initialize the starting index of a group.
10
           List<List<Integer>> largeGroups = new ArrayList<>(); // Initialize the list to store the result.
11
12
13
           // Iterate over the string to find large groups.
           while (startIndex < strLength) {</pre>
14
15
                int endIndex = startIndex; // Initialize the end index of the current group.
               // Increase the endIndex as long as the current character is the same as the start character.
16
               while (endIndex < strLength && s.charAt(endIndex) == s.charAt(startIndex)) {</pre>
17
18
                   ++endIndex;
19
20
               // Check if the current group is a large group (i.e., has a length of 3 or more).
```

// Add the start and end indices (end index is exclusive so subtract 1) of the large group to the result.

startIndex = endIndex; // Move the startIndex to the end of the current group to start checking the next group.

```
C++ Solution
 1 #include <vector>
 2 #include <string>
   class Solution {
   public:
       std::vector<std::vector<int>> largeGroupPositions(std::string s) {
            int length = s.size();  // Get the size of the input string
                                          // Initialize the starting index of a group
           int startIndex = 0;
           std::vector<std::vector<int>> largeGroups; // Result vector for storing large group positions
10
           // Iterate over the string to identify groups
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           while (startIndex < length) {</pre>
13
               int endIndex = startIndex; // Initialize the end index of the group to the current start index
14
               // Advance the end index as long as characters match the character at the start index
15
               while (endIndex < length && s[endIndex] == s[startIndex]) {</pre>
16
                   ++endIndex;
17
18
19
               // Check if the identified group is a large group (3 or more characters)
20
21
               if (endIndex - startIndex >= 3) {
                   // Add the start and end indexes (exclusive of the end) of the large group to the result
                   largeGroups.push_back({startIndex, endIndex - 1});
23
24
25
26
               startIndex = endIndex; // Move the start index to the end of the current group for the next iteration
27
28
29
            return largeGroups; // Return the vector of large groups
30
31 };
32
```

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Typescript Solution

function largeGroupPositions(s: string): number[][] {

let length: number = s.length; // Get the length of the input string

let startIndex: number = 0; // Initialize the starting index of a group

loop skipping groups of the same character, every character is visited at most once.

The space complexity is determined by the space needed to store the output, which is the list ans.

let largeGroups: number[][] = []; // Result array for storing large group positions

```
// Iterate over the string to identify groups
       while (startIndex < length) {</pre>
           let endIndex: number = startIndex; // Initialize the end index of the group to the current start index
           // Advance the end index as long as characters match the character at the start index
           while (endIndex < length && s.charAt(endIndex) === s.charAt(startIndex)) {</pre>
               endIndex++;
13
14
           // Check if the identified group is a large group (3 or more characters)
15
           if (endIndex - startIndex >= 3) 
16
               // Add the start and end indices (end is exclusive) of the large group to the result
17
               largeGroups.push([startIndex, endIndex - 1]);
19
20
           startIndex = endIndex; // Move the start index to the end of the current group for the next iteration
21
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23
       return largeGroups; // Return the array of large groups
24
25 }
26
Time and Space Complexity
Time Complexity
The given Python function scans through the string s once. The inner while-loop advances the index j as long as consecutive
```

characters are equal to s[i], and the outer while-loop is responsible for iterating over each character, but thanks to the inner while-

each character. Therefore, the overall time complexity is O(n).

Space Complexity

In the worst case, if we have a string where every three characters form a group (for example, "aaabbbccc"), the number of groups (and thus the number of sublists in ans) will be approximately n/3. Each group is represented by a pair of indices (which takes

The time complexity for scanning the string is O(n), where n is the length of the string, since in the worst case, we have to look at

constant space), so the space required grows linearly with the number of groups.

Thus, the space complexity is O(n) where n is the length of the string, as the size of the ans list is proportional to the number of large groups in the string.