



Problem Description

with the dimensions (n - 2) x (n - 2). For each cell in maxLocal, its value should be the largest number found in the corresponding 3 x 3 submatrix of grid. To clarify, the submatrix corresponds to a sliding window centered around the cell (i + 1, j + 1) in the original grid, with i and j denoting row and column indices, respectively, in maxLocal. Essentially, it can be visualized as overlaying a 3 x 3 grid on top of the main grid, moving it one cell at a time, and at each position, recording the maximum value found in this overlay into the new matrix.

The given problem involves an n x n integer matrix named grid. The objective is to generate a new integer matrix called maxLocal

Intuition

3 x 3 submatrix. We start this process one row and one column in from the top-left of the grid (since the edge rows and columns don't have enough neighbors for a full 3 x 3 submatrix) and end one row and one column before the bottom-right. At each iteration of the loop, we find the largest value from the current 3 x 3 submatrix and store it in the corresponding cell in the maxLocal matrix. More specifically, the iteration will have two loops: one for traversing rows i and another for columns j. For each position (i, j), we

To solve this challenge, the intuitive approach is to perform a nested loop traversal across the main grid that captures every possible

look at rows i to i + 2 and columns j to j + 2, creating a 3 x 3 region. Making use of a list comprehension and the max() function, we find the largest value in that region and save it as the value of maxLocal[i][j]. This two-step process - finding the 3 x 3 submatrix and then the maximum value within it - is repeated until the entire maxLocal matrix is filled.

The implementation utilizes a straightforward brute-force algorithm to address the problem. This strategy leverages the capability of

Solution Approach

Here's a step-by-step walkthrough of the algorithm, using Python as the reference language:

2. Initialize the maxLocal matrix filled with zeros, with the size $(n - 2) \times (n - 2)$, which will eventually store the largest values.

nested loops and list comprehension for easy iteration through the matrix.

1. Determine the size n of the input grid.

- 3. Start with two nested loops, where i iterates through rows from 0 to n 2, and j iterates through columns from 0 to n 2. These ranges are chosen to ensure we can always extract a 3×3 submatrix centered around grid[i + 1][j + 1].
- 4. For each position (i, j), extract the contiguous 3 x 3 submatrix. This is done by another nested loop, or in this case, a list comprehension, that iterates through all possible x and y coordinates in this submatrix, with x ranging from i to i + 2 and y ranging from j to j + 2.
- 5. Calculate the maximum value within this 3×3 submatrix using the max() function applied to the list comprehension, which iterates over the range of x and y and accesses the values in grid[x][y]. 6. Assign this maximum value to the corresponding cell in maxLocal[i][j].
- By using a list comprehension within the nested loops to calculate the maximum, the implementation avoids the need for an explicit
- inner loop for exploring the 3×3 submatrix. This makes the code more concise and readable.

Additionally, the approach does not require any extra data structures aside from the maxLocal matrix which is being filled in, indicating an in-place algorithm with no additional space complexity.

The final result will be the maxLocal matrix, which now contains the largest number from every contiguous 3 x 3 submatrix of the

1 class Solution: def largestLocal(self, grid: List[List[int]]) -> List[List[int]]: n = len(grid) ans = $[[0] * (n - 2) for _ in range(n - 2)]$

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for j in range(n - 2):
                   ans[i][j] = max(
                       grid[x][y] for x in range(i, i + 3) for y in range(j, j + 3)
10
           return ans
The above code snippet corresponds to the entire algorithm discussed, simplified into a Python method definition within a Solution
class.
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Example Walkthrough Let's assume we have the following 4×4 grid:

7 8 6 2 4 4 3 2 0

We want to apply the algorithm to create a maxLocal matrix with dimensions $(4 - 2) \times (4 - 2)$, which is 2×2 .

for i in range(n - 2):

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Here are the steps of the algorithm in action:
 1. Start with i=0 and j=0. The 3 x 3 submatrix is:
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3 7 8 6

1 9 9 8 1

original grid.

1 9 9 8 2 5 7 5

1 9 8 1

2. Move to i=0 and j=1. The 3 x 3 submatrix is:

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The maximum value in this submatrix is 9. So, maxLocal[0][1] will be 9.
3. Next, i=1 and j=0. The 3 x 3 submatrix is:
```

2 7 8 6

3 4 3 2

3 8 6 2

The maximum value in this submatrix is 8. So, maxLocal[1][0] will be 8.

The maximum value in this submatrix is 9. So, maxLocal[0][0] will be 9.

1 7 5 1 2 8 6 2 3 3 2 0

After filling all the cells in maxLocal, we get the final matrix:

for row in range(grid_size - 2):

for col in range(grid_size - 2):

4. Finally, i=1 and j=1. The 3 x 3 submatrix is:

```
The maximum value in this submatrix is 8. So, maxLocal[1][1] will be 8.
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2 8 8

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Python Solution

This 2x2 matrix represents the maximum values from each 3×3 sliding window within the original 4×4 grid.

3x3 local grids and each local grid reduces the dimension by 2 on each axis.

largest_values_grid = [[0] * (grid_size - 2) for _ in range(grid_size - 2)]

Iterating over each cell that will be the top-left corner of a 3x3 grid.

Finding the largest value in the current 3x3 local grid.

// Initialize max value for the current 3x3 subgrid

// Nested loops to go through each element in the 3x3 subgrid

// Update maxVal if a larger value is found

// Assign the largest value in the 3x3 subgrid to the answer grid

maxVal = Math.max(maxVal, grid[x][y]);

// Return the answer grid containing the largest values of each subgrid

int maxVal = Integer.MIN_VALUE;

maxLocalValues[i][j] = maxVal;

return maxLocalValues;

for (int x = i; x <= i + 2; ++x) {

for (int y = j; $y \le j + 2$; ++y) {

vector<vector<int>> result(size - 2, vector<int>(size - 2));

// Loop through each cell where a 3x3 grid can start.

for (let z = col; z < col + 3; z++) {

// Find the maximum value in the local 3x3 grid

// Assign the maximum value found to the corresponding position in resultGrid

localMax = Math.max(localMax, grid[k][z]);

for (int col = 0; col < size - 2; ++col) {

for (int row = 0; row < size - 2; ++row) {</pre>

def largestLocal(self, grid: List[List[int]]) -> List[List[int]]: # Determining the size of the given grid. grid_size = len(grid) # Preparing the answer grid with the reduced size (n-2) since we are looking for

class Solution:

from typing import List

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largest_value = max(
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18
                       grid[i][j] for i in range(row, row + 3) for j in range(col, col + 3)
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21
                   # Storing the largest value found in the corresponding cell of the answer grid.
22
                   largest_values_grid[row][col] = largest_value
23
24
           # Returning the answer grid containing all the largest values found in each 3x3 local grid.
25
           return largest_values_grid
26
Java Solution
   class Solution {
       // Method to find the largest element in every 3x3 subgrid
       public int[][] largestLocal(int[][] grid) {
           // Determine the size of the grid
           int gridSize = grid.length;
           // Initialize the answer grid with a reduced size
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           // because the border elements can't form a complete 3x3 subgrid
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           int[][] maxLocalValues = new int[gridSize - 2][gridSize - 2];
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           // Iterate through the grid, considering each 3x3 subgrid
           for (int i = 0; i <= gridSize - 3; ++i) {
13
               for (int j = 0; j <= gridSize - 3; ++j) {
14
```

class Solution { 5 public: vector<vector<int>> largestLocal(vector<vector<int>>& grid) { int size = grid.size(); // Get the size of the input grid.

C++ Solution

1 #include <vector>

2 using namespace std;

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                   // Iterate through each cell within the current 3x3 window.
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                    for (int localRow = row; localRow <= row + 2; ++localRow) {</pre>
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                        for (int localCol = col; localCol <= col + 2; ++localCol) {</pre>
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                            // Update the corresponding cell in the result grid with
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                            // the maximum value seen so far in the 3x3 window.
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                            result[row][col] = max(result[row][col], grid[localRow][localCol]);
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           // Return the populated result grid after processing the entire input grid.
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           return result;
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26 };
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Typescript Solution
   function largestLocal(grid: number[][]): number[][] {
       const gridSize = grid.length; // The size of the input grid
       // Initialize the result grid with dimensions (gridSize - 2) x (gridSize - 2)
       const resultGrid = Array.from({ length: gridSize - 2 }, () => new Array(gridSize - 2).fill(0));
       // Iterate through the grid to fill the resultGrid with the maximum values from each local 3x3 grid
       for (let row = 0; row < gridSize - 2; row++) {
            for (let col = 0; col < gridSize - 2; col++) {
                let localMax = 0; // The maximum value in the current 3x3 grid
10
               // Iterate over the 3x3 grid starting at (row, col)
11
               for (let k = row; k < row + 3; k++) {
12
```

// Initialize the result grid with dimensions (size -2) x (size -2) to accommodate the 3x3 window reduction.

20 resultGrid[row][col] = localMax; 21 22 23 24 // Return the result grid which contains the largest values found in each local 3x3 grid

return resultGrid;

Time and Space Complexity The provided code snippet is used to find the maximum local element in all 3×3 subgrids of a given 2D grid. Here is an analysis of its

Time Complexity: The time complexity of the code is determined by the nested loops and the operations within them. The outer two for loops iterate

time and space complexity:

over (n - 2) * (n - 2) elements since we're examining 3×3 subgrids and thus can't include the last two columns and rows for starting points of our subgrids. For each element in the answer grid, we find the maximum value within the 3×3 subgrid, which involves iterating over 9 elements (3 rows by 3 columns).

Therefore, the time complexity is: 0((n-2)*(n-2)*9), which simplifies to $0(n^2)$, assuming that the max operation within a constantsized subgrid takes constant time.

Space Complexity:

The space complexity is determined by the additional space used by the algorithm, not including the input. In this case, we're creating an output grid ans of size (n - 2) * (n - 2) to store the maxima of each 3×3 subgrid, which represents the additional

space used. Thus, the space complexity is: 0((n-2)*(n-2)) which simplifies to $0(n^2)$ as the size of the ans grid grows quadratically with the input size n.