

Problem Description

element repeats itself, and all these elements are present within arr1. Our goal is to sort arr1 so that the order of elements matches how they appear in arr2, while any elements that are not found in arr2 should be appended to the end of arr1 in ascending order. For example, if we have arr1 = [2,3,1,3,2,4,6,7,9,2,19] and arr2 = [2,1,4,3,9,6], the first step is to arrange the elements of

In this problem, we are provided with two integer arrays, arr1 and arr2. The elements within arr2 are distinct, which means no

arr1 that are also in arr2 in the order they appear in arr2: [2,2,2,1,4,3,3,9,6]. After that, we take the remaining elements that are not in arr2 (which are 7, 19, and 7 in this case) and place them at the end in sorted order: [7,19,7]. The final output array would be the concatenation of these two processes, giving us [2,2,2,1,4,3,3,9,6,7,7,19]. The challenge lies in figuring out a strategy that allows us to sort arr1 with these constraints in mind.

Intuition

The key to solving this problem efficiently is to recognize that we can use a custom sorting strategy. The default sort mechanisms

found in arr2.

Given that the elements in arr2 are the first to be sorted and they are in a specific order, we need to have this order play a significant role in our custom sorting strategy. For elements that are in arr2, we want them to appear first and in the order they are

most programming languages provide can often be customized with a user-defined key function that determines the sort order.

To achieve this, we can map each element of arr2 to its position (or index) within arr2. This can be done using a dictionary or a hash map that takes the element as a key and the position of that element in arr2 as the value.

In Python, this mapping can be created as follows: $pos = \{x: i \text{ for } i, x \text{ in enumerate(arr2)}\}$ For the elements of arr1 that are not found in arr2, we need to ensure they come after any element that is in arr2. One way to do

this is to assign them an index that is larger than any index that would be assigned to the elements that are in arr2. Since indices in

an index based on a large number (for example, 1000) plus their value, which ensures they will be sorted in ascending order after the

Python are zero-based and arr2 can have, at most, the same number of elements as arr1, all the elements not in arr2 can be given

elements of arr2. Now, using Python's sorted function, we can sort arr1 with a custom key. This key is defined by a lambda function that takes an element x and returns its corresponding index from the dictionary pos if it's in arr2, or 1000 + x otherwise. Here's the implementation in Python, which reflects our approach:

1 sorted(arr1, key=lambda x: pos.get(x, 1000 + x)) By using this custom sorting key, we align arr1 elements per arr2's order and then append the remaining elements in ascending

order right after. The overall runtime complexity of the solution is O(NlogN), where N is the length of arr1.

Solution Approach

The solution is implemented using a dictionary for direct mapping and the sorting algorithm that Python provides. Here's a step-bystep explanation, referencing the provided code:

1. Create a Mapping of arr2: The solution starts by creating a dictionary called pos. This dictionary maps each element x of arr2 to

its index i in the array. This is done using a dictionary comprehension: 1 pos = {x: i for i, x in enumerate(arr2)}

during sorting:

By doing this, we are setting up a quick reference so that we can look up the index of any element from arr2 instantly. This index acts as a custom priority value for the sorting algorithm.

1 sorted(arr1, key=lambda x: pos.get(x, 1000 + x))

The key parameter is crucial here. It takes a lambda function which determines how the elements of arr1 should be compared

2. Custom Sort with a Lambda Function: We need to sort arr1, but not by its elements' natural order. We will use Python's sorted

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    pos.get(x, 1000 + x) is the function's body that returns a value for each x that will be used to compare during sorting.

\circ If x is found in arr2 (and, as a result, the post dictionary), postget(x) returns the index of x from arr2.
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element x itself. This ensures that:

function, which allows us to specify a custom key function:

lambda x: is an anonymous function that takes x, an element from arr1.

■ The elements not in arr2 are ordered at the end because their sorting key is greater than any index assigned from arr2. ■ They are sorted in ascending order amongst themselves because if two elements are not found in arr2, their sorting key

By using this sorting strategy, we can maintain the relative ordering based on another array arr2 while also logically appending and

Thus, using a combination of dictionary for direct access and sorting algorithm with a custom key function, this solution efficiently

 \circ If x is not found in arr2, the method pos.get(x, 1000 + x) provides a default value which is the sum of 1000 and the

is simply their value (1000 + their own value respects the natural ascending order).

achieves the required array manipulation adhering to the conditions of the problem statement.

- sorting elements not found in arr2.
- Let's illustrate the solution approach using a smaller example. Suppose arr1 contains [8, 4, 5, 4, 6] and arr2 contains [4, 6]. We

end. Step 1: Create a Mapping for arr2 Elements:

want to sort arr1 such that the order of arr2 is preserved for common elements, and the rest are sorted in ascending order at the

2 # pos will be {4: 0, 6: 1}

Example Walkthrough

Here, 4 is at index 0 in arr2 and 6 is at index 1.

1 sorted(arr1, key=lambda x: pos.get(x, 1000 + x))

Step 2: Sort arr1 with a Custom Sort Key:

pos = {x: i for i, x in enumerate(arr2)}

Breaking down how the key function works for each element in arr1: 8 is not found in arr2, so the key would be 1000 + 8 = 1008.

The first 4 is found in arr2 and its index is 0; hence, the key would be 0.

We now sort arr1 using Python's sorted function with the custom key:

First, we create a dictionary that maps each element of arr2 to its index:

• 6 is found and its index is 1, so the key is 1. The sorted order using these keys would be [4, 4, 6, 5, 8].

• 5 is not found in arr2, so its key is 1000 + 5 = 1005.

The second 4 is found and its key is again 0.

Step 3: Review the Final Sorted Array: So, the elements from arr1 that were also in arr2 are now sorted in arr2 order: 4 comes first, followed by 6. The remaining elements,

not in arr2, follow at the end in ascending order: 5 and then 8.

In conclusion, after applying the custom key sorting strategy to arr1, we get the final sorted array: [4, 4, 6, 5, 8], successfully matching the required conditions.

public int[] relativeSortArray(int[] arr1, int[] arr2) {

// Fill the map with the positions of the elements

elementToIndexMap.put(arr2[index], index);

// Function to sort arr1 with respect to the order in arr2

std::unordered_map<int, int> elementToIndex;

// putting it at the end of the sorted array.

std::vector<std::pair<int, int>> tempArray;

for (int value : arr1) {

for (int i = 0; i < arr2.size(); ++i) {</pre>

std::vector<int> relativeSortArray(std::vector<int>& arr1, std::vector<int>& arr2) {

elementToIndex[arr2[i]] = i; // Map the element to its index in arr2

// We will use an unordered_map to store the position (index) of

// each element in arr2; this position is used when sorting arr1

// from original array arrl. If the element does not exist in arr2,

// Create a pair of (index, value) for each element in arr1

// Sort the array: first by the index position and then by the element values

sortedPairs.sort((a, b) => a[0] - b[0] || a[1] - b[1]);

// Map the sorted array back to a single array of values

return sortedPairs.map(pair => pair[1]);

// we assign the index as the size of arr2 which is effectively

// Vector to temporarily store pairs of position and value

for (int index = 0; index < arr2.length; ++index) {</pre>

// Create a hashmap to store the positions of each element in arr2

Map<Integer, Integer> elementToIndexMap = new HashMap<>(arr2.length);

def relative_sort_array(self, arr1: List[int], arr2: List[int]) -> List[int]:

Sort arr1 based on the condition that elements of arr2 should come first

The lambda function inside sorted uses the 'get' method to find the position

in the order they appear in arr2, followed by the remaining elements in

sorted_arr1 = sorted(arr1, key=lambda x: position_map.get(x, 1000 + x))

Create a dictionary to map elements of arr2 to their indices.

position_map = {value: index for index, value in enumerate(arr2)}

of x from position_map if x exists in arr2. Otherwise, it assigns a number 13 # greater than 1000 to ensure that x is positioned after all elements of arr2. # This is based on the assumption that the elements of arr1 do not exceed 1000. 14 15 16

return sorted_arr1 17

ascending order.

Python Solution

class Solution:

Java Solution

class Solution {

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           // Create a new 2D array to hold elements and their corresponding positions
           int[][] elementPositionPairs = new int[arr1.length][0];
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            for (int index = 0; index < elementPositionPairs.length; ++index) {</pre>
               // Use the position from arr2 if present, or else use the value from arr1 plus the length of arr2
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               elementPositionPairs[index] = new int[]{arr1[index], elementToIndexMap.getOrDefault(arr1[index], arr2.length + arr1[index
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17
           // Sort the 2D array based on the positions
           Arrays.sort(elementPositionPairs, (pair1, pair2) -> pair1[1] - pair2[1]);
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           // Place the sorted elements back into arr1
           for (int index = 0; index < elementPositionPairs.length; ++index) {</pre>
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                arr1[index] = elementPositionPairs[index][0];
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           // Return the sorted arr1
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           return arr1;
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C++ Solution
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1 #include <vector>

class Solution {

public:

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2 #include <unordered_map>

#include <algorithm>

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                tempArray.emplace_back(index, value);
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           // Sort the tempArray based on the previously stored positions (indices)
           // If two elements have the same position, they're compared by their values
           std::sort(tempArray.begin(), tempArray.end());
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           // Reassign sorted values back to arr1
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           for (int i = 0; i < arr1.size(); ++i) {
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               arr1[i] = tempArray[i].second; // Set arr1[i] with the value from the sorted pair
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37
           // Return the sorted arrl
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           return arr1;
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40 };
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Typescript Solution
 1 // Given two arrays arrl and arr2, the function sorts elements of arrl
 2 // such that the relative ordering of items in arr1 are the same as in arr2.
  // Elements not present in arr2 will be placed at the end of arr1 in ascending order.
   function relativeSortArray(arr1: number[], arr2: number[]): number[] {
       // Create a mapping of element to its position for quick access
       const elementToIndex: Map<number, number> = new Map();
       // Fill the map with elements of arr2 and their corresponding indices
       for (let index = 0; index < arr2.length; ++index) {</pre>
           elementToIndex.set(arr2[index], index);
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       // Initialize an array to hold pairs of position in arr2 and the value
       const sortedPairs: [number, number][] = [];
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14
       // Map each element of arr1 to its corresponding index in arr2, with a fallback index
15
       for (const element of arr1) {
16
           // Get the index position of the element from the map or use default if not present
           const indexInArr2 = elementToIndex.get(element) ?? arr2.length;
           sortedPairs.push([indexInArr2, element]);
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```

// Find the index if value exists in arr2; for elements not in arr2, set index to arr2's size

int index = elementToIndex.count(value) ? elementToIndex[value] : arr2.size();

Time and Space Complexity

Time Complexity The time complexity of the given code primarily depends on the sorting algorithm used by the sorted function in Python.

arr2 once.

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- The sorted function has a time complexity of O(N log N) where N is the length of arr1. This is because Python uses the TimSort algorithm for sorting, which is a hybrid sorting algorithm derived from merge sort and insertion sort.
- still O(N log N). Therefore, the total time complexity of the algorithm is $O(M + N \log N)$.

element, together with the condition to handle elements not in pos, the overall time complexity of the sorting with these operations is

However, since the key function in the sorted algorithm uses a lookup in the post dictionary, which is 0(1) time complexity for each

Creating the post dictionary has a time complexity of O(M) where M is the length of arr2, since we iterate over each element of

Space Complexity

The space complexity of the algorithm consists of the space needed for the post dictionary and any additional space used by the sorting algorithm:

• The post dictionary has a space complexity of O(M) because it contains as many entries as there are elements in arr2.

- The sorted function returns a new list and internally uses additional space which depends on the specifics of the implementation. In the worst case, this could be O(N) space.
- The total space complexity is therefore O(M + N).