## **Problem Description**

bikes, with the condition that n <= m, meaning there are at least as many bikes as there are workers. Each worker and bike has a unique position on the plane, represented by their coordinates.

The task involves assigning bikes to workers on a campus that is laid out on an X-Y coordinate plane. There are n workers and m

The goal is to assign each worker exactly one bike based on the shortest Manhattan distance between them. The Manhattan distance between two points p1 and p2 is defined as |p1.x - p2.x| + |p1.y - p2.y|, which essentially measures the distance between the points if you could only move along the grid horizontally or vertically.

if one worker has the option of multiple bikes at the same distance, the bike with the lower index is assigned.

The output should be an array where each element represents the index of the bike assigned to the worker at the corresponding index.

If multiple workers could be assigned to the same bike due to equal distances, the worker with the lower index gets priority. Similarly,

#### These combinations are then sorted primarily by the distance, then by the worker index, and finally by the bike index. This guarantees that the shortest and prioritized combinations (based on worker and bike indices) are considered first in the sorted list.

Once sorted, the algorithm iterates through the list of combinations and assigns bikes to workers one by one. During this process, both workers and bikes are marked as 'visited' to ensure that each worker gets only one bike, and each bike is assigned only once.

By sorting the combination of distances and applying the rules of priority in indexing, we ensure an efficient distribution of bikes to

The solution approach requires creating a list of all possible combinations of workers and bikes along with their Manhattan distances.

workers based on the proximity and predefined rules. This approach guarantees that the problem constraints will be respected and the correct assignments will be made.

**Solution Approach** The reference solution provided implements a straightforward but effective algorithm to solve the problem with a greedy approach.

### 1. Collect Pair Information: Initially, the solution collects all possible combinations of worker and bike pairs along with their Manhattan distances. This is done using a nested loop to iterate over each worker and bike to calculate their distance: dist =

Here is a walkthrough of the solution step by step:

abs(workers[i][0] - bikes[j][0]) + abs(workers[i][1] - bikes[j][1]). 2. Distance-Pair Array: These pairs, along with the distances, are stored in an array where each element is a tuple comprising the

- distance, the worker index i, and the bike index j: (dist, i, j). 3. Sorting Pair Array: Next, the tuples are sorted based on the first, second, and third elements (distance, worker index, and bike index, respectively). This orders the pairings primarily by the shortest distance, then by the smallest worker index, and lastly by the smallest bike index if there's a tie on the distance.
- 4. Initialize Visited Flags: Two arrays vis1 for workers and vis2 for bikes are created to keep track of whether a worker has been assigned a bike and whether a bike has been taken by a worker.

5. Assignment: The solution iterates through the sorted array and checks the visited flags for both workers and bikes. If both are

assigned to the worker: ans [i] = j and the visited flags are updated to True. 6. Return Assignment: Once all workers have been assigned a bike, the ans array, which holds the mapping of workers to their assigned bike indices, is returned.

It is worth noting that this approach leverages the Python's tuple sorting behavior which automatically sorts the tuples based on

their first element, and in case of ties, it uses the second element, and so on. The algorithm also makes efficient use of space by

using boolean arrays to keep track of the workers and bikes that have already been assigned. This algorithm achieves a balance

unvisited, (i.e., the worker has not yet been assigned a bike and the bike has not been assigned to any worker), the bike is

between performance and clarity, allowing for an effective solution to the problem.

 Workers' coordinates: [(0, 0), (2, 1)] (Worker 0 at (0, 0) and Worker 1 at (2, 1)) Bikes' coordinates: [(1, 0), (2, 2), (2, 3)] (Bike 0 at (1, 0), Bike 1 at (2, 2), and Bike 2 at (2, 3)) Step 1: Collect Pair Information We begin by calculating the Manhattan distance for all possible worker-bike pairs:

#### Worker 0 & Bike 1: abs(0 - 2) + abs(0 - 2) = 4 • Worker 0 & Bike 2: abs(0 - 2) + abs(0 - 3) = 5

• Worker 0 & Bike 0: abs(0 - 1) + abs(0 - 0) = 1

• Worker 1 & Bike 2: abs(2-2) + abs(1-3) = 2

vis1 = [False, False] (for workers)

is assigned Bike 1.

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Python Solution

vis2 = [False, False, False] (for bikes)

num workers = len(workers)

num\_bikes = len(bikes)

distances.sort()

Example Walkthrough

**Initial Setup:** 

• Worker 1 & Bike 0: abs(2 - 1) + abs(1 - 0) = 2• Worker 1 & Bike 1: abs(2-2) + abs(1-2) = 1

Step 4: Initialize Visited Flags Create boolean arrays to track the visited workers and bikes:

Step 2: Distance-Pair Array Construct an array with the pairs and their distances:

Let's consider a small example to illustrate the solution approach with 2 workers and 3 bikes.

- [(1, 0, 0), (4, 0, 1), (5, 0, 2), (2, 1, 0), (1, 1, 1), (2, 1, 2)]
- [(1, 0, 0), (1, 1, 1), (2, 1, 0), (2, 1, 2), (4, 0, 1), (5, 0, 2)]

Step 3: Sorting Pair Array Sort the array of tuples by distance, then by worker index, and then by bike index:

• Worker 0 is assigned Bike 0 because the pair (1, 0, 0) comes first and both Worker 0 and Bike 0 are unvisited. Mark Worker 0 and Bike 0 as visited.

Remaining pairs either contain a worker or a bike that has already been assigned, so we skip them.

+ abs(workers[worker\_index][1] - bikes[bike\_index][1])

Skip to the next unvisited pair (1, 1, 1). Assign Worker 1 to Bike 1 and mark them as visited.

Step 6: Return Assignment The ans array after the assignments is [0, 1], indicating that Worker 0 is assigned Bike 0 and Worker 1

Step 5: Assignment Iterate through the sorted pairs and assign bikes to workers:

Now the assignment mapping of workers to bike indices is returned, completing the algorithm execution for the given example.

- 1 from typing import List from itertools import product class Solution: def assignBikes(self, workers: List[List[int]], bikes: List[List[int]]) -> List[int]:
- distances = [] # List to hold the distance and associated worker and bike indices 8 9 # Calculate the manhattan distances between each worker-bike pair and store them 10 for worker index, bike index in product(range(num workers), range(num bikes)): 11 distance = abs(workers[worker index][0] - bikes[bike index][0]) \

# Mark the bike as visited

21 visited\_bikes = [False] \* num\_bikes 22 assignments = [-1] \* num workers # -1 to indicate unassigned workers 23 24 # Assign bikes to workers based on sorted distances

# Initialize the visited lists for workers and bikes

for distance, worker\_index, bike\_index in distances:

visited\_bikes[bike\_index] = True

# Sort the distances list in ascending order

visited\_workers = [False] \* num\_workers

distances.append((distance, worker\_index, bike\_index))

# Assign a bike to a worker if both are unvisited/unassigned

// Keep track of which workers and bikes have been visited/assigned

if (!workersVisited[workerIdx] && !bikesVisited[bikeIdx]) {

return assignments; // Return the final assignment of bikes to workers

int[] assignments = new int[numWorkers]; // Result array with assignments

// If the current worker and bike have not been visited yet, assign the bike to the worker

boolean[] workersVisited = new boolean[numWorkers];

// Assign bikes to workers based on the sorted distances

boolean[] bikesVisited = new boolean[numBikes];

workersVisited[workerIdx] = true;

assignments[workerIdx] = bikeIdx;

bikesVisited[bikeIdx] = true;

for (int[] arr : distances) {

int bikeIdx = arr[2];

// Function to assign bikes to workers

int workerIdx = arr[1];

if not visited\_workers[worker\_index] and not visited\_bikes[bike\_index]:

visited\_workers[worker\_index] = True # Mark the worker as visited

assignments[worker\_index] = bike\_index # Assign the bike to the worker

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           return assignments
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Java Solution
    class Solution {
         public int[] assignBikes(int[][] workers, int[][] bikes) {
             int numWorkers = workers.length;
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             int numBikes = bikes.length;
             // Create an array to store distance, worker index, and bike index
             int[][] distances = new int[numWorkers * numBikes][3];
  8
             // Calculate the Manhattan distance between each worker and bike pair
             int index = 0;
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             for (int i = 0; i < numWorkers; ++i) {</pre>
                 for (int j = 0; j < numBikes; ++j) {
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                     int dist = Math.abs(workers[i][0] - bikes[j][0]) +
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                                Math.abs(workers[i][1] - bikes[j][1]);
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                     distances[index++] = new int[] {dist, i, j};
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             // Sort the distances array by distance first, then worker index, then bike index
             Arrays.sort(distances, (a, b) -> {
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                 if (a[0] != b[0]) {
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                     return a[0] - b[0]; // Compare distances
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                 if (a[1] != b[1]) {
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                     return a[1] - b[1]; // Compare worker indices
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                 return a[2] - b[2]; // Compare bike indices
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             });
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C++ Solution

#include <vector>

#include <algorithm>

using namespace std;

2 #include <tuple>

6 class Solution {

public:

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vector<int> assignBikes(vector<vector<int>>& workers, vector<vector<int>>& bikes) {
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             int numWorkers = workers.size();
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             int numBikes = bikes.size();
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             // Create an array to store the distances and indices
             vector<tuple<int, int, int>> allPairs(numWorkers * numBikes);
             // Iterate over each worker and each bike to calculate distances
             for (int workerIndex = 0, pairIndex = 0; workerIndex < numWorkers; ++workerIndex) {</pre>
                 for (int bikeIndex = 0; bikeIndex < numBikes; ++bikeIndex) {</pre>
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                     int distance = abs(workers[workerIndex][0] - bikes[bikeIndex][0]) +
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                                    abs(workers[workerIndex][1] - bikes[bikeIndex][1]);
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                     // Store the tuple of distance, worker index, and bike index
 22
                     allPairs[pairIndex++] = {distance, workerIndex, bikeIndex};
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             // Sort all pairs by distance, then by worker index, then by bike index
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             sort(allPairs.begin(), allPairs.end());
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             // Vectors to keep track of which workers and bikes have been visited
 30
             vector<bool> workerVisited(numWorkers, false);
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             vector<bool> bikeVisited(numBikes, false);
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             // Initialize the answer vector to store the bike index for each worker
 34
             vector<int> answer(numWorkers, -1); // Default to -1 for unassigned
 35
 36
             // Iterate over each sorted pair
             for (auto& [distance, workerIndex, bikeIndex] : allPairs) {
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 38
                 // If the current worker and bike have not been visited
 39
                 if (!workerVisited[workerIndex] && !bikeVisited[bikeIndex]) {
 40
                     // Assign the bike to the worker
                     workerVisited[workerIndex] = true;
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                     bikeVisited[bikeIndex] = true;
 43
                     answer[workerIndex] = bikeIndex;
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             // Return the final assignment of bikes to workers
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             return answer;
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 50 };
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Typescript Solution
   // Import necessary elements from array functions
    import { abs, sort } from 'math';
    // Type definition for a pair representing distance, worker index, and bike index
    type WorkerBikePair = [number, number, number];
    // Function to calculate the Manhattan distance between a worker and a bike
    function calculateManhattanDistance(worker: number[], bike: number[]): number {
         return abs(worker[0] - bike[0]) + abs(worker[1] - bike[1]);
 10 }
 11
    // Function to assign bikes to workers based on their Manhattan distances
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       // Return the final assignments
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        return answer;
63 }
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```

Time and Space Complexity

1. The nested loop with product (range(n), range(m)) generates all possible pairs of workers and bikes. There are n workers and m bikes, so there are n \* m pairs. This operation has a time complexity of O(n \* m).

log n) complexity, and here we are sorting n \* m elements.

The time complexity and space complexity of the given code can be analyzed as follows:

function assignBikes(workers: number[][], bikes: number[][]): number[] {

// Array to store all worker-bike pairs with their corresponding distances

const distance: number = calculateManhattanDistance(workers[workerIndex], bikes[bikeIndex]);

for (let workerIndex = 0; workerIndex < numWorkers; ++workerIndex) {</pre>

for (let bikeIndex = 0; bikeIndex < numBikes; ++bikeIndex) {</pre>

// Sort the pairs by distance, then by worker index, then by bike index

// If distance and worker index are the same, compare by bike index

// Arrays to keep track of whether workers or bikes are already matched

// Proceed if both the worker and the bike are not yet visited

if (!workerVisited[workerIndex] && !bikeVisited[bikeIndex]) {

// Assign the bike to the worker in the answer array

const workerVisited: boolean[] = new Array(numWorkers).fill(false);

const bikeVisited: boolean[] = new Array(numBikes).fill(false);

// Array to keep the results of bike assignments for each worker

const answer: number[] = new Array(numWorkers).fill(-1);

// Assign bikes to workers based on the sorted allPairs

workerVisited[workerIndex] = true;

bikeVisited[bikeIndex] = true;

answer[workerIndex] = bikeIndex;

for (const [distance, workerIndex, bikeIndex] of allPairs) {

allPairs.push([distance, workerIndex, bikeIndex]);

// If distance is the same, compare by worker index

const numWorkers: number = workers.length;

// Populate the array with distances and indices

const numBikes: number = bikes.length;

const allPairs: WorkerBikePair[] = [];

allPairs.sort((a, b) => {

// Compare by distance

return a[0] - b[0];

return a[1] - b[1];

// Mark as matched

if (a[0] !== b[0]) {

if (a[1] !== b[1]) {

return a[2] - b[2];

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});

# complexity from O(n \* m).

**Time Complexity** 

3. Appending each tuple (dist, i, j) to the arr list is done n \* m times, which is also 0(n \* m). 4. The sort operation on arr has a time complexity of  $0(n * m * \log(n * m))$  because sorting is typically implemented with an 0(n \* m)

2. Calculating the Manhattan distance for each worker-bike pair takes constant time, so this does not change the overall time

- 5. The final loop goes over the sorted array, which has n \* m elements, and it assigns bikes to workers. Since each worker and bike is visited at most once, the loop will run at most n + m times - because once every worker has a bike, the process stops. The
- time complexity for this loop is thus O(n + m). Therefore, the overall time complexity of the code is dominated by the sorting step, which is  $0(n * m * \log(n * m))$ .
- Space Complexity 1. The arr list contains n \* m tuples, so it has a space complexity of O(n \* m).

2. The vis1 and vis2 arrays each have a space complexity of O(n) and O(m) respectively. Adding them together gives O(n + m).

3. The answer array ans has a space complexity of O(n) because it must hold an assignment for each worker. The additional space required is for the tuples and the arrays to track the visited workers and bikes. Since these do not depend on

the size of the input apart from the number of workers and bikes, we can consider them all in aggregate. Hence, the overall space complexity is 0(n \* m) as it is the term that grows more significantly with the input size.