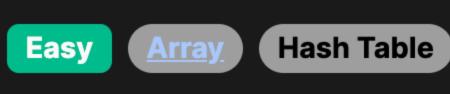
# 2869. Minimum Operations to Collect Elements



## **Problem Description**

You are given an array of positive integers, named nums, and a target number represented by the integer k. Your goal is to collect a list of elements labeled from 1 to k through a specific operation. The operation is defined as removing the last element from the array nums and placing that element in your collection. You must determine the fewest number of operations needed to collect all elements from 1 to k. If an element is not present in nums or cannot be obtained by performing the allowed operations, it implies that it is impossible to collect all elements from 1 to k.

For example, if your array nums is [1,3,2,4,3] and k is 3, you can perform the following operations:

• Remove 3 from nums and add it to the collection. Now your collection has [3] and nums is [1,3,2,4].

• Remove 4 from nums but do not add to the collection since 4 is not needed (we want elements from 1 to k=3). • Remove 2 from nums and it goes to the collection. Your collection is now [3,2] and nums is [1,3].

is to figure out this minimum number of operations for any given array nums and integer k.

- Finally, remove 3 from nums and since it's already in the collection, you can ignore it. Remove 1 from nums and add it to your collection.
- After these 4 operations, your collection has the elements [1, 2, 3], and thus, the minimum number of operations is 4. The task

Intuition

### The intuition behind the solution is to address the problem efficiently by working backwards, starting from the end of the array -

elements are getting added to the collection and ensure the following: • We only care about elements that are less than or equal to k, because we want to collect elements 1 to k. • We avoid adding duplicates to our collection because each number from 1 to k should only be collected once.

since it's the only place where we can remove elements - and moving towards the front. By doing this, we keep a check on what

We use a list called is added that keeps track of whether an element has been added to the collection. This array is of size k, where each index represents an element from 1 to k, and the value at each index represents whether the corresponding element

If we encounter an element, while traversing from the end, that is less than or equal to k and has not been added

has been added to the collection.

(is\_added[element - 1] is False), we mark it as added and increase our count of unique elements. We continue this process until our count reaches k, which means we have all the elements from 1 to k. Solution Approach

The solution involves a single pass through the given array in reverse order, beginning from the last element and moving towards the first. This strategy is chosen because elements can only be removed from the end of the array. The language of choice for

#### the implementation is Python. Here's a step-by-step explanation of the solution with reference to the provided code snippet:

values in is\_added are set to False, indicating that no elements have been collected yet. We define a count variable count, which keeps the count of unique elements that we have collected so far, starting the count from 0.

An array is\_added of size k is created to keep track of elements from 1 to k that have been added to our collection. Initially, all

For each element nums [i] encountered during the traversal:

We iterate through the array nums from the last element to the first, using a reverse loop. This is implemented by the loop for

We check if nums [i] is greater than k or if is\_added[nums[i] - 1] is True (the element has already been added to the

collection). If either condition is true, we continue to the next iteration without performing any operations since we either

- don't need the element or have already collected it. If nums [i] is needed and has not been added to the collection yet, we set is\_added [nums [i] - 1] to True and increment
- our count.

i in range(n - 1, -1, -1); where n is the length of the nums array.

collection. Once we've collected all required elements, we immediately return the result.

solution with a time complexity of O(n), where n is the number of elements in nums.

- As soon as our count equals k, we know we have collected all elements from 1 up to k. The minimum number of operations required is then the total length of the array minus the current index i, which gives us n - i. This approach efficiently ensures that we do not collect unnecessary elements and simultaneously avoid duplication in our
- **Example Walkthrough**

Using this algorithm, we leverage simple data structures such as an array (is\_added) and a counter variable to reach an optimal

The is\_added array will be initialized with values [False, False, False] which signifies that none of the numbers 1, 2, or 3 have been added to our collection yet.

○ Since 2 is less than or equal to k and hasn't been added to the collection yet (is\_added[2 - 1] is False), we add it by setting is\_added[1] to

3 is less than or equal to k and is not present in the collection (is\_added[3 - 1] is False), so we add it. Our is\_added array

6 (length of nums) - 3 (current index) = 3 operations to collect all necessary elements. Therefore, the fewest number of

now becomes [True, True, True]. At this point, we have all elements from 1 to k in our collection, and we stop the iteration.

Let's consider a small example to illustrate the solution approach with the array nums = [5,4,2,3,1,2] and target k = 3.

#### True. Our collection becomes [False, True, False], indicating that 2 has been added. Moving to nums [4] which is 1:

**Python** 

Next, we look at nums [3] which is 3:

operations required in this example is 3.

for i in range(n - 1, -1, -1):

continue

if count == k:

return n - i

• Following our rules, we add 1 to the collection and is\_added becomes [True, True, False].

Since we have collected all the numbers from 1 to k after reaching the 3rd index from the right (inclusive), we have done a total of

# Start iterating over the list from the end to the beginning

# If it is not possible to perform the operation, return -1

// This function calculates the minimum number of operations required

int countDistinct = 0; // Variable to count distinct integers

// Mark the number as added because it is distinct

vector<bool> isAdded(n, false); // Create a boolean vector to track added numbers

// Start from the end of the vector and look for distinct integers until 'k' are found

// If current number is greater than k or already counted as distinct, skip it

// If we have found 'k' distinct numbers, return the number of operations,

// which is the difference between array length and starting index

// The problem's constraints should ensure that this situation doesn't happen.

// Note that the loop is missing an exit condition and might lead

// to an out of bounds access in the nums vector or an infinite loop

// to reduce array 'nums' such that there are 'k' distinct integers

int minOperations(vector<int>& nums, int k) {

isAdded[nums[i] - 1] = true;

if (countDistinct == k) {

// if 'k' distinct numbers are not found.

return n - i;

for (int i = n - 1; ; --i) {

continue;

countDistinct++;

**}**;

**TypeScript** 

int n = nums.size(); // Obtain the size of nums

if (nums[i] > k || isAdded[nums[i] - 1]) {

// Increase the count of distinct numbers

# Skip if the number is greater than k or already added

# If we have added k unique numbers, return the number of operations

if nums[i] > k or is\_added[nums[i] - 1]:

We iterate over nums starting from the last element, so our loop begins at nums [5] which is 2.

Solution Implementation

class Solution: def min\_operations(self, nums: List[int], k: int) -> int: # Create a list to track if the required numbers have been added  $is\_added = [False] * k$ count = 0 # Counter for unique numbers added n = len(nums) # Calculate the length of the nums list

```
# Mark the number as added
is_added[nums[i] - 1] = True
count += 1 # Increment the counter by 1
```

return -1

import java.util.List;

class Solution {

Java

from typing import List

```
// Method to find the minimum number of operations to add the first k positive integers into the list
    public int minOperations(List<Integer> nums, int k) {
       // Array to keep track of which numbers between 1 to k have been added
       boolean[] isNumberAdded = new boolean[k];
       // Get the size of the input list
       int n = nums.size();
       // Counter for unique numbers added to the list
       int count = 0;
       // Iterate in reverse through the list
        for (int i = n - 1; i >= 0; i--) {
            int currentValue = nums.get(i);
           // If the current value is greater than k or already marked as added, skip it
           if (currentValue > k || isNumberAdded[currentValue - 1]) {
               continue;
            // Mark this number as added
            isNumberAdded[currentValue - 1] = true;
           // Increment the count of unique numbers
           count++;
           // If we have added k unique numbers, return the number of operations
           if (count == k) {
               return n - i;
       // If we exit the loop without returning, there's an error, so return —1 as it shouldn't happen
       // Each number between 1 and k should exist in a properly-sized list
       return -1;
C++
#include <vector>
using namespace std;
class Solution {
public:
```

```
function minOperations(nums: number[], k: number): number {
   // Get the length of the input array.
   const arrayLength: number = nums.length;
   // Initialize an array to keep track of which numbers have been added to the sequence.
   const isAdded: boolean[] = Array(k).fill(false);
   // Initialize count to keep track of unique numbers encountered that are not more than k.
    let uniqueCount: number = 0;
   // Iterate backwards through the array.
   for (let i: number = arrayLength - 1;; --i) {
       // If the current number is greater than k or it has already been counted, skip it.
       if (nums[i] > k || isAdded[nums[i] - 1]) {
           continue;
       // Mark the current number as added.
       isAdded[nums[i] - 1] = true;
       // Increment the count of unique numbers.
       ++uniqueCount;
       // If we have encountered k unique numbers, return the size of the sequence.
       if (uniqueCount === k) {
            return arrayLength - i;
```

```
// The loop was intentionally constructed to run indefinitely, control exits from within the loop.
      // If the function has not returned within the loop, it's unexpected as per the problem statement,
      // and may indicate an issue with the inputs. The following return statement is technically unreached.
      return -1; // Return an impossible count as indication of an error.
from typing import List
class Solution:
   def min_operations(self, nums: List[int], k: int) -> int:
       # Create a list to track if the required numbers have been added
        is_added = [False] * k
        count = 0 # Counter for unique numbers added
       n = len(nums) # Calculate the length of the nums list
       # Start iterating over the list from the end to the beginning
       for i in range(n - 1, -1, -1):
            if nums[i] > k or is_added[nums[i] - 1]:
               # Skip if the number is greater than k or already added
               continue
           # Mark the number as added
            is_added[nums[i] - 1] = True
            count += 1 # Increment the counter by 1
           if count == k:
               # If we have added k unique numbers, return the number of operations
               return n - i
```

## **Time Complexity** The time complexity of the given code is O(n), where n is the length of the input array nums. This is because there is a single for

Time and Space Complexity

# If it is not possible to perform the operation, return -1

therefore do not contribute to the space complexity beyond a constant factor.

## loop that iterates backwards over the array nums, and in each iteration, it performs a constant time check and assignment operation. Since these operations do not depend on the size of k and there are no nested loops, the iteration will occur n times,

return -1

leading to a linear time complexity with respect to the size of the array. **Space Complexity** The space complexity of the given code is O(k). The is\_added list is the only additional data structure whose size scales with the input parameter k. Since it is initialized to have k boolean values, the amount of memory used by this list is directly proportional to

the value of k. The rest of the variables used within the function (like count, n, and i) use a constant amount of space, and