1115. Print FooBar Alternately

Medium

## In this concurrency problem, we have a code snippet defining a class FooBar with two methods: foo() and bar(). These

**Problem Description** 

Concurrency

methods are called by two separate threads, the output is "foobar" repeated n times, with "foo" and "bar" alternating properly. To achieve this, we need to implement a mechanism that enforces the order of execution, such that the foo() method must print "foo" before the bar() method prints "bar," and this pattern is repeated for the entire loop.

methods print the strings "foo" and "bar", respectively, in a loop up to n times. The challenge is to ensure that when these

Intuition Concurrency problems often require synchronization techniques to ensure that multiple threads can work together without

conflicts or race conditions. In this case, we need to make sure that "foo" is always printed before "bar."

### The intuition behind the solution is to use semaphores—a classic synchronization primitive—to coordinate the actions of the

threads. In the context of this problem, we use two semaphores: one that controls the printing of "foo" (self.f) and one that controls the printing of "bar" (self.b). The self.f semaphore is initially set to 1, allowing the foo() method to print immediately. After printing, it releases the self.b

semaphore, which is initially set to 0, thus preventing bar() from printing until foo() is printed first. Once self.b is released by foo(), the bar() method can print "bar" and then release the self.f semaphore to allow the next "foo" to be printed. This

**Solution Approach** The solution utilizes the Semaphore class from Python's threading module as the primary synchronization mechanism, allowing us to enforce the strict alternation between foo and bar.

#### printed. Two Semaphore objects are created: self.f for "foo" with an initial value of 1, and self.b for "bar" with an initial

value of 0. The initial values are critical: self.f is set to 1 to allow foo() to proceed immediately, and self.b is set to 0 to block bar() until foo() signals it by releasing self.b.

Initialization: The FooBar class is initialized with an integer n, which represents the number of times "foobar" should be

The foo Method: It contains a loop that runs n times.

Once the semaphore is acquired, printBar() is executed to print "bar".

again, hence ensuring the sequence starts with "foo".

alternating process continues until the loop completes n times.

Here's a step-by-step explanation of the code implementation:

- Each iteration begins with self.f.acquire(), which blocks if the semaphore's value is 0. Since self.f is initialized to 1, foo() can start immediately on the first iteration. The printFoo() function is executed, printing "foo". • After printing "foo", the self.b.release() is called. This increments the count of the self.b semaphore, signaling the bar() method (if it is waiting) that it can proceed to print "bar".
- Each iteration starts by calling self.b.acquire(), which waits until the self.f semaphore is released by a previous foo() call, ensuring that "foo" has been printed before "bar" can proceed.

**Example Walkthrough** 

The bar Method:

It's also a loop running n times.

This alternating semaphore pattern locks each method in a waiting state until the other method signals that it has finished its task by releasing the semaphore. Since acquire() decreases the semaphore's value by 1 and release() increases it by 1, this careful incrementing and decrementing of semaphore values guarantees that the print order is preserved and that the strings "foo" and

"bar" are printed in the correct sequence to form "foobar" without getting mixed up or overwritten.

Semaphore self.b is set to 0 (locked), preventing bar() from being called until foo() is done.

self.f.acquire() is called, which succeeds immediately because self.f is 1 (unlocked).

• After printing "bar", it invokes self.f.release() to increment the semaphore count for foo, allowing the next iteration of foo() to print

Here's how the synchronization using semaphores will facilitate this: Initialization:  $\circ$  FooBar object is created with n = 2. Semaphore self.f is set to 1 (unlocked), allowing foo() to be called immediately.

Let's walk through a simple example with n = 2. We want our output to be "foobarfoobar" with "foo" always preceding "bar".

### printFoo() is executed, so "foo" is printed.

**Python** 

Java

C++

#include <semaphore.h>

#include <functional>

class FooBar {

int n ;

private:

public:

**TypeScript** 

let n: number;

n = count;

});

/\*\*

/\*\*

// The number of times to print "foo" and "bar"

let fooPromiseResolver: (() => void) | null = null;

let barPromiseResolver: (() => void) | null = null;

\* @param {number} count - The number of iterations to run the sequence.

\* @param {() => void} printFoo - A callback function that prints "foo".

if (fooPromiseResolver) fooPromiseResolver();

\* Initializes the synchronization primitives.

// Start with the ability to print "foo"

if (canPrintBar) canPrintBar();

\* Prints "foo" to the console or another output.

async function foo(printFoo: () => void): Promise<void> {

await new Promise<void>((resolve) => {

if (canPrintFoo) canPrintFoo();

fooPromiseResolver = resolve;

// The provided callback prints "foo"

barPromiseResolver = null;

\* Prints "bar" to the console or another output.

// Block until 'bar' is printed

// Allow "bar" to be printed

canPrintBar = (() => {

canPrintFoo = null;

// Prevent "bar" from printing until "foo" is printed

function initFooBar(count: number): void {

fooPromiseResolver = null;

for (let i = 0; i < n; i++) {

// Promises and callbacks for signaling

// Deferred promise resolvers

canPrintFoo = (() => {

canPrintBar = null;

printFoo():

}):

let canPrintFoo: (() => void) | null = null;

let canPrintBar: (() => void) | null = null;

class FooBar {

class FooBar:

from threading import Semaphore

def init (self, n: int):

self.sem foo = Semaphore(1)

self.sem\_bar = Semaphore(0)

for in range(self.n):

from typing import Callable

First Iteration:

Second Iteration:

1. foo() method is called by Thread 1.

printBar() is executed, printing "bar" after "foo".

printFoo() is executed, and another "foo" is printed.

self.b.release() is called, incrementing self.b to 1, which unlocks bar(). 2. bar() method is called by Thread 2. self.b.acquire() is called, which succeeds because self.b was released by foo().

self.f.release() is called, setting self.f back to 1 (unlocked) and allowing the next foo() to proceed.

1. Again, foo() method is called by Thread 1. ■ This time, since self.f was released by the previous bar() call, self.f.acquire() succeeds again.

self.b.release() is called, incrementing self.b and allowing bar() to be called.

printing functions, ensuring the correct order despite the concurrent execution of threads.

With self.b released, self.b.acquire() allows the thread to proceed.

self.n = n # Number of times "foo" and "bar" are to be printed.

"""Print "bar" n times, ensuring it alternates with "foo"."""

self.sem\_foo.release() # Unlock semaphore for "foo".

print bar() # Provided print function for "bar".

public void foo(Runnable printFoo) throws InterruptedException {

public void bar(Runnable printBar) throws InterruptedException {

// Output "foo"

// Output "bar"

// The number of times to print "foobar"

sem\_t sem\_foo\_, sem\_bar\_; // Semaphores used to coordinate the printing order

self.sem bar.acquire() # Wait for semaphore to be unlocked.

private final int loopCount: // The number of times "foo" and "bar" should be printed.

# Semaphore for "foo" is initially unlocked.

def bar(self, print bar: Callable[[], None]) -> None:

# Semaphore for "bar" is initially locked.

printBar() prints "bar", following the "foo" printed by the last foo() call.

■ Finally, self.f.release() is called, although in this case, it's unnecessary because we've reached our loop condition (n times) and no further foo() calls are needed.

2. bar() method is again called by Thread 2.

- Solution Implementation
- def foo(self, print foo: Callable[[], None]) -> None: """Print "foo" n times, ensuring it alternates with "bar".""" for in range(self.n): self.sem foo.acquire() # Wait for semaphore to be unlocked. print foo() # Provided print function for "foo". self.sem\_bar.release() # Unlock semaphore for "bar".

fooSemaphore.acquire(); // Acquire a permit before printing "foo", ensuring "foo" has the turn to print

barSemaphore.acquire(): // Acquire a permit before printing "bar", ensuring "bar" has the turn to print

fooSemaphore.release(); // Release a permit for "foo" after "bar" is printed, allowing "foo" to print next

barSemaphore.release(); // Release a permit for "bar" after "foo" is printed, allowing "bar" to print next

By the end of the two iterations, we've successfully printed "foobarfoobar". Each foo() preceded a bar() thanks to our

semaphore controls, and at no point could bar() leapfrog ahead of foo(). The semaphores effectively serialized access to the

#### private final Semaphore fooSemaphore = new Semaphore(1); // A semaphore for "foo", allowing "foo" to print first. private final Semaphore barSemaphore = new Semaphore(0); // A semaphore for "bar", initially locked until "foo" is printed. public FooBar(int n) {

this.loopCount = n;

// The method for printing "foo"

printFoo.run();

// The method for printing "bar"

printBar.run();

for (int i = 0; i < loopCount; i++) {

for (int i = 0; i < loopCount; i++) {

```
// Constructor that initializes the semaphores and count
FooBar(int n) : n (n) {
   // Initialize sem foo with a count of 1 to allow "foo" to print first
   sem init(&sem foo , 0, 1);
   // Initialize sem bar with a count of 0 to block "bar" until "foo" is printed
    sem_init(&sem_bar_, 0, 0);
// Deconstructor that destroys the semaphores
~FooBar() {
   sem destroy(&sem foo );
   sem_destroy(&sem_bar_);
// Method for printing "foo"
void foo(std::function<void()> printFoo) {
    for (int i = 0; i < n; ++i) {
        // Wait on sem foo to ensure "foo" is printed first
        sem wait(&sem foo );
        // printFoo() calls the provided lambda function to output "foo"
        printFoo();
        // Post (increment) sem_bar_ to allow "bar" to be printed next
        sem_post(&sem_bar_);
// Method for printing "bar"
void bar(std::function<void()> printBar) {
    for (int i = 0; i < n; ++i) {
        // Wait on sem bar to ensure "bar" is printed after "foo"
        sem wait(&sem bar );
        // printBar() calls the provided lambda function to output "bar"
        printBar();
        // Post (increment) sem_foo_ to allow the next "foo" to be printed
        sem_post(&sem_foo_);
```

#### \* @param {() => void} printBar - A callback function that prints "bar". async function bar(printBar: () => void): Promise<void> { for (let i = 0; i < n; i++) {

await new Promise<void>((resolve) => { barPromiseResolver = resolve; if (canPrintBar) canPrintBar(); }); // The provided callback prints "bar" printBar(); // Allow "foo" to be printed again canPrintFoo = (() => { fooPromiseResolver = null; if (barPromiseResolver) barPromiseResolver(); }): // Block until 'foo' is printed again canPrintBar = null; // Example usage: initFooBar(3); // Initialize printing "foo" and "bar" three times each foo(() => console.log('foo')); bar(() => console.log('bar')); from threading import Semaphore from typing import Callable class FooBar: def init (self, n: int): self.n = n # Number of times "foo" and "bar" are to be printed. # Semaphore for "foo" is initially unlocked. self.sem foo = Semaphore(1) # Semaphore for "bar" is initially locked. self.sem\_bar = Semaphore(0) def foo(self, print foo: Callable[[], None]) -> None: """Print "foo" n times, ensuring it alternates with "bar".""" for in range(self.n): self.sem foo.acquire() # Wait for semaphore to be unlocked. print foo() # Provided print function for "foo". self.sem\_bar.release() # Unlock semaphore for "bar".

# **Time Complexity**

for in range(self.n):

Time and Space Complexity

def bar(self, print bar: Callable[[], None]) -> None:

"""Print "bar" n times, ensuring it alternates with "foo"."""

self.sem\_foo.release() # Unlock semaphore for "foo".

print bar() # Provided print function for "bar".

that the memory usage is constant irrespective of the size of n.

self.sem bar.acquire() # Wait for semaphore to be unlocked.

times, where n is the input that represents the number of times the "foo" and "bar" functions should be called, respectively. The methods invoke acquire and release on semaphores, but the acquire/release operations are constant-time 0(1) operations, assuming that there is no contention (which should not happen here given the strict alternation). The printFoo and printBar functions are also called n times each, and if we consider these functions to have 0(1) time complexity, which is a reasonable assumption for a simple print operation, then this does not change the overall time complexity of the foo and bar methods. **Space Complexity** The space complexity of the FooBar class is 0(1) since the space required does not grow with n. The class maintains fixed resources: two semaphores and one integer variable. No additional space is allocated that would scale with the input n, meaning

The time complexity of the FooBar class methods foo and bar are both O(n). Each method contains a loop that iterates n