573. Squirrel Simulation

Problem Description

Medium Array

location.

travel to collect all the nuts and bring them to the tree, one at a time. The squirrel can move in four directions - up, down, left, and right, moving to the adjacent cell with each move. The distance is measured in the number of moves the squirrel makes. You are provided with: • an array tree with two integers indicating the position (tree_r, tree_c), where tree_r is the row and tree_c is the column of the tree's

In this problem, you are given the dimensions of a garden that is height units tall and width units wide. There is a tree, a squirrel,

and multiple nuts located at different positions within this garden. Your task is to calculate the minimal distance the squirrel must

• an array squirrel with two integers indicating the position (squirrel_r, squirrel_c), where squirrel_r is the row and squirrel_c is the column of the squirrel's initial location. • an array of nuts, where each element is an array [nut_ir, nut_ic] that provides the position of each nut in the garden.

The goal is to return a single integer that is the minimal total distance for all of this back-and-forth movement.

Intuition

since the squirrel must go to the nut and then return to the tree.

initially calculated sum, gives the result of the minimal total distance.

the path the squirrel takes if it chooses that nut first).

d = abs(i - a) + abs(j - b) + c

ans = min(ans, s + d - c * 2)

traveled. One key observation is that, except for the first nut, the squirrel will be traveling from the tree to a nut and then back to the tree for every other nut. This means that for every nut except the first, it will cost double the distance from the tree to the nut

The intuition behind the solution is to recognize that the problem is about optimizing the squirrel's path to minimize the distance

by picking a nut that's closer to its starting location, even if that nut is slightly further from the tree. To arrive at the solution approach, the solution calculates initially the sum of the distances for the squirrel to go from the tree to all the nuts and back to the tree, which would be double the sum of the distances from the tree to each nut. Then, for each nut, it calculates the actual distance the squirrel would travel if it picked up that particular nut first, which involves going from its initial

position to the nut, then to the tree, and subsequently to all other nuts and back to the tree (as already calculated).

However, for the first nut, the squirrel doesn't start at the tree; it starts at its initial position. If the first nut that the squirrel

collects is further from the tree than the squirrel's initial location, the squirrel could potentially reduce the total distance traveled

The additional distance for the first nut is, therefore, the distance from the squirrel to the nut plus the distance from that nut to the tree. The optimization comes from finding the minimum of these additional distances by checking each nut to determine which one should be the first nut that reduces the total travel distance the most. The solution iterates over all the nuts and calculates the difference between the optimized path (going first to each nut) and the

non-optimized path (doubling the distance from the tree to each nut). The minimum of these differences, when added to the

Solution Approach The solution's main strategy is a greedy algorithm, which involves selecting the first nut to pick up based on which choice will

minimize the overall distance traveled. To understand this strategy, let's break down the implementation step-by-step. **Initialize variables:** We store the tree's coordinates in x and y, and the squirrel's starting coordinates in a and b. x, y, a, b = *tree, *squirrel

the squirrel must travel to each nut and then back to the tree, this distance is multiplied by 2.

Precompute and Double the Constant Distances: We compute the sum of all the distances from each nut to the tree. Since

s = sum(abs(i - x) + abs(j - y) for i, j in nuts) * 2

Calculate the Minimum Additional Distance: For each nut, we calculate two distances: c is the distance from the current nut to the tree. od is the distance from the squirrel's initial position to the current nut plus the distance from the current nut to the tree (which represents

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those trips with the trip from the squirrel's starting position).
for i, j in nuts:
    c = abs(i - x) + abs(i - y)
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We then compute the total distance if the squirrel chooses that nut first by adding d to the precomputed sum s and

subtracting c * 2 (since we had previously doubled the distance from this nut to the tree but now we need to replace one of

Find the Optimal First Nut: The loop iterates over all nuts, finding the minimum possible total distance (ans) for the squirrel to

collect all nuts and bring them to the tree by possibly choosing each nut as the first pickup. The first nut to be picked is

Return the Result: Finally, the algorithm returns the minimum additional distance found, which, when added to the double sum of the constant distances, gives the minimum total distance needed to solve the problem.

implicitly chosen by finding the minimum value of ans over all possibilities.

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nut collection process.
Example Walkthrough
  Consider a scenario where our garden has dimensions height = 5 and width = 5, the tree is located at position (2, 2), the
  squirrel is at position (4, 4), and there are three nuts located at positions (0, 1), (1, 3), and (3, 2).
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By the end of the loop, ans holds the minimum total distance that factors in the optimal starting nut to pick up. This optimal path

ensures that the squirrel does not waste additional distance on its first trip, thereby reducing the total distance over the entire

The distance for each nut to the tree would be: ■ For nut at (0, 1): abs(0 - 2) + abs(1 - 2) = 3■ For nut at (1, 3): abs(1 - 2) + abs(3 - 2) = 2■ For nut at (3, 2): abs(3 - 2) + abs(2 - 2) = 1

$\mathbf{c} = 3$ • d = abs(0 - 4) + abs(1 - 4) + c = 7 (distance from squirrel to nut) + 3 (nut to tree) = 10

Initialize variables:

x, y = 2, 2 # Tree's position

a, b = 4, 4 # Squirrel's position

Sum of the distances = 3 + 2 + 1 = 6

s = 12 # Precomputed sum

• For nut at (0, 1):

• For nut at (1, 3):

• For nut at (3, 2):

Find the Optimal First Nut:

less than picking the nut at (0, 1) first).

■ c = 1

Solution Implementation

tree x, tree y = tree

min_distance = inf

for nut x, nut y in nuts:

squirrel_x, squirrel_y = squirrel

Initialize the minimum distance as infinity

int minimumDistance = Integer.MAX VALUE;

// instead of going to the tree first.

int totalDistanceToAllNuts = 0;

for (int[] nut : nuts) {

return minimumDistance;

Distance from the tree to the current nut

class Solution:

c = 2

Precompute and Double the Constant Distances:

Calculate the Minimum Additional Distance:

 \blacksquare s + d - c * 2 = 12 + 10 - 3 * 2 = 16

 \blacksquare s + d - c * 2 = 12 + 6 - 2 * 2 = 14

 \blacksquare s + d - c * 2 = 12 + 4 - 1 * 2 = 14

 \circ Double the sum as the squirrel goes back and forth = 6 * 2 = 12

For each nut, calculate c (tree to nut) and d (squirrel to nut + nut to tree):

= d = abs(1 - 4) + abs(3 - 4) + c = 4 (distance from squirrel to nut) + 2 (nut to tree) = 6

= d = abs(3 - 4) + abs(2 - 4) + c = 3 (distance from squirrel to nut) + 1 (nut to tree) = 4

result, which is the minimum total distance for the squirrel to collect all the nuts and bring them to the tree, is 14.

def minDistance(self, height: int, width: int, tree: List[int], squirrel: List[int], nuts: List[List[int]]) -> int:

Calculate the sum of distances from the tree to all nuts and back (doubled because of the round trip)

Distance from the squirrel to the current nut plus the distance from the current nut to the tree

squirrel_to_nut_plus_nut_to_tree = abs(nut_x - squirrel_x) + abs(nut_y - squirrel_y) + tree_to_nut

Calculate the difference when the squirrel goes to this nut first instead of going to the tree

total_distance = sum(abs(nut_x - tree_x) + abs(nut_y - tree_y) for nut_x, nut_y in nuts) * 2

Iterate through all the nuts to find the one with the minimum extra distance for the squirrel

We replace one of the tree-to-nut round trips with squirrel-to-nut then nut-to-tree

• Thus, the optimal first nut to pick up is either the one at (1, 3) or at (3, 2) since both result in a minimum additional distance of 14 (which is

The minimum distances calculated from picking each nut first are 16, 14, and 14, respectively.

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Return the Result:

    Since both of these nuts yield a minimum additional distance of 14, the squirrel can choose either of these as the first nut. Therefore, the
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Python
from typing import List
from math import inf
```

Deconstruct the tree and squirrel coordinates for ease of use

tree_to_nut = abs(nut_x - tree_x) + abs(nut_y - tree_y)

totalDistanceToAllNuts += calculateDistance(nut, tree);

minimumDistance = Math.min(minimumDistance, currentDistance);

// Helper method to calculate the Manhattan distance between two points a and b

distance_diff = squirrel_to_nut_plus_nut_to_tree - tree_to_nut * 2

public int minDistance(int height, int width, int[] tree, int[] squirrel, int[][] nuts) {

Update the minimum distance if we found a nut that results in a smaller extra distance for the squirrel min_distance = min(min_distance, total_distance + distance_diff) # The result is the smallest distance the squirrel needs to collect all nuts and put them in the tree return min_distance

// Calculate the total distance to collect all nuts and returning them to the tree, done twice (back and forth)

// Subtract the distance saved by the squirrel going directly to the nut, and then to the tree,

// Update the minimum distance if the current distance is less than what we have seen so far

int currentDistance = totalDistanceToAllNuts + distanceToSquirrel - distanceToTree * 2;

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totalDistanceToAllNuts *= 2;
// Try each nut as the first nut to calculate the minimum distance needed
for (int[] nut : nuts) {
    int distanceToTree = calculateDistance(nut, tree); // Distance from the current nut to the tree
    int distanceToSquirrel = calculateDistance(nut, squirrel) + distanceToTree; // Full trip from squirrel to nut and then to
```

Java

};

TypeScript

}):

class Solution {

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private int calculateDistance(int[] a, int[] b) {
        return Math.abs(a[0] - b[0]) + Math.abs(a[1] - b[1]);
C++
class Solution {
public:
    // Calculate the minimum total distance the squirrel must travel to collect all nuts
    // and put them in the tree, starting from the squirrel's initial position.
    int minDistance(int height, int width, vector<int>& tree, vector<int>& squirrel, vector<vector<int>>& nuts) {
        int minTotalDistance = INT MAX;
        int totalDistanceToTree = 0;
        // Calculate the total distance for all nuts to the tree (each trip accounts for going to and returning from the tree).
        for (auto& nut : nuts) {
            totalDistanceToTree += distance(nut, tree);
        totalDistanceToTree *= 2;
        // Evaluate each nut, calculating the distance savings if the squirrel starts from that nut.
        for (auto& nut : nuts) {
            int distanceToTree = distance(nut, tree);
            int distanceToSquirrel = distance(nut, squirrel);
            int currentDistance = distanceToSquirrel + distanceToTree; // Squirrel's first trip distance for this nut
            int distanceSaved = distanceToTree * 2; // The saved distance for fetching this nut last
            int totalCurrentDistance = totalDistanceToTree + currentDistance - distanceSaved;
            minTotalDistance = min(minTotalDistance, totalCurrentDistance);
```

return minTotalDistance; // Example usage: // let minHeight = 5; // let minWidth = 7; // let treePos = [2, 3]; // let squirrelStart = [4, 4];

Update the minimum distance if we found a nut that results in a smaller extra distance for the squirrel min_distance = min(min_distance, total_distance + distance_diff) # The result is the smallest distance the squirrel needs to collect all nuts and put them in the tree return min_distance

distance_diff = squirrel_to_nut_plus_nut_to_tree - tree_to_nut * 2

Calculate the sum of distances from the tree to all nuts and back (doubled because of the round trip)

Distance from the squirrel to the current nut plus the distance from the current nut to the tree

squirrel_to_nut_plus_nut_to_tree = abs(nut_x - squirrel_x) + abs(nut_y - squirrel_y) + tree_to_nut

Main Loop Over Nuts: The code then iterates over all the nuts to determine the minimum extra distance the squirrel has to

travel for going to one of the nuts first before going to the tree. Each iteration includes constant time computations (addition,

Calculate the difference when the squirrel goes to this nut first instead of going to the tree

total_distance = sum(abs(nut_x - tree_x) + abs(nut_y - tree_y) for nut_x, nut_y in nuts) * 2

Iterate through all the nuts to find the one with the minimum extra distance for the squirrel

We replace one of the tree-to-nut round trips with squirrel-to-nut then nut-to-tree

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The code performs a few distinct operations which each contribute to the overall time complexity:
   Initial Summation of Distances (s calculation): The code first calculates the sum of the double distances from the tree to all
   nuts. Since this operation is performed for each nut only once, it runs in O(n), where n is the number of nuts.
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Space Complexity

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No additional data structures: There are no extra data structures that grow with the input size.
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return minTotalDistance; // Calculate the Manhattan distance between two points represented as vectors. int distance(vector<int>& pointA, vector<int>& pointB) {

function distance(pointA: number[], pointB: number[]): number {

// Calculate the total distance for all nuts to the tree

totalDistanceToTree += distance(nut, tree);

const distanceToTree = distance(nut, tree);

// Squirrel's first trip distance for this nut

Initialize the minimum distance as infinity

Distance from the tree to the current nut

subtraction, and comparison). This loop runs in O(n).

tree_to_nut = abs(nut_x - tree_x) + abs(nut_y - tree_y)

// The saved distance for fetching this nut last

const distanceToSquirrel = distance(nut, squirrel);

const currentDistance = distanceToSquirrel + distanceToTree;

// Each trip accounts for going to and returning from the tree.

let minTotalDistance = Number.MAX_SAFE_INTEGER;

let totalDistanceToTree = 0;

nuts.forEach((nut) => {

totalDistanceToTree *= 2;

nuts.forEach((nut) => {

// let nutsPositions = [[3, 0], [2, 5]];

min_distance = inf

for nut x, nut y in nuts:

return abs(pointA[0] - pointB[0]) + abs(pointA[1] - pointB[1]);

// Calculate the Manhattan distance between two points represented as arrays.

// and put them in the tree, starting from the squirrel's initial position.

return Math.abs(pointA[0] - pointB[0]) + Math.abs(pointA[1] - pointB[1]);

// Calculate the minimum total distance the squirrel must travel to collect all nuts

function minDistance(height: number, width: number, tree: number[], squirrel: number[], nuts: number[][]): number {

// Evaluate each nut, calculating the difference in distance if the squirrel goes to this nut first.

const distanceSaved = distanceToTree * 2; // Total distance if starting from this nut const totalCurrentDistance = totalDistanceToTree - distanceSaved + currentDistance; minTotalDistance = Math.min(minTotalDistance, totalCurrentDistance); });

// let minDist = minDistance(minHeight, minWidth, treePos, squirrelStart, nutsPositions); // console.log(minDist); from typing import List from math import inf class Solution: def minDistance(self, height: int, width: int, tree: List[int], squirrel: List[int], nuts: List[List[int]]) -> int: # Deconstruct the tree and squirrel coordinates for ease of use tree x, tree y = tree squirrel_x, squirrel_y = squirrel

Time and Space Complexity **Time Complexity**

Therefore, the total time complexity of the given code is O(n), dominated by the iterative operations done for each nut. The space complexity of the code can be analyzed based on the space used by variables that are not input-dependent:

Constant space for variables: The variables x, y, a, b, s, and ans use constant space, 0(1). Hence, the space complexity of the code is 0(1).