2531. Make Number of Distinct Characters Equal String) Medium Hash Table Counting

Leetcode Link

In this problem, we are given two zero-indexed (meaning indexing starts from 0) strings word1 and word2. We are allowed to perform

Problem Description

a move which involves choosing two indices i from wordl and j from word2 (both indices must be within the bounds of their respective strings) and swapping the characters at these positions, i.e., word1[i] and word2[j]. The objective is to determine if it is possible to equalize the number of distinct characters in both word1 and word2 using exactly one

such move. If this is possible, our function should return true. Otherwise, it should return false.

Intuition

The intuition behind the solution is based on counting characters in each string and then considering every possible swap to see if it

To keep track of the characters, we use two arrays cnt1 and cnt2 of size 26 (assuming the input strings consist of lowercase

Then, we iterate through the counts of both cnt1 and cnt2. For every pair of characters (i, j) where cnt1[i] and cnt2[j] are

can lead us to the goal of equalizing the distinct number of characters in both strings with just one move.

We first count how many times each character appears in both word1 and word2.

both non-zero (indicating that character i is available to swap in word1 and character j in word2), we emulate a swap. After the virtual swap, we calculate the number of distinct characters currently present in cnt1 and cnt2.

alphabetic characters only). Each array corresponds to counting occurrences of characters in word1 and word2, respectively.

3 for c in word2:

word2 to word1.

return True

space besides the two counting arrays.

both words by performing exactly one swap.

1, ..., 0] also corresponding to [a, b, c, ..., z].

cnt1 becomes [0, 1, 1, 1, 0, 0, ..., 0]

cnt2 becomes [1, 0, 0, 0, 1, 1, ..., 0]

cnt2 has 3 distinct characters (a, e, f)

Now, we attempt every possible swap between characters in word1 and word2:

Example Walkthrough

cnt2[ord(c) - ord('a')] += 1

respective strings), a virtual swap is performed:

1 if sum(v > 0 for v in cnt1) == sum(v > 0 for v in cnt2):

Reverting Swap: If the numbers are not equal, the swap is reverted:

1 cnt1[i], cnt2[j] = cnt1[i] - 1, cnt2[j] - 1

2 cnt1[j], cnt2[i] = cnt1[j] + 1, cnt2[i] + 1

1 cnt1[i], cnt2[j] = cnt1[i] + 1, cnt2[j] + 1

Here's the thinking process to arrive at the solution:

- We check if the number of distinct characters is now equal in both. If it is, we return true. If it isn't, we revert the swap back to its original state and continue with the next possible swap pair. The reason we try every possible swap is that the distinct number of characters is influenced by both the characters being swapped
- Solution Approach
- The implementation of the solution follows a simple yet efficient brute-force approach to verify whether a single swap can make the number of distinct characters equal in both strings. Here's a walk-through:

in and out. So, to find out if any swap can achieve our goal, we have to examine each possibility.

• Data Structures: Two arrays cnt1 and cnt2 are utilized, each with a length of 26 corresponding to the 26 letters in the English alphabet. These arrays are used to count the occurrences of each character in word1 and word2 respectively.

Counting Characters: Iterate over word1 and word2 to count each character. This is done by converting the character to its ASCII

value with ord(c), subtracting the ASCII value of 'a' to normalize the index to 0-25, and incrementing the count at that index in

cnt1 or cnt2. This looks like the following: 1 for c in word1: cnt1[ord(c) - ord('a')] += 1

 Attempting Swaps: The next step is to consider every possible swap. For this, nested loops are used to go over every index i in cnt1 and every index j in cnt2. If cnt1[i] and cnt2[j] are both greater than 0 (meaning both characters are present in their

This emulates moving one occurrence of the i-th character from word1 to word2 and one occurrence of the j-th character from

• Checking Distinct Characters: After emulating the swap, a check is performed to determine if the number of distinct characters

in both arrays is the same. This is done by summing up the count of indices greater than 0 in both arrays:

is found after all possibilities have been considered, the function eventually returns False.

2 cnt1[j], cnt2[i] = cnt1[j] - 1, cnt2[i] - 1 And the algorithm continues to the next possible swap.

This algorithm is efficient in the sense that it goes through a limited set of possible swaps (at most 26 * 26) and requires no extra

• Result: If a successful swap that balances the number of distinct characters is found, the function returns True. If no such swap

Let's illustrate the solution approach with a small example: Suppose we have word1 = "abc" and word2 = "def". Our goal is to determine if we can equalize the number of distinct characters in

cnt1 (for word1) would be [1, 1, 1, 0, 0, 0, ..., 0] corresponding to [a, b, c, ..., z]. cnt2 (for word2) would be [0, 0, 0, 1, 1,

First, we initialize two arrays to count the occurrences of each character, cnt1 and cnt2. After iterating through each word:

1. Trying to swap a from word1 with d from word2 We adjust our count arrays to reflect this potential swap:

We now check if the number of distinct characters in both is the same: cnt1 has 3 distinct characters (b, c, d)

Revert back to original counts:

cnt1 becomes [1, 1, 1, 0, 0, 0, ..., 0]

cnt2 becomes [0, 0, 0, 1, 1, 1, ..., 0]

2. Trying to swap a from word1 with e from word2

Since the number of distinct characters is equal, we return True.

Update the character frequency for word2

count2[ord(char) - ord('a')] += 1

for j, frequency2 in enumerate(count2):

if frequency1 and frequency2:

return True

public boolean isItPossible(String word1, String word2) {

// Count the frequency of each character in wordl

// Count the frequency of each character in word2

bool isItPossible(std::string word1, std::string word2) {

// Populate frequency arrays for wordl

// Populate frequency arrays for word2

for (int j = 0; j < 26; ++j) {

--count1[i];

--count2[j];

++count1[j];

++count2[i];

int difference = 0;

// Iterate over each letter in the alphabet

// Iterate over each letter in the alphabet

if (count1[i] > 0 && count2[j] > 0) {

for (int k = 0; k < 26; ++k) {

if (count1[k] > 0) {

++difference;

if (count2[k] > 0) {

if (difference == 0) {

// Undo the simulated swap

return true;

++count1[i];

++count2[j];

--count1[j];

--count2[i];

--difference;

// Check if current letters are present in both words

// Simulate swapping the letters by updating counts

// Check if the frequencies match for each letter

// Variable to track the sum of differences in frequencies

// If the sum of differences is zero, words can be made identical

// Arrays to store letter frequencies for both words

for (int i = 0; i < word1.length(); ++i) {</pre>

for (int i = 0; i < word2.length(); ++i) {</pre>

countWord1[word1.charAt(i) - 'a']++;

int[] countWord1 = new int[26];

int[] countWord2 = new int[26];

for i, frequency1 in enumerate(count1):

... and so on for every character in word1 vs every character in word2. For the sake of this example, we already found a swap that works, so there's no need to continue. The function would now return

True. If a swap could not equalize the number of distinct characters, then we would finally return False after exhausting all the

However, if we need to continue to illustrate the rest of the process, we would revert this virtual swap and try other possibilities:

Update the character frequency for wordl for char in word1: 8 count1[ord(char) - ord('a')] += 1 9

for char in word2:

Python Solution

combinations.

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1 class Solution:
      def isItPossible(self, word1: str, word2: str) -> bool:
3
          # Count the frequency of each character in both words
          count1 = [0] * 26
          count2 = [0] * 26
```

Try swapping frequencies and check if both words can have the same character set

Decrement and increment the frequencies at position i and j respectively

If both words have the same number of distinct characters, return True

Calculate the number of distinct characters current in each word

Only proceed with the swap if both frequencies are non-zero

count1[i], count2[j] = count1[i] - 1, count2[j] - 1

count1[j], count2[i] = count1[j] + 1, count2[i] + 1

Revert the changes if the above condition is not met

count1[i], count2[j] = count1[i] + 1, count2[j] + 1

count1[j], count2[i] = count1[j] - 1, count2[i] - 1

// Create arrays to count the frequency of each character in both strings

distinct_in_word1 = sum(v > 0 for v in count1)

distinct_in_word2 = sum(v > 0 for v in count2)

if distinct_in_word1 == distinct_in_word2:

36 # If no swaps can result in both words having the same character set, return False return False 37 38

Java Solution

1 class Solution {

class Solution {

int count1[26] = {0};

int count2[26] = {0};

for (char c : word1) {

for (char c : word2) {

++count1[c - 'a'];

++count2[c - 'a'];

for (int i = 0; i < 26; ++i) {

public:

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                 countWord2[word2.charAt(i) - 'a']++;
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 17
             // Iterate over all pairs of characters
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             for (int i = 0; i < 26; ++i) {
 19
                 for (int j = 0; j < 26; ++j) {
 20
                     // If both characters are present in their respective words
 21
                     if (countWord1[i] > 0 && countWord2[j] > 0) {
 22
                         // Simulate swapping the characters
 23
                         countWord1[i]--;
 24
                         countWord2[j]--;
 25
                         countWord1[j]++;
 26
                         countWord2[i]++;
 27
 28
                         // Check if the frequency distribution matches after the swap
 29
                         int delta = 0; // Delta will store the net difference in frequencies
 30
                         for (int k = 0; k < 26; ++k) {
 31
                             if (countWord1[k] > 0) {
 32
                                 delta++;
 33
 34
                             if (countWord2[k] > 0) {
 35
                                 delta--;
 36
 37
 38
                         // If delta is zero, it means that the frequency distribution matches
 39
 40
                         if (delta == 0) {
 41
                             return true;
 42
 43
 44
                         // Undo the swap operation as it did not lead to a match
 45
                         countWord1[i]++;
 46
                         countWord2[j]++;
 47
                         countWord1[j]--;
 48
                         countWord2[i]--;
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 53
             // If no swaps resulted in a match, return false
 54
             return false;
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 56 }
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C++ Solution
    #include <string>
```

56 57 58 // If no swap can make the words identical, return false 59 return false; 60 61 };

```
Typescript Solution
     function isItPossible(word1: string, word2: string): boolean {
         // Arrays to store letter frequencies for both words
         let count1: number[] = new Array(26).fill(0);
         let count2: number[] = new Array(26).fill(0);
  6
         // Populate frequency arrays for wordl
         for (let c of word1) {
             count1[c.charCodeAt(0) - 'a'.charCodeAt(0)]++;
  8
  9
 10
 11
         // Populate frequency arrays for word2
         for (let c of word2) {
 12
             count2[c.charCodeAt(0) - 'a'.charCodeAt(0)]++;
 13
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 16
         // Iterate over each letter in the alphabet
 17
         for (let i = 0; i < 26; ++i) {
 18
             // Iterate over each letter in the alphabet
 19
             for (let j = 0; j < 26; ++j) {
                 // Check if current letters are present in both words
 20
                 if (count1[i] > 0 && count2[j] > 0) {
 21
 22
                     // Simulate swapping the letters by updating counts
 23
                     count1[i]--;
 24
                     count2[j]--;
 25
                     count1[j]++;
 26
                     count2[i]++;
 27
 28
                     // Variable to track the sum of differences in frequencies
 29
                     let difference = 0;
 30
                     // Check if the frequencies match for each letter
 31
                     for (let k = 0; k < 26; ++k) {
 32
                         if (count1[k] > 0) {
 33
                             difference++;
 34
 35
                         if (count2[k] > 0) {
 36
                             difference--;
 37
 38
 39
 40
                     // If the sum of differences is zero, words can be made identical
                     if (difference === 0) {
 41
 42
                         return true;
 43
 45
                     // Undo the simulated swap
 46
                     count1[i]++;
                     count2[j]++;
 47
 48
                     count1[j]--;
 49
                     count2[i]--;
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 54
         // If no swap can make the words identical, return false
 55
         return false;
 56
 57
```

The provided code iterates over each character in word1 and word2, counting the frequency of every character. Then it has nested loops where it iterates over the counts of characters in cnt1 and cnt2 arrays (size 26 for each letter of the alphabet) and performs operations to check if swapping characters could make the frequency of non-zero characters in the count arrays equal.

Time and Space Complexity

0(26^2) or simply 0(1) since 26 is a constant and does not change with the input size. Combining these, the overall time complexity is primarily affected by the character counting, so 0(n) for word1 plus 0(m) for word2, where n and m are the lengths of word1 and word2 respectively. Since the 26*26 operations are constant time and do not scale with n

Time Complexity

The total time complexity is O(n + m), where n is the length of word1 and m is the length of word2. Space Complexity

The time complexity of counting characters is O(n) for each word, where n is the length of the word. However, the nested loops

or m, the insignificant additional constant time doesn't affect the overall complexity.

create a bigger time complexity issue. There are 26 possible characters, leading to 26*26 comparisons in the worst case which, is

Putting it all together, the space complexity is 0(1) because it only requires fixed space for the frequency counts and a few variables for iteration and comparison, which does not scale with the input size.

The space complexity is much simpler to analyze. The space required by the algorithm is the space for the two count arrays cnt1

and cnt2 which hold the frequency of each character. As these arrays have a fixed size of 26, regardless of the input size, the space

complexity is 0(1).