1460. Make Two Arrays Equal by Reversing Subarrays

Easy Array Hash Table Sorting

Problem Description

In this problem, you are provided with two integer arrays target and arr. Both arrays have the same number of elements. Your goal is to determine if it's possible to make the array arr identical to the array target by performing a series of operations. For each operation, you can select any non-empty subarray from arr and reverse it. A subarray is a contiguous part of an array. You can reverse subarrays as many times as you need. You need to return true if arr can be made equal to target or false otherwise.

Intuition

The key insight to solve this problem is to understand that reversing a subarray does not change the overall content (numbers and their frequencies) of the array; it only changes the order of elements. If two arrays contain the same elements with the same frequencies (multisets are equal), it is always possible to make one array equal to another by reversing subarrays because you can always rearrange the elements to match. Therefore, the solution does not actually involve performing the subarray reversals; instead, it involves checking whether both arrays contain the same set of elements with the same frequency.

The approach to arrive at the solution is straightforward:

1. Sort both the target and arr arrays. Sorting brings elements of the same value next to each other and thus makes it easy to compare the

- arrays.

 2. After sorting, if both arrays are equal, it means that arr can be transformed into target through reversing subarrays. Hence, return true.
- 3. If the arrays do not match after sorting, it means arr cannot be made equal to target, so return false.

The implementation of the solution is quite simple and uses basic algorithms and data structures. Here is a detailed explanation

Solution Approach

of the approach:

1. **Sorting Algorithm:** The core of the solution uses a sorting algorithm. Both Python's sort() method on lists uses TimSort,

- which is a hybrid sorting algorithm derived from merge sort and insertion sort. It is a stable, adaptive, and iterative merge sort that requires fewer than n log(n) comparisons when running on partially sorted arrays, which makes it very efficient.

 2. **Comparison:** After <u>sorting</u>, the elements in both <u>target</u> and <u>arr</u> are in the same order if they are comprised of the same set of elements. The algorithm then simply compares the two sorted arrays to check for equality. This is done using the '=='
- 3. **Returning the Result:** If the comparison evaluates to true, it means that arr can be rearranged to match target by reversing subarrays; the function thus returns true. If the comparison is false, there are elements in arr that do not match those in target, indicating that no series of reversals will make the two arrays equal. In this case, the function returns false.

The solution does not require any additional data structures; it works with the input arrays themselves and returns a boolean

value. It is also important to note that since <u>sorting</u> changes the original arrays, if maintaining the original order is needed for any reason, one could sort copies of the arrays instead.

Here is the final solution encapsulated in a class, as provided in the reference code:

target.sort() arr.sort() return target == arr

def canBeEqual(self, target: List[int], arr: List[int]) -> bool:

operator in Python, which compares corresponding elements in both lists.

```
The above solution is concise as well as efficient due to the use of <u>sorting</u>, which is more time and space-efficient than other methods of comparison that could involve using hash maps or multiset data structures to compare frequencies of elements.
```

Let us consider a small example to illustrate the solution approach.

Suppose the target array is [1, 2, 3, 4] and the arr array is [2, 4, 1, 3]. We want to find out if we can make arr identical to

target by reversing subarrays.

Following the solution approach:

class Solution:

After sorting, target remains [1, 2, 3, 4] because it was already sorted. The arr array after sorting becomes [1, 2, 3, 4].

Sorting the arrays: We start by sorting both target and arr.

from typing import List # Import List from typing module for type annotations

target (List[int]): The target list that arr should match.

arr (List[int]): The list to compare with the target list.

return Arrays.equals(target, arr);

#include <vector> // Include vector header for using vectors

function canBeEqual(target: number[], arr: number[]): boolean {

// Determine the length of the 'target' array.

const occurrenceCount = new Array(1001).fill(0);

occurrenceCount[target[index]]++;

// Iterate over each element of 'target' and 'arr'.

for (let index = 0; index < arrayLength; index++) {</pre>

// Check if every value in our counting array is zero.

// Increment the count for the current element in 'target'.

After sorting, if target is equal to arr, it means arr can be made equal

to target by sorting. Otherwise, it's not possible.

const arrayLength = target.length;

Check if two lists, target and arr, can be made equal through sorting.

Comparing the sorted arrays: Now, we compare the sorted target array with the sorted arr array.

Since sorted_target == [1, 2, 3, 4] and sorted_arr == [1, 2, 3, 4], comparison shows that both arrays are identical.

Result: As the sorted arrays are identical, we can conclude that it is possible to make arr identical to target by reversing

- subarrays.

 Hence, according to the solution approach outlined in the problem content, the function would return true for these arrays.
- Solution Implementation

Python

class Solution: def canBeEqual(self, target: List[int], arr: List[int]) -> bool:

Args:

```
Returns:
       bool: True if arr can be made equal to target by sorting, False otherwise.
       # Sort both the target list and arr list in place
        target.sort()
       arr.sort()
       # After sorting, if target is equal to arr, it means arr can be made equal
       # to target by sorting. Otherwise, it's not possible.
        return target == arr
Java
class Solution {
    public boolean canBeEqual(int[] target, int[] arr) {
       // Sort the target array in-place
       Arrays.sort(target);
       // Sort the arr array in-place
       Arrays.sort(arr);
       // Check if the sorted arrays are equal
       // The equals method checks if the two arrays have the same elements in the same order
```

// Initialize an array for counting occurrences with fixed size 1001, filled with zeros.

// This is based on the constraint that the elements in 'target' and 'arr' are integers between 1 and 1000.

```
// Decrement the count for the current element in 'arr'.
occurrenceCount[arr[index]]--;
}
```

C++

The given code consists of two sort operations and one equality check operation. The sort operations on both target and arr are the dominant factors in the time complexity of this function.

Time Complexity

return target == arr

Time and Space Complexity

Assuming that the sort function is based on an algorithm with O(n log n) time complexity such as Timsort (which is the sorting

algorithm used by Python's built-in sort method), the time complexity for sorting both lists would be $O(n \log n)$, where n is the length of the lists.

The time complexity would be the sum of sorting the two lists:

Since the question implies that both lists should be of the same length for them to be possibly equal, we can assume

target.length == arr.length and use n as the length for both:

Total time for sorting: 2 * 0(n log n)

First sort: O(target.length log(target.length))

Second sort: 0(arr.length log(arr.length))

The equality check operation that follows (target == arr) compares each element between the two lists, which has a time complexity of O(n).

However, since the time for sorting $(0(n \log n))$ is greater than the time for comparison (0(n)), the overall time complexity of the

function is dominated by the sorting time:Overall time complexity: O(n log n)

Space Complexity

Considering the space complexity, the sort operations are done in-place in Python, which means that no additional space

proportional to the input size is required beyond a constant amount used by the sorting algorithm itself.

Thus, the space complexity of the function is:

• Space complexity: 0(1) (constant space complexity, not counting the input and output)