2507. Smallest Value After Replacing With Sum of Prime Factors **Leetcode Link** 

Medium Math Number Theory

2 \* 3 \* 3). So the sum of its prime factors is 2 + 3 + 3 = 8.

**Problem Description** The problem involves manipulating a given positive integer n by continuously replacing it with the sum of its prime factors. A key

detail is that if a prime factor is repeated in the factorization (that is, n is divisible by a prime number multiple times), that prime factor must be counted multiple times in the sum. For example, if our given number n is 18, its prime factors are 2 and 3, but 3 is counted twice because 18 is divisible by 3 twice (18 =

The goal is to repeat this process: take the sum of the prime factors of n and then replace n with this sum, and continue this until n no longer changes—it reaches its smallest possible value.

The challenge is to write an algorithm that performs this computation efficiently and find when the number n stops changing, which is the result to be returned.

Intuition

The solution makes use of a loop that continuously replaces the number n with the sum of its prime factors until n cannot be reduced

of n.

and should be included in the sum).

any further.

We start with a loop that runs indefinitely, checking if we can find prime factors of n and their sums. In each iteration, we initialize three variables - t to keep track of the original value of n at the start of the iteration, s to calculate the sum of prime factors, and i to iterate over potential prime factors starting from 2.

The inner while loop checks if i is a divisor of n, and if so, it keeps dividing n by i and adds i to the sum s until n is no longer divisible by i. The variable i is incremented to check for the next possible prime factor. Once we have tried all possible divisors up to n // i, we check if n itself is greater than 1 (which would mean n is a prime number

At this point, if the calculated sum of prime factors s is equal to the temporary variable t, it means n can no longer be reduced any further, and we return t.

The key intuition behind the algorithm is that the prime factorization of a number can help find a smaller representation of the number itself and by repeating this process and summing these factors repeatedly, we can converge to the smallest possible value

Solution Approach

The solution makes use of a basic factorization algorithm and control flow to reduce the number n to its smallest possible value by continuously summing its prime factors.

If s is not equal to t, we replace n with s to continue the process with a new reduced value of n.

2. Inside the loop, a temporary variable t is assigned the current value of n to keep track of its value through each iteration. This is important to identify when there is no further change possible. 3. Another variable s is initialized to 0; this will be used to calculate the sum of the prime factors of n.

5. The first inner while loop runs as long as i is less than or equal to n // i (since a factor larger than the square root of n would

4. An index i is set to 2, which is the first prime number. This variable is used to test potential factors of n.

continues until n is no longer divisible by i, accommodating for all instances of i as a factor.

have already been identified by its corresponding smaller factor except when n is prime). 6. A nested inner while loop checks if i divides n perfectly. If it does, n is divided by i, and i is added to the sum s. This loop

7. The index i is then incremented to check the next potential factor.

number to its smallest form based on prime factor accumulation.

Step 2: Assign the value of n to a temporary variable t, so t = 12.

Step 3: Initialize a variable s to 0. This will hold the sum of the prime factors of n.

Step 5: As  $i \ll n // i$  (since 2  $\ll 12 // 2$ ), we proceed with factorization.

Step 10: Because s is not equal to t, we set n to s; hence n is now 7.

Now n has changed from 12 to 7. We repeat the entire process again with n = 7.

Step 8: Since n is still greater than 1, and 7 is a prime number, we add 7 to s. Now s = 7.

# Check for factors of the number starting with the smallest prime factor

# Set num to sum\_of\_factors for the next iteration to check the new number

// Method to find the smallest value according to specified conditions

// Loop indefinitely until we find the smallest value

n /= i; // Divide n by the factor

// If there is a remaining factor greater than 1, add it to the sum

// If it matches, return the sum (as it is the smallest value)

// If the sum of factors is the same as the original value, return it

const originalValue: number = n; // Preserve the original value of 'n'

n /= i; // Divide 'n' by the factor for further factorization

// Check for a remaining prime factor greater than sqrt(n)

// Otherwise, set 'n' to the calculated sum of factors for the next iteration

let sumOfFactors: number = 1; // Initialize the sum of factors to 1, since 1 is a factor of all numbers

sumOfFactors += n / i; // Add the complement factor to the sum if it's different

// Check if the sum of the factors equals the original number

// Store the original value of n

// Initialize the sum of the factors

Step 1: Start with n = 12 and enter the while loop.

Step 7: Increment i to the next integer, which is 3.

Step 1: n is 7, so we enter the while loop.

Here's a step-by-step breakdown of the approach used in the solution:

1. The solution uses a while loop that will run until n no longer changes.

8. After all possible factors up to n // i have been tested, a final check outside the inner loops evaluates if n itself is greater than 1, implying n is a prime number. If it is prime, it is added to the sum s.

9. At the end of the iteration, the algorithm checks if the sum of the prime factors s is equal to the starting value t. If so, n has

reached its smallest value as no primes other than itself can be extracted and summed, and the algorithm returns the value of t.

10. If s is not equal to t, the algorithm replaces the value of n with s to repeat the factorization process on this new reduced number.

indivisible and signaling the end of the reduction process. This algorithm does not employ any complex data structures and follows a straightforward but effective pattern to reduce the

The algorithm essentially terminates when a number only comprises its prime self or when it's reduced to 1, both of which are

Let's walk through the solution approach with a small example where our given number n is 12. The goal is to continuously replace n with the sum of its prime factors and repeat this process until n is no longer reduced.

Step 4: Start with i = 2, which is the smallest prime factor.

Step 6: The inner loop checks if 12 is divisible by 2. It is, so we divide n by 2 to get 6 and add 2 to s. Now, s = 2 and n = 6.

We continue with the inner loop since 6 is still divisible by 2. We divide 6 by 2 to get 3 and add 2 to s. Now, s = 4 and n = 3.

## Step 8: Since 3 is a prime number larger than n // i but still divides n perfectly, we add 3 to s. Now, s = 4 + 3 giving us s = 7 and n

becomes 1.

Step 2: Set t to 7.

Step 3: Initialize s to 0.

Step 4: Start with i = 2.

**Python Solution** 

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Example Walkthrough

Step 9: Since n is now 1, we exit the inner loop. We check if s is equal to the temporary variable t. As s = 7 and t = 12, they are not equal.

Step 5: Begin inner loop. Because no i exists such that  $i \ll n // i$  and i divides 7, skip to step 8.

Step 9: Compare the sum of the prime factors s with t. They are equal (s = t = 7), indicating that we cannot reduce n any further.

The algorithm would then terminate and return the value 7 as the result, as n can no longer be reduced by the process described.

class Solution: def smallestValue(self, num: int) -> int: # Continue the loop until we find the smallest value while True: # Initialize temp variable to store original number, sum\_of\_factors, and start divisor from 2

temp, sum\_of\_factors, divisor = num, 0, 2

while divisor <= num // divisor:</pre>

sum\_of\_factors += num

if sum\_of\_factors == temp:

divisor += 1

return temp

num = sum\_of\_factors

public int smallestValue(int n) {

if (n > 1) {

int originalValue = n;

sumOfFactors += n;

if (sumOfFactors == originalValue) {

if (sumOfFactors == originalValue) {

1 // Function to find the smallest integer value with the same

2 // sum of factors (including 1 and the number itself) as 'n'

// Factorize and sum up the factors

if (n > 1 && n < originalValue) {</pre>

// Loop indefinitely until we find the smallest value

return sumOfFactors;

n = sumOfFactors;

function smallestValue(n: number): number {

while (true) {

if num > 1:

# If the divisor is a factor, divide num by the divisor and add to the sum\_of\_factors while num % divisor == 0: num //= divisor sum\_of\_factors += divisor 13 # Move to the next potential factor

# If there's a remaining number greater than one, it's a prime factor; add it to the sum\_of\_factors

# If the sum of factors is equal to the original number, we found the smallest value

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int sumOfFactors = 0;
               // Start dividing the number from 2 onwards to find its factors
                for (int i = 2; i <= originalValue / i; ++i) {</pre>
                    // Divide by i as long as it is a factor of n
                    while (n % i == 0) {
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                        sumOfFactors += i; // Add factor to the sum
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```

Java Solution

class Solution {

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                   return sumOfFactors;
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               // If it does not match, set n to the sumOfFactors for another iteration
               n = sumOfFactors;
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31 }
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C++ Solution
1 class Solution {
  public:
       // Function to find the smallest integer value with the same
       // sum of factors (including 1 and the number itself) as 'n'
       int smallestValue(int n) {
           // Loop indefinitely until we find the smallest value
           while (true) {
               int originalValue = n; // Preserve the original value of 'n'
               int sumOfFactors = 0; // Initialize the sum of factors to 0
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               // Factorize and sum up the factors
               for (int i = 2; i \le n / i; ++i) { // Only need to check up to sqrt(n)
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                   // While 'n' is divisible by 'i'
                   while (n % i == 0) {
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                       sumOfFactors += i; // Add the factor to the sum
                       n /= i; // Divide 'n' by the factor for further factorization
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               // If there is a remaining prime factor greater than sqrt(n), add it to the sum
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               if (n > 1) sumOfFactors += n;
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### for (let i = 2; $i \le Math.sqrt(n)$ ; ++i) { // Only need to check up to sqrt(n)// While 'n' is divisible by 'i' while (n % i === 0) { sumOfFactors += i; // Add the factor to the sum **if** (i != n / i) { 14

Typescript Solution

while (true) {

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sumOfFactors += n; // Add it to the sum only if it's different from the current value
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           // If the sum of factors is the same as the original value, return the smallest integer
           if (sumOfFactors == originalValue) {
               return originalValue;
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           n = sumOfFactors; // Set 'n' to the calculated sum of factors for the next iteration
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33 }
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Time and Space Complexity
Time Complexity
The time complexity of this code is governed by two nested loops. The outer loop runs indefinitely until a number is found where the
sum of its prime factors is equal to itself. The inner loop, on the other hand, is used for factorization and runs at most sqrt(n) times
because it checks for factors from 2 to at most sqrt(n).
```

## • In the worst-case scenario for a given n, we may have to factorize n, then s, and so on if n is initially not a prime or semiprime (a product of exactly two primes). Each new s will be smaller than n as we are adding up the factors. Thus, the inner factorization

The factorization loop:

process will take O(sqrt(m)) time for each number m we factorize, where m starts from n and gets smaller. The outer while loop: The code does not have a clear stopping condition within predictable bounds, due to the uncertain nature of the sum s

approaching the target condition s == t. However, for every iteration, the sum of the prime factors of n (s) gets closer to being a

The number of iterations can be considered proportional to the number of distinct prime factors of n in a sense but is not easy to

prime number itself. Once s equals a prime number, it will become equal to t in the next iteration, and the function will return t.

Considering the worst-case scenario, where n is a large composite number with many small prime factors, the time complexity can be high, but the exact upper bound is intricate to determine without more constraints on n.

# The space complexity is 0(1). There are only a few integer variables (t, s, i, and n) being used, which do not depend on the input

**Space Complexity** 

express in standard 0() notation.

size n, unless considering the size n itself needs. The space used by these variables is constant and does not grow with n.