1585. Check If String Is Transformable With Substring Sort Operations String] Greedy **Sorting** Hard

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Problem Description
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The goal is to determine if it's possible to convert string s into string t by performing several operations. In each operation, you can select a non-empty substring from s and rearrange its characters in ascending order. A substring is defined as a contiguous block of characters within the string. The operation can be applied any number of times to any substring within s. The problem

asks to return true if s can be transformed into t using these operations, otherwise return false. Intuition

To solve the problem, the solution uses the intuition that each character in string s can only move towards the left in each operation. This is due to the fact that sorting a substring in an ascending order will not move any character to a position greater than its original position.

Given that constraint, we check if we can transform s into t by considering the characters in t from left to right. We make use of a data structure pos, which is a dictionary of deques, where the keys are the digits (characters represented as integers) found in s, and the values are deques holding the indices of these digits. As we iterate through each character in t, we check if the character can be moved to its correct position in s. This involves

checking if there is an occurrence of the character left in s (pos[x] is not empty), and ensuring that there are no smaller characters (digits) to the left of the character's current position in s (any(pos[i] and pos[i][0] < pos[x][0] for i in range(x)) returns false). If these conditions are not met, it means that it is not possible to transform s into t and we return false. Otherwise, we "move" the character in s to its next position (which is simulated by popping from the left of the deque

corresponding to that character) and continue with the next character in t.

true. **Solution Approach** The implementation of the solution takes a strategic approach using the following algorithmic concepts and data structures:

Deque Data Structure: The solution uses a defaultdict to create a mapping of deques which are essentially double-ended

Index Tracking: Each character's position in s is tracked by storing its indices in order in the corresponding deque. This helps

queues to store the indices of corresponding characters. The deque allows efficient removal of elements from both ends

If we successfully go through the entire string t without returning false, it means the transformation is possible and we return

to determine the possibility of moving a character to a certain position. **Greedy Approach:** A greedy approach is applied by attempting to match characters from t with characters in s from left to

right. This checks the feasibility of reaching the desired pattern.

which is essential since we are trying to simulate the movement of characters within the string.

- Initialize pos: A defaultdict(deque) called pos is created to store the indices of characters in s. **Store Indices**: Iterating through s, the indices are stored in pos. For example, if s = '3421' then pos[1] will be deque([3]), pos [2] will be deque([2]), and so on.
- **Transformation Check**: The algorithm runs through each character c in t and performs the following: ∘ It converts c to an integer x. • It checks if pos[x] is not empty, which means there's still an occurrence of that character in s.

• It uses a combination of list comprehension and the any function to determine if there's any smaller character to the left of the current one

in s. This is done with any (pos[i] and pos[i][0] < pos[x][0] for i in range(x)). If true, then it's not possible to move x to the desired position without violating the sorting constraint (since a smaller number cannot be passed). Simulate Character Movement: If the checks pass, the character's index is moved (deque is popped from the left) as if the

return true. If at any point an inconsistency arises, the function returns false immediately.

Transformation Check: Now, we check if we can transform s into t character by character.

• We then "move" the '1' by popping from pos [1], simulating the sorting operation. pos now looks like this:

Return the Result: If no inconsistencies are found during the whole iteration, the transformation is possible, and therefore we

it's possible to convert s into t using the defined operations.

To understand the step-by-step algorithm:

only allows characters in s to move leftward (smaller index) and not rightward. **Example Walkthrough**

character has been sorted to its place in s. This is effectively simulating the operation defined in the problem.

Following the solution steps: Initialize pos: We create a defaultdict(deque) to hold the indices of characters from string s. The pos will look like this after

Let's illustrate the solution approach with a small example. Suppose we have s = "4321" and t = "1234". We want to find out if

The algorithm's correctness hinges on the property that you can only sort substrings in ascending order, which consequently

Store Indices: As we've initialized pos, each character from s has its index stored in the corresponding deque.

○ We start with t[0] which is '1', convert it into an integer x = 1, and see if pos[x] is not empty (which means '1' is still in s). pos[1] is

• We check if there's a smaller character to the left of the '1' in s using the any statement. Since '1' is the smallest character, no smaller character can be to the left, so we pass this check.

pos[4] = deque([0])

pos[3] = deque([1])

pos[2] = deque([2])

pos[4] = deque([0])

pos[3] = deque([1])

pos[1] = deque([])

transformation.

from collections import defaultdict

for digit in t:

def isTransformable(self, s: str, t: str) -> bool:

digit_positions[int(digit)].append(index)

public boolean isTransformable(String s, String t) {

Deque<Integer>[] positions = new Deque[10];

for (int i = 0; i < s.length(); i++) {</pre>

return false;

Arrays.setAll(positions, k -> new ArrayDeque<>());

if (positions[targetDigit].isEmpty()) {

for (int j = 0; j < targetDigit; j++) {</pre>

// If all checks pass, s is transformable into t

positions[targetDigit].poll();

// Import the Queue implementation if you don't have one.

function isTransformable(src: string, target: string): boolean {

import { Queue } from 'some-queue-library';

// Initialize queues for each digit.

for (let i = 0; i < src.length; i++) {</pre>

const digit = parseInt(src[i], 10);

if (digitPositions[digit].isEmpty()) {

for (let j = 0; j < digit; ++j) {

def isTransformable(self, s: str, t: str) -> bool:

digit_positions[int(digit)].append(index)

digit_positions[digit_value].popleft()

digit_positions = defaultdict(deque)

for index, digit in enumerate(s):

Iterate over each digit in t

return False

Time and Space Complexity

digit_value = int(digit)

return false;

digitPositions[digit].enqueue(i);

for (let i = 0; i < 10; i++) {

for (const char of target) {

return false;

const digitPositions: Queue<number>[] = [];

digitPositions[i] = new Queue<number>();

// This is just a placeholder, as TypeScript does not have a built-in queue.

// Array to store the positions of digits 0-9 in the string 'src'.

// Fill the queues with the positions of each digit in 'src'.

// Iterate through each character in the target string 'target'.

const digit = parseInt(char, 10); // Convert character to digit

// If there is no such digit in 'src', transformation is impossible.

// Check if any smaller digit is positioned before this digit in 'src'.

// If there is a smaller digit in front of this digit, it's not transformable.

// If all characters can be transformed without violating the constraints, return true.

Initialize a dictionary to store queues of indices for each digit in s

Fill the dictionary with the positions of each digit in s

if (!digitPositions[j].isEmpty() && digitPositions[j].front() < digitPositions[digit].front()) {</pre>

If there are no positions left for the current digit or if there is any digit with a smaller value

if not digit_positions[digit_value] or any(digit_positions[smaller_digit] and digit_positions[smaller_digit][0] < digit_positions[smaller_

in front of the current digit's earliest position, the transformation is not possible

If the current digit can be transformed, remove its earliest occurrence from the queue

To determine the time complexity, let us analyze the various operations being performed in the function.

If all digits in t can be reached by transforming s successfully, return True

// You can replace this with your own Queue class or any library implementation.

// Function to check if the string 'src' can be transformed into the string 'target'.

// Fill each queue with indices of digits in string s

// Initialize an array of queues representing each digit from 0 to 9

int digit = s.charAt(i) - '0'; // Get digit at char position i

return false; // s cannot be transformed into t

// If position is valid, remove it as it is 'used' for transformation

positions[digit].offer(i); // Add index to corresponding digit queue

// If there are no positions left for this digit, s cannot be transformed into t

// If there is a smaller digit and it comes before the current digit in source s

if (!positions[j].isEmpty() && positions[j].peek() < positions[targetDigit].peek()) {</pre>

// Check if any smaller digit appears after the current digit in the source

digit_positions = defaultdict(deque)

for index, digit in enumerate(s):

Iterate over each digit in t

return False

digit_value = int(digit)

Python

Java

class Solution {

class Solution:

we've finished:

pos[4] = deque([0])

pos[3] = deque([1])

pos[2] = deque([2])

pos[1] = deque([3])

deque([3]), so it's not empty.

pos[1] = deque([]) // '1' has been moved Repeat for the Next Character: We repeat the same check for the next character in t, which is '2'.

• With the any statement, we check for smaller characters. Since '1' has already been moved, and there is no '0' in s, the check passes.

Return the Result: Since we were able to move all characters from s to match the sorted order of t without encountering any

Continue With Remaining Characters: We continue this process for '3' and '4'. Each character passes the condition checks and gets moved.

Initialize a dictionary to store queues of indices for each digit in s

Fill the dictionary with the positions of each digit in s

x becomes 2 and pos [2] is deque([2]), so we have a '2' to work with.

• We "move" the '2' by popping from pos [2], and pos now looks like:

pos[2] = deque([]) // '2' has been moved

Solution Implementation

blocking condition, we conclude that it's possible to transform s into t. The function would return true.

This walkthrough demonstrates that given these operations and their constraints, we can determine the possibility of

transforming one string into another algorithmically, using deques for index tracking and a series of checks to simulate the

If the current digit can be transformed, remove its earliest occurrence from the queue digit_positions[digit_value].popleft() # If all digits in t can be reached by transforming s successfully, return True return True

If there are no positions left for the current digit or if there is any digit with a smaller value

if not digit_positions[digit_value] or any(digit_positions[smaller_digit] and digit_positions[smaller_digit][0] < dig

in front of the current digit's earliest position, the transformation is not possible

// Process the target string to see if it is transformable from source string for (int i = 0; i < t.length(); i++) {</pre> int targetDigit = t.charAt(i) - '0'; // Target digit to search for

return true;

```
C++
class Solution {
public:
   // Function to check if the string 's' can be transformed into the string 't'.
    bool isTransformable(string s, string t) {
       // Queue to store the positions of digits 0-9 in the string 's'.
       queue<int> digit_positions[10];
       // Fill the queues with the positions of each digit in 's'.
        for (int i = 0; i < s.size(); ++i) {
            digit_positions[s[i] - '0'].push(i);
       // Iterate through each character in the target string 't'.
        for (char& c : t) {
            int digit = c - '0'; // Convert character to digit
           // If there is no such digit in 's', transformation is impossible.
            if (digit_positions[digit].empty()) {
                return false;
            // Check if any smaller digit is positioned before our digit in 's'.
            for (int j = 0; j < digit; ++j) {
                // If there is a smaller digit in front of our digit, it's not transformable.
                if (!digit_positions[j].empty() && digit_positions[j].front() < digit_positions[digit].front()) {</pre>
                    return false;
            // Since this digit can be safely transformed, we pop it from the queue.
            digit_positions[digit].pop();
       // If all characters can be transformed without violating the constraints, return true.
       return true;
};
TypeScript
```

// Since this digit can be safely transformed, we dequeue it from the corresponding queue. digitPositions[digit].dequeue();

class Solution:

return true;

from collections import defaultdict

for digit in t:

return True

once.

Space Complexity:

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The given Python function isTransformable checks if it's possible to transform the string s into the string t by repeatedly moving
  a digit to the leftmost position if there are no smaller digits to its left.
Time Complexity:
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Initializing pos: Building the pos dictionary takes O(n) time, where n is the length of s, as we iterate over all characters in s

Checking Transformability: We iterate over each character in t, and for each character, perform a check that could iterate

over a maximum of 10 possible values (since the digits range from 0 to 9) - this is the any() function call. The inner check pos[i] and pos[i][0] < pos[x][0] is 0(1) for every digit i for each character of t, because it's merely

- indexing and comparison. Therefore, the total time for transforming, in the worst case, would be 0(10 * n) which simplifies to 0(n), since we don't count constants in Big O notation.
- **Popping from pos:** The popleft() operation is 0(1) for a deque. Therefore, the overall time complexity of the function is O(n).
- Space for pos: The pos dictionary stores deque objects for each unique digit found in s, and each deque can grow to the size of the number of occurrences of that digit. The total size of all deques combined will not exceed n. Therefore, the space

complexity contributed by pos is O(n).

Thus, the total space complexity is also O(n). In conclusion, the time complexity is O(n) and the space complexity is O(n).

Miscellaneous Space: Constant additional space used by iterators and temporary variables.