2410. Maximum Matching of Players With Trainers

Sorting

Problem Description

Greedy Array Two Pointers

represents the ability of a particular player, and each entry in the trainers array represents the training capacity of a particular trainer. A player can be matched with a trainer if the player's ability is less than or equal to the trainer's capacity. It's important to note that each player can only be matched with one trainer, and each trainer can match with only one player. The task is to find the maximum number of such player-trainer pairings where the conditions of matching ability to capacity are met.

In this LeetCode problem, you are given two 0-indexed integer arrays: players and trainers. Each entry in the players array

Intuition

The intuition behind the solution is straightforward: we want to pair as many players with trainers as we can, prioritizing the pairing of less able players with trainers who have just enough capacity to train them. This ensures that we do not 'waste' a

Medium

To achieve this, we sort both the players and trainers arrays. Sorting helps us easily compare the least able player with the least capable trainer and move up their respective arrays. If a match is found, both the player and the trainer are effectively

removed from the potential pool by moving to the next elements in the arrays (incrementing the index). If no match is found, we

move up the trainers array to find the next trainer with a higher capacity that might match the player's ability.

trainer with high capacity on a player with low ability when that trainer could be matched with a more able player instead.

The result is incrementally built by adding a match whenever we find a capable trainer for a player. We stop when we've either paired all players or run out of trainers. The sum of matches made gives us the maximum number of possible pairings. Solution Approach

The implementation provided uses a simple but effective greedy approach, heavily relying on sorting. The steps involve:

The algorithm works with two pointers, one (i) iterating through the players array, and the other (j) through the trainers array.

Sorting the players array in ascending order, which ensures we start with the player with the lowest ability.

Sorting the trainers array in ascending order, which allows us to start with the trainer having the lowest capacity.

- Here is the implementation broken down:
- Initialize the trainer index pointer j to 0.

Iterate over each player p in the sorted players list using a for-loop.

Within the loop, proceed with an inner while-loop which continues as long as j is less than the length of trainers and

the current trainer's capacity trainers[j] is less than the ability of the player p. The purpose of this loop is to find the

Greedy approach: The algorithm uses a greedy method by matching each player with the "smallest" available trainer that can

Sorting: By sorting the arrays, the solution leverages the property that once a player cannot be paired with a current trainer,

Algorithm and Data Structures:

train them.

process.

Example Walkthrough

• players: [2, 3, 4]

trainers: [1, 2, 5]

After sorting,

Following the solution approach:

After the while-loop, if j is still within the bounds of the trainers array, it means a suitable trainer has been found for

first trainer in the sorted list whose capacity is sufficient to train the player.

Initialize a variable ans to 0 to keep track of the number of matches made.

- the player p. We increment ans to count the match and also increment j to move to the next trainer. Once all players have been considered, return the value of ans. •
- Two pointers: It uses two pointers to pair players with trainers without revisiting or re-comparing them, thus optimizing the

they cannot be paired with any trainers before that one in the sorted list either.

Lists: The main data structure used here are lists (players and trainers), which are sorted.

- Patterns used: • Sorting and Two-pointers: This is a common pattern for efficiently pairing elements from two different sorted lists.
- Let's say we have the following players and trainers arrays:

We first sort both the players and trainers arrays:

players: [2, 3, 4] (already sorted) trainers: [1, 2, 5] (already sorted)

We initialize ans to 0. This variable will keep track of the successful player-trainer matches.

 We check the trainers in order: trainer 1 cannot train the player because the trainer's capacity is too low. Move to the next trainer (2), and we find a match. Trainer 2 can train the player with ability 2.

successfully.

Python

class Solution:

players.sort()

trainers.sort()

for player in players:

 There are no more trainers to compare since we exceeded the length of the trainers array. At the end of this process, the value of ans is 2, which signifies that we were able to match two pairs of players and trainers

We also initialize the trainer index pointer j to 0.

Now, let's go through each player and try to find a trainer match:

We increment ans to 1 and move to the next trainer (j becomes 2).

• The trainer in position 2 in the sorted trainers list has a capacity of 5, which is sufficient.

• We increment ans to 2, and since there are no more trainers, we move on to the final player.

Player with ability 2 is the first in the players array.

Player with ability 3 is the next.

We match this player with the trainer.

Player with ability 4:

Solution Implementation

from typing import List # Import List from typing module for type hints

Sort the list of players in ascending order

Sort the list of trainers in ascending order

Iterate over each player in the sorted list

trainer_index += 1

match count += 1

trainer_index += 1

Return the total number of matches

if trainer index < len(trainers):</pre>

Increment the match count

// Sort the players array in ascending order

// Sort the trainers array in ascending order

if (trainersIndex < trainers.size()) {</pre>

// Import the necessary module for sorting

// Function to sort an array in ascending order

return sort(array, (a, b) => a - b);

function sortArrayAscending(array: number[]): number[] {

// Sort the players and trainers in ascending order

// Function to match players with trainers based on their strength.

// Each player can only be matched with a trainer that is equal or greater in strength.

function matchPlayersAndTrainers(players: number[], trainers: number[]): number {

matches++: // Increment match count

return matches; // Return the total number of matches made

// Initialize the count of matches to 0

// Initialize the index for trainers to 0

def matchPlayersAndTrainers(self, players: List[int], trainers: List[int]) -> int:

Skip trainers which have less capacity than the current player's strength

// Increment the trainerIndex until we find a trainer that can match the player

trainersIndex++; // Move to the next trainer for the following players

while (trainerIndex < trainers.length && trainers[trainerIndex] < player) {</pre>

while trainer index < len(trainers) and trainers[trainer_index] < player:</pre>

If there is a trainer that can match the current player

Move to the next trainer for the next player

Initialize the count of matched players and trainers match count = 0# Initialize the pointer for trainers list trainer_index = 0

Java class Solution { public int matchPlayersAndTrainers(int[] players, int[] trainers) {

return match_count

Arrays.sort(players);

Arrays.sort(trainers);

int trainerIndex = 0;

// Iterate over each player

for (int player: players) {

int matches = 0;

```
trainerIndex++;
            // If there is a trainer that can match the player, increment the match count
            if (trainerIndex < trainers.length) {</pre>
                matches++; // A match is found
                trainerIndex++; // Move to the next trainer.
        // Return the total count of matches
        return matches;
C++
#include <vector>
#include <algorithm> // Include the necessary header for std::sort
class Solution {
public:
    // Function to match players with trainers based on their strength.
    // Each player can only be matched with a trainer that is equal or greater in strength.
    int matchPlayersAndTrainers(vector<int>& players, vector<int>& trainers) {
        std::sort(players.begin(), players.end()); // Sort the players in ascending order
        std::sort(trainers.begin(), trainers.end()); // Sort the trainers in ascending order
        int matches = 0; // Initialize the number of matches to zero
        int trainersIndex = 0; // Initialize the trainers index to the start of the array
        // Iterate through each player
        for (int playerStrength : players) {
            // Find a trainer that can match the current player's strength
            while (trainersIndex < trainers.size() && trainers[trainersIndex] < playerStrength) {</pre>
                trainersIndex++; // Increment trainers index if the current trainer is weaker than the player
            // If a suitable trainer is found, make the match
```

import { sort } from 'some-sorting-module'; // Please replace 'some-sorting-module' with the actual module you would use for sorting

const sortedPlayers: number[] = sortArrayAscending(players); const sortedTrainers: number[] = sortArrayAscending(trainers); let matches: number = 0; // Initialize the number of matches to zero let trainersIndex: number = 0; // Initialize the trainers index to the start of the array

};

TypeScript

```
// Iterate through each player and match with trainers based on strength
    for (const playerStrength of sortedPlayers) {
        // Find a trainer that can match the current player's strength
       while (trainersIndex < sortedTrainers.length && sortedTrainers[trainersIndex] < playerStrength) {</pre>
            trainersIndex++; // Increment trainers index if the current trainer is weaker than the player
       // If a suitable trainer is found, make the match
        if (trainersIndex < sortedTrainers.length) {</pre>
            matches++: // Increment match count
            trainersIndex++; // Move to the next trainer for the following players
    return matches; // Return the total number of matches made
// Export the function if this is part of a module
export { matchPlayersAndTrainers };
from typing import List # Import List from typing module for type hints
class Solution:
    def matchPlayersAndTrainers(self, players: List[int], trainers: List[int]) -> int:
       # Sort the list of players in ascending order
       players.sort()
       # Sort the list of trainers in ascending order
        trainers.sort()
       # Initialize the count of matched players and trainers
       match count = 0
       # Initialize the pointer for trainers list
        trainer index = 0
       # Iterate over each player in the sorted list
        for player in players:
            # Skip trainers which have less capacity than the current player's strength
           while trainer index < len(trainers) and trainers[trainer_index] < player:</pre>
                trainer_index += 1
           # If there is a trainer that can match the current player
            if trainer index < len(trainers):</pre>
                # Increment the match count
```

Return the total number of matches return match_count Time and Space Complexity

Space Complexity

match count += 1

trainer_index += 1

Move to the next trainer for the next player

Time Complexity The time complexity of the code is determined by the sorting operations and the subsequent for-loop with the inner while loop.

and m is the number of trainers. The for-loop iterates over each player, which is O(n). Within this loop, the while loop progresses without resetting, which

Sorting both the players and trainers list takes O(nlogn) and O(mlogm), respectively, where n is the number of players

makes it linear over m elements of trainers in total. In the worst case, all players are checked against all trainers, hence it contributes a maximum of O(m) time complexity.

The combined time complexity of these operations would be 0(nlogn) + 0(mlogm) + 0(n + m). However, since 0(nlogn) and O(mlogm) are the dominant terms, the overall time complexity simplifies to O(nlogn + mlogm).

The space complexity of the code is 0(1), which is the additional space required. The sorting happens in place, and no additional data structures are used aside from a few variables for keeping track of indexes and the answer.