3071. Minimum Operations to Write the Letter Y on a Grid

We have to transform this grid such that it forms the letter 'Y'. A grid forms the letter 'Y' when:

Matrix

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Problem Description
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Hash Table

Counting

1. All cells that are part of the letter 'Y' (the two diagonals converging at the center cell from the top left and top right and the vertical line from the center down to the bottom) contain the same value.

In this problem, we are provided with an odd-sized, 0-indexed $n \times n$ grid containing three possible values in each cell: 0, 1, or 2.

2. All cells that are not part of the letter 'Y' also contain the same value, which must be distinct from the value of the 'Y' cells. 3. A cell's value can be changed to either 0, 1, or 2 in one operation.

Our goal is to find the *minimum* number of such operations needed to write the letter 'Y' on the grid according to these rules.

Intuition

The solution approach leverages the concept of counting and enumeration. Knowing that all cells of the 'Y' must share one value

We build two count dictionaries: cnt1 keeps track of how many times each value appears within the cells that make up the letter

Medium

<u>Array</u>

(let's call it 'i'), and all cells not a part of 'Y' must share a different value ('j'), there are limited possibilities since we have only three different values to choose from (0, 1, or 2). Moreover, when choosing a value for the 'Y' part and a different value for the non-'Y' part of the grid, these values can't be the same.

'Y', and cnt2 does the same for the cells outside of the letter 'Y'. Then, for each pair of values (i for the 'Y', j for the rest) we calculate the number of operations that would be needed if we chose this pair of values. This is done simply by subtracting the

the 'Y':

number of already correctly valued cells from the total cell count, which gives us the number of cells we need to change. **Solution Approach** The implementation of the solution follows a counting and enumeration approach as indicated in the Reference Solution

Approach. Firstly, two Counter objects from Python's collections module, cnt1 and cnt2, are used. They serve as the data structures to keep track of the frequency of each value in the cells that belong to the 'Y' (cnt1) and those that do not (cnt2). The solution iterates over each cell in the grid using nested loops. For each cell at position (i, j), it determines if the cell

belongs to the letter 'Y'. This classification is done using three conditions, indicating the cell is part of one of the three parts of

• b = i + j == n - 1 and i <= n // 2: Checks if the cell is on the diagonal from the top-right to the center cell. • c = j == n // 2 and i >= n // 2: Checks if the cell is part of the vertical line from the center down. Using logical OR (a or b or c), we can ascertain if a cell is part of the 'Y'. If it is, increment the count for the value in cnt1, otherwise in cnt2.

In the minimizing step, we generate permutations of two different values, i and j (with i!= j), where i is a candidate value

for 'Y' cells, and j is for the non-'Y' cells. We subtract the sum of the count of already correct 'Y' cells (cnt1[i]) and the count of

• a = i == j and i <= n // 2: Checks if the cell is on the diagonal from the top-left to the center cell.

changes needed to form the letter 'Y' on the grid according to the problem's conditions.

Counting Values: We use two Counter objects, cnt1 for the 'Y' and cnt2 for the non-'Y'.

already correct non-'Y' cells (cnt2[j]) from the total number of cells in the grid (n * n). This gives us the count of operations for this specific combination of values i and j. As we want to minimize the number of operations, we take the minimum over all (i, j) pairs using a generator expression:

min(n * n - cnt1[i] - cnt2[j] for i in range(3) for j in range(3) if i != j

By iterating over all value combinations and taking the minimum of the computed operations, we determine the least amount of

Let's illustrate the solution approach using a small 3×3 grid example. Our grid looks like this: 1 0 1 0 2 0 1 0 1

Initially, we determine the 'Y' cells are located at indices (0,0), (0,2), (1,1), (2,1). The rest of the cells are not part of the 'Y'.

For cnt1 ('Y' cells), the grid has:

Now, following the solution approach:

Value 1 at indices (0,0) and (0,2), so cnt1[1] = 2.

No other values are present in the non-'Y' cells.

Proceeding similarly for other combinations:

Value 0 at indices (0,1), (1,0), (1,2), and (2,0), (2,2), so cnt2[0] = 5.

• Operations needed: 3 * 3 - cnt1[1] - cnt2[0] = 9 - 2 - 5 = 2.

Value 2 at index (1,1), so cnt1[2] = 1.

For cnt2 (non-'Y' cells), the grid has:

Example Walkthrough

ensuring i != j. For each combination, we calculate the needed operations. For i = 1 (value for 'Y') and j = 0 (value for non-'Y'):

Determining Operations: We enumerate the permutations of different values for 'Y' (i) and non-'Y' (j) from the set {0, 1, 2},

• i = 1, j = 2: Operations = 9 - cnt1[1] - cnt2[2] = 9 - 2 - 0 = 7 (since cnt2[2] is not present). • i = 2, j = 0: Operations = 9 - cnt1[2] - cnt2[0] = <math>9 - 1 - 5 = 3.

from collections import Counter

def minimumOperationsToWriteY(self, grid):

Get the size of the grid

non_Y_counter = Counter()

vert horiz counter = Counter()

grid (List[List[int]]): A square grid of integers.

int: The minimum number of operations needed.

Initialize counters for each part of the "Y"

Count the occurrences

vert_horiz_counter[x] += 1

non_Y_counter[x] += 1

public int minimumOperationsToWriteY(int[][] grid) {

int minimumOperationsToWriteY(vector<vector<int>>& grid) {

// Iterate through the grid and categorize elements

for (int i = 0; i < gridSize; ++i) {</pre>

} else {

if (i != i) {

for (int i = 0; i < gridSize; ++i) {</pre>

int gridSize = grid.size(); // total size of the grid (n x n)

int countGroupOne[3] = $\{0\}$; // frequency count of numbers in the first group (Y shape)

// Increase the correct frequency counter based on the booleans above

// Find the minimum number of operations by checking all combinations of colours between the groups

if (isDiagonalOne || isDiagonalTwo || isVerticalLine) {

++countGroupTwo[grid[i][j]]; // for non-Y shape

for (int i = 0; i < 3; ++i) { // iterate over possible colours for group one

// Ensure we are not counting the same colour for both groups

// And update minimumOperations if this count is lower

return minimumOperations; // Return the minimum number of operations found

* Calculates the minimum number of operations to write the integer 'Y'.

* @return {number} The minimum number of operations to write 'Y'.

function minimumOperationsToWriteY(grid: number[][]): number {

* @param {number[][]} grid - A square grid of numbers.

// Get the size of the arid.

for (int i = 0; i < 3; ++i) { // iterate over possible colours for group two

// Calculate operations needed to paint the current combination

++countGroupOne[grid[i][j]]; // for Y shape

int countGroupTwo[3] = $\{0\}$; // frequency count of numbers in the second group (not Y shape)

is vert or first diag = i == i and i <= n // 2

is_horiz_middle = j == n // 2 and i >= n // 2

is second diag = i + j == n - 1 and i <= n // 2

if is vert or first diag or is_second_diag or is_horiz_middle:

// Method to calculate the minimum number of operations needed to write 'Y' on the grid

int gridSize = grid.length; // Dimension of the grid (since it's n x n)

Python

class Solution:

Args:

Returns:

n = len(grid)

```
\circ min(2, 7, 3, 8) = 2.
  Therefore, the minimum number of operations needed to transform this 3×3 grid into the letter 'Y' is 2. We can achieve this by
  changing the middle element (1,1) from 2 to 1, and any of the non-'Y' elements from 0 to 1.
Solution Implementation
```

• i = 2, j = 1: Operations = 9 - cnt1[2] - cnt2[1] = 9 - 1 - 0 = 8 (since cnt2[1] is not present).

Calculate the minimum number of operations required to write the letter Y on the grid.

Minimizing Operations: We take the smallest number of operations from the calculations. In this case, it is:

Iterate over the grid to count the occurrences in the "Y" shape and elsewhere for i, row in enumerate(grid): for j, x in enumerate(row): # Calculate which part of the grid is the 'Y' shape

Compute the minimum operations by finding max occurrence in 'Y' and outside 'Y', excluding the same number in both places

n * n - vert horiz counter[i] - non Y counter[j] for i in range(3) for j in range(3) if i != j

Java

class Solution {

else:

min operations = min(

return min_operations

```
// Arravs to store the count of each number (0, 1, 2) in the regions that form 'Y'
        int[] countRegionY = new int[3];
        int[] countOutsideRegionY = new int[3];
        // Loop through each cell in the grid to separate the cells that will form 'Y'
        // and those that will not
        for (int i = 0; i < gridSize; ++i) {</pre>
            for (int i = 0; i < gridSize; ++i) {</pre>
                // Conditions to detect cells that are part of 'Y'
                boolean isDiagonalFromTopLeft = i == j && i <= gridSize / 2;</pre>
                boolean isDiagonalFromTopRight = i + i == gridSize - 1 && i <= gridSize / 2;</pre>
                boolean isVerticalMiddleLine = j == gridSize / 2 && i >= gridSize / 2;
                if (isDiagonalFromTopLeft || isDiagonalFromTopRight || isVerticalMiddleLine) {
                    // Increment the count for the region that forms 'Y'
                    ++countRegionY[grid[i][j]];
                } else {
                    // Increment the count for the region outside 'Y'
                    ++countOutsideRegionY[grid[i][j]];
        // We'll assume the worst case where we need to change every cell initially
        int minOperations = gridSize * gridSize;
        // Evaluate all combinations where the number for region 'Y' and outside are different
        for (int numberForY = 0: numberForY < 3: ++numberForY) {</pre>
            for (int numberOutsideY = 0; numberOutsideY < 3; ++numberOutsideY) {</pre>
                if (numberForY != numberOutsideY) {
                    // Calculate the number of operations needed if these two numbers were used
                    int operations = gridSize * gridSize - countRegionY[numberForY] - countOutsideRegionY[numberOutsideY];
                    // Check if the current operation count is lower than the minimum found so far
                    minOperations = Math.min(minOperations, operations);
        // Return the minimum number of operations found
        return minOperations;
C++
class Solution {
public:
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bool isDiagonalOne = (i == j) && (i <= gridSize / 2); // condition for main diagonal and upper half

bool isDiagonalTwo = (i + j == gridSize - 1) & (i <= gridSize / 2); // condition for anti-diagonal and upper half

int minimumOperations = gridSize * gridSize; // maximum number of operations (each cell made to write an element not already

minimumOperations = min(minimumOperations, gridSize * gridSize - countGroupOne[i] - countGroupTwo[j]);

bool is Vertical Line = (j == grid Size / 2) & (i >= grid Size / 2); // condition for the vertical middle line and lower

```
const gridSize = grid.length;
// Initialize counters for the number collections.
const crossCounters: number[] = Array(3).fill(0);
```

*/

};

TypeScript

```
const otherCounters: number[] = Array(3).fill(0);
    // Loop through the grid to count occurrences.
    for (let i = 0; i < gridSize; ++i) {</pre>
        for (let i = 0; i < gridSize; ++i) {</pre>
            // Check if the current cell is part of the cross (Y shape).
            const onFirstDiagonal = i === j && i <= gridSize >> 1;
            const onSecondDiagonal = i + i === gridSize - 1 && i <= gridSize >> 1;
            const onMiddleColumn = j === gridSize >> 1 && i >= gridSize >> 1;
            // Update the corresponding counters based on the cell's position.
            if (onFirstDiagonal || onSecondDiagonal || onMiddleColumn) {
                ++crossCounters[grid[i][j]];
            } else {
                ++otherCounters[grid[i][j]];
    // Initialize the answer with the maximum possible number of operations.
    let minimumOperations = gridSize * gridSize;
    // Determine the minimum operations required by trying all combinations of numbers.
    for (let i = 0; i < 3; ++i) {
        for (let i = 0; i < 3; ++i) {
            // We want to use different numbers for the cross and the other cells.
            if (i !== i) {
                // Calculate the minimum operations by subtracting the already correct cells.
                minimumOperations = Math.min(
                    minimumOperations.
                    gridSize * gridSize - crossCounters[i] - otherCounters[j]
    // Return the calculated minimum operations.
    return minimumOperations;
from collections import Counter
class Solution:
    def minimumOperationsToWriteY(self, grid):
        Calculate the minimum number of operations required to write the letter Y on the grid.
        Aras:
        grid (List[List[int]]): A square grid of integers.
        Returns:
        int: The minimum number of operations needed.
        # Get the size of the grid
        n = len(grid)
        # Initialize counters for each part of the "Y"
        vert horiz counter = Counter()
        non_Y_counter = Counter()
        # Iterate over the grid to count the occurrences in the "Y" shape and elsewhere
        for i, row in enumerate(grid):
            for i, x in enumerate(row):
                # Calculate which part of the grid is the 'Y' shape
                is vert or first diag = i == i and i <= n // 2
                is second diag = i + j == n - 1 and i <= n // 2
                is_horiz_middle = j == n // 2 and i >= n // 2
                # Count the occurrences
```

Time Complexity

n x n. The operations inside the for-loops are constant time operations, causing the overall time to be proportional to the number of elements in the grid. **Space Complexity**

The time complexity of the given code is $0(n^2)$. This is because the nested for-loops iterate over the entire grid, which is of size

Compute the minimum operations by finding max occurrence in 'Y' and outside 'Y', excluding the same number in both places

if is vert or first diag or is_second_diag or is_horiz_middle:

vert_horiz_counter[x] += 1

n * n - vert horiz counter[i] - non Y counter[j]

for i in range(3) for j in range(3) if i != j

non_Y_counter[x] += 1

else:

min operations = min(

return min_operations

Time and Space Complexity

The space complexity of the code is higher than 0(1) mentioned in the reference answer. It actually depends on the range of the elements in the grid. The two counters cnt1 and cnt2 are used to count occurrences of elements that satisfy certain conditions. Because Python's Counter is essentially a hash map, the space used by these counters will increase with the variety of elements in the grid. Assuming the range of numbers in the grid is k, the space complexity would be 0(k). If we were to consider k as being a constant because the problem might define a limit on the values of grid elements, then we could consider the space complexity to be 0(1). Without such a constraint being specified, the space complexity is not strictly constant.