

Problem Description

either num + 1 or num + 2. These two integers should be as close to each other as possible in terms of their absolute difference - meaning the difference between the two numbers without considering whether it is positive or negative should be minimal. The

The problem presents a scenario where you are given a single integer num. Your task is to find two integers whose product is

final solution does not require the integers to be in any specific order, thus either of the integers can come first.

Intuition

without leaving a remainder. In this case, our target numbers are num + 1 and num + 2. We look for the pair of factors, for each of these numbers, that are closest to each other. The intuition here is that the pair of factors that are closest to each other will have the smallest absolute difference. To efficiently find such a pair for a given target number, we can start checking from the square root of the target number and

In solving this problem, the concept of factors of a number is key. For any given number, its factors are the numbers that divide it

move downwards. The square root gives us a good starting point since it is the largest number that can multiply by itself to not exceed the target. Therefore, any factor larger than the square root would result in a product larger than our target number, disqualifying it from being a correct answer. Once we have the factors for both num + 1 and num + 2, we compare their absolute differences. The pair with the smaller

absolute difference is the closest pair – this represents our final answer. The reason for checking both num + 1 and num + 2 is to fulfill the task's requirement of finding the closest integers in absolute difference.

num + 1 or num + 2) and finds the closest divisors of x.

Solution Approach

Here's how the f(x) function works in detail: 1. Start by calculating the integer square root of x using int(sqrt(x)). This is the starting point for finding our factors.

The solution uses a straightforward approach by defining a helper function f(x) which takes a number x (which will either be

2. Iterate downwards from this square root to 1 (inclusive), as any factor greater than the square root of x would result in a product larger than x when multiplied with another integer.

- 3. In each iteration, check if x % i == 0 which is the condition to confirm if i is a factor of x.
- 4. Once a factor is found, calculate the pair by dividing x by the factor i. The pair will be [i, x // i]. 5. Return the pair of factors found.
- The main function closestDivisors(num: int) -> List[int] calls this helper function twice: once with num + 1 and once with
- 1. The function stores the pairs returned by f(num + 1) and f(num + 2) in variables a and b respectively.

• Starting the search from the square root of these numbers and iterating downwards to find these pairs.

2. It then compares the absolute differences of the two integers in each pair.

4. The result pair is returned.

num + 2.

iterating through all possible divisors, which significantly reduces the number of operations, especially for very large numbers. To summarize, the solution algorithm involves:

Algorithmically, this is an efficient approach because it only searches up to the square root of the target numbers rather than

Returning the pair with the smallest absolute difference between the two factors.

• Finding pairs of factors (divisors) for two numbers num + 1 and num + 2.

3. The pair with the smaller absolute difference is chosen as the result.

- Data structures used in this solution are basic and include primarily lists to store the pairs of factors. The pattern utilized here is an optimization over brute force where only necessary divisors are considered, which is made possible by the mathematical property that a number's divisors are symmetrical around the square root.
- Let's assume num = 8. We want to find two integers whose product is num + 1 or num + 2, which here would be 9 and 10. These two integers should have the smallest possible absolute difference. Let's apply our solution approach:

Then, we check if there's any number from 3 down to 1 that evenly divides 9:

point.

Example Walkthrough

○ 3 is a divisor of 9 since 9 % 3 == 0. \circ The pair is formed by 3 (the divisor) and 9 // 3 = 3. So the factor pair for num + 1 is [3, 3]. Next, we calculate the square root of num + 2, which is slightly over sqrt(10) but we use int(sqrt(10)) = 3 as our starting

We check divisors from 3 to 1 for the number 10:

We first calculate the square root of num + 1 which is sqrt(9) = 3.

 3 is not a divisor of 10 since 10 % 3 != 0. ○ 2 is a divisor since 10 % 2 == 0.

The pair for num + 2 is [2, 10 // 2], which is [2, 5].

- Now we have two pairs: [3, 3] for num + 1 and [2, 5] for num + 2. We compare their absolute differences:
- \circ The absolute difference for [3, 3] is 3 3 = 0. • The absolute difference for [2, 5] is 5 - 2 = 3.

The final result is [3, 3], indicating that these are the two integers (which in this case happen to be identical) that, when

We choose the pair with the smallest absolute difference, which is [3, 3], with an absolute difference of 0.

multiplied together, yield num + 1 (which is 9) and have the smallest absolute difference.

Start by finding the square root of 'x' and iterate backwards

closest divisors num plus one = find closest_divisors(num + 1)

closest_divisors_num_plus_two = find_closest_divisors(num + 2)

// Helper function that calculates the two closest divisors of 'x'

// Start from the square root of 'x' and check for the closest divisors by moving downwards

// If 'i' divides 'x' with no remainder, 'i' and 'x / i' are divisors of 'x'

// Found the closest divisors, return them in an array

private int[] findClosestDivisors(int x) {

return new int[] {i, x / i};

 $if (x \% i == 0) {$

we quickly identify the closest pair of numbers satisfying the condition without unnecessary computation.

This example walk-through demonstrates the solution approach. By checking only the divisors from the square root and below,

class Solution: def closestDivisors(self, num: int) -> List[int]: # Define a helper function to find the pair of divisors # of a number 'x' that are closest to each other.

if x % i == 0: # Return the divisor pair [i, x // i] return [i, x // i]

def find closest divisors(x):

for i in range(int(sgrt(x)), 0, -1):

If 'i' is a divisor of 'x'

Find the closest divisors for 'num + 1'

Find the closest divisors for 'num + 2'

Solution Implementation

from typing import List

from math import sqrt

Python

```
# Compare which pair of divisors has the smallest difference
        # and return that pair.
        if abs(closest divisors num plus one[0] - closest_divisors_num_plus_one[1]) < abs(closest_divisors_num_plus_two[0] - closest_
            return closest_divisors_num_plus_one
        else:
            return closest_divisors_num_plus_two
Java
class Solution {
    // This function finds two closest divisors of the input number 'num'
    public int[] closestDivisors(int num) {
        // Find the closest divisors for the number 'num + 1'
        int[] divisorsNumPlusOne = findClosestDivisors(num + 1);
        // Find the closest divisors for the number 'num + 2'
        int[] divisorsNumPlusTwo = findClosestDivisors(num + 2);
        // Compare abs difference of divisors pairs and return the pair with the smallest difference
        if (Math.abs(divisorsNumPlusOne[0] - divisorsNumPlusOne[1]) <</pre>
            Math.abs(divisorsNumPlusTwo[0] - divisorsNumPlusTwo[1])) {
            return divisorsNumPlusOne;
        } else {
            return divisorsNumPlusTwo;
```

for (int i = (int) Math.sqrt(x); true; --i) { // the condition is always true, it breaks inside if a divisor pair is found

TypeScript

};

// Import sart function from Math module

if (x % i === 0) {

// Function to find closest divisors of an integer 'num'

const findClosestDivisors = (x: number): number[] => {

// Lambda function to find the divisors closest to the square root of 'x'

// If i is a divisor, return the pair [i, x/i]

const divisorsForNumPlusOne: number[] = findClosestDivisors(num + 1);

// Start from the largest possible factor that could be closest to sart of x

for (let $i = Math.floor(sqrt(x)); i > 0; --i) { // Ensure i is always positive}$

// Since every number has at least one pair of divisors, this line should not be reached.

throw new Error("No divisors found"); // To handle edge cases theoretically unreachable

function closestDivisors(num: number): number[] {

return [i, x / i];

import { sqrt, abs } from 'math';

```
C++
#include <vector>
#include <cmath> // Include cmath for sqrt function
class Solution {
public:
    // Function to find the closest divisors of an integer 'num'
    vector<int> closestDivisors(int num) {
        // Lambda function to find the divisors closest to the square root of 'x'
        auto findClosestDivisors = [](int x) -> vector<int> {
            // Start from the largest possible factor that could be closest to sqrt of x
            for (int i = sart(x); i > 0; --i) { // Ensure i is always positive
                if (x \% i == 0) {
                    // If i is a divisor, return the pair (i, x/i)
                    return vector<int>{i, x / i};
            // This code should never reach here since every number has at least one pair of divisors
        };
        // Find the closest divisors for both num+1 and num+2
        vector<int> divisorsForNumPlusOne = findClosestDivisors(num + 1);
        vector<int> divisorsForNumPlusTwo = findClosestDivisors(num + 2);
        // Determine which pair of divisors has the smallest difference
        // Return the pair with the smallest difference
        return abs(divisorsForNumPlusOne[0] - divisorsForNumPlusOne[1]) < abs(divisorsForNumPlusTwo[0] - divisorsForNumPlusTwo[1]) ?
};
// Remember to include necessary headers before using this code.
```

```
const divisorsForNumPlusTwo: number[] = findClosestDivisors(num + 2);
// Determine which pair of divisors has the smallest difference, and return the pair with the smallest difference
```

// Find the closest divisors for both num+1 and num+2

```
return abs(divisorsForNumPlusOne[0] - divisorsForNumPlusOne[1]) < abs(divisorsForNumPlusTwo[0] - divisorsForNumPlusTwo[1])
        ? divisorsForNumPlusOne
        : divisorsForNumPlusTwo;
from typing import List
from math import sqrt
class Solution:
   def closestDivisors(self, num: int) -> List[int]:
        # Define a helper function to find the pair of divisors
       # of a number 'x' that are closest to each other.
       def find closest divisors(x):
           # Start by finding the square root of 'x' and iterate backwards
            for i in range(int(sqrt(x)), 0, -1):
                # If 'i' is a divisor of 'x'
                if x % i == 0:
                    # Return the divisor pair [i, x // i]
                    return [i, x // i]
       # Find the closest divisors for 'num + 1'
        closest divisors num plus one = find closest_divisors(num + 1)
       # Find the closest divisors for 'num + 2'
        closest_divisors_num_plus_two = find_closest_divisors(num + 2)
       # Compare which pair of divisors has the smallest difference
       # and return that pair.
        if abs(closest divisors num plus one[0] - closest_divisors_num_plus_one[1]) < abs(closest_divisors_num_plus_two[0] - closest_
```

Time and Space Complexity

numbers starting from int(sqrt(x)) to 1.

num has a negligible effect for large numbers.

used is for a handful of variables that store the divisor pairs and their differences.

return closest_divisors_num_plus_one else: return closest_divisors_num_plus_two

Time Complexity The time complexity of the function primarily depends on the for loop within the nested function f(x), which iterates over

Since the square root function essentially reduces the number of iterations to the square root of x, the time complexity for finding the divisors would be O(sqrt(x)).

Given that the function f(x) is called twice—once for num + 1 and once for num + 2—the overall time complexity is then

O(sqrt(num + 1) + sqrt(num + 2)), which simplifies to O(sqrt(num)), as the higher order term dominates and adding 1 or 2 to

Space Complexity

The space complexity of this algorithm is 0(1), as the space used does not grow with the input size num. The only extra space