2096. Step-By-Step Directions From a Binary Tree Node to Another

You are given the root of a binary tree with {" "} n nodes. Each node is uniquely assigned a value from {" "} 1 to n. You are also given an integer{" "} startValue representing the value of the start node{" "} s, and a different integer destValue representing the value of the destination node t. Find the shortest path starting from node s {" "} and ending at node t. Generate step-by-step directions of such path as a string

consisting of only the uppercase letters{" "} 'L', 'R', and 'U'. Each letter indicates a specific direction:

Return{" "} the step-by-step directions of the shortest path from node{" "} s to node t.

"U" means to go from a node to its parent{" "} node.

• 'L' means to go from a node to its{" "} left child node.

means to go from a node to its{" "} right child node.

Example 1:

Input: root = [5,1,2,3,null,6,4], startValue = 3, destValue $= 6{"}n"}$ Output: "UURL"{"\n"} **Explanation:** The shortest path is: $3 \rightarrow 1 \rightarrow 5 \rightarrow 2 \rightarrow 6.\{"\n"\}$

Input: root = [2,1], startValue = 2, destValue = $1{"\n"}$

Explanation: The shortest path is: 2 → 1.{"\n"}

Example 2:

Output: "L"{"\n"}

Constraints: • The number of nodes in the tree is n. • 2 <= n <= 10^5

• 1 <= startValue, destValue <= n startValue != destValue Solution

1 <= Node.val <= n

All the values in the tree are unique.

path(root, start)

We can perform a DFS ($\frac{\text{depth first search}}{\text{depth first search}}$) to get path(root, end). This path consists of 'L's and 'R's. We can do another DFS to get path(root, start). Replacing the 'L's and 'R's of path(root, start) with 'U's gives us path(start, root). Now we can

concatenate path(start, root) and path(root, end) to get the answer.

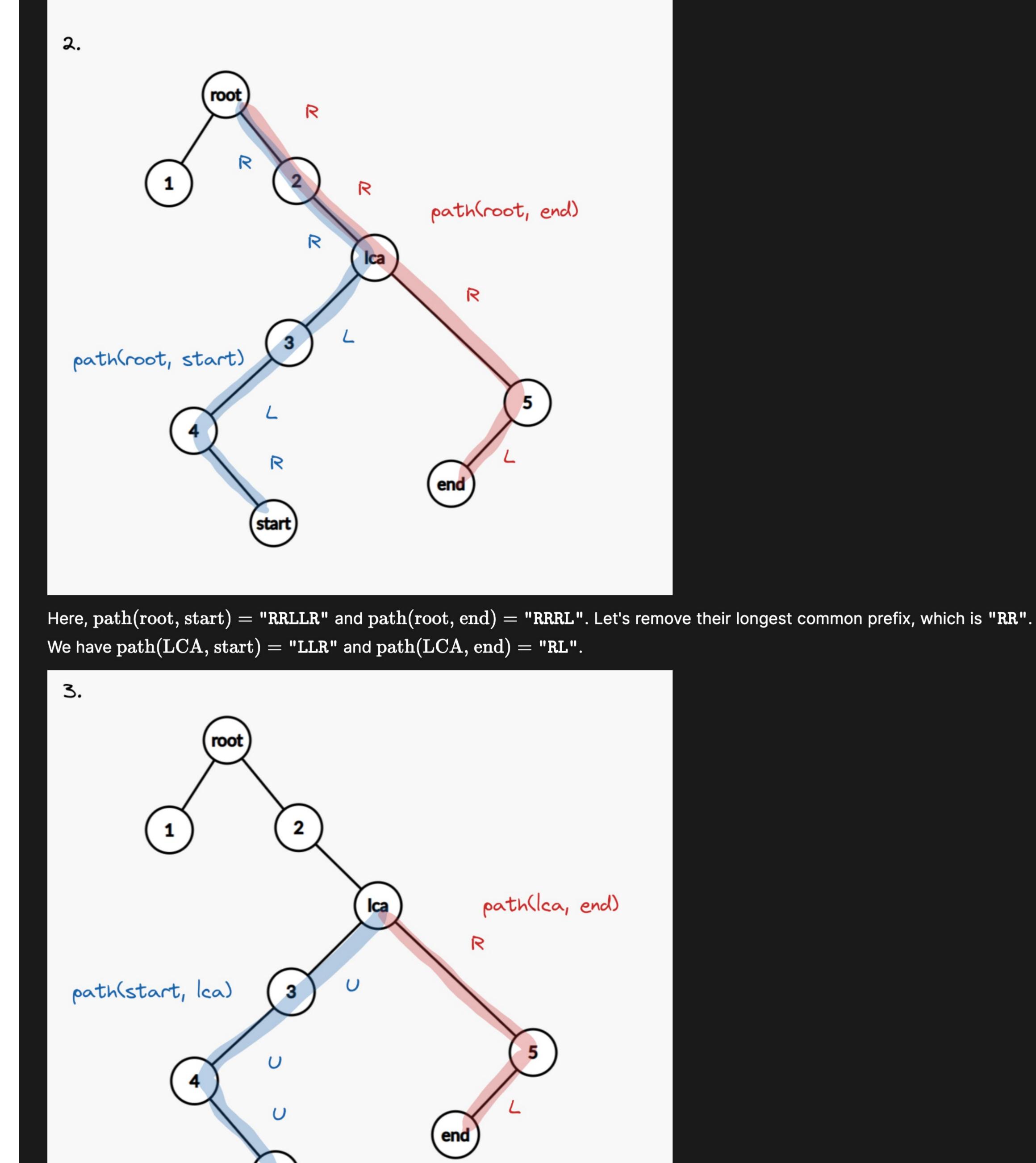
("U"s) before going down a non-negative number of times ("L"s or "R"s). The highest node in this path is known as the LCA (lowest common ancestor) of start and end.

Let path(node1, node2) denote the path from node1 to node2.

path(root, end)

First consider the case where path(start, end) goes through the root. Let's split this into path(start, root) + path(root, end).

In the general case, path(start, end) may not go through the root. Notice that this path goes up a non-negative number of times



TreeNode *left; TreeNode *right; TreeNode() : val(0), left(nullptr), right(nullptr) {} TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}

Then we replace all characters in path(root, start) with "U"s to obtain path(start, lca) = "UUU". Finally, we get

Each DFS takes $\mathcal{O}(n)$ and our string operations never happen on strings exceeding length $\mathcal{O}(n)$. The time complexity is $\mathcal{O}(n)$.

path(start, end) = path(start, LCA) + path(LCA, end) = "UUU" + "RL" = "UUURL".

TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left),

void getPath(TreeNode *cur, int targetValue, string &path, string &ans) {

The strings never exceed length $\mathcal{O}(n)$. The space complexity is $\mathcal{O}(n)$.

```
ans = path;
path.push back('L');
qetPath(cur->left, targetValue, path, ans);
path.back() = 'R';
getPath(cur->right, targetValue, path, ans);
```

Time complexity

Space complexity

C++ Solution

* struct TreeNode {

int val;

* right(right) {}

if (!cur)

int val;

*/

class Solution {

TreeNode left;

TreeNode() {}

TreeNode right;

this.val = val;

if (cur == null)

return;

path.append("L");

int i = 0;

i++;

this.left = left;

this.right = right;

if (cur.val == targetValue)

int strLen = path.length();

path.delete(strLen, strLen+1);

ans[0] = path.toString();

path.replace(strLen, strLen+1, "R");

TreeNode(int val) { this.val = val; }

TreeNode(int val, TreeNode left, TreeNode right) {

// so appending a character takes O(length of string).

qetPath(cur.left, targetValue, path, ans);

getPath(cur.right, targetValue, path, ans);

StringBuilder tmpPath = new StringBuilder();

String[] startPath = {""}, destPath = {" "};

qetPath(root, destValue, tmpPath, destPath);

getPath(root, startValue, tmpPath, startPath);

// Find the first point at which the paths diverge

// We use StringBuilder instead of String because String is immutable,

// ans is a String[1] instead of String because in Java, arrays are passed by reference.

void getPath(TreeNode cur, int targetValue, StringBuilder path, String[] ans) {

public String getDirections(TreeNode root, int startValue, int destValue) {

return "U".repeat(startPath[0].length()-i) + destPath[0].substring(i);

return;

path.pop_back();

if (cur->val == targetValue)

class Solution {

* Definition for a binary tree node.

/**

* };

public:

*/

string getDirections(TreeNode *root, int startValue, int destValue) { string tmpPath, startPath, destPath; qetPath(root, startValue, tmpPath, startPath); getPath(root, destValue, tmpPath, destPath); // Find the first point at which the paths diverge auto [itr1, itr2] = mismatch(startPath.begin(), startPath.end(), destPath.begin(), destPath.end()); return string(startPath.end() - itr1, 'U') + string(itr2, destPath.end()); **Java Solution** /** * Definition for a binary tree node. * public class TreeNode {

Python Solution # Definition for a binary tree node. # class TreeNode: def ___init___(self, val=0, left=None, right=None):

self.val = val self.left = left self.right = right class Solution: def getDirections(self, root: Optional[TreeNode], startValue: int, destValue: int) -> str: # ans is a list so that it passes by reference def getPath(cur, targetValue, path, ans): if cur is None: return if cur.val == targetValue: ans.append(''.join(path)); path.append('L'); getPath(cur.left, targetValue, path, ans) path[-1] = 'R'qetPath(cur.right, targetValue, path, ans) path.pop(-1) tmpPath = []startPath = [] destPath = []getPath(root, startValue, tmpPath, startPath) qetPath(root, destValue, tmpPath, destPath) startPath = startPath[0] destPath = destPath[0] # Find the first point at which the paths diverge i = 0while i < min(len(startPath), len(destPath)) and startPath[i] == destPath[i]:</pre> i += 1 return 'U'*(len(startPath)-i) + destPath[i:]

while (i < Math.min(startPath[0].length(), destPath[0].length()) && startPath[0].charAt(i) == destPath[0].ch</pre>