Sorting

Two Pointers

Problem Description

The problem presents us with an array named nums containing an equal number of odd and even integers. The task is to arrange this array in such a way that the parity of the elements matches the parity of their indices. In essence, all the even-indexed positions must contain even numbers, and all the odd-indexed positions must be filled with odd numbers. The challenge is to figure out a way to reorder the array to meet this requirement, knowing that there is at least one way to achieve it because of the even distribution of odd and even numbers.

Intuition

Since we know half the numbers are even and half are odd, we understand that there must be a correct place for each number in the 'reordered' array. The intuition is to iterate over the even indices and whenever we encounter an odd number at an even index, we search for an even number at an odd index to swap with. This is why we have a secondary pointer j starting at index 1, the first odd index, and it moves in steps of 2 (to remain at odd indices). Every time we find an even number at an odd index, we swap it with the mispositioned odd number we found at the even index. This swapping continues until all even indices hold even numbers (and thus, all odd indices hold odd numbers) since we are

The approach to solving this problem involves in-place reordering of the elements to meet our even-odd position requirements.

swapping odd for even, maintaining the parity balance. Once we're done with the loop, every even index i will have an even

The solution provided uses a two-pointer approach to rearrange the array such that even-indexed elements are even, and odd-

number and ever odd index will have an odd number, resulting in an array that satisfies the conditions of the problem. **Solution Approach**

indexed elements are odd. The algorithm iterates through the array with these key steps:

We initialize two pointers, i and j. Here i starts at index 0 (the first even index) and will be incremented by 2 in each iteration to ensure it only points to even indices. Pointer j starts at index 1 (the first odd index) and will be moved in

- We loop over the array with i up to n (n being the length of the array), making steps of size 2. At each iteration, we check if the number at index i is odd by observing the last bit nums[i] & 1. If it is equal to 1, it
- indicates that the number is odd and is in the wrong position since i is even. When an odd number is observed at an even index i, we look for an even number at an odd index j. We perform an inner

Python's tuple unpacking feature: nums[i], nums[j] = nums[j], nums[i].

increments of 2 to jump to the next odd index when needed.

- while loop to move j forward until we find an even number (also checking the last bit nums[j] & 1). As soon as an even number is found at an odd index j, we swap the elements at indexes i and j. This is done using
- The loop continues until i reaches the end of the array. At this point, all the even indices have even numbers due to our swapping mechanism, which ensures that every time we encounter a number out of place, we switch it with its corresponding pair.
- No additional data structures are needed, and the in-place algorithm performs swaps only when necessary. This leads to a space complexity of O(1) and a time complexity that is O(n) since each element is looked at once and each odd index is evaluated at most once.

The elegance of this solution comes from recognizing that, given the constraints, for every mispositioned odd number at an even

index, there must exist a mispositioned even number at an odd index that we can swap it with. Through careful stepping of the 🗓

and j pointers and efficient swapping, we're able to rearrange the array as needed without the use of any extra space.

0, 2, 4, and they should have even numbers, while the odd indices 1, 3, 5 should have odd numbers.

Begin looping over the array with i going from 0 to the length of the array, incremented by 2.

Example Walkthrough Let's use the following small example to illustrate the solution approach:

Consider the array nums = [3, 6, 1, 4, 7, 2]. Here, we have an equal number of odd and even integers. The even indices are

Following the steps from the solution approach:

On the first iteration: i is 0 and j is 1.

Initialize two pointers: i (starting at 0) and j (starting at 1).

nums [i] is 3, which is odd (3 & 1 equals 1), so it's out of place.

Now, we need to find an even number at an odd index j to swap with nums[i]. Since j is already at an odd index and

Increment i by 2. Now, i is 2.

nums[j] is 4, which is even.

Swap the elements at indices i and j. After swapping, nums becomes [6, 3, 1, 4, 7, 2].

nums [i] is 7, which is odd and incorrectly placed at an even index.

Initialize 'odd_index' to the first odd position

odd index += 2

public int[] sortArrayByParityII(int[] nums) {

int i = 0; // i should index even elements

while ((nums[j] & 1) != 0) {

int j = 1; // j should index odd elements

nums[j] is 6, which is even (6 & 1 equals 0), we don't need to move j.

- o nums [i] is 1, which is odd and out of place at an even index. Increment j by 2 to find the next even number. Now j is 3.
- Swap the elements at indices i and j. The array nums is now [6, 3, 4, 1, 7, 2]. Increment i by 2. Now, i is 4.
- 10. Increment j by 2 to move to the next odd index. Now j is 5.

nums[j] is 2, which is even.

- Swap the elements at indices i and j. The array nums is now [6, 3, 4, 1, 2, 7].
- Now each even index (0, 2, 4) has even numbers and each odd index (1, 3, 5) has odd numbers. The process is complete and the array nums is successfully reordered. The result is an array where the parity of the elements matches the parity of their indices,
- Solution Implementation

length = len(nums)

odd index = 1

thus [6, 3, 4, 1, 2, 7].

Python

class Solution: def sortArrayByParityII(self, nums: List[int]) -> List[int]: # Get the length of the input list

Increase 'odd_index' by 2 to check the next odd position

nums[even_index], nums[odd_index] = nums[odd_index], nums[even_index]

// i points to the next even index, j points to the next odd index

// Find the next number that's out of place at an odd index

* Sorts an array such that every even-indexed element is even and every odd-indexed element is odd.

// Initialize two pointers, `evenIndex` for even indexes and `oddIndex` for odd indexes.

// Swap the values at `evenIndex` and `oddIndex` to satisfy the condition.

[nums[evenIndex], nums[oddIndex]] = [nums[oddIndex], nums[evenIndex]];

* @param {number[]} nums - The input array to be sorted in place.

// Check if the current even index contains an odd value.

// Move the `evenIndex` pointer to the next even position.

// If so, look for an even value at the odd index positions.

function sortArrayByParityII(nums: number[]): number[] {

// Iterate over the array through the even indexes.

while ((nums[oddIndex] & 1) === 1) {

* @returns {number[]} The sorted array.

while (evenIndex < nums.length) {</pre>

if ((nums[evenIndex] & 1) === 1) {

oddIndex += 2;

let evenIndex: number = 0;

let oddIndex: number = 1;

// Iterate over the array, considering only even indices for i

// If the number at the current even index is odd,

// Swap the numbers at indices i and j

// we need to swap it with a number at an odd index.

Swap the odd number at the even index with the even number at the odd index

Iterate through all even indices of the list for even index in range(0, length, 2): # Check if the current even position contains an odd number if nums[even index] % 2 == 1: # Look for an even number at odd positions while nums[odd index] % 2 == 1:

Return the sorted by parity list return nums Java

while (i < nums.length) {</pre>

if ((nums[i] & 1) != 0) {

i += 2;

int temp = nums[i];

nums[i] = nums[j];

nums[j] = temp;

class Solution {

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// Move to the next even index
            i += 2;
        // Return the sorted array with even indices containing even numbers
        // and odd indices containing odd numbers
        return nums;
C++
#include <vector>
using namespace std;
class Solution {
public:
    // Function to sort the array so that at every even index we have an even number,
    // and at every odd index, we have an odd number.
    vector<int> sortArravBvParitvII(vector<int>& nums) {
        // Traverse through the elements at even indices.
        for (int evenIndex = 0, oddIndex = 1; evenIndex < nums.size(); evenIndex += 2) {</pre>
            // Check if the current even index has an odd number.
            if ((nums[evenIndex] & 1) == 1) {
                // Find the odd index that has an even number.
                while ((nums[oddIndex] & 1) == 1) {
                    oddIndex += 2;
                // Swap the odd number at an even index with the even number at an odd index.
                swap(nums[evenIndex], nums[oddIndex]);
        // Return the sorted array.
        return nums;
};
TypeScript
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/**

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evenIndex += 2;
    // Return the sorted array.
    return nums;
// Usage example:
// const sorted = sortArrayByParityII([4, 2, 5, 7]);
// console.log(sorted);
class Solution:
    def sortArrayByParityII(self, nums: List[int]) -> List[int]:
        # Get the length of the input list
        length = len(nums)
        # Initialize 'odd index' to the first odd position
        odd_index = 1
        # Iterate through all even indices of the list
        for even index in range(0, length, 2):
           # Check if the current even position contains an odd number
            if nums[even index] % 2 == 1:
               # Look for an even number at odd positions
               while nums[odd index] % 2 == 1:
                   # Increase 'odd_index' by 2 to check the next odd position
                   odd index += 2
                # Swap the odd number at the even index with the even number at the odd index
                nums[even_index], nums[odd_index] = nums[odd_index], nums[even_index]
        # Return the sorted by parity list
        return nums
Time and Space Complexity
Time Complexity
  The time complexity of the given code is O(n), where n is the length of the input list nums. This is because the outer loop runs
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for every second element (i.e., each even-indexed position), which results in n/2 iterations. Inside the loop, we check if the current even-indexed number is odd. If it is odd, we perform a while loop to find the next odd-indexed position that contains an odd number to swap with. The key point to note here is that each odd number is moved at most once since after it is moved to an even index, it doesn't need to be moved again. Therefore, the while loop altogether will also run for a maximum of n/2 odd numbers throughout the

entire execution of the algorithm. Combining the outer and inner loop we still have a linear relationship with the number of elements n, leading to a total time complexity of O(n).

Space Complexity

The space complexity of the code is 0(1) since only a fixed number of extra variables (n, j) are used, and these do not grow

with the input size. The swaps are done in-place, which means no additional space proportional to the input size is required.