356. Line Reflection

Hash Table

**Math** 

# **Problem Description**

Medium Array

The problem deals with a set of points on a 2D plane and the goal is to determine if there is a line parallel to the y-axis that can reflect all points symmetrically. In essence, we have to find out whether we can draw a vertical line such that every point has a mirrored counterpart on the other side of the line. The set of points before and after reflecting them across this line should be identical, even if some points are repeated.

## Intuition

have a mirrored point (x', y) such that both points are equidistant from the line of reflection. This implies that the sum of the xcoordinates of the two points (x + x') will be equal to twice the x-coordinate of the reflection line.

To solve this problem, we must recognize that a reflection over a line parallel to the y-axis means that each point (x, y) will

We approach the solution by first finding the minimal  $(min_x)$  and maximal  $(max_x)$  x-coordinates among all given points. The line of reflection, if it exists, would be exactly halfway between these two values, so the sum s for the reflection condition would be

 $min_x + max_x$ . Next, we store all given points in a set for constant-time lookups. Then, for every point in our input, we check if its symmetric

counterpart with respect to the potential line of reflection ((s - x, y)) exists in our set. If all points satisfy this condition, the points are symmetric with respect to a line, and the function should return true. Otherwise, no such line exists and we return false.

## To implement the solution, we use a set data structure to store the points for efficient querying, as well as simple variables to keep track of the minimum and maximum x-coordinates.

**Solution Approach** 

The solution is straightforward in terms of algorithms and does not require any complex data structure manipulation or intricate

patterns. It essentially involves the following steps: 1. Initialize min\_x to infinity and max\_x to negative infinity. These will be used to track the minimum and maximum x-values of the given points. 2. Traverse all points in the input list.

For each point (x, y), update min\_x and max\_x to keep track of the smallest and largest x-coordinates.

- Add the point (x, y) to the set point\_set for efficient lookups later. 3. Calculate the sum s which is the potential reflection line's x-coordinate times two (s = min\_x + max\_x). This is based on the principle that the
- 4. Finally, confirm that each point (x, y) has its mirrored counterpart. This is done by checking if (s x, y) is present in point\_set for every (x, y) in the list of points.

 $\circ$  For point (1, 1), we update min\_x to 1 (as 1 < infinity) and max\_x to 1 (as 1 > -infinity).

For point (3, 1), min\_x remains 1 and max\_x is updated to 3.

reflection's x-coordinate is  $s = min_x + max_x = 1 + 9 = 10$ .

line of reflection would be exactly in the middle of the leftmost and rightmost points.

If all points have their mirrored counterparts in the set, return true.

If there is even a single point without a mirrored counterpart, return false.

- The use of a set is crucial here as it allows us to verify the existence of the mirrored pairs in constant time (0(1)), thus keeping
- the entire algorithm's time complexity to O(n), where n is the number of points. There are no particular algorithms used beyond basic iteration and set operations. Also, there are no complex logical patterns or
- **Example Walkthrough**

mathematical formulas involved, just the straightforward application of the definition of a reflection over a vertical line.

Let's consider an example with the following set of points on a 2D plane: [(1, 1), (3, 1), (7, 1), (9, 1)]. Our goal is to determine if there exists a line parallel to the y-axis that reflects all these points symmetrically.

# the points.

Now we traverse our points list:

First, we initialize min\_x to infinity and max\_x to negative infinity to keep track of the minimum and maximum x-coordinates of

 For point (7, 1), min\_x remains 1 and max\_x is updated to 7. For the last point (9, 1), min\_x remains 1 and max\_x is updated to 9. Throughout the iteration, we also add each point to a set point\_set for efficient lookups.

After the first traversal, we have  $min_x = 1$  and  $max_x = 9$ . The sum s which represents twice the possible line of

 $\circ$  Point (1, 1) has a counterpart (9, 1) because s - x = 10 - 1 = 9, and (9, 1) is in point\_set.  $\circ$  Point (3, 1) should be mirrored with point (7, 1) because s - x = 10 - 3 = 7, and (7, 1) is indeed in point\_set.

min x, max x = float('inf'), float('-inf')

minX = Math.min(minX, point[0]);

maxX = Math.max(maxX, point[0]);

// of the line of reflection.

set<pair<int, int>> pointSet;

for (auto& point : points) {

for (auto& point : points) {

return false;

import { Set } from "typescript-collections";

function isReflected(points: number[][]): boolean {

const sumOfMinAndMaxX: number = minX + maxX;

for (const point of points) {

Time and Space Complexity

**Time Complexity** 

**Space Complexity** 

return true;

**}**;

**TypeScript** 

int sumOfMinAndMaxX = minX + maxX;

pointSet.add(List.of(point[0], point[1]));

// Calculate the sum of minX and maxX which is twice the X coordinate

// A set to store unique points (pair of X and Y coordinates).

// Loop through all points to populate pointSet and find the minX and maxX.

// Iterate through all points to check if the reflected point exists.

int reflectedX = sumOfMinAndMaxX - point[0];

if (!pointSet.count({reflectedX, point[1]})) {

// Importing arrays and sets from ECMAScript 6 (ES6) standard.

// Return true if all points have their reflected counterparts.

// Function to determine if a set of points reflects across a single vertical axis.

minX = min(minX, point[0]); // Update minX if current point's X is smaller.

 $\max X = \max(\max X, point[0]); // Update \max X if current point's X is larger.$ 

// Calculate the reflection of the current point across the reflection axis.

// Sum of minX and maxX is the total which, when halved, gives the X-coordinate of the reflection axis.

// Check if the reflected point is not in the set. If it's not, the point array is not reflected.

pointSet.insert({point[0], point[1]}); // Insert the point into the set.

point\_set = set() # Create a set to store unique points.

# Iterate over all points to find the min and max X values and add points to the set.

For the symmetry check, we traverse the list again:

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    Since points (7, 1) and (9, 1) have already been paired with their counterparts, the conditions are satisfied for all points.

All points have a mirrored counterpart with respect to the line whose x-coordinate is $/2, which in this case is 5. Since each
point (x, y) has a matching point (s - x, y) in the set, we can conclude that it's possible to draw a vertical line parallel to the
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y-axis at x = 5 that reflects all the points symmetrically. Therefore, the function should return true for this example.

**Python** class Solution: def isReflected(self, points: List[List[int]]) -> bool: # Initialize minimum and maximum X to positive and negative infinity respectively.

#### for x, v in points: min x = min(min x, x)max x = max(max x, x)

point\_set.add((x, y))

Solution Implementation

```
# Calculate the sum of min and max X, which should be equal to twice the X value of the reflection line.
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```
reflection_sum = min_x + max_x
        # Check if for each point (x, v), the reflected point across the Y-axis
        # given by (reflection sum -x, y) exists in the point set.
        # The reflection across the Y-axis is defined by the line X = (min \ x + max_x) / 2.
        return all((reflection_sum - x, y) in point_set for x, y in points)
Java
class Solution {
    public boolean isReflected(int[][] points) {
        // Initialize the max and min X coordinates to extreme values.
        final int MAX VALUE = Integer.MAX_VALUE;
        int minX = MAX VALUE;
        int maxX = Integer.MIN_VALUE;
        // Using a set to store unique point representations.
        Set<List<Integer>> pointSet = new HashSet<>();
        // Iterate over all points to find the minX and maxX values,
        // and add each point to the set.
        for (int[] point : points) {
```

```
int sum = minX + maxX;
        // Check if each point has its reflected counterpart in the set.
        for (int[] point : points) {
            // If the reflected point is not found in the set, return false.
            if (!pointSet.contains(List.of(sum - point[0], point[1]))) {
                return false;
       // If all points have their reflected counterpart, return true.
        return true;
C++
#include <vector>
#include <set>
#include <algorithm>
using std::vector;
using std::set;
using std::min;
using std::max;
using std::pair;
class Solution {
public:
    // Function to determine if a set of points reflects across a single vertical axis.
    bool isReflected(vector<vector<int>>& points) {
        // Constants to represent infinity and negative infinity.
        const int INF = 1 << 30;
        // Variables to store the minimum and maximum X-coordinates.
        int minX = INF, maxX = -INF;
```

### // Constants to represent infinity and negative infinity. const INF: number = 1 << 30;</pre> // Variables to store the minimum and maximum X-coordinates. let minX: number = INF; let maxX: number = -INF; // A set to store unique points (tuple of X and Y coordinates). let pointSet: Set<string> = new Set(); // Loop through all points to populate pointSet and find the minX and maxX. for (const point of points) { minX = Math.min(minX, point[0]); // Update minX if current point's X is smaller.

// Sum of minX and maxX is the total, which, when halved, gives the X-coordinate of the reflection axis.

maxX = Math.max(maxX, point[0]); // Update maxX if current point's X is larger.

// Insert the point into the set as a string to ensure uniqueness.

pointSet.add(JSON.stringify({x: point[0], y: point[1]}));

// Iterate through all points to check if the reflected point exists.

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const reflectedX: number = sumOfMinAndMaxX - point[0];
// Check if the reflected point is not in the set. If it's not, the point array is not reflected.
if (!pointSet.contains(JSON.stringify({x: reflectedX, y: point[1]}))) {
    return false;
```

// Calculate the reflection of the current point across the reflection axis.

```
// Return true if all points have their reflected counterparts.
   return true;
class Solution:
   def isReflected(self, points: List[List[int]]) -> bool:
       # Initialize minimum and maximum X to positive and negative infinity respectively.
       min x, max x = float('inf'), float('-inf')
       point_set = set() # Create a set to store unique points.
       # Iterate over all points to find the min and max X values and add points to the set.
       for x, v in points:
           min x = min(min x, x)
           max x = max(max x, x)
           point_set.add((x, y))
       # Calculate the sum of min and max X, which should be equal to twice the X value of the reflection line.
       reflection sum = min x + max x
       # Check if for each point (x, y), the reflected point across the Y-axis
       # given by (reflection sum -x, y) exists in the point set.
       # The reflection across the Y-axis is defined by the line X = (min \ x + max \ x) / 2.
       return all((reflection sum - x, y) in point set for x, y in points)
```

Then we check if each point has a reflected point in the point\_set using the formula (s - x, y) where s is the sum of min\_x and max\_x. This step also iterates over all n points and each lookup in a set is expected to be 0(1) on average,

leading to another O(n) time operation. Therefore, the overall time complexity is O(n) + O(n) = O(n) where n is the number of input points.

We iterate over all n points once to find the minimum  $(min_x)$  and maximum  $(max_x)$  x-coordinates, and to add each point to

We create a set, point\_set, which in the worst case, will contain all input points if all points are unique. This leads to a space complexity of O(n).

The space complexity of the code can be considered as follows:

The time complexity of the given code can be analyzed as follows:

the point\_set. This requires O(n) time.

No additional significant space is used in the algorithm. Hence, the overall space complexity of the algorithm is O(n) where n is the number of input points.