624. Maximum Distance in Arrays

Problem Description

Medium Greedy Array

In this problem, we are provided m separate arrays, each being sorted in ascending order. Our task is to find the maximum distance between any two integers where each integer is taken from a different array. The distance is defined as the absolute difference between those two integers. That is, if we take one integer a from one array and another integer b from a different array, the distance is |a - b|. We need to compute and return the maximum such distance possible.

Intuition

To solve this problem efficiently, we leverage the fact that each of the marrays is sorted in ascending order. Given this property, the smallest element of each array will be at the 0-th index and the largest element at the last index.

largest distance by taking the smallest element from one array and the largest from another array. So to maximize the distance, we always consider the minimum element from one array and the maximum element from the other. The approach is to iterate through the arrays while tracking the smallest and largest elements we have seen so far. For each new

To find the maximum distance, we need to make the difference as large as possible. This means we should consider the potential

array, we consider the distance between the current array's smallest element and the maximum element seen so far, and vice versa - the distance between the current array's largest element and the smallest element seen so far. The answer is the largest of all these distances. Here is the step-by-step reasoning:

3. Iterate through the remaining arrays starting from the second one. For each array:

Calculate the absolute difference between the current array's smallest element and mx (the largest element seen so far).

1. Initialize ans to 0, which will keep track of the maximum distance.

 Calculate the absolute difference between the current array's largest element and mi (the smallest element seen so far). Update ans with the largest value between ans, the first calculated difference, and the second calculated difference.

2. Take the smallest (mi) and largest (mx) elements from the first array to initiate tracking.

 Update mi with the smallest value between mi (the smallest value found so far) and the smallest element of the current array. Update mx with the largest value between mx (the largest value found so far) and the largest element of the current array.

4. After finishing the iteration, return ans as the maximum distance found.

- By only comparing the extreme ends of each array, we ensure that we are always considering the likely pairs to give us the maximum distance while maintaining a linear time complexity.
- The solution follows a greedy strategy by keeping track of the smallest and the largest element found so far across all arrays. Here is how the implementation reflects the approach:

• An initial answer ans is set to 0, which will hold the maximum distance. • The smallest element mi and the largest element mx from the first array are used to initialize the tracking.

Solution Approach

 Then, for every subsequent array, the implementation performs the following steps: 1. Calculates two distances: the distance between the smallest element of the current array and the largest element mx seen so far (a = abs(arr[0] - mx)), and the distance between the largest element of the current array and the smallest element mi seen so far (b =

abs(arr[-1] - mi). 2. The answer ans is updated to be the maximum of ans, a, and b.

largest element found up to the current point in iteration.

statements and math operations.

ans = max(ans, a, b)

mi = min(mi, arr[0])

- 3. The current array's smallest element is considered to update mi using mi = min(mi, arr[0]) if it is smaller than the current mi.
- This approach guarantees each array is only visited once, making the overall time complexity 0(m), where m is the number of arrays. There is no use of additional data structures, keeping space complexity to 0(1) as it only uses a few variables for tracking the minimum, maximum, and the current answer.

4. Similarly, the current array's largest element is used to update mx using mx = max(mx, arr[-1]) if it is larger than the current mx.

By using min() and max() functions, the elements mi and mx are continuously updated to always represent the smallest and

Here's a snippet of the main logic using the algorithm described: class Solution: def maxDistance(self, arrays: List[List[int]]) -> int: ans = 0mi, mx = arrays[0][0], arrays[0][-1]for arr in arrays[1:]:

The code directly follows this strategy and does not use any complex data structures or algorithms beyond basic conditional

mx = max(mx, arr[-1])return ans

```
Let's say we have the following 3 sorted arrays:
• Array 1: [1, 2, 3]
• Array 2: [4, 5]
• Array 3: [1, 5, 9]
Our goal is to find the maximum distance between any two integers from two different arrays.
```

two integers from two different arrays. Thus, we return 8.

def maxDistance(self, arrays: List[List[int]]) -> int:

min_value = min(min_value, array[0])

min = Math.min(min, array.get(0));

// Return the largest distance found.

// Start from the second array

for (int i = 1; i < arrays.size(); ++i) {</pre>

auto& currentArray = arrays[i];

// Reference to the current array to avoid copying

int distanceToMin = abs(currentArray[0] - maxElement);

int distanceToMax = abs(currentArray.back() - minElement);

maxDistance = max({maxDistance, distanceToMin, distanceToMax});

max = Math.max(max, array.get(array.size() - 1));

max_value = max(max_value, array[-1])

Return the computed maximum distance between any two pairs

min_value, max_value = arrays[0][0], arrays[0][-1]

Iterate over the arrays starting from the second one

Initialize the maximum distance and the minimum and maximum values

Compute the potential distances between the current array's first or last element

max_distance = max(max_distance, distance_min_to_max, distance_max_to_min)

Update the overall min_value and max_value with the current array's values

1. Initialize ans to 0, which will be used to keep track of the maximum distance found so far.

a, b = abs(arr[0] - mx), abs(arr[-1] - mi)

is 3.

Example Walkthrough

Next, we go to the third array [1, 5, 9]. We calculate the distance between 1 (the smallest element of this array) and the

mi or mx. Since 1 is equal to current mi and 9 is greater than current mx, we update mx to 9.

Now, we iterate through the remaining arrays to calculate potential distances and update mi and mx.

update mi and mx if necessary. Since 4 is greater than mi and 5 is greater than mx, we only update mx to 5.

current largest mx, which is 5. The distance is abs(1 - 5) = 4. Then, we calculate the distance between 9 (the largest element of this array) and the current smallest mi, which is 1. The distance here is abs(9 - 1) = 8. We update ans to the maximum of ans and these new distances. ans is now updated to 8, which is the maximum distance found. We check if we need to update

Finally, after iterating through all arrays, we have the largest distance ans as 8, which is the correct maximum distance between

2. We start with the first array and take its smallest and largest elements as the initial values for mi and mx respectively. In this case, mi is 1 and mx

Moving to the second array [4, 5], we calculate the distance between the smallest element in this array 4 and the current mx

3 which is abs (4 - 3) = 1. We also calculate the distance between the largest element in this array 5 and the current mi 1

which is abs(5 - 1) = 4. We update ans to the maximum of ans, which is 0, and these new distances, ans is now 4. We also

Solution Implementation **Python** from typing import List class Solution:

with the opposite ends of the known range (max_value and min_value) distance_min_to_max = abs(array[0] - max_value) distance_max_to_min = abs(array[-1] - min_value) # Update max_distance to be the maximum of the current max_distance and the new distances

max_distance = 0

for array in arrays[1:]:

return max_distance

```
Java
class Solution {
    // Method to find the maximum distance between any pair of elements
    // from different arrays within the list of arrays.
    public int maxDistance(List<List<Integer>> arrays) {
       // Initializing 'ans' as 0 to hold the maximum distance encountered.
        int maxDistance = 0;
       // Assign the first element of the first array as the minimum known value 'min'.
       int min = arrays.get(0).get(0);
       // Assign the last element of the first array as the maximum known value 'max'.
       int max = arrays.get(0).get(arrays.get(0).size() - 1);
       // Iterate over the remaining arrays in the list starting from the second array.
        for (int i = 1; i < arrays.size(); ++i) {</pre>
            List<Integer> array = arrays.get(i);
           // Calculate the absolute difference between the first element of the current array
           // and the known maximum. This represents a potential max distance.
            int distanceWithMax = Math.abs(array.get(0) - max);
           // Calculate the absolute difference between the last element of the current array
           // and the known minimum. This represents a potential max distance.
            int distanceWithMin = Math.abs(array.get(array.size() - 1) - min);
            // Update 'maxDistance' with the greater of the two newly calculated distances
            // if either is larger than the current 'maxDistance'.
```

```
return maxDistance;
#include <vector>
#include <algorithm> // include the algorithm header for using min, max functions
class Solution {
public:
    int maxDistance(vector<vector<int>>& arrays) {
        int maxDistance = 0; // maximum distance found so far
        int minElement = arrays[0][0]; // initialize with the first element of the first array
        int maxElement = arrays[0].back(); // initialize with the last element of the first array
```

maxDistance = Math.max(maxDistance, Math.max(distanceWithMax, distanceWithMin));

// Update the known minimum value 'min' if the first item of the current array is smaller.

// Update the known maximum value 'max' if the last item of the current array is larger.

```
// Update minElement with the smaller between minElement and the first element of the current array
           minElement = min(minElement, currentArray[0]);
            // Update maxElement with the larger between maxElement and the last element of the current array
           maxElement = max(maxElement, currentArray.back());
       // Return the maximum distance found
       return maxDistance;
};
TypeScript
// Importing the first and last functions from lodash for array manipulation
import { first, last } from 'lodash';
// Function to calculate the maximum distance between any pair of elements from different arrays
function maxDistance(arrays: number[][]): number {
    // Maximum distance found so far
    let maxDistance = 0;
    // Initialize with the first element of the first sub-array
    let minElement = first(arrays[0]) as number;
    // Initialize with the last element of the first sub-array
    let maxElement = last(arrays[0]) as number;
    // Start from the second sub-array
    for (let i = 1; i < arrays.length; ++i) {</pre>
       // Reference to the current sub-array to avoid copying
       const currentArray = arrays[i];
       // Calculate distance between the smallest element so far and the last element of the current sub-array
       const distanceToMin = Math.abs(first(currentArray) as number - maxElement);
```

// Calculate distance between the largest element so far and the first element of the current sub-array

// Update minElement with the smaller between minElement and the first element of the current sub-array

// Update maxElement with the larger between maxElement and the last element of the current sub—array

// Calculate distance between the smallest element so far and the last element of the current array

// Calculate distance between the largest element so far and the first element of the current array

// Update maxDistance with the largest of the three: itself, distanceToMin, distanceToMax

```
maxElement = Math.max(maxElement, last(currentArray) as number);
      // Return the maximum distance found
      return maxDistance;
from typing import List
```

// Update maxDistance with the largest of the three: itself, distanceToMin, distanceToMax

const distanceToMax = Math.abs(last(currentArray) as number - minElement);

maxDistance = Math.max(maxDistance, distanceToMin, distanceToMax);

minElement = Math.min(minElement, first(currentArray) as number);

Initialize the maximum distance and the minimum and maximum values

Return the computed maximum distance between any two pairs

def maxDistance(self, arrays: List[List[int]]) -> int:

```
max distance = 0
min_value, max_value = arrays[0][0], arrays[0][-1]
# Iterate over the arrays starting from the second one
for array in arrays[1:]:
    # Compute the potential distances between the current array's first or last element
    # with the opposite ends of the known range (max_value and min_value)
    distance_min_to_max = abs(array[0] - max_value)
    distance_max_to_min = abs(array[-1] - min_value)
    # Update max_distance to be the maximum of the current max_distance and the new distances
    max_distance = max(max_distance, distance_min_to_max, distance_max_to_min)
    # Update the overall min_value and max_value with the current array's values
    min_value = min(min_value, array[0])
    max_value = max(max_value, array[-1])
```

The time complexity of the given code is primarily determined by the loop that iterates over all elements of the arrays list, except the first element. In the worst case, where there are n arrays inside the arrays list, we iterate n - 1 times. Inside the loop, we

return max_distance

Time and Space Complexity

execute a constant number of operations for each array: calculating the absolute difference between the first and last elements with mi and mx, updating ans, and updating mi and mx with the current array's first and last elements, respectively. Since all these

Time Complexity

class Solution:

operations inside the loop have a constant time complexity, the overall time complexity of the loop is O(n). Hence, the total time complexity of the code is O(n). **Space Complexity** The space complexity is determined by the additional memory used by the program which is not part of the input. In the given

code, we use a fixed number of variables (ans, mi, mx, a, b) that do not depend on the size of the input. No other dynamic data structures or recursive calls which could use additional space are involved. Therefore, the space complexity is 0(1), which represents constant space usage.