213. House Robber II Medium **Dynamic Programming Array**

Problem Description

and each house contains a certain amount of money that can be stolen. However, if the robber tries to rob two adjacent houses, an alarm will be triggered and the police will be notified. The objective is to determine the maximum amount of money that can be stolen without triggering any alarms. Constraints include:

In this problem, we imagine a situation where a professional thief is planning to rob houses. The houses are laid out in a circle,

One cannot rob two adjacent houses on the same night.

• The arrangement of the houses in a circle, which means that the first and last houses are considered adjacent.

Given the circular arrangement, we need to find a way to calculate the maximum loot without robbing adjacent houses and also

constraint imposed by the circular arrangement of the houses.

accounting for the circular layout.

The problem can be seen as a variation of the classic dynamic programming problem "House Robber", with the additional

The intuition for solving this problem is to use dynamic programming to make decisions at each step. We need to figure out, with each house, whether it's more profitable to rob it or skip it. However, the circular arrangement makes it tricky; we cannot simply

start from one end and go to the other, as we may miss out on the optimal solution that involves wrapping around the end. To resolve this issue with circular arrangement, we make use of the following observations: • If we rob the first house, we cannot rob the last house due to the circular layout.

• If we ignore the first house, we can treat the problem as a non-circular arrangement and apply standard dynamic programming as in the original

"House Robber" problem.

So our approach is:

considering both possibilities – starting from the first house or skipping it.

1. Calculate the maximum money that can be robbed bypassing the first house (making the problem linear instead of circular). 2. Calculate the maximum money that can be robbed if we skip the last house for the same reason.

After calculating these two scenarios, the maximum of both will give us the solution to the original circular problem. The function

considering the previous decisions. The rob function then calls the helper function <u>rob</u> twice – once for the array of houses excluding the first house and once for the array excluding the last house – and returns the maximum of the two results. This accommodates the circular constraint by

<u>rob</u> implements this approach of robbing or skipping a house. It iterates through the given list of house values and, at each

step, decides whether to rob the current house or not. The choice is made based on the maximum money that can be gained

In summary, the solution splits the circular problem into two linear sub-problems and applies dynamic programming to each. The maximum of the solutions to these sub-problems is the answer to the original problem. **Solution Approach**

The solution provided uses dynamic programming, a method for solving complex problems by breaking them down into simpler sub-problems. It is particularly well-suited for optimization problems like ours, where we seek the maximum amount of money that can be robbed.

<u>Dynamic programming</u> often involves creating an array to store the sub-problems' solutions; however, in this implementation, only

two variables are maintained, reducing the space complexity. The algorithm employs a "rob-or-skip" strategy, where we look at

each house and decide whether to rob it or skip it based on which action would yield a higher profit. Here's how the variables f

and g correspond to this strategy:

maximum amount robbed up to the previous house where we did not rob, to ensure no two adjacent houses are robbed). The main function rob distinguishes between two cases due to the circular layout of houses. If there's only one house, it simply returns its value, as there's no constraint to consider. In other cases, the function utilizes a helper function _rob. The _rob

• g: The maximum amount of money that can be robbed including the current house (which is the sum of the money in the current house and the

f: The maximum amount of money that can be robbed up to the current house without robbing it (essentially, the skip option).

house and one excluding the last house. Let's understand how _rob works: 1. Initialize f, g to 0. These are the maximum amounts with the aforementioned conditions. 2. Loop through each value x in the nums list, where x represents the amount of money in the current house:

Calculate the new value of f as the maximum between the old f and g. This choice represents skipping the current house, so we take the

• Update g to be the sum of the old f (which represents the total money robbed up to the previous house without robbing it) and x (the

3. After the loop, the function returns the maximum amount between f and g which represents the maximum money the robber can achieve by

function implements the dynamic programming approach and is applied to two slices of the original list: one excluding the first

either robbing or not robbing the last house in the list. Back in the rob function, we then return the maximum value obtained from calling _rob with either of the list slices, thus

The thief has the following options:

Start with f and g both at 0.

Solution Implementation

prev max = curr max = 0

for value in sub nums:

return curr_max

return nums[0]

public int rob(int[] nums) {

if (houseCount == 1) {

return nums[0];

int houseCount = nums.length;

if len(nums) == 1:

Python

best previous outcome.

concluding the approach and ensuring we consider the circular property of the problem. The time complexity of the solution is O(n), where n is the number of houses, because it processes each house exactly once. The

amount in the current house). This represents robbing the current house.

- space complexity is O(1), because it uses a constant amount of extra space, despite the input size.
- **Example Walkthrough** Let's use a small example to illustrate the solution approach. Suppose we have a circular arrangement of houses with the following amounts of money: [2, 3, 2].

1. Rob the first house (with 2), and then can only rob the third house (with 2), totaling 4. 2. Skip the first house, rob the second house (with 3), and then they have to skip the third house because they can't rob adjacent houses.

If we ignore the first house, we treat the remaining houses [3, 2] as a linear problem. Using our dynamic programming

Start with f and g both at 0.

approach:

the current house's value, which is 0 + 2. Now f is 3 and g is 2. • The maximum of f and g is 3, which is the most we can rob from houses [3, 2].

• For the second house (with 2), we update f to the maximum of previous f and g, which is 3. We then set g to the previous f value plus

If we ignore the last house, we treat the remaining houses [2, 3] as a linear problem. Again, we apply dynamic programming:

o For the second house (with 3), we update f to the maximum of previous f and g, which is 2. We then set g to the previous f value plus

After comparing the two scenarios, we see the maximum money the thief can rob without triggering any alarms is 3. Thus, the

The maximum of f and g is 3, the most we can rob from houses [2, 3].

Applying the two scenarios from our solution approach:

robber robs the second house with 3 in both cases, and the solution to the problem is 3. This demonstrates how the circular constraint is handled by considering the problem in a linear fashion for two slices of the array.

∘ For the first house (with 2), we update g to 0 + 2. Now f is 0 and g is 2.

the current house's value, which is 0 + 3. Now f is 2 and g is 3.

Calculate the maximum amount of money that can be

Update the prev max and curr max to include

the current house in the robbery plan or not.

prev max, curr max = curr max, max(curr max, prev max + value)

The current maximum will be the solution for this subsequence.

// The main function to compute the maximum amount of money that can be robbed.

robbed from a subsequence of houses (sub_nums).

If there's only one house, return its value.

private int robMaxMoney(int[] nums, int left, int right) {

// Iterate through the range of houses.

int inclCurr = exclPrev + nums[i];

exclPrev = Math.max(inclPrev, exclPrev);

// Edge case: If there's only one house, return its value.

const robRange = (start: number, end: number): number => {

for (; start <= end; ++start) {</pre>

robPrevious = robNext;

prev max = curr max = 0

for value in sub nums:

1. Excluding the first house.

2. Excluding the last house.

return curr_max

return nums[0]

Time and Space Complexity

if len(nums) == 1:

return Math.max(robPrevious, robCurrent);

// Helper function to rob a range of houses from index l to r, inclusive.

// Iterate from the start index to the end index, inclusive.

// Return the maximum amount robbed between the last two houses.

return Math.max(robRange(0, numLength - 2), robRange(1, numLength - 1));

Calculate the maximum amount of money that can be

Update the prev max and curr max to include

the current house in the robbery plan or not.

prev max, curr max = curr max, max(curr max, prev max + value)

The current maximum will be the solution for this subsequence.

return max(_rob_subsequence(nums[1:]), _rob_subsequence(nums[:-1]))

robbed from a subsequence of houses (sub_nums).

If there's only one house, return its value.

Rob the houses by considering two scenarios:

Take the maximum from these two options.

// or from the second house to the last one since we cannot rob adjacent houses.

if (numLength === 1) {

return nums[0];

for (int i = left; i <= right; ++i) {</pre>

inclPrev = inclCurr;

∘ For the first house (with 3), we update g to 0 + 3 (since f is still 0). Now f is 0 and g is 3.

class Solution: def rob(self, nums: List[int]) -> int: def rob_subsequence(sub_nums):

Rob the houses by considering two scenarios: # 1. Excluding the first house. # 2. Excluding the last house. # Take the maximum from these two options. return max(_rob_subsequence(nums[1:]), _rob_subsequence(nums[:-1]))

// Calculate the new maximum excluding the current house by taking the maximum between inclPrev and exclPrev.

// Otherwise, find the maximum of two scenarios: excluding the first house or excluding the last house.

return Math.max(robMaxMoney(nums, 0, houseCount -2), robMaxMoney(nums, 1, houseCount -1));

// Helper function to calculate the max money that can be robbed in a specific range of houses.

int inclPrev = 0; // This will hold the maximum amount including the previous house.

int exclPrev = 0; // This will hold the maximum amount excluding the previous house.

// Initialize two variables to keep track of the two previous maximum values.

// Calculate the new maximum including the current house.

// Update inclPrev to be used in the next iteration.

```
// If there's only one house, the maximum money is what's in that house.
```

Java

class Solution {

```
// Return the maximum money that can be robbed within the specified range.
        // Must consider the final value of inclPrev and exclPrev, since one of them will have the maximum.
        return Math.max(inclPrev, exclPrev);
C++
class Solution {
public:
    // Main function to calculate the maximum amount of money that can be robbed.
    int rob(vector<int>& nums) {
        int n = nums.size();
        // If there's only one house, return its value.
        if (n == 1) {
            return nums[0];
        // Compare and return the maximum amount between robbing the first house or the second one.
        return max(robRange(nums, 0, n - 2), robRange(nums, 1, n - 1));
    // Helper function to calculate max amount from a range of houses.
    int robRange(vector<int>& nums, int left, int right) {
        int inclusive = 0; // This will store the max amount including the current house.
        int exclusive = 0; // This will store the max amount excluding the current house.
        // Iterate from the left to the right indices of the house range.
        for (; left <= right; ++left) {</pre>
            // Compute the new exclusive amount, which is the max of the previous inclusive and exclusive amounts.
            int newExclusive = max(inclusive, exclusive);
            // Update inclusive to be the sum of the old exclusive and current house value.
            inclusive = exclusive + nums[left];
            // Update exclusive to the newly computed value.
            exclusive = newExclusive;
        // Return the maximum of inclusive and exclusive amounts as the result.
        return max(inclusive, exclusive);
};
TypeScript
function rob(nums: number[]): number {
    // Determine the length of the array.
    const numLength = nums.length;
```

let robPrevious = 0; // Initialize to store the maximum amount that can be robbed up to the previous house

// The maximum amount robbed will be the maximum of either robbing from the first house to the second-to-last house,

let robNext = Math.max(robPrevious, robCurrent); // The next 'previous' is the max of current and previous robbed amounts

robCurrent = robPrevious + nums[start]; // Update the current to the sum of previous robbed amount and current h

// Move the 'next' value to 'previous' for the next iteration.

let robCurrent = 0; // Initialize to store the maximum amount that can be robbed up to the current house

class Solution: def rob(self, nums: List[int]) -> int: def rob_subsequence(sub_nums):

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The given Python code defines a method rob that solves the house robber problem for a neighborhood arranged in a circle, meaning the first and last houses cannot be robbed together. **Time Complexity**

once and performs a constant time operation of $\max(f, g)$ and an addition operation f + x for each element.

The rob function calls the _rob function twice: once with nums[1:] and once with nums[:-1]. Since slicing a list is O(n) and the <u>rob</u> function is also O(n), the total time complexity for each call is O(2n). Because the two calls do not depend on each other,

The time complexity of the helper function <u>rob</u> is O(n), where n is the number of elements in nums. It iterates through the list

Final time complexity: 0(n) **Space Complexity**

However, the main function rob creates two slices of the nums input list: nums[1:] and nums[:-1]. Each of these slices is of size

the overall time complexity of rob remains O(n).

n-1, where n is the length of the original nums. The space taken up by these lists is significant. Final space complexity: O(n), accounting for the space used by the list slices.

The helper function <u>rob</u> uses O(1) extra space, as it only keeps track of the two variables f and g.