3038. Maximum Number of Operations With the Same Score I

Simulation **Easy** Array

In this problem, we have an array of integers called nums. We are allowed to remove pairs of elements from the start of the array.

Problem Description

perform as many operations as possible with the condition that all operations must yield the same score. This means that each pair of elements we remove must have the same sum as the first pair we removed. To find the maximum number of operations we can perform under these conditions, we need to continuously check the first two elements of the remaining array, removing them only if they match the desired score. If at any point the sum of the two elements

Each time we remove a pair, we calculate the sum of the two elements and consider this the score of the operation. The goal is to

doesn't match the score, we must stop the operations, and that's the maximum number we can perform. Intuition

To solve this problem, the solution begins by presuming that the score we need every operation to match is the sum of the first

two elements in the array, denoted as s. The solution then proceeds to traverse the array in steps of 2, meaning it looks at every consecutive pair of elements.

During this traversal, the solution checks if the sum of each pair matches the initial score s. As soon as it encounters a pair that does not have a sum equal to s, or if there are no more pairs to check (e.g., after removing pairs the array has only one element left), the traversal stops. This is grounded in the rule that all operations must yield the same score, and since the score is defined by the first operation, any pair that doesn't match this score would break the condition.

Solution Approach The implemented solution follows a straightforward and direct approach. It doesn't employ complex algorithms or data structures but relies on a simple traversal of the array using a for loop.

The solution initializes a variable s with the sum of the first two elements in nums. This s will be the score that every

array.

break

for i in range(0, n, 2):

Here's a step-by-step explanation of the implementation:

if i + 1 == n or nums[i] + nums[i + 1] != s:

Let's illustrate the solution approach with a simple example.

Suppose we have the following array of integers as our nums:

ans, n = 0, len(nums) # ans = 0, n = 6

nums[0] + nums[1] = 10, which is equal to s.

nums[2] + nums[3] = 10, which is equal to s.

remove or the condition of maintaining the same score can no longer be met.

operation must match. s = nums[0] + nums[1]

It initializes an answer variable, ans, to count the number of operations performed and a variable n to store the length of the

ans, n = 0, len(nums)

The solution uses a for loop to iterate over the elements of the array in steps of two, looking at every pair of elements.

For each pair, it checks two conditions: whether the current index i + 1 equals the length of the array (which would mean there's no pair left to process) and whether the sum of the current pair does not match the score s.

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If none of the conditions for breaking the loop are met, it means that the current pair matches the score, and an operation can
 be performed. The ans variable is incremented to reflect this.
ans += 1
```

If either condition is true, the loop breaks, effectively stopping any further operations, as either there are no more pairs to

same score. return ans

This solution assumes that the input array nums has already been set up such that every valid operation (pair with matching sum

The loop continues until it has either checked all elements in the array or encountered a break. At the end of the loop, the

solution returns the count ans, which represents the maximum number of operations that have been performed with the

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s) is adjacent and that there are no possible operations (pairs with the desired sum) after encountering the first invalid one.
Therefore, it's essential to confirm that the input array will conform to these conditions; otherwise, the solution might not
correctly calculate the maximum number of operations.
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We will walk through the steps of the implementation using this array. First, we initialize the variable s with the sum of the first two elements in nums, which is 4 + 6 = 10. s = nums[0] + nums[1] # s = 10

We initialize our answer variable, ans, to 0 and a variable n to store the length of the array, which is 6 in this case.

The for loop begins and will iterate over the pairs of elements.

for i in range(0, n, 2):

First iteration:

ans += 1 # ans = 1

Second iteration:

ans += 1 # ans = 2

Third iteration:

ans += 1 # ans = 3

Python

Java

class Solution:

from typing import List

if len(nums) < 2:

return 0

operations count = 0

break

return operations_count

total_numbers = len(nums)

def maxOperations(self, nums: List[int]) -> int:

sum of pair = nums[0] + nums[1]

operations_count += 1

Return the total pairs formed.

public int maxOperations(int[] nums) {

int targetSum = nums[0] + nums[1];

for (int i = 0: i + 1 < n: i += 2) {

++operationsCount;

++operations_count;

function maxOperations(nums: number[]): number {

// Iterate over the 'nums' array with a step of 2.

if (nums[i] + nums[i + 1] === requiredSum) {

doesn't scale with the size of the input array nums.

const requiredSum = nums[0] + nums[1];

for (let i = 0; i + 1 < n; i += 2) {

operationsCount++;

const n = nums.length;

let operationsCount = 0;

return operations_count;

// Return the total number of valid operations

// Initialize a sum 'requiredSum' with the sum of the first two elements.

// 'operationsCount' will hold the number of valid operations performed.

// Check if the current and next element sum up to the 'requiredSum'.

// Increment the count of valid operations if the condition is met.

// 'n' represents the total number of elements in the 'nums' array.

for i in range(0, total numbers, 2):

Check if there are enough numbers to form pairs.

Initialize the sum of the first pair and the answer counter.

If we found a matching pair, increment the counter.

// Calculate the sum of the first pair of elements in the array

// Check if the current pair sums up to the target sum

// If it does, increment the operations counter

if (nums[i] + nums[i + 1] == targetSum) {

// Initialize a counter for the number of valid operations

if i + 1 == total_numbers or nums[i] + nums[i + 1] != sum_of_pair:

// Loop through the array in pairs, up to the second-to-last element (i + 1 < n)

Iterate through the list in steps of 2 to form pairs.

return ans # returns 3

well. We increment ans.

Example Walkthrough

nums = [4, 6, 4, 6, 8, 2]

At first, i = 0, and we look at the first pair nums [0] + nums [1] which is 4 + 6. This equals s (10), so we can remove this pair. No break is encountered, and we increment ans.

nums[4] + nums[5] = 10, which is equal to s. However, this is the last pair.

```
Finally, i = 4, and we look at the last pair nums[4] + nums[5] which is 8 + 2. This equals 10, the same as s. Even though it
matches s, we're at the end of the array, so the loop naturally concludes.
```

The loop finishes because there are no more elements in nums. The solution will return the count ans, which is 3 in this case.

Now i = 2, and we look at the next pair nums[2] + nums[3] which is 4 + 6. This also equals s = (10), so we remove this pair as

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in the array.
Solution Implementation
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Check whether we reached the end or if the sum of the current pair doesn't match sum_of_pair.

This example shows that the array $\frac{10}{100}$ allowed for 3 pairs to be removed with the sum matching $\frac{10}{100}$, hence 3 operations

were performed. The example follows the method of sequential checking and early termination if the condition is not met. In this

instance, the conditions were satisfied by all element pairs, so the maximum number of operations is equal to the number of pairs

int operationsCount = 0; // Length of the nums array int n = nums.length;

} else {

public class Solution {

// If a pair doesn't sum up to the target sum, break out of the loop // No need to continue as subsequent operations will not be valid break; // Return the total number of valid operations return operationsCount; C++ #include <vector> using namespace std; class Solution { public: // Function to determine the maximum number of operations where each operation consists of // finding a pair of elements from the array that add up to the same value int maxOperations(vector<int>& nums) { // Initialize the sum "s" with the sum of the first two elements int target sum = nums[0] + nums[1]; // Initialize the answer, which stores the number of valid operations int operations count = 0; // Get the size of the nums array int n = nums.size(); // Iterate over the elements of the vector in pairs for (int i = 0; i + 1 < n & nums[i] + nums[i + 1] == target sum; <math>i += 2) { // If the sum of the current pair of elements equals the target sum, increment the operations count

} else {

};

TypeScript

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// If the condition is not met, break the loop as it's assumed that nums were initially arranged in pairs with the same s
            break;
    // Return the total number of operations performed.
    return operationsCount;
from typing import List
class Solution:
    def maxOperations(self, nums: List[int]) -> int:
        # Check if there are enough numbers to form pairs.
        if len(nums) < 2:</pre>
            return 0
        # Initialize the sum of the first pair and the answer counter.
        sum of pair = nums[0] + nums[1]
        operations count = 0
        total_numbers = len(nums)
        # Iterate through the list in steps of 2 to form pairs.
        for i in range(0, total numbers, 2):
           # Check whether we reached the end or if the sum of the current pair doesn't match sum_of_pair.
            if i + 1 == total_numbers or nums[i] + nums[i + 1] != sum_of_pair:
                break
           # If we found a matching pair, increment the counter.
            operations_count += 1
       # Return the total pairs formed.
        return operations_count
Time and Space Complexity
  The time complexity of the code is 0(n/2) since the loop increments by 2 each time, resulting in looping through half of the array
  size in the worst case. However, this simplifies to 0(n) because constant factors are dropped in Big O notation.
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The space complexity of the code is 0(1) because it uses a fixed amount of additional space (variables s, ans, and n) that