

Problem Description

In this problem, we are working with a record of tasks completed by different employees, each of whom has a unique identifier ranging from 0 to n - 1. There is a given 2D integer array named logs, where each element is a pair consisting of an employee id and a leave time. This array essentially tracks which employee finished which task and the time they finished it. The leave times are unique for each task.

Each task begins immediately after the previous one ends, and the first task (0th task) starts at time 0. The goal is to determine the employee id that worked on the task that took the longest time to complete. In case there are tasks that took the same maximum amount of time and were worked on by different employees, the employee with the smaller id is the one that should be returned.

Intuition

the finish time of the current task. The oth task starts at time 0, so for this first task, the duration is simply the finishing time of the task since it started at time 0. Now, we keep track of the maximum duration observed and the associated employee id. As we iterate through the logs array, we

compare the current task's duration with what we've recorded as the maximum duration so far. If the current duration is greater, we

To solve this problem, we need to determine the duration that each employee spent on their task. Since the tasks are completed in a

serial manner, the duration an employee spent on a task can be calculated by subtracting the finish time of the previous task from

update our records to now consider this as the max duration, and we update the "hardest worker" to be the current employee's id. If the current duration matches the maximum duration seen so far (a tie), we check the employee ids involved and update our "hardest worker" to be the employee with the smaller id. The managing of the "last" task finish time is critical because it marks the start of the next task (other than the first task, which starts at time 0). After determining the duration for the current task, we update the "last" finish time to match the current task's

The given solution employs a for-loop to scan the logs list once, maintaining the current maximum duration and the employee with the smallest id that worked for this duration. Once the loop is complete, it returns the identifier of the "hardest worker" as the result.

Solution Approach

The solution approach uses a simple linear scan algorithm to determine which employee worked the longest on a single task. It

makes use of basic comparison and arithmetic operations to track and update the task with the longest duration and the

## corresponding employee id.

Here is the step-by-step explanation of the solution algorithm: 1. Initialize three variables: last, mx, and ans. Set both last and mx (maximum task duration so far) to 0, and ans (the id of the employee who worked the longest) to the first employee's id.

2. Iterate over each entry in the logs list. Each entry is a pair [uid, t] which stands for an employee id (uid) and the time the task was completed (t).

finish time, effectively setting it up for the next comparison.

- 3. For each entry, update the duration of the current task by calculating the difference between the current task's completion time t and the last task's completion time (stored in last). The expression for this is:
- 1 t -= last

4. Now check if the current duration t is greater than the maximum duration x or if there is a tie and the current employee's id is

Update mx to match the current duration t.

Update ans to the current employee's id (uid).

smaller than the one previously stored. If either condition is true:

- 5. Update the last task completion time to account for the subsequent task's duration calculation: 1 last += t
- 6. After processing all entries, return ans as the solution. This is the id of the employee who worked the most extended task. The solution uses no additional data structures; it relies only on a few variables to track the information needed and operates with an
- Example Walkthrough

the space complexity is constant O(1), as no extra space is needed regardless of the size of the input list logs.

O(n) time complexity where n is the number of entries in the logs list. The algorithm is efficient as it scans the logs only once, and

Let's walk through a small example to illustrate the solution approach: Suppose we have these logs for employees with durations in which they completed the tasks: 1 logs = [[0,3],[2,5],[0,10],[1,15]]

## In this logs array, the first element [0,3] indicates that employee 0 finished a task at time 3. The second element [2,5] indicates that employee 2 finished the following task at time 5; and so on.

Now, let's apply the solution approach: 1. We initialize last, mx, and ans to 0 since no task has been completed yet, and we need a starting point.

which is 3, so we don't change mx or ans. We then update last to 5.

def hardestWorker(self, n: int, logs: List[List[int]]) -> int:

# 'last\_event\_time' to capture the end time of the last event.

# or if it's equal and the current user ID is lower than the one stored.

# Update the hardest worker ID and the max task duration

hardest\_worker\_id, max\_duration = user\_id, current\_duration

# 'max\_duration' to store the duration of the longest task.

# Update 'last\_event\_time' for the next iteration

// Function to find the hardest worker where 'n' is the number of workers,

int hardestWorkerId = 0; // Store the ID of the hardest worker.

// 'maxWorkDuration' to keep the maximum continuous work duration.

let [longestWorkId, maxWorkDuration, lastTimestamp] = [0, 0, 0];

// 'lastTimestamp' to keep the last timestamp we calculated work duration from.

// Calculate the work duration by subtracting lastTimestamp from current timestamp.

int hardestWorker(int n, vector<vector<int>>& logs) {

// and 'logs' is the vector of vectors containing the worker ID and the time they logged out.

last\_event\_time = current\_time

# Return the ID of the hardest worker

return hardest\_worker\_id

# 'hardest\_worker\_id' to store the ID of the hardest worker.

which is 3, so we update mx to 5 and ans to 0. Then we update the last to 10.

since this is the longest duration we have seen so far, and it belongs to employee 0. last = 0

2. We begin iteration with the first entry [0,3]. This task took 3 time units because it started at time 0. We set mx to 3 and ans to 0,

3. We continue to the next entry [2,5]. The duration for this task is 5 - last which is 2 (5 - 3). This duration is not longer than mx

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4. The next entry is [0,10]. The duration for this task is 10 - last which is 5 (10 - 5). This duration is greater than the current mx
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1 last = 10

3 ans = 0

2 mx = 53 ans = 0

5. The last entry is [1,15]. The duration for this task is 15 - last which is 5. This is equal to our current mx. Now, we compare the

employee ids; since 1 is larger than the current ans which is 0, we keep ans as 0. We update last to 15.

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3 ans = 0
6. We have exhausted all entries, so we finish the iteration, and the ans is the final result. Employee 0 is the one who worked the
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longest on a single task.

# Initialize variables:

last\_event\_time = 0

hardest\_worker\_id = 0

# Iterate through the logs

max\_duration = 0

Python Solution

class Solution:

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2 mx = 5

1 last = 15

units to complete.

In this example, the final output would be 0, which indicates that employee 0 worked the longest time on a task, which took 5 time

12 for user\_id, current\_time in logs: 13 # Calculate the duration of the current task 14 current\_duration = current\_time - last\_event\_time 15 # Check if the current task duration is greater than the max duration found so far, 16

if max\_duration < current\_duration or (max\_duration == current\_duration and hardest\_worker\_id > user\_id):

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Java Solution
   class Solution {
       public int hardestWorker(int n, int[][] logs) {
           int employeeIdWithMaxWork = 0; // To hold the ID of the employee who worked the most in a single log
           int lastEndTime = 0; // To keep track of the end time of the last task
           int maxDuration = 0; // To hold the maximum duration of work done continuously
           // Iterate through the work logs
           for (int[] log : logs) {
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               int currentEmployeeId = log[0]; // ID of the current employee
 9
               int currentTime = log[1]; // Time when the current task was finished
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               // Calculate the duration the current employee worked
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               int currentDuration = currentTime - lastEndTime;
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               // Determine if the current employee worked more than the previous max,
               // or if it's equal to the max duration and the current employee's ID is smaller.
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               if (maxDuration < currentDuration || (maxDuration == currentDuration && employeeIdWithMaxWork > currentEmployeeId)) {
17
                   employeeIdWithMaxWork = currentEmployeeId; // Update the employee ID with the max work
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                   maxDuration = currentDuration; // Update the max duration
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               // Update the end time of the last task to the current task's end time
               lastEndTime = currentTime;
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           return employeeIdWithMaxWork; // Return the ID of the employee who worked the hardest
```

### int maxWorkDuration = 0; // Store the maximum work duration. 9 10 11 12

C++ Solution

#include <vector>

class Solution {

public:

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int lastLogTime = 0; // Store the last log time to calculate the duration.
           // Loop through the logs to identify the hardest worker.
           for (auto& log : logs) {
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               int workerId = log[0]; // Get the worker's ID from the log.
14
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               int logTime = log[1]; // Get the log time from the log.
16
               // Calculate the work duration for the current worker.
17
               int workDuration = logTime - lastLogTime;
19
20
               // Check if this worker has worked more or equal than the current maximum work duration
21
               // and if this worker has a smaller ID than the current hardest worker's ID in case of a tie.
               if (maxWorkDuration < workDuration || (maxWorkDuration == workDuration && hardestWorkerId > workerId)) {
23
                   maxWorkDuration = workDuration; // Update the maximum work duration.
                   hardestWorkerId = workerId; // Update the hardest worker ID.
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               // Update the last log time for the next iteration.
               lastLogTime = logTime;
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           // Return the ID of the hardest worker after all logs have been processed.
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           return hardestWorkerId;
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34 };
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Typescript Solution
1 // Define the function hardestWorker which takes a total number of workers (n)
2 // and an array of logs with worker IDs and timestamps, and
  // returns the ID of the worker who worked the hardest.
   function hardestWorker(totalWorkers: number, logs: number[][]): number {
       // Initialize variables:
       // 'longestWork' to track the longest amount of work done.
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## // Update the longestWorkId and maxWorkDuration with the current worker's ID and work duration. 19 longestWorkId = workerId; 20 maxWorkDuration = workDuration; // Update lastTimestamp to calculate next work duration correctly.

// Iterate through each log entry.

for (let [workerId, timestamp] of logs) {

let workDuration = timestamp - lastTimestamp;

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# Time and Space Complexity

lastTimestamp = timestamp; 25 26 27 // Return the ID of the worker who worked the hardest. 28 return longestWorkId; 29 } 30

// Check if the current work duration exceeds our maxWorkDuration or if it is equal and the workerId is lower.

if (maxWorkDuration < workDuration || (maxWorkDuration === workDuration && longestWorkId > workerId)) {

# comparisons). Space Complexity

**Time Complexity** 

The space complexity of the code is 0(1) because the code uses a fixed amount of space that does not depend on the input size. Variables last, mx, ans, and any other temporary variables used during iteration do not scale with the size of the input, as they are simply storage for current and maximum time differences as well as the corresponding worker ID.

The time complexity of the code is O(n), where n is the number of entries in the logs list. The reason is that the code iterates

through each log entry exactly once, performing a constant amount of work for each entry (updating variables and simple