## 2997. Minimum Number of Operations to Make Array XOR Equal to K

Medium

Bit Manipulation

Array

## **Problem Description**

In this problem, we are given an array nums of integers and a separate integer k. Each number in the array is represented in binary and we can flip any of the bits in its binary representation (i.e., change a 0 to a 1 or vice versa) as many times as we want. Our goal is to perform the minimum number of such bit flips so that when we compute the bitwise XOR of all the numbers in the array, the result equals k. It is also mentioned that we can flip the leading zeros—which are usually omitted in binary notation. Our task is to find out the minimum number of flips needed to achieve this.

## To approach this problem, we don't really need to flip any bits initially. Instead, we can directly compute the bitwise XOR of all the

Intuition

elements in nums. Since XOR is an associative operation, the order in which we perform it doesn't matter, and since XORing a number with itself yields zero, any duplicate numbers effectively cancel each other out. After this computation, we'll have a single number which is the result of the XOR of all the array elements. The problem now simplifies to finding the minimum number of bits we need to flip in this result to make it turn into k. To do this,

we can XOR this result with k. Why does that work? XORing two numbers highlights the bits where they are different (since XOR gives us 1 when we compare

bits that are different and 0 when bits are the same). The final step is to count the number of bits that are set to 1 in this XOR result. Since these 1's represent differing bits, they are the exact bits we need to flip to turn our array XOR result into k. The number of such bits is the minimum operations required,

which is precisely what we need. Solution Approach

The solution relies on bit manipulation and the properties of the XOR operation. The approach is straightforward with four key

XOR all elements with k: The reduce function from Python's functools is utilized here with the xor operator from the operator

## steps:

module. reduce applies the xor function cumulatively to the items of nums, from left to right, so as to reduce the array to a single value. This value represents the XOR of all elements of nums. Then, by XORing this result with k, we find out which bits

differ between the XOR of the array elements and k. Calculate the number of differing bits: In the final result, a bit set to 1 indicates a position where a bit must be flipped to obtain k. The Python method bit\_count() is used on the resulting number to get the count of bits that are set to 1.

Minimum operations required: The count obtained from bit\_count() is the minimum number of operations required. This is

**Return the result:** The final count of set bits is returned as the solution to the problem: the minimum number of bit flips needed.

because each bit that is set to 1 can be flipped individually, and no unnecessary bit flips are required.

In terms of data structures, this solution is efficient as it operates with integers directly and does not require any additional data structures. The complexity of the algorithm is primarily dependent on the number of bits required to represent the numbers in the

This solution leverages the elegance and efficiency of bit operations to solve the problem in a crisp and concise manner.

Here's the consolidated approach in code: from operator import xor

from functools import reduce

class Solution:

array, which is at most  $O(\log(\max(nums, k)))$ .

```
In this code:

    reduce(xor, nums, k) computes the XOR of all elements in nums along with k.

• .bit_count() then counts the number of set bits (bits that are 1) of the result from the reduce function, thereby counting the number of
 operations needed.
```

Let's illustrate the solution approach with a small example.

Suppose we have the following array of integers nums and the integer k:

return reduce(xor, nums, k).bit\_count()

def minOperations(self, nums: List[int], k: int) -> int:

The binary representations of the numbers are:

 1 → 01 •  $5 \rightarrow 101$ 

**Example Walkthrough** 

• nums = [2, 3, 1]

• k = 5

• 2 → 10

• 3 → 11

1. XOR all elements with k: First, calculate the XOR of all the elements in nums:

• 2 XOR 3 XOR 1 = 10 XOR 11 XOR 01 = 00 (binary for 0)

Here's how we apply the solution approach step by step:

2. Calculate the number of differing bits: We use bit\_count() to count how many bits are set to 1 in 101:

# that counts the number of set bits (1s) in the integer

operations to transform the array's XOR result into k.

**Return the result**: The final result is 2, so our function minOperations(nums, k) will return 2. This example demonstrates how to apply the proposed algorithm step by step, showing that we can achieve the desired result

the validity of the algorithm.

from functools import reduce

return xor\_result.bit\_count()

return bin(xor\_result).count('1')

return Integer.bitCount(k);

from operator import xor

**Python** 

`python

class Solution {

class Solution {

int xorResult = k;

for (int num : nums) {

xorResult ^= num;

public:

**}**;

/\*\*

Java

Now we XOR the result with k:

• 0 XOR 5 = 000 XOR 101 = 101

• 101 has two bits set to 1.

Solution Implementation

Minimum operations required: The number of set bits is the minimum number of operations required. Thus, we need 2

with a minimum number of bit flips. Each bit that differs between the XOR result of the array and k represents a flip we need to

make. In this case, flipping the second and third bits (from the right) of the number 0 (the XOR of nums) will result in the binary

number 101, which represents the integer 5. The result aligns with the expected output of the minOperations method, confirming

class Solution: def minOperations(self, nums: List[int], k: int) -> int: # This method calculates the minimum number of operations # using bitwise XOR and counts the number of set bits. # Perform bitwise XOR of all numbers in the list with k xor\_result = reduce(xor, nums, k) # Return the count of set bits in the XOR result # .bit\_count() is a method available from Python 3.10 onwards

Note: In the original code, `reduce(xor, nums, k)` applies the `xor` operation consecutively to the elements of `nums` list start

If you are working in an environment that does not support Python 3.10 or newer, you could replace the `bit\_count()` method with

```
* This method calculates the minimum number of operations to reduce the XOR of all elements in the array to zero.
* An operation is defined as changing any element of the array to any other value.
* @param nums The array of integers to be processed.
* @param k The initial value to be XORed with elements of the array.
* @return The minimum number of operations required.
public int minOperations(int[] nums, int k) {
   // Iterate through each element in the array
   for (int num : nums) {
       // XOR the current element with k
       k ^= num;
```

// Return the count of bits set to 1 in the result of the XOR operations, as each bit set to 1 requires one operation to

// Group bytes and sum them up.

// Sum up halves.

// Sum up words.

```
C++
#include <vector>
#include <algorithm> // For std::count
```

// Function to find the minimum number of operations to convert a vector of integers

// \_\_builtin\_popcount returns the number of set bits (1-bits) in an integer

// which represents the number of different bits from the desired XOR result.

// This will be the minimum number of operations required to achieve the XOR result k.

// so that the XOR of all its elements is equal to a given integer k.

// Each operation can flip a bit in an integer of the vector.

// Calculate the XOR of all elements in nums with k

int minOperations(std::vector<int>& nums, int k) {

return builtin popcount(xorResult);

num = (num + (num >>> 4)) & 0x0f0f0f0f;

# Return the count of set bits in the XOR result

return xor\_result.bit\_count()

Time and Space Complexity

# .bit\_count() is a method available from Python 3.10 onwards

# that counts the number of set bits (1s) in the integer

num += num >>> 8;

num += num >>> 16;

# Counts the number of set bits ('1's) in the binary representation of the integer

```
TypeScript
// Calculates the minimum number of operations to reduce the XOR of all elements to zero.
// Each operation consists of incrementing or decrementing an element in the array.
function minOperations(nums: number[], targetXOR: number): number {
   // Calculate the cumulative XOR of all numbers in the array and the initial targetXOR.
   for (const num of nums) {
        targetXOR ^= num; // XOR current number with accumulator.
   // Return the count of set bits in the resulting XOR to determine
   // the minimum number of operations required.
   return countSetBits(targetXOR);
// Counts the number of set bits (1s) in the binary representation of the integer.
function countSetBits(num: number): number {
   // Apply bit manipulation techniques to count the set bits.
```

num = num - ((num >>> 1) & 0x55555555); // Group bit pairs and sum them up.

// Mask out the irrelevant bits and return the number of set bits which

num = (num & 0x33333333) + ((num >>> 2) & 0x33333333); // Group nibbles and sum them up.

```
// is contained in the least significant 6 bits of the result.
      return num & 0x3f;
from functools import reduce
from operator import xor
class Solution:
    def minOperations(self, nums: List[int], k: int) -> int:
        # This method calculates the minimum number of operations
        # using bitwise XOR and counts the number of set bits.
        # Perform bitwise XOR of all numbers in the list with k
        xor_result = reduce(xor, nums, k)
```

If you are working in an environment that does not support Python 3.10 or newer, you could replace the `bit\_count()` method with the ```python # Counts the number of set bits ('1's) in the binary representation of the integer return bin(xor\_result).count('1')

xor operation to each element of the array exactly once. The space complexity of the code is 0(1) since the xor operation and bit\_count() method do not require any additional space

that grows with the size of the input array. The operation is performed in constant space, regardless of the length of nums.

The time complexity of the code is O(n), where n is the length of the array nums. This is because the reduce function applies the

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