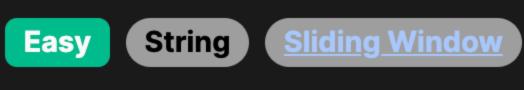
### 2379. Minimum Recolors to Get K Consecutive Black Blocks



### Problem Description

respectively. This string is indexed starting at 0, meaning the first element is at index 0. Our objective is to find the minimum number of operations required to ensure that there is at least one segment in the string consisting of k consecutive black blocks. An operation consists of changing a white block into a black block.

Imagine we have a block arrangement like "WBWBBW" and we want at least 3 consecutive black blocks. We need to determine

In this problem, we have a string blocks composed of characters 'W' and 'B', which represent white and black blocks,

the fewest changes needed to reach that goal; in this case, it's 1 operation by changing the second block from white to black resulting in "WBBWBB".

• We need to count the number of operations, which correlates to the number of white blocks ('W') that we need to recolor into black blocks ('B').

The problem specifies two key elements central to our task:

- ('B').
   We are seeking to create a continuous segment of k consecutive black blocks.
- Intuition

#### To arrive at the solution, we need to consider a <u>sliding window</u> approach. The reason a sliding window works well here is that

string, we'll look at the number of white blocks within each window, because those are our candidates for recoloring.

Initial thought might involve checking every possible window of k blocks in the string, and for each window, counting the number of white blocks and keeping track of the minimum count seen. But this is inefficient because we would be doing repetitive

we're interested in a continuous sequence of blocks — specifically, subsequences of length k. As we slide the window across the

counting. A more optimized approach involves only doing the counting once and then updating the count as the window slides.

The solution code implements this efficient approach:

• First, it counts the number of white blocks within the first k blocks. This is our initial answer — the worst-case number of operations, if the first k

• Then, the algorithm enters a loop where it slides the window one block at a time to the right. As the window slides, it adjusts the count by

blocks contain the most white blocks.

**Solution Approach** 

- adding one if a white block enters the window on the right, and subtracting one if a white block leaves the window on the left.

   After each adjustment, it updates the answer with the new count if it's lower than the previous answer.
- This allows us to go through the entire blocks string only once (after the initial count), adjusting the white block count and minimum number as we slide the window. The result is the minimum number of operations needed to ensure a sequence of k
- black blocks.

step walkthrough of the algorithm, as implemented in the provided solution code:

shifts the fixed-size window one block to the right with each iteration.

white blocks than any window seen before), update ans to cnt.

1. Initialize an integer variable ans (short for answer) to hold the minimum number of recoloring operations needed and another variable cnt (short for count) to keep track of the number of white blocks within the current window. Calculate the number of white blocks in the first window of size k using the count method on the substring blocks [:k]. Set ans to this initial count.

The implementation of this solution uses a simple but effective pattern known as the sliding window technique. Here's a step-by-

As the window slides, update cnt:
 Increment cnt by 1 if the new block coming into the right side of the window is a white block (blocks[i] == 'W').
 Decrement cnt by 1 if the block that is sliding out of the left side of the window is a white block (blocks[i - k] == 'W').

After updating cnt, compare it with the current value of ans. If cnt is smaller (which means the window now contains fewer

Use a loop to iterate through the string blocks, starting from the kth index up to the end of the string (not inclusive). This loop

No complex data structures are needed for this approach; only integer counters are used. The algorithm runs in linear time, O(n),

After the loop has iterated over the whole string, ans holds the minimum number of operations required. Return ans.

where n is the length of the blocks string. It only requires a constant amount of additional space, resulting in O(1) space complexity.

Example Walkthrough

Let's illustrate the solution approach with a smaller example. Consider the string blocks = "WWBWB" and let's say we want at least

By maintaining the count cnt as the window slides over the string, the solution avoids the need to recount white blocks in each

new window. This is much more efficient than a naive approach where you might recount white blocks for every new window of

#### first window of size 3, which is blocks[:3] = "WWB". This has 2 white blocks, so ans = 2. We also initialize the count variable

cnt = 2.

k = 3 consecutive black blocks.

**Update count and answer:** 

Following the steps outlined in the solution approach:

size k.

2. Iterate through blocks with a sliding window: We start our loop from index 3, which is the fourth block in blocks, and go up to the end of the string.

Initialize variables and count white blocks in the first window: The variable ans is set to the number of white blocks in the

initial window is white (blocks [0] = 'W'), which means we need to decrement cnt by 1. Now, cnt = 1.

Therefore, after the first iteration, the window looks at "WBW", with cnt = 1.

• At index 3, the block is 'W', so no new black block is added to the window. Since we're not adding a white block on the right, cnt remains 2.

However, we should also slide the window, which means we remove the leftmost block from our initial count. The leftmost block within our

Final return value: In the next and final iteration, the window would look at "BWB", and since there's 1 white block entering the

window on the right (blocks[4] == 'W'), cnt is incremented back to 2. However, the block that leaves on the left (blocks[1])

Following this process, we've completed our sweep of the block string, and ans holds the value 1, which represents the minimum

Thus, the output would be 1 because we need to change one white block to achieve a segment of 3 consecutive black blocks ("WBWB" would become "WBBB").

# Initialize minimum white blocks to be recolored with the count from the first window

# If the block that is exiting the window is white, decrement the count

// Initialize the count of white blocks within the first window of size k

# Update the minimum if the current count is less than the previous minimum

// Method to find the minimum number of recolors needed to get at least k consecutive black blocks

number of operations required to ensure there is a segment of 3 consecutive black blocks in the string.

**Update the answer:** Our new cnt is less than the previous ans, so we set ans to cnt. Now, ans = 1.

is a black block, so cnt remains 2. There's no need to update ans since cnt is not smaller than ans.

Solution Implementation

Python

for index in range(k, len(blocks)):
 # If the newly included block in the window is white, increment the count
 if blocks[index] == 'W':
 white\_count += 1

```
# Return the minimum number of white blocks that need to be recolored return min_recolors

Java
```

int whiteCount = 0;

class Solution {

public:

class Solution:

def minimumRecolors(self, blocks: str, k: int) -> int:

white\_count = blocks[:k].count('W')

if blocks[index - k] == 'W':

public int minimumRecolors(String blocks, int k) {

// Function to find the minimum number of recolors needed

// of white blocks in the first window of size 'k'

// Count the number of white blocks in the first 'k' block segment

// Increase whiteBlockCount if the current block is white

// Update minRecolors to the smallest number of white blocks

// Decrease whiteBlockCount if the leftmost block of

minRecolors = min(minRecolors, whiteBlockCount);

// The minimum recolors needed is initially set to the number

// Iterate over the blocks starting from the 'k'th block

int whiteBlockCount = count(blocks.begin(), blocks.begin() + k, 'W');

// to get at least 'k' consecutive black blocks

int minimumRecolors(string blocks, int k) {

int minRecolors = whiteBlockCount;

for (int i = k; i < blocks.size(); ++i) {</pre>

// the previous window was white

// seen in any window of size 'k'

def minimumRecolors(self, blocks: str, k: int) -> int:

white\_count = blocks[:k].count('W')

for index in range(k, len(blocks)):

if blocks[index - k] == 'W':

if blocks[index] == 'W':

white\_count += 1

white\_count -= 1

min recolors = white count

# Initial count of 'W' in the first window of size k

# Slide the window of size k through the blocks string

min recolors = min(min recolors, white count)

# Initialize minimum white blocks to be recolored with the count from the first window

# If the newly included block in the window is white, increment the count

# If the block that is exiting the window is white, decrement the count

# Return the minimum number of white blocks that need to be recolored

# Update the minimum if the current count is less than the previous minimum

whiteBlockCount += blocks[i] == 'W';

whiteBlockCount -= blocks[i - k] == 'W';

white\_count -= 1

min\_recolors = white\_count

# Initial count of 'W' in the first window of size k

# Slide the window of size k through the blocks string

min\_recolors = min(min\_recolors, white\_count)

```
for (int i = 0; i < k; ++i) {
            if (blocks.charAt(i) == 'W') {
                whiteCount++;
        // Initialize the answer with the count of white blocks within the first window
        int minRecolors = whiteCount;
        // Slide the window of size k across the string and update the minimum recolors requird
        for (int i = k; i < blocks.length(); ++i) {</pre>
            // If the entering character is white, increment white count
            if (blocks.charAt(i) == 'W') {
                whiteCount++;
           // If the exiting character is white, decrement white count
            if (blocks.charAt(i - k) == 'W') {
                whiteCount--;
           // Update the minimum recolors if the current count is less than the previous minimum
           minRecolors = Math.min(minRecolors, whiteCount);
       // Return the minimum number of recolors required
        return minRecolors;
C++
class Solution {
```

```
// Return the minimum number of recolors needed
       return minRecolors;
};
TypeScript
// Function to find the minimum number of recolors to get at least k consecutive black blocks.
// @param blocks - String representing the arrangement of black (B) and white (W) blocks.
// @param k - Number representing the desired consecutive black blocks length.
// @returns The minimum number of recolors required (changing 'W' to 'B').
function minimumRecolors(blocks: string, k: number): number {
    // Initialize a counter for white blocks in the first window of size k.
    let whiteCount = 0;
    for (let i = 0; i < k; ++i) {
       whiteCount += blocks[i] === 'W' ? 1 : 0;
    // The answer starts off as the number of white blocks in the first window.
    let minRecolors = whiteCount;
    // Slide the window of size k across the blocks string while updating the count.
    for (let i = k; i < blocks.length; ++i) {</pre>
       // If the new block is white, increase the count.
       whiteCount += blocks[i] === 'W' ? 1 : 0;
       // If the block exiting the window is white, decrease the count.
       whiteCount -= blocks[i - k] === 'W' ? 1 : 0;
       // Update the answer with the minimum count seen so far.
       minRecolors = Math.min(minRecolors, whiteCount);
    // Return the minimum number of recolors needed.
    return minRecolors;
```

# return min\_recolors Time and Space Complexity

The time complexity of the given code can be analyzed based on the operations performed within the loop. The loop runs from k to len(blocks), which indicates that the loop runs len(blocks) – k times. Within each iteration, the code performs constant-time operations such as comparison and increment/decrement operations. Therefore, the overall time complexity is 0(len(blocks) - k), which simplifies to 0(n) where n is the length of the blocks string.

## Space Complexity

**Time Complexity** 

class Solution:

The space complexity is determined by the amount of additional space used by the algorithm relative to the input size. The given code uses a fixed number of variables (ans, cnt, i) that do not depend on the size of the input. Therefore, the space complexity remains constant, regardless of the input size. Consequently, the space complexity of the code is 0(1), indicating constant space complexity.