1222. Queens That Can Attack the King

Medium Array Matrix Simulation

Problem Description

grid, akin to a coordinate plane, where the intersection of each horizontal row and vertical column denotes a square on the chessboard.

On an 8×8 chessboard, the game setup includes multiple black queens and one white king. The chessboard aligns with a 0-indexed

Leetcode Link

A 2D integer array queens, where each element queens[i] = [xQueen_i, yQueen_i] represents the position of the i-th black

We are provided with two inputs:

- queen on the chessboard. An integer array king of length 2, with king = [xKing, yKing] standing for the position of the white king on the chessboard.
- The challenge lies in identifying all the black queens that are in a direct position to attack the white king. The queens can attack in straight lines horizontally, vertically, or diagonally. The positions of the black queens capable of launching an attack on the king must

be returned and can be listed in any sequence. Intuition

The essence of finding a solution resides in understanding how queens move and attack in chess. A queen attacks along rows, columns, and diagonals until it's obstructed by another piece or the edge of the chessboard.

So, the intuitive approach is to inspect the chessboard row by row, column by column, and along each diagonal emanating from the king's position. We want to locate the closest queen in each of these directions, as only those queens would be able to attack the king without being blocked by other queens.

The solution involves the following steps: Transform the list of queen positions into a set for efficient presence checking. 2. Define the 8 possible directions in which the king can be attacked based on queen movements: vertical, horizontal, and the four

diagonals.

- 3. Iterate over these directions, "walking" from the king's position step by step in one direction until an edge of the board is reached or a queen is found.
- 4. When a queen is located in one of these directions, we record her position as she can attack the king, then move on to the next direction.
- Solution Approach
- The implementation of the solution leverages simple yet effective programming concepts to iterate through the grid and find the attacking queens. Here's a walk-through of the solution:

2. A set s is created to store the positions of the queens. This allows for O(1) look-up times when checking if a given position on

1. n is set to 8, which represents the size of the chessboard.

By proceeding as such, we collect the positions of all queens that pose a direct threat to the king.

the chessboard has a queen or not.

horizontally, and diagonally.

king in any given direction is considered a direct threat.

Let's consider a small example based on the solution approach described:

2. The queen's position would be added to a set: $s = \{(3, 5)\}$.

Now, we look at all 8 possible directions one by one:

3. We initiate an empty list ans to keep track of attacking queens.

Horizontally and Vertically (four directions): left, right, up, down

4. The algorithm uses two nested loops, each iterating from -1 to 1, to check all 8 possible directions a queen can attack from (excluding the stationary position (0, 0)). These two numbers (a, b) represent the "step" in each direction — vertically,

5. For each direction, we start at the king's position and "move" one step at a time in that direction by incrementing x by a and y by

6. During each step, we check if the new position (x, y) has a queen (if (x, y) in s). Once we find a queen, we add her position

b. We continue this step-wise movement until we reach the edge of the board (0 <= x + a < n and 0 <= y + b < n).

3. We define an empty list ans to eventually hold the positions of the queens that can attack the king.

- to the ans list using ans append ([x, y]). We then break out of the inner while loop to stop looking in that direction because it's not possible for another queen to attack the king from the same direction without being obstructed by the first one found. 7. After inspecting all 8 directions, we return the list ans which contains the positions of all queens that can directly attack the king.
- Example Walkthrough

By using sets for constant-time look-ups and only searching along each direction once, we ensure that the algorithm runs efficiently.

It avoids redundant checks and follows the movement of queens in chess closely to guarantee that only the queen closest to the

Suppose the 2D board has only one queen located at [3, 5] and the white king is at [4, 3]. We want to use the algorithm to determine if this queen can actually attack the king. 1. The chessboard size n is 8 (0-indexed from 0 to 7).

Diagonally (four directions): upper-left, upper-right, lower-left, lower-right For each direction, we start from the king's position [4, 3] and move in steps.

Let's pick two directions to illustrate:

Direction (1, 1) - This represents checking diagonally down-right from the king.

Direction (0, 1) - This represents checking to the right of the king.

squares. At [4, 5] we exit the board, and since the queen is not in that direct path, we do not find a threat from that direction.

2], [2, 1], [1, 0], but we don't find a queen in this path using our set s.

Size of the chess board

return attack_positions

final int SIZE = 8;

queen_positions = {(i, j) for i, j in queens}

Skip (0, 0) to prevent standing still

// Function that returns a list of queens that can attack the king.

public List<List<Integer>> queensAttacktheKing(int[][] queens, int[] king) {

for delta_col in range(-1, 2):

if delta_row or delta_col:

// Define the size of the chessboard

for (int[] queen : queens) {

// Board to track the positions of the queens

// Place queens on the board based on given positions

boolean[][] board = new boolean[SIZE][SIZE];

board[queen[0]][queen[1]] = true;

board_size = 8

class Solution:

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Java Solution

1 import java.util.ArrayList;

2 import java.util.List;

class Solution {

Now, the step-by-step process is as follows:

signifies that no queens can attack the king in the given board setup. Python Solution

1. Starting with the right direction, we move from [4, 3] to [4, 4], [4, 5], ... and we keep checking if there's a queen on these

2. Checking diagonally down-right, we go from [4, 3] to [5, 4] to [6, 5] and see there is no queen at these positions (since the

3. Continuing this for all 8 directions, we eventually check the up-left diagonal direction (-1, -1). We would step from [4, 3] to [3,

As we can see from our direction checks, the single queen at [3, 5] is not in a position to attack the king at [4, 3] in any of the 8

directions investigated by our algorithm. Therefore, the ans list remains empty, indicating that the king is not under attack from the

given position of the queen. At the end of the iterating through all possible directions, we would return the unchanged ans list, which

exit the board boundaries. Since we neither find a queen nor exit the boundaries, we move no further.

def queensAttacktheKing(self, queens: List[List[int]], king: List[int]) -> List[List[int]]:

Set representation of queen positions for constant time accessibility

List to hold the final positions of queens that can attack the king

set s only contains [3, 5], and this does not match our current position [6, 5]), and we continue until we either find a queen or

attack_positions = [] # Explore all 8 directions from the king's position: # These directions are combinations of moving vertically (a), horizontally (b), 12 # and diagonally on the chess board. for delta_row in range(-1, 2): 13

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current_row, current_col = king
                        # Move in the direction until hitting the edge of the board or finding a queen
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                        while 0 <= current_row + delta_row < board_size \</pre>
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                                and 0 <= current_col + delta_col < board_size:</pre>
21
                            # Update current position
22
                            current_row, current_col = current_row + delta_row, current_col + delta_col
23
                            # Check if current position is occupied by a queen
                            if (current_row, current_col) in queen_positions:
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25
                                # If found, add the position to the answer list and stop searching this direction
26
                                attack_positions.append([current_row, current_col])
27
                                break
28
           # Return all the unique positions from which the queens can attack the king directly
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1 #include <vector>

2 using namespace std;

class Solution {

public:

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36 }

17 18 // List to store positions of queens that can attack the king 19 List<List<Integer>> attackPositions = new ArrayList<>(); 20 21 // Directions - The pair (a, b) represents all 8 possible directions from a cell 22 11 (0, -1) (-1, -1) (-1, 0)23 11 (0, 1)(1, 0)24 (1, -1)(1, 1)25 for (int rowDir = -1; rowDir <= 1; rowDir++) {</pre> 26 for (int colDir = -1; colDir <= 1; colDir++) {</pre> 27 28 // Skip standstill direction as the king is not moving 29 if (rowDir != 0 || colDir != 0) { 30 // Starting position for searching in a specific direction int x = king[0] + rowDir, y = king[1] + colDir; // Traverse in the direction until a queen is found or edge of board is reached while $(x >= 0 \&\& x < SIZE \&\& y >= 0 \&\& y < SIZE) {$ 35 // Check if queen is found if (board[x][y]) { 36 37 attackPositions.add(List.of(x, y)); 38 break; // Break out of the loop as we're only looking for the closest queen 39 40 // Move to next position in the direction 41 x += rowDir; 42 y += colDir; 43 44 45 46 47 48 return attackPositions; 49 50 } 51 C++ Solution

vector<vector<int>> queensAttacktheKing(vector<vector<int>>& queens, vector<int>& king) {

bool squares[boardSize][boardSize]{ false }; // Initialize all squares to 'false'

// Directions that need to be checked: Horizontal, Vertical, Diagonal (8 directions)

// Start from the king's position and check in the current direction

while $(x \ge 0 \&\& x < boardSize \&\& y \ge 0 \&\& y < boardSize) {$

vector<vector<int>> attackedQueens; // To hold the positions of queens that can attack the king

// Skip checking the direction where both deltaX and deltaY are zero (the king's position)

// Add the queen's position to the result vector and break out of the loop to check the next direction

const int boardSize = 8; // The size of the chessboard

for (const auto& queen : queens) {

continue;

squares[queen[0]][queen[1]] = true;

for (int deltaX = -1; deltaX <= 1; ++deltaX) {</pre>

if (deltaX == 0 && deltaY == 0) {

// Continue until out of bounds

if (squares[x][y]) {

break;

for (int deltaY = -1; deltaY <= 1; ++deltaY) {</pre>

int x = king[0] + deltaX, y = king[1] + deltaY;

// If a queen is found on the current square

attackedQueens.push_back({x, y});

// Mark positions where queens are located on the chessboard

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                         // Move to the next square in the current direction
 35
                         x += deltaX;
 36
                         y += deltaY;
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             return attackedQueens; // Return the list of queens that can attack the king
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 43 };
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Typescript Solution
   function queensAttacktheKing(queens: number[][], king: number[]): number[][] {
       const boardSize = 8; // Chessboard size is 8x8
       // Create a boolean matrix to record the positions of queens
       const board: boolean[][] = Array.from({ length: boardSize }, () => Array.from({ length: boardSize }, () => false));
       // Mark the positions of all queens on the board
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       queens.forEach(([row, col]) => (board[row][col] = true));
       // Initialize an array to store queens that can attack the king
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       const attackingQueens: number[][] = [];
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       // Iterate over all possible directions from the king: 8 surrounding cells (horizontal, vertical, and diagonal)
       for (let rowDirection = -1; rowDirection <= 1; ++rowDirection) {</pre>
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           for (let colDirection = -1; colDirection <= 1; ++colDirection) {</pre>
14
               // Skip checking the same cell where the king is positioned
15
               if (rowDirection || colDirection) {
16
                   // Start from the king's position and check in the direction specified by rowDirection and colDirection
17
                   let [currentRow, currentCol] = [king[0] + rowDirection, king[1] + colDirection];
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19
                   // Continue checking the cells within the board limits
20
                   while (currentRow >= 0 && currentRow < boardSize && currentCol >= 0 && currentCol < boardSize) {</pre>
21
                       // If a queen is found in the current direction, add it to the attackingQueens array and break the loop
                       if (board[currentRow][currentCol]) {
                           attackingQueens.push([currentRow, currentCol]);
                           break;
                       // Move to the next cell in the current direction
28
                       currentRow += rowDirection;
                       currentCol += colDirection;
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Time and Space Complexity

return attackingQueens;

// Return the array containing the positions of the attacking queens

queen. Let's analyze the time and space complexity: **Time Complexity**

The code performs a search in each of the 8 directions from the king's position until it either hits the end of the board or finds a

The time complexity is O(n), where n is the number of cells in the board (in this case, n is 64, since it's an 8×8 chessboard). We

iterate over each direction only once and in the worst case, we traverse the entire length of the board. However, since the board size

is fixed, we can consider this time complexity to be 0(1) in terms of the input size, because the board doesn't grow with the input.

Space Complexity

The space complexity is O(q), where q is the number of queens, because we store the positions of the queens in a set s. The board size is fixed and does not influence the space complexity beyond the storage of queens. In the worst case where every cell contains a queen, space complexity would be O(n), but since the board's size is constant and doesn't scale with the input, this can also be referred to as 0(1) from a practical perspective.