# Problem Description

find the number of distinct excellent pairs in the array, where a pair (num1, num2) is considered excellent if it satisfies two conditions: Both num1 and num2 exist in nums.

2. The sum of the number of set bits (bits with value 1) in num1 OR num2 and num1 AND num2 must be greater than or equal to k.

- In this problem, we're presented with an array nums consisting of 0-indexed positive integers and a positive integer k. Our goal is to
- enough number of set bits to meet or exceed the threshold k.

We count the number of set bits using bitwise OR and AND operations, and we are looking for pairs that collectively have a large

Also, it's important to note that a pair where num1 is equal to num2 can also be considered excellent if the number of its set bits is sufficient and at least one occurrence of the number exists in the array. Moreover, we want the count of distinct excellent pairs, so the order matters—(a, b) is considered different from (b, a).

Intuition To find the number of excellent pairs efficiently, we need to think about the problem in terms of set bits because the conditions for

an excellent pair depend on the sum of set bits in num1 OR num2 and num1 AND num2.

### One intuitive approach is to use the properties of bitwise operations. For any two numbers, num1 OR num2 will have the highest

k.

possible set bit count of the two numbers because OR operation results in a 1 for each bit that is 1 in either num1 or num2. On the other hand, num1 AND num2 will have set bits only in positions where both num1 and num2 have set bits. Since only the set bits matter, and we're looking for pairs (num1, num2) that meet a certain combined set bit count, we can reduce

the complexity by avoiding the direct computation of num1 OR num2 and num1 AND num2 for all pairs. Instead, we preprocess by

counting the set bits for each unique number in nums. To ensure that we count each distinct pair only once, we eliminate duplicates in the array by converting it into a set. We then create a counter to keep track of how many numbers have a specific set bit count. This preprocessing step simplifies our task to just

combining counts from the counter, avoiding the need to directly compute the set bit sum for every possible pair. When we iterate over the unique numbers in nums set, we determine the number of set bits for each unique value using bit\_count(). We then iterate through our set bit count counter and add the count of numbers that have enough set bits to complement the current number's set bit count to reach at least k. This way, we find all the pairs that, when combined, meet or exceed the threshold

problem of having to generate and check every possible pair, which would be computationally inefficient. Solution Approach

The given solution employs a combination of bit manipulation and hash mapping to efficiently compute the number of excellent pairs.

By adding the counts for all such complementing set bit count pairs, we calculate the total number of excellent pairs, bypassing the

1. Eliminating Duplicates: The solution begins by converting the original list of numbers into a set. This serves two purposes: Ensures that each number is considered only once, thereby eliminating redundant pairs like (num1, num1) when num1

Helps later in avoiding double counting of excellent pairs with the same numbers but in different positions (since (a, b) and

#### (b, a) are distinct). 1 s = set(nums)

1 cnt = Counter()

Let's break down the implementation step-by-step:

appears multiple times in the input list.

will be equivalent to bit\_count(num1) + bit\_count(num2).

sum of bit\_count(v) and i is greater than or equal to k.

t holds the number of set bits for the current unique number v.

2. Counting Set Bits: Then, the algorithm uses a Counter from the collections module to keep track of how many numbers share the same set bit count.

- 2 for v in s: cnt[v.bit\_count()] += 1
  - The method bit\_count() is used to determine the number of set bits in each number. These counts are stored such that cnt[i]
- reflects the number of unique numbers with exactly i set bits. 3. Finding Excellent Pairs: The core of the solution is to find all the pairs that satisfy the condition of having a combined set bit sum (via bitwise OR and AND) greater than or equal to k. Since an OR operation can never reduce the number of set bits, and an

AND operation can only produce set bits that are already set in both numbers, the combined set bit count for a pair (num1, num2)

The solution iterates over each unique value v from the set of numbers and then checks for every bit count i stored in cnt if the

```
t = v.bit_count()
for i, x in cnt.items():
    if t + i >= k:
        ans += x
```

1 ans = 0

1 return ans

Example Walkthrough

· 1 has 1 set bit.

2 has 1 set bit.

Let's say we have an array nums = [3, 1, 2, 2] and k = 3.

Next, we count set bits in step 2. Every number will be processed as follows:

• For v = 2 also with 1 set bit, the situation is the same as for v = 1.

the set bit counts to find excellent pairs, which is computationally faster.

def countExcellentPairs(self, nums: List[int], k: int) -> int:

# Count the frequency of the bit count for each number

for bit\_count, freq in bit\_count\_freq.items():

if current\_bit\_count + bit\_count >= k:

excellent\_pairs\_count += freq

# Return the total count of excellent pairs

// Method to count the number of excellent pairs

public long countExcellentPairs(int[] nums, int k) {

Set<Integer> uniqueNumbers = new HashSet<>();

// Use a set to eliminate duplicate values from 'nums'

std::unordered\_set<int> unique\_numbers(nums.begin(), nums.end());

// Variable to store the final count of excellent pairs

if (current\_bit\_count + i >= k) {

// Return the final count of excellent pairs

function countExcellentPairs(nums: number[], k: number): number {

// Variable to store the final count of excellent pairs

// Set to eliminate duplicates from the input array

const uniqueNumbers = new Set<number>(nums);

const bitCount = new Array(32).fill(0);

uniqueNumbers.forEach(number => {

let excellentPairsCount = 0;

std::vector<int> bit\_count(32, 0);

for (int number : unique\_numbers) {

long long excellent\_pairs\_count = 0;

for (int number : unique\_numbers) {

return excellent\_pairs\_count;

function countSetBits(num: number): number {

for (int i = 0; i < 32; ++i) {

// There are at most 32 bits in an int, so we create an array of size 32

// Count the number of times a number with a particular bit count appears

// Iterate over each unique number and find the count of numbers that

excellent\_pairs\_count += bit\_count[i];

// Array to count the frequency of set bits (1s) in the binary representation of numbers

// have the required number of bits for forming an excellent pair with the current number

// If the sum of the bit counts of both numbers meets or exceeds k,

// add the frequency of the corresponding bit count to the answer

// Function to count the number of set bits in the binary representation of a number

// Logical Right Shift to process the next bit

// Array to count the frequency of set bits (1s) in the binary representation of numbers

bitCount[countSetBits(number)]++; // Increment the frequency of the bit count

// Count the number of times a number with a particular bit count appears

// Since integers in JavaScript are represented using 32 bits, we create an array of size 32

count += num & 1; // Increment count if the least significant bit is set

int current\_bit\_count = \_\_builtin\_popcount(number); // Get bit count of the current number

++bit\_count[\_builtin\_popcount(number)]; // \_builtin\_popcount returns the number of bits set to 1

return excellent\_pairs\_count

for (int num : nums) {

long totalPairs = 0;

uniqueNumbers.add(num);

bit\_count\_freq[num.bit\_count()] += 1

# Use a set to eliminate duplicates as each number contributes uniquely

excellent\_pairs\_count = 0 # Initialize the count of excellent pairs

numbers with 1 set bit, we add 2 to ans, resulting in ans = 2 so far.

For v = 3 with 2 set bits, we compare it against cnt entries.

the conditions are (3, 1), (3, 2), (1, 3), (2, 3).

from collections import Counter

for num in unique\_nums:

unique set bit numbers to quickly find the number of complements needed to meet the set bit threshold k. 4. Returning the result: Finally, after iterating through all unique numbers and their possible pairings, the solution returns ans, which holds the total number of distinct excellent pairs.

If t + 1 is greater than or equal to k, all numbers with a set bit count 1 can form an excellent pair with v. The count of such

numbers is x, and we add this to our total count of excellent pairs (ans). This step leverages the precomputed counts of

redundant operations and directly navigates to the crux of the problem, which significantly improves the computational efficiency. By employing a hash map to store unique set bit counts and identify complementing pairs, the solution reduces what would be a quadratic-time complexity task to a complexity proportional to the product of unique numbers and unique set bit counts present in the input.

First, we'll apply step 1 and eliminate duplicates by converting nums into a set, which will give us  $s = \{1, 2, 3\}$ .

In summary, the solution follows a smart preprocessing step to calculate and use set bit counts for optimization. This eliminates

· 3 has 2 set bits. The Counter object cnt turns out to be {1: 2, 2: 1}, indicating that there are 2 numbers with 1 set bit and 1 number with 2 set bits.

• For v = 1 with 1 set bit, cnt[1] = 2. Since 1 (set bits of v) + 1 (set bits of another number) is not >= k, no excellent pairs

o 2 (set bits of v) + 1 (set bits of another number) is >= k, thus excellent pairs are (3, 1) and (3, 2). Since there are 2

After processing all the unique numbers, we follow step 4 and conclude that there are 4 distinct excellent pairs. The pairs that satisfy

This example demonstrates the solution's efficiency, as it avoids checking all possible combinations of nums and directly focuses on

Now for step 3, we find excellent pairs. We go through each number in s and compare the set bit count with our k value:

### Finally, since the pairs (1, 3) and (2, 3) are also excellent pairs (order matters), we count them again. This adds another 2 to ans, making ans = 4.

are formed with v = 1.

Python Solution

unique\_nums = set(nums) # Counter to store the frequency of bit counts bit\_count\_freq = Counter()

# If the sum of bit counts is greater than or equal to k, add to the count

# Calculate the number of excellent pairs for num in unique\_nums: current\_bit\_count = num.bit\_count() # Get the bit count of current number 19 20 21 # Iterate over the bit count frequencies

Java Solution

class Solution {

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class Solution:

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int[] bitCounts = new int[32]; // Array to store how many numbers have a certain bit count
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           // Count the occurrence of each bit count for the unique elements
           for (int num : uniqueNumbers) {
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               int bits = Integer.bitCount(num); // Count the 1-bits in the binary representation of 'num'
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               ++bitCounts[bits];
                                                  // Increase the count for this number of 1-bits
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           // Iterate over the unique numbers to find pairs
           for (int num : uniqueNumbers) {
23
               int bits = Integer.bitCount(num); // Count the 1-bits for this number
24
               // Check for each possible bit count that could form an excellent pair with 'num'
25
               for (int i = 0; i < 32; ++i) {
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27
                   // Check if the sum of 1-bits is at least 'k'
28
                   if (bits + i >= k) {
                       totalPairs += bitCounts[i]; // If it is, add the count of numbers with 'i' bits to the total
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           return totalPairs; // Return the total count of excellent pairs
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36 }
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C++ Solution
 1 #include <vector>
   #include <unordered_set>
   class Solution {
   public:
        long long countExcellentPairs(std::vector<int>& nums, int k) {
           // Create a set to eliminate duplicates from the input vector
```

// To store the total number of excellent pairs

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num >>>= 1;
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        return count;
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});

Typescript Solution

let count = 0;

while (num > 0) {

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       // Iterate over each unique number and find the count of numbers that
       // have the required number of bits for forming an excellent pair with the current number
       uniqueNumbers.forEach(number => {
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           const currentBitCount = countSetBits(number); // Get bit count of the current number
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           for (let i = 0; i < 32; i++) {
31
               if (currentBitCount + i >= k) {
                   // If the sum of the bit counts of both numbers meets or exceeds k,
                   // add the frequency of the corresponding bit count to the answer
                   excellentPairsCount += bitCount[i];
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       });
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39
       // Return the final count of excellent pairs
       return excellentPairsCount;
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42 }
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Time and Space Complexity
The given code counts the number of "excellent" pairs in an array, with a pair (a, b) being excellent if the number of bits set to 1 in
their binary representation (a OR b) is at least k.
Time Complexity
```

## 1. Creating a set s from nums: This operation has a time complexity of O(N) where N is the number of elements in nums, as each element is added to the set once.

3. Populating the Counter: This again is O(U) since we're iterating over the set s once.

The time complexity of the function involves several steps:

- 4. Running the nested loops, where the outer loop is over each unique value v in s and the inner loop is over the counts in Counter:
- Since the outer loop runs O(U) times and the inner loop runs a maximum of O(U) times, the nested loop part has a worse-case complexity of O(U^2).

2. Counting the bit\_count t of each distinct number in s: Each call to v.bit\_count() is 0(1). Since we do this for each number in the

The total time complexity is, therefore,  $O(N + U^2)$ .

set s, this step has a complexity of O(U), where U is the number of unique numbers in nums.

The space complexity consists of: 1. The set s, which takes O(U) space.

Hence, the total space complexity of the function is O(U), where U is the count of unique numbers in nums.

**Space Complexity** 

2. The Counter object cnt, which takes another O(U) space. The additional space for the variable ans and loop counters is O(1).