1540. Can Convert String in K Moves

String

Problem Description

Hash Table

Medium

In this problem, you are given two strings s and t, and an integer k. The goal is to determine if you can transform string s into string t by performing up to k moves. A move consists of picking an index j in string s and shifting the character at that index a certain number of times. Each character in s may only be shifted once, and the shift operation wraps around the alphabet (e.g., shifting 'z' by 1 results in 'a'). The key constraints are:

• Each move involves shifting one character of s by i positions in the alphabet, where $1 \ll i \ll k$.

You can make no more than k moves in total.

- The index chosen for shifting in each move should not have been used in a previous move. You can only shift a character to the next one in the alphabet, with 'z' changing to 'a'.
- The task is to return true if s can be transformed into t using at most k moves under these rules; otherwise, return false.
- ntuition

The solution approach starts by realizing that in order to change s into t, for each position where s and t differ, we need to

calculate the number of shifts required to convert the character in s to the corresponding character in t. This can be done by finding the difference in their ASCII values and taking the modulus with 26 to handle the wrapping around the alphabet. The next step is to keep track of the number of times each shift amount is needed. We use a list, cnt, to record the frequency of

already the same, and no move is needed. Once we have this frequency array, we check if any of these shifts can be performed within the move limit k. For each non-zero shift amount, we calculate when the last shift can occur. Since each shift amount can be used every 26 moves, we determine the maximum moves needed for each shift (i + 26 * (cnt[i] - 1)). If this exceeds k, it's impossible to transform s into t, and we

each required shift amount from 0 to 25 (since there are 26 letters in the alphabet). A shift amount of 0 means the characters are

return false. The reason for subtracting 1 from cnt[i] is that the first occurrence of each shift amount does not have to wait for a full cycle it can happen immediately. Hence, only subsequent occurrences (if any) need to be delayed by 26 moves each.

The implementation of the solution involves a few key steps that use simple data structures and algorithms.

that only characters in s can be shifted, implying both strings must be of the same length to be convertible.

Length Check: First, we check if s and t are of equal length. If not, return False immediately because the problem states

Frequency Array: We create an array, cnt, of length 26 initialized to zero, which will hold the frequency of each shift amount

increment occurs at cnt[0].

we return False.

Example Walkthrough

Let's consider a simple example:

Solution Approach

required.

- Calculate Shifts: We loop over the characters of s and t simultaneously using Python's zip function. For each corresponding pair of characters (a, b), we calculate the shift required to turn a into b. The shift is calculated using the
- addition of 26 before the modulus operation is to ensure a positive result for cases where b comes before a in the alphabet. Count the Shifts: We increment the count of the appropriate shift amount in the cnt array. If no shift is needed, the

formula (ord(b) - ord(a) + 26) % 26, where ord() is a Python function that returns the ASCII value of a character. The

- Move Validation: After calculating all the necessary shifts, we iterate over the cnt array (starting from index 1 since index 0 represents no shift). For each shift amount i, we find out how late this shift can occur by multiplying the number of full cycles (cnt[i] - 1) by the cycle length 26, and adding the shift amount i itself. The result tells us at which move number the last shift could feasibly happen. If this move number is greater than k, the transformation is not possible within the move limit, so
- By using a frequency array and calculating the maximum move number needed for each shift, the algorithm efficiently determines the possibility of transformation without actually performing the moves, resulting in a less complex and time-efficient solution.

Result: If none of the calculated shift timings exceed k, then it is verified that all characters from s can be shifted to match

• s = "abc" • t = "bcd" • k = 2

• Both strings s and t are of equal length, which is 3. We can proceed with the transformation process.

which exceeds k = 2, implying that it's impossible to conduct all required shifts within k moves.

Calculate the shift difference and take modulo 26 to wrap around the alphabet

For multiple shifts 'i', we need to wait 26 more for each additional use

// If all shifts are possible within the allowed maximum, return true.

// If the lengths of source and target are not the same, then conversion is not possible

// Initialize an array to count the number of times a particular shift is needed

// +26 before % 26 takes care of negative shifts, turning them positive

// Check if the number of shifts 'i' can be performed within the maxShifts

int shiftCount[26] = {}; // There are 26 possible shifts (for 26 letters)

// Calculate the shift needed to convert source[i] to target[i]

bool canConvertString(string source, string target, int maxShifts) {

// % 26 ensures the shift is in the range [0, 25]

// Iterate over all possible shifts except 0 (no shift needed)

// If all shifts can be performed given the constraints, return true

If the lengths of the input strings differ, conversion is not possible

Check if the shifts can be achieved within the allowed operations 'k'

print(result) # Output: True or False depending on whether the conversion is possible

Calculate the maximum number of operations needed for this shift

For multiple shifts 'i', we need to wait 26 more for each additional use

If max operations exceeds 'k', then it's not possible to convert the string

Initialize an array to keep count of the number of shifts for each character

Calculate the shift difference and take modulo 26 to wrap around the alphabet

def can convert string(self, s: str, t: str, k: int) -> bool:

shift = (ord(char t) - ord(char s) + 26) % 26

Increment the count of the corresponding shift

 $max_{operations} = i + 26 * (shift_{counts}[i] - 1)$

result = solution instance.can convert string("input1", "input2", 10)

Iterate over the characters of both strings

int shift = (target[i] - source[i] + 26) % 26;

// Increment the count of the shift needed

if (source.size() != target.size()) {

// Iterate over the characters of the strings

for (int i = 0; i < source.size(); ++i) {</pre>

return false;

++shiftCount[shift];

for (int i = 1; i < 26; ++i) {

If max operations exceeds 'k', then it's not possible to convert the string

needed exceeds k. Thus, the false outcome would be the correct answer for this example.

Step 2: Frequency Array

Step 1: Length Check

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Step 3: Calculate Shifts
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For s[0]: 'a' to t[0]: 'b' requires 1 shift.

For s[1]: 'b' to t[1]: 'c' requires 1 shift.

∘ For s[2]: 'c' to t[2]: 'd' requires 1 shift.

• We iterate through cnt starting from index 1:

- We create an array cnt with length 26, initialized to zero: cnt = [0] * 26.

 The shift calculations: Shift for 'a' to 'b' is (ord('b') - ord('a') + 26) % 26 = 1.

Step 4: Count the Shifts For each shift calculated, increment the corresponding index in cnt:

Shift for 'b' to 'c' is (ord('c') - ord('b') + 26) % 26 = 1.

Shift for 'c' to 'd' is (ord('d') - ord('c') + 26) % 26 = 1.

We loop over s and t in parallel and calculate the shift required:

t within the move limit. Thus, the function returns True.

- For shift 1: cnt[1] becomes 3 because we need to shift by 1 three times.
- Solution Implementation

Python

class Solution:

Step 5: Move Validation

def can convert string(self, s: str, t: str, k: int) -> bool: # If the lengths of the input strings differ, conversion is not possible if len(s) != len(t): return False

for char s, char t in zip(s, t):

shift_counts[shift] += 1

if max operations > k:

shift = (ord(char t) - ord(char s) + 26) % 26

Increment the count of the corresponding shift

 $max_{operations} = i + 26 * (shift_{counts}[i] - 1)$

Initialize an array to keep count of the number of shifts for each character shift counts = [0] * 26# Iterate over the characters of both strings

∘ For shift amount i = 1 and frequency cnt[1] = 3, calculate maximum move needed: i + 26 * (cnt[1] - 1) = 1 + 26 * (3 - 1) = 53

As we can see, for the sample s and t, it would not be possible to transform s into t with only k moves since the latest move

```
# Check if the shifts can be achieved within the allowed operations 'k'
for i in range(1, 26):
    # Calculate the maximum number of operations needed for this shift
```

return true;

```
return False
        # If all shifts are possible within the operations limit, return True
        return True
# Example usage:
# solution instance = Solution()
# result = solution instance.can convert string("input1", "input2", 10)
# print(result) # Output: True or False depending on whether the conversion is possible
Java
class Solution {
    // Method to determine if it's possible to convert string s to string t
    // by shifting each character in s, at most k times.
    public boolean canConvertString(String s, String t, int k) {
        // Return false if the lengths of the strings are different.
        if (s.length() != t.length()) {
            return false;
        // Arrav to keep track of how many shifts for each letter are required.
        int[] shiftCounts = new int[26];
        // Calculate the shift required for each character to match the target string.
        for (int i = 0; i < s.length(); ++i) {</pre>
            int shift = (t.charAt(i) - s.charAt(i) + 26) % 26;
            ++shiftCounts[shift];
        // Check for each shift if it's possible within the allowed maximum of k shifts.
        for (int i = 1; i < 26; ++i) {
            // If the maximum shift needed for any character is more than k, return false.
            if (i + 26 * (shiftCounts[i] - 1) > k) {
                return false;
```

C++

public:

class Solution {

```
if (i + 26 * (shiftCount[i] - 1) > maxShifts) {
                return false; // If not possible, return false
       // If all shifts can be performed, return true
       return true;
};
TypeScript
function canConvertString(source: string, target: string, maxShifts: number): boolean {
   // If the lengths of source and target are not the same, then conversion is not possible
   if (source.length !== target.length) {
        return false;
   // Initialize an array to count the number of times a particular shift is needed
   // There are 26 possible shifts (for 26 letters in the alphabet)
    let shiftCount: number[] = new Array(26).fill(0);
   // Iterate over the characters of the strings
   for (let i = 0; i < source.length; i++) {</pre>
       // Calculate the shift needed to convert source[i] to target[i]
       // % 26 ensures the shift is in the range [0, 25]
       // Adding 26 before % 26 takes care of negative shifts by making them positive
        let shift = (target.charCodeAt(i) - source.charCodeAt(i) + 26) % 26;
       // Increment the count of the shift needed
        shiftCount[shift]++;
   // Iterate over all possible shifts except 0 (no shift needed)
   for (let i = 1; i < 26; i++) {
       // Check if the number of shifts 'i' can be performed within the maxShifts
       // This considers every 26th shift because the same letter can't be shifted again until 26 other shifts have occurred
       if (i + 26 * (shiftCount[i] - 1) > maxShifts) {
            return false; // If not possible, return false
```

// This considers every 26th shift because the same letter can't be shifted until 26 others have occurred

If all shifts are possible within the operations limit, return True return True # Example usage:

solution instance = Solution()

if len(s) != len(t):

return False

 $shift_counts = [0] * 26$

for i in range(1, 26):

for char s, char t in zip(s, t):

shift_counts[shift] += 1

if max operations > k:

return False

return true;

class Solution:

Time and Space Complexity The time complexity of the code is O(n), where n is the length of the strings s and t. This is because there is a single loop that iterates over the characters of s and t only once, and the operations within the loop are of constant time. The second loop is

not dependent on n and iterates up to a constant value (26), which does not affect the time complexity in terms of n. The space complexity of the code is 0(1), as there is a fixed-size integer array cnt of size 26, which does not scale with the input size. The rest of the variables use constant space as well.