2475. Number of Unequal Triplets in Array

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Hash Table
Easy
        <u>Array</u>
```

Problem Description

You are provided with an array of positive integers called nums. Each element of the array has a 0-based index. The task is to find

out how many groups of three indices (i, j, k) exist such that the following conditions are met: 1. The indices are in strictly increasing order, i.e., i < j < k.

- 2. The values of nums at these indices are pairwise distinct. This means that nums[i], nums[j], and nums[k] should all be different from each other.
 - In other words, for a triplet (i, j, k) to be counted, it must consist of three unique numbers from the nums array, each coming from a unique position in the array, where i, j, and k represent the positions (indices) of these numbers. The goal is to return the count of all such triplets.
- Intuition

To solve this problem, we need to find a way to calculate the number of distinct triplets without having to manually check each

possible combination, which would be inefficient especially for large arrays. We can use the following approach:

1. Count the frequency of each number in the array using a data structure like a dictionary (Counter in Python). 2. Iterate through each unique number's frequency, calculating how many distinct triplets can be formed with it.

- The intuition stems from the fact that if we know how many times a particular number occurs, and how many numbers are before
- and after it in the array (that haven't been used in a triplet yet), we can calculate all the valid triplets it can form. Here are the steps taken in the solution code:

• Use Counter to get the frequency of each distinct number in nums. • Initialize two counters, ans and a. ans will hold the final count of triplets, and a will keep track of the count of numbers processed so far.

• The number of triplets we can form with the current number as the middle element is a * b * c. This product comes from selecting one of

Let b represent the frequency of the current number.

Loop through the frequency count of each distinct number:

- Calculate c as the number of remaining elements that come after the current number, which is n a b (where n is the total length of nums).
- Add this triplet count to the ans. Update a to include the current number as well by adding b to it.
- The final answer is contained in ans, so we return it. **Solution Approach**

• We use the Counter class from Python's collections module to efficiently count the frequency of each element in the nums array.

the numbers before the current one (a choices), the current number itself (b choices), and one of the numbers after (c choices).

- The solution approach is based on the concept of counting and combinatorics. To implement the solution, the following
- components are used:

Data Structures:

Variables:

• cnt: A Counter object that holds the frequency of each unique element in nums. • n: The total number of elements in nums.

a: Keeps track of the number of elements processed so far (left side of the current element).

Algorithm:

processed.

n - a - b.

 b: The frequency of the current element being considered as the middle element of a triplet. o c: Represents the number of elements that are available to be picked after the current element.

3. Iterate over each count b in the values of cnt (since keys represent the unique numbers and values their respective counts):

Calculate the number of triplets possible with the current number as the middle element as a * b * c and add this to ans.

Each step of the algorithm is designed to avoid checking each triplet individually by leveraging the frequency counts and

1. Count the frequency of each unique number using cnt = Counter(nums). 2. Initialize a variable ans to accumulate the total number of valid triplets and a variable a to keep track of the count of numbers already

• ans: An accumulator to sum the number of valid triplets found.

• A dictionary-like container maps unique elements to their respective frequencies.

- Update a to include the element just processed by adding b to it. 4. After the loop ends, ans holds the total number of valid triplets, which is returned as the final answer.
- positions of numbers to calculate the number of potential triplets directly. This leads to a more efficient implementation that can handle large arrays without incurring a performance penalty typical of brute-force solutions.

Calculate the number of available elements after the current element (which could potentially be the third element of the triplet) as c =

Example Walkthrough Let's consider an example nums array to illustrate the solution approach:

From the problem statement, we want to count distinct triplets (i, j, k) such that nums[i], nums[j], and nums[k] are all unique. 1. First, we use Counter to get the frequency of each unique number in nums. from collections import Counter cnt = Counter(nums) # cnt will be Counter($\{1: 2, 2: 2, 3: 1\}$)

Now, we iterate over each count b in the values of cnt. The dictionary cnt has elements along with their frequencies: {1: 2,

ans = 0

2: 2, 3: 1}.

• First iteration (for number 1):

Second iteration (for number 2):

■ b for number 2 is 2.

• b for number 3 is 1.

ans remains 4.

Solution Implementation

Python

a = 0

nums = [1, 1, 2, 2, 3]

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■ b for number 1 is 2.
 c = n - a - b  which is  5 - 0 - 2 = 3 . There are three elements that can come after the first 1 to form a triplet.
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c = n - a - b which is 5 - 2 - 2 = 1. There is one element that can come after 2 to form a triplet.

■ Triplet count for 2 is a * b * c = 2 * 2 * 1 = 4. We can form 4 triplets with 2 as the middle element.

■ Update a = a + b to 2 (since we've now processed two occurrences of 1).

Initialize the variables ans for the total number of valid triplets and a for counting processed numbers.

We set n to the total number of elements in nums, which is 5.

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■ Update ans by adding the triplet count, ans = ans + 4 to 4.
    ■ Update a = a + b to 4 (as we've now processed two occurrences of 2).

    Third iteration (for number 3):
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c = n - a - b which is 5 - 4 - 1 = 0. There are no elements left to form triplets with 3 as the middle element.

■ Triplet count for 1 is a * b * c = 0 * 2 * 3 = 0 (since a is 0, no triplets can be formed with 1 as the middle element yet).

After the iteration ends, ans holds the total number of valid triplets. For this example, there are 4 valid triplets, and thus we return 4.

Initialize the answer and the count for the first number in the triplet

Increment count_first_number for the next iteration of the loop

// If the number is already in the map, increment its frequency,

// Initialize the answer variable to store the count of unequal triplets

// Variable 'prefixCount' is used to keep track of the count of numbers processed so far

// Calculate the count of numbers remaining after excluding the current number

// Update the answer by adding the number of unequal triplets that can be formed

// Update the prefix count by adding the frequency of the current number

Count of the third number in the triplet is total_numbers

minus the count_first_number and count_second_number

triplets without having to check each possible combination of three numbers.

def unequalTriplets(self, nums: List[int]) -> int:

number_counts = Counter(nums)

answer = count_first_number = 0

total_numbers = len(nums)

Count the frequency of each number in nums

Iterating through the counts of each unique number

for count_second_number in number_counts.values():

count_first_number += count_second_number

Return the total number of unequal triplets

// otherwise insert it with frequency 1

frequencyMap.merge(num, 1, Integer::sum);

// 'n' is the total number of elements in the input array

int suffixCount = n - prefixCount - frequency;

answer += prefixCount * frequency * suffixCount;

■ Triplet count for 3 is a * b * c = 4 * 1 * 0 = 0. No additional triplets can be formed.

from collections import Counter class Solution:

To conclude, using this example with nums = [1, 1, 2, 2, 3], the algorithm efficiently found the total number of 4 distinct

count_third_number = total_numbers - count_first_number - count_second_number # Calculate the number of triplets where the first, second, and third numbers # are all different answer += count_first_number * count_second_number * count_third_number

```
class Solution {
    public int unequalTriplets(int[] nums) {
       // Create a map to store the frequency of each number in the array
       Map<Integer, Integer> frequencyMap = new HashMap<>();
```

Java

return answer

for (int num : nums) {

int answer = 0;

int prefixCount = 0;

int n = nums.length;

// Iterate through the frequency map

prefixCount += frequency;

const lengthOfNums = nums.length;

for (const num of nums) {

let result = 0;

return result;

let accumulated = 0;

const countMap = new Map<number, number>();

// Initialize variable to hold the result

// Update the accumulated count

accumulated += currentCount;

number_counts = Counter(nums)

total_numbers = len(nums)

for (const currentCount of countMap.values()) {

// Return the total number of unequal triplets

def unequalTriplets(self, nums: List[int]) -> int:

Count the frequency of each number in nums

// Initialize a map to keep count of each number's occurrences

// Iterate over the map to calculate the answer using the formula

// Calculate the count for the third element in the triplet

const remaining = lengthOfNums - accumulated - currentCount;

// Calculate the number of unequal triplets for the current iteration

// 'accumulated' will keep track of the accumulated counts for each processed number

// Count the occurrences of each number in the nums array

countMap.set(num, (countMap.get(num) ?? 0) + 1);

result += accumulated * currentCount * remaining;

for (int frequency : frequencyMap.values()) {

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// Return the final count of unequal triplets
       return answer;
C++
#include <vector>
#include <unordered_map>
using namespace std;
class Solution {
public:
   // This function counts the number of unequal triplets within the input vector.
   int unequalTriplets(vector<int>& nums) {
       // Create a hash table to keep track of the count of each number in the vector.
       unordered_map<int, int> numCounts;
       // Increment the count for each value in the nums vector.
        for (int value : nums) {
           ++numCounts[value];
        int result = 0; // Variable to store the result.
        int accumulated = 0; // Accumulated counts of previous numbers.
       // Calculate the number of unequal triplets.
        for (auto& [value, count] : numCounts) {
            int remaining = nums.size() - accumulated - count; // Count of numbers that are not equal to current number.
            result += accumulated * count * remaining; // Multiply by count to form unequal triplets.
            accumulated += count; // Update accumulated with count of the current number.
       // Return the total number of unequal triplets.
       return result;
};
TypeScript
function unequalTriplets(nums: number[]): number {
   // Find the length of the nums array
```

```
from collections import Counter
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class Solution:

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# Initialize the answer and the count for the first number in the triplet
       answer = count_first_number = 0
       # Iterating through the counts of each unique number
       for count_second_number in number_counts.values():
           # Count of the third number in the triplet is total_numbers
           # minus the count_first_number and count_second_number
           count_third_number = total_numbers - count_first_number - count_second_number
           # Calculate the number of triplets where the first, second, and third numbers
           # are all different
           answer += count_first_number * count_second_number * count_third_number
           # Increment count_first_number for the next iteration of the loop
           count_first_number += count_second_number
       # Return the total number of unequal triplets
        return answer
Time and Space Complexity
Time Complexity
```

The time complexity of the code is primarily determined by the number of operations within the for loop which iterates over the values of the Counter object cnt. The Counter object itself is created by iterating over the nums list once, which takes O(n) time

where n is the number of elements in nums. Inside the for loop, each operation is constant time, so the overall time complexity of the for loop is 0(k), where k is the number

Therefore, the total time complexity is O(n) + O(k) = O(n) as the first iteration to create Counter and the second iteration over unique values can be bounded by n.

of unique elements in nums. Since k can vary from 1 to n, in the worst case where all elements are unique, k is equal to n.

elements in nums.

Thus, the overall space complexity of the code is O(n).

Space Complexity The space complexity is affected by the storage used for the Counter object cnt which stores the frequency of each unique element in nums. In the worst case, if all elements of nums are unique, the Counter would take O(n) space where n is the number of