1609. Even Odd Tree Medium Tree Breadth-First Search Binary Tree Leetcode Link

Problem Description The problem provides a binary tree and introduces a special condition called Even-Odd. A binary tree meets the Even-Odd condition

if: The root node is at level 0, its children are at level 1, their children at level 2, and so on.

- Nodes at levels with an even index (like 0, 2, 4, ...) must have odd integer values, and these values must be in strictly increasing
- order from left to right. Nodes at levels with an odd index (like 1, 3, 5, ...) must have even integer values, and these values must be in strictly decreasing order from left to right.
- The goal is to write a function that takes the root node of a binary tree as an input and returns true if the tree satisfies the Even-Odd condition; otherwise, it should return false.

Intuition

To solve this problem, we use the Breadth-First Search (BFS) algorithm. BFS allows us to traverse the tree level by level. This fits our needs perfectly since the Even-Odd condition applies to individual levels of the tree.

Here's how we approach the solution:

easily check the strictly increasing or decreasing order requirement.

 Maintain a prev value to compare with the current node's value ensuring the strictly increasing or decreasing order. After processing all nodes in the current level, toggle the even/odd level flag using XOR with 1. 3. If at any point a node does not meet the requirements, return false.

2. Traverse each level of the tree using a while loop that runs as long as there are nodes left in the queue:

1. Initialize a queue and add the root node to it. The queue is used to keep track of nodes to visit in the current level.

For each node, check if the value meets the required criteria based on whether the current level is even or odd.

- 4. If the loop completes without finding any violations, return true since the tree meets the Even-Odd condition. By using BFS, we are able to check the values level by level, and by maintaining a variable to track the previous node's value, we can
- Solution Approach
- The solution approach is based on a Breadth-First Search (BFS) pattern. Here's a walkthrough of the implementation details: 1. Initialize Variables: A flag variable even, which is initially set to 1, is used to represent the current parity of the level being

traversed (odd or even). A queue q is used to keep track of nodes in the current level. = deque([root])

2. BFS Loop: The solution iterates through the tree by using a while loop that continues as long as there are nodes in the queue.

1 while q:

1 for _ in range(len(q)):

1 prev = 0 if even else inf

1 if root.left:

3 if root.right:

1 even ^= 1

Example Walkthrough

q.append(root.left)

q.append(root.right)

done by using the XOR operator (even ^= 1).

the tree against the provided conditions level by level.

Suppose we have the following binary tree:

Let's take a small binary tree to illustrate the solution approach.

3. Level Traversal: Within the loop, a for loop iterates over each node in the current level. The number of nodes is determined by the current size of the queue (len(q)).

- 4. Initialization Per Level: At the beginning of each level, initialize prev to 0 if the current level is supposed to contain even values, and to inf for odd values.
- 5. Node Validation: For each node, the algorithm checks if the current node's value meets the appropriate conditions:

For even levels, node values must be odd and increasing, checked by root.val % 2 == 0 or prev >= root.val.

For odd levels, node values must be even and decreasing, checked by root.val % 2 == 1 or prev <= root.val.

- 1 if even and (root.val % 2 == 0 or prev >= root.val): 3 if not even and (root.val % 2 == 1 or prev <= root.val):</pre> return False
- 1 prev = root.val 7. Queue Update: The left and right children of the current node are added to the queue if they exist.

If a node does not meet its corresponding condition, the function returns False.

6. Track Previous Value: The prev value is updated to the current node's value for the next iteration.

satisfies the Even-Odd condition. By using a queue to manage nodes, this approach maintains a clear delineation between levels and employs BFS efficiently to verify

9. Completion: After the entire tree has been traversed without returning False, the function returns True, confirming that the tree

8. Toggle Even/Odd Level: After each level is processed, the even flag is toggled to 1 if it was 0, or to 0 if it was 1. The toggle is

Let's walk through the steps:

1. Initialize Variables: Set even to 1 and put the root node with value 5 in the queue. 2. BFS Loop: Start the while loop since the queue is not empty.

Since it's an even level, we expect an odd value and as it's the first value, it doesn't need to be increasing. The root's value is

We move on to the next node in the same level, which is 7. We expect an even value and a value less than prev (which is 10),

5, which is odd, so it passes. The function continues. 6. Track Previous Value: The prev value becomes the current node's value, which is 5.

Python Solution

class TreeNode:

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from collections import deque

level = 0

while queue:

Definition for a binary tree node.

queue = deque([root])

Traverse the tree by levels

for _ in range(len(queue)):

node = queue.popleft()

10. Level Traversal: There are two nodes in this level 11. Initialization Per Level: Set prev to 0 because we are at an odd level now.

7. Queue Update: Nodes 10 and 7 are children of 5, so add them to the queue.

9. BFS Loop: Now we have 10 and 7 in the queue; the while loop continues.

8. Toggle Even/Odd Level: At the end of the level, toggle even to 0.

3. Level Traversal: Process nodes in the current level. We have one node, which is 5.

4. Initialization Per Level: At level 0 (even level), we initialize prev to inf.

5. Node Validation: We check the value of the root node:

12. Node Validation: Evaluate the nodes in this odd level:

but 7 is odd, so it fails the condition.

We start with the node 10. We expect an even value and decreasing order (however, as it's the first node of this level, we

only check for an even value). Node 10 has an even value, so it passes for now.

This small example contradicts the Even-Odd condition in the second level and illustrates the checking mechanism per level using BFS traversal.

def __init__(self, val=0, left=None, right=None):

def isEvenOddTree(self, root: TreeNode) -> bool:

Initialize the queue with the root node

Initialize the level to 0 (0-indexed, so even)

Since we detected a violation at node 7, the function should return False.

self.val = val self.left = left self.right = right class Solution:

Depending on the current level, set the previous value accordingly

Check the Even-Odd Tree condition for the current node

Update the previous value for the next comparison

At even levels, values must be odd and strictly increasing

previous_value = 0 if level % 2 == 0 else float('inf')

For even levels, we start comparing from the smallest possible value (0)

For odd levels, we start comparing from the largest possible value (infinity)

31 if level % 2 == 0 and (node.val % 2 == 0 or previous_value >= node.val): return False 32 33 # At odd levels, values must be even and strictly decreasing if level % 2 == 1 and (node.val % 2 == 1 or previous_value <= node.val): 34 35 return False 36

previous_value = node.val

If all conditions are met, return True

if node.left:

if node.right:

Move to the next level

return false;

previousValue = root.value;

queue.offer(root.left);

queue.offer(root.right);

// If all levels meet the condition, return true

TreeNode() : val(0), left(nullptr), right(nullptr) {}

* Checks whether a binary tree is an even-odd tree.

increasing from left to right.

decreasing from left to right.

bool isEvenOddTree(TreeNode* root) {

bool isEvenLevel = true;

* @param root The root of the binary tree.

// Queue for level order traversal.

queue<TreeNode*> nodeQueue{{root}};

nodeQueue.pop();

isEvenLevel = !isEvenLevel;

root = nodeQueue.front();

TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}

TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left), right(right) {}

* A binary tree is an even-odd tree if it meets the following conditions:

* - At even-indexed levels, all nodes' values are odd, and they are

* - At odd-indexed levels, all nodes' values are even, and they are

* @return True if the binary tree is an even-odd tree, otherwise false.

// 'isEvenLevel' indicates whether the current level is even or odd.

// Check whether the current node's value is appropriately odd or even.

previousValue = root->val; // Update the 'previousValue'.

// Add child nodes to the queue for the next level.

// Toggle the level indicator after finishing each level.

if (root->left) nodeQueue.push(root->left);

if (root->right) nodeQueue.push(root->right);

return true; // The tree meets the even-odd tree conditions.

if (isEvenLevel && (root->val % 2 == 0 || previousValue >= root->val)) return false;

if (!isEvenLevel && (root->val % 2 == 1 || previousValue <= root->val)) return false;

if (root.left != null) {

if (root.right != null) {

isLevelEven = !isLevelEven;

return true;

* Definition for a binary tree node.

C++ Solution

4 struct TreeNode {

13 class Solution {

/**

*/

int val;

TreeNode *left;

TreeNode *right;

1 /**

3 */

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14 public:

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// Update the 'previousValue' with the current node's value

// Toggle the level indication for the next level traversal

// Add the left and right children, if they exist, to the queue for the next level

level += 1

return True

Add child nodes to the queue

queue.append(node.left)

queue.append(node.right)

Process all nodes in the current level

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Java Solution
   // Definition of the binary tree node.
   class TreeNode {
       int value;
       TreeNode left;
       TreeNode right;
       TreeNode() {}
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       TreeNode(int value) {
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           this.value = value;
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       TreeNode(int value, TreeNode left, TreeNode right) {
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           this.value = value;
14
           this.left = left;
15
           this.right = right;
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18 }
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   class Solution {
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       public boolean isEvenOddTree(TreeNode root) {
           // Tracks whether the current level is even (starting with the root level as even).
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           boolean isLevelEven = true;
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           // Queue for performing level order traversal
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           Deque<TreeNode> queue = new ArrayDeque<>();
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            queue.offer(root);
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           // Loop while there are nodes in the queue
           while (!queue.isEmpty()) {
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               // 'previousValue' will store the last value seen at the current level to check the strictly increasing or decreasing orc
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                int previousValue = isLevelEven ? 0 : 1000001;
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               // Number of nodes at the current level
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                for (int levelSize = queue.size(); levelSize > 0; --levelSize) {
                    root = queue.poll();
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                    // For even levels, values must be odd and strictly increasing
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                    if (isLevelEven && (root.value % 2 == 0 || previousValue >= root.value)) {
                        return false;
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                    // For odd levels, values must be even and strictly decreasing
                    if (!isLevelEven && (root.value % 2 == 1 || previousValue <= root.value)) {</pre>
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32 while (!nodeQueue.empty()) { 33 // 'previousValue' holds the previously encountered value in the current level. 34 int previousValue = isEvenLevel ? 0 : INT_MAX; // Initialize based on the level. // Process all nodes at the current level. 35 36 for (int i = nodeQueue.size(); i > 0; --i) {

Typescript Solution 1 interface TreeNode { val: number; left: TreeNode | null; right: TreeNode | null; 6 function isEvenOddTree(root: TreeNode | null): boolean { // 'isEvenLevel' indicates whether the current level is even or odd. let isEvenLevel = true; // Queue for level order traversal. 10 let nodeQueue: (TreeNode | null)[] = [root]; 11 12 13 while (nodeQueue.length > 0) { // 'previousValue' holds the previously encountered value in the current level. 14 let previousValue: number = isEvenLevel ? 0 : Number.MAX_SAFE_INTEGER; // Initialize based on the level. 15 16 // Process all nodes at the current level. 17 let levelSize = nodeQueue.length; 18 for (let i = 0; i < levelSize; i++) {</pre> let node = nodeQueue.shift()!; // '!' asserts that 'node' won't be null. 19 20 // Check whether the current node's value is appropriately odd or even. if (isEvenLevel && (node.val % 2 === 0 || previousValue >= node.val)) return false; 21 22 if (!isEvenLevel && (node.val % 2 === 1 || previousValue <= node.val)) return false;</pre> 23 24 previousValue = node.val; // Update the 'previousValue'. // Add child nodes to the queue for the next level. 25 if (node.left) nodeQueue.push(node.left); 26 27 if (node.right) nodeQueue.push(node.right); 28 29 // Toggle the level indicator after finishing each level. isEvenLevel = !isEvenLevel; 30 31 32 33 return true; // The tree meets the even-odd tree conditions. 34 } 35 Time and Space Complexity

traverses each node exactly once. Each node is popped from the queue once, and its value is checked against the conditions for an even-odd tree, which takes constant time. The enqueue operations for the children of the nodes also happen once per node,

maintaining the overall linear complexity with respect to the number of nodes. // The space complexity of the code is O(m) where m is the maximum number of nodes at any level of the binary tree, or the maximum

// The time complexity of the provided code is O(n) where n is the number of nodes in the binary tree. This is because the code

breadth of the tree. This is because the queue q can at most contain all the nodes at the level with the maximum number of nodes at some point during execution, which determines the maximum amount of space needed at any time.