1052. Grumpy Bookstore Owner

use this secret technique only once in the day.

Medium Array Sliding Window

Problem Description

enter the store. The number of customers entering the store at the beginning of the ith minute is given by customers[i]. After each minute, all customers for that particular minute leave.

The bookstore owner's mood fluctuates, and this affects customer satisfaction. The mood is captured by an array grumpy. If

In this problem, we have a bookstore owner who has their store open for n minutes. During each of those minutes, customers

The bookstore owner's mood fluctuates, and this affects customer satisfaction. The mood is captured by an array grumpy. If grumpy[i] is 1, the owner is grumpy in the ith minute; if 0, they are not grumpy. When the owner is grumpy, customers leaving in that minute are not satisfied.

The task is to determine the maximum number of customers that can be satisfied throughout the day by smartly applying the secret technique to a certain time window during the day.

However, the owner knows a secret technique that allows them to be not grumpy for minutes consecutive minutes, but they can

Intuition

The solution aims to identify the best time segment where using the secret technique will result in the maximum number of satisfied customers. We can always satisfy customers when the owner is not grumpy, so these customers are not our focus.

to s.

To achieve this, we observe that the key is to apply the secret technique during the consecutively grumpy moments when the number of customers is highest. Here's how the solution approaches the problem:

1. We track the number of customers that could be satisfied if the owner were not grumpy at all by adding up the number of

Instead, we want to maximize the number of satisfied customers during the owner's grumpy periods.

customers during grumpy minutes. This gives us the total potential satisfaction if the owner were not grumpy.

- 2. Then, we create a <u>sliding window</u> with the width equal to <u>minutes</u> the duration for which the secret technique can be applied. As we slide this window across the grumpy minutes, we calculate the number of dissatisfied customers in this window.
- We keep track of the maximum number of these potentially dissatisfied customers that can be turned into satisfied customers by applying the secret technique in the current <u>sliding window</u>.
 We use a variable t to calculate the temporary sum of customers that are currently not satisfied due to the owner being
- window (if the owner was grumpy at that minute) and add the customers at the new end of the window (if the owner is grumpy at the new end).

We keep adjusting the best maximum number of satisfied customers found so far (ans) each time we slide the window one

grumpy within the window. Whenever we move the window forward by a minute, we deduct the customers at the start of the

- minute forward by subtracting the number of unsatisfied customers in the window from the total potential satisfaction.

 In the end, we return the maximum number of customers that can be satisfied which will include the number of customers that were satisfied when the owner was not grumpy and the additional customers that were satisfied when using the secret technique during the grumpy periods.
- Solution Approach

 The solution approach utilizes a <u>sliding window</u> to find the optimal <u>minutes</u> to apply the secret technique so as to maximize the

We first calculate s, the sum of customers that are not satisfied initially due to the owner being grumpy. This is done by taking

the dot product of the customers and grumpy arrays. Only the minutes when the owner is grumpy (grumpy [i] == 1) contribute

cs registers the total number of customers that visit the store throughout the day, which is the sum of the entire customers

satisfaction of customers. Here's how the implementation details map to this strategy:

by incrementing t by customers[i] only if grumpy[i] == 1.

that minute was a grumpy one (customers[j] * grumpy[j]).

customers, but we skip the 1st minute since the owner is not grumpy.

subtracting those at the starting minute if it was a grumpy period.

customers for that day, resulting in a total satisfaction of 33 customers.

Calculate the total number of customers

temp_improved += customer * grumpiness

for (int i = 0; i < numCustomers; ++i) {</pre>

totalGrumpyCustomers += customers[i];

if (grumpy[i] == 1) {

for (int i = 0; i < n; ++i) {

for (int i = 0; i < n; ++i) {

if (endTime >= 0) {

totalUnsatisfied = 0;

windowUnsatisfied = 0;

const n: number = customers.length;

totalCustomers += customers[i];

Calculate the total number of customers

total_customers = sum(customers)

return max_improved

for (let i = 0; i < n; i++) {

totalCustomers = 0;

maxIncrease = 0;

class Solution:

total_customers = sum(customers)

temp_improved = max_improved = 0

Enumerate through each minute

def maxSatisfied(self, customers: List[int], grumpy: List[int], minutes: int) -> int:

Calculate the initial unsatisfied customer count (when the owner is grumpy)

Tracking variable for the current window of grumpy period that can be improved

for minute_index, (customer, grumpiness) in enumerate(zip(customers, grumpy), 1):

temp_improved -= customers[start_index] * grumpy[start_index]

int totalCustomersSatisfied = 0; // Sum of all customers that are satisfied

int numCustomers = customers.length; // Total number of customers

The answer is the maximum potential customers satisfied in any "X" minute window

Increase the count of potentially satisfied customers when the owner is grumpy

Remove the front of the window (the oldest minute moving out of the window)

unsatisfied = sum(customer * grumpiness for customer, grumpiness in zip(customers, grumpy))

1 customers.

Python

class Solution:

array without taking the owner's mood into account.

3. Two variables t and ans are initiated to zero. t will capture the total number of dissatisfied customers within the current

window, and ans will store the maximum number of customers that can be satisfied during the best window of minutes.

We then iterate through each minute i, simultaneously calculating the sliding window sum t for grumpy minutes. We do this

In each iteration, we check if the <u>sliding window</u> has extended beyond the prescribed <u>minutes</u>. This is done using a condition
 (j := i - minutes) >= 0. If true, it means we have a full window and we are moving it one step forward.
 We calculate the potential maximum number of satisfied customers <u>ans</u> by subtracting from <u>cs</u> the total dissatisfaction

without the current window (s - t). If it exceeds our previous best (ans), we update ans.

given problem that runs in linear time, based on the number of minutes the store is open (0(n)).

0], and the owner can be not grumpy for minutes = 3 consecutive minutes using the secret technique.

8. Once the loop is done, all possible windows of size minutes have been checked. The ans variable holds the maximum number of customers that can be satisfied by applying the secret technique during the best possible window. This value is returned.

Through the use of a simple loop and the sliding window technique, the implementation finds the optimal moment to use the

secret technique without the need for complex data structures or algorithms. This is an efficient and effective solution for the

After updating ans, we adjust t by subtracting the customers at the starting point of the previous window frame, but only if

Example Walkthrough

Let us consider the bookstore is open for n = 7 minutes, and the number of customers entering at each minute is given by

We first calculate the total potential satisfaction without using the secret technique, which is the sum of customers that visit when the owner is not grumpy. From our example, this is 1 (1st minute) + 1 (5th minute) + 10 (7th minute) = 12.
 Next, using a sliding window of size minutes = 3, we want to find a period where applying the secret technique can maximize customer satisfaction. We start with the window at the beginning:

• For the first window (minutes 1-3), the owner is in a good mood in the 1st minute, but grumpy in the 2nd and 3rd minute, where we have 10 + 10

We slide the window to the right by one minute (minutes 2-4), the owner is grumpy in all these minutes, resulting in 10 + 10 +

Continuing to slide our window to the right (minutes 3-5, minutes 4-6), we find the maximum number of grumpy customers

that would be satisfied by the secret technique. The window minutes 4-6 is where the owner is constantly grumpy with customers [10, 1, 10].

As we slide, we update the temporary sum t as we go, adding new customers who entered during a grumpy minute and

After sliding through all possible windows, the window minutes 4-6 has the highest number of grumpy customers (1 + 10 + 10 = 21) that could be turned into satisfied customers.
 The result is the sum of the initial satisfaction (12) and the additional satisfaction during the owner's grumpy period (21) when

Therefore, by applying the secret technique from the 4th to the 6th minute, the owner can maximize the number of satisfied

using the secret technique, which gives us the maximum customer satisfaction with the technique applied being 33.

- Solution Implementation
- # Slide the window of minutes for the "X" minute period
 if (start_index := minute_index minutes) >= 0:
 # Update the max potential customers satisfied for any "X" minute window
 max_improved = max(max_improved, total_customers (unsatisfied temp_improved))

// Calculate the initial total number of customers that are satisfied and the total affected by grumpiness

// If the owner is grumpy during customer i's visit, add to totalGrumpyCustomers

```
class Solution {
   public int maxSatisfied(int[] customers, int[] grumpy, int minutes) {
    int totalGrumpyCustomers = 0; // This will hold the total number of customers affected by grumpiness
```

Java

return max_improved

```
// Add to the total number of satisfied customers regardless of grumpiness
            totalCustomersSatisfied += customers[i];
        int temporaryGrumpyCustomers = 0; // Temporary variable to keep track of grumpy customers during the 'minutes' window
        int maxGrumpyCustomersConverted = 0; // Maximum number of grumpy customers that can be converted to satisfied customers
        for (int i = 0; i < numCustomers; ++i) {</pre>
            // If the owner is grumpy, temporarily add this customer to the count
            if (grumpy[i] == 1) {
                temporaryGrumpyCustomers += customers[i];
           // If we are beyond the 'minutes' window, subtract the customers at the start of the window
            if (i >= minutes) {
                if (grumpy[i - minutes] == 1) {
                    temporaryGrumpyCustomers -= customers[i - minutes];
            // Store the higher value between the current and previous maximum number of convertible grumpy customers
            maxGrumpyCustomersConverted = Math.max(maxGrumpyCustomersConverted, temporaryGrumpyCustomers);
       // Calculate final answer: Subtract the initial grumpy customers from the total customers satisfied and add the convertik
       return totalCustomersSatisfied - totalGrumpyCustomers + maxGrumpyCustomersConverted;
C++
class Solution {
public:
    int maxSatisfied(vector<int>& customers, vector<int>& grumpy, int minutes) {
        int totalUnsatisfied = 0; // To keep track of the total unsatisfied customers if the owner is grumpy
        int totalCustomers = 0; // To keep track of the total number of customers served satisfactorily whether the owner is grun
        int windowUnsatisfied = 0; // To keep track of unsatisfied customers in the current time window of size 'minutes'
        int maxIncrease = 0; // Variable to store the maximum increase in satisfied customers achievable by using the technique 1
```

int n = customers.size(); // Total number of customers, assuming customers and grumpy are of same size

int endTime = i - minutes; // Calculate end time of window (i - minutes + 1) - 1 for 0-based indexing

// Shrink the window from the left by removing the first element's impact if it was unsatisfied

maxIncrease = max(maxIncrease, totalCustomers - (totalUnsatisfied - windowUnsatisfied));

// After the loop, we must do one last check as the sliding window does not check the last window position

let maxIncrease: number = 0; // Maximum increase in satisfied customers possible by not being grumpy for 'minutes' duration

windowUnsatisfied += customers[i] * grumpy[i]; // Add to window unsatisfied customers

// If we've filled the first window, start calculating and updating maxIncrease

maxIncrease = max(maxIncrease, totalCustomers - (totalUnsatisfied - windowUnsatisfied));

// The answer is the maximum customers possible by using the technique in the best time window

// Calculate the maximum satisfied customers by using the technique

totalUnsatisfied += customers[i] * grumpy[i]; // Only add customers that are unsatisfied because the owner is grumpy

// Calculate the initial state, considering the owner is grumpy all the time

// Use a sliding window to determine when to utilize the owner's technique

windowUnsatisfied -= customers[endTime] * grumpy[endTime];

totalCustomers += customers[i]; // Add all customers

```
return maxIncrease;
}

};

TypeScript

let totalUnsatisfied: number = 0; // Track total unsatisfied customers if the owner is grumpy
let totalCustomers: number = 0; // Track total number of customers served satisfactorily
```

let windowUnsatisfied: number = 0; // Track unsatisfied customers in the current 'minutes' window

function maxSatisfied(customers: number[], grumpy: number[], minutes: number): number {

def maxSatisfied(self, customers: List[int], grumpy: List[int], minutes: int) -> int:

Calculate the initial unsatisfied customer count (when the owner is grumpy)

Tracking variable for the current window of grumpy period that can be improved

temp_improved -= customers[start_index] * grumpy[start_index]

The answer is the maximum potential customers satisfied in any "X" minute window

// Initial calculations without utilizing the special technique

totalUnsatisfied += customers[i] * grumpy[i];

```
// Using a sliding window to find the best time to apply the technique
for (let i = 0; i < n; i++) {
    windowUnsatisfied += customers[i] * grumpy[i];
    let endTime: number = i - minutes;

    if (endTime >= 0) {
        maxIncrease = Math.max(maxIncrease, totalCustomers - (totalUnsatisfied - windowUnsatisfied));
        windowUnsatisfied -= customers[endTime] * grumpy[endTime];
    }
}
// Handle the last possible window to make sure we don't miss the chance of optimizing at the end
maxIncrease = Math.max(maxIncrease, totalCustomers - (totalUnsatisfied - windowUnsatisfied));

// The result is the best scenario where the technique is applied in the optimal window
return maxIncrease;
```

```
temp_improved = max_improved = 0

# Enumerate through each minute
for minute_index, (customer, grumpiness) in enumerate(zip(customers, grumpy), 1):
    # Increase the count of potentially satisfied customers when the owner is grumpy
    temp_improved += customer * grumpiness

# Slide the window of minutes for the "X" minute period
    if (start_index := minute_index - minutes) >= 0:
        # Update the max potential customers satisfied for any "X" minute window
        max_improved = max(max_improved, total_customers - (unsatisfied - temp_improved))
```

Remove the front of the window (the oldest minute moving out of the window)

unsatisfied = sum(customer * grumpiness for customer, grumpiness in zip(customers, grumpy))

Time and Space Complexity

The time complexity of the provided code is O(N), where N is the length of the customers or grumpy lists. This is because the code

consists of a single loop that iterates over the entire length of the input lists exactly once. Within the loop, there are constant-time operations being performed, such as arithmetic operations, comparisons, and an assignment inside an if-statement that triggers based on the loop index.

The space complexity of the code is 0(1), meaning it requires constant additional space. The extra space used in the algorithm is for variables space, the area and if the latter found in the conditional if-statement). These variables do not depend on the size of the

for variables s, cs, t, ans, and j (the latter found in the conditional if-statement). These variables do not depend on the size of the input, hence the constant space complexity. No additional data structures or variable-sized containers are used that would require space proportional to the input size.