953. Verifying an Alien Dictionary String Hash Table Array **Leetcode Link** Easy

In the given problem, we're dealing with a hypothetical alien language that uses the same lowercase letters as English but has a

Problem Description

already sorted as per the alien dictionary.

the alien alphabetical order. We're given a list of words called words, which are supposedly written in this alien language. Our task is to determine if these words are sorted according to the alien lexicographical order. To put it simply, we need to confirm if each word would appear before the

different ordering of these letters. This order is represented by the string order, where each character's position indicates its rank in

following word in a dictionary organized by the provided alien alphabet. The lexicographical order is similar to what we're familiar with in an English dictionary—'apple' comes before 'banana' because 'a' is before 'b' in the English alphabetical order. However, for this problem, we need to ignore the English alphabetical order and instead

use the provided order string to determine the sorting sequence. To solve this, we have to read the words list using the alien alphabet and confirm that each subsequent word is either the same up to

a point and then greater according to the alien alphabet, or entirely greater if compared at the same position in each word.

Intuition

In essence, we need to write a function that checks this sorting without reordering the list—a simple 'yes' or 'no' to whether the list is

The intuition behind the provided solution approach is to use a dictionary named m to represent the mapping of each character in the alien language to its corresponding index according to order. Once we have this mapping, it is much easier to compare the words.

During the comparison, we want to ensure that for each position i up to the length of the longest word (in this case, arbitrarily

chosen as 20, assuming no word is longer than 20 characters), the following holds true:

order, we can return True as this means the list is sorted correctly up to this character.

• If the character at position i in the previous word has a higher index in our alien dictionary than the character at position i in the current word, we should immediately return False, as this indicates that the list is not sorted.

• If characters at position i are the same, we continue our checks without concluding anything just yet because the subsequent characters might determine the order.

If at this position i, we do not encounter any characters that are out of order, and all comparisons were either equal or in sorted

- If after checking all positions i, we haven't returned False, it is safe to assume that the list is sorted according to the alien language, and we return True.
- beyond the length of the shortest word, due to the handling of current and previous character index (curr and prev) with default values ensuring proper comparison even if one of the words is shorter.

The code smartly exits early when it finds evidence the words are not in order, and avoids unnecessary comparisons for characters

Solution Approach The solution provided makes use of a hash table (or dictionary in Python) and single-pass comparisons to implement the algorithm.

1. Mapping the Order: We start by creating a dictionary m where the key is each character of the alien alphabet and the value is its

corresponding rank in the alien language. 1 m = {c: i for i, c in enumerate(order)} Using enumerate, each character c is paired with its index i, effectively creating a rank ordering for the alien alphabet.

2. Comparing Characters in Words: The code uses a for loop to iterate over the positions of the characters in the words, up to a

predefined limit of 20 (assuming that none of the words will be longer than that). For each position i, it compares characters

position in the order string. This mapping allows us to efficiently translate each character in the alien words into its

1 for i in range(20):

1 prev = -1

2 if prev > curr:

1 if prev == curr:

we continue.

1 if valid:

1 return True

valid = False

return True

Here's a step-by-step breakdown of the implementation:

from different words at the same index.

(prev) would mean the list is not sorted.

1 curr = -1 if i >= len(x) else m[x[i]]

3. Tracking and Comparing Ranks of Characters: Inside the loop, two variables prev and curr are used to track the rank of the

4. Checking Sorted Order: As we iterate over each word x in words, we check if the current character's rank is less than the

previous one. Since we are performing lexicographical comparison, any current character (curr) with a lower rank than previous

current and previous characters, respectively. Initially, prev is set to -1 to ensure that the first comparison is valid.

return False 5. Validating Sequential Characters: If the previous character is the same as the current one, valid is set to False, valid is used to track whether all characters compared up to that point have been the same, in which case we need to check further.

6. Early Termination: If at any point, we successfully pass through all the words and their characters at a certain index without

7. Final Verdict: If we reach the end of our positional checks (the for loop) without finding any characters out of order, we

Key algorithms and data structures used in this implementation include the hash table for constant-time look-up, and a comparison

Let's consider a small example to illustrate the solution approach. Imagine we are given the following alien order and a list of words:

Using the provided order, we create a mapping m that will let us compare the characters based on their alien language order.

framework that examines each word character-by-character using its index. The process of early termination after successful

conclude that the words list is sorted according to the alien dictionary.

We will walk through checking if the words are sorted according to this alien language.

○ The mapping would result in: {'h': 0, 'l': 1, 'a': 2, ... 'z': 25}.

We compare 'l' (at position 0) of "leetcode" with 'h' of "hello".

We set up a loop that will go up to 20, which is more than enough for our example.

returning False, and valid remains True, it means we found a lexicographical order and can return True immediately. Otherwise,

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sequential checks leverages a common pattern in sorting algorithms to optimize performance by avoiding unnecessary work.
Example Walkthrough
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1 order = "hlabcdefgijkmnopqrstuvwxyz"

2. Comparing Characters in Words:

Next is "leetcode":

6. Early Termination:

False.

dictionary.

Python Solution

for i in range(20):

class Solution:

alphabet given.

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38 }

};

private:

return true;

return len1 <= len2;</pre>

for (const char of order) {

Typescript Solution

5. Validating Sequential Characters:

characters to determine the order.

3. Tracking and Comparing Ranks of Characters:

words = ["hello", "leetcode"]

1. Mapping the Order:

 \circ We initialize prev = -1 to indicate the start of comparisons. We will compare the ranks of characters at each position within the words. 4. Checking Sorted Order:

■ In the mapping m, 'I' has a rank of 1 and 'h' has a rank of 0. Since 1 is greater than 0, the order is maintained.

When comparing characters at the same position, if they are equal, we proceed with further positions as we need more

∘ In our example, if we found 'g' in "hello" greater than 'e' in "leetcode" at any given position, we would terminate and return

If all characters are equal up to the minimum length of the two words, we also proceed to the next word in the list.

After checking all positions and finding no disorder, we return True, meaning the list is sorted according to the alien

In practice, the comparison process stops at the first instance where prev > curr. As "hello" and "leetcode" are in correct order in

this alien language (since 'h' < 'l'), the result of our checks would be True. The words are indeed sorted according to the alien

Create a mapping from each alien character to its order index for easy comparison.

Go through all the words to compare the characters at the current index i.

If all words are sorted until now, no need to check further characters.

// Method to check if the words are sorted as per the given alien dictionary order

int previousIndex = -1; // Initialize previous character index as -1

// Subtract 'a' to convert char into an index (0-25), then assign its order

for (int i = 0; i < 20; ++i) { // Iterate at most up to 20 times, the maximum word length

// If previous index is greater than the current index, words are not sorted

previousIndex = currentIndex; // Update previousIndex for next iteration

// If all words had distinct characters or reached the end, exit from the loop

previous_index = -1 # Start with a comparison index of -1 for the first character.

Iterate up to 20 characters, assuming no word is longer than 20 characters.

Continue the loop until either the words are found to be in disorder or all characters are successfully compared.

Starting with "hello", we get the rank of 'h' as 0 for position 0. There's no previous word, so we move on.

 However, since 'h' is less than 'l' at the first position and there are no other characters in disorder, we do not terminate early and continue. 7. Final Verdict:

def isAlienSorted(self, words: List[str], order: str) -> bool:

if previous_index > current_index:

if previous_index == current_index:

is_sorted_until_now = False

public boolean isAlienSorted(String[] words, String order) {

charIndexMapping[order.charAt(i) - 'a'] = i;

if (previousIndex > currentIndex) {

if (previousIndex == currentIndex) {

// If all the words are in the correct order, return true

int char1Index = charIndexMap[word1[i] - 'a'];

int char2Index = charIndexMap[word2[i] - 'a'];

// If characters are not in the correct order, return false

// Create a mapping of each character to its position in the new language order

const minLength = Math.min(firstWord.length, secondWord.length);

if (!areWordsEqual && firstWord.length > secondWord.length) {

The time complexity of the given code consists of a few different components:

if (charPositionMap.get(firstWord[j]) > charPositionMap.get(secondWord[j])) {

if (charPositionMap.get(firstWord[j]) < charPositionMap.get(secondWord[j])) {</pre>

// The first word comes before the second, so we can break the comparison

// If all characters are equal, but the length of the first word is greater, the list is not sorted

// The first word is greater than the second, so the list is not sorted

if (char1Index < char2Index) return true;</pre>

if (char1Index > char2Index) return false;

function isAlienSorted(words: string[], order: string): boolean {

charPositionMap.set(char, charPositionMap.size);

// Find the minimum length to compare both words

const charPositionMap = new Map<string, number>();

// Iterate through the pairs of adjacent words

for (let i = 1; i < words.length; i++) {</pre>

// Compare characters of both words

for (let j = 0; j < minLength; j++) {</pre>

areWordsEqual = true;

const firstWord = words[i - 1];

const secondWord = words[i];

let areWordsEqual = false;

return false;

break;

Time and Space Complexity

int len1 = word1.size(), len2 = word2.size();

for (int i = 0; i < min(len1, len2); ++i) {</pre>

isDistinct = false;

int[] charIndexMapping = new int[26];

return false;

if (isDistinct) {

break searchLoop;

for (int i = 0; i < order.length(); ++i) {</pre>

// Mapping from character to its position in the alien language

previous_index = current_index

return False

if is_sorted_until_now:

return True

return True

searchLoop:

class Solution {

order_index = {char: index for index, char in enumerate(order)}

for word in words: 12 13 # If the current index is greater than the word length, use a default index of -1. # Else, use the mapped order index of the current character. 14 15 current_index = -1 if i >= len(word) else order_index[word[i]] 16 # If the previous character's index is greater than the current character's index, the list is not sorted.

Update the previous_index to the current character's order index for the next iteration.

If no problems were found in the above loop, the words list is sorted according to the alien language.

is_sorted_until_now = True # Flag to keep track of whether the order is sorted up to the current character index.

If the previous character's index is the same as the current, we can't be sure if it is sorted and need to check th

- 34 Java Solution
- 15 boolean isDistinct = true; // Flag to check if all chars at position i were unique so far 16 // Iterate over words to compare characters at index i 17 for (String word : words) { 18 // If i is greater than word length, currIndex is -1; else, get char's alien order index 19 20 int currentIndex = i >= word.length() ? -1 : charIndexMapping[word.charAt(i) - 'a']; 21

// If previous index is equal to the current index, we can't confirm order; isDistinct is false

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           return true; // Return true if words are sorted as per alien language
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43 }
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C++ Solution
  1 class Solution {
  2 public:
         // Function to check if the words vector is sorted according to the provided alien dictionary order.
         bool isAlienSorted(vector<string>& words, string order) {
             // Create a mapping from character to its index in the alien language
             vector<int> charIndexMap(26);
  6
             for (int i = 0; i < 26; ++i) {
                 // order[i] - 'a' calculates the offset index from 'a' for the character in the alien language
                 charIndexMap[order[i] - 'a'] = i;
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             // Compare each word with the next one to see if they are in the correct order
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             for (auto it = words.begin(); it != prev(words.end()); ++it) {
                 // If the current word is greater than the next word, return false
 14
                 if (!isCorrectOrder(*it, *next(it), charIndexMap)) {
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                     return false;
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 19
```

// Helper function to determine if two words are in the correct order according to the alien dictionary.

// If characters are different and correctly ordered, wordl is less than word2, stop comparing

bool isCorrectOrder(const string &word1, const string &word2, const vector<int>& charIndexMap) {

// Compare character by character according to the alien dictionary order

// If words are equal for min(len1, len2), the shorter word should come first

32 return false; 33 34 35 // If all words are properly sorted or equal, return true 36

return true;

• The outer loop runs a fixed 20 times which is a constant factor and thus does not impact the scalability of the algorithm. The

Space Complexity

Time Complexity

value 20 probably assumes that no word will be longer than 20 characters, which seems to be an arbitrary limit. However, this loop should ideally be in terms of the length of the longest word in words. It would be more accurate to consider it as O(L) where L is the length of the longest word in words.

us a complexity of O(N) for the inner loop. Within the inner loop, the operations are constant time, involving index lookup and comparison.

• Inside the outer loop, there is an inner loop that iterates over each word in the words list. If the number of words is N, this gives

Creating a mapping (dictionary) m from characters to their alien order indices has a complexity of 0(1) because the number of

characters in the order string is constant (assuming a fixed alphabet size, e.g., 26 for the English alphabet).

- Putting it all together, even with the arbitrary limit of 20 iterations, the actual scalable factor is the length of the words and the number of words. So the time complexity is O(L * N).
 - The space complexity for creating the m dictionary is 0(1), again because the alphabet size is constant. The code does not use any additional significant memory that scales with the input size, as all other variables are of fixed size regardless of the input.
- Thus, the overall space complexity of the code is 0(1).