123. Best Time to Buy and Sell Stock III Array **Dynamic Programming** Hard

# We are presented with an array called prices, where each element prices[i] represents the price of a stock on day i. The challenge

**Problem Description** 

then selling one share of the stock. It is important to note that you cannot hold more than one share at a time, which means you must sell the share you hold before buying another one. Therefore, the goal is to strategically choose two periods of time to buy and sell stocks to maximize your profit. Intuition

is to determine the maximum profit that can be made through at most two stock transactions. A transaction consists of buying and

The intuition behind the solution involves dynamic programming to keep track of profits across four different states, which represent the four actions you can take: 1. Buying the first stock (f1): For this action, we want to minimize the cost, so we keep track of the lowest price we've seen so far.

2. Selling the first stock (f2): Here, we calculate the profit from the first transaction. We aim to maximize the profit by selling at the highest price after buying at the lowest price.

3. Buying the second stock (f3): For this, we want the net cost to be minimal, which is purchasing a second stock at the lowest

effective price after accounting for the profit from the first sale. This means we subtract the price of the second stock from the

- profit made from selling the first stock. 4. Selling the second stock (f4): Finally, we want to maximize the total profit, which comes from selling the second stock at the highest possible price.
- By the end of the iteration, f4 will represent the maximum profit achievable with up to two trades.

The algorithm operates by iterating through the price array and updating these four states. Each price offers a potential to update

price (after adjusting for profit from f2) for the second transaction (f3), or a better final sell price for a higher overall profit (f4).

these states—either you get a better buy price (f1), a better sell price resulting in higher first transaction profit (f2), a lower net buy

**Solution Approach** 

The solution leverages dynamic programming to break down the problem into smaller subproblems and builds up the answer from

1. Initializing Variables: We initialize four variables, f1, f2, f3, and f4, which correspond to the states of buying the first stock,

selling the first stock, buying the second stock, and selling the second stock, respectively. Initially, we set f1 and f3 to prices [0] since we "buy" the stock on day zero (the maximum profit after buying is the negative value of the stock's price), and

the lowest effective price by considering the profits from the first sale.

those subproblems. Here's how the implementation carries this out:

f2 and f4 to 0 as we have not made any profit yet.

at the lowest price seen so far.

## Update f1 by taking the maximum of the current f1 and the negative current price, indicating the purchase of the first stock

f1 = max(f1, -price)

2. Iterating through Prices: We iterate through the prices array starting from day 1, updating the four states with each new price.

- Update f2 by taking the maximum of the current f2 and the current price plus f1. This reflects the sale of the first stock and captures the highest profit possible up to this point after the first transaction. f2 = max(f2, f1 + price)
- Update f3 by taking the maximum of the current f3 and f2 minus the current price. This reflects buying the second stock at
- f3 = max(f3, f2 price)• Update f4 by taking the maximum of the current f4 and the current price plus f3. This updates the total profit after the

3. Returning Result: After iterating through all prices, f4 will hold the maximum profit after at most two transactions, which we

Each of the updates done in the loop effectively simulates all possible buy and sell actions that could be taken on that day and

By the end of the loop, we have considered all possible scenarios of one or two transactions through the days, ensuring the

maximum profit is calculated. The key aspect of this solution is it avoids the need for nested loops, which would significantly

carries forward the best decision. The reason we update in the order  $f1 \rightarrow f2 \rightarrow f3 \rightarrow f4$  is due to the dependencies between the

transactions: selling the first stock must consider its purchase and buying the second stock must consider the sale of the first, etc.

return as the final result.

increase the computational complexity.

Example Walkthrough

most two transactions.

Following the solution approach:

1. Initializing Variables:

• f2 = 0 (no sale yet)

Day 1 (price = 4):

f4 = max(f4, f3 + price)

second sale is the highest we can get.

Let's walk through the solution approach using a small example. Consider the following array for prices: 1 prices = [3, 4, 5, 1, 3, 2, 10]

In this array, prices[i] represents the price of the stock on day i. We want to calculate the maximum profit that can be made with at

• f4 = 0 (no second sale yet) 2. Iterating through Prices:

 $\circ$  f1 = max(-3, -5) = -3  $\circ$  f2 = max(1, -3 + 5) = 2

 $\circ$  f3 = max(-3, 2 - 5) = -3

 $\circ$  f4 = max(1, -3 + 5) = 2

 $\circ$  f1 = max(-3, -3) = -3

 $\circ$  f1 = max(-3, -2) = -3

 $\circ$  f3 = max(1, 2 - 2) = 1

 $\circ$  f4 = max(4, 1 + 2) = 4

Day 6 (price = 10):

**Python Solution** 

class Solution:

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C++ Solution

1 #include <vector>

class Solution {

from typing import List

 $\circ$  f2 = max(2, -3 + 2) = 2

 $\circ$  f2 = max(2, -3 + 3) = 2

Day 2 (price = 5):

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Day 3 (price = 1):
  \circ f1 = max(-3, -1) = -3
  \circ f2 = max(2, -3 + 1) = 2
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 $\circ$  f4 = max(0, -3 + 4) = 1 (after possible second sale)

f1 = -3 (since we "buy" the stock on day 0 at price 3)

• f3 = -3 (after "buying" the first stock, we haven't bought the second one yet)

 $\circ$  f1 = max(-3, -4) = -3 (we keep the previous buy since it's cheaper)

 $\circ$  f2 = max(0, -3 + 4) = 1 (we can sell the first stock bought at -3 for 4)

∘ f3 = max(-3, 2 - 1) = 1 (we now buy the second stock since we have a net profit after first sale)  $\circ$  f4 = max(2, 1 + 1) = 2 Day 4 (price = 3):

∘ f3 = max(-3, 1 - 4) = -3 (considering profit from first sale, buying second stock at -3 is still better)

- $\circ$  f3 = max(1, 2 3) = 1 ∘ f4 = max(2, 1 + 3) = 4 (we sold the second stock we bought on day 3 for 3, making a profit) Day 5 (price = 2):
- $\circ$  f2 = max(2, -3 + 10) = 7  $\circ$  f3 = max(1, 7 - 10) = 1  $\circ$  f4 = max(4, 1 + 10) = 11

Thus, we conclude the maximum profit possible is \$11.

def maxProfit(self, prices: List[int]) -> int:

second\_buy, second\_sell = -prices[0], 0

first\_buy = max(first\_buy, -price)

# and that's the result to return

public int maxProfit(int[] prices) {

int secondBuyProfit = -prices[0];

for (int i = 1; i < prices.length; ++i) {</pre>

int secondSellProfit = 0;

return secondSellProfit;

2 #include <algorithm> // for std::max function

return second\_sell

for price in prices[1:]:

# Loop through each price starting from the second price

# Update first\_buy to be either its previous value

# Update first\_sell to be either its previous value

# Update second\_buy to be either its previous value

# Update second\_sell to be either its previous value

# The maximum profit after two transactions is located in second\_sell,

second\_sell = max(second\_sell, second\_buy + price)

// fl represents the maximum profit after the first buy

// f3 represents the maximum profit after the second buy

// f4 represents the maximum profit after the second sell

// Update the maximum profit after the first buy

// Update the maximum profit after the first sell

// Update the maximum profit after the second buy

// Update the maximum profit after the second sell

// Return the maximum profit after the second sell

firstBuyProfit = Math.max(firstBuyProfit, -prices[i]);

// Iterate through the list of prices starting from the second price

// Equivalent to buying at 'firstBuyProfit' and selling at prices[i]

firstSellProfit = Math.max(firstSellProfit, firstBuyProfit + prices[i]);

secondBuyProfit = Math.max(secondBuyProfit, firstSellProfit - prices[i]);

secondSellProfit = Math.max(secondSellProfit, secondBuyProfit + prices[i]);

// Equivalent to buying at 'secondBuyProfit' and selling at prices[i]

// This is the maximum profit we can make with at most two transactions

// We subtract 'prices[i]' because it's the price we are buying at after the first sell

first\_sell = max(first\_sell, first\_buy + price)

second\_buy = max(second\_buy, first\_sell - price)

# or the negative of the current price (representing a buy)

# or the sum of first\_buy and the current price (representing a sell)

# or the difference of first\_sell and the current price (representing a second buy)

# or the sum of second\_buy and the current price (representing a second sell)

 $\circ$  f1 = max(-3, -10) = -3

# Initialize the four states of profits # first\_buy represents the profit after the first buy # first\_sell represents the profit after the first sell # second\_buy represents the profit after the second buy # second\_sell represents the profit after the second sell first\_buy, first\_sell = -prices[0], 0 10

3. Returning Result: Lastly, after iterating through all prices, f4 holds the maximum profit of 11 after at most two transactions.

By following these steps using dynamic programming, we have efficiently computed the maximum profit possible with up to two

transactions without needing to iterate through every possible pair of buy and sell days, which would be far less efficient.

35 # Example usage: **36** # sol = Solution() 37 # max\_profit = sol.maxProfit([3,3,5,0,0,3,1,4]) 38 # print(max\_profit) # Output would be 6, representing the maximum profit possible with two transactions 39

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int firstBuyProfit = -prices[0];
          // f2 represents the maximum profit after the first sell
          int firstSellProfit = 0;
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Java Solution

1 public class Solution {

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public:
       int maxProfit(vector<int>& prices) {
           // Initialization of buyl and sell1 corresponds to the first transaction.
           // buyl is the maximum profit we can achieve by buying a stock on day 0,
           // which is simply the negative value of the stock price on that day.
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           int buy1 = -prices[0];
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           // sell1 is the maximum profit we can achieve by selling a stock after buying on day 0,
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           // initially this would be zero because we have just bought a stock and yet to sell it.
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           int sell1 = 0;
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           // Initialization of buy2 and sell2 corresponds to the second transaction.
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           // buy2 is the maximum profit we can potentially hold for a second buy,
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           // initially it's the same as buyl because we have not executed the first sell yet.
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           int buy2 = -prices[0];
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           // sell2 is the maximum profit we can get after completing the second sell,
           // initially, this would also be zero because we haven't sold any stock twice.
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           int sell2 = 0;
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            for (int i = 1; i < prices.size(); ++i) {</pre>
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               // For each day, calculate the maximum profit if we were to buy the stock
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               // for the first transaction, by comparing the previous buy profit and
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               // the negative of today's price (as if we bought the stock today)
                buy1 = std::max(buy1, -prices[i]);
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               // Calculate the maximum profit if we were to sell the stock
               // for the first transaction. Compare the previous sell profit and
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               // the profit from selling today (today's price plus the current buy profit).
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               sell1 = std::max(sell1, buy1 + prices[i]);
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               // For the second buy, calculate the maximum profit by comparing the previous second buy profit
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               // and the current overall profit minus today's price (buying the second stock today).
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               buy2 = std::max(buy2, sell1 - prices[i]);
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               // Finally, calculate the profit if we were to make the second sell today.
               // It compares the current second sell profit and the profit from selling today,
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               // which includes the potential second buy profit and today's price.
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               sell2 = std::max(sell2, buy2 + prices[i]);
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           // The result should be the maximum profit after two sales.
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           return sell2;
47 };
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// firstBuyProfit assumes the first transaction hasn't happened yet, hence negative value

// firstSellProfit assumes no transaction has been made, hence 0 profit

// secondBuyProfit takes into account profit from the first transaction

// secondSellProfit represents the cumulative profit from both transactions

// The maximum profit of the first buy is either the previous value or

// the current profit from the first sell minus the current price

// Return the cumulative maximum profit from both allowed transactions

secondBuyProfit = Math.max(secondBuyProfit, firstSellProfit - prices[i]);

secondSellProfit = Math.max(secondSellProfit, secondBuyProfit + prices[i]);

• The number of iterations in the loop, which is n-1 where n is the length of the prices list.

The constant time operations within each iteration, which don't depend on the size of the input.

// The maximum profit of the second sell is the greater of the previous value or

// the profit after selling at the current price (including the profit from the first sell)

// the negative of the current price (which means buying at the current price)

## // The maximum profit of the first sell is either the previous value or 23 // the profit after selling at the current price firstSellProfit = Math.max(firstSellProfit, firstBuyProfit + prices[i]); 24 25 26 // The maximum profit of the second buy is the greater of the previous value or

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Typescript Solution

function maxProfit(prices: number[]): number {

let firstBuyProfit = -prices[0];

let secondBuyProfit = -prices[0];

let firstSellProfit = 0;

let secondSellProfit = 0;

return secondSellProfit;

Time and Space Complexity

The time complexity of the code is determined by:

// Initialize profit states for the two transactions

// At start, no transaction has happened, hence 0 profit

firstBuyProfit = Math.max(firstBuyProfit, -prices[i]);

// Loop to calculate profits at each price point

for (let i = 1; i < prices.length; ++i) {</pre>

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to make at most two transactions to maximize profit. The variables f1, f2, f3, and f4 represent the four states of profits after each
transaction or waiting period.
Time Complexity:
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The given Python code represents a dynamic programming approach to solve a stock trading problem, where an individual is allowed

# There are four assignments in each iteration that happen in 0(1) time each.

Thus, the total time complexity is 0(n-1) \* 0(1), which simplifies to 0(n). **Space Complexity:** 

The space complexity of the code is determined by:

- The space taken by the four variables f1, f2, f3, f4, which is constant and does not scale with input size.
- No additional data structures are used that grow with input size. Hence, the space complexity of the code is 0(1), representing the constant space used by the variables.