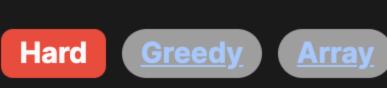
## 1121. Divide Array Into Increasing Sequences



#### **Problem Description**

The LeetCode problem in question requires us to determine if a sorted integer array nums can be divided into one or more disjoint increasing subsequences where each subsequence is of at least length k. A subsequence is a sequence that can be derived from another sequence by deleting some or no elements without changing the order of the remaining elements. Disjoint subsequences are such that they do not share any common elements. The nums array is given to be sorted in non-decreasing order (i.e., it can have duplicates, but the sequence is not decreasing). The function should return true if such a division is possible or false otherwise.

#### Intuition

most frequent number in nums limits the number of these subsequences we can create. This is because each subsequence can have at most one occurrence of any number, and hence the number of subsequences

To arrive at the solution, we must understand that if we can form disjoint increasing subsequences of at least length k, then the

cannot exceed the number of occurrences of the most frequent number. To find the solution, we follow these steps:

We calculate the frequency of the most frequent number, mx, by using the groupby function from the itertools module in

Python, which groups consecutive identical elements. We convert each group to a list and get its length to find out how many times that particular number appears in the array. We then multiply this frequency mx by k to find the minimum array length needed to have mx disjoint subsequences of length

k. This is because each subsequence needs at least one occurrence of the most frequent number and each must be of length

- k or more. We compare this minimum required length mx \* k to the actual length of nums. If mx \* k is less than or equal to the length of nums, then we can divide nums into the required disjoint increasing subsequences and we return true. Otherwise, we return
- false because there would not be enough numbers to form subsequences of length k with the most frequent number appearing in all of them. By following this approach, the solution effectively ensures whether there are enough elements in the array to distribute among subsequences of the required length without any overlap.

**Solution Approach** 

#### The groupby method: This function is used for grouping elements in an iterable. If you pass a sorted iterable to it, it groups all

consecutive duplicate elements together. In Python, it is available in the itertools module.

The implementation of the solution uses:

- The max function: After grouping the elements with groupby, the code calculates the maximum group size, which represents the frequency of the most common element in the sorted array nums. This is done by mapping each grouped sequence to its
- length and taking the maximum of these lengths. List comprehension: This is used for compactly applying operations to sequences. In the reference solution, list comprehension is utilized within the max function to create a list of lengths of the groups returned by groupby.

Comparison: Finally, the decision to return true or false is decided by comparing the product of the maximum group size

- and k with the length of the nums array. If this product is less than or equal to the length of nums, then true is returned. Otherwise, false is returned.
- The code snippet works as follows: mx = max(len(list(x))) for \_, x in groupby(nums): This line creates grouped sequences of identical consecutive elements

these lengths, the maximum length is taken, which is stored in mx.

return mx \* k <= len(nums): This checks whether the array has enough elements to form the required increasing subsequences. If the most frequent element's count (mx) times the minimum subsequence length (k) is less than or equal to the total number of elements in nums, then the condition is satisfied and true is returned. If not, false is returned.

In this implementation, the use of groupby is key because it allows us to easily find the frequency of the most common element in

in the sorted array nums using groupby. For each group, it takes the length by converting the group iterator to a list. Among

the sorted array, which is crucial for determining whether the subsequences can be formed. **Example Walkthrough** 

Let's illustrate the solution approach using a small example. Consider the sorted integer array nums = [1, 2, 3, 3, 4, 4, 5, 5, 5] and let k = 3. We want to determine if this array can be

### Following the described approach:

We use the groupby method to group identical consecutive elements. Thus, we get the groups [(1), (2), (3, 3), (4, 4),

divided into disjoint increasing subsequences each of at least length 3.

We then calculate mx \* k. Here, mx is 3 and k is 3, so mx \* k is 9.

(5, 5, 5)]. We calculate the maximum frequency mx. The groups' lengths are [1, 1, 2, 2, 3], so mx is 3 because 5 occurs three times.

Since  $mx * k \le the length of nums, the function would return true. This means we can divide nums into one or more disjoint$ increasing subsequences of length at least 3. One such division could be [[1, 3, 5], [2, 4, 5], [3, 4, 5]].

We compare mx \* k with the length of nums. The length of nums is 9, which is equal to mx \* k (both are 9).

# The number of subsequences of size k we can create is equal to the maximum frequency.

# If the total length of nums is at least as large as this number, then we can

// Method to check if an array can be divided into subsequences each of length k

int lastValue = 0; // to store the value of the last element processed

public boolean canDivideIntoSubsequences(int[] nums, int k) {

// Iterate over each element in the array

for (int currentValue : nums) {

for (int currentValue : nums) {

return false;

previousValue = currentValue;

int currentCount = 0; // to hold count of current element

// Iterate through each number in the given 'nums' vector

if (currentStreak \* k > nums.size()) {

// Otherwise, reset the streak count to 1 for the new number

// If the current value is the same as the previous value, increment the streak

// If the number of times a particular element needs to be repeated ('currentStreak' times 'k')

// exceeds the total length of the 'nums' vector, it is not possible to divide into subsequences.

currentStreak = (previousValue == currentValue) ? currentStreak + 1 : 1;

// Update the previous value to the current value for the next iteration

subsequence length k, we can determine if a sorted array can be divided into the required subsequences.

This example illustrates how by using the groupby function and checking the maximum frequency of a number times the minimum

**Python** 

#### def canDivideIntoSubsequences(self, nums: List[int], k: int) -> bool: # Calculate the maximum frequency of any number in the list max\_frequency = max(len(list(group)) for \_, group in groupby(nums))

Solution Implementation

from itertools import groupby

from typing import List

class Solution:

class Solution {

```
# divide nums into subsequences of size k, where each subsequence is strictly increasing.
        return max frequency * k <= len(nums)</pre>
Java
```

```
// If the current element is the same as the last, increment the count
           // Otherwise, reset the count for a new value
            currentCount = (lastValue == currentValue) ? currentCount + 1 : 1;
           // If the number of times an element appears multiplied by k exceeds the array length,
           // it's not possible to divide the array into subsequences of length k
            if (currentCount * k > nums.length) {
                return false;
            // Update the last processed value to the current value
            lastValue = currentValue;
       // If we didn't return false during the loop, it's possible to divide the array
       return true;
C++
class Solution {
public:
   bool canDivideIntoSubsequences(vector<int>& nums, int k) {
        int currentStreak = 0; // Count of how many times the current number has appeared consecutively
        int previousValue = 0; // The value of the previous element in the array
```

```
// If the loop completes without returning false, it means the 'nums' vector can be
       // divided into subsequences of length 'k' without violating the rules
       return true;
};
TypeScript
function canDivideIntoSubsequences(nums: number[], k: number): boolean {
    let currentStreak = 0; // Count of consecutive appearances of the current number
    let previousValue = 0; // The value of the previous element in the array
   // Iterate through each number in the 'nums' array
   for (let currentValue of nums) {
       // If the current value is the same as the previous value,
       // increment the streak, otherwise reset streak to 1
        if (previousValue === currentValue) {
           currentStreak++;
       } else {
            currentStreak = 1;
       // If the current streak times k exceeds the array size,
       // it's not possible to divide into subsequences
       if (currentStreak * k > nums.length) {
            return false; // Return early as the condition is violated
       // Update the previousValue to the currentValue for the next iteration
       previousValue = currentValue;
```

```
# If the total length of nums is at least as large as this number, then we can
# divide nums into subsequences of size k, where each subsequence is strictly increasing.
return max_frequency * k <= len(nums)</pre>
```

**Time and Space Complexity** 

to group adjacent elements together.

return true;

from typing import List

class Solution:

from itertools import groupby

# **Time Complexity**

// If we can iterate through the array without returning false,

def canDivideIntoSubsequences(self, nums: List[int], k: int) -> bool:

max\_frequency = max(len(list(group)) for \_, group in groupby(nums))

# Calculate the maximum frequency of any number in the list

// it means the array can be divided into subsequences of length 'k'

The time complexity of the code is primarily determined by the groupby operation from the itertools library and the calculation of the maximum group size (mx). The groupby operation is O(n), where n is the length of the list nums. This is because groupby iterates through the list once in order

# The number of subsequences of size k we can create is equal to the maximum frequency.

elements are the same and thus belong to a single group. Thus, the time complexity of the function is O(n) overall, as both the groupby operation and the calculation of the maximum group

The calculation of the maximum group size involves iterating over each group generated by groupby and finding the length of the

longest group created. This step is also 0(n) because in the worst case, it will iterate through all elements of nums once, if all

size are linear in terms of the input size. **Space Complexity** 

groups are not stored in memory all at once but instead one at a time as they are iterated over, the space complexity does not

The space complexity of this code involves the extra space required to store the groups formed by the groupby function. For each unique element in nums, a new group is formed, and a generator object is created for each group. However, since these

grow with the number of unique elements. The variable mx only requires constant space, and since groupby does not create a list of all groups but just an iterable, the space

complexity for the storage of groups is also constant.

Thus, the space complexity is 0(1), because the space required does not grow with the size of the input list nums.