1560. Most Visited Sector in a Circular Track

Easy Array Simulation

Problem Description

In this problem, we have a circular track with n sectors, each labeled from 1 to n. A marathon is taking place on this track, consisting of m rounds. Each round starts and ends at specified sectors, with the starting sector for each round given in an array `rounds`. Specifically, the ith round starts at sector `rounds[i - 1]` and ends at sector `rounds[i]`. The goal is to find out which sectors are the most visited throughout the marathon and return these sectors sorted in ascending order. It's important to notice that the sectors must be visited in the counter-clockwise direction and we wrap around the track each time we reach sector n.

Intuition

The solution to this problem relies on the observation that if the start sector of the marathon (`rounds[0]`) is less than or equal to the end sector (`rounds[-1]`), then all sectors in between are visited more frequently than any other sector. This is because each round always moves from a lower to a higher sector number, covering all sectors in between.

However, if the start sector is greater than the end sector, this indicates that the runners have passed the sector marked n and started a new lap, visiting the sectors from 1 to `rounds[-1]` again. Therefore, in this case, the most visited sectors will be from sector 1 to `rounds[-1]` and from `rounds[0]` to sector n.

The solution approach, therefore, is to check the relative positions of the start and end sectors and return the appropriate range of sectors:

1. If `rounds[0]` is less than or equal to `rounds[-1]`, return a list of sectors ranging from `rounds[0]` to `rounds[-1]`.
2. If `rounds[0]` is greater than `rounds[-1]`, there are two ranges of sectors to consider. Combine the range from 1 to `rounds[-1]` with the range from `rounds[0]` to n into a single list and return it.

Solution Approach

The solution for finding the most visited sectors on the circular track follows a straightforward approach based on conditional logic rather than elaborate algorithms or data structures.

Here's a step-by-step explanation of the implementation:

1. **Condition Check**: First, we check if `rounds[0]`, the starting sector of the first round, is less than or equal to `rounds[-1]`, the ending sector of the last round.
   - If this condition is `True`, it means that the runners have not crossed the sector marked n from the last round to the first round, so we only need to consider the sectors that lie between `rounds[0]` and `rounds[-1]`.
2. **Range Construction (No Wrap-Around)**: When `rounds[0] <= rounds[-1]`:
   - We use the `range` function to generate a list starting from `rounds[0]` to `rounds[-1]`, inclusive. This list represents the sectors that are covered in a direct path from the starting sector to the ending sector without any wrap-around.
3. **Range Construction (With Wrap-Around)**: When `rounds[0] > rounds[-1]`:
   - This indicates that during the marathon, runners have wrapped around the track, passing by sector n and back to sector 1. We need to consider two ranges of sectors.
   - The first range is from sector 1 to `rounds[-1]`, inclusive, which represents the sectors visited after the last wrap-around.
   - The *second* range is from `rounds[0]` to sector n, inclusive, which represents the sectors visited before the wrap-around occurs.
   - These two ranges are combined to form a complete list of the most visited sectors.
4. **List Concatenation**: To combine the two ranges when there is a wrap-around:
   - We utilize list concatenation to join the lists resulting from the two `range` calls into a single list, which is returned as the final result.

The code implementation uses only basic Python constructs such as conditional statements and list manipulations, thereby avoiding the need for complex algorithms or data structures. It leverages Python's `range` function to create lists directly corresponding to the ranges of sectors that are most visited, as deduced from the initial condition checks.

```
class Solution:
    def mostVisited(self, n: int, rounds: List[int]) -> List[int]:
        if rounds[0] <= rounds[-1]:
            return list(range(rounds[0], rounds[-1] + 1))
        else:
            return list(range(1, rounds[-1] + 1)) + list(range(rounds[0], n + 1))
```

The function `mostVisited` takes two parameters, n and `rounds`, which represent the number of sectors and the array of rounds, respectively. Within the function, we apply the logic outlined above to determine the most visited sectors and return them as a list.

The beauty of this solution lies in its simplicity and efficiency. Since it avoids additional computation by directly identifying the range of most visited sectors, it executes in constant time, making it suitable for even large values of n and numbers of rounds.

Example Walkthrough

Let's walk through a small example to illustrate the solution approach using a circular track with 5 sectors ($n = 5$) and an array of rounds `[1, 3, 5, 2]`. We apply the steps from the solution approach to determine the most visited sectors.

Given:

- $n = 5$ (sectors are 1, 2, 3, 4, 5)
- `rounds = [1, 3, 5, 2]`

We start by applying the first step in our solution approach.

1. **Condition Check**: We compare `rounds[0]` to `rounds[-1]` (we check whether $1 <= 2$).
   - In this case, the condition is `True` because 1 is less than 2.
   - Normally, this would imply that we only need to consider the sectors that lie between 1 and 2. However, since there is a wrap-around implied by the progression of sectors in the rounds from 5 back to 2, we need to consider the full marathon sequence.
2. **Range Construction (With Wrap-Around)**: Since `rounds[0]` is not greater than `rounds[-1]` but there is a wrap-around from the third to the fourth round (from 5 back to 2), we consider the entire sequence of sectors visited.
   - The first range is from sector 1 to sector 5, inclusive.
   - The second range starts again at 1 after wrapping around and ends at 2.
3. **List Concatenation**: We combine the ranges from the first round and after the wrap-around:
   - The sectors visited are initially `[1, 2, 3, 4, 5]`, and following the wrap-around, the sectors `[1, 2]` are visited again, leading to the sequence `[1, 2, 3, 4, 5, 1, 2]`.
   - Since sectors 1 and 2 were visited in the last part of the marathon, they would be added to the most visited sectors list.

Hence, after evaluating the rounds in order, we determine that the most visited sectors would be `[1, 2, 3, 4, 5]`, but `[1, 2]` have the highest visit count due to the wrap-around. We need to sort them to return the result as `[1, 2]`.

Using the given solution code, the `mostVisited` function would thus return `[1, 2, 3, 4, 5] + [1, 2]`, which simplifies to `[1, 2, 3, 4, 5]` since we are asked to return unique sectors visited. After sorting, the final answer remains `[1, 2, 3, 4, 5]`, since all sectors are most visited due to the marathon sequence that was provided in the `rounds` array.

Solution Implementation

Python

```
class Solution:
    def mostVisited(self, sector_count: int, rounds: List[int]) -> List[int]:
        # The first sector visited in the first round
        start_sector = rounds[0]
        # The last sector visited in the last round
        end_sector = rounds[-1]

        # If the race started (start_sector) and ended (end_sector) on the same or a directly subsequent sector,
        # or if the race direction was such that there was no wraparound
        if start_sector <= end_sector:
            # The most visited sectors are all sectors from start_sector to end_sector, inclusive
            return list(range(start_sector, end_sector + 1))
        else:
            # If there was a wraparound, the most visited sectors are from the beginning (sector 1)
            # to the end_sector, and from the start_sector to the end of the track (sector_count)
            return list(range(1, end_sector + 1)) + list(range(start_sector, sector_count + 1))

# Note: The List type needs to be imported from the typing module, so we should add the following line
# at the beginning of the file:
# from typing import List
```

```
Make sure to import the `List` type from the `typing` module for type annotations if it is not already imported:

```python
from typing import List
```
```

Java

```
import java.util.ArrayList;
import java.util.List;

class Solution {
    public List<Integer> mostVisited(int sectors, int[] rounds) {
        // The final round (m is now finalRoundIndex for clarity)
        int finalRoundIndex = rounds.length - 1;
        // List to store the most visited sectors
        List<Integer> mostVisitedSectors = new ArrayList<>();

        // If the first round is less than or equals to the last round, it means the path doesn't cross sector 1
        if (rounds[0] <= rounds[finalRoundIndex]) {
            // Iterate from start sector to end sector and add to the list
            for (int i = rounds[0]; i <= rounds[finalRoundIndex]; ++i) {
                mostVisitedSectors.add(i);
            }
        } else { // The path crosses sector 1
            // Add sectors from 1 to the final round's sector
            for (int i = 1; i <= rounds[finalRoundIndex]; ++i) {
                mostVisitedSectors.add(i);
            }
            // Add sectors from the starting round's sector to the last sector (wrapping around)
            for (int i = rounds[0]; i <= sectors; ++i) {
                mostVisitedSectors.add(i);
            }
        }

        // Return the list of most visited sectors
        return mostVisitedSectors;
    }
}
```

C++

```
#include <vector>

class Solution {
public:
    // This function returns the most visited sectors in a circular track after all rounds are completed.
    std::vector<int> mostVisited(int totalSectors, std::vector<int>& rounds) {
        int lastRoundIndex = rounds.size() - 1; // Get the index of the last round
        std::vector<int> mostVisitedSectors; // Vector to store the most visited sectors

        // Check if the starting sector number is less than or equal to the ending sector number
        if (rounds[0] <= rounds[lastRoundIndex]) {
            // If it is, add all sectors from the start sector to the end sector to the answer
            for (int i = rounds[0]; i <= rounds[lastRoundIndex]; ++i) {
                mostVisitedSectors.push_back(i);
            }
        } else {
            // If it's not, the path has lapped around the track
            // Add all sectors from 1 to the end sector
            for (int i = 1; i <= rounds[lastRoundIndex]; ++i) {
                mostVisitedSectors.push_back(i);
            }
            // Also add all sectors from the start sector to the total number of sectors
            for (int i = rounds[0]; i <= totalSectors; ++i) {
                mostVisitedSectors.push_back(i);
            }
        }
        // Return the list of the most visited sectors after all the rounds
        return mostVisitedSectors;
    }
};
```

TypeScript

```
// Define the function to calculate the most visited sectors.
// Take the total number of sectors and an array of rounds as inputs.
function mostVisited(totalSectors: number, rounds: number[]): number[] {
    // Get the index of the last round
    const lastRoundIndex = rounds.length - 1;
    // Initialize an array to store the most visited sectors
    let mostVisitedSectors: number[] = [];

    // Check if the starting sector number is less than or equal to the ending sector number
    if (rounds[0] <= rounds[lastRoundIndex]) {
        // Add all sectors from the start sector to the end sector to the most visited sectors
        for (let i = rounds[0]; i <= rounds[lastRoundIndex]; i++) {
            mostVisitedSectors.push(i);
        }
    } else {
        // The path has lapped around the track, so cover both segments of the lap

        // Add all sectors from 1 to the end sector number
        for (let i = 1; i <= rounds[lastRoundIndex]; i++) {
            mostVisitedSectors.push(i);
        }

        // Add all sectors from the start sector to the total number of sectors
        for (let i = rounds[0]; i <= totalSectors; i++) {
            mostVisitedSectors.push(i);
        }
    }

    // Return the array of the most visited sectors
    return mostVisitedSectors;
}


class Solution:
    def mostVisited(self, sector_count: int, rounds: List[int]) -> List[int]:
        # The first sector visited in the first round
        start_sector = rounds[0]
        # The last sector visited in the last round
        end_sector = rounds[-1]

        # If the race started (start_sector) and ended (end_sector) on the same or a directly subsequent sector,
        # or if the race direction was such that there was no wraparound
        if start_sector <= end_sector:
            # The most visited sectors are all sectors from start_sector to end_sector, inclusive
            return list(range(start_sector, end_sector + 1))
        else:
            # If there was a wraparound, the most visited sectors are from the beginning (sector 1)
            # to the end_sector, and from the start_sector to the end of the track (sector_count)
            return list(range(1, end_sector + 1)) + list(range(start_sector, sector_count + 1))

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Time and Space Complexity

The time complexity of the code is primarily determined by the creation of the list that gets returned. This involves generating a range of integers, which can be done in constant time for each integer in the range, and then converting that range into a list.

- In the best-case scenario, where `rounds[0] <= rounds[-1]`, we create a single range from `rounds[0]` to `rounds[-1]`, which includes at most n elements. Thus, the time complexity for this case is $O(n)$.
- In the worst-case scenario, where `rounds[0] > rounds[-1]`, we create two ranges. The first is from 1 to `rounds[-1]` and the second is from `rounds[0]` to n. In the worst case, these two ranges can also include up to n elements combined, so the time complexity remains $O(n)$.

The space complexity of the code is determined by the space needed to store the output list.

- In both cases mentioned above, the space complexity depends on the number of elements in the output list, which can be at most n. Thus, the space complexity for the algorithm is $O(n)$.

Therefore, the final time complexity is $O(n)$ and the space complexity is $O(n)$.