



**Problem Description** 

Math

String

In this problem, we are given two strings \$1 and \$2. Both strings have the same length and contain only the characters 'x' and 'y'. Our objective is to make the two strings identical by performing swaps between the characters of the two strings. A swap involves taking one character from s1 and one character from s2 and exchanging them. The goal is to determine the minimum number of swaps required to make the two strings identical. If it's not possible to make the strings equal, we must return -1.

### Intuition

The solution is based on counting how many characters are out of place when comparing both strings. We look for characters that are 'x' in s1 and 'y' in s2, and vice versa, since these are the only ones that need to be swapped to make the strings equal. There are two types of mismatches:

A 'yx' mismatch: where s1 has a 'y' and s2 has an 'x' at the same position.

An 'xy' mismatch: where s1 has an 'x' and s2 has a 'y' at the same position.

- 1. If we have an even number of 'xy' and 'yx' mismatches, we can always solve the problem. 2. We can swap odd mismatches within themselves, for instance, an odd count of 'xy' mismatches can be made even by
- performing one swap inside the 'xy' set (changing two 'xy' to two 'yx'). This adds one more swap to our answer for each odd count. 3. Each pair of mismatches ('xy' with 'yx') can be solved with one swap, so the pair count contributes directly to the total number of
- swaps. The algorithm checks if there is an even total number of mismatches. If it's odd, we return -1 because we would always end up with

one character that cannot be matched. If it's even, we sum the half of each count (because a swap fixes two mismatches) and case for odd counts, which contributes one additional swap for each. **Solution Approach** 

# The implementation of the solution uses a simple counter-based approach without any sophisticated data structures or patterns.

The core of the solution revolves around counting the discrepancies between the two strings and categorizing them as xy or yx, followed by calculating the number of swaps needed based on these counts. Let's walk through the steps present in the Python code provided:

1. Initialize two counters xy and yx to 0. These will hold the counts for mismatches where s1[i] is 'x' and s2[i] is 'y', and where

counter if a > b, which indicates a 'y' in s1 and an 'x' in s2.

- s1[i] is 'y' and s2[i] is 'x', respectively. 2. Iterate over both strings in parallel by using the zip function, which allows us to obtain pairs of characters (a, b) from \$1 and \$2.
- 3. For each pair (a, b), increment the xy counter if a < b, which means we have an 'x' in s1 and a 'y' in s2. Similarly, increment the yx
- 4. Once we've counted all mismatches, we check whether the sum of xy and yx is even. If it's not ((xy + yx) % 2 is truthy), we return -1 because this means it's impossible to make the strings equal.
- swapped directly, which also applies to yx // 2 pairs of 'yx' mismatches. Additionally, for each type of mismatch, if there's an odd one out (checked by xy % 2 and yx % 2), it requires an extra swap. 6. The minimum number of swaps required is the sum of these values, which is xy // 2 + yx // 2 + xy % 2 + yx % 2, meaning

5. If the sum is even, we calculate the number of swaps. We calculate xy // 2 to see how many pairs of 'xy' mismatches can be

the number of direct swaps for pairs plus the extra swaps for any remaining mismatch. The logic applied here uses basic arithmetic and logic to deduce the number of required actions to align the two strings. It's a

pragmatic approach that is both efficient and clear in its purpose. Example Walkthrough

## Let's illustrate the solution approach with a small example. Consider two strings s1 = "xyyx" and s2 = "yxyx". Following the steps of the solution:

1. Initialize counters xy and yx to 0.

2. Iterate over both strings s1 and s2 and compare the corresponding characters using zip:

- For the first pair (s1[0] = 'x', s2[0] = 'y'), since s1[0] < s2[0], increment xy to 1.</li>
- $\circ$  For the second pair (s1[1] = 'y', s2[1] = 'x'), s1[1] > s2[1], so increment yx to 1.
  - The third pair (s1[2] = 'y', s2[2] = 'y') matches, so no counter is incremented.  $\circ$  For the fourth pair (s1[3] = 'x', s2[3] = 'x'), again they match, so no counter is increased.
  - After this step, xy equals 1, and yx equals 1.
- 3. Check if the sum of xy and yx is even: 1 + 1 = 2 which is even, so continue. 4. Calculate the minimum number of swaps:
- Pairs of 'xy' mismatches that can be swapped directly: xy // 2 = 1 // 2 = 0. Pairs of 'yx' mismatches that can be swapped directly: yx // 2 = 1 // 2 = 0.

An extra swap is needed for the one 'yx' mismatch: yx % 2 = 1 % 2 = 1.

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character of s2. After these swaps, both s1 and s2 will be "xyxy".

# Counters for "x-y" and "y-x" pairs

if char1 == 'x' and char2 == 'y':

char charS1 = s1.charAt(i), charS2 = s2.charAt(i);

// If s1 has 'x' and s2 has 'y', increment countXY

// If s1 has 'y' and s2 has 'x', increment countYX

if (charS1 == 'x' && charS2 == 'y') {

if (charS1 == 'y' && charS2 == 'x') {

countXY++;

countYX++;

if ((countXY + countYX) % 2 == 1) {

if ((countXY + countYX) % 2 != 0) {

// The number of swaps is calculated by adding:

// - Half of countYX for the same reason as above

return -1; // Return -1 to indicate the impossibility

count\_xy = count\_yx = 0

5. The total number of swaps needed is: xy // 2 + yx // 2 + xy % 2 + yx % 2 = 0 + 0 + 1 + 1 = 2. Therefore, for s1 = "xyyx" and s2 = "yxyx", it takes a minimum of 2 swaps to make the strings identical. The swaps that could be performed would be between the first character of s1 with the second character of s2, and the second character of s1 with the first

Python Solution class Solution: def minimumSwap(self, s1: str, s2: str) -> int:

#### # Iterate over characters of both strings simultaneously for char1, char2 in zip(s1, s2): # If we find a "x-y" pair, increment count\_xy

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count_xy += 1
               # If we find a "y-x" pair, increment count_yx
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               elif char1 == 'y' and char2 == 'x':
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                   count_yx += 1
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           # If the sum of both counts is odd, we can't make them equal, thus return -1
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           if (count_xy + count_yx) % 2 != 0:
               return -1
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           # For pairs "xy" or "yx" that match directly, each pair takes one swap
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           swaps = count_xy // 2 + count_yx // 2
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           # For the remaining unpaired "xy" or "yx", they take two swaps to balance
           # As there would be one "xy" and one "yx" remaining for an even total of mismatches
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           swaps += 2 * (count_xy % 2)
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           return swaps
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Java Solution
   class Solution {
       public int minimumSwap(String s1, String s2) {
           // Counters for 'x' in s1 and 'y' in s2, and 'y' in s1 and 'x' in s2
           int countXY = 0, countYX = 0;
           // Loop through the strings to count 'xy' and 'yx' pairs
           for (int i = 0; i < s1.length(); i++) {
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               return -1;
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           // The minimum swaps is the sum of:
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           // - Half the pairs of 'xy' and 'yx', as two swaps can solve two pairs,
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           // - and one swap for each of the remaining 'xy' or 'yx' if there is an odd count.
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           return countXY / 2 + countYX / 2 + countXY % 2 + countYX % 2;
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29 }
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C++ Solution
 1 class Solution {
 2 public:
       // Function to calculate the minimum number of swaps to make two strings equal
       int minimumSwap(string s1, string s2) {
           int countXY = 0; // Number of occurrences where 'x' in s1 is aligned with 'y' in s2
           int countYX = 0; // Number of occurrences where 'y' in s1 is aligned with 'x' in s2
           // Loop through both strings character by character
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           for (int i = 0; i < s1.size(); ++i) {
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               char charFromS1 = s1[i], charFromS2 = s2[i];
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               if (charFromS1 == 'x' && charFromS2 == 'y') {
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                   // Increment count for 'x' in s1 and 'y' in s2
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                   countXY++;
               } else if (charFromS1 == 'y' && charFromS2 == 'x') {
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                   // Increment count for 'y' in s1 and 'x' in s2
                   countYX++;
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// If the sum of countXY and countYX is odd, return -1, as it's impossible to swap to equality

// If the sum of countXY and countYX is odd, it's not possible to make the strings equal

// - Half of countXY because two 'xy' pairs can be corrected by a single swap

#### 28 // - The remainder of countXY divided by 2, because one 'xy' cannot be solved without an extra 'yx' pair 29 30

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Similarly, one cannot solve an extra 'yx' without an extra 'xy' so countYX % 2 is also needed
           return countXY / 2 + countYX / 2 + countXY % 2 + countYX % 2;
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32 };
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Typescript Solution
   function minimumSwap(s1: string, s2: string): number {
        let countXY = 0, // count of 'x' in s1 and 'y' in s2 at the same position
           countYX = 0; // count of 'y' in s1 and 'x' in s2 at the same position
       // Iterate through the strings to count 'xy' and 'yx' pairs
       for (let i = 0; i < s1.length; ++i) {</pre>
           const charS1 = s1[i],
                 charS2 = s2[i];
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           if (charS1 === 'x' && charS2 === 'y') {
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               ++countXY; // Increase count for 'xy' pair
12
           if (charS1 === 'y' && charS2 === 'x') {
               ++countYX; // Increase count for 'yx' pair
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       // Check if the sum of 'xy' and 'yx' pairs is odd, return -1 since it's not possible to swap to equality
       if ((countXY + countYX) % 2 === 1) {
           return -1;
       // Calculate minimum swaps:
       // Each pair of 'xy' or 'yx' requires one swap
```

# Time and Space Complexity Time Complexity

18 19 21 22 23 24 // If there's an odd number of 'xy' or 'yx', it takes two additional swaps to fix both return Math.floor(countXY / 2) + Math.floor(countYX / 2) + (countXY % 2) + (countYX % 2); 26 27 } 28

## The time complexity of the given function is O(n), where n is the length of the strings s1 and s2. This is because the function consists of a single loop that iterates through all characters of s1 and s2 in a pairwise manner using the zip function. Within the loop, only

constant time operations are performed (comparison and addition).

remains constant regardless of the input size.

Space Complexity The space complexity of the given function is 0(1). There are only a fixed number of integer variables (xy and yx) initialized and used for counting the occurrences, which do not depend on the size of the input strings. Thus, the amount of additional memory required