Dynamic Programming

String

Array

#### **Problem Description**

Medium Greedy

colors[i] is the color of the ith balloon. Alice desires the string of balloons to be "colorful," which means no two adjacent balloons should share the same color. To help achieve this, Bob can remove balloons, but it takes time. The time it takes to remove each balloon is specified in an integer array neededTime, where neededTime[i] represents the seconds required to remove the ith balloon.

The goal is to find the minimum amount of time Bob needs to spend to turn the string of balloons into one where no two consecutive

Alice has n balloons tied in a row and each balloon is colored some color. This sequence of colors is given as a string colors where

balloons are of the same color.

### Given that we want to make the rope of balloons colorful with the least amount of time needed, we can intuit that when we encounter a sequence of balloons of the same color, we should remove all but one of them. Instead of randomly choosing which

Intuition

balloons to remove, we should aim to remove the balloons with the least time cost. The optimal balloon to keep is the one with the highest neededTime value because removing it would contribute the most to the total time.

Throughout the string of balloons, we iteratively look for sequences of consecutive balloons with the same color. For each of these sequences, we calculate the total time needed to remove all balloons (s) and the maximum time needed to remove a single balloon within that sequence (mx). To make the sequence colorful, we will keep the balloon with the maximum neededTime (the most costly to

remove) and remove all others. Hence, we subtract the mx from the total s and add this to our answer (ans).

The algorithm is efficient because it goes through the string once, using a while loop, checking each sequence of same-colored balloons and calculating the time cost on the fly. It has a linear time complexity relative to the length of the colors string, making it suitable even for a large number of balloons.

Solution Approach

The solution implementation follows a greedy approach, which entails iterating through the colors string and grouping balloons with

the same color. During this grouping, the algorithm also keeps track of the sum of neededTimes of these balloons (s) as well as the

### maximum neededTime for a single balloon in the group (mx). To do this, two pointers are used. The first pointer i marks the beginning of a group of same-colored balloons, and a second pointer j is used to find the end of this group.

Initialize ans as zero, which will store the minimum total time required to make the rope colorful.
 Loop through the string colors using the index i, which serves as the starting point of each group of same-colored balloons:

Initialize s as zero, which will hold the sum of the neededTime for all balloons currently in the sequence.

The approach can be broken down into the following steps:

o Initialize my as zero, which will keen track of the maximum neededTime of a halloon in the cu

Within this loop, start a nested loop with index j, beginning at the same position as i.

Initialize mx as zero, which will keep track of the maximum neededTime of a balloon in the current sequence.

For each balloon that has the same color as the one at position i (i.e., colors[j] == colors[i]), accumulate its neededTime

in s and update mx if the current balloon's neededTime[j] is greater than the current mx.

Set i to the value of j to begin processing the next sequence.

remove based on the criteria of the lowest time cost, which aligns with our greedy strategy.

1. We initialize ans as zero. This will keep track of the minimum total time needed.

- Increment j for as long as it does not exceed the length of the colors string and the color of the jth balloon is the same as the ith balloon.
- neededTime (mx) from the sum of neededTime (s) to get the time required to make the current sequence colorful and add this value to ans.

□ If the sequence length (j - i) is greater than 1, it means there are duplicate balloons. In this case, subtract the maximum

3. Once the loop is complete, all the neededTime for removing necessary balloons to avoid consecutive same colors has been added to ans. Return ans as the answer.
It is important to note that because this approach always chooses the most expensive balloon to keep in each consecutive sequence of same-colored balloons, the accumulated removal time is as small as possible. Essentially, you're selecting which balloons to

sequence of balloons colorful by removing consecutive balloons of the same color in the least total time possible.

Using the solution approach outlined earlier, we can work through this step by step:

2. We will loop through the string colors using index i. For simplicity, we consider each character as a group of same-colored

Let's consider an example where colors = "abcccbad" and neededTime = [1, 2, 3, 4, 5, 6, 7, 8]. Our goal is to make the

#### Start with i = 0; this is the first balloon with color 'a'.

balloons, even if it's a group of one.

Now at i = 1; this balloon has color 'b'.

At i = 2; this balloon has the color 'c'.

 $\circ$  Initialize s = 0 and mx = 0.

becomes 3).

example.

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C++ Solution

1 class Solution {

2 public:

**Python Solution** 

class Solution:

from typing import List

Example Walkthrough

• There are no consecutive balloons with the same color, so i is incremented to 1.

- We set j = 2 and start the nested loop.
- We go through the sequence, updating s and mx as follows:

There are no consecutive balloons with the same color, so i is incremented to 2.

We find there is a sequence of balloons with the same color 'c' starting from index 2 to 4.

For j = 4, colors[j] = 'c', s += neededTime[4] (s becomes 12) and mx = max(mx, neededTime[4]) (mx becomes 5).
 Since j - i (4 - 2) is greater than 1, we have duplicate balloons and we calculate the time to make the sequence colorful:

ans += s - mx (ans becomes 12 - 5 = 7).

def minCost(self, colors: str, neededTime: List[int]) -> int:

while i < n and colors[i] == colors[start]:</pre>

total\_time += sum\_time - max\_time

max\_time = max(max\_time, neededTime[i])

# If there was more than one balloon in the sequence,

// Process consecutive balloons of the same color

totalMinCost += sumNeededTime - maxNeededTime;

sumNeededTime += neededTime[endIndex];

// Move to the next balloon

if (endIndex - startIndex > 1) {

++endIndex;

return totalMinCost;

// Add the needed time of the current balloon to the sum

// Update the maximum needed time for the current sequence

maxNeededTime = Math.max(maxNeededTime, neededTime[endIndex]);

// If there is more than one balloon of the same color consecutively,

// Return the total minimum cost to make all balloons have distinct colors

sum\_time += neededTime[i]

# Move to the next balloon.

i += 1

**if** i - start > 1:

return total\_time

# Add the time for painting this balloon to the sum.

# Initialize the answer and the iterator i.

- We set i to 5 as the next starting point.
- 3. Moving on, loop continues, but no more sequences of consecutive colors are found.

  The final ans represents the minimum total time to remove all the consecutive same-colored balloons and is equal to 7 in this

Continuing with i = 5, there are no more consecutive balloons with the same color, so nothing gets added to ans.

■ For j = 2, colors[j] = 'c', so we update s += neededTime[2] (s becomes 3) and mx = max(mx, neededTime[2]) (mx

■ For j = 3, colors[j] = 'c', s += neededTime[3] (s becomes 7) and mx = max(mx, neededTime[3]) (mx remains 3).

total\_time = 0
i = 0
n = len(colors)

# Update the max\_time if this balloon's time is greater than the current max.

# Return the total minimum time to repaint the balloons so that no adjacent balloons have the same color.

# add the total time minus the time of the most expensive to repaint balloon.

# No need for updating i here, as it's already one past the current sequence.

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# Process the string until we reach the end of colors.

while i < n:

# Start of the same color sequence.

start = i

# Initialize the sum and the maximum time for the current sequence.

sum_time = 0

max_time = 0

# Iterate over the sequence of same colors.
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37 # Example usage:
38 # sol = Solution()
39 # colors = "abaac"
40 # neededTime = [1,2,3,4,5]
41 # print(sol.minCost(colors, neededTime)) # Should output 3
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Java Solution
   class Solution {
       public int minCost(String colors, int[] neededTime) {
           // Initialize the total minimum cost to 0
           int totalMinCost = 0;
           // Obtain the total number of balloons
           int balloonsCount = neededTime.length;
           // Iterate through the array of needed time to check for consecutive balloons of the same color
           for (int startIndex = 0, endIndex = 0; startIndex < balloonsCount; startIndex = endIndex) {</pre>
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               // Move the endIndex to the start of the current sequence
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               endIndex = startIndex;
               // Initialize the sum of needed time for all balloons in the current sequence and the maximum needed time
               int sumNeededTime = 0, maxNeededTime = 0;
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while (endIndex < balloonsCount && colors.charAt(endIndex) == colors.charAt(startIndex)) {</pre>

// increment the total minimum cost by the sum of needed times minus the maximum needed time

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int minCost(string colors, vector<int>& neededTime) {
           // This variable will hold the minimum total time
           int totalTime = 0;
           // The size of the colors string
           int n = colors.size();
           // Loop through the colors string
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           for (int i = 0, j = 0; i < n; i = j) {
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               // Look for a sequence of the same color
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               j = i;
               // 'sumTimes' holds the sum of time for balloons of the same color group
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               int sumTimes = 0;
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               // 'maxTime' holds the maximum time needed for a single balloon in the same color group
               int maxTime = 0;
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               // While we haven't reached the end of the string and the current color is
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               // the same as the one we started this sequence with, add to sumTimes and
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               // update maxTime if necessary
               while (j < n && colors[j] == colors[i]) {</pre>
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                    sumTimes += neededTime[j];
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                   maxTime = max(maxTime, neededTime[j]);
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                   ++j;
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               // If we have more than one balloon with the same color in sequence,
               // add to the total time the sum of needed time minus the maximum time
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               // needed for a single balloon (we are keeping the one which takes longest to remove)
               if (j - i > 1) {
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                    totalTime += sumTimes - maxTime;
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           // Return the total time calculated
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           return totalTime;
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38 };
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Typescript Solution
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# // as the one we started this sequence with, add to sumTimes and update maxTime while (j < n && colors[j] === colors[i]) { sumTimes += neededTime[j]; maxTime = Math.max(maxTime, neededTime[j]); j++;</pre>

j = i;

let sumTimes = 0;

let maxTime = 0;

if (j - i > 1) {

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return totalTime;

Time and Space Complexity

Time Complexity

the inner loop does not lead to a quadratic time complexity, because each character in the string colors is visited exactly once. After the inner loop, the index i jumps to the position j, skipping all the characters that have already been accounted for.

Thus, each character in the string causes an iteration of the outer loop, and the inner loop only runs for characters of a sequence of

The provided code consists of a single while loop that iterates through the string colors. Inside the loop, there's another while loop

that also iterates through the string but only when it finds consecutive characters that are the same. Despite this nested structure,

identical colors once. This leads to a linear time complexity with respect to the length of the string colors. Hence, the time complexity is O(n), where n is the length of colors.

function minCost(colors: string, neededTime: number[]): number {

// Loop through the colors string

for (let i = 0, j = 0; i < n; i = j) {

// Start of the same color sequence

totalTime += sumTimes - maxTime;

// Return the total minimum time calculated

let totalTime = 0; // This will hold the minimum total time

// Hold the sum of times for balloons of the same color group

// Hold the maximum time needed for a single balloon in the same color group

// This is because we're removing all but one balloon in the sequence

// While we haven't reached the end of the string and the current color is the same

// Add to total time the sum of needed times minus the max needed time for this sequence

const n = colors.length; // The size of the colors string

### Space Complexity

The code uses a constant amount of extra space: variables ans, i, s, mx, and j. When iterating over the input, no additional space that scales with the size of the input is used. The input itself (colors and neededTime) is not counted towards the space complexity.

Therefore, the space complexity is 0(1) since it does not allocate any additional space that grows with the input size.