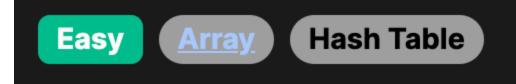
## 1. Two Sum



**Problem Description** 

In this problem, we have an array of integers called nums and a target integer called target. Our task is to find two distinct numbers within the array that when added together, equal the target. One important rule is that we cannot use the same element from the array twice in our sum. The output will be the indices (positions in the array) of these two numbers, and it does not matter in which order we provide these indices. Each valid input to the problem is guaranteed to have exactly one solution.

### Intuition

To solve this problem efficiently, we avoid the brute force approach of trying every possible pair of numbers to see if they add up to the target. Instead, we employ a hash table (dictionary in Python), which offers fast lookup times that average out to O(1) complexity. The basic idea is to iterate through the list once, and for each number, we check if the number required to reach the target (target - current\_number) has already been seen in the array. If it has, then we have found the two numbers whose sum is equal to the target.

# Solution Approach

ups which are critical for efficiency. The algorithm proceeds as follows: 1. Initialize an empty hash table (dictionary in Python dialect), we'll call it m.

The solution provided uses a hash table to map each element's value to its index in the array. This allows for constant-time look-

- 2. Iterate over the nums array, enumerating both the value x and its index i. Enumeration provides a convenient way of getting both the value and the index without additional overhead.
- 4. Check if y is present as a key in the hash table. If it is found, it means we had already seen the necessary pair earlier in the array. We then

3. For every value x, calculate its complement y by subtracting x from target (y = target - x).

- retrieve m[y], which is the index of y we had stored, and return a list containing the indices of y and x ([m[y], i]). This satisfies the requirement as their sum is equal to the target.
- 5. If y is not in the hash table, add the current value x along with its index i to the hash table (m[x] = i). This stores x for future reference if we later come across its complement y.

By only traversing the array once, the overall time complexity is O(n), where n is the number of elements in nums. The space complexity is also 0(n) as in the worst case, we could potentially store all elements in the hash table before finding a match.

class Solution:

Here's the provided solution approach step in the code:

```
def twoSum(self, nums: List[int], target: int) -> List[int]:
      m = {} # Step 1: Initialize the hash table
      for i, x in enumerate(nums): # Step 2: Enumerate through nums
          y = target - x # Step 3: Calculate complement y
         if y in m: # Step 4: Check if y is present in the hash table
              return [m[y], i] # Return the indices of the two elements adding up to target
         m[x] = i # Step 5: Store x with its index in hash table for future reference
The strength of this approach is that it smartly utilizes the hash table to avoid nested loops and thus reducing the time
```

complexity. The algorithm is linear time as it eliminates the need to examine every possible pair by keeping a record of what is needed to reach the target with the current number.

### Let's work through a small example to illustrate the solution approach.

**Example Walkthrough** 

Suppose our input array nums is [2, 7, 11, 15] and the target is 9.

Initialize an empty hash table (dictionary), which we'll name m.

Following the solution steps:

Now we start iterating over the nums array:

- First iteration (i = 0):
- - We take the first element nums [0] which is 2. ■ Calculate its complement y = target - nums[0], which gives us y = 9 - 2 = 7.
    - Check if 7 is present in the hash table m. It isn't, since m is empty.
    - Store the current value 2 along with its index 0 in the hash table: m[2] = 0.
    - Second iteration (i = 1):

def twoSum(self, nums: List[int], target: int) -> List[int]:

// Create a hashmap to store the value and its index

int current = nums[i]; // Current element value

Map<Integer, Integer> indexMap = new HashMap<>();

// Iterate over the array elements

for (int i = 0; i < nums.length; i++) {</pre>

std::unordered\_map<int, int> num\_to\_index;

for (int i = 0; i < nums.size(); ++i) {</pre>

// Iterate through each number in the vector.

return [numIndexMap.get(complement)!, i];

for index, number in enumerate(nums):

index\_map[number] = index

adding an element to m—is O(1).

# Calculate the complement of the current number

return [index\_map[complement], index]

int current\_num = nums[i]; // Current number in the iteration.

# Create a dictionary to store numbers and their indices

- Take the next element nums [1], which is 7. ■ Calculate its complement y = target - nums[1], which gives us y = 9 - 7 = 2.
- - Check if 2 is in the hash table m. Yes, it is! The complement 2 was stored during the first iteration.
  - Since the complement is found, we retrieve m[2] which is 0, the index of the complement 2. This gives us the indices [0, 1].
- The sum of the numbers at these indices (nums[0] + nums[1]) equals the target (2 + 7 = 9). As expected, the original problem statement guarantees that there is exactly one solution, so the problem is now solved with the output [0, 1].

By using this hash table approach, we efficiently found the two numbers that add up to the target in a single pass through the array, thereby using 0(n) time complexity instead of 0(n^2) which would result from using a brute force approach with nested

loops. Solution Implementation

## **Python**

index\_map = {}

### from typing import List # Import List for type annotation class Solution:

```
# Enumerate through the list of numbers
        for index, number in enumerate(nums):
           # Calculate the complement of the current number
            complement = target - number
           # If complement is in the index_map, a solution is found
            if complement in index_map:
               # Return the indices of the two numbers
               return [index_map[complement], index]
            # Otherwise, add the current number and its index to the index_map
            index_map[number] = index
       # If no solution is found, this return will not be reached due to guaranteed solution.
Java
import java.util.Map;
import java.util.HashMap;
class Solution {
    public int[] twoSum(int[] nums, int target) {
```

```
// Check if the complement is already in the map
            if (indexMap.containsKey(complement)) {
                // If complement is found, return the indices of the two numbers
                return new int[] {indexMap.get(complement), i};
            // Store the current value and its index in the map
            indexMap.put(current, i);
       // Note: The problem statement guarantees that there will always be exactly one solution,
       // so no need to return null or throw an exception here.
        throw new IllegalArgumentException("No two sum solution found");
C++
#include <vector>
#include <unordered_map>
class Solution {
public:
    // Function to find the indices of the two numbers that add up to a specific target.
    std::vector<int> twoSum(std::vector<int>& nums, int target) {
       // Create a hash map to store the value and its index.
```

int complement = target - current; // The complement which, when added to 'current', equals 'target'

```
int complement = target - current_num; // Find the complement of the current number.
           // If the complement is found in the map, return the pair of indices.
            if (num_to_index.count(complement)) {
                return {num_to_index[complement], i};
            // If complement is not found, add the current number and its index to the map.
            num_to_index[current_num] = i;
       // In case no solution is found (this part is unreachable if input is guaranteed to have one solution as stated in the pr
       return {}; // Return an empty vector.
};
TypeScript
// Define the function signature with input types and return type
function twoSum(nums: number[], target: number): number[] {
    // Initialize a map to store the numbers and their indices
   const numIndexMap = new Map<number, number>();
    for (let i = 0; i < nums.length; i++) {</pre>
        const currentNum = nums[i]; // The current number in the array
       const complement = target - currentNum; // The number that complements the current number to reach the target
       // Check if the complement is already in the map
       if (numIndexMap.has(complement)) {
```

```
// If complement is not found, add current number and its index to the map
          numIndexMap.set(currentNum, i);
      // Since the problem guarantees exactly one solution, the loop should never finish without returning.
      // If no solution is found (which violates the problem's constraints), throw an error.
      throw new Error('No two sum solution exists');
  // The function `twoSum` can now be used as per its signature with TypeScript type checking.
from typing import List # Import List for type annotation
class Solution:
   def twoSum(self, nums: List[int], target: int) -> List[int]:
       # Create a dictionary to store numbers and their indices
        index_map = {}
       # Enumerate through the list of numbers
```

#### complement = target - number # If complement is in the index\_map, a solution is found if complement in index\_map: # Return the indices of the two numbers

# Otherwise, add the current number and its index to the index\_map

# If no solution is found, this return will not be reached due to guaranteed solution.

// If complement is found, return its index along with the current index

Time and Space Complexity The time complexity of the code is O(n), where n is the length of the array nums. This is because the code loops through each

element in nums exactly once, and each operation within the loop—including checking if an element exists in the map m and

The space complexity of the code is also 0(n), since in the worst case, the code may insert each element of the array nums into

the map m. Therefore, the space used by the map m grows linearly with the number of elements in nums.