2310. Sum of Numbers With Units Digit K

Medium Greedy Math Dynamic Programming Enumeration

Problem Description

In this problem, we are given two integers, num and k. We need to construct a set of positive integers that satisfy two conditions:

1. The units digit (rightmost digit) of each integer in the set is k.

- 2. The sum of all the integers in the set is num.
- 2. The sum of all the integers in the set is num

we should return -1. It's also important to note that a set can contain multiple instances of the same integer, and an empty set has a sum of 0.

For example, if num = 58 and k = 9, one possible set satisfying the conditions could be $\{9, 9, 9, 9, 9, 9, 9, 4\}$, which has a sum of

We are asked to find the minimum possible size of such a set. If it's not possible to find any such set that satisfies the conditions,

58 and each integer has a 9 as its units digit. The goal is to find the smallest such set.

To come up with a solution, we need to first understand a few points about numbers and their units digits:

Intuition

The unit digit of a sum depends solely on the unit digits of the summands.
Since we can have multiple instances of the same number in our set, we can think of this problem as trying to reach num by adding a certain

- number of k's to the multiple of 10 (since the units digit of a multiple of 10 is always 0).
- The solution starts with an observation that we can keep adding the number k until we either reach num or we surpass it. When we add k to itself, we keep the unit digit the same and increase the tens digit. The key insight is that if num can be expressed as

a sum of integers with a unit digit k, then there must be a combination where the remainder when num is divided by 10 is k, or num is a multiple of 10.

With this in mind, we iterate over the range from 1 to num. For each index i, we check if num - k * i is a non-negative number and is also a multiple of 10 (to ensure the rest of the digits form a multiple of 10). If both conditions are met, i is the minimum

and is also a multiple of 10 (to ensure the rest of the digits form a multiple of 10). If both conditions are met, i is the minimum number of times we have to add k to form part of a sum that equals num. Therefore, i is the smallest possible size of our desired set.

If after iterating through the possible values of i we do not find a suitable number that meets the conditions, it means that it is

Solution Approach

The solution follows a straightforward brute-force approach. Since the problem requires finding the minimum number of integers

required to create the sum num, where each integer has the unit digit k, the algorithm performs a simple iteration to check if a

set can be constructed for each possible size incrementally.

Here is a breakdown of the implementation:

1. **Early Return for Zero:** If num is 0, we return 0 immediately since no number is needed to sum up to zero.

On each iteration, we calculate t, which is num - k * i. This represents the part of the sum that is not contributed by adding numbers

Iteration Over Possible Set Sizes:

with unit digit k.

We use the walrus operator := to assign the value to t while checking the condition in the same statement.

• We start a loop where i runs from 1 to num inclusive. This loop represents the iteration over possible sizes of the set.

- 3. Checking Validity of the Set:
 - Finding the Minimum Set Size:

 o If a valid set is found (the conditions in step 3 are met), the current value of i represents the minimum size of the set needed, and i is

We check if t is non-negative and is also a multiple of 10. Both conditions are required for t to potentially be the sum of integers with the

returned. 5. **Return -1 if No Set Exists:**

Let's illustrate the solution approach using a small example where num = 37 and k = 8.

unit digit 0, which when combined with i instances of k would sum up to num.

To check if t is a multiple of 10, we simply verify if t % 10 == 0.

not possible to form num exclusively with numbers ending in k, and we return -1.

- ∘ If no such i is found after completing the loop, it indicates that it is not possible to construct such a set. Therefore, the function returns -1.
- used can be identified as a brute-force search, checking all possible combinations until a valid one is found or until all options have been exhausted.

No complex data structures are required for this solution as it relies entirely on integer operations and a single loop. The pattern

One important algorithmic insight is the leveraging of modular arithmetic to identify potential matches for the conditions set by

the problem. This implementation efficiently approaches the problem with a time complexity of O(num), where num is the integer

representing the sum we are trying to reach.

Example Walkthrough

Following the provided steps:

I. Early Return for Zero: Since num is not 0, we do not return immediately and continue with the iterative process.

We start the loop with i=1 and go up to i=37 (because num = 37). In each iteration, we compute the value t as num - k * i. Checking Validity of the Set:

•

For each value of i, we want to check if t is a non-negative multiple of 10. In other words, t % 10 should be equal to 0.

• i = 1: t = 37 - 8 * 1 = 29. Since 29 % 10 ≠ 0, the set {8} with size 1 is not enough.

Return -1 if No Set Exists:

Solution Implementation

if num == 0:

return 0

for i in range(1, num + 1):

return i

remaining_value = num - k * i

public int minimumNumbers(int num, int k) {

Iteration Over Possible Set Sizes:

Let's walk through a few iterations in our example:

- i = 2: t = 37 8 * 2 = 21. Again, 21 % 10 ≠ 0, so the set {8, 8} with size 2 is not enough.
- i = 4: t = 37 8 * 4 = 5. Since $5 \% 10 \ne 0$, even the set $\{8, 8, 8, 8\}$ with size 4 does not meet the requirement.
- add up to 37 does not exist. We will never find a non-negative t that is a multiple of 10 because num itself is not a multiple of

If num is 0, no number needs to be added, so return 0

Iterate through i from 1 up to the given number

Continuing this process, we finally reach:

10 and k does not have a unit digit that could make up the difference.

4. Finding the Minimum Set Size:

• Since none of our checks for a valid set has passed, we don't have a minimum size to return.

∘ Since the loop has finished and no non-negative multiple of 10 was found for t, we return -1.

Calculate the remaining value after subtracting i times k from num

If both conditions are met, i is the minimum count of numbers needed

If no such number exists that the remaining value is divisible by 10, return -1

Check if remaining value is non-negative and divisible by 10

if remaining value >= 0 and remaining value % 10 == 0:

// Method to find the minimum number of non-negative integers needed

// If the sum required is 0, no numbers are needed.

// If num is 0, no numbers are needed, so return 0.

if (remainder >= 0 && remainder % 10 == 0) {

for (int count = 1; count <= num; ++count) {</pre>

int remainder = num - k * count;

// with last digit 'k'

// Iterate over possible counts of numbers with last digit 'k'

// Calculate the remaining value after subtracting 'count' numbers

// it means the number can be formed by 'count' numbers ending with 'k'

// If it's not possible to sum to 'num' with integers ending in 'k', return -1

// If the remainder is non-negative and is a multiple of 10,

return count; // Return the count as the answer

// where the last digit of each integer is 'k', and their sum is 'num'.

Thus, in this example, we can't find a set consisting of positive integers with the unit digit k (which is 8) that sums up to num (which is 37). The smallest possible set does not exist, hence the return value is -1.

Python

class Solution:
 def minimumNumbers(self, num: int, k: int) -> int:

We reach i = 5: t = 37 - 8 * 5 = 37 - 40 = -3. This is a negative number, so it is evident that the set with multiple 8s that

Java

 $if (num == 0) {$

if (num == 0) {

return -1;

return 0;

return -1

class Solution {

```
return 0;
        // Iterate from 1 to 'num' to find the minimum count of integers needed.
        for (int i = 1; i \le num; ++i) {
            // Calculate the remaining value after subtracting k * i'.
            int remainder = num - k * i;
            // If the remainder is non-negative and ends with a 0.
            // it means we found a valid group of 'i' numbers that sum up to 'num'.
            if (remainder >= 0 && remainder % 10 == 0) {
                return i; // Return the count 'i' as the result.
        // If no valid combination is found, return -1.
        return -1;
C++
class Solution {
public:
    // Function to find the minimum number of integers needed whose last digit is 'k'
    // that add up to the number 'num'
    int minimumNumbers(int num, int k) {
```

```
// This funct
```

};

```
// This function finds the minimum number of integers needed to sum up to a given number 'num',
// where each integer must end with the digit 'k'.
function minimumNumbers(num: number, k: number): number {
   // If 'num' is 0, no numbers are needed so return 0.
   if (num === 0) {
        return 0;
   // Determine the last digit of 'num'.
    let lastDigit = num % 10;
   // Check every possible integer with the last digit 'k'.
    for (let i = 1; i <= 10; i++) {
        let currentNumber = i * k:
       // The maximum number of integers with the last digit 'k' that we would need is 10,
       // as every digit from 0 to 9 will have appeared with 'k' at least once.
       // If 'currentNumber' is less than or equal to 'num' and ends with the same digit as 'num',
       // it's possible to sum up to 'num' with 'i' numbers ending with 'k'.
       if (currentNumber <= num && currentNumber % 10 === lastDigit) {</pre>
            return i;
   // If no such number could be found within the range, return -1 to indicate failure.
   return -1;
class Solution:
   def minimumNumbers(self, num: int, k: int) -> int:
       # If num is 0, no number needs to be added, so return 0
       if num == 0:
```

Check if remaining value is non-negative and divisible by 10 if remaining value >= 0 and remaining value % 10 == 0: # If both conditions are met, i is the minimum count of numbers needed return i

Time and Space Complexity

return -1

class Solution:
 def minimumNumbers(self, num: int, k: int) -> int:
 # If num is 0, no number needs to be added, so return 0
 if num == 0:
 return 0

Iterate through i from 1 up to the given number
 for i in range(1. num + 1):
 # Calculate the remaining value after subtracting i times k from num
 remaining_value = num - k * i

performs constant-time operations such as arithmetic subtraction, multiplication, and modulus operation. Therefore, the upper bound of loop iterations directly scales with the value of num.

The time complexity of the given code is O(num) because the loop runs a maximum of num iterations. In each iteration, it

The space complexity of the code is 0(1) because the space used does not depend on the size of the input num. The algorithm uses a fixed number of variables (i and t) that do not scale with input size.

If no such number exists that the remaining value is divisible by 10, return -1