594. Longest Harmonious Subsequence

Hash Table Sorting **Easy**

Problem Description

This problem asks for the length of the longest harmonious subsequence within an integer array nums. A harmonious array is defined as one where the difference between the maximum and the minimum values is exactly 1. It's important to note that a subsequence is different from a subset, as a subsequence maintains the original order of elements, and can be formed by removing zero or more elements from the array.

For example, given an array nums = [1,3,2,2,5,2,3,7], the longest harmonious subsequence is [3,2,2,2,3] with a length of 5.

Intuition

the approach is to check for each element num, whether there exists an element num + 1 in the array. If it exists, then a harmonious subsequence can potentially be formed using the elements num and num + 1. Since the elements need not be adjacent, we only care about the counts of the elements num and num + 1. We can find the total

The solution to this problem leverages hashing to keep track of the frequency of each number in the given array. The essence of

count of such pairs and keep updating the maximum found so far if the current count exceeds the previous maximum. Using a hash map or, in Python, a Counter object is ideal for this task as it allows us to efficiently store and access the frequencies of each element. Here's the intuitive breakdown of the process:

2. Iterate through each element num in the array.

3. Check if there's an element num + 1 in the counter hashmap.

1. Create a frequency counter (hash map) from the array to count occurrences of each element.

- 4. If so, calculate the sum of the count of num and num + 1 since those two form a harmonious sequence.
- 5. Compare this sum with the current longest harmonious subsequence length, and update it if this one is longer. 6. Continue this process until you have checked all elements.
- 7. Return the longest length obtained.
- This approach is efficient because it eliminates the need to consider actual subsequences. Instead, it relies on counts, which simplifies the process of verifying the harmonious condition.
- The implementation of the solution follows these steps:

used to count objects and store them as dictionary keys and their counts as dictionary value. Here it is used to build a hash

Solution Approach

map of each number (num) in nums array to its frequency.

subsequence we've seen so far.

Here's the key part of the code explained:

return ans # Step 5: Return the result

harmonious subsequence.

Iterate Through Each Element: We iterate through each element in the nums array. For each element num, we are going to check if num + 1 is a key in the counter. Check and Calculate: If num + 1 exists in our counter, we then know we can form a harmonious subsequence with num and

num + 1, since our Counter contains the counts of all numbers and we only want the difference between numbers to be

Counter Data Structure: We use Python's Counter class from the collections module, which is a subclass of a dictionary. It is

Maximize Answer: Now, having the sum of counts of num and num + 1, we compare it with our current answer (initially zero). If our sum is greater, we update the answer to this larger sum. Essentially, we are keeping track of the largest harmonious

exactly 1. To calculate the length of this potential subsequence we add the count of num and the count of num + 1.

- **Return the Result**: Finally, after iterating over all the elements in the nums array, we end up with the length of the longest harmonious subsequence in the ans variable. The algorithm uses O(N) space due to the counter, where N is the number of elements in the nums array. The time complexity is
- counter = Counter(nums) # Step 1: Build the counter hash map ans = 0for num in nums: # Step 2: Iterate through each element

Example Walkthrough

This code snippet illustrates how we use the algorithms and data structures mentioned to find the length of the longest

We now start to iterate through each element in nums. We take the first element 1 and check if 1 + 1 (which is 2) is a key in the counter. It is, so we find the sum of the counts of 1 and 2, which is 2 + 2 = 4. We compare this sum 4 with our current

example is 4.

Python

Java

class Solution {

#include <algorithm>

int findLHS(vector<int>& nums) {

for (int num : nums) {

nums.forEach(num => {

HashMap.forEach((count, num) => {

});

for (auto& [num, count] : frequencyMap) {

if (frequencyMap.count(num + 1)) {

// Count the frequency of each number in the given nums array

++frequencyMap[num]; // Increment the count for the number

class Solution {

public:

Solution Implementation

from collections import Counter

for num in nums:

longest_harmonious_subseq_length = 0

if num + 1 in frequency_map:

return longest_harmonious_subseq_length

Iterate through each number in the nums list

Check if the number has a companion number that is one greater

Calculate its length: count of num + count of num + 1

Harmonious subsequence found with num and num + 1

answer (which is 0 since we just started) and update the answer to 4.

Let's consider a small example to illustrate the solution approach with an array nums = [1, 1, 2, 2, 3].

also O(N), since we iterate over the array once to build the counter and once again to find the ans.

ans = max(ans, counter[num] + counter[num + 1]) # Step 4: Maximize the answer

if num + 1 in counter: # Step 3: Check if `num + 1` is present

from nums, and each value is how many times that element appears.

We move to the next element, which is also 1. Since we've already considered this number and the counter hasn't changed, the calculation would be the same, thus no change in the answer.

Next, we consider number 2. We check if 2 + 1 (which is 3) is a key in the counter. It is, so we find the sum of the counts of 2

First, we create a counter from the array. Our counter will look like this: {1: 2, 2: 2, 3: 1}, where each key is the element

- and 3, which is 2 + 1 = 3. We compare this sum 3 with our current answer 4, and since 3 is less than 4, we don't update the answer. Then, we take the next 2, and just like the previous 2, it yields the same calculation, so no change occurs.
- harmonious subsequence involving the number 3. Having examined all the elements, we end up with the answer 4, which is the length of the longest harmonious subsequence

Finally, we consider the last element 3 and check for 3 + 1 (which is 4). This is not a key in the counter, so we don't have a

[1, 1, 2, 2]. This example validates our solution approach: using a counter to efficiently compute the length of the longest harmonious subsequence by only considering the frequencies of num and num + 1 for each number in the array. The final answer for this

class Solution: def findLHS(self, nums: List[int]) -> int: # Count the frequency of each number in the list using Counter frequency_map = Counter(nums) # Initialize the variable to store the length of the longest harmonious subsequence

```
current_length = frequency_map[num] + frequency_map[num + 1]
        # Update the answer with the maximum length found so far
        longest_harmonious_subseq_length = max(longest_harmonious_subseq_length, current_length)
# Return the length of the longest harmonious subsequence found
```

public int findLHS(int[] nums) {

```
// Create a HashMap to keep track of the frequency of each number
       Map<Integer, Integer> frequencyMap = new HashMap<>();
       // Count the occurrences of each number in the array.
        for (int num : nums) {
            frequencyMap.put(num, frequencyMap.getOrDefault(num, 0) + 1);
       // Initialize variable to keep track of the longest harmonious subsequence
        int longestHarmoniousSubsequence = 0;
        // Iterate through the numbers in the array
        for (int num : nums) {
            // Check if the number that is one more than the current number exists in the map
            if (frequencyMap.containsKey(num + 1)) {
                // If it exists, calculate the sum of the frequencies of the current number
                // and the number that is one more than the current number
                int currentLength = frequencyMap.get(num) + frequencyMap.get(num + 1);
                // Update the longest harmonious subsequence if the current sum is greater
                longestHarmoniousSubsequence = Math.max(longestHarmoniousSubsequence, currentLength);
       // Return the length of the longest harmonious subsequence found
       return longestHarmoniousSubsequence;
C++
#include <vector>
#include <unordered_map>
```

std::unordered_map<int, int> frequencyMap; // Map to keep track of the frequency of each number in nums

int longestHarmoniousSequence = 0; // Variable to hold the length of the longest harmonious sequence

// Check if there is a number in the map which is exactly one more than the current number

// longest and the total count of the current number and the number that is one more.

// If found, update the longestHarmoniousSequence with the larger value between the previous

longestHarmoniousSequence = std::max(longestHarmoniousSequence, count + frequencyMap[num + 1]);

// Iterate through the numbers in the array to find the longest harmonious sequence

```
// Return the length of the longest harmonious sequence found
       return longestHarmoniousSequence;
};
TypeScript
// Importing necessary types from 'collections' module
import { HashMap } from "collectable";
// Declare a HashMap to keep track of the frequency of each number in nums
let frequencyMap: HashMap<number, number> = HashMap.empty();
// Function to find the length of the longest harmonious subsequence in the nums array
function findLHS(nums: number[]): number {
   // Count the frequency of each number in the given nums array
```

frequencyMap = HashMap.update<number, number>(n => (n || 0) + 1, num, frequencyMap);

// Iterate through the numbers in the map to find the longest harmonious sequence

Check if the number has a companion number that is one greater

Calculate its length: count of num + count of num + 1

Update the answer with the maximum length found so far

current_length = frequency_map[num] + frequency_map[num + 1]

Harmonious subsequence found with num and num + 1

Return the length of the longest harmonious subsequence found

```
// Check if there is a number in the map which is exactly one more than the current number
          if (HashMap.has(num + 1, frequencyMap)) {
              // If found, update the longestHarmoniousSequence with the larger value between the previous
              // longest and the total count of the current number and the number that is one more.
              const nextCount = HashMap.get(num + 1, frequencyMap) || 0;
              longestHarmoniousSequence = Math.max(longestHarmoniousSequence, count + nextCount);
      }, frequencyMap);
      // Return the length of the longest harmonious sequence found
      return longestHarmoniousSequence;
from collections import Counter
class Solution:
   def findLHS(self, nums: List[int]) -> int:
       # Count the frequency of each number in the list using Counter
        frequency_map = Counter(nums)
       # Initialize the variable to store the length of the longest harmonious subsequence
        longest_harmonious_subseq_length = 0
       # Iterate through each number in the nums list
        for num in nums:
```

longest_harmonious_subseq_length = max(longest_harmonious_subseq_length, current_length)

let longestHarmoniousSequence: number = 0; // Variable to hold the length of the longest harmonious sequence

```
Time Complexity
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Time and Space Complexity

if num + 1 in frequency_map:

return longest_harmonious_subseq_length

The function involves calculating the frequency of each number in the list, which can be done in O(n) time where n is the number of elements in nums. The for loop iterates through each element in nums once, and each lookup and update operation within the loop can be considered to have an average-case time complexity of 0(1) due to the hash table (dictionary) operations in Python. Therefore, the total time complexity of this function is O(n).

Space Complexity

The space complexity comes from the use of a counter (essentially a hash table), which stores each unique number and its corresponding frequency. In the worst case, if all elements in nums are unique, the space required would be O(n). Thus, the overall space complexity of the function is O(n).