406. Queue Reconstruction by Height

Medium Greedy Binary Indexed Tree Segment Tree Array Sorting

Problem Description

In this problem, we are given an array of people's attributes, where each person is represented by a pair of values `[h, k]`. The `h` represents the height of the person and `k` represents the number of people who are in front of this person in a queue and are taller or of the same height. The array `people` does not necessarily represent the order of people in a queue. Our task is to reconstruct the queue in such a way that the attributes of people in the reconstructed queue align with the original array `people`. The reconstructed queue should also be represented as an array where each element follows the same pair representation `[h, k]`, with each pair denoting the attributes of the people in the queue.

To be more concrete, if `people[i] = [h_i, k_i]`, we need to place the `i`th person in such a position in the queue that `k_i` is the exact count of people ahead of this person with heights greater than or equal to `h_i`.

Intuition

The intuition behind the solution lies in the fact that it's easier to place the tallest people first and then insert the shorter people at their appropriate positions based on the `k` value. Here's how we can think of the approach:

1. Sort `people` by their height in descending order; if two people have the same height, sort them by their `k` value in ascending order. This way, we start with the tallest people, and for those with the same height, we place those with fewer people in front first.
2. Initialize an empty list `ans` that will represent our queue.
3. Iterate over the sorted list, inserting each person into the `ans` list at the index equal to their `k` value. It works because by inserting the tallest people first, their `k` value corresponds exactly to the position they should be in the queue. As we insert the shorter people, we displace the taller ones only if needed, which maintains the integrity of the `k` values.

By following this approach, we ensure that each person is placed in accordance with the number of taller or equally tall people that are supposed to be in front of them according to the original array `people`.

Solution Approach

The solution uses a greedy algorithm approach along with list operations to reconstruct the queue. Here is a step-by-step walkthrough of the implementation based on the solution code provided:

1. Sort the `people` array in a way such that we first sort the people by height in descending order and then, if the heights are equal, sort by the `k` value in ascending order. The lambda function in the `sort` method accomplishes this by using `-x[0]` for descending height and `x[1]` for ascending `k` values. Sorting is done with this line:

   ```
   people.sort(key=lambda x: (-x[0], x[1]))
   ```

   By sorting in this manner, we prioritize taller individuals and, among those of identical height, prioritize those with a smaller `k` value, which is the count of people in front of them.

2. Initialize an empty list `ans` which will be used to construct the queue.

   ```
   ans = []
   ```

3. Iterate over the sorted `people` array. For each person `p`, we insert them into the `ans` list at the index specified by their `k` value. This is possible because at the time of their insertion, there will be `p[1]` (which is `k`) taller or equal height people already in the list, since those have been inserted earlier due to the sorting strategy.

   ```
   for p in people:
       ans.insert(p[1], p)
   ```

   The `insert` operation places the person at the correct position in the queue, shifting elements that are already there.

4. Return the `ans` list which now represents the reconstructed queue.

   ```
   return ans
   ```

Using this greedy strategy, we are able to re-order the array so that each person is positioned in accordance with their height and the count of people in front of them with equal or greater height. The use of sorting and list insertions are key aspects of the algorithm, ensuring that the resulting `ans` list satisfies the requirements that are stated in the problem.

Example Walkthrough

Let's consider a small example to illustrate the solution approach.

Suppose we are given the following array `people`:

```
people = [[7, 0], [4, 4], [7, 1], [5, 0], [6, 1], [5, 2]]
```

Following the steps outlined in the solution approach:

1. We first sort the array `people` by height in descending order and `k` value in ascending order when the heights are equal. After sorting, our `people` array looks like this:

   ```
   people = [[7, 0], [7, 1], [6, 1], [5, 0], [5, 2], [4, 4]]
   ```

2. We initialize an empty list `ans`:

   ```
   ans = []
   ```

3. We then iterate over the sorted `people` array and insert each person into the `ans` list at the index specified by their `k` value:
   - We insert `[7, 0]` at index `0`: `ans = [[7, 0]]`
   - We insert `[7, 1]` at index `1`: `ans = [[7, 0], [7, 1]]`
   - We insert `[6, 1]` at index `1`: `ans = [[7, 0], [6, 1], [7, 1]]`
   - We insert `[5, 0]` at index `0`: `ans = [[5, 0], [7, 0], [6, 1], [7, 1]]`
   - We insert `[5, 2]` at index `2`: `ans = [[5, 0], [7, 0], [5, 2], [6, 1], [7, 1]]`
   - Finally, we insert `[4, 4]` at index `4`: `ans = [[5, 0], [7, 0], [5, 2], [6, 1], [4, 4], [7, 1]]`
4. The final `ans` list is our reconstructed queue which aligns with the original `people` array attributes:

   ```
   ans = [[5, 0], [7, 0], [5, 2], [6, 1], [4, 4], [7, 1]]
   ```

Using the greedy solution approach of sorting the individuals by height and then inserting them into the queue based on their `k` value, we have successfully reconstructed a queue that satisfies the requirements of the problem. In this queue, for every person, the `k` value is the exact count of people ahead of them with heights greater than or equal to their own.

Solution Implementation

Python

```python
from typing import List

class Solution:
    def reconstructQueue(self, people: List[List[int]]) -> List[List[int]]:
        # Sort the 'people' array in descending order of height (h),
        # and in the case of a tie, in ascending order of the number of people in front (k).
        people.sort(key=lambda person: (-person[0], person[1]))

        # Initialize an empty list to hold the final queue reconstruction
        queue = []

        # Iterate over the sorted 'people' list
        for person in people:
            # Insert each person into the queue. The index for insertion is the k-value,
            # which represents the number of people in front of them with equal or greater height.
            queue.insert(person[1], person)

        # Return the reconstructed queue
        return queue
```

Java

```java
import java.util.Arrays;
import java.util.ArrayList;
import java.util.List;

class Solution {

    public int[][] reconstructQueue(int[][] people) {
        // Sort the array. First, sort by height in descending order.
        // If heights are equal, sort by the number of people in front (k-value) in ascending order.
        Arrays.sort(people, (person1, person2) ->
            person1[0] == person2[0] ? person1[1] - person2[1] : person2[0] - person1[0]);

        // Initialize a list to hold the final queue reconstruction,
        // which will allow us to insert people at specific indices.
        List<int[]> reconstructedQueue = new ArrayList<>(people.length);

        // Iterate over the sorted array, and insert each person into the list
        // at the index specified by their k-value.
        for (int[] person : people) {
            // The second value of each person array (person[1])
            // specifies the index at which this person should be added in the queue
            reconstructedQueue.add(person[1], person);
        }

        // Convert the list back to an array before returning it.
        return reconstructedQueue.toArray(new int[reconstructedQueue.size()][]);
    }
}
```

C++

```cpp
#include <vector>
#include <algorithm>

class Solution {
public:
    std::vector<std::vector<int>> reconstructQueue(std::vector<std::vector<int>>& people) {
        // Sort the people to arrange them according to their height in descending order.
        // If two people are of the same height, order them by their "k-value" in ascending order.
        std::sort(people.begin(), people.end(), [](const std::vector<int>& person1, const std::vector<int>& person2) {
            return person1[0] > person2[0] || (person1[0] == person2[0] && person1[1] < person2[1]);
        });

        // Initialize an empty vector for the final arrangement.
        std::vector<std::vector<int>> arrangedQueue;

        // Insert each person into the arrangedQueue.
        // The index to insert the person is determined by the person's "k-value",
        // which indicates the number of people in front of this person who have a height greater than or equal to this person's
        for (const std::vector<int>& person : people) {
            arrangedQueue.insert(arrangedQueue.begin() + person[1], person);
        }

        // Return the final arrangement of the queue.
        return arrangedQueue;
    }
};
```

TypeScript

```typescript
function reconstructQueue(people: number[][]): number[][] {
    // Sort the `people` array to arrange them by their height in descending order.
    // If two people have the same height, order them by their k-value in ascending order.
    people.sort((a, b) => {
        return b[0] - a[0] || a[1] - b[1];
    });

    // Initialize an empty array for the final arrangement.
    let arrangedQueue: number[][] = [];

    // Insert each person into the `arrangedQueue`.
    // The index at which to insert is determined by the person's k-value,
    // which indicates the number of people in front of them with a height greater than or equal.
    for (let person of people) {
        arrangedQueue.splice(person[1], 0, person);
    }

    // Return the final arrangement of the queue.
    return arrangedQueue;
}

// Example of usage:
// let peopleInput: number[][] = [[7, 0], [4, 4], [7, 1], [5, 0], [6, 1], [5, 2]];
// let queue = reconstructQueue(peopleInput);
// console.log(queue); // Outputs arranged queue based on the given conditions.
```

```python
from typing import List

class Solution:
    def reconstructQueue(self, people: List[List[int]]) -> List[List[int]]:
        # Sort the 'people' array in descending order of height (h),
        # and in the case of a tie, in ascending order of the number of people in front (k).
        people.sort(key=lambda person: (-person[0], person[1]))

        # Initialize an empty list to hold the final queue reconstruction
        queue = []

        # Iterate over the sorted 'people' list
        for person in people:
            # Insert each person into the queue. The index for insertion is the k-value,
            # which represents the number of people in front of them with equal or greater height.
            queue.insert(person[1], person)

        # Return the reconstructed queue
        return queue
```

Time and Space Complexity

Time Complexity

The time complexity of the code can be broken down into two major operations:

1. Sorting the `people` array:

   This operation uses a sorting algorithm, typically Timsort in Python, which has a time complexity of `O(n log n)` for an array of n elements.

2. Reconstructing the queue by inserting elements:

   The `insert` operation within a list has a time complexity of `O(n)` in the worst case because it may require shifting the elements of the list.

   Since `insert` is called for each person in the sorted array and the list can grow up to n elements, this operation will have a worst-case complexity of `O(n^2)` as it performs n insertions into a structure that can be up to n in size.

By combining these, the overall time complexity of the code is `O(n log n) + O(n^2)`, which simplifies to `O(n^2)` since `O(n^2)` is the dominating term when n is large.

Space Complexity

The space complexity of the code is the amount of extra space used by the algorithm, not counting the space occupied by the input itself. In this case:

1. Sorted `people` array:

   Since the sorting operation is done in-place, it does not contribute to additional space complexity beyond what is taken by the input `people` array.

2. The `ans` list:

   This list is created to store the output and grows to contain n elements, which are the same elements from the input `people` array. Thus, the space required is `O(n)`.

Therefore, the overall space complexity of the code is `O(n)`.