2514. Count Anagrams String Hash Table Math Combinatorics Counting **Leetcode Link** Hard

The problem gives you a string s which consists of one or more words separated by single spaces. A string t is considered an

Problem Description

anagram of string s if for every word in t, there is a corresponding word in s that has exactly the same letters, but possibly in a different order. The goal is to calculate the number of distinct anagrams of s. Since the result could be a very large number, you're asked to return this number modulo 10^9 + 7, which is often used in programming contests to avoid dealing with exceedingly large numbers that could cause overflow errors.

length of the word.

Intuition To solve this problem, we need to think about what an anagram actually is. An anagram of a word is a permutation of its letters. If we consider each word separately, the number of different permutations (and thus anagrams) for a word is given by the factorial of the

That's because permuting the duplicates among themselves doesn't create unique anagrams. For example, for the word "aabb", it has 4 letters, so naïvely we might think there are 4! (which is 24) anagrams. However, both 'a'

However, there's a catch: if a word has duplicate letters, we must divide the total count by the factorials of the counts of each letter.

and 'b' are repeated twice. We need to divide 4! by the number of permutations of 'a' (which is 2!) and by the number of permutations of 'b' (which is 2!) to get the right answer: 4! / (2! * 2!), which equals 6.

The solution code first builds a list f of precomputed factorials modulo 10^9 + 7 to avoid recalculating factorials for each word, thereby saving time.

Then, for each word in the given string, we calculate the factorial of the length of the word, multiply this by the resulting ans so far, and then modify ans by multiplying it by the modular inverses of the factorials of the counts of each letter in the word. The function pow(f[v], -1, mod) is used to compute the modular inverse of f[v] under the modulo $10^9 + 7$.

modulo $10^9 + 7$. **Solution Approach**

Finally, we return ans, which—because we've taken the modulo at each step—will be the correct count of distinct anagrams of s

The provided solution code implements the following approach: 1. Precompute Factorials: Initially, the code precomputes factorials of numbers up to 10^5 and stores them in a list f. This list will

done modulo 10^9 + 7 to prevent integer overflow and to comply with the requirement of returning the answer modulo 10^9 + 7.

2. Iterate Over Words: The core logic of the solution iterates over each word within the input string s. We use the split() method to break the string into a list of words.

where it might cause overflow errors.

3. Count Letter Frequencies: For each word, a Counter (from Python's collections module) is used to calculate the frequencies of each letter within the word. The counter creates a dictionary that maps each letter to its frequency.

be used to look up the factorial of any number during the computation without having to recalculate it. This precomputation is

number of permutations of the letters without considering duplicate letters.

5. Adjust for Duplicate Letters: Since we need to account for repeated letters, we loop over the values of our letter frequency

4. Calculate Anagrams: For every word, we first multiply ans with the factorial of the length of the word. This represents the total

these factorials using the pow function, which allows us to compute inverse modulo operations. We then multiply ans by this inverse to correctly account for the permutations of the duplicate letters.

6. Keep Result within Modulo: Throughout these operations, we ensure that ans is kept within the modulo 10^9 + 7 by taking the

modulo operation after each multiplication. This step is crucial because it ensures that the value of ans never exceeds the limit

counter and get the factorial of each letter count. Using the precomputed factorials in f, we calculate the modular inverse of

7. Return the Result: After iterating through all the words and updating ans accordingly, we return ans, which is the total number of distinct anagrams of the original string s modulo 10^9 + 7. It's worth noting that the pow(x, -1, mod) computation is a way of finding the modular multiplicative inverse of x mod mod when x and mod are coprime (which they are in this case because the mod is prime and the factorial is not zero).

Using this approach, the solution code efficiently calculates the number of distinct anagrams of the given string s by breaking the

Let's walk through a small example to illustrate the solution approach using the string s = "abc cab".

Factorials up to the maximum length 10^5 are precomputed and stored in a list f modulo 10^9 + 7. For this example, we would have f[0] = 1, f[1] = 1, f[2] = 2, f[3] = 6, ... and so on.

For each word, we count letter frequencies. For abc, it's {'a': 1, 'b': 1, 'c': 1}. For cab, it's the same since cab is a permutation

Step 4: Calculate Anagrams

of abc.

Example Walkthrough

Step 1: Precompute Factorials

Step 2: Iterate Over Words

Step 3: Count Letter Frequencies

We start with ans = 1. For each word:

need to adjust for duplicates because all letters are distinct.

characters, steps 4 and 5 would show more adjustment for duplicates.

Count the frequency of each character in the word

answer *= pow(factorials[count], -1, MOD)

// Iterating over the character frequencies,

// Finally, return the product casted to an integer

// Calculate the number of anagrams for words in a string

for (int i = 1; i <= word.size(); ++i) {</pre>

int charIndex = word[i - 1] - 'a';

// Use modular inverse to divide ans by the factorial part

// Calculate x raised to the power of n mod MOD using fast exponentiation.

return (ans * pow(factorial, MOD - 2)) % MOD;

let result: number = 1; // Start with result as 1

// If n is odd, multiply the result by x.

result = (result * x) % MOD;

// Square x and reduce n by half.

function pow(x: number, n: number): number {

// Continue looping until n becomes zero.

// Increment the count for this character

// Break down the string into words and compute the result

for (int count : characterCount) {

if (count > 0) {

return (int) product;

// taking the modInverse of their factorial, and updating the product

product = (product * modInverse.intValue()) % MOD;

// We declare 'mod' as a static constant because it's a property of the class and

// to store each word individually

// to store the factorial part

// to store the number of permutations

std::stringstream ss(s); // used to break the string into words

int count[26] = {0}; // to count occurrences of each letter

// Calculate zero-based index for character in alphabet

BigInteger factorial = BigInteger.valueOf(factorials[count]);

BigInteger modInverse = factorial.modInverse(BigInteger.valueOf(MOD));

for count in character_counts.values():

the multiplicative inverse of the factorial of the character's count

We split the string s into words using the split() method, resulting in the list ["abc", "cab"].

problem down into calculations involving factorials and considering the impact of duplicate letters.

• For the word cab, we repeat the process. Multiplying ans by f[3] which is 6, we get ans = 36 modulo 10^9 + 7. **Step 5: Adjust for Duplicate Letters**

Step 6: Keep Result within Modulo

modulo $10^9 + 7$. So we would return ans = 36.

Pre-calculate factorial values modulo MOD

Split the string into words

answer %= MOD

for word in s.split():

Both words in our example do not have duplicate letters, so we don't need to adjust for duplicates. If there were duplicates, we would multiply ans by the modular inverses of the factorials of the letter counts.

words, there is no issue with overflow, and we don't need further modulo operations here. **Step 7: Return the Result**

We have completed the iteration over all words, and the final ans represents the total number of distinct anagrams of the string s

This step ensures that ans remains within the range allowed by modulo 10^9 + 7. In our case, since ans is 36 after processing both

• For the word abc, the length is 3, so we multiply ans by f[3] which is 6. Now, ans=ans * 6 % (10^9 + 7)->ans` = 6. We don't

That number, 36, says that there are 36 different permutations of s considering each word and its anagrams, but this is only for the

length and uniqueness of the chosen words "abc" and "cab". In a more complex example with longer words and repeating

Python Solution from collections import Counter

for i in range(1, 10**5 + 1):

answer = 1

return answer

MOD = 10**9 + 7

factorials = [1]

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C++ Solution

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#include <sstream>

#include <string>

class Solution {

// should not change.

static const int MOD = 1e9 + 7;

int countAnagrams(std::string s) {

std::string word;

long ans = 1;

long factorial = 1;

while (ss >> word) {

// Process each word in the string

++count[charIndex];

new_value = factorials[-1] * i % MOD factorials.append(new_value) 10 class Solution: 12 def countAnagrams(self, s: str) -> int: 13 # Initialize the answer as 1

character_counts = Counter(word) 19 # Multiply the answer by the factorial of the length of the word 20 answer *= factorials[len(word)] answer %= MOD 23 24 # For each character, update the answer by multiplying with

```
33 # sol = Solution()
34 # result = sol.countAnagrams("the quick brown fox")
35 # print(result)
36
```

32 # Example usage:

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Java Solution
    import java.math.BigInteger;
    class Solution {
         // Define the modulus constant for taking the mod of large numbers to prevent overflow
         private static final int MOD = (int) 1e9 + 7;
  6
         // Method to count the number of anagrams for groups of words separated by spaces
         public int countAnagrams(String s) {
  8
             int length = s.length();
  9
 10
             // Factorial array to store precomputed factorials under MOD
 11
             long[] factorials = new long[length + 1];
 12
             factorials[0] = 1; // 0! is 1
 13
 14
             // Calculate factorial for each number up to length, with mod
 15
             for (int i = 1; i <= length; ++i) {</pre>
 16
                 factorials[i] = (factorials[i - 1] * i) % MOD;
 17
 18
 19
             long product = 1; // Initialize the product to 1
 20
             // Split the given string by spaces and process each group of characters
 21
             for (String word : s.split(" ")) {
 22
                 int[] characterCount = new int[26];
 23
 24
                 // Count the frequency of each character in the current word
 25
                 for (int i = 0; i < word.length(); ++i) {</pre>
 26
                     characterCount[word.charAt(i) - 'a']++;
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                 // Multiply the product by the factorial of the length of the word
 30
                 product = (product * factorials[word.length()]) % MOD;
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26 27 28 29

```
// Update the permutation count: ans = ans * i
 30
                     ans = (ans * i) % MOD;
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 32
                     // Update the factorial part using the letter's occurrence
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                     factorial = (factorial * count[charIndex]) % MOD;
 35
 36
 37
             // Use modular inverse to divide ans by the factorial part
 38
             return static_cast<int>((ans * pow(factorial, MOD - 2)) % MOD);
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 41
         // Calculate x raised to the power of n mod MOD using fast exponentiation
 42
         long pow(long x, int n) {
 43
             long result = 1L; // start with result as 1
 44
             // Continue looping until n becomes zero
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 46
             while (n > 0) {
                 // If n is odd, multiply the result by x
 47
                 if (n % 2 == 1) {
 48
 49
                     result = (result * x) % MOD;
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 51
 52
                 // Square x and reduce n by half
 53
                 x = (x * x) % MOD;
 54
                 n /= 2;
 55
 56
 57
             return result;
 58
 59 };
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Typescript Solution
   // Define the modulo constant.
    const MOD: number = 1e9 + 7;
    // Calculate the number of anagrams for words in a string.
  5 // Break down the string into words and compute the result.
  6 function countAnagrams(s: string): number {
        const words: string[] = s.split(/\s+/); // Split the string into words
         let ans: number = 1; // To store the number of permutations
         let factorial: number = 1; // To store the factorial part
  9
 10
 11
         // Process each word in the string
 12
         for (let word of words) {
 13
             let count: number[] = new Array(26).fill(0); // To count occurrences of each letter
 14
             for (let i: number = 0; i < word.length; i++) {</pre>
 15
 16
                 // Calculate zero-based index for character in alphabet
 17
                 let charIndex: number = word.charCodeAt(i) - 'a'.charCodeAt(0);
 18
 19
                 // Increment the count for this character
 20
                 count[charIndex]++;
 21
 22
                 // Update the permutation count: ans = ans * (i + 1)
 23
                 ans = (ans * (i + 1)) % MOD;
 24
 25
                 // Update the factorial part using the letter's occurrence
 26
                 factorial = (factorial * count[charIndex]) % MOD;
 27
```

Time and Space Complexity

return result;

while (n > 0) {

if (n % 2 === 1) {

x = (x * x) % MOD;

n = Math.floor(n / 2);

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32 }

run only once, and each operation is a multiplication and a modulo operation, which are both 0(1). When analyzing the countAnagrams method, for each word in the input string s, the code counts the frequency of characters, which is O(K) where K is the

Time Complexity

length of the word. Multiplying by the factorial of the length of the word is 0(1) since we have precomputed the factorials. The for loop for dividing by the factorial of the count of each character runs at most 26 times (since there are only 26 letters in the alphabet), and each iteration does a modular inverse and multiplication, which can both be considered 0(1) under modular arithmetic. Therefore, for each word, the complexity is dominated by O(K). Since this is done for every word in the string, and if there are W words in s, the overall time complexity of countAnagrams is O(W*K). **Space Complexity**

The time complexity of the precomputation part (where the factorials are computed up to 10**5) is 0(N) where N is 10**5. This part is

The space complexity for storing the precomputed factorials is O(N), where N is 10**5. Aside from that, the space used by the

function countAnagrams depends on the size of the counter object, which is at most 26 for each word. Therefore, the space complexity for the function countAnagrams is 0(1), making the total space complexity of the program 0(N).