Medium

Problem Description

In this problem, we have two digits that are considered "lucky" – 4 and 7. A number is termed as a "lucky number" if and only if every digit in the number is either a 4 or a 7. We are given a task to find the k^th^ lucky number when all the lucky numbers are sorted in increasing order. We need to provide this lucky number as a string.

The challenge here is to calculate the k^th^ lucky number without generating all previous lucky numbers.

For instance, the first few lucky numbers are: 1st lucky number is "4" 2nd lucky number is "7" 3rd lucky number is "44" (and so on)

Intuition

digit 1 with 7, we get a system where each binary number corresponds to a lucky number. For example, binary 0 (which in our modified system would be "4"), binary 1 ("7"), binary 10 ("44"), binary 11 ("47") and so on. To find the k^th^ lucky number, we follow a similar approach as we would finding the k^th^ number in binary. However, unlike a

To solve this problem, we can draw an analogy from binary numbers. If we replace every binary digit 0 with 4 and every binary

standard binary system where each position can be 0 or 1, we only have two digits 4 and 7, representing the two possible states at each position of a lucky number. We determine the length of our lucky number (in digits) by finding the smallest number in such that k is less than 2ⁿ. This represents the level at which the kth lucky number exists if we were to visualize all lucky numbers in a binary tree form. Once we have the number of digits n, we can construct the lucky number from most significant digit to least significant digit, by checking if k is in the lower half (which would correspond to '4') or the upper half (which would correspond to '7') of the values for

that digit's position—this is akin to deciding between 0 and 1 in binary representation. If it's in the upper half, we know this digit is '7', and we subtract the size of the lower half (2^(n - 1)) from k to continue finding the rest of the digits. If it's in the lower half, we simply assign '4' to this current digit. We iterate this process until we have all n digits of our lucky number. Solution Approach

The solution implements the intuition discussed above with a focus on optimizing the process. Here's the detailed explanation of

the code step-by-step: We start by initializing n to 1. This n variable will eventually indicate the number of digits in our k^th^ lucky number.

sits.

in the upper half of the range for the next digit.

generating all previous lucky numbers.

The first while loop checks how many digits the k^th^ lucky number will have. It works under the principle that there are 2^n lucky numbers with n digits (2ⁿ - 1 possible combinations plus the all-4s combination). So, if k is greater than 2ⁿ, k is

- not within the range of lucky numbers that have n digits. Therefore, we subtract 2ⁿ from k and increment n by 1, and iterate until k is less than or equal to 2ⁿ. This locates the correct 'level' of the binary-tree-like structure where our number
- Once we have the number of digits n, we initialize an empty list ans, which will store each digit of the k^th^ lucky number as we compute it. The second while loop executes n times, decreasing n with each iteration. Each iteration of the loop decides one digit of the lucky number, starting from most significant to least significant. The conditional within this loop if k <= 1 << n is checking whether k fits in the lower half of the range for the current digit (which would correspond to '4'). If so, '4' is

appended to $\frac{1}{2}$ or $\frac{1}$

After the loop completes, ans holds the digits of our k'th' lucky number in order. We finally return the number as a string by joining each element of the list with "".join(ans).

This implementation uses a list (ans) as the primary data structure to build the lucky number and follows a binary search pattern

to decide each digit of the lucky number, improving the efficiency by avoiding unnecessary computation that would come from

Example Walkthrough Let's illustrate the solution approach using an example. We want to find the 5th lucky number. Initialize n to 1, because every lucky number has at least one digit.

Begin the first while loop to find the number of digits our 5th lucky number will have. We know there are 2ⁿ lucky numbers

for each digit length. Initially, n is 1, and 2ⁿ is 2, which is less than 5. So we increment n to 2, and now 2ⁿ is 4, which is still

less than 5. Incrementing n once more gives us 2ⁿ as 8, which is greater than 5. So, we stop here and know that our 5th

not change k.

lucky number has 3 digits.

We now start with k = 5 and n = 3. Initialize an empty list ans to store the digits.

Enter the second while loop, which will run three times (since our lucky number has 3 digits):

1 << (n-1) (which is 2) from k. k becomes 3, and we append '7' to ans.

We join these to form the 5th lucky number: "".join(ans) equals "477".

∘ In the third and final iteration, n is decremented to 1, making 1 << n equal to 2. k is 3, which is greater; therefore, we are again in the upper half. We subtract $1 \ll (n-1)$ from k, which is 1, making k now equal to 2. We append '7' to ans. The second while loop completes and our list ans has the values ['4', '7', '7'].

o In the first iteration, 1 << n equals 1 << 3 which equals 8. This is greater than k (5), so k is in the lower half. We append '4' to ans and do

o In the second iteration, we decrement n to 2. Now 1 << n is 4, which is less than k. Hence, k is in the upper half, and we need to subtract

- Therefore, the 5th lucky number is "477". This example clearly shows how the binary-like approach works by deciding one digit at a time, based on whether k is in the lower or upper half of the range for that digit.
- **Python**

Increment the digit count as we're moving on to numbers with more digits

If k is in the second half, append 7 and adjust k accordingly

If k is within the range of the first half, append 4, as it is the smaller digit

num digits = 1# Find the number of digits the kth lucky number must have while k > (1 << num digits):</pre> # Decrease k by the count of lucky numbers with num_digits digits

Construct the kth lucky number by going through each digit place while num digits: num digits -= 1 # Decrement digits count as we build the number from high to low

else:

answer_digits = []

Solution Implementation

def kth lucky number(self, k: int) -> str:

if k <= (1 << num digits):

answer_digits.append("4")

 $k = 1 \ll num diaits$

num digits += 1

Initialize the number of digits to be considered to 1

Initialize the answer as an empty list to hold the digits

// Function to find the kth lucky number where lucky numbers are

// positive integers whose decimal representation contains only the digits 4 and 7.

// Find the number of digits in the kth lucky number by using powers of 2

// If 'k' is greater, the digit is '7' and we adjust 'k' accordingly.

// Join the digits and return the resulting lucky number as a string.

k -= 1 << numDigits: // Subtract the number of numbers with `numDigits` digits

// Each additional digit doubles the count of lucky numbers available

++numDigits; // Move to the next digit length

class Solution:

```
answer digits.append("7")
                k -= 1 << num_digits # Decrement k as we have used one of the 7s
        # Join all the individual digits to form the kth lucky number and return it
        return "".join(answer_digits)
Java
class Solution {
    public String kthLuckyNumber(int k) {
        // 'n' represents the number of digits in the lucky number
        int n = 1;
        // Find the number of digits in the kth lucky number by comparing k with powers of 2
        while (k > (1 << n)) {
            k = (1 << n);
            ++n;
        // Build the kth lucky number starting with the most significant digit
        StringBuilder ans = new StringBuilder();
        while (n > 0) {
            // Check the kth bit of n to decide whether to append '4' or '7'
            // If k is in the first half of the range for the current digit length, append '4'
            if (k <= (1 << (n - 1))) {
                ans.append('4');
            } else {
                // If k is in the second half, append '7' and update k
                ans.append('7');
                k = (1 << (n - 1));
            n--;
        return ans.toString();
```

C++

public:

#include <string>

class Solution {

std::string kthLuckyNumber(int k) {

int numDigits = 1;

// Start counting digits from 1

while (k > (1 << numDigits)) {</pre>

luckyNumbers.push('4');

luckyNumbers.push('7');

k -= 1 << counter;

} else {

```
// Initialize an empty string to store the kth lucky number
        std::string luckyNumber;
        // Construct the lucky number digit by digit
        while (numDigits--) {
            // If the remaining k is less or equal to the number of lucky numbers with the current number of digits
            if (k <= (1 << numDigits)) {</pre>
                luckyNumber.push\_back('4'); // A '4' is appended when k is in the first half within the current digit's range
            } else {
                luckyNumber.push back('7'); // Otherwise, a '7' is appended
                k = 1 \ll numDigits; // And we adjust k to reflect that we're now considering the second half range
        // Return the constructed lucky number as a string
        return luckyNumber;
};
TypeScript
function kthLuckyNumber(k: number): string {
    // Initialize counter 'n' which represents the number of binary digits.
    let counter = 1;
    // As long as 'k' is greater than '2^n', decrease 'k' by '2^n' and increment 'n'.
    while (k > (1 << counter)) {</pre>
        k -= 1 << counter;
        ++counter;
    // Initialize an array 'luckyNumbers' to store the lucky number digits.
    const luckyNumbers: string[] = [];
    // Build the lucky number by determining if each digit is a '4' or a '7'.
    while (counter-- > 0) {
        if (k <= (1 << counter)) {
            // If 'k' is less than or equal to '2^n', the digit is '4'.
```

```
return luckyNumbers.join('');
class Solution:
    def kth lucky number(self, k: int) -> str:
       # Initialize the number of digits to be considered to 1
       num_digits = 1
       # Find the number of digits the kth lucky number must have
       while k > (1 << num digits):</pre>
           # Decrease k by the count of lucky numbers with num_digits digits
           k = 1 \ll num digits
           # Increment the digit count as we're moving on to numbers with more digits
           num_digits += 1
       # Initialize the answer as an empty list to hold the digits
        answer digits = []
       # Construct the kth lucky number by going through each digit place
        while num digits:
           num digits -= 1 # Decrement digits count as we build the number from high to low
           if k <= (1 << num digits):
                # If k is within the range of the first half, append 4, as it is the smaller digit
                answer_digits.append("4")
           else:
                # If k is in the second half, append 7 and adjust k accordingly
                answer digits.append("7")
                k -= 1 << num_digits # Decrement k as we have used one of the 7s
       # Join all the individual digits to form the kth lucky number and return it
        return "".join(answer_digits)
Time and Space Complexity
  The provided code calculates the k-th lucky number where lucky numbers are composed only of the digits 4 and 7.
```

Time Complexity

The main component of the code involves a while loop that runs until k is less than 1 << n, where n starts at 1 and gets incremented. Inside the loop, k is decremented by 1 << n. After finding the value of n such that k is within bounds, the code

the length of the binary representation of k, which can be described as $O(\log k)$.

executes another while loop that decreases n on each iteration and constructs the lucky number. In the worst-case scenario, n will be proportional to the number of digits in the k-th lucky number.

The construction of the lucky number (ans.append("4") or ans.append("7")) depends on the number of digits n. Since the loops' maximum number of iterations is equal to the number of digits, the overall time complexity of constructing the lucky number is $O(\log k)$.

The actual number of loop iterations is related to the bit length of the input k. Therefore, the time complexity is determined by

Space Complexity

The extra space used in the solution is allocated for the list ans to construct the return string. In the worst case, the length of ans will be equal to n, the number of digits in the k-th lucky number. Since the value of n is dependent on the logarithm of k, the space complexity is O(log k) for storing the resulting string.