402. Remove K Digits

<u>Greedy</u>

String]

Monotonic Stack

Problem Description

This LeetCode problem asks you to find the smallest possible integer after removing exactly k digits from a string num that represents a non-negative integer. The goal is to reduce the size of the number while keeping the remaining digits in the same order as they were in the original number.

Intuition

Medium

should ensure that the higher place values (like tens, hundreds etc.) have the smallest possible digits. Therefore, while parsing the string from left to right, if we encounter a digit that is larger than the digit following it, we remove the larger digit (which is at a higher place value). This decision is greedy because it makes the best choice at each step, aiming to keep the smallest digits at higher place values. To efficiently perform removals and keep track of the digits, a stack is an excellent choice. Each time we add a new digit to the

The intuition behind the solution is to use a greedy algorithm. If we want the resulting number to be the smallest possible, we

stack, we compare it to the element on top of the stack (which represents the previous digit in the number). If the new digit is smaller, it means we can make the number smaller by popping the larger digit off the stack. This process is repeated up to k times as required by the problem statement. The stack represents the digits of the resulting number with the smallest digits at the bottom (higher place values). When k

removals are done, or the string is fully parsed, we take the bottom n - k digits from the stack (where n is the length of num),

since k digits have been removed, and that forms our result. Leading zeroes are removed as they do not affect the value of the number. If all digits are removed, we must return '0', which is the smallest non-negative integer. **Solution Approach**

We create an empty list called stk, which we will use as a stack to keep track of the valid digits of the smallest number we

are constructing.

We need to retain len(num) - k digits to form the new number after we have removed k digits. The variable remain holds

The implementation of this algorithm is straightforward once you understand the intuition:

- this value.
- While we still have more digits k to remove, and the stack stk is not empty, and the digit at the top of the stack stk[-1] is greater than the current digit c, we pop the top of the stack. This is because keeping c, which is smaller, will yield a smaller number.
- We also decrement k by 1 each time we pop a digit off the stack since that counts as one removal.

any leading zeros with .lstrip('0').

remove the top element of the stack.

Here's the step-by-step process:

Iterate over each digit in "1432219":

Add '2' to the stack. stk = ['1', '2'].

We iterate over each character c in the string num:

- After the check (and potential removal), we append the current digit c to the stack. This digit is now part of the new number. After we finish iterating over num, the stack contains the digits of the resulting number, but it might have more digits than necessary if we didn't need to remove k digits. Thus, we slice the stack up to remain digits.
- Next, we need to convert the list of digits into a string. We join the digits in stk up to the remain index and then we remove
- because we must return a valid number and 0 is the smallest non-negative integer. In any other case, we return the joined string of digits that now represents the smallest possible integer after the removal of k digits.

The last step is to handle the case where all digits are removed, resulting in an empty string. If that happens, we return '0'

ideal choice to store the digits of the new number because it allows for easy removal of the last added digit when a smaller digit comes next. The process is greedy and makes local optimum choices by preferring smaller digits in the higher place values. Remember, in Python, a list can act as a stack with the append method to push elements onto the stack and the pop method to

This algorithm makes use of a stack, which is a classic data structure that operates on a Last In, First Out (LIFO) principle. It's an

Let's consider a small example to illustrate the solution approach. Suppose the input string num is "1432219" and k is 3. We want to remove 3 digits to make the number as small as possible.

Initialize an empty list stk to represent the stack. The number of digits we want to remain in the final number is remain =

len(num) - k = 7 - 3 = 4.

Example Walkthrough

 Next is '4'. '4' is greater than '1', so we keep it and push '4' to stk. \circ Then comes '3'. '3' is smaller than '4' and k > 0, so we pop '4' out of the stack. Now stk = ['1'] and k = 2. \circ Now we have '2'. '2' is smaller than '3', so we pop '3'. Now, stk = ['1'] and k = 1.

- ∘ Finally, '1' is smaller than '2', so we pop the last '2' from stk. stk = ['1', '2'], and k = 0 (no more removals allowed). Since we've already removed 3 digits, just push '1' and then '9' to the stack. Now, stk = ['1', '2', '1', '9'].

Start with the first digit '1'. Since the stack is empty, we add '1' to stk.

- We've finished processing each digit and our stack stk represents the smallest number we could make. However, we need to make sure we have the right number of digits, which should be remain = 4. Since stk already contains 4 digits, there's no
- need to slice.

We return '1219', which is the smallest number possible after removing 3 digits from "1432219".

Another '2' comes, which is the same as the last digit, so we push '2' to stk. stk = ['1', '2', '2'].

This example illustrates how the stack helps efficiently manage the digits of the new number, ensuring that smaller digits remain at the higher place values whenever possible.

Join the stack to form a number and strip leading zeros (if any). result = ''.join(stk).lstrip('0'). In this case, '1219'.

- **Python**
- # Initialize a stack to keep track of the digits stack = []

Remove the last digit from the stack as it's greater than the current one stack.pop() # Decrease the count of digits we can remove k -= 1

for digit in num:

stack.append(digit)

Solution Implementation

def removeKdigits(self, num: str, k: int) -> str:

remaining_digits_count = len(num) - k

Number of digits to remain in the final number

Iterate over each character in the input string

while k and stack and stack [-1] > digit:

Add the current digit to the stack

return result.isEmpty() ? "0" : result;

string removeKdigits(string num, int k) {

for (char& digit : num) {

while (k > 0) {

// Iterate through each character in the input number

// and whether we have still digits to remove

while (k > 0 && !stack.empty() && stack.back() > digit) {

--k; // Decrement the count of digits to remove

// This is necessary when the sequence was initially increasing

stack.pop back(); // Remove the last digit from 'stack'

--k; // Decrement the count of digits to remove

int startIndex = 0; // Index to keep track of leading zeros

while (startIndex < stack.size() && stack[startIndex] == '0') {</pre>

// Remove leading zeros from the 'stack'

stack += digit; // Add the current digit to 'stack'

final_number = ''.join(stack[:remaining_digits_count])

While we can still remove digits, and the stack is not empty,

and the current digit is smaller than the last digit in the stack:

Build the final number string from the stack up to the remaining digits

// If the resulting string is empty, return "0" instead; otherwise, return the string.

// Function to remove 'k' digits from the string 'num' to get the smallest possible number.

// Check if the current digit is smaller than the last digit in 'stack'

// Further remove digits from the end if we haven't removed enough 'k' digits

string stack; // Using 'stack' to store the characters representing the smallest number

stack.pop back(); // Remove the last digit from 'stack' to maintain the smallest number

class Solution:

```
# Strip leading zeros from the final number and return it, or return '0' if empty
        return final_number.lstrip('0') or '0'
Java
class Solution {
    public String removeKdigits(String num, int k) {
        // Create a StringBuilder to use as a stack to keep track of digits.
        StringBuilder stack = new StringBuilder();
        // Iterate through each character in the input string.
        for (char digit : num.toCharArray()) {
            // While the current digit is smaller than the last digit in the stack
            // and we still have digits to remove (k > 0), remove the last digit.
            while (k > 0 \&\& stack.length() > 0 \&\& stack.charAt(stack.length() - 1) > digit) {
                stack.deleteCharAt(stack.length() - 1);
                k--;
            // Append the current digit to the stack (StringBuilder).
            stack.append(digit);
        // If after the iteration we still need to remove more digits, remove from the end.
        while (k > 0) {
            stack.deleteCharAt(stack.length() - 1);
            k--;
        // Remove leading zeros by finding the index of the first non-zero digit.
        int nonZeroIndex = 0;
        while (nonZeroIndex < stack.length() && stack.charAt(nonZeroIndex) == '0') {</pre>
            nonZeroIndex++;
        // Create a new string starting from the first non-zero digit.
        String result = stack.substring(nonZeroIndex);
```

C++

public:

class Solution {

```
++startIndex; // Increment index to skip the leading zero
        string result = stack.substr(startIndex); // Extract the non-zero starting substring as result
        return result.empty() ? "0" : result; // If result is empty, return "0"; otherwise, return the result
};
TypeScript
function removeKdigits(numString: string, k: number): string {
    // Convert the string to an array of characters for easier manipulation
    let digitArray = [...numString];
    // Keep removing digits until we have removed k digits
    while (k > 0) {
        let indexToDelete = 0; // Initialize deletion index
        // Find where the digit is greater than the one following it; that's our deletion target
        while (indexToDelete < digitArray.length - 1 && digitArray[indexToDelete + 1] >= digitArray[indexToDelete]) {
            indexToDelete++;
        // Remove the digit at the identified deletion index
        digitArray.splice(indexToDelete, 1);
        // Decrement the count of digits we still need to remove
        k--;
    // Join the array back into a string and strip leading zeroes, if any
    let result = digitArray.join('').replace(/^0*/g, '');
    // If the result is an empty string, return '0', otherwise return the processed number string
    return result || '0';
class Solution:
    def removeKdigits(self, num: str, k: int) -> str:
        # Initialize a stack to keep track of the digits
```

Time and Space Complexity

stack = []

for digit in num:

stack.pop()

stack.append(digit)

k -= 1

Number of digits to remain in the final number

Iterate over each character in the input string

while k and stack and stack [-1] > digit:

Add the current digit to the stack

return final_number.lstrip('0') or '0'

While we can still remove digits, and the stack is not empty,

Decrease the count of digits we can remove

final_number = ''.join(stack[:remaining_digits_count])

and the current digit is smaller than the last digit in the stack:

Build the final number string from the stack up to the remaining digits

Strip leading zeros from the final number and return it, or return '0' if empty

Remove the last digit from the stack as it's greater than the current one

remaining_digits_count = len(num) - k

The time complexity of the given code can be analyzed based on the operations performed. The code iterates over each character in the string num which has a length of n. In the worst case, each character may be pushed to and popped from the

Time Complexity

stack stk once. Pushing and popping from the stack are 0(1) operations, but since the inner while loop could run up to k times for each character, it might appear at first as if the complexity is O(nk). However, each element is pushed and popped at most once, resulting in a time complexity of O(n) overall because the while loop can't execute more than n times over the course of the entire function. Therefore, the total time complexity of the algorithm is: 0(n)

Space Complexity

The space complexity is determined by the space used by the stack stk, which in the worst case could contain all characters if k is zero or if all characters are in increasing order. Therefore, the space complexity is proportional to the length of the input string num. Thus, the space complexity of the algorithm is:

0(n)