2740. Find the Value of the Partition

Sorting **Medium** Array

Problem Description

In this problem, we are presented with an array of positive integers named nums. Our task is to divide this array into two nonempty sub-arrays, nums1 and nums2, in such a way that the difference between the maximum value in nums1 and the minimum value in nums2 is as small as possible. This difference is referred to as the "value of the partition," expressed mathematically as |max(nums1) - min(nums2)|. We need to find and return the smallest possible value of this partition.

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minimum of another is achieved when the elements in both subsets are as close to each other as possible. Since nums contains only positive integers, the optimal partition would place consecutive elements from the sorted order of nums into nums1 and nums2. Sorting nums ensures that any two adjacent numbers have the smallest possible difference. After sorting, we can iterate through nums, taking one element to nums1 and the next to nums2, while keeping track of the minimum value of |nums[i+1] - nums[i] | for all i. This minimum value will occur between two consecutive elements in the sorted array, forming the optimal boundary between nums1 and nums2. Using the pairwise function from Python, we can easily iterate through adjacent pairs of elements to find this boundary, resulting in the minimum partition value.

The intuition behind the solution is based on the fact that the smallest difference between the maximum of one subset and the

The solution to this problem uses sorting and a simple iteration to determine the minimum partition value. The algorithm involves

the following steps:

Solution Approach

Sorting: The first action is to sort the input array nums. Sorting is a fundamental step because it allows us to consider elements in their natural order to find the smallest difference between adjacent elements. Python's built-in .sort() method

sorts the array in place in ascending order. The time complexity of the sorting operation is O(N log N), where N is the number of elements in nums. Iterating with Adjacent Pairs: Once the array is sorted, the next step is to iterate through the array and consider each pair of adjacent elements. This is where the minimum possible partition value will occur. To do this, we can use the pairwise

function, which is a handy tool for iterating over a list in overlapping pairs. This function simplifies the process of comparing

- Calculating the Minimum Partition Value: For each pair of adjacent elements (a, b) generated by pairwise(nums), we calculate the difference b - a. The partition value is the absolute difference between the maximum of one partition and the minimum of the second partition, so by considering differences between consecutive elements, we are directly calculating
- potential partition values. The minimum partition value is then found by taking the minimum of all these differences using the min function. **Return the Result**: Finally, the function returns the minimum partition value that was calculated. No additional data structures are required other than the space needed for sorting. The pairwise function generates a tuple for

each pair of elements, which uses only a constant amount of additional space. Overall, the algorithm is efficient and

straightforward, with the sorting step dominating the overall time complexity.

each element with its subsequent neighbor without the need for complex index management.

return min(b - a for a, b in pairwise(nums)) In this code, pairwise(nums) is an iterable that gives us each adjacent pair (a, b) from the sorted nums, and the generator expression min(b - a for a, b in pairwise(nums)) calculates the minimum difference between any two consecutive elements

Example Walkthrough Let's go through a simple example to illustrate the solution approach described above. Consider the following array of numbers:

Following the provided solution approach, here is how we would find the smallest possible value of the partition:

After sorting: nums = [1, 2, 4, 5]

Adjacent pairs: (1, 2), (2, 4), (4, 5)

 \circ For the pair (2, 4), the difference is 4 - 2 = 2.

 \circ For the pair (4, 5), the difference is 5 - 4 = 1.

Sort the input list in ascending order

The 'pairwise' utility is not available in Python standard library

// Update the minValueOfPartition with the smallest difference found

minValueOfPartition = Math.min(minValueOfPartition, nums[i] - nums[i - 1]);

// between adjacent elements in the sorted array.

// Return the minimum value of the partition found.

return minValueOfPartition;

// Return the minimum difference found.

function findValueOfPartition(nums: number[]): number {

let minDifference: number = Number.MAX_SAFE_INTEGER;

// Iterate over the sorted array to find the smallest difference

// Update minDifference with the smallest difference found so far.

minDifference = Math.min(minDifference, nums[i] - nums[i - 1]);

// Sort the array in non-decreasing order.

return minDifference;

nums.sort((a, b) => a - b);

// between consecutive elements.

for (let i = 1; i < nums.length; i++) {</pre>

nums = [4, 2, 5, 1]

Here's how the solution might look in Python:

def findValueOfPartition(self, nums: List[int]) -> int:

in the array, which corresponds to the minimum partition value.

from itertools import pairwise

nums.sort()

class Solution:

Sorting: We start by sorting the array. Before sorting: nums = [4, 2, 5, 1]

Sorting ensures that we consider the elements in increasing order to find the smallest difference between adjacent elements.

 \circ For the pair (1, 2), the difference is 2 - 1 = 1.

Calculating the Minimum Partition Value: We calculate the difference between each pair of adjacent numbers:

Iterating with Adjacent Pairs: We use the pairwise function to iterate through the sorted array in adjacent pairs:

We are looking for the minimum of these differences. **Return the Result**: The smallest difference from our pairs is 1 (which occurs between the pairs (1, 2) and (4, 5)).

Therefore, the smallest possible value of the partition for the input nums = [4, 2, 5, 1] is 1. When we apply the provided Python code to this array, the function findValueOfPartition returns 1 as the result of the calculation. Solution Implementation

from itertools import tee class Solution: def find_value_of_partition(self, nums: List[int]) -> int:

```
# until Python 3.10. A custom implementation is needed for earlier versions.
# Here is a custom implementation of pairwise utility using tee and zip.
def pairwise(iterable):
    "s \rightarrow (s0,s1), (s1,s2), (s2, s3), ..."
```

a, b = tee(iterable)

next(b, None)

return zip(a, b)

nums.sort()

Python

```
# Calculate the minimum difference between consecutive elements in the sorted list
        # This difference represents the smallest partition value
        return min(b - a for a, b in pairwise(nums))
Please note that you need to import the `List` typing from the `typing` module to use list type hints. Here's how you include it:
```python
from typing import List
Java
class Solution {
 public int findValueOfPartition(int[] nums) {
 // Sort the array to bring similar values closer.
 Arrays.sort(nums);
 // Initialize the minimum difference to a large value.
 // Instead of 1 << 30, Integer.MAX_VALUE is used for readability.
 int minValueOfPartition = Integer.MAX_VALUE;
 // Loop through the sorted array starting from the second element
 for (int i = 1; i < nums.length; ++i) {</pre>
```

```
C++
#include <vector> // Necessary for using std::vector
#include <algorithm> // Necessary for using std::sort and std::min
class Solution {
public:
 // Function to find the minimum difference between any two elements after sorting the array.
 int findValueOfPartition(std::vector<int>& nums) {
 // Sort the array in non-decreasing order.
 std::sort(nums.begin(), nums.end());
 // Initialize the answer with a large value.
 // INT_MAX from limits.h could also be used for maximum allowable integer.
 int minDifference = INT_MAX;
 // Iterate over the sorted array to find the smallest difference
 // between consecutive elements.
 for (int i = 1; i < nums.size(); ++i) {</pre>
 // Update minDifference with the smallest difference found so far.
 minDifference = std::min(minDifference, nums[i] - nums[i - 1]);
```

#### // Return the minimum difference found. return minDifference;

from itertools import tee

**TypeScript** 

```
class Solution:
 def find_value_of_partition(self, nums: List[int]) -> int:
 # Sort the input list in ascending order
 nums.sort()
 # The 'pairwise' utility is not available in Python standard library
 # until Python 3.10. A custom implementation is needed for earlier versions.
 # Here is a custom implementation of pairwise utility using tee and zip.
 def pairwise(iterable):
 "s \rightarrow (s0,s1), (s1,s2), (s2, s3), ..."
 a, b = tee(iterable)
 next(b, None)
 return zip(a, b)
 # Calculate the minimum difference between consecutive elements in the sorted list
 # This difference represents the smallest partition value
 return min(b - a for a, b in pairwise(nums))
Please note that you need to import the `List` typing from the `typing` module to use list type hints. Here's how you include it:
```python
from typing import List
Time and Space Complexity
Time Complexity
```

// Function to find the minimum difference between any two elements after sorting the array.

// Initialize the answer with a large value. TypeScript's maximum safe integer can be used.

The given code has two main operations that dictate the time complexity: sorting the list nums and then computing the minimum difference between consecutive elements.

The sorting of a list of n elements has a time complexity of O(n log n) using the Timsort algorithm, which is the default sorting algorithm in Python.

complexity is: 0(n log n)

The pairwise function generates tuples containing pairs of consecutive elements in the sorted list. Iterating through the nums list

to find the minimum difference is a linear operation with a time complexity of 0(n - 1), which simplifies to 0(n). Combining both, the overall time complexity is $0(n \log n) + 0(n)$. As $0(n \log n)$ is the dominating factor, the final time

Space Complexity

pairwise iterator. The sort operation can be done in-place, so it does not significantly add to the space complexity, thus being 0(1).

The space complexity of the code is determined by the additional space required for sorting and the space required for the

The pairwise function, however, creates a new iterator that generates pairs of consecutive elements without generating all pairs at once. Therefore, the space complexity due to pairwise is 0(1).

Combining both the space complexities from sorting and pairwise, the overall space complexity is: 0(1)