

The problem focuses on finding the largest variance within any substring of a given string s, which contains only lowercase English letters. The variance of a string is defined as the largest difference in the number of occurrences between any two distinct characters in it. It's important to consider that the characters could be the same or different.

To clarify, let's say we're given a string like "abcde". In this string, each character appears exactly once, so the variance between any two characters is zero. However, if we had a string like "aabbcc", the variance could be higher. If we choose the substring "aab", the variance would be 2 because there are two 'a' characters and none of the 'b' character (or vice versa). We need to consider all possible substrings to find the one with the largest variance.

The intuition behind the solution lies in realizing that the variance could only be maximized if we consider pairs of different characters in the string, as those pairs are the ones that can actually have a non-zero variance. So the first step is to iterate over all

Intuition

possible pairs of characters (permutations of two distinct letters from the 26 lowercase English characters). Once a pair of characters is selected (let's say 'a' and 'b'), we apply a two-pass scanning algorithm on the string to track the difference in the occurrence counts between 'a' and 'b'. We keep two counters: one (f[0]) resets every time we encounter an 'a' after a 'b', while the other (f[1]) continues to accumulate or starts over from f[0] - 1 when a 'b' is encountered after at least one 'a'. We

look to maximize the difference (f[1]), which is essentially the variance, while scanning. In the end, we take the maximum recorded value during all scans as the answer. This approach works because by focusing on pairs of characters at a time, we are able to isolate the potential substrings that can provide the maximum variance, rather than getting lost in the overwhelming number of all possible substrings.

Along with tracking the occurrences of 'a' and 'b', there's also an understanding that we include a subtraction operation when we encounter a 'b', making sure we are actually counting a valid substring with at least one 'a' when 'b' appears, since variance requires both characters to be present. A -inf is used to handle cases where a valid substring does not exist yet. By considering every single position in the string as a potential starting or ending point for substrings with maximum variance between two specific characters,

the algorithm guarantees that no possible substring is overlooked, thus ensuring the correctness of the solution.

Solution Approach The implementation employs a brute force strategy enhanced with a smart iteration mechanism and memory optimization to calculate the maximum variance for all substrings of the given string. The key elements of this approach involve permutations, two

Firstly, the algorithm generates all possible permutations of two distinct characters from the set of lowercase English letters. This is done using permutations (ascii_lowercase, 2) which loops through all pairs of two distinct characters since we're looking for the

Permutations of Characters

Tracking Variance with Two Counters In the solution, an array f with two elements is used:

This dichotomy allows the algorithm to keep track of two scenarios simultaneously:

Maximal Variance Calculation

variance achievable among all substrings of s.

largest variance within any substring of s.

Initialize the counters f[0] and f[1] to 0.

Start with 'a': f[0] = 1, f[1] = 1.

2. Loop through each character c in s = "abbab":

Next, 'a': f[0] = 1, f[1] = max(f[1] + 1, f[0]) = 1.

comparison variables, and a linear scan through the string.

difference in occurrences between two potentially different characters.

f[1] keeps track of the accumulated variance between a and b.

Accumulating occurrences of a when no b has interfered (handled by f[0]).

Both f[0] and f[1] are reset to zero if a b appears directly after a b. However, when an a is followed by a b, f[1] takes the role of capturing the current variance by taking either the current accumulated variance minus one, or the variance just after the last a

Capturing the 'variance' until a reset is needed due to a b occurrence (handled by f[1]).

As we scan the string s, for each character c, the algorithm performs the following actions:

f[0] serves as a resettable counter for occurrences of character a since the last occurrence of b.

minus one, determined by max(f[1] - 1, f[0] - 1). This handling ensures that at least one a is present before a b is counted, which is necessary for a valid substring.

• If c is a, we increment both f[0] and f[1]. If c is b, we decrement f[1] and choose the larger of the two: f[1] or f[0] - 1, effectively either continuing with the current

After each update, we check if f[1] exceeds the current ans. If it does, we update ans to the value of f[1]. The algorithm avoids

the string ensures that we consider every potential substring for each character pair. Eventually, ans will contain the maximum

substring or starting a new substring that excludes the last a. f [0] is reset to zero as we can't start a new substring with just b.

false starts with -inf, ensuring that f[1] only records valid substrings. The loop through permutations of characters ensures that we consider every potential character disparity. The nested loop through

0(26*26*n) where n is the length of the string, due to the nested loops—26 for each character in the permutations, and n for the iterations over the string. Example Walkthrough

Let's illustrate the solution approach using a small example. Suppose we are given the string s = "abbab". Our task is to find the

The space complexity of the solution is 0(1) as we only maintain a few variables and a pair of counters, and the time complexity is

We initiate a variable ans to store the maximum variance found. Initially, ans = -inf. Pair ('a', 'b'):

consider only a few pairs here, such as ('a', 'b') and ('b', 'a'), since these are the only characters in our example string.

Firstly, we generate all possible permutations of two distinct characters from the set of lowercase English letters. For simplicity, let's

○ For c = 'b': Decrement f[1] and then set f[1] to the maximum of f[1] and f[0] - 1. Reset f[0] to 0.

the pair ('a', 'b').

Let's go through the string "abbab":

For c = 'a': Increment both f[0] and f[1] (since we initially consider the occurrence of 'a' without 'b' involved).

 Next, 'b': f[1] = max(f[1] - 1, f[0] - 1) = 0, f[0] = 0. Update ans = max(ans, f[1]) = 1. • Next, 'b': $f[1] = \max(f[1] - 1, f[0] - 1) = -1$, since there are no 'a' in front, we don't update ans.

Following the same steps but reversing the roles of 'a' and 'b', we would not find a variance larger than what we already have from

After considering these permutations, we find that the maximum variance is 1, which we may have obtained from the substring "ab"

Pair ('b', 'a'):

Last, 'b': f[1] = max(f[1] - 1, f[0] - 1) = 0. Update ans = max(ans, f[1]) = 1.

present, it is unnecessary.

1 from itertools import permutations

max_variance = 0

from string import ascii_lowercase

def largestVariance(self, s: str) -> int:

Iterate over each character in the string

if current_char == char_a:

freq_before_b += 1

freq_before_b = 0

max_freq_after_b += 1

elif current_char == char_b:

Initialize the answer to 0

for current_char in s:

public int largestVariance(String s) {

// Length of the input string.

continue;

// Variable to store the maximum variance found so far.

if (firstChar == secondChar) {

for (char firstChar = 'a'; firstChar <= 'z'; ++firstChar) {</pre>

int[] frequency = new int[] {0, -length};

if (s.charAt(i) == firstChar) {

for (int i = 0; i < length; ++i) {

frequency[0]++;

frequency[1]++;

frequency[0] = 0;

// Iterate over each character in the string.

} else if (s.charAt(i) == secondChar) {

// Update the maximum variance found.

maxVariance = Math.max(maxVariance, frequency[1]);

// Iterate over all possible pairs of characters (a and b are different).

// If both characters are the same, skip this iteration.

// f[0] is the streak of 'a's, f[1] is the max variance for the current window.

// If the current character is 'a', increase both frequencies.

// If the current character is 'b', calculate the variance.

// Reset the streak of 'a's because 'b' is encountered.

frequency[1] = Math.max(frequency[0] - 1, frequency[1] - 1);

// Initialize with 0 for streak and -n for variance because variance cannot be less than -n.

for (char secondChar = 'a'; secondChar <= 'z'; ++secondChar) {</pre>

// Array to keep track of the frequency of character 'a'

int length = s.length();

int maxVariance = 0;

or "ba" in the string "abbab".

Python Solution

from math import inf

class Solution:

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Java Solution

class Solution {

In a complete implementation, we would evaluate all permutations of character pairs, but in this small example with only 'a' and 'b' The process effectively scans through every substring of s and calculates the variance for pairs of characters, updating ans when a higher variance is found. The time complexity is mainly driven by the nested loops: there are 26*26 permutations of character pairs,

and for each pair, we scan the entire string s of length n, resulting in an 0(26*26*n) time complexity.

Increment the frequency counter of 'a' when 'a' is found

max_freq_after_b = max(max_freq_after_b - 1, freq_before_b - 1)

Iterate through all ordered pairs of distinct lowercase ascii characters 10 for char_a, char_b in permutations(ascii_lowercase, 2): 11 12 13 # Initialize the frequency counters for the current pair, where:

- # freq_before_b records the number of 'a' characters seen so far without any 'b' in between 14 # max_freq_after_b records the maximum frequency difference after at least one 'b' was found 16 freq_before_b = 0 17 max_freq_after_b = -inf 18
- 30 # Update the max_variance if we find a new maximum 31 if max_variance < max_freq_after_b:</pre> 32 max_variance = max_freq_after_b 33 34 # Return the largest variance found 35 return max_variance 36

Decrement the frequency counter of 'b' when 'b' is found; reset freq_before_b

36 37 38 39 // Return the largest variance found. return maxVariance; 40 41

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C++ Solution
  1 class Solution {
  2 public:
        // Method to find the largest variance of a substring for a given string
         int largestVariance(string s) {
             // Length of the input string
             int n = s.size();
             // Initialize the maximum variance to 0
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             int maxVariance = 0;
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             // Iterate over all possible character pairs (a, b) where a != b
             for (char a = 'a'; a <= 'z'; ++a) {
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                 for (char b = 'a'; b <= 'z'; ++b) {
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                     // Skip if both characters are the same
                     if (a == b) continue;
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                     // Initialize frequency counters for characters a and b
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                     // f[0] will track frequency of a, f[1] tracks max variance
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                     int frequencies [2] = \{0, -n\};
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                     // Iterate over each character in the string
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                     for (char currentChar : s) {
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                         // Increment frequency count if a matches the current character
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                         if (currentChar == a) {
                             frequencies[0]++;
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                             frequencies[1]++;
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                         // Decrement frequency count and update variance if b matches
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                         else if (currentChar == b) {
                             frequencies[1] = max(frequencies[1] - 1, frequencies[0] - 1);
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                             frequencies[0] = 0;
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                         // Update the maximum variance found so far
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                         maxVariance = max(maxVariance, frequencies[1]);
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             // Return the final maximum variance
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             return maxVariance;
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    };
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Typescript Solution
  1 // Function to find the largest variance of a substring for a given string
  2 function largestVariance(s: string): number {
         // Length of the input string
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26 27 frequencies[0] = 0; 28 29 // Update the maximum variance found so far maxVariance = Math.max(maxVariance, frequencies[1]); 30

return maxVariance;

Time and Space Complexity

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let n: number = s.length;

let maxVariance: number = 0;

// Initialize the maximum variance to 0

if (a === b) continue;

// Return the final maximum variance

22 frequencies[1]++; 23 24 // Decrement frequency count and update variance if 'b' matches 25 else if (currentChar.charCodeAt(0) === b) { frequencies[1] = Math.max(frequencies[1] - 1, frequencies[0] - 1);

if (currentChar.charCodeAt(0) === a) {

// Iterate over each character in the string

// Iterate over all possible character pairs (a, b) where a != b

for (let b = 'a'.charCodeAt(0); b <= 'z'.charCodeAt(0); ++b) {</pre>

// Initialize frequency counters for characters 'a' and 'b'

// Increment frequency count if 'a' matches the current character

for (let a = 'a'.charCodeAt(0); a <= 'z'.charCodeAt(0); ++a) {</pre>

// Skip if both characters are the same

let frequencies: number[] = [0, -n];

for (let currentChar of s) {

frequencies[0]++;

- Time Complexity The time complexity of the code can be analyzed based on the operations it performs:
- English alphabet. Since there are 26 letters in the English alphabet, there are 26 * 25 such permutations because after choosing the first letter, there are 25 remaining letters to form a pair.

Thus, the space complexity of the algorithm is 0(1) as it only uses a constant amount of space.

2. The function then iterates through each character in the string s for every pair of letters (a, b). The length of the string s is n. 3. Within the nested loop, there are constant-time operations (comparisons, assignments, max, addition, and subtraction).

1. The function permutations (ascii_lowercase, 2) generates all possible permutations of two distinct lowercase letters from the

- Given the above steps, the time complexity is 0(26 * 25 * n), which simplifies to 0(n) because the number of permutations (26 * 25) is a constant.
- Space Complexity The space complexity of the code is determined by:
 - 1. Variables ans, f[0], and f[1], which are used to keep track of the current and maximum variance. 2. No additional data structures that grow with input size are used.