2903. Find Indices With Index and Value Difference I

Easy <u>Array</u>

In this problem, you're given an integer array called nums, with a length of n. You're also given two integers: indexDifference

Problem Description

they meet the following criteria: 1. The absolute difference between i and j must be at least indexDifference, 2. The absolute difference between the values at nums[i] and nums[j] must be at least valueDifference.

and valueDifference. Your task is to find two indices i and j such that both i and j are within the range from 0 to n - 1 and

In terms of the outcome, you need to return an array called answer. This array should consist of the indices [i, j] if such a pair

of indices exists. If there are multiple valid pairs, you can return any one of them. If no valid pairs are found, then return [-1, -1].

An interesting point to note is that according to the problem statement, i and j can be the same index, which implies that indexDifference could potentially be 0.

The primary intuition behind the solution is the usage of a sliding window technique, combined with the maintenance of the

minimum and maximum values within the window. The sliding window is defined by two pointers, i and j, that maintain a

distance apart specified by indexDifference. The pointers are used as markers to capture a subarray within nums to check against our two conditions.

At the outset, i starts at the position indexDifference, and j starts at 0. By doing this, the gap between i and j reflects the indexDifference requirement of our problem. We maintain two variables, mi and mx, to keep track of the indices where the minimum and maximum values are found within our sliding window that ends at the current j index. While sliding i further along the array, we update mi and mx to account for the

entry of new values into the window and the exit of old values. When updating mi and mx, if nums[j] is less than nums[mi], we reassign mi to j, because we have found a new minimum.

Conversely, if nums[j] is greater than nums[mx], we reassign mx to j due to identifying a new maximum. After every movement of i and update of mi and mx, we check our two conditions against nums[i] (the current value at i). If

the difference between nums[i] and the value at nums[mi] is greater than or equal to valueDifference, we have found a valid

pair [mi, i]. Alternatively, if the difference between the maximum value (nums [mx]) and nums [i] is also greater than or equal to

If we reach the end of the array without finding a pair that satisfies both conditions, we conclude that no such pair exists, and we return the default output [-1, -1].

valueDifference, then [mx, i] is a valid pair. In this situation, we immediately return the pair as it meets our requirements.

sections one at a time. This allows for checking the conditions over smaller segments in a single pass through the array, making the solution more efficient than a brute force approach that would involve examining all possible index pairs.

The solution uses a sliding window approach, which involves moving a fixed-size window across the array to examine sub-

To implement this technique, the algorithm maintains two pointers: i and j. These pointers define the bounds of the sliding window. The pointer i starts at the index equal to indexDifference while j starts at 0, thus immediately satisfying the condition abs(i - j) >= indexDifference because i - j is initialized to indexDifference.

As the algorithm iterates over the array, starting from i = indexDifference, it keeps track of the indices of the minimum and the maximum values found so far to the left of j. These indices are stored in variables mi and mx, respectively.

mx = j

array.

Example Walkthrough

indexDifference.

Here's how we proceed step by step:

continue the iteration, we would:

from typing import List

min_idx = max_idx = 0

Initialize min and max index pointers

for current_idx in range(idx_diff, len(nums)):

if nums[compare idx] > nums[max_idx]:

return [min_idx, current_idx]

return [max_idx, current_idx]

if nums[current idx] - nums[min idx] >= val diff:

if nums[max idx] - nums[current idx] >= val_diff:

int minIndex = 0: // Initializing the minimum value index

int maxIndex = 0; // Initializing the maximum value index

if (nums[currentIndex] < nums[minIndex]) {</pre>

minIndex = currentIndex;

if (nums[i] > nums[maxIndex]) {

// nums: The array of numbers to search within

if (nums[i] - nums[minIndex] >= valueDiff) {

if (nums[maxIndex] - nums[i] >= valueDiff) {

return $\{-1, -1\}$; // If no pair found, return $\{-1, -1\}$

// indexDifference: The maximum allowed difference between the indices

maxIndex = j; // Update maxIndex if a new maximum is found

// This function finds two indices such that the difference of the elements at these indices

// Returns an array with two numbers representing the indices, or [-1, -1] if no such pair exists

function findIndices(nums: number[], indexDifference: number, valueDifference: number): number[] {

for (let currentIndex = indexDifference; currentIndex < nums.length; currentIndex++) {</pre>

// Initialize the indices for the minimum value (minIndex) and maximum value (maxIndex) found.

// is at least the given valueDifference and the indices are separated by at most the given

// valueDifference: The minimum required difference between the values at the indices

// Iterate over the array, starting from the element at the indexDifference.

// Calculate the index of the element we are comparing against,

return {minIndex, i}; // Pair found, return indices

return {maxIndex, i}; // Pair found, return indices

max_idx = compare_idx

class Solution:

return [mi, i]

return [mx, i]

for i in range(indexDifference, len(nums)):

if nums[i] - nums[mi] >= valueDifference:

if nums[mx] - nums[i] >= valueDifference:

j = i - indexDifference

Solution Approach

Within the loop, we first check whether the current element at index j changes the minimum or maximum: if nums[j] < nums[mi]:</pre> mi = iif nums[i] > nums[mx]:

After updating mi and mx, we check if nums[i] differs enough from the minimum or maximum value to satisfy the valueDifference condition:

```
thus it returns [-1, -1].
In terms of data structures, no additional structures are needed beyond the use of a few variables to keep track of the indices
and values encountered. This algorithm is space-efficient because it operates directly on the input array without requiring extra
space proportional to the input size.
The choice of a sliding window and keeping track of minimum and maximum values eliminates the need to compare every
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element with every other element, thereby reducing the time complexity from O(n^2) to O(n), where n is the length of the input

If either of these checks succeeds, the function immediately returns the corresponding pair of indices, as they meet both

prescribed conditions. If the function reaches the end of the array without returning, this means no valid pairs were found, and

4, 5], with indexDifference = 3 and valueDifference = 3. Our task is to find indices i and j such that abs(i - j) >= indexDifference and abs(nums[i] - nums[j]) >=valueDifference.

According to the given solution approach, we initialize the sliding window by setting i to indexDifference and j to 0. This

immediately satisfies the first condition as the difference between the indices i = 3 and j = 0 is 3, which is equal to

Let's walk through an example to illustrate the solution approach described above. Consider the integer array nums = [1, 2, 3,

1. On the first iteration where i = 3: \circ We have j = 0• We initiate mi to j since there are no previous values to compare with, and similarly, mx is also initiated to j ○ The elements under consideration are [nums[0], nums[3]] i.e., [1, 4]

In this example, we directly found a pair that met both conditions, and thus we would return [0, 3]. However, if we needed to

2. Increment i to the next position and decrement j to keep the window size constant while satisfying the indexDifference. If nums[j] changes

• Since nums[3] - nums[0] fulfills the valueDifference condition, we return [0, 3] as the indices that satisfy both conditions.

Solution Implementation **Python**

We then maintain mi and mx to keep track of the minimum and maximum values within the window.

 \circ We see that nums[3] - nums[0] = 4 - 1 = 3, which is equal to valueDifference

the minimum or maximum within the new window, update mi or mx accordingly.

3. Check if nums[i] differs enough from nums[mi] or nums[mx] as explained previously.

4. If we find a pair, we return it. If not, we continue iterating until i reaches the end of the array.

If no valid pairs are found by the end of the array, we return [-1, -1] as specified.

def findIndices(self, nums: List[int], idx diff: int, val_diff: int) -> List[int]:

Compute the comparison index that matches the index difference

public int[] findIndices(int[] nums, int indexDifference, int valueDifference) {

// Update the maximum value index if a new maximum is found

compare_idx = current_idx - idx_diff # Check and update the min and max indices based on the values at compare_idx if nums[compare idx] < nums[min_idx]:</pre> min idx = compare idx

If the value difference requirement is met with the minimum, return the indices

If the value difference requirement is met with the maximum, return the indices

If the required value difference isn't found, return [-1, -1] as per problem statement

Traverse the array, starting from the index that enables the required index difference

```
// Loop through the array starting from the index equal to the indexDifference to the end of the array
for (int i = indexDifference; i < nums.length; ++i) {</pre>
    int currentIndex = i - indexDifference; // Calculate the index to compare with
    // Update the minimum value index if a new minimum is found
```

return $\begin{bmatrix} -1, & -1 \end{bmatrix}$

Java

class Solution {

```
if (nums[currentIndex] > nums[maxIndex]) {
                maxIndex = currentIndex;
            // Check if the difference between the current value and the minimum value found so far is at least valueDifference
            if (nums[i] - nums[minIndex] >= valueDifference) {
                return new int[] {minIndex, i}; // Return the indices if condition is met
           // Check if the difference between the maximum value found so far and the current value is at least valueDifference
            if (nums[maxIndex] - nums[i] >= valueDifference) {
                return new int[] {maxIndex, i}; // Return the indices if condition is met
        // Return [-1, -1] if no such pair of indices is found
        return new int[] \{-1, -1\};
C++
#include <vector>
class Solution {
public:
    // Method to find the two indices in the array nums such that the difference
    // between their values is at least valueDifference and their index difference is exactly indexDifference
    // Args:
        nums: The input vector of integers
        indexDiff: The required difference between the indices of the two elements
        valueDiff: The minimum required value difference between the two elements
    // Returns:
    // A vector with two elements: the indices of the elements in nums that satisfy the above criteria
        If no such pair exists, returns \{-1, -1\}
    std::vector<int> findIndices(std::vector<int>& nums, int indexDiff, int valueDiff) {
        int minIndex = 0, maxIndex = 0; // Initialized to store the index of minimum and maximum values seen so far
        for (int i = indexDiff; i < nums.size(); ++i) {</pre>
            int j = i - indexDiff; // Calculate the corresponding index
            if (nums[i] < nums[minIndex]) {</pre>
                minIndex = j; // Update minIndex if a new minimum is found
```

// Check if the difference between the current value and the minimum value seen so far is at least valueDiff

// Check if the difference between the maximum value seen so far and the current value is at least valueDiff

// Method to find indices in an array such that the difference between their values is at least a given value and their positions

};

TypeScript

// indexDifference.

let minIndex = 0;

let maxIndex = 0;

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// which is indexDifference behind the current index.
        const compareIndex = currentIndex - indexDifference;
        // Update minIndex if the current compareIndex points to a new minimum value
        if (nums[compareIndex] < nums[minIndex]) {</pre>
            minIndex = compareIndex;
       // Update maxIndex if the current compareIndex points to a new maximum value
        if (nums[compareIndex] > nums[maxIndex]) {
           maxIndex = compareIndex;
        // Check if the difference between the current element and the minimum value is large enough.
        if (nums[currentIndex] - nums[minIndex] >= valueDifference) {
            return [minIndex, currentIndex]; // Return the indices if the condition is met.
       // Check if the difference between the maximum value and the current element is large enough.
        if (nums[maxIndex] - nums[currentIndex] >= valueDifference) {
            return [maxIndex, currentIndex]; // Return the indices if the condition is met.
   // If no suitable pair of indices is found, return [-1, -1].
   return [-1, -1];
from typing import List
class Solution:
   def findIndices(self, nums: List[int], idx diff: int, val_diff: int) -> List[int]:
       # Initialize min and max index pointers
        min_idx = max_idx = 0
       # Traverse the array, starting from the index that enables the required index difference
       for current_idx in range(idx_diff, len(nums)):
           # Compute the comparison index that matches the index difference
           compare_idx = current_idx - idx_diff
           # Check and update the min and max indices based on the values at compare_idx
           if nums[compare idx] < nums[min_idx]:</pre>
               min idx = compare idx
            if nums[compare idx] > nums[max_idx]:
                max_idx = compare_idx
```

return [max_idx, current_idx] # If the required value difference isn't found, return [-1, -1] as per problem statement return [-1, -1]

Time and Space Complexity

Time Complexity The time complexity of the provided code is O(n). This is achieved by iterating over the array once from indexDifference to the length of the array len(nums). Only constant time checks and updates are performed within the loop, leading to a linear time complexity relative to the array's size.

If the value difference requirement is met with the minimum, return the indices

If the value difference requirement is met with the maximum, return the indices

if nums[current idx] - nums[min idx] >= val_diff:

if nums[max idx] - nums[current idx] >= val_diff:

return [min_idx, current_idx]

Space Complexity The space complexity of the code is 0(1). No additional space proportional to the input size is used. Only a fixed number of

variables mi, mx, and j are used, which occupy constant space regardless of the input array size.