2280. Minimum Lines to Represent a Line Chart

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Medium
           Geometry
                      <u>Array</u>
                               <u>Math</u>
                                         Number Theory
                                                          Sorting
```

We have a 2D integer array named stockPrices, where each element consists of two integers representing a day and the corresponding stock price on that day. The task is to determine the minimum number of straight lines required to connect all the

points plotted on an XY plane, with the day as the X-axis and the price as the Y-axis, in the given order after sorting the data by the day. Essentially, we want to find the fewest lines that connect each adjacent point (each day to the next) to form a line chart. Intuition

**Problem Description** 

To solve this problem, the solution leverages the concept of slopes of lines. When connecting points to draw the line chart, we can continue using the same line as long as consecutive segments share the same slope. When the slopes differ, we need a new line. To find the slope of the line segment between two points (x, y) and (x1, y1), we calculate (y1 - y) / (x1 - x), which gives us the

rate of change in price with respect to days. Now, if we compare this slope with the next segment's slope between the points (x1,

y1) and (x2, y2), we would normally compute (y2 - y1) / (x2 - x1) and compare. To avoid division and possible precision issues with floating-point numbers, we instead compare cross-products: (y1 - y) \* (x2 - x1) and (y2 - y1) \* (x1 - x). If these two products are not equal, the slope has changed, and we need an additional line. The approach first sorts the stockPrices based on the day to ensure the points are in the correct order. Then, iterating through each pair of adjacent points, we calculate the aforementioned cross-product to decide whether the current segment can be

connected with the prior one or if a new line has to start. The variable ans keeps track of the number of lines needed. Remember, we increment the number of lines when a change in slope is detected between segments, and after calculating the

Solution Approach The implementation given in the reference solution approach follows these steps:

Sort the Input: We start by sorting the stockPrices array. This ensures that the stock prices are ordered by day, which is

slopes, we update the dx and dy to be the differences between the current and next points to be ready for the next iteration.

## Initialize Variables: Two variables dx and dy are initialized to store the differences in the x (day) and y (price) coordinates of

successive points. Additionally, ans is initialized to count the number of lines required.

essential as we need to connect points from one day to the next.

yields pairs of elements from the input list.

needed to represent the chart.

without redundant calculations or comparisons.

Suppose we have the following 2D array stockPrices:

- Iterate Through Points: We use a loop to go through each pair of adjacent points in the sorted array. To do this efficiently, the code leverages the pairwise utility, which is not a standard Python function but likely a user-defined or library function that
- Calculate and Compare Slopes: For each pair (x, y) and (x1, y1), we calculate the change in x as dx1 = x1 x and the change in y as dy1 = y1 - y. We then compare the products of cross-multiplying these with the previous segment's dx and dy
- Increment Line Count: When a slope change is detected, we increment ans by 1. **Update Differences**: Whether a new line is added or not, we update dx and dy with the differences dx1 and dy1, to use them in the next iteration for slope comparison.

**Return Result**: After all points have been processed, the variable ans is returned. This gives us the minimum number of lines

to check if the slopes match. If dy \* dx1 does not equal dx \* dy1, the slope has changed, indicating a need for a new line.

• Looping and Iteration: The use of a loop along with pairwise comparison is a common pattern when we need to check adjacent elements in a sequence.

Here is the explanation of the algorithm, data structure, or pattern used in each step:

• Sorting: It's the first step and is critical for comparing the slopes of segments in the correct order.

• ans is initialized to 1 because we need at least one line to connect the first point to the next one.

**Example Walkthrough** Let's walk through a small example to illustrate the solution approach.

• Slope Comparison: This step avoids division by comparing slopes using cross-multiplication, which is a mathematical technique to avoid

By systematically comparing each pair of points, the solution efficiently calculates the number of lines needed to form the chart

floating-point precision errors and ensure that comparisons are accurate when checking for collinearity between points.

stockPrices = [[1, 100], [2, 200], [3, 300], [4, 200], [5, 100], [6, 200]]

**Sort the Input:** Since stockPrices is already sorted by the day (the first element of each sub-array), we don't need to sort it again.

○ On the next iteration, we compare the point [2, 200] with [3, 300]. We calculate the changes dx1 = 3 - 2 and dy1 = 300 - 200, resulting

## ∘ Imagine comparing the first point [1, 100] with the second [2, 200]. This sets our initial dx and dy to [2–1, 200–100] = [1, 100].

**Initialize Variables:** 

**Iterate Through Points:** 

**Increment Line Count:** 

**Update Differences:** 

**Return Result:** 

Solution Implementation

from itertools import pairwise

stock\_prices.sort()

num\_lines = 0

return num\_lines

prev\_diff\_x, prev\_diff\_y = 0, 1

num\_lines += 1

# Initialize the count of lines needed to 0

# Iterate over each pair of adjacent stock prices

prev\_diff\_x, prev\_diff\_y = diff\_x, diff\_y

# Return the total number of lines needed

public int minimumLines(int[][] stockPrices) {

if prev\_diff\_y \* diff\_x != prev\_diff\_x \* diff\_y:

# Update the previous differences for the next iteration

Arrays.sort(stockPrices, (a, b) -> Integer.compare(a[0], b[0]));

// Iterate over the stock prices starting from the second price

int lineCount = 0; // Counter for the minimum lines needed

if (prevDeltaY \* deltaX != prevDeltaX \* deltaY) {

for (int i = 1; i < stockPrices.length; ++i) {</pre>

// Get the previous day's stock price

// Get the current day's stock price

prevDeltaX = deltaX;

prevDeltaY = deltaY;

// Initialize variables to keep track of previous slope components

10.

**Python** 

**Calculate and Compare Slopes:** 

in [1, 100]. • Now we compare the slopes. Since dy \* dx1 = dx \* dy1 (100 \* 1 = 1 \* 100), the slope hasn't changed, and we stay on the same line.

We set dx and dy to dx1 and dy1 respectively, for the next comparison.

• We don't increment ans as the line didn't change.

Continue this process for the remaining points:

• We set dx and dy to 0 as there is no previous point to compare with yet.

Comparing [4, 200] with [5, 100], dx1 = 5 - 4 and dy1 = 100 - 200, resulting in [1, -100].  $\circ$  Slopes match (as they're both [1, -100]), so we don't increment ans and just update dx and dy.

Comparing [3, 300] with [4, 200], we calculate dx1 = 4 - 3 and dy1 = 200 - 300, resulting in [1, -100].

 $\circ$  Slopes differ (dy \* dx1 = 100 \* 1!= -1 \* 100 = dx \* dy1), so we increment ans to 2 and update dx and dy.

Finally, comparing [5, 100] with [6, 200], dx1 = 6 - 5 and dy1 = 200 - 100, resulting in [1, 100].

With this hands-on example, we've illustrated how the solution approach can be applied to determine the minimum number of lines required to connect a series of points in a line chart.

After all points have been processed, we conclude that we need ans = 3 lines to connect all the points.

 $\circ$  Slopes differ again (dy \* dx1 = -100 \* 1!= 1 \* 100 = dx \* dy1), so we increment ans to 3.

from typing import List class Solution: def minimumLines(self, stock\_prices: List[List[int]]) -> int: # Sort the stock prices based on the days

# Initialize the previous differences in x and y to 0 and 1 respectively as placeholders.

# then the points are not on the same line. Increment the number of lines.

// Sort the stock prices array based on the day (the 0th element of each sub-array);

// if the days are equal, sort based on the price (the 1st element of each sub-array)

int previousX = stockPrices[i - 1][0], previousY = stockPrices[i - 1][1];

// Check if the current line has a different slope than the previous one

++lineCount; // A different slope means a new line is needed

// Update previous slope components for the next iteration

int currentX = stockPrices[i][0], currentY = stockPrices[i][1];

int deltaX = currentX - previousX, deltaY = currentY - previousY;

int prevDeltaX = 0, prevDeltaY = 1; // Initiated to 0 and 1 to handle the starting calculation

// Calculate the change in x and y (i.e., the slope components) for the current line

// by comparing cross products of the slope components (to avoid floating-point division)

for (current\_x, current\_y), (next\_x, next\_y) in pairwise(stock\_prices): # Compute the differences in x and y for the current pair diff\_x, diff\_y = next\_x - current\_x, next\_y - current\_y # If the cross product of the differences does not equal to zero,

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Java
class Solution {
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// Return the total number of lines needed to connect all points with straight lines
        return lineCount;
C++
#include <vector>
#include <algorithm>
class Solution {
public:
    int minimumLines(std::vector<std::vector<int>>& stockPrices) {
       // Sort stockPrices by their x-value (date)
        std::sort(stockPrices.begin(), stockPrices.end());
       // Initialize previous difference in x and y to compare slopes.
       // Starting with an impossible value for the very first comparison.
       int prevDeltaX = 0, prevDeltaY = 1;
       // Begin with no lines drawn
        int numLines = 0;
       // Iterate through stockPrices starting from the second point
        for (int i = 1; i < stockPrices.size(); ++i) {</pre>
           // Get x and y of the current and previous points
            int currX = stockPrices[i - 1][0], currY = stockPrices[i - 1][1];
            int nextX = stockPrices[i][0], nextY = stockPrices[i][1];
           // Calculate differences in x and y for the current segment
            int deltaX = nextX - currX, deltaY = nextY - currY;
            // If the cross product of the two segments is not zero,
           // they are not collinear, so the slope has changed and a new line is needed.
            if (static_cast<long long>(prevDeltaY) * deltaX != static_cast<long long>(prevDeltaX) * deltaY) {
                ++numLines;
                // Update previous differences to current one for next iteration comparison.
                prevDeltaX = deltaX;
                prevDeltaY = deltaY;
       // Return the total number of lines needed
       return numLines;
```

```
class Solution:
   def minimumLines(self, stock_prices: List[List[int]]) -> int:
       # Sort the stock prices based on the days
        stock_prices.sort()
```

from typing import List

**}**;

**TypeScript** 

function minimumLines(stockPrices: number[][]): number {

// Initialize the number of lines needed to connect all points

let previousSlope: [BigInt, BigInt] = [BigInt(0), BigInt(0)];

const [previousX, previousY] = stockPrices[i - 1];

const differenceX = BigInt(currentX - previousX);

const differenceY = BigInt(currentY - previousY);

const [currentX, currentY] = stockPrices[i];

previousSlope = [differenceX, differenceY];

// Return the total number of lines required

// Start from the second point and compare slopes of consecutive points

// Get the coordinates of the current point and the previous point

// Sort the stock prices based on the first value of each pair (the dates are sorted)

// Initialize a variable to store the previous slope as a pair of BigInts for accurate calculation

// Calculate the difference in x and y using BigInt for accuracy (to handle very large numbers)

// If it's the first line or the slope differs from the previous slope, increment the line count

if (i == 1 || differenceX \* previousSlope[1] !== differenceY \* previousSlope[0]) {

# Initialize the previous differences in x and y to 0 and 1 respectively as placeholders.

// Update the previous slope to the current slope for the next iteration

// Get the number of stock price entries

stockPrices.sort((a, b) => a[0] - b[0]);

let linesRequired = 0;

const stockPricesCount = stockPrices.length;

for (let i = 1; i < stockPricesCount; i++) {</pre>

linesRequired++;

prev\_diff\_x, prev\_diff\_y = 0, 1

num lines += 1

# Initialize the count of lines needed to 0

# Iterate over each pair of adjacent stock prices

prev\_diff\_x, prev\_diff\_y = diff\_x, diff\_y

operations do not increase the overall complexity beyond O(n).

# Return the total number of lines needed

for (current\_x, current\_y), (next\_x, next\_y) in pairwise(stock\_prices):

# If the cross product of the differences does not equal to zero,

# then the points are not on the same line. Increment the number of lines.

# Compute the differences in x and y for the current pair

diff\_x, diff\_y = next\_x - current\_x, next\_y - current\_y

# Update the previous differences for the next iteration

if prev\_diff\_y \* diff\_x != prev\_diff\_x \* diff\_y:

return linesRequired;

from itertools import pairwise

num\_lines = 0

return num\_lines

Time and Space Complexity

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Time Complexity
  The given Python code first sorts the stockPrices. Sorting takes 0(n log n) time where n is the number of stockPrices.
  Then, it iterates through the sorted prices to determine the number of lines required. This iteration involves a single pass through
  the stockPrices list, which takes O(n) time.
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Since sorting is the most time-consuming operation and is followed by a linear pass, the total time complexity of this algorithm is  $O(n \log n) + O(n)$ , which simplifies to  $O(n \log n)$ .

Inside the loop, it performs constant-time arithmetic operations to determine if a new line segment is required. Therefore, these

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Here, the algorithm uses a fixed amount of space to store temporary values like dx, dy, dx1, and dy1, so it uses 0(1) additional
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space.

Sorting is typically done in-place (especially if the Python sort() method is used), therefore, the space complexity remains 0(1).

In summary, the space complexity of the algorithm is 0(1).

**Space Complexity** The space complexity is determined by the additional space required by the algorithm besides the input.