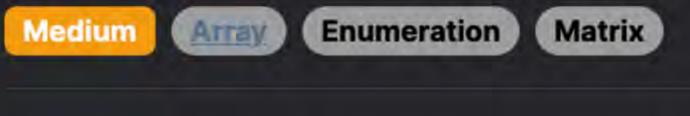
2018. Check if Word Can Be Placed In Crossword



Problem Description

This problem involves a matrix representing a crossword puzzle, which includes lowercase English letters, spaces (represented by 1) for empty cells, and the 1#1 character for blocked cells. The goal is to determine if a given word can be placed in the puzzle

Leetcode Link

following certain conditions. The word can be placed either horizontally or vertically and must adhere to the following rules: 1. The word cannot be placed in cells that contain the '#' character (blocked cells). 2. Each letter of the word must either fill an empty cell (designated by spaces ' ') or match an existing letter on the board.

- 3. If the word is placed horizontally, there should not be any empty cells or other letters immediately to the left or right of the word.
- 4. If the word is placed vertically, there should not be any empty cells or other letters immediately above or below the word.
- The task is to return true if the word can be placed on the board according to the rules, or false otherwise.

To decide whether the given word can be positioned within the board, we should check each cell of the matrix where the first letter

### of the word could potentially be placed. This checking has to take into account the orientation (horizontal and vertical) and also the direction (from left to right, right to left, top to bottom, and bottom to top).

letters. This involves:

Intuition

The intuition behind the provided solution is to systematically iterate over every cell in the matrix and try to match the word considering all possible starting positions and directions that adhere to the crossword rules. For every potential starting position, we examine whether the word would fit without violating any constraints such as running into blocked cells or mismatching existing

 Checking horizontally to the right (left\_to\_right) and to the left (right\_to\_left). Checking vertically downwards (up\_to\_down) and upwards (down\_to\_up). If any of these checks succeed, indicating that the word fits without issue, the function will return true. If no such fitting place is found across the entire board, the function will return false.

The solution efficiently prunes the search by ensuring the word's placement does not start or end next to an empty cell or a different

- letter when placing words horizontally or vertically. As such, it starts the placement from the border of the board or next to a blocked
- cell and checks if every letter of the word can be placed in a suitable position.

**Solution Approach** The solution uses a nested loop to iterate through every cell in the board. For each cell, it checks if this cell could be a potential starting point for the word by following these steps:

1. It checks if the current cell and its immediate neighbor in the opposite direction of the check are either on the edge of the board or blocked by a '#'. This ensures that we only start at valid positions according to the rules of the puzzle. 2. If the starting position is valid, it then invokes the check function which will attempt to place the word starting from that position,

# The check function is designed to validate the placement of the word by iterating over each letter of the word and checking the

The current board cell is not blocked (i.e., not '#').

The board matrix which stores the characters as a 2D list.

following conditions: The current position is within the bounds of the board.

The current board cell is either empty (i.e., ' ') or matches the corresponding letter in the word.

Variables m and n which represent the number of rows and columns in the board, respectively.

 $left_{to} = (j == 0 \text{ or board}[i][j - 1] == '#') and check(i, j, 0, 1)$ 

right\_to\_left = (j == n - 1 or board[i][j + 1] == '#') and check(i, j, 0, -1)

moving either horizontally or vertically and either forwards or backwards depending on the check being performed.

- The check function also ensures that the letter after the last one of the word (calculated by x, y = i + a \* k, j + b \* k) is either out of bounds or blocked. This ensures that the word does not end next to a cell that could violate the horizontal or vertical
- placement rules. The data structures used in the solution are:

The solution approach does not use any additional complex algorithms or patterns. It simply leverages careful iteration and checking of board states to determine if the word can be placed.

The and operator in left\_to\_right, right\_to\_left, up\_to\_down, and down\_to\_up checks combines the start position validation and

the check function call. If any of these conditions return True, it means the word can be placed in the board following the puzzle

 $up_{to}down = (i == 0 \text{ or board}[i - 1][j] == '#') \text{ and } check(i, j, 1, 0)$ down to up = (i == m - 1 or board[i + 1][j] == '#') and check(i, j, -1, 0)if left\_to\_right or right\_to\_left or up\_to\_down or down\_to\_up:

return True

1 for i in range(m):

return False

Example Walkthrough

rules.

positions.

for j in range(n):

Variable k which is the length of the word.

Here's a code snippet encapsulating that logic:

```
Consider a 3 x 3 crossword puzzle board and the word "cat":
1 board = [
      ['#', 'c', '#'],
      [' ', 'a', ' '],
```

see that 't' can match 't'. Thus, "cat" can be placed vertically from (0,1) to (2,1).

would return true, and the algorithm would confirm that "cat" can be placed on the board.

We want to check if we can place the word "cat" on this board.

efficiently designed to stop once a valid position is found.

return False

i < 0 or i >= rows or

j < 0 or j >= cols or

i, j = i + delta\_i, j + delta\_j

word\_length = len(word) # Get the length of the word

return False

return True

# If no valid placement is found, return False

int endCol = j + colIncrement \* wordLength;

// Check each character to see if the word fits

if (i < 0 | | i >= rows | | j < 0 | | j >= cols

for (int p = 0; p < wordLength; ++p) {</pre>

return false;

i += rowIncrement;

j += colIncrement;

return false;

// Check if the word goes out of bounds or is not terminated properly

|| (board[i][j] != ' ' && board[i][j] != word.charAt(p))) {

if (endRow < 0 || endRow >= rows || endCol < 0 || endCol >= cols || board[endRow][endCol] != '#') {

for char in word:

if (

):

return True

The cell (0,1) contains the first letter of the word "cat", 'c', and cell (0,0) is blocked, which is a valid start position. However, we cannot place "cat" horizontally to the right because there is no space to fit the entire word "cat". Next, we check horizontally to the left (right\_to\_left). This direction is not applicable in this case as we are looking for starting

Then, we check vertically downwards (up\_to\_down). From (0,1), we realize (0,0) is blocked, providing a potential starting position

for "cat". We can successfully match 'c' with 'c', then move to the next cell (1,1) and match 'a' with 'a', and finally move to (2,1) and

Since we found a valid placement for "cat", we do not need to check vertically upwards (down\_to\_up) from (0,1). The check function

So in our algorithm, as soon as it runs the up\_to\_down check starting at (0,1), it will return true, indicating that the word "cat" can

indeed be placed on the board vertically. The result is obtained without having to check other cells or directions, as the solution is

First, we check horizontally to the right (left\_to\_right). Starting from (0,0) we find it's a blocked cell ('#'), so we move to (0,1).

Using our solution approach, we will check each cell starting from (0,0) to (2,2) to find a valid placement for "cat".

Python Solution from typing import List

def place\_word\_in\_crossword(self, board: List[List[str]], word: str) -> bool: # Function to check if the word fits starting from position (i, j) in the direction specified by (delta\_i, delta\_j) def is\_valid\_placement(i, j, delta\_i, delta\_j): 6 # Move to the end of the word in the specified direction to check if it is within bounds or blocked by '#' end\_i, end\_j = i + delta\_i \* word\_length, j + delta\_j \* word\_length

left to right = (j == 0 or board[i][j - 1] == '#') and is\_valid\_placement(i, j, 0, 1)

 $top_{to}bottom = (i == 0 or board[i - 1][j] == '#') and is_valid_placement(i, j, 1, 0)$ 

right\_to\_left =  $(j == cols - 1 or board[i][j + 1] == '#') and is_valid_placement(i, j, 0, -1)$ 

bottom\_to\_top = (i == rows - 1 or board[i + 1][j] == '#') and is\_valid\_placement(i, j, -1, 0)

# Check for out of bounds or if the current board cell is blocked or does not match the word character

if not (0 <= end\_i < rows and 0 <= end\_j < cols) or (board[end\_i][end\_j] == '#'):</pre>

# Iterate through each character of the 'word' to check for a valid placement

(board[i][j] != ' ' and board[i][j] != char)

# Move to the next cell in the specified direction

rows, cols = len(board), len(board[0]) # Get the dimensions of the board

# If the word can be placed in any direction, return True

if left\_to\_right or right\_to\_left or top\_to\_bottom or bottom\_to\_top:

#### 29 30 # Iterate over every cell in the board 31 for i in range(rows): for j in range(cols): 32 33 # Check all four directions from the current cell: left to right, right to left, top to bottom, bottom to top

class Solution:

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```
return False
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 46 # Example usage:
 47 # solution = Solution()
 48 # result = solution.place_word_in_crossword(board=[['#',' ','#'],[' ',' ','#'],['#','c',' ']], word="abc")
 49 # print(result) # Output will be True or False based on if the word can be placed on the board
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Java Solution
   class Solution {
         private int rows;
         private int cols;
         private char[][] board;
  4
         private String word;
  5
         private int wordLength;
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  8
         // Method to check if the word can be placed in the crossword
         public boolean placeWordInCrossword(char[][] board, String word) {
  9
 10
             rows = board.length;
             cols = board[0].length;
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 12
             this.board = board;
 13
             this.word = word;
 14
             wordLength = word.length();
 15
             // Traverse the board to check every potential starting point
 16
             for (int i = 0; i < rows; ++i) {
 17
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                 for (int j = 0; j < cols; ++j) {
 19
                     // Check four possible directions from the current cell
 20
                     // Left to right
 21
                     boolean leftToRight = (j == 0 \mid | board[i][j - 1] == '#') && canPlaceWord(i, j, 0, 1);
 22
                     // Right to left
 23
                     boolean rightToLeft = (j == cols - 1 \mid | board[i][j + 1] == '#') && canPlaceWord(i, j, 0, -1);
 24
                     // Up to down
 25
                     boolean upToDown = (i == 0 \mid | board[i - 1][j] == '#') && canPlaceWord(i, j, 1, 0);
 26
                     // Down to up
 27
                     boolean downToUp = (i == rows - 1 || board[i + 1][j] == '#') && canPlaceWord(i, j, -1, 0);
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 29
                     // If any direction is possible, return true
 30
                     if (leftToRight || rightToLeft || upToDown || downToUp) {
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                         return true;
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             // If no direction is possible, return false
             return false;
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         // Helper method to check if the word can be placed starting from (i, j) in the specified direction (a, b)
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         private boolean canPlaceWord(int i, int j, int rowIncrement, int colIncrement) {
             int endRow = i + rowIncrement * wordLength;
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```

#### 58 return true; 59 60 61

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C++ Solution
  1 class Solution {
  2 public:
         // Function to determine if a word can be placed in a crossword
         bool placeWordInCrossword(vector<vector<char>>& board, string word) {
             int numRows = board.size(), numCols = board[0].size(); // board dimensions
             int wordLen = word.size(); // length of the word to be placed
  6
             // Lambda function to check if the word fits in the given direction
             auto check = [&](int row, int col, int deltaRow, int deltaCol) {
  9
 10
                 int endRow = row + deltaRow * wordLen, endCol = col + deltaCol * wordLen;
 11
                 // Check if the end position is not blocked by '#'
                 if (endRow >= 0 && endRow < numRows && endCol >= 0 && endCol < numCols && board[endRow][endCol] != '#') {
 12
 13
                     return false;
 14
 15
                 // Iterate over each character in the word
 16
                 for (char& c : word) {
                     // Check boundaries and match the character with the board or wildcard
 17
 18
                     if (row < 0 || row >= numRows || col < 0 || col >= numCols || (board[row][col] != ' ' && board[row][col] != c)) {
 19
                         return false;
 20
 21
                     row += deltaRow;
 22
                     col += deltaCol;
 23
 24
                 return true;
 25
             };
 26
             // Iterate over each cell in the board
 27
 28
             for (int i = 0; i < numRows; ++i) {</pre>
 29
                 for (int j = 0; j < numCols; ++j) {
                     // Check four possible directions where the word can be placed
 30
 31
                     bool leftToRight = (j == 0 || board[i][j - 1] == '#') && check(i, j, 0, 1);
 32
                     bool rightToLeft = (j == numCols - 1 \mid | board[i][j + 1] == '#') && check(i, j, 0, -1);
                     bool upToDown = (i == 0 \mid | board[i - 1][j] == '#') && check(i, j, 1, 0);
 33
                     bool downToUp = (i == numRows - 1 \mid | board[i + 1][j] == '#') && check(i, j, -1, 0);
 34
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 36
                     // If the word can be placed in any direction, return true
 37
                     if (leftToRight || rightToLeft || upToDown || downToUp) {
 38
                         return true;
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 42
             // Return false if the word can't be placed on the board in any direction
 43
             return false;
 44
 45 };
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Typescript Solution
  1 // Function to check if a word can be placed in a crossword
    function placeWordInCrossword(board: char[][], word: string): boolean {
         const numRows = board.length; // Number of rows in the board
         const numCols = board[0].length; // Number of columns in the board
         const wordLen = word.length; // Length of the word to be placed
  6
  7
         // Helper function to check if the word fits in the given direction
         const check = (row: number, col: number, deltaRow: number, deltaCol: number): boolean => {
  8
             const endRow = row + deltaRow * wordLen;
  9
 10
             const endCol = col + deltaCol * wordLen;
             // Check if the end position is outside the boundary or blocked by '#'
 11
```

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// Iterate over each character in the word

const char = word[index];

return false;

row += deltaRow;

col += deltaCol;

// Iterate over each cell in the board

for (let i = 0; i < numRows; ++i) {

return true;

return false;

for (let index = 0; index < wordLen; ++index) {</pre>

// Check boundaries and match the character with the board or wildcard

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};

Time Complexity

30 for (let j = 0; j < numCols; ++j) {</pre> // Check four possible directions where the word can be placed 31 32 const leftToRight =  $(j === 0 \mid | board[i][j - 1] === '#') && check(i, j, 0, 1);$ 33 const rightToLeft =  $(j === numCols - 1 \mid | board[i][j + 1] === '#') && check(i, j, 0, -1);$ const upToDown = (i === 0 || board[i - 1][undefined] === '#' || board[i - 1][j] === '#') && check(i, j, 1, 0); 34 const downToUp =  $(i === numRows - 1 \mid | board[i + 1] === undefined \mid | board[i + 1][j] === '#') && check(i, j, -1, 0);$ // If the word can be placed in any direction, return true if (leftToRight || rightToLeft || upToDown || downToUp) { 39 40 41 42 // Return false if the word can't be placed on the board in any direction 43 44 return false; 45 46

if (endRow >= 0 && endRow <= numRows && endCol >= 0 && endCol <= numCols && (board[endRow] === undefined || board[endRow][e

if (row < 0 || row >= numRows || col < 0 || col >= numCols || (board[row][col] !== ' ' && board[row][col] !== char)

## loop up to the length of the word (k). Additionally, when evaluating the board for possible placements, the algorithm checks the perpendicular cells to ensure placement is

Time and Space Complexity

do not impact the linear relationship between the time complexity and the number of cells times the length of the word. Space Complexity

The time complexity of the given code is 0 (m \* n \* k), where m is the number of rows in the board, n is the number of columns in the

board, and k is the length of the word to be placed. This complexity arises because the code iterates over all cells of the board (m \*

n) and for each cell, it attempts to place the word in all four directions. The check function, which is called for each direction, runs a

at the beginning or end of a word sequence (which is a constant time check). Due to these operations being constant in time, they

The space complexity of the code is 0(1) (constant space complexity). This is because the algorithm only uses a fixed amount of extra space for variables that store the dimensions of the board and indices during the checks regardless of the input size. No additional space proportional to the input size is required beyond what is used to store the board and word, which are inputs and not counted towards the space complexity.