2507. Smallest Value After Replacing With Sum of Prime Factors Leetcode Link

Medium (Math)

Problem Description

The problem involves manipulating a given positive integer n by continuously replacing it with the sum of its prime factors. A key detail is that if a prime factor is repeated in the factorization (that is, n is divisible by a prime number multiple times), that prime factor must be counted multiple times in the sum.

For example, if our given number n is 18, its prime factors are 2 and 3, but 3 is counted twice because 18 is divisible by 3 twice (18 = 2 * 3 * 3). So the sum of its prime factors is 2 + 3 + 3 = 8.

The challenge is to write an algorithm that performs this computation efficiently and find when the number n stops changing, which is the result to be returned.

The goal is to repeat this process: take the sum of the prime factors of n and then replace n with this sum, and continue this until n

Intuition

The solution makes use of a loop that continuously replaces the number n with the sum of its prime factors until n cannot be reduced

any further.

and should be included in the sum).

Solution Approach

three variables - t to keep track of the original value of n at the start of the iteration, s to calculate the sum of prime factors, and 1 to iterate over potential prime factors starting from 2.

We start with a loop that runs indefinitely, checking if we can find prime factors of n and their sums. In each iteration, we initialize

The inner while loop checks if i is a divisor of n, and if so, it keeps dividing n by i and adds i to the sum s until n is no longer divisible by i. The variable i is incremented to check for the next possible prime factor. Once we have tried all possible divisors up to n // i, we check if n itself is greater than 1 (which would mean n is a prime number

At this point, if the calculated sum of prime factors s is equal to the temporary variable t, it means n can no longer be reduced any further, and we return t.

If s is not equal to t, we replace n with s to continue the process with a new reduced value of n. The key intuition behind the algorithm is that the prime factorization of a number can help find a smaller representation of the

number itself and by repeating this process and summing these factors repeatedly, we can converge to the smallest possible value of n.

The solution makes use of a basic factorization algorithm and control flow to reduce the number n to its smallest possible value by continuously summing its prime factors.

3. Another variable s is initialized to 0; this will be used to calculate the sum of the prime factors of n.

continues until n is no longer divisible by i, accommodating for all instances of i as a factor.

4. An index i is set to 2, which is the first prime number. This variable is used to test potential factors of n.

2. Inside the loop, a temporary variable t is assigned the current value of n to keep track of its value through each iteration. This is important to identify when there is no further change possible.

5. The first inner while loop runs as long as i is less than or equal to n // i (since a factor larger than the square root of n would have already been identified by its corresponding smaller factor except when n is prime).

1, implying n is a prime number. If it is prime, it is added to the sum s.

indivisible and signaling the end of the reduction process.

Step 4: Start with i = 2, which is the smallest prime factor.

number to its smallest form based on prime factor accumulation.

with the sum of its prime factors and repeat this process until n is no longer reduced.

Step 3: Initialize a variable s to 0. This will hold the sum of the prime factors of n.

Step 5: As $i \ll n // i$ (since 2 $\ll 12 // 2$), we proceed with factorization.

Step 10: Because s is not equal to t, we set n to s; hence n is now 7.

Here's a step-by-step breakdown of the approach used in the solution:

1. The solution uses a while loop that will run until n no longer changes.

7. The index i is then incremented to check the next potential factor.

8. After all possible factors up to n // i have been tested, a final check outside the inner loops evaluates if n itself is greater than

6. A nested inner while loop checks if i divides n perfectly. If it does, n is divided by i, and i is added to the sum s. This loop

reached its smallest value as no primes other than itself can be extracted and summed, and the algorithm returns the value of t. 10. If s is not equal to t, the algorithm replaces the value of n with s to repeat the factorization process on this new reduced number.

The algorithm essentially terminates when a number only comprises its prime self or when it's reduced to 1, both of which are

This algorithm does not employ any complex data structures and follows a straightforward but effective pattern to reduce the

9. At the end of the iteration, the algorithm checks if the sum of the prime factors s is equal to the starting value t. If so, n has

Example Walkthrough

Let's walk through the solution approach with a small example where our given number n is 12. The goal is to continuously replace n

Step 1: Start with n = 12 and enter the while loop. Step 2: Assign the value of n to a temporary variable t, so t = 12.

We continue with the inner loop since 6 is still divisible by 2. We divide 6 by 2 to get 3 and add 2 to s. Now, s = 4 and n = 3.

Step 7: Increment i to the next integer, which is 3.

Step 8: Since 3 is a prime number larger than n // i but still divides n perfectly, we add 3 to s. Now, s = 4 + 3 giving us s = 7 and n

Step 9: Since n is now 1, we exit the inner loop. We check if s is equal to the temporary variable t. As s = 7 and t = 12, they are not

Step 6: The inner loop checks if 12 is divisible by 2. It is, so we divide n by 2 to get 6 and add 2 to 5. Now, s = 2 and n = 6.

Step 1: n is 7, so we enter the while loop.

Step 9: Compare the sum of the prime factors s with t. They are equal (s = t = 7), indicating that we cannot reduce n any further. The algorithm would then terminate and return the value 7 as the result, as n can no longer be reduced by the process described.

Step 5: Begin inner loop. Because no i exists such that i <= n // i and i divides 7, skip to step 8.

Check for factors of the number starting with the smallest prime factor

Set num to sum_of_factors for the next iteration to check the new number

Step 8: Since n is still greater than 1, and 7 is a prime number, we add 7 to s. Now s = 7.

divisor += 1 14 # If there's a remaining number greater than one, it's a prime factor; add it to the sum_of_factors 15 if num > 1: 16 sum_of_factors += num # If the sum of factors is equal to the original number, we found the smallest value

If the divisor is a factor, divide num by the divisor and add to the sum_of_factors

Now n has changed from 12 to 7. We repeat the entire process again with n = 7.

Step 2: Set t to 7.

Step 3: Initialize s to 0.

Step 4: Start with i = 2.

Python Solution

class Solution:

Java Solution

class Solution {

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C++ Solution

1 class Solution {

public:

becomes 1.

equal.

def smallestValue(self, num: int) -> int: # Continue the loop until we find the smallest value while True: # Initialize temp variable to store original number, sum_of_factors, and start divisor from 2 temp, sum_of_factors, divisor = num, 0, 2

while divisor <= num // divisor:</pre>

while num % divisor == 0:

sum_of_factors += divisor

Move to the next potential factor

num //= divisor

if sum_of_factors == temp:

return temp

public int smallestValue(int n) {

int originalValue = n;

int sumOfFactors = 0;

while (true) {

22 num = sum_of_factors 23

```
sumOfFactors += i; // Add factor to the sum
14
15
                        n /= i; // Divide n by the factor
16
               // If there is a remaining factor greater than 1, add it to the sum
               if (n > 1) {
19
                    sumOfFactors += n;
20
21
22
               // Check if the sum of the factors equals the original number
               if (sumOfFactors == originalValue) {
```

// If it matches, return the sum (as it is the smallest value)

// If it does not match, set n to the sumOfFactors for another iteration

// Start dividing the number from 2 onwards to find its factors

// Method to find the smallest value according to specified conditions

// Loop indefinitely until we find the smallest value

for (int i = 2; i <= originalValue / i; ++i) {</pre>

// Function to find the smallest integer value with the same

// sum of factors (including 1 and the number itself) as 'n'

// Loop indefinitely until we find the smallest value

int originalValue = n; // Preserve the original value of 'n'

int sumOfFactors = 0; // Initialize the sum of factors to 0

// Divide by i as long as it is a factor of n

// Store the original value of n

while $(n \% i == 0) {$

return sumOfFactors;

n = sumOfFactors;

int smallestValue(int n) {

while (true) {

while (true) {

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// Factorize and sum up the factors

while (n % i === 0) {

if (i != n / i) {

if (n > 1 && n < originalValue) {

if (sumOfFactors == originalValue) {

return originalValue;

// While 'n' is divisible by 'i'

// Initialize the sum of the factors

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11
               // Factorize and sum up the factors
               for (int i = 2; i \le n / i; ++i) { // Only need to check up to sqrt(n)
12
                   // While 'n' is divisible by 'i'
                   while (n \% i == 0) {
14
15
                       sumOfFactors += i; // Add the factor to the sum
                       n /= i; // Divide 'n' by the factor for further factorization
16
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20
               // If there is a remaining prime factor greater than sqrt(n), add it to the sum
               if (n > 1) sumOfFactors += n;
21
22
23
               // If the sum of factors is the same as the original value, return it
24
               if (sumOfFactors == originalValue) {
                   return sumOfFactors;
25
26
27
28
               // Otherwise, set 'n' to the calculated sum of factors for the next iteration
29
               n = sumOfFactors;
30
31
32 };
33
Typescript Solution
 1 // Function to find the smallest integer value with the same
2 // sum of factors (including 1 and the number itself) as 'n'
   function smallestValue(n: number): number {
       // Loop indefinitely until we find the smallest value
```

const originalValue: number = n; // Preserve the original value of 'n'

sumOfFactors += i; // Add the factor to the sum

// Check for a remaining prime factor greater than sqrt(n)

for (let i = 2; i <= Math.sqrt(n); ++i) { // Only need to check up to sqrt(n)

n /= i; // Divide 'n' by the factor for further factorization

let sumOfFactors: number = 1; // Initialize the sum of factors to 1, since 1 is a factor of all numbers

sumOfFactors += n / i; // Add the complement factor to the sum if it's different

sumOfFactors += n; // Add it to the sum only if it's different from the current value

// If the sum of factors is the same as the original value, return the smallest integer

n = sumOfFactors; // Set 'n' to the calculated sum of factors for the next iteration

Time Complexity The time complexity of this code is governed by two nested loops. The outer loop runs indefinitely until a number is found where the

because it checks for factors from 2 to at most sqrt(n).

Time and Space Complexity

The factorization loop:

 In the worst-case scenario for a given n, we may have to factorize n, then s, and so on if n is initially not a prime or semiprime (a product of exactly two primes). Each new s will be smaller than n as we are adding up the factors. Thus, the inner factorization process will take O(sqrt(m)) time for each number m we factorize, where m starts from n and gets smaller. The outer while loop:

The code does not have a clear stopping condition within predictable bounds, due to the uncertain nature of the sum s

be high, but the exact upper bound is intricate to determine without more constraints on n.

sum of its prime factors is equal to itself. The inner loop, on the other hand, is used for factorization and runs at most sqrt(n) times

The number of iterations can be considered proportional to the number of distinct prime factors of n in a sense but is not easy to express in standard 0() notation. Considering the worst-case scenario, where n is a large composite number with many small prime factors, the time complexity can

approaching the target condition s == t. However, for every iteration, the sum of the prime factors of n (s) gets closer to being a

prime number itself. Once s equals a prime number, it will become equal to t in the next iteration, and the function will return t.

The space complexity is 0(1). There are only a few integer variables (t, s, i, and n) being used, which do not depend on the input

Space Complexity

size n, unless considering the size n itself needs. The space used by these variables is constant and does not grow with n.

Number Theory

no longer changes—it reaches its smallest possible value.