1975. Maximum Matrix Sum

Matrix

# **Problem Description**

Greedy Array

Medium

matrix. An operation consists of selecting any two adjacent elements (elements that share a border) and multiplying both by -1. This toggling between positive and negative can be done as many times as desired. The goal is to determine the highest possible sum of all elements in the matrix after performing zero or more of these operations.

Given an n x n matrix filled with integers, the task is to perform strategic operations to maximize the sum of all elements in the

Intuition

pairs of adjacent elements, it may not always be possible to convert all negatives to positives. The strategy is as follows: Calculate the total sum of all the absolute values of the elements in the matrix. This is the maximum sum the matrix can

The key to solving this problem is realizing that each operation can be used to negate the effect of negative numbers in the

matrix. It's generally advantageous to have all numbers positive for the highest sum. However, since operations are limited to

possibly have if we could individually toggle each element to be positive.

sum.

- Track the smallest absolute value of the elements in the matrix. This value is important because if there's an odd number of negative elements that can't be paired, the smallest element's negativity will have the smallest negative impact on the total
- Count the number of negative elements in the matrix. If this count is even, it means every negative element can be paired with another negative to turn both into positives through one operation.

If the count of negative elements is odd, there will remain one unpaired negative element affecting the total sum. In this case,

- subtract twice the smallest absolute value found in the matrix from the total sum to account for the impact of the remaining negative number after all possible pairs have been negated. If the smallest value is zero, it means an operation can be done without impacting the total sum, as multiplying zero by -1 does not change the value.
- The solution leverages these intuitions to determine the maximum sum. Solution Approach

To implement the solution, we perform a series of steps that follow the intuition described earlier. Here's a breakdown of the

Initialize a sum variable s to 0, which will store the total sum of absolute values of the matrix elements. Also, create a counter

algorithm based on the code provided:

each row). For every element v in the matrix:

summing all absolute values. Return the sum s.

the sum. So, return s.

Let's consider a small 3x3 matrix as an example:

# cnt to keep track of the number of negative elements and a variable mi to keep the minimum absolute value observed in the

matrix. Iterate through each element of the matrix using a nested loop (iterating first through rows, then through elements within

- $\circ$  Add the absolute value of the element to s (s += abs(v)), progressively building the total sum of absolute values. Update the minimum absolute value mi if the absolute value of the current element is less than the current mi (mi = min(mi, abs(v))). ∘ If the current element v is negative, increment the counter cnt (if v < 0: cnt += 1). This helps keep track of how many negative numbers are in the matrix, crucial for deciding the following steps.
- After iterating through the matrix: Check if the number of negative elements cnt is even or the smallest value mi is zero:
  - If cnt is even, it's possible to negate all negative elements by using the operation, and the highest sum can be obtained by simply
  - If the number of negative elements cnt is odd and mi is not zero, it means that not all negative elements can be negated, and there will be one negative element affecting the total sum. Return s - mi \* 2, subtracting twice the smallest absolute

value to account for the remaining negative element's impact on the total sum.

■ If mi is zero, it means one of the elements is zero, and no subtraction is needed as multiplying zero by -1 multiple times will not change

calculation of sums, tracking the smallest value, and counting occurrences of a particular condition (negativity in this case), which are fairly common operations in matrix manipulation problems.

This approach uses simple data structures (just integers and loops over the 2D matrix), with the main pattern being the

Following the solution approach:

o cnt is 4, which is even. Therefore, we don't need to subtract anything from s, because we can pair all negatives and multiply by -1 to turn

### Iterate through the matrix and calculate: ∘ s (the sum of absolute values), which is | −1 | + | −2 | + | 3 | + | 4 | + | 5 | + | −6 | + | −7 | + | 8 | + | 9 | = 1 + 2 + 3 + 4

them positive.

+ 5 + 6 + 7 + 8 + 9 = 45.

**Example Walkthrough** 

Since cnt is even, we can return s which equals 45 as our answer.

Therefore, the highest possible sum after applying the operations is 45.

According to the strategy, the operations performed would be:

Initialize s as 0, cnt as 0, and mi as infinity (or a very large number).

 $\circ$  mi (the smallest absolute value), which is min(infinity, 1, 2, 3, 4, 5, 6, 7, 8, 9) = 1.

 $\circ$  cnt (the count of negative numbers), there are 4 negatives: -1, -2, -6, -7 so cnt = 4.

 Multiplying -1 and -2 by -1 to turn them positive. Multiplying –6 and –7 by –1 to turn them positive.

After the iteration, check the conditions:

**Python** 

total\_sum = negative\_count = 0

if value < 0:</pre>

return total\_sum

# Iterate over each row in the matrix

total\_sum += abs(value)

negative\_count += 1

long long maxMatrixSum(vector<vector<int>>& matrix) {

// Loop over each element in the row

negativeCount += value < 0;</pre>

sumOfAbsoluteValues += abs(value);

if (negativeCount % 2 == 0 || minAbsValue == 0) {

minAbsValue = min(minAbsValue, abs(value));

int minAbsValue = INT\_MAX;

// Loop over each row in the matrix

return sumOfAbsoluteValues;

for (int& value : row) {

for (auto& row : matrix) {

Solution Implementation

class Solution:

for row in matrix: # Iterate over each value in the row for value in row: # Add the absolute value of the current element to the total sum

# Initialize total sum, negative count, and minimum positive value

# Record the minimum positive value (smallest absolute value)

# If there's an odd number of negatives, subtract twice the minimum positive

# value to account for the one value that will remain negative

minimum\_positive\_value = min(minimum\_positive\_value, abs(value))

# If the current value is negative, increment the negative count

minimum\_positive\_value = float('inf') # Represents infinity

def maxMatrixSum(self, matrix: List[List[int]]) -> int:

```
# If there is an even number of negatives (or a zero element, which can flip sign without penalty),
# the result is simply the total sum of absolute values
if negative_count % 2 == 0 or minimum_positive_value == 0:
```

class Solution {

public:

```
return total_sum - minimum_positive_value * 2
Java
class Solution {
   public long maxMatrixSum(int[][] matrix) {
        long sum = 0; // Initialize a sum variable to hold the total sum of matrix elements
        int negativeCount = 0; // Counter for the number of negative elements in the matrix
        int minAbsValue = Integer.MAX_VALUE; // Initialize to the maximum possible value to track the smallest absolute value see
       // Loop through each row of the matrix
       for (int[] row : matrix) {
           // Loop through each value in the row
           for (int value : row) {
               sum += Math.abs(value); // Add the absolute value of the element to the sum
               // Find the smallest absolute value in the matrix
               minAbsValue = Math.min(minAbsValue, Math.abs(value));
               // If the element is negative, increment the negativeCount
               if (value < 0) {
                   negativeCount++;
       // If the count of negative numbers is even or there's at least one zero, return the sum of absolute values
       if (negativeCount % 2 == 0 || minAbsValue == 0) {
           return sum;
       // Since the negative count is odd, we subtract twice the smallest absolute value to maximize the matrix sum
       return sum - (minAbsValue * 2);
C++
```

long long sumOfAbsoluteValues = 0; // This variable will store summation of absolute values of all elements in matrix

// This variable will keep track of the smallest absolute value encountered

int negativeCount = 0; // Counter for the number of negative elements in the matrix

// Update the smallest absolute value encountered if current absolute value is smaller

// Add the absolute value of the current element to the total sum

// If the number of negative values is even or the minimum absolute value is 0,

// If the current element is negative, increment the negative counter

// we can make all elements non-negative without decreasing the sum of absolute values.

```
// If the number of negative values is odd, we subtract twice the smallest
       // absolute value to compensate for the one element that will remain negative.
       return sumOfAbsoluteValues - minAbsValue * 2;
};
TypeScript
/**
 * Calculates the maximum absolute sum of any submatrix of the given matrix.
 * @param {number[][]} matrix - The 2D array of numbers representing the matrix.
 * @return {number} - The maximum absolute sum possible by potentially negating any submatrix element.
function maxMatrixSum(matrix: number[][]): number {
    let negativeCount = 0; // Count of negative numbers in the matrix
    let sum = 0; // Sum of the absolute values of the elements in the matrix
    let minAbsValue = Infinity; // Smallest absolute value found in the matrix
    // Iterate through each row of the matrix
    for (const row of matrix) {
       // Iterate through each value in the row
        for (const value of row) -
            sum += Math.abs(value); // Add the absolute value to the sum
           minAbsValue = Math.min(minAbsValue, Math.abs(value)); // Update min absolute value if necessary
            negativeCount += value < 0 ? 1 : 0; // Increment count if the value is negative</pre>
    // If the count of negative numbers is even, the sum is already maximized
    if (negativeCount % 2 == 0) {
        return sum;
    // Otherwise, subtract double the smallest absolute value to negate an odd count of negatives
    return sum - minAbsValue * 2;
// Example usage:
// const matrix = [[1, -1], [-1, 1]];
```

#### # If there's an odd number of negatives, subtract twice the minimum positive # value to account for the one value that will remain negative return total\_sum - minimum\_positive\_value \* 2

Time and Space Complexity

return total\_sum

// const result = maxMatrixSum(matrix);

class Solution:

// console.log(result); // Output should be 4

total\_sum = negative\_count = 0

for row in matrix:

for value in row:

if value < 0:

# Iterate over each row in the matrix

total\_sum += abs(value)

def maxMatrixSum(self, matrix: List[List[int]]) -> int:

# Iterate over each value in the row

negative\_count += 1

# the result is simply the total sum of absolute values

if negative\_count % 2 == 0 or minimum\_positive\_value == 0:

# Initialize total sum, negative count, and minimum positive value

# Add the absolute value of the current element to the total sum

# Record the minimum positive value (smallest absolute value)

minimum positive value = min(minimum positive value, abs(value))

# If the current value is negative, increment the negative count

# If there is an even number of negatives (or a zero element, which can flip sign without penalty),

minimum\_positive\_value = float('inf') # Represents infinity

# **Time Complexity**

The given code iterates through all elements of the matrix with dimension n x n. For each element, it performs a constant number of operations: calculating the absolute value, updating the sum s, comparing with the minimum value mi, and incrementing a counter cnt if the value is negative. • Iterating through all elements takes 0(n^2) time, where n is the dimension of the square matrix (since there are n rows and n columns).

 All the operations inside the nested loops are constant time operations. Hence, combining these, the overall time complexity of the code is  $0(n^2)$ .

**Space Complexity** 

• The variables s, cnt, and mi use constant space. There are no additional data structures that grow with the size of the input.

Therefore, the space complexity of the code is 0(1), which indicates constant space usage regardless of the input size.

The space complexity is determined by the additional space required by the code, not including the space taken by the inputs.