2120. Execution of All Suffix Instructions Staying in a Grid String Medium Simulation

robot's position is defined by the startPos, which is an array of two integers representing the row and column. Additionally, we are provided with a string s of length m, representing a sequence of instructions ('L', 'R', 'U', 'D') that tell the robot how to move within the grid.

The task is to find out, for each instruction starting from the ith position in the string s, how many instructions the robot can execute

Leetcode Link

before one of two conditions occurs: 1. The robot is instructed to move off the grid (out of bounds).

The final output should be an array answer where each element answer[i] represents the number of instructions the robot can

Check if the next move stays within the grid boundaries.

by considering every possible execution path from every starting point.

b. Check if the movement is within grid bounds:

position and increment the t counter.

execute when starting from the ith instruction in s.

2. The robot reaches the end of the instruction string without moving off the grid.

The solution to this problem lies in simulating the path of the robot's movement on the grid from each possible starting point in the

instruction string s.

For each instruction starting point i, we simulate the robot's movement by iterating over instructions from s[i] to s[m - 1] and tracking its current position. We use a dictionary mp to translate the instruction characters ('L', 'R', 'U', 'D') into corresponding row

and column movements.

With each step, we:

 If it does, we update the robot's position and increment the counter for executed instructions. If the next move would take the robot out of bounds, we stop and consider the number of executed instructions as the result for the starting point i. We repeat this process for every instruction as a potential starting point, obtaining the total number of instructions the robot can

The approach relies on a brute-force simulation for each starting point, which simplifies the implementation and ensures correctness

Solution Approach

point in the instruction string s until it either moves off the grid or runs out of instructions. Here are the steps and the constructs used in the implementation:

2. Create a mapping (mp) which translates instruction characters ('L', 'R', 'U', 'D') to their corresponding movements on the grid

The provided solution approach utilizes the simulation strategy, where we follow the robot's instructions from each possible start

as row (x) and column (y) changes.

Algorithm:

point.

3. Loop through each instruction i as a start point in the instruction string s. For each start point:

a. Initial Position: Set the robot's current position (x, y) to the starting position (startPos).

1. Initialize an empty list (ans) to collect the number of executable instructions for each starting point.

execute for each case. These counts are added to the ans list in the order of instruction starting points.

4. Iterate over the instructions from the current start point i to the end of the string. For each instruction j: a. Retrieve the corresponding movement from the mp dictionary.

 \circ If the resulting position (x + a, y + b) is within the grid (0 <= x < n and 0 <= y < n), apply the movement to the current

b. Execution Counter (t): Initialize a counter to keep track of how many instructions have been executed from that starting

c. After the loop, append the value of t (the number of instructions executed) to the ans list. 5. Once we have simulated starting from each instruction and collected the executable instruction counts, return ans.

If the move would go off-grid, break out of the loop since the robot cannot execute this instruction.

• List: To store the final count of executable instructions (ans) and the starting position (startPos). • **Dictionary**: To translate instructions into actual x and y movements (mp).

Let's walk through a small example to illustrate the solution approach described above.

are trying to query, thus avoiding the need for more complex algorithms or optimizations. By using this brute-force simulation, the solution effectively handles any given n x n grid and instruction set s, ensuring a correct

• Simulation: This problem is directly addressed through simulatory execution of instructions, mimicking the exact behavior we

Suppose we have a grid of size 3×3 , and the robot starts at startPos = [1,1] (the center of the grid). Imagine our string of

and complete answer.

Example Walkthrough

Start from startPos (1,1).

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Start from startPos (1, 1).

from typing import List

results = []

class Solution:

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After the loop, we have: ans = [3, 2, 1].

instruction string before moving off the grid.

Now let's simulate according to the algorithm:

1. Initialize an empty list (ans) to store results.

Data Structures:

Patterns:

3. Loop through each i in s — here, i will take values 0 for 'R', 1 for 'U', and 2 for 'L'. For i = 0 ('R'):

2. Create a mapping (mp) of {'L': (0, -1), 'R': (0, 1), 'U': (-1, 0), 'D': (1, 0)} to represent movements.

instructions, s, is "RUL". We want to find out how many instructions the robot can execute starting at each character of the

 Move right (0, 1) to (1,2). Grid boundary not exceeded, increment t to 1. Next is 'U', move up (-1, ∅) to (0,2). Grid boundary not exceeded, increment t to 2. Last is 'L', move left (∅, -1) to (0,1). Grid boundary not exceeded, increment t to 3. • Add t = 3 to ans.

 Next is 'L', move left (∅, −1) to (0,0). Grid boundary not exceeded, increment t to 2. There are no more instructions, so add t = 2 to ans.

For i = 1 ('U'):

For i = 2 ('L'):

Move up (-1, ∅) to (0, 1). Grid boundary not exceeded, increment t to 1.

 Move left (∅, −1) to (1,0). Grid boundary not exceeded, increment t to 1. No more instructions, so add t = 1 to ans.

instructions; from 'U', 2 instructions; and from 'L', 1 instruction before going off the grid or running out of instructions.

def execute_instructions(self, n: int, start_pos: List[int], instructions: str) -> List[int]:

Mapping for each instruction to its corresponding row and column changes

Get the row and column changes for the current instruction

row_change, col_change = direction_map[instructions[j]]

Update the current position based on the instruction

public int[] executeInstructions(int n, int[] startPos, String instructions) {

// Loop through each position in the instructions as a starting point

// A map to associate directions with their respective 2D coordinate changes

int x = startPos[0], y = startPos[1]; // Current position to check from

// Loop through the instructions starting from the current position

for (int current = start; current < instructionsLength; ++current) {</pre>

x += deltaX; // Update the current x-coordinate

y += deltaY; // Update the current y-coordinate

vector<int> executeInstructions(int n, vector<int>& startPos, string s) {

int instructionCount = s.size(); // number of instructions to execute

++validMoves; // Increment the valid move counter

results[start] = validMoves; // Store the valid moves for this start position

return results; // Return the array containing valid moves for each start position

// Retrieve the 2D coordinate change for the current direction

int instructionsLength = instructions.length(); // Length of the instructions string

int[] results = new int[instructionsLength]; // Array to store the result for each position

int validMoves = 0; // Counter to count the valid moves from the current starting position

char currentInstruction = instructions.charAt(current); // Get the current instruction

// A move took us out of the grid bounds, stop checking further instructions

int deltaX = directionMap.get(currentInstruction)[0], deltaY = directionMap.get(currentInstruction)[1];

Check if the new position is still within bounds

Increase the count of valid moves

Return the list of valid moves for each starting point

Map<Character, int[]> directionMap = new HashMap<>(4);

directionMap.put('L', new int[] {0, -1});// Left move

directionMap.put('D', new int[] {1, 0}); // Down move

for (int start = 0; start < instructionsLength; ++start) {</pre>

directionMap.put('U', new int[] {-1, 0});// Up move

directionMap.put('R', new int[] {0, 1}); // Right move

direction_map = {"L": [0, -1], "R": [0, 1], "U": [-1, 0], "D": [1, 0]}

Execute instructions starting at index i and onwards

4. Return ans, which gives us the count of executable instructions for the robot starting from each position in the string s.

So, the output would be [3, 2, 1] for our example, indicating that starting from the first instruction 'R', the robot can execute 3

Python Solution

Initialize the result list

Length of the instruction string

for i in range(num_instructions):

x += row_change

y += col_change

return results

Java Solution

class Solution {

num_instructions = len(instructions)

Iterate over each instruction starting position

for j in range(i, num_instructions):

if $0 \le x \le n$ and $0 \le y \le n$:

valid_moves += 1

Starting position for this sequence of instructions 16 x, y = start_pos 17 19 # Count of valid moves from the current starting instruction 20 valid_moves = 0

35 else: 36 # If out of bounds, no more valid moves, break the loop 37 break 38 39 # Add the count of valid moves for this sequence to the results list results.append(valid_moves) 40 41

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// Check if the move within the bounds of the grid
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                    if (0 \le x + deltaX \& x + deltaX < n \& 0 \le y + deltaY \& y + deltaY < n) {
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C++ Solution

#include <vector>

#include <unordered_map>

using namespace std;

2 #include <string>

6 class Solution {

7 public:

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} else {

break;

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unordered_map<char, vector<int>> directionMap; // map to hold the direction vectors
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            // Direction vectors for L, R, U, D
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            directionMap['L'] = \{0, -1\};
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            directionMap['R'] = \{0, 1\};
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            directionMap['U'] = \{-1, 0\};
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            directionMap['D'] = \{1, 0\};
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            // Iterate over each instruction to see how many steps we can take starting from that instruction
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            for (int i = 0; i < instructionCount; ++i) {</pre>
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                int posX = startPos[0]; // current X position
                int posY = startPos[1]; // current Y position
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                int validSteps = 0; // counter for the number of valid steps we can take for current instruction
23
                // Starting from the i-th instruction, execute instructions until out of bounds
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                for (int j = i; j < instructionCount; ++j) {</pre>
26
                    // Get the move's direction vector (dx, dy) based on the j-th instruction
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                    int dx = directionMap[s[j]][0];
28
                    int dy = directionMap[s[j]][1];
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                    // Check if the new position after the move is within bounds
                    if (0 \le posX + dx \& posX + dx < n \& 0 \le posY + dy \& posY + dy < n) {
31
32
                        // Update position and increment valid step count
33
                        posX += dx;
34
                        posY += dy;
35
                        ++validSteps;
36
                    } else {
37
                        // If the move is out of bounds, break the loop
38
                        break;
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                // Store the number of valid steps for the i-th starting instruction
42
                answer[i] = validSteps;
43
            return answer; // return the vector with the results
44
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46 };
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```

vector<int> answer(instructionCount); // vector to hold the number of valid instructions for each index

// Update the current position based on the instruction. 18 19 if (currentInstruction === 'U') { 20 } else if (currentInstruction === 'D') { 21 22

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Typescript Solution

1 // This function takes a grid size (n), a starting position (startPos),

const instructionsLength = instructions.length;

// Loop through each instruction in the string.

let [currentRow, currentCol] = startPos;

for (let i = 0; i < instructionsLength; i++) {</pre>

let stepCount: number;

currentRow--;

currentRow++;

currentCol--;

currentCol++;

answerArray[i] = stepCount - i;

Therefore, the time complexity is $0(m^2)$.

input string s.

} else {

return answerArray;

break;

const answerArray = new Array(instructionsLength);

2 // and a string of instructions (s) and calculates the number of valid positions

// Initialize the robot's current position to the starting position.

// Iterate over the remaining instructions from the current position.

for (stepCount = i; stepCount < instructionsLength; stepCount++) {</pre>

const currentInstruction = instructions[stepCount];

} else if (currentInstruction === 'L') {

3 // the robot can move to following the instructions starting at each position within the string.

function executeInstructions(gridSize: number, startPos: number[], instructions: string): number[] {

// If the new position is out of bounds, stop processing further instructions.

// Return the array containing the number of valid steps for each starting position in the string.

// The number of valid steps taken is the difference between the current

// step count and the starting step index for this round of instructions.

Time and Space Complexity The provided Python code implements a function that calculates the number of valid instructions from each starting point in the string s. For calculation, it simulates the movement on an $n \times n$ grid based on the instructions given in s.

• Time Complexity: The outer loop runs m times, where m is the length of the string s. An inner loop runs for each character in s

starting at the position i, where i is the index of the outer loop. In the worst case, the inner loop runs m times (when i is 0), and

as i increases, the number of iterations decreases. Hence, the number of total operations can be approximated by the sum of

if (currentRow === -1 || currentRow === gridSize || currentCol === -1 || currentCol === gridSize) {

the first m natural numbers. This sum is (m*(m+1))/2 which has a complexity of $0(m^2)$. The precise sum of operations for the inner loop is: 1 m + (m - 1) + (m - 2) + ... + 2 + 1 = m*(m+1)/2

 Space Complexity: The space complexity is determined by the storage used by the algorithm besides the input. Here, the variables used (ans, mp, x, y, a, b, t) do not grow with the size of the input n or the length of the string s; they either are constants or only store single values. The ans list, however, grows with the length of s since an entry is added for each character in s.

Thus, space complexity is 0(m) because the only variable that scales with input size is the ans list, with m being the length of the

Problem Description In this problem, we are given a scenario where we have an n x n grid, and a robot is situated at a starting position on this grid. The