2410. Maximum Matching of Players With Trainers

Sorting

Two Pointers

Problem Description

Greedy Array

represents the ability of a particular player, and each entry in the trainers array represents the training capacity of a particular trainer. A player can be matched with a trainer if the player's ability is less than or equal to the trainer's capacity. It's important to note that each player can only be matched with one trainer, and each trainer can match with only one player. The task is to find the maximum number of such player-trainer pairings where the conditions of matching ability to capacity are met.

In this LeetCode problem, you are given two 0-indexed integer arrays: players and trainers. Each entry in the players array

Intuition

Medium

pairing of less able players with trainers who have just enough capacity to train them. This ensures that we do not 'waste' a trainer with high capacity on a player with low ability when that trainer could be matched with a more able player instead. To achieve this, we sort both the players and trainers arrays. Sorting helps us easily compare the least able player with the

The intuition behind the solution is straightforward: we want to pair as many players with trainers as we can, prioritizing the

least capable trainer and move up their respective arrays. If a match is found, both the player and the trainer are effectively removed from the potential pool by moving to the next elements in the arrays (incrementing the index). If no match is found, we move up the trainers array to find the next trainer with a higher capacity that might match the player's ability. The result is incrementally built by adding a match whenever we find a capable trainer for a player. We stop when we've either paired all players or run out of trainers. The sum of matches made gives us the maximum number of possible pairings.

Solution Approach

The implementation provided uses a simple but effective greedy approach, heavily relying on sorting. The steps involve:

Sorting the players array in ascending order, which ensures we start with the player with the lowest ability.

Sorting the trainers array in ascending order, which allows us to start with the trainer having the lowest capacity.

- The algorithm works with two pointers, one (i) iterating through the players array, and the other (j) through the trainers array. Here is the implementation broken down:
- Initialize a variable ans to 0 to keep track of the number of matches made.

Initialize the trainer index pointer j to 0. • Iterate over each player p in the sorted players list using a for-loop.

player p. We increment ans to count the match and also increment j to move to the next trainer.

Algorithm and Data Structures:

train them.

Within the loop, proceed with an inner while-loop which continues as long as j is less than the length of trainers and the current trainer's capacity trainers[j] is less than the ability of the player p. The purpose of this loop is to find the first

Once all players have been considered, return the value of ans.

- trainer in the sorted list whose capacity is sufficient to train the player.
- Lists: The main data structure used here are lists (players and trainers), which are sorted.

they cannot be paired with any trainers before that one in the sorted list either.

Two pointers: It uses two pointers to pair players with trainers without revisiting or re-comparing them, thus optimizing the process.

Sorting: By sorting the arrays, the solution leverages the property that once a player cannot be paired with a current trainer,

Greedy approach: The algorithm uses a greedy method by matching each player with the "smallest" available trainer that can

After the while-loop, if j is still within the bounds of the trainers array, it means a suitable trainer has been found for the

- Patterns used: • Sorting and Two-pointers: This is a common pattern for efficiently pairing elements from two different sorted lists.
- **Example Walkthrough** Let's say we have the following players and trainers arrays:

• trainers: [1, 2, 5]

• players: [2, 3, 4]

We first sort both the players and trainers arrays: After sorting,

We initialize ans to 0. This variable will keep track of the successful player-trainer matches.

We check the trainers in order: trainer 1 cannot train the player because the trainer's capacity is too low.

Move to the next trainer (2), and we find a match. Trainer 2 can train the player with ability 2.

• The trainer in position 2 in the sorted trainers list has a capacity of 5, which is sufficient.

players: [2, 3, 4] (already sorted)

- We also initialize the trainer index pointer j to 0.
- Player with ability 2 is the first in the players array.

Player with ability 4:

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successfully.

Python

Player with ability 3 is the next.

Following the solution approach:

trainers: [1, 2, 5] (already sorted)

 We match this player with the trainer. • We increment ans to 2, and since there are no more trainers, we move on to the final player.

We increment ans to 1 and move to the next trainer (j becomes 2).

Now, let's go through each player and try to find a trainer match:

There are no more trainers to compare since we exceeded the length of the trainers array.

from typing import List # Import List from typing module for type hints

Initialize the count of matched players and trainers

If there is a trainer that can match the current player

Initialize the pointer for trainers list

if trainer_index < len(trainers):</pre>

Increment the match count

#include <algorithm> // Include the necessary header for std::sort

// Iterate through each player

for (int playerStrength : players) {

// Function to match players with trainers based on their strength.

int matches = 0; // Initialize the number of matches to zero

// If a suitable trainer is found, make the match

return matches; // Return the total number of matches made

if (trainersIndex < trainers.size()) {</pre>

matches++; // Increment match count

int matchPlayersAndTrainers(vector<int>& players, vector<int>& trainers) {

// Find a trainer that can match the current player's strength

// Each player can only be matched with a trainer that is equal or greater in strength.

std::sort(players.begin(), players.end()); // Sort the players in ascending order

int trainersIndex = 0; // Initialize the trainers index to the start of the array

std::sort(trainers.begin(), trainers.end()); // Sort the trainers in ascending order

while (trainersIndex < trainers.size() && trainers[trainersIndex] < playerStrength) {</pre>

trainersIndex++; // Move to the next trainer for the following players

trainersIndex++; // Increment trainers index if the current trainer is weaker than the player

trainer_index += 1

- Solution Implementation
 - class Solution: def matchPlayersAndTrainers(self, players: List[int], trainers: List[int]) -> int: # Sort the list of players in ascending order players.sort() # Sort the list of trainers in ascending order

At the end of this process, the value of ans is 2, which signifies that we were able to match two pairs of players and trainers

Iterate over each player in the sorted list for player in players: # Skip trainers which have less capacity than the current player's strength while trainer_index < len(trainers) and trainers[trainer_index] < player:</pre>

trainers.sort()

match_count = 0

trainer_index = 0

```
match_count += 1
                # Move to the next trainer for the next player
                trainer_index += 1
       # Return the total number of matches
       return match_count
Java
class Solution {
    public int matchPlayersAndTrainers(int[] players, int[] trainers) {
       // Sort the players array in ascending order
       Arrays.sort(players);
       // Sort the trainers array in ascending order
       Arrays.sort(trainers);
       // Initialize the count of matches to 0
        int matches = 0;
       // Initialize the index for trainers to 0
       int trainerIndex = 0;
       // Iterate over each player
        for (int player: players) {
           // Increment the trainerIndex until we find a trainer that can match the player
           while (trainerIndex < trainers.length && trainers[trainerIndex] < player) {</pre>
                trainerIndex++;
           // If there is a trainer that can match the player, increment the match count
            if (trainerIndex < trainers.length) {</pre>
                matches++; // A match is found
                trainerIndex++; // Move to the next trainer.
       // Return the total count of matches
       return matches;
```

TypeScript

};

C++

public:

#include <vector>

class Solution {

```
// Import the necessary module for sorting
  import { sort } from 'some-sorting-module'; // Please replace 'some-sorting-module' with the actual module you would use for sort
  // Function to sort an array in ascending order
  function sortArrayAscending(array: number[]): number[] {
      return sort(array, (a, b) => a - b);
  // Function to match players with trainers based on their strength.
  // Each player can only be matched with a trainer that is equal or greater in strength.
  function matchPlayersAndTrainers(players: number[], trainers: number[]): number {
      // Sort the players and trainers in ascending order
      const sortedPlayers: number[] = sortArrayAscending(players);
      const sortedTrainers: number[] = sortArrayAscending(trainers);
      let matches: number = 0; // Initialize the number of matches to zero
      let trainersIndex: number = 0; // Initialize the trainers index to the start of the array
      // Iterate through each player and match with trainers based on strength
      for (const playerStrength of sortedPlayers) {
          // Find a trainer that can match the current player's strength
          while (trainersIndex < sortedTrainers.length && sortedTrainers[trainersIndex] < playerStrength) {</pre>
              trainersIndex++; // Increment trainers index if the current trainer is weaker than the player
          // If a suitable trainer is found, make the match
          if (trainersIndex < sortedTrainers.length) {</pre>
              matches++; // Increment match count
              trainersIndex++; // Move to the next trainer for the following players
      return matches; // Return the total number of matches made
  // Export the function if this is part of a module
  export { matchPlayersAndTrainers };
from typing import List # Import List from typing module for type hints
class Solution:
   def matchPlayersAndTrainers(self, players: List[int], trainers: List[int]) -> int:
       # Sort the list of players in ascending order
        players.sort()
       # Sort the list of trainers in ascending order
        trainers.sort()
       # Initialize the count of matched players and trainers
       match_count = 0
       # Initialize the pointer for trainers list
        trainer index = 0
       # Iterate over each player in the sorted list
```

Return the total number of matches return match_count

Time and Space Complexity

for player in players:

trainer_index += 1

match_count += 1

trainer_index += 1

if trainer_index < len(trainers):</pre>

Increment the match count

The time complexity of the code is determined by the sorting operations and the subsequent for-loop with the inner while loop.

Time Complexity

Sorting both the players and trainers list takes O(nlogn) and O(mlogm), respectively, where n is the number of players and m is the number of trainers.

Skip trainers which have less capacity than the current player's strength

while trainer_index < len(trainers) and trainers[trainer_index] < player:</pre>

If there is a trainer that can match the current player

Move to the next trainer for the next player

- The for-loop iterates over each player, which is O(n). Within this loop, the while loop progresses without resetting, which makes it linear over m elements of trainers in total. In the worst case, all players are checked against all trainers, hence it contributes a maximum of O(m) time complexity.
- The combined time complexity of these operations would be 0(nlogn) + 0(mlogm) + 0(n + m). However, since 0(nlogn) and O(mlogm) are the dominant terms, the overall time complexity simplifies to O(nlogn + mlogm).

The space complexity of the code is 0(1), which is the additional space required. The sorting happens in place, and no additional

Space Complexity

data structures are used aside from a few variables for keeping track of indexes and the answer.