1219. Path with Maximum Gold

Medium Array Backtracking Matrix Leetcode Link

### **Problem Description** In this challenge, you are given a two-dimensional grid that represents a gold mine, with each cell containing a certain amount of

gold or being empty (indicated by a 0). The objective is to collect as much gold as possible with the following constraints:

- You can collect gold from any cell and start from there. You may move one step at a time in any of the four cardinal directions (left, right, up, or down). Once you leave a cell, you cannot enter it again; you collect the gold once only.
- You cannot step into a cell with zero gold. You can end your collection route at any cell containing gold.
- Your task is to determine the maximum amount of gold that can be collected following these rules.
- To solve this problem, we need to explore every possible path that maximizes gold collection, which is a classic Depth-First Search

## keep track of the gold collected along each path and then backtracking, which means restoring the state before the DFS call, to

Intuition

The process involves exploring a cell, collecting its gold, and then recursively checking all other paths from its adjacent cells. We

Here's the approach: 1. Iterate through each cell in the grid. 2. If a cell has gold, initiate a DFS search from that cell.

(DFS) scenario. The intuition behind using DFS is that we want to explore all potential routes from each cell containing gold.

5. After exploring all directions from the current cell, backtrack by restoring the cell's gold value. 6. Keep track of the maximum gold collected during each DFS call.

- 7. Once all cells have been explored, return the maximum gold collected.

ensure that we can explore other paths from the current cell.

This approach exhaustively searches all possible paths and ensures that we find the maximum amount of gold that can be collected in the mine.

3. Mark the cell as visited by setting its value to 0 to prevent revisiting, and collect the gold.

4. Explore all four directions recursively, summing up the gold from subsequent cells.

- Solution Approach
- grid. Here's a breakdown of how the DFS is implemented in the provided solution: A helper function dfs(i, j) is defined to perform the DFS search from the cell at row i and column j. • The function first checks if the current cell (i, j) is out of bounds or if it contains 0 gold, in which case it returns 0 since no gold

The implementation uses Depth-First Search (DFS), a recursive algorithm that exhaustively searches through all possible paths in the

# gold from these recursive calls.

search from that cell.

Example Walkthrough

maximum is updated accordingly.

The input grid itself, which is a 2-dimensional list.

can be collected.

as visited to avoid revisiting it in this path. The function then recursively explores the four adjacent cells (left, right, up, and down) by calling dfs(i + a, j + b) for each direction and adds the result to the current gold value. The max() function is used to choose the path that yields the maximum

After exploring all directions, the function backtracks by resetting the gold value of the current cell (grid[i][j] = t) to ensure

If the cell contains gold, the function temporarily sets the gold amount to zero (grid[i][j] = 0). This serves as marking the cell

- that future DFS calls can visit this cell from a different starting point. • The main function iterates over every cell in the grid and calls dfs(i, j) if the cell contains gold, thereby initiating the DFS
- The built-in max() function is applied across all cells to find the single highest gold value collected. The data structures used in this approach are:

Each call to dfs(i, j) will return the maximum amount of gold that can be collected starting from that cell, and the overall

By using recursive DFS, the solution is able to explore all paths optimally without the need for additional data structures to track visited cells or paths, since the grid is modified in place and later restored during backtracking.

3. From here, we can move right or down. Moving right, we find 2 gold. We collect it and mark the cell as visited, updating the grid:

4. Since we can't move into a cell with zero gold, we backtrack and restore the gold value of the cell (0,1). Our grid is back to:

6. Now, we cannot move any further from cell (1,0), so we finish this path and record that we've collected 7 gold as our current

Since there's only one other cell with gold we haven't started from (1,0), we move there and repeat the DFS process.

8. We try moving from cell (1,0) with 3 gold, but since the adjacent cells with gold have already been visited in previous iterations,

9. Our record shows the maximum gold collected was 7 gold from the path that started at (0,0), moved to (1,0), and then to (0,1).

So, the maximum amount of gold that can be collected in this example, following the DFS approach we outlined in the solution, is 7.

7. After that, we backtrack, restoring the grid, and continue the process for the remaining cells. The grid is restored to its previous

1 4 2 2 3 0

Temporary variables for storing the current gold amount and for carrying out the DFS.

Let's illustrate the solution approach using a small example of a gold mine grid:

We want to collect the maximum amount of gold following the given rules.

Here is a step-by-step walkthrough of the solution implemented with DFS:

1. We start our search with the cell at the top-left corner (0,0) which contains 4 gold. We initiate DFS from this cell. 2. We collect the 4 gold from this cell and mark it as visited by setting its value to 0. Our grid now looks like this:

The total gold collected is 4.

The total gold now is 4 + 3 = 7.

there's no further gold to collect.

def get\_maximum\_gold(self, grid):

return 0

gold = grid[x][y]

grid[x][y] = gold

return total\_gold

total\_gold = gold + max(

rows, cols = len(grid), len(grid[0])

grid[x][y] = 0

def dfs(x, y):

1 0 2 2 3 0

1 0 0

2 3 0

1 0 2

1 0 2

2 0 0

maximum.

class Solution:

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Now we have 4 + 2 = 6 gold.

And we explore the next direction, which is down, finding 3 gold in cell (1,0).

5. We repeat the process for cell (1,0): collect the 3 gold, mark the cell as visited, and update the grid:

state: 1 4 2 2 3 0

Performs a depth-first search to find the maximum gold that can

:param grid: List[List[int]] representing the grid of cells with gold.

Depth-first search helper function starting from cell (x, y).

if not  $(0 \le x \le rows and 0 \le y \le cols and grid[x][y])$ :

# Calculate the gold gathered in all four directions

dfs(x + dx, y + dy) for dx, dy in directions

# Initialize the number of rows and columns of the grid

directions = [(0, 1), (0, -1), (-1, 0), (1, 0)]

private int[][] grid; // Holds the grid with gold amounts

// Computes the maximum gold that can be collected

public int getMaximumGold(int[][] grid) {

// Number of rows in the grid

// Number of columns in the grid

// Directions array representing the 4 possible movements: up, right, down, left.

// Compute the maximum gold recursively starting from the current cell.

if  $(x < 0 \mid | x >= grid.size() \mid | y < 0 \mid | y >= grid[0].size() \mid | grid[x][y] == 0) {$ 

// Mark the current cell as visited by setting it to 0.

std::vector<int> directions =  $\{-1, 0, 1, 0, -1\}$ ;

// Iterate over each cell in the grid.

for (int i = 0; i < grid.size(); ++i) {</pre>

// Main function to call for getting the maximum gold.

int getMaximumGold(std::vector<std::vector<int>>& grid) {

for (int j = 0; j < grid[0].size(); ++j) {</pre>

maxGold = std::max(maxGold, dfs(i, j, grid));

// Helper DFS function to explore the grid for gold collection.

int gold = grid[x][y]; // Store the gold in the current cell.

// Recursively collect the gold by going in one direction.

// Update the maximum gold collected from the current path.

grid[x][y] = gold; // Reset the cell to its original gold value.

\* Retrieves the maximum amount of gold by visiting cells in the grid.

\* Each cell can be visited only once, and we must avoid cells with 0 gold.

// Temporarily set the current cell to 0 to mark as visited

\* @param {number[][]} grid - The grid of cells containing gold amounts.

\* @return {number} - The maximum amount of gold that can be collected.

function getMaximumGold(grid: number[][]): number {

function dfs(row: number, col: number): number {

const currentGold: number = grid[row][col];

const directions: number[] = [-1, 0, 1, 0, -1];

const nextRow: number = row + directions[k];

// Accumulate the max gold along this path

const nextCol: number = col + directions[k + 1];

// Reset the current cell's value after exploring all paths

// Initialize the answer to 0; it will store the maximum gold found

// Start from each cell in the grid to find the maximum gold

const rowCount: number = grid.length;

return 0;

grid[row][col] = 0;

const colCount: number = grid[0].length;

let maxGoldFromHere: number = 0;

// Explore in all four directions

// Directions: up, right, down, left

for (let k: number = 0; k < 4; ++k) {

maxGoldFromHere = std::max(maxGoldFromHere, gold + nextGold);

int nextGold = dfs(x + directions[i], y + directions[i + 1], grid);

return maxGoldFromHere; // Return the maximum gold obtained starting from (x, y).

// The depth-first search function explores all possible paths from the current cell.

// Boundary check and gold check, return 0 if it's out of bounds or cell is empty

if (row < 0 || row >= rowCount || col < 0 || col >= colCount || grid[row][col] === 0) {

maxGoldFromHere = Math.max(maxGoldFromHere, currentGold + dfs(nextRow, nextCol));

int dfs(int x, int y, std::vector<std::vector<int>>& grid) {

// Boundary check and cell with zero gold check

The idea is to try all possible paths and get the maximum gold.

:param x: int representing the x-coordinate of the current cell.

:return: int representing the maximum amount of gold that can be collected.

be collected starting from any point in the grid.

- Python Solution
- 17 :param y: int representing the y-coordinate of the current cell. :return: int representing the collected gold amount via the path. 18 19 20 # Check if the current position is out of bounds or has no gold

# Store the gold at the current position and mark this cell as visited by setting the gold to 0

# Backtrack and reset the gold to original since we might visit this cell again via a different path

# Define all four possible directions to move on the grid, which are right, left, up, and down

#### 43 # Using a grid comprehension, perform depth-first search starting from each cell # and return the maximum gold found in all the paths 44 return max(dfs(x, y) for x in range(rows) for y in range(cols)) 45 46 # The modified class and function names now follow PEP 8 naming conventions.

Java Solution

class Solution {

class Solution {

int maxGold = 0;

return maxGold;

return 0;

grid[x][y] = 0;

int maxGoldFromHere = 0;

// Explore all 4 adjacent cells.

for (int i = 0; i < 4; ++i) {

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private:

private int rows;

private int cols;

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// Initialize rows and cols based on the input grid size
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             rows = grid.length;
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             cols = grid[0].length;
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             this.grid = grid; // Assign the grid
             int maxGold = 0; // Initialize maximum gold collected
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             // Iterate over all cells of the grid
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             for (int i = 0; i < rows; ++i) {
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                 for (int j = 0; j < cols; ++j) {
                     // Update maxGold with the maximum of the current maxGold or the gold collected by DFS from this cell
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                     maxGold = Math.max(maxGold, dfs(i, j));
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             return maxGold; // Return the maximum gold collected
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         // Helper method for depth-first search
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         private int dfs(int row, int col) {
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             // Base case: If the cell is out of the grid bounds or has 0 gold, return 0
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             if (row < 0 || row >= rows || col < 0 || col >= cols || grid[row][col] == 0) {
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                 return 0;
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             // Store gold value of the current cell
             int gold = grid[row][col];
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             // Mark the current cell as visited by setting gold to 0
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             grid[row][col] = 0;
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             // Array to facilitate iteration over the adjacent cells (up, right, down, left)
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             int[] directions = \{-1, 0, 1, 0, -1\};
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             int maxGold = 0; // Initialize max gold collected from this cell
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             // Iterate over all adjacent cells
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             for (int k = 0; k < 4; ++k) {
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                 // Calculate the gold collected by DFS of the adjacent cells and update maxGold accordingly
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                 maxGold = Math.max(maxGold, gold + dfs(row + directions[k], col + directions[k + 1]));
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             // Backtrack: reset the value of the cell to the original gold amount
             grid[row][col] = gold;
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             // Return the maximum gold that can be collected from this cell
             return maxGold;
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C++ Solution
  1 #include <vector>
    #include <algorithm> // include algorithm library for using max function
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**Typescript Solution** 

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        for (let i: number = 0; i < rowCount; ++i) {
            for (let j: number = 0; j < colCount; ++j) {</pre>
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                maxGold = Math.max(maxGold, dfs(i, j));
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```

grid[row][col] = currentGold;

return maxGoldFromHere;

let maxGold: number = 0;

Time and Space Complexity

return maxGold;

Let's denote m as the number of rows and n as the number of columns in the grid. In the worst case, the entire grid could be filled with positive gold values, requiring an exhaustive search from each cell.

The time complexity of the provided code is calculated based on several factors:

4 directions to explore next. Therefore, an upper bound of the time complexity is 0(4^(m\*n)), which is exponential due to the recursive exploration with

The algorithm performs a depth-first search (DFS) starting from each cell that has a positive gold value.

• Each DFS explores four possible directions to go next (except when on the border or if a cell is already visited).

branching. For space complexity:

• The space used by the call stack during the recursive DFS would be equal to the maximum depth of the recursion which, in the worst case, is O(m\*n) since we might have to traverse the entire grid if it is all filled with gold. • We are also modifying the input grid to mark the cells as visited by temporarily setting grid[i][j] = 0 and resetting it back to its

For each DFS, the maximum depth could be min(m\*n) if the path is allowed to traverse every cell. However, due to path restrictions

(not revisiting a cell with no gold), the maximum depth will be less than or equal to m\*n. Each time we explore a cell, there are at most

- - original value before returning from the DFS. This in-place modification means we don't use extra space proportional to the grid size (except the implicit recursion stack). Thus, the space complexity is 0(m\*n).