2323. Find Minimum Time to Finish All Jobs II

Medium Greedy Array Sorting

Problem Description

You are provided with two integer arrays, jobs and workers, each indexed starting from 0. The arrays have the same lengths, meaning they contain an equal number of elements. jobs[i] represents the time required to complete the i-th job, while workers[j] represents the time a j-th worker can dedicate to working each day. The task is to assign each job to one worker so that every worker gets exactly one job. Your goal is to calculate the minimum number of days necessary to finish all the jobs after they have been allocated to the workers.

The intuition behind the solution involves two observations:

Intuition

Sorting Pattern: If we sort both the jobs and workers arrays in ascending order, we ensure the biggest jobs are matched

- with the workers who can work the most per day. This helps us in minimizing the overall time since we are using the maximum capacity of the most efficient workers. Minimum Days Calculation: After sorting, we pair each job with a worker. To calculate the number of days a worker needs to
- Since we are dealing with whole days, we need to round up, which is achieved by (a + b 1) // b in Python (integer division with ceiling). The max function applied to the result of these calculations across all job-worker pairs ensures that we find the scenario that takes the longest. This is the minimum number of days needed to complete all jobs, since all jobs are being worked on

complete a job, we can take the time needed for the job (a) and divide it by the time the worker can put in each day (b).

simultaneously, and once the job taking the longest is done, all other jobs would have also been completed. Solution Approach

The solution uses a greedy algorithm approach which is implemented through sorting and pairing, leading to an efficient way to

Firstly, we sort both the jobs and workers arrays in ascending order. By doing this, we can pair the largest job with the

worker who can work the longest hours, and so on down the line. Sorting is a common algorithmic pattern for problems related to optimization, as it often helps to align data in a way that simplifies pairing or comparison.

calculate the minimum time necessary to complete all jobs. Here's a walkthrough of the implementation:

using Python's built-in zip() function, which pairs elements with the same index from each of the input sequences. For each pair (a, b) taken from zip(jobs, workers), we calculate the number of days it would take for worker b to complete job a. This is where the formula (a + b - 1) // b comes into play. It ensures that we perform the integer division

We then iterate over the pairs of sorted jobs and workers. These pairs are generated by zipping the two arrays together

with a ceiling effect since we are dealing with full days. If a job doesn't require a full extra day, we still count the partial day as

parallel, the overall completion time for all jobs is dictated by the job-worker pair that takes the longest. This longest time will

- a full one, since a worker cannot be partially assigned to a job each day. The max() function is then used on the sequence of these calculations. Because each worker is working on different jobs in
- be the maximum value out of the individual times calculated. In summary, by sorting the jobs and workers, we facilitate an optimal pairing strategy. Then we calculate the time each optimal job-worker pair takes, ensuring to account for partial days as full. Lastly, we identify the longest duration out of these pairs as our

Data Structures used: Arrays: to store the input jobs and workers data and manipulate it through sorting. • Integers: to represent individual job times and worker capacities, and to calculate and represent the minimum number of days required.

Calculation: to determine the number of days for each pair and identify the maximum completion time.

Patterns used:

solution.

Example Walkthrough

Here's how we implement our solution:

paired = zip([2, 3, 7], [1, 3, 5])

paired = [(2, 1), (3, 3), (7, 5)]

Sorting: to optimally align the jobs to workers.

Let's walk through a small example to illustrate the solution approach. Consider the following arrays for jobs and workers: jobs = [3, 2, 7]

• Greedy: to pair the highest capacity workers with the largest jobs to minimize overall time.

Iteration: to go through the pairs of jobs and workers and calculate the time taken for each.

workers = [1, 5, 3]

First, we sort both arrays to align the largest jobs with the workers who can dedicate the most time per day. After sorting:

iobs = [2, 3, 7]

workers = [1, 3, 5]

Next, we pair each job with a worker by zipping the two sorted arrays.

```
Now, for each paired tuple, calculate the number of days it takes for the worker to finish the job using (a + b - 1) // b.
Here, a is the time needed for the job, and b is the time the worker can work per day.
```

Last, we find the maximum number of days from all the pairs; this will be the minimum number of days required to complete

def minimumTime(self, jobs: List[int], workers: List[int]) -> int:

the workers who take the least amount of time per unit.

Sort the lists to ensure that the shortest jobs are matched with

The -1 is used to ensure that we do not add an extra unnecessary

time unit if iob duration is perfectly divisible by worker speed.

For each pair of job and worker, the time required for that worker

// This function calculates the minimum time required to assign jobs to workers

int minimumTime(vector<int>& jobs, vector<int>& workers) {

// Iterate over the matched pairs of jobs and workers.

// Import the necessary functions from lodash for sorting and max operations

sort(jobs.begin(), jobs.end());

return minimumTimeRequired;

sort(workers.begin(), workers.end());

for (int i = 0; i < jobs.size(); ++i) {</pre>

// such that each job is assigned to one worker, and the time taken is minimized.

// The time taken for each job-worker pair is the ceil of division of job by worker.

int minimumTimeRequired = 0; // Initialize the minimum time required as 0.

// Sort the jobs and workers vectors to match the smallest job with the fastest worker.

// Return the maximum time out of all job-worker pairs as that would be the bottleneck.

// Calculate the time taken for each job-worker pair, which is jobs[i] divided by workers[i],

// rounded up to the nearest integer, and update minimumTimeRequired if this time is larger.

minimumTimeRequired = max(minimumTimeRequired, (jobs[i] + workers[i] - 1) / workers[i]);

 \circ For the third pair (7, 5), the calculation is (7 + 5 - 1) // 5 = 2 days.

 \circ For the first pair (2, 1), the calculation is (2 + 1 - 1) // 1 = 2 days.

 \circ For the second pair (3, 3), the calculation is (3 + 3 - 1) // 3 = 1 day.

```
all the jobs.
\max_{days} = \max([2, 1, 2]) = 2
```

Python from typing import List

Calculate the time taken for each worker to complete their job. # We're using (iob duration + worker speed - 1) // worker speed to # determine the minimum number of full time units needed for completion.

So, with the given job and worker assignments, the minimum number of days required to finish all the jobs is 2 days.

```
times = [(job_duration + worker_speed - 1) // worker_speed for job_duration, worker_speed in zip(jobs, workers)]
# The maximum time among all the workers is what will determine
# the total time required to finish all jobs, since we have to wait for
# the slowest iob-worker pair to finish.
```

return max(times)

to complete their iob is calculated.

import java.util.Arrays; // Necessary import for Arrays.sort()

Solution Implementation

class Solution:

Java

iobs.sort()

workers.sort()

```
class Solution {
    // Method to calculate the minimum time needed to assign jobs to workers such that each worker gets exactly one job.
    public int minimumTime(int[] jobs, int[] workers) {
        Arrays.sort(jobs): // Sort the array of jobs in non-decreasing order.
        Arrays.sort(workers); // Sort the array of workers in non-decreasing order.
        int maximumTime = 0; // Initialize maximum time to 0.
        for (int i = 0; i < iobs.length; ++i) {
            // Calculate the time it takes for worker i to complete job i.
            // Divide the iob[i] work units by the workers[i] efficiency, rounding up to the nearest whole number.
            int currentTime = (jobs[i] + workers[i] - 1) / workers[i];
            // Update maximumTime to be the maximum of itself and the time calculated for the current worker.
            maximumTime = Math.max(maximumTime, currentTime);
        // Return the overall maximum time it takes for all jobs to be completed.
        return maximumTime;
C++
#include <vector>
#include <algorithm>
using namespace std;
class Solution {
```

};

TypeScript

public:

```
import { sortBy, max } from 'lodash';
// This function calculates the minimum time required to assign jobs to workers
// such that each iob is assigned to one worker, and the time taken is minimized.
// The time taken for each job-worker pair is the ceil of the division of the job by the worker.
function minimumTime(jobs: number[], workers: number[]): number {
    // Sort the jobs and workers arrays to match the smallest job with the fastest worker.
    jobs = sortBy(jobs);
    workers = sortBy(workers);
    let minimumTimeRequired = 0; // Initialize the minimum time required as 0.
    // Iterate over the matched pairs of jobs and workers.
    for (let i = 0; i < iobs.length; ++i) {
       // Calculate the time taken for each job-worker pair, which is jobs[i] divided by workers[i],
        // rounded up to the nearest integer, by adding workers[i] -1 before the division.
        // Update minimumTimeRequired if this time is larger.
        minimumTimeRequired = max([minimumTimeRequired, Math.ceil(jobs[i] / workers[i])]);
    // Return the maximum time out of all job-worker pairs as that would be the bottleneck.
    return minimumTimeRequired;
// Usage example (if you want to test the function):
// const iobsExample = [3, 2, 10];
// const workersExample = [1, 2, 3];
// const minTime = minimumTime(jobsExample, workersExample);
// console.log(minTime); // Output will be the minimum time required based on the job and worker assignments.
from typing import List
class Solution:
   def minimumTime(self, jobs: List[int], workers: List[int]) -> int:
```

times = $[(job_duration + worker_speed - 1) // worker_speed for job_duration, worker_speed in zip(jobs, workers)]$

Time and Space Complexity **Time Complexity**

return max(times)

iobs.sort()

workers.sort()

2. Sorting the workers list: This is another sorting operation also with O(n log n), assuming workers list has the same length as jobs. Since these are two consecutive operations, the combined time complexity for both sorts will still be 0(n log n) because the

the code is $O(n \log n)$.

constants are dropped in Big O notation. 3. The zip function and the comprehension: Creating pairs of jobs and workers with zip is O(n) and the comprehension iterates over each pair

Sort the lists to ensure that the shortest jobs are matched with

Calculate the time taken for each worker to complete their iob.

We're using (job duration + worker speed - 1) // worker speed to

The -1 is used to ensure that we do not add an extra unnecessary

The maximum time among all the workers is what will determine

time unit if iob duration is perfectly divisible by worker speed.

For each pair of job and worker, the time required for that worker

the total time required to finish all jobs, since we have to wait for

The time complexity of the given code can be broken down into the following parts:

1. Sorting the jobs list: Sorting an array of n items has a time complexity of $O(n \log n)$.

determine the minimum number of full time units needed for completion.

the workers who take the least amount of time per unit.

to complete their job is calculated.

the slowest job-worker pair to finish.

once, making it O(n) as well.

the input size. Therefore, the space complexity for the sort operation is 0(1).

- Combining all the above, we see that the most time-consuming operations are the sorts. Therefore, the overall time complexity of
- **Space Complexity** For space complexity:

Sorting the lists in-place: Since Python's sort method on lists sorts the list in place, it doesn't use extra space proportional to

zip and the comprehension: zip takes 0(1) additional space as it returns an iterator, and the comprehension also takes 0(1) space because it only computes the maximum time without storing intermediate results in an array or list.

Hence, the total space complexity of the code is 0(1).