2251. Number of Flowers in Full Bloom

Hard Array Hash Table Binary Search Ordered Set Prefix Sum Sorting

Problem Description

In this problem, we are dealing with a scenario related to flowers blooming. The input includes two arrays:

1. A 2D array `flowers`, where each sub-array contains two elements indicating the start and end of the full bloom period for a particular flower. The flower blooms inclusively from `start_i` to `end_i`.
2. An array `persons`, where each element represents the time a person arrives to see the flowers.

The goal is to determine how many flowers are in full bloom for each person when they arrive. The output should be an array `answer`, where `answer[i]` corresponds to the number of flowers in full bloom at the time `i` th person arrives.

Intuition

To solve this problem, we can use a two-step strategy involving sorting and binary search:

1. **Sorting**: We separate the start and end times of the bloom periods into two lists and sort them. The sorted start times help us determine how many flowers have started blooming at a given point, and the sorted end times indicate how many flowers have finished blooming.
2. **Binary Search**: When a person arrives, we want to count the flowers that have begun blooming but haven't finished. We use the binary search algorithm to find:
   - The index of the *first* end time that is *strictly greater* than the arrival time of the person, which indicates how many flowers have finished blooming. We can use the number `bisect_left` on the sorted end times.
   - The index of the *first* start time that is *greater than or equal to* the arrival time, which tells us how many flowers have started to bloom. We use `bisect_right` for this on the sorted start times.

By subtracting the number of flowers that have finished blooming from those that have started, we get the count of flowers in full bloom when a person arrives. We repeat this process for each person and compile the results into the final answer array.

Solution Approach

The solution approach uses a combination of sorting and binary search to efficiently determine how many flowers are in full bloom for each person's arrival time. Here's the implementation explained step by step:

1. **Sort Starting and Ending Times**: First, we extract all the start times and end times from the `flowers` array into separate lists and sort them:

```
start = sorted(a for a, _ in flowers)
end = sorted(b for _, b in flowers)
```

   Sorting these lists allows us to use binary search later on. The `start` list will be used to determine how many flowers have started blooming by a certain time, and the `end` list will help determine how many flowers have ended their bloom.

2. **Binary Search for Bloom Count**: The next step is to iterate over each person's arrival time `p` in the `persons` list and find out the count of flowers in bloom at that particular time. For each `p`:

```
bisect_right(start, p) - bisect_left(end, p)
```

   - `bisect_right(start, p)` finds the index in the sorted `start` list where `p` would be inserted to maintain the order. This index represents the count of all flowers that have started blooming up to time `p` (including `p`).
   - `bisect_left(end, p)` finds the index in the sorted `end` list where `p` could be inserted to maintain the order. This index signifies the count of flowers that have not finished blooming by time `p`.

   By subtracting the numbers obtained from `bisect_left` on the `end` list from `bisect_right` on the `start` list, we obtain the total number of flowers in bloom at the arrival time of `p`.

3. **Compile Results**: The above operation is repeated for each person's arrival time, and the results are compiled into the answer list. This list comprehends the count of flowers in bloom for each person, as per their arrival times:

```
return [bisect_right(start, p) - bisect_left(end, p) for p in persons]
```

In the end, the `answer` list is returned, which provides the solution, i.e., the number of flowers in full bloom at the time of each person's arrival.

The algorithms and data structures used here, like sorting and binary search (`bisect` module in Python), enable us to solve the problem in a time-efficient manner, taking advantage of the ordered datasets for quick lookups.

Example Walkthrough

Imagine we have an array of flowers where the blooms are represented as `flowers = [[1,3], [2,5], [4,7]]` and an array of persons with arrival times as `persons = [1, 3, 5]`. We want to find out how many flowers are in full bloom each person sees when they arrive.

First, we need to process the flowers' bloom times. We sort the start times `[1, 2, 4]` and the end times `[3, 5, 7]` of the blooming periods.

Now let's walk through the steps to get the number of flowers in bloom for each person:

1. Person at time `1`:
   - Using `bisect_right` for the sorted start times: `bisect_right([1, 2, 4], 1)` gives us index `1`, indicating one flower has started blooming.
   - Using `bisect_left` for the end times: `bisect_left([3, 5, 7], 1)` gives us index `0`, indicating no flowers have finished blooming.
   - The difference `1 (started) - 0 (ended)` tells us that exactly one flower is in full bloom.
2. Person at time `3`:
   - `bisect_right([1, 2, 4], 3)` results in index `2`, as two flowers have bloomed by time 3.
   - `bisect_left([3, 5, 7], 3)` gives us index `1`, as one flower has stopped blooming.
   - The difference `2 (started) - 1 (ended)` is 1, so one flower is blooming for this person.
3. Person at time `5`:
   - `bisect_right([1, 2, 4], 5)` gives an index of `3` - all three flowers have started blooming by time 5.
   - `bisect_left([3, 5, 7], 5)` yields index `2`, as two flowers have finished blooming strictly before time 5.
   - The difference `3 (started) - 2 (ended)` is 1, indicating that one flower is in bloom when this person arrives.

Thus, for the persons arriving at times `1`, `3`, and `5`, the function will return `[1, 1, 1]` as the number of flowers in full bloom at each of their arrival times.

Solution Implementation

Python

```
from bisect import bisect_right, bisect_left

class Solution:
    def fullBloomFlowers(self, flowers: List[List[int]], persons: List[int]) -> List[int]:
        # Sort the start times and end times of the flowers' blooming periods
        start_times = sorted(start for start, _ in flowers)
        end_times = sorted(end for _, end in flowers)

        # Calculate the number of flowers in full bloom for each person's visit
        bloom_counts = [
            # The total number of flowers that have started blooming by person p's visit time
            bisect_right(start_times, p) -
            # Subtracting the number of flowers that have finished blooming by person p's visit time
            bisect_left(end_times, p)
            for p in persons
        ]

        return bloom_counts

# Example usage:
# sol = Solution()
# print(sol.fullBloomFlowers([[1, 10], [3, 3]], [4, 5]))
```

Java

```
import java.util.Arrays;

public class Solution {
    public int[] fullBloomFlowers(int[][] flowers, int[] people) {
        int flowerCount = flowers.length; // Number of flowers
        int[] bloomStart = new int[flowerCount];
        int[] bloomEnd = new int[flowerCount];

        // Extract the start and end bloom times for each flower into separate arrays
        for (int i = 0; i < flowerCount; ++i) {
            bloomStart[i] = flowers[i][0];
            bloomEnd[i] = flowers[i][1];
        }

        // Sort the start and end bloom times arrays
        Arrays.sort(bloomStart);
        Arrays.sort(bloomEnd);

        int peopleCount = people.length; // Number of people
        int[] answer = new int[peopleCount]; // Array to store the answers

        // For each person, calculate the number of flowers in full bloom
        for (int i = 0; i < peopleCount; ++i) {
            // Number of flowers that have started blooming minus
            // the number of flowers that have already ended blooming
            answer[i] = findInsertionPoint(bloomStart, people[i] + 1) - findInsertionPoint(bloomEnd, people[i]);
        }

        return answer; // Return the array containing the number of flowers in full bloom for each person
    }

    private int findInsertionPoint(int[] times, int value) {
        int left = 0;  // Start of the search range
        int right = times.length;  // End of the search range

        // Binary search to find the insertion point of 'value'
        while (left < right) {
            int mid = (left + right) / 2; // Midpoint of the current search range
            if (times[mid] >= value) {
                right = mid;  // Adjust the search range to the left half
            } else {
                left = mid + 1;  // Adjust the search range to the right half
            }
        }
        return left;  // The insertion point is where we would add 'value' to keep the array sorted
    }
}
```

C++

```
#include <vector>
#include <algorithm>

using namespace std;

class Solution {
public:
    vector<int> fullBloomFlowers(vector<vector<int>>& flowers, vector<int>& people) {
        // Number of flower intervals
        int n = flowers.size();

        // Separate vectors to hold the start and end times for each flower
        vector<int> starts;
        vector<int> ends;

        // Loop over all flowers to populate the start and end vectors
        for (auto& flower : flowers) {
            starts.push_back(flower[0]);
            ends.push_back(flower[1]);
        }

        // Sort the start and end vectors to prepare for binary search
        sort(starts.begin(), starts.end());
        sort(ends.begin(), ends.end());

        // Vector to hold the number of flowers in full bloom for each person
        vector<int> bloomCounts;

        // Loop through each person to determine how many flowers are in full bloom
        for (auto& person : people) {
            // Find the position of the first flower that starts after the person's time (exclusive)
            // This gives us the number of flowers that have started blooming
            auto flowersStarted = upper_bound(starts.begin(), starts.end(), person) - starts.begin();

            // Find the position of the first flower that ends at or before the person's time (inclusive)
            // This gives us the number of flowers that have already ceased blooming
            auto flowersEnded = lower_bound(ends.begin(), ends.end(), person) - ends.begin();

            // Subtract flowersEnded from flowersStarted to get the number of flowers in full bloom
            bloomCounts.push_back(flowersStarted - flowersEnded);
        }

        // Return the counts of flowers in full bloom for each person
        return bloomCounts;
    }
};
```

TypeScript

```
function fullBloomFlowers(flowers: number[][], people: number[]): number[] {
    const flowerCount = flowers.length; // Number of flowers
    // Arrays to store the start and end times of each flower's bloom.
    const bloomStarts = new Array(flowerCount).fill(0);
    const bloomEnds = new Array(flowerCount).fill(0);

    // Split the flowers' bloom times into start and end times.
    for (let i = 0; i < flowerCount; ++i) {
        bloomStarts[i] = flowers[i][0];
        bloomEnds[i] = flowers[i][1];
    }

    // Sort the start and end times.
    bloomStarts.sort((a, b) => a - b);
    bloomEnds.sort((a, b) => a - b);

    // Array to store the result for each person.
    const results: number[] = [];
    for (const person of people) {
        // Find the number of flowers blooming at the time person visits.
        // Start of blooms minus end of blooms before person arrives.
        const flowersBloomingStart = search(bloomStarts, person + 1); // Start of blooms after person
        const flowersBloomingEnd = search(bloomEnds, person); // End of blooms by the time person visits
        results.push(flowersBloomingStart - flowersBloomingEnd); // Number of flowers currently in bloom
    }
    return results;
}

// Binary search helper function to find the index at which a flower's start or end time is greater than or equal to x.
function search(nums: number[], x: number): number {
    let left = 0;
    let right = nums.length;
    while (left < right) {
        const mid = left + ((right - left) >> 1); // Prevents potential overflow
        if (nums[mid] >= x) {
            right = mid; // Look in the left half
        } else {
            left = mid + 1; // Look in the right half
        }
    }
    return left; // Left is the index where nums[left] >= x
}
```

```
from bisect import bisect_right, bisect_left

class Solution:
    def fullBloomFlowers(self, flowers: List[List[int]], persons: List[int]) -> List[int]:
        # Sort the start times and end times of the flowers' blooming periods
        start_times = sorted(start for start, _ in flowers)
        end_times = sorted(end for _, end in flowers)

        # Calculate the number of flowers in full bloom for each person's visit
        bloom_counts = [
            # The total number of flowers that have started blooming by person p's visit time
            bisect_right(start_times, p) -
            # Subtracting the number of flowers that have finished blooming by person p's visit time
            bisect_left(end_times, p)
            for p in persons
        ]

        return bloom_counts

# Example usage:
# sol = Solution()
# print(sol.fullBloomFlowers([[1, 10], [3, 3]], [4, 5]))
```

Time and Space Complexity

Time Complexity

The given code consists of three main parts:

1. Sorting the start times of the flowers: `sorted(a for a, _ in flowers)`
2. Sorting the end times of the flowers: `sorted(b for _, b in flowers)`
3. Iterating through each person and using binary search to find the count of bloomed flowers: `[bisect_right(start, p) - bisect_left(end, p) for p in persons]`

Let's consider n as the number of flowers and m as the number of persons. Here's a breakdown of the time complexity:

- **Sorting the start and end times**: Sorting takes $O(n \log n)$ time for both the start and end lists. Hence the combined sorting time is $2 * O(n \log n)$.
- **Binary search for each person**: For each person, `bisect_right` and `bisect_left` are performed once. These operations have a time complexity of $O(\log n)$. Since these operations are performed for m persons, the total time for this part is $O(m \log n)$.

Adding these up, the overall time complexity of the code is $O(n \log n + m \log n)$.

Space Complexity

The space complexity comes from the additional lists used to store start and end times:

- **Start and end lists**: Two lists are created to store start and end times, each of size n. Hence, the space taken by these lists is $2 * O(n)$.

Therefore, the overall space complexity of the code is $O(n)$.