1375. Number of Times Binary String Is Prefix-Aligned

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Medium <u>Array</u>
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Problem Description

a sequence of steps defined by an integer array flips, we flip the bits of this binary string from 0 to 1. The ith element of flips represents the bit position in the binary string that will be flipped during the ith step of the process.

In this LeetCode problem, we are given an initially zeroed binary string of length n, where the indexes are 1-indexed. Throughout

A binary string is considered prefix-aligned after the ith step if all the bits from the beginning of the string to position i are set

to 1, while the rest of the string remains at 0. The task is to calculate the total number of times the binary string is prefix-aligned during the entire flipping process. For example:

Initial binary string of length 5: 00000 - flips sequence: [3,2,4,1,5]

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After each flip:
- Step 1, flip position 3: 00100 (Not prefix-aligned as the first bit is still 0)
- Step 2, flip position 2: 01100 (Not prefix-aligned as the first two bits are not all 1)
- Step 3, flip position 4: 01110 (Not prefix-aligned as the first four bits are not all 1)
- Step 4, flip position 1: 11110 (Not prefix-aligned as the first four bits are not all 1)
- Step 5, flip position 5: 11111 (Prefix-aligned as all the bits are 1)
Hence, the binary string is prefix-aligned 1 time during the flipping process.
  Understanding the problem is crucial before attempting to create a solution.
Intuition
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The intuition behind the solution is to keep track of the highest bit position (let's call it mx) that has been flipped at each step. For

each step i, we compare mx with the current step index i: if mx is equal to i, it means that all bits up to the current step index have been flipped to 1, and hence the string is prefix-aligned.

Here is how this intuition is applied to the solution: 1. Initialize the counter ans to keep track of the number of times the string is prefix-aligned, and mx to record the highest bit position flipped so far.

3. Update mx to be the maximum between its current value and the bit position x of the current flip. 4. If mx is equal to i after flipping the bit at position x, increment ans by 1 since the string is currently prefix-aligned.

- 5. Continue the loop until all flips are processed. 6. Return ans, the total count of prefix-aligned occurrences.
- Solution Approach

to loop over flips, starting from index 1 because the problem is 1-indexed.

2. Iterate through each flip in the sequence, keeping track of the current step number i (starting from 1).

- The solution provided follows an iterative approach with two primary variables in play: ans and mx. The variable ans serves as a counter for the instances when the binary string is prefix-aligned, and mx keeps track of the maximum index that has been
- flipped so far. 1. Initialize ans to 0. This variable will count the number of times the binary string is prefix-aligned during the flip sequence.
- 2. Initialize mx to 0. This variable represents the maximum position of the flipped bit encountered up to the current step.

x), which ensures mx always reflects the farthest position that was flipped till the current step. 5. The check mx == i will be True if all bits from the start to the current position i are 1 (thus, the string is prefix-aligned). If so, we increment ans by 1. This leverages the fact that a binary string is prefix-aligned if the maximum flipped position at step i is equal to i itself. Any flip sequence that has the largest flip within the range [1, i] at i ensures that all preceding bits are already flipped to 1.

6. Finally, the loop concludes once all the elements in flips have been iterated, and the ans value is returned. ans now contains the number of

4. For each flip, update mx to be the maximum between the current mx and the flip position x. This is achieved by the expression mx = max(mx,

We use a for loop to iterate through each flip in the provided flips list. The loop uses the built-in enumerate function in Python

times the binary string was prefix-aligned during the process. This algorithm doesn't use any complex data structures and requires no additional space besides the two variables ans and mx,

3. During each iteration of the loop, we get two values: i, the step number starting from 1, and x, the current position to flip.

making it space-efficient. The time complexity of the algorithm is O(n) where n is the number of flips since we are going through each flip exactly once. It efficiently solves the problem by keeping track of only the current state necessary to determine prefix alignment at each step without reconstructing the binary string at each instance.

Suppose we have an initially zeroed binary string of length 4 and the following flips sequence: [1,3,2,4]. • Initial binary string: 0000 We iterate through each flip, updating the maximum position mx flipped, and checking after each step if the string is prefixaligned.

Step 1, flip position 1: The binary string after the flip becomes 1000. mx is updated to 1. Since mx (1) equals the step number

New binary string: 1000

o ans still = 1

o ans still = 1

 \circ mx = 4

answer ans is 2.

the number of flips.

 \circ mx = 1

Example Walkthrough

 \circ ans = 1 Step 2, flip position 3: The binary string becomes 1010. The new mx is 3, which is now the highest position flipped so far.

i (1), we increment ans by 1. The string is prefix-aligned.

Let's walk through a small example to illustrate the solution approach.

New binary string: 1010 o mx = 3

Step 3, flip position 2: After this flip, the binary string is 1110, and mx remains 3. As mx (3) is not equal to the step number i

Step 4, flip position 4: The final binary string is 1111. mx is updated to 4, which now equals the step number i (4). We

Since mx (3) does not equal the step number i (2), and is not incremented. The string is not prefix-aligned.

(3), ans remains the same. The string is **not prefix-aligned**. New binary string: 1110 • mx still = 3

increment ans by 1 again because the string is prefix-aligned.

- New binary string: 1111
- o ans = 1 + 1 = 2 At the end of the flipping process, we have encountered 2 instances where the binary string is prefix-aligned, which means our

Using the solution approach, we keep track of the maximum flipped position and check for prefix alignment in each step

efficiently. The variables ans and mx enable us to do this without needing additional data structures or performing complex

operations. This makes the algorithm space-efficient and straightforward to implement with a time complexity of O(n), where n is

Initialize the number of moments when all bulbs are blue and the current maximum flipped bulb position.

If the maximum flipped position equals the current moment index, all the bulbs are blue.

Update the maximum flipped position if the current flip position is greater.

- `moments all blue`: The count of times when all lights up to the current moment (included) turned blue.

max_flipped_position = max(max_flipped_position, flip_position)

- `flips`: The list of light bulb positions to flip during each moment, starting from 1.

Python # This class contains a method to determine the number of times all light bulbs are blue. class Solution: def numTimesAllBlue(self, flips: List[int]) -> int:

moments all blue = 0

max_flipped_position = 0

return moments_all_blue

Loop through each flip by index and value.

for moment, flip position in enumerate(flips, 1):

moments_all_blue += max_flipped_position == moment

Return the total count of moments when all the bulbs are blue.

Solution Implementation

It's important to import the `List` typing from the `typing` module to ensure the type hint is recognized by the Python interpreter: ```python from typing import List

In the code above:

```
Java
class Solution {
    public int numTimesAllBlue(int[] flips) {
        int numMomentsAllBlue = 0; // This will store the number of moments when all bulbs are blue
        int maxTurnedOnBulb = 0; // This will keep track of the highest numbered bulb that has been turned on
        // Iterate through the flips array. Each flip represents turning on the bulb at that index.
        for (int moment = 1; moment <= flips.length; ++moment) {</pre>
            // Update the maxTurnedOnBulb with the maximum value between the current max and the bulb flipped at this moment
            maxTurnedOnBulb = Math.max(maxTurnedOnBulb, flips[moment - 1]);
            // If the maximum turned—on bulb number equals the current moment, increment the numMomentsAllBlue counter
            // It means all bulbs up to that point are blue
            if (maxTurnedOnBulb == moment) {
                ++numMomentsAllBlue;
        // Return the total number of moments when all flipped—on bulbs are blue
        return numMomentsAllBlue;
```

- `max flipped position`: The maximum position (index) among the flipped light bulbs. If the maximum position we have flipped so far

- The `enumerate` function is used to loop through each flip with its corresponding moment, beginning with 1. The `enumerate` function

};

C++

public:

#include <vector>

class Solution {

#include <algorithm> // For using the max function

int numTimesAllBlue(vector<int>& flips) {

return countBlueMoments;

// Loop through each flip in the flips vector

Loop through each flip by index and value.

for moment, flip position in enumerate(flips, 1):

moments_all_blue += max_flipped_position == moment

Return the total count of moments when all the bulbs are blue.

Update the maximum flipped position if the current flip position is greater.

- `moments all blue`: The count of times when all lights up to the current moment (included) turned blue.

If the maximum flipped position equals the current moment index, all the bulbs are blue.

max_flipped_position = max(max_flipped_position, flip_position)

- `flips`: The list of light bulb positions to flip during each moment, starting from 1.

maxFlipped = max(maxFlipped, flips[i - 1]);

for (int i = 1; i <= flips.size(); ++i) {</pre>

// Function to count the number of moments when all bulbs are blue

int countBlueMoments = 0; // Initialize a counter for the blue moments

// it means all bulbs up to that point are on (and hence blue)

// Return the total count of moments when all turned-on bulbs are blue

int maxFlipped = 0; // This will keep track of the maximum bulb number flipped so far

// Update the maximum flipped bulb number if the current flip is greater

// If the maximum flipped bulb number equals the number of flips so far,

countBlueMoments += (maxFlipped == i) ? 1 : 0; // Use the ternary operator for the condition

```
TypeScript
function numTimesAllBlue(lightSwitches: number[]): number {
    let countAllBlue = 0; // This will hold the number of times all lights turned blue
    let maxSwitchedOn = 0; // Tracks the maximum light switch number turned on so far
    // Loop through each light switch flip
    for (let i = 1; i <= lightSwitches.length; ++i) {</pre>
        // Update maxSwitchedOn to be the highest of the current max or the current light switch flipped
        maxSwitchedOn = Math.max(maxSwitchedOn, lightSwitches[i - 1]);
        // If maxSwitchedOn is equal to the number of flips so far, all lights are blue
        if (maxSwitchedOn === i) {
            countAllBlue += 1;
    // Return the total number of times all lights have turned blue
    return countAllBlue;
# This class contains a method to determine the number of times all light bulbs are blue.
class Solution:
    def numTimesAllBlue(self, flips: List[int]) -> int:
        # Initialize the number of moments when all bulbs are blue and the current maximum flipped bulb position.
        moments all blue = 0
        max_flipped_position = 0
```

Time and Space Complexity

return moments_all_blue

In the code above:

from typing import List

Time Complexity

```python

The given algorithm consists of a single for loop that iterates through every element in the flips list exactly once. The enumeration of flips does not change the overall time complexity. Inside the loop, the algorithm performs a constant amount of work on each iteration, including a comparison and an assignment. Therefore, the time complexity of this function is O(n) where n is the number of elements in flips.

- `max flipped position`: The maximum position (index) among the flipped light bulbs. If the maximum position we have flipped so far

It's important to import the `List` typing from the `typing` module to ensure the type hint is recognized by the Python interpreter:

- The `enumerate` function is used to loop through each flip with its corresponding moment, beginning with 1. The `enumerate` function

## **Space Complexity**

The space complexity of the algorithm is defined by the amount of additional memory used by the algorithm as a function of the input size. Within the provided algorithm, and and mx are the only variables that occupy extra space, and their memory footprint does not depend on the size of the input. As a result, the space complexity is 0(1) because the amount of extra memory used is constant, irrespective of the input size.