

1122. Relative Sort Array

Easy Array Hash Table Counting Sort Sorting

[Leetcode Link](#)

Problem Description

In this problem, we are provided with two integer arrays, `arr1` and `arr2`. The elements within `arr2` are distinct, which means no element repeats itself, and all these elements are present within `arr1`. Our goal is to sort `arr1` so that the order of elements matches how they appear in `arr2`, while any elements that are not found in `arr2` should be appended to the end of `arr1` in ascending order.

For example, if we have `arr1 = [2,3,1,3,2,4,6,7,9,2,19]` and `arr2 = [2,1,4,3,9,6]`, the first step is to arrange the elements of `arr1` that are also in `arr2` in the order they appear in `arr2`: `[2,2,2,1,4,3,3,9,6]`. After that, we take the remaining elements that are not in `arr2` (which are 7, 19, and 7 in this case) and place them at the end in sorted order: `[7,19,7]`. The final output array would be the concatenation of these two processes, giving us `[2,2,2,1,4,3,3,9,6,7,7,19]`.

The challenge lies in figuring out a strategy that allows us to sort `arr1` with these constraints in mind.

Intuition

The key to solving this problem efficiently is to recognize that we can use a custom sorting strategy. The default sort mechanisms most programming languages provide can often be customized with a user-defined key function that determines the sort order.

Given that the elements in `arr2` are the first to be sorted and they are in a specific order, we need to have this order play a significant role in our custom sorting strategy. For elements that are in `arr2`, we want them to appear first and in the order they are found in `arr2`.

To achieve this, we can map each element of `arr2` to its position (or index) within `arr2`. This can be done using a dictionary or a hash map that takes the element as a key and the position of that element in `arr2` as the value.

In Python, this mapping can be created as follows: `pos = {x: i for i, x in enumerate(arr2)}`

For the elements of `arr1` that are not found in `arr2`, we need to ensure they come after any element that is in `arr2`. One way to do this is to assign them an index that is larger than any index that would be assigned to the elements that are in `arr2`. Since indices in Python are zero-based and `arr2` can have, at most, the same number of elements as `arr1`, all the elements not in `arr2` can be given an index based on a large number (for example, 1000) plus their value, which ensures they will be sorted in ascending order after the elements of `arr2`.

Now, using Python's `sorted` function, we can sort `arr1` with a custom key. This key is defined by a lambda function that takes an element `x` and returns its corresponding index from the dictionary `pos` if it's in `arr2`, or `1000 + x` otherwise.

Here's the implementation in Python, which reflects our approach:

```
1 sorted(arr1, key=lambda x: pos.get(x, 1000 + x))
```

By using this custom sorting key, we align `arr1` elements per `arr2`'s order and then append the remaining elements in ascending order right after. The overall runtime complexity of the solution is $O(N \log N)$, where `N` is the length of `arr1`.

Solution Approach

The solution is implemented using a dictionary for direct mapping and the sorting algorithm that Python provides. Here's a step-by-step explanation, referencing the provided code:

- Create a Mapping of `arr2`:** The solution starts by creating a dictionary called `pos`. This dictionary maps each element `x` of `arr2` to its index `i` in the array. This is done using a dictionary comprehension:

```
1 pos = {x: i for i, x in enumerate(arr2)}
```

By doing this, we are setting up a quick reference so that we can look up the index of any element from `arr2` instantly. This index acts as a custom priority value for the sorting algorithm.

- Custom Sort with a Lambda Function:** We need to sort `arr1`, but not by its elements' natural order. We will use Python's `sorted` function, which allows us to specify a custom key function:

```
1 sorted(arr1, key=lambda x: pos.get(x, 1000 + x))
```

The `key` parameter is crucial here. It takes a lambda function which determines how the elements of `arr1` should be compared during sorting:

- `lambda x:` is an anonymous function that takes `x`, an element from `arr1`.
- `pos.get(x, 1000 + x)` is the function's body that returns a value for each `x` that will be used to compare during sorting.
- If `x` is found in `arr2` (and, as a result, the `pos` dictionary), `pos.get(x)` returns the index of `x` from `arr2`.
- If `x` is not found in `arr2`, the method `pos.get(x, 1000 + x)` provides a default value which is the sum of 1000 and the element `x` itself. This ensures that:
 - The elements not in `arr2` are ordered at the end because their sorting key is greater than any index assigned from `arr2`.
 - They are sorted in ascending order amongst themselves because if two elements are not found in `arr2`, their sorting key is simply their value (`1000 + their own value` respects the natural ascending order).

By using this sorting strategy, we can maintain the relative ordering based on another array `arr2` while also logically appending and sorting elements not found in `arr2`.

Thus, using a combination of dictionary for direct access and sorting algorithm with a custom key function, this solution efficiently achieves the required array manipulation adhering to the conditions of the problem statement.

Example Walkthrough

Let's illustrate the solution approach using a smaller example. Suppose `arr1` contains `[8, 4, 5, 4, 6]` and `arr2` contains `[4, 6]`. We want to sort `arr1` such that the order of `arr2` is preserved for common elements, and the rest are sorted in ascending order at the end.

Step 1: Create a Mapping for `arr2` Elements:

First, we create a dictionary that maps each element of `arr2` to its index:

```
1 pos = {x: i for i, x in enumerate(arr2)}
2 # pos will be {4: 0, 6: 1}
```

Here, 4 is at index 0 in `arr2` and 6 is at index 1.

Step 2: Sort `arr1` with a Custom Sort Key:

We now sort `arr1` using Python's `sorted` function with the custom key:

```
1 sorted(arr1, key=lambda x: pos.get(x, 1000 + x))
```

Breaking down how the key function works for each element in `arr1`:

- 8 is not found in `arr2`, so the key would be `1000 + 8 = 1008`.
- The first 4 is found in `arr2` and its index is 0; hence, the key would be 0.
- 5 is not found in `arr2`, so its key is `1000 + 5 = 1005`.
- The second 4 is found and its key is again 0.
- 6 is found and its index is 1, so the key is 1.

The sorted order using these keys would be `[4, 4, 6, 5, 8]`.

Step 3: Review the Final Sorted Array:

So, the elements from `arr1` that were also in `arr2` are now sorted in `arr2` order: 4 comes first, followed by 6. The remaining elements, not in `arr2`, follow at the end in ascending order: 5 and then 8.

In conclusion, after applying the custom key sorting strategy to `arr1`, we get the final sorted array: `[4, 4, 6, 5, 8]`, successfully matching the required conditions.

Python Solution

```
1 class Solution:
2     def relativeSortArray(self, arr1: List[int], arr2: List[int]) -> List[int]:
3         # Create a dictionary to map elements of arr2 to their indices.
4         position_map = {value: index for index, value in enumerate(arr2)}
5
6         # Sort arr1 based on the condition that elements of arr2 should come first
7         # in the order they appear in arr2, followed by the remaining elements in
8         # ascending order.
9         sorted_arr1 = sorted(arr1, key=lambda x: position_map.get(x, 1000 + x))
10
11        # The lambda function inside sorted uses the 'get' method to find the position
12        # of x from position_map if x exists in arr2. Otherwise, it assigns a number
13        # greater than 1000 to ensure that x is positioned after all elements of arr2.
14        # This is based on the assumption that the elements of arr1 do not exceed 1000.
15
16        return sorted_arr1
17
```

Java Solution

```
1 class Solution {
2     public int[] relativeSortArray(int[] arr1, int[] arr2) {
3         // Create a hashmap to store the positions of each element in arr2
4         Map<Integer, Integer> elementToIndexMap = new HashMap<>(arr2.length);
5         // Fill the map with the positions of the elements
6         for (int index = 0; index < arr2.length; ++index) {
7             elementToIndexMap.put(arr2[index], index);
8         }
9
10        // Create a new 2D array to hold elements and their corresponding positions
11        int[][] elementPositionPairs = new int[arr1.length][2];
12        for (int index = 0; index < elementPositionPairs.length; ++index) {
13            // Use the position from arr2 if present, or else use the value from arr1 plus the length of arr2
14            elementPositionPairs[index] = new int[]{arr1[index], elementToIndexMap.getOrDefault(arr1[index], arr2.length + arr1[index])};
15        }
16
17        // Sort the 2D array based on the positions
18        Arrays.sort(elementPositionPairs, (pair1, pair2) -> pair1[1] - pair2[1]);
19
20        // Place the sorted elements back into arr1
21        for (int index = 0; index < elementPositionPairs.length; ++index) {
22            arr1[index] = elementPositionPairs[index][0];
23        }
24
25        // Return the sorted arr1
26        return arr1;
27    }
28 }
29
```

C++ Solution

```
1 #include <vector>
2 #include <unordered_map>
3 #include <algorithm>
4
5 class Solution {
6 public:
7     // Function to sort arr1 with respect to the order in arr2
8     std::vector<int> relativeSortArray(std::vector<int>& arr1, std::vector<int>& arr2) {
9         // We will use an unordered_map to store the position (index) of
10        // each element in arr2; this position is used when sorting arr1
11        std::unordered_map<int, int> elementToIndex;
12        for (int i = 0; i < arr2.size(); ++i) {
13            elementToIndex[arr2[i]] = i; // Map the element to its index in arr2
14        }
15
16        // Vector to temporarily store pairs of position and value
17        // from original array arr1. If the element does not exist in arr2,
18        // we assign the index as the size of arr2 which is effectively
19        // putting it at the end of the sorted array.
20        std::vector<std::pair<int, int>> tempArray;
21        for (int value : arr1) {
22            // Find the index if value exists in arr2; for elements not in arr2, set index to arr2's size
23            int index = elementToIndex.count(value) ? elementToIndex[value] : arr2.size();
24            // Create a pair of (index, value) for each element in arr1
25            tempArray.emplace_back(index, value);
26        }
27
28        // Sort the tempArray based on the previously stored positions (indices)
29        // If two elements have the same position, they're compared by their values
30        std::sort(tempArray.begin(), tempArray.end());
31
32        // Reassign sorted values back to arr1
33        for (int i = 0; i < arr1.size(); ++i) {
34            arr1[i] = tempArray[i].second; // Set arr1[i] with the value from the sorted pair
35        }
36
37        // Return the sorted arr1
38        return arr1;
39    };
40 };
41
```

Typescript Solution

```
1 // Given two arrays arr1 and arr2, the function sorts elements of arr1
2 // such that the relative ordering of items in arr1 are the same as in arr2.
3 // Elements not present in arr2 will be placed at the end of arr1 in ascending order.
4 function relativeSortArray(arr1: number[], arr2: number[]): number[] {
5     // Create a mapping of element to its position for quick access
6     const elementToIndex: Map<number, number> = new Map();
7     // Fill the map with elements of arr2 and their corresponding indices
8     for (let index = 0; index < arr2.length; ++index) {
9         elementToIndex.set(arr2[index], index);
10    }
11
12    // Initialize an array to hold pairs of position in arr2 and the value
13    const sortedPairs: [number, number][] = [];
14
15    // Map each element of arr1 to its corresponding index in arr2, with a fallback index
16    for (const element of arr1) {
17        // Get the index position of the element from the map or use default if not present
18        const indexInArr2 = elementToIndex.get(element) ?? arr2.length;
19        sortedPairs.push([indexInArr2, element]);
20    }
21
22    // Sort the array: first by the index position and then by the element values
23    sortedPairs.sort((a, b) => a[0] - b[0] || a[1] - b[1]);
24
25    // Map the sorted array back to a single array of values
26    return sortedPairs.map(pair => pair[1]);
27 }
28
```

Time and Space Complexity

Time Complexity

The time complexity of the given code primarily depends on the sorting algorithm used by the `sorted` function in Python.

- Creating the `pos` dictionary has a time complexity of $O(M)$ where `M` is the length of `arr2`, since we iterate over each element of `arr2` once.
- The `sorted` function has a time complexity of $O(N \log N)$ where `N` is the length of `arr1`. This is because Python uses the TimSort algorithm for sorting, which is a hybrid sorting algorithm derived from merge sort and insertion sort.

However, since the key function in the sorted algorithm uses a lookup in the `pos` dictionary, which is $O(1)$ time complexity for each element, together with the condition to handle elements not in `pos`, the overall time complexity of the sorting with these operations is still $O(N \log N)$.

Therefore, the total time complexity of the algorithm is $O(M + N \log N)$.

Space Complexity

The space complexity of the algorithm consists of the space needed for the `pos` dictionary and any additional space used by the sorting algorithm:

- The `pos` dictionary has a space complexity of $O(M)$ because it contains as many entries as there are elements in `arr2`.
- The `sorted` function returns a new list and internally uses additional space which depends on the specifics of the implementation. In the worst case, this could be $O(N)$ space.

The total space complexity is therefore $O(M + N)$.