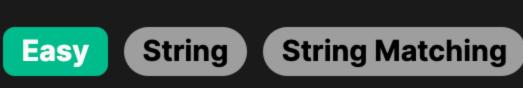
1668. Maximum Repeating Substring



Problem Description

The problem involves finding the highest number of times (k) a word can repeat itself as a contiguous subsequence within a larger string. For example, if the sequence is "ababc" and the word is "ab", then the word "ab" repeats once as a substring. However, it does not repeat twice in a row, so it cannot be considered 2-repeating. The task is to find the maximum k-repeating value, which means we want to find the highest number of consecutive repetitions of word within sequence.

If word does not appear in sequence even once, then the maximum k-repeating value is 0. The goal is to return this maximum k value.

Intuition

the number of times word can fit into sequence, which is len(sequence) // len(word). We start checking from the maximum possible repetition (k value) and work backwards. In each iteration, we construct a string by repeating word k times and check if that construction is a substring of sequence. If it is, that means word is k-repeating in sequence, and we have found our maximum k-repeating value.

The intuition behind the solution is a straightforward search approach. We know that the maximum k value cannot be more than

The given solution method maxRepeating defines a simple approach, utilizing a control structure to find the maximum k-

Solution Approach

repeating value. It uses a for loop and string multiplication operations to achieve this. Here's a step-by-step overview of the algorithm:

First, we need to establish the upper bound of our search for k. The most number of times word can fit into sequence is found through integer division (//). This is calculated by len(sequence) // len(word). It gives us the maximum number of

times word can be repeated before it would exceed the length of sequence.

- We then use a for loop to iterate from this maximum possible value of k down to 0. The reason for iterating backwards is to find the maximum k-repeating value first. If we were to iterate forwards, we would need to check every possible value of k
- upto the maximum, which is less efficient. In each iteration of the loop, we create a new string by multiplying word by k. The multiplication operator * used on strings in Python creates a new string by repeating the operand string k times.
- if the first operand is found within the second operand string. If the condition word * k in sequence is true, the loop breaks and returns the current value of k, which is the maximum k-

We then check if this newly created string is a substring of sequence by using the in operator in Python, which returns True

If the loop completes without finding any k value for which word * k is a substring of sequence, the function implicitly returns None, which is not relevant to our problem statement since we expect an integer.

The choice of using this specific pattern is due to its simplicity and effectiveness for the given problem. No additional data

structures are needed, and the control flow along with built-in string operations is sufficient to arrive at the correct result. With this approach, we minimize the number of checks we need to perform and ensure that we can return the highest k as soon

Example Walkthrough

Let's go through an example to illustrate the solution approach. For this example, let's take the sequence as

Establish the upper bound: According to step 1, we find the maximum number of times word can fit into sequence. We get

"aabbabbaabbaabb" and the word as "abb".

with k decremented by 1.

as it is found.

repeating value.

len(sequence) // len(word) which is 18 // 3 = 6. Thus, the word "abb" can be repeated at most 6 times in a row without exceeding the length of the sequence.

- Iterate from the maximum possible value of k down to 0: We start a for loop from 6 down to 1. String multiplication and substring check: In the first iteration, k is 6, so we multiply "abb" by 6. We check if "abbabbabbabbabbabb" (word * k) is a substring of the sequence "aabbabbaabbaabbaabb". It isn't, so we continue the loop
- Repeat step 3 with k = 5, 4, 3, ..., constructing the string "abb" * k each time and checking if it's a substring of "aabbabbabbaabb". When k = 3, we build "abbabbabb" (which is "abb" repeated 3 times) and check if it's a substring of the sequence. We find
- that "aabbabbaabbaabb" does contain "abbabbabb". Thus, word does k-repeat for k = 3. The function then breaks from the loop and returns 3, as this is the highest k value for which word * k is a subsequence of the sequence.

Using this approach, we efficiently found the highest number of contiguous subsequences of "abb" within the larger string

Solution Implementation

"aabbabbaabbaabb" by starting from the largest possible repetition (k) and reducing it until we found a match.

:param sequence: The string in which to search for repeating 'word'. :param word: The word to look for in 'sequence'. :return: The maximum number of times 'word' can be repeated.

in 'sequence' as a substring.

def max_repeating(self, sequence: str, word: str) -> int:

max possible repeats = len(sequence) // len(word)

Find the maximum number of times 'word' can be consecutively repeated

Calculate the maximum possible repetitions of 'word' within 'sequence'

// This function finds the maximum number of times 'word' can be repeated

int maxCount = 0; // Stores the maximum count of repeating 'word'

// Loop repeats until 'word' can no longer fit into 'sequence'

if (sequence.find(repeatedWord) != string::npos) {

string repeatedWord = word: // Starts with one 'word' and we'll append more

// Check if the current 'repeatedWord' is a substring of 'sequence'

maxCount = k; // Update the maxCount to the current repetition number

// consecutively in the string 'sequence'.

int maxRepeating(string sequence, string word) {

for (int k = 1; k <= possibleRepeats; ++k) {</pre>

Python

class Solution:

```
# Iterate over the possible repetitions, starting from the most and descending
        for k in range(max possible repeats, -1, -1):
            # If the repeat sequence of 'word' is within 'sequence', return that repeat count
            if word * k in sequence:
                return k
        # If no repetition of 'word' is found, we return 0
        return 0 # Technically this line is not necessary due to the loop's range
Java
class Solution {
    // Defines a method to find the maximum number of times 'word' repeats in 'sequence'
    public int maxRepeating(String sequence, String word) {
        // Start from the maximum possible repetitions and decrement
        for (int k = sequence.length() / word.length(); k > 0; --k) {
            String repeatedWord = word.repeat(k); // Construct the word repeated 'k' times
            if (sequence.contains(repeatedWord)) { // Check if 'sequence' contains the repeated 'word'
                return k; // If found, return the current repetition count 'k'
        return 0; // If no repetition is found, return 0
C++
```

#include <string>

class Solution {

public:

using namespace std;

} else {

```
// If the 'repeatedWord' is not in 'sequence', break out of the loop
                break;
            // Append 'word' to 'repeatedWord' for the next iteration to check for the next number of repeats
            repeatedWord += word;
        // Return the maxCount, which is the maximum number of times 'word' repeats consecutively in 'sequence'
        return maxCount;
TypeScript
/**
 * Find the maximum number of times the word can be repeated consecutively in the sequence.
 * @param {string} sequence - The string to search within.
 * @param {string} word - The word to look for in the sequence.
 * @returns {number} The maximum number of times the word is repeated.
function maxRepeating(sequence: string, word: string): number {
    // Determine the lengths of the sequence and the word
    let sequenceLength: number = sequence.length;
    let wordLength: number = word.length;
    // Start from the maximum possible repetition (sequence length divided by word length)
    // and go down to find the maximum repeating consecutive occurrences
    for (let repeatCount: number = Math.floor(sequenceLength / wordLength); repeatCount > 0; repeatCount--) {
        // Generate the word repeated 'repeatCount' times
        let repeatedWord: string = word.repeat(repeatCount);
        // Check if the sequence includes the repeated word
        if (sequence.includes(repeatedWord)) {
            // If the sequence contains the word repeated 'repeatCount' times, return 'repeatCount' as the result
            return repeatCount;
```

int possibleRepeats = sequence.size() / word.size(); // Calculates the maximum possible times 'word' could repeat

```
// If no repetition is found, return 0
   return 0;
class Solution:
   def max_repeating(self, sequence: str, word: str) -> int:
       Find the maximum number of times 'word' can be consecutively repeated
       in 'sequence' as a substring.
        :param sequence: The string in which to search for repeating 'word'.
        :param word: The word to look for in 'sequence'.
       :return: The maximum number of times 'word' can be repeated.
       # Calculate the maximum possible repetitions of 'word' within 'sequence'
       max_possible_repeats = len(sequence) // len(word)
       # Iterate over the possible repetitions, starting from the most and descending
       for k in range(max possible repeats, -1, -1):
           # If the repeat sequence of 'word' is within 'sequence', return that repeat count
           if word * k in sequence:
               return k
       # If no repetition of 'word' is found, we return 0
       return 0 # Technically this line is not necessary due to the loop's range
```

Time and Space Complexity **Time Complexity:**

The time complexity of checking if a string is a substring of another string is O(n), where n is the length of the string. Here, it is checking a substring of maximum length len(word) * k in a string of length len(sequence). This needs to be done for each k

The given code finds the maximum number of times the word word can be repeated in the sequence string. It checks for every k

from len(sequence) // len(word) down to 0 to see if word * k is a substring of the sequence.

grows linearly with the size of this generated string, which is at most the length of sequence itself.

is len(word). This simplifies to $O((n^2/m) * (n/m)) = O(n^3/m^2)$.

from len(sequence) // len(word) to 0. Therefore, the time complexity can be approximated as O((n/m) * (n + n - m + n - 2m + ... + m)), where n is len(sequence) and m

Space Complexity: The space complexity of the code is O(n), with n being the maximum length of word * k which can be generated for the comparison with sequence. Essentially, the space required is for the substring that is created during word * k. The memory used