1616. Split Two Strings to Make Palindrome

Medium **Two Pointers** String

Problem Description

into two substrings each. For string a, this produces a_prefix and a_suffix, and for string b, this results in b_prefix and b_suffix. After the split, the goal is to check whether concatenating a_prefix with b_suffix, or b_prefix with a_suffix, results in a palindrome. A palindrome is a string that reads the same forward as it does backward, like "racecar" or "madam".

You are provided with two strings a and b that are of equal length. Your task is to pick an index at which you will split both a and b

being an empty string and a_suffix being the entire string a, or vice versa. The challenge is to return true if you can create a

palindrome by doing such a split and concatenation, or false otherwise.

Keep in mind that when you split a string into a prefix and suffix, either part can be empty. So, you can end up with a_prefix

Intuition

1. If either a or b is already a palindrome, you don't need to split them; the answer is true. 2. If the beginning of a matches the end of b (and vice versa) until a certain point, then there's a chance that the remainder of the string is a

palindrome by itself; we must check this.

Let's try to simplify the problem with the following insights:

- 3. If we can't form a palindrome from the start, perhaps we could if we split the strings later, so we can check from the opposite direction as well; start from the ends and move towards the middle.
- Now, to the solution approach:
- Check if the current prefix of a and suffix of b can form a palindrome: Iterate through string a from the start and string b from the end simultaneously. As soon as the characters at the current indices don't match, stop.

Check the inner substrings: When the characters don't match, you have to check if the inner substring a[i:j+1] or b[i:j+1] is a palindrome. If either of these substrings is a palindrome, it means that the whole string can be a palindrome, because the

- previous characters on both ends were matching. Check the opposite case: Repeat steps 1 and 2, but this time try to match the beginning of b and the end of a.
- Combine the checks: If any check returns true, a palindrome can be formed. This approach works because starting from the ends and moving towards the middle ensures that by the time a mismatch occurs, if the remaining unchecked substring is a palindrome itself, the previous matching parts would make the whole string a
- palindrome. This holds true for both beginning from the start and checking towards the end, as well as the opposite. **Solution Approach**

The solution makes use of two helper functions check1 and check2 to determine if a palindrome can be formed by combining the

The check1 function is called twice, first with a and b in their original order and then with their order switched. This ensures

starting from the beginning of a and the other (j) starting from the end of b. This continues as long as the characters at these

that we check for palindromes by making a split before and also after the midpoint of the strings. Inside check1, a two-pointer approach is used. This approach involves iterating over the strings with two indices: one (i)

prefixes and suffixes of two strings a and b.

The implementation involves the following steps:

part, left unchecked by check1, forms a palindrome.

b_suffix or b_prefix + a_suffix forms a palindrome.

Step 1: Check for Palindrome Using the Two-Pointer Approach

Step 2: Invoking check2 for non-matching characters

• The checks from check1 and check2 have all returned false.

def checkPalindromeFormation(self, a: str, b: str) -> bool:

def is_palindrome_substring(s: str, start: int, end: int) -> bool:

Compare the substring with its reverse to check for a palindrome

// Move towards the middle from both ends if corresponding characters match

// Check if we have reached the middle or if either half is a palindrome

return left >= right || isPalindrome(prefix, left, right) || isPalindrome(suffix, left, right);

// Helper method to check if the substring of a is a palindrome within the given interval [left, right]

// Helper function to check if switching between 'a' and 'b' strings can form a palindrome.

Check if it is possible to create a palindrome by taking a prefix from one string and

Loop until the characters at i and j are equal and i is less than j

def is_palindrome_substring(s: str, start: int, end: int) -> bool:

Compare the substring with its reverse to check for a palindrome

Check if the substring s[start:end+1] is a palindrome.

Check if the strings a and b can form a palindrome by switching prefix and suffix at any point.

If we've scanned the entire string and it's a palindrome, or if the substring a[i:j+1]

return i >= j or is_palindrome_substring(a, i, j) or is_palindrome_substring(b, i, j)

Return True if either combination of a prefix from a and suffix from b, or a prefix from b

while (left < right && prefix[left] == suffix[right]) {</pre>

// Continue checking for palindrome within the interval

bool isPalindrome(string& str, int left, int right) {

while (left < right && str[left] == str[right]) {</pre>

// If we reached the middle, then it's a palindrome

function checkPalindromeFormation(a: string, b: string): boolean {

function isPalindromeAfterSwitch(a: string, b: string): boolean {

++left;

++left;

};

TypeScript

--right;

return left >= right;

let leftIndex = 0;

let rightIndex = b.length - 1;

--right;

Return True if either combination of a prefix from a and suffix from b, or a prefix from b

Check if the substring s[start:end+1] is a palindrome.

return s[start:end+1] == s[start:end+1][::-1]

return check_palindrome(a, b) or check_palindrome(b, a)

public boolean checkPalindromeFormation(String a, String b) {

return checkCombination(a, b) || checkCombination(b, a);

// Check both combinations, a with b and b with a

and suffix from a, forms a palindrome

def check_palindrome(a: str, b: str) -> bool:

a suffix from the other string.

A couple of scenarios can lead to a successful palindrome formation:

indices are equal. If they're not equal, this is where we might have a chance to split and form a palindrome, and we then move on to step 3.

Once a mismatch is found, the check2 function comes into play. It takes the substring from a or b starting at index i and

ending at j, and checks if it is a palindrome by comparing it to its reverse (a[i : j + 1][::-1]). This checks if the middle

• The looping in check1 proceeds until i is greater than or equal to j, meaning all characters matched properly, so the strings are already palindromic considering the parts we are checking. • The check2 function returns true, which means that even though there was a mismatch, the remaining unmatched part is a palindrome itself, thereby allowing for a successful palindrome formation when the outer matched parts are included.

Finally, the checkPalindromeFormation function returns true if either call to check1 returns true. Otherwise, it returns false,

indicating that no palindromic combination could be found. The solution efficiently uses the two-pointer technique to iterate and compare characters, which is a common pattern used in

palindrome cannot be formed, leading to an efficient algorithm both in terms of time and space complexity.

string manipulation algorithms. The method of checking the remaining inner substring separately helps to stop early when a

Example Walkthrough Let us walk through an example to illustrate the solution approach.

Assume we are given the strings a = "xayb" and b = "abzx". The goal is to determine if there is a split index where $a_prefix +$

We begin by using check1 to compare a and the reverse of b using a two-pointer technique. Pointers i and j start at the start of a and the end of b respectively. • We compare a[i] with b[j]. For i=0 and j=3, a[0] is 'x' and b[3] is 'x' as well, they match, so we continue.

• Increment i and decrement j, and compare a[1] with b[2]. We have a[1] is 'a' and b[2] is 'z'. They do not match, so we proceed to step 3.

Now we call check2 for substring a[i : j + 1], which is "ay". • Check if "ay" is a palindrome by comparing it to its reverse "ya". It is not, so we go back to check1 and also invoke check2 on b[i : j + 1] which

Step 3: Check the Opposite Case

Step 4: Combine the checks

described.

Python

Java

class Solution {

class Solution:

is "bz".

• Now we switch a and b and compare b and the reverse of a using a two-pointer technique initiated by calling check1 again.

Check if "bz" is a palindrome by comparing it to its reverse "zb". It is not a palindrome either.

• Calling check2 for b[0:4], checking if "abzx" is a palindrome, which it is not. • Similarly, calling check2 for a [0:4], checking if "xayb" is a palindrome, which it is not.

• This time b[0] is 'a' and the reversed a[3] is 'b', they do not match. Directly proceed to check the inner substrings.

Solution Implementation

Check if it is possible to create a palindrome by taking a prefix from one string and

• This means that in this example, we cannot form a palindrome from any split index by concatenating substrings from a and b.

Thus, the final result returned by the checkPalindromeFormation function for strings a = "xayb" and b = "abzx" would be false.

No palindromic combination is possible with the given strings following the checks performed according to the method

i, j = 0, len(b) - 1# Loop until the characters at i and j are equal and i is less than j while i < j and a[i] == b[j]:</pre> i, j = i + 1, j - 1# If we've scanned the entire string and it's a palindrome, or if the substring a[i:j+1] # or b[i:j+1] is a palindrome, return True return i >= j or is_palindrome_substring(a, i, j) or is_palindrome_substring(b, i, j)

// Main method to check if a palindrome can be formed by replacing a prefix of one string with a suffix of another

Check if the strings a and b can form a palindrome by switching prefix and suffix at any point.

```
// Helper method to check if a palindrome can be formed with a particular combination
    private boolean checkCombination(String prefixString, String suffixString) {
        int startIndex = 0;
        int endIndex = suffixString.length() - 1;
       // Move from both ends towards the center comparing characters of prefixString and suffixString
       while (startIndex < endIndex && prefixString.charAt(startIndex) == suffixString.charAt(endIndex)) {</pre>
            startIndex++;
            endIndex--;
       // If indices have crossed, the resulting string is already a palindrome
        // or check if substring of prefixString or suffixString from startIndex to endIndex is a palindrome
        return startIndex >= endIndex || isSubStrPalindrome(prefixString, startIndex, endIndex) || isSubStrPalindrome(suffixString)
    // Method to check if the substring from startIndex to endIndex is a palindrome in string str
    private boolean isSubStrPalindrome(String str, int startIndex, int endIndex) {
        while (startIndex < endIndex && str.charAt(startIndex) == str.charAt(endIndex)) {</pre>
            startIndex++;
            endIndex--;
       // If indices have crossed, substring is a palindrome
        return startIndex >= endIndex;
C++
class Solution {
public:
   // Public method to check if a palindrome can be formed from two strings
    bool checkPalindromeFormation(string a, string b) {
       // It checks both combinations: a|b and b|a
        return checkPalindrome(a, b) || checkPalindrome(b, a);
private:
   // Helper method to check if a palindrome can be formed by prefixing and suffixing substrings from two strings
    bool checkPalindrome(string& prefix, string& suffix) {
        int left = 0;
        int right = suffix.size() - 1;
```

```
// Check from the outside towards the center if characters match
          while (leftIndex < rightIndex && a.charAt(leftIndex) === b.charAt(rightIndex)) {</pre>
              leftIndex++;
              rightIndex--;
          // If true, the substring is already a palindrome or can form a palindrome with the rest of 'a' or 'b'
          return leftIndex >= rightIndex || isPalindromeSubstring(a, leftIndex, rightIndex) || isPalindromeSubstring(b, leftIndex,
      // Helper function to check if a substring of 'a' is a palindrome.
      function isPalindromeSubstring(a: string, leftIndex: number, rightIndex: number): boolean {
          // Check from the outside towards the center if characters match
          while (leftIndex < rightIndex && a.charAt(leftIndex) === a.charAt(rightIndex)) {</pre>
              leftIndex++;
              rightIndex--;
          // If true, then the substring from 'leftIndex' to 'rightIndex' is a palindrome
          return leftIndex >= rightIndex;
      // Check both combinations of strings to see if a palindrome can be formed.
      return isPalindromeAfterSwitch(a, b) || isPalindromeAfterSwitch(b, a);
class Solution:
```

```
Time and Space Complexity
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return check_palindrome(a, b) or check_palindrome(b, a)

return s[start:end+1] == s[start:end+1][::-1]

def checkPalindromeFormation(self, a: str, b: str) -> bool:

def check_palindrome(a: str, b: str) -> bool:

or b[i:j+1] is a palindrome, return True

a suffix from the other string.

i, j = 0, len(b) - 1

while i < j and a[i] == b[j]:</pre>

and suffix from a, forms a palindrome

i, j = i + 1, j - 1

string. If the characters match, it moves the pointers inwards. When they stop matching, it calls check2 to check if the substring between the pointers in either string is a palindrome. check2 checks if a substring is a palindrome by comparing the substring to its reverse.

1. The while loop in check1 runs in O(n) time in the worst case, where n is the length of the strings, since each character is compared at most

The worst-case overall time complexity is the sum of these operations, which is O(n) + O(n) + O(n) + O(n), simplifying to O(n).

check1 is a function that uses two-pointers to compare the characters from the beginning of one string and the end of the other

The given Python code defines a function checkPalindromeFormation which takes two strings a and b and checks if a palindrome

can be formed by taking a prefix of one string and the suffix of the other string. There are two helper functions, check1 and

2. The check2 function, when called, takes O(k) time, where k is the length of the substring being checked. In the worst case, k is n (the entire string), so check2 has a worst-case time complexity of O(n). 3. The slicing a[i:j+1] and reversing a[i:j+1][::-1] operations in check2 both take 0(n) time.

4. The checkPalindromeFormation function calls check1 twice, once with (a, b) and once with (b, a).

Space Complexity:

check2.

Time Complexity:

once.

2. The check2 function uses additional space for the substring and its reverse, which in the worst case could be the entire string, so it is 0(n). 3. No additional data structures are used.

1. The two-pointer approach in check1 only uses a constant amount of extra space for the pointers and indices, so it is 0(1).

Therefore, the space complexity of the code is primarily determined by the space needed to store the substring and its reverse in check2, which is 0(n).