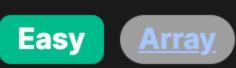
3000. Maximum Area of Longest Diagonal Rectangle



Problem Description

The given problem requires us to find the area of a rectangle with the longest diagonal among all rectangles in the provided 2D integer array dimensions. Each element of this array is itself an array of two elements; the first element represents the length and the second represents the width of a rectangle. If multiple rectangles have diagonals of the same maximum length, we need to find the area of the rectangle that has the maximum area among them.

Intuition

diagonal is the hypotenuse of the right-angled triangle formed by the length and width of the rectangle. To find the longest diagonal, we will square the lengths and widths of the rectangles to avoid working with floating-point numbers that arise from square roots, which simplifies comparisons. The solution iterates over each rectangle, calculating the square of its diagonal (1**2 + w**2). We also track the maximum

The primary step is to understand that the diagonal of a rectangle can be calculated using the Pythagorean theorem, where the

diagonal length found so far (mx) and the area of the corresponding rectangle (ans). If a rectangle's diagonal squared is greater than the current maximum, we update both mx and ans. If it's equal to the current maximum, we only update ans if the rectangle's area is larger than the current largest area.

The implementation of the solution adopts a straightforward approach consistent with the intuition previously described:

Solution Approach

We initialize two variables, ans and mx, both set to 0. ans will hold the maximum area found, and mx will hold the square of the

- longest diagonal found so far. We then loop through each rectangle's dimensions in the dimensions list. For each rectangle, we calculate t, the square of the
- equivalent to applying the Pythagorean theorem to find the diagonal length without taking the square root. Within the loop, we perform a check to see if the calculated t is greater than our current maximum mx. If it is, we have found a new rectangle with a longer diagonal, so we update mx to this new maximum, and we update ans to the area of the current

length of the diagonal, using the formula 1**2 + w**2 where 1 is the length and w is the width of the rectangle. This step is

- rectangle, which is 1 * w. If t is equal to the current mx, it means we have found another rectangle with a diagonal equal to the longest one seen so far. Since the problem requires us to return the area of the rectangle with the maximum area in such cases, we update ans with
- the larger of its current value or the current rectangle's area. After the loop has finished executing, ans contains the correct maximum area, which we return. This approach ensures that the algorithm only passes through the array of rectangles once, making it efficient. The algorithm's
- time complexity is O(n), where n is the number of rectangles given.

beyond the two tracking variables, making it a straightforward but effective approach to solve the problem.

In summary, this solution uses simple iteration and comparison, with no need for complex data structures or additional memory

Example Walkthrough

Let's consider a small example with the dimensions array containing the following dimensions of rectangles: [[3, 4], [5, 12],

[8, 15]].

Let's begin the loop:

We start by initializing ans = 0 (to keep track of the area with the maximum diagonal) and mx = 0 (to keep track of the square of the longest diagonal).

- For the first rectangle [3, 4], compute t = 3**2 + 4**2 = 9 + 16 = 25. \circ Compare t (25) with mx (0). Since 25 is greater than 0, update mx = 25 and ans = 3 * 4 = 12.
- Move to the third rectangle [8, 15]:

Proceed to the second rectangle [5, 12]:

 \circ Compute t = 5**2 + 12**2 = 25 + 144 = 169.

```
\circ Compute t = 8**2 + 15**2 = 64 + 225 = 289.
\circ Compare t (289) with mx (169). Since 289 is greater, update mx = 289 and ans = 8 * 15 = 120.
```

Initialize the maximum diagonal length and area seen so far

// Calculate the square of the diagonal for current rectangle

// If a larger diagonal is found, update maxDiagonalSquare and maxArea

// If the current diagonal equals the maximum diagonal found so far,

// update the maximum area if the current area is greater

int diagonalSquare = length * length + width * width;

maxDiagonalSq = diagonalSq;

} else if (maxDiagonalSq == diagonalSq) {

function areaOfLargestDiagonal(dimensions: number[][]): number {

if (maxDiagonalSquared < diagonalSquared) {</pre>

// Iterate through each dimension pair in the dimensions array

let maxArea = 0; // Initialize maxArea to store the largest area found

maxDiagonalSquared = diagonalSquared; // Update maxDiagonalSquared

maxArea = max(maxArea, length * width);

maxArea = length * width;

// Return the maximum area found

return maxArea;

Loop through each dimension pair in the dimensions list

 \circ Compare t (169) with mx (25). Since 169 is greater, update mx = 169 and ans = 5 * 12 = 60.

In this example walkthrough, we see how the algorithm successfully identifies the rectangle with the longest diagonal and, consequently, the largest area among the rectangles with such a diagonal.

The loop has finished. The final values are mx = 289 and ans = 120. Thus, the maximum area of the rectangle with the longest

Python

class Solution: def areaOfMaxDiagonal(self, dimensions: List[List[int]]) -> int:

max_area = 0

max_diagonal_length = 0

for length, width in dimensions:

from typing import List

Solution Implementation

diagonal is 120.

```
# Calculate the diagonal length's square based on Pythagorean theorem
           diagonal_length_sq = length**2 + width**2
           # If a new maximum diagonal length is found
           if diagonal_length_sq > max_diagonal_length:
               # Update maximum diagonal length and area
               max_diagonal_length = diagonal_length_sq
               max_area = length * width
           # If current diagonal is equal to the maximum found but the area is larger
           elif diagonal_length_sq == max_diagonal_length:
               # Update the area to the larger one
               max_area = max(max_area, length * width)
       # Return the maximum area corresponding to the largest diagonal length
       return max area
Java
class Solution {
   public int areaOfMaxDiagonal(int[][] dimensions) {
        int maxArea = 0; // Holds the area of rectangle with largest diagonal so far
        int maxDiagonalSquare = 0; // Square of the largest diagonal seen so far
       // Loop through each dimensions array, where 'dimension' represents [length, width]
        for (int[] dimension : dimensions) {
           int length = dimension[0]; // Length of current rectangle
           int width = dimension[1]; // Width of current rectangle
```

```
if (maxDiagonalSquare < diagonalSquare) {</pre>
                maxDiagonalSquare = diagonalSquare;
                maxArea = length * width; // Update area to the area of the current rectangle
            } else if (maxDiagonalSquare == diagonalSquare) {
                // If the same diagonal is found, update maxArea if current area is larger
                maxArea = Math.max(maxArea, length * width);
       // Return the area of the rectangle with the largest diagonal encountered
       return maxArea;
C++
class Solution {
public:
    // Function to calculate the area of the rectangle with the maximum diagonal length
    int areaOfMaxDiagonal(vector<vector<int>>& dimensions) {
                             // Variable to store the maximum area found so far
        int maxArea = 0;
        int maxDiagonalSq = 0; // Variable to store the square of the maximum diagonal length
       // Loop through the list of dimensions
        for (auto& dimension : dimensions) {
            int length = dimension[0]; // Store the length of the current rectangle
            int width = dimension[1]; // Store the width of the current rectangle
            int diagonalSq = length * length + width * width; // Calculate the square of the diagonal length
            // If the current diagonal is greater than the maximum found so far,
            // update the maximum diagonal and the maximum area accordingly
            if (maxDiagonalSq < diagonalSq) {</pre>
```

```
for (const [length, width] of dimensions) {
    const diagonalSquared = length * length + width * width; // Calculate the square of the diagonal for current dimensions
   // Compare the squared diagonal length with the maximum found so far
```

TypeScript

};

```
maxArea = length * width; // Update maxArea with the area of the current dimensions
          } else if (maxDiagonalSquared === diagonalSquared) {
              maxArea = Math.max(maxArea, length * width); // If diagonals are equal, choose the larger area
      // Return the area of the rectangle with the longest diagonal
      return maxArea;
from typing import List
class Solution:
   def areaOfMaxDiagonal(self, dimensions: List[List[int]]) -> int:
       # Initialize the maximum diagonal length and area seen so far
       max_diagonal_length = 0
       max area = 0
       # Loop through each dimension pair in the dimensions list
        for length, width in dimensions:
           # Calculate the diagonal length's square based on Pythagorean theorem
            diagonal length sq = length**2 + width**2
           # If a new maximum diagonal length is found
```

let maxDiagonalSquared = 0; // Initialize maxDiagonalSquared to store the square of the longest diagonal found

If current diagonal is equal to the maximum found but the area is larger elif diagonal_length_sq == max_diagonal_length: # Update the area to the larger one $max_area = max(max_area, length * width)$

max_area = length * width

if diagonal_length_sq > max_diagonal_length:

Update maximum diagonal length and area

Return the maximum area corresponding to the largest diagonal length

max_diagonal_length = diagonal_length_sq

Time and Space Complexity

return max_area

Time Complexity

The provided function area0fMaxDiagonal processes each pair of dimensions in the input list exactly once. In each iteration, the function calculates the square of the length and width (1**2 + w**2) and compares it to the maximum diagonal length found so far (mx < t) and mx == t). The calculations and comparisons inside the loop are constant time operations. Let n be the number of pairs in the input list dimensions. Since the loop iterates over all pairs, the time complexity is O(n), where n

is the length of the input list.

Space Complexity

The function uses a fixed number of variables (ans, mx, t, l, w) that do not depend on the size of the input list. Therefore, the space complexity is O(1), meaning it requires a constant amount of additional space.