982. Triples with Bitwise AND Equal To Zero

```
Bit Manipulation
                       Array
                                Hash Table
Hard
```

Given an array of integers nums, we are tasked with finding the count of all unique "AND triples." An "AND triple" is defined as a combination of three indices (i, j, k) within the bounds of the array nums such that the following conditions are met:

- 0 <= j < nums.length • 0 <= k < nums.length

• 0 <= i < nums.length

**Problem Description** 

- The key condition for an "AND triple" is that the bitwise AND operation (&) applied to nums[i], nums[j], and nums[k] results in zero, i.e., nums[i] & nums[j] & nums[k] == 0.

The bitwise AND operation takes two numbers as operands and performs the AND operation on every pair of corresponding bits.

The goal is to write a function that will return the total number of such "AND triples." Intuition

The operation results in a 1 in each bit position for which the corresponding bits of both operands are 1s.

To solve the problem efficiently, we focus on the definition of the bitwise AND operator. When performing an AND operation

result in a time complexity of O(n^3), which is not efficient for large arrays.

where the result is 0, any bit in the resulting number must have been 0 in at least one of the operands. This means that as long as there is a zero at the same position in either i, j, or k index of nums, their AND operation will yield a zero. An intuitive approach might be trying all possible triples (i, j, k) and checking their AND operation one by one, but this would

occurrences for every possible result of nums[i] & nums[j]. We can precompute and store these counts in a Counter dictionary, which is a type of dictionary provided by Python's collections module that counts the occurrences of each element. After counting all the possible AND results of pairs (i, j), we iterate through each value z in nums and for each pair (xy, v) in

Instead, we can take advantage of the commutative property of the AND operation (i.e., a & b = b & a) and count the number of

our Counter dictionary (where xy is a result of nums[i] & nums[j] and v is the count of how many times this result occurred). If xy & z == 0, it means this z paired with all v occurrences of xy will contribute v to the total count of "AND triples". **Solution Approach** 

The solution employs an efficient approach that involves a combination of precomputation and hash mapping through the Counter class, which is part of Python's collections module. The idea is to leverage the properties of the bitwise AND operation to reduce the computational complexity. Here's a breakdown of the implementation steps using the provided solution code:

Precomputation of Pairwise AND results: We create a Counter to precompute and store all possible bitwise AND results x &

## y for each pair in the array nums. This precomputation is done in a nested loop where x and y iterate over all elements in nums.

cnt = Counter(x & y for x in nums for y in nums) cnt now holds the frequency of each result that can be obtained by performing the bitwise AND operation on any two elements in nums.

every combination of xy (a precomputed AND result from cnt) and z (an element in nums), we check if the bitwise AND of xy & z equals 0. return sum(v for xy, v in cnt.items() for z in nums if xy & z == 0)

Iterate Through the Counter and the Array: We iterate through both the Counter dictionary cnt and the array nums. For

In this line of code, xy represents a possible result of nums[i] & nums[j] and v is the number of times this result occurs (the

count). We are interested in the cases where xy & z is zero because that means we've found a valid "AND triple". In the above implementation, the complex part of the problem, which is finding pairs whose AND is 0, is handled by the Counter, which amortizes that work over all element pairs with a single pass through <a href="nums">nums</a>. Then, the final loop provides a multiplication

factor (v) for each zero-resulting AND pair when checked against all elements of nums. This clever use of hash mapping paired

• Counter (hash map): A specialized dictionary for counting hashable objects, quite handy for counting occurrences of results (precomputed AND

with the bitwise operation properties results in an elegant solution that is more computationally efficient than brute force.

Algorithm Patterns: Hash mapping: To store and access the count of occurrences of pairwise AND results efficiently. • Bitwise operations: Core part of the logic, determining the property of the "AND triple".

• Time Complexity: O(n^2), since we iterate over all pairs once to compute the AND results and store them in the Counter, and then iterate over

• Space Complexity: O(n^2), primarily due to the space required to store the Counter dictionary, which in the worst case could hold a distinct

Let's consider a small array nums = [2, 1, 4] to illustrate the solution approach:

**Example Walkthrough** 

**4 & 4 = 4** 

Data Structures:

operations).

Complexity Analysis:

this Counter and **nums** in a nested fashion.

**Precomputation of Pairwise AND results:** 

2 & 4 = 0 (binary 10 AND 100 results in 000)

Iterate Through the Counter and the Array:

count for every pair combination.

nums, the pairs and their AND results are: **2** & 2 = 2

■ 2 & 1 = 0 (since the binary representation of 2 is 10 and 1 is 01, the AND operation results in 00)

• The frequency of AND results will then be stored in a Counter: cnt becomes {2: 1, 0: 3, 1: 1, 4: 1}.

■ For xy = 1: The AND operation with all elements in nums (2, 1, 4) does not yield 0, so no count is added.

■ For xy = 2: The AND operation with any of nums does not yield 0, so no count is added.

**1** & 1 = 1 ■ 1 & 4 = 0 (binary 01 AND 100 results in 000)

With the Counter ready, we now iterate through each unique AND result stored in cnt and each element z in nums, and count if the AND of

■ For xy = 0: The AND operation with all elements in nums (2, 1, 4) will yield 0, and since cnt[0] = 3, we have 3 valid "AND triples" for

We will calculate the bitwise AND for every pair of numbers in the array and count the occurrences of each result. For our example array

■ For xy = 4: The AND operation with any of nums does not yield 0, so no count is added. ○ The count for this iteration will be 0 + 9 + 0 + 0 = 9.

Solution Implementation

from collections import Counter

triplet\_count = 0

for z in nums:

return triplet\_count

**Python** 

Java

class Solution {

class Solution:

each element in nums, totaling 9.

def countTriplets(self, nums: List[int]) -> int:

if bitwise result & z == 0:

triplet\_count += frequency

for bitwise result, frequency in frequency counter.items():

the stored result with z equals 0:

Therefore, the total count of all unique "AND triples" for the array nums = [2, 1, 4] is 9.

# We calculate bitwise AND for all possible pairs of numbers in the nums list. frequency\_counter = Counter(x & y for x in nums for y in nums) # Initialize a result variable to store the count of triplets.

# For each bitwise AND result, we check if there's a number z in nums such that

# to the triplet\_count since each occurrence contributes to a valid triplet.

# After iterating through all items and nums, we return the count of triplets.

// Counts the triplets (x, y, z) such that (x & y & z) == 0 from the given nums array.

// Initialize the count array that will hold the frequency of occurrences of x & y.

// If x & y & z == 0, add the count of x & y to answer.

# bitwise result & z is equal to 0. If ves, we add the frequency of the bitwise result

# Iterate through the items in the frequency counter. Each item is a tuple (bitwise result, frequency).

# Create a counter (dictionary) to store the frequency of each bitwise AND result.

public int countTriplets(int[] nums) { // Find the maximum value in nums to determine the size of the count array. int maxVal = 0; for (int num : nums) { maxVal = Math.max(maxVal, num);

int[] count = new int[maxVal + 1];

for (int xy = 0; xy <= maxVal; ++xy) {</pre>

answer += count[xy];

// Return the total count of valid triplets.

 $if ((xy \& z) == 0) {$ 

// Increment the count of x & y.

```
// Initialize the answer to count the number of valid triplets.
int answer = 0:
// Iterate over all possible combinations of x \& y and check with each z in nums.
```

return answer;

int answer = 0;

return answer;

**}**;

#include <vector>

#include <cstring>

#include <algorithm>

for (int x : nums) {

for (int v : nums) {

for (int z : nums) {

count[x & y]++;

```
class Solution {
public:
   // Function to count the number of triplets (x, y, z) from the nums array such that
   // \times \& v \& z == 0, where & is the bitwise AND operator.
   int countTriplets(vector<int>& nums) {
       // Find the maximum element to determine the size of the count array.
        int maxElement = *max_element(nums.begin(), nums.end());
       // Create and initialize the count array to store the frequency
       // of bitwise AND results of all possible pairs (x, y).
        int count[maxElement + 1];
        memset(count, 0, sizeof(count));
       // Populate the count array with the frequency of all possible (x & y) results.
        for (int x : nums) {
            for (int y : nums) {
                count[x & y]++;
```

// Initialize the answer to count the number of valid triplets.

for (int xy = 0; xy <= maxElement; ++xy) {</pre>

answer += count[xy];

// Return the total count of valid triplets.

 $if ((xy \& z) == 0) {$ 

for (int z : nums) {

// Return the total count of triplets.

def countTriplets(self, nums: List[int]) -> int:

if bitwise result & z == 0:

triplet\_count += frequency

return tripletCount;

class Solution:

from collections import Counter

triplet count = 0

for z in nums:

return triplet\_count

Time and Space Complexity

// Iterate over all possible values of (x & y) and every element z in nums.

// Check if the current combination satisfies x & y & z == 0.

// If so, add the frequency of (x & y) to the answer.

```
TypeScript
function countTriplets(nums: number[]): number {
    // Find the maximum number in the 'nums' array.
    const maxNum = Math.max(...nums);
    // Create an array to store the count of all possible 'x & y' results.
    const count: number[] = Array(maxNum + 1).fill(0);
    // Calculate the frequency of each 'x & y' result.
    for (const x of nums) {
        for (const y of nums) {
            count[x & y]++;
    // Initialize a variable to store the total count of triplets.
    let tripletCount = 0;
    // Iterate through all possible 'x & y' results.
    for (let xy = 0; xy \le maxNum; ++xy) {
        // For each element 'z' in 'nums',
        // if 'xv & z' is zero, increment the total count of triplets by the pre-calculated frequency of 'xy'.
        for (const z of nums) {
            if ((xy \& z) === 0) {
                tripletCount += count[xy];
```

# Create a counter (dictionary) to store the frequency of each bitwise AND result.

# For each bitwise AND result, we check if there's a number z in nums such that

# to the triplet\_count since each occurrence contributes to a valid triplet.

# After iterating through all items and nums, we return the count of triplets.

# bitwise result & z is equal to 0. If yes, we add the frequency of the bitwise result

# We calculate bitwise AND for all possible pairs of numbers in the nums list.

frequency\_counter = Counter(x & y for x in nums for y in nums)

# Initialize a result variable to store the count of triplets.

for bitwise result, frequency in frequency counter.items():

## **Time Complexity** The time complexity of the provided code can be analyzed by breaking down the nested loops and operations:

**Space Complexity** 

this results in  $n^2$  iterations. • For each iteration, the bitwise AND operation is performed once, which takes constant time, so the nested loops contribute 0(n^2) to the time complexity.

# Iterate through the items in the frequency counter. Each item is a tuple (bitwise result, frequency).

unique values. Each update to the Counter (a form of dictionary) takes an average of 0(1) time. • The final loop iterates v times over the elements z in nums for each xy key-value pair in cnt. If v represents the count of occurrences of xy,

m is at most  $n^2$ , and therefore m \* n is at most  $n^3$ , which is dominated by the  $n^2$  term when n is large.

and there are m such unique xy pairs (with m being at most  $n^2$ ), the final loop could iterate up to m \* n times. The bitwise AND operation and comparison inside the final loop also take constant time.

• The count (cnt) is updated during these iterations, which, in the worst case, will store n^2 unique values if all possible AND operations result in

• There are two nested loops that iterate over the array nums to compute bitwise AND (&) for every pair of elements. If n is the length of nums,

- Based on the information above, the total time complexity of this code is  $0(n^2) + 0(m * n)$ , which simplifies to  $0(n^2)$  because
- For space complexity, we consider the additional space used apart from the input:

No other significant additional space is used that grows with the input size.

- The Counter object (cnt) holds the result of bitwise AND operations for each pair of elements in nums, which can grow up to n^2 unique keyvalue pairs in the worst case. Hence, space complexity due to cnt is 0(n^2).
- Thus, the overall space complexity of the code is  $O(n^2)$ .