

**Problem Description** 

In this problem, you have a sorted list of integers from 1 to n. You apply a specific elimination algorithm to this list that consists of repeatedly removing numbers in a particular order until only one number remains. The rules for the removal process are as follows:

- 1. Begin with the first element in the list and then remove every other element. This is considered a left-to-right pass.
- 2. On the next iteration, start from the end of the remaining list and remove every other element, moving right to left. 3. Continue this process, alternating between left-to-right and right-to-left, until there is only one number left in the list.
- The task is to determine what that last remaining number will be given the integer n.

Intuition

To find the solution without simulating the entire elimination process (which would be inefficient for large n), we can observe patterns in the list as we remove numbers. We notice that on each pass, the number of elements in the list is halved. Therefore, we can track the following variables throughout the process: • a1: The first number left in the list after each pass.

- an: The last number left in the list after each pass.
- i: The index of the current pass, starting from 0.
- step: The distance between remaining numbers after each pass. • cnt: The count of remaining numbers after each pass.
- The critical insight is that we can mathematically determine the new values of a1 and an without explicitly simulating each pass. Specifically, a1 increments by step on every pass, and an decrements by step when the pass starts from the right and the count

of remaining numbers is odd. We continue this process until cnt (the count of remaining numbers) is reduced to 1, which indicates only one number is left. At that point, a1 will be the last remaining number, and we can return it. By tracking these variables and updating them after each pass in an alternating fashion, we arrive at the last number without having to actually remove elements from the list, which significantly reduces the time complexity and allows us to solve the

problem with a high efficiency. **Solution Approach** 

## actually removing elements from a list. Instead, we use variables to keep track of the state of the list. Below is the detailed explanation of how the implementation corresponds to the solution approach:

2).

We initially set a1 to 1, an to n, i to 0, step to 1, and cnt to n. These variables represent the start and end of the list, the index of the pass (which determines the direction of elimination), the gap between numbers after each iteration, and the count of numbers remaining, respectively.

The solution implemented in Python uses a straightforward loop to simulate the process of elimination described above, without

• i: Used to track the number of completed passes, and it helps to determine the direction of the current pass. • step: The distance between consecutive remaining numbers. On each pass, the step doubles since every second element is removed. • cnt: The count of remaining numbers in the list. Initially the same as n, and it is halved after each pass (cnt >>= 1 is equivalent to cnt = cnt //

The loop continues until cnt becomes 1, which means there's only one number left. Inside the loop: • On each iteration, we check if the pass is odd (i % 2) or even. For even passes (i is 0, 2, 4, ...), we always increase a1 by step. For odd passes

• a1: This variable represents the value of the first element in the remaining list.

• an: Represents the value of the last element in the remaining list.

(i is 1, 3, 5, ...), we reduce an by step only. • If the count of remaining numbers is odd (cnt % 2), then on every pass an is decremented for even i and a1 is incremented for odd i.

• After processing the current pass, we double the step (step <<= 1 is equivalent to step = step \* 2) because every other number was removed

- during the pass. • The count of remaining elements is halved (cnt >>= 1) as we effectively remove every second element in the list on each pass.
- We increment i after each pass to alternate the direction for the next pass. At the end of the loop, a1 is the last number remaining, which is what we return. The solution makes use of bitwise operations for
- efficiency, such as >>= for division by 2 and <<= for multiplication by 2. The algorithm efficiently tracks and updates the bounds of the remaining list without the need to maintain the list itself, resulting in a time complexity of O(log n).
- **Example Walkthrough**

Let's walk through an example using the solution approach with n = 10. The sorted list from 1 to n will initially be [1, 2, 3, 4, 5, 6, 7, 8, 9, 10]. • Initialization: We set a1 to 1, an to 10, i to 0, step to 1, and cnt to 10.

### • First Pass (i = 0, left-to-right): We increment a1 by step since it's an even pass, so a1 becomes 2. The step doubles to 2, and cnt is halved to 5. Updated list would look like [2, 4, 6, 8, 10] if we were removing elements, but we're not.

and cnt is halved to 1.

• Second Pass (i = 1, right-to-left): Now, it's an odd pass and cnt is odd, so an (last number) is decremented by step, becoming 10 - 2 = 8. We double the step to 4, and halve cnt to 2. • The conceptual list would now be [2, 6].

• Third Pass (i = 2, left-to-right): For this even pass, we increase a1 again by step, resulting in a1 = 2 + 4 = 6. The step is doubled to become 8,

Since cnt is now 1, there's only one number left, and a1 is that last remaining number which is 6. Hence, the last remaining number when using this elimination algorithm on a sorted list from 1 to 10 is 6.

Let's illustrate the solution approach with a real example using `n = 10`. Initially, `a1` is set to 1, representing

- In the first pass (left-to-right), we eliminate every other number, starting with the second one. We update `a1` to

- The second pass (right-to-left) starts with an odd count, so we decrease `an` by the current `step` value, which is

- On the third pass (left-to-right), since it's even, we increase `a1` by the step (now 4) to get `a1` = 6. Doubling Therefore, the last standing number is `6` which we find by intelligently keeping track of the eliminated numbers' ra

Solution Implementation

#### first\_element, last\_element = 1, n # Initialize iteration count, step increment, and count of remaining elements iteration, step, remaining\_count = 0, 1, n

def lastRemaining(self, n: int) -> int:

while remaining\_count > 1:

if iteration % 2:

# Initialize the first and last elements

# Loop until only one element remains

last element -= step

# If the iteration is odd, operate from the right (last to first)

# If count is odd, decrement the last element

// If the count is odd, increment the first element

firstElement += step; // Increment first element

// If the count is odd, decrement the last element

if (remainingElements % 2 == 1) {

// In even rounds, move from left to right

if (remainingElements % 2 == 1) {

lastElement -= step;

// Loop until there's only one number remaining

if (remaining % 2 === 1) {

if (remaining % 2 === 1) {

return first; // Return the remaining number

if (round % 2 === 1) {

// let result = lastRemaining(10);

} else {

// Example usage

// In every odd round, update the last element

// In every even round, update the first element

firstElement += step;

# Adjust first element if the count is odd

**Python** 

class Solution:

```
if remaining count % 2:
                   first_element += step
           # If the iteration is even, operate from the left (first to last)
           else:
               # Increment the first element
               first_element += step
               # Adjust the last element if the count is odd
               if remaining_count % 2:
                   last element -= step
           # Halve the remaining count as we remove half of the elements each iteration
           remaining_count >>= 1
           # Double the step size since every round we skip one more element
           step <<= 1
           # Move to the next iteration
           iteration += 1
       # The first element is the last remaining number
       return first_element
Java
class Solution {
   public int lastRemaining(int n) {
       int firstElement = 1;  // Initialize the first element of the sequence to 1
       int lastElement = n;  // Initialize the last element of the sequence to n
       int step = 1;
                                 // Step size, which doubles each iteration
       // Loop until only one element remains. cnt keeps track of the remaining elements.
       for (int round = 0, remainingElements = n; remainingElements > 1;
            remainingElements >>= 1, step <<= 1, ++round) {
           // In odd rounds, move from right to left
           if (round % 2 == 1) {
                lastElement -= step; // Subtract step from the last element
```

else {

return firstElement;

```
C++
class Solution {
public:
    int lastRemaining(int n) {
        int first = 1, last = n; // Initialize the first and last elements
       int step = 1; // This will represent the step size after each round
       // Loop until there's only one number remaining
        for (int round = 0, remaining = n; remaining > 1; remaining >>= 1, step <<= 1, ++round) {</pre>
           // In every odd round, update the last element
            if (round % 2) {
                last -= step; // Move the last element backwards by the step size
                if (remaining % 2) {
                    first += step; // If there's an odd number of elements, move the first element forwards
            } else {
                // In every even round, update the first element
                first += step; // Move the first element forwards by the step size
                if (remaining % 2) {
                    last -= step; // If there's an odd number of elements, move the last element backwards
       return first; // Return the remaining number
};
TypeScript
// Function to determine the last remaining number after eliminating every second element in rounds
function lastRemaining(n: number): number {
    let first: number = 1; // Initialize the first element
    let last: number = n; // Initialize the last element
    let step: number = 1; // This will represent the step size after each round
```

// Once the loop is finished, there's only one element left, which is the first element.

```
class Solution:
   def lastRemaining(self, n: int) -> int:
       # Initialize the first and last elements
        first_element, last_element = 1, n
       # Initialize iteration count, step increment, and count of remaining elements
        iteration, step, remaining_count = 0, 1, n
       # Loop until only one element remains
       while remaining_count > 1:
           # If the iteration is odd, operate from the right (last to first)
            if iteration % 2:
               # If count is odd, decrement the last element
                last element -= step
               # Adjust first element if the count is odd
               if remaining_count % 2:
                    first_element += step
           # If the iteration is even, operate from the left (first to last)
           else:
               # Increment the first element
               first_element += step
               # Adjust the last element if the count is odd
               if remaining_count % 2:
                    last_element -= step
            # Halve the remaining count as we remove half of the elements each iteration
            remaining_count >>= 1
            # Double the step size since every round we skip one more element
           step <<= 1
            # Move to the next iteration
            iteration += 1
       # The first element is the last remaining number
        return first_element
```

for (let round: number = 0, remaining: number = n; remaining > 1; remaining >>= 1, step <<= 1, round++) {</pre>

first += step; // If there's an odd number of remaining elements, move the first element forwards

last -= step; // If there's an odd number of remaining elements, move the last element backwards

last -= step; // Move the last element backwards by the step size

first += step; // Move the first element forwards by the step size

// console.log(result); // This would compute the result for n=10 and print it to the console

# Time and Space Complexity

the start and end, and every second number from the list is removed at each step. **Time Complexity** 

The time complexity of the function is determined by the while loop which runs until cnt > 1. The number cnt represents the

remaining numbers in the sequence and is halved (cnt >>= 1) at each iteration of the while loop due to removing every second

number. As a result, the time complexity of the loop, and hence the function, is O(log n), as it reduces the sequence size by a

numbers from a list that has been recreated at each step of the problem. In the problem, numbers are removed alternately from

The given Python code defines a method lastRemaining that calculates the last number remaining after repeatedly removing

## **Space Complexity**

factor of two at each step until only one number is left.

The space complexity of the provided code is 0(1). This is because the algorithm uses a fixed number of variables (a1, an, i, step, cnt) and doesn't require any additional space that would grow with the input size n. It does not utilize any complex data structures that would require space proportional to the input size.