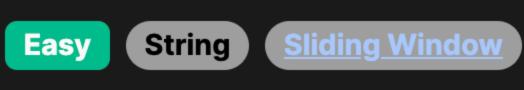
# 2379. Minimum Recolors to Get K Consecutive Black Blocks



# **Problem Description**

respectively. This string is indexed starting at 0, meaning the first element is at index 0. Our objective is to find the minimum number of operations required to ensure that there is at least one segment in the string consisting of k consecutive black blocks. An operation consists of changing a white block into a black block. Imagine we have a block arrangement like "WBWBBW" and we want at least 3 consecutive black blocks. We need to determine

In this problem, we have a string blocks composed of characters 'W' and 'B', which represent white and black blocks,

the fewest changes needed to reach that goal; in this case, it's 1 operation by changing the second block from white to black resulting in "WBBWBB".

The problem specifies two key elements central to our task: • We need to count the number of operations, which correlates to the number of white blocks ('W') that we need to recolor into black blocks

('B').

- We are seeking to create a continuous segment of k consecutive black blocks.
- Intuition

### To arrive at the solution, we need to consider a sliding window approach. The reason a sliding window works well here is that

string, we'll look at the number of white blocks within each window, because those are our candidates for recoloring. Initial thought might involve checking every possible window of k blocks in the string, and for each window, counting the number of white blocks and keeping track of the minimum count seen. But this is inefficient because we would be doing repetitive

we're interested in a continuous sequence of blocks — specifically, subsequences of length k. As we slide the window across the

The solution code implements this efficient approach: • First, it counts the number of white blocks within the first k blocks. This is our initial answer — the worst-case number of operations, if the first

counting. A more optimized approach involves only doing the counting once and then updating the count as the window slides.

• Then, the algorithm enters a loop where it slides the window one block at a time to the right. As the window slides, it adjusts the count by adding one if a white block enters the window on the right, and subtracting one if a white block leaves the window on the left.

white blocks than any window seen before), update ans to cnt.

k blocks contain the most white blocks.

- After each adjustment, it updates the answer with the new count if it's lower than the previous answer. This allows us to go through the entire blocks string only once (after the initial count), adjusting the white block count and
- minimum number as we slide the window. The result is the minimum number of operations needed to ensure a sequence of k
- black blocks.

The implementation of this solution uses a simple but effective pattern known as the sliding window technique. Here's a step-bystep walkthrough of the algorithm, as implemented in the provided solution code: Initialize an integer variable ans (short for answer) to hold the minimum number of recoloring operations needed and another

#### variable cnt (short for count) to keep track of the number of white blocks within the current window. Calculate the number of white blocks in the first window of size k using the count method on the substring blocks[:k]. Set ans to this initial count.

complexity.

Solution Approach

Use a loop to iterate through the string blocks, starting from the kth index up to the end of the string (not inclusive). This loop shifts the fixed-size window one block to the right with each iteration.

As the window slides, update cnt: Increment cnt by 1 if the new block coming into the right side of the window is a white block (blocks[i] == 'W'). Decrement cnt by 1 if the block that is sliding out of the left side of the window is a white block (blocks[i - k] == 'W').

After the loop has iterated over the whole string, ans holds the minimum number of operations required. Return ans.

No complex data structures are needed for this approach; only integer counters are used. The algorithm runs in linear time, O(n),

where n is the length of the blocks string. It only requires a constant amount of additional space, resulting in O(1) space

new window. This is much more efficient than a naive approach where you might recount white blocks for every new window of

After updating cnt, compare it with the current value of ans. If cnt is smaller (which means the window now contains fewer

By maintaining the count cnt as the window slides over the string, the solution avoids the need to recount white blocks in each

size k. **Example Walkthrough** 

Let's illustrate the solution approach with a smaller example. Consider the string blocks = "WWBWB" and let's say we want at least

k = 3 consecutive black blocks. Following the steps outlined in the solution approach: Initialize variables and count white blocks in the first window: The variable ans is set to the number of white blocks in the

first window of size 3, which is blocks[:3] = "WWB". This has 2 white blocks, so ans = 2. We also initialize the count variable

## Iterate through blocks with a sliding window: We start our loop from index 3, which is the fourth block in blocks, and go up

2.

**Python** 

class Solution:

to the end of the string.

cnt = 2.

**Update count and answer:** At index 3, the block is 'W', so no new black block is added to the window. Since we're not adding a white block on the right, cnt remains

However, we should also slide the window, which means we remove the leftmost block from our initial count. The leftmost block within our

**Update the answer:** Our new cnt is less than the previous ans, so we set ans to cnt. Now, ans = 1. Final return value: In the next and final iteration, the window would look at "BWB", and since there's 1 white block entering the window on the right (blocks[4] == 'W'), cnt is incremented back to 2. However, the block that leaves on the left

Following this process, we've completed our sweep of the block string, and lans holds the value 1, which represents the

(blocks[1]) is a black block, so cnt remains 2. There's no need to update ans since cnt is not smaller than ans.

minimum number of operations required to ensure there is a segment of 3 consecutive black blocks in the string.

Solution Implementation

# If the newly included block in the window is white, increment the count

# If the block that is exiting the window is white, decrement the count

// Initialize the answer with the count of white blocks within the first window

// If the entering character is white, increment white count

// If the exiting character is white, decrement white count

// Update minRecolors to the smallest number of white blocks

// Function to find the minimum number of recolors to get at least k consecutive black blocks.

# Update the minimum if the current count is less than the previous minimum

# Return the minimum number of white blocks that need to be recolored

// @param blocks - String representing the arrangement of black (B) and white (W) blocks.

minRecolors = min(minRecolors, whiteBlockCount);

// seen in any window of size 'k'

return minRecolors;

**}**;

**TypeScript** 

// Return the minimum number of recolors needed

// Slide the window of size k across the string and update the minimum recolors requird

initial window is white (blocks[0] = 'W'), which means we need to decrement cnt by 1. Now, cnt = 1.

Therefore, after the first iteration, the window looks at "WBW", with cnt = 1.

Thus, the output would be 1 because we need to change one white block to achieve a segment of 3 consecutive black blocks ("WWBWB" would become "WWBBB").

def minimumRecolors(self, blocks: str, k: int) -> int:

if blocks[index] == 'W':

white count += 1

int whiteCount = 0;

for (int i = 0; i < k; ++i) {

whiteCount++;

int minRecolors = whiteCount;

whiteCount++;

if (blocks.charAt(i) == 'W') {

for (int i = k; i < blocks.length(); ++i) {</pre>

if (blocks.charAt(i - k) == 'W') {

if (blocks.charAt(i) == 'W') {

if blocks[index - k] == 'W':

- # Initial count of 'W' in the first window of size k white count = blocks[:k].count('W') # Initialize minimum white blocks to be recolored with the count from the first window min\_recolors = white\_count # Slide the window of size k through the blocks string for index in range(k, len(blocks)):
- white count -= 1 # Update the minimum if the current count is less than the previous minimum min\_recolors = min(min\_recolors, white\_count) # Return the minimum number of white blocks that need to be recolored return min\_recolors
- Java class Solution { // Method to find the minimum number of recolors needed to get at least k consecutive black blocks public int minimumRecolors(String blocks, int k) { // Initialize the count of white blocks within the first window of size k
- whiteCount--; // Update the minimum recolors if the current count is less than the previous minimum minRecolors = Math.min(minRecolors, whiteCount); // Return the minimum number of recolors required return minRecolors; C++ class Solution { public: // Function to find the minimum number of recolors needed // to get at least 'k' consecutive black blocks int minimumRecolors(string blocks, int k) { // Count the number of white blocks in the first 'k' block segment int whiteBlockCount = count(blocks.begin(), blocks.begin() + k, 'W'); // The minimum recolors needed is initially set to the number // of white blocks in the first window of size 'k' int minRecolors = whiteBlockCount; // Iterate over the blocks starting from the 'k'th block for (int i = k; i < blocks.size(); ++i) {</pre> // Increase whiteBlockCount if the current block is white whiteBlockCount += blocks[i] == 'W'; // Decrease whiteBlockCount if the leftmost block of // the previous window was white whiteBlockCount -= blocks[i - k] == 'W';
- // @param k Number representing the desired consecutive black blocks length. // @returns The minimum number of recolors required (changing 'W' to 'B'). function minimumRecolors(blocks: string, k: number): number { // Initialize a counter for white blocks in the first window of size k. let whiteCount = 0; for (let i = 0; i < k; ++i) { whiteCount += blocks[i] === 'W' ? 1 : 0; // The answer starts off as the number of white blocks in the first window. let minRecolors = whiteCount; // Slide the window of size k across the blocks string while updating the count. for (let i = k; i < blocks.length; ++i) {</pre> // If the new block is white, increase the count. whiteCount += blocks[i] === 'W' ? 1 : 0; // If the block exiting the window is white, decrease the count. whiteCount -= blocks[i - k] === 'W' ? 1 : 0;
- // Update the answer with the minimum count seen so far. minRecolors = Math.min(minRecolors, whiteCount); // Return the minimum number of recolors needed. return minRecolors; class Solution: def minimumRecolors(self, blocks: str, k: int) -> int: # Initial count of 'W' in the first window of size k white count = blocks[:k].count('W') # Initialize minimum white blocks to be recolored with the count from the first window min\_recolors = white\_count # Slide the window of size k through the blocks string for index in range(k, len(blocks)): # If the newly included block in the window is white, increment the count if blocks[index] == 'W': white count += 1 # If the block that is exiting the window is white, decrement the count if blocks[index - k] == 'W':
- Time and Space Complexity

return min\_recolors

white count -= 1

min\_recolors = min(min\_recolors, white\_count)

**Time Complexity** The time complexity of the given code can be analyzed based on the operations performed within the loop. The loop runs from k to len(blocks), which indicates that the loop runs len(blocks) - k times. Within each iteration, the code performs constanttime operations such as comparison and increment/decrement operations. Therefore, the overall time complexity is O(len(blocks) - k), which simplifies to O(n) where n is the length of the blocks string.

# **Space Complexity**

The space complexity is determined by the amount of additional space used by the algorithm relative to the input size. The given code uses a fixed number of variables (ans, cnt, i) that do not depend on the size of the input. Therefore, the space complexity remains constant, regardless of the input size. Consequently, the space complexity of the code is 0(1), indicating constant space complexity.