

1111. Maximum Nesting Depth of Two Valid Parentheses Strings

MediumStackString

Leetcode Link

Problem Description

The problem requires us to operate on a string that is said to be a valid parentheses string (VPS). A VPS is defined by the following properties:

- It consists of "(" and ")" characters only.
- It can be an empty string.
- It can be the concatenation of two valid parentheses strings, **A** and **B**.
- It can be a string that encapsulates a valid parentheses string within parentheses, like "(A)".

The concept of nesting depth `depth(S)` is also introduced, which is calculated as follows:

- `depth("")` (the depth of an empty string) is 0.
- `depth(A + B)` (the depth of the concatenation of two VPS's) is the maximum of the depths of **A** and **B**.
- `depth("(" + A + ")")` (the depth when a VPS **A** is within parentheses) is $1 + \text{depth}(A)$.

The task is to split a given VPS `seq` into two disjoint subsequences **A** and **B**. Both **A** and **B** should themselves be VPS's, and the total length of **A** and **B** should be equal to the length of `seq`. Out of all possible **A** and **B**, we need to find the pair that minimizes the maximum of their depths.

The output should be an array `answer` where each element is either 0 or 1. `answer[i]` should be 0 if the character `seq[i]` is part of the subsequence **A**, otherwise it's part of **B**. Though there can be multiple correct **A** and **B** pairs, any valid solution is acceptable.

Intuition

The intuition behind the solution is related to the observation that in a valid parentheses string, a subsequence **A** can be formed by taking some of the "left parentheses" (and the matching "right parentheses") to keep **A** as a VPS, while the rest can form the subsequence **B**. By doing so, we can potentially balance the depth between **A** and **B**.

A straightforward approach to split the pairs between **A** and **B** is to alternate assignments of the parentheses as we traverse the given `seq`. Each time we encounter a "left parentheses" (, we assign it to **A** if we are at an "even level" and to **B** if we are at an "odd level". When we encounter a "right parentheses"), we step back one level and then perform the similar alternating assignment.

This alternating assignment helps ensure that we don't create too deep a nesting in any of the subsequences, hence minimizing the maximum depth between them.

To implement this, we keep a counter `x` to track the current level of depth during the traversal, and toggle the assignment of parentheses based on the current level's parity (even or odd) indicated by `x & 1`. An even `x` represents that we assign to sequence **A** (`answer[i]` is 0), and an odd `x` represents that we assign to sequence **B** (`answer[i]` is 1).

Solution Approach

The solution uses a simple algorithm that iterates through the given VPS string `seq` and assigns each parenthesis to either subsequence **A** or **B** based on the current depth level. Here's how it's implemented:

- Initialize an array `ans` with the same length as `seq` to store the assignment of each parenthesis to subsequence **A** (represented by 0) or **B** (represented by 1).
- Initialize a variable `x` that will keep track of the current depth level as we walk through the string. This variable will increment when a "left parentheses" (is encountered and decrement when a "right parentheses") is encountered.
- Iterate over the characters of `seq` using `enumerate` to have both index `i` and character `c`. For each character:
 - If `c` is a "left parentheses" (, determine if the current depth level `x` is even or odd by checking `x & 1`. If it's even (`x & 1` is 0), it means we are at an even depth level, so `ans[i]` is set to 0, indicating the parenthesis belongs to **A**. If it's odd (`x & 1` is 1), `ans[i]` is set to 1, meaning it belongs to **B**. After assigning the parenthesis, increment the depth level `x`.
 - If `c` is a "right parentheses"), first decrement the depth level `x` because we are closing a parenthesis. Then check `x & 1` and assign `ans[i]` in the same way as if it was a "left parentheses".
- Once the iteration is complete, return the `ans` array which now represents the VPS split between **A** and **B** in such a way that the maximum depth of both is minimized.

This algorithmic approach ensures that the split of parentheses between **A** and **B** keeps their respective depths balanced, thus minimizing the maximum depth of any VPS that could be formed by them. No additional data structures are needed beyond the `ans` array for the output and a single integer variable for depth tracking.

Example Walkthrough

Let's consider the VPS sequence `seq = "()()())"`

We wish to split this sequence into two disjoint subsequences **A** and **B** such that the maximum depth of **A** and **B** is minimized.

The algorithm goes as follows:

- Initialize `ans` to be an array of the same length as `seq`. In this case, `ans = [0, 0, 0, 0, 0, 0]`.
- Initialize our depth counter, `x`, to 0.
- Now, we iterate through the `seq` character by character.

Here's a step-by-step walkthrough:

- `i = 0, c = '('`. Since `x` is 0 (even), `ans[i]` is set to 0. `x` increments to 1. Now, `ans = [0, 0, 0, 0, 0, 0]`.
- `i = 1, c = '('`. `x` is 1 (odd), `ans[i]` is set to 1. `x` increments to 2. Now, `ans = [0, 1, 0, 0, 0, 0]`.
- `i = 2, c = '('`. Before assignment, we decrement `x` to 1. Since `x` is now 1 (odd), `ans[i]` is set to 1. Now, `ans = [0, 1, 1, 0, 0, 0]`.
- `i = 3, c = '('`. `x` is 1 (odd), `ans[i]` is set to 1. `x` increments to 2. Now, `ans = [0, 1, 1, 0, 0, 0]`.
- `i = 4, c = '('`. We decrement `x` to 1 before assigning. `x` is 1 (odd), `ans[i]` is set to 1. Now, `ans = [0, 1, 1, 0, 1, 0]`.
- `i = 5, c = '('`. We decrement `x` to 0 before assigning. Since `x` is now 0 (even), `ans[i]` is set to 0. Now, `ans = [0, 1, 1, 0, 1, 1]`.

So the final `ans` representing the split is `[0, 1, 1, 0, 1, 0]`. Here, **A** consists of the parentheses at positions 0, 3, and 5 corresponding to the VPS "(()", and **B** consists of the parentheses at positions 1, 2, and 4 corresponding to the VPS "()()". Both **A** and **B** are valid parentheses strings and the overall maximum depth is minimized.

Python Solution

```
1 from typing import List
2
3 class Solution:
4     def maxDepthAfterSplit(self, seq: str) -> List[int]:
5         # Initialize an array to store the assignment of depths to each parenthesis.
6         assigned_depth = [0] * len(seq)
7
8         # Declare a variable to keep track of the current depth level.
9         depth_level = 0
10
11        # Iterate over each character in the sequence alongside its index.
12        for index, char in enumerate(seq):
13            if char == "(":
14                # If the parenthesis is an opening one, determine the depth.
15                # We use bitwise AND with 1 to alternate between 0 and 1.
16                assigned_depth[index] = depth_level & 1
17                # Increment depth level since we've encountered an opening parenthesis.
18                depth_level += 1
19            else:
20                # If the parenthesis is a closing one, first decrement the depth level.
21                depth_level -= 1
22                # After decreasing the depth level, determine the depth for the closing parenthesis.
23                assigned_depth[index] = depth_level & 1
24
25        # Return the final array with the assigned depths.
26        return assigned_depth
27
```

Java Solution

```
1 class Solution {
2     // This method splits the maximum depth of balanced sub-sequences from the given sequence
3     public int[] maxDepthAfterSplit(String seq) {
4         // Get the length of the sequence
5         int n = seq.length();
6         // Initialize the answer array to hold the group assignment of each character in 'seq'
7         int[] answer = new int[n];
8         // Loop through the characters of 'seq'
9         for (int i = 0, depthLevel = 0; i < n; ++i) {
10            // If the current character is an opening parenthesis '('
11            if (seq.charAt(i) == '(') {
12                // Assign the current group based on the parity of 'depthLevel' and increment 'depthLevel'
13                answer[i] = depthLevel++ & 1;
14            } else {
15                // If it's a closing parenthesis ')', decrement 'depthLevel' first
16                // Then assign the current group based on the parity of 'depthLevel'
17                answer[i] = --depthLevel & 1;
18            }
19        }
20        // Return the array containing group assignments for each character
21        return answer;
22    }
23 }
24
```

C++ Solution

```
1 #include <vector>
2 #include <string>
3
4 class Solution {
5 public:
6     // Function to assign depths after split based on the sequence of parentheses
7     std::vector<int> maxDepthAfterSplit(std::string seq) {
8         int size = seq.size(); // Get the size of the sequence
9         std::vector<int> answer(size); // Initialize a vector to store the answer with the same size as the input
10        for (int i = 0, depth = 0; i < size; ++i) {
11            if (seq[i] == '(') {
12                // If the character is '(', increment the depth and assign the current depth modulo 2 to the answer.
13                // This is because depth alternates for better balancing the depth after splitting.
14                answer[i] = depth++ & 1;
15            } else {
16                // For ')', decrement the depth first because we are ending a level of depth
17                // before assigning to answer.
18                answer[i] = --depth & 1;
19            }
20            // " & 1" effectively takes the least significant bit of the depth,
21            // which is equivalent to depth % 2 (but faster), to alternate between 0 and 1.
22        }
23        return answer; // Return the computed vector of depths.
24    }
25 };
26
```

Typescript Solution

```
1 function maxDepthAfterSplit(seq: string): number[] {
2     const sequenceLength = seq.length; // The length of the input string
3     const answer: number[] = new Array(sequenceLength); // Initialize the answer array with the same length as input
4
5     // Initialize variables for iteration and depth calculation
6     let depth = 0; // Used to track the depth of nested parentheses
7
8     // Loop through each character in the input string
9     for (let index = 0; index < sequenceLength; ++index) {
10        if (seq[index] === '(') {
11            // If the current character is an opening parenthesis, determine which group it belongs to
12            // Use bitwise AND with 1 to alternate groups (0 and 1) as depth increases
13            answer[index] = depth++ & 1;
14        } else {
15            // If the current character is a closing parenthesis, decrement depth first before determining the group
16            // Again, use bitwise AND with 1 to alternate groups
17            answer[index] = --depth & 1;
18        }
19    }
20
21    return answer; // Return the populated answer array
22 }
23
```

Time and Space Complexity

The given code snippet takes a string `seq`, representing a sequence of parentheses, and returns a list of integers that correspond to the depth of each parenthesis in a way such that the sequence is split into two sequences **A** and **B**, with both being valid parentheses strings and depth ideally balanced.

Time Complexity

The time complexity of this code is $O(n)$. This is because there is a single loop that iterates through the string `seq`, which is of length `n`. Within this loop, all operations (assigning `ans[i]`, incrementing or decrementing `x`, and bitwise comparison with 1) are constant time operations, $O(1)$. Since these constant time operations are repeated `n` times, the overall time complexity remains linear.

Space Complexity

The space complexity is also $O(n)$. The additional space used by the algorithm is primarily for the output list `ans`, which has the same length as the input string `seq`. There are no other data structures that grow with the size of the input, and variables like `x` and `i` use constant space. Therefore, the space complexity is directly proportional to the input size.