# **Problem Description**

find the number of distinct excellent pairs in the array, where a pair (num1, num2) is considered excellent if it satisfies two conditions: 1. Both num1 and num2 exist in nums.

In this problem, we're presented with an array nums consisting of 0-indexed positive integers and a positive integer k. Our goal is to

2. The sum of the number of set bits (bits with value 1) in num1 OR num2 and num1 AND num2 must be greater than or equal to k. We count the number of set bits using bitwise OR and AND operations, and we are looking for pairs that collectively have a large

an excellent pair depend on the sum of set bits in num1 OR num2 and num1 AND num2.

enough number of set bits to meet or exceed the threshold k.

the order matters—(a, b) is considered different from (b, a). Intuition

Also, it's important to note that a pair where num1 is equal to num2 can also be considered excellent if the number of its set bits is

sufficient and at least one occurrence of the number exists in the array. Moreover, we want the count of distinct excellent pairs, so

To find the number of excellent pairs efficiently, we need to think about the problem in terms of set bits because the conditions for

## One intuitive approach is to use the properties of bitwise operations. For any two numbers, num1 OR num2 will have the highest

hand, num1 AND num2 will have set bits only in positions where both num1 and num2 have set bits. Since only the set bits matter, and we're looking for pairs (num1, num2) that meet a certain combined set bit count, we can reduce

the complexity by avoiding the direct computation of num1 OR num2 and num1 AND num2 for all pairs. Instead, we preprocess by

combining counts from the counter, avoiding the need to directly compute the set bit sum for every possible pair.

problem of having to generate and check every possible pair, which would be computationally inefficient.

possible set bit count of the two numbers because OR operation results in a 1 for each bit that is 1 in either num1 or num2. On the other

counting the set bits for each unique number in nums. To ensure that we count each distinct pair only once, we eliminate duplicates in the array by converting it into a set. We then create a counter to keep track of how many numbers have a specific set bit count. This preprocessing step simplifies our task to just

When we iterate over the unique numbers in nums set, we determine the number of set bits for each unique value using bit\_count(). We then iterate through our set bit count counter and add the count of numbers that have enough set bits to complement the current number's set bit count to reach at least k. This way, we find all the pairs that, when combined, meet or exceed the threshold k.

Solution Approach The given solution employs a combination of bit manipulation and hash mapping to efficiently compute the number of excellent pairs.

By adding the counts for all such complementing set bit count pairs, we calculate the total number of excellent pairs, bypassing the

1. Eliminating Duplicates: The solution begins by converting the original list of numbers into a set. This serves two purposes: • Ensures that each number is considered only once, thereby eliminating redundant pairs like (num1, num1) when num1

• Helps later in avoiding double counting of excellent pairs with the same numbers but in different positions (since (a, b) and

### (b, a) are distinct). 1 s = set(nums)

1 cnt = Counter()

cnt[v.bit\_count()] += 1

2 for v in s:

1 ans = 0

Let's break down the implementation step-by-step:

appears multiple times in the input list.

will be equivalent to bit\_count(num1) + bit\_count(num2).

sum of bit\_count(v) and i is greater than or equal to k.

which holds the total number of distinct excellent pairs.

Let's say we have an array nums = [3, 1, 2, 2] and k = 3.

Next, we count set bits in step 2. Every number will be processed as follows:

• For v = 2 also with 1 set bit, the situation is the same as for v = 1.

the set bit counts to find excellent pairs, which is computationally faster.

def countExcellentPairs(self, nums: List[int], k: int) -> int:

# Count the frequency of the bit count for each number

# Counter to store the frequency of bit counts

bit\_count\_freq[num.bit\_count()] += 1

# Calculate the number of excellent pairs

# Iterate over the bit count frequencies

# Return the total count of excellent pairs

# Use a set to eliminate duplicates as each number contributes uniquely

excellent\_pairs\_count = 0 # Initialize the count of excellent pairs

current\_bit\_count = num.bit\_count() # Get the bit count of current number

numbers with 1 set bit, we add 2 to ans, resulting in ans = 2 so far.

• For v = 3 with 2 set bits, we compare it against cnt entries.

the conditions are (3, 1), (3, 2), (1, 3), (2, 3).

from collections import Counter

unique\_nums = set(nums)

for num in unique\_nums:

for num in unique\_nums:

return excellent\_pairs\_count

uniqueNumbers.add(num);

for (int num : uniqueNumbers) {

++bitCounts[bits];

long totalPairs = 0;

bit\_count\_freq = Counter()

class Solution:

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38 };

t holds the number of set bits for the current unique number v.

2. Counting Set Bits: Then, the algorithm uses a Counter from the collections module to keep track of how many numbers share the same set bit count.

- The method bit\_count() is used to determine the number of set bits in each number. These counts are stored such that cnt[i]
- reflects the number of unique numbers with exactly i set bits. 3. Finding Excellent Pairs: The core of the solution is to find all the pairs that satisfy the condition of having a combined set bit

sum (via bitwise OR and AND) greater than or equal to k. Since an OR operation can never reduce the number of set bits, and an

AND operation can only produce set bits that are already set in both numbers, the combined set bit count for a pair (num1, num2)

The solution iterates over each unique value v from the set of numbers and then checks for every bit count i stored in cnt if the

2 for v in s: t = v.bit\_count() for i, x in cnt.items(): if t + i >= k: ans += x

∘ If t + i is greater than or equal to k, all numbers with a set bit count i can form an excellent pair with v. The count of such

numbers is x, and we add this to our total count of excellent pairs (ans). This step leverages the precomputed counts of

unique set bit numbers to quickly find the number of complements needed to meet the set bit threshold k.

4. Returning the result: Finally, after iterating through all unique numbers and their possible pairings, the solution returns ans,

1 return ans In summary, the solution follows a smart preprocessing step to calculate and use set bit counts for optimization. This eliminates

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redundant operations and directly navigates to the crux of the problem, which significantly improves the computational efficiency.
By employing a hash map to store unique set bit counts and identify complementing pairs, the solution reduces what would be a
quadratic-time complexity task to a complexity proportional to the product of unique numbers and unique set bit counts present in
the input.
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**Example Walkthrough** 

• 1 has 1 set bit.

making ans = 4.

• 2 has 1 set bit. • 3 has 2 set bits.

The Counter object cnt turns out to be {1: 2, 2: 1}, indicating that there are 2 numbers with 1 set bit and 1 number with 2 set bits.

First, we'll apply step 1 and eliminate duplicates by converting nums into a set, which will give us  $s = \{1, 2, 3\}$ .

Now for step 3, we find excellent pairs. We go through each number in s and compare the set bit count with our k value: • For v = 1 with 1 set bit, cnt[1] = 2. Since 1 (set bits of v) + 1 (set bits of another number) is not >= k, no excellent pairs are formed with v = 1.

Finally, since the pairs (1, 3) and (2, 3) are also excellent pairs (order matters), we count them again. This adds another 2 to ans,

After processing all the unique numbers, we follow step 4 and conclude that there are 4 distinct excellent pairs. The pairs that satisfy

This example demonstrates the solution's efficiency, as it avoids checking all possible combinations of nums and directly focuses on

o 2 (set bits of v) + 1 (set bits of another number) is >= k, thus excellent pairs are (3, 1) and (3, 2). Since there are 2

- Python Solution
- for bit\_count, freq in bit\_count\_freq.items(): 22 # If the sum of bit counts is greater than or equal to k, add to the count 24 if current\_bit\_count + bit\_count >= k: 25 excellent\_pairs\_count += freq 26

// To store the total number of excellent pairs

int bits = Integer.bitCount(num); // Count the 1-bits in the binary representation of 'num'

// Increase the count for this number of 1-bits

int[] bitCounts = new int[32]; // Array to store how many numbers have a certain bit count

// Count the occurrence of each bit count for the unique elements

// Variable to store the final count of excellent pairs

if (current\_bit\_count + i >= k) {

// Return the final count of excellent pairs

// Iterate over each unique number and find the count of numbers that

excellent\_pairs\_count += bit\_count[i];

// have the required number of bits for forming an excellent pair with the current number

// If the sum of the bit counts of both numbers meets or exceeds k,

// add the frequency of the corresponding bit count to the answer

// Function to count the number of set bits in the binary representation of a number

int current\_bit\_count = \_\_builtin\_popcount(number); // Get bit count of the current number

long long excellent\_pairs\_count = 0;

for (int number : unique\_numbers) {

return excellent\_pairs\_count;

function countSetBits(num: number): number {

for (int i = 0; i < 32; ++i) {

#### // Method to count the number of excellent pairs public long countExcellentPairs(int[] nums, int k) { // Use a set to eliminate duplicate values from 'nums' 6 Set<Integer> uniqueNumbers = new HashSet<>(); for (int num : nums) {

class Solution {

Java Solution

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           // Iterate over the unique numbers to find pairs
           for (int num : uniqueNumbers) {
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                int bits = Integer.bitCount(num); // Count the 1-bits for this number
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               // Check for each possible bit count that could form an excellent pair with 'num'
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               for (int i = 0; i < 32; ++i) {
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                   // Check if the sum of 1-bits is at least 'k'
                   if (bits + i >= k) {
                       totalPairs += bitCounts[i]; // If it is, add the count of numbers with 'i' bits to the total
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           return totalPairs; // Return the total count of excellent pairs
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36 }
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C++ Solution
 1 #include <vector>
   #include <unordered_set>
   class Solution {
   public:
        long long countExcellentPairs(std::vector<int>& nums, int k) {
           // Create a set to eliminate duplicates from the input vector
           std::unordered_set<int> unique_numbers(nums.begin(), nums.end());
           // Array to count the frequency of set bits (1s) in the binary representation of numbers
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           // There are at most 32 bits in an int, so we create an array of size 32
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           std::vector<int> bit_count(32, 0);
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           // Count the number of times a number with a particular bit count appears
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           for (int number : unique_numbers) {
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                ++bit_count[__builtin_popcount(number)]; // __builtin_popcount returns the number of bits set to 1
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### while (num > 0) { count += num & 1; // Increment count if the least significant bit is set 6 num >>>= 1;

Typescript Solution

let count = 0;

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// Logical Right Shift to process the next bit
       return count;
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   function countExcellentPairs(nums: number[], k: number): number {
       // Set to eliminate duplicates from the input array
       const uniqueNumbers = new Set<number>(nums);
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       // Array to count the frequency of set bits (1s) in the binary representation of numbers
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       // Since integers in JavaScript are represented using 32 bits, we create an array of size 32
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       const bitCount = new Array(32).fill(0);
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       // Count the number of times a number with a particular bit count appears
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       uniqueNumbers.forEach(number => {
           bitCount[countSetBits(number)]++; // Increment the frequency of the bit count
       });
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       // Variable to store the final count of excellent pairs
25
       let excellentPairsCount = 0;
26
       // Iterate over each unique number and find the count of numbers that
       // have the required number of bits for forming an excellent pair with the current number
29
       uniqueNumbers.forEach(number => {
           const currentBitCount = countSetBits(number); // Get bit count of the current number
30
           for (let i = 0; i < 32; i++) {
31
               if (currentBitCount + i >= k) {
                   // If the sum of the bit counts of both numbers meets or exceeds k,
                   // add the frequency of the corresponding bit count to the answer
                   excellentPairsCount += bitCount[i];
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       });
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39
       // Return the final count of excellent pairs
       return excellentPairsCount;
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42 }
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Time and Space Complexity
The given code counts the number of "excellent" pairs in an array, with a pair (a, b) being excellent if the number of bits set to 1 in
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# **Time Complexity** The time complexity of the function involves several steps:

element is added to the set once.

their binary representation (a OR b) is at least k.

set s, this step has a complexity of O(U), where U is the number of unique numbers in nums. 3. Populating the Counter: This again is O(U) since we're iterating over the set s once.

- 4. Running the nested loops, where the outer loop is over each unique value v in s and the inner loop is over the counts in Counter:
- Since the outer loop runs O(U) times and the inner loop runs a maximum of O(U) times, the nested loop part has a worse-case complexity of O(U^2).

1. Creating a set s from nums: This operation has a time complexity of O(N) where N is the number of elements in nums, as each

2. Counting the bit\_count t of each distinct number in s: Each call to v.bit\_count() is 0(1). Since we do this for each number in the

- The total time complexity is, therefore,  $O(N + U^2)$ .
- The space complexity consists of: 1. The set s, which takes O(U) space.

2. The Counter object cnt, which takes another O(U) space.

**Space Complexity** 

The additional space for the variable ans and loop counters is 0(1). Hence, the total space complexity of the function is O(U), where U is the count of unique numbers in nums.