2198. Number of Single Divisor Triplets

Medium Math

## **Problem Description** You are tasked with finding the number of special triplets in a given array of positive integers, where the array is indexed starting at

sum of the numbers at these indices, specifically nums[i] + nums[j] + nums[k], is divisible by exactly one of the three numbers nums [i], nums [j], or nums [k]. The goal is to count how many such triplets exist in the provided array. Intuition

0. A triplet consists of three distinct indices (i, j, k). This triplet is considered special - termed as a single divisor triplet - if the

**Leetcode Link** 

To solve this problem, we should first understand that checking every possible triplet would be quite inefficient if we did so naïvely, since that requires looking at all possible combinations of numbers in the array. However, by utilizing the Counter class from Python's

collections module, we can improve the efficiency when certain numbers are repeated in nums. The solution approach involves the following steps: 1. We count the occurrences of each number in the array utilizing the Counter class.

2. Once we have the counts of each number, we iterate through the array thrice – simulating the selection of three distinct

elements, denoted by a, b, and c, for our triplet. If a, b, and c are the same number, we skip to avoid checking the same

combination again. It's important to note that a, b, or c could be the same number but from different indices. 3. For each triplet, we check if the sum of a, b, and c (s = a + b + c) is divisible by only one of the three numbers. The check

number of possible triplets.

function helps in this regard by returning True if exactly one divisibility condition is met. 4. If the check function returns True, we must account for the occurrences of each number. Since the numbers might be duplicated in nums, we need to calculate the number of triplets that we can form with the given counts of a, b, and c. We do this by considering the different scenarios:

∘ If two numbers are the same and the third is different, we use combinations like cnt1 \* (cnt1 - 1) \* cnt3 to reflect the

especially when dealing with an array containing many repeated numbers. **Solution Approach** 

Using the Counter allows us to avoid unnecessary duplicate computations and vastly reduces the number of required operations,

If all three numbers are different, we simply multiply the counts: cnt1 \* cnt2 \* cnt3.

5. We sum up these counts to get the total number of single divisor triplets.

The implementation of the solution leverages the Counter class from Python's collections module to efficiently track the number of times each number appears in the input array nums. This is crucial to avoid duplicate computations when numbers are repeated. Here's a step-by-step breakdown of the algorithm:

1. A Counter object is created for the array nums to get the count of each unique number. For example, if nums is [1, 2, 2, 3], the Counter object would be {1: 1, 2: 2, 3: 1}.

2. We define a helper function check(a, b, c) which calculates the sum s of the parameters a, b, and c. It then checks if s is

divisible by exactly one of the values a, b, or c by using the modulo operation s % x for each and counts the occurrences of 0

# (true divisibility) using a generator expression. If this count is 1, it returns True; otherwise, it returns False.

different. For example:

divisor triplets.

**Example Walkthrough** 

**Step 1: Counter Object Creation** 

Step 3: Iterate over unique combinations

• First we pick a = 1 from the counter.

Step 4: Apply the check function and calculate the triplet count

pick c. This gives us a count of (2 \* (2 - 1) / 2) \* 1 = 1.

Step 5: Add valid triplets count to the accumulator

cnt2 ways to choose b, and cnt3 ways to choose c.

exhaustively examining each distinct triplet in the original array.

We create a Counter object from nums which gives us {1: 1, 2: 2, 3: 1, 4: 1}.

We start iterating over every unique combination using the counts from the Counter object:

3. The solution then iterates through every possible combination of the counts of each unique number, looking for ways to form triplets. The outer loop picks a number, denoted as a, from the Counter. The nested loops pick additional numbers, b and c.

- 4. For each triplet (a, b, c), it calls the check function to verify whether it's a single divisor triplet. If it is, the count of possible single divisor triplets that can be made with a, b, and c is calculated. This depends on whether a, b, and c are the same or
- ∘ If a is the same as b but different from c, the number of triplets is cnt1 \* (cnt1 1) / 2 \* cnt3 because there are cnt1 \* (cnt1 - 1) / 2 ways to choose two numbers a and b from cnt1 occurrences, and cnt3 ways to choose c.

5. It's essential to handle overcounting carefully; the solution ensures that each triplet of numbers is counted exactly once, with

attention to the distinction between using combinations vs. permutations when selecting elements from counts.

If all three numbers are different, the number of triplets will be cnt1 \* cnt2 \* cnt3 since there are cnt1 ways to choose a,

7. After checking all possible combinations of numbers from the Counter, the final result is returned, which is the accumulated count of all valid triplets.

Thus, the algorithm makes smart use of combinatorics and the properties of divisibility to compute the desired count without

6. As each valid triplet is found, its count is added to an accumulator variable, ans, which finally holds the total count of single

**Given Array** Consider the array nums = [1, 2, 2, 3, 4].

Let's walk through an example using the provided solution approach to get a clearer understanding of how the algorithm works.

**Step 2: Define the check function** The check(a, b, c) function takes three numbers and returns True if the sum s = a + b + c is divisible by exactly one of a, b, or c.

## • Then we pick b = 2. Since we have two 2's, cnt2 = 2. • Next, we pick c = 4. Since there is only one 4, cnt3 = 1.

Using the check function, we find that 1 + 2 + 4 = 7 is not divisible by any number exactly once, so we discard this triplet.

Next, we try a = 2, b = 2, and c = 4 (where a and b are the same). The sum is 2 + 2 + 4 = 8, and this sum is divisible by 4 only,

which makes it a valid triplet. Since a and b are the same, there are cnt2 \* (cnt2 - 1) / 2 ways to pick a and b, and cnt3 ways to

After iterating through all unique number combinations from the Counter, we sum up all valid triplet counts we have found. Suppose

we found another valid triplet using different numbers, let's say we found 1 more valid triplet in the process. Our total count ans

# We continue this process for other possible combinations.

Step 6: Final result

would be 1 + 1 = 2.

exactly by one of the three numbers.

1 from collections import Counter

from typing import List

class Solution:

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C++ Solution

1 #include <vector>

4 class Solution {

public:

2 using namespace std;

Thus, for the given array [1, 2, 2, 3, 4], our final result is 2, meaning there are two single divisor triplets whose sum is divisible

Each time we find a valid triplet combination, we increment our accumulator ans with the count of its occurrences.

Python Solution

sum\_triplet = a + b + c

for a, count\_a in num\_counter.items():

for b, count\_b in num\_counter.items():

elif b == c:

else:

public long singleDivisorTriplet(int[] nums) {

if (divisorCount != 1) {

**if** (i == j && i == k) {

} else if (i == j) {

} else if (i == k) {

} else if (j == k) {

long long singleDivisorTriplet(vector<int>& nums) {

vector<int> frequency(101, 0);

for (int number : nums) {

frequency[number]++;

for (int i = 1; i <= 100; ++i) {

for (int j = 1;  $j \le 100$ ; ++j) {

for (int k = 1;  $k \le 100$ ; ++k) {

int freq\_i = frequency[i];

int freq\_j = frequency[j];

int freq\_k = frequency[k];

int sum = i + j + k;

**if**  $(i == j \&\& i == k) {$ 

} else if (i == j) {

} else if (i == k) {

} else if (j == k) {

// Return the total count of single divisor triplets

// Return the total count of single divisor triplets

61 const sampleResult: number = singleDivisorTriplet(sampleNums);

performs a fixed series of modulo operations and comparisons.

cnt3, and space for the check function's return value).

} else {

// Return the computed answer.

return answer;

// Calculate the number of valid configurations.

// i and j are the same, but k is different.

// i and k are the same, but j is different.

// j and k are the same, but i is different.

answer += (long) countI \* countJ \* countK;

answer += (long) countI \* (countI - 1) \* (countI - 2) / 6;

answer += (long) countI \* (countI - 1) / 2 \* countK;

answer += (long) countI \* (countI - 1) / 2 \* countJ;

answer += (long) countI \* countJ \* (countJ - 1) / 2;

// All three numbers are the same.

// All three numbers are different.

// Counter array to keep the frequency of each number from 1 to 100

// Fill the counter with frequencies of numbers in the input vector nums

long long answer = 0; // Variable to store the final count of triplets

// Fetch the frequencies of the current combination

// Count the number of divisors the sum has from the triplet

// Now, handle the cases depending on the uniqueness of i, j, k

int divisorCount = (sum % i == 0) + (sum % j == 0) + (sum % k == 0);

// We are only interested in triplets where exactly one of the numbers divides the sum

answer += (long long) freq\_i \* freq\_j \* freq\_k; // Just the product of the counts

answer += (long long) freq\_i \* (freq\_i - 1) / 2 \* freq\_k; // Combination of nC2 times the count of k

answer += (long long) freq\_i \* (freq\_i - 1) / 2 \* freq\_j; // Combination of nC2 times the count of j

answer += (long long) freq\_j \* (freq\_j - 1) / 2 \* freq\_i; // Combination of nC2 times the count of i

answer += freq\_i \* (freq\_i - 1) \* (freq\_i - 2) / 6; // Combination of nC3

// Check all possible combinations of three numbers (i, j, k)

// Calculate the sum of the triplet

if (divisorCount != 1) continue;

// All numbers are the same

// i and j are the same, k is different

// i and k are the same, j is different

// j and k are the same, i is different

// i, j, and k are all different

continue;

return count\_triplets // 6

40 # result = solution.singleDivisorTriplet([nums])

int[] count = new int[101];

for (int num : nums) {

count[num]++;

38 # The above code can be used as follows:

def singleDivisorTriplet(self, nums: List[int]) -> int:

def is\_single\_divisor\_triplet(a: int, b: int, c: int) -> bool:

# Iterate over all possible combinations of a, b, and c

for c, count\_c in num\_counter.items():

if a == b and b == c:

if is\_single\_divisor\_triplet(a, b, c):

# All numbers are different

# Since each triplet is counted 6 times (all permutations), we divide by 6

// Initialize a counter array to count the occurrences of each number up to 100.

return sum(sum\_triplet % x == 0 for x in [a, b, c]) == 1

10 11 # Count the occurrences of each number in the list 12 num\_counter = Counter(nums) 13 count\_triplets = 0 # Initialize the count of valid triplets

# Check if the current combination is a single divisor triplet

count\_triplets += count\_a \* count\_b \* count\_c

# If any two numbers are equal, adjust the count accordingly

count\_triplets += count\_a \* (count\_a - 1) // 2 \* count\_b

count\_triplets += count\_a \* count\_b \* (count\_b - 1) // 2

# A helper function to check if only one of a, b, or c is a divisor of their sum

23 # All numbers are the same 24 count\_triplets += count\_a \* (count\_a - 1) \* (count\_a - 2) // 6 25 elif a == b: 26 count\_triplets += count\_a \* (count\_a - 1) // 2 \* count\_c 27 elif a == c:

# Java Solution

class Solution {

39 # solution = Solution()

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// Initialize a variable to store the answer.
            long answer = 0;
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            // Iterate through all possible trios and calculate the sum.
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            for (int i = 1; i \le 100; i++) {
                for (int j = 1; j \le 100; j++) {
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                    for (int k = 1; k \le 100; k++) {
                        // Get the count of each number in the trio.
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                        int countI = count[i];
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                        int countJ = count[j];
                        int countK = count[k];
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                        int sum = i + j + k;
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                        // Check if exactly one number divides the sum.
24
                        int divisorCount = 0;
25
                        divisorCount += sum % i == 0 ? 1 : 0;
26
                        divisorCount += sum % j == 0 ? 1 : 0;
27
                        divisorCount += sum % k == 0 ? 1 : 0;
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                        // If not exactly one divisor, then continue to the next trio.
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### 38 39 40 41 42 43

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             return answer;
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 57 };
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Typescript Solution
    const MAX_NUM = 100;
    // Function to count single divisor triplets
    function singleDivisorTriplet(nums: number[]): number {
         // Initialize a counter array to keep the frequency of each number from 1 to 100
         const frequency: number[] = new Array(MAX_NUM + 1).fill(0);
  8
         // Fill the counter with frequencies of numbers in the input array nums
         nums.forEach(number => {
  9
             frequency[number]++;
 10
         });
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 13
         let answer: number = 0; // Variable to store the final count of triplets
 14
 15
         // Check all possible combinations of three numbers (i, j, k)
 16
         for (let i = 1; i <= MAX_NUM; ++i) {</pre>
 17
             for (let j = 1; j <= MAX_NUM; ++j) {</pre>
                 for (let k = 1; k <= MAX_NUM; ++k) {</pre>
 18
 19
                     // Fetch the frequencies of the current combination
                     const freqI: number = frequency[i];
 21
                     const freqJ: number = frequency[j];
 22
                     const freqK: number = frequency[k];
 23
 24
                     // Calculate the sum of the triplet
 25
                     const sum: number = i + j + k;
 26
                     // Count the number of divisors the sum has from the triplet
                     const divisorCount: number = (sum % i === 0 ? 1 : 0)
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                                                   + (sum % j === 0 ? 1 : 0)
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                                                   + (sum % k === 0 ? 1 : 0);
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 31
                     // We are only interested in triplets where exactly one of the numbers divides the sum
 32
                     if (divisorCount !== 1) continue;
 33
 34
                     // Handle the cases depending on the uniqueness of i, j, k
 35
                     if (i === j && i === k) {
 36
                         // All numbers are the same (choose 3 from freqI)
 37
                         answer += freqI * (freqI - 1) * (freqI - 2) / 6;
 38
                     } else if (i === j) {
 39
                         // i and j are the same, k is different (choose 2 from freqI and multiply by freqK)
                         answer += freqI * (freqI - 1) / 2 * freqK;
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                     } else if (i === k) {
 42
                         // i and k are the same, j is different (choose 2 from freqI and multiply by freqJ)
 43
                         answer += freqI * (freqI - 1) / 2 * freqJ;
 44
                     } else if (j === k) {
 45
                         // j and k are the same, i is different (choose 2 from freqJ and multiply by freqI)
 46
                         answer += freqJ * (freqJ - 1) / 2 * freqI;
                     } else {
 47
 48
                         // i, j, and k are all different (multiply counts of freqI, freqJ, freqK)
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                         answer += freqI * freqJ * freqK;
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```

# The time complexity of the algorithm can be analyzed based on the nested loops and the operations performed within them. We iterate over all unique elements in nums up to three times due to the nested loops. If there are k unique elements, then the triple-

**Time Complexity:** 

**Space Complexity:** 

the triplet.

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return answer;

const sampleNums: number[] = [1, 2, 3];

Time and Space Complexity

62 console.log(sampleResult); // Output the result

59 // Sample usage

increments to ans involve multiplication, which is an O(1) operation. Therefore, the overall time complexity is  $0(k^3)$ .

nested loop runs 0(k^3) times. The check function is called inside the innermost loop, which operates in 0(1) time because it simply

For each unique triplet (a, b, c), we calculate the number of such triplets considering their counts in the original list. The

The given Python code aims to count the number of triplets in a list where exactly one of the three numbers is a divisor of the sum of

• The counter variable is a Counter object, which stores the frequency of each unique number in nums. It occupies O(k) space, where k is the number of unique numbers in nums. • Other than that, the space usage is constant 0(1) because we only use a fixed number of extra variables (ans, a, b, c, cnt1, cnt2,

The space complexity can be determined by the additional space used by the algorithm, which is not part of the input.

Thus, the overall space complexity is O(k).