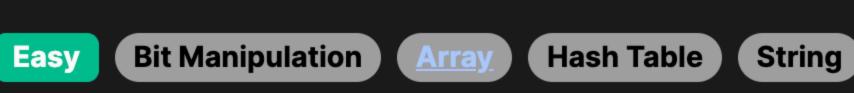
1684. Count the Number of Consistent Strings



Problem Description

The problem provides us with a string allowed which is made up of distinct characters, meaning no character is repeated in the string allowed. Additionally, we are given an array of strings called words. Our task is to determine how many strings in the array words are "consistent."

A string is defined as consistent if every character in the string appears in the allowed string. In other words, if there's even one character in a word that is not present in allowed, then that word is not considered consistent.

The goal is to return the number of consistent strings in the array words.

Intuition

the allowed string. To optimize this process, we convert the allowed string to a set. Sets in Python are collections that provide O(1) time complexity for checking if an element is present, which is faster than if allowed were a list or a string, where the time complexity would be O(n). Once we have the set of allowed characters, we iterate over each word in words. For each word, we check whether every

To solve this problem, the intuitive approach is to check each word in words and ensure all of its characters are contained within

character is in our set of allowed characters. We use the all() function which returns True if all elements of the iterable (the characters in the word) are True (or if the iterable is empty). The expression (c in s for c in w) is a generator that yields True or False for each character c in the word w depending on

whether c is in the set s or not. The all() function takes this generator as input and evaluates to True only if every character in the word is in the set of allowed characters.

Finally, we use the sum() function to count how many words are consistent by adding up 1 for every True result from the all()

check. **Solution Approach**

The first step is to convert the allowed string into a set:

s = set(allowed)

The solution provided uses Python's set and comprehension features to perform the task efficiently.

- Converting to a set is significant because set operations in Python are usually faster than list or string operations. Specifically,
 - checking for membership using the in operator is very fast for sets, taking O(1) time on average.

of consistent strings.

single, readable line of code.

Now, let's apply the solution step by step:

sum(all(c in s for c in w) for w in words) This comprehensive line does several things:

o for w in words: We start by going over each word in the words list. o all(c in s for c in w): For each word, we create a generator expression that yields True for each character c that is in the set s. The

Next, we use a list comprehension to iterate over each word in words:

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all() function checks if all values provided by this generator are True, meaning every character in the word is in the allowed set.
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• The entire all() expression will be True if the word is consistent and False otherwise. o sum(...): We are then using sum() to add up all the True values. Since True is equivalent to 1 in Python, sum() effectively counts the number

- This algorithm is very concise and takes advantage of Python's high-level operations to solve the problem with both simplicity and efficiency. The use of all() combined with a generator expression and sum() allows us to perform the entire operation in a
- **Example Walkthrough** Let's walk through a small example to illustrate the solution approach. Suppose the string allowed contains the characters "abc" and we have an array words containing the words ["ab", "ac", "de", "abc", "xy"].

We convert the allowed string into a set to expedite membership checks:

•

• "ab": Consistent

• "ac": Consistent

Python

s = set('abc') # This creates the set {'a', 'b', 'c'} Next, we iterate over each word in words and use the all() function to check if every character of a word is in the set s:

Let's break this down for each word: For the word "ab": the characters are 'a' and 'b', both of which are in the set s. So all(c in s for c in 'ab') will yield

True for both characters, and all() will return True.

result = sum(all(c in s for c in w) for w in words)

For the word "de": the character 'd' is not in the set s, so all(c in s for c in 'de') will yield False when checking 'd', and

all() will return False. For the word "abc": all the characters 'a', 'b', and 'c' are in the set s, and all() will return True.

For the word "xy": neither 'x' nor 'y' is in the set s, so all(c in s for c in 'xy') will yield False immediately when

For the word "ac": the characters are 'a' and 'c', both of which are in the set s. The all() function will also return True.

- checking 'x', and all() will return False. So, we get the following results for our words array:
- "de": Not consistent
- "abc": Consistent • "xy": Not consistent

Finally, the sum() function would add up all the True values (represented as 1 in Python):

Therefore, the output would be 3, as there are three consistent strings in the array words.

Solution Implementation

// Method to count number of consistent strings

boolean[] isAllowed = new boolean[26];

for (char c : allowed.toCharArray()) {

// Initialize count for consistent strings

// Iterate through each word in the array 'words'

// If the word is consistent, increment the count

if (!allowedLetters.test(ch - 'a')) return false;

// Iterate through each word and increment the consistent count if the word is consistent

// Counts the number of words from the 'words' array that contain only characters from the 'allowed' string.

// @param words - An array of strings, where each string consists only of lowercase English letters.

@param allowed - A string consisting of distinct lowercase English letters, representing the allowed characters.

// @return The count of words from the 'words' array that are "consistent" - contain only characters from 'allowed'.

// All characters in the word are allowed

// Return the final count of consistent strings

return true;

return consistentCount;

for (std::string& word : words) {

if (isConsistent(word)) {

++consistentCount;

};

isAllowed[c - 'a'] = true;

for (String word : words) {

// Array to keep track of allowed characters

result = sum([True, True, False, True, False]) # This equals 3

class Solution: def countConsistentStrings(self, allowed: str, words: list[str]) -> int: # Convert the allowed string into a set for O(1) lookup times allowed_chars = set(allowed)

Count the words in which all the characters are in the allowed set

public int countConsistentStrings(String allowed, String[] words) {

Summing up the boolean values where True is counted as 1, False as 0

// Populate the isAllowed array with characters from the 'allowed' string

consistent_words_count = sum(all(char in allowed_chars for char in word) for word in words)

return consistent_words_count # Return the total count of consistent words Java

int count = 0;

class Solution {

```
if (isConsistent(word, isAllowed)) {
                count++;
        // Return the total count of consistent strings
        return count;
    // Helper method to check if a word is consistent
    private boolean isConsistent(String word, boolean[] isAllowed) {
        // Iterate through each character in the word
        for (int i = 0; i < word.length(); i++) {</pre>
           // If the character is not in the allowed list, return false
            if (!isAllowed[word.charAt(i) - 'a']) {
                return false;
        // Return true if all characters of the word are allowed
        return true;
C++
#include <vector>
#include <string>
#include <bitset>
class Solution {
public:
    int countConsistentStrings(std::string allowed, std::vector<std::string>& words) {
        // Initialize a bitset to store whether each letter in the alphabet is allowed
        std::bitset<26> allowedLetters;
        for (char ch : allowed) {
            allowedLetters.set(ch - 'a');
       // Counter for the number of consistent strings
        int consistentCount = 0;
        // Lambda function to check if all characters in a word are allowed
        auto isConsistent = [&](std::string& word) {
           for (char ch : word) {
                // If any character is not allowed, return false immediately
```

TypeScript

};

```
function countConsistentStrings(allowed: string, words: string[]): number {
      // Converts a string to a bitmask where each bit set represents the presence of a character in the string.
      // @param str - A string to convert to a bitmask.
      // @return A bitmask representing the characters of the input string.
      const convertToBitmask = (str: string): number => {
          let bitmask = 0;
          for (const char of str) {
              bitmask |= 1 << (char.charCodeAt(0) - 'a'.charCodeAt(0));</pre>
          return bitmask;
      };
      // Create a bitmask for the allowed characters.
      const allowedMask = convertToBitmask(allowed);
      let consistentCount = 0; // Initialize the count of consistent strings.
      // Loop through each word in the array.
      for (const word of words) {
          // If the bitmask of the word OR'd with the allowedMask equals the allowedMask,
          // it means the word contains only characters from 'allowed'.
          if ((allowedMask | convertToBitmask(word)) === allowedMask) {
              consistentCount++; // Increment count as the word is consistent.
      // Return the total count of consistent strings.
      return consistentCount;
   def countConsistentStrings(self, allowed: str, words: list[str]) -> int:
       # Convert the allowed string into a set for O(1) lookup times
        allowed_chars = set(allowed)
       # Count the words in which all the characters are in the allowed set
       # Summing up the boolean values where True is counted as 1, False as 0
        consistent_words_count = sum(all(char in allowed_chars for char in word) for word in words)
        return consistent_words_count # Return the total count of consistent words
Time Complexity
  The time complexity of the function depends on two factors: the length of the words list meaning the number of words it contains
```

class Solution: Time and Space Complexity

(let's denote this number as n) and the average length of the words (let's denote it as k). The function iterates over each word

and then over each character in the word to check if it is in the set s. Checking membership in a set is an 0(1) operation on

average. Therefore, the time complexity for checking one word is O(k), and for all the words it's O(n * k).

Space Complexity

The space complexity is O(s) where s is the number of unique characters in the allowed string because these are stored in a set. The other variable that could contribute to space complexity is win the generator expression, but since the space required for wis reused as the iteration proceeds and since the strings in words are input to the function and not duplicated by it, this does not add to the space complexity of the solution itself. The space complexity for the output of the sum function is 0(1) since it's accumulating the result into a single integer. Therefore, the overall space complexity is 0(s) + 0(1) which simplifies to 0(s).