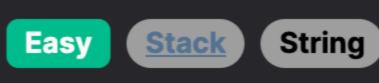
#### 20. Valid Parentheses



### **Problem Description**

The problem presents a scenario in which we are given a string s consisting of six possible characters: the opening and closing brackets of three types—parentheses (), square brackets [], and curly braces {}. The challenge is to determine whether this string represents a sequence of brackets that is considered valid based on certain rules. A string of brackets is deemed valid if it satisfies the following conditions:

- 1. Each opening bracket must be closed by a closing bracket of the same type.
- 2. Opening brackets must be closed in the correct order. That means no closing bracket should interrupt the corresponding pair of an opening bracket. 3. Each closing bracket must have an associated opening bracket of the same type before it.
- For example, a string "(())" is valid because each opening parenthesis (has a corresponding closing parenthesis) that occurs later

in the string, and they are properly nested and ordered.

#### The intuition behind the solution utilizes a common data structure known as a stack, which operates on a last-in, first-out (LIFO)

Intuition

principle. This principle mimics the necessary behavior for tracking opening brackets and ensuring they are closed in the correct order. The steps followed in the solution are: 1. Initialize an empty stack to keep track of opening brackets.

3. When an opening bracket is encountered ('(', '[', or '{'), push it onto the stack. This represents waiting for a closing bracket

2. Traverse the input string character by character.

- to match. 4. When a closing bracket is encountered (')', ']', or '}'), check if it forms a pair with the last opening bracket added to the stack
- (i.e., on top of the stack). 5. If the stack is empty (no opening bracket available to match) or the closing bracket does not form a valid pair with the opening bracket on top of the stack, we know the string is invalid, and we immediately return False.
- 6. If we successfully traverse the entire input string and the stack is empty, meaning all opening brackets have been matched
- correctly, we return True. If the stack is not empty, it indicates there are unmatched opening brackets, and therefore, the string is invalid. The concise implementation of this algorithm ensures that both the correct type and order of brackets are validated for the string to

**Solution Approach** 

The solution utilizes a stack data structure to track opening brackets and ensure they have the appropriate closing brackets. Let's

#### walk through the implementation step by step:

be considered valid.

2. Pairs Set: We define a set d containing string representations of the valid bracket pairs: '()', '[]', and '{}'. This helps quickly check if an encountered closing bracket correctly pairs with the last opening bracket on the stack.

o Opening Bracket: If c is an opening bracket ('(', '[', or '{'}), it is pushed onto stk, waiting for the corresponding closing

3. Iterate Over the String: The algorithm iterates over each character c in the input string s.

1. Initialize a Stack: A list named stk is created to serve as a stack.

implies all opening brackets were properly matched and closed.

- bracket. Closing Bracket: If c is a closing bracket (')', ']', or '}'):
  - The stack is checked to ensure it's not empty, which would mean there's no opening bracket to match the closing one.
    - If the stack is not empty, the top element is popped. We concatenate it with c to check if they form a valid pair by checking against set d.
    - If either condition fails the stack was empty or the concatenation of the popped element with c does not form a valid pair – the function immediately returns False because we've detected an invalid bracket sequence.
- If the stack is empty (return not stk), it indicates a valid bracket sequence and returns True.

If the stack is not empty, some opening brackets were not closed, thus the function returns False, representing an invalid

4. Final Stack Check: After processing all characters in the string, the algorithm checks if the stack is empty. An empty stack

- bracket sequence. The solution is efficient with a linear time complexity, O(n), where n is the length of the string s. It only requires a single pass through
- the string and constant-time operations for each character. The space complexity is also O(n), in the worst case where the string consists entirely of opening brackets, which would all be pushed onto the stack.

Example Walkthrough Let's consider a small example to clearly understand the solution approach. Suppose we are given the string  $s = "\{[()()]\}$ ". We want to determine if this string represents a valid sequence of brackets. Following the solution steps:

#### 1. Initialize a Stack: We start with an empty list stk that will be used as our stack. 2. Pairs Set: We have a predefined set d that contains '()', '[]', and '{}' to represent valid bracket pairs.

pair (). So far so good.

3. Iterating Over the String: We iterate through each character in the string s.

• For the second character [, it's also an opening bracket, so we push it onto stk. For the third character (, another opening bracket gets pushed onto stk.

For the first character {, it's an opening bracket, so we push it onto stk.

- The fourth character ) is a closing bracket, so we pop the last element (, which matches the closing bracket, forming a valid
- The fifth character ( is once again an opening bracket and is pushed onto stk. • The sixth character ) is a closing bracket, and popping from stk gives us the matching (, forming another valid pair. We

# If the character is an opening parenthesis, push it onto the stack

# If the stack is empty, all parentheses were valid and correctly nested

(leftBracket == '{' && rightBracket == '}') ||

(leftBracket == '[' && rightBracket == ']');

// Function to check if a given string has valid parentheses

// Using a string as a stack to store opening brackets

// Iterate through all characters in the input string

if (c == '(' || c == '{' || c == '[') {

stack.push\_back(c);

return false;

stack.pop\_back();

// If the character is an opening bracket, push it onto the stack

// If the stack is empty or characters don't match, return false

// If the characters match, pop the opening bracket from the stack

else if (stack.empty() || !match(stack.back(), c)) {

continue this process for the remaining characters.

• The seventh character ] is a closing bracket, and popping from stk gives us [, which is the correct complement, forming the pair [].

• Finally, the last character } is a closing bracket, and after popping from stk we get {, its correct opening pair, forming {}.

closing bracket in the proper order.

4. Final Stack Check: At the end of iteration, stk is empty because every opening bracket has been matched with the correct

has correctly determined the validity by using a stack to manage the ordering and pairing of the brackets.

Since the stack is empty, we return True, indicating that the given string  $s = "\{[()()]\}$ " is a valid sequence of brackets. Our solution

stack = [] # Create a set with valid parentheses pairs valid\_pairs = {'()', '[]', '{}'} # Iterate over each character in the string

```
if char in '({[':
11
                    stack.append(char)
12
               # If the stack is empty or the formed pair is not valid, return False
13
               elif not stack or stack.pop() + char not in valid_pairs:
14
                    return False
15
```

**Java Solution** 

for char in s:

return not stack

Python Solution

def isValid(self, s: str) -> bool:

# Initialize an empty list to use as a stack

1 class Solution:

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C++ Solution

1 class Solution {

bool isValid(string s) {

for (char c : s) {

string stack;

else {

2 public:

```
1 class Solution {
       // Method to determine if an input string has valid parentheses
       public boolean isValid(String s) {
           // Use a deque as a stack to keep track of the opening brackets
           Deque<Character> stack = new ArrayDeque<>();
           // Iterate over each character in the input string
 6
           for (char c : s.toCharArray()) {
               // If the current character is an opening bracket, push it onto the stack
8
               if (c == '(' || c == '{' || c == '[') {
9
10
                   stack.push(c);
               } else {
11
12
                   // If the stack is empty or the current closing bracket doesn't match
13
                   // the last opening bracket, the string is not valid
                   if (stack.isEmpty() || !match(stack.pop(), c)) {
14
15
                       return false;
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           // If the stack is empty, all brackets were matched correctly
20
           return stack.isEmpty();
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       // Helper method to determine if two brackets are a matching pair
24
       private boolean match(char leftBracket, char rightBracket) {
25
           // Return true if pairs match, false otherwise
26
           return (leftBracket == '(' && rightBracket == ')') ||
```

#### 24 // If the stack is empty, all brackets were properly closed, return true 25 26 27 28

```
return stack.empty();
       // Utility function to check if the opening and closing brackets match
29
       bool match(char left, char right) {
            return (left == '(' && right == ')') |
30
                   (left == '[' && right == ']') ||
31
                   (left == '{' && right == '}');
32
33
34 };
35
Typescript Solution
 1 // Map that associates opening braces with their corresponding closing braces.
   const bracketPairs = new Map<string, string>([
        ['(', ')'],
       ['[', ']'],
['{', '}'],
    1);
    /**
    * Function to determine if the given string has valid bracket pairing.
    * @param {string} str - The input string containing brackets to be checked.
    * @return {boolean} - Returns true if the string is valid, otherwise false.
12
    */
   function isValid(str: string): boolean {
       // A stack to keep track of the expected closing brackets.
14
       const expectedBracketsStack: string[] = [];
15
16
       // Iterate over each character in the string.
17
       for (const char of str) {
18
19
           // Check if the character is an opening bracket and get its pair.
           if (bracketPairs.has(char)) {
20
               // Push the expected closing bracket onto the stack.
21
22
                expectedBracketsStack.push(bracketPairs.get(char)!);
23
           } else {
24
               // If the character is not a matching closing bracket,
25
               // or the stack is empty (mismatched brackets), return false.
26
               if (expectedBracketsStack.pop() !== char) {
27
                    return false;
```

### 34 } 35

return expectedBracketsStack.length === 0;

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each character in the input string exactly once.

• Best Case: When the string is empty or consists of a single pair of brackets, the time complexity is 0(1) because it takes a constant amount of time.

The space complexity of the code is also O(n), as in the worst case (when all characters in the input string are opening brackets), the

# **Time Complexity:**

Time and Space Complexity The time complexity of the given code is O(n), where n is the length of the input string. This is because the algorithm iterates over

// If the stack is empty, all brackets were properly closed; otherwise, return false.

- Average Case: For a typical string with a mix of opening and closing brackets, the time complexity remains 0(n) because each
  - character is processed once. • Worst Case: In the worst case scenario, where there are n characters and the string is properly nested to the deepest level, each character is still processed exactly once, giving us a time complexity of O(n).

## stack stk will contain all characters in the input string.

- **Space Complexity:**
- The space complexity is O(n) which occurs when all characters are opening brackets and hence, all are pushed onto the stack. This represents the maximum space utilized by the algorithm. As the stack grows with each opening bracket and shrinks with each closing bracket, the actual space used depends on the

string or a string with just one pair of brackets) to 0(n) (when all characters are opening brackets).

number of unmatched opening brackets at any point in the algorithm. Therefore, space usage can vary from 0(1) (for an empty