2086. Minimum Number of Food Buckets to Feed the Hamsters



Greedy

Medium

String

**Dynamic Programming** 

In this problem, we're given a string hamsters, which represents a row of positions that can either contain a hamster, denoted by 'H', or be empty, denoted by '.'. We need to place food buckets at some of these empty positions. There are two conditions to ensure that a hamster can be fed: • There must be at least one food bucket immediately to the left (i - 1) or immediately to the right (i + 1) of the hamster's position.

• We want to place the minimum number of food buckets necessary to feed all the hamsters.

Our task is to find this minimum number of food buckets or determine that it's impossible to feed all hamsters, in which case we return -1.

Intuition

number while ensuring all hamsters can access at least one bucket. This optimality comes from the following rule: whenever we have a choice to place a food bucket either to the left or the right of a hamster, we prefer to place it to the right. Doing so might enable us to feed the next hamster with the same bucket, hence reducing the total number.

The solution approach is to iterate through the string and decide the placement of the food buckets optimally to minimize their

• If the right side is not available, we check the left side. If it's empty and we haven't already placed a bucket there while placing it for a previous

Here's the approach in detail:

hamster, we place a bucket there and increment the answer by 1. If neither side is available for placing a bucket, we return −1 as it's impossible to feed this hamster.

right. We increment the answer by 1 and skip checking the next position by advancing two steps in the iteration.

• We iterate through each position in the string. If we find a hamster ('H'), we check its neighbouring positions.

The critical insight is that by considering the right side first, we utilize each bucket to its maximum potential, thereby ensuring the

• If we have an empty position to the right of the hamster, we always place a bucket there, as it could potentially feed another hamster to the

- least number of buckets is used. **Solution Approach**
- The solution uses a greedy algorithm, which is a common approach for optimization problems where we aim to make the locally

Start iterating through the string street using variable i that begins from index 0.

optimal choice at each stage with the hope of finding a global optimum.

Here's a step-by-step walkthrough of the implementation based on the reference solution: Initialize a counter ans to 0, which will hold the minimum number of food buckets required.

Check if the current character at index i is a hamster ('H'). If there is a hamster at index i and the next index i + 1 is within bounds and empty (contains '.'), place a food bucket

at i + 1. This is because placing the bucket to the right could potentially also feed a hamster at index i + 2.

■ If a bucket can be placed at i - 1, increase ans by 1. There's no need to move i forward by two positions in this case since we are

If neither the left nor the right side of the current hamster can have a bucket placed, it is impossible to feed this hamster,

■ Move i two positions forward to i + 2, as we have already dealt with hamster at i and potentially at i + 1 if there is one. If you can't place the bucket to the right because i + 1 is out of bounds or not empty, check if the left of the hamster at

only dealing with the current hamster.

minimum number of buckets required. We return ans.

space complexity since we only use a fixed number of extra variables.

Assume the input string representing the row of positions is: "H..HH..H.."

At index 3, again, there's no hamster; thus, we move to index 4.

Following the solution approach, we'll determine where to place the food buckets:

At index 2, there's no hamster, just a food bucket, so we move to the next index.

so we return -1.

• After placing the bucket, increase ans by 1 to count the newly placed bucket.

Continue this process until we have iterated over the entire string. The ans variable by the end of iteration corresponds to the

If the current character is not a hamster, we simply move to the next index by incrementing i.

i is available (make sure you're not out of bounds and the index i - 1 is empty).

**Example Walkthrough** Let's go through an example to illustrate the solution approach.

No additional data structures are used apart from variables for indexing and counting. This makes the algorithm efficient both in

terms of time and space complexity, achieving (O(n)) time complexity, where (n) is the length of the street string, and (O(1))

Starting at index 0, we see 'H', the position contains a hamster.  $\circ$  **i** + 1 is within bounds and empty, so we place a bucket at **i** + 1 (position 1). • We increment ans by 1 (now ans is 1). We skip the next position by moving two positions forward to index 2.

## At index 4, there's a hamster 'H'. $\circ$ **i** + 1 is within bounds and empty, so we put a bucket at **i** + 1 (position 5).

Hence, we return 3.

class Solution:

Solution Implementation

We move two positions forward to index 6.

At index 6, there's no hamster; we move to index 7.

• We increment ans by 1 (now ans is 2).

At index 7, there's another hamster 'H'.

We place a bucket at i - 1 (position 8).

def minimumBuckets(self, street: str) -> int:

# If we encounter a house 'H'

if street[index] == 'H':

index += 2

return -1

public int minimumBuckets(String street) {

int streetLength = street.length();

for (int i = 0; i < streetLength; ++i) {</pre>

if (street.charAt(i) == 'H') {

++bucketCount:

++bucketCount;

// Return the minimum number of buckets needed

return -1;

function minimumBuckets(street: string): number {

let streetSize: number = street.length;

// Iterate through the 'street' string

for (let i = 0; i < streetSize; i++) {</pre>

bucketCount++;

bucketCount++;

index, length = 0, len(street)

# If we encounter a house 'H'

if street[index] == 'H':

index += 2

return -1

bucket count += 1

bucket\_count += 1

# Move to the next position in street

while index < length:</pre>

else:

index += 1

return -1;

let bucketCount: number = 0;

if (street[i] === 'H') {

i += 2;

} else {

return bucketCount;

// Variable 'streetSize' holds the size of the 'street' string

if (i + 1 < streetSize && street[i + 1] === '.') {</pre>

} else if (i > 0 && street[i - 1] === '.') {

// Return the total minimum number of buckets required

// position has already been considered

# Initialize the index and get the length of the street

if index + 1 < length and street[index + 1] == '.':</pre>

# If no space on the right, check for space on the left

# Return the total number of buckets needed after placing them optimally

# Loop through each character in the street string

# Increment the bucket count

# Increment the bucket count

elif index > 0 and street[index - 1] == '.':

// Initialize 'bucketCount' to count the minimum number of buckets required

// Check if there is a spot for a bucket to the right of the house

// Check if there is a spot for a bucket to the left of the house

// If there are no spots to place a bucket around the house,

// it's not possible to water all houses, so return -1

// Place a bucket to the right of the house, increase the bucket count,

// Place a bucket to the left of the house, increase the bucket count

// This does not require skipping a position because the previous

// and skip the next position because it is covered by the bucket

i += 2:

} else {

bucket count += 1

bucket\_count += 1

# Move to the next position in street

// Get the length of the street represented as a string

// Initialize the count of buckets required to water flowers

// Check if there is a house in the current position

index, length = 0, len(street)

while index < length:</pre>

else:

index += 1

return bucket\_count

int bucketCount = 0;

# Initialize the number of buckets needed

 i + 1 is within bounds but contains another hamster, so we can't place a bucket there. ○ We look to the left (i - 1), but there's already a food bucket there placed for the previous hamster, so no new bucket is needed here.

At index 9, there's a hamster 'H'.

 $\circ$  i + 1 is out of bounds, so we look at i - 1, which is within bounds and empty.

We move to the next index 8 and find it empty. Move to index 9.

- We increment ans by 1 (now ans is 3).
- **Python**

# Check if there is space on the right to place a bucket and avoid out of range error

# If no space available to place a bucket, it's not possible to water the house

# Skip the next house as the current bucket will water this house

// If there is an empty space next to the house, place the bucket there

// Skip the next position as we've already placed the bucket

// If there is an empty space before the house, place the bucket there

// If there are no empty spaces around the house, it's impossible to water

if (i + 1 < streetLength && street.charAt(i + 1) == '.') {</pre>

} else if (i > 0 && street.charAt(i - 1) == '.') {

Having gone through the string, we needed a minimum of 3 food buckets to feed all hamsters in the given row of positions.

bucket\_count = 0 # Initialize the index and get the length of the street

if index + 1 < length and street[index + 1] == '.':</pre>

# Return the total number of buckets needed after placing them optimally

# Loop through each character in the street string

# Increment the bucket count

# If no space on the right, check for space on the left elif index > 0 and street[index - 1] == '.': # We don't skip the index here because the bucket waters the house on the left # Increment the bucket count

```
// Iterate over each position in the street
```

Java

class Solution {

```
return bucketCount;
C++
class Solution {
public:
    int minimumBuckets(string street) {
        // Variable 'n' holds the size of the street string
        int streetSize = street.size();
        // Initialize 'bucketCount' to count the minimum number of buckets required
        int bucketCount = 0;
        // Iterate through the street string
        for (int i = 0; i < streetSize; ++i) {</pre>
            // Check if the current position is a house 'H'
            if (street[i] == 'H') {
                // Check if there is a spot for a bucket to the right of the house
                if (i + 1 < streetSize && street[i + 1] == '.') {</pre>
                    // Increase the bucket count and skip the next spot since it's covered by the bucket
                    ++bucketCount;
                    i += 2;
                // Check if there is a spot for a bucket to the left of the house
                } else if (i > 0 && street[i - 1] == '.') {
                    // Increase the bucket count
                    ++bucketCount;
                } else {
                    // If there are no spots to place a bucket around the house, return —1 indicating it's not possible
                    return -1;
        // Return the total minimum number of buckets required
        return bucketCount;
};
TypeScript
// Function to compute the minimum number of buckets required
```

## class Solution: def minimumBuckets(self, street: str) -> int: # Initialize the number of buckets needed bucket\_count = 0

return bucket\_count Time and Space Complexity **Time Complexity** The given code involves a single loop through the length of the string street. Within this loop, the operations are constant time

checks and assignments. The loop iterates at most n times, where n is the length of street. In the worst case, every character

# Check if there is space on the right to place a bucket and avoid out of range error

# We don't skip the index here because the bucket waters the house on the left

# If no space available to place a bucket, it's not possible to water the house

# Skip the next house as the current bucket will water this house

is visited once, and update operations are constant time 0(1). Therefore, the time complexity is 0(n).

**Space Complexity** The space complexity of the code is 0(1). This is because the code only uses a fixed number of extra variables (ans, i, and n) that do not depend on the size of the input street. There are no data structures used that grow with the input size. Hence, the space required does not scale with the size of the input.