2342. Max Sum of a Pair With Equal Sum of Digits

Sorting Heap (Priority Queue) Medium <u>Array</u> Hash Table

Problem Description

where i is not equal to j, such that the sum of the digits of nums[i] is equal to the sum of the digits of nums[j]. Once such indices are found, we need to return the maximum value of the sum nums[i] + nums[j] that we can obtain by considering all possible pairs of indices that meet the given condition. To clarify, if we pick nums[i] as "123" and nums[j] as "51", both have the sum of digits equal to 6 and thus meet the criteria

This problem presents us with an array nums filled with positive integers, and our goal is to find two different indices i and j,

since 1+2+3 = 5+1.

The immediate brute-force approach would be to compute the sum of digits for every pair of numbers in the array and then find

Intuition

for large arrays. The given solution leverages hashing to optimize the process. It uses a dictionary (or hash map) to keep track of the highest value number v for each unique sum of digits y. This way, when a new number is processed, we can easily check if there is

the maximum pair that meets the criteria. However, this would be inefficient with a time complexity of O(n^2) which is not optimal

already a stored number with the same sum of digits using the hash map. While iterating through the **nums** array, for each number: 1. We calculate the sum of digits y.

 Calculate a potential new maximum by adding the current number v and the stored number d[y]. Compare it with the current known maximum ans and update ans if the new potential maximum is greater.

- 3. Regardless of whether y was already present or not, we update the dictionary with the highest value v for the sum of digits y. We use max(d[y], v) to ensure we always keep the larger number for that sum of digits, which is crucial for maximizing the eventual sum of nums[i] +

2. If y is already in the dictionary d, it means we have encountered another number previously with the same sum of digits. We can then:

nums[j]. The result of this process is the maximum sum that can be formed under the given constraints, which the function returns. If no

default values for non-existing keys. The algorithm proceeds as follows:

loop effectively extracts and sums up each digit of the number.

is assuming that the hash map operations (insert and lookup) are O(1) on average.

such pair exists, the answer remains as the initialized value of -1. **Solution Approach**

The implementation of the solution makes use of a hash map to efficiently track the maximum number encountered for each unique sum of digits. The Python code utilizes a defaultdict from the collections module which simplifies the management of

1. Initialize ans as -1. This will store the eventual maximum sum nums[i] + nums[j] if a valid pair is found. If no valid pairs are found, ans will

earlier. In this case:

return -1. 2. Initialize a defaultdict named d to store pairs of (sum of digits: maximum number with that sum). 3. Iterate over each number v in the input nums array. Calculate the sum of the digits y of the current number v using a while loop that adds v % 10 to y and then floor divides v by 10. This

o If the sum of digits y is already a key in the dictionary d, then there exists a different number with the same sum of digits encountered

Compute the possible new maximum sum d[y] + v and compare it with ans. Update ans with this new maximum if it is greater.

down to O(n * k), where n is the number of elements in nums and k is the average number of digits in numbers within nums. This

At the end of the loop, ans contains the maximum sum of nums[i] + nums[j] where sums of their digits are equal or remains -1

- Update the dictionary with the key y by ensuring it stores the maximum v: d[y] = max(d[y], v). This step is crucial since it is possible to encounter multiple numbers with the same sum of digits and we are only interested in storing the largest one for the pair-wise comparison.
- The data structure used is crucial for optimizing the time complexity of this problem. By using a hash map, accessing and updating the maximum number for a sum of digits is done in constant time, bringing the overall time complexity of the algorithm

if no such pair exists. This value is then returned as the answer. Example Walkthrough Let's consider a small example using the array nums = [42, 33, 60].

2. Initialize the hash map d as a defaultdict with the default type as int. It will hold the sum of digits as the key and the respective maximum number as the value. 3. Start with the first number 42. The sum of its digits is 4 + 2 = 6. Since there is no entry in d with 6 as key, add it with d[6] = 42.

(which is 42) and 33 is greater than ans. The current sum is 42 + 33 = 75, so we update ans = 75. Then we update d[6] to be the max of

4. Move to the next number $\frac{33}{3}$. The sum of its digits is $\frac{3}{7} + \frac{3}{7} = \frac{6}{3}$. There is already an entry with the sum 6, so we check if the sum of $\frac{d[6]}{d[6]}$

new maximum number 60.

Solution Implementation

1. Initialize ans as -1.

5. Proceed to the last number 60. The sum of its digits is 6 + 0 = 6. Again, there's an entry for sum 6. We compute the potential new maximum

Initialize the maximum sum as -1 (assuming no answer is found yet)

check if the current number contributes to a larger max sum

// Return the maximum pair sum of numbers with the same digit sum, else -1.

// Function to calculate the maximum sum of a pair of numbers with the same sum of digits

int maxPairSum = -1; // Initialize max pair sum as -1 to handle cases with no valid pair

// Update the maxPairSum with the sum of the top two numbers in the current group

// Check if there are at least two numbers in the current digit sum group

// Sort the numbers within the current group in descending order

// Iterate through each number to calculate their digit sums and group them

// Calculate the sum of digits for the current number

// Add the number to its corresponding digit sum group

maxPairSum = max(maxPairSum, group[0] + group[1]);

for (int value = number; value > 0; value /= 10) {

digitSumGroups[digitSum].emplace_back(number);

sort(group.rbegin(), group.rend());

Calculate the sum of digits for the current number

If the sum of digits has been seen before,

if digit sum in digit sum max num:

42 and 33, which is 42, so no change is required.

def maximumSum(self, nums: List[int]) -> int:

- which is d[6] (42) plus 60 equals 102, and since it is greater than the current ans (75), we update ans = 102. We then update d[6] with the
- At the end of this process, ans holds the value 102, which is the maximum sum of nums[i] + nums[j] with equal digit sums. Since no other pairs are left to be considered, we would return 102 as the result.
- **Python** from collections import defaultdict class Solution:

Initialize a dictionary to store the maximum number for each digit sum digit_sum_max_num = defaultdict(int) # Iterate through each number in the given list

while temp num: digit sum += temp_num % 10 temp_num //= 10

return maxPairSum;

int maximumSum(vector<int>& numbers) {

for (int& number : numbers) {

int digitSum = 0;

vector<vector<int>> digitSumGroups(100);

digitSum += value % 10;

// Iterate through all digit sum groups

for (auto& group : digitSumGroups) {

// Return the maximum pair sum found

return maxPairSum;

class Solution:

from collections import defaultdict

for num in nums:

return max_sum

Time and Space Complexity

complexity for this part is 0(d * n).

Space Complexity

digit sum = 0

temp num = num

while temp num:

def maximumSum(self, nums: List[int]) -> int:

digit_sum_max_num = defaultdict(int)

Iterate through each number in the given list

digit sum += temp_num % 10

};

if (group.size() > 1) {

for num in nums:

digit sum = 0

temp num = num

 $\max sum = -1$

```
max_sum = max(max_sum, digit_sum_max_num[digit_sum] + num)
            # Update the maximum number for the current digit sum
            digit_sum_max_num[digit_sum] = max(digit_sum_max_num[digit_sum], num)
        # Return the maximum sum of two numbers having the same sum of digits
        return max_sum
Java
class Solution {
    public int maximumSum(int[] nums) {
        // This variable will hold the answer, initialized to -1 as per the problem statement.
        int maxPairSum = -1;
        // This arrav will store the maximum number encountered for each digit sum.
        int[] maxNumWithDigitSum = new int[100];
        // Iterate through all the numbers in the input array.
        for (int number : nums) {
            int sumOfDigits = 0;
            // Calculate the sum of digits of the current number.
            for (int tempNumber = number; tempNumber > 0; tempNumber /= 10) {
                sumOfDigits += tempNumber % 10;
            // If there's already a number with the same digit sum encountered,
            // check if the two numbers form a larger pair sum.
            if (maxNumWithDigitSum[sumOfDigits] > 0) {
                maxPairSum = Math.max(maxPairSum, maxNumWithDigitSum[sumOfDigits] + number);
            // Update the array with the maximum number for the current digit sum.
            maxNumWithDigitSum[sumOfDigits] = Math.max(maxNumWithDigitSum[sumOfDigits], number);
```

// Initialize a vector of vectors to group numbers by their digit sums (up to 81 for 99999, which is maximum 5*9)

C++

public:

class Solution {

```
return maxPairSum;
};
TypeScript
// Define the maxDigits value representing the maximum possible sum of digits which is 9*5=45
const MAX_DIGITS_SUM = 45;
// Function to calculate the sum of the digits of a number
const sumOfDigits = (number: number): number => {
    let digitSum = 0;
    while (number > 0) {
        digitSum += number % 10;
        number = Math.floor(number / 10);
    return digitSum;
// Function to calculate the maximum sum of a pair of numbers with the same sum of digits
const maximumSum = (numbers: number[]): number => {
    // Initialize an array to group numbers by their digit sums
    const digitSumGroups: number[][] = Array.from({ length: MAX_DIGITS_SUM + 1 }, () => []);
    // Group numbers by their digit sum
    numbers.forEach(number => {
        const digitSum = sumOfDigits(number);
        digitSumGroups[digitSum].push(number);
    });
    let maxPairSum = -1; // Initialize max pair sum as -1 to indicate no valid pair yet
    // Iterate through all digit sum groups to find the maximum pair sum
    digitSumGroups.forEach(group => {
        if (group.length > 1) {
            // Sort the aroup in descending order
            group.sort((a, b) => b - a);
            // Update the maxPairSum with the sum of the two largest numbers in the current group
            maxPairSum = Math.max(maxPairSum, group[0] + group[1]);
    });
    // Return the maximum pair sum found, or -1 if no valid pairs exist
```

temp_num //= 10 # If the sum of digits has been seen before, # check if the current number contributes to a larger max sum if digit sum in digit sum max num:

Update the maximum number for the current digit sum

Calculate the sum of digits for the current number

// The functions can now be used globally as part of the TypeScript codebase

Initialize the maximum sum as -1 (assuming no answer is found yet)

Initialize a dictionary to store the maximum number for each digit sum

max_sum = max(max_sum, digit_sum_max_num[digit_sum] + num)

Return the maximum sum of two numbers having the same sum of digits

digit_sum_max_num[digit_sum] = max(digit_sum_max_num[digit_sum], num)

```
Time Complexity
  The time complexity of the given code can be analyzed by considering two major parts:
 1. Iterating through each number in the nums list.
 2. Summing the digits of each number and updating the dictionary.
  For the first part, iterating through the nums list is O(n) where n is the length of nums.
```

Combining both parts, the total time complexity is 0(n * d). Note that if the input numbers have a bounded size (for example, they are all 32-bit integers), d can be considered a constant, and the time complexity can be viewed as O(n).

For the second part, we must consider the number of digits in each number for the summation of digits and updating the

dictionary. The number of digits d in a number v is proportional to log10(v). Hence, summing the digits of a number takes 0(d)

time where d is the number of digits. Since for each number, we perform a digit sum and update the dictionary, the overall time

The space complexity is determined by the space required to store the d dictionary which holds the maximum number encountered for each digit sum.

• The number of unique digit sums is at most 9*m, where m is the maximum number of digits in a number of the nums list since the largest digit sum for a number with m digits would be 9 * m (if each digit is 9). Each entry in the dictionary stores an integer value (which is constant space).

Thus, the space complexity is O(m), where m represents the maximum number of digits across all numbers in nums.

If we again assume the input numbers are all 32-bit integers, m is a constant (10, since 2^31 - 1 has 10 digits), making the space complexity 0(1).