128. Longest Consecutive Sequence

Hash Table

Problem Description

Union Find Array

The problem asks us to find the length of the longest sequence of consecutive numbers in an unsorted array called nums. A

Medium

The tricky part of this problem is that we need to come up with a solution that has a time complexity of O(n). This means we cannot afford the luxury of sorting the array as it would typically require 0(n * log n) time. Thus, we must find a way to keep track of sequences efficiently, despite the order of elements being arbitrary.

is a consecutive sequence, but [1, 3, 4] is not. Our task is to find the longest such sequence in the given array.

sequence of numbers is considered consecutive if every number follows the previous one without any gaps. For example, [1, 2,

Intuition

To solve this problem in O(n) time, we need to think of a data structure that allows us to quickly check if an element exists in the set and if we can extend a consecutive sequence. A hash table, or in Python a set, is an ideal candidate because it allows us to

Here's the intuition for the solution approach:

the list to a set. We iterate through each number x in the original array. For each x, we have two conditions: \circ If x - 1 is not in the set, x could be the start of a new sequence.

Convert the nums array into a set to eliminate duplicates and allow for 0(1) existence checks. It takes 0(n) time to convert

∘ If x - 1 is in the set, x is part of a sequence that started before x and we don't need to check it as it will already be covered when we

query the existence of an element in constant 0(1) time.

check the beginning of its sequence. When we find a number x that is the start of a new sequence (because x - 1 is not in the set), we then proceed to check how long this sequence is by continuously incrementing y (initialized as x + 1) as long as y is present in the set.

Each time we extend the sequence, we update the length of the current sequence and update the answer ans if the current

This approach guarantees we only make a constant number of passes through the array and that we only consider each

sequence is longer than the previously recorded longest sequence.

sequence from its beginning, ensuring our algorithm runs in O(n) time.

Solution Approach The solution approach can be decomposed into key steps that align with the algorithmic design and utilize data structures such

Step 1: Building a Set Firstly, we convert the given list nums into a set s. This process removes any duplicate elements and facilitates constant time checks for the presence of integers. This is critical as it allows for the linear time algorithm we're aiming for.

Step 2: Iteration and Sequence Detection We iterate through each number x in the list nums. For each number, we check if its predecessor (x - 1) is in the set.

if x - 1 not in s:

for x in nums:

v = x + 1

while y in s:

y += 1

s = set(nums)

as hash tables effectively.

Step 3: Extension of the Sequence When we find that x could be the start of a sequence, we try to find out where the sequence ends. We initialize a variable y as x + 1 and while y is in the set, we keep incrementing y by one to extend the sequence.

If x - 1 is not in the set, it implies that x could potentially be the start of a new consecutive sequence.

represents the length of the longest consecutive elements sequence found.

efficiently and satisfy the problem's constraints.

find the longest sequence of consecutive numbers in this array.

Duplications are removed and we can check for existence in constant time.

We iterate through nums. Assume our iteration order is the same as the array's order.

```
Step 4: Update Longest Sequence Length After we find a sequence starting with x and ending before y, the length of this
 sequence is y - x. If this length is greater than any previously found sequences, we update our answer ans.
ans = max(ans, y - x)
```

Step 5: Return the Result Once we've considered each number in the array, we return ans as the answer to the problem, which

This approach takes advantage of the hash table pattern via the set s, which provides us with the constant time lookups needed

to achieve an overall O(n) time complexity. Thus, we harness the capability of hash tables to manage our computations

Let's illustrate the solution approach using a small example. Consider the unsorted array nums = [4, 1, 3, 2, 6]. Our goal is to

Example Walkthrough

Step 1: Building a Set We transform the **nums** array into a set: $s = set([4, 1, 3, 2, 6]) # s = \{1, 2, 3, 4, 6\}$

if 3 not in s: # False, hence we skip

Iteration 2: x = 1

v = 2

while y in s:

Step 5: Return the Result

Python

Java

C++

#include <vector>

#include <unordered set>

returned as the solution.

Solution Implementation

longest_streak = 0

for number in nums:

return longest_streak

// Add all elements to the set.

// Go through each element in the array.

if (!numSet.contains(num - 1)) {

int currentNum = num;

int currentStreak = 1;

currentNum += 1:

// Return the longest streak length.

return longestStreak;

currentStreak += 1;

for (int num : nums) {

int longestStreak = 0;

for (int num : nums) {

numSet.add(num);

Iteration 1: x = 4

Step 2: Iteration and Sequence Detection

We check if 3 (x - 1) is in the set:

We check if 0 (x - 1) is in the set:

Since 3 is present, 4 is not the start of a new sequence.

if 0 not in s: # True, thus 1 might be a sequence start

We extend the sequence from 1 onwards to find its length:

We calculate the length of the current sequence 4 - 1 which is 3.

We update ans to the length of this sequence if it is the longest found so far.

y += 1 # y becomes 3, then 4, stops at 5

ans = max(ans, 4 - 1) # if ans was 0, it becomes 3

Since 0 is not present, 1 is a start of a new sequence.

Step 3: Extension of the Sequence

The sequence we found is 1, 2, 3, 4. **Step 4: Update Longest Sequence Length**

Iterations continue with 2, 3, and 6 but no other new sequence is found with a length greater than 3.

Initialize the current number as the possible start of a sequence

Increment the current num to find the length of the streak

Update the longest streak with the maximum streak found

longest_streak = max(longest_streak, current_streak)

// Initialize the variable for the longest consecutive sequence.

// Check if current number is the beginning of a sequence.

// Initialize the current streak length.

while (numSet.contains(currentNum + 1)) {

// Update the longest streak found so far.

// Initialize the current number as the potential start of the sequence.

// Expand the current streak if consecutive numbers are found.

longestStreak = Math.max(longestStreak, currentStreak);

After iterating through all numbers, the longest sequence found is from 1 to 4, which has a length of 4. Thus, ans = 4 and is

class Solution: def longestConsecutive(self, nums: List[int]) -> int: # Create a set from the list for O(1) lookups num set = set(nums)

Iterate over each number in the list

if number - 1 not in num set:

current num = number

current num += 1

current_streak += 1

current_streak = 1

Check if it's the start of a sequence

while current num + 1 in num_set:

Return the length of the longest consecutive sequence

class Solution { public int longestConsecutive(int[] nums) { // Create a hash set to store the unique elements of the array. Set<Integer> numSet = new HashSet<>();

```
#include <algorithm>
class Solution {
public:
   int longestConsecutive(vector<int>& nums) {
       // Create an unordered set to hold unique elements for constant-time lookups.
       unordered set<int> numbersSet(nums.begin(), nums.end());
        int longestStreak = 0; // Variable to store the length of the longest consecutive sequence found.
       // Iterate over each element in the vector.
        for (int num : nums) {
            // Check if the current number is the beginning of a sequence by looking for num -1.
            if (!numbersSet.count(num - 1)) {
                // If num is the start of a sequence, look for all consecutive numbers starting with num + 1.
               int currentNum = num + 1;
               // Continue checking for the next consecutive number in the sequence.
               while (numbersSet.count(currentNum)) {
                    currentNum++;
                // Update the longest streak with the length of the current sequence.
                longestStreak = max(longestStreak, currentNum - num);
       // Return the longest length of consecutive sequence found.
        return longestStreak;
};
TypeScript
// Function to find the length of the longest consecutive elements sequence.
function longestConsecutive(nums: number[]): number {
   // Initialising a set to store unique numbers from the input.
   const numSet: Set<number> = new Set(nums);
    let longestStreak = 0; // Stores the length of the longest consecutive sequence.
   // Iterate over each number in the set.
```

// Starting number of the current sequence.

// Initializing current streak length.

Time and Space Complexity The given code is designed to find the length of the longest consecutive elements sequence in an unsorted array. It utilizes a set

return longest_streak

for (const num of numSet) {

return longestStreak;

num set = set(nums)

longest_streak = 0

for number in nums:

class Solution:

if (!numSet.has(num - 1)) {

currentNum++;

currentStreak++;

let currentNum = num;

let currentStreak = 1;

// Check if current number is the beginning of a sequence.

// Update the longest streak if current one is longer.

longestStreak = Math.max(longestStreak, currentStreak);

Initialize the current number as the possible start of a sequence

Increment the current num to find the length of the streak

Update the longest streak with the maximum streak found

longest_streak = max(longest_streak, current_streak)

// Incrementally check consecutive numbers.

// Return the length of the longest consecutive sequence.

def longestConsecutive(self, nums: List[int]) -> int:

Create a set from the list for O(1) lookups

Check if it's the start of a sequence

while current num + 1 in num_set:

Return the length of the longest consecutive sequence

Iterate over each number in the list

if number - 1 not in num set:

current num = number

current num += 1

to achieve an average time complexity of O(n).

1. Creating a set from the list of numbers, which takes O(n) time.

current_streak += 1

current_streak = 1

while (numSet.has(currentNum + 1)) {

Time Complexity: The algorithm has two main parts:

2. Looping through each number in the array and extending the consecutive sequence if the current number is the start of a sequence. This part is

Combining these two parts still results in a total of O(n) time complexity since other operations inside the loop are constant time

on average, such as checking for membership in the set and updating the ans variable. **Space Complexity:**

The space complexity is O(n) because a set is created to store the elements of the array, and no other data structures that depend on the size of the input are used.

also 0(n) on average because each number is visited only once during the sequence extension process.