2279. Maximum Bags With Full Capacity of Rocks

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Problem Description
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Medium

Greedy Array Sorting

number we can achieve with the given additionalRocks.

reach full capacity. This is done using a list comprehension:

ans by 1 and reduce additionalRocks by the amount used:

number of additionalRocks that we can distribute across these bags. The objective is to maximize the number of bags that are filled to their full capacity using the additionalRocks available. A bag is considered to be at full capacity if the number of rocks it contains equals its capacity.

In this problem, we're dealing with a set of n bags, each with a certain capacity. The capacity array represents the maximum

number of rocks each bag can hold, while the rocks array shows the current number of rocks in each bag. We also have a certain

Intuition

The intuition behind the solution is based on the idea that to maximize the number of bags at full capacity, we should fill the bags

that need the fewest additional rocks first. This is a greedy approach where we prioritize the bags that are closest to being full

because adding rocks to them will quickly increase the count of fully filled bags. We calculate the difference between the capacity and rocks for each bag, which represents the number of additional rocks needed to fill each bag to capacity. We then sort this list to get the bags that need the least additional rocks at the beginning.

Stepping through this sorted list, we distribute the additionalRocks as long as we have enough to fill a bag. Each time we fill a bag, we decrement the number of additionalRocks by the number used and increment the count of fully-filled bags by one. This process continues until we run out of additionalRocks or fill all the bags. The final count of fully-filled bags is the maximum

Solution Approach The solution implements a simple greedy strategy using Python lists and sorting algorithm. Here's a step-by-step walkthrough of the implementation:

Calculate the difference between the capacity and the rocks for each bag to find out how many more rocks are needed to

d = [a - b for a, b in zip(capacity, rocks)]

for v in d:

this value:

return ans

if v <= additionalRocks:</pre>

additionalRocks -= v

Following the steps outlined in the solution approach:

d.sort() # The sorted list d will remain [1, 1, 2]

Using the list comprehension:

ans = 0 # Starts at 0

Solution Implementation

remaining_capacity.sort()

filled_bags = 0

else:

break

return filled_bags

from typing import List

class Solution:

Python

if v <= additionalRocks:</pre>

ans += 1

Here, a represents an element from the capacity array and b represents the corresponding element from the rocks array. The zip function pairs each element from capacity with the corresponding element from rocks.

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The problem is now reduced to filling the bags with the least difference first. To do this efficiently, sort the list d in non-
 decreasing order:
d.sort()
```

Initialize a variable ans to count the number of bags that can be filled to full capacity:

ans = 0Iterate through the sorted list d and try to fill each bag. If the additionalRocks is enough to fill the current bag, increase the

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    v represents the number of additional rocks needed for the current bag.

• If additionalRocks is at least v, it means we can fill this bag. Then we update additionalRocks to reflect the rocks used.
• The loop continues either until there are no more rocks left (additionalRocks is less than the next v in the list) or all bags are checked.
Once the loop is complete, ans is the maximum number of bags that can be filled to full capacity, and the function returns
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time complexity of this approach is $0(n \log n)$. The space complexity is 0(n) due to the creation of the list d. **Example Walkthrough**

Let's illustrate the solution approach with a small example. Suppose we have the following inputs:

First, we calculate the difference between capacity and rocks for each bag.

d = [a - b for a, b in zip(capacity, rocks)] # d will become [1, 1, 2]

We initialize ans to keep track of the bags that can be filled to full capacity:

Now, we iterate through the sorted list d and distribute additionalRocks:

additionalRocks -= v # Decrease the additionalRocks by v

Loop through [1, 1, 2] with additionalRocks starting at 3

This approach uses the built-in sorting function which typically has a time complexity of 0(n log n) where n is the number of

elements in the list. The subsequent iteration through the sorted list has a linear time complexity of O(n). As a result, the overall

• capacity array: [5, 3, 7] • rocks array: [4, 2, 5] • additionalRocks: 3 We want to find out the maximum number of bags that can be filled to their full capacity using these 3 additional rocks.

We then sort this list to consider the bags that need the fewest additional rocks first:

for v in d:

The first bag needs 1 rock to be full, which we have, so ans becomes 1 and additionalRocks becomes 2.

The second bag also needs 1 rock, so ans is incremented to 2 and additionalRocks is reduced to 1.

Therefore, we have maximized the number of full bags (2 out of 3) using the 3 additionalRocks provided.

def maximumBags(self, capacity: List[int], rocks: List[int], additionalRocks: int) -> int:

Calculate the remaining capacity of each bag by subtracting the number of rocks

Initialize the counter for the number of bags that can be completely filled

Sort the remaining capacities to prioritize bags that need fewer rocks to reach capacity

currently in each bag from the bag's total capacity.

Iterate through the sorted remaining capacities

Increment the filled bags counter

additionalRocks -= required_rocks

Return the total number of completely filled bags

as no further bags can be completely filled

public int maximumBags(int[] capacity, int[] rocks, int additionalRocks) {

// Get the number of bags by checking the length of the capacity array.

// Initialize a counter for the maximum number of bags that can be filled.

// Subtract the used rocks from the available additional rocks.

for required_rocks in remaining_capacity:

filled_bags += 1

int numBags = capacity.length;

for (int i = 0; i < numBags; ++i) {

Arrays.sort(remainingCapacity);

maxFilledBags++;

int maxFilledBags = 0;

} else {

break;

int[] remainingCapacity = new int[numBags];

// Calculate the remaining capacity for each bag.

// Iterate over the sorted remaining capacities.

for (int requiredRocks : remainingCapacity) {

if (requiredRocks <= additionalRocks) {</pre>

additionalRocks -= requiredRocks;

// 'rocks' array represents the current number of rocks in each bag.

// Iterate over the sorted bags and try to fill them

filledBags++; // Increment filled bags count

// Return the number of bags that have been filled to capacity

currently in each bag from the bag's total capacity.

additionalRocks -= required rocks

Return the total number of completely filled bags

as no further bags can be completely filled

// Get the number of bags

let filledBags = 0;

return filledBags;

from typing import List

else:

Time Complexity

Space Complexity

break

Time and Space Complexity

return filled bags

class Solution:

const numBags = capacity.length;

requiredRocks.sort((a, b) => a - b);

// 'additionalRocks' represents the total number of additional rocks available.

const requiredRocks = capacity.map((cap, index) => cap - rocks[index]);

// Calculate the difference between bag capacity and current number of rocks

// Initialize a counter to keep track of the number of bags that can be filled

function maximumBags(capacity: number[], rocks: number[], additionalRocks: number): number {

// Sort the required rocks in ascending order — to prioritize bags that need fewer rocks to reach capacity

for (let i = 0; i < numBags && (requiredRocks[i] === 0 || requiredRocks[i] <= additionalRocks); i++) {</pre>

additionalRocks -= requiredRocks[i]; // Subtract the used rocks from the additional rocks

def maximumBags(self, capacity: List[int], rocks: List[int], additionalRocks: int) -> int:

Calculate the remaining capacity of each bag by subtracting the number of rocks

// Increment the count of filled bags.

remainingCapacity[i] = capacity[i] - rocks[i];

Increment the count of full bags

```
The process stops here as we have distributed all additional rocks that we can. The final count of fully-filled bags is the value
 of ans:
return ans # Returns 2 as the maximum number of full bags
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The third bag needs 2 rocks to be full. However, we only have 1 additionalRock left, so we cannot fill this bag to capacity.

If the current bag requires fewer or equal rocks than we have available, # use those rocks to fill the bag if required_rocks <= additionalRocks:</pre>

remaining_capacity = [total_cap - current_rocks for total_cap, current_rocks in zip(capacity, rocks)]

Java class Solution {

// Create an array to store the difference between capacity and current rocks in each bag.

// Function to determine the maximum number of bags that can be filled given capacities, current rocks, and additional rocks.

// Sort the remaining capacities in ascending order; to fill as many bags as possible starting with the ones requiring the le

// If the required rocks to fill a bag is less than or equal to the available additional rocks...

// If the remaining rocks are not sufficient to fill the next bag, break out of the loop.

Decrement the available rocks by the number of rocks used for the current bag

If the current bag requires more rocks than available, break the loop,

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// Return the maximum number of bags that can be filled.
        return maxFilledBags;
C++
class Solution {
public:
    // Function to find the maximum number of bags that can be filled given the remaining capacity.
    int maximumBags(vector<int>& capacity, vector<int>& rocks, int additionalRocks) {
        int numBags = capacity.size(); // Get the number of bags.
        vector<int> remainingCapacity(numBags); // Vector to hold remaining capacities of the bags.
        // Calculate the remaining capacity for each bag.
        for (int i = 0; i < numBags; ++i) {
            remainingCapacity[i] = capacity[i] - rocks[i];
        // Sort the remaining capacities in ascending order.
        sort(remainingCapacity.begin(), remainingCapacity.end());
        int maxFilledBags = 0; // Counter for maximum number of bags that can be completely filled.
        // Iterate over each bag's remaining capacity.
        for (int& remaining: remainingCapacity) {
            // If there are not enough rocks to fill the next bag, break the loop.
            if (remaining > additionalRocks) break;
            // If we have enough rocks to fill the current bag:
                                             // Increment the count of filled bags.
            maxFilledBags++;
            additionalRocks -= remaining; // Use the rocks to fill the bag.
        return maxFilledBags; // Return the maximum number of bags that can be filled.
};
TypeScript
// Function to determine the maximum number of bags that can be filled to capacity
// with a given number of additional rocks.
// 'capacity' array represents the capacity of each bag.
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remaining_capacity.sort()
# Initialize the counter for the number of bags that can be completely filled
filled bags = 0
# Iterate through the sorted remaining capacities
for required rocks in remaining capacity:
   # If the current bag requires fewer or equal rocks than we have available,
   # use those rocks to fill the bag
    if required_rocks <= additionalRocks:</pre>
        # Increment the filled bags counter
        filled_bags += 1
        # Decrement the available rocks by the number of rocks used for the current bag
```

If the current bag requires more rocks than available, break the loop,

Sort the remaining capacities to prioritize bags that need fewer rocks to reach capacity

remaining_capacity = [total_cap - current_rocks for total_cap, current_rocks in zip(capacity, rocks)]

1. The list comprehension d = [a - b for a, b in zip(capacity, rocks)] takes O(n) time, where n is the number of elements in capacity and rocks. 2. Sorting the list d.sort() has a time complexity of O(n log n) because it uses the Timsort algorithm which is Python's standard sorting

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algorithm.
3. The loop for v in d: iterates through each element of the list d once, giving a time complexity of O(n).
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The time complexity of the provided code consists of several parts:

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Since sorting the list is the most expensive operation, the overall time complexity of the code is 0(n \log n).
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The space complexity of the code also involves a few components:

1. The list comprehension generates a new list d of size n, resulting in O(n) space complexity.

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2. Sorting the list is done in-place in Python, so it doesn't require additional space other than some constant workspace, hence 0(1).
3. The variables ans and additionalRocks use constant space, 0(1).
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Combining these, the total space complexity of the code is O(n) because the new list d is the dominant factor.