Problem Description

The problem focuses on finding the largest variance within any substring of a given string s, which contains only lowercase English letters. The variance of a string is defined as the largest difference in the number of occurrences between any two distinct characters in it. It's important to consider that the characters could be the same or different.

To clarify, let's say we're given a string like "abcde". In this string, each character appears exactly once, so the variance between any two characters is zero. However, if we had a string like "aabbcc", the variance could be higher. If we choose the substring "aab", the variance would be 2 because there are two 'a' characters and none of the 'b' character (or vice versa). We need to consider all possible substrings to find the one with the largest variance.

The intuition behind the solution lies in realizing that the variance could only be maximized if we consider pairs of different characters in the string, as those pairs are the ones that can actually have a non-zero variance. So the first step is to iterate over all

Intuition

possible pairs of characters (permutations of two distinct letters from the 26 lowercase English characters). Once a pair of characters is selected (let's say 'a' and 'b'), we apply a two-pass scanning algorithm on the string to track the difference in the occurrence counts between 'a' and 'b'. We keep two counters: one (f [0]) resets every time we encounter an 'a' after a 'b', while the other (f[1]) continues to accumulate or starts over from f[0] - 1 when a 'b' is encountered after at least one 'a'. We

look to maximize the difference (f[1]), which is essentially the variance, while scanning. In the end, we take the maximum recorded value during all scans as the answer. This approach works because by focusing on pairs of characters at a time, we are able to isolate the potential substrings that can provide the maximum variance, rather than getting lost in the overwhelming number of all possible substrings.

Along with tracking the occurrences of 'a' and 'b', there's also an understanding that we include a subtraction operation when we encounter a 'b', making sure we are actually counting a valid substring with at least one 'a' when 'b' appears, since variance requires both characters to be present. A -inf is used to handle cases where a valid substring does not exist yet. By considering every single position in the string as a potential starting or ending point for substrings with maximum variance between two specific characters,

the algorithm guarantees that no possible substring is overlooked, thus ensuring the correctness of the solution. **Solution Approach** The implementation employs a brute force strategy enhanced with a smart iteration mechanism and memory optimization to calculate the maximum variance for all substrings of the given string. The key elements of this approach involve permutations, two

Permutations of Characters Firstly, the algorithm generates all possible permutations of two distinct characters from the set of lowercase English letters. This is

done using permutations (ascii_lowercase, 2) which loops through all pairs of two distinct characters since we're looking for the difference in occurrences between two potentially different characters.

• f[1] keeps track of the accumulated variance between a and b. This dichotomy allows the algorithm to keep track of two scenarios simultaneously:

Tracking Variance with Two Counters

In the solution, an array f with two elements is used:

Capturing the 'variance' until a reset is needed due to a b occurrence (handled by f[1]).

comparison variables, and a linear scan through the string.

Both f[0] and f[1] are reset to zero if a b appears directly after a b. However, when an a is followed by a b, f[1] takes the role of capturing the current variance by taking either the current accumulated variance minus one, or the variance just after the last a

• f[0] serves as a resettable counter for occurrences of character a since the last occurrence of b.

minus one, determined by max(f[1] - 1, f[0] - 1). This handling ensures that at least one a is present before a b is counted, which is necessary for a valid substring.

false starts with -inf, ensuring that f[1] only records valid substrings.

Accumulating occurrences of a when no b has interfered (handled by f[0]).

As we scan the string s, for each character c, the algorithm performs the following actions: If c is a, we increment both f[0] and f[1]. If c is b, we decrement f[1] and choose the larger of the two: f[1] or f[0] - 1, effectively either continuing with the current

After each update, we check if f[1] exceeds the current ans. If it does, we update ans to the value of f[1]. The algorithm avoids

the string ensures that we consider every potential substring for each character pair. Eventually, ans will contain the maximum

substring or starting a new substring that excludes the last a. f[0] is reset to zero as we can't start a new substring with just b.

The loop through permutations of characters ensures that we consider every potential character disparity. The nested loop through

Example Walkthrough

Maximal Variance Calculation

variance achievable among all substrings of s.

The space complexity of the solution is 0(1) as we only maintain a few variables and a pair of counters, and the time complexity is

largest variance within any substring of s.

1. Initialize the counters f[0] and f[1] to 0.

2. Loop through each character c in s = "abbab":

• Next, 'a': f[0] = 1, f[1] = max(f[1] + 1, f[0]) = 1.

iterations over the string.

Let's illustrate the solution approach using a small example. Suppose we are given the string s = "abbab". Our task is to find the

0(26*26*n) where n is the length of the string, due to the nested loops—26 for each character in the permutations, and n for the

Firstly, we generate all possible permutations of two distinct characters from the set of lowercase English letters. For simplicity, let's consider only a few pairs here, such as ('a', 'b') and ('b', 'a'), since these are the only characters in our example string. We initiate a variable ans to store the maximum variance found. Initially, ans = -inf.

∘ For c = 'a': Increment both f[0] and f[1] (since we initially consider the occurrence of 'a' without 'b' involved).

∘ For c = 'b': Decrement f[1] and then set f[1] to the maximum of f[1] and f[0] - 1. Reset f[0] to 0.

Let's go through the string "abbab":

• Last, 'b': f[1] = max(f[1] - 1, f[0] - 1) = 0. Update ans = max(ans, f[1]) = 1.

Iterate through all ordered pairs of distinct lowercase ascii characters

Initialize the frequency counters for the current pair, where:

Update the max_variance if we find a new maximum

// Method to find the largest variance of a substring for a given string

// Iterate over all possible character pairs (a, b) where a != b

// Initialize frequency counters for characters a and b

// f[0] will track frequency of a, f[1] tracks max variance

// Increment frequency count if a matches the current character

// Skip if both characters are the same

// Iterate over each character in the string

for char_a, char_b in permutations(ascii_lowercase, 2):

• Start with 'a': f[0] = 1, f[1] = 1. • Next, 'b': f[1] = max(f[1] - 1, f[0] - 1) = 0, f[0] = 0. Update ans = max(ans, f[1]) = 1.

Pair ('a', 'b'):

- Pair ('b', 'a'):
- Following the same steps but reversing the roles of 'a' and 'b', we would not find a variance larger than what we already have from the pair ('a', 'b').

In a complete implementation, we would evaluate all permutations of character pairs, but in this small example with only 'a' and 'b'

The process effectively scans through every substring of s and calculates the variance for pairs of characters, updating ans when a

After considering these permutations, we find that the maximum variance is 1, which we may have obtained from the substring "ab" or "ba" in the string "abbab".

• Next, 'b': $f[1] = \max(f[1] - 1, f[0] - 1) = -1$, since there are no 'a' in front, we don't update ans.

1 from itertools import permutations from string import ascii_lowercase from math import inf

21 if current_char == char_a: 22 # Increment the frequency counter of 'a' when 'a' is found 23 freq_before_b += 1 24 max_freq_after_b += 1 elif current_char == char_b:

Decrement the frequency counter of 'b' when 'b' is found; reset freq_before_b

max_freq_after_b = max(max_freq_after_b - 1, freq_before_b - 1)

```
higher variance is found. The time complexity is mainly driven by the nested loops: there are 26*26 permutations of character pairs,
and for each pair, we scan the entire string s of length n, resulting in an 0(26*26*n) time complexity.
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present, it is unnecessary.

Python Solution

class Solution:

def largestVariance(self, s: str) -> int:

freq_before_b = 0

Return the largest variance found

public int largestVariance(String s) {

// Length of the input string.

int length = s.length();

return max_variance

if max_variance < max_freq_after_b:</pre>

max_variance = max_freq_after_b

Initialize the answer to 0

max_variance = 0

```
# freq_before_b records the number of 'a' characters seen so far without any 'b' in between
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               # max_freq_after_b records the maximum frequency difference after at least one 'b' was found
16
               freq_before_b = 0
17
               max_freq_after_b = -inf
18
               # Iterate over each character in the string
19
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               for current_char in s:
```

Java Solution

class Solution {

```
// Variable to store the maximum variance found so far.
           int maxVariance = 0;
           // Iterate over all possible pairs of characters (a and b are different).
           for (char firstChar = 'a'; firstChar <= 'z'; ++firstChar) {</pre>
                for (char secondChar = 'a'; secondChar <= 'z'; ++secondChar) {</pre>
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                    if (firstChar == secondChar) {
                        // If both characters are the same, skip this iteration.
13
                        continue;
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15
                    // Array to keep track of the frequency of character 'a'
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                   // f[0] is the streak of 'a's, f[1] is the max variance for the current window.
17
                    // Initialize with 0 for streak and -n for variance because variance cannot be less than -n.
18
                    int[] frequency = new int[] {0, -length};
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                    // Iterate over each character in the string.
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                    for (int i = 0; i < length; ++i) {</pre>
23
                        if (s.charAt(i) == firstChar) {
                            // If the current character is 'a', increase both frequencies.
                            frequency[0]++;
26
                            frequency[1]++;
27
                        } else if (s.charAt(i) == secondChar) {
                            // If the current character is 'b', calculate the variance.
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                            frequency[1] = Math.max(frequency[0] - 1, frequency[1] - 1);
                            // Reset the streak of 'a's because 'b' is encountered.
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                            frequency [0] = 0;
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                        // Update the maximum variance found.
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                        maxVariance = Math.max(maxVariance, frequency[1]);
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           // Return the largest variance found.
            return maxVariance;
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42 }
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C++ Solution
  1 class Solution {
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2 public:

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int largestVariance(string s) {

int n = s.size();

int maxVariance = 0;

// Length of the input string

// Initialize the maximum variance to 0

for (char b = 'a'; b <= 'z'; ++b) {

int frequencies[2] = $\{0, -n\}$;

for (char currentChar : s) {

if (currentChar == a) {

frequencies[0]++;

frequencies[1]++;

for (char a = 'a'; a <= 'z'; ++a) {

if (a == b) continue;

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 27
                         // Decrement frequency count and update variance if b matches
                         else if (currentChar == b) {
                              frequencies[1] = \max(\text{frequencies}[1] - 1, \text{frequencies}[0] - 1);
                              frequencies[0] = 0;
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 32
                         // Update the maximum variance found so far
 33
                         maxVariance = max(maxVariance, frequencies[1]);
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 36
             // Return the final maximum variance
 37
             return maxVariance;
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    };
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Typescript Solution
  1 // Function to find the largest variance of a substring for a given string
  2 function largestVariance(s: string): number {
         // Length of the input string
         let n: number = s.length;
         // Initialize the maximum variance to 0
         let maxVariance: number = 0;
         // Iterate over all possible character pairs (a, b) where a != b
         for (let a = 'a'.charCodeAt(0); a <= 'z'.charCodeAt(0); ++a) {</pre>
  9
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             for (let b = 'a'.charCodeAt(0); b <= 'z'.charCodeAt(0); ++b) {</pre>
                 // Skip if both characters are the same
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                 if (a === b) continue;
 13
 14
                 // Initialize frequency counters for characters 'a' and 'b'
 15
                 let frequencies: number[] = [0, -n];
 16
 17
                 // Iterate over each character in the string
                 for (let currentChar of s) {
 18
                     // Increment frequency count if 'a' matches the current character
 19
                     if (currentChar.charCodeAt(0) === a) {
 20
                         frequencies[0]++;
 21
 22
                         frequencies[1]++;
 23
 24
                     // Decrement frequency count and update variance if 'b' matches
 25
                     else if (currentChar.charCodeAt(0) === b) {
```

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Time Complexity

return maxVariance;

Time and Space Complexity

// Return the final maximum variance

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36 }

26 frequencies [1] = Math.max(frequencies [1] - 1, frequencies [0] - 1);27 frequencies[0] = 0; 28 // Update the maximum variance found so far 29 maxVariance = Math.max(maxVariance, frequencies[1]);

1. The function permutations (ascii_lowercase, 2) generates all possible permutations of two distinct lowercase letters from the English alphabet. Since there are 26 letters in the English alphabet, there are 26 * 25 such permutations because after choosing

the first letter, there are 25 remaining letters to form a pair.

The space complexity of the code is determined by:

The time complexity of the code can be analyzed based on the operations it performs:

- 2. The function then iterates through each character in the string s for every pair of letters (a, b). The length of the string s is n. 3. Within the nested loop, there are constant-time operations (comparisons, assignments, max, addition, and subtraction).
- 25) is a constant. **Space Complexity**

Given the above steps, the time complexity is 0(26 * 25 * n), which simplifies to 0(n) because the number of permutations (26 *

- 2. No additional data structures that grow with input size are used.
- 1. Variables ans, f[0], and f[1], which are used to keep track of the current and maximum variance.
- Thus, the space complexity of the algorithm is 0(1) as it only uses a constant amount of space.