## 799. Champagne Tower

Medium <u>Dynamic Programming</u>

## **Problem Description**

pyramid has a number of glasses equal to the row number, starting from 1 for the first row, 2 for the second row, and continuing up to the 100th row. Each glass can hold one cup of champagne. When we pour champagne into the top glass, it fills glasses below it in the following manner: once the current glass is full, any extra champagne splits equally between the two glasses directly below it. This process continues down the pyramid until the champagne either fills glasses or spills on the floor if it reaches the last row.

The actual problem is to determine how full a specific glass, given by its row (query\_row) and position in the row (query\_glass), is after pouring a certain number of cups (poured) of champagne into the top glass of the tower.

This problem presents a scenario where glasses are stacked in a pyramid fashion to form a champagne tower. Each row in the

Intuition

To solve this problem, we can use a dynamic programming approach where we simulate the pouring process. We will create a 2D

## array that represents the champagne glasses and track the amount of champagne in each glass. When we pour champagne into

the glass at the top, we continue to distribute it down the rows. If a glass receives more than one cup of champagne, it overflows, and the excess amount is shared equally between the two glasses directly below it.

We're only interested in the distribution of champagne up to and including the query\_row, so we can limit the simulation to that. By initializing the overflow process from the top and proceeding row by row, we can efficiently calculate how much champagne is in each glass when the overflow process ceases - either when we reach the desired row or when no more champagne can

overflow.

When representing the overflow, any glass that has more than one cup of champagne will distribute the excess. This is calculated by subtracting 1 cup (the glass's capacity) from its champagne amount, dividing the remainder by 2, and adding the result to the two glasses below. This is iterated for each row until we get to the row we need to query, which contains the glass of interest.

Finally, we query the amount of champagne in the specific glass. If the glass is full, it will have one cup of champagne; if it is not

full, it will have an amount equal to whatever was poured into it through the overflow process. Since glasses can't hold more than one cup, the answer will be 1 if the calculated amount is over 1 cup, or the calculated amount itself if it's less than or equal to 1 cup.

The solution uses <u>dynamic programming</u> to solve the problem efficiently. Here's how the implementation goes:

1. Create a 2D Array: We initialize a 2D array | f | with dimensions 101×101, which represents the champagne glasses. We use a

## 3. Simulate Pouring Process: We iterate through the rows up to query\_row + 1. For each row i, we iterate through the glasses

since that's how full the glass is.

f[0][0] = poured

 $f = [[0] * 101 for _ in range(101)]$ 

and want to find out how full the glass at row 2, position 1 (zero-indexed) is.

then divide this by 2, so each glass below will receive 1 cup.

Step 2 (Top Glass): We pour 3 cups into the top glass f[0][0], so it becomes 3.

for i in range(query\_row + 1):

 $\mathbf j$  in that row up to  $\mathbf i + \mathbf 1$ , because the number of glasses in each row is equal to the row number.

Base Condition: We set the champagne in the top glass f[0][0] to be equal to the poured amount.

size of 101 because that accounts for the 100th row potentially overflowing into an additional row below it.

is greater than 1.

5. **Distribute Excess Champagne**: If the glass overflows, we calculate the excess amount that will flow to the glasses below by

Check for Overflow: Inside the nested loop, we check if a glass has more than one cup of champagne by checking if f[i][j]

subtracting one cup and dividing the remainder by two (this is half).

6. **Update Adjacent Glasses**: We update the two glasses below the current glass (positions [i + 1][j] and [i + 1][j + 1] in

the f array) by adding the half to both. This represents the flow of excess champagne to lower-level glasses.

8. **Query Result**: After completing the simulation up to the query\_row, we simply return the value at f[query\_row] [query\_glass]. If it's more than 1, it means the glass is overflowing, and we return 1. Otherwise, we return the actual value

Cap Glass Fullness: We set the current glass f[i][j] to 1 because a glass can't hold more than one cup of champagne.

The patterns used in the solution are <u>dynamic programming</u> (memoization of intermediate results to avoid recomputation) and simulation (simulating the behavior of champagne pouring over the glasses).

manner. The <u>dynamic programming</u> pattern reduces the overall time complexity as compared to a naive recursive approach.

The algorithm iteratively updates the state of the champagne tower as champagne is poured until it reaches the query condition.

This allows us to only perform necessary calculations and retrieve the exact fullness of the requested glass in an efficient

Here's the code snippet that performs these steps:

class Solution:
 def champagneTower(self, poured: int, query\_row: int, query\_glass: int) -> float:

for j in range(i + 1):
 if f[i][j] > 1:
 half = (f[i][j] - 1) / 2

f[i][i] = 1
 f[i + 1][j] += half
 f[i + 1][j + 1] += half
 return f[query\_row][query\_glass]

Let's illustrate the solution approach with a small example. Suppose we pour 3 cups of champagne into the top glass of a tower

Step 3 (First Row): As there is an overflow (since  $\frac{3}{2} > \frac{1}{2}$ ), we calculate the excess champagne, which is  $\frac{3}{2} - \frac{1}{2} = \frac{2}{2}$  cups. We

```
• Step 1 (Initialization): We create the 2D array f to represent our glasses, all initially with 0 champagne.
```

Row 0: 3

Row 1: 1 1

Row 1: 1 1

Row 0: 1

Row 1: 1 1

Row 2: 0 0 0

Row 2: 0 0 0

Row 2: 0 0 0

Here's how the process works step-by-step:

At this point, the array will look like this:

**Example Walkthrough** 

The array now:

Row 0: 1

Step 4 (Second Row): Now we look at each glass in the first row. Each glass has 1 cup, so no overflow occurs.

Step 5 (Third Row): Since there's no overflow from the first row, the second row's glasses remain as they are.

Step 6 (Result Query): Now we query the fullness of the glass at row 2, position 1, which corresponds to f[2][1]. As we

In conclusion, after pouring 3 cups of champagne, the glass at row 2, position 1 is empty. The code provided efficiently

simulates the entire pouring process, accounts for any overflows, and gives us the correct amount of champagne in any given

The final state of the array up to row 2:

can see, it is 0, meaning no champagne has reached this glass. Therefore, the return value would be 0.

# Initialize a Pascal's triangle with 101 rows and columns as the champagne tower, and only 0 is filled.

# Simulating the pouring process up to the query\_row + 1 because we need values from the given query row.

def champagneTower(self, poured: int, query row: int, query glass: int) -> float:

# Calculate the champagne that flows to each glass below.

# Return the amount of champagne in the glass at query\_row and query\_glass, capped at 1.

# If there's an overflow in the current glass...

overflow = (tower[row][col] - 1) / 2.0

tower =  $[[0.0] * 101 for _ in range(101)]$ 

# The champagne poured into the top glass.

for row in range(query row + 1):

for col in range(row + 1):

if tower[row][col] > 1:

return min(1, tower[query\_row][query\_glass])

for (int row = 0: row <= quervRow: row++) {</pre>

for (int glass = 0; glass <= row; glass++) {</pre>

if (champagneLevel[row][glass] > 1) {

champagneLevel[row][glass] = 1;

// Return the amount of champagne in the gueried glass

return champagneLevel[queryRow][queryGlass];

// Check if the current glass is overflowing

champagneLevel[row + 1][glass] += overflow;

// (it will be at most 1 since any excess would have overflowed)

champagneLevel[row + 1][glass + 1] += overflow;

// Initialize the first row with the amount of champagne poured into the first glass

// Initialize the next row with zeros which will represent the empty glasses

for (let glassIndex = 0; glassIndex < rowIndex; glassIndex++) {</pre>

// equally to the two glasses below it in the next row

def champagneTower(self, poured: int, query row: int, query glass: int) -> float:

# Calculate the champagne that flows to each glass below.

# Distribute the overflowed champagne to the two glasses below.

# Return the amount of champagne in the glass at query\_row and query\_glass, capped at 1.

# Reset the current glass to full after overflow.

# If there's an overflow in the current glass...

overflow = (tower[row][col] - 1) / 2.0

tower[row + 1][col] += overflow

tower[row + 1][col + 1] += overflow

const overflow = (currentRow[glassIndex] - 1) / 2;

// Set the nextRow to be the currentRow for the next iteration

for (let rowIndex = 1; rowIndex <= queryRow; rowIndex++) {</pre>

const nextRow = new Array(rowIndex + 1).fill(0);

nextRow[qlassIndex] += overflow;

nextRow[glassIndex + 1] += overflow;

// Iterate over each glass in the current row

if (currentRow[glassIndex] > 1) {

// Iterate over the rows of the champagne tower until the specified query row is reached

// If the current glass has more than one unit of champagne, it overflows

// Return the amount of champagne in the specified glass, capped at 1 because glasses can't hold

// Calculate the amount of champagne that overflows from the current glass and distribute it

# Initialize a Pascal's triangle with 101 rows and columns as the champagne tower, and only 0 is filled.

# Simulating the pouring process up to the query\_row + 1 because we need values from the given query row.

double overflow = (champagneLevel[row][glass] - 1) / 2.0;

// Distribute the overflowing champagne to the two glasses below it

glass.

Solution Implementation

tower[0][0] = poured

**Python** 

Java

class Solution {

class Solution:

- # Reset the current glass to full after overflow.
  tower[row][col] = 1
  # Distribute the overflowed champagne to the two glasses below.
  tower[row + 1][col] += overflow
  tower[row + 1][col] += overflow
- // Computes the amount of champagne in a glass located at (guervRow, queryGlass)
  // after pouring a certain amount into the top glass of the tower.
  public double champagneTower(int poured, int queryRow, int queryGlass) {
   // Initialize a 2D array to hold the quantity of champagne in each glass
   double[][] champagneLevel = new double[101][101];

   // Pour champagne into the top glass
   champagneLevel[0][0] = poured;

  // Fill the glasses for each row till the queried row

// Calculate the amount of champagne that overflows, to be divided between the two glasses below

// Current glass should not have more than 1 unit of champagne after overflowing

```
class Solution {
public:
    double champagneTower(int poured, int queryRow, int queryGlass) {
        // Initialize the array that will store the quantity of champagne in each glass.
        // Only need 100 rows according to the problem statement.
        double glasses[100][100] = {0.0};
        // Pour the champagne into the top glass.
        glasses[0][0] = poured;
        // Start from the top and fill down to the queried row.
        for (int row = 0; row <= queryRow; ++row) {</pre>
            for (int glass = 0; glass <= row; ++glass) {</pre>
                // Check if there is any overflow in the current glass.
                if (glasses[row][glass] > 1) {
                    // Calculate the amount of champagne that flows to the glasses below.
                    double overflow = (glasses[row][glass] - 1.0) / 2.0;
                    // Ensure the current glass is filled to its capacity.
                    glasses[row][glass] = 1;
                    // Distribute the overflow to the two glasses below equally.
                    glasses[row + 1][glass] += overflow;
                    glasses[row + 1][glass + 1] += overflow;
        // Return the amount of champagne in the gueried glass.
        // If the glass is full or under, it will contain the exact amount.
        // If it did not receive enough champagne, it will contain the remaining amount.
        return glasses[queryRow][queryGlass];
};
TypeScript
// Defines a function to simulate pouring champagne into a tower and queries the amount of champagne
// in a specific glass at a specific row after the champagne is poured.
function champagneTower(poured: number, queryRow: number, queryGlass: number): number {
```

```
tower = [[0.0] * 101 for _ in range(101)]

# The champagne poured into the top glass.
tower[0][0] = poured
```

class Solution:

let currentRow = [poured];

currentRow = nextRow;

// more than one unit of champagne

return Math.min(1, currentRow[queryGlass]);

for row in range(query row + 1):

Time and Space Complexity

for col in range(row + 1):

if tower[row][col] > 1:

tower[row][col] = 1

return min(1, tower[query\_row][query\_glass])

```
Time Complexity

The time complexity of the code is primarily determined by the nested for loops. In the worst case, the loop runs for query_row + 1 iterations, and for each iteration i, it runs i + 1 times, because the inner loop ranges up to the current row number. Thus, the number of operations can be approximated as a sum of an arithmetic series, which calculates to roughly 1 + 2 + ... + (query_row + 1), which is ((query_row + 1) * (query_row + 2)) / 2. Simplifying this arithmetic series results in
```

O(query\_row^2). Hence, the time complexity is <code>O(query\_row^2)</code>.

Space Complexity

The space complexity is determined by the size of the 2D list <code>f</code>, which is created to store the quantity of champagne in each glass. Since the list is initialized with dimensions 101×101, this is essentially a constant space allocation, not varying with the size of the input. Therefore, the space complexity is <code>O(1)</code> in terms of the input poured, <code>query\_row</code>, and <code>query\_glass</code>; however, if considering the size of the grid as part of the complexity, it would be <code>O(101 \* 101)</code> which simplifies to <code>O(1)</code> under Big O notation since 101 is a constant.