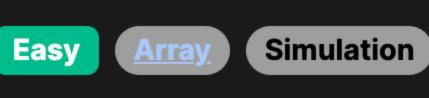
2460. Apply Operations to an Array



Problem Description

In this problem, you are given an array of non-negative integers. Your task is to perform a series of operations on this array, specifically n - 1 operations where n is the size of the array. Each operation is performed on the ith element of the array (considering 0-indexing), following these rules:

• If they are equal, double the ith element (nums[i] = nums[i] * 2) and set the i + 1th element to 0.

array manipulation techniques. The key steps in this approach are as follows:

Check if the ith element of the array (nums[i]) is equal to the next element (nums[i + 1]).

- If they are not equal, move on to the next operation without making changes to the current ith element.
- After you have performed all the operations, you need to shift all the 0's in the array to the end, while preserving the order of the non-zero elements. The challenge is to perform these operations sequentially and then return the modified array.

Example: Given an array [1,0,2,0,0,1], after performing all the operations and shifting the 0's, the resulting array would be

[1,2,1,0,0,0]. Intuition

The solution involves a two-step approach. First, we perform the n-1 operations as specified, inspecting each pair of adjacent

elements set to 0 that now should be moved to the end. The intuition behind the first step is straightforward: loop through the array, compare each element with its neighbor, and if they are the same, apply the operation. We have to remember that these operations should be applied sequentially, meaning the result

elements and applying the doubling and zeroing rules. After all operations are complete, we're left with an array with some

of one operation may affect subsequent operations. Therefore, careful in-place manipulation of the array is necessary. For the second step, the intuition is to keep track of where the next non-zero element should be placed. Essentially, this involves a second pass through the array, where we move each non-zero element leftwards to "fill in" the non-zero portion of the array.

This is why a separate counter (i in the provided solution) is maintained to keep track of the index at which the next non-zero element should be inserted. Non-zero elements are placed in the array in their original order until all non-zero elements have been accounted for. Finally, the rest of the array is filled with 0s. **Solution Approach**

The provided solution follows a straightforward two-pass approach which efficiently addresses the requirements with simple

Doubling and Zeroing in Place: The first pass goes through the array from the start to the second-to-last element. At each

doubling integers.

index i, the algorithm checks if nums[i] equals nums[i + 1]. If they are equal, it doubles nums[i] using the left shift operator (<<= 1 is equivalent to multiplying by 2) and sets nums [i + 1] to 0. Using the bitwise shift here is a more efficient way of

Shifting Non-zero Elements: In the second pass, the algorithm traverses the array only once more and maintains a separate index i which keeps track of the position where the next non-zero element should go. Hence, for each element x in the array, if x is non-zero, it is placed at nums[i] and the index i is incremented. This effectively compacts all non-zero elements towards the beginning of the array.

Filling Remaining with Zeros: Because the original array is modified in place during step 1, and the non-zero elements are

moved forward in step 2, the remaining elements in the array (from the current index i to the end) are already implicitly 0. If

there was any need to explicitly set them to 0, it could be done in a final pass; however, the code efficiency is improved by realizing this step is not necessary given the initial array manipulation. This approach takes O(n) time due to the two sequential passes through the array, where n is the number of elements in nums. No additional data structures are needed, so the space complexity is 0(1) as the solution uses only a fixed amount of extra space to

Example Walkthrough Let's illustrate the solution approach with a small example. Suppose we have the following array: Input array: [2, 2, 3, 3, 3]

Doubling and Zeroing in Place:

Shifting Non-zero Elements:

We start at the first element and compare it with the next element. 1. nums [0] is 2, and nums [1] is also 2. They are equal, so we double nums [0] (2 becomes 4) and set nums [1] to 0. Now the array looks like this: [4,

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0, 3, 3, 3].
2. We move to the next non-zero pair. nums [2] is 3, and nums [3] is also 3. Doubling nums [2] we get 6, and set nums [3] to 0: [4, 0, 6, 0, 3].
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store the counters and temporary values.

Now, we make a second pass through the array and move all non-zero elements to the front.

5. Next, the 3 is placed at nums [i] (now nums [2]), and i is incremented to 3.

the array while preserving the order of the non-zero elements.

def apply_operations(self, nums: List[int]) -> List[int]:

Initialize a pointer for the index of 'result'.

return result; // Return the resulting array

// Loop through each pair of adjacent numbers

for (int idx = 0; idx < size -1; ++idx) {

if (nums[idx] == nums[idx + 1]) {

// Get the size of the nums array

int size = nums.size();

vector<int> result(size);

for (int& num : nums) {

if (num) {

double its value and set the next element to 0.

Create a new list 'result' with the same size filled with zeros.

Iterate over 'nums' to populate non-zero elements in the 'result'.

Get the length of the list 'nums'.

if nums[i] == nums[i + 1]:

nums[i] *= 2

result = [0] * length

result_index = 0

nums[i + 1] = 0

Here is how the solution approach would be applied to this array:

- 1. We set up a separate index i starting at 0 to track where to place non-zero elements. 2. Starting from the beginning of the array, when we find a non-zero element, we move it to the nums [i] position and increment i.
- 3. For the first non-zero element, 4 stays in its original position, and i is incremented to 1. 4. Skipping over the zero, we come to 6. 6 is placed at nums[i] (which is nums[1] now), and i becomes 2.

3. nums [4] is the last element and doesn't have a pair to compare with, so the array remains [4, 0, 6, 0, 3] after the first pass.

So, the non-zero part of the array is now [4, 6, 3]. Since the array's length is 5 and we already have the remaining elements implicitly set to 0 due to the first pass, we get:

Output array: [4, 6, 3, 0, 0]

Solution Implementation **Python**

We have now successfully completed all operations as defined by the problem statement, and the zeros are shifted to the end of

Iterate over the list elements, except for the last element. for i in range(length - 1): # If the current element is the same as the next element,

length = len(nums)

from typing import List

class Solution:

```
for num in nums:
           # If the element is non-zero, put it in the next position of 'result'.
            if num:
                result[result_index] = num
                result_index += 1
       # Return the 'result' list containing the processed numbers.
        return result
Java
class Solution {
    // Method to apply operations on an array of integers
    public int[] applyOperations(int[] nums) {
        int length = nums.length; // Get the length of the array
       // Loop through each element, except the last one
        for (int i = 0; i < length - 1; ++i) {
           // Check if the current element is equal to the next element
            if (nums[i] == nums[i + 1]) {
                // If so, double the current element
                nums[i] \ll 1; // Same as <math>nums[i] = nums[i] * 2
                // And set the next element to zero
                nums[i + 1] = 0;
       int[] result = new int[length]; // Create a new array to store the results
        int index = 0; // Initialize result array index
       // Iterate through the original array
        for (int num : nums) {
           // Copy non-zero elements to the result array
           if (num > 0) {
                result[index++] = num; // Assign and then increment the index
```

```
class Solution {
public:
    vector<int> applyOperations(vector<int>& nums) {
```

C++

```
return result;
  };
  TypeScript
  function applyOperations(nums: number[]): number[] {
      const length = nums.length; // The total number of elements in the array 'nums'
      // Double the current number and set the next one to 0 if they're equal
      for (let index = 0; index < length - 1; ++index) {</pre>
          if (nums[index] === nums[index + 1]) {
              nums[index] *= 2; // Double the current number
              nums[index + 1] = 0; // Set the next number to 0
      // Initialize a new array 'result' with the same length as 'nums' and fill it with 0s
      const result: number[] = Array(length).fill(0);
      // Pointer to the position in 'result' where the next non-zero element will be placed
      let resultIndex = 0;
      // Move all non-zero elements to the 'result' array
      for (const number of nums) {
          if (number !== 0) {
              result[resultIndex++] = number; // Assign non-zero element and move to the next index
      return result; // Return the transformed array
from typing import List
class Solution:
   def apply_operations(self, nums: List[int]) -> List[int]:
       # Get the length of the list 'nums'.
        length = len(nums)
```

// If adjacent numbers are equal, double the current number and set next number to zero

int resultIndex = 0; // Initiate a result index to populate result vector with non-zero values

result[resultIndex++] = num; // Add to result and increment the position

// Return the result vector (which doesn't contain zeros between non-zero numbers)

nums[idx] <<= 1; // double the number (same as <math>nums[idx] *= 2)

// Create a new vector to store the resulting numbers after applying operations

// If the current number is non-zero, add it to the result vector

nums[idx + 1] = 0; // set the next number to zero

// Iterate over the modified nums array to filter out the zeros

```
return result
Time and Space Complexity
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for i in range(length - 1):

nums[i] *= 2

result = [0] * length

result index = 0

for num in nums:

Time Complexity

nums[i + 1] = 0

if nums[i] == nums[i + 1]:

The given function applyOperations consists of two separate for-loops that are not nested. The first for-loop iterates through the list nums, except for the last element, performing constant-time operations. The iteration

Iterate over the list elements, except for the last element.

double its value and set the next element to 0.

Initialize a pointer for the index of 'result'.

result[result_index] = num

result_index += 1

If the current element is the same as the next element,

Create a new list 'result' with the same size filled with zeros.

Iterate over 'nums' to populate non-zero elements in the 'result'.

Return the 'result' list containing the processed numbers.

If the element is non-zero, put it in the next position of 'result'.

complexity for this portion is O(n-1) which simplifies to O(n). The second for-loop iterates through each element in nums once. It fills in the non-zero elements to the list ans. The assignment and increment operations are constant time operations, and since this loop iterates n times, the time complexity for this loop is also O(n).

occurs exactly n - 1 times, where n is the length of nums. Since no other operations are nested inside this loop, the time

Combining both loops, which are sequential and not nested, the overall time complexity of the function is 0(n) + 0(n) which simplifies to O(n).

Space Complexity

Regarding space complexity, a new list ans of the same size as the input list nums is created, which denotes the extra space used by the algorithm. This implies a space complexity of O(n).

No additional data structures are used that grow with the input size, hence the total space complexity of the function is O(n).