581. Shortest Unsorted Continuous Subarray

differ will give us the boundaries of the subarray that needs to be sorted.

Two Pointers Sorting Monotonic Stack

Problem Description

Medium

<u>Greedy</u> <u>Array</u>

subarray) of this array that, if we sort it, the entire array becomes sorted in non-decreasing order (i.e., ascending order with duplicates allowed). The goal is to find the length of this smallest segment.

We are looking for the subarray where once sorted, the values before and after this selected part fit in order and don't disrupt the

In this LeetCode problem, we are given an array of integers (nums). The task is to identify the smallest segment (continuous

non-decreasing sequence. To solve this problem, we have to distinguish between the parts of the array which are already in the correct position and those that are not.

Intuition

There are different strategies that we could use to solve this problem, some being more optimal than others. The initial, more

straightforward approach is to sort a copy of the array and to compare it to the original. The first and last positions where they

However, the actual implemented solution follows a more optimal and intuitive approach that avoids using extra space and reduces the time complexity. The solution involves iterating over the array while keeping track of the maximum and minimum values seen from the left and the right, respectively.

While scanning from the left, we maintain the maximum value encountered so far - let's call it mx. If at any index i, the value x is less than mx, this suggests x is out of its required sorted position because everything before it is already part of a non-decreasing sequence. Hence, x should be included in the subarray that needs sorting. The rightmost index x where this occurs

marks the endpoint of the unsorted subarray. Simultaneously, we scan from the right and maintain the minimum value encountered so far - referenced as mi. Upon finding an element greater than mi at index n - i - 1 from the end, it indicates this element is also out of order. The leftmost index 1

where this situation arises becomes the beginning of the unsorted subarray.

Initially, l and r are set to -1, and if they remain unchanged after the iterations, it means the entire array is already sorted, and we return length 0. If they have been updated, the length of the minimum subarray required to be sorted is calculated by the difference r - l + 1.

The intuition behind updating r and 1 only when we find values out of their required position is that only these values define the

Solution Approach

The solution involves a single pass approach that cleverly makes use of two properties, namely tracking the maximum and

minimum values from the left and right, respectively. By doing this, we can identify the bounds of the unsorted subarray without

Initialize Variables: Two pointers 1 and r are initialized to -1 which will represent the left and right bounds of the unsorted

boundaries of the shortest unsorted subarray—and only by sorting them can we achieve a fully sorted array.

subarray. We also initialize two variables, mi to +infinity, representing the minimum value seen from the right; and mx to - infinity, representing the maximum value seen from the left.

the right.

Example Walkthrough

of O(n) and space complexity of O(1).

having to sort the array or use extra space.

Let's understand this algorithm step by step:

2. Iterate Over the Array: We loop over all elements in the array using a variable i that goes from 0 to n - 1, where n is the length of the array. While iterating, we perform the following:
Update Right Boundary (r): For the i-th element x, we compare it with the maximum value seen so far mx. If x is smaller than mx, it means x is out of place and should be part of the subarray to be sorted, so we update r to be the current index i.

• Update Maximum (mx): If x is greater than or equal to mx, we update mx to be x because it's the new maximum value seen so far.

Update Left Boundary (1): We compare nums [n - i - 1] with the minimum value seen from the right mi. If nums [n - i - 1] is bigger than mi, it's out of place and should be included in the subarray to be sorted, so we update 1 to be n - i - 1.
 Update Minimum (mi): If nums [n - i - 1] is less than or equal to mi, we update mi to this value because it's now the minimum seen from

 \circ At each iteration of i, we also consider the mirror position from the right end of the array: index n - i - 1.

sorted, which is r - l + 1.

By iterating from both ends towards the center and updating l and r on-the-go, we are leveraging the information of where the sequence breaks from being non-decreasing. mx and mi act as sentinels to detect any number that doesn't fit into the non-decreasing order up to that point, thus efficiently finding the shortest unsorted window in one single pass with a time complexity

Calculate the Subarray Length: After iterating through the array, we check if r was updated. If r is not updated (r == -1), it

means the array was already sorted, and we return 0. Otherwise, we return the length of the subarray that needs to be

Let's walk through a small example using the described solution approach. Consider the array nums = [2, 6, 4, 8, 10, 9, 15].

Step 1: Initialize Variables

• l = -1

• r = -1

i. When i = 0, element nums[i] = 2.

Let's see this in action:

• mi = +infinity

mx = -infinity

Step 2: Iterate Over the Array

mi = min(+infinity, 15) = 15.

• mi becomes min(15, 9) = 9.

ii. Moving to i = 1, nums [i] = 6.

We start iterating from the first element to the last.

• Element 6 is not less than mx (2). Thus, mx becomes 6.

For each element, we perform the checks and updates as specified.

mx = max(-infinity, 2) = 2.
 The mirror index from the right is n - 0 - 1 = 6. Element nums[6] = 15.

As we continue, mx and mi will not update any further because the remaining elements (nums[5] = 9 and nums[6] = 15) are

• Correcting for the left boundary due to the two-pointer check in the algorithm, 1 should be 1, not 2. The element nums [1] = 6 initiated the

iii. Next, i = 2, nums[i] = 4.

The mirror index from the right is n - 2 - 1 = 4. Element nums[4] = 10.
mi is now the min(9, 10) = 9.

mi stays at 9 since 8 is less than mi.

• 10 is not less than mx. Update mx to 10.

• 4 is less than mi (9), so we update 1 = 2.

• 4 is less than mx (6), hence we update r = 2.

• 8 is not less than mx (6). So, mx becomes 8.

• The mirror index is n - 3 - 1 = 3. Element nums[3] = 8.

• The mirror index is n - 4 - 1 = 2. Element nums [2] = 4.

After completing the iteration, 1 is found to be 2 and r is also 2.

We need to find the full range of the unsorted segment.

iv. For i = 3, nums[i] = 8.

• The mirror index from the right is n - 1 - 1 = 5. Element nums[5] = 9.

- v. When i = 4, nums[i] = 10.
- greater than all seen max values and less than all seen min values, respectively.

 Step 3: Calculate the Subarray Length

descending sequence.

Solution Implementation

from typing import List

class Solution:

Python

Thus, the length of the subarray is r - l + 1 = 5 - 1 + 1 = 5.
 The subarray [6, 4, 8, 10, 9] from indices 1 to 5 is the smallest segment that, once sorted, the entire array nums becomes

sorted. The length of this segment is 5.

def findUnsortedSubarrav(self. nums: List[int]) -> int:
 # Initializing minimum and maximum values found while traversing
 minimum, maximum = float('inf'), float('-inf')

Initializing the left and right boundaries of the subarray

Found a new right boundary if current number is less than the maximum seen so far

Found a new left boundary if current number is more than the minimum seen so far

Iterate over the list to find the right boundary of the subarray

• They both referred to the same element because we performed the 1 update after the r update.

We already know r should be at least 5, because nums [5] = 9 was out of place.

else:
 # Update the maximum value seen so far
 maximum = nums[i]

Checking the left boundary using the mirrored index (from the end)

if minimum < nums[length - i - 1]:</pre>

 $left_boundary = length - i - 1$

minimum = nums[length - i - 1]

Update the minimum value seen so far

right_boundary = i

left_boundary, right_boundary = -1, -1

Length of the input list

for i in range(length):

if maximum > nums[i]:

length = len(nums)

else:

class Solution {

#include <vector>

class Solution {

public:

#include <algorithm>

using namespace std;

If the right boundarv is still -1, the array is already sorted, hence return 0
Otherwise return the length of the unsorted subarray
return 0 if right_boundary == -1 else right_boundary - left_boundary + 1
Java

int n = nums.length; // Length of the input array

public int findUnsortedSubarray(int[] nums) {

// Initialize variables

} else {

for (int i = 0; i < n; ++i) {

if (maxUnsorted > nums[i]) {

int findUnsortedSubarray(vector<int>& nums) {

for (int i = 0; i < n; ++i) {

right = i;

} else {

} else {

* @return The length of such subarray.

let minOutOfOrder = Infinity;

if (maxVal > nums[i]) {

maxVal = nums[i];

left = n - i - 1;

if (minVal < nums[n - i - 1]) {

minVal = nums[n - i - 1];

return right == -1 ? 0 : right - left + 1;

function findUnsortedSubarray(nums: number[]): number {

const int INF = 1e9; // Use uppercase for constant values

int n = nums.size(); // Size of the input vector

// Simultaneously, check the unsorted subarray from the end of the array
if (minUnsorted < nums[n - i - 1]) {
 leftIndex = n - i - 1; // Update the left boundary
} else {
 minUnsorted = nums[n - i - 1]; // Update the current min for the sorted part
}

// If rightIndex is not updated, the array is fully sorted, return 0
// Otherwise, return the length of the subarray that must be sorted
return rightIndex == -1 ? 0 : rightIndex - leftIndex + 1;
}
</pre>
C++

int left = -1, right = -1; // Initialize the left and right pointers of the subarray

int minVal = INF, maxVal = -INF; // Initialize the minimum and maximum values seen so far

// Iterate forward through the vector to find the right boundary of the unsorted subarray

// Iterate backward through the vector to find the left boundary of the unsorted subarray

// If right is -1, the arrav is alreadv sorted; otherwise, calculate the length of the unsorted subarray

maxUnsorted = nums[i]; // Update the current max for the sorted part

final int MAX INT = Integer.MAX VALUE; // Use MAX_VALUE constant for clarity

// Loop through the array to find the unsorted subarray's boundaries

rightIndex = i; // Update the right boundary

int leftIndex = -1, rightIndex = -1; // Track the left and right boundaries of the subarray

// If current max is greater than the current element, it might belong to the unsorted part

int minUnsorted = MAX_INT, maxUnsorted = Integer.MIN_VALUE; // Set initial values for min and max of the unsorted subarray

/** * Finds the length of the smallest contiguous subarray that, if sorted, results in the whole array being sorted. * @param nums Array of numbers to be checked.

*/

let left = -1:

let right = -1;

};

```
let maxOutOfOrder = -Infinity;
    const length = nums.length;
    // Traverse the array to find the boundaries where the order is incorrect.
    for (let i = 0; i < length; ++i) {</pre>
        // From the left, identify the rightmost index that is out of order.
        if (maxOutOfOrder > nums[i]) {
            right = i;
        } else {
            maxOutOfOrder = nums[i];
        // From the right, identify the leftmost index that is out of order.
        if (minOutOfOrder < nums[length - i - 1]) {</pre>
            left = length -i - 1;
        } else {
            minOutOfOrder = nums[length - i - 1];
    // If no out-of-order elements are found, the array is already sorted (return 0).
    // Otherwise, return the length of the unsorted subarray.
    return right === -1 ? 0 : right - left + 1;
from typing import List
class Solution:
    def findUnsortedSubarray(self, nums: List[int]) -> int:
        # Initializing minimum and maximum values found while traversing
        minimum, maximum = float('inf'), float('-inf')
```

Otherwise return the length of the unsorted subarray return 0 if right_boundary == -1 else right_boundary - left_boundary + 1 Time and Space Complexity

if minimum < nums[length - i - 1]:</pre>

 $left_boundary = length - i - 1$

minimum = nums[length - i - 1]

left_boundary, right_boundary = -1, -1

Length of the input list

for i in range(length):

if maximum > nums[i]:

right_boundary = i

maximum = nums[i]

length = len(nums)

else:

else:

depend on the size of the array.

Initializing the left and right boundaries of the subarray

Update the maximum value seen so far

Update the minimum value seen so far

Iterate over the list to find the right boundary of the subarray

Checking the left boundary using the mirrored index (from the end)

If the right boundary is still -1, the array is already sorted, hence return 0

Found a new right boundary if current number is less than the maximum seen so far

Found a new left boundary if current number is more than the minimum seen so far

The time complexity of the given code is O(n). This is because the code contains only one for loop that iterates over all elements of the array, nums, exactly once. Inside the loop, it performs a constant number of operations for each element, which does not

The space complexity of the code is 0(1) which means constant space. No additional space is utilized that is dependent on the input size n, besides a fixed number of individual variables (mi, mx, l, r, and n) that occupy space which does not scale with the size of the input array.