## 1836. Remove Duplicates From an Unsorted Linked List

Medium Hash Table **Linked List** 

**Problem Description** 

**Leetcode Link** 

than once. Once we have identified such values, we are then to delete all the nodes containing any of those values from the linked list. The goal is to return the linked list after we've made all the necessary deletions.

In this problem, we are given the head of a singly linked list. Our task is to find all the values within this linked list that appear more

Intuition

To solve this problem, we need a way to track the frequency of each value present in the linked list. A common and efficient way to do this is by using a hash table, also known as a dictionary in Python, where the keys are the values from the linked list and the corresponding values are the counts of occurrences for each key.

The solution approach involves two main steps. First, we need to traverse the entire linked list to populate the hash table with the

list. In the second step, we need to traverse the linked list again and this time, delete nodes that have a count greater than one. This

correct counts for each value. With the completed hash table, we can then identify which values appear more than once in the linked

amounts to checking the hash table for the count of the current node's value. If it's greater than one, it means this value appears multiple times and hence the node should be deleted. We need to carefully update the next pointers of the nodes that are not deleted to ensure we have a properly linked list at the end.

A dummy node is typically used as an anchor to manage the head of the list during deletion, especially in cases where the head node itself might need to be deleted. This dummy node initially points to the head of the list, and we start our iteration from the head while keeping track of the previous node as well. If a node needs to be deleted, we can bypass it by setting the next pointer of the previous node to the current node's next. If a node doesn't need to be deleted, we just move the previous pointer to the current

node. After iterating through the entire list and making necessary deletions, the dummy's next pointer points to the head of the modified list, which we return as the final result. **Solution Approach** 

The implementation of the solution begins with importing Counter from the collections module in Python. The Counter class provides a convenient way to count hashable objects. It is essentially a dictionary where elements are stored as dictionary keys and

# Here's the step-by-step breakdown of the implementation:

their counts are stored as dictionary values.

1. Initialize the Counter: An instance of the Counter is created, which will keep track of the number of occurrences of each element in the list. This is done by the line cnt = Counter(). 2. First Pass - Count the Occurrences: The first traversal of the list occurs here, fulfilling the responsibility of counting

the count of the current value cur.val by doing cnt[cur.val] += 1. This forms the frequency mapping required to identify

3. Setup the Dummy Node and Pointers: A dummy node is created with ListNode(0, head). This node is a placeholder to help

occurrences of each value. We continue traversing the list until we reach the end (cur is None). During the traverse, we increment

duplicates.

to the next of the current, effectively bypassing the current node in the list.

cur is reset to head. 4. Second Pass - Delete Duplicates: The list is traversed again. This time we have our frequency map ready, and hence for each node, we check if its value appears more than once by verifying cnt[cur.val] > 1. If this condition holds, it means the node is a

manage deletions, especially when the head of the list might need to be deleted. The pre pointer is set to the dummy node and

duplicate and should be removed from the list. To delete the current node cur, we set the next pointer of the previous node pre

However, if the current node's value does not appear more than once, we need to keep it, and simply update the pre pointer to reference the current node. After either of these checks is performed, we move the current pointer cur to the next node in the list (cur.next).

5. Return Modified List: After the traversal is complete and we have erased all the nodes that had duplicate values, the list is now

This algorithm effectively solves the problem using O(n) time complexity for the two traversals and O(n) space complexity for the

modified. The dummy next holds the reference to the new head of the list, which is returned as the final result.

**Example Walkthrough** Let's assume we have a singly linked list with the following values:  $3 \rightarrow 4 \rightarrow 4 \rightarrow 2 \rightarrow 3 \rightarrow 1 \rightarrow 2$ .

1. Initialize the Counter: We start by creating an empty Counter object: cnt = Counter().

2. First Pass - Count the Occurrences: We traverse the list and count the occurrences of each value. The counter after this pass

Our objective is to identify all values that appear more than once and then remove all nodes containing any such values. Following

### 3. Setup the Dummy Node and Pointers: We create the dummy node and set up our pointers. Now, we have dummy -> 3 -> 4 -> 4 -> 2 -> 3 -> 1 -> 2, with pre pointing to dummy and cur pointing to the first node with value 3.

Python Solution

class ListNode:

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44 }

from collections import Counter

self.val = val

self.next = next

current\_node = head

while current\_node:

current = head

while current:

ListNode current = head;

while (current != null) {

while (current != null) {

} else {

return dummy.next;

current = current.next;

ListNode dummy = new ListNode(0, head);

if (valueCount.get(current.val) > 1) {

previous.next = current.next;

previous = current;

# Definition for singly-linked list.

def \_\_init\_\_(self, val=0, next=None):

value\_counter[current\_node.val] += 1

# Start with a dummy node that points to the head of the list.

# Traverse the linked list again, this time to remove duplicates.

# This simplifies edge cases such as deleting the head node.

current\_node = current\_node.next

if value\_counter[current.val] > 1:

cur now points to the first 4.

List traversal is now complete.

our solution approach, here's an example walkthrough:

will look like this:  $cnt = \{3: 2, 4: 2, 2: 2, 1:1\}.$ 

counter.

4. Second Pass - Delete Duplicates: We begin traversing the list again. Here's how we process each node:

cur now points to the second 3, and we bypass it, as cnt[3] > 1.

cur moves to the second 2, which is bypassed, as cnt[2] > 1.

 Look at cur node with value 4, cnt [4] > 1, it is a duplicate, so we bypass it. cur now points to the second 4, we repeat the above step. cur now points to the first 2, and we bypass it, as cnt[2] > 1.

Look at cur node with value 3, cnt[3] > 1, it is a duplicate, so we update pre.next to cur.next, bypassing the current 3.

5. Return Modified List: At this point of the walkthrough, only the node with value 1 remains, and the updated list points to it. So, the dummy next is pointing to the node with value 1, which is now the head of our resulting list.

cur now points to 1, since cnt[1] == 1, it's not a duplicate, pre is updated to reference this node.

Hence, after the algorithm finishes, the linked list that we return will only contain the node with value 1.

class Solution: def deleteDuplicatesUnsorted(self, head: ListNode) -> ListNode: # Create a Counter to keep track of the frequency of each value in the linked list. value\_counter = Counter()

22 dummy\_node = ListNode(0, head) 23 24 # Initialize two pointers, 'previous' and 'current'. # 'previous' will lag one behind 'current' as we traverse the list. 26 previous = dummy node

# If the current node's value has a count greater than 1, it's a duplicate.

# Traverse the linked list to populate the counter with the frequencies of each value.

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                   # Therefore, skip this node by setting the previous node's next to be the next node.
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                   previous.next = current.next
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               else:
                   # If it's not a duplicate, move the 'previous' pointer up to be the 'current' node.
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                   previous = current
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               # Move the current pointer to the next node in the list.
39
               current = current.next
40
           # Return the modified list, starting from the dummy head's next value
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           return dummy_node.next
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Java Solution
   /**
    * Definition for singly-linked list.
    */
   class ListNode {
       int val;
       ListNode next;
       ListNode() {}
       ListNode(int val) { this.val = val; }
       ListNode(int val, ListNode next) { this.val = val; this.next = next; }
9
10 }
11
   class Solution {
       public ListNode deleteDuplicatesUnsorted(ListNode head) {
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           // HashMap to store the frequency of each value in the list
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           Map<Integer, Integer> valueCount = new HashMap<>();
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           // First pass: count the occurrences of each value
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valueCount.put(current.val, valueCount.getOrDefault(current.val, 0) + 1);

// Dummy node to simplify edge cases at the head of the list

current = head; // Start again from the head of the list

// Second pass: remove nodes with values that appear more than once

// If current node's value count is more than 1, skip it

ListNode previous = dummy; // Maintain the node before the current node

// Only move the previous pointer if current node is unique

// Return the next node of the dummy, which is the new head of the modified list

current = current.next; // Move to the next node in the list

// Return the new list starting at the node after the dummy node.

\* Function to delete all duplicates from an unsorted singly-linked list

\* @returns {ListNode | null} - The modified list with duplicates removed

function deleteDuplicatesUnsorted(head: ListNode | null): ListNode | null {

\* @param {ListNode | null} head - The head of the singly-linked list

// Map to store the frequency count of each value in the list

const frequencyCount: Map<number, number> = new Map();

### C++ Solution // Definition for a singly-linked list node. 2 struct ListNode { int val; ListNode \*next; // Constructor to initialize a node with a value and next pointer (default nullptr). ListNode(int x = 0, ListNode \*nextNode = nullptr) : val(x), next(nextNode) {} **}**; 9 class Solution { 10 public: // Function to delete nodes with duplicate values from an unsorted linked list. 11 12 ListNode\* deleteDuplicatesUnsorted(ListNode\* head) { 13 // An unordered map to store the count of each value in the list. 14 unordered\_map<int, int> valueCounts; 15 // First pass: count occurrences of each value in the list. 16 for (ListNode\* currentNode = head; currentNode != nullptr; currentNode = currentNode->next) { 17 valueCounts[currentNode->val]++; 18 19 20 21 // Dummy node to facilitate deletion from the head of the list. ListNode dummyNode; 22 23 24 // Previous node pointer starts from dummy node; current node starts from head. 25 ListNode \*previousNode = &dummyNode, \*currentNode = head; 26 27 // Second pass: remove nodes with values that have more than one occurrence. 28 while (currentNode != nullptr) { 29 // If the current value exists more than once, skip the current node. if (valueCounts[currentNode->val] > 1) { 30 31 previousNode->next = currentNode->next; 32 } else { 33 // Only move the previous node pointer if the current value isn't duplicated. 34 previousNode = currentNode; 35 36 // Move to the next node in the list. 37 currentNode = currentNode->next;

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Typescript Solution

return dummyNode.next;

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// Count the occurrences of each value by traversing the list
       for (let currentNode = head; currentNode !== null; currentNode = currentNode.next) {
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           const value = currentNode.val;
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           frequencyCount.set(value, (frequencyCount.get(value) ?? 0) + 1);
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       // Create a dummy node that points to the head of the list
       const dummyHead = new ListNode(0, head);
       // Traverse the list with two pointers, `previousNode` and `currentNode`, connected as: previousNode -> currentNode
       for (
           let previousNode = dummyHead, currentNode = head;
22
           currentNode !== null;
23
           currentNode = currentNode.next
24
25
           // Check the frequency count of the currentNode's value
26
           if (frequencyCount.get(currentNode.val)! > 1) {
27
               // If count is more than 1, it is a duplicate, remove it by updating the next pointer of the previous node
28
               previousNode.next = currentNode.next;
           } else {
29
               // If current value is not a duplicate, move previousNode pointer to the current node
30
31
               previousNode = currentNode;
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       // Return the modified list, omitting the dummy head
36
       return dummyHead.next;
37 }
38
   // ListNode class definition for reference
   class ListNode {
       val: number;
       next: ListNode | null;
42
43
       constructor(val?: number, next?: ListNode | null) {
44
           this.val = (val === undefined ? 0 : val);
45
           this.next = (next === undefined ? null : next);
46
47
48 }
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Time and Space Complexity
```

The provided code aims to remove all nodes that have duplicate values from an unsorted singly-linked list. The algorithm works in two passes. The first pass counts the occurrences of each value using a counter (cnt), and the second pass removes nodes with values that occur more than once. • Time complexity: The time complexity of the code is O(n) where n is the length of the linked list. This is because each node in

- the list is visited exactly twice once while counting the occurrences (first while loop) and once while removing duplicates (second while loop). Both operations for each node take constant time, so the total time is linear with respect to the number of nodes in the list. • Space complexity: The space complexity of the code is also O(n). This is due to the use of a counter (cnt) to store the
- occurrence count for each value present in the linked list. In the worst case, if all n nodes have unique values, the counter will need to store an entry for each value, resulting in O(n) space used.