

242. Valid Anagram

EasyHash TableStringSorting

Problem Description

The problem presents two input strings, `s` and `t`, and asks to determine whether `t` is an anagram of `s`. An **anagram** is defined as a word or phrase that is created by rearranging the letters of another word or phrase, using all the original letters exactly once. The goal is to return `true` if `t` is an anagram of `s`, and `false` otherwise. This implies that if `t` is an anagram of `s`, both strings should have the same letters with the exact amount of occurrences for each letter.

Intuition

To solve this problem efficiently, we can use a counting approach. The intuition comes from the definition of an anagram – the strings must have the same characters in the same quantities. First, we can quickly check if `s` and `t` are of the same length. If they're not, `t` cannot be an anagram of `s` and we can immediately return `false`.

Once we've established that `s` and `t` are of equal length, the next step is to count the occurrences of each character in `s`. We can do this efficiently by using a hash table (in Python, this could be a `Counter` from the `collections` module). The hash table will map each character to the number of times it appears in `s`.

After setting up the counts for `s`, we traverse the string `t`. For each character in `t`, we decrease its corresponding count in the hash table by one. If at any point we find that a character's count goes below zero, this means that `t` contains more of that character than `s` does, which violates the anagram property. In this case, we can return `false`.

If we can successfully traverse all characters in `t` without any counts going below zero, it means that `t` has the same characters in the same quantities as `s`. Hence, `t` is an anagram of `s`, and we return `true`.

Solution Approach

The solution approach utilizes a hash table to keep track of the characters in string `s`. The use of a hash table allows us to efficiently map characters to the number of times they occur in `s`. This is a common pattern when we need to count instances of items, such as characters in a string.

The solution proceeds with the following steps:

- Check Length:** We compare the lengths of `s` and `t`. If they are different, we immediately return `false` because an anagram must have the same number of characters.
- Initialize Counter:** We initialize a `Counter` from the `collections` module in Python for the first string `s`. This `Counter` object effectively creates the hash table that maps each character to its count in the string.
- Character Traversal & Counting:** We then iterate over each character `c` in the second string `t` and decrement the count of `c` in our Counter by one for each occurrence. The `Counter` object allows us to do this in constant time for each character.
- Check Negative Counts:** We check if the decremented count of any character goes below zero. If this happens, it means that `t` has an extra occurrence of a character that does not match the count in `s`, and we return `false`.
- Return True:** If we go through the entire string `t` without encountering a negative count, this means that `t` has the exact same characters in the same amounts as `s`, confirming that `t` is an anagram of `s`. Thus, we return `true`.

By using a hash table, we achieve a linear time complexity of  $O(n)$ , where  $n$  is the length of the strings, since we go through each string once. This makes the algorithm efficient for large strings. The solution combines the steps of setting up the Counter, iterating over `t`, and verifying the counts into a concise and effective approach.

Example Walkthrough

Let's consider an example where `s = "listen"` and `t = "silent"`. We want to determine if `t` is an anagram of `s` using the solution approach described above.

- Check Length:** We check the lengths of both `s` and `t`. In this case, both strings are of length 6. As they are equal, we proceed to the next step.
- Initialize Counter:** We utilize a `Counter` from the `collections` module in Python to tally the characters in string `s`. So, we get the counts as follows:

{'l': 1, 'i': 1, 's': 1, 't': 1, 'e': 1, 'n': 1}

- Character Traversal & Counting:** We iterate over each character in string `t`.
  - Decrement count for 's': {'l': 1, 'i': 1, 's': 0, 't': 1, 'e': 1, 'n': 1}
  - Decrement count for 'i': {'l': 1, 'i': 0, 's': 0, 't': 1, 'e': 1, 'n': 1}
  - Decrement count for 'l': {'l': 0, 'i': 0, 's': 0, 't': 1, 'e': 1, 'n': 1}
  - Decrement count for 'e': {'l': 0, 'i': 0, 's': 0, 't': 1, 'e': 0, 'n': 1}
  - Decrement count for 'n': {'l': 0, 'i': 0, 's': 0, 't': 1, 'e': 0, 'n': 0}
  - Decrement count for 't': {'l': 0, 'i': 0, 's': 0, 't': 0, 'e': 0, 'n': 0}
- Check Negative Counts:** Throughout the traversal, we check the counts after each decrement. No count goes below zero, meaning there are no extra occurrences of any character.
- Return True:** After traversing `t`, all the counts are exactly zero, which means that `t` has precisely the same characters in the same amounts as `s`. Therefore, according to the solution approach, we return `true`, confirming that `t` is an anagram of `s`.

Solution Implementation

Python

```
from collections import Counter

class Solution:
    def isAnagram(self, s: str, t: str) -> bool:
        # If the lengths of the strings are not equal, they cannot be anagrams
        if len(s) != len(t):
            return False

        # Create a counter dictionary for all characters in the first string 's'
        char_count = Counter(s)

        # Iterate over each character in the second string 't'
        for char in t:
            # Decrease the count of the character in the counter dictionary
            char_count[char] -= 1

            # If the count of a character goes below zero,
            # it means 't' has more occurrences of that character than 's',
            # so 's' and 't' cannot be anagrams
            if char_count[char] < 0:
                return False

        # If we've not returned False till now, 's' and 't' are anagrams
        return True
```

Java

```
class Solution {
    public boolean isAnagram(String s, String t) {
        // Check if both strings are of equal length
        if (s.length() != t.length()) {
            return false; // If not equal, they can't be anagrams
        }

        // Create an array to store the frequency of each letter
        int[] letterCount = new int[26];

        // Iterate over each character of both strings
        for (int i = 0; i < s.length(); i++) {
            // Increment the count for each letter in string s
            letterCount[s.charAt(i) - 'a']++;
            // Decrement the count for each letter in string t
            letterCount[t.charAt(i) - 'a']--;
        }

        // Check if all counts are zero, indicating anagrams
        for (int count : letterCount) {
            if (count != 0) {
                return false; // If any count is not zero, s and t are not anagrams
            }
        }

        // If all counts are zero, then s and t are anagrams
        return true;
    }
}
```

C++

```
#include <string>
#include <vector>
#include <algorithm>

class Solution {
public:
    // Function to determine if two strings are anagrams of each other
    bool isAnagram(std::string s, std::string t) {
        // If the strings are not of the same size, they cannot be anagrams
        if (s.size() != t.size()) {
            return false;
        }

        // Create a vector of 26 elements to count the occurrences of each letter
        std::vector<int> charCounts(26, 0);

        // Increment and decrement counts for each character in both strings respectively
        for (int i = 0; i < s.size(); ++i) {
            ++charCounts[s[i] - 'a']; // Increment count for the current char in string s
            --charCounts[t[i] - 'a']; // Decrement count for the current char in string t
        }

        // Check if all counts are zero, if so, strings are anagrams, otherwise, they are not
        return std::all_of(charCounts.begin(), charCounts.end(), [](int count) {
            return count == 0;
        });
    }
};
```

TypeScript

```
function isAnagram(source: string, target: string): boolean {
    // Check if both strings are of equal length; if not, they cannot be anagrams
    if (source.length !== target.length) {
        return false;
    }

    // Create an array of 26 elements to represent counts of each letter in the alphabet
    const letterCounts = new Array(26).fill(0);

    for (let i = 0; i < source.length; ++i) {
        // Calculate the count for each letter in the source string (increment)
        letterCounts[source.charCodeAt(i) - 'a'.charCodeAt(0)]++;

        // Calculate the count for each letter in the target string (decrement)
        letterCounts[target.charCodeAt(i) - 'a'.charCodeAt(0)]--;
    }

    // Check if all counts return to zero; if so, the strings are anagrams of each other
    return letterCounts.every(count => count === 0);
}
```

```
from collections import Counter

class Solution:
    def isAnagram(self, s: str, t: str) -> bool:
        # If the lengths of the strings are not equal, they cannot be anagrams
        if len(s) != len(t):
            return False

        # Create a counter dictionary for all characters in the first string 's'
        char_count = Counter(s)

        # Iterate over each character in the second string 't'
        for char in t:
            # Decrease the count of the character in the counter dictionary
            char_count[char] -= 1

            # If the count of a character goes below zero,
            # it means 't' has more occurrences of that character than 's',
            # so 's' and 't' cannot be anagrams
            if char_count[char] < 0:
                return False

        # If we've not returned False till now, 's' and 't' are anagrams
        return True
```

Time and Space Complexity

The time complexity of the provided code is  $O(n)$  where  $n$  represents the length of the string `s` (and `t`, since they are of the same length for the comparison to be valid). This complexity arises because the code iterates over each character in both `s` and `t` exactly once.

The space complexity of the code is  $O(C)$  where  $C$  is the size of the character set used in the strings. In this context, since we are typically dealing with lowercase English letters,  $C$  is equal to 26. The space complexity comes from the use of a counter to store the frequency of each character in string `s`.