977. Squares of a Sorted Array

Two Pointers Sorting

Problem Description The problem provides us with an array of integers nums that is sorted in non-decreasing order. Our task is to return an array

containing the squares of each element from the nums array, and this resulting array of squares should also be sorted in nondecreasing order. For example, if the input is [-4, -1, 0, 3, 10], after squaring each number, we get [16, 1, 0, 9, 100] and then we need to sort this array to get the final result [0, 1, 9, 16, 100].

Intuition

Given that the input array is sorted in non-decreasing order, we realize that the smallest squares might come from the absolute

magnitudes whether they are positive or negative. Therefore, we can use a two-pointer approach: 1. We create two pointers, i at the start of the array and j at the end. 2. We also create an array res of the same length as nums to store our results. 3. We iterate from the end of the res array backward, deciding whether the square of nums[i] or nums[j] should occupy the current position based on which one is larger. This ensures that the largest square is placed at the end of res array first.

values of negative numbers at the beginning of the array as well as the positive numbers at the end. The key insight is that the

squares of the numbers will be largest either at the beginning or at the end of the array since squaring emphasizes larger

4. We move pointer i to the right if nums[i] squared is greater than nums[j] squared since we have already placed the square of nums[i] in the result array.

steps of the algorithm are implemented in the following way:

- 5. Similarly, we move pointer j to the left if the square of nums[j] is greater to ensure we are always placing the next largest square. 6. We continue this process until all positions in res are filled with the squares of nums in non-decreasing order.
- 7. Finally, the res array is returned.
- Solution Approach

Initialize pointers and an array: A pointer i is initialized to the start of the array (0), a pointer j is initialized to the end of the array (n - 1 where n is the length of the array), and an array res of the same length as nums is created to store the results.

the start.

moving towards the start.

in a sorted array of squares which is then returned.

Iterate to fill result array: A loop is used, where the index k starts from the end of the res array (n - 1) and decrements with each iteration. The while loop continues until i is greater than j, which means all elements have been considered.

The solution makes use of a two-pointer approach, an efficient algorithm when dealing with sorted arrays or sequences. The

- Compare and place squares: During each loop iteration, the solution compares the squares of the values at index i and j (nums[i] * nums[i] vs nums[j] * nums[j]) to decide which one should be placed at the current index k of the res array.
- The larger square is placed at res[k], and the corresponding pointer (i or j) is moved. If nums[i] * nums[i] is greater than nums[j] * nums[j], this means that the square of the number pointed to by i is

currently the largest remaining square, so it's placed at res[k], and i is incremented to move to the next element from

Conversely, if nums[j] * nums[j] is greater than or equal to nums[i] * nums[i], then nums[j] squared is placed at

- res [k], and j is decremented to move to the next element from the end. **Decrement k:** After each iteration, k is decremented to fill the next position in the res array, starting from the end and
- This approach uses no additional data structures other than the res array to produce the final sorted array of squares. It is space-optimal, requiring O(n) additional space, and time-optimal with O(n) time complexity because it avoids the need to sort the

squares after computation, which would take O(n log n) if a sorting method was used after squaring the elements.

goal is to compute the squares of each number and get a sorted array as a result. Here's how it works:

Return the result: Once the while loop is done, all elements have been squared and placed in the correct position, resulting

Example Walkthrough

Let's walk through a small example to illustrate the solution approach. We will use the input array nums = [-3, -2, 1, 4]. Our

Initialize pointers and an array: We initialize i to 0, j to 3 (since there are four elements, n - 1 = 3), and an array res with the length of 4 to store the results: res = [0, 0, 0, 0]. **Iterate to fill result array**: We start a while loop with k = 3, which is the last index in the rest array.

• Squares: nums[i] * nums[i] = 9, nums[j] * nums[j] = 16. Since 16 is greater than 9, we place 16 at res[k]: res = [0, 0, 0, 16].

Compare and place squares:

○ First iteration: nums[i] is -3, nums[j] is 4.

Increment i to 1 and decrement k to 1.

 Decrement j to 2 and k to 2. **Next iteration:**

- Now i = 0 (with nums[i] = -3), j = 2 (with nums[j] = 1). • Squares: nums[i] * nums[i] = 9, nums[j] * nums[j] = 1. \circ 9 is greater than 1, so we place 9 at res[k]: res = [0, 0, 9, 16].
- Now i = 1 (with nums[i] = -2), j = 2 (with nums[j] = 1).

• Now i = 2 (with nums[i] = 1), j = 2 (with nums[j] = 1), and k = 0.

Initialize a result array of the same length as the input array

Initialize pointers for the start and end of the input array,

Loop through the array from both ends towards the middle

and a pointer for the position to insert into the result array

start_pointer, end_pointer, result_pointer = 0, length - 1, length - 1

// returns a new array with the squares of each number sorted in non-decreasing order.

int[] sortedSquares = new int[length]; // Create a new array to hold the result

// Compare the squares to decide which to place next in the result array

// and a pointer 'k' for the position to insert into the result array, starting from the end.

// If the start square is greater, place it in the next open position at 'k',

int length = nums.length; // Store the length of the input array

// Initialize pointers for the start and end of the input array,

// Calculate the square of the start and end elements

result[position] = nums[right] * nums[right];

--right; // Move the right pointer one step left

// Return the result which now contains the squares in non-decreasing order

* Returns an array of the squares of each element in the input array, sorted in non-decreasing order.

// Two pointers approach: start from the beginning (i) and the end (i) of the nums array.

// Square of nums[i] is greater, so store it at index k in res and move i forward.

Compare the squared values and add the larger one to the end of the result array

// Square of nums[i] is greater or equal, so store it at index k in res and move j backward.

--position; // Move the position pointer one step left

* @param {number[]} nums - The input array of integers.

const sortedSquares = (nums: number[]): number[] => {

// n is the length of the input array nums.

const n: number = nums.length;

const res: number[] = new Array(n);

* @return {number[]} - The sorted array of squares of the input array.

// res is the resulting array of squares, initialized with the size of nums.

// The index k is used for the current position in the resulting array res.

// The larger square is placed at the end of array res, at index k.

for (let i: number = 0, j: number = n - 1, k: number = n - 1; $i \le j$;) {

// Compare squares of current elements pointed by i and i.

if (nums[i] * nums[i] > nums[i] * nums[i]) {

res[k--] = nums[i] * nums[i];

res[k--] = nums[j] * nums[j];

end_square = nums[end_pointer] ** 2

result[result pointer] = start_square

result[result pointer] = end_square

Move the result pointer to the next position

if start square > end square:

start_pointer += 1

end_pointer -= 1

Return the sorted square array

result pointer -= 1

Time and Space Complexity

else:

return result

int startSquare = nums[start] * nums[start];

// then increment the start pointer.

int endSquare = nums[end] * nums[end];

if (startSquare > endSquare) {

for (int start = 0, end = length - 1, k = length - 1; start <= end;) {

• Squares: nums[i] * nums[i] = 4, nums[j] * nums[j] = 1.

 4 is greater than 1, so we place 4 at res[k]: res = [0, 4, 9, 16]. Increment i to 2 and decrement k to 0.

Next iteration:

Final iteration:

sorted positions.

from typing import List

Python

Solution Implementation

result = [0] * length

while start pointer <= end pointer:</pre>

end_pointer -= 1

Return the sorted square array

public int[] sortedSquares(int[] nums) {

result_pointer -= 1

Square the values at both pointers

end_square = nums[end_pointer] ** 2

start square = nums[start pointer] ** 2

result[result pointer] = end_square

// Method that takes an array of integers as input and

Move the result pointer to the next position

res becomes [1, 4, 9, 16]. **Return the result**: At the end of the loop, we have the final sorted array of squares res = [1, 4, 9, 16]. Following these steps, we have successfully transformed the nums array into a sorted array of squares without needing to sort

There's only one element left, so we square it and place it at res[k]: nums[i] * nums[i] = 1.

class Solution: def sortedSquares(self, nums: List[int]) -> List[int]: # Get the length of the input array length = len(nums)

them again after squaring. This approach efficiently uses the original sorted order to place the squares directly in the correct

Compare the squared values and add the larger one to the end of the result array if start square > end square: result[result pointer] = start_square start_pointer += 1

else:

return result

Java

class Solution {

```
sortedSquares[k--] = startSquare;
                ++start;
            } else {
                // If the end square is greater or equal, place it in the next open position at 'k',
                // then decrement the end pointer.
                sortedSquares[k--] = endSquare;
                --end;
        // Return the array with sorted squares
        return sortedSquares;
C++
#include <vector>
using namespace std;
class Solution {
public:
    vector<int> sortedSquares(vector<int>& nums) {
        int size = nums.size();
        vector<int> result(size); // This will store the final sorted squares of numbers
        // Use two pointers to iterate through the array from both ends
                       // Start pointer for the array
        int right = size - 1; // End pointer for the array
        int position = size - 1; // Position to insert squares in the result array from the end
        // While left pointer does not surpass the right pointer
        while (left <= right) {</pre>
            // Compare the square of the elements at the left and right pointer
            if (nums[left] * nums[left] > nums[right] * nums[right]) {
                // If the left square is larger, place it in the result array
                result[position] = nums[left] * nums[left]:
                ++left; // Move the left pointer one step right
            } else {
                // If the right square is larger or equal, place it in the result array
```

} else {

++i:

--j;

return result;

};

/**

TypeScript

```
// Return the sorted array of squares.
   return res;
};
from typing import List
class Solution:
   def sortedSquares(self, nums: List[int]) -> List[int]:
       # Get the length of the input array
        length = len(nums)
       # Initialize a result array of the same length as the input array
        result = [0] * length
       # Initialize pointers for the start and end of the input array,
        # and a pointer for the position to insert into the result array
        start_pointer, end_pointer, result_pointer = 0, length - 1, length - 1
       # Loop through the array from both ends towards the middle
       while start pointer <= end pointer:</pre>
            # Square the values at both pointers
            start square = nums[start pointer] ** 2
```

Time Complexity

The time complexity of the code is O(n). Each element in the nums array is accessed once during the while loop. Despite the fact that there are two pointers (i, j) moving towards each other from opposite ends of the array, each of them moves at most n steps. The loop ends when they meet or cross each other, ensuring that the total number of operations does not exceed the number of elements in the array.

Space Complexity

The space complexity of the code is O(n). Additional space is allocated for the res array, which stores the result. This array is of the same length as the input array nums. No other additional data structures are used that grow with the size of the input, so the total space required is directly proportional to the input size n.