1229. Meeting Scheduler

Two Pointers Sorting

Problem Description

Medium Array

meeting duration. Each person's schedule is represented by a list of non-overlapping time slots where a time slot is an array [start, end] showing availability from start to end. The goal is to find the earliest starting time slot that is available in both schedules and lasts at least for the given duration. If there's no such common time slot, we return an empty array.

The problem is about finding a mutual meeting time slot between two people given their individual schedules and a required

Intuition To solve this problem, we can use the two-pointer technique. Since no individual time slots overlap within a single person's schedule, we can sort both schedules by the starting times of the slots. We then compare the slots from both schedules to find overlapping slots.

We use two pointers i and j to traverse slots1 and slots2 respectively. In each iteration, we find the latest start time by taking the maximum of slots1[i][0] and slots2[j][0] and the earliest end time by taking the minimum of slots1[i][1] and slots2[j][1]. If the overlapping time between the latest start and earliest end is greater than or equal to the required duration, we have found a suitable time slot and return the start and end of this meeting slot.

If the overlap is not sufficient, we move the pointer forward in the list which has the earlier ending time slot, hoping to find a longer slot that might overlap with the other person's next slot. This process continues until we either find a suitable slot or

exhaust all available slots in either list.

considering the earliest available slots first and eliminates the need for checking past time slots. Sorting is crucial as it sets up the structure for the two-pointer technique to work effectively.

these values:

if end - start >= duration:

slots after the initial sort.

Solution Approach

Two-pointer Technique: Two pointers, i and j, are used to iterate through slots1 and slots2 respectively. At each step, i refers to the current slot in the first person's schedule, and j refers to the current slot in the second person's schedule.

corresponding to the slot with the earlier end time is incremented:

Let's consider an example to illustrate the solution approach:

Person A's schedule (slots1): [[10, 50], [60, 120], [140, 210]]

slots1[i] = [10, 50] and slots2[j] = [0, 15]

slots1[i] = [10, 50] and slots2[j] = [25, 50]

Required duration for the meeting: 8 minutes

operation in this example.

Person B's schedule (slots2): [[0, 15], [25, 50], [60, 70], [80, 100]]

The provided Python solution follows a straightforward two-pointer approach to solve the problem:

Finding Overlaps: For the current pair of time slots pointed by i and j, we calculate the overlap by determining the maximum of the two start times and the minimum of the two end times. The variables start and end are used to record

Sorting the time slots: Both slots1 and slots2 are sorted based on their starting times. This ensures that we are always

- start = max(slots1[i][0], slots2[i][0]) end = min(slots1[i][1], slots2[j][1])
- Checking Duration: We then check if the overlapping duration is greater than or equal to the required duration by subtracting start from end:

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If the condition is met, we have found a valid time slot and can return [start, start + duration] as the solution.
Advancing the Pointers: If the overlapping time slot is not long enough, we need to discard the time slot that ends earlier and
move forward. This decision is made by comparing the end times of the current time slots pointed by i and j. The pointer
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if slots1[i][1] < slots2[j][1]:</pre> i += 1else: j += 1

```
end of slots1 or j reaches the end of slots2). If no common time slots with sufficient duration are found by the end of
   either list, we return an empty array [] as specified.
This algorithm makes efficient use of the sorted structure of the time slots and the two-pointer technique to minimize the number
of comparisons and quickly find the first suitable time slot, achieving a time complexity that is linear in the size of the input time
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Continuation and Termination: The above steps are continued in a loop until one of the lists is exhausted (i.e., i reaches the

Example Walkthrough

This step ensures that we're always trying to find overlap with the nearest possible future slot.

Two-pointer Technique: Initialize two pointers, i for slots1 and j for slots2. Initially, i = 0 and j = 0. **Finding Overlaps:** First Comparison:

■ The overlap is from 10 (max of 10 and 0) to 15 (min of 50 and 15), which is 5 minutes long. This is not enough for the 8-minute

Sorting the time slots: Both schedules are already sorted based on their starting times, eliminating the need for a sort

■ We have found a suitable slot, so we can return the result [25, 25 + 8], which is [25, 33]. This result indicates that Person A and Person B can successfully schedule a meeting starting at 25 minutes past the hour and

■ The overlap is from 25 to 50, which is 25 minutes long and suffices for the 8-minute duration.

Move the pointer j forward because slots2[j][1] is less than slots1[i][1].

The method described above is efficient because it continuously seeks the earliest possible meeting time by looking for overlaps in the schedules and moves forward based on the end times of the current slots. As soon as a fitting time slot is found, it returns

the result without unnecessary comparisons of later time slots, saving time and computation.

from typing import List

slots1.sort()

slots2.sort()

class Solution:

Solution Implementation

duration.

Second Comparison:

lasting until 33 minutes past the hour.

Python

If the overlapping period is greater than or equal to the required duration

Otherwise, increase the index for slots2 to check the next slot

result = sol.minAvailableDuration([[10, 50], [60, 120], [140, 210]], [[0, 15], [60, 70]], 8)

// Sort both sets of time slots in ascending order based on the start times

// Loop through both sets of time slots to find a common free duration

// Find the latest start time and earliest end time between two slots

int latestStart = std::max(slots1[index1][0], slots2[index2][0]);

int earliestEnd = std::min(slots1[index1][1], slots2[index2][1]);

// Move to the next slot in the set with the earlier ending time

std::sort(slots1.begin(), slots1.end());

std::sort(slots2.begin(), slots2.end());

// Use indices to track the current slot in each set

if (earliestEnd - latestStart >= duration) {

if (slots1[index1][1] < slots2[index2][1]) {</pre>

while (index1 < slots1.size() && index2 < slots2.size()) {</pre>

return {latestStart, latestStart + duration};

// If no common free duration is found, return an empty vector

If no overlapping period long enough for the meeting is found, return an empty list

Return the start time and start time plus the duration

return [overlap_start, overlap_start + duration]

Sort the time slots for both people to allow for easy comparison

Initialize variables to track current index in slots1 and slots2

if overlap end - overlap start >= duration:

def minAvailableDuration(self, slots1: List[List[int]], slots2: List[List[int]], duration: int) -> List[int]:

index1 = index2 = 0# Get the total number of slots for both people total slots1 = len(slots1) total_slots2 = len(slots2)

Iterate over the slots until at least one person's slots are fully checked while index1 < total slots1 and index2 < total slots2:</pre> # Find the overlapping start time of the current slots overlap start = max(slots1[index1][0], slots2[index2][0]) # Find the overlapping end time of the current slots overlap_end = min(slots1[index1][1], slots2[index2][1])

If the end time in slots1 is before the end time in slots2, # increase the index for slots1 to check the next slot if slots1[index1][1] < slots2[index2][1]:</pre> index1 += 1

else:

return []

Example usage:

sol = Solution()

index2 += 1

```
# print(result) # This will output: [60, 68]
Java
class Solution {
    public List<Integer> minAvailableDuration(int[][] slots1, int[][] slots2, int duration) {
        // Sort the time slots for both people based on the start times
        Arrays.sort(slots1, (a, b) -> a[0] - b[0]);
        Arrays.sort(slots2, (a, b) \rightarrow a[0] - b[0]);
        int index1 = 0; // Index for navigating person 1's time slots
        int index2 = 0; // Index for navigating person 2's time slots
        int len1 = slots1.length; // Total number of slots for person 1
        int len2 = slots2.length; // Total number of slots for person 2
        // Iterate through both sets of slots
        while (index1 < len1 && index2 < len2) {</pre>
            // Calculate the overlap start time
            int overlapStart = Math.max(slots1[index1][0], slots2[index2][0]);
            // Calculate the overlap end time
            int overlapEnd = Math.min(slots1[index1][1], slots2[index2][1]);
            // Check if the overlapping duration is at least the required duration
            if (overlapEnd - overlapStart >= duration) {
                // If so, return the start time of the meeting and start time plus duration
                return Arrays.asList(overlapStart, overlapStart + duration);
            // Move to the next slot in the list that has an earlier end time
            if (slots1[index1][1] < slots2[index2][1]) {</pre>
                index1++;
            } else {
                index2++;
        // If no common slot is found that fits the duration, return an empty list
        return Collections.emptyList();
C++
#include <vector>
#include <algorithm>
```

std::vector<int> minAvailableDuration(std::vector<std::vector<int>>& slots1, std::vector<std::vector<int>>& slots2, int duration)

// Check if the overlap duration between two slots is greater than or equal to the desired duration

// If so, return the start time and start time plus duration as the available duration

};

class Solution {

int index1 = 0;

int index2 = 0;

++index1;

++index2;

} else {

return {};

public:

```
TypeScript
function minAvailableDuration(slots1: number[][], slots2: number[][], duration: number): number[] {
    // Sort both sets of time slots by start time in ascending order
    slots1.sort((a, b) => a[0] - b[0]);
    slots2.sort((a, b) => a[0] - b[0]);
    // Initialize pointers to traverse the slots
    let index1 = 0:
    let index2 = 0;
    while (index1 < slots1.length && index2 < slots2.length) {</pre>
        // Calculate the maximum start time and the minimum end time between two slots
        const latestStart = Math.max(slots1[index1][0], slots2[index2][0]);
        const earliestEnd = Math.min(slots1[index1][1], slots2[index2][1]);
        // Check if there is a slot sufficient for the required duration
        if (earliestEnd - latestStart >= duration) {
            // If found, return the time range starting from latestStart and continuing for the duration
            return [latestStart, latestStart + duration];
        // Move to the next slot in the list with the earlier end time
        if (slots1[index1][1] < slots2[index2][1]) {</pre>
            index1++; // Move the pointer in the first list forward
        } else {
            index2++; // Move the pointer in the second list forward
    // If no matching slots were found, return an empty array
    return [];
// Example usage:
// const result = minAvailableDuration([[10.50].[60.120].[140.210]].[[0.15].[60.70]].8);
// console.log(result); // Should log a valid time slot with at least 8 minutes of duration, or an empty array if there's no slot.
from typing import List
class Solution:
    def minAvailableDuration(self, slots1: List[List[int]], slots2: List[List[int]], duration: int) -> List[int]:
        # Sort the time slots for both people to allow for easy comparison
        slots1.sort()
        slots2.sort()
        # Initialize variables to track current index in slots1 and slots2
```

```
return []
# Example usage:
# sol = Solution()
# result = sol.minAvailableDuration([[10. 50], [60, 120], [140, 210]], [[0, 15], [60, 70]], 8)
```

else:

index1 += 1

index2 += 1

print(result) # This will output: [60, 68]

Time and Space Complexity

slots1 and O(n log n) for slots2.

index1 = index2 = 0

total slots1 = len(slots1)

total_slots2 = len(slots2)

Get the total number of slots for both people

while index1 < total slots1 and index2 < total slots2:</pre>

if overlap end - overlap start >= duration:

if slots1[index1][1] < slots2[index2][1]:</pre>

Find the overlapping start time of the current slots

overlap_end = min(slots1[index1][1], slots2[index2][1])

return [overlap_start, overlap_start + duration]

increase the index for slots1 to check the next slot

Find the overlapping end time of the current slots

overlap start = max(slots1[index1][0], slots2[index2][0])

Iterate over the slots until at least one person's slots are fully checked

Return the start time and start time plus the duration

If the end time in slots1 is before the end time in slots2,

If the overlapping period is greater than or equal to the required duration

Otherwise, increase the index for slots2 to check the next slot

If no overlapping period long enough for the meeting is found, return an empty list

The given algorithm primarily consists of two parts: sorting the time slots and then iterating through the sorted lists to find a common available duration.

Time Complexity

First, we sort both slots1 and slots2. Assuming that slots1 has m intervals and slots2 has n intervals, the time taken to sort these lists using a comparison-based sorting algorithm, like Timsort (Python's default sorting algorithm), is 0(m log m) for

pair of slots once, which leads to a complexity of 0(m + n) since each list is traversed at most once. The overall time complexity is the sum of the complexities of sorting and iterating through the slot lists, which is 0(m log m + n log n + m + n). However, the log factor dominates the linear factor in computational complexity analysis for large values.

After sorting, we have a while loop that runs until we reach the end of one of the slot lists. In the worst case, we'll compare each

Hence, the time complexity simplifies to $0(m \log m + n \log n)$. **Space Complexity**

The space complexity of the algorithm is determined by the space we use to sort slots1 and slots2. Since the Timsort algorithm can require a certain amount of space for its operation, the space complexity is 0(m+n) due to the auxiliary space used in sorting. All other operations use a constant amount of space, and no additional space is used that depends on the input size, so the space complexity remains O(m+n).