# 2644. Find the Maximum Divisibility Score

<u>Array</u> Easy

### **Problem Description**

The problem provides us with two arrays: nums (containing integers) and divisors (containing possible divisors). We are tasked with finding the integer from divisors which has the highest divisibility score, where the divisibility score of an element in divisors is the count of elements in nums that it can divide evenly (without leaving a remainder). In case more than one divisor has the same highest score, we need to return the smallest among them.

To put it simply, we want to answer the question: "Which number in divisors divides the most numbers in nums with no remainder?" and if there's a tie, we select the smallest number.

Intuition

its divisibility score. We begin by setting a variable mx to 0, to keep track of the maximum score found so far, and another variable ans to keep the

This problem is suited for a brute force approach that checks each number in divisors against every number in nums to calculate

- divisor with the maximum score. We initialize ans with the first element of divisors since we need to return an element from divisors in case all scores are zero. We iterate over each divisor in divisors. For each divisor, we count how many numbers in nums are divisible by it. This is done
- using a summation with a conditional check, where we use the modulo operator % to test if the remainder is zero (meaning divisibility). If the current divisor's score is greater than the maximum score we've seen (mx), we update mx with this new score and also
- update ans with the current divisor. If the current divisor's score is equal to the maximum score but less than the current ans, we update ans to this divisor since it
- **Solution Approach**

The solution provided in the reference code implements the brute force approach described in the intuition section. Here is a

## detailed breakdown of the approach:

 We have a for loop that goes through each element div in the divisors array. • Inside the loop, we have a key expression sum(x % div == 0 for x in nums). This expression counts the number of elements in nums that are divisible by div (the x % div == 0 part checks if x is divisible by div without a remainder).

is smaller and we want the smallest divisor in case of a tie.

- This is a generator comprehension within the sum function that goes over each element x in nums and yields 1 if x is divisible by div, and 0 otherwise. The sum function then adds up these 0s and 1s to get the total count. • The cnt variable is used to hold this count, which represents the divisibility score of the current divisor div.
- We compare cnt against mx to determine if we have found a higher divisibility score: o If cnt is greater than mx, we have indeed found a new divisor with a higher score. We assign cnt to mx, and div to ans.
- o If cnt is equal to mx but div is smaller than the current ans, then we update ans with div. This ensures that among divisors with the same highest score, we will return the smallest divisor.

scenarios where the array sizes are manageable and high efficiency is not a critical requirement.

- No additional data structures are required, making this approach efficient in terms of space complexity. The time complexity can be considered as O(n \* m), where 'n' is the length of the nums array and 'm' is the length of the divisors array, since each divisor is checked against all nums.
- The algorithm employs a simple comparison-based technique, and its strength lies in its straightforwardness and direct mapping

to the problem statement without any need for optimization tricks or complicated data structures. This is highly suitable for

**Example Walkthrough** Let's consider two small arrays to illustrate the solution approach:

# Now, we want to figure out which number in divisors can divide the most numbers in nums without leaving a remainder, and if

• divisors = [2, 3, 4]

• nums = [4, 8, 12]

there's a tie, we choose the smallest number.

```
We start with an initial maximum score mx set to 0 and the current answer ans set to the first element in divisors, which is 2.
```

∘ When div = 2: 4 % 2 == 0 (true, score is 1)

■ 12 % 2 == 0 (true, score is 3) ■ Total score for 2 is 3. mx is updated to 3, ans is updated to 2.

After the iteration, the highest score is 3 with ans being 2. Since there's a tie between 2 and 4, we take the smaller number

■ 8 % 3 != 0 (false, score remains 0)

■ Total score for 3 is 1, which is less than mx(3). We don't update mx or ans.

We then iterate through each number in divisors to calculate its divisibility score:

```
■ 12 % 3 == 0 (true, score is 1)
```

∘ When div = 4:

 $\circ$  When div = 3:

4 % 4 == 0 (true, score is 1) ■ 8 % 4 == 0 (true, score is 2)

4 % 3 != 0 (false, score remains 0)

■ 12 % 4 == 0 (true, score is 3) ■ Total score for 4 is also 3, equal to mx. However, since ans is smaller (2 < 4), we do not update ans.

■ 8 % 2 == 0 (true, score is 2)

which is 2. Thus, the final answer is 2, because it can divide all three numbers in nums with no remainder and is the smallest among the

divisors with the highest divisibility score.

Solution Implementation

# max\_count with zero to keep track of the highest score and count.

count divisible = sum(num % divisor == 0 for num in nums)

max\_count, max\_score = count\_divisible, divisor

// Return the divisor with the highest divisibility score.

// Function to find the divisor that has the maximum divisibility score.

// The score for a divisor is defined as the number of elements in 'nums'

// Initialize the maximum count of divisible numbers to zero.

// Initialize the answer with the first divisor as a starting point.

// Count how many numbers in 'nums' are divisible by 'divisor'.

#include <algorithm> // include necessary headers for std::min

int maxDivScore(vector<int>& nums, vector<int>& divisors) {

// that are divisible by this divisor.

int maxDivisibleCount = 0;

// Iterate over each divisor.

for (int divisor : divisors) {

int divisibleCount = 0;

for (int num : nums) {

if (num % divisor == 0) {

++divisibleCount;

# Loop through each divisor in the divisors list.

if max\_count < count\_divisible:</pre>

# Count how many numbers in nums are divisible by the current divisor.

# If the current count is higher than the max\_count found so far,

# update max\_count and max\_score with the current count and divisor

count\_divisible = sum(num % divisor == 0 for num in nums)

for divisor in divisors:

int maxScoreDivisor = divisors[0];

return maxDivisor;

elif max\_count == count\_divisible and max\_score > divisor:

if max\_count < count\_divisible:</pre>

max\_score = divisor

# to the smallest divisor in case of a tie.

# Count how many numbers in nums are divisible by the current divisor.

# If the current count is higher than the max\_count found so far,

# update max\_count and max\_score with the current count and divisor

# than the max\_score, update the max\_score to the current divisor.

# Return the divisor that has the highest divisibility score, giving preference

# If the current count is equal to the max\_count, but the divisor is smaller

class Solution: def max\_div\_score(self, nums: List[int], divisors: List[int]) -> int: # Initialize max\_score with the first element in divisors and

### max\_score, max\_count = divisors[0], 0 # Loop through each divisor in the divisors list. for divisor in divisors:

**Python** 

```
return max_score
Java
class Solution {
   // Method to find the divisor with the highest divisibility score.
   // Divisibility score is defined by how many numbers in nums are divisible by the divisor.
    public int maxDivScore(int[] nums, int[] divisors) {
       // Initialize the answer with the first divisor, assuming it has the maximum score initially.
       int maxDivisor = divisors[0];
       // Initialize the maximum count of divisible numbers for any divisor.
       int maxCount = 0;
       // Iterate through all the divisors
        for (int divisor : divisors) {
            // Initialize count for the current divisor.
            int count = 0;
           // Count how many numbers in nums are divisible by this divisor.
            for (int num : nums) {
                if (num % divisor == 0) {
                    count++;
           // Update the maxDivisor and maxCount if the current divisor has a higher count.
            if (maxCount < count) {</pre>
                maxCount = count;
                maxDivisor = divisor;
            } else if (maxCount == count) {
                // If the current divisor has the same count, choose the smallest one.
                maxDivisor = Math.min(maxDivisor, divisor);
```

C++

public:

#include <vector>

class Solution {

```
// If the current count is greater than the maximum found so far, update the maximum count
            // and change the answer to the current divisor.
            if (maxDivisibleCount < divisibleCount) {</pre>
                maxDivisibleCount = divisibleCount;
                maxScoreDivisor = divisor;
            // If the current count equals the maximum found so far, select the smaller divisor.
            } else if (maxDivisibleCount == divisibleCount) {
                maxScoreDivisor = std::min(maxScoreDivisor, divisor);
        // Return the divisor with the maximum divisibility score.
        // In case of a tie, the smallest such divisor is returned.
        return maxScoreDivisor;
};
TypeScript
// Function that calculates the maximum division score for a given array of numbers and divisors.
// The division score is defined by the number of times the numbers in the array can be evenly divided by the divisors.
// The function returns the divisor that gives the highest division score. In case of a tie, it returns the smallest divisor.
function maxDivScore(nums: number[], divisors: number[]): number {
    let bestDivisor: number = divisors[0]; // Initialize bestDivisor as the first divisor
    let maxDivisibleCount: number = 0; // Initialize maximum divisible count (division score) as 0
    // Loop through each divisor
    for (const divisor of divisors) {
        // Calculate the division score for the current divisor by reducing the nums array
        const divisibleCount = nums.reduce((count, num) => count + (num % divisor === 0 ? 1 : 0), 0);
        // Update the bestDivisor and maxDivisibleCount if this divisor has a higher division score
        if (maxDivisibleCount < divisibleCount) {</pre>
            maxDivisibleCount = divisibleCount;
            bestDivisor = divisor;
        } else if (maxDivisibleCount === divisibleCount && bestDivisor > divisor) {
            // If the division score is the same but the current divisor is smaller, update the bestDivisor
            bestDivisor = divisor;
    // Return the divisor with the highest division score (bestDivisor)
    return bestDivisor;
```

```
def max_div_score(self, nums: List[int], divisors: List[int]) -> int:
   # Initialize max_score with the first element in divisors and
    # max_count with zero to keep track of the highest score and count.
    max score, max count = divisors[0], 0
```

class Solution:

```
max_count, max_score = count_divisible, divisor
           # If the current count is equal to the max_count, but the divisor is smaller
           # than the max_score, update the max_score to the current divisor.
           elif max_count == count_divisible and max_score > divisor:
               max_score = divisor
       # Return the divisor that has the highest divisibility score, giving preference
       # to the smallest divisor in case of a tie.
       return max_score
Time and Space Complexity
  The given Python code defines the method maxDivScore which finds the divisor from the list divisors that maximizes the number
  of elements in nums that can be evenly divided by it. In case of ties, the smallest such divisor is returned.
```

**Time Complexity:** 

iterates once for each element in divisors. Inside the loop, the sum function iterates over each element in nums: • Let n be the length of nums. • Let d be the length of divisors.

The time complexity of the method can be determined by analyzing the for loop and the sum function within it. The for loop

For each divisor, we perform n modulus operations and n equality checks. Therefore, for all divisors, we perform this operation d times. The time complexity is 0(n\*d) because we have two nested loops: one iterating over divisors and the other over nums.

**Space Complexity:** The space complexity is determined by the extra space used by the function beyond the input lists. In this case, the only extra space used are a few variables (ans, mx, div, and cnt) that remain constant regardless of input size. Hence, the space complexity is 0(1) as there are no data structures used that scale with the size of the input.