1849. Splitting a String Into Descending Consecutive Values

## **Problem Description**

String ]

Medium

**Backtracking** 

substrings such that: 1. The numerical values of these substrings, when converted from strings to numbers, form a sequence in descending order.

This problem asks us to check whether a given string s, comprised exclusively of digits, can be split into two or more non-empty

For example, if s = "543210", it is possible to split it into ["54", "3", "2", "10"], where each number is 1 less than the previous, and

so the condition is met.

2. Backtracking: If at any point the condition is not met, the function backtracks, trying a different split.

2. The difference between the numerical values of each pair of adjacent substrings is exactly 1.

There cannot be any rearrangement of the substrings.

The key constraints are that the substrings must be non-empty and must appear in the same order they do in the original string.

Intuition

# 1. Recursion: Recursively split the string and check if the current split creates a number that is exactly 1 less than the preceding number.

different points and recursively check if the resulting substrates satisfy the conditions. The solution uses three key ideas:

3. Early Stopping: If the given string has been entirely traversed and more than one substring meeting the conditions has been found, we can stop

To solve this problem, we can use a depth-first search (DFS) algorithm. The DFS approach will try to partition the string s at

To implement the solution:

the search and return true.

• A recursive function dfs is defined, which takes parameters i (the current starting index in s), x (the last numerical value of the substring), and k (the count of valid splits found so far).

• Initially, dfs is called with starting index 0, x set as -1 (as there is no previous number), and k as 0 (no splits found yet). • Within dfs, we iterate over the string starting from index i, attempting to form a new substring from s[i] to s[j]. • We check if our substring satisfies the conditions; if it does, we continue the search by recursively calling dfs with the next starting index

• If at any point, a full pass through the string is made (i equals the length of s) and more than one valid substring (k > 1) has been found, the

**Solution Approach** 

(j+1), the new number y, and increment k.

- function returns true, indicating the splitting is possible. Otherwise, the function eventually returns false. Using DFS, this solution effectively searches for all possible substrings that might fit the criteria, and upon finding the first valid
- series of splits, it returns true. If no valid series is found by the time the entire string has been traversed, it returns false.
- The implementation of the solution revolves around a classic DFS algorithm. Here's a step-by-step explanation of the approach using the given Python code:

Initialization: A helper function, dfs, is defined inside the main class Solution. This recursive function is at the heart of the

Recursive Search: Inside the dfs function, starting from index i, the code attempts to build a substring piece by piece by

Conditions Check: After each additional character is included in y, we check if y is exactly one less than the previous

Continuation and Backtracking: If the condition fails or the recursive call doesn't result in a valid solution, the function

continues to the next iteration of the loop, effectively trying a longer substring (more characters included from s) or

Main Function Call: To start the process, the splitString method in the Solution class calls dfs(0, -1, 0) indicating it

starts looking for splits from the beginning of the string (i = 0), with no prior number (x = -1), and no splits found yet (k = 0)

DFS implementation. It aims to try all possible splits of the string s beginning from index i. Base Case: If dfs reaches the end of the string (i = len(s)), it checks whether at least two substrings with the required properties have been found (k > 1). If so, the function returns True; otherwise, False.

#### iterating from i to each j within the loop. The goal is to form a numerical value y of the current substring being considered (s[i:j+1]).

number x. If x is -1 (indicating this is the first number), any y is acceptable (since there is no prior number to compare with). o If this condition is satisfied, the function immediately proceeds with a recursive call, dfs(j + 1, y, k + 1), where j + 1 is the new

- starting index for the next substring, y is now the last number, and k + 1 indicates one more valid split has been found. ∘ If the returned value from this recursive call is True, indicating that proceeding from index j + 1 has led to a valid sequence, the current dfs also returns True.
- backtracking if all possibilities starting from index i have been exhausted. **Return Result**: The dfs function will return False if none of the iterations and recursive calls result in a valid split sequence.

On the other hand, the function will exit with a True at the first instance of finding a valid sequence of substrings.

0). This algorithm effectively explores all potential ways to split the string through DFS while avoiding a full search when the first

valid sequence of splits or exhausts all possibilities and determines it's not possible.

valid solution is found. It also uses recursion and backtracking systematically to traverse the decision tree until it either finds a

split into a sequence of descending numbers that are each 1 less than the previous, specifically into ["3", "2", "1"]. Initial Call: • The splitString method calls dfs(0, -1, 0) to start the recursive search.

Let's walk through a small example to illustrate the solution approach with the string s = "321". This string should be possible to

 We start by picking s[0:1] which is "3", and convert it to number y = 3. • Since the last number x was -1 (indicating no previous number), "3" is accepted as the first valid substring. • We then call dfs(1, 3, 1) to proceed to the next character. Action in dfs - Second Iteration from index 1:

## • Now x = 2, and i = 2.

Base Case Reached:

class Solution:

Java

class Solution {

**Example Walkthrough** 

Action in dfs - First Iteration from index 0:

• The loop runs from i = 0 to the end of the string.

• Now x = 3 and we start from i = 1.

We try to form substrings starting from the first character.

We pick s[1:2] which is "2", and convert it to number y = 2.

• We then call dfs(2, 2, 2) since we have found two substrings "3" and "2".

• We now call dfs(3, 1, 3) because we have found a third substring "1".

• The count of valid splits k is 3 (greater than 1), so we return True.

• The call dfs(3, 1, 3) has i == len(s), meaning we have reached the end of the string.

• We check if y is exactly 1 less than  $\times$  (3 - 1 = 2), and it is.

Action in dfs - Third Iteration from index 2:

• We pick s[2:3] which is "1", and convert it to number y = 1. Again, we check if y is exactly 1 less than x (2 - 1 = 1), which it is.

• The valid sequence resulting from the above splits is ["3", "2", "1"], and therefore the initial call of splitString returns True.

• We have successfully split the string into a descending sequence where each number is exactly 1 less than the previous substring.

This simple example illustrates the DFS-based recursive process of finding valid substrings and how backtracking occurs if a

substring does not fit the required criteria. The process continues until the entire string is either successfully split or determined

Result:

Solution Implementation **Python** 

# Helper function to perform depth-first search

for j in range(current index, len(s)):

to be unsplittable according to the conditions.

def split string(self, s: str) -> bool:

current number = 0

return False

return dfs(0, -1, 0)

return True

// i: starting index for the next number

return splitCount > 1;

return true;

return false;

bool splitString(string s) {

if (start == s.size()) {

return count > 1;

long long currentValue = 0;

// The current value being formed.

for (int j = start; j < s.size(); ++j) {</pre>

// Forming the current value.

**if** (i == str.length()) {

// lastNumber: the last number found in the split

for (int j = i; j < str.length(); ++j) {</pre>

// If no valid split is found, return false

// and `count` is the number of splits made so far.

// Iterate through the string starting from `start`.

depthFirstSearch(i + 1, currentValue, count + 1)) {

// If no valid split could be found, return false.

// can be split according to the problem requirements.

// If dfs call is successful, return true indicating that the string

// Function to check if the input string can be split into consecutive decreasing integers.

console.log(splitString(s)); // It will log the output of the splitString function.

# Base condition: Check if we have reached the end of the string

# Valid split if at least one number has been split before

# Check if it's the start (-1) or if the current number is exactly

# one less than the previous number, and then recursively call dfs

# If the recursive call returns True, propagate it upwards

// Call the depthFirstSearch starting from index 0 with an undefined previous value (-1) and no splits (0 count).

// splitCount: number of valid splits found so far

private boolean dfs(int i, long lastNumber, int splitCount) {

// Iterate over the string to form the next number

// Append the next digit to the current number

dfs(i + 1, currentNumber, splitCount + 1)) {

// If a valid split is found, return true

def dfs(current index, previous number, split count): # Base condition: Check if we have reached the end of the string if current index == len(s): # Valid split if at least one number has been split before return split\_count > 1

# Variable to store the current number being formed

# If we cannot find a valid split, return False

# Initiate the depth-first search with initial values

private String str; // Variable to hold the input string

# Accumulate the digits to form the current number

# Check if it's the start (-1) or if the current number is exactly

# one less than the previous number, and then recursively call dfs

# If the recursive call returns True, propagate it upwards

// Method to check if we can make a split where each number is one less than the previous

// If we have reached the end of the string, check if we made more than one valid split

// Check if the current number is 1 less than the last number or if this is the first number

// Continue the DFS search from the next index, with the current number as the last number

long currentNumber = 0; // Variable to store the current number being formed

// Function to check if the input string can be split into consecutive decreasing integers.

// Lambda function to perform Depth First Search (DFS) to find if the string can be split.

// Base case: if we have reached the end of the string and made at least one split.

function<bool(int, long long, int)> dfs = [&](int start, long long prevValue, int count) -> bool {

// `start` is the index to start searching from, `prevValue` is the previous value found,

currentNumber = currentNumber \* 10 + (str.charAt(j) - '0');

if ((lastNumber == -1 || lastNumber - currentNumber == 1) &&

current\_number = current\_number \* 10 + int(s[j])

public boolean splitString(String s) { this.str = s; // Assign the given string to the class variable // Begin depth-first search from the start of the string, with initial value -1 and no splits return dfs(0, -1, 0); // Recursive DFS method to check for valid splits

if (previous number ==-1 or previous number - current number ==1) and dfs(j + 1, current\_number, split\_count + 1):

C++

public:

class Solution {

```
currentValue = currentValue * 10 + (s[i] - '0');
                // To prevent overflow, break if the number is too large.
                if (currentValue > 1e10) {
                    break;
                // If this is the first number or if the current number is exactly
                // one less than the previous, recursively call dfs.
                if ((prevValue == -1 || prevValue - currentValue == 1) &&
                    dfs(i + 1, currentValue, count + 1)) {
                    // If dfs call is successful, return true indicating that the string
                    // can be split according to the problem requirements.
                    return true;
            // If no valid split could be found, return false.
            return false;
        };
        // Call the dfs starting from index 0 with an undefined previous value (-1) and no splits (0 \ count) .
        return dfs(0, -1, 0);
};
TypeScript
// Type alias for the depth-first search function type.
type DFSFunction = (start: number, prevValue: number, count: number) => boolean;
// Function to perform the Depth First Search (DFS) to find if the string can be split
// into consecutive decreasing integers.
const depthFirstSearch: DFSFunction = (start, prevValue, count) => {
    // Base case: if we have reached the end of the string and made at least one split.
    if (start === s.length) {
        return count > 1;
    // The current value being formed.
    let currentValue = 0;
    // Iterate through the string starting from start.
    for (let i = start; i < s.length; j++) {</pre>
        // Forming the current value.
        currentValue = currentValue * 10 + (s[i].charCodeAt(0) - '0'.charCodeAt(0));
        // To prevent overflow, break if the number is too large.
        if (currentValue > 1e10) {
            break;
        // If this is the first number or if the current number is exactly
        // one less than the previous, recursively call depthFirstSearch.
        if ((prevValue === -1 || prevValue - currentValue === 1) &&
```

```
return False
# Initiate the depth-first search with initial values
return dfs(0, -1, 0)
```

Time and Space Complexity

return true;

const splitString = (s: string): boolean => {

let s: string = "1234"; // Example input string

def split string(self, s: str) -> bool:

current number = 0

if current index == len(s):

return True

return split\_count > 1

# Helper function to perform depth-first search

for j in range(current index, len(s)):

def dfs(current index. previous number. split count):

# Variable to store the current number being formed

# If we cannot find a valid split, return False

# Accumulate the digits to form the current number

uses a depth-first search (DFS) approach to recursively try out different splits.

n choices at most once for each of the 2<sup>n</sup> potential splits of the string.

current\_number = current\_number \* 10 + int(s[j])

return depthFirstSearch(0, -1, 0);

return false;

// Example usage

class Solution:

**}**;

**}**;

**Time Complexity** The DFS function, dfs, is called recursively, at each step it tries to construct a new number, y, by adding digits one by one from the input string s. For each number y formed, a check is made whether the previous number x is exactly one more than y (i.e., x

- y == 1). The function dfs is potentially called once for each new number y formed, which can happen for each starting

The given Python code is a solution to the problem of splitting a string into a sequence of decreasing consecutive numbers. It

if (previous number ==-1 or previous number - current number ==1) and dfs(j + 1, current\_number, split\_count + 1):

### position i in the string s. The worst case time complexity is $0(n * 2^n)$ , where n is the length of the string. This is because at each step, we have two choices: either include the current digit in the current number y or start a new number beginning with the current digit. We make

**Space Complexity** The space complexity of the code is O(n). This comes from two factors:

for a split starting at each digit of the string. 2. The variable y in the inner loop, which is re-constructed for each recursive call but doesn't add to the space complexity as it is just an integer and not a recursive structure.

1. The recursive call stack, which at maximum depth could be O(n), as in the worst case the recursion could go as deep as the length of the string

Hence, the space requirement grows linearly with the input size, dominated by the depth of the recursive calls.