7. Reverse Integer

Medium Math

Problem Description

The task is to take a signed 32-bit integer `x` and reverse the order of its digits. For example, if the input is `123`, the output should be `321`. If the input is `-123`, the output should be `-321`. The tricky part comes with the boundaries of a 32-bit signed integer, which ranges from `-2^31` to `2^31 - 1`. If reversing the digits of `x` would cause the number to fall outside this range, the function should return `0` instead. This means we need to be careful with overflow—an issue that occurs when the reversed integer is too large or too small to be represented by a 32-bit signed integer.

Intuition

To solve this problem, we first set up two boundaries, `mi` and `mx`, which represent the minimum and maximum values of a 32-bit signed integer, respectively. These values are `-2^31` and `2^31 - 1`.

We want to build the reversed number digit by digit. We can isolate the last digit of `x` by taking `x % 10` (the remainder when `x` is divided by 10). This last digit, referred to as `y` in our code, is the next digit to be placed in the reversed number.

However, we need to be careful not to cause an overflow when we add this new digit to the reversed number. Before we add `y` to the reversed number `ans`, we check if adding the digit would cause an overflow. To do this, we check if `ans` is either less than `mi / 10 + 1` or greater than `mx / 10`. If it's outside this range, we return `0`.

If it's safe to add the digit, we proceed. We add the digit to `ans` by multiplying `ans` by 10 (which "shifts" the current digits to the left) and then adding `y`. This process effectively reverses the digits of `x`.

For the next iteration, we need to remove the last digit from `x`. We do this by subtracting `y` from `x` and then dividing by 10.

We repeat this process until `x` has no more digits left. The result is a reversed number that fits within the 32-bit signed integer range, or `0` if an overflow would have occurred.

The time complexity is `O(\log |x|)` because the process continues for as many digits as `x` has, and the space complexity is `O(1)` as there is a constant amount of memory being used regardless of the size of `x`.

Solution Approach

The implementation uses a straightforward algorithm that iterates through the digits of the input number `x` and constructs the reversed number without using additional data structures or complex patterns. Let's detail the steps using the provided Reference Solution Approach:

1. **Initialization**: We start by setting the initial reversed number `ans` to 0. We also define the minimum and maximum values `mi` and `mx` for a 32-bit signed integer, which are `-2^31` and `2^31 - 1`.
2. **Reversing Digits**: The `while` loop runs as long as there are digits left in `x`. Within the loop, we take the following steps:
   - Isolate the last digit `y` of `x` by computing `x % 10`.
   - If `x` is negative and `y` is positive, adjust `y` by subtracting 10 to make it negative.
3. **Checking for Overflow**: Before appending `y` to `ans`, we must confirm that `ans * 10 + y` will not exceed the boundaries set by `mi` and `mx`. To avoid overflow, we check:
   - If `ans` is less than `mi/10 + 1` or greater than `mx/10`, we return `0` immediately, as adding another digit would exceed the 32-bit signed integer limits.
4. **Building the Reversed Number**: If it is safe to proceed, we multiply `ans` by 10 (which shifts the reversed number one place to the left) and add `y` to `ans`. This action reverses `y` from its position in `x` to its new reversed position in `ans`.
5. **Updating the Original Number `x`**: We update `x` by removing its last digit. This is done by subtracting `y` from `x` and then dividing by 10.
6. **Completion**: The loop repeats this process, accumulating the reversed number in `ans` until all digits are processed.

The core of this approach is predicated on the mathematical guarantees regarding integer division and modulus operations in Python. The guard checks for overflow by considering both scale (multiplication by 10) and addition (adding the digit) separately and only proceeds if the operation stays within bounds.

By following the constraints of a 32-bit signed integer at every step and efficiently using arithmetic operations, the `reverse` function achieves the reversal of digits robustly and efficiently.

Example Walkthrough

To illustrate the solution approach, let's take `x = 1469` as our example.

1. **Initialization**:
   - `ans` is initialized to 0.
   - `mi` is set to `-2^31` (-2147483648).
   - `mx` is set to `2^31 - 1` (2147483647).
2. **Reversing Digits**: Begin while loop with `x = 1469`.
   - Isolate the last digit `y` by computing `1469 % 10 = 9`.
   - `x` is positive so we keep `y = 9`.
3. **Checking for Overflow**:
   - `ans` is currently 0, which is greater than `mi/10 + 1` (-214748364) and less than `mx/10` (214748364), so continue without returning `0`.
4. **Building the Reversed Number**:
   - We multiply `ans` by 10, which is still 0, and add `y` to get `ans = 9`.
5. **Updating the Original Number `x`**:
   - Update `x` to remove the last digit: `x = (1469 - 9) / 10` which simplifies to `x = 146`.

Next iteration of the loop with `x = 146`:

- Isolate `y = 146 % 10 = 6`.
- Check for overflow: `ans = 9` is still within bounds.
- Update `ans`: `ans = 9 * 10 + 6 = 96`.
- Update `x`: `x = (146 - 6) / 10` which simplifies to `x = 14`.

Next iteration with `x = 14`:

- Isolate `y = 14 % 10 = 4`.
- Check for overflow: `ans = 96` is still within bounds.
- Update `ans`: `ans = 96 * 10 + 4 = 964`.
- Update `x`: `x = (14 - 4) / 10` which simplifies to `x = 1`.

Final iteration with `x = 1`:

- Isolate `y = 1 % 10 = 1`.
- Check for overflow: `ans = 964` is still within bounds.
- Update `ans`: `ans = 964 * 10 + 1 = 9641`.
- Update `x`: `x = (1 - 1) / 10` which simplifies to `x = 0`.

Now `x = 0`, the while loop terminates, and the reversed number `ans = 9641` is returned. There were no issues with overflow throughout the process, so the result `9641` is the correct reversed integer for our example of `x = 1469`.

This process demonstrates that by evaluating the overflow conditions before each digit is added to the reversed number, and by building the reversed number step by step, it's possible to safely reverse the digits of `x` without using extra space or complex data structures. Additionally, due to the use of modulo and division operations, the solution efficiently handles the reversal process for each digit.

Python Solution

```python
class Solution:
    def reverse(self, x: int) -> int:
        # This variable will hold the reversed number
        reversed_number = 0

        # These define the range of acceptable 32-bit signed integer values
        min_int, max_int = -2**31, 2**31 - 1

        while x:
            # Check if the reversed_number will overflow when multiplied by 10
            if reversed_number < min_int // 10 + 1 or reversed_number > max_int // 10:
                # Return 0 on overflow as per problem constraints
                return 0

            # Extract the least significant digit of the current number
            digit = x % 10

            # Adjustments for negative numbers when the extracted digit is non-zero
            if x < 0 and digit > 0:
                digit -= 10

            # Shift reversed_number digits to the left and add the new digit
            reversed_number = reversed_number * 10 + digit

            # Remove the least significant digit from x
            x = (x - digit) // 10

        # Return the reversed number within the 32-bit signed integer range
        return reversed_number

```

Java Solution

```java
class Solution {
    public int reverse(int x) {
        // Initialize answer to hold the reversed number
        int reversedNumber = 0;

        // Loop until x becomes 0
        while (x != 0) {
            // Check for overflow/underflow condition, return 0 if violated
            // Integer.MIN_VALUE is -2^31 and Integer.MAX_VALUE is 2^31 - 1
            if (reversedNumber < Integer.MIN_VALUE / 10 || reversedNumber > Integer.MAX_VALUE / 10) {
                return 0;
            }

            // Add the last digit of x to reversedNumber
            reversedNumber = reversedNumber * 10 + x % 10;

            // Remove the last digit from x
            x /= 10;
        }

        // Return the reversed number
        return reversedNumber;
    }
}

```

C++ Solution

```cpp
#include <climits> // For INT_MIN and INT_MAX

class Solution {
public:
    int reverse(int x) {
        int reversedNumber = 0;

        // Loop until all digits are processed
        while (x != 0) {
            // Check if multiplying by 10 will cause overflow
            if (reversedNumber < INT_MIN / 10 || reversedNumber > INT_MAX / 10) {
                return 0;  // Return 0 if overflow would occur
            }

            // Pop the last digit from 'x' using modulus and add it to 'reversedNumber'
            reversedNumber = reversedNumber * 10 + x % 10;

            // Remove the last digit from 'x' by dividing it by 10
            x /= 10;
        }

        return reversedNumber;  // Return the reversed number
    }
};

```

Typescript Solution

```typescript
/**
 * Reverse an integer.
 * @param {number} x - The integer to be reversed.
 * @return {number} - The reversed integer, or 0 if the reversed integer overflows 32-bit signed integer range.
 */
const reverseInteger = (x: number): number => {
    // Define the minimum and maximum values for 32-bit signed integer.
    const INT_MIN: number = -(2 ** 31);
    const INT_MAX: number = 2 ** 31 - 1;

    let reversed: number = 0;

    while (x !== 0) {
        // Check for potential overflow by comparing with pre-divided limits.
        if (reversed < Math.floor(INT_MIN / 10) || reversed > Math.floor(INT_MAX / 10)) {
            return 0;
        }

        // Perform the reverse by multiplying the current reversed by 10 and adding the last digit of x.
        reversed = reversed * 10 + (x % 10);

        // Floor division by 10 to get the next digit (in TypeScript `~~` is replaced by Math.trunc).
        // Since x can be negative, we use trunc instead of floor to correctly handle negative numbers.
        x = Math.trunc(x / 10);
    }

    return reversed;
};

```

Time and Space Complexity

Time Complexity

The time complexity of the given code is dependent on the number of digits in the integer `x`. Since we are handling the integer digit by digit, the number of operations is linearly proportional to the number of digits. If the integer `x` has `n` digits, then the time complexity is `O(n)`.

Space Complexity

The space complexity of the provided code is `O(1)`. This is because we are only using a fixed amount of additional space (`ans`, `mi`, `mx`, `y`, and a few variables for control flow) regardless of the input size.