3043. Find the Length of the Longest Common Prefix

Medium Hash Table Array String

Problem Description

common prefix between any two integers, where one integer comes from arr1 and the other comes from arr2. A "prefix" in this context means the digits at the start of an integer, up until any length of that integer. For instance, 12 is a prefix

In this problem, we are given two arrays of positive integers, arr1 and arr2. We are tasked with finding the length of the longest

of 123. If a prefix is found in both integers from arr1 and arr2, it is considered a common prefix. Note that if the prefix doesn't exist in both integers, it's not common.

Our goal is to determine the size of the utmost common prefix discovered among all possible pairs made by taking one integer from each of the two arrays. If no common prefixes are found, then the answer we return should be 0.

Intuition

To solve this problem, we apply the concept of storing the prefixes in a hash table. The solution revolves around two main steps:

Firstly, we enumerate all possible prefixes from the integers in arr1 and store each one in a set (which acts as our hash table).

Secondly, for every integer in arr2, we start with the integer itself and keep dividing it by 10 (which effectively removes the least

significant digit) to generate its prefixes. Each generated prefix is checked against our hash table to see if it's a common prefix.

interested in the longest common prefix – so we start with checking the entire number and then reduce it from the right. While we are checking for prefixes, we remember the length of the longest common prefix we've encountered by using a variable.

The logic behind gradually reducing the integer starting from its original value down to its high-order digits is that we are

This variable is updated anytime we find a longer common prefix than the current record. This approach allows us to efficiently find the longest common prefix among all pairs by leveraging the quick lookup capabilities

Solution Approach

The solution can be broken down into two main phases, using the Python programming language:

• For each number x in arr1, we start a loop where we keep adding x to our set s. Right after adding, we divide x by 10 using an integer division (x //= 10). This effectively removes the least significant digit of x, leaving us with a new prefix. We repeat this until x becomes 0,

of a hash table.

which means we've added all possible prefixes of x to the set.

Prefix Generation for arr1:

Finding Longest Common Prefix with arr2:

longest common prefix in an efficient manner.

 \circ Add 12345 to s, now s = {12345}.

• We initialize an answer variable ans to 0. This will hold the length of the longest common prefix we find. • We iterate through each number x of arr2. Similar to the first phase, we keep dividing x by 10 to get its prefixes. ∘ For each generated prefix of x, we check if this prefix exists in the set s. If it does, it means we have found a common prefix. We then compare the length of this prefix (by converting it to a string and getting the length) to our variable ans. If the length is greater, we update

ans to the new length. As soon as we find a matching prefix, we break out of the loop for the current number x because we're looking for the longest common

prefix, and any further divisions will only produce shorter prefixes which are not of interest. In summary, the solution utilizes a hash set for quick existence checks of prefixes, a loop to generate all possible prefixes of

• We initiate by creating an empty set s, which will be used to store the prefixes of the integers present in arr1.

- integers in arr1, and another loop to check prefixes of integers in arr2 against this set. We also use the max function to update the length of the longest common prefix as we iterate through the elements of arr2. This approach guarantees that we find the
- **Example Walkthrough** Let's illustrate the solution with a small example using the following arrays:

• arr1 = [12345, 678] • arr2 = [123, 458]Prefix Generation for arr1 • Initialize set s = {}.

∘ Iterate by dividing by 10 and adding to s: 1234, 123, 12, 1, until x becomes 0. s becomes {12345, 1234, 123, 12, 1}. • For 678 in arr1:

• For 12345 in arr1:

```
\circ Add 678 to s, now s = {12345, 1234, 123, 12, 1, 678}.
```

- Continue dividing by 10 and adding to s: 67, 6. \circ Final s = {12345, 1234, 123, 12, 1, 678, 67, 6}.
- Finding Longest Common Prefix with arr2

• Initialize ans = 0.

458 is not in s.

• For 123 in arr2: • 123 is in set s. A common prefix of length 3 is found.

Update ans = 3 because it is greater than the current ans value.

 No need to keep dividing 123 by 10 because we found the longest prefix with it. • For 458 in arr2:

Divide by 10, get 45. 45 not in s.

1. Prefix Generation for `arr1`:

- Initialize `ans = 0`.

Divide by 10, get 4. 4 not in s. All prefixes checked, move to next integer.

- Now, $s = \{12345, 1234, 123, 12, 1, 678, 67, 6\}$

- Create an empty set `s = {}`. Add all prefixes of `12345` to `s`. We add `12345`, `1234`, `123`, `12`, and `1`. Add all prefixes of `678` to `s`. We add `678`, `67`, and `6`.

2. Finding Longest Common Prefix with `arr2`:

Solution Implementation

Look at `123` from `arr2`. We find `123` in `s`. Update `ans` to `3` since `123` has length `3`.

- Look at `458` from `arr2`. We do not find `458`, `45`, or `4` in `s`. We do not update `ans`.

The longest common prefix found is of length `3`, making our final answer `3`.

Initialize 'max prefix_length' to keep track of the longest prefix length

Again, generate prefixes by continuously dividing the number by 10

If the current prefix is in 'prefixes' set, it is a common prefix

Loop through each number in the second list (nums2)

Update 'num' by removing the last digit

After iterating through both arrays, the longest common prefix we found was of length 3. Hence, our answer is 3.

Given `arr1 = [12345, 678]` and `arr2 = [123, 458]`, let's walk through the solution step by step.

```
Python
class Solution:
    def longestCommonPrefix(self, nums1: List[int], nums2: List[int]) -> int:
        # Initialize an empty set to store unique prefixes
        prefixes = set()
        # Loop through each number in the first list (nums1)
        for num in nums1:
            # Generate prefixes by continuously dividing the number by 10
            while num:
                # Add the current prefix to the set
                prefixes.add(num)
                # Update 'num' by removing the last digit
```

Update 'max prefix length' if this prefix is longer than the current max max prefix length = max(max prefix length, len(str(num))) # Break the loop as we've found the longest prefix for this number break

num //= 10

num //= 10

return longestPrefixLength;

unordered_set<int> prefixes;

for (: num: num /= 10) {

int longestPrefixLength = 0;

for (; num; num /= 10) {

for (int num : nums2) {

prefixes.insert(num);

// Iterating through each number in nums2

// Return the length of the longest common prefix.

Initialize an empty set to store unique prefixes

Loop through each number in the first list (nums1)

Add the current prefix to the set

Loop through each number in the second list (nums2)

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Again, generate prefixes by continuously dividing the number by 10

If the current prefix is in 'prefixes' set, it is a common prefix

max prefix length = max(max prefix length, len(str(num)))

Update 'max prefix length' if this prefix is longer than the current max

Break the loop as we've found the longest prefix for this number

return longestPrefix;

prefixes = set()

for num in nums1:

while num:

prefixes.add(num)

if num in prefixes:

break

num //= 10

max_prefix_length = 0

for num in nums2:

while num:

class Solution:

for (int num : nums1) {

// Required to use std::vector

// Required for log10 function

int longestCommonPrefix(vector<int>& nums1, vector<int>& nums2) {

// Insert all prefixes of all numbers in nums1 into the set

// Create an unordered set to store all the prefixes of elements in nums1

// Break down the number into prefixes and check against our set

// If the current prefix is in the set, it's a common prefix

// Initialize 'longestPrefixLength' to store the length of the longest prefix found

// A prefix here is obtained by repeatedly dividing the number by 10

#include <unordered set> // Required for std::unordered set

using namespace std; // To refrain from using 'std::' prefix

if num in prefixes:

max_prefix_length = 0

for num in nums2:

while num:

```
# Return the length of the longest common prefix
        return max_prefix_length
Java
class Solution {
    // Method to find the longest common prefix length represented in two arrays
    public int longestCommonPrefix(int[] arr1, int[] arr2) {
        // Create a HashSet to store unique prefixes
        Set<Integer> prefixes = new HashSet<>();
        // Iterate through every number in the first array
        for (int num : arr1) {
            // Loop to add all prefixes of the current number to the set
            for (int prefix = num; prefix > 0; prefix /= 10) {
                prefixes.add(prefix);
        // Initialize the variable 'longestPrefixLength' to store the length of the longest common prefix
        int longestPrefixLength = 0;
        // Iterate through every number in the second array
        for (int num : arr2) {
            // Loop to check if any prefix of the current number exists in the set
            for (int prefix = num; prefix > 0; prefix /= 10) {
                // If a common prefix is found
                if (prefixes.contains(prefix)) {
                    // Update 'longestPrefixLength' with the length of the current prefix if it's the longest found so far
                    longestPrefixLength = Math.max(longestPrefixLength, String.valueOf(prefix).length());
                    // Break the loop since we only need the longest common prefix
                    break;
        // Return the length of the longest common prefix
```

C++

public:

#include <vector>

#include <cmath>

class Solution {

```
if (prefixes.count(num)) {
                    // Logarithm base 10 of num gives us the number of digits - 1
                    longestPrefixLength = max(longestPrefixLength, (int)log10(num) + 1);
                    // once the longest prefix for current num is found, no need to look for shorter ones
                    break;
        // Return the length of the longest common prefix
        return longestPrefixLength;
};
TypeScript
function longestCommonPrefix(arr1: number[], arr2: number[]): number {
    // Create a new Set to store unique digits from arr1.
    const uniqueDigits: Set<number> = new Set<number>();
    // Extract digits from each number in arrl and add them to the Set.
    for (let num of arr1) {
        // Iterate through digits of 'num' by continuously dividing by 10.
       while (num) {
            // Add the rightmost digit to the Set and truncate 'num' to remove that digit.
            uniqueDigits.add(num % 10);
            num = Math.floor(num / 10);
    // Initialize 'longestPrefix' to 0, which will keep track of
    // the length of the longest prefix found that matches a digit in 'uniqueDigits'.
    let longestPrefix: number = 0;
    // Look for common digits in arr2 that exist in the 'uniqueDigits' Set.
    for (let num of arr2) {
       // Again, iterate through digits of 'num'.
       while (num) {
            // If the rightmost digit is found in 'uniqueDigits', we have a common prefix.
            if (uniqueDigits.has(num % 10)) {
                // Update 'longestPrefix' with the maximum of current value
                // and the current number 'num' length.
                longestPrefix = Math.max(longestPrefix, Math.floor(Math.log10(num)) + 1);
            // Truncate 'num' to remove its rightmost digit.
            num = Math.floor(num / 10);
```

num //= 10 # Return the length of the longest common prefix return max_prefix_length

Time and Space Complexity

The time complexity of the algorithm is 0(m * log M + n * log N), where m is the length of arr1, n is the length of arr2, M is the maximum value in arr1, and N is the maximum value in arr2. This is because for each element in arr1 and arr2, the algorithm potentially divides the element by 10 until it reaches 0, which occurs in 0(log M) or 0(log N) time for each element respectively.

The space complexity of the algorithm is 0(m * log M) because the algorithm stores each prefix of the numbers from arr1 in a set. The number of prefixes for a number with a logarithmic number of digits is also logarithmic, so each number contributes O(log M) space, and there are m numbers in arr1.