

# Problem Description

The problem is to find out how many different permutations of the numbers 1 to n can be constructed in such a way that all prime numbers are at prime indices, with indices considered to be 1-indexed (meaning they start from 1, not 0). Note that prime numbers are integers greater than 1 that have no divisors other than 1 and themselves.

to consider two separate conditions for the permutation to meet the prime placement criteria: (1) The numbers themselves must be prime, and (2) Their positions in the permutation must also be prime index locations. Given a number n, there will be a specific count of prime numbers within the range from 1 to n, and those prime numbers must be

It is important to understand that indices in a list and the actual numbers placed at those indices can be prime. Therefore, we need

arranged in the prime indices. All the non-prime numbers, therefore, will go into the remaining positions. Since the answer could potentially be a very large number, the problem asks us to return the result modulo 10^9 + 7, a common

technique in algorithms to avoid integer overflow and to keep the numbers within a manageable range.

# The solution is based on combinatorics. Specifically, we can break down the problem into two separate tasks:

1. Counting the prime numbers from 1 to n. This tells us how many numbers need to be placed in prime index positions. 2. Calculating permutations of the prime numbers in the prime indices and the permutations of the non-prime numbers in the non-

- prime indices. The Sieve of Eratosthenes algorithm, which is an efficient way to find all primes smaller than a given number, is a good approach to
- get the count of prime numbers. With this algorithm, we iterate over each number from 2 to n and mark multiples of each number (which cannot be prime) as non-prime. The numbers that remain unmarked at the end of this process are the prime numbers.

Once we count the number of primes within the range of 1 to n (let's say there are cnt primes), we then must calculate the total number of permutations of these cnt prime numbers, which is simply the factorial of cnt (factorial(cnt)). Simultaneously, we must also calculate the permutations for the remaining n - cnt non-prime numbers, which can be arranged in

any order in the remaining positions. The total permutations for these non-prime numbers are factorial(n - cnt). The total number of prime arrangements is then the product of these two permutations: factorial(cnt) for the prime numbers and

The final step is to return this product modulo 1009 + 7 to ensure the output fits within the expected range of values as required by the problem.

The solution code encloses the prime-counting logic within a function count, then calculates the factorials, multiplies them, and

Solution Approach

The implementation of the solution can be broken down into two main parts: counting prime numbers and computing factorials for the permutations.

## The count function initiates a list primes of boolean values, initially set to True, representing the numbers from 0 to n. The boolean values will represent if a number is prime (True) or not (False).

factorial(n - cnt) for the non-prime numbers.

applies the modulo operation to return the answer.

1. Counting Prime Numbers using Sieve of Eratosthenes:

 It then iterates over the list starting from the first prime number 2. For each number 1, if 1 is marked as True (indicating it is still considered prime), we: Increment the cnt which counts the number of primes.

- 2. Calculating Factorials: The number of permissible arrangements of prime numbers is calculated by taking the factorial of the count of prime
- Similarly, for non-prime numbers, we use the factorial of the difference between the total count n and the count of prime numbers (factorial(n - cnt)).

specified by the problem and to avoid integer overflow.

numbers (factorial(cnt)).

3. Calculating the Answer:

Proceed to mark all multiples of i as False since they are not primes.

After the loop completes, cnt holds the number of prime numbers between 1 and n.

ensures that the algorithm runs in a relatively fast time frame appropriate for the size of input n.

Now, apply Step 1: Counting Prime Numbers using the Sieve of Eratosthenes:

Multiply the results of these factorials to find the total permutations: 6 \* 6 = 36.

So, for n = 6, there are 36 permutations where prime numbers occupy the prime indices.

Take the result modulo 10^9 + 7 gives 36 mod (10^9 + 7) = 36.

def count\_primes(n: int) -> int:

for i in range(2, n + 1):

prime\_count = count\_primes(n)

return arrangements % (10\*\*9 + 7)

for (int i = 2;  $i \le n$ ; ++i) {

1 using ll = long long; // Alias for long long type

int numPrimeArrangements(int n) {

2 const int MOD = 1e9 + 7; // Constants should be in uppercase

result = (result \* BigInt(i)) % MOD;

\* Counts the number of prime numbers up to `n`.

\* @returns The count of prime numbers up to `n`.

isPrime[j] = false;

function countPrimes(n: number): number {

let primeCount: number = 0;

if (isPrime[i]) {

return primeCount;

for (let i = 2; i <= n; ++i) {

primeCount++;

Time and Space Complexity

\* @param n The number up to which to count primes.

const isPrime: boolean[] = new Array<boolean>(n + 1).fill(true);

for (let j = i \* 2; j <= n; j += i) {

complexity: O(n \* log(log(n)) + n + n), simplifying to O(n \* log(log(n))).

 The answer to how many unique permutations of numbers are available such that primes are at prime indices is given by the product of these two factorials: factorial(cnt) \* factorial(n - cnt).

Since the product of these factorials can be very large, the final answer is taken modulo 10^9 + 7 to fit within the bounds

Throughout the implementation, we see the use of basic data structures like lists (primes list for the sieve algorithm) and the use of the range function to traverse numbers and multiples. Due to Python's built-in modulo operation and factorial computation, no additional custom data structures or algorithms are required for computing factorials and their modulo.

In programming terms, this implementation leverages memoization implicitly by pre-computing the factorials of the numbers from 1

to n only once, thereby avoiding redundant calculations. This approach, combined with the efficiency of the Sieve of Eratosthenes,

**Example Walkthrough** Let us illustrate the solution approach by walking through a small example. Consider the number n = 6. We want to find out how

many permutations of the numbers 1 to 6 can be formed where prime numbers occupy prime indices (1-indexed). Firstly, we need to identify the prime numbers from 1 to 6 and also the prime indices. The prime numbers will be 2, 3, and 5, and the prime indices are 2, 3, and 5 since 1 is not considered prime.

Initialize the primes list to track prime numbers up to n (ignoring 0 and 1). This looks like [True, True, True, True, True, True].

• Starting from 2, mark non-prime multiples as False. The updated primes list becomes [True, True, True, True, False, True].

# • Count the number of True values excluding the first position which corresponds to 1. We have three primes 2, 3, 5 (which is cnt

prime numbers.

**Python Solution** 

= 3).Next, for Step 2: Calculating Factorials:

Finally, in Step 3: Calculating the Answer:

• Calculate the factorial of the count of prime numbers (which is 3): factorial(3) = 3! = 6 to account for permutations among

class Solution: def numPrimeArrangements(self, n: int) -> int: from math import factorial

is\_prime = [True] \* (n + 1) # Initialize a list to track prime numbers

# Calculate the number of arrangements as the factorial of the prime count,

if (isPrime[i]) { // Check if the number is marked as a prime

++count; // Increment the count of prime numbers

int primeCount = countPrimes(n); // Count the number of primes up to n

// Mark all multiples of i as non-prime

for (int j = i + i;  $j \le n$ ; j += i) {

return count; // Return the count of prime numbers

isPrime[j] = false;

# multiplied by the factorial of the count of non-prime numbers

# Return the number of arrangements modulo (10\*\*9 + 7)

arrangements = factorial(prime\_count) \* factorial(n - prime\_count)

Count the number of prime numbers less than or equal to n using the Sieve of Eratosthenes algorithm.

Also calculate the factorial of non-prime numbers (n - cnt): factorial(6 - 3) = factorial(3) = 3! = 6.

#### 13 if is\_prime[i]: # If i is a prime number 14 count += 1 15 for j in range(i \* 2, n + 1, i): 16 is\_prime[j] = False # Mark multiples of i as not prime 17 return count

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count = 0

```
Java Solution
   class Solution {
         private static final int MOD = (int) 1e9 + 7; // Constant for the modulo operation
         // Counts the number of prime arrangements possible up to n
  4
         public int numPrimeArrangements(int n) {
             int primeCount = countPrimes(n); // Counts the number of prime numbers up to n
  6
             long arrangements = factorial(primeCount) * factorial(n - primeCount);
             // Computes the arrangements as prime! * (n - prime)!
  8
             return (int) (arrangements % MOD); // Returns the result modulo MOD
  9
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 12
         // Calculates the factorial of a number using the modulo operation
 13
         private long factorial(int n) {
 14
             long result = 1;
 15
             for (int i = 2; i <= n; ++i) {
 16
                 result = (result * i) % MOD; // Calculates factorial with modulo at each step
 17
 18
             return result;
 19
 21
         // Counts the number of prime numbers up to n
 22
         private int countPrimes(int n) {
 23
             int count = 0; // Initialize count of primes to 0
             boolean[] isPrime = new boolean[n + 1]; // Create an array to mark non-prime numbers
 24
 25
             Arrays.fill(isPrime, true); // Assume all numbers are prime initially
```

### // Solution class containing methods for prime arrangements calculation class Solution { public: // Calculates the number of prime arrangements for a given number n

C++ Solution

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// Calculate the factorial of prime count and non-prime count respectively,
           // then multiply them together modulo MOD
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           ll arrangements = factorial(primeCount) * factorial(n - primeCount) % MOD;
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           return static_cast<int>(arrangements);
14
       // Function to calculate factorial of a number n modulo MOD
       ll factorial(int n) {
17
18
           ll result = 1;
           for (int i = 2; i \le n; ++i) {
19
               result = (result * i) % MOD;
           return result;
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25
       // Function to count the number of prime numbers up to n
26
       int countPrimes(int n) {
27
           vector<bool> isPrime(n + 1, true); // Create a sieve initialized to true
           int primeCount = 0;
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           for (int i = 2; i \le n; ++i) {
               if (isPrime[i]) {
30
                   ++primeCount; // Increment count if i is a prime
31
                   // Mark all multiples of i as non-prime
33
                   for (int j = i * 2; j <= n; j += i) {
                       isPrime[j] = false;
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           return primeCount;
39
40 };
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Typescript Solution
  // Define type alias for bigint
2 type BigIntAlias = bigint;
   // Constant for modular arithmetic
   const MOD: BigIntAlias = BigInt(1e9 + 7);
   * Calculates the factorial of a number `n` modulo `MOD`
    * @param n The number to calculate the factorial of.
    * @returns The factorial of `n` modulo `MOD`.
11
    */
   function factorial(n: number): BigIntAlias {
       let result: BigIntAlias = BigInt(1);
       for (let i = 2; i <= n; ++i) {
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```

### const primeCount: number = countPrimes(n); 45 const arrangements: BigIntAlias = (factorial(primeCount) \* factorial(n - primeCount)) % MOD; 46 return Number(arrangements);

**Time Complexity** 

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/\*\*

\*/

return result;

38 39 \* Calculates the number of prime arrangements for a given number `n`. \* @param n The number to calculate prime arrangements for. \* @returns The number of prime arrangements for `n`. 43 \*/ function numPrimeArrangements(n: number): number {

prime numbers up to n, has a time complexity of O(n \* log(log(n))). The factorial function (which is not shown here, but its complexity can be derived from typical implementations) typically has a

The time complexity of the count function is O(n \* log(log(n))). This is because the sieve of Eratosthenes, which is used to find all

prime count cnt and one for the non-prime count n - cnt. Therefore, the overall time complexity of the code is dominated by the sieve in the count function, which gives us the total time

time complexity of O(n). Given that the maximum value for factorial calculation is n, two such calculations are performed - one for the

# For space complexity, the count function allocates an array primes of size n + 1, which leads to O(n) space complexity. The space

**Space Complexity** 

requirement for the calculation of the factorial depends on the implementation, but typically, it can be calculated in O(1) space.

Therefore, the overall space complexity of the code is O(n) for the sieve of Eratosthenes array storage.