2300. Successful Pairs of Spells and Potions

Medium Array Two Pointers Binary Search Sorting

In this problem, you are provided with two arrays called spells and potions. The spells array contains elements representing the strength of each spell, and similarly, the potions array contains elements representing the strength of each potion. The task is to determine for each spell, how many potions can successfully combine with it to achieve or exceed a given success threshold. A spell and potion pair is deemed successful if the product (multiplication) of their strengths is at least equal to the success value provided. You should return an array where each value corresponds to the number of successful potion combinations for each spell in the

Leetcode Link

multiply each spell with each potion: For the first spell (strength 10), it pairs successfully with two potions (potions of strength 10 and 8) since both products are at

To clarify with an example, suppose spells = [10, 20], potions = [5, 8, 10], and success = 100. To find successful pairs, you

- least 100. For the second spell (strength 20), it pairs successfully with all three potions, as all products will be 100 or greater. The result array will be [2, 3] as the first spell has 2 successful combinations, and the second has 3.
- Intuition

The goal is to find an efficient way to pair spells with potions such that the product of their strengths is at least the success target. A

brute-force approach would be to try all possible pairs, but this can be inefficient, especially with large arrays.

spells array.

Problem Description

Instead, by first sorting the potions array, we can take advantage of binary search to quickly find the number of successful potions for each spell. Binary search operates by repeatedly dividing the search interval in half, which is much faster than scanning every element.

Once the potions array is sorted, for each spell, we want to find the position where the potion strength is just enough or more to meet the success target when paired with the spell. We need to find the smallest potion that, when multiplied by the spell's strength, equals or exceeds the success threshold. All potions to the right in the sorted array are guaranteed to also form successful pairs

The solution leverages the bisect_left function from Python's bisect module, which uses binary search to find the insertion point for a given element to maintain sorted order. In this context, it finds the first index in potions where the potion strength is sufficient for the success criteria with a given spell. We then subtract this index from the total number of potions (m) to get the count of successful potions for that spell.

For each spell in spells, we apply this logic and generate the final result array. Solution Approach

The solution implements a binary search mechanism to optimize the process of finding successful spell-potion pairs. Here's a step-

1. Sorting the Potions: Begin by sorting the potions array. This is crucial for binary search to work since binary search requires a sorted array to function correctly. Sorting is done in ascending order.

at the minimum, when multiplied by the spell's strength, is equal to success.

and index is the position returned by bisect_left.

Here's the critical code snippet with explanations for clarity:

2 m = len(potions) # Storing the length of the 'potions' array.

3 # Step 2, 3, and 4: List comprehension iterating over 'spells'.

4 return [m - bisect_left(potions, success / v) for v in spells]

1 potions.sort() # Step 1: Sorting the 'potions' array.

spells representing the n in the complexity.

wise explanation of the approach used in the solution:

2. Using Binary Search: For each spell's strength value, we apply a binary search to determine the number of potions that, when

because they are equal or stronger.

potion index that can achieve the required target when combined with the current spell. 3. Calculating Successful Pairs:

multiplied by this spell's strength, will at least be equal to the success value. This is done by finding the leftmost (smallest)

 We use the bisect_left function which finds the index at which we could insert the value success / spell_strength into potions to maintain sorted order. • The value we are searching for is calculated by success / spell_strength because we want to find the potion strength that,

Since potions is sorted, every potion at and beyond the index returned by bisect_left would result in a successful pair

- when combined with the current spell. 4. Storing Results: ∘ For each spell, we calculate m - index to find out the number of successful potions, where m is the total number of potions
- successful pairs due to the sorted property of the array.
- 5. Generating the Output: Finally, we return a list comprehension that iterates over every spell in spells, applies the above logic using bisect_left, and calculates the number of successful pairs for that spell.

This operation gives us the count because all elements from the position index to the end of the potions array will be

Through sorting and binary searching, this algorithm achieves a complexity of O(n log m), where n is the number of spells, and m is

the number of potions. Sorting takes $0 (m \log m)$, and then each binary search operation takes $0 (\log m)$, with the linear iteration over

Consider we have spells = [15, 10], potions = [1, 5, 20, 8], and success = 120. We aim to determine for each spell, how many

Example Walkthrough

1. Sorting the Potions: The first step is sorting the potions array. Before: [1, 5, 20, 8] After sorting: [1, 5, 8, 20]

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2. Using Binary Search: We then apply binary search to find the count of potions that can pair with a spell to achieve the success.
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2.

3. Calculating Successful Pairs:

• For spell 15: Using bisect_left, we find the index where 8 could fit in the sorted potions array [1, 5, 8, 20], which is index

For spell 10: We need a potion of at least 120 / 10 = 12.

Let's walk through a small example based on the solution approach given above.

potions achieve or exceed the success threshold when multiplied by the spell.

o For spell 10: Using bisect_left, we find the index where 12 could fit, which is after 8 and before 20. Hence, the index would be 3. 4. Storing Results:

5. Generating the Output: According to the steps described in the solution approach, we create a list with the counts of successful

∘ For spell 15: m - index is 4 - 2 = 2. So, there are 2 successful potion combinations for the spell 15.

∘ For spell 10: m - index is 4 - 3 = 1. Therefore, there is 1 successful potion combination for the spell 10.

For spell 15: We search for the smallest potion that when multiplied is at least 120. That potion should be 120 / 15 = 8.

combinations for each spell. The result for our example will be [2, 1]. For better understanding, here's how the binary search part of the code would operate in a step-by-step manner:

index = 2, total potions m = 4, successful combinations: 4 - 2 = 2.

Spell is 15: bisect_left([1, 5, 8, 20], 120/15)

index = 3, successful combinations: 4 - 3 = 1.

It finds the position of 8 because 8 * 15 is exactly 120.

 Spell is 10: bisect_left([1, 5, 8, 20], 120/10) It passes over 8 (since 8 * 10 is less than 120) and settles before 20.

Thus, the final returned value for this particular example with the spells and potions given would be [2, 1]. Each spell has a

Python Solution

corresponding count of potions that, when multiplied, meet or exceed the success threshold.

15 Returns: 16 List containing the count of successful potion combinations for each spell. 17 18 # Sort the potions list in ascending order for binary searching

in the sorted 'potions' list that meets or exceeds the 'success' threshold when combined with the spell.

Using list comprehension to generate the list of counts for successful combinations

int mid = (left + right) / 2; // Calculate the mid index.

if ((long) spells[i] * potions[mid] >= successThreshold) {

return successfulPairs; // Return the array of successful pair counts.

// Check if the current spell and potion at mid index is a successful pair.

right = mid; // If successful, narrow the search to the left half.

// of potions minus the number of potions that did not meet the success threshold.

// The number of successful pairs for the current spell is the total number

vector<int> successfulPairs(vector<int>& spells, vector<int>& potions, long long success) {

left = mid + 1; // If not successful, narrow the search to the right half.

For each spell strength 'spell_strength' in 'spells', we find the index of the first potion

def successful_pairs(self, spells: List[int], potions: List[int], success: int) -> List[int]: Find the number of potions for each spell that when combined result in a product equal to or greater than the success threshold. 8 9 10 Args: spells: List of integers representing the strength of spells. 11 12 potions: List of integers representing the volume of potions. 13 success: An integer representing the minimum success threshold for a spell-potion combination. 14

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33 }

1 from bisect import bisect_left

potions.sort()

Length of the potions list

num_potions = len(potions)

} else {

successfulPairs[i] = nPotions - left;

// Function to find the number of successful pairs

// Sort the potions array for binary search

sort(potions.begin(), potions.end());

// Iterate over each spell

from typing import List

class Solution:

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           # The count of successful pairings for each spell is then the total number of potions
28
           # minus this index, giving the number of potions that meet or exceed the product threshold.
29
           return [num_potions - bisect_left(potions, success / spell_strength) for spell_strength in spells]
30
Java Solution
   import java.util.Arrays; // Ensure to import the necessary Arrays class for sorting.
   class Solution {
       // This function determines the number of successful pairs of spells and potions.
       public int[] successfulPairs(int[] spells, int[] potions, long successThreshold) {
           Arrays.sort(potions); // Sort the potions array for binary search.
           int nSpells = spells.length; // Number of spells.
           int nPotions = potions.length; // Number of potions.
           int[] successfulPairs = new int[nSpells]; // Array to store the answer.
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12
           // Iterate over each spell.
           for (int i = 0; i < nSpells; ++i) {</pre>
13
               int left = 0, right = nPotions; // Set search boundaries for binary search.
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15
               // Binary search to find the first potion that results in a successful pair.
16
               while (left < right) {</pre>
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// ans will hold the answer to the problem 9 vector<int> ans; 10 // m is the size of the potions array 12 int m = potions.size();

C++ Solution

1 class Solution {

2 public:

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           for (int spell : spells) {
               // Find the first potion that, when multiplied with the spell,
16
               // is greater than or equal to the success threshold.
17
               // We use 1.0 to cast success to double for correct division.
19
               int index = lower_bound(potions.begin(), potions.end(), (success + spell - 1) / spell) - potions.begin();
20
               // Push the number of successful pairs into the answer vector.
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22
               // Subtract the found index from the size to determine the number of successful potions.
23
               ans.push_back(m - index);
24
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26
           // Return the final answer
27
           return ans;
28
29 };
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Typescript Solution
   function successfulPairs(spells: number[], potions: number[], successThreshold: number): number[] {
       // Sort the potions array in ascending order
       potions.sort((a, b) => a - b);
       // The number of potions available
       const potionCount = potions.length;
 6
       // Initialize an array to store the number of successful pairs for each spell
 8
       const successfulPairsCount: number[] = [];
 9
10
11
       // Iterate over each spell
12
       for (const spell of spells) {
13
           // Set the initial bounds for binary search
           let left = 0;
14
           let right = potionCount;
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16
           // Perform binary search to find the index where
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35 successfulPairsCount.push(potionCount - left); 36 37 38

Time and Space Complexity

while (left < right) {

} else {

right = mid;

left = mid + 1;

// Calculate the middle index

const mid = Math.floor((left + right) / 2);

if (spell * potions[mid] >= successThreshold) {

Time Complexity The time complexity of the given code snippet consists of two parts: sorting the potions list and performing binary searches using

the bisect_left function.

// Return the array containing the counts of successful spell-potion pairs return successfulPairsCount; 39

// If it does, we need to find if there's a smaller index that also satisfies the condition

- 1. Sorting: The potions.sort() operation has a time complexity of O(n log n) where n is the number of potions. 2. Binary Search: The binary search inside the list comprehension using bisect_left is performed for each spell. If there are s spells, and the binary search has a time complexity of O(log n) for each spell, the entire list comprehension has a time
- complexity of O(s log n). Therefore, combining both operations, the total time complexity is $0(n \log n + s \log n)$.

Space Complexity

1. Sorting Space: The sorting is done in place, which does not require additional space, so it's 0(1). 2. Output List: There is a list comprehension that generates a list of the same length as the number of spells. Therefore, it has a

space complexity of O(s) where s is the number of spells.

The space complexity is determined by the additional space required beyond the input data.

// spell * potionValue is greater than or equal to success threshold

// Check if the current combination meets the success requirement

// If not, move the left pointer to narrow the search

// The number of potions that meet the success criteria with the current spell

// is the length of the potions array minus the left bound found by binary search

Given that O(s) is the larger term between O(1) and O(s), the overall space complexity is O(s).