

Problem Description

In this problem, you're given an integer array called nums, with a length of n. You're also given two integers: indexDifference and valueDifference. Your task is to find two indices i and j such that both i and j are within the range from 0 to n - 1 and they meet the following criteria:

2. The absolute difference between the values at nums[i] and nums[j] must be at least valueDifference.

The absolute difference between i and j must be at least indexDifference,

In terms of the outcome, you need to return an array called answer. This array should consist of the indices [i, j] if such a pair of

indices exists. If there are multiple valid pairs, you can return any one of them. If no valid pairs are found, then return [-1, -1]. An interesting point to note is that according to the problem statement, i and j can be the same index, which implies that indexDifference could potentially be 0. Intuition

The primary intuition behind the solution is the usage of a sliding window technique, combined with the maintenance of the minimum and maximum values within the window. The sliding window is defined by two pointers, 1 and 1, that maintain a distance apart

conditions. At the outset, i starts at the position indexDifference, and j starts at 0. By doing this, the gap between i and j reflects the indexDifference requirement of our problem.

specified by indexDifference. The pointers are used as markers to capture a subarray within nums to check against our two

We maintain two variables, mi and mx, to keep track of the indices where the minimum and maximum values are found within our sliding window that ends at the current j index. While sliding i further along the array, we update mi and mx to account for the entry of new values into the window and the exit of old values.

When updating mi and mx, if nums[j] is less than nums[mi], we reassign mi to j, because we have found a new minimum. Conversely, if nums[j] is greater than nums[mx], we reassign mx to j due to identifying a new maximum.

After every movement of i and update of mi and mx, we check our two conditions against nums[i] (the current value at i). If the difference between nums[i] and the value at nums[mi] is greater than or equal to valueDifference, we have found a valid pair [mi, i]. Alternatively, if the difference between the maximum value (nums [mx]) and nums [i] is also greater than or equal to

valueDifference, then [mx, i] is a valid pair. In this situation, we immediately return the pair as it meets our requirements.

If we reach the end of the array without finding a pair that satisfies both conditions, we conclude that no such pair exists, and we return the default output [-1, -1]. The approach's efficiency comes from the fact that it avoids checking every possible pair of indices, which would otherwise lead to a

Solution Approach

The solution uses a sliding window approach, which involves moving a fixed-size window across the array to examine sub-sections one at a time. This allows for checking the conditions over smaller segments in a single pass through the array, making the solution more efficient than a brute force approach that would involve examining all possible index pairs.

To implement this technique, the algorithm maintains two pointers: 1 and 1. These pointers define the bounds of the sliding window.

The pointer i starts at the index equal to indexDifference while j starts at 0, thus immediately satisfying the condition abs(i - j) >= indexDifference because i - j is initialized to indexDifference.

mx = 1

return [mx, i]

proportional to the input size.

condition:

less efficient solution with a higher time complexity.

As the algorithm iterates over the array, starting from i = indexDifference, it keeps track of the indices of the minimum and the maximum values found so far to the left of j. These indices are stored in variables mi and mx, respectively. for i in range(indexDifference, len(nums)): j = i - indexDifference

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Within the loop, we first check whether the current element at index j changes the minimum or maximum:
  if nums[j] < nums[mi]:</pre>
      mi = i
  if nums[j] > nums[mx]:
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After updating mi and mx, we check if nums [i] differs enough from the minimum or maximum value to satisfy the valueDifference

If either of these checks succeeds, the function immediately returns the corresponding pair of indices, as they meet both prescribed

conditions. If the function reaches the end of the array without returning, this means no valid pairs were found, and thus it returns

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1 if nums[i] - nums[mi] >= valueDifference:
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[-1, -1].
In terms of data structures, no additional structures are needed beyond the use of a few variables to keep track of the indices and
values encountered. This algorithm is space-efficient because it operates directly on the input array without requiring extra space
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5], with indexDifference = 3 and valueDifference = 3.

if nums[mx] - nums[i] >= valueDifference:

with every other element, thereby reducing the time complexity from O(n^2) to O(n), where n is the length of the input array. **Example Walkthrough**

Let's walk through an example to illustrate the solution approach described above. Consider the integer array nums = [1, 2, 3, 4,

Our task is to find indices i and j such that abs(i - j) >= indexDifference and abs(nums[i] - nums[j]) >= valueDifference.

The choice of a sliding window and keeping track of minimum and maximum values eliminates the need to compare every element

According to the given solution approach, we initialize the sliding window by setting i to indexDifference and j to 0. This immediately satisfies the first condition as the difference between the indices i = 3 and j = 0 is 3, which is equal to

Here's how we proceed step by step:

class Solution:

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Java Solution

class Solution {

indexDifference.

1. On the first iteration where i = 3: ○ We have j = 0

• Since nums [3] - nums [0] fulfills the valueDifference condition, we return [0, 3] as the indices that satisfy both conditions.

We initiate mi to j since there are no previous values to compare with, and similarly, mx is also initiated to j

In this example, we directly found a pair that met both conditions, and thus we would return [0, 3]. However, if we needed to

2. Increment i to the next position and decrement j to keep the window size constant while satisfying the indexDifference. If

continue the iteration, we would:

4. If we find a pair, we return it. If not, we continue iterating until i reaches the end of the array.

def findIndices(self, nums: List[int], idx_diff: int, val_diff: int) -> List[int]:

3. Check if nums [i] differs enough from nums [mi] or nums [mx] as explained previously.

If no valid pairs are found by the end of the array, we return [-1, -1] as specified.

Initialize min and max index pointers

if nums[compare_idx] < nums[min_idx]:</pre>

if nums[compare_idx] > nums[max_idx]:

return [max_idx, current_idx]

min_idx = compare_idx

min_idx = max_idx = 0

nums [j] changes the minimum or maximum within the new window, update mi or mx accordingly.

○ We see that nums[3] - nums[0] = 4 - 1 = 3, which is equal to valueDifference

o The elements under consideration are [nums [0], nums [3]] i.e., [1, 4]

We then maintain mi and mx to keep track of the minimum and maximum values within the window.

Python Solution from typing import List

for current_idx in range(idx_diff, len(nums)): 10 # Compute the comparison index that matches the index difference 11 12 compare_idx = current_idx - idx_diff 13

// Method to find indices in an array such that the difference between their values is at least a given value and their positions

// Check if the difference between the current value and the minimum value found so far is at least valueDifference

// Check if the difference between the maximum value found so far and the current value is at least valueDifference

// Loop through the array starting from the index equal to the indexDifference to the end of the array

max_idx = compare_idx 18 19 20 # If the value difference requirement is met with the minimum, return the indices if nums[current_idx] - nums[min_idx] >= val_diff: 21 22 return [min_idx, current_idx]

If the value difference requirement is met with the maximum, return the indices

Check and update the min and max indices based on the values at compare_idx

Traverse the array, starting from the index that enables the required index difference

If the required value difference isn't found, return [-1, -1] as per problem statement return [-1, -1]

if nums[max_idx] - nums[current_idx] >= val_diff:

int currentIndex = i - indexDifference; // Calculate the index to compare with 10 11 // Update the minimum value index if a new minimum is found 12 if (nums[currentIndex] < nums[minIndex]) {</pre> minIndex = currentIndex; 14

return new int[] {minIndex, i}; // Return the indices if condition is met

return new int[] {maxIndex, i}; // Return the indices if condition is met

// Update the maximum value index if a new maximum is found

public int[] findIndices(int[] nums, int indexDifference, int valueDifference) {

int minIndex = 0; // Initializing the minimum value index

int maxIndex = 0; // Initializing the maximum value index

for (int i = indexDifference; i < nums.length; ++i) {</pre>

if (nums[currentIndex] > nums[maxIndex]) {

if (nums[i] - nums[minIndex] >= valueDifference) {

if (nums[maxIndex] - nums[i] >= valueDifference) {

maxIndex = currentIndex;

31 32 33 // Return [-1, -1] if no such pair of indices is found 34 return new int[] {-1, -1}; 35 36 }

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C++ Solution
 1 #include <vector>
   class Solution {
   public:
       // Method to find the two indices in the array nums such that the difference
       // between their values is at least valueDifference and their index difference is exactly indexDifference
       // Args:
            nums: The input vector of integers
            indexDiff: The required difference between the indices of the two elements
            valueDiff: The minimum required value difference between the two elements
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       // Returns:
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            A vector with two elements: the indices of the elements in nums that satisfy the above criteria
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            If no such pair exists, returns {-1, -1}
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       std::vector<int> findIndices(std::vector<int>& nums, int indexDiff, int valueDiff) {
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            int minIndex = 0, maxIndex = 0; // Initialized to store the index of minimum and maximum values seen so far
            for (int i = indexDiff; i < nums.size(); ++i) {</pre>
16
                int j = i - indexDiff; // Calculate the corresponding index
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                if (nums[j] < nums[minIndex]) {</pre>
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                    minIndex = j; // Update minIndex if a new minimum is found
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               if (nums[j] > nums[maxIndex]) {
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                    maxIndex = j; // Update maxIndex if a new maximum is found
23
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               // Check if the difference between the current value and the minimum value seen so far is at least valueDiff
25
               if (nums[i] - nums[minIndex] >= valueDiff) {
26
                    return {minIndex, i}; // Pair found, return indices
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28
               // Check if the difference between the maximum value seen so far and the current value is at least valueDiff
29
               if (nums[maxIndex] - nums[i] >= valueDiff) {
30
                    return {maxIndex, i}; // Pair found, return indices
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           return \{-1, -1\}; // If no pair found, return \{-1, -1\}
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35 };
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6 // valueDifference: The minimum required difference between the values at the indices
 7 // Returns an array with two numbers representing the indices, or [-1, -1] if no such pair exists
   function findIndices(nums: number[], indexDifference: number, valueDifference: number): number[] {
       // Initialize the indices for the minimum value (minIndex) and maximum value (maxIndex) found.
       let minIndex = 0;
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       let maxIndex = 0;
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       // Iterate over the array, starting from the element at the indexDifference.
       for (let currentIndex = indexDifference; currentIndex < nums.length; currentIndex++) {</pre>
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           // Calculate the index of the element we are comparing against,
           // which is indexDifference behind the current index.
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           const compareIndex = currentIndex - indexDifference;
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           // Update minIndex if the current compareIndex points to a new minimum value
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           if (nums[compareIndex] < nums[minIndex]) {</pre>
               minIndex = compareIndex;
21
22
23
24
           // Update maxIndex if the current compareIndex points to a new maximum value
25
           if (nums[compareIndex] > nums[maxIndex]) {
               maxIndex = compareIndex;
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           // Check if the difference between the current element and the minimum value is large enough.
30
           if (nums[currentIndex] - nums[minIndex] >= valueDifference) {
31
                return [minIndex, currentIndex]; // Return the indices if the condition is met.
32
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           // Check if the difference between the maximum value and the current element is large enough.
35
           if (nums[maxIndex] - nums[currentIndex] >= valueDifference) {
                return [maxIndex, currentIndex]; // Return the indices if the condition is met.
36
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       // If no suitable pair of indices is found, return [-1, -1].
40
       return [-1, -1];
42 }
Time and Space Complexity
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1 // This function finds two indices such that the difference of the elements at these indices

2 // is at least the given valueDifference and the indices are separated by at most the given

5 // indexDifference: The maximum allowed difference between the indices

Typescript Solution

4 // nums: The array of numbers to search within

3 // indexDifference.

Time Complexity The time complexity of the provided code is O(n). This is achieved by iterating over the array once from indexDifference to the length of the array len(nums). Only constant time checks and updates are performed within the loop, leading to a linear time

complexity relative to the array's size.

The space complexity of the code is 0(1). No additional space proportional to the input size is used. Only a fixed number of variables

Space Complexity

mi, mx, and j are used, which occupy constant space regardless of the input array size.