2900. Longest Unequal Adjacent Groups Subsequence I

Dynamic Programming



Problem Description

The goal is to find the longest subsequence from an array of indices [0, 1, ..., n - 1] such that for any two consecutive indices in the subsequence i_j and i_{j+1} , the elements in the binary array groups at those indices are not the same, i.e., groups $[i_j]$!= groups [i_{j + 1}]. Each index in the subsequence corresponds to a word in the words array. The task is to return an array of words that represents this longest subsequence.

Leetcode Link

A subsequence is a sequence that can be derived from another sequence by deleting some or no elements without changing the order of the remaining elements. Importantly, the words in the words array may have different lengths, which doesn't impact the selection of the subsequence.

Intuition

The approach to finding the longest subsequence where consecutive elements in groups are different is a greedy one. This means we can make local, optimal choices at each step without needing to consider the rest of the array.

For every element at index i in groups, we have two scenarios - either i is the first index (i.e., i == 0), or groups [i] is different from the previous element groups [i - 1]. If either of these conditions is met, we can include the corresponding word words [i] in our subsequence.

words have their corresponding groups elements equal, effectively giving us the longest subsequence by the definition provided.

By iterating over the entire groups array and including words that meet the criteria, we ensure that no two consecutive selected

long as they meet the aforementioned condition. Since the condition only depends on the current and previous groups elements, we only need a simple iteration to build our solution without needing to backtrack or look ahead.

The intuition comes from the fact that to maximize the length of the subsequence, we want to include as many words as possible as

The provided Python solution uses a list comprehension to create and return the subsequence of words, which is essentially a

Solution Approach

single-pass greedy algorithm. The algorithm iterates through the groups array and applies a selection criteria to each element to determine if the corresponding element from the words array should be included in our final output or not.

Algorithm:

built-in enumerate function to obtain each element in groups and its index simultaneously. 2. In this comprehension, for every element x and index i, the following condition is checked: i == 0 or x != groups[i - 1]. This

1. Initialize a list comprehension that will evaluate each element x in groups along with its corresponding index i. It uses Python's

- condition says that the first element (i == 0) should always be included, and then every subsequent element should be included only if it is different than the one preceding it (x != groups[i - 1]). This ensures that no two adjacent elements in the subsequence have the same groups value. 3. If the condition is true for a given i, we select words[i] for inclusion in the final output.
- 4. Once the list comprehension finishes iterating over all elements in groups, it will have produced a list of words that constitutes
- the longest subsequence satisfying the problem's constraints. 5. The final step is to return this list of words.
- **Data Structures:**

Patterns used:

We utilize Python's list data structure to store the words and groups.

word) based only on the current and immediate previous elements from groups, which guarantees the finding of the global optimum (the longest subsequence under the given conditions). • List comprehension, a concise way to create lists, is used for its readability and efficiency in selecting the appropriate words.

• The solution applies a greedy approach to the problem: at each step, it makes a local optimum choice (whether to include a

This simple yet effective approach leverages the characteristics of the problem's constraints to arrive at an optimal solution without needing a complex algorithm.

Example Walkthrough

words arrays:

• groups = [1, 0, 0, 1, 0, 1]

Let's walk through a small example to illustrate the solution approach described above. Suppose we have the following groups and

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in the groups array. Let's apply the algorithm step by step:
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consecutive elements must be different.

• ["apple", "banana", "cherry", "mango", "peach"]

longest_subsequence_words = []

words = ["apple", "banana", "grape", "cherry", "mango", "peach"]

1. Start iterating through the groups array, comparing the element at index i with the one at index i - 1. 2. The first element in groups is 1 (at index 0). Since i == 0, we don't have a previous element to compare with, so we include

Our task is to find the longest subsequence such that no two consecutive indices in the subsequence correspond to equal elements

- "apple" from the words array in our subsequence.
- subsequence. 4. At index 2, groups has another 0. This time, groups [2] == groups [1], so "grape" does not get included in our subsequence since

3. The next element in groups is 0 (at index 1). Since groups[1] != groups[0], we include "banana" from the words array in our

5. Now at index 3, we have a 1 in groups. Since groups [3] != groups [2], "cherry" gets included in our subsequence.

6. Moving to index 4, the element in groups is 0. Because groups [4] != groups [3], we include "mango" in our subsequence.

- 7. Lastly, at index 5, groups contains a 1. As groups [5] != groups [4], we include "peach" in our subsequence. Following the steps of our algorithm, the final subsequence of words is:
- Thus, by iterating through each element of the groups array and checking our defined condition, we successfully construct the

greedy approach to solve the problem, as we made local optimal selections to achieve a global optimal solution.

Initialize an empty list to store the words in the longest subsequence

longest_subsequence_words.append(words[i])

// Initialize an ArrayList to store the resulting words.

List<String> result = new ArrayList<>();

// Iterate through each group by index.

for (int index = 0; index < n; ++index) {</pre>

// then it is a part of the longest non-repeating subsequence.

// Return the answer vector containing the words in the longest non-repeating subsequence.

if (index == 0 || groups[index] != groups[index - 1]) {

// Add the current word to the answer vector.

answer.emplace_back(words[index]);

Iterate through each index and corresponding group number in the groups list

public List<String> getWordsInLongestSubsequence(int n, String[] words, int[] groups) {

Import the List type from typing module for type hints from typing import List class Solution: def getWordsInLongestSubsequence(self, n: int, words: List[str], groups: List[int]) -> List[str]:

longest subsequence of words without having identical consecutive elements from groups. This illustrates the effectiveness of the

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for i, group_number in enumerate(groups):
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               # Check if it's the first word or if the current group number is different from the previous one
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               if i == 0 or group_number != groups[i - 1]:
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                   # If yes, append the corresponding word to the longest_subsequence_words list
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Python Solution

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# Return the list of words that form the longest subsequence
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           return longest_subsequence_words
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Java Solution
 1 import java.util.ArrayList;
   import java.util.List;
   class Solution {
        * Finds the words in the longest subsequence with alternating groups.
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                        the number of words.
        * @param n
        * @param words an array of words.
10
        * @param groups an array of group identifiers corresponding to each word.
11
        * @return a list of words in the longest subsequence with alternating groups.
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// Iterate over the words to find the longest subsequence. 18 for (int i = 0; i < n; ++i) { 19 20

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// Add the first word and any word that starts a new group (compared to the previous word).
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               if (i == 0 || groups[i] != groups[i - 1]) {
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                   result.add(words[i]);
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           // Return the list of words in the longest subsequence.
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           return result;
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28 }
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C++ Solution
1 #include <vector>
2 #include <string>
   class Solution {
   public:
       // Function that generates a vector of strings, which consists of the words
       // in the longest non-repeating subsequence based on the given groups.
       // Parameters:
       // n - the number of elements in the words and groups vectors.
       // words - a vector of strings representing the words.
       // groups - a vector of integers, where each integer corresponds to the group of the word at the same index.
       std::vector<std::string> getWordsInLongestSubsequence(int n, std::vector<std::string>& words, std::vector<int>& groups) {
13
           // Answer vector to store the resulting words.
           std::vector<std::string> answer;
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// If we are at the first word, or the current word belongs to a different group than the previous one,

27 28 }; 29

return answer;

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Typescript Solution
   function getWordsInLongestSubsequence(totalWords: number, wordsArray: string[], wordGroups: number[]): string[] {
       // Initialize an array to hold the resulting longest subsequence of words.
       const longestSubsequence: string[] = [];
       // Iterate through the array of words to identify the longest subsequence.
       for (let index = 0; index < totalWords; ++index) {</pre>
           // If we are at the first word or the current word's group is different from the previous word's group,
           // it is a part of the longest subsequence, so we add it to the result array.
           if (index === 0 || wordGroups[index] !== wordGroups[index - 1]) {
               longestSubsequence.push(wordsArray[index]);
       // Return the longest subsequence found.
14
       return longestSubsequence;
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Time and Space Complexity

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once and performs a constant time check for each element. The space complexity is O(k), where k is the number of unique subsequences identified, which in the worst case could be equal to n. This would happen if no two consecutive elements in groups are the same, resulting in each word from words being added to the

output list. Thus, the space used for the output list is proportional to the number of selected words.

The time complexity of the code is O(n), where n is the length of the list groups. This is because the code iterates over the list groups