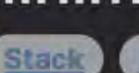
## 921. Minimum Add to Make Parentheses Valid







Problem Description

String Medium Stack Greedy

The problem presents a situation where we have a string s which consists only of the parentheses characters "(" and ")". A valid parentheses string is defined as such if it meets one of the following criteria:

Leetcode Link

- It's an empty string.
- It can be separated into two substrings A and B such that both A and B are themselves valid parentheses strings.
- It can be enclosed within a pair of parentheses meaning if A is a valid string, then (A) is also valid.

The task at hand is to determine the fewest number of moves required to make the initial string s into a valid parentheses string. Each move consists of inserting exactly one parenthesis character (either "(" or ")") at any position within the string.

### Intuition

The intuition behind the solution stems from the understanding of how a valid parentheses string is structured. Fundamentally, for every opening parenthesis "(", there should be a corresponding closing parenthesis ")" to make it valid. If we traverse the string from left to right, at any point, the number of closing parentheses should not exceed the number of opening parentheses.

When encountering an opening parenthesis, it suggests the start of a potentially valid substring, so we increment a counter. Upon finding a closing parenthesis, if there is a previously unmatched opening parenthesis (our counter is greater than zero), we can pair this closing parenthesis with that opening parenthesis, decrementing the counter. If the counter is zero (indicating no unmatched opening parentheses), we require an additional opening parenthesis to match the current closing parenthesis, thus incrementing the answer-keeping track of moves needed.

matching closing parentheses. So, we add the number of remaining unmatched opening parentheses to the answer. By accumulating the moves required to insert missing opening or closing parentheses, we can calculate the minimum number of

Finally, after running through the string, if there's any unmatched opening parenthesis remaining (counter is not zero), those need

moves to make the string valid.

# The algorithm implemented in the solution is relatively straightforward and efficient, using a single pass through the string, and it

Solution Approach

relies on the use of counters to track the state of the parentheses. No additional data structures are needed, which makes the space complexity constant, 0(1), as we only use a couple of integer variables to keep track of counts. The algorithm can be described as follows:

cnt will keep track of the balance of opening and closing parentheses as we iterate through the string. 2. Iterate through each character c of the given string s:

1. Initialize two variables, and and cont to zero. Here, and will hold the total number of moves required to make the string valid, and

- If c is an opening parenthesis "(", increment the cnt counter, since we have an additional unmatched opening parenthesis.
  - If c is a closing parenthesis ")": If cnt is greater than zero, it means there is a preceding unmatched opening parenthesis which can be paired with this
  - closing parenthesis, so we decrement cnt.
    - If cnt is zero, it means there are no unmatched opening parentheses to pair with this closing parenthesis, therefore, we need to add an opening parenthesis, thus we increment the ans counter.
- 3. After the iteration is complete, if there is a non-zero cnt, this means there are cnt number of unmatched opening parentheses remaining. These will all need matching closing parentheses, so we add cnt to ans.
- 4. The final value of ans is the minimum number of moves required to make the string valid.

This approach works well as it leverages the inherent structure of valid parentheses strings. Since we only look at the balance

between opening and closing brackets without necessarily considering their exact positions, the order in which we would insert the

additional parentheses doesn't change the number of moves we need to make, making this approach both simple and effective. The time complexity of the solution is linear, O(n), where n is the length of the input string because we go through the string exactly once. Example Walkthrough

### 1. We initialize two variables, ans = 0 and cnt = 0.

2. We start iterating through the string s from left to right.

Let's consider a small example here with the string s = ")) ((" to illustrate the solution approach."))

- We begin with the first character ):
- balance it out, so we increment ans to 1. Now, ans = 1 and cnt remains 0.

  - The second character is ): the same logic applies, so ans is incremented again. Now, ans = 2, and cnt is still 0. For the third character ), we repeat the same increment on ans. So, ans = 3, and cnt = 0.

• cnt is at 0, meaning there are no unmatched opening (for this). We would need one opening (before this one to

- We now encounter an opening parenthesis (. We increment cnt by 1 because we have one unmatched opening parenthesis, so cnt = 1.
- The fifth character is another (, so now cnt = 2.

three opening (at the beginning and three closing) at the end, resulting in a valid parentheses string ((())).

# If the character is an open parenthesis, increment the unmatched count.

# If it's a close parenthesis and there's an unmatched open, pair it and decrement.

unmatched (left, which means we need 3 matching closing) to make the string valid.

- The sixth character is yet again (, bringing cnt to 3. 3. We have finished iterating through the string. Now we must account for the unmatched opening parentheses. We have cnt = 3
- 4. We add cnt to ans. So, ans becomes 3 + 3 = 6.

def minAddToMakeValid(self, s: str) -> int:

unmatched\_open += 1

} else if (balanceCount > 0) {

additionsRequired++;

balanceCount--;

} else {

if char == '(':

# keep track of the unmatched parentheses.

# Initialize variables to count the necessary additions and

Python Solution

The minimum number of moves required to make the string )))((( valid by insertions is ans = 6. We can achieve this by inserting

#### additions\_needed = unmatched\_open = 0 # Iterate through each character in the string. for char in s:

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class Solution:

```
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               elif unmatched_open:
                   unmatched_open -= 1
14
               # Otherwise, if there is no unmatched open, we need an addition (an open parenthesis).
15
16
               else:
                   additions_needed += 1
17
18
           # Add any remaining unmatched open parentheses to the total additions needed.
19
           additions_needed += unmatched_open
20
21
22
           # Return the total number of additions needed to make the string valid.
23
           return additions_needed
24
Java Solution
   class Solution {
       public int minAddToMakeValid(String s) {
           int additionsRequired = 0; // Count of parentheses to add to make the string valid
           int balanceCount = 0;  // Keep track of the balance between opening and closing brackets
           // Loop through each character in the string
           for (char character : s.toCharArray()) {
               if (character == '(') {
                   // An opening parenthesis increments the balance count
9
                   balanceCount++;
```

// A closing parenthesis decrements the balance count if there are unmatched opening ones

// If there are no unmatched opening, we need an opening parenthesis

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18
19
           // Add the remaining unmatched opening parentheses to the count of required additions
20
           additionsRequired += balanceCount;
22
23
           // Return the total number of additions required to make the string valid
24
           return additionsRequired;
25
26 }
27
C++ Solution
 1 class Solution {
2 public:
       int minAddToMakeValid(string s) {
           int balance = 0; // This will keep track of the balance between '(' and ')'
           int additions = 0; // Counter for the required additions to make the string valid
           // Loop through each character in the string
           for (char c : s) {
               // If it's an opening bracket, increase the balance
9
               if (c == '(') {
                   balance++;
11
12
13
               // If it's a closing bracket, check if there is a matching opening bracket
               else {
14
                   // If there is a matching opening bracket, decrement the balance
15
                   if (balance > 0) {
16
                       balance--;
17
```

#### 26 // Add outstanding balance to the additions. These are unmatched opening brackets. 27 additions += balance;

else {

additions++;

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28
29
           // Return the total number of additions required to make the string valid
30
           return additions;
31
32 };
33
Typescript Solution
   // Function to calculate the minimum number of additions needed to make the parentheses string valid
   function minAddToMakeValid(s: string): number {
        let balance = 0; // This will keep track of the balance between '(' and ')'
       let additions = 0; // Counter for the additions required to make the string valid
       // Loop through each character in the string
       for (let i = 0; i < s.length; i++) {
           const c = s[i];
 8
 9
           // If it's an opening bracket, increase the balance
10
           if (c === '(') {
               balance++;
           } else { // Implicitly c is ')', as it's not '('
13
14
               // If there is a matching opening bracket, decrement the balance
               if (balance > 0) {
15
                   balance--;
16
               } else {
17
                   // If there is no matching opening bracket, increment additions
                   additions++;
19
20
21
22
23
24
       // Add any unmatched opening brackets to the additions
```

// If there is no matching opening bracket, increment additions

#### additions += balance; 26 // Return the total number of additions required to make the string valid 27

# Time and Space Complexity **Time Complexity**

return additions; 29 }

# through each character of the string exactly once, executing a constant number of operations for each character.

**Space Complexity** The space complexity of the provided code is 0(1), which signifies constant space. This is due to the fact that the space required

The time complexity of the provided code is O(n), where n is the length of the input string s. This is because the code iterates

does not scale with the size of the input string. The variables ans and cnt require a fixed amount of space regardless of the length of s.