



Problem Description

The problem presents a scenario in which we are given a string s consisting of six possible characters: the opening and closing brackets of three types—parentheses (), square brackets [], and curly braces {}. The challenge is to determine whether this string represents a sequence of brackets that is considered valid based on certain rules. A string of brackets is deemed valid if it satisfies the following conditions:

- Each opening bracket must be closed by a closing bracket of the same type. 2. Opening brackets must be closed in the correct order. That means no closing bracket should interrupt the corresponding pair of
- an opening bracket. 3. Each closing bracket must have an associated opening bracket of the same type before it.
- For example, a string "(())" is valid because each opening parenthesis (has a corresponding closing parenthesis) that occurs later

in the string, and they are properly nested and ordered.

Intuition

The intuition behind the solution utilizes a common data structure known as a stack, which operates on a last-in, first-out (LIFO) principle. This principle mimics the necessary behavior for tracking opening brackets and ensuring they are closed in the correct order. The steps followed in the solution are:

3. When an opening bracket is encountered ('(', '[', or '{'}), push it onto the stack. This represents waiting for a closing bracket

1. Initialize an empty stack to keep track of opening brackets.

(i.e., on top of the stack).

2. Traverse the input string character by character.

- to match. 4. When a closing bracket is encountered (')', ']', or '}'), check if it forms a pair with the last opening bracket added to the stack
- 5. If the stack is empty (no opening bracket available to match) or the closing bracket does not form a valid pair with the opening bracket on top of the stack, we know the string is invalid, and we immediately return False.
- 6. If we successfully traverse the entire input string and the stack is empty, meaning all opening brackets have been matched correctly, we return True. If the stack is not empty, it indicates there are unmatched opening brackets, and therefore, the string
- is invalid. The concise implementation of this algorithm ensures that both the correct type and order of brackets are validated for the string to be considered valid.

Solution Approach

The solution utilizes a stack data structure to track opening brackets and ensure they have the appropriate closing brackets. Let's

walk through the implementation step by step:

2. Pairs Set: We define a set d containing string representations of the valid bracket pairs: '()', '[]', and '{}'. This helps quickly check if an encountered closing bracket correctly pairs with the last opening bracket on the stack.

 Opening Bracket: If c is an opening bracket ('(', '[', or '{'}), it is pushed onto stk, waiting for the corresponding closing bracket.

3. Iterate Over the String: The algorithm iterates over each character c in the input string s.

- Closing Bracket: If c is a closing bracket (')', ']', or '}'):
- If the stack is not empty, the top element is popped. We concatenate it with c to check if they form a valid pair by

implies all opening brackets were properly matched and closed.

1. Initialize a Stack: A list named stk is created to serve as a stack.

checking against set d.

■ If either condition fails – the stack was empty or the concatenation of the popped element with c does not form a valid

• The stack is checked to ensure it's not empty, which would mean there's no opening bracket to match the closing one.

- pair the function immediately returns False because we've detected an invalid bracket sequence. 4. Final Stack Check: After processing all characters in the string, the algorithm checks if the stack is empty. An empty stack
- o If the stack is empty (return not stk), it indicates a valid bracket sequence and returns True.

If the stack is not empty, some opening brackets were not closed, thus the function returns False, representing an invalid

bracket sequence. The solution is efficient with a linear time complexity, O(n), where n is the length of the string s. It only requires a single pass through

the string and constant-time operations for each character. The space complexity is also O(n), in the worst case where the string

consists entirely of opening brackets, which would all be pushed onto the stack.

Example Walkthrough Let's consider a small example to clearly understand the solution approach. Suppose we are given the string $s = "\{[()()]\}$ ". We want to determine if this string represents a valid sequence of brackets. Following the solution steps:

2. Pairs Set: We have a predefined set d that contains '()', '[]', and '{}' to represent valid bracket pairs.

closing bracket in the proper order.

Initialize an empty list to use as a stack

Create a set with valid parentheses pairs

3. Iterating Over the String: We iterate through each character in the string s.

The fourth character) is a closing bracket, so we pop the last element (, which matches the closing bracket, forming a valid

For the first character {, it's an opening bracket, so we push it onto stk.

For the third character (, another opening bracket gets pushed onto stk.

For the second character [, it's also an opening bracket, so we push it onto stk.

The fifth character (is once again an opening bracket and is pushed onto stk.

If the stack is empty or the formed pair is not valid, return False

elif not stack or stack.pop() + char not in valid_pairs:

// Helper method to determine if two brackets are a matching pair

(leftBracket == '{' && rightBracket == '}') ||

(leftBracket == '[' && rightBracket == ']');

private boolean match(char leftBracket, char rightBracket) {

return (leftBracket == '(' && rightBracket == ')') ||

// Return true if pairs match, false otherwise

1. Initialize a Stack: We start with an empty list stk that will be used as our stack.

- pair (). So far so good.
- The sixth character) is a closing bracket, and popping from stk gives us the matching (, forming another valid pair. We continue this process for the remaining characters.
- The seventh character 1 is a closing bracket, and popping from stk gives us [, which is the correct complement, forming the pair []. Finally, the last character } is a closing bracket, and after popping from stk we get {, its correct opening pair, forming {}.

4. Final Stack Check: At the end of iteration, stk is empty because every opening bracket has been matched with the correct

- Since the stack is empty, we return True, indicating that the given string s = "{[()()]}" is a valid sequence of brackets. Our solution has correctly determined the validity by using a stack to manage the ordering and pairing of the brackets.
- Python Solution class Solution: def isValid(self, s: str) -> bool:

valid_pairs = {'()', '[]', '{}'} # Iterate over each character in the string for char in s: 9 # If the character is an opening parenthesis, push it onto the stack 10

```
15
                    return False
16
           # If the stack is empty, all parentheses were valid and correctly nested
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           return not stack
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30 }

stack = []

if char in '({[':

stack.append(char)

```
Java Solution
 1 class Solution {
       // Method to determine if an input string has valid parentheses
       public boolean isValid(String s) {
           // Use a deque as a stack to keep track of the opening brackets
           Deque<Character> stack = new ArrayDeque<>();
           // Iterate over each character in the input string
 6
           for (char c : s.toCharArray()) {
               // If the current character is an opening bracket, push it onto the stack
               if (c == '(' || c == '{' || c == '[') {
 9
                   stack.push(c);
10
11
               } else {
                   // If the stack is empty or the current closing bracket doesn't match
12
13
                   // the last opening bracket, the string is not valid
                   if (stack.isEmpty() || !match(stack.pop(), c)) {
14
                       return false;
15
16
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19
           // If the stack is empty, all brackets were matched correctly
20
           return stack.isEmpty();
```

16 17 18 19

C++ Solution

```
1 class Solution {
   public:
       // Function to check if a given string has valid parentheses
       bool isValid(string s) {
           // Using a string as a stack to store opening brackets
            string stack;
           // Iterate through all characters in the input string
           for (char c : s) {
 9
               // If the character is an opening bracket, push it onto the stack
10
                if (c == '(' || c == '{' || c == '[') {
11
12
                    stack.push_back(c);
13
               // If the stack is empty or characters don't match, return false
14
15
                else if (stack.empty() || !match(stack.back(), c)) {
                    return false;
               // If the characters match, pop the opening bracket from the stack
               else {
20
                    stack.pop_back();
21
22
23
           // If the stack is empty, all brackets were properly closed, return true
24
            return stack.empty();
25
26
27
28
       // Utility function to check if the opening and closing brackets match
29
       bool match(char left, char right) {
            return (left == '(' && right == ')') |
30
                   (left == '[' && right == ']') ||
31
                   (left == '{' && right == '}');
32
33
34 };
35
```

1 // Map that associates opening braces with their corresponding closing braces.

16 17 18

Typescript Solution

['(', ')'],

['[', ']'], ['{', '}'],

1);

2 const bracketPairs = new Map<string, string>([

```
/**
    * Function to determine if the given string has valid bracket pairing.
    * @param {string} str - The input string containing brackets to be checked.
    * @return {boolean} - Returns true if the string is valid, otherwise false.
    */
12
   function isValid(str: string): boolean {
       // A stack to keep track of the expected closing brackets.
14
       const expectedBracketsStack: string[] = [];
15
       // Iterate over each character in the string.
       for (const char of str) {
           // Check if the character is an opening bracket and get its pair.
           if (bracketPairs.has(char)) {
20
               // Push the expected closing bracket onto the stack.
21
22
               expectedBracketsStack.push(bracketPairs.get(char)!);
           } else {
23
               // If the character is not a matching closing bracket,
24
               // or the stack is empty (mismatched brackets), return false.
               if (expectedBracketsStack.pop() !== char) {
26
27
                   return false;
28
29
30
31
32
       // If the stack is empty, all brackets were properly closed; otherwise, return false.
33
       return expectedBracketsStack.length === 0;
34 }
35
Time and Space Complexity
The time complexity of the given code is O(n), where n is the length of the input string. This is because the algorithm iterates over
each character in the input string exactly once.
```

The space complexity of the code is also O(n), as in the worst case (when all characters in the input string are opening brackets), the stack stk will contain all characters in the input string.

Time Complexity:

• Best Case: When the string is empty or consists of a single pair of brackets, the time complexity is 0(1) because it takes a

- Average Case: For a typical string with a mix of opening and closing brackets, the time complexity remains O(n) because each character is processed once. • Worst Case: In the worst case scenario, where there are n characters and the string is properly nested to the deepest level,
- each character is still processed exactly once, giving us a time complexity of O(n).

Space Complexity:

constant amount of time.

- The space complexity is O(n) which occurs when all characters are opening brackets and hence, all are pushed onto the stack. This represents the maximum space utilized by the algorithm.
- As the stack grows with each opening bracket and shrinks with each closing bracket, the actual space used depends on the number of unmatched opening brackets at any point in the algorithm. Therefore, space usage can vary from 0(1) (for an empty string or a string with just one pair of brackets) to O(n) (when all characters are opening brackets).