



Problem Description

calculate the sum of the diagonal elements, which includes elements from both the primary diagonal and the secondary diagonal. The primary diagonal is the one that starts from the top left corner and ends at the bottom right corner. The secondary diagonal starts from the top right corner and ends at the bottom left corner. However, there's a catch: If any element is common between the primary and secondary diagonals (which would be the case for the central element in a matrix with odd dimensions), we must include it only once in our sum.

The problem provides us a square matrix mat, which means the number of rows and columns in the matrix are equal. Our task is to

Intuition

The primary diagonal elements have the same index for their row and column. In terms of indexes, these are the elements mat [1]

To approach this problem, we consider the primary and secondary diagonals of the matrix.

[i] where i ranges from 0 to the n-1, wherein n is the size of one dimension since it's a square matrix.

The secondary diagonal elements have row and column indexes that sum up to n-1. In other words, the indexes are of the form

mat[i][j] where j is n-i-1. Now, the challenge is to make sure that we don't double-count the element in the case of an odd-dimension matrix where the

primary and secondary diagonals intersect. To avoid this, we simply check if the index 1 is equal to the index 1. If they are equal, this means we're looking at the central element in the case of an odd-size matrix and we shouldn't add it again. Therefore, the sum ans starts at 0, and we iterate through each row with its index i. While iterating, we calculate j as n-i-1 for each

row to pinpoint the element in the secondary diagonal. We then add to ans the sum of elements mat[i][i] and mat[i][j], unless i j, in which case we only add the element mat[i][i] once. By iterating over all rows, we calculate the sum required by the problem without having to deal with two separate loops or

Solution Approach

The solution provided is written in Python and follows a straightforward approach that efficiently computes the sum of the diagonals of a square matrix.

Here's a detailed walkthrough of the implementation:

2. We need the size of the matrix, which is the length of one of its sides (since it is square), so we take the length of the matrix n = len(mat).

We start by initializing a variable ans to 0. This will hold the cumulative sum of the diagonal elements.

maintaining additional data structures, therefore making our approach both efficient and straightforward.

3. The for loop is used to iterate over each row of the matrix. The loop variable i serves as the index for both rows and the primary

diagonal elements. enumerate is used so we can have both the index and the row elements available during each iteration.

- 4. We then calculate the column index j for the secondary diagonal element corresponding to the current row. As explained earlier, j is given by n - i - 1.
- 5. Inside the loop, we add to ans the value of the primary diagonal element row[i]. 6. We also want to add the secondary diagonal element row[j] unless i is equal to j (which is the case when we are at the central

element of a matrix with an odd number of rows/columns). To do this succinctly, we add row[j] only if j != i, otherwise, we add

- 0. This conditional addition is achieved by the expression (0 if j == i else row[j]).
- 7. After the loop has processed all the rows, ans now contains the total sum of both diagonals, with no duplicates for the intersecting element (if any).
- There are no complex algorithms, data structures, or patterns involved here. The solution effectively uses index manipulation to address the specific elements required for the sum, making it a simple yet effective approach to solving this problem.
- 1 class Solution: def diagonalSum(self, mat: List[List[int]]) -> int:

Here is the core logic encapsulated in Python code:

8. Finally, the method returns the computed ans.

n = len(mat)for i, row in enumerate(mat): i = n - i - 1ans += row[i] + (0 if j == i else row[j])

```
By carefully choosing and manipulating indexes, we maintain a clear and concise solution without additional memory usage for
storing intermediate results or redundant computations.
Example Walkthrough
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Let's walk through an example to illustrate the solution approach. Suppose we have the following 3×3 square matrix: 1 mat = [

[1, 2, 3],

[4, 5, 6],

[7, 8, 9]

3. Start the for loop with i iterating from 0 to 2 (since the matrix has 3 rows and columns).

Here's the step-by-step process of how the solution works:

4. For each row i, calculate the index j for accessing the secondary diagonal's element, which is n - i - 1. This will give us 2, 1, 0 for i = 0, 1, 2 respectively.

1. Initialize ans to 0. This will keep track of the sum of diagonal elements.

2. Determine the size n of the matrix. In this case, n = len(mat) = 3.

5. Now, add the primary diagonal element to ans. In the first iteration with i = 0, add mat [0] [0] which is 1. 6. Check if i equals j. If they are not the same, add the secondary diagonal element to ans. With i = 0 and j = 2, they aren't

1), we don't add the secondary diagonal element because that would be double-counting.

9. After adding these elements through the loop, the sum ans becomes 1 + 3 + 5 + 9 + 7 = 25.

(because j = 0), so add the secondary diagonal element mat [2] [0] which is 7.

equal, so add mat[0][2] which is 3. 7. For the second iteration where i = 1, we add mat[1][1] (the middle element) to ans. Since i and j are equal here (both equal to

8. In the third iteration with i = 2, add the primary diagonal element mat [2] [2] which is 9 to ans. Index i and j are not the same

- 10. The method returns the value of ans, which is 25. This is the required sum of the elements on the primary and secondary diagonals of the matrix, without double-counting the center element.
 - class Solution: def diagonalSum(self, mat: List[List[int]]) -> int: ans = 0

When we call diagonalSum(mat) with our example matrix, it will return 25, which is the correct answer. This demonstrates the efficiency and simplicity of the provided solution approach.

j = n - i - 1

if j != i:

return total_sum

total_sum += row[i]

Return the computed sum

Putting it all into the Python code, we have:

j = n - i - 1

for i, row in enumerate(mat):

ans += row[i] + (0 if j == i else row[j])

n = len(mat)

return ans

Python Solution

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```
class Solution:
   def diagonalSum(self, matrix: List[List[int]]) -> int:
        # Initialize the sum of the diagonals
        total_sum = 0
       # Get the size of the matrix (assuming it's square)
        n = len(matrix)
       # Loop over each row and calculate the diagonal sum
        for i, row in enumerate(matrix):
            # Calculate the index for the secondary diagonal
```

total_sum += row[j]

Add the primary diagonal element

```
21
Java Solution
   class Solution {
       public int diagonalSum(int[][] matrix) {
           int totalSum = 0; // This will hold the sum of the diagonal elements
           int size = matrix.length; // The matrix is size x size
           // Loop through each row of the matrix
           for (int i = 0; i < size; ++i) {
               int reverseIndex = size - i - 1; // Calculate the corresponding column index for the secondary diagonal
9
               // Add the primary diagonal element
10
               totalSum += matrix[i][i];
11
12
13
               // If it's not the same element (which would be the case in the middle of an odd-sized matrix)
               // then add the secondary diagonal element
14
               if (i != reverseIndex) {
15
                   totalSum += matrix[i][reverseIndex];
16
```

// Return the sum of primary and secondary diagonals, excluding the middle element if counted twice

Add the secondary diagonal element if it's not the same as the primary diagonal

C++ Solution

return totalSum;

```
#include<vector>
   class Solution {
   public:
       // Function to calculate the sum of the elements on the diagonals of a square matrix
       int diagonalSum(std::vector<std::vector<int>>& mat) {
            int total = 0; // Used to store the sum of the diagonal elements
            int size = mat.size(); // Get the size of the square matrix
           // Iterate through each row of the matrix
10
           for (int rowIndex = 0; rowIndex < size; ++rowIndex) {</pre>
12
                int colIndex = size - rowIndex - 1; // Calculate the column index for the secondary diagonal
13
               // Sum the primary diagonal element
14
               total += mat[rowIndex][rowIndex];
15
16
17
               // Sum the secondary diagonal element only if it's not the same as the primary diagonal element
               if (rowIndex != colIndex) {
19
                   total += mat[rowIndex][colIndex];
20
21
22
           return total; // Return the sum of the diagonal elements
24 };
25
Typescript Solution
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// Initialize 'sum' to zero, which will store the final diagonal sum let sum = 0;

function diagonalSum(matrix: number[][]): number {

const matrixSize = matrix.length;

// 'matrixSize' stores the size of the matrix (number of rows/columns)

```
// Iterate through each row of the matrix
       for (let row = 0; row < matrixSize; row++) {
           // Add the elements from both the primary and secondary diagonals for the current row
           sum += matrix[row][row] + matrix[row][matrixSize - 1 - row];
11
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       // If the matrix size is odd, subtract the central element once to correct the sum
14
       if (matrixSize % 2 === 1) {
           // The central element is at the position ['matrixSize' / 2] in both dimensions
16
           sum -= matrix[matrixSize >> 1] [matrixSize >> 1];
17
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       // Return the calculated sum of the diagonal elements
20
       return sum;
21
22 }
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Time and Space Complexity
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Time Complexity The provided code traverses each row only once, and within each row, it accesses two elements directly by their index, which is an (O(1)) operation. Since there are n rows in a square matrix with size (n \times n), the overall time complexity of the code is (O(n)) where (n) is the number of rows (and also the number of columns) in the matrix.

Space Complexity

The code uses a fixed number of variables (ans, n, i, j, row). It does not depend on the size of the input matrix, therefore the space complexity is (O(1)), that is, constant space complexity.