2343. Query Kth Smallest Trimmed Number Quickselect Radix Sort Heap (Priority Queue)

Divide and Conquer

Problem Description

String

Medium Array

In this LeetCode problem, we're given an array of strings called nums, with each string representing a number with leading zeros if required so that all numbers are of the same length. Along with this array, we're given a 2D array of queries. Each query is a pair [k, trim]. For each query, we need to do the following:

Sorting

Leetcode Link

beginning (left side) of the string until we are left with trim number of digits. 2. Find the kth smallest number in this trimmed list. If two numbers are the same, the one whose index is smaller in the original nums

1. Trim each number in the nums array to keep only the rightmost trim digits. This is equivalent to removing digits from the

- array is considered the smaller one. 3. After the answer is determined for each individual query, we reset each number in the nums array to its original value before
- The result is an array of indices indicating the positions of the kth smallest trimmed numbers in the original nums array for each query.

Intuition

The given problem can be approached by looking at each query and processing the nums array based on that query. This involves transforming the original array into a form where it can be sorted, and then picking the kth smallest element.

Here's the intuitive breakdown of the solution:

2. After the array is trimmed, each number along with its original index must be kept together so that after sorting, we can refer

1. For each query, begin by trimming all the numbers in nums. This task involves taking the last trim digits of each number.

back to where the number came from. Tuple pairs can be used for this, with the trimmed number first for sorting purposes,

kth smallest number along with its index.

2. Looping through each [k, trim] pair in the queries list:

proceeding with the next query.

- followed by the original index. 3. We sort the collection of these pairs to arrange the trimmed numbers in ascending order along with their original indices. Sorting
- in Python is stable, which means that if two elements are equal, their order relative to each other will remain the same as it was in the input. 4. Once the array is sorted, we simply pick the element that is at the k-1 position (due to the 0-indexing) which corresponds to the
- 5. This index is added to the result list, which is returned at the end after processing all queries. Utilizing the sort functionality of Python in this way allows us to address the problem in a straightforward manner.
- **Solution Approach**
- The implementation follows directly from the intuitive approach that was discussed. Here's a step-by-step walkthrough using the provided Python solution:

1. Initializing an empty list named ans to store the indices of the kth smallest trimmed numbers for each query.

practice.

during each query.

• 102 becomes 2

473 becomes 3

query, the result is 0.

the result is 0.

class Solution:

Java Solution

import java.util.Arrays;

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answer = []

return answer

modify the original array.

• The key operation in the loop is to create a list of tuples for each number in nums. The tuple consists of the trimmed number (v[-trim:]) and its original index (i). This is facilitated by Python's list comprehension and slicing features. • For the trimming operation, v[-trim:] is used to get the substring of v from the -trim position (counting from the right) to

 Along with the trimmed number, the original index is also stored in the tuple to keep track of its position in the original nums array.

the end. This essentially gives us the rightmost trim digits of each number.

- 3. The list of tuples is then sorted. As mentioned earlier, Python's sorting is stable, so if two trimmed numbers are equal, their
- indexing. From this tuple, [1] is used to extract the second element, which is the original index of the trimmed number in the nums array. 5. The extracted index is appended to the ans list.

6. This process is repeated for each query, and the original nums array remains unaltered as we only use substrings and do not

4. After the list is sorted, the k-1 indexed tuple in the sorted list gives us the kth smallest number because Python uses 0-based

relative order will remain the same and thus the number from the lower index in the original list will be deemed smaller. The

sorting algorithm typically used in Python is Timsort, which is efficient for the kind of partial orderings that appear often in

The solution's time complexity is dominated by the sorting step inside each query loop. If n is the length of nums and m is the length of queries, the expected time complexity is O(m * n * log(n)), with additional space complexity of O(n) for the list of tuples created

7. After processing all queries, the ans list is returned, containing the indices of the kth smallest trimmed number for each query.

- **Example Walkthrough**
- Let's illustrate the solution approach with a small example. Suppose nums is ["102", "473", "251", "814"] and the queries array is [[2,1],[1,2]].
- o 251 becomes 1 ○ 814 becomes 4 The array of trimmed numbers paired with their indices is [('2', 0), ('3', 1), ('1', 2), ('4', 3)].

3. The second smallest number after trimming is at index 0 because the pair ('2', 0) is the second element. Thus, for the first

4. For the next query [1,2], the trim value is 2, so we don't actually trim since all nums elements are already of length 2 or smaller. 5. The sorted array of numbers paired with their indices is [('02', 0), ('14', 3), ('51', 2), ('73', 1)].

3. The index of the kth smallest trimmed number is then found and appended to the result list.

2. We then sort these pairs to get [('1', 2), ('2', 0), ('3', 1), ('4', 3)].

1. For the first query [2,1], trim value is 1. We trim each element in nums to the last digit:

As a result, after processing both queries, we get an answer array of [0, 0]. This walkthrough matches the steps provided in the solution approach:

6. The smallest number (after trimming to two digits, which in this case changes nothing) is at index 0. So, for the second query,

- 1. A list of tuples is created for each number and its index with the numbers trimmed to the correct size. 2. The list of tuples is sorted based on the trimmed numbers.
- **Python Solution** from typing import List

Initialize an empty list to store the answer

String.valueOf(j)

// Compare trimmed strings.

// Return the array of answers.

// Sort the array of trimmed strings and indices.

// If equal, compare their original indices.

// Get the index of the k-th smallest trimmed string.

answers[i] = Integer.parseInt(trimmedAndIndices[k - 1][1]);

int comparison = a[0].compareTo(b[0]);

Arrays.sort(trimmedAndIndices, (a, b) -> {

};

});

return answers;

4. This process is repeated for each query without altering the original nums array.

Iterate through each query in the list of queries for k, trim in queries: # Create a sorted list of tuples where each tuple contains

trimmed_sorted_list = sorted((num[-trim:], index) for index, num in enumerate(numbers))

def smallestTrimmedNumbers(self, numbers: List[str], queries: List[List[int]]) -> List[int]:

a trimmed number (last 'trim' characters) and its original index

From the sorted list, select the k-th smallest trimmed number

Hence the final indices of the kth smallest trimmed numbers for each query are determined and returned in a list.

- # (as sorting starts with 0, we use k-1 as the index) # and append its original index to the answer list answer.append(trimmed_sorted_list[k - 1][1]) 17 18 # Return the list of original indices of the k-th smallest trimmed numbers 19
- class Solution { public int[] smallestTrimmedNumbers(String[] nums, int[][] queries) { // Get the number of strings in the nums array and the total number of queries. int numOfStrings = nums.length; int numOfQueries = queries.length; // Create an array to store the answers for each query. int[] answers = new int[numOfQueries]; 11 12 13 // Initialize a 2D array to store both the trimmed string and its original index. String[][] trimmedAndIndices = new String[numOfStrings][2]; 14 15 // Iterate over each query 16 for (int i = 0; i < numOfQueries; ++i) {</pre> // Extract the k-th value and the trim length from the current query. int k = queries[i][0]; 19 int trimLength = queries[i][1]; 20 21 22 for (int j = 0; j < numOfStrings; ++j) {</pre> // Trim each string in nums from the end by the given trim length and store the result along with the original index. 24 trimmedAndIndices[j] = new String[] { 25 nums[j].substring(nums[j].length() - trimLength),

return comparison == 0 ? Integer.compare(Integer.parseInt(a[1]), Integer.parseInt(b[1])) : comparison;

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C++ Solution

1 #include <vector>

2 #include <string>

#include <algorithm>

using namespace std;

class Solution { 7 public: // Function to return the indices of the smallest trimmed numbers for each query vector<int> smallestTrimmedNumbers(vector<string>& numbers, vector<vector<int>>& queries) { int numOfNumbers = numbers.size(); // Total number of strings in the numbers vector 10 vector<pair<string, int>> trimmedNumbers(numOfNumbers); // Pair to hold trimmed strings and original indices vector<int> answer; // Vector to hold the final results 12 13 14 // Iterate through each query for (auto& query : queries) { int k = query[0]; // kth smallest number to find after trimming int trimLength = query[1]; // Length of the number to consider after trimming 19 // Prepare the trimmed numbers along with their original indices for (int i = 0; i < numOfNumbers; ++i) {</pre> 20 // Trim the number keeping the last 'trimLength' digits and store the original index trimmedNumbers[i] = {numbers[i].substr(numbers[i].size() - trimLength), i}; 25 // Sort the trimmed numbers, (since we're using pairs, it sorts by the first element (trimmed string), and uses the secon 26 sort(trimmedNumbers.begin(), trimmedNumbers.end()); 27 28 // Add the original index of the kth smallest trimmed number to the answer vector answer.push_back(trimmedNumbers[k - 1].second); 29 30 31 32 // Return the final result after processing all queries 33 return answer; 34 35 }; 36 Typescript Solution // In TypeScript, we can simply use arrays and strings, and TypeScript has built—in types for these. // Function to return the indices of the smallest trimmed numbers for each query function smallestTrimmedNumbers(numbers: string[], queries: number[][]): number[] { const numOfNumbers = numbers.length; // Total number of strings in the numbers array const trimmedNumbers: { trimmed: string; originalIndex: number }[] = []; // Array to hold trimmed strings and original indices const answer: number[] = []; // Array to hold the final results

const trimLength = query[1]; // The number of characters to consider from the end of the string after trimming

// Trim the number, keeping the last 'trimLength' characters, and store along with the original index

// Sort the trimmed numbers (the array gets sorted by the 'trimmed' property, and uses the 'originalIndex' for tie-breaking)

26 ? a.trimmed.localeCompare(b.trimmed) 27 : a.originalIndex - b.originalIndex); 29 30 // Add the original index of the kth smallest trimmed number to the answer array

return answer;

Time Complexity

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// Iterate through each query

for (const query of queries) {

for (let i = 0; i < numOfNumbers; ++i) {</pre>

trimmedNumbers[i] = {

trimmedNumbers.sort((a, b) =>

a.trimmed !== b.trimmed

originalIndex: i

const k = query[0]; // kth smallest number to find after trimming

// Prepare the trimmed numbers along with their original indices

trimmed: numbers[i].slice(-trimLength),

Time and Space Complexity

The time complexity of the solution consists of two main operations:

answer.push(trimmedNumbers[k - 1].originalIndex);

// Return the final result after processing all queries

- 1. Trimming and sorting the strings: For each query, we trim the strings to the last trim characters and sort them. This operation has a complexity of O(n * m * log(n)), where n is the length of nums and m is the length of the trimmed string. The m factor comes from the time to create the substring for each number, and log(n) is for the sorting.
- Thus, the total time complexity is O(q * n * m * log(n)).

2. The loop runs for each query, adding a factor of the number of queries q.

Space Complexity

- Space complexity pertains to the extra space required by the algorithm. Here's what it includes: 1. The trimmed and tuples list t, which has a size O(n) for each query since we store the trimmed strings and their original indices.
 - 2. The answer list ans that holds one value for each query, giving it a size of O(q).
- Thus, the total space complexity would be 0(n + q).