2225. Find Players With Zero or One Losses Counting Hash Table Sorting Medium Array

Leetcode Link

The problem provides you with an integer array matches, where each element represents the outcome of a match as [winner_i,

loser_i]. The task is to find all players who have not lost any matches and all players who have lost exactly one match, ensuring that all listed players have participated in at least one match.

It's important to list the players in ascending order. If a player does not appear in the matches array, they are not considered since no match records for them exist. Additionally, players will not lose more than one match because match outcomes are unique.

The first list (answer[0]) should contain all players who have never lost a match.

The second list (answer[1]) should include all players who have lost exactly one match.

Intuition The intuition behind the solution is to track the number of losses for each player. This is achievable by counting the occurrences of

To do this systematically, we can use the following steps:

1. Utilize a Counter to keep a tally of the losses for each player. 2. Iterate over the matches list, incrementing the count for losers. If a player wins a match, ensure they're in the Counter with a loss count of 0 (since winning means they have not lost that match).

 Players with 0 losses go into answer [0]. Players with exactly 1 loss go into answer [1].

- 4. Sort both lists to satisfy the ascending order requirement.
- 5. Return the sorted lists as the final result. Using the Counter, we abstract the complexity of tracking individual losses and make it simple to determine which list a player
- belongs to based on their number of losses after processing all match outcomes.
- it allows us to keep track of how many times each player appears as the loser and does not count their wins. 2. Iterate through the matches list, and for each match [winner, loser]:

Increment the loser's count by one to signify their loss in this match.

haven't lost it.

3. After tallying the losses for each player, create a two-dimensional list ans, where ans [0] will eventually contain players with 0

losses, and ans [1] will contain players with exactly one loss. 4. Iterate through the items in the Counter (for u, v in cnt.items():). For each player (u) and their loss count (v):

o If v is less than 2 (meaning the player has either won all their matches or lost just one), add the player to ans [v]. This works

Check if the winner is already in the Counter. If not, initialize their loss count to 0 because winning a match implies they

6. Finally, return the list ans as the result, where ans [0] contains all players that have not lost any matches and ans [1] contains all players that have lost exactly one match.

because v can only be 0 or 1, as we are only interested in players with no losses or exactly one loss.

- Example Walkthrough
- matches = [[1, 3], [2, 1], [4, 2], [5, 2]]In this array, the subarrays represent match outcomes with winners and losers respectively.

After initializing: Counter({}) 2. Start iterating through matches: Match [1, 3]: Counter({'3': 1}) Match [2, 1]: Counter({'3': 1, '1': 1})

Before proceeding with the loser of the next match, ensure that winner 4 has a loss count of 0: Counter({'3': 1, '1': 1,

Match [4, 2]: Counter({'3': 1, '1': 1, '2': 1, '4': 0}) ∘ Again, ensure that winner 5 has a loss count of 0 before the next loser is processed: Counter({'3': 1, '1': 1, '2': 1, '4': 0, '5': 0})

'4': 0})

Match [5, 2]: Counter({'3': 1, '1': 1, '2': 2, '4': 0, '5': 0}) 3. Now we have all loss counts, create list ans with two sublists for 0 losses and 1 loss:

Player 2 is not added to either sublist since they've lost more than one match.

There's no need to explicitly check if 'winner' is in 'count_losses'

since getting count_losses[winner] will default to 0 if not present

If a player has less than 2 losses, add them to the respective list

Players with no losses are served first & with exactly one loss served after

Combine the two lists into a single list of lists for the result

result = [winners_with_zero_losses, winners_with_one_loss]

o ans = [[], []] 4. Iterate through the counts and populate ans accordingly:

- Player 1 has 1 loss: ans = [[], [3, 1]] Player 4 has 0 losses: ans = [[4], [3, 1]] Player 5 has 0 losses: ans = [[4, 5], [3, 1]]
- o ans[0].sort():ans = [[4, 5], [3, 1]] o ans[1].sort():ans = [[4, 5], [1, 3]]

6. The final sorted ans list is returned from the function:

5. Sort both ans [0] and ans [1] as required:

Player 3 has 1 loss: ans = [[], [3]]

ans = [[4, 5], [1, 3]]

class Solution:

from collections import Counter

def findWinners(self, matches): count_losses = Counter()

for winner, loser in matches:

winners_with_zero_losses = []

winners_with_zero_losses.sort()

winners_with_one_loss.sort()

// Process all match results

for (int[] match : matches) {

int loser = match[1];

int winner = match[0];

winnersList.add(new ArrayList<>());

return result

winners_with_one_loss = []

count_losses[loser] += 1

Count the number of losses for each player

Iterate over the players and their loss counts

for player, loss_count in count_losses.items():

Sort the lists to meet the output criteria

public List<List<Integer>> findWinners(int[][] matches) {

lossCountMap.putIfAbsent(winner, 0);

// Increment the loss count for the loser

List<List<Integer>> winnersList = new ArrayList<>();

// and those who have lost exactly one match (one-match-losers)

vector<vector<int>> findWinners(vector<vector<int>>& matches) {

// Make sure every player is included in the map

// Increment the loss count for the loser of the match

vector<vector<int>> answer(2); // To hold the final result

for (auto& match : matches) {

int winner = match[0];

int loser = match[1];

++lossCount[loser];

if (losses < 2) {

type MatchResult = [number, number];

if (!lossCount.has(winner)) {

lossCount.set(winner, 0);

for (const match of matches) {

const winners: number[] = [];

const [winner, loser] = match;

const oneMatchLosers: number[] = [];

oneMatchLosers.push(player);

oneMatchLosers.sort((a, b) => a - b);

return [winners, oneMatchLosers];

const matchResults: MatchResult[] = [

const results = findWinners(matchResults);

Time and Space Complexity

// Return the sorted lists as a 2D array

winners.sort((a, b) => a - b);

if (!lossCount.count(winner)) {

for (auto& playerLossPair : lossCount) {

int player = playerLossPair.first;

int losses = playerLossPair.second;

answer[losses].push_back(player);

// Sort the list of winners and one-match-losers

sort(answer[0].begin(), answer[0].end());

sort(answer[1].begin(), answer[1].end());

return answer; // Return the sorted lists

// Increment the loss count for the loser of the match

function findWinners(matches: MatchResult[]): number[][] {

// Process each match result and update the loss count for the players

// Usage of the findWinners function can be demonstrated with an example:

console.log(`Winners: \${results[0]}, One-Match Losers: \${results[1]}`);

const currentLossCount = lossCount.get(loser) || 0;

lossCount.set(loser, currentLossCount + 1);

processMatch(lossCount, winner, loser);

lossCount[winner] = 0;

unordered_map<int, int> lossCount; // Map storing the loss count for each player

// Process each match result and update the loss count for the players

// Iterate through the map to classify players based on their loss counts

// If the player has less than 2 losses, add them to the respective list

// Import statements are not required in TypeScript for data structures like arrays and maps.

function processMatch(lossCount: Map<number, number>, winner: number, loser: number): void {

// Function to sort and categorize players based on their loss counts into winners and one-match-losers

const lossCount: Map<number, number> = new Map(); // Map to store the loss count for each player

// Arrays to hold the list of players who never lost and those who lost exactly one match

// Type Alias for readability, representing a match result with winner and loser.

// Ensure every player is present in the map; if not add them with zero losses

// A map to keep track of loss counts for each player

Map<Integer, Integer> lossCountMap = new HashMap<>();

// Initialize the winner's loss count if not already present in map

lossCountMap.put(loser, lossCountMap.getOrDefault(loser, 0) + 1);

if loss_count < 2:</pre> if loss_count == 0: winners_with_zero_losses.append(player) elif loss count == 1: winners_with_one_loss.append(player)

This holds our result, where ans [0] contains players 4 and 5 who have never lost any matches, and ans [1] contains players 1 and 3

Java Solution

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           winnersList.add(new ArrayList<>());
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           // Iterate through each entry in the loss count map
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           for (Map.Entry<Integer, Integer> entry : lossCountMap.entrySet()) {
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               int player = entry.getKey();
               int losses = entry.getValue();
               // If the player has lost fewer than 2 matches, include them in the appropriate sublist (0 losses or 1 loss)
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               if (losses < 2) {
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                   winnersList.get(losses).add(player);
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           // Sort the sublists of players with 0 losses and exactly 1 loss
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           Collections.sort(winnersList.get(0));
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           Collections.sort(winnersList.get(1));
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           // Return the result list
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           return winnersList;
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44 }
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C++ Solution
   #include <vector>
  2 #include <unordered_map>
    #include <algorithm>
  5 class Solution {
  6 public:
         // Function to find the players who have never lost a match (winners)
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// Create the result list containing two lists: one for all the players who have never lost, and one for the players who have

44 45 };

31 // Iterate through the map to classify players based on their loss counts for (const [player, losses] of lossCount) { // If the player has no losses, they are a winner. If only one loss, they are a one-match-loser. 35 **if** (losses === 0) { winners.push(player);

Time Complexity 1. Creating the counter: O(N) - Creating the counter object requires iterating over the list of matches where N is the number of matches.

The given Python code seeks to find all players who never lost a game (who are undefeated), and all players who lost exactly one

- 3. Traversing the counter for players with losses less than 2: 0(P) We go through the counter which contains P unique players. 4. Sorting winners and players with one loss: O(PlogP) - Sorting is performed on the players' list for those who have not lost and those who have lost one match. In the worst case, all players could either be winners or have one loss, hence the sorting can be
- The overall time complexity is the sum of these operations, dominated by the sorting steps: O(N) + O(P) + O(PlogP). Since the sorting term is usually the most significant for large lists, we can simplify this expression to: O(PlogP).
- 1. Counter object: 0(P) The counter object holds at most P unique players, which is the space required. 2. Answer List: 0(P) - The answer list is a list of two lists, which in the worst case would hold all P players, resulting in 0(P) space.
- Therefore, the final complexity of the provided code is:

- Here's how you need to organize the output:
- **Problem Description**

- each player in the loser position of every match.
- 3. After processing all matches, go through the Counter and categorize players based on their loss count.

- **Solution Approach**

- The solution's implementation follows an efficient approach to categorize players based on their match outcomes:
- 1. A Counter object is utilized to maintain the tally of losses for each player. This data structure is optimal for this purpose because

- 5. The lists need to be sorted in increasing order, as per the problem specifications. Thus, both ans [0] and ans [1] are sorted using the sort() method.
- This algorithm efficiently uses a hash map (provided by Counter in Python) to count occurrences, which is ideal for frequency counting tasks. The overall time complexity of the algorithm is determined by the number of matches and the sorting step, and the space complexity is largely influenced by the Counter used to store losses per player. Let's go through a small example to illustrate the solution approach. Assume we have the following matches array:
- 1. Initialize a Counter to count losses:

- who have lost exactly one match, listed in ascending order. Python Solution
- 12 13 17 18 19 20 23 24

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- import java.util.ArrayList; 2 import java.util.HashMap; import java.util.List; import java.util.Map; import java.util.Collections; class Solution {
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- 19 20 21 22

- 9 10 11 12 13

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Typescript Solution 6 // Function to update the loss count map with match results

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- 2. Initializing player counts: 0(N) We iterate through all matches, and for each match we perform a check and counter increment which operates in constant time, resulting in O(N).
- The overall space complexity combines both aspects, remaining O(P) since it is not multiplied by any factor.
- O(PlogP) for P unique players. **Space Complexity**

36 } else if (losses === 1) { 37 38 39 40 41 42 // Sort the list of winners and one-match-losers

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[1, 3],

[2, 3],

[3, 6],

[5, 6],

[5, 7]

- game. Let's break down its time and space complexity.
- - Time Complexity: O(PlogP) Space Complexity: 0(P)