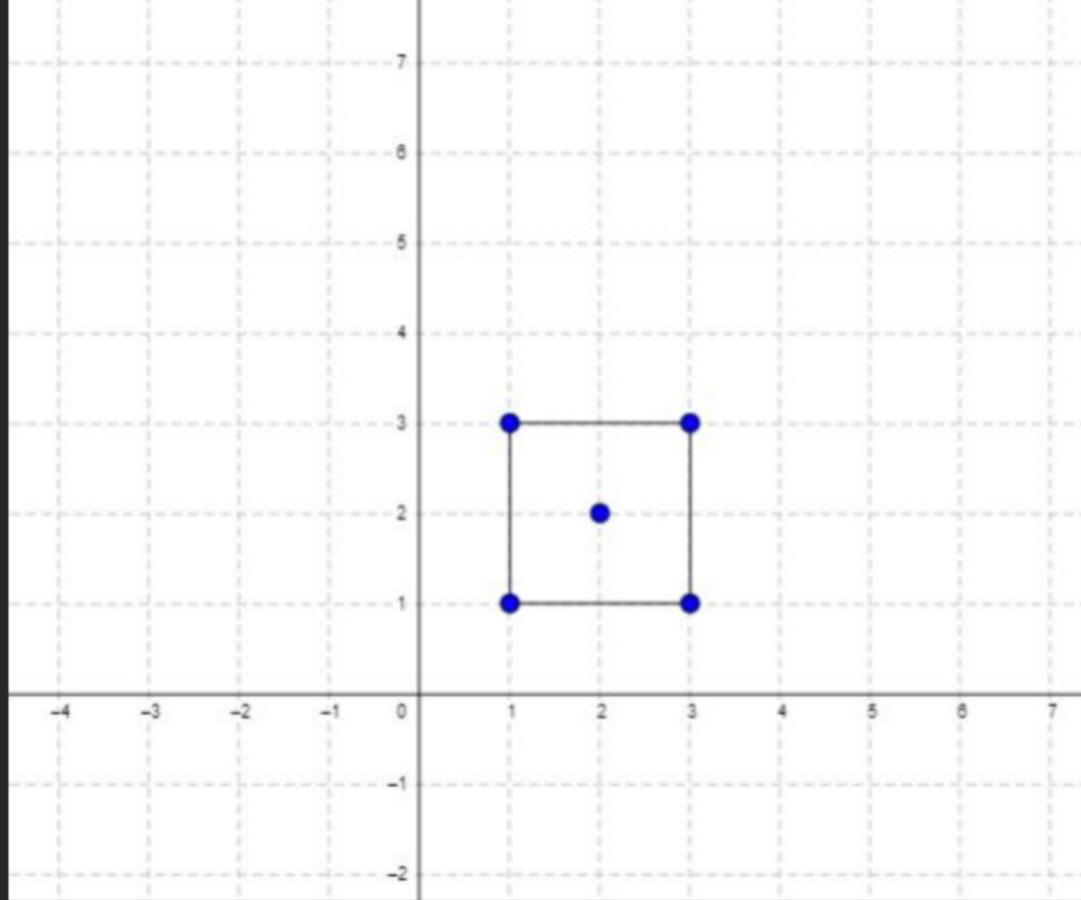
939. Minimum Area Rectangle

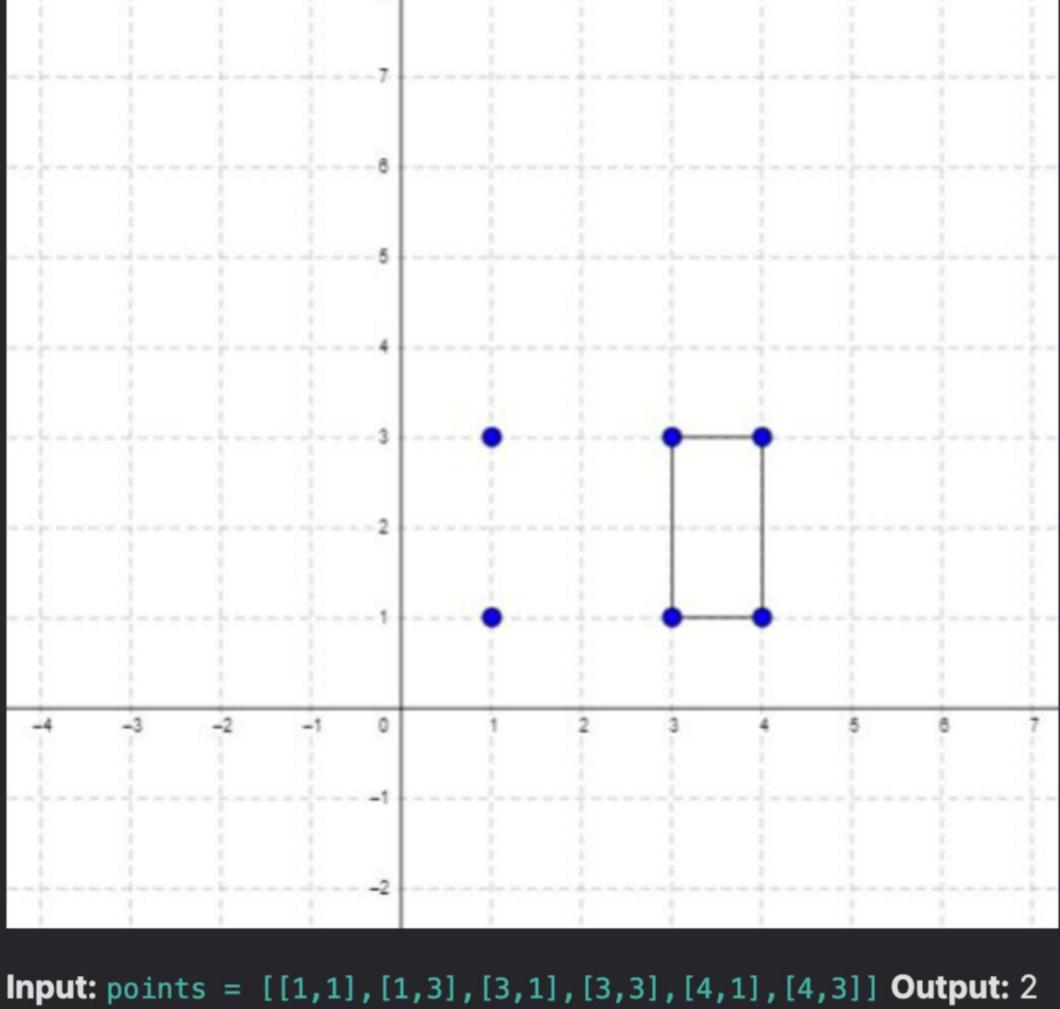
Leetcode Link

You are given an array of points in the X-Y plane points where points [i] = $[x_i, y_i]$.

Return the minimum area of a rectangle formed from these points, with sides parallel to the X and Y axes. If there is not any such rectangle, return 0.

Example 1:





Solution

Brute Force Solution

Then, we return the minimum area from a rectangle formed with these points. One key point is that we need to make sure the rectangle has positive area.

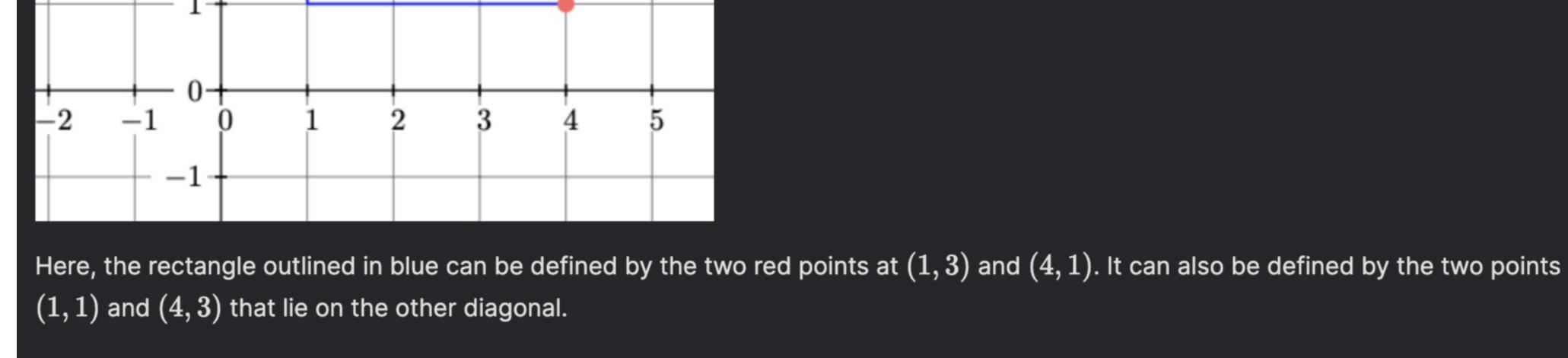
Let N denote the size of points. This algorithm runs in $\mathcal{O}(N^4)$.

Since we need to form a rectangle from 4 different points, we can check all combinations of 4 points to see if it forms a rectangle.

Full Solution

Let's try to optimize our algorithm to find all possible rectangles faster. One observation we can make is that a rectangle can be

Example



rectangle has positive area (i.e. $x_1 \neq x_2, y_1 \neq y_2$).

defined by two points that lie on one of the two diagonals.

Now, instead of trying all combinations of 4 different points from points, we'll try all combinations of 2 different points from points to be the two defining points of the rectangle.

The two defining points have to be a part of points for the rectangle to exist. In addition, we need to check if the two other points in

the rectangle exist in points. Specifically, let's denote the two defining points as (x_1,y_1) and (x_2,y_2) . We'll need to check if (x_1,y_2)

and (x_2,y_1) exist in points. This is where we can use a hashmap to do this operation in $\mathcal{O}(1)$. We'll also need to make sure the

Time Complexity In our algorithm, we check all combinations of 2 different points in points. Since each check runs in $\mathcal{O}(1)$ and there are $\mathcal{O}(N^2)$ combinations, this algorithm runs in $\mathcal{O}(N^2)$.

Since we store $\mathcal{O}(N)$ integers in our <u>hashmap</u>, our space complexity is $\mathcal{O}(N)$. Space Complexity: $\mathcal{O}(N)$.

C++ Solution

class Solution {

int ans = INT_MAX;

public:

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Space Complexity

Time Complexity: $\mathcal{O}(N^2)$.

unordered_map<int, unordered_map<int, bool>> hashMap; for (vector<int> point : points) { // add all points into hashmap hashMap[point[0]][point[1]] = true;

index1++) { // iterate through first defining point

for (int index2 = index1 + 1; index2 < points.size();</pre>

for (int index1 = 0; index1 < points.size();</pre>

int x2 = points[index2][0];

int y2 = points[index2][1];

int x1 = points[index1][0];

int y1 = points[index1][1];

if (x1 == x2 ||

continue;

int minAreaRect(vector<vector<int>>& points) {

index2++) { // iterate through second defining point

y1 == y2) { // rectangle doesn't have positive area

```
if (hashMap[x1].count(y2) &&
22
                        hashMap[x2].count(
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                            y1)) { // check if other points in rectangle exist
24
                        ans = min(ans, abs(x1 - x2) * abs(y1 - y2));
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           if (ans == INT_MAX) { // no solution
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                return 0;
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           return ans;
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33 };
Java Solution
1 class Solution {
       public int minAreaRect(int[][] points) {
           HashMap<Integer, HashMap<Integer, Boolean>> hashMap = new HashMap<>();
            for (int[] point : points) { // add all points into hashmap
                if (!hashMap.containsKey(point[0])) {
                    hashMap.put(point[0], new HashMap<>());
                hashMap.get(point[0]).put(point[1], true);
 9
            int ans = Integer.MAX_VALUE;
            for (int index1 = 0; index1 < points.length;</pre>
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                 index1++) { // iterate through first defining point
13
                int x1 = points[index1][0];
                int y1 = points[index1][1];
14
                for (int index2 = index1 + 1; index2 < points.length;</pre>
15
```

index2++) { // iterate through second defining point

if $(x1 == x2 \mid | y1 == y2)$ { // rectangle doesn't have positive area

ans = Math.min(ans, Math.abs(x1 - x2) * Math.abs(y1 - y2));

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return 0;

return ans;

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Small note: You can use a set in python which acts as a hashset and essentially serves the same purpose as a hashmap for this

&& hashMap.get(x2).containsKey(y1)) { // check if other points in rectangle exist

```
1 class Solution:
       def minAreaRect(self, points: List[List[int]]) -> int:
           min_area = 10 ** 9
           points_table = {}
           for x, y in points: # add all points into hashset
               points_table[(x, y)] = True
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           for x1, y1 in points: # iterate through first defining point
               for x2, y2 in points: # iterate through second defining point
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                   if x1 > x2 and y1 > y2: # Skip looking at same point
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                       if (x1, y2) in points_table and (x2, y1) in points_table: # check if other points in rectangle exist
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                           area = abs(x1 - x2) * abs(y1 - y2)
                           if area:
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                               min_area = min(area, min_area)
16
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```

Example 2:

Input: points = [[1,1],[1,3],[3,1],[3,3],[2,2]] Output: 4

Constraints: $1 \leq \text{points.length} \leq 500 \text{ points[i].length} == 2.0 \leq x_i, y_i \leq 4*10^4$ All the given points are unique.

Python Solution solution.

int x2 = points[index2][0];

int y2 = points[index2][1];

if (hashMap.get(x1).containsKey(y2)

if (ans == Integer.MAX_VALUE) { // no solution

continue;

return 0 if min_area == 10 ** 9 else min_area

Got a question? Ask the Teaching Assistant anything you don't understand.