In this problem, you are given a binary tree and need to find all the "lonely" nodes within it. A lonely node is defined as a node that is

Problem Description

right child with no left sibling. It is important to note that the root of the tree is not considered lonely since it does not have a parent. The goal is to return a list of the values of these lonely nodes. The order of values in the output list is not important.

the only child of its parent, meaning it has no sibling. The node, in this context, could be either a left child with no right sibling or a

Intuition The solution to this problem uses a classic tree traversal approach using a Depth-First Search (DFS) algorithm. The idea is to

traverse the tree starting from the root and check at each node if it has any lonely children. If a node has exactly one child (either left

Here's the intuition behind the DFS approach for this problem: Start DFS from the root of the tree. Upon visiting each node, check if the node has only one child.

 If the node only has a right child (meaning the left child is None), record the value of the right child. Recur for both the left and right children of the current node.

or right), that child is a lonely node, and its value is added to the answer list.

- Once DFS is complete, the answer list will contain the values of all the lonely nodes.
- · Return the answer list.

If the node only has a left child (meaning the right child is None), record the value of the left child.

- This method ensures that every node is visited, and no lonely nodes are missed. It uses the recursive nature of DFS to backtrack and traverse all paths within the tree efficiently.
- Solution Approach

parameter of this function refers to the current node being visited.

is None when a right child exists, or root, right is None when a left child exists.

7. The dfs function is called with the root of the binary tree as its argument, starting the traversal.

8. Once the full tree has been traversed, the ans list is complete, and it is returned as the final result.

approach for solving tree traversal problems. Here's a step-by-step explanation of how the code works:

The provided Python solution implements a recursive Depth-First Search (DFS) strategy on the binary tree, which is a common

1. A helper function named dfs is defined, which will be invoked on the root node and recursively on each child. The root

2. The base case of our recursive function checks for two conditions:

 If the root is None, meaning we have reached a leaf node's child (which doesn't exist), in which case the function returns immediately without doing anything further.

If both the left and right children of the current node are None, meaning the current node is a leaf, there's no need to proceed

further as leaf nodes cannot have lonely nodes. 3. For each non-null node visited by dfs, the function checks if the node has a single child. This is verified by checking if root. left

Let's consider a binary tree with the following structure:

Node 4 is the only child of node 2, and it's also lonely.

allows the function to traverse the whole tree thoroughly.

- 4. If the node has only one child, the value of that lonely child node (root.right.val or root.left.val) is added to the ans list. 5. After checking for loneliness, dfs is called recursively on both the left and right children of the current node, if they exist. This
- 6. Before calling the dfs function, an empty list named ans is created. This will be used to collect the values of the lonely nodes that are encountered during the recursion.
- exploring all nodes of a tree systematically. This specific problem doesn't require maintaining any additional data structures aside from the ans list that accumulates the result. The simplicity and elegance of recursion make the solution concise and highly readable.

The underlying concepts used in this solution include recursive functions, tree traversal, and DFS, which is a fundamental pattern for

 Node 1 is the root and has two children (2 and 3), so it's not lonely. Node 2 has one child (4), but no right sibling, so it's lonely. Node 3 has one child missing (right child), so its left child (5) is lonely.

2. We check if root is None. Since 1 has children, it's not None, we proceed. It's also not a leaf, so no base case conditions are met.

We want to find all the lonely nodes, which in this case are 2, 4, 5, and 6.

Now, following the solution approach:

4. We call dfs(2) and dfs(3) recursively.

In this tree:

Example Walkthrough

It's not None, and it's not a leaf.

Node 5 has a left child (6), but no right sibling, so node 5 is also lonely.

1. We define dfs (root), which we'll call on the root node (1 in our example).

3. We check if node 1 has a single child. It has two, so no nodes are added to the ans list.

For node 4:

Node 5 has no right sibling, so we add 6 (its left child) to ans.

All children are None. It's a leaf node, so no further action is taken.

on the DFS strategy laid out in the problem's solution approach.

def __init__(self, val=0, left=None, right=None):

Helper function to perform DFS

It is None when checking for its children, so no further actions are taken (leaf node).

We would have just walked through the algorithm to correctly identify nodes 4, 5, and 6 as the lonely nodes in the given tree based

 Node 3 has one child missing (right child is None), so we add 5 (its left child) to ans. We call dfs(5).

We call dfs(6).

It's not None, and it's not a leaf.

It's not None, and it's not a leaf.

For node 6:

5. After all recursions complete, we've added all the values of the lonely nodes - 4, 5, and 6 - to the ans list.

self.left = left self.right = right class Solution:

Perform a depth-first search to find all nodes that have only one child (lonely nodes)

If the node is None or is a leaf node (no children), there's nothing to do

if node is None or (node.left is None and node.right is None):

24 # If the node has a left child but no right child, add the left child's value 25 if node.right is None: lonely_nodes.append(node.left.val) 26 27 # Recursively apply DFS to the left and right children 28

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    Check if node 2 has a single child: Node 2 has only one child (4), so we add 4 to ans.

    We call dfs(4) recursively.
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For node 3:

For node 5:

For node 2:

Python Solution

class TreeNode:

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1 # Definition for a binary tree node.

self.val = val

def dfs(node):

return

dfs(node.left)

lonely_nodes = []

dfs(node.right)

If the node has a right child but no left child, add the right child's value 20 if node.left is None: 21 22 lonely_nodes.append(node.right.val) 23

Initialize an empty list to store lonely node values

Trigger DFS from the root of the tree

// Public method to find all lonely nodes

dfs(root);

return lonelyNodes;

return;

dfs(node.left);

dfs(node.right);

* Definition for a binary tree node.

private void dfs(TreeNode node) {

if (node.left == null) {

if (node.right == null) {

public List<Integer> getLonelyNodes(TreeNode root) {

// Private helper method to perform depth-first search

lonelyNodes.add(node.right.val);

lonelyNodes.add(node.left.val);

// Recursively apply DFS to the left subtree

// Recursively apply DFS to the right subtree

// Start depth-first search traversal from root to find lonely nodes

// Base case: if the node is null or it's a leaf (no children)

if (node == null || (node.left == null && node.right == null)) {

// If the node has no left child, the right child is a lonely node

// If the node has no right child, the left child is a lonely node

def get_lonely_nodes(self, root: Optional[TreeNode]) -> List[int]:

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           dfs(root)
           # Return the list of lonely node values
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           return lonely_nodes
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  # Note: The 'Optional' type and 'List' need to be imported from 'typing' module.
  # Example:
41 # from typing import List, Optional
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Java Solution
  // Class to define a binary tree node
   class TreeNode {
       int val;
       TreeNode left;
       TreeNode right;
       TreeNode() {}
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       TreeNode(int val) {
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           this.val = val;
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       TreeNode(int val, TreeNode left, TreeNode right) {
           this.val = val;
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           this.left = left;
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           this.right = right;
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   class Solution {
       // List to store the values of lonely nodes
       private List<Integer> lonelyNodes = new ArrayList<>();
23
```

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C++ Solution

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* struct TreeNode {
          int val;
          TreeNode *left;
          TreeNode *right;
          TreeNode() : val(0), left(nullptr), right(nullptr) {}
          TreeNode(int x) : val(x), left(nullptr), right(nullptr) {}
          TreeNode(int x, TreeNode *left, TreeNode *right) : val(x), left(left), right(right) {}
    * };
11 */
12 class Solution {
13 public:
       // Function to collect all lonely nodes in a binary tree
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       vector<int> getLonelyNodes(TreeNode* root) {
           vector<int> lonelyNodes; // This will hold the lonely nodes, which are nodes that have only one child
           // Define the DFS function to traverse the tree
           function<void(TreeNode*)> dfs = [&](TreeNode* node) {
19
               // Base case: If the node is null or it is a leaf node (no children), return
20
               if (!node || (!node->left && !node->right)) return;
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               // If the node only has a right child
               if (!node->left) {
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                    lonelyNodes.push_back(node->right->val); // Add the value of the right child to our answer
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               // If the node only has a left child
               if (!node->right) {
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                    lonelyNodes.push_back(node->left->val); // Add the value of the left child to our answer
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               // Recursively call the DFS on the left and right children
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               dfs(node->left);
35
               dfs(node->right);
36
           };
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           // Start DFS traversal from the root
           dfs(root);
39
           // Return the list of lonely nodes found
           return lonelyNodes;
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43 };
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Typescript Solution
   // Definition for a binary tree node
   class TreeNode {
       val: number;
       left: TreeNode | null;
       right: TreeNode | null;
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28 // If the node only has a left child 29 if (!node.right) // Add the value of the left child to our answer 31 lonelyNodes.push(node.left.val); 32

dfs(node.left);

dfs(node.right);

if (!node.left) {

this.val = val;

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12 }

this.left = left;

const lonelyNodes: number[] = [];

17 // The DFS function to traverse the tree

const dfs = (node: TreeNode | null): void => {

// If the node only has a right child

lonelyNodes.push(node.right.val);

if (!node || (!node.left && !node.right)) return;

// Add the value of the right child to our answer

// Recursively call the DFS on the left and right children

this.right = right;

```
// Function to collect all lonely nodes in a binary tree
  const getLonelyNodes = (root: TreeNode): number[] => {
      // Start DFS traversal from the root
       dfs(root);
       // Return the list of lonely nodes found
       return lonelyNodes;
45 };
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Time and Space Complexity
Time Complexity
The time complexity of the code is O(n), where n is the number of nodes in the given binary tree. This complexity arises because the
algorithm needs to visit each node exactly once to determine if it has a lonely node (a node with only one child).
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constructor(val: number = 0, left: TreeNode | null = null, right: TreeNode | null = null) {

// This array will hold the lonely nodes, which are nodes that have only one child

// Base case: If the node is null or it is a leaf node (no children), return

Space Complexity The space complexity of the code is O(h), where h is the height of the binary tree. This is because the depth of the recursive call stack will go as deep as the height of the tree in the worst case. Note that this doesn't include the space taken by the output list ans which, in the worst case, can have up to n - 1 elements if all nodes have only one child. If including the output, the space complexity would be 0(n).