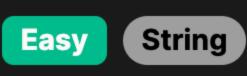
1869. Longer Contiguous Segments of Ones than Zeros



Problem Description

The problem presents us with a binary string s, which consists only of 0s and 1s. Our goal is to determine whether the longest sequence of consecutive 1s is longer than the longest sequence of consecutive 0s in the binary string. If the longest sequence of 1s is longer, we should return true; otherwise, we return false. A sequence of consecutive 1s or 0s means they are directly next to each other without being interrupted by the other digit.

For example, let's take a binary string s = "110100010". Here, the longest segment of 1s has length 2, and the longest segment of

0s has a length of 3. So in this case, the function should return false because the segment of 0s is longer than the segment of 1s.

If there are no 0s in the string, the length of the longest segment of 0s is 0, and similarly, if there are no 1s, the length of the longest segment of 1s is 0.

The algorithm needs to go through the string s and track the length of the current segment of consecutive 1s and 0s, as well as

Intuition

• If the character is a 0, we increase the temporary count of 0s (t0) by 1 and reset the temporary count for 1s (t1) to 0 because the sequence of 1s is interrupted.

the length of the longest segments found so far. The idea is to iterate through the string once and, for each character:

- Similarly, if the character is a 1, we do the opposite: increase the count for 1s and reset the count for 0s.
- While performing these operations, we also keep track of the maximum lengths seen so far for both 0s and 1s. This is done by comparing the current temporary counts (to and t1) with the max counts (no and n1) respectively and updating the max counts if

At the end of the iteration, we simply compare the longest segment of 1s to the longest segment of 0s and return true if the segment of 1s is longer; otherwise, we return false.

Solution Approach

find the longest contiguous segments of 1s and 0s in a single pass through the string.

respectively.

the current is greater.

Here is a step-by-step walkthrough of the implementation from the provided solution code: Initialize four variables: no and n1 are set to 0 to keep track of the maximum lengths of the segments of os and 1s

The solution to the "Check if Binary String Has at Most One Segment of Ones" problem relies on an approach that allows us to

Iterate through each character c in the binary string s:

counting variables and updating maximums, which only employ basic operations.

∘ If c is '0', increment the count of t0 because we are in a segment of 0s, and reset t1 to 0 as we're no longer in a segment of 1s.

∘ If c is '1', increment the count of t1 because we are in a segment of 1s, and reset t0 to 0 as we're no longer in a segment of 0s.

respectively; to and t1 are set to 0 to keep the count of the current length of the segments of 0s and 1s.

After every iteration (for each character c), refresh the maximum values no and n1 by comparing them with to and t1. This ensures that after processing the complete string, no and n1 hold the lengths of the longest segments of os and 1s,

Update n0 by evaluating n0 = max(n0, t0).

contiguous segment of 0s.

 Update n1 by evaluating n1 = max(n1, t1). Finally, once the iteration is finished, return the result of comparing n1 and n0 by returning True if n1 > n0 and False otherwise. This comparison dictates whether the longest contiguous segment of 1s is strictly longer than the longest

The algorithm, thus, performs in O(n) time complexity, where n is the length of the binary string, because it requires just one scan through the string. The space complexity is O(1), as the amount of extra space used does not grow with the size of the input.

No additional data structures or complex patterns are needed for the implementation. The efficiency lies in the simplicity of using

in the solution approach, here's a step-by-step explanation of what happens: Initialize variables: n0 = 0, n1 = 0 to keep track of the maximum lengths of 0s and 1s, and t0 = 0, t1 = 0 to keep track of the

Let us consider a small example with the binary string s = "0110011". When this string is passed through the algorithm described

current lengths.

Example Walkthrough

Iteration starts: \circ c = '0': t0 becomes 1, t1 is reset to 0. We compare n0 with t0, now n0=1 and n1=0.

- ∘ c = '1': t1 becomes 1, t0 is reset to 0. Since n1 (which is 0) is less than t1 (which is 1), we update n1=1. Now n0=1 and n1=1. ∘ c = '1': t1 increases to 2, t0 remains 0. Update n1 to 2, as it is greater than the current n1. Now n0=1 and n1=2. o c = '0': t0 becomes 1, t1 is reset to 0, do not update n0 or n1 as current values of t0 and t1 are not greater than n0 and n1.
- \circ c = '1': t1 becomes 1, t0 is reset to 0. Do not update n0 or n1.

```
\circ c = '1': t1 increases to 2, t0 remains 0. Do not update n1 as t1 equals n1. Final values are n0=2 and n1=2.
     Once the iteration is over, we compare n1 and n0. Since n1 is not greater than n0, we return False.
  No segment of 1s is strictly longer than the longest segment of 0s in the binary string s = "0110011", so our algorithm correctly
  returns False based on the provided method. The complexity remains linear since we only went through the string a single time,
  and the space complexity is constant as we only kept track of a fixed number of variables.
Solution Implementation
```

Python

// Update the longest sequence lengths if the current sequences are longer

longestZeroSeq = Math.max(longestZeroSeq, currentZeroSeq);

longestOneSeq = Math.max(longestOneSeq, currentOneSeq);

return longestOneSeq > longestZeroSeq;

longestZero = max(longestZero, currentZero);

* Determines if the longest contiguous segment of 1s is longer than the

* @param {string} s - A binary string containing only '0's and '1's.

* @return {boolean} - True if the longest segment of 1s is longer than

the longest segment of 0s; otherwise, false.

let maxZeros: number = 0; // Tracks the length of the longest segment of 0s.

let max0nes: number = 0; // Tracks the length of the longest segment of 1s.

let currentZeros: number = 0; // Tracks the length of the current segment of 0s.

let currentOnes: number = 0; // Tracks the length of the current segment of 1s.

* longest contiguous segment of 0s in the binary string.

longestOne = max(longestOne, currentOne);

return longestOne > longestZero;

const checkZeroOnes = (s: string): boolean => {

// Update the longest sequence length of '1's if necessary

// Return true if the longest sequence of '1's is greater than the longest sequence of '0's

 \circ c = '0': t0 increases to 2, t1 remains 0. Update n0 to 2, as it is equal to the current t0. Now n0=2 and n1=2.

class Solution: def checkZeroOnes(self, s: str) -> bool: # Initialize variables to keep track of the longest consecutive sequence of '0's and '1's max_zeros = max_ones = 0

for char in s: **if** char == '0': # Increase the sequence count of '0's, reset the sequence of '1's current_zeros += 1

else:

current_zeros = current_ones = 0

current_ones = 0

current_zeros = 0

Iterate over the characters in the string

Do the opposite when we encounter '1's

```
current ones += 1
           # Update the maximum sequence lengths for '0's and '1's
           max zeros = max(max zeros, current zeros)
           max_ones = max(max_ones, current_ones)
       # Check if the maximum sequence of '1's is greater than that of '0's
       return max_ones > max_zeros
Java
class Solution {
    public boolean checkZeroOnes(String s) {
        int longestZeroSeq = 0; // Variable to track the length of the longest contiguous sequence of '0's
        int longestOneSeq = 0; // Variable to track the length of the longest contiguous sequence of '1's
        int currentZeroSeq = 0; // Variable to track the length of the current contiguous sequence of '0's
        int currentOneSeq = 0; // Variable to track the length of the current contiguous sequence of '1's
       // Iterate over each character in the string
        for (int i = 0; i < s.length(); ++i) {</pre>
            if (s.charAt(i) == '0') { // If the current character is '0'
                currentZeroSeq++; // Increase the length of the current '0' sequence
                currentOneSeq = 0; // Reset the length of the '1' sequence
            } else { // If the current character is '1'
                currentOneSeq++; // Increase the length of the current '1' sequence
                currentZeroSeq = 0; // Reset the length of the '0' sequence
```

};

/**

*/

TypeScript

```
C++
class Solution {
public:
   // Function to check if the longest continuous sequence of 1s is longer than the longest continuous sequence of 0s
   bool checkZeroOnes(string s) {
        int longestZero = 0; // Stores the length of the longest sequence of '0's
        int longestOne = 0; // Stores the length of the longest sequence of '1's
        int currentZero = 0; // Tracks the current sequence length of '0's
        int currentOne = 0; // Tracks the current sequence length of '1's
       // Loop through the input string character by character
       for (char c : s) {
           if (c == '0') {
                // If the current character is '0', increase the count of the current sequence of '0's
               ++currentZero;
               // Reset the current sequence of '1's since it's interrupted by a '0'
               currentOne = 0;
            } else { // c is '1'
               // If the current character is '1', increase the count of the current sequence of '1's
               ++currentOne;
               // Reset the current sequence of '0's since it's interrupted by a '1'
                currentZero = 0;
           // Update the longest sequence length of '0's if necessary
```

// After going through the string, check if the longest sequence of '1's is strictly greater than '0's

```
for (let char of s) {
          if (char === '0') {
              currentZeros++; // Increment the count of contiguous 0s.
              currentOnes = 0; // Reset the count of contiguous 1s.
          } else {
              currentOnes++; // Increment the count of contiguous 1s.
              currentZeros = 0; // Reset the count of contiguous 0s.
         // Update the maximum counts if the current counts are greater.
         maxZeros = Math.max(maxZeros, currentZeros);
          maxOnes = Math.max(maxOnes, currentOnes);
      // Return true if longest segment of 1s is longer than longest segment of 0s.
      return maxOnes > maxZeros;
  // The function can be used as follows:
  // const result: boolean = checkZeroOnes("11001");
  // console.log(result); // Output would be true if the longest segment of 1s is longer than the longest segment of 0s.
class Solution:
   def checkZeroOnes(self, s: str) -> bool:
       # Initialize variables to keep track of the longest consecutive sequence of '0's and '1's
       max_zeros = max_ones = 0
       current_zeros = current_ones = 0
       # Iterate over the characters in the string
       for char in s:
           if char == '0':
               # Increase the sequence count of '0's, reset the sequence of '1's
               current zeros += 1
               current_ones = 0
           else:
               # Do the opposite when we encounter '1's
               current_zeros = 0
               current_ones += 1
           # Update the maximum sequence lengths for '0's and '1's
           max_zeros = max(max_zeros, current_zeros)
           max_ones = max(max_ones, current_ones)
       # Check if the maximum sequence of '1's is greater than that of '0's
       return max_ones > max_zeros
Time and Space Complexity
```

The time complexity of the given code is O(n), where n is the length of the input string s. This is because the code consists of a

single for-loop that iterates over each character in the string exactly once. During each iteration, the code does a constant amount of work: updating counters to and t1, resetting one counter depending

operations are constant time, the loop represents a linear time complexity relative to the length of the string. The space complexity of the code is 0(1). The amount of extra space used by the algorithm does not depend on the size of the

on the character, and calculating the maximum length of consecutive '0's and '1's seen so far using max(). Since all these

input string. The variables n0, n1, t0, and t1 use a fixed amount of space. No additional space that grows with input size is required. Therefore, the space complexity is constant.