2001. Number of Pairs of Interchangeable Rectangles

Medium Array Hash Table Math Counting Number Theory

Problem Description

In this problem, we are given a list of n rectangles, each represented by their width and height as an element in a 2D integer array `rectangles`. For example, the i-th rectangle is represented as `rectangles[i] = [width_i, height_i]`.

The task here is to find how many pairs of these rectangles are interchangeable. Two rectangles are considered interchangeable if they have the equivalent width-to-height ratio. In mathematical terms, rectangle i is interchangeable with rectangle j if `width_i/height_i` equals `width_j/height_j`. Note that we will be using decimal division to compare ratios.

The expected result is to return the total count of such interchangeable pairs among the rectangles provided.

Intuition

To find pairs of interchangeable rectangles, we focus on calculating the width-to-height ratio for each rectangle. However, directly comparing floating-point ratios might lead to precision issues. To avoid this, we normalize each width-to-height ratio by dividing both width and height by their greatest common divisor (GCD). This way, two rectangles that have the same width-to-height ratio will have the same normalized width and height, thus can be represented by the same key when stored.

To efficiently track how many rectangles have the same normalized width and height, we use a `Counter` data structure (a type of dictionary that is optimized for counting hashable objects). When we normalize a rectangle, we check how many we've seen with that same normalized width and height so far (using the `Counter`), and we increment our answer by that count—this is because each of those could form an interchangeable pair with the current rectangle.

For each rectangle, we then increment its count in the `Counter`. This step ensures that we keep track of all interchangeable rectangles encountered.

Combining these strategies allows us to count the pairs without comparing every rectangle to every other rectangle, which would be inefficient. The use of the greatest common divisor (GCD) ensures that we are accurately identifying interchangeable rectangles without floating-point arithmetic issues.

Solution Approach

The solution uses a combination of hashing and the greatest common divisor (GCD) algorithm to efficiently count interchangeable rectangle pairs.

Here is a step-by-step breakdown of the implementation:

- We initialize the variable `ans` to store the total count of interchangeable rectangle pairs and a `Counter` data structure called `cnt`, which will be used to keep track of the occurrences of normalized width-to-height ratios.
- We iterate over each rectangle given in the `rectangles` list.
  - For the current rectangle, w represents the width and h represents the height.
  - We calculate the greatest common divisor of w and h using the function `gcd(w, h)`. This step is crucial as it allows us to normalize the width and height to their simplest ratio form.
  - We then normalize w and h by dividing them by the calculated GCD. As a result, the tuple `(w // g, h // g)` represents the unique key of the current ratio in its simplest form.
  - The expression `cnt[(w, h)]` retrieves the current count of rectangles that have been seen with the same normalized ratio. This count directly corresponds to the number of new pairs that can be formed with the current rectangle since each of those previous rectangles could form a pair with it.
  - We increment `ans` by the count retrieved from `cnt[(w, h)]`.
  - Finally, we increment `cnt[(w, h)]` by 1 because we have encountered another rectangle with the same normalized ratio.

By hashing the normalized ratios, the algorithm ensures constant-time lookups for previous occurrences, bypassing the need for nested loops, which would increase the time complexity significantly.

The final result stored in `ans` is returned, which represents the total number of pairs of interchangeable rectangles within the input array.

Using this approach, the solution is both efficient and precise, utilizing the mathematical properties of ratios and the efficiencies gained through hashing.

Example Walkthrough

Let's consider an example with the following list of rectangles: `rectangles = [[4, 8], [3, 6], [10, 20], [15, 30]]`.

We want to find the total count of interchangeable rectangle pairs that can be formed from these rectangles.

1. First, we initialize `ans` to 0, which will hold our answer, and `cnt`, a `Counter` to keep track of normalized width-to-height ratios.
2. For the first rectangle `[4, 8]`, we compute its GCD to normalize the ratio. `gcd(4, 8)` is 4, so when we divide both by 4, the normalized ratio is (1, 2). We have not seen this ratio before, so `cnt[(1, 2)]` is 0. We then increment `cnt[(1, 2)]` by 1.
3. Next, rectangle `[3, 6]` comes down to the same normalized ratio (1, 2) after dividing by the `gcd(3, 6)` which is 3. We check `cnt[(1, 2)]` which is 1 (from the previous step), so we increase `ans` by 1 and `cnt[(1, 2)]` becomes 2.
4. Then, with rectangle `[10, 20]`, after dividing by `gcd(10, 20)` which is 10, the ratio is again (1, 2). The `cnt[(1, 2)]` is currently 2, meaning we can form 2 more pairs with each of the rectangles we've seen before, so we add 2 to `ans`, making it a total of 3 so far, and `cnt[(1, 2)]` becomes 3.
5. Lastly, for the rectangle `[15, 30]`, we divide by `gcd(15, 30)` which is 15, and we get the normalized ratio (1, 2). Since `cnt[(1, 2)]` is 3, we add 3 to `ans` and increment `cnt[(1, 2)]`.

Now, `ans` is 6, representing the total number of interchangeable rectangle pairs, and `cnt[(1, 2)]` is 4, indicating we saw the same ratio 4 times.

By applying our solution approach, we have efficiently found that there are 6 pairs of interchangeable rectangles in our example list without having to compare each rectangle to every other rectangle. This process has bypassed potential floating-point issues and kept the computation time to a minimum.

Solution Implementation

Python

```python
from math import gcd
from collections import Counter

class Solution:
    def interchangeableRectangles(self, rectangles: List[List[int]]) -> int:
        # Initialize a variable to keep count of pairs
        pair_count = 0

        # Initialize a counter to keep track of the occurrences of each ratio
        ratio_counter = Counter()

        # Loop through each rectangle
        for width, height in rectangles:
            # Calculate the greatest common divisor of width and height
            gcd_value = gcd(width, height)

            # Normalize the width and height by dividing them by the gcd to obtain the ratio
            normalized_width, normalized_height = width // gcd_value, height // gcd_value

            # The number of rectangles with the same width-to-height ratio so far
            # will be the number of interchangeable rectangle pairs we can form with the current one
            pair_count += ratio_counter[(normalized_width, normalized_height)]

            # Increment the counter for the current ratio
            ratio_counter[(normalized_width, normalized_height)] += 1

        return pair_count
```

Java

```java
class Solution {
    public long interchangeableRectangles(int[][] rectangles) {
        long countInterchangeablePairs = 0; // this will hold the final answer
        int n = rectangles.length + 1; // using length + 1 to ensure a unique mapping when multiplied
        Map<Long, Integer> ratioCountMap = new HashMap<>(); // this map stores the counts of each unique rectangle ratio

        // Loop through each rectangle
        for (int[] rectangle : rectangles) {
            int width = rectangle[0];
            int height = rectangle[1];
            int gcdValue = gcd(width, height); // compute the greatest common divisor of width and height
            width /= gcdValue; // normalize the width by the gcd
            height /= gcdValue; // normalize the height by the gcd
            long ratioHash = (long) width * n + height; // create a unique hash for the width/height ratio

            // Update the number of interchangeable pairs count
            countInterchangeablePairs += ratioCountMap.getOrDefault(ratioHash, 0);
            // Increase the count for this ratio in the map
            ratioCountMap.merge(ratioHash, 1, Integer::sum);
        }
        return countInterchangeablePairs; // return the total count of interchangeable rectangle pairs
    }

    // Helper function to compute the greatest common divisor (GCD) of two numbers
    private int gcd(int a, int b) {
        return b == 0 ? a : gcd(b, a % b);
    }
}
```

C++

```cpp
#include <vector>
#include <unordered_map>
using namespace std;

class Solution {
public:
    long long interchangeableRectangles(vector<vector<int>>& rectangles) {
        long long answer = 0;  // This will hold the final count of interchangeable rectangle pairs
        int numRectangles = rectangles.size();  // Total number of rectangles given
        unordered_map<long long, int> countMap;  // Hashmap to store the count of rectangles with the same ratio

        for (auto& rectangle : rectangles) {
            int width = rectangle[0];  // Extract the width and height of the current rectangle
            int height = rectangle[1];  // Find greatest common divisor of width and height to get the ratio
            int gcdValue = gcd(width, height);  // Simplify width by dividing by the gcd
            height /= gcdValue;  // Similarly, simplify height by dividing by the gcd

            // 'key' is the unique identifier for the ratio of width to height.
            // Multiplying width by a large number (greater than the maximum height)
            // and then adding the height ensures a unique value for each width to height ratio.
            long long key = 1LL * width * (numRectangles + 1) + height;

            // If a rectangle with the same ratio was already encountered,
            // then increase the count of interchangeable pairs by the count of those rectangles.
            answer += countMap[key];

            // Increment the count of the current ratio by 1.
            countMap[key]++;
        }

        return answer;  // Return the total count of interchangeable rectangle pairs.
    }

private:
    // Utility function to calculate greatest common divisor
    int gcd(int a, int b) {
        while (b != 0) {
            int temp = b;
            b = a % b;
            a = temp;
        }
        return a;
    }
};
```

TypeScript

```typescript
/**
 * Function to find the number of pairs of rectangles that are interchangeable.
 * Rectangles are interchangeable if one can become the other by resizing.
 * We normalize each rectangle's dimensions by their greatest common divisor (GCD)
 * and count occurrences of each unique pair of normalized dimensions.
 * @param {number[][]} rectangles - A list of rectangle specifications, given by [width, height].
 * @return {number} - The number of interchangeable rectangle pairs.
 */
function interchangeableRectangles(rectangles: number[][]): number {
    const countMap: Map<number, number> = new Map();
    let pairCount: number = 0;

    for (let [width, height] of rectangles) {
        // Calculate the greatest common divisor of width and height
        const gcdValue: number = gcd(width, height);

        // Normalize the width and height using gcdValue
        width = Math.floor(width / gcdValue);
        height = Math.floor(height / gcdValue);

        // Create a unique hash (or map key) for the normalized rectangle dimensions
        const hash: number = width + (rectangles.length + 1) * height;

        // Increment the pair count by the number of occurrences already found
        pairCount += (countMap.get(hash) || 0);

        // Update the map with the new count for the current normalized dimensions
        countMap.set(hash, (countMap.get(hash) || 0) + 1);
    }

    return pairCount;
}

/**
 * Helper function to calculate the greatest common divisor of two numbers.
 * @param {number} a - First number
 * @param {number} b - Second number
 * @return {number} - The greatest common divisor of a and b.
 */
function gcd(a: number, b: number): number {
    if (b === 0) {
        return a;
    }
    return gcd(b, a % b);
}
```

```python
from math import gcd
from collections import Counter

class Solution:
    def interchangeableRectangles(self, rectangles: List[List[int]]) -> int:
        # Initialize a variable to keep count of pairs
        pair_count = 0

        # Initialize a counter to keep track of the occurrences of each ratio
        ratio_counter = Counter()

        # Loop through each rectangle
        for width, height in rectangles:
            # Calculate the greatest common divisor of width and height
            gcd_value = gcd(width, height)

            # Normalize the width and height by dividing them by the gcd to obtain the ratio
            normalized_width, normalized_height = width // gcd_value, height // gcd_value

            # The number of rectangles with the same width-to-height ratio so far
            # will be the number of interchangeable rectangle pairs we can form with the current one
            pair_count += ratio_counter[(normalized_width, normalized_height)]

            # Increment the counter for the current ratio
            ratio_counter[(normalized_width, normalized_height)] += 1

        return pair_count
```

Time and Space Complexity

Time Complexity

The time complexity of the given code is primarily determined by the for loop that iterates over every rectangle within the `rectangles` list. During each iteration, the following operations occur:

- Calculating the greatest common divisor (GCD) for the width and height of the rectangle, which is done using the `gcd` function. The time complexity of the `gcd` function is generally $O(\log(\min(w, h)))$, where w and h are the width and height of the rectangle.
- Reducing the width and height by dividing by their GCD.
- Checking and updating the count of a particular ratio (width to height) in the hash map implemented by the `Counter` class.

Assuming n represents the number of rectangles, the time complexity for the `gcd` operation, which is the most expensive operation inside the loop, adds up to $O(n * \log(\min(w, h)))$. However, since this is a reduction to the simplest form, the value of $\log(\min(w, h))$ is small compared to n. Therefore, we often approximate this to $O(n)$ when w and h are not extreme values.

The lookup and update in the hash map `cnt` have an average-case time complexity of $O(1)$ for each operation since these operations are constant time in a hash map (dictionary in Python) on average.

Thus, the overall average time complexity of the code is $O(n)$.

Space Complexity

The space complexity of the code is influenced by the space required to store the reduced width and height pairs in a hash map. In the worst case, if all rectangles have different width-to-height ratios, the space complexity would be $O(n)$ because each rectangle would be represented in the hash map after reduction.

Furthermore, the counter `ans` is using $O(1)$ space and the `gcd` call does not use additional space proportional to n.

Therefore, the total space complexity of the given code is $O(n)$.