1358. Number of Substrings Containing All Three Characters

Sliding Window

Problem Description

Hash Table

String)

The given problem requires us to find the number of substrings in a string s that contains at least one occurrence of the characters 'a', 'b', and 'c'. The string s consists only of these three characters. A substring is any sequence of consecutive characters from the string. The goal is to count all such possible substrings where each of the three characters appears at least once.

Intuition

Medium

string. The key intuition here is to understand that once we have found a substring containing all three characters, extending this substring to the right (by adding more characters in sequence) will also form valid substrings containing all three characters.

Here's the step-by-step intuition:

To solve this problem, we use a sliding window approach to track the latest positions of 'a', 'b', and 'c' while iterating through the

1. Initialize a dictionary to store the latest index of 'a', 'b', and 'c'. By default, they are set to -1, indicating that they haven't been found yet.

- 2. Iterate through the string character by character, updating the dictionary with the new index of each character encountered.
- 3. At each step, determine the smallest of the three indices because the smallest index indicates the rightmost position up to which we have seen all three characters together.
- 4. For the current index i, the count of valid substrings ending at i will be the minimum index among the latest indices of 'a', 'b', and 'c' plus one. This is because any substring starting from index 0 to the minimum index will have all three characters up to the current position i.
- 5. Sum up all these counts to get the total number of substrings containing all three characters.

 This simple yet elegant solution effectively counts all the necessary substrings by considering each valid end position and how

many substrings it can generate based on the earlier occurrences of 'a', 'b', and 'c'.

The solution approach uses the concept of pointers and a hash map (in Python, a dictionary) to efficiently keep track of the latest occurrence of each character 'a', 'b', and 'c'.

$d = {\text{"a": } -1, \text{"b": } -1, \text{"c": } -1}$ # Dictionary to store the latest index for 'a', 'b', and 'c' ans = 0 # This will hold the total count of substrings

Solution Approach

The dictionary d serves as our hash map, holding the most recent indices of each character. Initially, all characters are set to -1,

```
indicating they have not been encountered yet.

Then, we start iterating over each character in the string:
```

for i, c in enumerate(s):
 d[c] = i # Update the latest index for the character `c`
 ans += min(d["a"], d["b"], d["c"]) + 1 # Count substrings ending at the current index `i`

As we iterate, for each character (c) at index i, we update its latest index in the dictionary. The min(d["a"], d["b"], d["c"]) finds the smallest index of the three, which, as previously mentioned, is the furthest right we can go while still having all three

characters. By adding 1 to this minimum value, we get the number of substrings ending at index i that has all three characters.

return ans

1. Initialize a dictionary with keys 'a', 'b', 'c', and values -1.

Each time we find a character, update its latest index in the dictionary.

Finally, the sum of all such substrings is returned as the final answer:

2. Start iterating over each character in the string using a for loop.

Algorithm

- 5. Accumulate this count to a running total.
- 6. After the loop finishes, return the total count as the answer.

Now we iterate through the string while keeping track of the latest index of each character in d.

in a time complexity of O(n) and a space complexity of O(1), where n is the length of the string.

Calculate the number of valid substrings that end at the current index (minimum index of 'a', 'b', 'c' + 1).

Example Walkthrough

We start by initializing the dictionary d with {"a": -1, "b": -1, "c": -1}.
 Set our running total of valid substrings ans to ∅.

This algorithm is efficient because it processes each character exactly once and uses constant space for the dictionary, resulting

For i = 0 (character = 'a'): Update d to {"a": 0, "b": -1, "c": -1}.

```
The minimum index among 'a', 'b', and 'c' is −1, so ans += −1 + 1, resulting in ans = 0.
For i = 1 (character = 'b'):
```

Let's use a string s = "abcabc" to illustrate the solution approach.

- Update d to {"a": 0, "b": 1, "c": -1}.
- The minimum index among 'a', 'b', and 'c' is still -1, so ans +=-1+1, resulting in ans =0.

For i = 2 (character = 'c'):

For i = 4 (character = 'b'):

For i = 5 (character = 'c'):

Solution Implementation

Python

Java

C++

public:

class Solution {

int numberOfSubstrings(string s) {

int substringCount = 0;

// Initialize the answer to 0.

// Iterate over the string.

int lastSeenPositions[3] = $\{-1, -1, -1\}$;

class Solution {

class Solution:

- Update d to {"a": 0, "b": 1, "c": 2}.
- For i = 3 (character = 'a'):
 Update d to {"a": 3, "b": 1, "c": 2}.

 \circ The minimum index among 'a', 'b', and 'c' is now 1, so ans += 1 + 1, resulting in ans = 3.

 \circ Now we have all three characters, and the minimum index is 0, so ans += 0 + 1, resulting in ans = 1.

- Update d to {"a": 3, "b": 4, "c": 2}.
 The minimum index among 'a', 'b', and 'c' is 2, so ans += 2 + 1, resulting in ans = 6.
- Update d to {"a": 3, "b": 4, "c": 5}.
 The minimum index among 'a', 'b', and 'c' is 3, so ans += 3 + 1, resulting in ans = 10.

Create a dictionary to keep track of the last seen index of 'a', 'b', and 'c'

This is because a valid substring must include at least one of each

Throughout this process, we efficiently tracked the latest positions of 'a', 'b', and 'c', calculated the substrings terminating at each character position, and accumulated this to find the total number of substrings containing all three characters.

After the loop finishes, we return the total count, which is 10.

def numberOfSubstrings(self, string: str) -> int:

last_seen_index[char] = index

last_seen_index = {"a": -1, "b": -1, "c": -1}

Update the last seen index for the current character

Initialize answer to store the number of valid substrings answer = 0 # Enumerate over the characters of the string for index, char in enumerate(string):

Increment the answer by one more than the smallest last seen index among 'a', 'b', and 'c'

```
answer += min(last_seen_index.values()) + 1

# Return the total count of valid substrings that contain at least one of each 'a', 'b', and 'c'
return answer
```

```
public int numberOfSubstrings(String s) {
    // Array to store the latest positions of characters 'a', 'b', and 'c'
    int[] latestPosition = new int[] {-1, -1, -1};

    // This will hold the count of valid substrings
    int answer = 0;

    // Iterate over each character in the string
    for (int i = 0; i < s.length(); ++i) {
        char currentChar = s.charAt(i);

        // Update the latest position of the current character
        latestPosition[currentChar - 'a'] = i;

        // Find the smallest index among the latest positions of 'a', 'b', and 'c'
        // and add 1 to get the count of valid substrings ending with the current character
        int minPosition = Math.min(latestPosition[0], Math.min(latestPosition[1], latestPosition[2]));
        answer += minPosition + 1;
    }

    return answer; // Return the total count of valid substrings
}</pre>
```

// Find the smallest index among the last seen positions of 'a', 'b', and 'c'.
// Add 1 because indices are 0-based, and we're interested in the number of element of the sint minusestSeenPosition = min(lastSeenPositions[0].

// Update the last seen position for the current character.

lastSeenPositions[s.charCodeAt(index) - 'a'.charCodeAt(0)] = index;

// Find the smallest index among the last seen positions of 'a', 'b', and 'c'.

for (int index = 0; index < s.size(); ++index) {</pre>

lastSeenPositions[s[index] - 'a'] = index;

```
// Add 1 because indices are 0-based, and we're interested in the number of elements.
            int minLastSeenPosition = min(lastSeenPositions[0],
                                          min(lastSeenPositions[1], lastSeenPositions[2])) + 1;
           // Add the number of valid substrings ending with the current character.
           // This is possible because any substring ending at the current index
            // and starting before or at the smallest last seen index will contain all three characters.
            substringCount += minLastSeenPosition;
       // Return the total count of valid substrings.
       return substringCount;
};
TypeScript
// Function to count the number of substrings containing all three characters 'a', 'b', and 'c'.
function numberOfSubstrings(s: string): number {
    // Initialize an array to store the last seen positions of 'a', 'b', and 'c'.
    const lastSeenPositions: number[] = [-1, -1, -1];
    // Initialize the answer to 0.
    let substringCount: number = 0;
    // Iterate over the string.
    for (let index = 0; index < s.length; ++index) {</pre>
       // Update the last seen position for the current character.
```

// Function to count the number of substrings containing all three characters 'a', 'b', and 'c'.

// Initialize an array to store the last seen positions of 'a', 'b', and 'c'.

```
# Update the last seen index for the current character
last_seen_index[char] = index
# Increment the answer by one more than the smallest last seen index among 'a', 'b', and 'c'
# This is because a valid substring must include at least one of each
answer += min(last_seen_index.values()) + 1

# Return the total count of valid substrings that contain at least one of each 'a', 'b', and 'c'
return answer
```

Create a dictionary to keep track of the last seen index of 'a', 'b', and 'c'

last_seen_index = {"a": -1, "b": -1, "c": -1}

Enumerate over the characters of the string

for index, char in enumerate(string):

answer = 0

Time and Space Complexity

Initialize answer to store the number of valid substrings

The time complexity of the given code is O(n), where n is the length of the string s. This is because the code iterates over each character in the string exactly once. Within the loop, updating the dictionary and calculating the minimum value and the cumulative sum is done in constant time.

The space complexity of the code is 0(1). The space is constant because the dictionary d only stores three key-value pairs regardless of the size of the input string, corresponding to the characters 'a', 'b', and 'c'.