

# 1496. Path Crossing

EasyHash TableString

[Leetcode Link](#)

## Problem Description

In this LeetCode problem, we are given a string `path` that represents a sequence of moves. Each character in `path` stands for a directional move: 'N' for north, 'S' for south, 'E' for east, and 'W' for west. Each move is one unit long. We start at the origin point (0, 0) on a two-dimensional plane and follow the moves indicated in the `path` string. The task is to determine whether or not our path ever crosses itself. In other words, if we ever visit the same point more than once during our walk, we return `true`. If our path does not cross and we never visit the same point more than once, we return `false`.

## Intuition

To solve this problem, the intuitive approach is to track every position we visit as we traverse the path defined by the string. We can use a set to store our visited positions because sets allow fast lookup times to check whether we have been to a position or not, as duplicates are not allowed in a set.

We start by initializing our position to the origin (0, 0) and create an empty set called `vis` (short for "visited") which will hold tuples of our coordinates on the 2D plane. As we iterate over each move in the path string, we update our current position by incrementing or decrementing our `i` (for the north-south axis) and `j` (for the east-west axis) accordingly.

After each move, we check whether the new coordinate (represented as a tuple `(i, j)`) is already present in our `vis` set. If it is, it means we've just moved to a spot we've previously visited, which means our path has crossed, and we return `true`. If the coordinate is not in the set, we add it to the set and continue onto the next move in the path.

We repeat this process for each move in the path. If we finish iterating over all moves without returning `true`, it means our path never crosses itself, and we return `false`.

## Solution Approach

The solution to the problem implemented in Python uses a set data structure and simple coordinate manipulation to track the movement on the path. Below is an overview of the approach, breaking down how the algorithm works.

1. Initialize the current position to the origin, `(i, j) = (0, 0)`.
2. Create a set named `vis` (short for visited) and add the initial position to it. Sets are chosen because they store unique elements, allowing us to quickly check if a position has been visited before.
3. Loop through each character in the `path` string:
  - The `for c in path:` loop iterates over each character in the `path` string.
  - The `match` statement (a feature available in Python 3.10 and above) works like a switch-case statement found in other languages. It matches the character `c` with one of the cases: 'N', 'S', 'E', or 'W'.
  - Based on the direction, we update our `(i, j)` coordinates:
    - For 'N', we decrement `i` to move north (`i -= 1`).
    - For 'S', we increment `i` to move south (`i += 1`).
    - For 'E', we increment `j` to move east (`j += 1`).
    - For 'W', we decrement `j` to move west (`j -= 1`).
4. After updating the coordinates, we check if the new position `(i, j)` is already present in the `vis` set:
  - If the condition `(i, j) in vis:` is `True`, we return `True` since the path has crossed a previously visited position.
  - If the position is not found in the set, we add the new position to the set with `vis.add((i, j))`.
5. If the loop completes without finding any crossing, the `return False` statement at the end of the function ensures we return `False`, as no path has been crossed.

This approach uses straightforward coordinate tracking and set membership checks to efficiently solve the problem. The time complexity is  $O(N)$ , where `N` is the length of the path, since we visit each character once, and the space complexity is also  $O(N)$ , due to the storage required for the set that holds the visited positions.

## Example Walkthrough

Let's assume our given `path` string is "NESWW".

Following the solution approach, here's a step-by-step illustration of how the algorithm will execute:

1. We initialize our current position at the origin (0, 0), so `(i, j) = (0, 0)`.
2. We create an empty set `vis` and add the initial position to it, so `vis = {(0, 0)}`.
3. We start looping through each character in the `path`:
  - The first character is 'N'. We decrement `i` because we're moving north, so `i = 0 - 1 = -1` and `j` remains 0. The new position is `(-1, 0)`, which is not in `vis`, so we add it: `vis = {(0, 0), (-1, 0)}`.
  - The second character is 'E'. We increment `j` to move east, so `i` remains -1, and `j = 0 + 1 = 1`. The new position is `(-1, 1)`, which is also not in `vis`, so we add it: `vis = {(0, 0), (-1, 0), (-1, 1)}`.
  - The third character is 'S'. We increment `i` to move south, so `i = -1 + 1 = 0` and `j` remains 1. The new position is `(0, 1)`, not in `vis`, so we add it: `vis = {(0, 0), (-1, 0), (-1, 1), (0, 1)}`.
  - The fourth character is 'W'. We decrement `j` to move west, so `i` remains 0, and `j = 1 - 1 = 0`. The position `(0, 0)` is already in `vis`, indicating we've returned to the origin. Since this position is revisited, we would return `True` as the path crosses itself.

Therefore, the function would return `True` based on the input path "NESWW", because we revisited the starting point, indicating a crossing path.

## Python Solution

```
1 class Solution:
2     def isPathCrossing(self, path: str) -> bool:
3         # initialize starting point
4         x, y = 0, 0
5         # set to keep track of visited coordinates
6         visited = {(0, 0)}
7
8         # iterate over each character in the path string
9         for direction in path:
10             # move in the corresponding direction
11             if direction == 'N':
12                 x -= 1
13             elif direction == 'S':
14                 x += 1
15             elif direction == 'E':
16                 y += 1
17             elif direction == 'W':
18                 y -= 1
19
20             # check if the new position has already been visited
21             if (x, y) in visited:
22                 # if we've been here before, path crosses. Return True
23                 return True
24
25             # add the new position to the set of visited coordinates
26             visited.add((x, y))
27
28             # if visited all positions without crossing, return False
29             return False
30
```

## Java Solution

```
1 class Solution {
2
3     public boolean isPathCrossing(String path) {
4         // Two variables to keep track of current position
5         int x = 0, y = 0;
6         // Use a HashSet to store visited coordinates.
7         Set<Integer> visited = new HashSet<>();
8         // Hash for the origin, adding it as the first visited coordinate
9         visited.add(0);
10
11         // Iterate over the path characters
12         for (int index = 0; index < path.length(); ++index) {
13             char direction = path.charAt(index);
14
15             // Move in the grid according to the current direction
16             switch (direction) {
17                 case 'N': // Moving north decreases the y-coordinate
18                     y--;
19                     break;
20                 case 'S': // Moving south increases the y-coordinate
21                     y++;
22                     break;
23                 case 'E': // Moving east increases the x-coordinate
24                     x++;
25                     break;
26                 case 'W': // Moving west decreases the x-coordinate
27                     x--;
28                     break;
29             }
30
31             // Calculate a unique hash for the current position.
32             // Multiplying by a large enough number to not mix coordinates
33             int hash = y * 20000 + x;
34
35             // Check if this position has been visited before, if so, path crosses
36             if (!visited.add(hash)) {
37                 return true; // early return if the path crosses itself
38             }
39         }
40
41         // If no crossing points were found, return false
42         return false;
43     }
44 }
45
```

## C++ Solution

```
1 #include <unordered_set>
2 #include <string>
3
4 class Solution {
5 public:
6     // Determines if a path crosses itself based on commands in a string
7     bool isPathCrossing(const std::string& path) {
8         // Initialize (i, j) as the starting position (0, 0)
9         int x = 0, y = 0;
10
11         // Create a hash set to track visited positions with a unique key
12         std::unordered_set<int> visitedPositions{{0}};
13
14         // Iterate through each character in the path string
15         for (const char &direction : path) {
16             // Update position based on direction
17             if (direction == 'N') {
18                 --x; // Move north
19             } else if (direction == 'S') {
20                 ++x; // Move south
21             } else if (direction == 'E') {
22                 ++y; // Move east
23             } else {
24                 --y; // Move west
25             }
26
27             // Calculate a unique key for the position
28             int key = x * 20001 + y; // Use prime number to reduce collisions
29
30             // Check if the position has been visited before
31             if (visitedPositions.count(key)) {
32                 // If visited before, path crosses itself
33                 return true;
34             }
35
36             // Add the new position to the set of visited positions
37             visitedPositions.insert(key);
38         }
39
40         // If no crossing occurred, return false
41         return false;
42     }
43 };
44
```

## Typescript Solution

```
1 function isPathCrossing(path: string): boolean {
2     // Initialize current position at the origin (0,0)
3     let position: [number, number] = [0, 0];
4     // Create a set to store visited coordinates as a unique identifier
5     const visited: Set<string> = new Set();
6
7     // Add the starting position (origin) to the visited set
8     visited.add(position.toString());
9
10    // Iterate through each character in the path string
11    for (const direction of path) {
12        // Update the position according to the direction
13        switch (direction) {
14            case 'N': // North decreases the x coordinate
15                position[0]--;
16                break;
17            case 'S': // South increases the x coordinate
18                position[0]++;
19                break;
20            case 'E': // East increases the y coordinate
21                position[1]++;
22                break;
23            case 'W': // West decreases the y coordinate
24                position[1]--;
25                break;
26        }
27
28        // Convert the tuple to a string to create a unique identifier for the position
29        const positionKey = position.toString();
30        // If the position has been visited, return true and exit
31        if (visited.has(positionKey)) {
32            return true;
33        }
34        // Add the new position to the visited set
35        visited.add(positionKey);
36    }
37    // If no crossing paths are detected, return false
38    return false;
39 }
40
```

## Time and Space Complexity

The given Python code checks if a path crosses itself based on a string of movement commands ('N', 'S', 'E', 'W' corresponding to North, South, East, and West movements). The code uses a set `vis` to track all the visited coordinates.

### Time Complexity:

The time complexity of the code is  $O(n)$ , where `n` is the length of the input string `path`. This is because the code iterates through each character of the `path` string exactly once.

For each character, the operations performed (updating coordinates and checking the set for the existence of the coordinates) are constant time operations, thus each character in the path requires a constant amount of time processing.

### Space Complexity:

The space complexity of the code is  $O(n)$ , where `n` is the length of the input string `path`. In the worst case, none of the positions will be revisited, so the set `vis` will contain a unique pair of coordinates for each move in the `path`. Thus, the maximum size of the set is proportional to the number of movements, which corresponds to the length of the `path`.

In summary, the code has a linear time and space complexity with respect to the length of the input `path`.