# **Problem Description**

an array called ages where ages [i] represents the age of the i-th person.

In this problem, we're modeling a simplified version of friend requests on a social media platform. Every user has an age, and there is

The task is to count the number of friend requests made on the platform under certain rules. Person x can send a friend request to person y only if they meet the following three conditions:

3. If person y is over 100 years old, then person x must also be over 100 (no one under 100 can send a friend request to someone

1. Person y's age is **not** less than or equal to 0.5 times the age of person x plus 7.

over 100).

2. Person y's age is less than or equal to the age of person x.

It's important to note that friend requests aren't reciprocal (if x requests y, y does not automatically request x), and no one can send a friend request to themselves. The goal is to figure out the total number of these possible friend requests.

Intuition

## To solve the problem, we analyze the conditions under which one person will send a friend request to another and apply these to all possible pairs of ages. We need to consider the age limitations and the fact that there could be multiple people with the same age.

n2).

Since the problem limits ages to be between 1 and 120, we can iterate over each possible age for both x and y (the sender and the receiver of the friend request). For each age pair, we calculate the number of possible friend requests based on the number of people who have those ages.

The solution uses a Counter to tally the number of people at each age which allows efficient lookups. For each pair of ages (i, j), if i can send a request to j according to the rules above, we increase our friend request count by n1 \* n2 where n1 is the number of people with age i and n2 is the number of people with age j. We need to account for the fact that people cannot friend request themselves, so if i equals j, we subtract the number of people who would have otherwise friend-requested themselves (which is

This approach uses two nested loops to check all possible age pairs and applies the given conditions to calculate the number of friend requests. **Solution Approach** 

The solution provided proceeds through a process that we can break down step by step: 1. Import the Counter Class: The Counter class from the collections module is used, which provides a way to count the number

of occurrences of each unique element in an iterable. In this case, it's used to keep track of how many people are of each age.

# 2. Initialize the Answer: A variable ans is initialized to 0. This variable keeps a running total of all valid friend requests.

to n2 people of age j.

answer.

•  $j \le 0.5 * i + 7$ 

(person y).

3. Double Loop Through Each Age: Two nested loops iterate over all possible ages from 1 to 120. The first loop uses i to represent the age of the person sending the request (person x), and the second loop uses j to represent the age of the potential recipient

4. Getting Count of Each Age Group: Using the Counter created, we retrieve n1, the count of people with age i and n2, the count of people with age j.

5. Apply the Conditions: Within the nested loop, for each (i, j) pair, the code checks whether the three conditions mentioned in

6. Update the Answer: If the conditions are satisfied (meaning that i can send a friend request to j), the number of friend requests

is updated by adding n1 \* n2 to the total. This accounts for all possible friend requests that can be made by n1 people of age i

- the problem statement are not met. If all three conditions are false, it means that a person with age i can send a friend request to the person with age j.
- self requests, which is n2. This is because when the ages are equal, each person has been counted as sending a request to themselves once, which is not allowed. 8. Return the Answer: After all pairs of ages have been checked, the total number of friend requests is returned as the final

7. Self-request Adjustment: Since no one can send a friend request to themselves, if i equals j, the code subtracts the number of

• j > i • j > 100 && i < 100

The solution iterates over all age combinations only once, resulting in an efficient approach with a time complexity that is constant—

not dependent on the number of people but rather on the fixed number range (which is 1 to 120). This is a key insight that greatly

simplifies an otherwise potentially complex problem, especially since a straightforward nested loop over the original list of ages

**Example Walkthrough** 

2. Initialize the Answer: We set ans = 0.

Here is how the conditions look like in code:

- Let's illustrate the solution approach: 1. Import the Counter Class: First, we import Counter from collections.
- So, n1 is the count of people per age on the outer loop and n2 is that on the inner loop. 5. Apply the Conditions: We check if person x (i) can send a friend request to person y (j) by negating the conditions given.

Person 14 cannot send to anyone (fails condition 1 with everyone else).

4. Getting Count of Each Age Group: We use a Counter to count age occurrences in our example.

would have a time complexity more directly tied to the number of persons.

Suppose we have an array of ages: ages = [14, 16, 16, 60].

ages present in our array). Here, they are: 14, 16, and 60.

Counter(ages) gives {14: 1, 16: 2, 60: 1}.

= 2 valid friend requests from ages 16 to 16.

def num\_friend\_requests(self, ages: List[int]) -> int:

# Counter object to store the frequency of each age

count\_a = age\_count[age\_a] # Number of people with age age\_a

friend\_requests += count\_a \* count\_b

if (senderAge == receiverAge) {

// Return the total number of friend requests

int numFriendRequests(vector<int>& ages) {

// Count the number of instances of each age

return friendRequests;

for (int age : ages) {

ageCounter[age]++;

friendRequests -= receiverCount;

// Create a counter array to store the number of people at each age.

vector<int> ageCounter(121, 0); // Initialized with 121 elements all set to 0

# Inner loop to iterate through all possible ages to find potential friends

# Check the conditions when a person A can send a friend request to B:

# If the conditions are met, increase the count of friend requests

count\_b = age\_count[age\_b] # Number of people with age age\_b

# Condition 1: B should not be less than or equal to 0.5 \* A + 7

# Condition 3: If B is over 100, then A must be over 100 as well

also over 100). 6. Update the Answer: We go through combinations, considering our array and other conditions.

Person 60 can send to everyone 16 and over, but since there are only people of age 16 and 60, they can send to the 2

Persons of age 16 can send to each other but not to age 14 (fails condition 1) or 60 (fails condition 2).

2 people of age 16 sending requests to 2 people of age 16 (including themselves initially) is 2 \* 2.

 $\circ$  Not (j > 100 && i < 100), which simplifies to i >= 100 || j <= 100 (since we can't send to someone over 100 unless i is

3. Double Loop Through Each Age: Now, we start two nested loops of ages between 1 and 120 (but in practice, we only consider

7. Self-request Adjustment: ∘ Since people can't friend request themselves, we subtract each person of age 16 from that count, which gives us 2 \* 2 - 2

 $\circ$  The final ans = 4.

through the use of a count for each age bracket.

age\_count = Counter(ages)

for age\_a in range(1, 121):

for age\_b in range(1, 121):

people of age 16.

 $\circ$  j > 0.5 \* i + 7

∘ j <= i

8. Return the Answer: ○ We have 2 requests from persons aged 16 to 16, and each person aged 60 can send to both aged 16, so 2 + 2 = 4.

• There is no need to adjust for age 14 or 60 as 14 cannot send any requests and there is only one person aged 60.

1 from collections import Counter # Import the Counter class from the collections module class Solution:

# Condition 2: B should be less than or equal to A (A can send to B of same age or younger)

if not (age\_b <= 0.5 \* age\_a + 7 or age\_b > age\_a or (age\_b > 100 and age\_a < 100)):

# If both ages are the same, we have to subtract the instances where A is B

This is the total number of valid friend requests that happen in our small example: 4. The process uses logical conditions and

arithmetic to carefully tally the requests between age groups, accounting for self-requests and avoids unnecessary computations

# Initialize the answer to 0 friend\_requests = 0 10 # Loop through all possible ages from 1 to 120

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Python Solution

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                        # because a person cannot friend request themselves
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                        if age_a == age_b:
                            friend_requests -= count_b
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31
32
            # Return the total friend requests that can be made
33
            return friend_requests
34
Java Solution
   class Solution {
        public int numFriendRequests(int[] ages) {
            // Array to keep count of each age (up to 120)
            int[] ageCount = new int[121];
            // Fill the ageCount array with the frequency of each age
            for (int age : ages) {
                ageCount[age]++;
 9
10
            // Variable to hold the final result
            int friendRequests = 0;
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            // Loop through each age for the sender
            for (int senderAge = 1; senderAge < 121; senderAge++) {</pre>
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                int senderCount = ageCount[senderAge];
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                // Only continue if there are people with this age
17
                if (senderCount > 0) {
                    // Loop through each age for the receiver
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20
                    for (int receiverAge = 1; receiverAge < 121; receiverAge++) {</pre>
21
                        int receiverCount = ageCount[receiverAge];
22
                        // Check if the friend request condition is satisfied
23
                        if (!(receiverAge \leftarrow 0.5 * senderAge + 7 || receiverAge \rightarrow senderAge || (receiverAge \rightarrow 100 && senderAge \leftarrow 100))) {
24
                             // Add the product of the counts of the respective ages
25
                             friendRequests += senderCount * receiverCount;
26
                             // Deduct the count when sender and receiver are of the same age
```

#### 10 11 12 int totalFriendRequests = 0; // Initialize the total number of friend requests to 0 13 // Loop through the ages from 1 to 120 (inclusive) 14

C++ Solution

1 class Solution {

2 public:

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for (int ageA = 1; ageA <= 120; ++ageA) {</pre>
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                int countAgeA = ageCounter[ageA]; // Number of people with age A
16
17
               // Loop through all possible ages for potential friends
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                for (int ageB = 1; ageB <= 120; ++ageB) {</pre>
                    int countAgeB = ageCounter[ageB]; // Number of people with age B
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22
                    // Check if a friend request can be sent according to the problem's conditions
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                    if (!(ageB \leftarrow 0.5 * ageA + 7 || // Age B should not be less than or equal to 0.5 * Age A + 7
                          ageB > ageA ||
24
                                                // Age B should be less than or equal to Age A
25
                          (ageB > 100 && ageA < 100) // If Age B is > 100, then Age A should also be > 100
26
                        totalFriendRequests += countAgeA * countAgeB; // Add the count product to total requests
28
29
                        // If ages are the same, subtract the self request count
30
                        if (ageA == ageB) {
                            totalFriendRequests -= countAgeB; // Subtract the diagonal
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37
           // Return the total number of friend requests
38
           return totalFriendRequests;
39
40 };
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Typescript Solution
   // Array to store the number of people at each age, initialized with 121 elements all set to 0
   const ageCounter: number[] = new Array(121).fill(0);
   // Function that calculates the number of friend requests
    function numFriendRequests(ages: number[]): number {
       // Count the number of instances of each age in the input array
       ages.forEach((age) => {
            ageCounter[age]++;
       });
10
       let totalFriendRequests: number = 0; // Initialize the total number of friend requests to 0
11
12
       // Iterate through the ages from 1 to 120 (inclusive) to calculate the number of friend requests
13
       for (let ageA = 1; ageA <= 120; ++ageA) {</pre>
14
            const countAgeA = ageCounter[ageA]; // Number of people with age A
16
17
           // Loop through all possible ages to check for potential friend matches
            for (let ageB = 1; ageB <= 120; ++ageB) {</pre>
               const countAgeB = ageCounter[ageB]; // Number of people with age B
20
```

# Time and Space Complexity

return totalFriendRequests;

ageB > ageA ||

// Return the total computed friend requests

complexity of the code is predominated by the Counter operation, resulting in O(N).

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(ageB > 100 && ageA < 100) // If Age B is over 100, then Age A must also be over 100 24 25 26 totalFriendRequests += countAgeA \* countAgeB; // Accumulate the product of the counts into total requests 28 // If Age A is the same as Age B, subtract the count for self requests 29 if (ageA === ageB) { totalFriendRequests -= countAgeB; // Adjust for self-request scenarios by subtracting the count 33 34

// Age B must be less than or equal to Age A

if (!(ageB  $\leftarrow$  0.5 \* ageA + 7 || // Age B should not be less than or equal to 0.5 times Age A + 7

// Check if a friend request can be sent according to the given conditions

## The given code consists of two nested loops each ranging from 1 to 121, resulting in a fixed number of iterations for the loops. Since the loops do not depend on the size of the input (ages list), they have a constant runtime. However, the overall time complexity does still depend on the Counter operation performed on the ages list earlier, which iterates over the entire list once.

Time Complexity

can be 120.

The Counter operation has a time complexity of O(N) where N is the number of elements in ages. The nested loops have a constant time complexity of 0(121 \* 121) because they iterate over a fixed range independent of the input size. Therefore, the total time

**Space Complexity** The space complexity of the code is influenced by the additional storage needed for the counter variable, which depends on the number of unique elements in ages. In the worst case, each age is unique, so the space complexity for counter is O(K) where K is the

number of unique ages. Given the constraints of the problem (ages are between 1 and 120), the maximum number of unique ages K

Therefore, the space complexity is O(K), yet since K is constant and limited to 120, this could also be considered O(1) in the context of this problem where K does not scale with the size of the input.