1272. Remove Interval

Medium Array

In this problem, we're dealing with a mathematical representation of sets using intervals of real numbers. Each interval is

Problem Description

represented as [a, b], which means it includes all real numbers x such that $a \ll x \ll b$. We are provided with two things:

• A sorted list of disjoint intervals, intervals, which together make up a set. The intervals are disjoint, meaning they do not overlap, and they are sorted in ascending order based on their starting points.

• Another interval, toBeRemoved, which we need to remove from the set represented by intervals.

Our objective is to return a new set of real numbers obtained by removing toBeRemoved from intervals. This set also needs to be represented as a sorted list of disjoint intervals. We need to consider that part of an interval might be removed, all of it might

be removed, or it might not be affected at all, depending on whether it overlaps with toBeRemoved.

The key to solving this problem is to examine each interval in intervals and figure out its relation with toBeRemoved. There are

three possibilities:

2. The interval is partially or completely inside the range of toBeRemoved and needs to be trimmed or removed. 3. The interval straddles the edges of toBeRemoved and might need to be split into two intervals.

1. The interval is completely outside the range of toBeRemoved and therefore remains unaffected.

• If the current interval ends before toBeRemoved starts or starts after toBeRemoved ends, it's disjoint and can be added to the result as is. • If there is overlap, we may need to trim the current interval. If the start of the current interval is before to BeRemoved, we can take the portion

Given that intervals is sorted, we can iterate over each interval and handle the cases as follows:

from the interval's start up to the start of toBeRemoved. Similarly, if the interval ends after toBeRemoved, we can take the portion from the end of toBeRemoved to the interval's end.

• If b <= x, then the interval [a, b] is completely before toBeRemoved and also remains unaffected.

Solution Approach

• We need to handle the edge cases where toBeRemoved completely covers an interval, in which case we add nothing to the result for that

- The provided solution employs a straightforward approach to tackle the problem by iterating through each interval in the given sorted list intervals and comparing it with the toBeRemoved interval. Here's a step by step process used in the implementation:
- The solution starts by initializing an empty list ans, which will eventually contain the resulting set of intervals after the

It then enters a loop over each interval [a, b] in the intervals list. For each interval, it checks whether there is an intersection with the toBeRemoved interval, [x, y]. It does this by verifying

removal process.

removal and is added to ans.

interval exactly once.

toBeRemoved interval:

Step 1: Initialize Result List

• Current interval [1, 4).

ans = [] (empty to begin with)

Step 2: Loop Over Each Interval in intervals

Step 3: Check for Intersection with toBeRemoved

Step 3: Check for Intersection with toBeRemoved

Step 3: Check for Intersection with toBeRemoved

• intervals = [[1, 4), [6, 8), [10, 13)]

the original set with the toBeRemoved interval excluded.

interval.

two conditions: If a >= y, then the interval [a, b] is completely after toBeRemoved and thus is unaffected.

When either of the above conditions is true, the current interval can be added directly to the ans list without modification

o If the start of the interval a is before x (the start of toBeRemoved), then the segment [a, x) of the original interval is unaffected by the

- since it doesn't intersect with toBeRemoved. If the interval does intersect with toBeRemoved, the solution needs to handle slicing the interval into potentially two parts:
- Similarly, if the end of the interval b is after y (the end of toBeRemoved), then the segment [y, b) remains after the removal and is also added to ans.
- The loops continue for all intervals in intervals, applying the above logic. After processing all intervals, the solution returns the ans list, which now contains the modified set of intervals, representing
- **Example Walkthrough** Let's consider the following small example to illustrate the solution approach. Assume we have the following intervals list and

The algorithm makes use of simple conditional checks and relies on the sorted nature of the input intervals for its correctness

and efficiency. The overall time complexity is O(n), where n is the number of intervals in intervals, since it processes each

toBeRemoved = [7, 12) Using the steps outlined in the solution approach:

• The interval [1, 4) does not intersect with [7, 12), as 4 < 7. • Since the interval is completely before toBeRemoved, add it to ans: ans = [[1, 4)].

Next, we take the interval [6, 8).

• The interval [6, 8) does intersect with [7, 12) since the interval starts before and ends in the range of toBeRemoved.

• The start of the interval 6 is before the start of toBeRemoved 7.

Step 5: Handle Slicing the Interval

Next, we take the interval [10, 13).

• Add the segment [6, 7) to ans: ans = [[1, 4), [6, 7)].

Step 5: Handle Slicing the Interval

13)].

Step 6: Continue the Loop

Step 7: Return the ans List

No more intervals to process.

Solution Implementation

updated_intervals = []

return updated_intervals

• The final result is ans = [[1, 4), [6, 7), [12, 13)].

The solution approach has efficiently handled the example intervals list by considering the toBeRemoved interval and has produced a result that correctly represents the set after removal.

This will store the final list of intervals after removing the specified interval

If the current interval doesn't overlap with the interval to be removed,

updated_intervals.append([interval_start, removal_start])

updated_intervals.append([removal_end, interval_end])

public List<List<Integer>> removeInterval(int[][] intervals, int[] toBeRemoved) {

// Function to remove the interval `toBeRemoved` from the list of `intervals`

// Check if part of the interval is before toBeRemoved

// Check if part of the interval is after toBeRemoved

updatedIntervals.push([removeEnd, end]);

// Return the final list of intervals after removal

removal_start, removal_end = toBeRemoved

for interval start, interval end in intervals:

we can add it to the updated list as-is

if interval start < removal start:</pre>

Return the updated list of intervals after removal

if interval end > removal end:

// Add the part of the interval after toBeRemoved

Extracting start and end points of the interval to be removed

Iterate through each interval in the provided list of intervals

def removeInterval(self. intervals: List[List[int]]. toBeRemoved: List[int]) -> List[List[int]]:

This will store the final list of intervals after removing the specified interval

If the current interval doesn't overlap with the interval to be removed,

if interval start >= removal end or interval end <= removal start:</pre>

If there is an overlap and the start of the current interval

updated_intervals.append([interval_start, removal_start])

updated_intervals.append([removal_end, interval_end])

Similarly, if the end of the current interval is after the end of

the interval to be removed, add the non-overlapping part to the result.

updated_intervals.append([interval_start, interval_end])

is before the start of the interval to be removed,

add the non-overlapping part to the result.

if (end > removeEnd) {

return updatedIntervals;

updated_intervals = []

else:

from typing import List

class Solution:

int removeStart = toBeRemoved[0], removeEnd = toBeRemoved[1];

if (start >= removeEnd || end <= removeStart) {</pre>

updatedIntervals.push_back(interval);

if (start < removeStart) {</pre>

// Iterate through all intervals

} else {

for (auto& interval : intervals) {

vector<vector<int>> removeInterval(vector<vector<int>>& intervals, vector<int>& toBeRemoved) {

// toBeRemoved[0] is the start of the interval to be removed, toBeRemoved[1] is the end

// Check if the current interval is completely outside the toBeRemoved interval

// Add interval to the result as it doesn't overlap with toBeRemoved

vector<vector<int>> updatedIntervals; // This will store the final intervals after removal

int start = interval[0], end = interval[1]; // Start and end of the current interval

Similarly, if the end of the current interval is after the end of

the interval to be removed, add the non-overlapping part to the result.

Iterate through each interval in the provided list of intervals

add the non-overlapping part to the result.

for interval start, interval end in intervals:

we can add it to the updated list as-is

if interval start < removal start:</pre>

Return the updated list of intervals after removal

// Function to remove a specific interval from a list of intervals

if interval end > removal end:

• The interval [10, 13) does intersect with [7, 12), because the interval starts inside and ends after the range of toBeRemoved.

• Since the end of the interval 13 is after the end of toBeRemoved 12, we add the segment [12, 13) to ans: ans = [[1, 4], [6, 7], [12,

def removeInterval(self, intervals: List[List[int]], toBeRemoved: List[int]) -> List[List[int]]: # Extracting start and end points of the interval to be removed removal_start, removal_end = toBeRemoved

class Solution:

from typing import List

Python

Java

class Solution {

if interval start >= removal end or interval end <= removal start:</pre> updated_intervals.append([interval_start, interval_end]) else: # If there is an overlap and the start of the current interval # is before the start of the interval to be removed,

```
// x and v represents the start and end of the interval to be removed
int removeStart = toBeRemoved[0];
int removeEnd = toBeRemoved[1];
// Preparing a list to store the resulting intervals after removal
List<List<Integer>> updatedIntervals = new ArrayList<>();
// Iterate through each interval in the input intervals array
for (int[] interval : intervals) {
    // a and b represents the start and end of the current interval
    int start = interval[0];
    int end = interval[1];
    // Check if the current interval is completely before or after the interval to be removed
    if (start >= removeEnd || end <= removeStart) {</pre>
        // Add to the result as there is no overlap
        updatedIntervals.add(Arrays.asList(start, end));
    } else {
        // If there's an overlap, we may need to add the non-overlapping parts of the interval
        if (start < removeStart) {</pre>
            // Add the part of the interval before the interval to be removed
            updatedIntervals.add(Arrays.asList(start, removeStart));
        if (end > removeEnd) {
            // Add the part of the interval after the interval to be removed
            updatedIntervals.add(Arrays.asList(removeEnd, end));
// Return the list of updated intervals
return updatedIntervals;
```

C++

public:

class Solution {

```
// Add the part of interval before toBeRemoved
                    updatedIntervals.push_back({start, removeStart});
                // Check if part of the interval is after toBeRemoved
                if (end > removeEnd) {
                    // Add the part of interval after toBeRemoved
                    updatedIntervals.push_back({removeEnd, end});
        // Return the final list of intervals after removal
        return updatedIntervals;
};
TypeScript
// Define the interval type as a tuple of two numbers
type Interval = [number, number];
// Function to remove the interval `toBeRemoved` from the list of `intervals`
function removeInterval(intervals: Interval[], toBeRemoved: Interval): Interval[] {
    // `toBeRemoved[0]` is the start of the interval to be removed, `toBeRemoved[1]` is the end
    const removeStart = toBeRemoved[0];
    const removeEnd = toBeRemoved[1];
    // This will store the final intervals after removal
    const updatedIntervals: Interval[] = [];
    // Iterate through all intervals
    for (const interval of intervals) {
        // `start` and `end` of the current interval
        const start = interval[0];
        const end = interval[1];
        // Check if the current interval is completely outside the toBeRemoved interval
        if (start >= removeEnd || end <= removeStart) {</pre>
            // Add the interval to the result as it doesn't overlap with toBeRemoved
            updatedIntervals.push(interval);
        } else {
            // Check if part of the interval is before toBeRemoved
            if (start < removeStart) {</pre>
                // Add the part of the interval before toBeRemoved
                updatedIntervals.push([start, removeStart]);
```

```
return updated_intervals
Time and Space Complexity
```

The primary operation in this function occurs within a single loop that iterates over all the original intervals in the list intervals. Within each iteration of the loop, the function performs constant-time checks and operations to possibly add up to two intervals to the ans list. Since there are no nested loops and the operations inside the loop are of constant time complexity, the overall

time complexity of the function is directly proportional to the number of intervals n in the input list. Therefore, the time

The code snippet provided is for a function that removes an interval from a list of existing intervals and returns the resulting list of

disjoint intervals after the removal. The computational complexity analysis for time and space complexity is as follows:

Space complexity: For space complexity, the function creates a new list ans to store the resulting intervals after the potential removal and

complexity is O(n).

Time complexity:

modification of the existing intervals. In the worst-case scenario, where no interval is completely removed and every interval needs to be split into two parts (one occurring before x and one after y of the toBeRemoved interval), the resulting list could potentially hold up to 2n intervals - doubling the input size. However, notice that this is a linear relationship with respect to the number of input intervals n. Therefore, the space complexity of the function is O(n) as well. In summary, both the time complexity and space complexity of the given code are O(n), where n is the number of intervals in the input list intervals.