

**Problem Description** 

Hash Table Counting

In this problem, you are given an integer array called nums. A number x is considered lonely if it meets two conditions:

1. It appears only once in the array (nums).

2. Neither x + 1 (the next consecutive number) nor x - 1 (the previous consecutive number) appear in the array. Your objective is to find all the lonely numbers within the nums array and return them. The order in which you return these lonely

numbers does not matter, meaning they can be in any sequence.

### The intuition behind the solution to this problem is to determine the occurrence of each number as well as the occurrence of its

Intuition

because it enables us to count the occurrences of each number in nums with ease. Once we have the counts of each number, we can iterate through the items in the counter, which consist of the number and its count. During the iteration, we apply the conditions that define loneliness:

immediate neighbors (i.e., num - 1 and num + 1). To efficiently manage this, a Counter from Python's collections module is useful

 The count of the number num should be exactly 1 (it appears only once). • The count of num - 1 should be 0 (no adjacent smaller number). The count of num + 1 should be 0 (no adjacent larger number).

- If a number satisfies all three conditions, it is a lonely number and is added to the result list (ans). After the iteration is complete, the list ans is returned as the list of all lonely numbers found in the array.
- **Solution Approach** The solution uses a hash table, specifically the Counter class from Python's collections module, to keep track of the frequency with

which each number appears in the nums array. This is a common pattern when dealing with problems that require you to track the

occurrences of elements because Counter provides an efficient way to store and query frequencies.

## Here's a step-by-step explanation of the implementation:

the values are the counts of each number. 1 counter = Counter(nums)

1. Initialize the Counter with the nums array. The Counter will create a hash table where the keys are the numbers from the array and

2. Create an empty list ans that will eventually contain all the lonely numbers. 1 ans = []

- 3. Iterate over the items in the counter. For each iteration, you have a num which is the number from the array and cnt which is the
- count of how many times that number appears. 1 for num, cnt in counter.items():
- Check if the current number num appears only once (cnt == 1). Check if the number immediately smaller than the current number (num - 1) does not appear in the array (counter [num - 1] == 0).

4. Inside the loop, you apply three checks for each number based on the definition of a lonely number:

5. If all three checks pass, it means that num is a lonely number, and you append it to the ans list.

6. After the loop finishes, the ans list contains all the lonely numbers and the function returns this list.

Let's consider an example to illustrate the solution approach. Suppose we have the following nums array:

• Check if the number immediately larger than the current number (num + 1) does not appear in the array (counter [num + 1] == 0).

1 return ans

the number of elements in nums.

1 nums = [4, 10, 5, 8, 20, 15, 11, 10]

Example Walkthrough

- 1 if cnt == 1 and counter[num 1] == 0 and counter[num + 1] == 0: ans.append(num)
- By using Counter, we achieve a solution that is simple and efficient with respect to time complexity since we look up the counts in constant time O(1) and we only iterate through the elements of the array once, making the overall time complexity O(N), where N is
- Follow these steps to find all the lonely numbers: 1. Initialize the Counter with the nums array:

### 1 ans = []

lonely numbers:

1 ans.append(20)

1 return ans

3. Iterate over the items in the counter:

In the first iteration, num is 4 and cnt is 1.

2 # First iteration: num = 4, cnt = 1

# Second iteration: num = 10, cnt = 2

1 for num, cnt in counter.items():

# and so on...

4. Apply the checks for loneliness:

1 counter = Counter([4, 10, 5, 8, 20, 15, 11, 10])

2. Create an empty list ans to hold the lonely numbers:

2 # counter now looks like: {4:1, 10:2, 5:1, 8:1, 20:1, 15:1, 11:1}

- Skipping to the next number meeting first condition which is 5. Check if 5 appears only once: Yes (cnt == 1).
- For num = 20, it appears only once, and neither 19 nor 21 appear in the array. So, 20 is lonely. ■ For num = 15, it appears only once, and neither 14 nor 16 appear in the array. So, 15 is lonely.

Check if 4 appears only once: Yes (cnt == 1).

Check if 3 appears in the array: No (counter[3] == 0).

Skipping ahead to numbers that meet the first condition (20 and 15).

Continue this process for the rest of the numbers.

5. After applying the loneliness checks and iterating through all the counter's items, find that 20 and 15 meet the criteria for being

The final output for our example is [20, 15], which means 20 and 15 are the lonely numbers in the given nums array. They appear only

■ Check if 5 appears in the array: Yes (counter[5]!= 0), so 4 is not lonely since its immediate larger neighbor exists.

• Check if 4 appears in the array: Yes (counter[4] != 0), so 5 is not lonely since its immediate smaller neighbor exists.

2 ans.append(15) 6. Finally, return the ans list, which contains all the lonely numbers found in nums:

once, and neither their immediate smaller nor larger neighbors appear in the array.

**Python Solution** 

def findLonely(self, nums: List[int]) -> List[int]:

# Iterate through the items in the counter

lonely\_numbers.append(num)

for num, count in num\_counter.items():

# Create a counter for all numbers in the list

# Initialize an empty list to store lonely numbers

2 # The final returned ans list is: [20, 15]

from collections import Counter

num\_counter = Counter(nums)

lonely\_numbers = []

from typing import List

class Solution:

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36 };

# Return the list of lonely numbers 18 return lonely\_numbers 19 20

# Check if the number appears only once and neither (num-1) nor (num+1) appear

# If conditions are met, append the number to the lonely\_numbers list

if count == 1 and num\_counter[num - 1] == 0 and num\_counter[num + 1] == 0:

// Create a HashMap to store the frequency of each number in the array

frequencyMap.put(num, frequencyMap.getOrDefault(num, 0) + 1);

// A number is lonely if it appears exactly once and neither of its

// adjacent numbers (number - 1, number + 1) are present in the array.

Map<Integer, Integer> frequencyMap = new HashMap<>();

// Initialize a list to store the lonely numbers

List<Integer> lonelyNumbers = new ArrayList<>();

// Iterate through the map to find lonely numbers

lonelyNumbers.push\_back(number);

// Initialize a hashMap to store the frequency of each number in the array

// Return the list of 'lonely' numbers.

let frequencyMap: Map<number, number> = new Map();

// Initialize an array to store lonely numbers

// Iterate through the entries of the frequency map

let lonelyNumbers: Array<number> = [];

return lonelyNumbers;

1 function findLonely(nums: number[]): number[] {

frequencyMap.forEach((number, count) -> {

// Iterate through the array and populate the frequency map

if (count == 1 && !frequencyMap.containsKey(number - 1)

&& !frequencyMap.containsKey(number + 1)) {

#### class Solution { // This method finds all elements in the array that stand alone without // any adjacent or duplicate elements public List<Integer> findLonely(int[] nums) { 9

**Java Solution** 

1 import java.util.Map;

2 import java.util.HashMap;

import java.util.List;

import java.util.ArrayList;

for (int num : nums) {

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27
                   lonelyNumbers.add(number);
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           });
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30
           // Return the list of lonely numbers
31
32
           return lonelyNumbers;
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34 }
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C++ Solution
1 #include <vector>
2 #include <unordered_map>
   using namespace std;
  class Solution {
  public:
       // Function to find 'lonely' integers in the array.
       // An integer is 'lonely' if its count is one and no adjacent integers (num - 1, num + 1) are present.
       vector<int> findLonely(vector<int>& nums) {
9
           // Create a hash map to keep track of the frequency of each number in the array.
10
           unordered_map<int, int> frequencyMap;
11
           // Iterate over the array and increment the count for each number in the frequencyMap.
           for (int num : nums) {
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               ++frequencyMap[num];
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16
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           // Prepare the answer vector to store the 'lonely' numbers.
           vector<int> lonelyNumbers;
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           // Iterate through the elements of frequencyMap
           for (const auto& element : frequencyMap) -
               int number = element.first;
                                                 // The number itself
               int count = element.second;
                                                 // Frequency of the number
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```

// Check if the count is 1 (unique number) and both adjacent numbers do not exist.

// If conditions are satisfied, add number to the 'lonely' numbers list.

if (count == 1 && !frequencyMap.count(number - 1) && !frequencyMap.count(number + 1)) {

#### // Populate the frequency map with counts of each number for (let num of nums) { frequencyMap.set(num, (frequencyMap.get(num) || 0) + 1);

Typescript Solution

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for (let [num, count] of frequencyMap.entries()) {
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          // Check if the current number is lonely, i.e., count is 1 and neither num-1 nor num+1 exist in the map
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          // If conditions are met, add the number to the lonelyNumbers array
           if (count === 1 && !frequencyMap.has(num - 1) && !frequencyMap.has(num + 1)) {
               lonelyNumbers.push(num);
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       // Return the array containing all lonely numbers
23
       return lonelyNumbers;
24 }
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Time and Space Complexity
Time Complexity
The time complexity of the code is primarily determined by three factors: the creation of the counter which counts the occurrences
of each element, the iteration through the counter dictionary, and the checks performed for each number.
 1. Creating the counter from the nums list takes O(n) time where n is the number of elements in the nums list because it requires one
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# pass through all elements.

2. Iterating through the counter dictionary's items also takes O(n) in the worst case, which is when all elements in nums are unique. We need to clarify this is "worst case" because the counter may have fewer than n keys if there are duplicate values in nums.

- 3. For each element in the counter, we check if the cnt is 1 and if the adjacent numbers num 1 and num + 1 are not present. These checks are constant-time operations, 0(1), because dictionary lookup is an 0(1) operation on average due to hash table
- implementation. Combining these factors, the overall time complexity is O(n) + O(n) \* O(1), which simplifies to O(n).

The space complexity involves the space taken up by the counter dictionary and the ans list.

**Space Complexity** 

1. The counter dictionary will hold at most n key-value pairs if all elements in nums are unique, resulting in a space complexity of

- O(n). 2. The ans list can also hold at most n elements in the worst case, where every element is lonely. Therefore, the space complexity
- for the ans list is also O(n). Therefore, the combined space complexity of the algorithm is also O(n), where n is the number of elements in the nums list.