529. Minesweeper **Breadth-First Search** Matrix Medium Depth-First Search Array **Leetcode Link**

Problem Description

triggering a mine. When a player clicks a square: 1. If a mine is revealed ('M'), the game ends and the mine becomes an 'X'.

The Minesweeper game consists of a board with mines ('M') and empty squares ('E'). The objective is to reveal all squares without

- 2. If an empty square ('E') without adjacent mines is revealed, it becomes a 'B', representing a blank space, and its adjacent unrevealed squares should also be revealed recursively. 3. If an empty square ('E') with adjacent mines is revealed, it shows a digit ('1' to '8') representing the number of mines in the
- adjacent squares. 4. The board state is returned when no more squares can be revealed.
- The task is to implement the logic that simulates this part of Minesweeper game behavior.

Intuition

To solve this problem, we use Depth-First Search (DFS) algorithm, where we start from the clicked position and explore adjacent squares. Here's a step-by-step breakdown:

1. Check the square at the clicked position: If it's a mine ('M'), change it to an 'X'.

- If it's an empty square ('E'), perform the following: Count the number of mines in the adjacent squares.
 - If there are adjacent mines, update the square with the mine count.
- If there are no adjacent mines, update the square to a 'B' and recursively reveal adjacent 'E' squares. 2. The base case for recursive DFS is when the current square is not an 'E' or when it is out of bounds, in which case we just return.
- 3. By following the steps above, we ensure that we reveal only the necessary squares to either show the number of adjacent mines or to expand the empty regions ('E' to 'B').
- numbered squares or the edge of the board.

4. The recursion naturally takes care of revealing all connected empty squares ('E' to 'B') and stopping when it encounters

- This DFS approach encapsulates the essence of Minesweeper's revealing mechanic, carefully unveiling the board according to the game's rules.
- Solution Approach

implementation steps, corresponding to the given Python code: 1. The nested dfs function is defined to perform a DFS starting from a particular cell on the board.

The implementation of the solution is a straightforward application of Depth-First Search (DFS). Here are the details of the

• First, we count the number of mines ('M') around the current square (i, j). This is done in a nested loop where we iterate

initial click.

Example Walkthrough

['E', 'E', 'E', 'M'], ['E', 'M', 'E', 'E'],

['E', 'E', 'M', 'E']]

1. Initialize – We note the player's click on cell (1, 2), which is 'E'.

2. The main logic for the DFS method is as follows:

over all adjacent cells including diagonals.

 \circ If there is at least one mine adjacent to (i, j), we update the current square with the number of mines (cnt). o If there are no adjacent mines, we update the current square to 'B'. We then recursively call dfs on all adjacent unrevealed squares ('E') to continue revealing the board.

3. Before DFS is initiated, we first check the cell at the click coordinates. If the clicked cell is a mine ('M'), we update it to 'X' to

indicate the game is over. Otherwise, we call our dfs function with the click coordinates. 4. The dfs function is designed to stop the recursive calls when it encounters a cell that is not 'E', which handles both the case

when revealing 'B' squares, stopping at the border, and stopping at numbered squares.

- 5. The dfs function also ensures that we never go out of bounds of the board by checking whether the coordinates (x, y) are within the range [0, m) for rows and [0, n) for columns, where m and n are the lengths of the board's rows and columns respectively.
- This implementation ensures that all recursive calls to reveal cells only act on valid, in-bound, and appropriate cells, preventing unnecessary work and ensuring that the board state is mutated correctly according to Minesweeper rules.

6. The updateBoard function then returns the updated board once all the possible and necessary cells are revealed based on the

Let's go through a small example to illustrate the solution approach. Suppose we have a Minesweeper board like this, and the player clicks on the cell at row 1, column 2 (indexes are zero-based): [['E', 'E', 'E', 'E'],

We will follow the solution steps for our DFS approach using this board:

4. Update Cell - Since there's one mine adjacent to (1, 2), we update the board as follows, replacing 'E' with '1' to indicate the

2. DFS Function Call – We check whether the clicked cell is 'M' or 'E'. Since it's 'E', we invoke our dfs function here.

3. Count Mines – We look at all the adjacent cells around (1, 2). The adjacent cells contain one mine at (2, 1). Therefore, the mine

number of adjacent mines:

from typing import List

def reveal(i: int, j: int):

board[i][j] = "B"

// Iterate through the adjacent cells

gameBoard[row][col] = 'B';

if (mineCount > 0) {

} else {

for (int x = row - 1; $x \le row + 1$; ++x) {

for (int y = col - 1; $y \le col + 1$; ++y) {

gameBoard[row][col] = (char) (mineCount + '0');

for (int y = col - 1; $y \le col + 1$; ++y) {

dfs(x, y); // Recursive call to dfs

for (int x = row - 1; $x \le row + 1$; ++x) {

int rows = board.size(), cols = board[0].size();

// Using lambda function to perform Depth-First Search (DFS)

function<void(int, int)> dfs = [&](int row, int col) {

int clickRow = click[0], clickCol = click[1];

mineCount++; // Increment the mine counter

// If there are mines found around the cell, display the mine count

vector<vector<char>> updateBoard(vector<vector<char>>& board, vector<int>& click) {

// If there are no mines around the cell, mark the cell as 'B' for blank

// Recursively reveal adjacent cells that are not mines and are unrevealed ('E')

if $(x >= 0 \&\& x < rows \&\& y >= 0 \&\& y < cols && gameBoard[x][y] == 'E') {$

mine_count = 0

The clicked position

Count mines around the current cell

if board[x][y] == "M":

class Solution:

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[['E', 'E', 'E', 'E'],

count cnt is 1.

['E', 'E', 'M', 'E']]

5. Adjacent Squares – Since cell (1, 2) is next to only one mine, we don't perform recursive calls on its neighbors as per standard Minesweeper rules.

6. Final Board - The updateBoard function returns the board since we cannot reveal any further cells based on the click position.

This updated board represents the state after the single click, adhering to the rules of Minesweeper where numbered cells should

not reveal their neighbors. If the initial click was on a different 'E' cell, especially one not adjacent to any mines, the recursive dfs

would take effect, revealing a larger area by turning 'E' cells to 'B' and uncovering all interconnected blank spaces.

for x in range(max(i - 1, 0), min(i + 2, rows)): # Limits the range on the board

def updateBoard(self, board: List[List[str]], click: List[int]) -> List[List[str]]:

for x in range(max(i - 1, 0), min(i + 2, rows)):

if board[x][y] == "E":

reveal(x, y)

for y in range(max(j - 1, 0), min(j + 2, columns)):

The recursive function to reveal the board starting from the clicked cell.

Python Solution

12 mine_count += 1 13 # If there are mines around the cell, update with mine count 14 15 if mine_count > 0: board[i][j] = str(mine_count) 16 17 else: # Otherwise, set the cell to "B" for blank and reveal surrounding cells 18

for y in range(max(j - 1, 0), min(j + 2, columns)): # Limits the range on the board

23 24 25 # Get the size of the board 26 rows, columns = len(board), len(board[0])

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click_row, click_col = click
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             # If the clicked cell contains a mine, game over
             if board[click_row][click_col] == "M":
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                 board[click_row][click_col] = "X"
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             else:
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                 # Start revealing from the clicked cell
 36
                 reveal(click_row, click_col)
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             # Return the updated board
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             return board
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Java Solution
  1 class Solution {
         private char[][] gameBoard; // Renamed board to gameBoard for clarity
                                    // Renamed m to rows for clarity
         private int rows;
         private int cols;
                                    // Renamed n to cols for clarity
  4
  5
         // Function updates the game board when a user clicks on a cell (i.e., uncovers the cell)
  6
         public char[][] updateBoard(char[][] board, int[] click) {
             rows = board.length;
                                       // Total number of rows in the board
             cols = board[0].length;
                                       // Total number of columns in the board
  9
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             gameBoard = board;
                                        // Assign board to gameBoard for internal use
 11
 12
             int clickRow = click[0], clickCol = click[1]; // Row and column indexes of the click position
 13
 14
             // Check if the clicked cell contains a mine
 15
             if (gameBoard[clickRow][clickCol] == 'M') {
                 gameBoard[clickRow][clickCol] = 'X'; // If there's a mine, mark the cell as 'X'
 16
 17
             } else {
 18
                 // If the clicked cell does not contain a mine, perform Depth First Search (DFS) from this cell
 19
                 dfs(clickRow, clickCol);
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 21
             return gameBoard; // Return the updated board
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         // Performs Depth First Search (DFS) to reveal cells
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         private void dfs(int row, int col) {
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             int mineCount = 0; // Counter for adjacent mines
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// Check if the adjacent cell is within the board and if it contains a mine

if $(x >= 0 \&\& x < rows \&\& y >= 0 \&\& y < cols \&\& gameBoard[x][y] == 'M') {$

55 C++ Solution

public:

class Solution {

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```
int mineCount = 0;
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                 // Count mines in adjacent cells
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                 for (int x = row - 1; x \le row + 1; ++x) {
                     for (int y = col - 1; y \le col + 1; ++y) {
 13
 14
                         if (x >= 0 \&\& x < rows \&\& y >= 0 \&\& y < cols && board[x][y] == 'M') {
 15
                             ++mineCount;
 16
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                 // If mines exist around the cell, update the cell with the mine count. Otherwise, mark the cell as 'B' and continue DF
 21
                 if (mineCount > 0) {
                     board[row][col] = mineCount + '0';
 22
                 } else {
 23
                     board[row][col] = 'B';
 24
                     for (int x = row - 1; x \le row + 1; ++x) {
 25
 26
                         for (int y = col - 1; y \le col + 1; ++y) {
 27
                             if (x >= 0 \&\& x < rows \&\& y >= 0 \&\& y < cols && board[x][y] == 'E') {
 28
                                 dfs(x, y);
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             };
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 35
             // Handle the initial click
 36
             if (board[clickRow][clickCol] == 'M') {
                 // If the clicked cell contains a mine, game over
 38
                 board[clickRow][clickCol] = 'X';
             } else {
 39
 40
                 // If the clicked cell is empty ('E'), perform DFS to reveal cells
                 dfs(clickRow, clickCol);
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             return board;
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 46 };
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Typescript Solution
    function updateBoard(board: string[][], click: number[]): string[][] {
         const rowCount = board.length; // Number of rows in the board
         const colCount = board[0].length; // Number of columns in the board
         const [rowClicked, colClicked] = click; // Destructure click into row and column indices
         // Deep-First Search function to reveal the board
         const revealBoard = (row: number, col: number) => {
             let mineCount = 0;
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             // Count the number of mines surrounding the current cell
  9
             for (let x = row - 1; x \le row + 1; ++x) {
 10
                 for (let y = col - 1; y \le col + 1; ++y) {
 11
                     if (x >= 0 \&\& x < rowCount \&\& y >= 0 \&\& y < colCount \&\& board[x][y] === 'M') {
 12
 13
                         mineCount++;
 14
 15
 16
 17
 18
             if (mineCount > 0) { // If mines are found around the cell, update the cell with the mine count
                 board[row][col] = mineCount.toString();
 19
 20
             } else { // If no mines are found, mark the cell as 'B' and continue the search
 21
                 board[row][col] = 'B';
```

return board; // Return the updated board state 42 43 44

Time and Space Complexity

// Click processing

for (let x = row - 1; $x \le row + 1$; ++x) {

revealBoard(x, y);

if (board[rowClicked][colClicked] === 'M') {

board[rowClicked][colClicked] = 'X';

revealBoard(rowClicked, colClicked);

for (let y = col - 1; y <= col + 1; ++y) {

// Recursively reveal the board for surrounding cells

// If the click is on a mine, game is over. Mark the clicked mine with 'X'

// If the click is not on a mine, reveal the board from the clicked spot

scenario, dfs performs a Depth-First Search on the entire board starting from the clicked cell.

The space complexity is hence 0(m * n) due to the recursion stack in the worst-case scenario.

if $(x >= 0 \&\& x < rowCount \&\& y >= 0 \&\& y < colCount \&\& board[x][y] === 'E') {$

The given Python code is a solution for the Minesweeper game where a player clicks on a cell on the board and depending on the board cell content, different actions may be taken.

The time complexity is determined by the number of times dfs is invoked and how many cells it inspects each time. In the worst-case

• Each cell is visited at most once due to the fact that once a cell is visited it changes from "E" (empty) to "B" (blank) or it becomes a digit representing the number of adjacent mines.

once.

Time Complexity:

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};

} else {

 For each cell visited, dfs inspects up to 8 adjacent cells. Given a board of size $m \times n$, the time complexity is thus $0(m \times n)$ since in the worst-case scenario, we might end up visiting each cell

- **Space Complexity:** The space complexity of this DFS solution is primarily determined by the recursion stack used by the dfs function.
- In the worst-case scenario, the DFS could recurse to a depth where the entire board is traversed, making the maximum recursion depth m * n in a situation where all cells are empty ("E") and connected, requiring the algorithm to visit every cell before hitting a base case.

When looking at auxiliary space (excluding space for inputs and outputs), the space complexity depends on the number of recursive calls placed on the stack, which remains 0(m * n).