Problem Description

the queries array match the pattern. We say that a query string matches the pattern if it's possible to insert lowercase English letters into the pattern such that after insertion, the pattern is equal to the query. Importantly, these insertions can happen at any index within the pattern, but we cannot add any new characters into the query string. The output of the problem is a boolean array answer where each element answer[i] is true if queries[i] matches the pattern, otherwise false. For example, if the query is "FoBar" and the pattern is "FB", we can see that by adding "o" between "F" and "B" and appending "ar"

The problem gives us an array of strings called queries and a single string called pattern. Our goal is to determine which strings in

after "B", the pattern matches the query (after the insertions, the pattern would be "FoBar"). Therefore, the answer for this query would be true. Intuition

queries[i] and pattern one by one. By using an iterative process, we can validate if we can morph the pattern into each query. The intuition for the solution is based on a couple of observations:

forward.

1. The characters in both queries[i] and pattern must match in order, and the matching characters in queries[i] must be in the same positions as in pattern. Uppercase letters in the query are "anchors" that must be present and in the correct order in the

The solution involves defining a match-checking function that uses two pointers to traverse and compare characters of the

pattern for a match to be possible. 2. Lowercase letters in queries[i] can be skipped because we're allowed to insert lowercase letters into the pattern. However,

- uppercase letters can't be skipped or inserted because they define the structure of the pattern that must be followed. This allows us to implement a pointer-based approach:
- Iterate over queries[i] with pointer i, and for each character, check if it matches the jth character of pattern. If the characters match, move both pointers forward.

• If the characters do not match, and the ith character of queries[i] is lowercase, we can attempt to skip it by moving pointer i

• If we encounter an uppercase letter that doesn't match or we reach the end of queries[i] without fully traversing the pattern,

Initialize two pointers, i and j, to traverse queries[i] and pattern, respectively.

- that query does not match the pattern. After pointer j has traversed the pattern, we need to ensure that all remaining characters in queries[i] are lowercase;
- otherwise, it means there are extra uppercase letters that make the query not match the pattern.
- desired boolean output array. The time complexity is linear with respect to the total number of characters in all queries and the length of the pattern, i.e.,

The match-checking function embodied the above logic, consistently applying it to every query in the array, thus creating the

effectively O(n*m), where n is the number of queries and m is the length of the longest query string.

The solution to the given problem involves creating a helper function, check(s, t), which is designed to determine if the string s (a query) matches the string t (the pattern). This function is crucial as it encapsulates the logic for matching based on the rules outlined in the problem description.

insertion of lowercase letters. o If s[i] matches t[j], increment both pointers i and j to check the next characters since a match dictates we move in both

queries and pattern.

2. Iterate through s using the pointer i:

Helper Function: check(s, t)

of the check function implementation using the s and t variables:

Solution Approach

o If the characters do not match and s[i] is not a lowercase letter, return false as we can't insert uppercase letters or skip non-matching uppercase letters. 3. Once j has reached the end of t, we check the remainder of s to ensure all leftover characters are lowercase. If we encounter an

uppercase letter or i hasn't yet reached the end of s, we return false as the additional characters can't be part of the pattern.

This function uses the two-pointer technique, which is a common pattern used in string manipulation algorithms. Here's a breakdown

o If s[i] is a lowercase letter that does not match t[j], increment i to skip it. This step embodies the rule that allows the

Main Logic

4. If i has reached the end of s, we return true, indicating a successful match.

case scenario, consuming linear time relative to the length of each individual query.

1. Use a list comprehension to iterate over all query strings in queries.

1. Initialize two pointers: i starts at the beginning of s while j starts at the beginning of t.

2. For each query, call the check function with the query and the pattern. 3. Convert the result of the check function into boolean values and store them in the final output list.

length of the longest query string. This is because we must potentially traverse each character in each query string in the worst-

In summary, this approach is simple yet effective. It leverages the two-pointer pattern to navigate two strings simultaneously and

verifies compliance with the pattern-matching rules. The algorithm is sharp in that it halts and decides as soon as a non-compliance

The main part of the solution consists of iterating over each string in the queries array and applying the check function:

Time Complexity Analysis The overall time complexity of the solution is O(n * m), where n is the number of query strings in the queries array, and m is the

Let's walk through an example to illustrate the solution approach described above. Suppose we have the following queries array and pattern:

is lowercase.

are lowercase.

1. "After":

2. "AFter":

Example Walkthrough

is detected, ensuring optimal performance.

1 queries = ["After", "AFter", "AFilter", "BFilter"]

∘ Initialize pointers i=0 and j=0. Since queries[i] = 'A' and pattern[j] = 'A', they match, and we increment both pointers.

the other is uppercase, but since queries [i] is lowercase, we can skip it by incrementing i.

Now i=1 and j=1, and we have queries[i] = 'f' and pattern[j] = 'F'. They do not match because one is lowercase and

With i=2 and j=1, we now have queries[i] = 't' and pattern[j] = 'F'. The letters do not match, so we increment i as 't'

• Now i=2, and since j has reached the end of pattern, we only need to check that the remaining characters in queries[i]

Since j has not moved and i reached the end of queries[i] without matching 'F' from pattern, we return false.

○ We set i=0 and j=0. They match ('A' with 'A'), so both pointers increment. At i=1, j=1, queries[i] = 'F' and pattern[j] = 'F', they match, so increment both pointers again.

○ At i=1, j=1, 'F' matches 'F'. Increment both again.

matching for each string in the queries array.

query_length = len(query)

pattern_length = len(pattern)

pattern_index = query_index = 0

Traverse the pattern characters.

query_index += 1

return False

query_index += 1

// Iterate over each query

for (String query : queries) {

pattern_index += 1

while pattern_index < pattern_length:</pre>

i=2 now points to 'i' which is lowercase, so we can skip it.

Skipping all lowercase 'l', 't', 'e', 'r' by incrementing i each time.

 As 't', 'e', and 'r' are all lowercase, we move i to the end of the string. Given that all conditions are satisfied, i has reached the end of s, and j has reached the end of t, we return true. 3. "AFilter":

 \circ Both pointers start at the beginning. i=0, j=0, and 'A' matches 'A'. Both increment.

To determine which query strings match the pattern, we apply the check function to each one:

○ At i=3, j=1, we have queries[i] = 'e', still lowercase. We skip it.

○ Now i=4, j=1, and queries[i] = 'r'. Another lowercase to skip.

 Having processed all characters of pattern and ensured no uppercase letters in the remaining queries[i], we return true. 4. "BFilter":

The resulting list based on the check function would be [false, true, true, false]. This methodically applies the rules from the

problem description to each query against the pattern. By utilizing this process, the check function can efficiently validate pattern

Start with i=0 and j=0. We find that 'B' doesn't match 'A' in pattern and it's uppercase, which means we cannot skip it or

Python Solution

Move to the next character in both query and pattern.

// Function to check the list of strings 'queries' match the camel case pattern 'pattern'

// Helper method to check if a single string 'query' matches the camel case 'pattern'

int queryIndex = 0, patternIndex = 0; // Indices to track position within query and pattern

public List<Boolean> camelMatch(String[] queries, String pattern) {

int queryLength = query.length(); // Length of query string

int patternLength = pattern.length(); // Length of the pattern

private boolean isMatch(String query, String pattern) {

++queryIndex;

return false;

return queryIndex == queryLength;

++queryIndex;

++queryIndex;

for (auto& query : queries) {

const queryLength = query.length;

let queryIndex = 0;

const patternLength = pattern.length;

++patternIndex;

// Check if query ended or characters don't match

if (queryIndex == queryLength || query[queryIndex] != pattern[patternIndex]) {

// Move to the next character in both the query and the pattern

while (queryIndex < queryLength && islower(query[queryIndex])) {</pre>

// If we reached the end of the query, all checks are passed

return answers; // Return a vector with the results for each query

// Iterate through all the queries to check against the pattern

answers.push_back(matchesPattern(query, pattern));

List<Boolean> matches = new ArrayList<>(); // List to hold the results

Skip lower case characters from query that are not a part of the pattern.

If query runs out of characters or does not match pattern character, return False.

if query_index == query_length or query[query_index] != pattern[pattern_index]:

def camelMatch(self, queries: List[str], pattern: str) -> List[bool]:

Helper function to check if the query matches the pattern.

def matches_pattern(query: str, pattern: str) -> bool:

All remaining characters in the query should be lowercase. while query_index < query_length and query[query_index].islower():</pre> query_index += 1 # True if the end of the query has been reached, False otherwise.

while query_index < query_length and query[query_index] != pattern[pattern_index] and query[query_index].islower():</pre>

change the pattern to match it. Immediately we return false, as it's not possible for this query to match the pattern.

class Solution:

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49 };

return query_index == query_length # Use list comprehension to apply the helper function to all queries in the input list. 31 return [matches_pattern(query, pattern) for query in queries]

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// Add the result of matching each query with the pattern to the matches list
               matches.add(isMatch(query, pattern));
10
           return matches;
11
```

Java Solution

1 class Solution {

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19
           // Iterate through the characters in 'pattern'
           while (patternIndex < patternLength) {</pre>
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21
               // Advance in 'query' while current character does not match the pattern character
               // and it's a lowercase letter
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23
               while (queryIndex < queryLength && query.charAt(queryIndex) != pattern.charAt(patternIndex)</pre>
                        && Character.isLowerCase(query.charAt(queryIndex))) {
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25
                    queryIndex++;
26
27
               // If the end of 'query' is reached or characters do not match, the query doesn't follow the pattern
28
               if (queryIndex == queryLength || query.charAt(queryIndex) != pattern.charAt(patternIndex)) {
29
                    return false;
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31
               // Move to the next character in both 'query' and 'pattern'
32
               queryIndex++;
33
                patternIndex++;
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36
           // After the end of 'pattern', all remaining characters in 'query' must be lowercase for it to be a match
37
           while (queryIndex < queryLength && Character.isLowerCase(query.charAt(queryIndex))) {</pre>
38
                queryIndex++;
39
40
           // Return true if the end of 'query' is reached, otherwise false
41
           return queryIndex == queryLength;
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C++ Solution
 1 #include <vector>
2 #include <string>
   using namespace std;
   class Solution {
   public:
       // Function to determine if each query matches the given pattern
       vector<bool> camelMatch(vector<string>& queries, string pattern) {
            vector<bool> answers;
           // Lambda function to check if a single query matches the pattern
10
            auto matchesPattern = [](const string& query, const string& pattern) {
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12
                int queryLength = query.size();
                int patternLength = pattern.size();
13
                int queryIndex = 0, patternIndex = 0;
14
15
               // Loop through each character in the pattern
16
               while (patternIndex < patternLength) {</pre>
17
                    // Advance through the query string to find the match for the current pattern character
18
```

while (queryIndex < queryLength && query[queryIndex] != pattern[patternIndex] && islower(query[queryIndex])) {</pre>

// All characters in the pattern are matched, check if the remaining characters in the query are all lowercase

1 // Function to match camel case patterns within a set of query strings function camelMatch(queries: string[], pattern: string): boolean[] { // Helper function to check if a single query matches the camel case pattern const checkPattern = (query: string, pattern: string): boolean => {

Typescript Solution

```
let patternIndex = 0;
           // Iterate through both query and pattern
           for (; patternIndex < patternLength; ++queryIndex, ++patternIndex) {</pre>
10
               // Skip lowercase characters in query that do not match the current pattern character
               while (queryIndex < queryLength &&</pre>
                      query[queryIndex] !== pattern[patternIndex] &&
                      query.charCodeAt(queryIndex) >= 97) { // ASCII 97 = 'a'
14
                   ++queryIndex;
15
16
               // Pattern character does not match, or query ended before pattern
17
               if (queryIndex === queryLength || query[queryIndex] !== pattern[patternIndex]) {
19
                   return false;
20
21
           // Skip all trailing lowercase characters in the query
           while (queryIndex < queryLength && query.charCodeAt(queryIndex) >= 97) {
23
               ++queryIndex;
           // If all characters of query are checked, it's a match
           return queryIndex == queryLength;
       };
28
29
       // Array to store results for each query
       const results: boolean[] = [];
       // Check each query and add the result to the results array
33
       for (const query of queries) {
           results.push(checkPattern(query, pattern));
34
35
       return results;
36
37 }
38
Time and Space Complexity
The given code implements a function to check if each query in a list of queries matches a given camel case pattern.
Time Complexity
```

• The check function goes through each character in the query string s and the pattern string t at most once. The inner while loop runs as long as there are lowercase characters in s that are not in t, and the outer while loop runs as long as there are characters in t.

Space Complexity

 In the worst case, each character of the query string will be checked exactly once against the pattern. This results in a linear complexity in terms of the length of the query string. The check function is called once for each query in the queries list.

The time complexity of the function is determined as follows:

If the average length of the queries is m and the length of the pattern is n, and there are k queries, then the time complexity is 0(k * (m + n)), since for every query we have a separate traversal through the pattern string and the query string.

• The check function uses a constant amount of extra space since it only uses a few integer variables to keep track of the indices.

• The output list has the same length as the number of queries. This is where the space is used. However, this is not counted towards the extra space complexity as this is required for the output of the function.

So, the space complexity of the function is 0(1) (constant space complexity) for the auxiliary space.

The space complexity of the function is considered as follows: