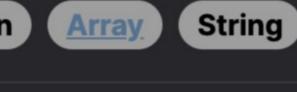


Bit Manipulation

Problem Description



Given an array of strings named words, the goal is to find the maximum product of the lengths of any two words in the array that do not have any letters in common. More formally, we want to find the maximum value of length(word[i]) * length(word[j]) where word[i] and word[j] are such that no letter appears in both words simultaneously. If there are no such pairs of words that satisfy this condition, the function should return 0.

Intuition

letters. However, this would require checking every pair which leads to a time complexity of O(n^2 * m), where n is the number of words and m is the average length of a word, rendering this approach inefficient for large input sizes. To optimize this, we can preprocess each word to create a bitmask that represents the set of characters of the word. In the bitmask,

The straightforward approach to solve this problem would be to compare each pair of words and check if they have any common

the i-th bit is set if the word contains the i-th letter of the alphabet. Since there are only 26 lowercase English letters, this bitmask can be represented by an integer. By comparing the bitmasks of two words, we can efficiently check if two words have common letters. If the result of a bitwise AND operation between two masks is zero, then the two words do not share any common letters. With the bitmask precomputation step, the time complexity is reduced significantly because the comparison of every pair now only

takes constant time, O(1), leading to an overall time complexity of O(n^2 + n * m). Here's a step-by-step breakdown of the solution approach: 1. Iterate through each word in the input, calculate its bitmask, and store these masks in an array.

2. Iterate through all pairs of words. For each pair, use the precomputed bitmasks to check if the words share common letters by performing a bitwise AND operation.

- 3. If the words do not share any common letters (bitwise AND result is 0), calculate the product of their lengths and update the
- answer with the maximum product found so far. 4. After checking all pairs, return the maximum product found.
- This bitmasking technique leverages the power of bitwise operations to significantly improve the time efficiency for checking if two words have common letters.

the bit position corresponding to each letter.

Solution Approach The solution uses a bit manipulation technique along with an array for storing the bitmasks. Here's a step-by-step breakdown of the

1. Initial Setup

implementation:

 Calculate the number n of words in the input list. Initialize an array mask of length n to store the bitmasks which will represent the unique characters of each word by setting

Iterate through the list of words with the index i. For each word word[i], do the following:

2. Creating Bitmasks

- For each character ch in word[i], shift 1 left by ord(ch) ord('a') bits to create a bitmask for the letter. ■ Use the bitwise OR | operation to update mask[i] by turning on the bit corresponding to each character in the word.
- 3. Finding the Maximum Product

Initialize mask[i] to 0.

 Initialize a variable ans to 0 for storing the maximum product found. Iterate through all unique pairs of words with indices i and j (with j > i to avoid duplication of pairs), to compare their

■ If the bitwise AND & of mask[i] and mask[j] is 0, it means the words have no common letters.

• If so, calculate the product of their lengths len(words[i]) * len(words[j]). Update ans with the maximum of itself and the calculated product.

bitmasks.

4. Returning the Result

Example Walkthrough

- After checking all pairs, return ans as the final result.
- The code uses bitmasks stored in an integer to efficiently represent and compare the characters of the words. Rather than checking
- each letter individually for every pair, the solution leverages bitwise operations to check for common letters in constant time, greatly
- enhancing performance.

Let's take an example array of strings words = ["abc", "de", "fg", "hi", "aeh"].

 The first step is to represent each word by a bitmask. Let's work through the words: For word "abc", the bitmask would be (set bit positions at 0, 1, 2): 0b111 which is 7 in decimal.

∘ For word "fg", the bitmask would be (set bit positions at 5, 6): 0b1100000 which is 96 in decimal. For word "hi", the bitmask would be (set bit positions at 7, 8): 0b110000000 which is 384 in decimal.

∘ For "aeh", the bitmask is (set bit positions at 0, 4, 7): 0b10010001 which is 145 in decimal. • As a result, the mask array after processing the words array would be [7, 24, 96, 384, 145].

• In the next step, we will look for pairs of words with no overlapping bits in their bitmasks:

Compare "de" (24) and "fg" (96): (24 & 96) equals 0; product is 2 * 2 = 4.

○ Compare "abc" (7) and "de" (24): (7 & 24) equals Ø which means no common letters. Product of lengths equals 3 * 2 = 6. Compare "abc" (7) and "fg" (96): (7 & 96) equals 0; product is 3 * 2 = 6.

∘ For word "de", the bitmask would be (set bit positions at 3, 4): 0b11000 which is 24 in decimal.

- Compare "abc" (7) and "hi" (384): (7 & 384) equals 0; product is 3 * 2 = 6. ○ Compare "abc" (7) and "aeh" (145): (7 & 145) is not 0 as they share letters, so do not calculate product.
- Compare "de" (24) and "hi" (384): (24 & 384) equals 0; product is 2 * 2 = 4.
 - Compare "de" (24) and "aeh" (145): (24 & 145) is not 0, so do not calculate product. Compare "fg" (96) and "hi" (384): (96 & 384) equals 0; product is 2 * 2 = 4.
 - Compare "fg" (96) and "aeh" (145): (96 & 145) equals 0; product of lengths is 2 * 3 = 6.
- Compare "hi" (384) and "aeh" (145): (384 & 145) equals 0; product is 2 * 3 = 6. The final step is to find the maximum product from these comparisons:

individually. The final solution would implement the process outlined in the "Solution Approach" section efficiently by leveraging

- The maximum product is 6. • Return the maximum product, which is 6. The words yielding this product are "abc" with "de", "abc" with "fg", "abc" with "hi", "fg"
- with "aeh", and "hi" with "aeh". This illustrates how the bitmask approach greatly simplifies the computation and avoids the overhead of checking each letter
- bitwise operations and precomputation of the bitmasks for each word. Python Solution

Create a list to store the bitmask representation of each word

Compare every pair of words to find the maximum product of lengths

of two words which have no characters in common (no common bits in the bitmask).

Generate a bitmask for each word where bit i is set if the

word contains the i-th letter of the alphabet

The calculated products are 6, 6, 6, 4, 4, 4, 6, 6.

class Solution: def max_product(self, words: List[str]) -> int: # Get the number of words in the list

12 for i, word in enumerate(words): for ch in word: 13 $masks[i] \mid = 1 \ll (ord(ch) - ord('a'))$ 14 15 # Initialize max_product to 0 max_product = 0

```
if masks[i] & masks[j] == 0: # No common characters
                       # Update max_product if this pair has a larger product
24
25
                       max_product = max(max_product, len(words[i]) * len(words[j]))
26
27
           # Return the maximum product found
```

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from typing import List

num_words = len(words)

 $masks = [0] * num_words$

for i in range(num_words - 1):

// Return the maximum product found

int maxProduct(vector<string>& words) {

int wordsCount = words.size(); // Count of the words in the vector

vector<int> masks(wordsCount); // Bitmasks for each word

return maxProduct;

return max_product

for j in range(i + 1, num_words):

```
Java Solution
   class Solution {
       public int maxProduct(String[] words) {
           // Get the length of the words array
           int length = words.length;
           // Create an array to store the bitmask representation of each word
           int[] bitMasks = new int[length];
           // Convert each word into a bitmask representation and store it in the bitMasks array
           for (int i = 0; i < length; ++i) {</pre>
               for (char c : words[i].toCharArray()) {
10
                   // Set the bit corresponding to the character 'c'
11
12
                   bitMasks[i] |= (1 << (c - 'a'));
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14
15
           // Initialize the maximum product to 0
16
           int maxProduct = 0;
17
18
           // Compare each pair of words to find the pair with the maximum product of lengths
19
20
           // where the words do not share any common characters
           for (int i = 0; i < length - 1; ++i) {
               for (int j = i + 1; j < length; ++j) {
22
                   // Check if the two words share any common characters using the '&' bitwise operator
24
                   if ((bitMasks[i] & bitMasks[j]) == 0) {
25
                       // Calculate the product of the lengths of the two words
26
                       int product = words[i].length() * words[j].length();
27
                       // Update maxProduct with the maximum value between the existing maxProduct and the current product
28
                       maxProduct = Math.max(maxProduct, product);
29
```

11 12 13 14

C++ Solution

1 #include <vector>

2 #include <string>

6 class Solution {

7 public:

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#include <algorithm>

using namespace std;

```
// Create bitmasks. Each bit in mask represents if a character ('a' to 'z') is in the word
           for (int i = 0; i < wordsCount; ++i) {</pre>
                for (char ch : words[i]) {
                    masks[i] |= 1 << (ch - 'a'); // Set the bit corresponding to the current character
15
16
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            int maxProduct = 0; // Initialize max product to be 0
           // Compare each pair of words
            for (int i = 0; i < wordsCount - 1; ++i) {
21
22
                for (int j = i + 1; j < wordsCount; ++j) {</pre>
23
                   // If two words have no common letters, their masks will not share any common bits
                   // The bitwise AND of their masks will be 0
24
                    if (!(masks[i] & masks[j])) {
25
                        // Update max product if this pair gives us a bigger product
26
27
                        maxProduct = max(maxProduct, (int)(words[i].size() * words[j].size()));
28
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32
           // Return the maximum product found
33
            return maxProduct;
34
35 };
36
Typescript Solution
   // Import statements are not necessary in plain TypeScript, and there are no direct equivalents for `vector` and `using namespace stc
   // Function to calculate the maximum product of lengths of two words that don't share common characters
   function maxProduct(words: string[]): number {
        const wordsCount: number = words.length; // Count of the words in the array
       const masks: number[] = new Array(wordsCount); // Bitmasks for each word
       // Initialize all masks to 0
 8
       masks.fill(0);
 9
10
       // Create bitmasks. Each bit in a mask represents if a character ('a' to 'z') is in the word
11
       for (let i = 0; i < wordsCount; ++i) {</pre>
            for (const char of words[i]) {
13
                masks[i] |= 1 << (char.charCodeAt(0) - 'a'.charCodeAt(0)); // Set the bit corresponding to the current character
14
```

maxProduct = Math.max(maxProduct, words[i].length * words[j].length); 27 28 29 30 31

return maxProduct;

// Compare each pair of words

for (let i = 0; i < wordsCount - 1; ++i) {

// Return the maximum product found

let maxProduct = 0; // Initialize max product to be 0

for (let j = i + 1; j < wordsCount; ++j) {</pre>

if ((masks[i] & masks[j]) === 0) {

characters. The analysis of time and space complexity is as follows:

// The bitwise AND of their masks will be 0

Time and Space Complexity The provided code calculates the maximum product of the lengths of two words such that the two words do not share any common

// If two words have no common letters, their masks will not share any common bits

// Update max product if this pair gives us a bigger product

bitmask. This process occurs in O(N * L) time, where N is the number of words, and L is the average length of the words.

Time Complexity

2. Comparing the bitmasks: The nested loop compares each pair of generated bitmasks. This results in 0(N^2) comparisons, since each word's bitmask is compared with every other word's bitmask.

1. Calculating the bitmask for each word: The first loop goes through each word and each character within those words to create a

- 3. Checking for common characters and updating the maximum product happens in constant time, 0(1), for each pair of words compared.
- Combining these steps, the overall time complexity is $0(N * L + N^2)$.

Since typically L <= 1000 and the main contributing factor for large inputs is N, and because N^2 grows faster than N * L, the N^2

Space Complexity

- 1. The space used to store the bitmasks: For N words, we store a bitmask for each, which uses O(N) space. 2. No additional significant space is used, since other variables use constant space.
- Hence, the space complexity of the code is O(N).

term dominates for large N, so the overall time complexity can be considered $O(N^2)$.