812. Largest Triangle Area

<u>Array</u>

Problem Description

Geometry

Easy

that can be formed by any three different points from the array. The resulting area should be approximated to within 10^-5 of the correct value to be considered acceptable. To solve this problem, we need to understand that a triangle can be formed by any three non-collinear points on a plane. Hence, we must consider all possible combinations of three points from our array and calculate the area of the triangle they form. This is

The given problem provides an array of points representing coordinates on the X-Y plane. Each point is defined by a pair of

integers, [xi, yi], which represent its X and Y coordinates, respectively. The task is to calculate the area of the largest triangle

essentially a combinatorial problem where we seek the maximum value from a set of results obtained from these combinations. Intuition

The intuitive solution for this problem leverages the mathematical formula for the area of a triangle when given three points. The formula to calculate the area T of a triangle with vertices at coordinates (x1, y1), (x2, y2), and (x3, y3) is: T = |(x1(y2-y3) + x2(y3-y1) + x3(y1-y2)) / 2|

```
space. Remember, in a 2D plane, the area A of a triangle formed by points (x1, y1), (x2, y2), and (x3, y3) can also be
```

computed using the following formula derived from the cross product:

However, the provided solution code uses the cross product method to find the area of a triangle formed by three points in 2D

A = |(x1(y2-y3) + x2(y3-y1) + x3(y1-y2)) / 2|

This is equivalent to:

A = |((x2-x1)*(y3-y1) - (x3-x1)*(y2-y1)) / 2|

maximum area (ans), it updates the ans variable.

implementation steps and the rationale behind them.

keep track of the maximum area found.

Here's a step-by-step explanation of the solution:

The formula is based on the fact that the absolute value of the cross product of two vectors can be interpreted as the area of the

area is found.

the triangle formed by the points.

parallelogram they span, and the area of the triangle is half of that parallelogram. The code implements a brute-force approach by considering every possible triplet of points:

1. Three nested loops iterate through all combinations of three points. It computes vectors [u1, v1] and [u2, v2] that represent two sides of a potential triangle. 2. The cross product of these vectors is calculated, which is essentially the determinant of a 2×2 matrix, and it corresponds to twice the area of

3. This value is halved and the absolute value is taken to get the area of the triangle. If the calculated area is greater than a previously recorded

4. After exhausting all possible combinations, the largest recorded triangle area ans is returned. The approach is exhaustive and straightforward but can be inefficient for large input sizes because it considers all possible point

Solution Approach The solution approach utilizes a straightforward brute-force algorithm to solve the problem. Here's an in-depth breakdown of the

Algorithm: The algorithm loops through each possible combination of three distinct points in the input array to compute the

Data Structures: No specific data structures are needed for this solution except for the input list of points and a variable to

triplets, resulting in a time complexity of $O(n^3)$, where n is the number of points. Nevertheless, it ensures that the maximum

area of the triangle formed by those points. It uses the cross product method to determine the area, which works effectively for this 2D geometric problem.

pattern is common in problems where all combinations of elements need to be considered.

Patterns: The pattern here is the use of multiple nested loops to generate all possible triplets of points from the list. This

CP = (u1 * v2 - u2 * v1)

t = abs(u1 * v2 - u2 * v1) / 2

iteration. **Nested Loops:** The three nested for loops are used to iterate over all possible combinations of three points from the points

array. These loops pick one point in each iteration, creating triplets such as (points[i], points[j], points[k]).

Initial Maximum Area: A variable called ans is initialized to zero, which will store the maximum triangle area found during the

- Vector Calculation: Within the innermost loop, two vectors are determined that represent the sides of the triangle from the point (x1, y1) to (x2, y2) and from (x1, y1) to (x3, y3).
- vectors in 2D space. Area Calculation: The actual area of the triangle is found by halving and taking the absolute value of the cross product CP,

Updating Maximum Area: If the computed triangle area t is greater than the current maximum ans, it is updated to t.

We consider only the 'z' component, which in a 3D cross product represents the area of the parallelogram formed by two

Result: Once all combinations have been considered, and the maximum area has been found, the variable ans holds the

largest triangle area, which is then returned as the final result.

the cross product approach explained in the content.

Vector Calculation: We calculate two vectors:

Since t is 6 and ans is 0, we update ans to 6.

to find the triangle with the maximum area.

Initialize the max area to 0

for x2, y2 in points:

for x3, y3 in points:

u1, v1 = x2 - x1, y2 - y1

u2, v2 = x3 - x1, y3 - y1

Return the largest found area of a triangle

public double largestTriangleArea(int[][] points) {

for (int[] point3 : points) {

// Return the largest area found.

points.forEach(point3 => {

const x3 = point3[0], y3 = point3[1];

const vector1X = x2 - x1, vector1Y = y2 - y1;

const vector2X = x3 - x1, vector2Y = y3 - y1;

// Update largestArea if a larger area is found

largestArea = Math.max(largestArea, area);

const examplePoints: number[][] = [[0, 0], [0, 1], [1, 0], [0, 2], [2, 0]];

def largest triangle area(self, points: List[List[int]]) -> float:

area = abs(u1 * v2 - u2 * v1) / 2

max area = max(max area, area)

Return the largest found area of a triangle

three nested loops each iterating over all the points.

Iterate over all possible combinations of three points to form triangles

(cross product of two vectors) divided by 2

Update the max area if the current area is larger

// Calculate the vectors from point1 to point2, and from point1 to point3

const area = Math.abs(vector1X * vector2Y - vector2X * vector1Y) / 2.0;

// Calculate the cross product of the vectors to get twice the area of the triangle

console.log(largestTriangleArea(examplePoints)); // Possible output: 2 (for the triangle formed by the last three points)

Calculate the triangle's area using the absolute value of the determinant

return maxArea;

C++

int x3 = point3[0], y3 = point3[1];

// which is $|x^2 - x^1| + |y^2 - y^1|$

int vector1X = $x^2 - x^1$, vector1Y = $y^2 - y^1$;

int vector2X = x3 - x1, vector2Y = y3 - y1;

// and (point1 to point3) divided by 2.

maxArea = Math.max(maxArea, area);

area = abs(u1 * v2 - u2 * v1) / 2

max area = max(max area, area)

From (0,0) to (0,4) which is (0 − 0, 4 − 0) or (0,4).

 \circ From (0,0) to (3,0) which is (3 - 0, 0 - 0) or (3,0).

Cross Product: The cross product of these two vectors is computed as:

which is the determinant of the 2×2 matrix formed by vectors u1, v1 and u2, v2:

```
requirements.
```

This elementary brute-force solution ensures that even though it might not be efficient for a large number of points (due to its

O(n^3) time complexity), it will always yield the correct maximum area of a triangle that can be formed, satisfying the problem

Let's use a small set of points to illustrate the solution approach provided: Suppose we have the following array of points: [(0,0), (0,4), (3,0)].

This array forms a right-angled triangle with the right angle at the origin (0,0). The area of this triangle can be calculated using

Initial Maximum Area: We initialize a variable ans to zero. This will store the maximum triangle area we find.

 The first loop starts with i = 0, with points[i] will be (0, 0). The second loop starts with j = 1, so points[j] will be (0, 4). • The third loop starts with k = 2, so points[k] will be (3, 0).

With only one combination of points, the final ans we get is the maximum area, which is 6 in this case. Hence, the largest triangle

Cross Product: \circ We compute the cross product CP of these vectors, which will be 0 * 0 - 3 * 4, resulting in -12. **Area Calculation:**

class Solution:

Java

class Solution {

max area = 0

return max_area

for x1, v1 in points:

Example Walkthrough

Nested Loops:

 \circ The absolute value of the cross product is |CP| = |-12| = 12. \circ The area of the triangle would be half of this, t = CP / 2 = 6. **Updating Maximum Area:**

that can be formed by any three points from the given array has an area equal to 6.

Iterate over all possible combinations of three points to form triangles

Calculate vector (u1, v1) from point 1 to point 2

Calculate vector (u2, v2) from point 1 to point 3

Update the max area if the current area is larger

(cross product of two vectors) divided by 2

By using the above steps on our small set of points, the brute-force algorithm effectively determines that the right-angled triangle formed by these points produces the largest area among all possible triangles formed from these points. This example demonstrates the underlying mechanics of the solution approach, where the algorithm evaluates all combinations of three points

def largest triangle area(self, points: List[List[int]]) -> float:

Solution Implementation Python

Calculate the triangle's area using the absolute value of the determinant

// Iterate through all the points to get the third point (p3) of the triangle.

// Calculate the vectors from point1 to point2 and from point1 to point3

// the absolute value of the cross product of vectors (point1 to point2)

double area = Math.abs(vector1X * vector2Y - vector2X * vector1Y) / 2.0;

// Update maxArea if the calculated area is larger than the current maxArea.

 $|x3 - x1 \ v3 - v1|$ divided by 2. Note: this represents

// Calculate the area of the triangle using the determinant |u1 v1|

double maxArea = 0; // Initialize maxArea as the largest triangle area found so far. // Iterate through all the points to get the first point (p1) of the triangle. for (int[] point1 : points) { int x1 = point1[0], y1 = point1[1]; // Iterate through all the points to get the second point (p2) of the triangle. for (int[] point2 : points) { int x2 = point2[0], y2 = point2[1];

|u2 v2|

```
#include <vector>
#include <cmath>
#include <algorithm>
class Solution {
public:
    double largestTriangleArea(vector<vector<int>>& points) {
        double largestArea = 0; // Initialize the variable to hold the largest triangle area
        // Iterate over all possible combinations of three points to find the largest area
        for (auto& point1 : points) {
            int x1 = point1[0], y1 = point1[1];
            for (auto& point2 : points) {
                int x2 = point2[0], y2 = point2[1];
                for (auto& point3 : points) {
                    int x3 = point3[0], y3 = point3[1];
                    // Calculate the vector from point1 to point2 and point1 to point3
                    int vector1X = x2 - x1, vector1Y = v2 - v1;
                    int vector2X = x3 - x1, vector2Y = y3 - y1;
                    // Use the cross product of vectors to find the area of the triangle
                    double area = std::abs(vector1X * vector2Y - vector2X * vector1Y) / 2.0;
                    // Update largestArea if a larger area is found
                    largestArea = std::max(largestArea, area);
        return largestArea; // Return the result
};
TypeScript
function largestTriangleArea(points: number[][]): number {
    let largestArea: number = 0; // Initialize the variable to hold the largest triangle area
    // Triple nested loop to iterate over all possible combinations of three points
    points.forEach(point1 => {
        const x1 = point1[0], y1 = point1[1];
        points.forEach(point2 => {
            const x2 = point2[0], y2 = point2[1];
```

```
for x3, v3 in points:
   # Calculate vector (u1, v1) from point 1 to point 2
   u1. v1 = x2 - x1. v2 - v1
   # Calculate vector (u2, v2) from point 1 to point 3
   u2. v2 = x3 - x1. v3 - v1
```

return max_area

for x1, y1 in points:

});

return largestArea;

max area = 0

// Return the largest area found

Initialize the max area to 0

for x2, y2 in points:

});

// Example usage:

class Solution:

});

Time and Space Complexity The time complexity of the code is $0(n^3)$ where n is the number of points provided in the points list. This is because there are

The space complexity of the code is 0(1) because the only extra space used is for variables to store the current point coordinates, the components of two vectors (u1, v1, u2, v2), and the area of the current triangle (t). These variables require a constant amount of space regardless of the input size.