2507. Smallest Value After Replacing With Sum of Prime Factors

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Medium Math
              Number Theory
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Problem Description The problem involves manipulating a given positive integer n by continuously replacing it with the sum of its prime factors. A key

factor must be counted multiple times in the sum. For example, if our given number n is 18, its prime factors are 2 and 3, but 3 is counted twice because 18 is divisible by 3 twice (18 = 2 * 3 * 3). So the sum of its prime factors is 2 + 3 + 3 = 8.

detail is that if a prime factor is repeated in the factorization (that is, n is divisible by a prime number multiple times), that prime

The goal is to repeat this process: take the sum of the prime factors of n and then replace n with this sum, and continue this until n no longer changes—it reaches its smallest possible value.

The challenge is to write an algorithm that performs this computation efficiently and find when the number n stops changing, which is the result to be returned.

The solution makes use of a loop that continuously replaces the number n with the sum of its prime factors until n cannot be

We start with a loop that runs indefinitely, checking if we can find prime factors of n and their sums. In each iteration, we initialize

reduced any further.

any further, and we return t.

by continuously summing its prime factors.

value of n.

three variables - t to keep track of the original value of n at the start of the iteration, s to calculate the sum of prime factors, and i to iterate over potential prime factors starting from 2.

The inner while loop checks if i is a divisor of n, and if so, it keeps dividing n by i and adds i to the sum s until n is no longer divisible by i. The variable i is incremented to check for the next possible prime factor. Once we have tried all possible divisors up to n // i, we check if n itself is greater than 1 (which would mean n is a prime

number and should be included in the sum). At this point, if the calculated sum of prime factors s is equal to the temporary variable t, it means n can no longer be reduced

If s is not equal to t, we replace n with s to continue the process with a new reduced value of n. The key intuition behind the algorithm is that the prime factorization of a number can help find a smaller representation of the number itself and by repeating this process and summing these factors repeatedly, we can converge to the smallest possible

Solution Approach

The solution makes use of a basic factorization algorithm and control flow to reduce the number n to its smallest possible value

Here's a step-by-step breakdown of the approach used in the solution: The solution uses a while loop that will run until n no longer changes. Inside the loop, a temporary variable t is assigned the current value of n to keep track of its value through each iteration.

An index i is set to 2, which is the first prime number. This variable is used to test potential factors of n.

number.

The first inner while loop runs as long as i is less than or equal to n // i (since a factor larger than the square root of n would have already been identified by its corresponding smaller factor except when n is prime).

A nested inner while loop checks if i divides n perfectly. If it does, n is divided by i, and i is added to the sum s. This

At the end of the iteration, the algorithm checks if the sum of the prime factors s is equal to the starting value t. If so, n has

reached its smallest value as no primes other than itself can be extracted and summed, and the algorithm returns the value of

The index i is then incremented to check the next potential factor. After all possible factors up to n // i have been tested, a final check outside the inner loops evaluates if n itself is greater

replace n with the sum of its prime factors and repeat this process until n is no longer reduced.

than 1, implying n is a prime number. If it is prime, it is added to the sum s.

indivisible and signaling the end of the reduction process.

Step 1: Start with n = 12 and enter the while loop.

Step 7: Increment i to the next integer, which is 3.

Step 1: n is 7, so we enter the while loop.

number to its smallest form based on prime factor accumulation.

Step 2: Assign the value of n to a temporary variable t, so t = 12.

Step 4: Start with i = 2, which is the smallest prime factor.

Step 3: Initialize a variable s to 0. This will hold the sum of the prime factors of n.

Now n has changed from 12 to 7. We repeat the entire process again with n = 7.

Step 5: Begin inner loop. Because no i exists such that i <= n // i and i divides 7, skip to step 8.

loop continues until n is no longer divisible by i, accommodating for all instances of i as a factor.

Another variable s is initialized to 0; this will be used to calculate the sum of the prime factors of n.

This is important to identify when there is no further change possible.

t. If s is not equal to t, the algorithm replaces the value of n with s to repeat the factorization process on this new reduced 10.

The algorithm essentially terminates when a number only comprises its prime self or when it's reduced to 1, both of which are

This algorithm does not employ any complex data structures and follows a straightforward but effective pattern to reduce the

Example Walkthrough Let's walk through the solution approach with a small example where our given number n is 12. The goal is to continuously

Step 5: As $i \le n // i$ (since 2 $\le 12 // 2$), we proceed with factorization. Step 6: The inner loop checks if 12 is divisible by 2. It is, so we divide n by 2 to get 6 and add 2 to s. Now, s = 2 and n = 6.

We continue with the inner loop since 6 is still divisible by 2. We divide 6 by 2 to get 3 and add 2 to s. Now, s = 4 and n = 3.

Step 8: Since 3 is a prime number larger than n // i but still divides n perfectly, we add 3 to s. Now, s = 4 + 3 giving us s =

Step 9: Since n is now 1, we exit the inner loop. We check if s is equal to the temporary variable t. As s = 7 and t = 12, they

Step 2: Set t to 7.

further.

described.

class Solution:

Python

Java

class Solution {

Step 3: Initialize s to 0.

Step 4: Start with i = 2.

Solution Implementation

while True:

def smallestValue(self, num: int) -> int:

while num % divisor == 0:

sum of factors += divisor

Move to the next potential factor

num //= divisor

divisor += 1

num = sum_of_factors

public int smallestValue(int n) {

int originalValue = n;

int sumOfFactors = 0;

while (n % i == 0) {

sumOfFactors += n;

return sumOfFactors;

while (true) {

7 and n becomes 1.

are not equal. Step 10: Because s is not equal to t, we set n to s; hence n is now 7.

Step 8: Since n is still greater than 1, and 7 is a prime number, we add 7 to s. Now s = 7. Step 9: Compare the sum of the prime factors s with t. They are equal (s = t = 7), indicating that we cannot reduce n any

The algorithm would then terminate and return the value 7 as the result, as n can no longer be reduced by the process

Initialize temp variable to store original number, sum_of_factors, and start divisor from 2 temp, sum of factors, divisor = num, 0, 2# Check for factors of the number starting with the smallest prime factor while divisor <= num // divisor:</pre> # If the divisor is a factor, divide num by the divisor and add to the sum_of_factors

// Method to find the smallest value according to specified conditions

// Start dividing the number from 2 onwards to find its factors

// If there is a remaining factor greater than 1, add it to the sum

// If it matches, return the sum (as it is the smallest value)

// If it does not match, set n to the sumOfFactors for another iteration

sumOfFactors += i; // Add factor to the sum

// Check if the sum of the factors equals the original number

// Loop indefinitely until we find the smallest value

for (int i = 2; i <= originalValue / i; ++i) {</pre>

// Divide by i as long as it is a factor of n

n /= i; // Divide n by the factor

// Store the original value of n

// Initialize the sum of the factors

if (sumOfFactors == originalValue) {

// Function to find the smallest integer value with the same

// sum of factors (including 1 and the number itself) as 'n'

// Loop indefinitely until we find the smallest value

// Factorize and sum up the factors

// While 'n' is divisible by 'i'

// Check for a remaining prime factor greater than sqrt(n)

Continue the loop until we find the smallest value

sumOfFactors += n; // Add it to the sum only if it's different from the current value

// If the sum of factors is the same as the original value, return the smallest integer

n = sumOfFactors; // Set 'n' to the calculated sum of factors for the next iteration

if $(n > 1 \&\& n < originalValue) {$

if (sumOfFactors == originalValue) {

return originalValue;

def smallestValue(self, num: int) -> int:

class Solution:

int originalValue = n; // Preserve the original value of 'n'

int sumOfFactors = 0; // Initialize the sum of factors to 0

for (int i = 2; $i \le n / i$; ++i) { // Only need to check up to sqrt(n)

Continue the loop until we find the smallest value

if num > 1: sum of factors += num # If the sum of factors is equal to the original number, we found the smallest value if sum of factors == temp: return temp # Set num to sum of factors for the next iteration to check the new number

If there's a remaining number greater than one, it's a prime factor; add it to the sum_of_factors

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C++
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int smallestValue(int n) {

while (true) {

class Solution {

public:

n = sumOfFactors;

if (n > 1) {

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while (n % i == 0) {
                    sumOfFactors += i; // Add the factor to the sum
                    n /= i; // Divide 'n' by the factor for further factorization
            // If there is a remaining prime factor greater than sqrt(n), add it to the sum
            if (n > 1) sumOfFactors += n;
            // If the sum of factors is the same as the original value, return it
            if (sumOfFactors == originalValue) {
                return sumOfFactors;
            // Otherwise, set 'n' to the calculated sum of factors for the next iteration
            n = sumOfFactors;
TypeScript
// Function to find the smallest integer value with the same
// sum of factors (including 1 and the number itself) as 'n'
function smallestValue(n: number): number {
    // Loop indefinitely until we find the smallest value
    while (true) {
        const originalValue: number = n; // Preserve the original value of 'n'
        let sumOfFactors: number = 1; // Initialize the sum of factors to 1, since 1 is a factor of all numbers
        // Factorize and sum up the factors
        for (let i = 2; i \le Math.sgrt(n); ++i) { // Only need to check up to sgrt(n)
            // While 'n' is divisible by 'i'
            while (n % i === 0) {
                sumOfFactors += i; // Add the factor to the sum
                if (i != n / i) {
                    sumOfFactors += n / i; // Add the complement factor to the sum if it's different
                n /= i; // Divide 'n' by the factor for further factorization
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while True:
           # Initialize temp variable to store original number, sum_of_factors, and start divisor from 2
            temp, sum of factors, divisor = num, 0, 2
           # Check for factors of the number starting with the smallest prime factor
           while divisor <= num // divisor:</pre>
                # If the divisor is a factor, divide num by the divisor and add to the sum_of_factors
               while num % divisor == 0:
                   num //= divisor
                    sum of factors += divisor
               # Move to the next potential factor
                divisor += 1
           # If there's a remaining number greater than one, it's a prime factor; add it to the sum_of_factors
            if num > 1:
                sum of factors += num
           # If the sum of factors is equal to the original number, we found the smallest value
            if sum of factors == temp:
                return temp
           # Set num to sum of factors for the next iteration to check the new number
            num = sum_of_factors
Time and Space Complexity
Time Complexity
```

times because it checks for factors from 2 to at most sqrt(n). The factorization loop:

• In the worst-case scenario for a given n, we may have to factorize n, then s, and so on if n is initially not a prime or semiprime (a product of exactly two primes). Each new s will be smaller than n as we are adding up the factors. Thus, the inner factorization process will take O(sqrt(m)) time for each number m we factorize, where m starts from n and gets smaller. The outer while loop:

• The code does not have a clear stopping condition within predictable bounds, due to the uncertain nature of the sum s approaching the target

a prime number, it will become equal to t in the next iteration, and the function will return t. The number of iterations can be considered

condition s == t. However, for every iteration, the sum of the prime factors of n (s) gets closer to being a prime number itself. Once s equals

The time complexity of this code is governed by two nested loops. The outer loop runs indefinitely until a number is found where

the sum of its prime factors is equal to itself. The inner loop, on the other hand, is used for factorization and runs at most sqrt(n)

proportional to the number of distinct prime factors of n in a sense but is not easy to express in standard 0() notation. Considering the worst-case scenario, where n is a large composite number with many small prime factors, the time complexity

can be high, but the exact upper bound is intricate to determine without more constraints on n.

Space Complexity

The space complexity is 0(1). There are only a few integer variables (t, s, i, and n) being used, which do not depend on the input size n, unless considering the size n itself needs. The space used by these variables is constant and does not grow with n.