2280. Minimum Lines to Represent a Line Chart

Number Theory

Sorting

Math

<u>Array</u>

Problem Description

Geometry

corresponding stock price on that day. The task is to determine the minimum number of straight lines required to connect all the points plotted on an XY plane, with the day as the X-axis and the price as the Y-axis, in the given order after sorting the data by the day. Essentially, we want to find the fewest lines that connect each adjacent point (each day to the next) to form a line chart.

We have a 2D integer array named stockPrices, where each element consists of two integers representing a day and the

Intuition To solve this problem, the solution leverages the concept of slopes of lines. When connecting points to draw the line chart, we can continue using the same line as long as consecutive segments share the same slope. When the slopes differ, we need a new

Medium

line. To find the slope of the line segment between two points (x, y) and (x1, y1), we calculate (y1 - y) / (x1 - x), which gives us the rate of change in price with respect to days. Now, if we compare this slope with the next segment's slope between the points (x1, y1) and (x2, y2), we would normally compute (y2 - y1) / (x2 - x1) and compare. To avoid division and possible precision

issues with floating-point numbers, we instead compare cross-products: (y1 - y) * (x2 - x1) and (y2 - y1) * (x1 - x). If

these two products are not equal, the slope has changed, and we need an additional line. The approach first sorts the stockPrices based on the day to ensure the points are in the correct order. Then, iterating through each pair of adjacent points, we calculate the aforementioned cross-product to decide whether the current segment can be connected with the prior one or if a new line has to start. The variable ans keeps track of the number of lines needed. Remember, we increment the number of lines when a change in slope is detected between segments, and after calculating the

slopes, we update the dx and dy to be the differences between the current and next points to be ready for the next iteration. Solution Approach

Sort the Input: We start by sorting the stockPrices array. This ensures that the stock prices are ordered by day, which is essential as we need to connect points from one day to the next.

Initialize Variables: Two variables dx and dy are initialized to store the differences in the x (day) and y (price) coordinates

Iterate Through Points: We use a loop to go through each pair of adjacent points in the sorted array. To do this efficiently, the

change in y as dy1 = y1 - y. We then compare the products of cross-multiplying these with the previous segment's dx and

Update Differences: Whether a new line is added or not, we update dx and dy with the differences dx1 and dy1, to use

Return Result: After all points have been processed, the variable ans is returned. This gives us the minimum number of lines

code leverages the pairwise utility, which is not a standard Python function but likely a user-defined or library function that

yields pairs of elements from the input list.

them in the next iteration for slope comparison.

needed to represent the chart.

without redundant calculations or comparisons.

Suppose we have the following 2D array stockPrices:

Sort the Input:

Initialize Variables:

Iterate Through Points:

Increment Line Count:

Return Result:

Solution Implementation

from itertools import pairwise

num_lines = 0

prev_diff_x, prev_diff_y = 0, 1

num_lines += 1

Initialize the count of lines needed to 0

Iterate over each pair of adjacent stock prices

prev_diff_x, prev_diff_y = diff_x, diff_y

Return the total number of lines needed

for (current x, current y), (next x, next y) in pairwise(stock_prices):

If the cross product of the differences does not equal to zero,

then the points are not on the same line. Increment the number of lines.

int prevDeltaX = 0, prevDeltaY = 1; // Initiated to 0 and 1 to handle the starting calculation

// Calculate the change in x and y (i.e., the slope components) for the current line

// by comparing cross products of the slope components (to avoid floating-point division)

int previousX = stockPrices[i - 1][0], previousY = stockPrices[i - 1][1];

// Check if the current line has a different slope than the previous one

// Return the total number of lines needed to connect all points with straight lines

++lineCount; // A different slope means a new line is needed

// Update previous slope components for the next iteration

int currentX = stockPrices[i][0], currentY = stockPrices[i][1];

int deltaX = currentX - previousX, deltaY = currentY - previousY;

Compute the differences in x and y for the current pair

diff_x, diff_y = next_x - current_x, next_y - current_y

Update the previous differences for the next iteration

// Initialize variables to keep track of previous slope components

// Iterate over the stock prices starting from the second price

int lineCount = 0; // Counter for the minimum lines needed

if (prevDeltaY * deltaX != prevDeltaX * deltaY) {

for (int i = 1; i < stockPrices.length; ++i) {</pre>

// Get the previous day's stock price

// Get the current day's stock price

prevDeltaX = deltaX;

prevDeltaY = deltaY;

return lineCount;

if prev diff v * diff_x != prev_diff_x * diff_y:

from typing import List

class Solution:

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Python

Calculate and Compare Slopes:

Increment Line Count: When a slope change is detected, we increment ans by 1.

Here is the explanation of the algorithm, data structure, or pattern used in each step:

• Sorting: It's the first step and is critical for comparing the slopes of segments in the correct order.

The implementation given in the reference solution approach follows these steps:

of successive points. Additionally, ans is initialized to count the number of lines required.

- Calculate and Compare Slopes: For each pair (x, y) and (x1, y1), we calculate the change in x as dx1 = x1 x and the
- dy to check if the slopes match. If dy * dx1 does not equal dx * dy1, the slope has changed, indicating a need for a new line.
- Looping and Iteration: The use of a loop along with pairwise comparison is a common pattern when we need to check adjacent elements in a sequence. • Slope Comparison: This step avoids division by comparing slopes using cross-multiplication, which is a mathematical technique to avoid floating-point precision errors and ensure that comparisons are accurate when checking for collinearity between points.

By systematically comparing each pair of points, the solution efficiently calculates the number of lines needed to form the chart

Example Walkthrough

stockPrices = [[1, 100], [2, 200], [3, 300], [4, 200], [5, 100], [6, 200]]

Let's walk through a small example to illustrate the solution approach.

• We set dx and dy to 0 as there is no previous point to compare with yet. ans is initialized to 1 because we need at least one line to connect the first point to the next one.

Since stockPrices is already sorted by the day (the first element of each sub-array), we don't need to sort it again.

∘ Imagine comparing the first point [1, 100] with the second [2, 200]. This sets our initial dx and dy to [2-1, 200-100] = [1, 100].

• Now we compare the slopes. Since dy * dx1 = dx * dy1 (100 * 1 = 1 * 100), the slope hasn't changed, and we stay on the same line.

- On the next iteration, we compare the point [2, 200] with [3, 300]. We calculate the changes dx1 = 3 2 and dy1 = 300 200, resulting in [1, 100].
 - We don't increment ans as the line didn't change. **Update Differences:**

 \circ Slopes differ (dy * dx1 = 100 * 1!= -1 * 100 = dx * dy1), so we increment ans to 2 and update dx and dy.

Comparing [4, 200] with [5, 100], dx1 = 5 - 4 and dy1 = 100 - 200, resulting in [1, -100].

 \circ Slopes differ again (dy * dx1 = -100 * 1 != 1 * 100 = dx * dy1), so we increment ans to 3.

After all points have been processed, we conclude that we need ans = 3 lines to connect all the points.

Comparing [3, 300] with [4, 200], we calculate dx1 = 4 - 3 and dy1 = 200 - 300, resulting in [1, -100].

With this hands-on example, we've illustrated how the solution approach can be applied to determine the minimum number of

∘ Slopes match (as they're both [1, -100]), so we don't increment ans and just update dx and dy. Finally, comparing [5, 100] with [6, 200], dx1 = 6 - 5 and dy1 = 200 - 100, resulting in [1, 100].

Continue this process for the remaining points:

lines required to connect a series of points in a line chart.

We set dx and dy to dx1 and dy1 respectively, for the next comparison.

def minimumLines(self, stock prices: List[List[int]]) -> int: # Sort the stock prices based on the days stock_prices.sort() # Initialize the previous differences in x and y to 0 and 1 respectively as placeholders.

public int minimumLines(int[][] stockPrices) { // Sort the stock prices array based on the day (the 0th element of each sub-array); // if the days are equal, sort based on the price (the 1st element of each sub-array) Arrays.sort(stockPrices, (a, b) -> Integer.compare(a[0], b[0]));

Java

class Solution {

return num_lines

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C++
#include <vector>
#include <algorithm>
class Solution {
public:
    int minimumLines(std::vector<std::vector<int>>& stockPrices) {
        // Sort stockPrices by their x-value (date)
        std::sort(stockPrices.begin(), stockPrices.end());
        // Initialize previous difference in x and y to compare slopes.
        // Starting with an impossible value for the very first comparison.
        int prevDeltaX = 0, prevDeltaY = 1;
        // Begin with no lines drawn
        int numLines = 0;
        // Iterate through stockPrices starting from the second point
        for (int i = 1; i < stockPrices.size(); ++i) {</pre>
            // Get x and y of the current and previous points
            int currX = stockPrices[i - 1][0], currY = stockPrices[i - 1][1];
            int nextX = stockPrices[i][0], nextY = stockPrices[i][1];
            // Calculate differences in x and y for the current segment
            int deltaX = nextX - currX, deltaY = nextY - currY;
            // If the cross product of the two segments is not zero,
            // they are not collinear, so the slope has changed and a new line is needed.
            if (static cast<long long>(prevDeltaY) * deltaX != static_cast<long long>(prevDeltaX) * deltaY) {
                ++numLines;
                // Update previous differences to current one for next iteration comparison.
                prevDeltaX = deltaX;
                prevDeltaY = deltaY;
        // Return the total number of lines needed
        return numLines;
};
```

TypeScript

function minimumLines(stockPrices: number[][]): number {

// Initialize the number of lines needed to connect all points

let previousSlope: [BigInt, BigInt] = [BigInt(0), BigInt(0)];

const [previousX. previousY] = stockPrices[i - 1];

const differenceX = BigInt(currentX - previousX);

const differenceY = BigInt(currentY - previousY);

const [currentX, currentY] = stockPrices[i];

previousSlope = [differenceX, differenceY];

// Return the total number of lines required

// Start from the second point and compare slopes of consecutive points

// Get the coordinates of the current point and the previous point

// Sort the stock prices based on the first value of each pair (the dates are sorted)

// Initialize a variable to store the previous slope as a pair of BigInts for accurate calculation

// Calculate the difference in x and v using BigInt for accuracy (to handle very large numbers)

// If it's the first line or the slope differs from the previous slope, increment the line count

if (i == 1 || differenceX * previousSlope[1] !== differenceY * previousSlope[0]) {

// Update the previous slope to the current slope for the next iteration

// Get the number of stock price entries

stockPrices.sort((a, b) => a[0] - b[0]);

let linesRequired = 0;

const stockPricesCount = stockPrices.length;

for (let i = 1; i < stockPricesCount; i++) {</pre>

linesRequired++;

return linesRequired;

from itertools import pairwise

```
from typing import List
class Solution:
    def minimumLines(self, stock prices: List[List[int]]) -> int:
        # Sort the stock prices based on the days
        stock_prices.sort()
        # Initialize the previous differences in x and y to 0 and 1 respectively as placeholders.
        prev_diff_x, prev_diff_y = 0, 1
        # Initialize the count of lines needed to 0
        num_lines = 0
        # Iterate over each pair of adjacent stock prices
        for (current x, current y), (next x, next y) in pairwise(stock_prices):
           # Compute the differences in x and y for the current pair
           diff_x, diff_y = next_x - current_x, next_y - current_y
           # If the cross product of the differences does not equal to zero,
            # then the points are not on the same line. Increment the number of lines.
            if prev diff v * diff_x != prev_diff_x * diff_y:
                num_lines += 1
           # Update the previous differences for the next iteration
            prev_diff_x, prev_diff_y = diff_x, diff_y
        # Return the total number of lines needed
        return num_lines
Time and Space Complexity
Time Complexity
  The given Python code first sorts the stockPrices. Sorting takes O(n log n) time where n is the number of stockPrices.
  Then, it iterates through the sorted prices to determine the number of lines required. This iteration involves a single pass through
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the stockPrices list, which takes O(n) time.

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operations do not increase the overall complexity beyond O(n).
Since sorting is the most time-consuming operation and is followed by a linear pass, the total time complexity of this algorithm is
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Inside the loop, it performs constant-time arithmetic operations to determine if a new line segment is required. Therefore, these

Space Complexity The space complexity is determined by the additional space required by the algorithm besides the input.

Here, the algorithm uses a fixed amount of space to store temporary values like dx, dy, dx1, and dy1, so it uses 0(1) additional

space. Sorting is typically done in-place (especially if the Python sort() method is used), therefore, the space complexity remains 0(1).

In summary, the space complexity of the algorithm is 0(1).

 $O(n \log n) + O(n)$, which simplifies to $O(n \log n)$.