3038. Maximum Number of Operations With the Same Score I

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Easy <u>Array</u> Simulation
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**Problem Description** 

Each time we remove a pair, we calculate the sum of the two elements and consider this the score of the operation. The goal is to perform as many operations as possible with the condition that all operations must yield the same score. This means that each pair of elements we remove must have the same sum as the first pair we removed.

To find the maximum number of operations we can perform under these conditions, we need to continuously check the first two

In this problem, we have an array of integers called nums. We are allowed to remove pairs of elements from the start of the array.

elements of the remaining array, removing them only if they match the desired score. If at any point the sum of the two elements doesn't match the score, we must stop the operations, and that's the maximum number we can perform.

## To solve this problem, the solution begins by presuming that the score we need every operation to match is the sum of the first two elements in the array, denoted as s. The solution then proceeds to traverse the array in steps of 2, meaning it looks at every

Intuition

consecutive pair of elements.

During this traversal, the solution checks if the sum of each pair matches the initial score s. As soon as it encounters a pair that does not have a sum equal to s, or if there are no more pairs to check (e.g., after removing pairs the array has only one element

left), the traversal stops. This is grounded in the rule that all operations must yield the same score, and since the score is defined by the first operation, any pair that doesn't match this score would break the condition.

Solution Approach

The implemented solution follows a straightforward and direct approach. It doesn't employ complex algorithms or data structures

## but relies on a simple traversal of the array using a for loop. Here's a step-by-step explanation of the implementation:

array.

s = nums[0] + nums[1]

1. The solution initializes a variable s with the sum of the first two elements in nums. This s will be the score that every operation must match.

2. It initializes an answer variable, ans, to count the number of operations performed and a variable n to store the length of the

ans, n = 0, len(nums)

3. The solution uses a for loop to iterate over the elements of the array in steps of two, looking at every pair of elements.

For each pair, it checks two conditions: whether the current index i + 1 equals the length of the array (which would mean

break

ans += 1

score.

return ans

for i in range(0, n, 2):

If either condition is true, the loop breaks, effectively stopping any further operations, as either there are no more pairs to

there's no pair left to process) and whether the sum of the current pair does not match the score s.

remove or the condition of maintaining the same score can no longer be met.

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If none of the conditions for breaking the loop are met, it means that the current pair matches the score, and an operation can
be performed. The ans variable is incremented to reflect this.
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if i + 1 == n or nums[i] + nums[i + 1] != s:

6. The loop continues until it has either checked all elements in the array or encountered a break. At the end of the loop, the solution returns the count ans, which represents the maximum number of operations that have been performed with the same

This solution assumes that the input array nums has already been set up such that every valid operation (pair with matching sum

s) is adjacent and that there are no possible operations (pairs with the desired sum) after encountering the first invalid one. Therefore, it's essential to confirm that the input array will conform to these conditions; otherwise, the solution might not correctly calculate the maximum number of operations.

Let's illustrate the solution approach with a simple example.

Suppose we have the following array of integers as our nums:

We initialize our answer variable, ans, to 0 and a variable n to store the length of the array, which is 6 in this case.

## s = nums[0] + nums[1] # s = 10

for i in range(0, n, 2):

# First iteration:

ans += 1 # ans = 1

# Second iteration:

# Third iteration:

ans += 1 # ans = 3

in the array.

**Python** 

return ans # returns 3

Solution Implementation

from typing import List

class Solution:

well. We increment ans.

nums = [4, 6, 4, 6, 8, 2]

3. The for loop begins and will iterate over the pairs of elements.

We will walk through the steps of the implementation using this array.

4. At first, i = 0, and we look at the first pair nums [0] + nums [1] which is 4 + 6. This equals s = (10), so we can remove this pair.

First, we initialize the variable s with the sum of the first two elements in nums, which is 4 + 6 = 10.

5. Now i = 2, and we look at the next pair nums [2] + nums [3] which is 4 + 6. This also equals s = (10), so we remove this pair as

No break is encountered, and we increment ans.

# nums[0] + nums[1] = 10, which is equal to s.

# nums[2] + nums[3] = 10, which is equal to s.

ans, n = 0, len(nums) # ans = 0, n = 6

ans += 1 # ans = 2

6. Finally, i = 4, and we look at the last pair nums[4] + nums[5] which is 8 + 2. This equals 10, the same as s. Even though it

The loop finishes because there are no more elements in nums. The solution will return the count ans, which is 3 in this case.

This example shows that the array nums allowed for 3 pairs to be removed with the sum matching s = 10, hence 3 operations

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were performed. The example follows the method of sequential checking and early termination if the condition is not met. In this instance, the conditions were satisfied by all element pairs, so the maximum number of operations is equal to the number of pairs
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def maxOperations(self, nums: List[int]) -> int:

for i in range(0, total\_numbers, 2):

break

return operations\_count

int operationsCount = 0;

int n = nums.length;

} else {

// Length of the nums array

++operations\_count;

function maxOperations(nums: number[]): number {

// Iterate over the 'nums' array with a step of 2.

const requiredSum = nums[0] + nums[1];

const n = nums.length;

let operationsCount = 0;

Time and Space Complexity

return operations\_count;

// Return the total number of valid operations

// Initialize a sum 'requiredSum' with the sum of the first two elements.

// 'operationsCount' will hold the number of valid operations performed.

// 'n' represents the total number of elements in the 'nums' array.

operations\_count += 1

# Return the total pairs formed.

public int maxOperations(int[] nums) {

int targetSum = nums[0] + nums[1];

for (int i = 0; i + 1 < n; i += 2) {

++operationsCount;

if (nums[i] + nums[i + 1] == targetSum) {

# Iterate through the list in steps of 2 to form pairs.

# If we found a matching pair, increment the counter.

// Calculate the sum of the first pair of elements in the array

// Check if the current pair sums up to the target sum

// If it does, increment the operations counter

// Initialize a counter for the number of valid operations

if i + 1 == total\_numbers or nums[i] + nums[i + 1] != sum\_of\_pair:

// Loop through the array in pairs, up to the second-to-last element (i + 1 < n)

matches s, we're at the end of the array, so the loop naturally concludes.

# nums[4] + nums[5] = 10, which is equal to s. However, this is the last pair.

# Check if there are enough numbers to form pairs.
if len(nums) < 2:
 return 0

# Initialize the sum of the first pair and the answer counter.
sum\_of\_pair = nums[0] + nums[1]
operations\_count = 0
total\_numbers = len(nums)</pre>

# Check whether we reached the end or if the sum of the current pair doesn't match sum\_of\_pair.

Java

public class Solution {

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// If a pair doesn't sum up to the target sum, break out of the loop
                // No need to continue as subsequent operations will not be valid
                break;
       // Return the total number of valid operations
        return operationsCount;
C++
#include <vector>
using namespace std;
class Solution {
public:
   // Function to determine the maximum number of operations where each operation consists of
   // finding a pair of elements from the array that add up to the same value
    int maxOperations(vector<int>& nums) {
       // Initialize the sum "s" with the sum of the first two elements
        int target_sum = nums[0] + nums[1];
        // Initialize the answer, which stores the number of valid operations
        int operations_count = 0;
       // Get the size of the nums array
        int n = nums.size();
       // Iterate over the elements of the vector in pairs
        for (int i = 0; i + 1 < n \&\& nums[i] + nums[i + 1] == target_sum; <math>i += 2) {
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// If the sum of the current pair of elements equals the target sum, increment the operations count

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for (let i = 0; i + 1 < n; i += 2) {
    // Check if the current and next element sum up to the 'requiredSum'.
    if (nums[i] + nums[i + 1] === requiredSum) {
        // Increment the count of valid operations if the condition is me</pre>
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**TypeScript** 

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if (nums[i] + nums[i + 1] === requiredSum) {
              // Increment the count of valid operations if the condition is met.
              operationsCount++;
          } else {
              // If the condition is not met, break the loop as it's assumed that nums were initially arranged in pairs with the sa
              break;
      // Return the total number of operations performed.
      return operationsCount;
from typing import List
class Solution:
   def maxOperations(self, nums: List[int]) -> int:
       # Check if there are enough numbers to form pairs.
       if len(nums) < 2:</pre>
            return 0
       # Initialize the sum of the first pair and the answer counter.
        sum_of_pair = nums[0] + nums[1]
        operations_count = 0
        total_numbers = len(nums)
       # Iterate through the list in steps of 2 to form pairs.
        for i in range(0, total_numbers, 2):
           # Check whether we reached the end or if the sum of the current pair doesn't match sum_of_pair.
            if i + 1 == total_numbers or nums[i] + nums[i + 1] != sum_of_pair:
                break
            # If we found a matching pair, increment the counter.
            operations count += 1
       # Return the total pairs formed.
        return operations count
```

The time complexity of the code is 0(n/2) since the loop increments by 2 each time, resulting in looping through half of the array size in the worst case. However, this simplifies to 0(n) because constant factors are dropped in Big O notation.

The space complexity of the code is 0(1) because it uses a fixed amount of additional space (variables s, ans, and n) that doesn't scale with the size of the input array nums.