1134. Armstrong Number

Easy Math

An Armstrong number (also known as a narcissistic number) is a number that is the sum of its own digits each raised to the power

Problem Description

number.

To explain it with an example, let's take the number 153 which is a 3-digit number. If we raise each digit to the power of 3 (since it's a 3-digit number) and sum them, we get: $1^3 + 5^3 + 3^3 = 1 + 125 + 27 = 153$, which is equal to the original number.

of the number of digits. In this problem, you are given an integer n, and the task is to determine whether n is an Armstrong

it's a 3-digit number) and sum them, we get: $1^3 + 5^3 + 3^3 = 1 + 125 + 27 = 153$, which is equal to the original number. Therefore, 153 is an Armstrong number.

This definition generalizes to any number of digits. For any given n, we wish to find out whether the number is equal to the sum

of its own digits each raised to the power of k, where k is the number of digits in n.

Intuition

The intuition behind the solution follows from the definition of an Armstrong number. We need to: 1. Determine the number of digits k in n, which can be done by converting the number n to a string and finding its length.

2. Sum up each digit raised to the kth power. This can be done by repeatedly taking the last digit of n using the modulo operator % 10, raising it to the kth power, and adding it to a sum.

number, and the function should return true. Otherwise, it should return false.

- to the kth power, and adding it to a sum.

 3. After processing a digit, we can remove it from n by using integer division by 10 (// 10) to shift all the digits to the right.
- 4. Continue this process until all digits have been processed (i.e., until n becomes 0).

 5. Finally, compare the obtained sum with the original number n, and if they are equal, n is an Armstrong number; otherwise, it is not.

This approach systematically deconstructs the number n to access its individual digits while keeping track of the sum of the

digits each raised to the kth power. If at the end, this sum is equal to n, then the number is confirmed to be an Armstrong

iterates through the digits of the number while computing a sum that follows the rule for Armstrong numbers.

Solution Approach

The implementation of the solution for determining if a number is an Armstrong number involves a straightforward approach that

1. First, find the number of digits k in our number n. This is done by converting n to a string and getting the length of that

s += (x % 10) ** k

indeed an Armstrong number:

the working copy of the original number.

elementary arithmetic and control flow constructs in Python.

manipulate without altering the original number 371:

x //= 10 # x becomes 37 after the first iteration

3), and add the result to the sum s:

x //= 10

return s == n

Example Walkthrough

string:

k = len(str(n))

2. Initialize two variables: one for the sum (s) which we will use to accumulate the k-th powers of the digits, and a copy of n (x) which we will use to manipulate and access each digit without altering the original number n:

Here is the step by step breakdown of the algorithm:

which we will use to manipulate and access each digit without altering the original number n: s, x = 0, n

while x:

Inside the loop, extract the last digit of x using % 10, raise it to the power of k, and add the result to the sum s:

5. Then remove the last digit from x by dividing x by 10 and keeping the integer part only (effectively shifting the digits to the right):

Use a while loop to iterate over each digit in the number. The loop continues as long as x is not zero:

After finishing the loop, we check if the calculated sum s is equal to the original number n. If they match, it means that n is

The pattern followed here is an iterative approach that decomposes the number digit by digit to perform the necessary arithmetic operations, a common technique when dealing with numbers in programming.

This solution does not require any complex algorithms or data structures and efficiently arrives at the answer by harnessing

The data structures used in this approach are quite simple - just integers for keeping track of the sum, the number of digits, and

Let us use the number 371 as an example to illustrate the solution approach.

1. First, we find the number of digits k in our number n, which is 371. We convert 371 to a string and get the length of that string:

k = len(str(371)) # k is 3 because there are 3 digits in the number

We initialize a variable for the sum s at 0 to accumulate the k-th powers of the digits, and we make a copy of n (x) to

3 We use a while

Python

Java

class Solution {

class Solution {

bool isArmstrong(int n) {

return sumOfPowers == n;

int numDigits = std::to string(n).size();

// Iterate over the digits of the number

for (int temp = n; temp != 0; temp /= 10) {

sumOfPowers += std::pow(temp % 10, numDigits);

public:

int sum = 0;

class Solution:

def is armstrong(self, n: int) -> bool:

sum_of_powers, temp_number = 0, n

temp_number //= 10

return sum_of_powers == n

public boolean isArmstrong(int number) {

Remove the last digit from temp_number

int numOfDigits = String.valueOf(number).length();

k = len(str(n))

s, x = 0, 371

3. We use a while loop to iterate over each digit in the number. The loop continues as long as \times (which begins as 371) is not zero:

while x: # True initially because x is 371

Inside the loop, we extract the last digit of x (which is 1 in the first iteration) using % 10, raise it to the power of k (which is

s += (x % 10) ** k # Adds 1**3 to s, so s becomes 1

Since they match, it means that 371 is indeed an Armstrong number:

return s == n # returns True, as 371 is indeed an Armstrong number

and added to s, resulting in s being 1 + 7**3 which is 1 + 343 making s now 344.

7. With x updated to 3 (37 // 10), the loop goes for one more iteration with 3 raised to the power of 3 and added to s, resulting in 344 + 3**3, which equals 344 + 27 making s 371.

The loop continues with x now being 37, and in the second iteration, the last digit 7 (37 % 10) will be raised to the power of 3

After finishing the loop (since x is now 0), we check if the calculated sum s (371) is equal to the original number n (371).

By following these steps, we have checked that 371 is an Armstrong number using the method described in the solution

approach. The while loop terminated after processing all digits of number 371, raising each digit to the power of the number of

We remove the last digit from x by dividing x by 10 and keeping the integer part only, which shifts all digits to the right:

digits, summing them, and comparing the result with the original number.

Solution Implementation

Convert the number to a string and calculate the length to determine k

Initialize the sum and a temporary variable with the original number

Check if the sum of the digits to the power of k is equal to the original number

// Initialize sum to store the sum of digits raised to the power of 'numOfDigits'

// Convert the integer n to a string to compute its length (number of digits)

int sumOfPowers = 0; // This will hold the sum of the digits raised to the power of numDigits

// Calculate the current digit to the power of numDigits and add it to sumOfPowers

// Return true if the sum of the powers equals the original number, false otherwise

Check if the sum of the digits to the power of k is equal to the original number

// Calculate the number of digits in 'number' by converting it to a string and getting its length

Loop through each digit of the number
while temp number > 0:
 # Get the last digit and raise it to the power of k
 last digit = temp number % 10
 sum_of_powers += last_digit ** k

```
// Temporary variable to hold the current number as we compute the Armstrong sum
        int tempNumber = number;
       // Loop to calculate the sum of each digit raised to the power numOfDigits
       while (tempNumber > 0) {
           // Get the last digit of tempNumber
            int digit = tempNumber % 10;
            // Add the digit raised to the power of 'numOfDigits' to the sum
            sum += Math.pow(digit, numOfDigits);
            // Remove the last digit from tempNumber
            tempNumber /= 10;
       // Check if the calculated sum is equal to the original number
        return sum == number;
C++
                        // Include cmath for the pow function
#include <cmath>
#include <string>
                        // Include string to convert integers to strings
```

function isArmstrong(number: number): boolean { // Calculate the length of the number to de const numberOfDigits = String(number).lengt

TypeScript

};

```
// Calculate the length of the number to determine the power to be used.
   const numberOfDigits = String(number).length;
   // Initialize the sum of the digits raised to the power of the number of digits.
   let sum = 0;
   // Initialize a variable 'value' to iterate through the digits of 'number'.
   for (let value = number; value; value = Math.floor(value / 10)) {
       // Add the current digit raised to the 'numberOfDigits' power to the sum.
       sum += Math.pow(value % 10, numberOfDigits);
   // Return true if the sum is equal to the original number (Armstrong number), otherwise return false.
   return sum === number;
class Solution:
   def is armstrong(self, n: int) -> bool:
       # Convert the number to a string and calculate the length to determine k
       k = len(str(n))
       # Initialize the sum and a temporary variable with the original number
       sum_of_powers, temp_number = 0, n
       # Loop through each digit of the number
       while temp number > 0:
           # Get the last digit and raise it to the power of k
           last digit = temp number % 10
```

The given Python code defines a method isArmstrong which determines whether a given integer n is an Armstrong number or

not. An Armstrong number (also known as a narcissistic number) of k digits is an integer such that the sum of its own digits each

raised to the power of k equals the number itself. To analyze its computational complexity, let's consider the following:

Time Complexity:

Space Complexity:

sum_of_powers += last_digit ** k

temp_number //= 10

return sum_of_powers == n

Time and Space Complexity

Remove the last digit from temp_number

The while loop runs as long as x is not zero. Since x is reduced by a factor of 10 in each iteration, there will be 0(d) iterations.
Within each iteration, the modulus operation x % 10, the power operation (x % 10) ** k and integer division x //= 10 are performed. The

• Determining the number of digits, k = len(str(n)), requires O(d) time.

• Let d be the number of digits in the integer n, which is equal to len(str(n)).

- power operation has a time complexity of O(k) because it involves multiplying the number x % 10 k times.

 Therefore, the total time complexity within the while loop is O(d * k).
- Since k is itself 0(d) because the number of digits d is log10(n), and k $\approx log10(n)$, the overall time complexity when combining these operations is $0(d^2)$.
- No additional data structures that grow with input size are used.
 The variables themselves only occupy constant space, 0(1).
 Thus, the overall space complexity of the method is 0(1).

• The space complexity is determined by the extra space used by the variables k, s, and x.

The provided Python code's time complexity is $0(d^2)$ and space complexity is 0(1).