



Problem Description

to left, handling carry as needed.

the array represents a single digit of the number in the correct order. For instance, if num = 1321, then the array form would be [1,3,2,1]. If we were to add k = 9 to this number, we're trying to find the array form of 1321 + 9, which would result in [1,3,3,0]. This is akin to how we perform addition by hand, starting from the rightmost digit (the least significant digit) and adding it to the

The problem is to implement the addition of a non-negative integer k to a number represented as an array, where each element of

corresponding digit of the other number (if available), carrying over the excess to the next digit on the left if the sum exceeds 9.

This solution approach leverages the simple idea of adding two numbers together just as we would on paper, digit by digit from right

Intuition

We initialize a pointer i to the last index of the num array and set carry to 0. In a loop, we perform the following until we have exhausted all digits in the num array and fully added k and accounted for any remaining carry:

• We calculate the total sum for the current digit by adding the appropriate digit from num (or 0 if we've passed the beginning of num), the last digit of k, and any carry from the previous step.

- We then use the divmod() function to both determine the digit to append to our answer in reverse (v) and the new carry value, ensuring that we only carry over when the sum is 10 or greater.
- Append v to the answer array ans.
- Update the value of k to reflect the leftover part after extracting the last digit by performing integer division by 10. Decrement the index i to move to the next digit to the left.
- Once the loop concludes (all digits and the carry have been processed), we reverse the ans array to represent the correct number as the digits have been appended in reverse order, starting with the least significant digit. This reversed ans array represents the num +

numbers. Here is the step-by-step break down:

k in array form, which is what we return as the solution.

Solution Approach The implementation of the solution is a straightforward simulation of the addition operation that we perform manually between

1. Initialize the index i to point to the last digit of the input array num which represents the least significant digit of the number we are given. Also, initialize a carry variable to 0, which will keep track of any carryover during the addition process.

2. Create an empty list ans that will store the result of the addition in reverse order.

4. The sum for the current position is calculated by adding the current digit of num array, if any (num[i] or 0 if i is less than 0), the last digit of k (k % 10), and the current carry. This is done using the line:

3. Execute a while loop with the condition to continue as long as there is a digit remaining in the num array (i >= 0), or there is still

1 carry += (0 if i < 0 else num[i]) + (k % 10)</pre>

5. The carry and the value to append to the result list, v, are determined by using divmod():

a part of k that hasn't been added (k), or there is a carry from the previous addition (carry).

- 1 carry, v = divmod(carry, 10) This operation returns the quotient that becomes the next carry and the remainder v which is the digit to be appended to the
- answer.
- 6. Append v to the ans list:

significant digit. Therefore, we reverse ans to correct the order:

to find the array form of 256 + 34, which would result in [2,9,0].

2. An empty list ans is created to store the result.

1 carry, v = divmod(10, 10) # carry = 1, v = 0

We append 0 to ans (since v = 0) and adjust k and i:

9. After the loop exits, the answer list ans contains the digits of the result in reversed order since we started adding from the least

10. Return the reversed ans list as the final output, which now represents the num + k in the correct array-form.

In summary, the algorithm uses simple arithmetic and an iterative approach to simulate the addition process with careful

8. Decrement the index i to move the pointer to the left, preparing for the next iteration to add the next higher order digit.

7. The next digit of k is obtained by reducing k by one decimal place using the floor division operation:

bookkeeping of carries. The data structure used is a list to store the digits of the result which is eventually reversed to match the expected output format.

Suppose we have the array num = [2,5,6] which represents the number 256 and we want to add k = 34 to this number. We expect

Example Walkthrough Let's take a small example to illustrate the solution approach.

1 ans.append(v)

1 k //= 10

1 return ans[::-1]

Here is how the algorithm would process this example: 1. Initialize the index i to 2 since the last digit of num is at index 2. Also, the carry variable is initialized to 0.

1 carry += num[2] + k % 10 # carry += 6 + 4 carry now becomes 10. The next step is to split this into carry and v:

1 k //= 10 # k becomes 3 2 i -= 1 # i becomes 1

1 carry += num[1] + k % 10 # carry += 5 + 3

5. The next iteration starts with i = 1, carry = 1, and k = 3. The sum is:

carry now becomes 9. Again, we split this into carry and v:

1 carry, v = divmod(9, 10) # carry = 0, v = 9

1 carry += num[0] + k % 10 # carry += 2 + 0

carry now becomes 2. Split into carry and v:

1 carry, v = divmod(2, 10) # carry = 0, v = 2

We append 9 to ans and adjust k and i:

3. Enter the while loop which continues as long as $i \ge 0$, k > 0, or carry is not 0.

4. In the first iteration, i is 2, there is no carry, and k is 34. So, we calculate:

1 k //= 10 # k becomes 0 2 i -= 1 # i becomes 0

6. The final iteration starts with i = 0, carry = 0, and k = 0. We perform the addition:

9. The final output is [2, 9, 0], which is the correct array representation of 256 + 34.

adds two numbers when one is given as an array and the other as an integer (k).

public List<Integer> addToArrayForm(int[] num, int k) {

// Also add the last digit of k to carry

carry += (n < 0 ? 0 : num[n--]) + k % 10;

while $(n >= 0 \mid \mid k > 0 \mid \mid carry > 0) {$

result.addFirst(carry % 10);

carry /= 10;

k /= 10;

return result;

int carry = 0; // To hold carry-over values during addition

// If n is within bounds, add num[n] to carry; else add 0.

// Add the unit's place of carry to the front of result

// Remove the unit's place digit from carry and k

We append 2 to ans: 1 ans.append(v) 7. The loop finishes since i is now less than 0, k is 0, and carry is also 0. The result list ans is currently [0, 9, 2].

def addToArrayForm(self, num: List[int], k: int) -> List[int]: current_index, carry_over = len(num) - 1, 0 result = [] # Loop until we've processed each digit of `num` and `k`

8. We reverse ans to get the correct order:

1 ans[::-1] # becomes [2, 9, 0]

Python Solution

class Solution:

class Solution {

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from typing import List

```
# as well as any remaining carry over value.
           while current_index >= 0 or k or carry_over:
               # If the current index is valid, extract the digit from `num`,
11
12
               # otherwise use 0. Add the last digit of `k` and the carry over.
13
               carry_over += (num[current_index] if current_index >= 0 else 0) + (k % 10)
               # Use divmod to get the new carry over and the digit to be added to the result.
14
15
               carry_over, digit = divmod(carry_over, 10)
               result.append(digit)
16
               # Update `k` and `current_index` for the next iteration.
17
18
               k //= 10
19
               current_index -= 1
20
21
           # The `result` list is currently in reverse, as digits were added from
22
           # the least significant digit (rightmost). We need to reverse it before returning.
23
           return result[::-1]
24
Java Solution
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int n = num.length - 1; // Initialize the index to the last element of the input array

LinkedList<Integer> result = new LinkedList<>(); // Using LinkedList to utilize addFirst method

// Loop until we've processed all elements of num, or until k becomes 0, or until there's no carry left

This walkthrough demonstrates each step of the algorithm on a small example, successfully showing how the solution approach

C++ Solution

1 #include <vector>

2 #include <algorithm>

class Solution {

5 public:

```
// Adds an integer k to a large integer represented as a vector of digits.
       vector<int> addToArrayForm(vector<int>& num, int k) {
           int numIndex = num.size() - 1; // Start from the last digit of the number.
           int carry = 0; // Initialize carry for addition.
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           vector<int> result; // The result vector to store the sum.
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12
           // Iterate until all digits are processed, or there is no carry, or k is not zero.
13
           while (numIndex >= 0 || k > 0 || carry > 0) {
               // If numIndex is within bounds, add the current digit to carry.
14
15
               // Otherwise, add 0 (when numIndex < 0).
               carry += (numIndex >= 0 ? num[numIndex] : 0) + (k % 10);
16
               result.push_back(carry % 10); // Extract the last digit of carry and add it to result.
17
               carry /= 10; // Update carry for the next iteration.
               k /= 10; // Move to the next digit in k.
19
20
               numIndex--; // Move to the previous digit in num.
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23
           // The result currently contains the digits in reverse order.
24
           reverse(result.begin(), result.end()); // Reverse the result to get the correct order.
25
26
           return result;
27
28 };
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Typescript Solution
   function addToArrayForm(numArray: number[], k: number(): number[] {
```

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// Convert 'k' to an array of its digits

let answer: number[] = [];

let kArray: number[] = [...String(k)].map(Number);

// Initialize the answer array to hold the result

// Sum is used to handle the current value while adding digits

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let carry: number = 0;
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       // Process while there are digits in 'numArray' or 'kArray', or if there is a carry value
11
       while (numArray.length | | kArray.length | | carry) {
           // Pop the last digit from 'numArray' or use 0 if empty
           let numDigit = numArray.pop() || 0;
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16
           // Pop the last digit from 'kArray' or use 0 if empty
           let kDigit = kArray.pop() || 0;
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18
19
           // Add the digits and the carry
20
           carry += numDigit + kDigit;
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22
           // Unshift (insert at the beginning) the last digit of the sum into 'answer'
           answer.unshift(carry % 10);
24
25
           // Floor division by 10 to get a new carry value
           carry = Math.floor(carry / 10);
26
27
28
29
       // Return the result stored in 'answer'
30
       return answer;
31 }
32
Time and Space Complexity
Time Complexity
```

Space Complexity

num even if k is already reduced to zero.

The space complexity of the code is $O(\max(N, \log K))$ because we create a list ans that stores the result, which in the worst case will have max(N, logK) digits (the maximum out of the number of digits in num and the number of digits in k after addition). Furthermore, space is used for variables i, carry, v, and the stack space for the recursive call when reversing the answer list, but these do not depend on the size of the inputs and thus contribute only a constant factor, which is negligible when analyzing space complexity.

The time complexity of the code is O(max(N, logK)) where N is the number of digits in num and logK represents the number of digits

in k. This is because the while loop runs once for each digit of the number num and for each digit of the number k. When k has more

digits than num, the loop will continue until all digits of k have been processed. In the reverse situation, it will process all the digits of