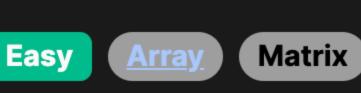
## 1572. Matrix Diagonal Sum



#### **Problem Description**

to calculate the sum of the diagonal elements, which includes elements from both the primary diagonal and the secondary diagonal. The primary diagonal is the one that starts from the top left corner and ends at the bottom right corner. The secondary diagonal starts from the top right corner and ends at the bottom left corner. However, there's a catch: If any element is common between the primary and secondary diagonals (which would be the case for the central element in a matrix with odd dimensions), we must include it only once in our sum.

The problem provides us a square matrix mat, which means the number of rows and columns in the matrix are equal. Our task is

Intuition

The primary diagonal elements have the same index for their row and column. In terms of indexes, these are the elements

To approach this problem, we consider the primary and secondary diagonals of the matrix.

- mat[i][i] where i ranges from 0 to the n-1, wherein n is the size of one dimension since it's a square matrix. The secondary diagonal elements have row and column indexes that sum up to n-1. In other words, the indexes are of the
- form mat[i][j] where j is n-i-1. Now, the challenge is to make sure that we don't double-count the element in the case of an odd-dimension matrix where the
- primary and secondary diagonals intersect. To avoid this, we simply check if the index i is equal to the index j. If they are equal, this means we're looking at the central element in the case of an odd-size matrix and we shouldn't add it again.

Therefore, the sum ans starts at 0, and we iterate through each row with its index i. While iterating, we calculate j as n-i-1 for each row to pinpoint the element in the secondary diagonal. We then add to ans the sum of elements mat[i][i] and mat[i][j], unless i == j, in which case we only add the element mat[i][i] once.

Solution Approach

### diagonals of a square matrix.

Here's a detailed walkthrough of the implementation: We start by initializing a variable ans to 0. This will hold the cumulative sum of the diagonal elements.

The solution provided is written in Python and follows a straightforward approach that efficiently computes the sum of the

= len(mat).

- The for loop is used to iterate over each row of the matrix. The loop variable i serves as the index for both rows and the primary diagonal elements. enumerate is used so we can have both the index and the row elements available during each iteration.

We then calculate the column index j for the secondary diagonal element corresponding to the current row. As explained

We need the size of the matrix, which is the length of one of its sides (since it is square), so we take the length of the matrix n

Inside the loop, we add to ans the value of the primary diagonal element row[i]. We also want to add the secondary diagonal element row[j] unless i is equal to j (which is the case when we are at the

central element of a matrix with an odd number of rows/columns). To do this succinctly, we add row[j] only if j = i,

After the loop has processed all the rows, ans now contains the total sum of both diagonals, with no duplicates for the intersecting element (if any).

otherwise, we add  $\emptyset$ . This conditional addition is achieved by the expression ( $\emptyset$  if j == i else row[j]).

- There are no complex algorithms, data structures, or patterns involved here. The solution effectively uses index manipulation to address the specific elements required for the sum, making it a simple yet effective approach to solving this problem.
- n = len(mat)for i, row in enumerate(mat):

Here is the core logic encapsulated in Python code:

def diagonalSum(self, mat: List[List[int]]) -> int:

class Solution:

[1, 2, 3],

[4, 5, 6],

[7, 8, 9]

ans = 0

Finally, the method returns the computed ans.

earlier, j is given by n - i - 1.

j = n - i - 1ans += row[i] + (0 if j == i else row[j]) return ans

By carefully choosing and manipulating indexes, we maintain a clear and concise solution without additional memory usage for

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storing intermediate results or redundant computations.
Example Walkthrough
  Let's walk through an example to illustrate the solution approach. Suppose we have the following 3×3 square matrix:
mat = [
```

Start the for loop with i iterating from 0 to 2 (since the matrix has 3 rows and columns). For each row i, calculate the index j for accessing the secondary diagonal's element, which is n - i - 1. This will give us 2,

1, 0 for i = 0, 1, 2 respectively.

ans = 0

n = len(mat)

return ans

Solution Implementation

total\_sum = 0

**Python** 

class Solution:

for i, row in enumerate(mat):

j = n - i - 1

Here's the step-by-step process of how the solution works:

Now, add the primary diagonal element to ans. In the first iteration with i = 0, add mat [0] [0] which is 1.

Initialize ans to 0. This will keep track of the sum of diagonal elements.

Determine the size n of the matrix. In this case, n = len(mat) = 3.

equal, so add mat[0][2] which is 3. For the second iteration where i = 1, we add mat [1] [1] (the middle element) to ans. Since i and j are equal here (both equal

diagonals of the matrix, without double-counting the center element.

ans += row[i] + (0 if j == i else row[j])

efficiency and simplicity of the provided solution approach.

def diagonalSum(self, matrix: List[List[int]]) -> int:

# Loop over each row and calculate the diagonal sum

int totalSum = 0; // This will hold the sum of the diagonal elements

int size = matrix.length; // The matrix is size x size

# Initialize the sum of the diagonals

to 1), we don't add the secondary diagonal element because that would be double-counting.

(because j = 0), so add the secondary diagonal element mat [2] [0] which is 7. After adding these elements through the loop, the sum ans becomes 1 + 3 + 5 + 9 + 7 = 25.

The method returns the value of ans, which is 25. This is the required sum of the elements on the primary and secondary

In the third iteration with i = 2, add the primary diagonal element mat [2] [2] which is 9 to ans. Index i and j are not the same

Check if i equals j. If they are not the same, add the secondary diagonal element to ans. With i = 0 and j = 2, they aren't

- Putting it all into the Python code, we have: class Solution: def diagonalSum(self, mat: List[List[int]]) -> int:
  - When we call diagonalSum(mat) with our example matrix, it will return 25, which is the correct answer. This demonstrates the

# Get the size of the matrix (assuming it's square) n = len(matrix)

int reverseIndex = size -i - 1; // Calculate the corresponding column index for the secondary diagonal

// If it's not the same element (which would be the case in the middle of an odd-sized matrix)

// Return the sum of primary and secondary diagonals, excluding the middle element if counted twice

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for i, row in enumerate(matrix):
           # Calculate the index for the secondary diagonal
             = n - i - 1
            # Add the primary diagonal element
            total_sum += row[i]
           # Add the secondary diagonal element if it's not the same as the primary diagonal
            if j != i:
                total_sum += row[j]
       # Return the computed sum
       return total_sum
Java
class Solution {
```

public int diagonalSum(int[][] matrix) {

for (int i = 0; i < size; ++i) {

totalSum += matrix[i][i];

if (i != reverseIndex) {

function diagonalSum(matrix: number[][]): number {

// Iterate through each row of the matrix

for (let row = 0; row < matrixSize; row++) {</pre>

const matrixSize = matrix.length;

// 'matrixSize' stores the size of the matrix (number of rows/columns)

// Initialize 'sum' to zero, which will store the final diagonal sum

// Add the elements from both the primary and secondary diagonals for the current row

return totalSum;

// Loop through each row of the matrix

// Add the primary diagonal element

// then add the secondary diagonal element

totalSum += matrix[i][reverseIndex];

```
C++
#include<vector>
class Solution {
public:
   // Function to calculate the sum of the elements on the diagonals of a square matrix
    int diagonalSum(std::vector<std::vector<int>>& mat) {
        int total = 0; // Used to store the sum of the diagonal elements
        int size = mat.size(); // Get the size of the square matrix
       // Iterate through each row of the matrix
        for (int rowIndex = 0; rowIndex < size; ++rowIndex) {</pre>
            int colIndex = size - rowIndex - 1; // Calculate the column index for the secondary diagonal
           // Sum the primary diagonal element
            total += mat[rowIndex][rowIndex];
           // Sum the secondary diagonal element only if it's not the same as the primary diagonal element
            if (rowIndex != colIndex) {
                total += mat[rowIndex][colIndex];
        return total; // Return the sum of the diagonal elements
};
TypeScript
```

let sum = 0;

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sum += matrix[row][row] + matrix[row][matrixSize - 1 - row];
      // If the matrix size is odd, subtract the central element once to correct the sum
      if (matrixSize % 2 === 1) {
          // The central element is at the position ['matrixSize' / 2] in both dimensions
          sum -= matrix[matrixSize >> 1][matrixSize >> 1];
      // Return the calculated sum of the diagonal elements
      return sum;
class Solution:
   def diagonalSum(self, matrix: List[List[int]]) -> int:
       # Initialize the sum of the diagonals
        total_sum = 0
       # Get the size of the matrix (assuming it's square)
       n = len(matrix)
       # Loop over each row and calculate the diagonal sum
       for i, row in enumerate(matrix):
           # Calculate the index for the secondary diagonal
           j = n - i - 1
           # Add the primary diagonal element
           total sum += row[i]
           # Add the secondary diagonal element if it's not the same as the primary diagonal
           if j != i:
               total_sum += row[j]
       # Return the computed sum
       return total_sum
Time and Space Complexity
```

#### **Time Complexity** The provided code traverses each row only once, and within each row, it accesses two elements directly by their index, which is

an (O(1)) operation. Since there are n rows in a square matrix with size (n \times n), the overall time complexity of the code is (O(n)) where (n) is the number of rows (and also the number of columns) in the matrix. **Space Complexity** 

# The code uses a fixed number of variables (ans, n, i, j, row). It does not depend on the size of the input matrix, therefore the

space complexity is (O(1)), that is, constant space complexity.