2610. Convert an Array Into a 2D Array With Conditions

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Problem Description
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<u>Array</u>

Hash Table

specific conditions. The 2D array must use all the integers from nums and arrange them so that each row contains only unique integers, without any repetitions. Additionally, we need to minimize the number of rows in the 2D array. The final 2D array can have varying numbers of elements per row, and if there are several correct ways to create it, any valid answer is acceptable.

In this problem, we're given an integer array nums, and our goal is to transform this array into a two-dimensional (2D) array with

number of times.

Intuition The solution leverages the idea of distribution: we want to spread out the occurrence of each integer in nums across different

Medium

rows of the resulting 2D array. Here's the thought process leading to the solution: • First, we need to know how many times each number appears in nums. To do this efficiently, a Counter is used, which is essentially a hash table (dictionary) where each key is a unique number from nums and its corresponding value is the count of how many times it appears.

• Since each row must contain distinct integers, the number of times an integer appears dictates how many rows it will have to be spread across.

• The next step is to iterate through each unique number and its count, and add it to a new row in the 2D array until we've placed it the required

• To maintain the minimum number of rows, we only add a new row when it's necessary – that is, when the number of occurrences of the current

number exceeds the current number of rows in the ans array. • We'll append the number to increasingly indexed rows until we've placed it accordingly (based on its count). Solution Approach

The implemented solution uses a hash table and array manipulation to satisfy the problem's requirements.

Here's a step-by-step breakdown of the implementation: Counting Elements: We start by creating a Counter from the nums array. This Counter acts as a hash table that maps each

unique integer in **nums** to the number of its occurrences. cnt = Counter(nums)

for x, v in cnt.items():

for i in range(v):

Preparing the Answer Array: We initialize an empty list ans, which will become the 2D array we must return.

ans = []

We check if the current row i exists in ans. If not, we create a new row (an empty list).

Finally, we return the ans list, which now represents the 2D array with the desired properties:

We create a Counter from the **nums** array that gives us the frequency of each unique number:

Let's walk through a small example to illustrate the solution approach. Suppose our input array nums is:

of occurrences), x is added to each subsequent row.

We leverage the hash table for efficient lookups and counting.

We use simple list manipulations to build the answer.

Row Management: During the iteration for each unique element:

Populating the 2D Array: We iterate over each unique element x and its occurrence count v in the Counter.

if len(ans) <= i:</pre> ans.append([])

• We ensure minimal rows by checking if a row exists before adding to it, and only when necessary do we create a new row. This approach ensures that each row will contain distinct elements from nums, and we only increase the total row count when the frequency of an integer necessitates an additional row, thus achieving the minimization of rows for our 2D array.

• We then append x to the proper row, filling the 2D array such that each row will have distinct integers. Since we iterate v times (the count

Example Walkthrough

nums = [5, 5, 6, 6, 6, 7]

2. Preparing the Answer Array

return ans

ans[i].append(x)

By following this algorithm:

We will apply each step of the solution approach to this array. 1. Counting Elements

This tells us that the number 6 appears 3 times, 5 appears 2 times, and 7 appears 1 time.

We iterate over the elements and their counts in the cnt dictionary.

We initialize an empty list ans to represent our 2D array.

For the number 6 (which appears 3 times):

Next, for the number 5 (which appears 2 times):

Lastly, for the number 7 (which appears 1 time):

Since ans has 0 rows, we add a new row and insert 6.

cnt = Counter(nums) # Output: Counter($\{6: 3, 5: 2, 7: 1\}$)

4. Row Management

ans = [[6]]

ans = [[6], [6]]

ans = [[6], [6], [6]]

ans = [[6, 5], [6], [6]]

ans = [[6, 5, 7], [6, 5], [6]]

ans = [[6, 5, 7], [6, 5], [6]]

Solution Implementation

from collections import Counter

only unique integers.

Python

class Solution:

3. Populating the 2D Array

ans = []

There are already 3 rows, so we place 7 into the first row that does not have 7.

For each unique element, we go through the count of its occurrences and manage rows accordingly.

Still needs to place 6 two more times, check next row which is empty, add a new row and insert 6.

Still needs to place 6 one more time, check next row which is empty, add a new row and insert 6.

ans = [[6, 5], [6, 5], [6]]

There are already 3 rows, so we place 5 into the first row.

Needs to place 5 one more time, we place it into the second row.

Final 2D Array The final ans array representing our 2D array is:

We have now placed each number from nums keeping the rows distinct and minimized the number of rows in the process.

We return this array as our solution. It has the minimum number of rows, contains all integers from nums, and each row contains

matrix = []# Iterate through the counted numbers and their counts for number, count in num counter.items(): # Loop for the count number of times for each number

for i in range(count):

Return the resulting matrix

int[] count = new int[n + 1];

return matrix

int n = nums.length;

for (int num : nums) {

++count[num];

if len(matrix) <= i:</pre>

matrix.append([])

matrix[i].append(number)

public List<List<Integer>> findMatrix(int[] nums) {

// Find the length of the input array.

for (int num = 1; num <= n; ++num) {</pre>

for (int i = 0; i < frequency; ++i) {</pre>

result.add(new ArrayList<>());

if (result.size() <= i) {</pre>

result.get(j).add(num);

if (matrix.size() <= j) {</pre>

matrix[j].push_back(num);

// Return the organized matrix

function findMatrix(nums: number[]): number[][] {

// Calculate the length of the input array.

// Each index corresponds to a value from the input array.

// Iterate through the possible numbers in the input array.

for (let i = 0; i < frequencyCounter[num]; ++i) {</pre>

def findMatrix(self, nums: List[int]) -> List[List[int]]:

Count the occurrences of each number in the input list

Initialize an empty list to store the resulting matrix

Iterate through the counted numbers and their counts

Loop for the count number of times for each number

If the matrix has fewer rows than needed, add a new row

for number, count in num counter.items():

matrix.append([])

const frequencyCounter: number[] = new Array(length + 1).fill(0);

// Initialize the answer matrix.

const answerMatrix: number[][] = [];

const length: number = nums.length;

for (const num of nums) {

++frequencyCounter[num];

for (let num = 1; num <= length; ++num) {</pre>

if (answerMatrix.length <= j) {</pre>

answerMatrix.push([]);

answerMatrix[j].push(num);

// Return the constructed matrix.

num_counter = Counter(nums)

for i in range(count):

if len(matrix) <= i:</pre>

return answerMatrix;

matrix = []

class Solution:

from collections import Counter

return matrix;

};

TypeScript

matrix.push_back(std::vector<int>());

// Initialize a counter array of length `length + 1` to keep track of the frequency of each number.

// Count the frequency of each number in input array and update the frequencyCounter array.

// For each number, we create a row in the matrix for the number of times it appears.

// If we don't have enough rows in the answerMatrix, add a new empty row.

// Append the number to the respective row in the matrix.

// Add the current number to the j-th row

int frequency = count[num];

List<List<Integer>> result = new ArrayList<>();

// Initialize a list to hold the final groups of numbers.

// Create an array to keep track of the count of each number.

// For each occurrence of the number, place it into a sub-list.

// If the current sub-list doesn't exist, create it.

// Add the current number to the appropriate sub-list.

// Count the occurrences of each number in the input array.

num_counter = Counter(nums)

def findMatrix(self, nums: List[int]) -> List[List[int]]:

Count the occurrences of each number in the input list

Initialize an empty list to store the resulting matrix

If the matrix has fewer rows than needed, add a new row

Append the current number to the i-th row of the matrix

// Iterate over the possible numbers from 1 to n and organize them into the result list.

Java

class Solution {

// Return the list of lists containing grouped numbers. return result; C++ #include <vector> class Solution { public: // Function that rearranges numbers into a sorted matrix based on their frequency std::vector<std::vector<int>> findMatrix(std::vector<int>& nums) { std::vector<std::vector<int>> matrix; // Will hold the final sorted matrix int n = nums.size(); std::vector<int> count(n + 1, 0); // Initialize counting vector with zeros // Count how many times each number appears in the input vector for (int num : nums) { ++count[num]; // Iterate through each unique number in the input array for (int num = 1; num <= n; ++num) {</pre> int frequency = count[num]; // Get the frequency of the current number // Construct rows of the matrix based on the frequency of the current number for (int j = 0; j < frequency; ++j) { // If there are not enough rows in the matrix, add a new row

Append the current number to the i-th row of the matrix matrix[i].append(number) # Return the resulting matrix return matrix

Time and Space Complexity

space complexity is also O(n).

count to build the ans list. Counting the frequency of each element using Counter(nums) can be associated with a time complexity of O(n), where n is the

length of the array nums. This operation involves going through all elements once to determine their frequencies.

The time complexity of the code is primarily determined by two factors: counting the elements in nums and then iterating over the

After counting, the code iterates over the count dictionary and, for each element x, it appends x to the lists in ans v times (where v is the frequency of x). Since the total number of append operations is equal to the length of the nums list (every

number from the list is appended exactly once), this also has a time complexity of O(n). Therefore, the overall time complexity of the code is O(n). As for space complexity, we are using additional data structures: a dictionary for the counts and a list of lists for the ans. Since

both the dictionary and the list of lists will at most store n entries (each corresponding to an element in the original nums list), the