1221. Split a String in Balanced Strings

Greedy **String** Counting Easy

Problem Description A balanced string is a string that contains an equal number of 'L' and 'R' characters. The problem asks to split a given balanced string s into the maximum number of balanced substrings. A substring is considered balanced if it contains an equal number of

'L' and 'R' characters. The challenge is to find the maximum number of balanced substrings that can be obtained from the input string.

To solve this problem, we need to iterate through the string and keep track of the number of 'L' and 'R' characters we've

Intuition

encountered. The idea is to increase a counter every time we encounter 'L' and decrease it every time we encounter 'R'. When this counter returns to zero, it means we have found an equal number of 'L' and 'R' characters, hence a balanced substring. By initializing a variable 1 to zero, we use it as a counter to keep track of the balance between 'L' and 'R' characters. If we

encounter 'L', we increment 1, and if we encounter 'R', we decrement 1. Each time 1 becomes zero, it's an indication that the substring from the start to the current position is balanced, and we increment our answer ans by one.

right, any time 1 is zero, we have found a balanced substring, and we can safely split the string at that point and start counting a new potential balanced substring. The algorithm ends when we've gone through the entire string, and the final value of ans is the maximum number of balanced

This approach works because we are given that the original string s is balanced. Therefore, as we process the string from left to

Solution Approach

The implementation of the solution uses a simple but effective algorithm that requires only a single pass through the string. Its

In Python, we define a class Solution with a method balancedStringSplit that takes a single argument, the string s. The

simplicity is derived from its reliance on a single integer counter and the characteristics of the string being balanced. Let's walk

substrings we can obtain.

through the code to understand this better.

substrings we could form.

if c == 'L':

else:

Example Walkthrough

l += 1

method is structured as follows: We initialize two integers ans and 1 to 0. ans will hold the final count of the balanced substrings, and 1 will be used as a counter to track the balance between 'L' and 'R' characters within a substring.

We iterate over each character c in the string s with a for loop.

- After updating 1 for the current character, we check if 1 has become 0. If it has, this indicates that we've encountered an
- equal number of 'L' and 'R' characters up to the current point in the string, forming a balanced substring. We increment ans to count this new balanced substring.
- Once the loop has finished, we have iterated over the entire string, and ans contains the maximum number of balanced

Inside the loop, we check if c is 'L'. If so, we increment 1, otherwise (meaning c is 'R'), we decrement 1.

- Here is the Python code that implements this algorithm: class Solution:
- def balancedStringSplit(self, s: str) -> int: ans = l = 0for c in s:

l -= 1 if l == 0:

 \circ We encounter 'R': 1 is decremented by 1 (1 = -1).

Next is 'L': 1 is incremented by 1 (1 = 0).

 \circ Then 'R': 1 is decremented by 1 (1 = 1).

Continuation of the sequence:

substrings.

Python

Java

class Solution {

class Solution:

Solution Implementation

balanced count = 0

balance = 0

for char in s:

Encounter 'L': 1 is incremented by 1 (1 = 1).

def balancedStringSplit(self, s: str) -> int:

Iterate over each character in the string

Check if the substring is balanced

Continuing the iteration:

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ans += 1
       return ans
There are no additional data structures needed for this solution, and it works in O(n) time, where n is the length of the string,
because we are making just one pass through it. The space complexity is 0(1) as we are only using a fixed amount of additional
space (the two integers ans and 1).
This algorithm is based on the pattern that a substring is balanced if and only if the number of 'L's is equal to the number of 'R's.
Given that the entire string is balanced, we know that each 'L' will eventually be matched by an 'R'. Therefore, the key insight is to
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We want to determine the maximum number of balanced substrings we can obtain from this string. We initialize $\frac{1}{2}$ ans = 0 (the count of balanced substrings) and $\frac{1}{2}$ = 0 (the balance counter). The string s is "RLLRRLRL".

Let's consider a simple example to illustrate the solution approach. Suppose we are given the balanced string s = "RLLRRLRL".

- \circ We encounter 'L': 1 is incremented by 1 (1 = 1). Next is 'L': 1 is incremented by 1 again (1 = 2).
- And 'R': 1 is decremented by 1 again (1 = 0). As 1 has returned to 0, we increment ans by 1 (ans = 2).

count the number of balanced substrings incrementally as we traverse the string.

As we iterate through the characters of s, we apply the following logic:

At this stage, we have identified another balanced substring: "LLRR". For the remaining characters:

Since 1 is now 0, we have found our first balanced substring "RL". ans is incremented by 1 (ans = 1).

 \circ We have 'L': increment 1 by 1 (1 = 1). \circ Then 'R': decrement 1 by 1 (1 = 0).

Increment ans (now ans = 3) since we've found another balanced substring "LR".

 And finally 'R': 1 is decremented back to 0 (1 = 0). • Increment ans once more (final ans = 4) as the last pair "LR" forms a balanced substring.

Initialize count of balanced substrings and a balance counter

Increment the count of balanced substrings

// Method to count the number of balanced strings in the input string 's'.

int balanceCount = 0; // To store the number of balanced strings found

int balance = 0; // A variable to track the balance between 'L' and 'R' characters

// A balanced string has an equal number of 'L' and 'R' characters.

// When balance is zero, a balanced string is found

characters, thus meeting the criteria for balanced substrings. The method balancedStringSplit would return 4 for the input "RLLRRLRL", and this concludes our example walkthrough using the solution approach.

Now that we've processed the entire string s, ans holds the value 4, which represents the maximum number of balanced

The substring splits we've found are "RL", "LLRR", "LR", and "LR". Each of these substrings has an equal number of 'L' and 'R'

If the character is 'L', increment the balance counter **if** char == 'L': balance += 1 # If the character is 'R', decrement the balance counter else: balance -= 1

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balanced_count += 1
# Return the total count of balanced substrings
return balanced_count
```

if balance == 0:

public int balancedStringSplit(String s) {

for (char c : s.toCharArray()) {

if (c == 'L') {

balance++;

balance--;

} else if (c == 'R') {

// Loop through each character in the string

// Increment balance when 'L' is found

// Decrement balance when 'R' is found

// Return the total number of balanced strings

* Function to count the number of balanced strings that can be split.

* @return {number} - The count of balanced strings that can be split.

// Iterate through the string to check for balanced segments.

// If balance is 0, a balanced segment is found.

* Balanced strings are those where the quantity of 'L' and 'R' characters is the same.

* @param {string} inputString - The string to be checked for balanced splits.

return countBalanced;

for (let character of inputString) {

balance += (character == 'L') ? 1 : -1:

```
if (balance == 0) {
                balanceCount++;
        // Return the total number of balanced strings
        return balanceCount;
C++
class Solution {
public:
    // This function counts how many times the input string can be split
    // into balanced strings, where "balanced" means the number of 'L's and 'R's
    // in the substring are equal.
    int balancedStringSplit(string s) {
        int countBalanced = 0; // To keep track of the count of balanced strings
        int balanceFactor = 0; // To keep track of the balance between 'L' and 'R'
        // Iterate over each character in the string
        for (char c : s) {
            // Increment balance factor for 'L' and decrement for 'R'
            if (c == 'L') {
                balanceFactor++;
            } else { // c == 'R'
                balanceFactor--;
            // When the balance factor is zero, we have a balanced string
            if (balanceFactor == 0) {
                countBalanced++; // Increment the number of balanced strings
```

*/ const balancedStringSplit = (inputString: string): number => { let balancedCount = 0; // Counter for balanced splits let balance = 0; // Helper to keep track of the current balance between 'L' and 'R'

};

/**

TypeScript

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if (balance === 0) {
            balancedCount++;
   return balancedCount; // Return the total number of balanced strings.
};
// Example usage:
// let result = balancedStringSplit("RLRRLLRLRL");
// console.log(result); // Outputs: 4, because there are four balanced substrings "RL", "RRLL", "RL", "RL"
class Solution:
   def balancedStringSplit(self, s: str) -> int:
       # Initialize count of balanced substrings and a balance counter
       balanced count = 0
       balance = 0
       # Iterate over each character in the string
       for char in s:
           # If the character is 'L', increment the balance counter
           if char == 'L':
                balance += 1
           # If the character is 'R', decrement the balance counter
           else:
                balance -= 1
           # Check if the substring is balanced
           if balance == 0:
               # Increment the count of balanced substrings
                balanced_count += 1
       # Return the total count of balanced substrings
       return balanced_count
```

Time and Space Complexity

Time Complexity

each character in the string exactly once, performing a constant amount of work for each character with simple arithmetic operations and a comparison.

The time complexity of the code is O(n), where n is the length of the input string s. This is because the code iterates through

Space Complexity The space complexity of the code is 0(1). The code uses a fixed number of integer variables (ans and 1) regardless of the input

size. These variables take up a constant amount of space and do not scale with the size of the input string.