Problem Description

pairs of indices (i, j) where $0 \ll i, j < arr1$, length. The expression to maximize is: 1 |arr1[i] - arr1[j]| + |arr2[i] - arr2[j]| + |i - j|

Given two integer arrays arr1 and arr2 of the same length, the task is to calculate the maximum value of a specific expression for all

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denotes the absolute value, and the goal is to find the maximum result possible by choosing different values of 1 and 1.
The
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Intuition

The intuitive approach to this problem would be to consider all possible pairs of indices (i, j) and calculate the value of the

must find a smarter way to ascertain the maximum value without directly examining every pair. The important realization is that the expression can be rewritten and analyzed in terms of its components. Notice that the expression comprises three terms: the absolute difference between arr1 elements, the absolute difference between arr2 elements, and the

expression for each pair, but this would lead to a rather inefficient solution with a quadratic runtime complexity. To optimize this, one

We can reformulate each term like this: |arr1[i] - arr1[j]| can be either (arr1[i] - arr1[j]) or -(arr1[i] - arr1[j])
|arr2[i] - arr2[j]| can be either (arr2[i] - arr2[j]) or -(arr2[i] - arr2[j])

Each of these can have two possible signs (+ or -), which gives us a total of 2 * 2 * 2 = 8 combinations. Each combination can be thought of representing a different "direction" or case. However, for this solution, we take into account that 📋 – ϳ is always added

- j| can be either (i - j) or -(i - j)

absolute difference between indices.

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to the expression, thus we only need to consider four combinations for the terms from the arrays arr1 and arr2.
The code uses a pairwise function in combination with a dirs tuple to go through these four combinations (-1 and 1 represent the
possible signs). For each case, it initializes minimum (mi) and maximum (mx) values that could be yielded by the expression when
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evaluating for all indices (i, j). At each step in the loop, the code updates the mx and mi to find the maximum possible value for the current combination and eventually computes the maximum possible value for the overall expression by comparing it against the maximum (ans) of previous cases.

Thus, instead of comparing all pairs of (i, j), we keep track of the max and min value of a * arr1[i] + b * arr2[i] + i for each

direction and use these values to calculate and update the running maximum value over the entire array. Solution Approach The solution utilizes the fact that every absolute difference between two numbers can be represented as either a positive or a

cases. The goal is to calculate the maximum value of our expression within each of these cases. Since we have two arrays and we need to consider positive and negative contributions of their elements independently, we need to

evaluate four cases. The dirs tuple (1, -1, -1, 1) along with pairwise is used to iterate over these cases. The pairwise function is

negative difference. Therefore, for each element in the arrays, we can apply either a positive or a negative sign, leading to different

not built into Python, but it would simply return pairs of elements from dirs, e.g., (1, -1), (-1, -1), (-1, 1). This means we evaluate the following cases for each index 1:

1. arr1[i] - arr2[i] + i

2. -arr1[i] - arr2[i] + i

3. -arr1[i] + arr2[i] + i

4. arr1[i] + arr2[i] + i For each of the four cases, we use a loop to iterate over all indices i of our input arrays.

Within each loop iteration, we apply the current case's signs to elements of arr1[i] and arr2[i] and add the index i. We update mx to be the maximum value seen so far and mi to be the minimum value seen so far. 1 mx = max(mx, a * arr1[i] + b * arr2[i] + i)2 mi = min(mi, a * arr1[i] + b * arr2[i] + i)

The maximum value for this case is then mx - mi, which represents the widest range we've seen between the expression values for

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ans = max(ans, mx - mi)
This line updates the overall maximum value ans with the maximum value obtained in the current case. By the end of the for-loop
that iterates over the pairwise elements, we've considered all possible signs for differences and added the absolute difference of
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any pair of indices within this particular case.

indices, yielding the maximum value of the original expression.

For these arrays, we want to maximize the expression:

1 |arr1[i] - arr1[j]| + |arr2[i] - arr2[j]| + |i - j|

Example Walkthrough Let's illustrate the solution approach with an example. Suppose we have the following arrays:

We notice that as i and j get farther apart, the expression tends to increase because of the |i - j| term, but the first two absolute

First, let's manually calculate the expression for some pairs to understand the problem:

• We don't actually need the pairwise function here because we know our cases in advance:

Finally, we return ans, which represents the maximum value of the expression across all pairs (i, j).

• For i = 0 and j = 2: \circ The expression would be |1 - 3| + |4 - 6| + |0 - 2| = 2 + 2 + 2 = 6

terms can also impact the result.

1 Case 1: arr1[i] - arr2[i] + i

2 Case 2: -arr1[i] - arr2[i] + i

3 Case 3: -arr1[i] + arr2[i] + i

4 Case 4: arr1[i] + arr2[i] + i

• For i = 0 and j = 1:

1 arr1 = [1, 2, 3]

2 arr2 = [4, 5, 6]

1. We will consider four cases by applying the pairwise concept to the dirs tuple.

Let's now use the optimized approach described in the solution:

3. Let's start with Case 1 (arr1[i] - arr2[i] + i)

 \circ For i = 0:1 - 4 + 0 = -3

4. Next is Case 2 (-arr1[i] - arr2[i] + i)

 \circ For i = 0: -1 - 4 + 0 = -5

 \circ For i = 1: -2 - 5 + 1 = -6

 \circ For i = 2: -3 - 6 + 2 = -7

higher than the previous cases.

from itertools import product

def maxAbsValExpr(self, arr1, arr2):

max_absolute_value_expr = -inf

from math import inf

class Solution:

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 \circ The expression would be |1-2|+|4-5|+|0-1|=1+1+1=3

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2. We will iterate over each element i of arr1 and arr2 and calculate the expressions for each of the four cases, updating the
  maximum (mx) and minimum (mi) as we go. We then use these to compute the range mx - mi in each case which is a potential
  maximum for our final answer.
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 \circ mx of Case 2 is -5, mi is -7, the range mx - mi is -5 - (-7) = 2

 \circ For i = 1: 2 - 5 + 1 = -2 \circ For i = 2:3 - 6 + 2 = -1 \circ mx of Case 1 is -1, mi is -3, and the range mx - mi is -1 - (-3) = 2

5. Proceed similarly for Case 3 and Case 4, updating mx and mi for each case, then compute the range mx - mi to see if the result is

6. After evaluating all four cases, the final answer ans would be the maximum range mx - mi found among all cases.

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For our example, the ranges were the same (2) for Case 1 and Case 2, and if we calculate it for Cases 3 and 4, we would get 10 and 8
respectively. Hence, the maximum value for the expression |arr1[i] - arr1[j]| + |arr2[i] - arr2[j]| + |i - j| across all i, j
with 0 <= i, j < arr1.length would be 10.
This example shows how to compute the maximum value effectively without checking each combination of i and j individually. By
considering the varied outcomes of applying both the positive and negative signs, we manage to find the optimal result with a more
efficient approach.
Python Solution
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Initialize maximum absolute value expression to a very small number

Iterate over all combinations of directions for arr1 and arr2.

Iterate over the pairs (i, (x, y)) where i is the index,

x is the element from arrl, and y is the element from arr2

Calculate the current value using the expression given by the problem

max_absolute_value_expr = max(max_absolute_value_expr, max_val - min_val)

// Initialize variable 'maxDifference' to keep track of the maximum absolute value

// Iterate over the four possible combinations of expressions represented by multipliers.

Each combination represents multiplying arr1 and arr2 with

Initialize maximum and minimum values observed

1 or -1, effectively checking all possible cases.

for i, (x, y) in enumerate(zip(arr1, arr2)):

Return the maximum absolute value expression found

// expression found. This value will be returned at the end.

// Using Integer.MIN_VALUE to handle negative overflow edge case.

// Length of the given arrays, assuming both arrays have the same length.

// Variables representing the direction for arr1[i] and arr2[i]

coefficientB = directions[k + 1]; // Coefficient for arr2

maxAnswer = max(maxAnswer, maxExprValue - minExprValue);

return maxAnswer; // Return the found maximum absolute value of an expression

for (int i = 0; i < arraySize; ++i) {</pre>

1 // Calculates the maximum absolute value expression for two arrays

// Define multipliers for different expression scenarios

// Initialize the maximum answer as the smallest integer possible

// Iterate through all possible expressions based on multipliers

2 function maxAbsValExpr(arr1: number[], arr2: number[]): number {

for (let expIndex = 0; expIndex < 4; ++expIndex) {</pre>

const coeffB = multipliers[expIndex + 1];

const coeffA = multipliers[expIndex];

// Select multipliers for current expression

const multipliers = [1, -1, -1, 1, 1];

let maxAnswer = Number.MIN_SAFE_INTEGER;

int maxExprValue = -INF, // Initialize max expression value in the current loop

// Loop through each element in the array to find max and min of expression

// Find the maximum value expression with the current coefficients

// Find the minimum value expression with the current coefficients

minExprValue = INF; // Initialize min expression value in the current loop

// Update the overall maximum answer with the difference between max and min

maxExprValue = max(maxExprValue, coefficientA * arr1[i] + coefficientB * arr2[i] + i);

minExprValue = min(minExprValue, coefficientA * arr1[i] + coefficientB * arr2[i] + i);

public int maxAbsValExpr(int[] arr1, int[] arr2) {

int maxDifference = Integer.MIN_VALUE;

int n = arr1.length;

for (int k = 0; k < 4; ++k) {

current_value = dir1 * x + dir2 * y + i

for dir1, dir2 in product((1, -1), repeat=2):

max_val, min_val = -inf, inf

return max_absolute_value_expr

for the current direction combination

23 # Update the maximum and minimum observed values max_val = max(max_val, current_value) 24 25 min_val = min(min_val, current_value) 26 27 # Update the maximum absolute value expression with the 28 # maximum difference between observed values

// Define the multipliers to represent the four possible combinations // of adding or subtracting arr1[i] and arr2[i]. // These multipliers will be applied inside the loop below. int[] multipliers = $\{1, -1, -1, 1\}$;

class Solution {

Java Solution

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// in the expression based on the current multiplier.
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               int a = multipliers[k], b = multipliers[k + 1];
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               // Initialize 'maxValue' and 'minValue' to track the maximum and minimum values
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               // of the expression for a given set of multipliers.
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               // Using Integer.MIN_VALUE and Integer.MAX_VALUE to start with the extreme possible values.
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               int maxValue = Integer.MIN_VALUE, minValue = Integer.MAX_VALUE;
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               // Iterate through the elements of the arrays to calculate the expressions.
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               for (int i = 0; i < n; ++i) {
29
                   // Calculate the current value of the expression.
                   int currentValue = a * arr1[i] + b * arr2[i] + i;
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                   // Update 'maxValue' and 'minValue' if the current value is greater than 'maxValue'
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                   // or less than 'minValue' respectively.
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                   maxValue = Math.max(maxValue, currentValue);
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                   minValue = Math.min(minValue, currentValue);
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                   // Update 'maxDifference' with the maximum difference found so far.
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                   maxDifference = Math.max(maxDifference, maxValue - minValue);
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           // Return the maximum absolute value expression found.
           return maxDifference;
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45 }
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C++ Solution
  1 #include <vector>
  2 #include <algorithm>
  3 using namespace std;
  5 class Solution {
    public:
         int maxAbsValExpr(vector<int>& arr1, vector<int>& arr2) {
             // Directions are the coefficients for arrl and arr2 in the combinations
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             // Last element is duplicated for easy loop termination
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             int directions [5] = \{1, -1, -1, 1, 1\};
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             const int INF = 1 << 30; // Define infinity as a large number</pre>
             int maxAnswer = -INF; // Initialize the maximum answer to negative infinity
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             int arraySize = arr1.size(); // Get the size of the input arrays
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             // Iterate over the 4 combinations of signs in the expression
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             for (int k = 0; k < 4; ++k) {
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                 int coefficientA = directions[k], // Coefficient for arr1
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Typescript Solution

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           // Initialize max and min variables for current expression scenario
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           let maxCurrent = Number.MIN_SAFE_INTEGER;
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            let minCurrent = Number.MAX_SAFE_INTEGER;
18
           // Iterate through elements of the arrays to compute expressions
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           for (let i = 0; i < arr1.length; ++i) {</pre>
21
               const val1 = arr1[i];
22
               const val2 = arr2[i];
23
24
               // Calculate the expression value with current i
25
               const expression = coeffA * val1 + coeffB * val2 + i;
26
27
               // Update the current maximum and minimum
28
               maxCurrent = Math.max(maxCurrent, expression);
29
               minCurrent = Math.min(minCurrent, expression);
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31
               // Update the global maximum answer
32
               maxAnswer = Math.max(maxAnswer, maxCurrent - minCurrent);
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36
       // Return the calculated maximum absolute value expression
37
       return maxAnswer;
38 }
Time and Space Complexity
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Space Complexity:

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The given Python code calculates the maximum absolute value expression for two arrays arr1 and arr2.

- Time Complexity: ○ There are 4 pairs of (a, b) which correspond to each combination of {1, -1} for each array element.
 - runs for n iterations where n is the length of the input arrays. Inside the loop, we perform constant-time operations such as comparisons and arithmetic operations. Therefore, the time complexity is O(4n) which simplifies to O(n).

For each of these 4 pairs, we iterate through both arr1 and arr2 simultaneously, using enumerate(zip(arr1, arr2)), which

- The additional space used by the algorithm is constant. It only needs a few variables for tracking maximum values, minimum values, and indices (mx, mi, and i), regardless of the input size.
- As such, the space complexity is 0(1).