Math

Medium Stack Array

Problem Description

You are asked to evaluate an arithmetic expression provided as an array of strings, tokens, which uses Reverse Polish Notation (RPN). This notation places the operator after the operands. For example, the expression "3 4 +" in RPN is equivalent to "3 + 4" in standard notation. Your task is to calculate the result of the expression and return the resulting integer value.

The expression only contains the operators +, -, *, and /.

Several points to consider for this problem are:

- Operands could be either integers or sub-expressions.
- When performing division, the result is always truncated towards zero.

The expression does not contain any division by zero.

The expression given is always valid and can be evaluated without error.

track of the operands as they appear and then apply the operators in the correct order.

- The result of the evaluated expression and any intermediate operations will fit within a 32-bit integer.
- Intuition
- To solve this problem, we need to understand the nature of Reverse Polish Notation. In RPN, every time we encounter an operator, it

The intuition for the solution is as follows: 1. We iterate through each string (token) in the tokens array.

applies to the last two operands that were seen. A stack is the perfect data structure for this evaluation because it allows us to keep

3. If the token is an operator, we pop the last two numbers from the stack and apply this operator; these two numbers are the

operands for any operator we come across.

- operands for the operator.
- 4. The result of the operation is then pushed back onto the stack.

5. After applying an operator, the stack should be in a state that is ready to process the next token.

2. If the token is a number (single digit or multiple digits), we convert it to an integer and push it onto a stack.

- 6. When we've processed all the tokens, the stack will contain only one element, which is the final result of the expression. Division in Python by default results in a floating-point number. Since the problem states that the division result should truncate
- toward zero, we explicitly convert the result to an int, which discards the decimal part.
- Solution Approach

The key here is to iterate through the tokens once and perform operations in order, ensuring the stack's top two elements are the

The solution makes use of a very simple yet powerful algorithm that utilizes a stack data structure to process the given tokens one

by one. Here are the steps it follows: Initialize an empty list nums that will act as a stack to store the operands.

2. Iterate over each token in the tokens array.

quotients.

Addition: nums [-2] += nums [-1]

Subtraction: nums [-2] -= nums [-1]

the expression.

 If the token is a numeric value (identified by either being a digit or having more than one character, which accounts for numbers like "-2"), we convert the token to an integer and push it onto the stack. If the token is an operator (+, -, *, /), we perform the following:

Pop the top two elements from the stack. Since the last element added is at the top of the stack, we'll refer to these as

towards zero. This is important as it handles the truncation towards zero for negative numbers correctly. The simple floor division operator // in Python truncates towards negative infinity, which can give incorrect results for negative

the second operand (at nums [-1]) and the first operand (at nums [-2]) in that order.

■ The result of the operation is then placed back into the stack at the position of the first operand (nums [-2]). ■ The second operand (nums [-1]), which has already been used, is removed from the stack. 3. After processing all the tokens, there should be a single element left in the nums stack. This element is the result of evaluating

Apply the operator to these two operands. For addition, subtraction, and multiplication, this is straightforward.

For division, we apply integer division which is the same as dividing and then applying the int function to truncate

the stack ensures that the operands for any operator are always readily available at the top of the stack. Here is a snippet of how the arithmetic operations are processed:

The algorithm used here is particularly efficient because it has a linear time complexity, processing each token exactly once. The

space complexity of this approach is also linear, as it depends on the number of tokens that are pushed into the stack. The use of

 Multiplication: nums [-2] *= nums [-1] Division: nums[-2] = int(nums[-2] / nums[-1]) Once finished, the program returns nums [0] as the result of the expression.

Let's use the following RPN expression as our example: "2 1 + 3 *" which, in standard notation, translates to (2 + 1) * 3.

2. We iterate through the tokens: ["2", "1", "+", "3", "*"].

a. The first token is "2", which is a number. We push it onto the stack: nums = [2].

b. The second token is "1", which is also a number. We push it onto the stack: nums = [2, 1].

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1. We initialize an empty list nums to serve as our stack: nums = [].
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Python Solution

class Solution:

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C++ Solution

#include <vector>

#include <stack>

class Solution {

public:

#include <string>

from typing import List

By following the solution approach:

Example Walkthrough

c. The third token is "+", which is an operator. We need to pop the top two numbers and apply the operator: - We pop the first

d. The fourth token is "3", a number. We push it onto the stack: nums = [3, 3].

We push the result onto the stack: nums = [3].

def evalRPN(self, tokens: List[str]) -> int:

if token == "+":

elif token == "-":

else: # Division

number_stack.pop()

return number_stack[0]

return stack.pop();

int evalRPN(vector<string>& tokens) {

if len(token) > 1 or token.isdigit():

number_stack.append(int(token))

Perform the operation based on the operator

number_stack[-2] += number_stack[-1]

number_stack[-2] *= number_stack[-1]

Return the result which is the only number left in the stack

// The final result is the only element in the stack, return it

// Create a stack to keep track of integers for evaluation

Stack for storing numbers

number_stack = []

elser

e. The fifth token is "*", an operator: - We pop the second operand, secondOperand = nums.pop(), which is 3, leaving nums = [3]. - We pop the first operand, firstOperand = nums.pop(), which is 3, leaving nums = []. - We multiply the two operands: stackResult = firstOperand * secondOperand, which equals 3 * 3 = 9. - We push the result onto the stack: nums = [9]. 3. After processing all the tokens, we are left with a single element in our stack nums = [9], which is our result.

So, the given RPN expression "2 1 + 3 *" evaluates to 9. Thus, the function would return 9 as the result of the expression.

operand: secondOperand = nums.pop(), which is 1. Now nums = [2]. - We pop the second operand: firstOperand = nums.pop(),

which is 2. Now nums = []. - We add the two operands: stackResult = firstOperand + secondOperand, which equals 2 + 1 = 3.

Loop over each token in the input list for token in tokens: # If the token represents a number (accounting for negative numbers)

Convert the token to an integer and push it onto the stack

Ensure integer division for negative numbers too

Pop the last two numbers, add them, and push the result back

number_stack[-2] = int(float(number_stack[-2]) / number_stack[-1])

Pop the last number (second operand) from the stack as it's been used

Pop the last two numbers, subtract the second from the first, and push back number stack[-2] -= number stack[-1] 21 22 elif token == "*": 23 # Pop the last two numbers, multiply, and push the result back

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Java Solution
   class Solution {
       public int evalRPN(String[] tokens) {
           Deque<Integer> stack = new ArrayDeque<>(); // Create a stack to hold integer values
           // Iterate over each token in the input array
           for (String token : tokens) {
6
               // Check if the token is a number (either single digit or multi-digit)
               if (token.length() > 1 || Character.isDigit(token.charAt(0))) {
                   // Push the number onto the stack
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                   stack.push(Integer.parseInt(token));
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               } else {
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                   // Pop the top two elements for the operator
                   int secondOperand = stack.pop();
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                   int firstOperand = stack.pop();
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                   // Apply the operator on the two operands based on the token
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                   switch (token) {
                       case "+":
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                           stack.push(firstOperand + secondOperand); // Add and push the result
20
                       case "-":
22
                           stack.push(firstOperand - secondOperand); // Subtract and push the result
23
                       case "*":
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25
                           stack.push(firstOperand * secondOperand); // Multiply and push the result
26
                           break;
                       case "/":
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                           stack.push(firstOperand / secondOperand); // Divide and push the result
29
                           break;
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stack<int> numbers;
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             // Iterate over each token in the Reverse Polish Notation expression
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             for (const string& token : tokens) {
                 // If the token represents a number (can be multiple digits or negative)
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                 if (token.size() > 1 || isdigit(token[0])) {
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                     // Convert the string token to an integer and push onto the stack
                     numbers.push(stoi(token));
                 } else { // If the token is an operator
                     // Pop the second operand from the stack
 18
 19
                     int operand2 = numbers.top();
 20
                     numbers.pop();
 21
 22
                     // Pop the first operand from the stack
                     int operand1 = numbers.top();
 24
                     numbers.pop();
 25
 26
                     // Perform the operation based on the type of operator
 27
                     switch (token[0]) {
 28
                         case '+': // Addition
 29
                             numbers.push(operand1 + operand2);
 30
                             break;
 31
                         case '-': // Subtraction
                             numbers.push(operand1 - operand2);
                             break;
                         case '*': // Multiplication
                             numbers.push(operand1 * operand2);
 36
                             break;
                         case '/': // Division
 37
 38
                             numbers.push(operand1 / operand2);
 39
                             break;
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             // The final result is the only number remaining on the stack
 45
             return numbers.top();
 46
 47
    };
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Typescript Solution
  1 // Evaluates the value of an arithmetic expression in Reverse Polish Notation.
  2 // Input is an array of tokens, where each token is either an operator or an operand.
  3 // Valid operators are: '+', '-', '*', and '/'.
  4 // Assumes the RPN expression is valid.
    function evalRPN(tokens: string[]): number {
         // Stack to hold the operands for evaluation.
```

39 break; case '/': 40 // Use truncation to conform to the requirements of integer division in RPN. 41 42 // The '~~' is a double bitwise NOT operator, used here as a substitute for Math.trunc 43 // which is to conform to the behavior specified in the problem statement. stack.push(~~(firstOperand / secondOperand));

const stack: number[] = [];

for (const token of tokens) {

} else {

if (isNumeric(token)) {

switch (token) {

case '+':

case '-':

case '*':

break;

break;

break;

default:

if (stack.length !== 1) {

return stack[0];

// Helper function to check if a string is a numeric value.

// If the token is a number, push it onto the stack.

return !isNaN(parseFloat(token)) && isFinite(Number(token));

stack.push(firstOperand + secondOperand);

stack.push(firstOperand - secondOperand);

stack.push(firstOperand * secondOperand);

// The result of the expression is the last element of the stack.

(append and pop), can be considered to have constant time complexity, 0(1).

// Safety check: there should only be one element left in the stack.

// Since we know the token is an operator, pop two operands from the stack.

// Safety check: Ensure operands are valid numbers to avoid runtime errors.

if (typeof firstOperand === 'undefined' || typeof secondOperand === 'undefined') {

throw new Error("Invalid token: Encountered an unknown operator.");

throw new Error("Invalid Expression: The final stack should only contain one element.");

throw new Error("Invalid Expression: Insufficient operands for the operator.");

// Perform the operation according to the current token and push the result onto the stack.

function isNumeric(token: string): boolean {

// Process each token in the RPN expression.

stack.push(Number(token));

const secondOperand = stack.pop();

const firstOperand = stack.pop();

Time and Space Complexity

Space Complexity

The given Python function evalRPN evaluates Reverse Polish Notation (RPN) expressions. The time complexity and space complexity analysis are as follows: **Time Complexity** The time complexity of this function is O(n), where n is the number of tokens in the input list tokens. This is because the function

iterates through each token exactly once. Each operation within the loop, including arithmetic operations and stack operations

The space complexity of the code is O(n) in the worst case, where n is the number of tokens in the list. This worst-case scenario occurs when all tokens are numbers and are thus pushed onto the stack. In the best-case scenario, where the input is balanced with numbers and operators, the space complexity could be better than O(n), but the upper bound remains O(n). The auxiliary space required is for the nums stack used to perform the calculations. There are no other data structures that use significant memory.