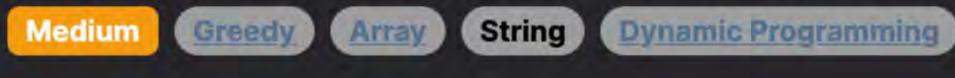
## 2900. Longest Unequal Adjacent Groups Subsequence I



Leetcode Link

**Problem Description** 

The goal is to find the longest subsequence from an array of indices [0, 1, ..., n - 1] such that for any two consecutive indices in the subsequence i\_j and i\_{j + 1}, the elements in the binary array groups at those indices are not the same, i.e., groups[i\_j] != groups [i\_{j + 1}]. Each index in the subsequence corresponds to a word in the words array. The task is to return an array of words that represents this longest subsequence.

A subsequence is a sequence that can be derived from another sequence by deleting some or no elements without changing the order of the remaining elements. Importantly, the words in the words array may have different lengths, which doesn't impact the selection of the subsequence.

#### Intuition

we can make local, optimal choices at each step without needing to consider the rest of the array.

The approach to finding the longest subsequence where consecutive elements in groups are different is a greedy one. This means

For every element at index i in groups, we have two scenarios - either i is the first index (i.e., i == 0), or groups [i] is different from the previous element groups [i - 1]. If either of these conditions is met, we can include the corresponding word words [i] in our subsequence.

words have their corresponding groups elements equal, effectively giving us the longest subsequence by the definition provided.

By iterating over the entire groups array and including words that meet the criteria, we ensure that no two consecutive selected

long as they meet the aforementioned condition. Since the condition only depends on the current and previous groups elements, we only need a simple iteration to build our solution without needing to backtrack or look ahead.

The intuition comes from the fact that to maximize the length of the subsequence, we want to include as many words as possible as

#### The provided Python solution uses a list comprehension to create and return the subsequence of words, which is essentially a

**Solution Approach** 

single-pass greedy algorithm. The algorithm iterates through the groups array and applies a selection criteria to each element to determine if the corresponding element from the words array should be included in our final output or not.

#### Algorithm:

built-in enumerate function to obtain each element in groups and its index simultaneously. 2. In this comprehension, for every element x and index i, the following condition is checked: i == 0 or x != groups [i - 1]. This

1. Initialize a list comprehension that will evaluate each element x in groups along with its corresponding index i. It uses Python's

- condition says that the first element (i == 0) should always be included, and then every subsequent element should be included only if it is different than the one preceding it  $(x \mid = groups[i - 1])$ . This ensures that no two adjacent elements in the subsequence have the same groups value. 3. If the condition is true for a given i, we select words[i] for inclusion in the final output.
- the longest subsequence satisfying the problem's constraints.
- 5. The final step is to return this list of words.

4. Once the list comprehension finishes iterating over all elements in groups, it will have produced a list of words that constitutes

**Data Structures:** 

#### We utilize Python's list data structure to store the words and groups.

Patterns used:

### The solution applies a greedy approach to the problem: at each step, it makes a local optimum choice (whether to include a

optimum (the longest subsequence under the given conditions). List comprehension, a concise way to create lists, is used for its readability and efficiency in selecting the appropriate words.

word) based only on the current and immediate previous elements from groups, which guarantees the finding of the global

This simple yet effective approach leverages the characteristics of the problem's constraints to arrive at an optimal solution without needing a complex algorithm.

Example Walkthrough

#### words arrays:

• groups = [1, 0, 0, 1, 0, 1]words = ["apple", "banana", "grape", "cherry", "mango", "peach"]

Our task is to find the longest subsequence such that no two consecutive indices in the subsequence correspond to equal elements

Let's walk through a small example to illustrate the solution approach described above. Suppose we have the following groups and

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in the groups array. Let's apply the algorithm step by step:
 1. Start iterating through the groups array, comparing the element at index i with the one at index i - 1.
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"apple" from the words array in our subsequence.

["apple", "banana", "cherry", "mango", "peach"]

longest\_subsequence\_words = []

if i == 0 or group\_number != groups[i - 1]:

longest subsequence words.append(words[i])

// Initialize an ArrayList to store the resulting words.

// Iterate over the words to find the longest subsequence.

List<String> result = new ArrayList<>();

// Iterate through each group by index.

for (int index = 0; index < n; ++index) {</pre>

// then it is a part of the longest non-repeating subsequence.

// Return the answer vector containing the words in the longest non-repeating subsequence.

if (index == 0 || groups[index] != groups[index - 1]) {

// Add the current word to the answer vector.

answer.emplace\_back(words[index]);

# Return the list of words that form the longest subsequence

consecutive elements must be different.

2. The first element in groups is 1 (at index 0). Since i == 0, we don't have a previous element to compare with, so we include

- 3. The next element in groups is 0 (at index 1). Since groups [1] != groups [0], we include "banana" from the words array in our
- subsequence. 4. At index 2, groups has another 0. This time, groups [2] == groups [1], so "grape" does not get included in our subsequence since
- 5. Now at index 3, we have a 1 in groups. Since groups [3] != groups [2], "cherry" gets included in our subsequence.

6. Moving to index 4, the element in groups is 0. Because groups [4] != groups [3], we include "mango" in our subsequence.

- Following the steps of our algorithm, the final subsequence of words is:
- Thus, by iterating through each element of the groups array and checking our defined condition, we successfully construct the longest subsequence of words without having identical consecutive elements from groups. This illustrates the effectiveness of the

greedy approach to solve the problem, as we made local optimal selections to achieve a global optimal solution.

# If yes, append the corresponding word to the longest\_subsequence\_words list

# Initialize an empty list to store the words in the longest subsequence

public List<String> getWordsInLongestSubsequence(int n, String[] words, int[] groups) {

7. Lastly, at index 5, groups contains a 1. As groups [5] != groups [4], we include "peach" in our subsequence.

# Import the List type from typing module for type hints from typing import List class Solution: def getWordsInLongestSubsequence(self, n: int, words: List[str], groups: List[int]) -> List[str]:

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# Iterate through each index and corresponding group number in the groups list
           for i, group_number in enumerate(groups):
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               # Check if it's the first word or if the current group number is different from the previous one
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Python Solution

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return longest_subsequence_words
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Java Solution
 1 import java.util.ArrayList;
   import java.util.List;
   class Solution {
        * Finds the words in the longest subsequence with alternating groups.
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                        the number of words.
        * @param n
        * @param words an array of words.
10
        * @param groups an array of group identifiers corresponding to each word.
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        * @return a list of words in the longest subsequence with alternating groups.
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```

#### for (int i = 0; i < n; ++i) { 19 20 21 22

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// Add the first word and any word that starts a new group (compared to the previous word).
               if (i == 0 || groups[i] != groups[i - 1]) {
                   result.add(words[i]);
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           // Return the list of words in the longest subsequence.
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           return result;
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28 }
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C++ Solution
1 #include <vector>
   #include <string>
   class Solution {
   public:
       // Function that generates a vector of strings, which consists of the words
       // in the longest non-repeating subsequence based on the given groups.
       // Parameters:
       // n - the number of elements in the words and groups vectors.
       // words - a vector of strings representing the words.
       // groups - a vector of integers, where each integer corresponds to the group of the word at the same index.
       std::vector<std::string> getWordsInLongestSubsequence(int n, std::vector<std::string>& words, std::vector<int>& groups) {
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           // Answer vector to store the resulting words.
           std::vector<std::string> answer;
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#### 28 }; 29

return answer;

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Typescript Solution
   function getWordsInLongestSubsequence(totalWords: number, wordsArray: string[], wordGroups: number[]): string[] {
       // Initialize an array to hold the resulting longest subsequence of words.
       const longestSubsequence: string[] = [];
       // Iterate through the array of words to identify the longest subsequence.
       for (let index = 0; index < totalWords; ++index) {</pre>
           // If we are at the first word or the current word's group is different from the previous word's group,
           // it is a part of the longest subsequence, so we add it to the result array.
           if (index === 0 || wordGroups[index] !== wordGroups[index - 1]) {
               longestSubsequence.push(wordsArray[index]);
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       // Return the longest subsequence found.
       return longestSubsequence;
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// If we are at the first word, or the current word belongs to a different group than the previous one,

# Time and Space Complexity

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once and performs a constant time check for each element. The space complexity is O(k), where k is the number of unique subsequences identified, which in the worst case could be equal to n. This would happen if no two consecutive elements in groups are the same, resulting in each word from words being added to the

output list. Thus, the space used for the output list is proportional to the number of selected words.

The time complexity of the code is O(n), where n is the length of the list groups. This is because the code iterates over the list groups