## 1624. Largest Substring Between Two Equal Characters



**Problem Description** 

**Leetcode Link** 

characters in the string, excluding the characters themselves. If no such substring exists (i.e., there aren't two identical characters in the string with anything between them), the function should return -1. The term "substring" indicates a consecutive sequence of characters within the string s.

The problem presents a string s and requires finding the length of the longest substring that is present between two identical

For example, in the string "abca", the longest substring between two equal characters 'a' is "bc", which has the length of 2. In a case where the string is "abc", there are no two equal characters with something between them, so the result would be -1.

# Intuition

them.

first and last occurrences, since that will determine the length of the substring between them. To efficiently do this, we should remember the index of the first occurrence of each character when we see it for the first time. With these considerations in mind, we use a dictionary d to store characters as keys and their first encountered index positions as

To solve this problem, we must find characters in the string that appear more than once and calculate the distance between their

values. We iterate over the string s and for each character c: If c is already in the dictionary d, this means we have previously seen it, and hence we calculate the distance between this

• If c is not in the dictionary, we record its index i in the dictionary d, marking its first occurrence.

occurrence and its first occurrence, which is i - d[c] - 1, and update the answer ans if this distance is larger.

The variable ans is used to keep track of the maximum length found so far. If no such length is found (meaning ans never gets updated), it remains -1, which is also the default value we return when there are no two identical characters with substrings between

Thus, by using a single pass over the string and a dictionary to track the indices of characters, we arrive at an efficient solution to the problem.

**Solution Approach** 

## choice for this problem due to its fast lookups and insertions, which operate on average in O(1) time complexity.

of c.

at hand.

Let's walk through the implementation step by step:

• A variable ans is initiated with a value of -1. This will eventually hold the maximum length of any substring found between two

The implementation uses a dictionary to store the first occurrence of each character in the given string. A dictionary is an excellent

equal characters, excluding the characters themselves.

• We iterate over the characters of the string s using a for loop, where i is the index and c is the character at that index.

A dictionary d is initialized to keep a record of the first index where each character appears.

- Inside the loop, we check if the character c is already in the dictionary d:  $\circ$  If c is in d, it means we have encountered c before. The substring length between the two c characters is i - d[c] - 1.
- We compare this length with the current maximum stored in ans. If it's larger, we update ans with this new value.
- If c is not in d, we add c to the dictionary with the current index i as its value. This marks the location of the first occurrence
- This approach is efficient because we maintain a sliding window between two occurrences of the same character by remembering only the first occurrence and using the current index to calculate the length of the in-between substring. As a result, the time
- complexity is O(n), where n is the length of the string, because we are making a single pass through the string. With this understanding, the provided Python code realizes the approach effectively and delivers the correct result for the problem

• Once the loop is complete, ans will have the maximum length found or remain -1 if no such substring exists.

Example Walkthrough Let's consider a string s with the value "character". We are trying to find the length of the longest substring present between two

Initialize ans as −1.

Following the solution approach, let's break it down step by step:

identical characters (excluding those characters).

1. For index 0, character c: it's not in d, so we add c: 0 to the dictionary. 2. For index 1, character h: it's also not in d, add h: 1.

Initialize a dictionary d for keeping the first index occurrence of each character.

- 3. For index 2, character a: it's not in d, add a: 2.
- 4. For index 3, character r: not in d, add r: 3. 5. For index 4, character a: a is already in d at index 2. Calculate the distance: 4 - 2 - 1 = 1. Update ans to 1.

Start iterating through the string:

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6. For index 5, character c: c is in d with index 0. The distance is 5 - 0 - 1 = 4. Since 4 is greater than the current ans value of
  1, update ans to 4.
```

# Initialize the answer with -1 as per problem constraints

# Store the first occurrence of the character

first\_occurrence[char] = index

// If the character has not been seen before,

// If we have seen the character before,

// calculate the length between the current and first index.

// Then, update the maxLength if the calculated length is larger.

maxLength = Math.max(maxLength, i - firstOccurrence[charIndex] - 1);

// store its index in firstOccurrence.

if (firstOccurrence[charIndex] == -1) {

firstOccurrence[charIndex] = i;

// Return the maximum length found.

# Iterate over the string to find the max length between equal characters

- 7. For index 6, character t: not in d, add t: 6. 8. For index 7, character e: not in d, add e: 7.
- End of iteration. The final ans is 4, which is the length of the longest substring ("aract") between two identical characters without including them ("c...c" and "r...r").

9. For index 8, character r: r is in d with index 3. The distance is 8 - 3 - 1 = 4. ans is already 4, so no update is needed.

The code returns 4 as the length of the longest substring between two repeating characters, excluding the characters themselves.

def maxLengthBetweenEqualCharacters(self, s: str) -> int: # Dictionary to store the first occurrence of each character first\_occurrence = {}

# If character is already seen, calculate the length between the current and first occurrence

### 12 if char in first\_occurrence: 13 length\_between = index - first\_occurrence[char] - 1 max\_length = max(max\_length, length\_between) 14 15 else:

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Python Solution

 $max_length = -1$ 

for index, char in enumerate(s):

class Solution:

```
19
           # Return the maximum length found
20
           return max_length
21
Java Solution
   class Solution {
       public int maxLengthBetweenEqualCharacters(String s) {
           // Array to store the first occurrence index of each character.
           // Initialized to -1 indicating that the character hasn't been seen yet.
           int[] firstOccurrence = new int[26];
           Arrays.fill(first0ccurrence, -1);
           // Variable to store the maximum length found.
           int maxLength = -1;
9
10
           // Loop through the string to find the maximum length.
           for (int i = 0; i < s.length(); ++i) {</pre>
12
               // Calculate the index for the character 'a' as 0, 'b' as 1, etc.
13
               int charIndex = s.charAt(i) - 'a';
14
```

### 31 } 32

} else {

return maxLength;

```
C++ Solution
 1 class Solution {
   public:
       int maxLengthBetweenEqualCharacters(string s) {
           // Create a vector 'firstIndex' to store the first occurrence index of each lowercase letter.
           // Initialized with -1, indicating we have not seen the character yet.
           vector<int> firstIndex(26, -1);
           // Variable to store the maximum length found between two equal characters.
           int maxLen = -1;
10
           // Iterate through the string to check each character.
11
12
           for (int i = 0; i < s.size(); ++i) {
               // Converting the current character to its corresponding index (0-25 for 'a'-'z').
13
               int index = s[i] - 'a';
14
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16
               // If we have not seen this character before, store its index.
17
               if (firstIndex[index] == -1) {
                   firstIndex[index] = i;
18
               } else {
19
                   // If we've seen this character before, calculate the length between the two equal characters.
20
                   // Update 'maxLen' if we find a longer length.
22
                   maxLen = max(maxLen, i - firstIndex[index] - 1);
23
24
25
26
           // Return the maximum length found.
27
           return maxLen;
28
29 };
30
```

Typescript Solution

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1 /**
    * Finds the maximum length of a substring between two identical characters
    * in the input string 's'.
    * @param s The string to be evaluated.
    * @return The maximum length of a substring.
    function maxLengthBetweenEqualCharacters(s: string): number {
       // 'n' holds the length of the string 's'.
       const n: number = s.length;
10
       // 'positionIndex' is an array to track the first occurrence position of each alphabet character.
       const positionIndex: number[] = new Array(26).fill(-1);
       // 'maxLength' will keep track of the maximum length found.
       let maxLength: number = -1;
14
15
       // Iterate over each character in the string.
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17
       for (let i = 0; i < n; i++) {
18
           // Compute the zero-based alphabet index of the current character.
           const charIndex: number = s[i].charCodeAt(0) - 'a'.charCodeAt(0);
20
           // If it is the first occurrence of the character, store its position.
           if (positionIndex[charIndex] === -1) {
21
               positionIndex[charIndex] = i;
           } else {
               // If a pair is found, update the 'maxLength', if necessary
               maxLength = Math.max(maxLength, i - positionIndex[charIndex] - 1);
26
27
28
       // Return the maximum length found, or -1 if no such length exists.
29
       return maxLength;
30
31 }
32
Time and Space Complexity
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### The given Python code iterates over the string s only once with a for loop. For each character in the string, it performs a constanttime operation to check if the character is in the dictionary d, update the maximum length using max(), and set the value in the

where n is the length of the string s.

complexity of maintaining the dictionary is:

where n is the length of the string s.

Time Complexity

dictionary. None of these operations depend on the size of the input string inside the loop, so each iteration of the loop runs in constant time, 0(1). The time complexity of the entire function is thus proportional to the number of characters n in the string s, as each character is processed just once. This results in an overall time complexity of: 1 0(n)

**Space Complexity** 

1 0(n)

However, if we consider the constraint of the problem that the input string s consists of only lowercase English letters, then there is a

constant maximum of 26 unique characters that can be stored in the dictionary d. This would imply a constant space complexity: 1 0(1)

Depending on the input constraints specified in the problem statement, you should use the appropriate analysis for the space complexity.

The space complexity of the code depends on the number of unique characters in the string s, as a dictionary d stores the first index at which each character appears. In the worst case, all characters of the string s are unique, which would require storing each character in the dictionary d. Assuming the input string s contains n characters and taking into account the possibility of n unique characters, the space