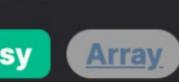
# 2562. Find the Array Concatenation Value









**Leetcode Link** 

# **Problem Description**

what you need to know:

This LeetCode problem involves manipulating an array of integers to calculate what is referred to as a "concatenation value". Here's

- You have an array of integers called nums, with each element having a 0-based index.
- The term "concatenation" of two numbers here refers to creating a new number by joining the digit sequences of both numbers end to end.
- For example, concatenating 15 and 49 yields 1549.
- Initially, the concatenation value is 0. To find the final concatenation value, you go through the following steps:
- 1. If nums contains more than one number, take the first and last elements.
  - 2. Concatenate those two elements and add their concatenation's numerical value to the concatenation value of nums.
  - 4. If only one number exists in nums, add its value to the concatenation value and remove it.
- This process repeats until nums is empty.

3. Remove the first and last elements from nums.

- Your goal is to return the final concatenation value after all elements are combined and removed according to the steps above.
- Intuition

no more elements left to process.

## To approach the solution to this problem, you start by understanding that you'll be reducing the array from both ends until there are

Here's how you arrive at the solution step-by-step:

 Recognize that array elements should be considered from both ends. At each step where there are at least two elements, you consider the first and last element for the operation.

- You convert each number to a string, concatenate those strings, then convert the resulting string back to an integer.
- Add this integer to the running total of the concatenation value.
- This process is repeated, updating two pointers (i and j) that represent the current first and last elements in the array. When the pointers meet or cross, you've processed all elements.
- The final concatenation value can be returned after considering all elements in the required manner.

If there's only one element left in the process, simply add its value to the total.

- **Solution Approach**
- The implementation of the solution uses a straightforward simulation algorithm with two pointers. Let's go through the specifics:
- We start by initializing an accumulator for the answer, ans, to start summing up the concatenation values.

• Two pointers, i and j, are used to traverse the nums array from both ends toward the center.

- i is initialized to 0 as the start of the array. j is initialized to the last index of the array, which is len(nums) − 1.
- We enter a while loop that continues as long as i < j indicating there are at least two elements in the current subrange.</li>
  - Inside the loop, we concatenate the numerical strings of nums[i] and nums[j], by doing str(nums[i]) + str(nums[j]), and

while loop because it has no pair element.

- then convert this string back to an integer using int(). The resultant integer is added to the ans accumulator.
- towards the center of the array. • If we exit the loop and find that i == j, it means there is one remaining element in the array which wasn't processed by the

• After concatenating and adding to ans, we advance i forward by one (i += 1) and move j backward by one (j -= 1) to move

simulation.

In that case, we simply add the value of this last remaining element (nums[i]) to ans.

operations. The core pattern here is the two-pointer technique that allows us to efficiently process elements from both ends of the array.

• The solution approach is completed by returning the ans variable which contains the final concatenation value after the

This problem does not require any complex data structures or algorithms, as it simply utilizes basic array manipulation and integer

Let's consider a small example to illustrate the solution approach with the nums array nums = [5,6,2,8]. 1. Initialize ans = 0 to keep track of the concatenation value.

## Now we start the while loop:

4. Second iteration:

Example Walkthrough

3. First iteration:

- Since i < j, concatenate the first and last elements: nums[i] is 5 and nums[j] is 8.</li> • Their concatenation as strings is '5' + '8' which is '58'. Convert this back to an integer to get 58.
  - Add this to ans: ans = ans + 58, which makes ans = 58.  $\circ$  Now, increment i so i = 1 and decrement j so j = 2.

2. Initialize two pointers: i = 0 for the start of the array and j = len(nums) - 1 = 3 for the end of the array.

- The condition i < j still holds true.</li> Concatenate nums[i] which is 6 and nums[j] which is 2 to get '62'.
- $\circ$  Move i to i = 2 and j to j = 1.

unpaired element, we would add its value directly to ans.

def findTheArrayConcVal(self, nums: List[int]) -> int:

# Loop until the pointers meet or cross

# Initialize the answer to 0

while left < right:</pre>

if left == right:

answer += nums[left]

6. Since all elements have been processed, the final concatenation value is the current value of ans, which is 120. We return ans.

 $\circ$  As an integer, it is 62. Update ans = 58 + 62, making ans = 120.

Python Solution from typing import List

Therefore, the final concatenation value for the input array [5,6,2,8] using the solution approach provided would be 120.

5. Now, i >= j, and the while loop condition is not met. However, we don't have an element that's unpaired. If there were an

answer = 0 # Set pointers for the start and end of the array left, right = 0, len(nums) - 19

# Concatenate the numbers at the pointers, convert to int and add to the answer

# If there is a middle element (odd number of elements), add it to the answer

// convert the result to integer, and add it to the result.

// Move the left index forward and the right index backward.

// Add the middle element directly to the result.

2 // Pairs the first and last elements moving inwards and adds their concatenation.

// If there is an unpaired middle element, it is added directly to the result.

Thus, the time complexity is  $0(n/2 \times \log M)$ , which simplifies to  $0(n \times \log M)$ .

function findTheArrayConcVal(nums: number[]): number {

// Get the length of the array

// Initialize the front index

// Initialize the back index

const arrayLength = nums.length;

// Initialize the answer to zero

let backIndex = arrayLength - 1;

result += Long.parseLong(nums[leftIndex] + "" + nums[rightIndex]);

// Check if there's a middle element left (in case of odd number of elements).

### answer += int(str(nums[left]) + str(nums[right])) 14 15 # Move the pointers towards the center 16 left, right = left + 1, right - 1

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class Solution:

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           # Return the final answer
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            return answer
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Java Solution
   class Solution {
        * Calculates the "array conc val" by concatenating pairs of elements
        * from the beginning and end of the array moving towards the center.
        * If there's a middle element (odd number of elements), it adds it as is.
        * @param nums An array of integers.
        * @return The calculated "array conc val".
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       public long findTheArrayConcVal(int[] nums) {
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            long result = 0; // Initialize the result variable.
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           int leftIndex = 0; // Start at the beginning of the array.
            int rightIndex = nums.length - 1; // Start at the end of the array.
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           // Loop through the array from both ends until indices meet or cross.
           while (leftIndex < rightIndex) {</pre>
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               // Concatenate the elements at current indices as strings,
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### 33 34 return result; // Return the computed "array conc val". 35 36 } 37

leftIndex++;

rightIndex--;

if (leftIndex == rightIndex) {

result += nums[leftIndex];

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C++ Solution
 1 class Solution {
 2 public:
       // Function to find the array concatenated value
       long long findTheArrayConcVal(vector<int>& nums) {
            long long concatenatedSum = 0; // Initialize sum of concatenated values
           int left = 0; // Starting index from the beginning of the array
           int right = nums.size() - 1; // Starting index from the end of the array
           // Loop through the array from both ends until the pointers meet or cross
           while (left < right) {</pre>
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               // Concatenate the values at the current indices, convert to number, and add to sum
               concatenatedSum += stoll(to_string(nums[left]) + to_string(nums[right]));
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               // Move the left pointer forward and the right pointer backward
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               ++left;
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               --right;
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           // If there is a middle element (odd number of elements), add it to the sum
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           if (left == right) {
               concatenatedSum += nums[left];
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           // Return the final concatenated sum
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           return concatenatedSum;
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27 };
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Typescript Solution
 1 // This function finds a value based on the concatenation of array elements.
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### 13 14 // Loop until frontIndex is less than backIndex 15 while (frontIndex < backIndex) {</pre> // Concatenate the elements by converting them to strings, adding them, and then converting back to a number 16 answer += Number(`\${nums[frontIndex]}\${nums[backIndex]}`);

let answer = 0;

let frontIndex = 0;

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// Increment the front index
           frontIndex++;
          // Decrement the back index
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          backIndex--;
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       // If there is a middle element, add it to the answer
24
       if (frontIndex === backIndex) {
26
           answer += nums[frontIndex];
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29
       // Return the calculated answer
30
       return answer;
31 }
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Time and Space Complexity
code.
Time Complexity
```

The provided Python code calculates a special value based on the input nums list. Let's analyze the time and space complexity of this The time complexity of the algorithm is determined by the while loop, which runs as long as i < j. Since the indices i and j move towards each other with each iteration, the loop executes approximately n/2 times, where n is the total number of elements in nums. Inside this loop, the algorithm concatenates the string representations of numbers at indices i and j, which takes 0(\log M) time, where M is the maximum value in the array (since the number of digits of a number x is proportional to  $\log x$ ).

**Space Complexity** The space complexity of the algorithm is determined by the extra space needed to store the intermediate string representations created during the concatenation operation. The longest possible string is the concatenation of the two largest numbers in nums.

Thus, the space complexity is proportional to the length of this string, which is 0(2 × \log M). This simplifies to 0(\log M) since constant factors are dropped in Big O notation.