# 1481. Least Number of Unique Integers after K Removals

Counting ) Sorting Medium **Greedy** Array Hash Table

## **Problem Description**

Given an array arr of integers and an integer k, the task is to determine the minimum number of unique integers that remain in the array after precisely k elements have been removed. The key to solving this problem lies in effectively selecting which elements to remove to achieve the least number of unique integers possible.

Intuition

unique count as quickly. Ideally, we want to remove the numbers that occur the least amount of times last. We can apply a strategy that consists of the following steps: 1. Count the frequency of each unique integer in the array using a hash table (or Counter in Python). 2. Sort these unique integers by their frequency in ascending order. This way, the elements that appear less frequently are at the beginning of the

To solve this problem, we should consider removing elements that appear more frequently first since this will not reduce the

- sorted list.
- 3. Traverse the sorted list, and with each step, decrease k by the frequency of the current element. This simulates removing that element's occurrences from the array.
- 4. If k becomes negative, it means we can't remove all occurrences of the current element without exceeding the allowed k removals. Therefore, the current count of unique integers minus the number of elements we've been able to fully remove by this point gives us the answer.
- up with 0 unique integers.

5. If we go through the entire list without k becoming negative, it means we managed to remove all occurrences of certain elements, and we end

Having this strategy in place allows us to reach the solution in an efficient manner. **Solution Approach** 

The implementation of the solution follows a straightforward approach outlined in previous steps, which uses common data structures and algorithms:

### Hash Table (Counter): We employ a hash table which is native to Python called Counter from the collections module. This

structure automatically counts the frequency of each unique integer, which is essential to our strategy. It simplifies the process of determining how many times each integer appears in the array.

cnt = Counter(arr) **Sorting:** After counting the occurrences of each integer, we sort these counts. This sorting is ascending, meaning that unique integers with fewer occurrences are considered first when we start removing elements.

```
sorted(cnt.values())
```

k -= v

if k < 0:

simulating the removal of v occurrences from the array. for i, v in enumerate(sorted(cnt.values())):

Traversal and Subtraction: With the sorted frequencies, we traverse the list. For each frequency value v, we subtract k by v,

```
If k becomes less than zero during the iteration, it indicates that we cannot remove all occurrences of the current element
as it would exceed k. Hence, the minimal number of unique integers is obtained by subtracting the current index i from
```

being linear traversals and basic arithmetic operations.

Let's use a small example to illustrate the solution approach:

return len(cnt) - i

unique integers remaining is minimal.

the total number of unique integers (the length of cnt):

to remove all occurrences of certain elements leading to zero unique integers left. return 0

The overall complexity of the solution is determined mainly by the sorting operation and the counting operation, with the rest

Return Result: If we are able to traverse the entire sorted list of frequencies without k becoming negative, we have managed

**Example Walkthrough** 

Suppose arr = [4, 3, 1, 1, 2, 3, 3] and k = 3. Our goal is to remove k elements from arr in such a way that the number of

### 2 appears 1 time 3 appears 3 times

1 appears 2 times

4 appears 1 time Using a hash table:

cnt =  $\{1: 2, 2: 1, 3: 3, 4: 1\}$ **Sort by Frequency:** We sort these counts in ascending order based on frequency:

∘ We start with the first element (frequency of 1). We remove one instance of the integer with a count of 1, which decreases k to 2.

∘ With the third element (frequency of 2), we can remove both instances of the integer 1, which would decrease k to -1.

• We move to the next element (another frequency of 1). We remove one instance of another integer with a count of 1, which decreases k to

```
However, since we can't go negative, it means we can't remove all instances of 1. At this point, we have removed two unique
```

Here's the traversal in action:

k -= v

if k < 0:

Solution Implementation

from collections import Counter

from typing import List

class Solution:

**Python** 

Sorted counts: [1, 1, 2, 3]

The number of unique integers is initially 4.

integers (those with the initial frequency of 1), resulting in a remaining unique integer count of 2 (those with frequencies of 2 and 3 are still present).

def findLeastNumOfUniqueInts(self, arr: List[int], k: int) -> int:

# because all elements can be removed to achieve k deletion

# Reduce the count of deletable elements by the current count value

# If k becomes negative, we can't delete anymore unique integers

# Create a counter for all elements in the array

# Go through the counts starting from the smallest

for index, value in enumerate(sorted\_counts):

int findLeastNumOfUniqueInts(vector<int>& arr, int k) {

for (auto& [number, count] : frequencyMap) {

sort(frequencies.begin(), frequencies.end());

for (int i = 0; i < frequencies.size(); ++i) {</pre>

// Extract the frequencies and sort them in ascending order

// Subtract the frequency of the current number from k

unordered\_map<int, int> frequencyMap;

frequencies.push\_back(count);

++frequencyMap[number];

for (int number : arr) {

vector<int> frequencies;

k -= frequencies[i];

// Create a hashmap to count the occurrence of each integer in the array

// Determine the least number of unique integers by removing k occurrences

int uniqueIntegers = frequencies.size(); // start with all unique integers

return 4 - i # number of unique integers initially minus the index

**Traverse and Remove:** We then traverse the sorted list and remove k elements.

Count Frequency: We first count how many times each integer appears in the array:

- cnt =  $\{1: 2, 2: 1, 3: 3, 4: 1\}$  $sorted_counts = [1, 1, 2, 3]$ k = 3for i, v in enumerate(sorted\_counts):
- Return Result: With the loop broken, we know we've removed instances of unique integers only until the array's k becomes negative. This means we return 4 - 2, which equals 2. Therefore, the minimum number of unique integers remaining in the array arr after removing exactly k elements is 2.
- counter = Counter(arr) # Sort the counts of each unique integer sorted\_counts = sorted(counter.values())

```
if k < 0:
        # Return the count of remaining unique integers
        return len(counter) - index
# If k is not negative after trying to remove all, return 0
```

return 0

#include <algorithm>

class Solution {

public:

Java

k -= value

```
class Solution {
    public int findLeastNumOfUniqueInts(int[] arr, int k) {
       // Create a hashmap to store the frequency of each integer in the array
       Map<Integer, Integer> frequencyMap = new HashMap<>();
       // Populate the frequency map
        for (int num : arr) {
            frequencyMap.merge(num, 1, Integer::sum); // Increment the count for each occurrence of a number
       // Create a list to store the frequencies only
       List<Integer> frequencies = new ArrayList<>(frequencyMap.values());
       // Sort the frequencies in ascending order
       Collections.sort(frequencies);
       // Iterate over the list of frequencies
        for (int i = 0, totalUniqueNumbers = frequencies.size(); i < totalUniqueNumbers; ++i) {</pre>
            k -= frequencies.get(i); // Subtract the frequency from 'k'
            if (k < 0) {
                // If 'k' becomes negative, the current frequency can't be fully removed
                // so return the number of remaining unique integers
                return totalUniqueNumbers - i;
       // If all frequencies have been removed with 'k' operations, return 0 as there are no unique integers left
       return 0;
C++
#include <vector>
#include <unordered_map>
```

return 0;

from typing import List

class Solution:

from collections import Counter

counter = Counter(arr)

if k < 0:

return 0

```
// If k becomes negative, we can't remove any more numbers
            if (k < 0) {
                return uniqueIntegers - i; // Return the remaining number of unique integers
       // If k is non-negative after all removals, we've removed all duplicates
       return 0;
};
TypeScript
function findLeastNumOfUniqueInts(arr: number[], k: number): number {
    // Create a map to hold the frequency of each integer in the array
    const frequencyMap: Map<number, number> = new Map();
    // Iterate over the array and populate the frequency map
    for (const number of arr) {
        frequencyMap.set(number, (frequencyMap.get(number) || 0) + 1);
    // Extract the frequency values from the map and store them in an array
    const frequencies: number[] = [];
    for (const frequency of frequencyMap.values()) {
        frequencies.push(frequency);
    // Sort the frequencies array in ascending order
    frequencies.sort((a, b) => a - b);
    // Iterate over the sorted frequencies
    for (let i = 0; i < frequencies.length; ++i) {</pre>
       // Decrement k by the current frequency
       k -= frequencies[i];
       // If k becomes negative, we've used up k removals, so we return
       // the number of unique integers left, which is the length of the
       // frequencies array minus the current index
       if (k < 0) {
            return frequencies.length - i;
```

```
# Sort the counts of each unique integer
sorted_counts = sorted(counter.values())
# Go through the counts starting from the smallest
for index, value in enumerate(sorted_counts):
    # Reduce the count of deletable elements by the current count value
    k -= value
   # If k becomes negative, we can't delete anymore unique integers
```

# Return the count of remaining unique integers

# If k is not negative after trying to remove all, return 0

# because all elements can be removed to achieve k deletion

def findLeastNumOfUniqueInts(self, arr: List[int], k: int) -> int:

# Create a counter for all elements in the array

return len(counter) - index

// If we've processed all frequencies and haven't used up k removals,

// all integers have been removed and 0 unique integers are left

Time and Space Complexity The time complexity of the given code is  $0(n \log n)$ . This complexity arises from the sorting operation sorted(cnt.values()) where cnt.values() represents the counts of unique integers in the array arr; sorting these counts requires 0(n log n) time since sorting is typically done using comparison-based algorithms like quicksort or mergesort which have 0(n log n) complexity

in the average and worst case. The space complexity of the code is O(n) because we are storing counts of the elements in the array in a dictionary cnt. In the worst case, if all elements are unique, it will contain n key-value pairs which relate directly to the size of the input array arr.