206. Reverse Linked List

Linked List

Problem Description

Recursion

Easy

value and a pointer/reference to the next node in the sequence. A singly linked list means that each node points to the next node and there is no reference to previous nodes. The problem provides a pointer to the head of the linked list, where the 'head' represents the first node in the list. Our goal is to take this linked list and return it in the reversed order. For instance, if the linked list is $1 \rightarrow 2 \rightarrow 3 \rightarrow \text{null}$, the reversed list should be $3 \rightarrow 2 \rightarrow 1 \rightarrow \text{null}$. Intuition

To reverse the <u>linked list</u>, we iterate over the original list and rearrange the next pointers without creating a new list. The intuition

behind this solution is to take each node and move it to the beginning of the new reversed list as we traverse through the original

The task is to reverse a singly <u>linked list</u>. A linked list is a data structure where each element (often called a 'node') contains a

list. We maintain a temporary node, often referred to as a 'dummy' node, which initially points to null, as it will eventually become the tail of the reversed list once all nodes are reversed. We iterate from the head towards the end of the list, and with each iteration, we do the following: • Temporarily store the next node (since we are going to disrupt the next reference of the current node).

• Set the next reference of the current node to point to what is currently the first node of the reversed list (initially, this is null or dummy next).

• Move to the next node in the original list using the reference we stored earlier.

- This process ensures that we do not lose track of the remaining parts of the original list while building the reversed list. After we

• Move the dummy's next reference to the current node, effectively placing the current node at the beginning of the reversed list.

have iterated through the entire original list, the dummy next will point to the new head of the reversed list, which we then return as the result.

Solution Approach The provided solution employs an iterative approach to go through each node in the linked list and reverse the links. Here's a step-by-step walk-through of the algorithm used:

A pointer called curr is initialized to point to the head of the original list. This pointer is used to iterate over the list.

The iteration starts with a while loop which continues as long as curr is not null. This ensures we process all nodes in the list. Inside the loop, next temporarily stores currenext, which is the pointer to the next node in the original list. This is crucial since

we are going to change currenext to point to the new list and we don't want to lose the reference to the rest of the original

iteration, the current node now points to the head of the reversed list.

The new reversed list referenced by dummy next is returned.

O(n), where n is the number of nodes in the list.

We want to reverse it to become:

3 -> 2 -> 1 -> null

dummy -> null

A new ListNode called dummy is created, which acts as the placeholder before the new reversed list's head.

- list. We then set curr.next to point to dummy.next. Since dummy.next represents the start of the new list, or null in the first
- dummy next is updated to curr to move the starting point of the reversed list to the current node. At this point, curr is effectively inserted at the beginning of the new reversed list.
- curr is updated to next to move to the next node in the original list, using the pointer we saved earlier. Once all nodes have been processed and the loop exits, dummy next will be the head of the new reversed list.

By updating the next pointers of each node, the solution reverses the direction of the list without allocating any additional nodes,

We create a ListNode called dummy that will initially serve as a placeholder for the reversed list. At the beginning, dummy next is

We store currenext in next, so next points to 2. next will help us move forward in the list after we've altered currenext.

5. We update curr.next to point to dummy.next, which is currently null. Now the first node (1) points to null, the start of our new reversed list.

which makes it an in-place reversal with a space complexity of O(1). Each node is visited once, resulting in a time complexity of

- **Example Walkthrough**
- Let's illustrate the solution approach with a small example. Suppose we have the following linked list: 1 -> 2 -> 3 -> null

set to null. We initialize a pointer curr to point to the head of the original list which is the node with the value 1.

Here's the step-by-step process to achieve that using the provided algorithm:

Starting the iteration, we enter the while loop since curr is not null.

```
dummy -> null <- 1 2 -> 3 -> null
```

dummy -> 1 -> null

dummy -> 1 -> null

curr -> 2 -> 3 -> null

dummy -> 2 -> 1 -> null

curr ----| 3 -> null next --^

dummy -> 3 -> 2 -> 1 -> null

then updated to the null we saved in next:

next -> 2 -> 3 -> null

curr -> 1 -> 2 -> 3 -> null

curr ----| 7. We update curr to next, moving forward in the original list. curr now points to 2.

8. The loop continues. Again, we save currinext to next, and update currinext to point to dummy next. Then we shift the start of the reversed list

9. In the final iteration, we perform similar steps. We save currenext to next, set currenext to dummy next, and move dummy next to curre curreis

Once curr is null, the while loop terminates, and we find that dummy next points to 3, which is the new head of the reversed

by setting dummy next to the current node and update curr to next. After this iteration, dummy points to the new head 2, and our reversed list grows:

6. We move the start of the reversed list to curr by setting dummy next to curr. The reversed list now starts with 1.

Lastly, we return the reversed list starting from $dummy_next$, which is $3 \rightarrow 2 \rightarrow 1 \rightarrow null$. And that completes the reversal of our linked list using the iterative approach described in the solution.

Solution Implementation

Definition for singly-linked list.

dummy_node = ListNode()

self.val = val

self.next = next

def __init__(self, val=0, next=None):

while current_node is not None:

next_node = current_node.next

dummy_node.next = current_node

current_node = next_node

return dummy_node.next

current_node.next = dummy_node.next

Move to the next node in the original list

// Dummy node that will help in reversing the list.

// Move to the next node in the original list.

// Move to the next node in the original list.

// The head of the new reversed list is 'dummy->next.'

current = nextNode;

// Definition for a node in a singly-linked list

return dummy->next;

};

/**

*/

TypeScript

interface ListNode {

val: number;

next: ListNode | null;

return previousNode;

Definition for singly-linked list.

dummy_node = ListNode()

current_node = head

self.val = val

self.next = next

def __init__(self, val=0, next=None):

Start from the head of the list

Iterate over the linked list

while current_node is not None:

next_node = current_node.next

dummy_node.next = current_node

current_node = next_node

current_node.next = dummy_node.next

Move to the next node in the original list

Save the next node

def reverseList(self, head: ListNode) -> ListNode:

class ListNode:

class Solution:

* Reverses a singly linked list.

// Return the reversed linked list which is pointed by dummy's next.

ListNode dummy = new ListNode();

ListNode current = head;

while (current != null) {

// Pointer to traverse the original list.

// Iterating through each node in the list.

// Temporary store the next node.

ListNode nextTemp = current.next;

current.next = dummy.next;

dummy.next = current;

current = nextTemp;

return dummy.next;

Save the next node

def reverseList(self, head: ListNode) -> ListNode:

Initialize a dummy node, which will be the new head after reversal

The dummy node's next now points to the head of the reversed list

Reverse the link so that current_node.next points to the node before it

// Reversing the link so that current.next points to the new head (dummy.next).

// Move the dummy's next to the current node making it the new head of the reversed list.

curr ----

list.

Python

class ListNode:

class Solution:

10.

Start from the head of the list current_node = head # Iterate over the linked list

```
Java
// Definition for singly-linked list.
class ListNode {
    int val;
    ListNode next;
    ListNode() {}
    ListNode(int val) { this.val = val; }
    ListNode(int val, ListNode next) { this.val = val; this.next = next; }
class Solution {
    /**
    * Reverses the given linked list.
    * @param head The head of the original singly-linked list.
    * @return The head of the reversed singly-linked list.
    public ListNode reverseList(ListNode head) {
```

```
C++
// Definition for singly-linked list node.
struct ListNode {
   int val;
                   // The value of the node.
                       // Pointer to the next node in the list.
   ListNode *next;
   // Default constructor initializes with default values.
   ListNode(): val(0), next(nullptr) {}
   // Constructor initializes with a given value and next pointer set to nullptr.
   ListNode(int x) : val(x), next(nullptr) {}
   // Constructor initializes with a given value and a given next node pointer.
   ListNode(int x, ListNode *next) : val(x), next(next) {}
class Solution {
public:
   // Function to reverse a singly-linked list.
   ListNode* reverseList(ListNode* head) {
       // The 'dummy' node acts as the new head of the reversed list.
       ListNode* dummy = new ListNode();
       // 'current' node will traverse the original list.
       ListNode* current = head;
       // Iterate through the list until we reach the end.
       while (current != nullptr) {
           // 'nextNode' temporarily stores the next node.
           ListNode* nextNode = current->next;
           // Reverse the 'current' node's pointer to point to the new list.
            current->next = dummy->next;
           // The 'current' node is prepended to the new list.
           dummy->next = current;
```

```
function reverseList(head: ListNode | null): ListNode | null {
   // Return immediately if the list is empty
   if (head === null) {
       return head;
   // Initialize pointers
    let previousNode: ListNode | null = null; // Previous node in the list
    let currentNode: ListNode | null = head; // Current node in the list
   // Iterate through the list
   while (currentNode !== null) {
       const nextNode: ListNode | null = currentNode.next; // Next node in the list
       // Reverse the current node's pointer
       currentNode.next = previousNode;
       // Move the previous and current pointers one step forward
       previousNode = currentNode;
       currentNode = nextNode;
```

// By the end, previousNode is the new head of the reversed linked list

Initialize a dummy node, which will be the new head after reversal

* @param {ListNode | null} head - The head node of the linked list to be reversed

* @return {ListNode | null} The new head of the reversed linked list

```
# The dummy node's next now points to the head of the reversed list
       return dummy_node.next
Time and Space Complexity
  The time complexity of the provided code is O(n), where n is the number of nodes in the linked list. This is because the code
```

Reverse the link so that current_node.next points to the node before it

iterates through all the nodes in the list a single time. The space complexity of the code is 0(1). The space used does not depend on the size of the input list, since only a finite number of pointers (dummy, curr, next) are used, which occupy constant space.