String

Recursion

Stack

Hard

one of several forms: a direct true or false value, a negation of a subexpression, a logical AND of one or more subexpressions, or a logical OR of one or more subexpressions. The different expressions are represented as: 't': evaluates to true

The problem at hand entails evaluating a boolean expression that can be either true ('t') or false ('f'). The expression can take

'f': evaluates to false

subExprn (with n >= 1)

- '!(subExpr)': evaluates to the logical NOT of the inner expression subExpr • '&(subExpr1, subExpr2, ..., subExprn)': evaluates to the logical AND of all the inner expressions subExpr1, subExpr2, ...,
 - subExprn (with n >= 1) • '|(subExpr1, subExpr2, ..., subExprn)': evaluates to the logical OR of all the inner expressions subExpr1, subExpr2, ...,
- The goal is to compute the result of the given boolean expression, assured by the problem's constraints to be valid and conform to the rules above.
- Intuition

The intuition behind the solution for evaluating the boolean expression is to process the expression in a structured way, making use of a stack to keep track of the characters that form a subexpression until it's complete and can be evaluated. This resembles the

classic stack-based evaluation of expressions which is commonly used in parsing algorithms.

Here's how the approach breaks down: 1. Initialize a Stack: Since the expression is nested with parentheses, we can use a stack to keep track of and evaluate the inner

∘ If we encounter a t, f, !, &, or |, we push it onto the stack as these characters indicate the beginning of an expression or a value.

which corresponds to the truth values.

equal to 't' to return True else it will be False.

data structure to keep track of the components of the expression.

logical operations as specified by the nested structure of the parentheses and operators.

• Initialize counters for true (t) and false (f) occurrences within the subexpression.

expressions in the correct order.

2. Iterate through Each Character:

∘ If a character is a), it means we have reached the end of an inner expression, so we begin to evaluate it. 3. Evaluate the Expression:

• We keep popping from the stack until we get t, f, or an operator. Here, we count how many t and f we have encountered

- Based on the operator before the sequence of t and f values (!, &, |), we apply the NOT, AND, or OR operation accordingly. For AND, the result is false if any f is found; for OR, the result is true if any t is found; for NOT, the result is true if f is found
- and vice versa.
- The result of the inner expression (resulting t or f) is then pushed back onto the stack as the new 'value' for that subexpression. 4. Result:

After evaluating the entire expression, the last element on the stack is the value of the whole expression. We check if it's

Solution Approach To decipher and evaluate the boolean expression provided in the problem statement, the solution leverages the use of a stack (stk)

This approach carefully constructs truth values by evaluating the smallest subexpressions first and combining them according to

 Iterate through every character (c) in the string expression. • When encountering characters that are 't', 'f', '!', '&', or '|', they are pushed onto the stack (stk), as these signify values or operators.

• If the character is a closing parenthesis ')', it indicates the end of a subexpression and triggers the evaluation sequence:

For !: Invert the boolean value, push 't' if there was an f in the subexpression, and 'f' if there was a t.

After the iteration is complete, the last character left in the stack is the result of the entire expression. Return True if this

Pop characters from the stack until an operator is found, counting the t and f values found in the way. ○ Based on the popped operator, apply the logical NOT (!), AND (६), or OR (|) operation:

conditions.

Example Walkthrough

Step 1: Initialize Stack

Step 2: Start Iterating

Initially, our stack (stk) is empty.

We encounter '(' and ignore it.

We encounter ',' and ignore it.

Step 4: Evaluate NOT Subexpression

on stk: stack = ['&', 't', 't', 't']

Here is how the implementation works:

■ For &: If any f was found, the result of the AND operation is 'f', otherwise, it's 't'. ■ For |: If any t was found, the result of the OR operation is 't', otherwise, it's 'f'. The result of this operation replaces the subexpression by being pushed onto the stack.

A key algorithmic pattern used here is stack-based parsing, which is a common approach for handling nested expressions. Using a

effectively evaluate the boolean expressions as required by the problem's specifications.

character is 't', otherwise return False.

available after inner expressions are evaluated.

the results into higher-level expressions—essentially a form of bottom-up evaluation. This pattern is effective because we process the expression linearly, maintaining only the information necessary for the current subexpression's evaluation. The stack implicitly handles the "waiting" operations, resuming them when their operands become

In Python, the solution makes use of the match statement (available from Python 3.10), which provides a clean and readable way to

perform actions based on the value of the operator. This increases the clarity of the code, especially when dealing with multiple

Overall, the combination of a simple stack along with a case-based action sequence allows the implementation to concisely and

stack allows us to respect the precedence of the logical operations, dealing with the innermost expressions first before combining

et's consider a small example that illustrates the solution approach: Given the expression: &(t, |(f,t),!(f)) Now, let's step through the process of evaluating it:

 We encounter '&' and push it onto stk: stack = ['&'] We encounter '(' and ignore it as it's not part of evaluation.

• We found a 't' in our count, so the result of the OR subexpression is 't'. We push this result on stk: stack = ['&', 't', 't']

• We start by popping from stk and we find 'f'. Once we hit '!', we apply NOT on our f, which makes it a 't'. We push this 't'

At this point, we've evaluated all inner expressions and are left with the AND operation to perform on the remaining stack elements.

We encounter ',' and ignore it. We encounter 't' and push it onto stk: stack = ['&', 't', '|', 'f', 't']

Step 3: Evaluate OR Subexpression

We encounter 't' and push it onto stk: stack = ['&', 't']

We encounter ',' (a delimiter between expression values) and ignore it.

We encounter 'f' and push it onto stk: stack = ['&', 't', '|', 'f']

We encounter '|' and push it onto stk: stack = ['&', 't', '|']

• We start popping from stk until we hit the | operator: we get 't' and 'f' (we count 1 true and 1 false). Once we hit '|', we check for the | operation rules.

We encounter '!' and push it onto stk: stack = ['&', 't', 't', '!']

• Once we've evaluated this final expression, our stack is now: stack = ['t'].

Iterate through each character in the expression

true_count += stack[-1] == 't'

false_count += stack[-1] == 'f'

result_char = 't' if false_count else 'f'

result_char = 'f' if false_count else 't'

result_char = 't' if true_count else 'f'

Push the result of the evaluation back onto the stack

// Stack to hold characters representing boolean values and operators

int countTrue = 0; // Count of 't' (true) characters

int countFalse = 0; // Count of 'f' (false) characters

true_count = false_count = 0

while stack[-1] in 'tf':

elif operator == '&':

elif operator == '|':

stack.append(result_char)

public boolean parseBoolExpr(String expression) {

for (char c : expression.toCharArray()) {

stack.push(c);

} else if (c == ')') {

Deque<Character> stack = new ArrayDeque<>();

// Iterate through each character in the expression

if (c != '(' && c != ')' && c != ',') {

for char in expression:

if char in 'tf!&|':

elif char == ')':

stack.append(char)

Our stack is left with a single element 't', which means our entire expression evaluates to True.

by evaluating inner subexpressions and progressively building up to the entire expression's value.

If it's an operator or a boolean value, push it onto the stack

When we encounter a closing parenthesis, we should evaluate the expression

Count the number of true ('t') and false ('f') until we hit an operator

For and, result is false if we have any false, otherwise it's true

For or, result is true if we have any true, otherwise it's false

// If the character is not a parenthesis or comma, push it onto the stack

// When a closing parenthesis is encountered, evaluate the expression inside

• We encounter ')' and now we have to evaluate the subexpression with OR.

 We encounter '(' and ignore it. We encounter 'f' and push it onto stk: stack = ['&', 't', 't', '!', 'f'] We encounter ')' and now evaluate the NOT subexpression.

• We keep popping and find all ts, meaning that our AND operation will result in true. The top-level operator '&' ensures we check all values.

Step 5: Final Evaluation

Python Solution

Step 6: Result

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class Solution: def parseBoolExpr(self, expression: str) -> bool: stack = []

This step-by-step process demonstrates how the stack-based evaluation method effectively breaks down and solves the expression

18 stack.pop() # Remove the boolean value as it's counted 19 20 # Pop the operator from the stack and apply the logic for each operator 21 operator = stack.pop() if operator == '!': 22 23 # For not, result is true if we have a false, otherwise it's false

35 # At the end, the stack should only have one element which is the result 36 # Return True if the result is 't', otherwise False 37 return stack[0] == 't' 38

Java Solution

class Solution {

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16
                     // Pop characters from the stack until the corresponding operator is found
                     while (stack.peek() == 't' || stack.peek() == 'f') {
 17
 18
                         if (stack.peek() == 't') {
 19
                             countTrue++;
 20
                         } else {
 21
                             countFalse++;
 22
 23
                         stack.pop();
 24
 25
                     // Pop the operator ('!', '&', or '|')
 26
 27
                     char operator = stack.pop();
 28
 29
                     // Evaluate the expression based on the operator
                     if (operator == '!' && countFalse > 0) {
 30
 31
                         // For '!', the expression is true if it has any 'f' (since '!' is a negation)
 32
                         c = 't';
 33
                     } else if (operator == '&' && countFalse == 0) {
 34
                         // For '&', the expression is true if all are 't'
 35
                         c = 't';
 36
                     } else if (operator == '|' && countTrue > 0) {
                         // For '|', the expression is true if at least one is 't'
 37
 38
                         c = 't';
 39
                     } else {
                         // In all other cases, the expression is false
 40
                         c = 'f';
 41
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 44
                     // Push the result of the evaluated expression back onto the stack
 45
                     stack.push(c);
 46
 47
             // The result of the entire expression is at the top of the stack
 48
             // Return true if the top of the stack is 't', otherwise false
 49
             return stack.peek() == 't';
 50
 51
 52 }
 53
C++ Solution
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// The top of the stack now contains the result, return 'true' if 't', 'false' otherwise. 43 44 return charStack.top() == 't'; 45 46 }; 47

Typescript Solution

2 // Supported operators: '!', '&', and '|'

3 // Operands can be 't' (true) or 'f' (false)

1 class Solution {

// Function to evaluate the boolean expression.

stack<char> charStack; // Create a stack to manage the characters.

// Process all operands up to the previous operator.

trueCount += charStack.top() == 't';

falseCount += charStack.top() == 'f';

// Evaluate the expression based on the operator.

// Now, pop the operator from the stack.

char operatorChar = charStack.top();

if (operatorChar == '!') {

if (operatorChar == '&') {

if (operatorChar == '|') {

1 // Parses the given boolean expression and returns the result.

4 // Expression examples: "|(f,t)", "&(t,f)", "!(t)", "&(!(&(t,f,t)),|(f,t))"

5 // Each expression is guaranteed to be valid and only consist of these characters.

charStack.push(c);

// Push the result onto the stack.

while (charStack.top() == 't' || charStack.top() == 'f') {

// If the character is not an operator or a comma, push it onto the stack.

// Increment trueCount or falseCount based on the operand.

charStack.pop(); // Remove the processed operand from the stack.

} else if (c == ')') { // When a closing bracket is encountered, evaluate the expression inside it.

int trueCount = 0, falseCount = 0; // Variables to count the number of true and false values.

c = falseCount > 0 ? 't' : 'f'; // Logical NOT: 'true' if any is 'false', else 'false'.

c = falseCount > 0 ? 'f' : 't'; // Logical AND: 'false' if any is 'false', else 'true'.

c = trueCount > 0 ? 't' : 'f'; // Logical OR: 'true' if any is 'true', else 'false'.

// Iterating over each character in the input expression.

if (c != '(' && c != ')' && c != ',') {

bool parseBoolExpr(string expression) {

charStack.push(c);

charStack.pop();

for (char c : expression) {

public:

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function parseBoolExpr(expression: string): boolean {
         // The current index pointer used while scanning through the expression
         let currentIndex = 0;
  8
         // The length of the expression
  9
         const expressionLength = expression.length;
 10
 11
 12
         // Recursive function to evaluate the expression
         const evaluate = (): boolean[] => {
 13
 14
             let resultStack: boolean[] = [];
 15
             while (currentIndex < expressionLength) {</pre>
 16
                 const currentChar = expression[currentIndex++];
 17
 18
                 if (currentChar === ')') {
 19
                     // End of the current expression
 20
                     break;
 21
 22
 23
                 if (currentChar === '!') {
 24
                     // Not operation: Flip the boolean value of the next expression
 25
                     resultStack.push(!evaluate()[0]);
                 } else if (currentChar === '|') {
 26
 27
                     // Or operation: Return true if any of the values in the next expression is true
 28
                     resultStack.push(evaluate().some(value => value));
 29
                 } else if (currentChar === '&') {
 30
                     // And operation: Return true only if all values in the next expression are true
 31
                     resultStack.push(evaluate().every(value => value));
 32
                 } else if (currentChar === 't') {
 33
                     // Literal 'true' value
                     resultStack.push(true);
                 } else if (currentChar === 'f') {
 35
                     // Literal 'false' value
 36
                     resultStack.push(false);
 37
 38
 39
 40
             return resultStack;
         };
 41
 42
 43
         // Start evaluating the expression and return the result of the top-level expression
         return evaluate()[0];
 44
 45
 46
Time and Space Complexity
The time complexity of the given code is O(n), where n is the length of the expression string. This is because the code iterates over
each character in the input string exactly once. During the iteration, each character is processed in a way that takes constant time.
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Operations like pushing to and popping from the stack, comparing characters, and assigning values all occur in 0(1) time.

The space complexity of the code is also O(n). This is due to the use of a stack that, in the worst case, might need to store an element for every character in the expression string before any reduction happens. The case that would result in this space usage is one where the string contains a long series of open parentheses before reaching any operator or closed parenthesis.