### 219. Contains Duplicate II



### **Problem Description**

The goal is to find out whether there is a pair of indices i and j in the given integer array nums such that nums[i] is equal to nums[j] and the absolute difference between i and j is less than or equal to k. This problem is checking for duplicates within a specified range of indices in an array. The crucial point here is that i and j must be distinct, which means you cannot compare an element with itself. If such a pair exists, then the correct output is true; otherwise, false.

#### Intuition

through the nums array, we keep updating the hash table with the elements as keys and their indices as values. At each step, we look at the current element x and check if it's already in the hash table. If it is, and the stored index is close enough to the current index (the difference is less than or equal to k), then we've found the required pair, and we can immediately return true. If there's no such element or the distance is too large, we simply keep the hash table updated by overwriting the index of the current element. After checking all elements, if no such pair is found, the function returns false.

The intuitive approach for this problem uses a hash table to track the indices of elements we've already seen. As we iterate

This approach is efficient as it only requires a single pass through the array, and the lookups and updates in the hash table are

performed in constant time. Hence, the total time complexity is O(n), where n is the number of elements in the nums array.

### The algorithm utilizes a hash table to implement the solution efficiently. The hash table stores elements from the array nums as

Solution Approach

keys, and the most recent index at which each element occurs as the corresponding value.

We initiate the iteration through the array with a for loop. At each step, we consider the current element and its index (i, x)

fetched from enumerate(nums)).

1. If the current element x already exists in the hash table d and the difference between the current index i and the previously

Here's the step-by-step implementation:

- stored index for x (d[x]) is less than or equal to k, then we have found two indices i and j (where j is d[x]) that satisfy both conditions: nums[i] == nums[j] and  $abs(i j) \ll k$ . In this case, the function immediately returns true.

  2. If the current element x does not satisfy the condition, or if it is not present in the hash table, it means we have not yet found
- a duplicate within the specified range. Therefore, we update the hash table with the current index i for element x: d[x] = i.

  This step is crucial because it keeps the hash table updated with the latest index where each number is found, ensuring that when checking for the range condition, we are always using the smallest possible index difference.

  3. After iterating through all the elements in the array, if we have not returned true, it means no two indices satisfy the given
- conditions, and the function concludes by returning false.

  In summary, the hash table is essential to keep track of the indices of the elements we've seen, enabling us to check the

length of the array nums.

Example Walkthrough

Let's take a small example to better understand the solution approach. Consider the array nums = [1, 2, 3, 1, 2, 3] and let's

existence of nearby duplicates in constant time. This approach leads to a linear time complexity, which is O(n), with n being the

#### say k = 2. We want to find if there are any duplicates within an index range of 2.

We start with an empty hash table d = {}.
 We iterate through the nums array with index i and element x.

b. At i = 1, x = 2. The hash table is now: d = 1

equal to k. So, the function returns false.

for index. value in enumerate(nums):

b. At i = 1, x = 2. The hash table is now:  $d = \{1: 0, 2: 1\}$ .

```
c. At i = 2, x = 3. The hash table updates to: d = \{1: 0, 2: 1, 3: 2\}.
```

a. At i = 0, x = 1. We add 1 to the hash table with its index:  $d = \{1: 0\}$ .

which is 0. The difference between the current index 3 and the last index 0 is 3 which is not less than or equal to k. We

def containsNearbyDuplicate(self, nums: List[int], k: int) -> bool:

# If the value is in index map, and the current index minus the

for (int currentIndex = 0; currentIndex < nums.length; ++currentIndex) {</pre>

// If so, we found a nearby duplicate and return true.

if (currentIndex - lastIndex <= k) {</pre>

indexMap.put(nums[currentIndex], currentIndex);

return true;

// Use getOrDefault to find the last index of the current number.

int lastIndex = indexMap.getOrDefault(nums[currentIndex], -1000000);

// This function checks if there are any duplicates within 'k' indices apart in the array 'nums'

// Check if the current number is in the map and if the difference between

function containsNearbyDuplicate(nums: number[], k: number): boolean {

// Create a map to keep track of the numbers and their indices

const indexMap: Map<number, number> = new Map();

for (let index = 0; index < nums.length; ++index) {</pre>

// the indices is less than or equal to 'k'

// Iterate through the 'nums' array

// If it's not found, use a default value that is far away from any possible index.

// Check if the current index and the last index of the same value are within k steps

// Update the map with the current value and its index for the next iteration checks

# stored index is less than or equal to k, a nearby duplicate exists

update d with the new index for 1:  $d = \{1: 3, 2: 1, 3: 2\}$ . e. At i = 4, x = 2. Again, 2 is already in the hash table. The last occurrence was at index 1. The difference between i = 4 and 1 is 3, which again, is not less than or equal to k. We update d:  $d = \{1: 3, 2: 4, 3: 2\}$ .

d. At i = 3, x = 1. The current element 1 is already in the hash table. We check the index of the last occurrence of 1 in d,

- f. At i = 5, x = 3. The 3 is found in the hash table at index 2. The difference between i = 5 and 2 is 3, which is also not less than or equal to k. Update d: d = {1: 3, 2: 4, 3: 5}.

  3. After completing the iteration, we have not found any elements where the absolute difference in their indices was less than or
- Through this example, it's clear how the hash table d is used to keep track of the indices of elements. Despite encountering duplicates, the conditions for k were not satisfied, which the algorithm correctly identified.

Solution Implementation

Python

# # Initialize a dictionary to store the value as key and its index as value index\_map = {} # Enumerate over the list to have both index and value

class Solution:

```
if value in index_map and index - index_map[value] <= k:</pre>
                return True
            # Update the index value in the index map for each value
            # It ensures that if the same value comes up again, the index_map stores
            # the latest index, which is useful for distance calculation
            index_map[value] = index
        # Return False if no nearby duplicates found within the given k distance
        return False
To ensure Python 3 syntax, especially for typing:
- Make sure `List` is imported from `typing` module, which will be used for the type hint of the `nums` argument.
Here is how you would include that import:
```python
from typing import List
class Solution:
    # the rest of the code remains the same as above
Java
import java.util.HashMap;
import java.util.Map;
class Solution {
    public boolean containsNearbyDuplicate(int[] nums, int k) {
        // Initialize a HashMap to store the value and its most recent index
        Map<Integer, Integer> indexMap = new HashMap<>();
```

```
// If no nearby duplicates are found, return false.
        return false;
C++
#include <vector>
#include <unordered map>
using namespace std;
class Solution {
public:
    bool containsNearbyDuplicate(vector<int>& nums, int k) {
        // Create a hashmap to store the most recent index of each value observed
        unordered_map<int, int> valueToIndexMap;
        for (int i = 0; i < nums.size(); ++i) {</pre>
            // Check if the current value exists in the map, i.e., it has been encountered before
            if (valueToIndexMap.count(nums[i])) {
                // If the previous occurrence is within k indices from the current index, return true
                if (i - valueToIndexMap[nums[i]] <= k) {</pre>
                    return true;
            // Update the hashmap with the current index for this value
            valueToIndexMap[nums[i]] = i;
        // If no duplicate elements are within the given k indices, return false
        return false;
};
```

# if (indexMap.has(nums[index]) && index - indexMap.qet(nums[index])! <= k) { return true; // Duplicate found within 'k' distance }</pre>

**TypeScript** 

```
// Update the index of the current number in the map
        indexMap.set(nums[index], index);
    // No duplicate found within 'k' distance
    return false;
class Solution:
   def containsNearbyDuplicate(self, nums: List[int], k: int) -> bool:
        # Initialize a dictionary to store the value as key and its index as value
        index_map = {}
        # Enumerate over the list to have both index and value
        for index, value in enumerate(nums):
           # If the value is in index map, and the current index minus the
           # stored index is less than or equal to k, a nearby duplicate exists
            if value in index_map and index - index_map[value] <= k:</pre>
                return True
           # Update the index value in the index map for each value
           # It ensures that if the same value comes up again, the index map stores
           # the latest index, which is useful for distance calculation
            index_map[value] = index
        # Return False if no nearby duplicates found within the given k distance
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## class Solution: # the rest of the code remains the same as above

Time and Space Complexity

once. Inside the loop, the code performs constant time operations including checks in a dictionary and assignment.

The space complexity is also O(n) as the code allocates a dictionary 'd' that could potentially store all the unique elements of 'nums' if there are no nearby duplicates. In the worst case, the dictionary will have as many entries as there are elements in 'nums'.

The time complexity of the provided code is O(n) because the code uses a single loop that iterates over the elements in 'nums'