123. Best Time to Buy and Sell Stock III Array Dynamic Programming Hard

Leetcode Link

Problem Description We are presented with an array called prices, where each element prices [i] represents the price of a stock on day i. The challenge

is to determine the maximum profit that can be made through at most two stock transactions. A transaction consists of buying and then selling one share of the stock. It is important to note that you cannot hold more than one share at a time, which means you must sell the share you hold before buying another one. Therefore, the goal is to strategically choose two periods of time to buy and sell

1. Buying the first stock (f1): For this action, we want to minimize the cost, so we keep track of the lowest price we've seen so far. 2. Selling the first stock (f2): Here, we calculate the profit from the first transaction. We aim to maximize the profit by selling at the

highest price after buying at the lowest price. 3. Buying the second stock (f3): For this, we want the net cost to be minimal, which is purchasing a second stock at the lowest

effective price after accounting for the profit from the first sale. This means we subtract the price of the second stock from the

- profit made from selling the first stock. 4. Selling the second stock (f4): Finally, we want to maximize the total profit, which comes from selling the second stock at the
- highest possible price. The algorithm operates by iterating through the price array and updating these four states. Each price offers a potential to update these states—either you get a better buy price (f1), a better sell price resulting in higher first transaction profit (f2), a lower net buy
- **Solution Approach**

price (after adjusting for profit from f2) for the second transaction (f3), or a better final sell price for a higher overall profit (f4).

The solution leverages dynamic programming to break down the problem into smaller subproblems and builds up the answer from those subproblems. Here's how the implementation carries this out:

By the end of the iteration, f4 will represent the maximum profit achievable with up to two trades.

at the lowest price seen so far.

- f1 = max(f1, -price) Update 12 by taking the maximum of the current 12 and the current price plus 11. This reflects the sale of the first stock and captures the highest profit possible up to this point after the first transaction.
- f2 = max(f2, f1 + price)
 - the lowest effective price by considering the profits from the first sale.
 - f4 = max(f4, f3 + price)

return as the final result.

1 prices = [3, 4, 5, 1, 3, 2, 10]

Following the solution approach:

most two transactions.

increase the computational complexity. Example Walkthrough

In this array, prices [i] represents the price of the stock on day i. We want to calculate the maximum profit that can be made with at

Let's walk through the solution approach using a small example. Consider the following array for prices:

1. Initializing Variables: f1 = -3 (since we "buy" the stock on day 0 at price 3) f2 = 0 (no sale yet)

 \circ f4 = max(0, -3 + 4) = 1 (after possible second sale) Day 2 (price = 5):

 \circ f1 = max(-3, -5) = -3

 \circ f2 = max(1, -3 + 5) = 2

 \circ f3 = max(-3, 2 - 5) = -3

Day 1 (price = 4):

 \circ f1 = max(-3, -1) = -3

• f1 = max(-3, -4) = -3 (we keep the previous buy since it's cheaper)

 \circ f2 = max(0, -3 + 4) = 1 (we can sell the first stock bought at -3 for 4)

- \circ f2 = max(2, -3 + 1) = 2 • f3 = max(-3, 2 - 1) = 1 (we now buy the second stock since we have a net profit after first sale) \circ f4 = max(2, 1 + 1) = 2
- Day 4 (price = 3):
- \circ f3 = max(1, 2 3) = 1

 \circ f1 = max(-3, -2) = -3

 \circ f1 = max(-3, -10) = -3

Python Solution

class Solution:

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from typing import List

- f4 = max(2, 1 + 3) = 4 (we sold the second stock we bought on day 3 for 3, making a profit) Day 5 (price = 2):
- \circ f3 = max(1, 2 2) = 1 \circ f4 = max(4, 1 + 2) = 4
- \circ f2 = max(2, -3 + 10) = 7 \circ f3 = max(1, 7 - 10) = 1 \circ f4 = max(4, 1 + 10) = 11

Thus, we conclude the maximum profit possible is \$11.

def maxProfit(self, prices: List[int]) -> int:

first_buy, first_sell = -prices[0], 0

for price in prices[1:]:

second_buy, second_sell = -prices[0], 0

first_buy = max(first_buy, -price)

Initialize the four states of profits

first_buy represents the profit after the first buy

first_sell represents the profit after the first sell

second_buy represents the profit after the second buy

second_sell represents the profit after the second sell

Loop through each price starting from the second price

second_buy = max(second_buy, first_sell - price)

Update first_buy to be either its previous value

or the negative of the current price (representing a buy)

18 19 # Update first_sell to be either its previous value 20 # or the sum of first_buy and the current price (representing a sell) first_sell = max(first_sell, first_buy + price) 21 22 # Update second_buy to be either its previous value

or the difference of first_sell and the current price (representing a second buy)

- 32 # and that's the result to return 33 return second_sell 34 35 # Example usage: 36 # sol = Solution()
- Java Solution public class Solution {

public int maxProfit(int[] prices) {

int firstSellProfit = 0;

return secondSellProfit;

int buy1 = -prices[0];

int buy2 = -prices[0];

int sell1 = 0;

int firstBuyProfit = -prices[0];

int secondBuyProfit = -prices[0];

26 27 28 // Update the maximum profit after the second sell 29 30 secondSellProfit = Math.max(secondSellProfit, secondBuyProfit + prices[i]); 31

// Return the maximum profit after the second sell

- class Solution { public: int maxProfit(vector<int>& prices) { // buyl is the maximum profit we can achieve by buying a stock on day 0,
- 28 29 30 31 32 sell1 = std::max(sell1, buy1 + prices[i]);

buy2 = std::max(buy2, sell1 - prices[i]);

sell2 = std::max(sell2, buy2 + prices[i]);

44 // The result should be the maximum profit after two sales. 45 return sell2; 47 }; 48

let secondBuyProfit = -prices[0];

// At start, no transaction has happened, hence 0 profit 13 let secondSellProfit = 0; 14 15 16 // Loop to calculate profits at each price point for (let i = 1; i < prices.length; ++i) { 17 // The maximum profit of the first buy is either the previous value or 18

// the profit after selling at the current price

firstBuyProfit = Math.max(firstBuyProfit, -prices[i]);

- 25 26 27 28 29
- Time and Space Complexity The given Python code represents a dynamic programming approach to solve a stock trading problem, where an individual is allowed

The time complexity of the code is determined by:

transaction or waiting period.

Time Complexity:

• The number of iterations in the loop, which is n-1 where n is the length of the prices list. The constant time operations within each iteration, which don't depend on the size of the input.

to make at most two transactions to maximize profit. The variables f1, f2, f3, and f4 represent the four states of profits after each

Thus, the total time complexity is 0(n-1) * 0(1), which simplifies to 0(n).

The space complexity of the code is determined by:

Space Complexity:

- The space taken by the four variables f1, f2, f3, f4, which is constant and does not scale with input size. No additional data structures are used that grow with input size.
- Hence, the space complexity of the code is 0(1), representing the constant space used by the variables.

stocks to maximize your profit.

Intuition The intuition behind the solution involves dynamic programming to keep track of profits across four different states, which represent the four actions you can take:

1. Initializing Variables: We initialize four variables, f1, f2, f3, and f4, which correspond to the states of buying the first stock, selling the first stock, buying the second stock, and selling the second stock, respectively. Initially, we set f1 and f3 to prices [0] since we "buy" the stock on day zero (the maximum profit after buying is the negative value of the stock's price), and f2 and f4 to 0 as we have not made any profit yet.

2. Iterating through Prices: We iterate through the prices array starting from day 1, updating the four states with each new price. Update f1 by taking the maximum of the current f1 and the negative current price, indicating the purchase of the first stock

Update f3 by taking the maximum of the current f3 and f2 minus the current price. This reflects buying the second stock at

- f3 = max(f3, f2 price) Update f4 by taking the maximum of the current f4 and the current price plus f3. This updates the total profit after the
 - second sale is the highest we can get.
- By the end of the loop, we have considered all possible scenarios of one or two transactions through the days, ensuring the maximum profit is calculated. The key aspect of this solution is it avoids the need for nested loops, which would significantly

3. Returning Result: After iterating through all prices, f4 will hold the maximum profit after at most two transactions, which we

Each of the updates done in the loop effectively simulates all possible buy and sell actions that could be taken on that day and

carries forward the best decision. The reason we update in the order $f1 \rightarrow f2 \rightarrow f3 \rightarrow f4$ is due to the dependencies between the

transactions: selling the first stock must consider its purchase and buying the second stock must consider the sale of the first, etc.

• f3 = -3 (after "buying" the first stock, we haven't bought the second one yet) f4 = 0 (no second sale yet) 2. Iterating through Prices:

○ f3 = max(-3, 1 - 4) = -3 (considering profit from first sale, buying second stock at -3 is still better)

 \circ f4 = max(1, -3 + 5) = 2 Day 3 (price = 1):

- \circ f1 = max(-3, -3) = -3 \circ f2 = max(2, -3 + 3) = 2
- \circ f2 = max(2, -3 + 2) = 2 Day 6 (price = 10):

3. Returning Result: Lastly, after iterating through all prices, f4 holds the maximum profit of 11 after at most two transactions.

By following these steps using dynamic programming, we have efficiently computed the maximum profit possible with up to two

transactions without needing to iterate through every possible pair of buy and sell days, which would be far less efficient.

- # Update second_sell to be either its previous value # or the sum of second_buy and the current price (representing a second sell) second_sell = max(second_sell, second_buy + price)
- 30 31 # The maximum profit after two transactions is located in second_sell,
- 37 # max_profit = sol.maxProfit([3,3,5,0,0,3,1,4]) 38 # print(max_profit) # Output would be 6, representing the maximum profit possible with two transactions
- 11 12 // f4 represents the maximum profit after the second sell 13 int secondSellProfit = 0; 14 15 // Iterate through the list of prices starting from the second price 16 for (int i = 1; i < prices.length; ++i) {

// Update the maximum profit after the first buy

firstBuyProfit = Math.max(firstBuyProfit, -prices[i]);

// fl represents the maximum profit after the first buy

// f2 represents the maximum profit after the first sell

// f3 represents the maximum profit after the second buy

- // Update the maximum profit after the first sell // Equivalent to buying at 'firstBuyProfit' and selling at prices[i] firstSellProfit = Math.max(firstSellProfit, firstBuyProfit + prices[i]); // Update the maximum profit after the second buy // We subtract 'prices[i]' because it's the price we are buying at after the first sell secondBuyProfit = Math.max(secondBuyProfit, firstSellProfit - prices[i]); // Equivalent to buying at 'secondBuyProfit' and selling at prices[i]
- C++ Solution 1 #include <vector> 2 #include <algorithm> // for std::max function // Initialization of buyl and sell1 corresponds to the first transaction.

// sell1 is the maximum profit we can achieve by selling a stock after buying on day 0,

// initially this would be zero because we have just bought a stock and yet to sell it.

// initially it's the same as buyl because we have not executed the first sell yet.

// which is simply the negative value of the stock price on that day.

// Initialization of buy2 and sell2 corresponds to the second transaction.

// buy2 is the maximum profit we can potentially hold for a second buy,

// This is the maximum profit we can make with at most two transactions

19 // sell2 is the maximum profit we can get after completing the second sell, 20 // initially, this would also be zero because we haven't sold any stock twice. 21 int sell2 = 0; 22 23 for (int i = 1; i < prices.size(); ++i) {</pre> 24 // For each day, calculate the maximum profit if we were to buy the stock 25 // for the first transaction, by comparing the previous buy profit and 26 // the negative of today's price (as if we bought the stock today) buy1 = std::max(buy1, -prices[i]); 27 // Calculate the maximum profit if we were to sell the stock // for the first transaction. Compare the previous sell profit and // the profit from selling today (today's price plus the current buy profit). 33

// Finally, calculate the profit if we were to make the second sell today.

// which includes the potential second buy profit and today's price.

// It compares the current second sell profit and the profit from selling today,

// For the second buy, calculate the maximum profit by comparing the previous second buy profit

// and the current overall profit minus today's price (buying the second stock today).

Typescript Solution function maxProfit(prices: number[]): number { // Initialize profit states for the two transactions // firstBuyProfit assumes the first transaction hasn't happened yet, hence negative value let firstBuyProfit = -prices[0]; // firstSellProfit assumes no transaction has been made, hence 0 profit let firstSellProfit = 0;

// secondBuyProfit takes into account profit from the first transaction

// secondSellProfit represents the cumulative profit from both transactions

// The maximum profit of the second buy is the greater of the previous value or // the current profit from the first sell minus the current price secondBuyProfit = Math.max(secondBuyProfit, firstSellProfit - prices[i]); // The maximum profit of the second sell is the greater of the previous value or 30 // the profit after selling at the current price (including the profit from the first sell) 31 secondSellProfit = Math.max(secondSellProfit, secondBuyProfit + prices[i]); 33 34 35 // Return the cumulative maximum profit from both allowed transactions 36 return secondSellProfit; 37 }

// the negative of the current price (which means buying at the current price)

// The maximum profit of the first sell is either the previous value or

firstSellProfit = Math.max(firstSellProfit, firstBuyProfit + prices[i]);

There are four assignments in each iteration that happen in 0(1) time each.