

## Longest Common Subsequence (LCS)

### Longest Common Subsequence (LCS)

- A subsequence of a sequence  $S$  is obtained by deleting zero or more symbols from  $S$ . For example, the following are all subsequences of “president”: **pred**, **sdn**, **predent**.
- The longest common subsequence problem is to **find a maximum length common subsequence between two sequences**.

### LCS

For instance,

Sequence 1: president

Sequence 2: providence

Its LCS is priden.

```

president
| | | | |
| | | | |
providence
  
```

### LCS

Sequence 1: algorithm

Sequence 2: alignment

One of its LCS is algm.

```

a l g o r i t h m
| | | | |
a l i g n m e n t
  
```

## Longest Common Subsequence (LCS)

- S1:ACCGGTCGAGTGCGCGGAAGCCGGC CGAA
- S2:GTCGTTCGGAATGCCGTTGCTCTGTAA
- S3:GTCGTCGGAAGCCGGCCGAA

4/26/2018

5

## LCS Algorithm

- First we'll find the length of LCS. Later we'll modify the algorithm to find LCS itself.
- Define  $X_i$ ,  $Y_j$  to be the prefixes of  $X$  and  $Y$  of length  $i$  and  $j$  respectively
- Define  $c[i,j]$  to be the length of LCS of  $X_i$  and  $Y_j$
- Then the length of LCS of  $X$  and  $Y$  will be  $c[m,n]$

$$c[i,j] = \begin{cases} c[i-1,j-1]+1 & \text{if } x[i] = y[j], \\ \max(c[i,j-1], c[i-1,j]) & \text{otherwise} \end{cases}$$

8

## LCS recursive solution

$$c[i, j] = \begin{cases} c[i-1, j-1] + 1 & \text{if } x[i] = y[j], \\ \max(c[i, j-1], c[i-1, j]) & \text{otherwise} \end{cases}$$

- We start with  $i = j = 0$  (empty substrings of  $x$  and  $y$ )
- Since  $X_0$  and  $Y_0$  are empty strings, their LCS is always empty (i.e.  $c[0, 0] = 0$ )
- LCS of empty string and any other string is empty, so for every  $i$  and  $j$ :  $c[0, j] = c[i, 0] = 0$

9

## LCS recursive solution

$$c[i, j] = \begin{cases} c[i-1, j-1] + 1 & \text{if } x[i] = y[j], \\ \max(c[i, j-1], c[i-1, j]) & \text{otherwise} \end{cases}$$

- When we calculate  $c[i, j]$ , we consider two cases:
- **First case:  $x[i] == y[j]$ :** one more symbol in strings  $X$  and  $Y$  **matches**, so the length of LCS  $X_i$  and  $Y_j$  **equals to the length of LCS of smaller strings  $X_{i-1}$  and  $Y_{j-1}$ , plus 1**

10

## LCS recursive solution

$$c[i, j] = \begin{cases} c[i-1, j-1] + 1 & \text{if } x[i] = y[j], \\ \max(c[i, j-1], c[i-1, j]) & \text{otherwise} \end{cases}$$

- **Second case:  $x[i] \neq y[j]$**
- As symbols **don't match**, our solution is not improved, and the length of  $\text{LCS}(X_i, Y_j)$  **is the same as before** (i.e. maximum of  $\text{LCS}(X_i, Y_{j-1})$  and  $\text{LCS}(X_{i-1}, Y_j)$ )

11

## LCS-LENGTH( $X, Y$ )

```

1  m = X.length
2  n = Y.length
3  let b[1..m, 1..n] and c[0..m, 0..n] be new tables
4  for i = 1 to m
5      c[i, 0] = 0
6  for j = 0 to n
7      c[0, j] = 0
8  for i = 1 to m
9      for j = 1 to n
10         if x_i == y_j
11             c[i, j] = c[i-1, j-1] + 1
12             b[i, j] = "↖"
13         elseif c[i-1, j] >= c[i, j-1]
14             c[i, j] = c[i-1, j]
15             b[i, j] = "↑"
16         else c[i, j] = c[i, j-1]
17             b[i, j] = "←"
18  return c and b

```

The b table returned by LCS-LENGTH enables us to quickly construct an LCS.

## LCS Length Algorithm

LCS-Length( $X, Y$ )

1.  $m = \text{length}(X)$  // get the # of symbols in  $X$
2.  $n = \text{length}(Y)$  // get the # of symbols in  $Y$
3. for  $i = 1$  to  $m$   $c[i, 0] = 0$  // special case:  $Y_0$
4. for  $j = 1$  to  $n$   $c[0, j] = 0$  // special case:  $X_0$
5. for  $i = 1$  to  $m$  // for all  $X_i$
6.     for  $j = 1$  to  $n$  // for all  $Y_j$
7.         if ( $X_i == Y_j$ )
8.              $c[i, j] = c[i-1, j-1] + 1$
9.         else  $c[i, j] = \max(c[i-1, j], c[i, j-1])$
10. return  $c$

13

## LCS Example

- $X = \text{ABCB}$
- $Y = \text{BDCAB}$

**What is the Longest Common Subsequence of  $X$  and  $Y$ ?**

$\text{LCS}(X, Y) = \text{BCB}$

$X = \text{A} \text{ B } \text{ C } \text{ B}$

$Y = \text{ B } \text{ D } \text{ C } \text{ A } \text{ B}$

14

### LCS Example (0)

$X = ABCB; m = |X| = 4$   
 $Y = BDCAB; n = |Y| = 5$   
 Allocate array  $c[4,5]$

ABCB  
BDCAB

j	0	1	2	3	4	5	
i	Yj	B	D	C	A	B	
0	Xi						
1	A						
2	B						
3	C						
4	B						

15

### LCS Example (1)

for  $i = 1$  to  $m$        $c[i,0] = 0$   
 for  $j = 1$  to  $n$        $c[0,j] = 0$

ABCB  
BDCAB

j	0	1	2	3	4	5	
i	Yj	B	D	C	A	B	
0	Xi	0	0	0	0	0	0
1	A	0					
2	B	0					
3	C	0					
4	B	0					

4/26/2018 16

### LCS Example (2)

if ( $X_i == Y_j$ )  
 $c[i,j] = c[i-1,j-1] + 1$   
 else  $c[i,j] = \max(c[i-1,j], c[i,j-1])$

ABCB  
BDCAB

j	0	1	2	3	4	5	
i	Yj	B	D	C	A	B	
0	Xi	0	0	0	0	0	0
1	A	0	0				
2	B	0					
3	C	0					
4	B	0					

4/26/2018 17

### LCS Example (3)

if ( $X_i == Y_j$ )  
 $c[i,j] = c[i-1,j-1] + 1$   
 else  $c[i,j] = \max(c[i-1,j], c[i,j-1])$

ABCB  
BDCAB

j	0	1	2	3	4	5	
i	Yj	B	D	C	A	B	
0	Xi	0	0	0	0	0	0
1	A	0	0	0	0		
2	B	0					
3	C	0					
4	B	0					

4/26/2018 18

### LCS Example (4)

if ( $X_i == Y_j$ )  
 $c[i,j] = c[i-1,j-1] + 1$   
 else  $c[i,j] = \max(c[i-1,j], c[i,j-1])$

ABCB  
BDCAB

j	0	1	2	3	4	5	
i	Yj	B	D	C	A	B	
0	Xi	0	0	0	0	0	0
1	A	0	0	0	0	1	
2	B	0					
3	C	0					
4	B	0					

4/26/2018 19

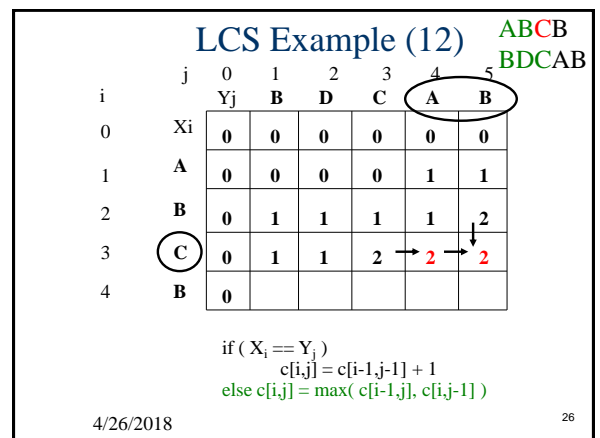
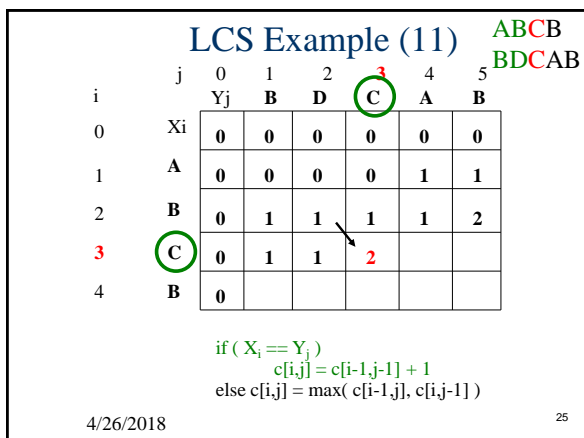
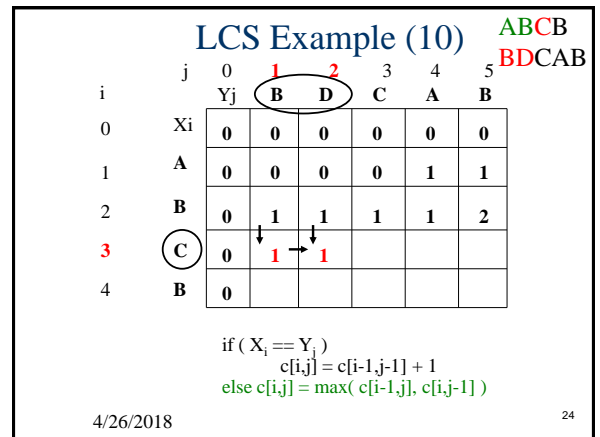
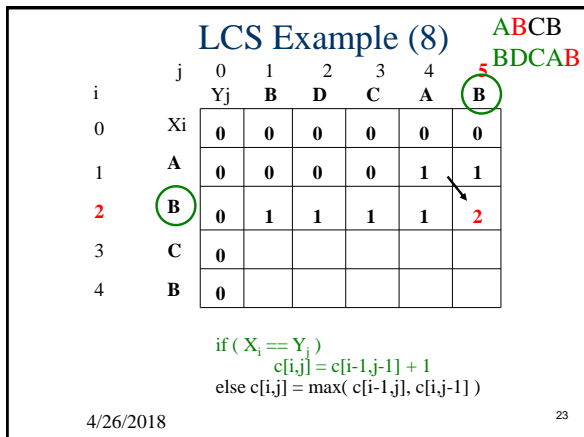
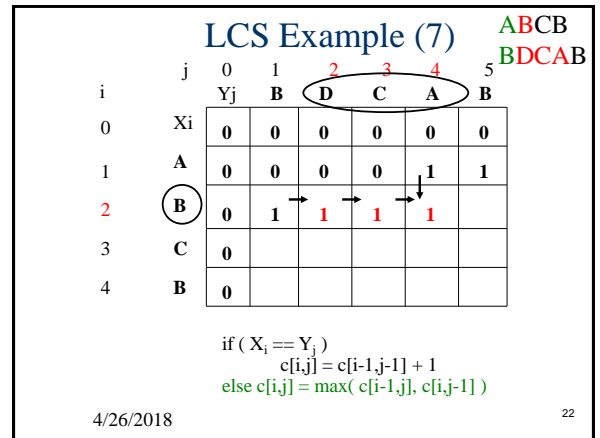
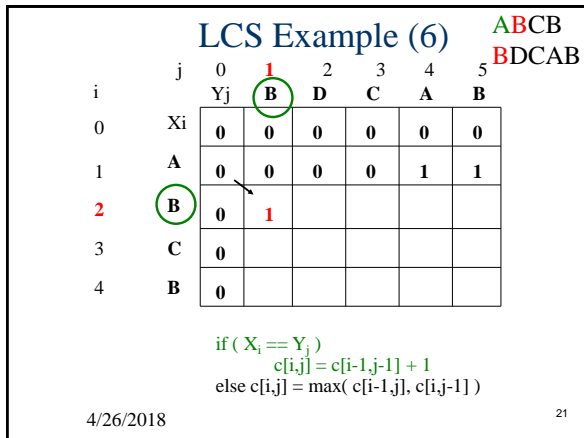
### LCS Example (5)

if ( $X_i == Y_j$ )  
 $c[i,j] = c[i-1,j-1] + 1$   
 else  $c[i,j] = \max(c[i-1,j], c[i,j-1])$

ABCB  
BDCAB

j	0	1	2	3	4	5	
i	Yj	B	D	C	A	B	
0	Xi	0	0	0	0	0	0
1	A	0	0	0	0	1	1
2	B	0					
3	C	0					
4	B	0					

4/26/2018 20



### LCS Example (13)

ABCB  
BDCAB

j	0	1	2	3	4	5
i	Yj	B	D	C	A	B
0	Xi	0	0	0	0	0
1	A	0	0	0	0	1
2	B	0	1	1	1	2
3	C	0	1	1	2	2
4	B	0	1			

if (  $X_i == Y_j$  )  
 $c[i,j] = c[i-1,j-1] + 1$   
else  $c[i,j] = \max( c[i-1,j], c[i,j-1] )$

4/26/2018 27

### LCS Example (14)

ABCB  
BDCAB

j	0	1	2	3	4	5
i	Yj	B	D	C	A	B
0	Xi	0	0	0	0	0
1	A	0	0	0	0	1
2	B	0	1	1	1	2
3	C	0	1	1	2	2
4	B	0	1	1	2	2

if (  $X_i == Y_j$  )  
 $c[i,j] = c[i-1,j-1] + 1$   
else  $c[i,j] = \max( c[i-1,j], c[i,j-1] )$

4/26/2018 28

### LCS Example (15)

ABCB  
BDCAB

j	0	1	2	3	4	5
i	Yj	B	D	C	A	B
0	Xi	0	0	0	0	0
1	A	0	0	0	0	1
2	B	0	1	1	1	2
3	C	0	1	1	2	2
4	B	0	1	1	2	3

if (  $X_i == Y_j$  )  
 $c[i,j] = c[i-1,j-1] + 1$   
else  $c[i,j] = \max( c[i-1,j], c[i,j-1] )$

4/26/2018 29

### LCS example

■  $X=\{A, B, C, B, D, A, B\}, Y=\{B, D, C, A, B, A\}$

4/26/2018 30

### LCS-LENGTH(X,Y)

```

1  m = X.length
2  n = Y.length
3  let b[1..m, 1..n] and c[0..m, 0..n] be new tables
4  for i = 1 to m
5    c[i, 0] = 0
6  for j = 0 to n
7    c[0, j] = 0
8  for i = 1 to m
9    for j = 1 to n
10     if  $x_i == y_j$ 
11        $c[i, j] = c[i-1, j-1] + 1$ 
12        $b[i, j] = "\nwarrow"$ 
13     elseif  $c[i-1, j] \geq c[i, j-1]$ 
14        $c[i, j] = c[i-1, j]$ 
15        $b[i, j] = "\uparrow"$ 
16     else  $c[i, j] = c[i, j-1]$ 
17        $b[i, j] = "\leftarrow"$ 
18  return c and b

```

The b table returned by LCS-LENGTH enables us to quickly construct an LCS.

4/26/2018 32

### How to find actual LCS

#### PRINT-LCS(b, X, i, j)

```

1  if i == 0 or j == 0
2    return
3  if b[i, j] == "\nwarrow"
4    PRINT-LCS(b, X, i-1, j-1)
5    print  $x_i$ 
6  elseif b[i, j] == "\uparrow"
7    PRINT-LCS(b, X, i-1, j)
8  else PRINT-LCS(b, X, i, j-1)

```

4/26/2018 32

## LCS Algorithm Running Time

- LCS algorithm calculates the values of each entry of the array  $c[m,n]$
- So what is the running time?

$O(m*n)$

since each  $c[i,j]$  is calculated in constant time, and there are  $m*n$  elements in the array

4/26/2018

33

## How to find actual LCS

- So far, we have just found the *length* of LCS, but not LCS itself.
  - We want to modify this algorithm to make it output Longest Common Subsequence of X and Y
- Each  $c[i,j]$  depends on  $c[i-1,j]$  and  $c[i,j-1]$  or  $c[i-1,j-1]$
- For each  $c[i,j]$  we can say how it was acquired:

2	2
2	3

For example, here  
 $c[i,j] = c[i-1,j-1] + 1 = 2 + 1 = 3$

34

## How to find actual LCS - continued

- Remember that

$$c[i,j] = \begin{cases} c[i-1,j-1] + 1 & \text{if } x[i] = y[j], \\ \max(c[i,j-1], c[i-1,j]) & \text{otherwise} \end{cases}$$

- So we can start from  $c[m,n]$  and go backwards
- Whenever  $c[i,j] = c[i-1,j-1] + 1$ , remember  $x[i]$  (because  $x[i]$  is a part of LCS)
- When  $i=0$  or  $j=0$  (i.e. we reached the beginning), output remembered letters in reverse order

35

## Finding LCS

		j	0	1	2	3	4	5
i		Yj	B	D	C	A	B	
0	Xi	0	0	0	0	0	0	
1	A	0	0	0	0	1	1	
2	B	0	1	1	1	1	2	
3	C	0	1	1	2	2	2	
4	B	0	1	1	2	2	3	

4/26/2018

36

## Finding LCS (2)

	j	0	1	2	3	4	5
i	Yj	B	D	C	A	B	
0	Xi	0	0	0	0	0	
1	A	0	0	0	0	1	1
2	B	0	1	1	1	1	2
3	C	0	1	1	2	2	2
4	B	0	1	1	2	2	3

LCS (reversed order): **B C B**

LCS (straight order): **B C B**

37

$X=\{A, B, C, B, D, A, B\}, Y=\{B, D, C, A, B, A\}$

	j	0	1	2	3	4	5	6
i	yj	B	D	C	A	B	A	
0	xj	0	0	0	0	0	0	
1	A	0	0	0	0	1	1	
2	B	0	1	1	1	2	2	
3	C	0	1	1	2	2	2	
4	B	0	1	1	2	2	3	
5	D	0	1	2	2	2	3	
6	A	0	1	2	2	3	3	
7	B	0	1	2	2	3	4	

4.

38