

## Greedy Algorithms

### Greedy algorithm

- A **greedy algorithm** is an algorithmic paradigm that follows the problem solving heuristic of making the locally optimal choice at each stage with the hope of finding a global optimum.

### Minimum Spanning Trees

- A spanning tree of a graph  $(G=V,E)$  with a weight is a tree
  - with no cycles.
  - has the same set of vertices of  $G$
  - Is a subgraph of graph  $G$ .
- A minimum spanning tree of a weighted graph  $G$  is the spanning tree of  $G$  whose edges sum to minimum weight.
- There can be more than one minimum spanning tree in a graph

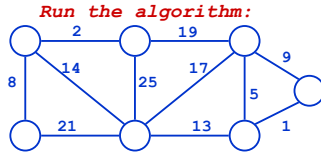
### Kruskal's Algorithm

```
Kruskal (G,w)
{
    T =  $\emptyset$ ;
    for each v  $\in$  V(G)
        MakeSet(v); //creates |V| trees
    sort E by increasing edge weight w
    for each (u,v)  $\in$  E (in sorted order)
        //checks if endpoints u, v belong to same tree
        if FindSet(u)  $\neq$  FindSet(v)
            T = T  $\cup$  {{u,v}};
            Union(u,v);
}
```

## Kruskal's Algorithm

Kruskal( $G, w$ )

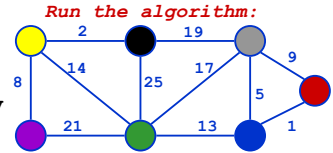
```
{
  T = ∅;
  for each v ∈ V
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  sort E by increasing edge weight w
  for each (u,v) ∈ E (in sorted order)
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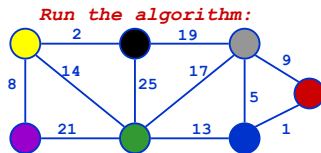
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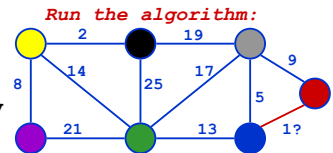
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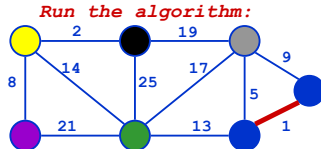
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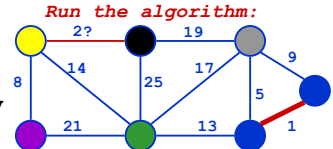
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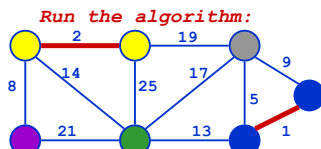
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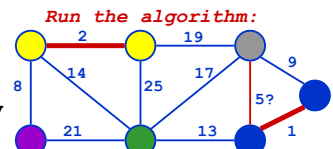
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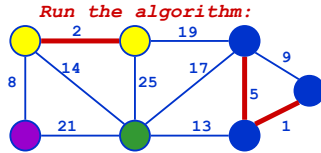
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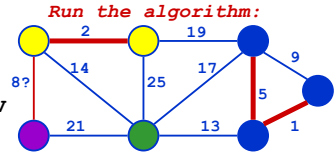
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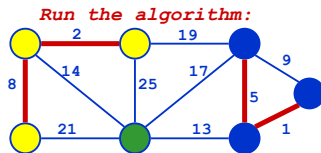
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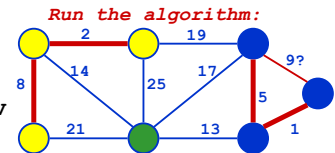
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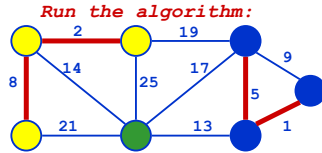
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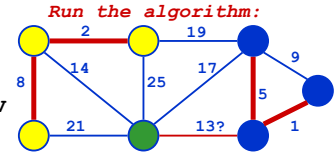
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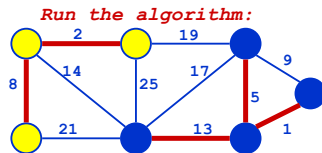
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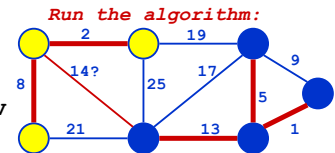
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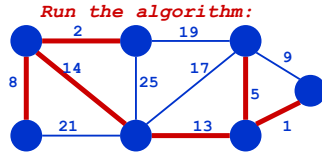
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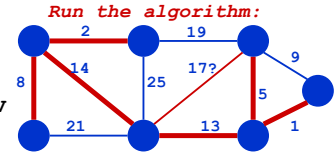
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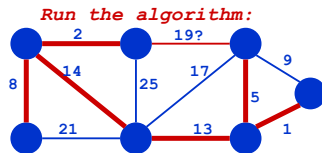
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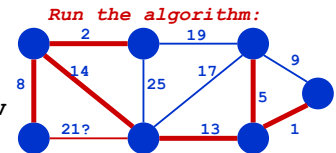
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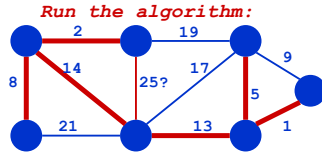
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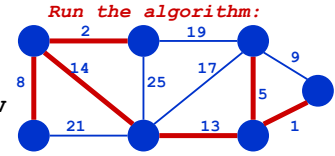
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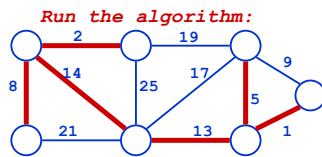
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## Prim's Algorithm

If  $G$  is connected, every vertex will appear in the minimum spanning tree.

Prim's algorithm starts from one vertex and grows the rest of the tree an edge at a time.

As a **greedy algorithm**, the cheapest edge with which can grow the tree by one vertex without creating a cycle is selected

## Prim's Algorithm

```

MST-Prim(G, w, r)
  Q ← V[G];
  for each u ∈ Q
    key[u] ← ∞;      //k:min value to connect to tree
    p[u] ← NIL      //p:parent
  key[r] ← 0;
  p[r] ← NIL;
  while (Q not empty)
    u = ExtractMin(Q);
    for each v ∈ Adj[u]
      if (v ∈ Q and w(u,v) < key[v])
        p[v] ← u;
        key[v] ← w(u,v);

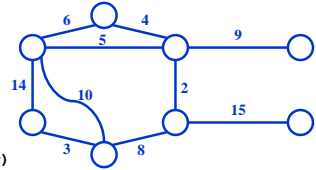
```

## Prim's Algorithm

```

MST-Prim(G, w, r)
  Q = V[G];
  for each u ∈ Q
    key[u] = ∞;
  key[r] = 0;
  p[r] = NULL;
  while (Q not empty)
    u = ExtractMin(Q);      Run on example graph
    for each v ∈ Adj[u]
      if (v ∈ Q and w(u,v) < key[v])
        p[v] = u;
        key[v] = w(u,v);

```

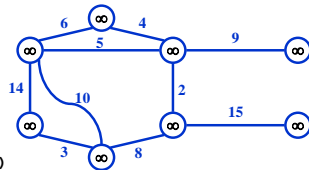


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  while (Q not empty)
    u = ExtractMin(Q);      Run on example graph
    for each v ∈ Adj[u]
      if (v ∈ Q and w(u,v) < key[v])
        p[v] = u;
        key[v] = w(u,v);

```

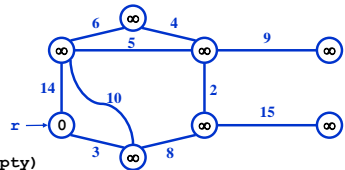


## Prim's Algorithm

```

MST-Prim(G, w, r)
  Q = V[G];
  for each u ∈ Q
    key[u] = ∞;
  key[r] = 0;
  p[r] = NULL;
  while (Q not empty)
    u = ExtractMin(Q);      Pick a start vertex r
    for each v ∈ Adj[u]
      if (v ∈ Q and w(u,v) < key[v])
        p[v] = u;
        key[v] = w(u,v);

```

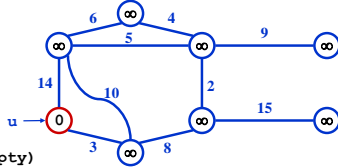




## Prim's Algorithm

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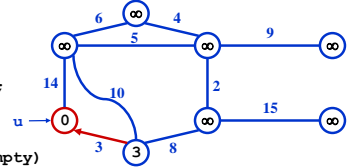
MST-Prim(G, w, r)
  Q = V[G];
  for each u ∈ Q
    key[u] = ∞;
  key[r] = 0;
  p[r] = NULL;
  while (Q not empty)
    u = ExtractMin(Q); Red vertices have been removed from Q
    for each v ∈ Adj[u]
      if (v ∈ Q and w(u,v) < key[v])
        p[v] = u;
        key[v] = w(u,v);
  
```



## Prim's Algorithm

```

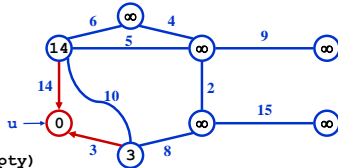
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  Q = V[G];
  for each u ∈ Q
    key[u] = ∞;
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  p[r] = NULL;
  while (Q not empty)
    u = ExtractMin(Q); Red arrows indicate parent pointers
    for each v ∈ Adj[u]
      if (v ∈ Q and w(u,v) < key[v])
        p[v] = u;
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## Prim's Algorithm

```

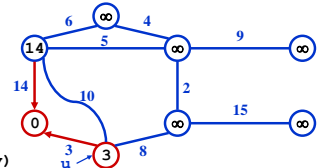
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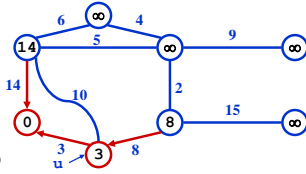


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MST-Prim( $G, w, r$ )

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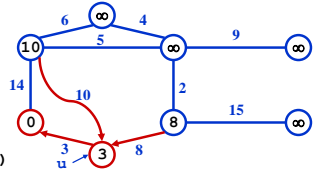


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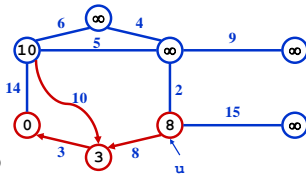


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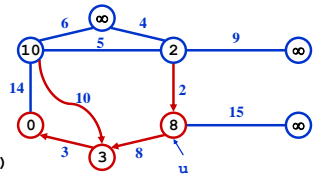


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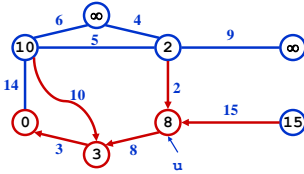
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            p[v] = u;
            key[v] = w(u,v);

```



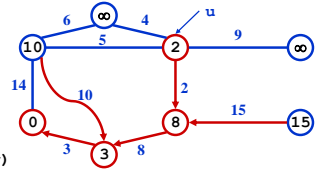
## Prim's Algorithm

MST-Prim( $G, w, r$ )

```

Q = V[G];
for each u ∈ Q
    key[u] = ∞;
key[r] = 0;
p[r] = NULL;
while (Q not empty)
    u = ExtractMin(Q);
    for each v ∈ Adj[u]
        if (v ∈ Q and w(u,v) < key[v])
            p[v] = u;
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```



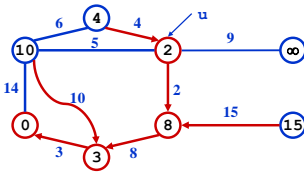
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        if (v ∈ Q and w(u,v) < key[v])
            p[v] = u;
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```



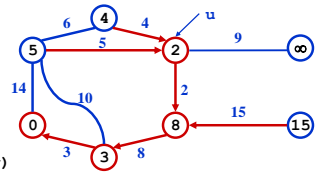
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```



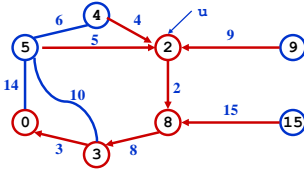
## Prim's Algorithm

MST-Prim( $G, w, r$ )

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        if (v ∈ Q and w(u,v) < key[v])
            p[v] = u;
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```



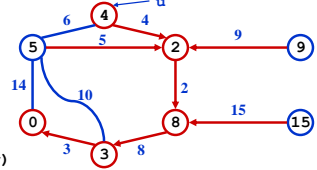
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MST-Prim( $G, w, r$ )

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        if (v ∈ Q and w(u,v) < key[v])
            p[v] = u;
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```



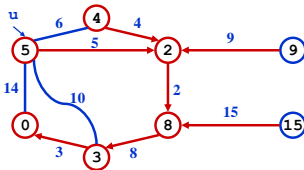
## Prim's Algorithm

MST-Prim( $G, w, r$ )

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        if (v ∈ Q and w(u,v) < key[v])
            p[v] = u;
            key[v] = w(u,v);

```



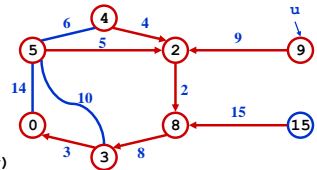
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MST-Prim( $G, w, r$ )

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        if (v ∈ Q and w(u,v) < key[v])
            p[v] = u;
            key[v] = w(u,v);

```

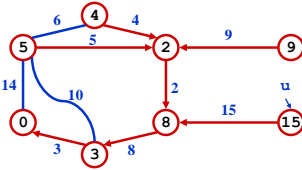


## Prim's Algorithm

```

MST-Prim(G, w, r)
  Q = V[G];
  for each u ∈ Q
    key[u] = ∞;
  key[r] = 0;
  p[r] = NULL;
  while (Q not empty)
    u = ExtractMin(Q);
    for each v ∈ Adj[u]
      if (v ∈ Q and w(u,v) < key[v])
        p[v] = u;
        key[v] = w(u,v);

```



## Kruskal's Algorithm

```

Kruskal(G, w)
{
  T = ∅;
  for each v ∈ V
    MakeSet(v);
  sort E by increasing edge weight w →
  for each (u,v) ∈ E (in sorted order)
    if FindSet(u) ≠ FindSet(v)
      T = T ∪ {(u,v)};
      Union(u, v);
}

```

$|V|$   
 $|E| \lg E$   
 $|V|$   
 $= |E| \lg E = O(E \lg V)$

## Analysis of Prim

```

MST-Prim(G, w, r)
  Q = V[G];
  for each u ∈ Q
    key[u] = ∞;
  key[r] = 0;
  p[r] = NULL;
  while (Q not empty)
    u = ExtractMin(Q);
    for each v ∈ Adj[u]
      if (v ∈ Q and w(u,v) < key[v])
        p[v] = u;
        key[v] = w(u,v);

```

$\Theta(V)$  for the initialization loop.  
 $|V|$  for the while loop.  
 $\text{Deg}(u)$  for the inner loop.  
 DecreaseKey (Called  $\Theta(E)$  times)

## Analysis of Prim

Time =  $\Theta(V) \cdot T_{\text{EXTRACT-MIN}} + \Theta(E) \cdot T_{\text{DECREASE-KEY}}$

$Q$	$T_{\text{EXTRACT-MIN}}$	$T_{\text{DECREASE-KEY}}$	Total
array	$O(V)$	$O(1)$	$O(V^2)$
binary heap	$O(\lg V)$	$O(\lg V)$	$O(E \lg V)$