

String Editing

String Editing

- Problem
 - Find The Edit Distance Between Two Strings

Edit Distance

Applications

- Approximate String Matching
- Spell checking
- Google – finding similar word variations
- DNA sequence comparison
- Pattern Recognition

String Editing

- We are given 2 strings
- $X = x_1, x_2, \dots, x_n$ and $Y = y_1, y_2, \dots, y_m$
where $x_i, 1 \leq i \leq n$ and $y_j, 1 \leq j \leq m$
- Problem: Transform $X \rightarrow Y$
- Edit operations:
 - Insert (I),
 - Delete (D),
 - Change (C): a symbol of X into another.

Cost (Transforming $X \rightarrow Y$)

- Cost = Sum of costs of the individual operation in the sequence
- $D(x_i) \rightarrow$ cost of deleting the symbol x_i from X
- $I(y_j) \rightarrow$ cost of inserting the symbol y_j into X
- $C(x_i, y_j) \rightarrow$ cost of changing the symbol x_i of X into y_j

Example

- $X = x_1, x_2, x_3, x_4, x_5 = a, a, b, a, b$
- $Y = y_1, y_2, y_3, y_4 = b, a, b, b$
- Cost of deleting and inserting be 1
- Cost of changing be 2
- $1 \leq i \leq 5$
- $1 \leq j \leq 4$
- **cost (i, j):** minimum cost of any edit sequence for transforming $X = x_1, x_2, x_3, x_4, x_5$ into y_1, y_2, y_3, y_4

cost(i, j)

1. $i=j=0$, $\text{cost}(i,j)=0$
2. $j=0$ and $i>0 \rightarrow$ transform X into Y by a sequence of deletes
 - $\text{cost}(i,0)=\text{cost}(i-1,0) + D(x_i)$
3. $i=0$ and $j>0 \rightarrow$ transform X into Y by a sequence of inserts
 - $\text{cost}(0,j)=\text{cost}(0,j-1) + I(y_j)$

cost(i, j)

4. If $i \neq 0$ and $j \neq 0$, transforming $X=x_1, x_2, x_3, x_4, x_5$ into y_1, y_2, y_3, y_4
 1. Transform x_1, x_2, \dots, x_{i-1} into y_1, y_2, \dots, y_j using a minimum-cost edit sequence and then delete x_i . The corresponding cost is $\text{cost}(i-1, j) + D(x_i)$.
 2. Transform x_1, x_2, \dots, x_{i-1} into y_1, y_2, \dots, y_{j-1} using a minimum-cost edit sequence and then change the symbol x_i to y_j . The associated cost is $\text{cost}(i-1, j-1) + C(x_i, y_j)$.
 3. Transform x_1, x_2, \dots, x_i into y_1, y_2, \dots, y_{j-1} using a minimum-cost edit sequence and then insert y_j . This corresponds to a cost of $\text{cost}(i, j-1) + I(y_j)$.

cost(i, j)

$$\text{cost}(i, j) = \begin{cases} 0 & i = j = 0 \\ \text{cost}(i-1, 0) + D(x_i) & j = 0, i > 0 \\ \text{cost}(0, j-1) + I(y_j) & i = 0, j > 0 \\ \text{cost}'(i, j) & i > 0, j > 0 \end{cases}$$

$$\text{where } \text{cost}'(i, j) = \min \left\{ \begin{array}{l} \text{cost}(i-1, j) + D(x_i), \\ \text{cost}(i-1, j-1) + C(x_i, y_j), \\ \text{cost}(i, j-1) + I(y_j) \end{array} \right\}$$

Example

- $X=x_1, x_2, x_3, x_4, x_5=a, a, b, a, b$
- $Y=y_1, y_2, y_3, y_4=b, a, b, b$
- Cost of deleting and inserting be 1
- Cost of changing be 2
- $1 \leq i \leq 5$
- $1 \leq j \leq 4$

Cost table

$j \rightarrow$	0	1	2	3	4
$i \downarrow$	0	1	2	3	4
1					
2					
3					
4					
5					

cost(1,1), cost(1,2)

$$\begin{aligned} \text{cost}(1,1) &= \min \{ \text{cost}(0,1) + D(x_1), \text{cost}(0,0) + C(x_1, y_1), \text{cost}(1,0) + I(y_1) \} \\ &= \min \{ 2, 2, 2 \} = 2 \end{aligned}$$

$$\begin{aligned} \text{cost}(1,2) &= \min \{ \text{cost}(0,2) + D(x_1), \text{cost}(0,1) + C(x_1, y_2), \text{cost}(1,1) + I(y_2) \} \\ &= \min \{ 3, 1, 3 \} = 1 \end{aligned}$$

$j \rightarrow$	0	1	2	3	4
$i \downarrow$	0	1	2	3	4
1		1			
2		2			
3		3			
4		4			
5		5			

Cost table

$j \rightarrow$	0	1	2	3	4	
$i \downarrow$	0	0	1	2	3	4
1	1	2	1	2	3	
2	2	3	2	3	4	
3	3	2	3	2	3	
4	4	3	2	3	4	
5	5	4	3	2	3	

cost(5,4)

- $X=x_1, x_2, x_3, x_4, x_5=a, a, b, a, b$
- $Y=y_1, y_2, y_3, y_4 = b, a, b, b$
- $\text{cost}(5,4)=3$.
- Possible minimum cost edit sequence
 1. delete x_1 , delete x_2 , and insert y_4
 2. Change $x_1 \rightarrow y_1$ and delete x_4

Edit Distance

- How many edits are needed to exactly match the Target with the Pattern
- Target: **TCGACGT** CA
- Pattern: T GACGTGC
- Three:
 - By Deleting C and A from the target, and by Deleting G from the Pattern

Edit Distance – Dynamic Programming

