



# Transactions and Recovery: Concurrency Control

# Example

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| T1:

- Interleaving:
- T1:  $A=A+100$ ,  $B=B-100$

| T2:

- Interleaving:
- T2:  $A=1.06*A$ ,  $B=1.06*B$

The DBMS's view:

T1:  $R(A)$ ,  $W(A)$ ,  $R(B)$ ,  $W(B)$

T2:  $R(A)$ ,  $W(A)$ ,  $R(B)$ ,  $W(B)$

# Scheduling Transactions

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## | Serial schedule

- Schedule that does not interleave the actions of different transactions.

## | Equivalent Schedule

- For any database state, the effect of executing the first schedule is identical to the effect of executing the second schedule.

## | Serializable Schedule

- A schedule that is equivalent to some serial execution of the transactions.

# Conflict Serializable Schedules

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| Two schedules are **conflict equivalent** if:

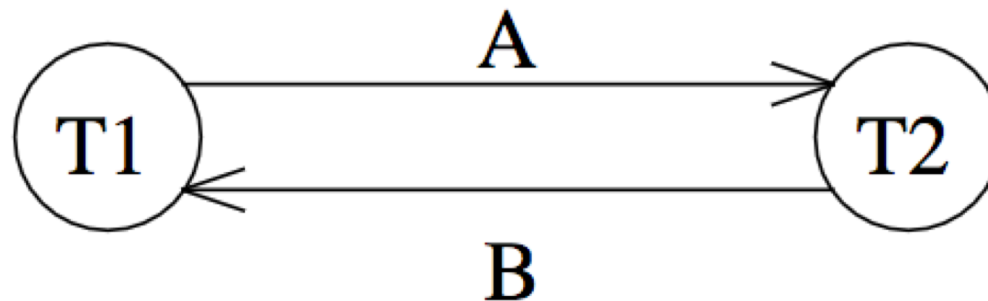
- Involve the same actions of the same transactions
- Every pair of conflicting actions is ordered the same way

| Schedule S is **conflict serializable** if S is conflict equivalent to some serial schedule

## Example: Not Conflict Serializable

T1: R(A), W(A), R(B), W(B)

T2: R(A), W(A), R(B), W(B)



Dependency Graph