Principles of Distributed and Parallel Database Systems: Part 2 Distributed Query Processing



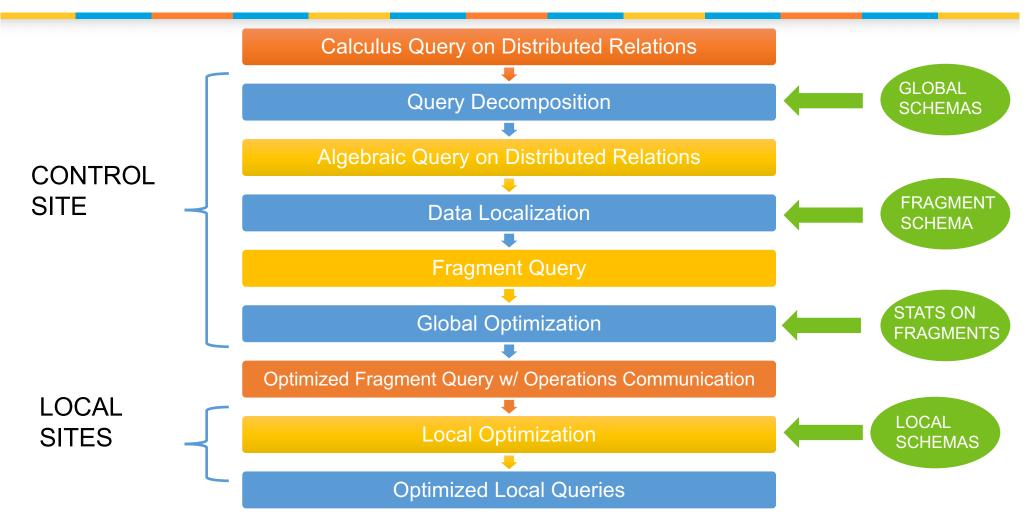
Objectives



Objective

Illustrate how concurrency control is achieved in distributed databases

Distributed Query Processing



Query Processing

Decomposition

Localization

Optimization

Decomposition

Same as in centralized system

Normalization

Eliminating redundancy

Algebraic rewriting

Normalization

Convert from general language to a "standard" form (e.g., Relational Algebra)

Example

```
Select A,C
From R,S
Where (R.B=1 and S.D=2) or (R.C>3 and S.D.=2)
    \sigma (R.B=1 v R.C>3) \wedge S.D.=2)
                                           Conjunctive
                                             normal
                                              form
```

Also: Detect invalid expressions

E.g.: Select * from R where R.A = 3

☐ R does not have "A" attribute

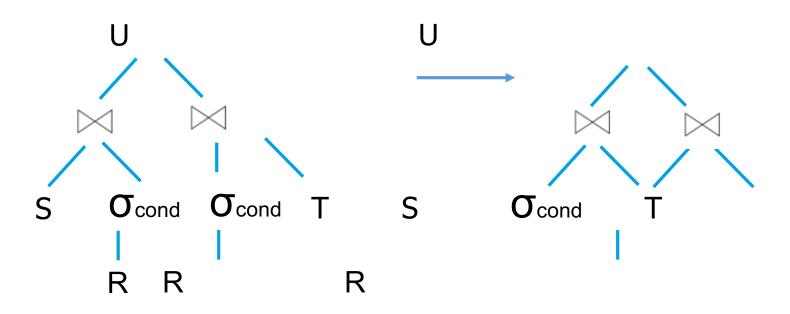
Eliminate redundancy

E.g. in conditions:

$$(S.A=1) \land (S.A>5) \Rightarrow False$$

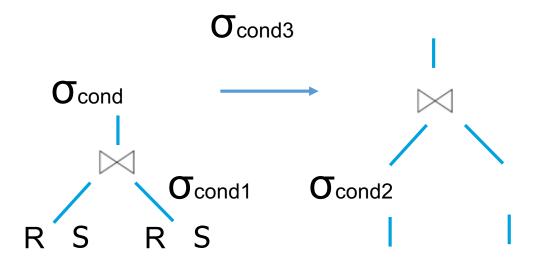
$$(S.A<10) \land (S.A<5) \Rightarrow S.A<5$$

E.g.: Common sub-expressions



Algebraic rewriting

E.g.: Push conditions down



Algebraic rewriting

After decomposition:

 One or more algebraic query trees on relations

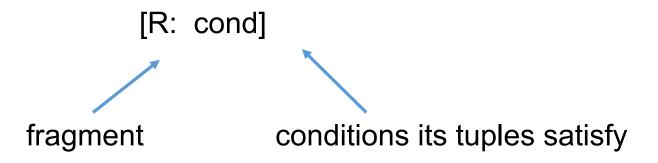
Localization:

Replace relations by corresponding fragments

Localization steps

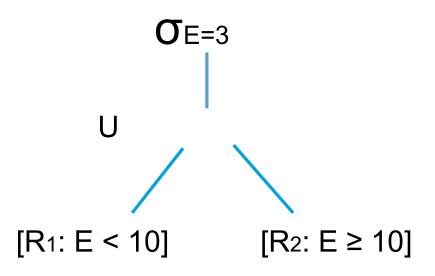
- Start with query
- Replace relations by fragments
- Push ∪: up ()
 - $-\pi,\sigma$: down
- Simplify eliminate unnecessary operations

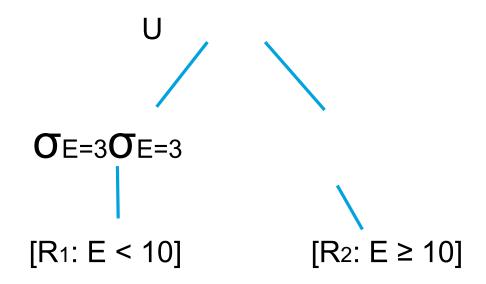
Notation for fragment



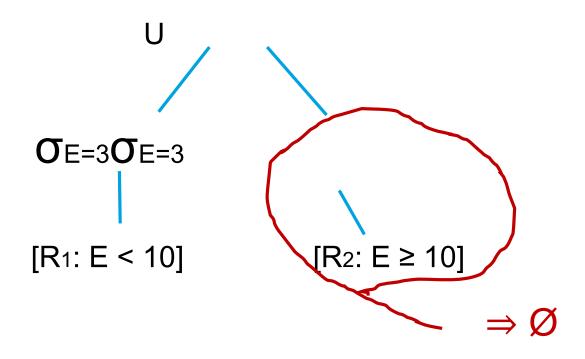
R

What if we have two Fragments [R₁: E < 10] [R₂: $E \ge 10$]





Anything Wrong here?



```
σE=3 | [R<sub>1</sub>: E < 10]
```