Team notebook

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1 Algorithms

1.1 sliding window

```
* Given an array ARR and an integer K, the problem boils down to
     computing for each index i: min(ARR[i], ARR[i-1], ..., ARR[i-K+1]).
* if mx == true, returns the maximun.
* http://people.cs.uct.ac.za/~ksmith/articles/sliding_window_minimum.html
* */
vector<int> sliding_window_minmax(vector<int> & ARR, int K, bool mx) {
 deque< pair<int, int> > window;
 vector<int> ans;
 for (int i = 0; i < ARR.size(); i++) {</pre>
     while (!window.empty() && window.back().first <= ARR[i])</pre>
       window.pop_back();
     while (!window.empty() && window.back().first >= ARR[i])
       window.pop_back();
   }
   window.push_back(make_pair(ARR[i], i));
   while(window.front().second <= i - K)</pre>
     window.pop_front();
   ans.push_back(window.front().first);
 }
 return ans;
```

2 Data structures

2.1 STL Treap

```
#include <ext/rope> //header with rope
using namespace std;
using namespace __gnu_cxx; //namespace with rope and some additional stuff
int main()
   ios_base::sync_with_stdio(false);
   rope <int> v; //use as usual STL container
   int n, m;
   cin >> n >> m;
   for(int i = 1; i <= n; ++i)</pre>
       v.push_back(i); //initialization
   int 1, r;
   for(int i = 0; i < m; ++i)</pre>
       cin >> 1 >> r;
       --1, --r;
       rope \langle int \rangle cur = v.substr(1, r - 1 + 1);
       v.erase(1, r - 1 + 1);
       v.insert(v.mutable_begin(), cur);
   for(rope <int>::iterator it = v.mutable_begin(); it !=
        v.mutable_end(); ++it)
       cout << *it << " ";
   return 0;
}
```

2.2 STL order statistics tree II

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>

using namespace std;
using namespace __gnu_pbds;

typedef tree<int,null_type,less<int>,rb_tree_tag,
tree_order_statistics_node_update> order_set;
```

```
order_set X;
int get(int y) {
  int l=0,r=1e9+1;
  while(l<r) {</pre>
    int m=l+((r-l)>>1);
    if(m-X.order_of_key(m+1)<y)</pre>
    else
      r=m;
 }
 return 1;
}
main(){
  ios::sync_with_stdio(0);
  cin.tie(0);
  int n,m;
  cin>>n>>m;
  for(int i=0;i<m;i++) {</pre>
    char a;
    int b;
    cin>>a>>b;
    if(a=='L')
      cout<<get(b)<<endl;</pre>
      X.insert(get(b));
  }
}
/***
Input
20 7
L 5
D 5
L 4
L 5
D 5
L 4
L 5
Output
5
4
```

```
6
4
7
***/
```

2.3 STL order statistics tree

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <bits/stdc++.h>
using namespace __gnu_pbds;
using namespace std;
typedef
tree<
  pair<int,int>,
 null_type,
 less<pair<int,int>>,
 rb_tree_tag,
  tree_order_statistics_node_update>
ordered_set;
main()
{
   ios::sync_with_stdio(0);
    cin.tie(0);
    int n;
    int sz=0;
   cin>>n;
    vector<int> ans(n,0);
   ordered_set t;
   int x,y;
   for(int i=0;i<n;i++)</pre>
    {
       cin>>x>>y;
       ans[t.order_of_key({x,++sz})]++;
       t.insert({x,sz});
   for(int i=0;i<n;i++)</pre>
       cout<<ans[i]<<'\n';</pre>
```

```
/***
Input
5
1 1
5 1
7 1
3 3
5 5
Output
1
2
1
1
0
***/
```

2.4 hash table

```
/**
 * Micro hash table, can be used as a set.
 * Very efficient vs std::set
 * */

const int MN = 1001;
struct ht {
  int _s[(MN + 10) >> 5];
  int len;
  void set(int id) {
    len++;
    _s[id >> 5] |= (1LL << (id & 31));
  }
  bool is_set(int id) {
    return _s[id >> 5] & (1LL << (id & 31));
  }
};</pre>
```

2.5 heavy light decomposition

```
// Heavy-Light Decomposition
struct TreeDecomposition {
 vector<int> g[MAXN], c[MAXN];
 int s[MAXN]; // subtree size
 int p[MAXN]; // parent id
 int r[MAXN]; // chain root id
 int t[MAXN]; // index used in segtree/bit/...
 int d[MAXN]; // depht
 int ts;
 void dfs(int v, int f) {
   p[v] = f;
   s[v] = 1;
   if (f != -1) d[v] = d[f] + 1;
   else d[v] = 0;
   for (int i = 0; i < g[v].size(); ++i) {</pre>
     int w = g[v][i];
     if (w != f) {
       dfs(w, v);
       s[v] += s[w];
 void hld(int v, int f, int k) {
   t[v] = ts++;
   c[k].push_back(v);
   r[v] = k;
   int x = 0, y = -1;
   for (int i = 0; i < g[v].size(); ++i) {</pre>
     int w = g[v][i];
     if (w != f) {
       if (s[w] > x) {
        x = s[w];
         y = w;
     }
   if (y != -1) {
     hld(y, v, k);
   for (int i = 0; i < g[v].size(); ++i) {</pre>
```

```
int w = g[v][i];
     if (w != f && w != y) {
       hld(w, v, w);
     }
   }
 }
  void init(int n) {
   for (int i = 0; i < n; ++i) {</pre>
     g[i].clear();
 }
 void add(int a, int b) {
   g[a].push_back(b);
   g[b].push_back(a);
 void build() {
   ts = 0;
   dfs(0, -1);
   hld(0, 0, 0);
 }
};
```

2.6 persistent array

```
struct node {
  node *1, *r;
  int val;

  node (int x) : 1(NULL), r(NULL), val(x) {}
  node () : 1(NULL), r(NULL), val(-1) {}
};

typedef node* pnode;

pnode update(pnode cur, int 1, int r, int at, int what) {
  pnode ans = new node();

  if (cur != NULL) {
    *ans = *cur;
}
```

```
if (1 == r) {
    ans-> val = what;
    return ans;
}
int m = (1 + r) >> 1;
if (at <= m) ans-> 1 = update(ans-> 1, 1, m, at, what);
else ans-> r = update(ans-> r, m + 1, r, at, what);
return ans;
}

int get(pnode cur, int 1, int r, int at) {
    if (cur == NULL) return 0;
    if (1 == r) return cur-> val;
    int m = (1 + r) >> 1;
    if (at <= m) return get(cur-> 1, 1, m, at);
else         return get(cur-> r, m + 1, r, at);
}
```

2.7 persistent seg tree

```
/**
 * Important:
 * When using lazy propagation remembert to create new
 * versions for each push_down operation!!!
 * */
struct node {
 node *1, *r;
 long long acc;
 int flip;
 node (int x) : 1(NULL), r(NULL), acc(x), flip(0) {}
 node () : 1(NULL), r(NULL), acc(0), flip(0) {}
}:
typedef node* pnode;
pnode create(int 1, int r) {
 if (1 == r) return new node();
 pnode cur = new node();
 int m = (1 + r) >> 1;
 cur \rightarrow 1 = create(1, m);
  cur-> r = create(m + 1, r);
```

```
return cur;
}
pnode copy_node(pnode cur) {
 pnode ans = new node();
 *ans = *cur;
 return ans;
}
void push_down(pnode cur, int 1, int r) {
 assert(cur):
 if (cur-> flip) {
   int len = r - l + 1;
   cur-> acc = len - cur-> acc;
   if (cur-> 1) {
     cur-> 1 = copy_node(cur-> 1);
     cur-> 1 -> flip ^= 1;
   if (cur-> r) {
     cur-> r = copy_node(cur-> r);
     cur-> r -> flip ^= 1;
   cur-> flip = 0;
 }
}
int get_val(pnode cur) {
 assert(cur);
 assert((cur-> flip) == 0);
 if (cur) return cur-> acc;
 return 0;
}
pnode update(pnode cur, int 1, int r, int at, int what) {
 pnode ans = copy_node(cur);
 if (1 == r) {
   assert(1 == at);
   ans-> acc = what;
   ans-> flip = 0;
   return ans;
 int m = (1 + r) >> 1;
 push_down(ans, 1, r);
 if (at <= m) ans-> 1 = update(ans-> 1, 1, m, at, what);
 else ans-> r = update(ans-> r, m + 1, r, at, what);
```

```
push_down(ans-> 1, 1, m);
  push_down(ans-> r, m + 1, r);
  ans-> acc = get_val(ans-> 1) + get_val(ans-> r);
 return ans;
}
pnode flip(pnode cur, int 1, int r, int a, int b) {
 pnode ans = new node();
 if (cur != NULL) {
   *ans = *cur;
 if (1 > b || r < a)
   return ans;
 if (1 >= a && r <= b) {</pre>
   ans-> flip ^= 1;
   push_down(ans, 1, r);
   return ans;
 int m = (1 + r) >> 1;
  ans-> 1 = flip(ans-> 1, 1, m, a, b);
 ans-> r = flip(ans-> r, m + 1, r, a, b);
 push_down(ans-> 1, 1, m);
 push_down(ans-> r, m + 1, r);
 ans-> acc = get_val(ans-> 1) + get_val(ans-> r);
 return ans;
}
long long get_all(pnode cur, int 1, int r) {
  assert(cur);
 push_down(cur, 1, r);
 return cur-> acc;
void traverse(pnode cur, int 1, int r) {
 if (!cur) return;
  cout << 1 << " - " << r << " : " << (cur-> acc) << " " << (cur-> flip)
      << endl:
 traverse(cur-> 1, 1, (1 + r) >> 1);
 traverse(cur-> 1, 1 + ((1 + r) >> 1), r);
```

2.8 segment tree

```
/**
 * Taken from: http://codeforces.com/blog/entry/18051
const int N = 1e5; // limit for array size
int n; // array size
int t[2 * N];
void build() { // build the tree
 for (int i = n - 1; i > 0; --i) t[i] = t[i << 1] + t[i << 1|1];
}
// Single modification, range query.
void modify(int p, int value) { // set value at position p
 for (t[p += n] = value; p > 1; p >>= 1) t[p>>1] = t[p] + t[p^1];
}
int query(int 1, int r) { // sum on interval [1, r)
 int res = 0;
 for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
   if (1\&1) res += t[1++];
   if (r&1) res += t[--r];
 }
 return res;
// Range modification, single query.
void modify(int 1, int r, int value) {
 for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
   if (l&1) t[l++] += value;
   if (r&1) t[--r] += value;
 }
}
int query(int p) {
 int res = 0;
 for (p += n; p > 0; p >>= 1) res += t[p];
 return res;
}
/**
* If at some point after modifications we need to inspect all the
```

```
* elements in the array, we can push all the modifications to the
 * leaves using the following code. After that we can just traverse
 * elements starting with index n. This way we reduce the complexity
 * from O(n \log(n)) to O(n) similarly to using build instead of n
     modifications.
 * */
void push() {
 for (int i = 1; i < n; ++i) {</pre>
   t[i<<1] += t[i];
   t[i << 1|1] += t[i];
   t[i] = 0;
 }
}
// Non commutative combiner functions.
void modify(int p, const S& value) {
 for (t[p += n] = value; p >>= 1;) t[p] = combine(t[p<<1], t[p<<1|1]);
S query(int 1, int r) {
 S resl, resr;
 for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
   if (1\&1) resl = combine(resl, t[1++]);
   if (r&1) resr = combine(t[--r], resr);
 return combine(resl, resr);
// To be continued ...
```

2.9 sparse table

```
// RMQ.
const int MN = 100000 + 10; // Max number of elements
const int ML = 18; // ceil(log2(MN));
struct st {
  int data[MN];
  int M[MN][ML];
  int n:
```

```
void read(int _n) {
   n = _n;
   for (int i = 0; i < n; ++i)
     cin >> data[i];
 }
 void build() {
   for (int i = 0; i < n; ++i)
     M[i][0] = data[i];
   for (int j = 1, p = 2, q = 1; p \le n; ++j, p \le 1, q \le 1)
     for (int i = 0; i + p - 1 < n; ++i)
       M[i][j] = max(M[i][j-1], M[i+q][j-1]);
 }
 int query(int b, int e) {
   int k = log2(e - b + 1);
   return max(M[b][k], M[e + 1 - (1<<k)][k]);</pre>
 }
};
```

2.10 splay tree

```
using namespace std;
#include<bits/stdc++.h>
#define D(x) cout<<x<<endl;</pre>
typedef int T;
struct node{
 node *left, *right, *parent;
 T key;
 node (T k) : key(k), left(0), right(0), parent(0) {}
};
struct splay_tree{
 node *root;
  void right_rot(node *x) {
   node *p = x->parent;
   if (x->parent = p->parent) {
     if (x->parent->left == p) x->parent->left = x;
     if (x->parent->right == p) x->parent->right = x;
```

```
if (p->left = x->right) p->left->parent = p;
 x->right = p;
 p->parent = x;
void left_rot(node *x) {
 node *p = x->parent;
 if (x->parent = p->parent) {
   if (x->parent->left == p) x->parent->left = x;
   if (x->parent->right == p) x->parent->right = x;
 if (p->right = x->left) p->right->parent = p;
 x \rightarrow left = p;
 p->parent = x;
void splay(node *x, node *fa = 0) {
 while( x->parent != fa and x->parent != 0) {
   node *p = x->parent;
   if (p->parent == fa)
     if (p->right == x)
      left_rot(x);
     else
       right_rot(x);
   else {
     node *gp = p->parent; //grand parent
     if (gp->left == p)
      if (p->left == x)
         right_rot(x),right_rot(x);
         left_rot(x),right_rot(x);
     else
       if (p->left == x)
         right_rot(x), left_rot(x);
       else
         left_rot(x), left_rot(x);
   }
 if (fa == 0) root = x;
void insert(T key) {
 node *cur = root;
```

```
node *pcur = 0;
   while (cur) {
     pcur = cur;
     if (key > cur->key) cur = cur->right;
     else cur = cur->left;
   cur = new node(key);
   cur->parent = pcur;
   if (!pcur) root = cur;
   else if (key > pcur->key ) pcur->right = cur;
   else pcur->left = cur;
   splay(cur);
 }
 node *find(T key) {
   node *cur = root:
   while (cur) {
     if (key > cur->key) cur = cur->right;
     else if(key < cur->key) cur = cur->left;
     else return cur;
   return 0;
 }
 splay_tree(){ root = 0;};
};
```

2.11 trie

```
const int MN = 26; // size of alphabet
const int MS = 100010; // Number of states.

struct trie{
    struct node{
        int c;
        int a[MN];
    };

    node tree[MS];
    int nodes;

void clear(){
        tree[nodes].c = 0;
```

```
memset(tree[nodes].a, -1, sizeof tree[nodes].a);
   nodes++;
  }
  void init(){
   nodes = 0;
   clear();
  int add(const string &s, bool query = 0){
   int cur_node = 0;
   for(int i = 0; i < s.size(); ++i){</pre>
     int id = gid(s[i]);
     if(tree[cur_node].a[id] == -1){
       if(query) return 0;
       tree[cur_node].a[id] = nodes;
       clear();
     cur_node = tree[cur_node].a[id];
   if(!query) tree[cur_node].c++;
   return tree[cur_node].c;
};
```

3 Geometry

3.1 closest pair problem

```
struct point {
  double x, y;
  int id;
  point() {}
  point (double a, double b) : x(a), y(b) {}
};

double dist(const point &o, const point &p) {
  double a = p.x - o.x, b = p.y - o.y;
  return sqrt(a * a + b * b);
}
```

```
double cp(vector<point> &p, vector<point> &x, vector<point> &y) {
 if (p.size() < 4) {</pre>
   double best = 1e100;
   for (int i = 0; i < p.size(); ++i)</pre>
     for (int j = i + 1; j < p.size(); ++j)</pre>
       best = min(best, dist(p[i], p[j]));
   return best;
 int ls = (p.size() + 1) >> 1;
 double 1 = (p[ls - 1].x + p[ls].x) * 0.5;
 vector<point> xl(ls), xr(p.size() - ls);
 unordered_set<int> left;
 for (int i = 0; i < ls; ++i) {</pre>
   x1[i] = x[i];
   left.insert(x[i].id);
 for (int i = ls; i < p.size(); ++i) {</pre>
   xr[i - ls] = x[i];
 vector<point> yl, yr;
  vector<point> pl, pr;
 yl.reserve(ls); yr.reserve(p.size() - ls);
 pl.reserve(ls); pr.reserve(p.size() - ls);
 for (int i = 0; i < p.size(); ++i) {</pre>
   if (left.count(y[i].id))
     yl.push_back(y[i]);
   else
     yr.push_back(y[i]);
   if (left.count(p[i].id))
     pl.push_back(p[i]);
   else
     pr.push_back(p[i]);
 }
 double dl = cp(pl, xl, yl);
 double dr = cp(pr, xr, yr);
 double d = min(dl, dr);
  vector<point> yp; yp.reserve(p.size());
 for (int i = 0; i < p.size(); ++i) {</pre>
   if (fabs(y[i].x - 1) < d)
     yp.push_back(y[i]);
 }
```

```
for (int i = 0; i < yp.size(); ++i) {
   for (int j = i + 1; j < yp.size() && j < i + 7; ++j) {
      d = min(d, dist(yp[i], yp[j]));
   }
}
return d;
}

double closest_pair(vector<point> &p) {
   vector<point> x(p.begin(), p.end());
   sort(x.begin(), x.end(), [](const point &a, const point &b) {
      return a.x < b.x;
   });
   vector<point> y(p.begin(), p.end());
   sort(y.begin(), y.end(), [](const point &a, const point &b) {
      return a.y < b.y;
   });
   return cp(p, x, y);
}</pre>
```

3.2 squares

```
typedef long double ld;
const ld eps = 1e-12;
int cmp(ld x, ld y = 0, ld tol = eps) {
   return ( x \le y + tol) ? (x + tol < y) ? -1 : 0 : 1;
}
struct point{
 ld x, y;
 point(ld a, ld b) : x(a), y(b) {}
 point() {}
};
struct square{
 ld x1, x2, y1, y2,
    a, b, c;
 point edges[4];
  square(ld _a, ld _b, ld _c) {
   a = _a, b = _b, c = _c;
   x1 = a - c * 0.5;
```

```
x2 = a + c * 0.5:
   v1 = b - c * 0.5;
   y2 = b + c * 0.5;
   edges[0] = point(x1, y1);
   edges[1] = point(x2, y1);
   edges[2] = point(x2, y2);
   edges[3] = point(x1, y2);
 }
};
ld min_dist(point &a, point &b) {
 1d x = a.x - b.x
    y = a.y - b.y;
 return sqrt(x * x + y * y);
}
bool point_in_box(square s1, point p) {
 if (cmp(s1.x1, p.x) != 1 && cmp(s1.x2, p.x) != -1 &&
     cmp(s1.y1, p.y) != 1 && cmp(s1.y2, p.y) != -1)
   return true;
 return false;
}
bool inside(square &s1, square &s2) {
 for (int i = 0; i < 4; ++i)
   if (point_in_box(s2, s1.edges[i]))
     return true:
 return false;
}
bool inside_vert(square &s1, square &s2) {
 if ((cmp(s1.y1, s2.y1) != -1 && cmp(s1.y1, s2.y2) != 1) ||
     (cmp(s1.y2, s2.y1) != -1 \&\& cmp(s1.y2, s2.y2) != 1))
   return true;
return false;
}
bool inside_hori(square &s1, square &s2) {
 if ((cmp(s1.x1, s2.x1) != -1 && cmp(s1.x1, s2.x2) != 1) ||
     (cmp(s1.x2, s2.x1) != -1 \&\& cmp(s1.x2, s2.x2) != 1))
   return true;
return false;
}
```

```
ld min_dist(square &s1, square &s2) {
 if (inside(s1, s2) || inside(s2, s1))
   return 0;
 ld ans = 1e100;
 for (int i = 0; i < 4; ++i)
   for (int j = 0; j < 4; ++j)
     ans = min(ans, min_dist(s1.edges[i], s2.edges[j]));
 if (inside_hori(s1, s2) || inside_hori(s2, s1)) {
   if (cmp(s1.v1, s2.v2) != -1)
     ans = min(ans, s1.y1 - s2.y2);
   if (cmp(s2.y1, s1.y2) != -1)
     ans = min(ans, s2.y1 - s1.y2);
 if (inside_vert(s1, s2) || inside_vert(s2, s1)) {
   if (cmp(s1.x1, s2.x2) != -1)
     ans = min(ans, s1.x1 - s2.x2);
   else
   if (cmp(s2.x1, s1.x2) != -1)
     ans = min(ans, s2.x1 - s1.x2);
 return ans:
```

3.3 triangles

Let a, b, c be length of the three sides of a triangle.

$$p = (a + b + c) * 0.5$$

The inradius is defined by:

$$iR = \sqrt{\frac{(p-a)(p-b)(p-c)}{p}}$$

The radius of its circumcircle is given by the formula:

$$cR = \frac{abc}{\sqrt{(a+b+c)(a+b-c)(a+c-b)(b+c-a)}}$$

4 Graphs

4.1 bridges

```
struct edge{
 int to, id;
 edge(int a, int b) : to(a), id(b) {}
};
struct graph {
 vector<vector<edge> > g;
 vector<int> vi, low, d, pi, is_b;
 int ticks, edges;
  graph(int n, int m) {
   g.assign(n, vector<edge>());
   is_b.assign(m, 0);
   vi.resize(n);
   low.resize(n);
   d.resize(n);
   pi.resize(n);
   edges = 0;
 }
  void add_edge(int u, int v) {
   g[u].push_back(edge(v, edges));
   g[v].push_back(edge(u, edges));
   edges++;
 }
  void dfs(int u) {
   vi[u] = true;
   d[u] = low[u] = ticks++;
   for (int i = 0; i < g[u].size(); ++i) {</pre>
     int v = g[u][i].to;
     if (v == pi[u]) continue;
     if (!vi[v]) {
       pi[v] = u;
       dfs(v);
       if (d[u] < low[v])</pre>
         is_b[g[u][i].id] = true;
```

```
low[u] = min(low[u], low[v]);
     } else {
       low[u] = min(low[u], d[v]);
   }
  // Multiple edges from a to b are not allowed.
  // (they could be detected as a bridge).
 // If you need to handle this, just count
  // how many edges there are from a to b.
  void comp_bridges() {
   fill(pi.begin(), pi.end(), -1);
   fill(vi.begin(), vi.end(), 0);
   fill(low.begin(), low.end(), 0);
   fill(d.begin(), d.end(), 0);
   ticks = 0;
   for (int i = 0; i < g.size(); ++i)</pre>
     if (!vi[i]) dfs(i);
 }
};
```

4.2 directed mst

```
const int inf = 1000000 + 10;

struct edge {
   int u, v, w;
   edge() {}
   edge(int a,int b,int c) : u(a), v(b), w(c) {}
};

/**

* Computes the minimum spanning tree for a directed graph

* - edges : Graph description in the form of list of edges.

* each edge is: From node u to node v with cost w

* - root : Id of the node to start the DMST.

* - n : Number of nodes in the graph.

* */

int dmst(vector<edge> &edges, int root, int n) {
   int ans = 0;
   int cur_nodes = n;
```

```
while (true) {
 vector<int> lo(cur_nodes, inf), pi(cur_nodes, inf);
 for (int i = 0; i < edges.size(); ++i) {</pre>
   int u = edges[i].u, v = edges[i].v, w = edges[i].w;
   if (w < lo[v] and u != v) {
     lo[v] = w:
     pi[v] = u;
   }
 }
 lo[root] = 0:
 for (int i = 0; i < lo.size(); ++i) {</pre>
   if (i == root) continue;
   if (lo[i] == inf) return -1;
 }
 int cur id = 0:
 vector<int> id(cur_nodes, -1), mark(cur_nodes, -1);
 for (int i = 0; i < cur_nodes; ++i) {</pre>
   ans += lo[i]:
   int u = i;
   while (u != root and id[u] < 0 and mark[u] != i) {</pre>
     mark[u] = i;
     u = pi[u];
   if (u != root and id[u] < 0) { // Cycle</pre>
      for (int v = pi[u]; v != u; v = pi[v])
        id[v] = cur_id;
      id[u] = cur_id++;
   }
 }
 if (cur id == 0)
   break:
 for (int i = 0; i < cur_nodes; ++i)</pre>
   if (id[i] < 0) id[i] = cur_id++;</pre>
 for (int i = 0; i < edges.size(); ++i) {</pre>
   int u = edges[i].u, v = edges[i].v, w = edges[i].w;
   edges[i].u = id[u];
   edges[i].v = id[v];
   if (id[u] != id[v])
     edges[i].w -= lo[v];
 }
 cur_nodes = cur_id;
```

```
root = id[root];
}
return ans;
}
```

4.3 eulerian path

```
// Taken from https://github.com/lbv/pc-code/blob/master/code/graph.cpp
// Eulerian Trail
struct Euler {
 ELV adj; IV t;
 Euler(ELV Adj) : adj(Adj) {}
 void build(int u) {
   while(! adj[u].empty()) {
     int v = adj[u].front().v;
     adj[u].erase(adj[u].begin());
     build(v);
   t.push_back(u);
}:
bool eulerian_trail(IV &trail) {
 Euler e(adj);
 int odd = 0, s = 0;
    for (int v = 0; v < n; v++) {
    int diff = abs(in[v] - out[v]);
    if (diff > 1) return false;
    if (diff == 1) {
    if (++odd > 2) return false;
    if (out[v] > in[v]) start = v;
    }
    }
    */
  e.build(s);
 reverse(e.t.begin(), e.t.end());
 trail = e.t;
 return true;
```

4.4 karp min mean cycle

```
/**
 * Finds the min mean cycle, if you need the max mean cycle
 * just add all the edges with negative cost and print
 * ans * -1
 * test: uva, 11090 - Going in Cycle!!
const int MN = 1000;
struct edge{
 int v;
 long long w;
 edge(){} edge(int v, int w) : v(v), w(w) {}
};
long long d[MN][MN];
// This is a copy of g because increments the size
// pass as reference if this does not matter.
int karp(vector<vector<edge> > g) {
 int n = g.size();
 g.resize(n + 1); // this is important
 for (int i = 0; i < n; ++i)
   if (!g[i].empty())
     g[n].push_back(edge(i,0));
  ++n;
 for(int i = 0; i < n; ++i)
   fill(d[i],d[i]+(n+1),INT_MAX);
 d[n - 1][0] = 0;
 for (int k = 1; k \le n; ++k) for (int u = 0; u \le n; ++u) {
   if (d[u][k - 1] == INT_MAX) continue;
   for (int i = g[u].size() - 1; i >= 0; --i)
     d[g[u][i].v][k] = min(d[g[u][i].v][k], d[u][k-1] + g[u][i].w);
 }
 bool flag = true;
 for (int i = 0; i < n && flag; ++i)</pre>
   if (d[i][n] != INT_MAX)
```

```
flag = false;

if (flag) {
    return true; // return true if there is no a cycle.
}

double ans = 1e15;

for (int u = 0; u + 1 < n; ++u) {
    if (d[u][n] == INT_MAX) continue;
    double W = -1e15;

    for (int k = 0; k < n; ++k)
        if (d[u][k] != INT_MAX)
        W = max(W, (double)(d[u][n] - d[u][k]) / (n - k));

    ans = min(ans, W);
}

// printf("%.2lf\n", ans);
cout << fixed << setprecision(2) << ans << endl;
return false;
}</pre>
```

4.5 konig's theorem

In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover

4.6 minimum path cover in DAG

Given a directed acyclic graph G = (V, E), we are to find the minimum number of vertex-disjoint paths to cover each vertex in V.

We can construct a bipartite graph $G' = (Vout \cup Vin, E')$ from G, where :

```
Vout = \{v \in V : v \text{ has positive out } - degree\}
Vin = \{v \in V : v \text{ has positive in } - degree\}
E' = \{(u, v) \in Vout \times Vin : (u, v) \in E\}
```

Then it can be shown, via König's theorem, that G' has a matching of size m if and only if there exists n-m vertex-disjoint paths that cover each vertex in

G, where n is the number of vertices in G and m is the maximum cardinality bipartite mathching in G'.

Therefore, the problem can be solved by finding the maximum cardinality matching in G' instead.

NOTE: If the paths are note necesarily disjoints, find the transitive closure and solve the problem for disjoint paths.

4.7 planar graph (euler)

Euler's formula states that if a finite, connected, planar graph is drawn in the plane without any edge intersections, and v is the number of vertices, e is the number of edges and f is the number of faces (regions bounded by edges, including the outer, infinitely large region), then:

$$f + v = e + 2$$

It can be extended to non connected planar graphs with c connected components:

$$f + v = e + c + 1$$

4.8 tarjan scc

```
const int MN = 20002;
struct tarjan_scc {
  int scc[MN], low[MN], d[MN], stacked[MN];
  int ticks, current_scc;
  deque<int> s; // used as stack.

  tarjan_scc() {}

  void init () {
    memset(scc, -1, sizeof scc);
    memset(d, -1, sizeof d);
    memset(stacked, 0, sizeof stacked);
    s.clear();
    ticks = current_scc = 0;
}

  void compute(vector<vector<int> > &g, int u) {
    d[u] = low[u] = ticks++;
```

```
s.push_back(u);
   stacked[u] = true;
   for (int i = 0; i < g[u].size(); ++i) {</pre>
     int v = g[u][i];
     if (d[v] == -1)
       compute(g, v);
     if (stacked[v]) {
       low[u] = min(low[u], low[v]);
     }
   }
   if (d[u] == low[u]) { // root
     int v;
     do {
       v = s.back();s.pop_back();
       stacked[v] = false;
       scc[v] = current_scc;
     } while (u != v);
     current_scc++;
 }
};
```

4.9 two sat (with kosaraju)

```
/**
 * Given a set of clauses (a1 v a2)^(a2 v a3)...
 * this algorithm find a solution to it set of clauses.
 * test: http://lightoj.com/volume_showproblem.php?problem=1251
 **/

#include<bits/stdc++.h>
using namespace std;
#define MAX 100000
#define endl '\n'
vector<int> G[MAX];
vector<int> GT[MAX];
vector<int> Ftime;
vector<vector<int> > SCC;
bool visited[MAX];
int n;
```

```
void dfs1(int n){
 visited[n] = 1;
 for (int i = 0; i < G[n].size(); ++i) {</pre>
   int curr = G[n][i]:
   if (visited[curr]) continue;
   dfs1(curr);
 }
 Ftime.push_back(n);
}
void dfs2(int n, vector<int> &scc) {
 visited[n] = 1;
 scc.push_back(n);
 for (int i = 0;i < GT[n].size(); ++i) {</pre>
   int curr = GT[n][i]:
   if (visited[curr]) continue;
   dfs2(curr, scc);
 }
}
void kosaraju() {
 memset(visited, 0, sizeof visited);
 for (int i = 0; i < 2 * n; ++i) {</pre>
   if (!visited[i]) dfs1(i);
 }
 memset(visited, 0, sizeof visited);
 for (int i = Ftime.size() - 1; i >= 0; i--) {
   if (visited[Ftime[i]]) continue;
   vector<int> _scc;
   dfs2(Ftime[i],_scc);
   SCC.push_back(_scc);
 }
}
 * After having the SCC, we must traverse each scc, if in one SCC are -b
     y b, there is not a solution.
```

```
* Otherwise we build a solution, making the first "node" that we find
     truth and its complement false.
bool two sat(vector<int> &val) {
 kosaraju();
 for (int i = 0; i < SCC.size(); ++i) {</pre>
   vector<bool> tmpvisited(2 * n, false);
   for (int j = 0; j < SCC[i].size(); ++j) {</pre>
     if (tmpvisited[SCC[i][j] ^ 1]) return 0;
     if (val[SCC[i][j]] != -1) continue;
     else {
       val[SCC[i][j]] = 0;
       val[SCC[i][j] ^ 1] = 1;
     tmpvisited[SCC[i][j]] = 1;
 }
 return 1;
// Example of use
int main() {
 int m, u, v, nc = 0, t; cin >> t;
 // n = "nodes" number, m = clauses number
 while (t--) {
   cin >> m >> n;
   Ftime.clear():
   SCC.clear();
   for (int i = 0; i < 2 * n; ++i) {
     G[i].clear();
     GT[i].clear();
   // (a1 v a2) = (a1 -> a2) = (a2 -> a1)
   for (int i = 0; i < m ; ++i) {</pre>
     cin >> u >> v;
     int t1 = abs(u) - 1;
     int t2 = abs(v) - 1;
     int p = t1 * 2 + ((u < 0)? 1 : 0);
     int q = t2 * 2 + ((v < 0)? 1 : 0);
```

```
G[p ^ 1].push_back(q);
    G[q ^ 1].push_back(p);
    GT[p].push_back(q ^ 1);
    GT[q].push_back(p ^ 1);
  vector < int > val(2 * n, -1);
  cout << "Case " << ++nc <<": ":
  if (two_sat(val)) {
    cout << "Yes" << endl;</pre>
    vector<int> sol:
    for (int i = 0; i < 2 * n; ++i)
      if (i % 2 == 0 and val[i] == 1)
        sol.push_back(i / 2 + 1);
    cout << sol.size() ;</pre>
    for (int i = 0; i < sol.size(); ++i) {</pre>
      cout << " " << sol[i];
    cout << endl;</pre>
  } else {
    cout << "No" << endl;</pre>
  }
}
return 0;
```

5 Math

}

5.1 Lucas theorem

For non-negative integers m and n and a prime p, the following congruence relation holds: :

$$\binom{m}{n} \equiv \prod_{i=0}^{k} \binom{m_i}{n_i} \pmod{p},$$

where:

$$m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0,$$

and:

$$n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$$

are the base p expansions of m and n respectively. This uses the convention that $\binom{m}{n} = 0$ if $m \le n$.

5.2 cumulative sum of divisors

```
The function SOD(n) (sum of divisors) is defined
as the summation of all the actual divisors of
an integer number n. For example,
 SOD(24) = 2+3+4+6+8+12 = 35.
The function CSOD(n) (cumulative SOD) of an integer n, is defined as
    below:
 csod(n) = \sum_{i=1}^{n} sod(i)
It can be computed in O(sqrt(n)):
*/
long long csod(long long n) {
 long long ans = 0;
 for (long long i = 2; i * i <= n; ++i) {
   long long j = n / i;
   ans += (i + j) * (j - i + 1) / 2;
   ans += i * (j - i);
 return ans;
```

5.3 fft

```
/**
 * Fast Fourier Transform.
 * Useful to compute convolutions.
 * computes:
 * C(f star g)[n] = sum_m(f[m] * g[n - m])
 * for all n.
 * test: icpc live archive, 6886 - Golf Bot
 * */

using namespace std;
#include<bits/stdc++.h>
#define D(x) cout << #x " = " << (x) << endl
#define endl '\n'</pre>
```

```
const int MN = 262144 << 1;</pre>
int d[MN + 10], d2[MN + 10];
const double PI = acos(-1.0);
struct cpx {
 double real, image;
 cpx(double _real, double _image) {
   real = _real;
   image = _image;
 }
 cpx(){}
};
cpx operator + (const cpx &c1, const cpx &c2) {
 return cpx(c1.real + c2.real, c1.image + c2.image);
}
cpx operator - (const cpx &c1, const cpx &c2) {
 return cpx(c1.real - c2.real, c1.image - c2.image);
}
cpx operator * (const cpx &c1, const cpx &c2) {
 return cpx(c1.real*c2.real - c1.image*c2.image, c1.real*c2.image +
      c1.image*c2.real);
}
int rev(int id, int len) {
 int ret = 0;
 for (int i = 0; (1 << i) < len; i++) {
   ret <<= 1;
   if (id & (1 << i)) ret |= 1;</pre>
 }
 return ret;
}
cpx A[1 << 20];
void FFT(cpx *a, int len, int DFT) {
 for (int i = 0; i < len; i++)</pre>
   A[rev(i, len)] = a[i];
 for (int s = 1; (1 << s) <= len; s++) {
   int m = (1 << s);
```

```
cpx wm = cpx(cos(DFT * 2 * PI / m), sin(DFT * 2 * PI / m));
   for(int k = 0; k < len; k += m) {
     cpx w = cpx(1, 0);
     for(int j = 0; j < (m >> 1); j++) {
       cpx t = w * A[k + j + (m >> 1)];
       cpx u = A[k + j];
       A[k + j] = u + t;
       A[k + j + (m >> 1)] = u - t;
       w = w * wm:
     }
   }
 if (DFT == -1) for (int i = 0; i < len; i++) A[i].real /= len,
      A[i].image /= len;
 for (int i = 0; i < len; i++) a[i] = A[i];</pre>
 return:
cpx in[1 << 20];
void solve(int n) {
 memset(d, 0, sizeof d);
 int t;
 for (int i = 0; i < n; ++i) {</pre>
   cin >> t;
   d[t] = true:
 int m;
 cin >> m;
 vector<int> q(m);
 for (int i = 0; i < m; ++i)</pre>
   cin >> q[i];
 for (int i = 0; i < MN; ++i) {</pre>
   if (d[i])
     in[i] = cpx(1, 0);
   else
     in[i] = cpx(0, 0);
 FFT(in, MN, 1);
 for (int i = 0; i < MN; ++i) {</pre>
   in[i] = in[i] * in[i];
 FFT(in, MN, -1);
```

```
int ans = 0;
for (int i = 0; i < q.size(); ++i) {
    if (in[q[i]].real > 0.5 || d[q[i]]) {
        ans++;
    }
}
cout << ans << endl;
}
int main() {
    ios_base::sync_with_stdio(false);cin.tie(NULL);
    int n;
    while (cin >> n)
        solve(n);
    return 0;
}
```

5.4 fibonacci properties

Let A, B and n be integer numbers.

$$k = A - B \tag{1}$$

$$F_A F_B = F_{k+1} F_A^2 + F_k F_A F_{A-1} \tag{2}$$

$$\sum_{i=0}^{n} F_i^2 = F_{n+1} F_n \tag{3}$$

ev(n) = returns 1 if n is even.

$$\sum_{i=0}^{n} F_i F_{i+1} = F_{n+1}^2 - ev(n) \tag{4}$$

$$\sum_{i=0}^{n} F_i F_{i-1} = \sum_{i=0}^{n-1} F_i F_{i+1} \tag{5}$$

5.5 sigma function

the sigma function is defined as:

$$\sigma_x(n) = \sum_{d|n} d^x$$

when x=0 is called the divisor function, that counts the number of positive divisors of n.

Now, we are interested in find

$$\sum_{d|n} \sigma_0(d)$$

if n is written as prime factorization:

$$n = \prod_{i=1}^{k} P_i^{e_k}$$

we can demonstrate that:

$$\sum_{d|n} \sigma_0(d) = \prod_{i=1}^k g(e_k + 1)$$

where g(x) is the sum of the first x positive numbers:

$$g(x) = (x * (x+1))/2$$

6 Matrix

6.1 matrix

```
const int MN = 111;
const int mod = 10000;

struct matrix {
   int r, c;
   int m[MN][MN];

matrix (int _r, int _c) : r (_r), c (_c) {
   memset(m, 0, sizeof m);
}

void print() {
   for (int i = 0; i < r; ++i) {
      for (int j = 0; j < c; ++j)
   }</pre>
```

```
cout << m[i][j] << " ";</pre>
     cout << endl;</pre>
   }
 }
  int x[MN][MN];
  matrix & operator *= (const matrix &o) {
   memset(x, 0, sizeof x);
   for (int i = 0; i < r; ++i)
     for (int k = 0; k < c; ++k)
       if (m[i][k] != 0)
         for (int j = 0; j < c; ++j) {
           x[i][j] = (x[i][j] + ((m[i][k] * o.m[k][j]) % mod)) % mod;
   memcpy(m, x, sizeof(m));
   return *this;
 }
};
void matrix_pow(matrix b, long long e, matrix &res) {
 memset(res.m, 0, sizeof res.m);
 for (int i = 0; i < b.r; ++i)</pre>
   res.m[i][i] = 1;
 if (e == 0) return;
 while (true) {
   if (e & 1) res *= b;
   if ((e >>= 1) == 0) break;
   b *= b;
 }
```

7 Misc

7.1 Template Java

```
import java.io.*;
import java.util.StringTokenizer;

public class Template {
    public static void main(String []args) throws IOException {
```

```
Scanner in = new Scanner(System.in);
       OutputWriter out = new OutputWriter(System.out);
       Task solver = new Task();
       solver.solve(in, out);
       out.close();
}
class Task{
   public void solve(Scanner in, OutputWriter out){
}
class Scanner{
   public BufferedReader reader;
   public StringTokenizer st;
   public Scanner(InputStream stream){
       reader = new BufferedReader(new InputStreamReader(stream));
       st = null;
   }
   public String next(){
       while(st == null || !st.hasMoreTokens()){
           trv{
              String line = reader.readLine();
              if(line == null) return null;
              st = new StringTokenizer(line);
          }catch (Exception e){
              throw (new RuntimeException());
          }
       }
       return st.nextToken();
   public int nextInt(){
       return Integer.parseInt(next());
   public long nextLong(){
       return Long.parseLong(next());
   public double nextDouble(){
       return Double.parseDouble(next());
```

```
}
}
class OutputWriter{
   BufferedWriter writer;
   public OutputWriter(OutputStream stream){
       writer = new BufferedWriter(new OutputStreamWriter(stream));
   }
   public void print(int i) throws IOException {
       writer.write(i);
   }
   public void print(String s) throws IOException {
       writer.write(s):
   }
   public void print(char []c) throws IOException {
       writer.write(c);
   }
   public void close() throws IOException {
       writer.close();
   }
```

7.2 dates

```
// number of days in blocks of years
const int p400 = 400*365 + leap_years(400);
const int p100 = 100*365 + leap_years(100);
const int p4 = 4*365 + 1;
const int p1 = 365;
int date_to_days(int d, int m, int y)
 return (y - 1) * 365 + leap_years(y - 1) + (is_leap(y) ? B[m] : A[m]) +
void days_to_date(int days, int &d, int &m, int &y)
 bool top100; // are we in the top 100 years of a 400 block?
 bool top4; // are we in the top 4 years of a 100 block?
 bool top1; // are we in the top year of a 4 block?
 v = 1;
 top100 = top4 = top1 = false;
 y += ((days-1) / p400) * 400;
 d = (days-1) \% p400 + 1;
 if (d > p100*3) top100 = true, d = 3*p100, y += 300;
 else y += ((d-1) / p100) * 100, d = (d-1) % p100 + 1;
 if (d > p4*24) top4 = true, d -= 24*p4, y += 24*4;
 else y += ((d-1) / p4) * 4, d = (d-1) % p4 + 1;
 if (d > p1*3) top1 = true, d = p1*3, y += 3;
 else y += (d-1) / p1, d = (d-1) % p1 + 1;
  const int *ac = top1 && (!top4 || top100) ? B : A;
 for (m = 1; m < 12; ++m) if (d \le ac[m + 1]) break;
 d = ac[m]:
```

7.3 fraction

```
struct frac{
  long long x, y;
  frac(long long a, long long b) {
    long long g = __gcd(a, b);
    x = a / g;
```

```
y = b / g;
}
bool operator < (const frac &o) const {
   return (x * o.y < y * o.x);
}
};</pre>
```

7.4 io

```
// taken from :
    https://github.com/lbv/pc-code/blob/master/solved/c-e/diablo/diablo.cpp
// this is very fast as well :
    https://github.com/lbv/pc-code/blob/master/code/input.cpp
typedef unsigned int u32;
#define BUF 524288
struct Reader {
  char buf[BUF]; char b; int bi, bz;
 Reader() { bi=bz=0; read(); }
 void read() {
   if (bi==bz) { bi=0; bz = fread(buf, 1, BUF, stdin); }
   b = bz ? buf[bi++] : 0; }
 void skip() { while (b > 0 && b <= 32) read(); }</pre>
 u32 next_u32() {
   u32 v = 0; for (skip(); b > 32; read()) v = v*10 + b-48; return v; }
 int next_int() {
   int v = 0; bool s = false;
   skip(); if (b == '-') { s = true; read(); }
   for (; 48 <= b \& \& b <= 57; read()) v = v*10 + b-48; return s? -v : v;}
  char next_char() { skip(); char c = b; read(); return c; }
};
```

8 Number theory

8.1 convolution

```
typedef long long int LL;
typedef pair<LL, LL> PLL;
inline bool is_pow2(LL x) {
```

```
return (x & (x-1)) == 0;
inline int ceil_log2(LL x) {
 int ans = 0;
 --x;
 while (x != 0) {
   x >>= 1:
   ans++;
 return ans:
/* Returns the convolution of the two given vectors in time proportional
    to n*log(n).
* The number of roots of unity to use nroots_unity must be set so that
     the product of the first
* nroots_unity primes of the vector nth_roots_unity is greater than the
     maximum value of the
* convolution. Never use sizes of vectors bigger than 2^24, if you need
     to change the values of
* the nth roots of unity to appropriate primes for those sizes.
vector<LL> convolve(const vector<LL> &a, const vector<LL> &b, int
    nroots_unity = 2) {
 int N = 1 \ll ceil_log2(a.size() + b.size());
 vector<LL> ans(N.O), fA(N), fB(N), fC(N):
 LL modulo = 1;
 for (int times = 0; times < nroots_unity; times++) {</pre>
   fill(fA.begin(), fA.end(), 0);
   fill(fB.begin(), fB.end(), 0);
   for (int i = 0; i < a.size(); i++) fA[i] = a[i];</pre>
   for (int i = 0; i < b.size(); i++) fB[i] = b[i];</pre>
   LL prime = nth_roots_unity[times].first;
   LL inv_modulo = mod_inv(modulo % prime, prime);
   LL normalize = mod_inv(N, prime);
   ntfft(fA, 1, nth_roots_unity[times]);
   ntfft(fB, 1, nth_roots_unity[times]);
   for (int i = 0; i < N; i++) fC[i] = (fA[i] * fB[i]) % prime;</pre>
   ntfft(fC, -1, nth_roots_unity[times]);
   for (int i = 0; i < N; i++) {</pre>
     LL curr = (fC[i] * normalize) % prime;
     LL k = (curr - (ans[i] % prime) + prime) % prime;
     k = (k * inv_modulo) % prime;
     ans[i] += modulo * k;
```

```
}
  modulo *= prime;
}
return ans;
}
```

8.2 crt

```
/**
 * Chinese remainder theorem.
 * Find z such that z % x[i] = a[i] for all i.
 * */
long long crt(vector<long long> &a, vector<long long> &x) {
 long long z = 0;
 long long n = 1;
 for (int i = 0; i < x.size(); ++i)
    n *= x[i];

for (int i = 0; i < a.size(); ++i) {
    long long tmp = (a[i] * (n / x[i])) % n;
    tmp = (tmp * mod_inv(n / x[i], x[i])) % n;
    z = (z + tmp) % n;
}

return (z + n) % n;
}</pre>
```

8.3 discrete logarithm

```
// Computes x which a ^ x = b mod n.
long long d_log(long long a, long long b, long long n) {
  long long m = ceil(sqrt(n));
  long long aj = 1;
  map<long long, long long> M;
  for (int i = 0; i < m; ++i) {
    if (!M.count(aj))
        M[aj] = i;
        aj = (aj * a) % n;
}</pre>
```

```
long long coef = mod_pow(a, n - 2, n);
coef = mod_pow(coef, m, n);
// coef = a ^ (-m)
long long gamma = b;
for (int i = 0; i < m; ++i) {
   if (M.count(gamma)) {
     return i * m + M[gamma];
   } else {
     gamma = (gamma * coef) % n;
   }
return -1;
}</pre>
```

8.4 ext euclidean

```
void ext_euclid(long long a, long long b, long long &x, long long &y,
    long long &g) {
    x = 0, y = 1, g = b;
    long long m, n, q, r;
    for (long long u = 1, v = 0; a != 0; g = a, a = r) {
        q = g / a, r = g % a;
        m = x - u * q, n = y - v * q;
        x = u, y = v, u = m, v = n;
    }
}
```

8.5 highest exponent factorial

```
int highest_exponent(int p, const int &n){
  int ans = 0;
  int t = p;
  while(t <= n){
    ans += n/t;
    t*=p;
  }
  return ans;
}</pre>
```

8.6 miller rabin

```
const int rounds = 20;
// checks whether a is a witness that n is not prime, 1 < a < n
bool witness(long long a, long long n) {
 // check as in Miller Rabin Primality Test described
 long long u = n - 1;
 int t = 0;
 while (u % 2 == 0) {
   t++;
   u >>= 1:
 long long next = mod_pow(a, u, n);
 if (next == 1) return false;
 long long last;
 for (int i = 0; i < t; ++i) {</pre>
   last = next:
   next = mod_mul(last, last, n);
   if (next == 1) {
     return last != n - 1;
   }
 }
 return next != 1;
}
// Checks if a number is prime with prob 1 - 1 / (2 ^ it)
// D(miller_rabin(999999999999997LL) == 1);
// D(miller_rabin(999999999971LL) == 1);
// D(miller_rabin(7907) == 1);
bool miller_rabin(long long n, int it = rounds) {
 if (n <= 1) return false;</pre>
 if (n == 2) return true;
 if (n % 2 == 0) return false;
 for (int i = 0; i < it; ++i) {</pre>
   long long a = rand() \% (n - 1) + 1;
   if (witness(a, n)) {
     return false;
   }
 }
 return true;
```

8.7 mod inv

```
long long mod_inv(long long n, long long m) {
  long long x, y, gcd;
  ext_euclid(n, m, x, y, gcd);
  if (gcd != 1)
    return 0;
  return (x + m) % m;
}
```

8.8 mod mul

```
// Computes (a * b) % mod
long long mod_mul(long long a, long long b, long long mod) {
  long long x = 0, y = a % mod;
  while (b > 0) {
    if (b & 1)
        x = (x + y) % mod;
        y = (y * 2) % mod;
        b /= 2;
  }
  return x % mod;
}
```

8.9 mod pow

```
// Computes ( a ^ exp ) % mod.
long long mod_pow(long long a, long long exp, long long mod) {
  long long ans = 1;
  while (exp > 0) {
    if (exp & 1)
      ans = mod_mul(ans, a, mod);
      a = mod_mul(a, a, mod);
      exp >>= 1;
  }
  return ans;
}
```

8.10 number theoretic transform

```
typedef long long int LL;
typedef pair<LL, LL> PLL;
/* The following vector of pairs contains pairs (prime, generator)
* where the prime has an Nth root of unity for N being a power of two.
 * The generator is a number g s.t g^(p-1)=1 (mod p)
 * but is different from 1 for all smaller powers */
vector<PLL> nth roots unity {
 {1224736769,330732430},{1711276033,927759239},{167772161,167489322},
  {469762049,343261969},{754974721,643797295},{1107296257,883865065}};
PLL ext_euclid(LL a, LL b) {
 if (b == 0)
   return make_pair(1,0);
 pair<LL,LL> rc = ext_euclid(b, a % b);
 return make_pair(rc.second, rc.first - (a / b) * rc.second);
}
//returns -1 if there is no unique modular inverse
LL mod_inv(LL x, LL modulo) {
 PLL p = ext_euclid(x, modulo);
 if ( (p.first * x + p.second * modulo) != 1 )
   return -1:
 return (p.first+modulo) % modulo;
}
//Number theory fft. The size of a must be a power of 2
void ntfft(vector<LL> &a, int dir, const PLL &root_unity) {
 int n = a.size():
 LL prime = root_unity.first;
 LL basew = mod_pow(root_unity.second, (prime-1) / n, prime);
 if (dir < 0) basew = mod_inv(basew, prime);</pre>
 for (int m = n: m \ge 2: m \ge 1) {
   int mh = m >> 1;
   LL w = 1:
   for (int i = 0; i < mh; i++) {</pre>
     for (int j = i; j < n; j += m) {
       int k = j + mh;
       LL x = (a[j] - a[k] + prime) % prime;
       a[j] = (a[j] + a[k]) \% prime;
       a[k] = (w * x) \% prime;
```

```
w = (w * basew) % prime;
}
basew = (basew * basew) % prime;
}
int i = 0;
for (int j = 1; j < n - 1; j++) {
  for (int k = n >> 1; k > (i ^= k); k >>= 1);
  if (j < i) swap(a[i], a[j]);
}
</pre>
```

8.11 pollard rho factorize

```
long long pollard_rho(long long n) {
 long long x, y, i = 1, k = 2, d;
 x = y = rand() \% n;
 while (1) {
   ++i;
   x = mod_mul(x, x, n);
   x += 2;
   if (x \ge n) x -= n;
   if (x == y) return 1;
   d = \_gcd(abs(x - y), n);
   if (d != 1) return d:
   if (i == k) {
     y = x;
     k *= 2;
 }
 return 1;
// Returns a list with the prime divisors of n
vector<long long> factorize(long long n) {
 vector<long long> ans;
 if (n == 1)
   return ans;
 if (miller_rabin(n)) {
   ans.push_back(n);
 } else {
   long long d = 1;
   while (d == 1)
```

```
d = pollard_rho(n);
vector<long long> dd = factorize(d);
ans = factorize(n / d);
for (int i = 0; i < dd.size(); ++i)
    ans.push_back(dd[i]);
}
return ans;
}</pre>
```

8.12 totient sieve

```
for (int i = 1; i < MN; i++)
   phi[i] = i;

for (int i = 1; i < MN; i++)
   if (!sieve[i]) // is prime
    for (int j = i; j < MN; j += i)
        phi[j] -= phi[j] / i;</pre>
```

8.13 totient

```
long long totient(long long n) {
   if (n == 1) return 0;
   long long ans = n;
   for (int i = 0; primes[i] * primes[i] <= n; ++i) {
      if ((n % primes[i]) == 0) {
        while ((n % primes[i]) == 0) n /= primes[i];
        ans -= ans / primes[i];
    }
   if (n > 1) {
      ans -= ans / n;
   }
   return ans;
}
```

9 Strings

9.1 Incremental Aho Corasick

```
class IncrementalAhoCorasic {
 static const int Alphabets = 26;
 static const int AlphabetBase = 'a';
 struct Node {
   Node *fail:
   Node *next[Alphabets];
   int sum;
   Node() : fail(NULL), next{}, sum(0) { }
 };
 struct String {
   string str;
   int sign;
 };
public:
 //totalLen = sum of (len + 1)
 void init(int totalLen) {
   nodes.resize(totalLen);
   nNodes = 0;
   strings.clear();
   roots.clear();
   sizes.clear();
   que.resize(totalLen);
  void insert(const string &str, int sign) {
   strings.push_back(String{ str, sign });
   roots.push_back(nodes.data() + nNodes);
   sizes.push_back(1);
   nNodes += (int)str.size() + 1;
   auto check = [&]() { return sizes.size() > 1 && sizes.end()[-1] ==
        sizes.end()[-2]; };
   if(!check())
     makePMA(strings.end() - 1, strings.end(), roots.back(), que);
   while(check()) {
     int m = sizes.back();
     roots.pop_back();
     sizes.pop_back();
     sizes.back() += m;
```

```
if(!check())
       makePMA(strings.end() - m * 2, strings.end(), roots.back(), que);
   }
 }
 int match(const string &str) const {
   int res = 0;
   for(const Node *t : roots)
     res += matchPMA(t, str);
   return res;
 }
private:
 static void makePMA(vector<String>::const_iterator begin,
      vector<String>::const_iterator end, Node *nodes, vector<Node*>
      &que) {
   int nNodes = 0;
   Node *root = new(&nodes[nNodes ++]) Node();
   for(auto it = begin; it != end; ++ it) {
     Node *t = root;
     for(char c : it->str) {
       Node *&n = t->next[c - AlphabetBase];
       if(n == nullptr)
        n = new(&nodes[nNodes ++]) Node();
       t = n;
     t->sum += it->sign;
   }
   int qt = 0;
   for(Node *&n : root->next) {
     if(n != nullptr) {
       n->fail = root;
       que[qt ++] = n;
     } else {
       n = root;
     }
   for(int qh = 0; qh != qt; ++ qh) {
     Node *t = que[qh];
     int a = 0;
     for(Node *n : t->next) {
       if(n != nullptr) {
         que[qt ++] = n;
         Node *r = t->fail;
         while(r->next[a] == nullptr)
```

```
r = r - > fail:
         n->fail = r->next[a];
         n->sum += r->next[a]->sum;
       }
       ++ a;
  static int matchPMA(const Node *t, const string &str) {
   int res = 0:
   for(char c : str) {
     int a = c - AlphabetBase;
     while(t->next[a] == nullptr)
       t = t->fail;
     t = t- next[a]:
     res += t->sum;
   return res;
  vector<Node> nodes;
  int nNodes:
  vector<String> strings;
  vector<Node*> roots:
 vector<int> sizes:
 vector<Node*> que;
};
int main() {
  int m:
  while(~scanf("%d", &m)) {
   IncrementalAhoCorasic iac;
   iac.init(600000);
   rep(i, m) {
     int ty;
     char s[300001];
     scanf("%d%s", &ty, s);
     if(ty == 1) {
       iac.insert(s, +1);
     } else if(ty == 2) {
       iac.insert(s, -1);
     } else if(ty == 3) {
       int ans = iac.match(s);
```

```
printf("%d\n", ans);
    fflush(stdout);
} else {
    abort();
}

return 0;
```

9.2 minimal string rotation

```
// Lexicographically minimal string rotation
int lmsr() {
 string s;
 cin >> s;
 int n = s.size();
 s += s;
 vector<int> f(s.size(), -1);
 int k = 0:
 for (int j = 1; j < 2 * n; ++j) {
   int i = f[j - k - 1];
   while (i != -1 && s[j] != s[k + i + 1]) {
     if (s[j] < s[k + i + 1])
       k = j - i - 1;
     i = f[i];
   if (i == -1 \&\& s[j] != s[k + i + 1]) {
     if (s[j] < s[k + i + 1]) {
       k = j;
     f[j - k] = -1;
   } else {
     f[j - k] = i + 1;
   }
 }
 return k;
```

9.3 suffix array

```
* 0 (n log^2 (n))
 * See http://web.stanford.edu/class/cs97si/suffix-array.pdf for reference
struct entry{
 int a, b, p;
  entry(){}
  entry(int x, int y, int z): a(x), b(y), p(z){}
  bool operator < (const entry &o) const {</pre>
   return (a == o.a) ? (b == o.b) ? (p < o.p) : (b < o.b) : (a < o.a);
};
struct SuffixArray{
  const int N;
  string s;
  vector<vector<int> > P;
  vector<entry> M;
  SuffixArray(const string &s): N(s.length()), s(s), P(1, vector<int>
      (N, 0)), M(N) {
    for (int i = 0; i < N; ++i)</pre>
     P[0][i] = (int) s[i];
   for (int skip = 1, level = 1; skip < N; skip *= 2, level++) {</pre>
     P.push_back(vector<int>(N, 0));
     for (int i = 0 ; i < N; ++i) {</pre>
       int next = ((i + skip) < N) ? P[level - 1][i + skip] : -10000;
       M[i] = entry(P[level - 1][i], next, i);
     sort(M.begin(), M.end());
     for (int i = 0; i < N; ++i)</pre>
       P[level][M[i].p] = (i > 0 \text{ and } M[i].a == M[i - 1].a \text{ and } M[i].b ==
            M[i - 1].b) ? P[level][M[i - 1].p] : i;
   }
  }
  vector<int> getSuffixArray(){
    vector<int> &rank = P.back();
    vector<pair<int, int> > inv(rank.size());
   for (int i = 0; i < rank.size(); ++i)</pre>
     inv[i] = make_pair(rank[i], i);
    sort(inv.begin(), inv.end());
    vector<int> sa(rank.size());
```

```
for (int i = 0; i < rank.size(); ++i)</pre>
     sa[i] = inv[i].second;
   return sa;
 }
 // returns the length of the longest common prefix of s[i...L-1] and
      s[j...L-1]
 int lcp(int i, int j) {
   int len = 0;
   if (i == j) return N - i;
   for (int k = P.size() - 1; k >= 0 && i < N && j < N; --k) {
     if (P[k][i] == P[k][j]) {
       i += 1 << k;
       j += 1 << k;
       len += 1 << k;
     }
   }
   return len;
 }
};
```

9.4 suffix automaton

```
/*
 * Suffix automaton:
 * This implementation was extended to maintain (online) the
 * number of different substrings. This is equivalent to compute
 * the number of paths from the initial state to all the other
 * states.
 * The overall complexity is O(n)
 * can be tested here:
     https://www.urionlinejudge.com.br/judge/en/problems/view/1530
 * */
struct state {
 int len, link;
 long long num_paths;
 map<int, int> next;
};
const int MN = 200011;
state sa[MN << 1];
```

```
int sz, last;
long long tot_paths;
void sa_init() {
 sz = 1;
 last = 0;
 sa[0].len = 0;
 sa[0].link = -1;
 sa[0].next.clear();
 sa[0].num_paths = 1;
 tot_paths = 0;
void sa_extend(int c) {
 int cur = sz++;
 sa[cur].len = sa[last].len + 1;
 sa[cur].next.clear();
 sa[cur].num_paths = 0;
 for (p = last; p != -1 \&\& !sa[p].next.count(c); p = sa[p].link) {
   sa[p].next[c] = cur;
   sa[cur].num_paths += sa[p].num_paths;
   tot_paths += sa[p].num_paths;
 if (p == -1) {
   sa[cur].link = 0;
 } else {
   int q = sa[p].next[c];
   if (sa[p].len + 1 == sa[q].len) {
     sa[cur].link = q;
   } else {
     int clone = sz++;
     sa[clone].len = sa[p].len + 1;
     sa[clone].next = sa[q].next;
     sa[clone].num_paths = 0;
     sa[clone].link = sa[q].link;
     for (; p!= -1 && sa[p].next[c] == q; p = sa[p].link) {
       sa[p].next[c] = clone;
       sa[q].num_paths -= sa[p].num_paths;
       sa[clone].num_paths += sa[p].num_paths;
     sa[q].link = sa[cur].link = clone;
   }
```

```
last = cur;
}
```

9.5 z algorithm

```
using namespace std;
#include<bits/stdc++.h>
vector<int> compute_z(const string &s){
 int n = s.size();
 vector<int> z(n,0);
 int l,r;
 r = 1 = 0;
 for(int i = 1; i < n; ++i){</pre>
   if(i > r) {
     1 = r = i;
     while(r < n and s[r - 1] == s[r])r++;
     z[i] = r - 1;r--;
   }else{
     int k = i-1;
     if(z[k] < r - i +1) z[i] = z[k];
     else {
      l = i;
       while(r < n and s[r - 1] == s[r])r++;
       z[i] = r - 1;r--;
     }
 return z;
int main(){
 //string line;cin>>line;
 string line = "alfalfa";
 vector<int> z = compute_z(line);
 for(int i = 0; i < z.size(); ++i ){</pre>
   if(i)cout<<" ";</pre>
   cout<<z[i];
 }
 cout<<endl;</pre>
```

```
// must print "0 0 0 4 0 0 1"

return 0;
}
```