Team notebook

August 22, 2016

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1 Algorithms

1.1 sliding window

```
/*
 * Given an array ARR and an integer K, the problem boils down to
     computing for each index i: min(ARR[i], ARR[i-1], ..., ARR[i-K+1]).
 * if mx == true, returns the maximun.
 * http://people.cs.uct.ac.za/~ksmith/articles/sliding_window_minimum.html
 * */

vector<int> sliding_window_minmax(vector<int> & ARR, int K, bool mx) {
  deque< pair<int, int> > window;
  vector<int> ans;
  for (int i = 0; i < ARR.size(); i++) {
    if (mx) {
      while (!window.empty() && window.back().first <= ARR[i])
        window.pop_back();
    } else {
      while (!window.empty() && window.back().first >= ARR[i])
        window.pop_back();
    }
    window.push_back(make_pair(ARR[i], i));

    while(window.front().second <= i - K)
      window.pop_front();

    ans.push_back(window.front().first);
  }
  return ans;
}
```

2 Data structures

2.1 hash table

```
/**
 * Micro hash table, can be used as a set.
 *  Very efficient vs std::set
 * */

const int MN = 1001;
struct ht {
  int _s[(MN + 10) >> 5];
  int len;
  void set(int id) {
    len++;
    _s[id >> 5] |= (1LL << (id & 31));
  }
  bool is_set(int id) {
    return _s[id >> 5] & (1LL << (id & 31));
  }
};
```

2.2 heavy light decomposition

```
// Heavy-Light Decomposition
struct TreeDecomposition {
  vector<int> g[MAXN], c[MAXN];
  int s[MAXN]; // subtree size
  int p[MAXN]; // parent id
  int r[MAXN]; // chain root id
  int t[MAXN]; // index used in segtree/bit/...
  int d[MAXN]; // depht
  int ts;

  void dfs(int v, int f) {
    p[v] = f;
    s[v] = 1;
    if (f != -1) d[v] = d[f] + 1;
    else d[v] = 0;

    for (int i = 0; i < g[v].size(); ++i) {
      int w = g[v][i];
      if (w != f) {
        dfs(w, v);
        s[v] += s[w];
      }
    }
  }
```

```
  void hld(int v, int f, int k) {
    t[v] = ts++;
    c[k].push_back(v);
    r[v] = k;

    int x = 0, y = -1;
    for (int i = 0; i < g[v].size(); ++i) {
      int w = g[v][i];
      if (w != f) {
        if (s[w] > x) {
          x = s[w];
          y = w;
        }
      }
    }
    if (y != -1) {
      hld(y, v, k);
    }

    for (int i = 0; i < g[v].size(); ++i) {
      int w = g[v][i];
      if (w != f && w != y) {
        hld(w, v, w);
      }
    }
  }

  void init(int n) {
    for (int i = 0; i < n; ++i) {
      g[i].clear();
    }
  }

  void add(int a, int b) {
    g[a].push_back(b);
    g[b].push_back(a);
  }

  void build() {
    ts = 0;
    dfs(0, -1);
    hld(0, 0, 0);
  }
};
```

2.3 persistent array

```
struct node {
  node *l, *r;
  int val;

  node (int x) : l(NULL), r(NULL), val(x) {}
  node () : l(NULL), r(NULL), val(-1) {}
};

typedef node* pnode;

pnode update(pnode cur, int l, int r, int at, int what) {
  pnode ans = new node();

  if (cur != NULL) {
    *ans = *cur;
  }
  if (l == r) {
    ans-> val = what;
    return ans;
  }
  int m = (l + r) >> 1;
  if (at <= m) ans-> l = update(ans-> l, l, m, at, what);
  else ans-> r = update(ans-> r, m + 1, r, at, what);
  return ans;
}

int get(pnode cur, int l, int r, int at) {
  if (cur == NULL) return 0;
  if (l == r) return cur-> val;
  int m = (l + r) >> 1;
  if (at <= m) return get(cur-> l, l, m, at);
  else        return get(cur-> r, m + 1, r, at);
}
```

2.4 persistent seg tree

```
/**
 * Important:
 * When using lazy propagation remembert to create new
 * versions for each push_down operation!!!
 * */
```

```cpp
struct node {
  node *l, *r;
  long long acc;
  int flip;

  node (int x) : l(NULL), r(NULL), acc(x), flip(0) {}
  node () : l(NULL), r(NULL), acc(0), flip(0) {}
};

typedef node* pnode;

pnode create(int l, int r) {
  if (l == r) return new node();
  pnode cur = new node();
  int m = (l + r) >> 1;
  cur-> l = create(l, m);
  cur-> r = create(m + 1, r);
  return cur;
}

pnode copy_node(pnode cur) {
  pnode ans = new node();
  *ans = *cur;
  return ans;
}

void push_down(pnode cur, int l, int r) {
  assert(cur);
  if (cur-> flip) {
    int len = r - l + 1;
    cur-> acc = len - cur-> acc;
    if (cur-> l) {
      cur-> l = copy_node(cur-> l);
      cur-> l -> flip ^= 1;
    }
    if (cur-> r) {
      cur-> r = copy_node(cur-> r);
      cur-> r -> flip ^= 1;
    }
    cur-> flip = 0;
  }
}

int get_val(pnode cur) {
  assert(cur);
  assert((cur-> flip) == 0);
  if (cur) return cur-> acc;
  return 0;
}

pnode update(pnode cur, int l, int r, int at, int what) {
  pnode ans = copy_node(cur);
  if (l == r) {
    assert(l == at);
    ans-> acc = what;
    ans-> flip = 0;
    return ans;
  }
  int m = (l + r) >> 1;
  push_down(ans, l, r);
  if (at <= m) ans-> l = update(ans-> l, l, m, at, what);
  else ans-> r = update(ans-> r, m + 1, r, at, what);

  push_down(ans-> l, l, m);
  push_down(ans-> r, m + 1, r);
  ans-> acc = get_val(ans-> l) + get_val(ans-> r);
  return ans;
}

pnode flip(pnode cur, int l, int r, int a, int b) {
  pnode ans = new node();

  if (cur != NULL) {
    *ans = *cur;
  }
  if (l > b || r < a)
    return ans;

  if (l >= a && r <= b) {
    ans-> flip ^= 1;
    push_down(ans, l, r);
    return ans;
  }

  int m = (l + r) >> 1;
  ans-> l = flip(ans-> l, l, m, a, b);
  ans-> r = flip(ans-> r, m + 1, r, a, b);
  push_down(ans-> l, l, m);
  push_down(ans-> r, m + 1, r);
```

```
  ans-> acc = get_val(ans-> l) + get_val(ans-> r);
  return ans;
}

long long get_all(pnode cur, int l, int r) {
  assert(cur);
  push_down(cur, l, r);
  return cur-> acc;
}

void traverse(pnode cur, int l, int r) {
  if (!cur) return;
  cout << l << " - " << r << " : " << (cur-> acc) << " " << (cur-> flip)
      << endl;
  traverse(cur-> l, l, (l + r) >> 1);
  traverse(cur-> l, 1 + ((l + r) >> 1), r);
}
```

2.5 segment tree

```
/**
 * Taken from: http://codeforces.com/blog/entry/18051
 * */

const int N = 1e5; // limit for array size
int n; // array size
int t[2 * N];

void build() { // build the tree
  for (int i = n - 1; i > 0; --i) t[i] = t[i<<1] + t[i<<1|1];
}

// Single modification, range query.
void modify(int p, int value) { // set value at position p
  for (t[p += n] = value; p > 1; p >>= 1) t[p>>1] = t[p] + t[p^1];
}

int query(int l, int r) { // sum on interval [l, r)
  int res = 0;
  for (l += n, r += n; l < r; l >>= 1, r >>= 1) {
    if (l&1) res += t[l++];
    if (r&1) res += t[--r];
  }
  return res;
}

// Range modification, single query.

void modify(int l, int r, int value) {
  for (l += n, r += n; l < r; l >>= 1, r >>= 1) {
    if (l&1) t[l++] += value;
    if (r&1) t[--r] += value;
  }
}

int query(int p) {
  int res = 0;
  for (p += n; p > 0; p >>= 1) res += t[p];
  return res;
}

/**
 * If at some point after modifications we need to inspect all the
 * elements in the array, we can push all the modifications to the
 * leaves using the following code. After that we can just traverse
 * elements starting with index n. This way we reduce the complexity
 * from O(n log(n)) to O(n) similarly to using build instead of n
    modifications.
 * */

void push() {
  for (int i = 1; i < n; ++i) {
    t[i<<1] += t[i];
    t[i<<1|1] += t[i];
    t[i] = 0;
  }
}

// Non commutative combiner functions.

void modify(int p, const S& value) {
  for (t[p += n] = value; p >>= 1; ) t[p] = combine(t[p<<1], t[p<<1|1]);
}

S query(int l, int r) {
  S resl, resr;
  for (l += n, r += n; l < r; l >>= 1, r >>= 1) {
    if (l&1) resl = combine(resl, t[l++]);
```

```
    if (r&1) resr = combine(t[--r], resr);
  }
  return combine(resl, resr);
}

// To be continued ...
```

2.6 sparse table

```
// RMQ.
const int MN = 100000 + 10; // Max number of elements
const int ML = 18; // ceil(log2(MN));

struct st {
  int data[MN];
  int M[MN][ML];
  int n;

  void read(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i)
      cin >> data[i];
  }

  void build() {
    for (int i = 0; i < n; ++i)
      M[i][0] = data[i];
    for (int j = 1, p = 2, q = 1; p <= n; ++j, p <<= 1, q <<= 1)
      for (int i = 0; i + p - 1 < n; ++i)
        M[i][j] = max(M[i][j - 1], M[i + q][j - 1]);
  }
  int query(int b, int e) {
    int k = log2(e - b + 1);
    return max(M[b][k], M[e + 1 - (1<<k)][k]);
  }
};
```

2.7 splay tree

```
using namespace std;
#include<bits/stdc++.h>
#define D(x) cout<<x<<endl;

typedef int T;

struct node{
  node *left, *right, *parent;
  T key;
  node (T k) : key(k), left(0), right(0), parent(0) {}
};

struct splay_tree{


  node *root;

  void right_rot(node *x) {
    node *p = x->parent;
    if (x->parent = p->parent) {
      if (x->parent->left == p) x->parent->left = x;
      if (x->parent->right == p) x->parent->right = x;
    }
    if (p->left = x->right) p->left->parent = p;
    x->right = p;
    p->parent = x;
  }

  void left_rot(node *x) {
    node *p = x->parent;
    if (x->parent = p->parent) {
      if (x->parent->left == p) x->parent->left = x;
      if (x->parent->right == p) x->parent->right = x;
    }
    if (p->right = x->left) p->right->parent = p;
    x->left = p;
    p->parent = x;
  }

  void splay(node *x, node *fa = 0) {

    while( x->parent != fa and x->parent != 0) {
      node *p = x->parent;
      if (p->parent == fa)
        if (p->right == x)
          left_rot(x);
        else
```

```
        right_rot(x);
    else {
      node *gp = p->parent; //grand parent
      if (gp->left == p)
        if (p->left == x)
          right_rot(x),right_rot(x);
        else
          left_rot(x),right_rot(x);
      else
        if (p->left == x)
          right_rot(x), left_rot(x);
        else
          left_rot(x), left_rot(x);
    }
  }
  if (fa == 0) root = x;
}

void insert(T key) {
  node *cur = root;
  node *pcur = 0;
  while (cur) {
    pcur = cur;
    if (key > cur->key) cur = cur->right;
    else cur = cur->left;
  }
  cur = new node(key);
  cur->parent = pcur;
  if (!pcur) root = cur;
  else if (key > pcur->key ) pcur->right = cur;
  else pcur->left = cur;
  splay(cur);
}

node *find(T key) {
  node *cur = root;
  while (cur) {
    if (key > cur->key) cur = cur->right;
    else if(key < cur->key) cur = cur->left;
    else return cur;
  }
  return 0;
}

splay_tree(){ root = 0;};
};
```

2.8 trie

```
const int MN = 26; // size of alphabet
const int MS = 100010; // Number of states.

struct trie{
  struct node{
    int c;
    int a[MN];
  };

  node tree[MS];
  int nodes;

  void clear(){
    tree[nodes].c = 0;
    memset(tree[nodes].a, -1, sizeof tree[nodes].a);
    nodes++;
  }

  void init(){
    nodes = 0;
    clear();
  }

  int add(const string &s, bool query = 0){
    int cur_node = 0;
    for(int i = 0; i < s.size(); ++i){
      int id = gid(s[i]);
      if(tree[cur_node].a[id] == -1){
        if(query) return 0;
        tree[cur_node].a[id] = nodes;
        clear();
      }
      cur_node = tree[cur_node].a[id];
    }
    if(!query) tree[cur_node].c++;
    return tree[cur_node].c;
  }

};
```

3 Geometry

3.1 closest pair problem

```
struct point {
  double x, y;
  int id;
  point() {}
  point (double a, double b) : x(a), y(b) {}
};

double dist(const point &o, const point &p) {
  double a = p.x - o.x, b = p.y - o.y;
  return sqrt(a * a + b * b);
}

double cp(vector<point> &p, vector<point> &x, vector<point> &y) {
  if (p.size() < 4) {
    double best = 1e100;
    for (int i = 0; i < p.size(); ++i)
      for (int j = i + 1; j < p.size(); ++j)
        best = min(best, dist(p[i], p[j]));
    return best;
  }

  int ls = (p.size() + 1) >> 1;
  double l = (p[ls - 1].x + p[ls].x) * 0.5;
  vector<point> xl(ls), xr(p.size() - ls);
  unordered_set<int> left;
  for (int i = 0; i < ls; ++i) {
    xl[i] = x[i];
    left.insert(x[i].id);
  }
  for (int i = ls; i < p.size(); ++i) {
    xr[i - ls] = x[i];
  }

  vector<point> yl, yr;
  vector<point> pl, pr;
  yl.reserve(ls); yr.reserve(p.size() - ls);
  pl.reserve(ls); pr.reserve(p.size() - ls);
  for (int i = 0; i < p.size(); ++i) {
    if (left.count(y[i].id))
      yl.push_back(y[i]);
    else
      yr.push_back(y[i]);

    if (left.count(p[i].id))
      pl.push_back(p[i]);
    else
      pr.push_back(p[i]);
  }

  double dl = cp(pl, xl, yl);
  double dr = cp(pr, xr, yr);
  double d = min(dl, dr);
  vector<point> yp; yp.reserve(p.size());
  for (int i = 0; i < p.size(); ++i) {
    if (fabs(y[i].x - l) < d)
      yp.push_back(y[i]);
  }
  for (int i = 0; i < yp.size(); ++i) {
    for (int j = i + 1; j < yp.size() && j < i + 7; ++j) {
      d = min(d, dist(yp[i], yp[j]));
    }
  }
  return d;
}

double closest_pair(vector<point> &p) {
  vector<point> x(p.begin(), p.end());
  sort(x.begin(), x.end(), [](const point &a, const point &b) {
    return a.x < b.x;
  });
  vector<point> y(p.begin(), p.end());
  sort(y.begin(), y.end(), [](const point &a, const point &b) {
    return a.y < b.y;
  });
  return cp(p, x, y);
}
```

3.2 squares

```
typedef long double ld;

const ld eps = 1e-12;
```

```
int cmp(ld x, ld y = 0, ld tol = eps) {
    return ( x <= y + tol) ? (x + tol < y) ? -1 : 0 : 1;
}

struct point{
  ld x, y;
  point(ld a, ld b) : x(a), y(b) {}
  point() {}
};

struct square{
  ld x1, x2, y1, y2,
     a, b, c;
  point edges[4];
  square(ld _a, ld _b, ld _c) {
    a = _a, b = _b, c = _c;
    x1 = a - c * 0.5;
    x2 = a + c * 0.5;
    y1 = b - c * 0.5;
    y2 = b + c * 0.5;
    edges[0] = point(x1, y1);
    edges[1] = point(x2, y1);
    edges[2] = point(x2, y2);
    edges[3] = point(x1, y2);
  }
};

ld min_dist(point &a, point &b) {
  ld x = a.x - b.x,
     y = a.y - b.y;
  return sqrt(x * x + y * y);
}

bool point_in_box(square s1, point p) {
  if (cmp(s1.x1, p.x) != 1 && cmp(s1.x2, p.x) != -1 &&
      cmp(s1.y1, p.y) != 1 && cmp(s1.y2, p.y) != -1)
    return true;
  return false;
}

bool inside(square &s1, square &s2) {
  for (int i = 0; i < 4; ++i)
    if (point_in_box(s2, s1.edges[i]))
      return true;
  return false;
}

bool inside_vert(square &s1, square &s2) {
  if ((cmp(s1.y1, s2.y1) != -1 && cmp(s1.y1, s2.y2) != 1) ||
      (cmp(s1.y2, s2.y1) != -1 && cmp(s1.y2, s2.y2) != 1))
    return true;
 return false;
}

bool inside_hori(square &s1, square &s2) {
  if ((cmp(s1.x1, s2.x1) != -1 && cmp(s1.x1, s2.x2) != 1) ||
      (cmp(s1.x2, s2.x1) != -1 && cmp(s1.x2, s2.x2) != 1))
    return true;
 return false;
}

ld min_dist(square &s1, square &s2) {
  if (inside(s1, s2) || inside(s2, s1))
    return 0;

  ld ans = 1e100;
  for (int i = 0; i < 4; ++i)
    for (int j = 0; j < 4; ++j)
      ans = min(ans, min_dist(s1.edges[i], s2.edges[j]));


  if (inside_hori(s1, s2) || inside_hori(s2, s1)) {
    if (cmp(s1.y1, s2.y2) != -1)
      ans = min(ans, s1.y1 - s2.y2);
    else
    if (cmp(s2.y1, s1.y2) != -1)
      ans = min(ans, s2.y1 - s1.y2);
  }

  if (inside_vert(s1, s2) || inside_vert(s2, s1)) {
    if (cmp(s1.x1, s2.x2) != -1)
      ans = min(ans, s1.x1 - s2.x2);
    else
    if (cmp(s2.x1, s1.x2) != -1)
      ans = min(ans, s2.x1 - s1.x2);
  }

  return ans;
}
```

3.3 triangles

Let a, b, c be length of the three sides of a triangle.

$$p = (a + b + c) * 0.5$$

The inradius is defined by:

$$iR = \sqrt{\frac{(p-a)(p-b)(p-c)}{p}}$$

The radius of its circumcircle is given by the formula:

$$cR = \frac{abc}{\sqrt{(a+b+c)(a+b-c)(a+c-b)(b+c-a)}}$$

4 Graphs

4.1 bridges

```
struct edge{
  int to, id;
  edge(int a, int b) : to(a), id(b) {}
};

struct graph {
  vector<vector<edge> > g;
  vector<int> vi, low, d, pi, is_b;

  int ticks, edges;

  graph(int n, int m) {
    g.assign(n, vector<edge>());
    is_b.assign(m, 0);
    vi.resize(n);
    low.resize(n);
    d.resize(n);
    pi.resize(n);
    edges = 0;
  }

  void add_edge(int u, int v) {
    g[u].push_back(edge(v, edges));
    g[v].push_back(edge(u, edges));
    edges++;
  }

  void dfs(int u) {
    vi[u] = true;
    d[u] = low[u] = ticks++;
    for (int i = 0; i < g[u].size(); ++i) {
      int v = g[u][i].to;
      if (v == pi[u]) continue;
      if (!vi[v]) {
        pi[v] = u;
        dfs(v);
        if (d[u] < low[v])
          is_b[g[u][i].id] = true;


        low[u] = min(low[u], low[v]);
      } else {
        low[u] = min(low[u], d[v]);
      }
    }
  }

  // Multiple edges from a to b are not allowed.
  // (they could be detected as a bridge).
  // If you need to handle this, just count
  // how many edges there are from a to b.
  void comp_bridges() {
    fill(pi.begin(), pi.end(), -1);
    fill(vi.begin(), vi.end(), 0);
    fill(low.begin(), low.end(), 0);
    fill(d.begin(), d.end(), 0);
    ticks = 0;
    for (int i = 0; i < g.size(); ++i)
      if (!vi[i]) dfs(i);
  }
};
```

4.2 directed mst

```
const int inf = 1000000 + 10;

struct edge {
```

```cpp
  int u, v, w;
  edge() {}
  edge(int a,int b,int c) : u(a), v(b), w(c) {}
};

/**
 * Computes the minimum spanning tree for a directed graph
 * - edges : Graph description in the form of list of edges.
 *    each edge is: From node u to node v with cost w
 * - root : Id of the node to start the DMST.
 * - n    : Number of nodes in the graph.
 * */

int dmst(vector<edge> &edges, int root, int n) {
  int ans = 0;
  int cur_nodes = n;
  while (true) {
    vector<int> lo(cur_nodes, inf), pi(cur_nodes, inf);
    for (int i = 0; i < edges.size(); ++i) {
      int u = edges[i].u, v = edges[i].v, w = edges[i].w;
      if (w < lo[v] and u != v) {
        lo[v] = w;
        pi[v] = u;
      }
    }

    lo[root] = 0;
    for (int i = 0; i < lo.size(); ++i) {
      if (i == root) continue;
      if (lo[i] == inf) return -1;
    }
    int cur_id = 0;
    vector<int> id(cur_nodes, -1), mark(cur_nodes, -1);
    for (int i = 0; i < cur_nodes; ++i) {
      ans += lo[i];
      int u = i;
      while (u != root and id[u] < 0 and mark[u] != i) {
        mark[u] = i;
        u = pi[u];
      }
      if (u != root and id[u] < 0) { // Cycle
         for (int v = pi[u]; v != u; v = pi[v])
           id[v] = cur_id;
         id[u] = cur_id++;
      }
    }

    if (cur_id == 0)
      break;

    for (int i = 0; i < cur_nodes; ++i)
      if (id[i] < 0) id[i] = cur_id++;

    for (int i = 0; i < edges.size(); ++i) {
      int u = edges[i].u, v = edges[i].v, w = edges[i].w;
      edges[i].u = id[u];
      edges[i].v = id[v];
      if (id[u] != id[v])
        edges[i].w -= lo[v];
    }
    cur_nodes = cur_id;
    root = id[root];
  }

  return ans;
}
```

4.3 eulerian path

```cpp
// Taken from https://github.com/lbv/pc-code/blob/master/code/graph.cpp
// Eulerian Trail

struct Euler {
  ELV adj; IV t;
  Euler(ELV Adj) : adj(Adj) {}
  void build(int u) {
    while(! adj[u].empty()) {
      int v = adj[u].front().v;
      adj[u].erase(adj[u].begin());
      build(v);
    }
    t.push_back(u);
  }
};
bool eulerian_trail(IV &trail) {
  Euler e(adj);
  int odd = 0, s = 0;
  /*
```

```
    for (int v = 0; v < n; v++) {
    int diff = abs(in[v] - out[v]);
    if (diff > 1) return false;
    if (diff == 1) {
    if (++odd > 2) return false;
    if (out[v] > in[v]) start = v;
    }
    }
    */
  e.build(s);
  reverse(e.t.begin(), e.t.end());
  trail = e.t;
  return true;
}
```

4.4 karp min mean cycle

```
/**
 * Finds the min mean cycle, if you need the max mean cycle
 * just add all the edges with negative cost and print
 * ans * -1
 *
 * test: uva, 11090 - Going in Cycle!!
 * */

const int MN = 1000;
struct edge{
  int v;
  long long w;
  edge(){} edge(int v, int w) : v(v), w(w) {}
};

long long d[MN][MN];
// This is a copy of g because increments the size
// pass as reference if this does not matter.
int karp(vector<vector<edge> > g) {
  int n = g.size();

  g.resize(n + 1); // this is important

  for (int i = 0; i < n; ++i)
    if (!g[i].empty())
      g[n].push_back(edge(i,0));
  ++n;

  for(int i = 0;i<n;++i)
    fill(d[i],d[i]+(n+1),INT_MAX);

  d[n - 1][0] = 0;

  for (int k = 1; k <= n; ++k) for (int u = 0; u < n; ++u) {
    if (d[u][k - 1] == INT_MAX) continue;
    for (int i = g[u].size() - 1; i >= 0; --i)
      d[g[u][i].v][k] = min(d[g[u][i].v][k], d[u][k - 1] + g[u][i].w);
  }

  bool flag = true;

  for (int i = 0; i < n && flag; ++i)
    if (d[i][n] != INT_MAX)
      flag = false;

  if (flag) {
    return true; // return true if there is no a cycle.
  }

  double ans = 1e15;

  for (int u = 0; u + 1 < n; ++u) {
    if (d[u][n] == INT_MAX) continue;
    double W = -1e15;

    for (int k = 0; k < n; ++k)
      if (d[u][k] != INT_MAX)
        W = max(W, (double)(d[u][n] - d[u][k]) / (n - k));

    ans = min(ans, W);
  }

  // printf("%.2lf\n", ans);
  cout << fixed << setprecision(2) << ans << endl;

  return false;
}
```

4.5 konig's theorem

In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover

4.6 minimum path cover in DAG

Given a directed acyclic graph $G = (V, E)$, we are to find the minimum number of vertex-disjoint paths to cover each vertex in V.

We can construct a bipartite graph $G' = (Vout \cup Vin, E')$ from G, where :

$$Vout = \{v \in V : v \ has \ positive \ out - degree\}$$

$$Vin = \{v \in V : v \ has \ positive \ in - degree\}$$

$$E' = \{(u, v) \in Vout \times Vin : (u, v) \in E\}$$

Then it can be shown, via König's theorem, that G' has a matching of size m if and only if there exists $n - m$ vertex-disjoint paths that cover each vertex in G, where n is the number of vertices in G and m is the maximum cardinality bipartite mathching in G'.

Therefore, the problem can be solved by finding the maximum cardinality matching in G' instead.

NOTE: If the paths are note necesarily disjoints, find the transitive closure and solve the problem for disjoint paths.

4.7 planar graph (euler)

Euler's formula states that if a finite, connected, planar graph is drawn in the plane without any edge intersections, and v is the number of vertices, e is the number of edges and f is the number of faces (regions bounded by edges, including the outer, infinitely large region), then:

$$f + v = e + 2$$

It can be extended to non connected planar graphs with c connected components:

$$f + v = e + c + 1$$

4.8 tarjan scc

```cpp
const int MN = 20002;

struct tarjan_scc {
  int scc[MN], low[MN], d[MN], stacked[MN];
  int ticks, current_scc;
  deque<int> s; // used as stack.

  tarjan_scc() {}

  void init () {
    memset(scc, -1, sizeof scc);
    memset(d, -1, sizeof d);
    memset(stacked, 0, sizeof stacked);
    s.clear();
    ticks = current_scc = 0;
  }

  void compute(vector<vector<int> > &g, int u) {
    d[u] = low[u] = ticks++;
    s.push_back(u);
    stacked[u] = true;
    for (int i = 0; i < g[u].size(); ++i) {
      int v = g[u][i];
      if (d[v] == -1)
        compute(g, v);
      if (stacked[v]) {
        low[u] = min(low[u], low[v]);
      }
    }

    if (d[u] == low[u]) { // root
      int v;
      do {
        v = s.back();s.pop_back();
        stacked[v] = false;
        scc[v] = current_scc;
      } while (u != v);
      current_scc++;
    }
  }
};
```

4.9 two sat (with kosaraju)

```
/**
 *  Given a set of clauses (a1 v a2)^(a2 v a3)....
 *  this algorithm find a solution to it set of clauses.
 *  test: http://lightoj.com/volume_showproblem.php?problem=1251
 **/

#include<bits/stdc++.h>
using namespace std;
#define MAX 100000
#define endl '\n'

vector<int> G[MAX];
vector<int> GT[MAX];
vector<int> Ftime;
vector<vector<int> > SCC;
bool visited[MAX];
int n;


void dfs1(int n){
  visited[n] = 1;

  for (int i = 0; i < G[n].size(); ++i) {
    int curr = G[n][i];
    if (visited[curr]) continue;
    dfs1(curr);
  }

  Ftime.push_back(n);
}

void dfs2(int n, vector<int> &scc) {
  visited[n] = 1;
  scc.push_back(n);

  for (int i = 0;i < GT[n].size(); ++i) {
    int curr = GT[n][i];
    if (visited[curr]) continue;
    dfs2(curr, scc);
  }
}

void kosaraju() {
  memset(visited, 0, sizeof visited);

  for (int i = 0; i < 2 * n ; ++i) {
    if (!visited[i]) dfs1(i);
  }

  memset(visited, 0, sizeof visited);
  for (int i = Ftime.size() - 1; i >= 0; i--) {
    if (visited[Ftime[i]]) continue;
    vector<int> _scc;
    dfs2(Ftime[i],_scc);
    SCC.push_back(_scc);
  }
}

/**
 * After having the SCC, we must traverse each scc, if in one SCC are -b
     y b, there is not a solution.
 * Otherwise we build a solution, making the first "node" that we find
     truth and its complement false.
 **/


bool two_sat(vector<int> &val) {
  kosaraju();
  for (int i = 0; i < SCC.size(); ++i) {
    vector<bool> tmpvisited(2 * n, false);
    for (int j = 0; j < SCC[i].size(); ++j) {
      if (tmpvisited[SCC[i][j] ^ 1]) return 0;
      if (val[SCC[i][j]] != -1) continue;
      else {
        val[SCC[i][j]] = 0;
        val[SCC[i][j] ^ 1] = 1;
      }
      tmpvisited[SCC[i][j]] = 1;
    }
  }
  return 1;
}

// Example of use

int main() {
```

```cpp
  int m, u, v, nc = 0, t; cin >> t;
  // n = "nodes" number, m = clauses number

  while (t--) {
    cin >> m >> n;
    Ftime.clear();
    SCC.clear();
    for (int i = 0; i < 2 * n; ++i) {
      G[i].clear();
      GT[i].clear();
    }

    // (a1 v a2) = (a1 -> a2) = (a2 -> a1)
    for (int i = 0; i < m ; ++i) {
      cin >> u >> v;
      int t1 = abs(u) - 1;
      int t2 = abs(v) - 1;
      int p = t1 * 2 + ((u < 0)? 1 : 0);
      int q = t2 * 2 + ((v < 0)? 1 : 0);
      G[p ^ 1].push_back(q);
      G[q ^ 1].push_back(p);
      GT[p].push_back(q ^ 1);
      GT[q].push_back(p ^ 1);
    }

    vector<int> val(2 * n, -1);
    cout << "Case " << ++nc <<": ";
    if (two_sat(val)) {
      cout << "Yes" << endl;
      vector<int> sol;
      for (int i = 0; i < 2 * n; ++i)
        if (i % 2 == 0 and val[i] == 1)
          sol.push_back(i / 2 + 1);
      cout << sol.size() ;

      for (int i = 0; i < sol.size(); ++i) {
        cout << " " << sol[i];
      }
      cout << endl;
    } else {
      cout << "No" << endl;
    }
  }
  return 0;
}
```

5 Math

5.1 cumulative sum of divisors

```cpp
/**
The function SOD(n) (sum of divisors) is defined
as the summation of all the actual divisors of
an integer number n. For example,

  SOD(24) = 2+3+4+6+8+12 = 35.

The function CSOD(n) (cumulative SOD) of an integer n, is defined as
    below:

  csod(n) = \sum_{i = 1}^{n} sod(i)

It can be computed in O(sqrt(n)):
*/

long long csod(long long n) {
  long long ans = 0;
  for (long long i = 2; i * i <= n; ++i) {
    long long j = n / i;
    ans += (i + j) * (j - i + 1) / 2;
    ans += i * (j - i);
  }
  return ans;
}
```

5.2 fft

```cpp
/**
 * Fast Fourier Transform.
 * Useful to compute convolutions.
 * computes:
 *    C(f star g)[n] = sum_m(f[m] * g[n - m])
 * for all n.
 * test: icpc live archive, 6886 - Golf Bot
 * */


using namespace std;
```

```cpp
#include<bits/stdc++.h>
#define D(x) cout << #x " = " << (x) << endl
#define endl '\n'

const int MN = 262144 << 1;
int d[MN + 10], d2[MN + 10];


const double PI = acos(-1.0);

struct cpx {
  double real, image;
  cpx(double _real, double _image) {
    real = _real;
    image = _image;
  }
  cpx(){}
};

cpx operator + (const cpx &c1, const cpx &c2) {
  return cpx(c1.real + c2.real, c1.image + c2.image);
}

cpx operator - (const cpx &c1, const cpx &c2) {
  return cpx(c1.real - c2.real, c1.image - c2.image);
}

cpx operator * (const cpx &c1, const cpx &c2) {
  return cpx(c1.real*c2.real - c1.image*c2.image, c1.real*c2.image +
      c1.image*c2.real);
}

int rev(int id, int len) {
  int ret = 0;
  for (int i = 0; (1 << i) < len; i++) {
    ret <<= 1;
    if (id & (1 << i)) ret |= 1;
  }
  return ret;
}

cpx A[1 << 20];

void FFT(cpx *a, int len, int DFT) {
  for (int i = 0; i < len; i++)
    A[rev(i, len)] = a[i];
  for (int s = 1; (1 << s) <= len; s++) {
    int m = (1 << s);
    cpx wm = cpx(cos( DFT * 2 * PI / m), sin(DFT * 2 * PI / m));
    for(int k = 0; k < len; k += m) {
      cpx w = cpx(1, 0);
      for(int j = 0; j < (m >> 1); j++) {
        cpx t = w * A[k + j + (m >> 1)];
        cpx u = A[k + j];
        A[k + j] = u + t;
        A[k + j + (m >> 1)] = u - t;
        w = w * wm;
      }
    }
  }
  if (DFT == -1) for (int i = 0; i < len; i++) A[i].real /= len,
      A[i].image /= len;
  for (int i = 0; i < len; i++) a[i] = A[i];
  return;
}

cpx in[1 << 20];

void solve(int n) {
  memset(d, 0, sizeof d);
  int t;
  for (int i = 0; i < n; ++i) {
    cin >> t;
    d[t] = true;
  }
  int m;
  cin >> m;
  vector<int> q(m);
  for (int i = 0; i < m; ++i)
    cin >> q[i];

  for (int i = 0; i < MN; ++i) {
    if (d[i])
      in[i] = cpx(1, 0);
    else
      in[i] = cpx(0, 0);
  }

  FFT(in, MN, 1);
  for (int i = 0; i < MN; ++i) {
```

```
    in[i] = in[i] * in[i];
  }
  FFT(in, MN, -1);

  int ans = 0;
  for (int i = 0; i < q.size(); ++i) {
    if (in[q[i]].real > 0.5 || d[q[i]]) {
      ans++;
    }
  }
  cout << ans << endl;
}

int main() {
  ios_base::sync_with_stdio(false);cin.tie(NULL);
  int n;
  while (cin >> n)
    solve(n);
  return 0;
}
```

5.3 fibonacci properties

Let A, B and n be integer numbers.

$$k = A - B \tag{1}$$

$$F_A F_B = F_{k+1} F_A^2 + F_k F_A F_{A-1} \tag{2}$$

$$\sum_{i=0}^{n} F_i^2 = F_{n+1} F_n \tag{3}$$

$ev(n)$ = returns 1 if n is even.

$$\sum_{i=0}^{n} F_i F_{i+1} = F_{n+1}^2 - ev(n) \tag{4}$$

$$\sum_{i=0}^{n} F_i F_{i-1} = \sum_{i=0}^{n-1} F_i F_{i+1} \tag{5}$$

6 Matrix

6.1 matrix

```
const int MN = 111;
const int mod = 10000;

struct matrix {
  int r, c;
  int m[MN][MN];

  matrix (int _r, int _c) : r (_r), c (_c) {
    memset(m, 0, sizeof m);
  }

  void print() {
    for (int i = 0; i < r; ++i) {
      for (int j = 0; j < c; ++j)
        cout << m[i][j] << " ";
      cout << endl;
    }
  }

  int x[MN][MN];
  matrix & operator *= (const matrix &o) {
    memset(x, 0, sizeof x);
    for (int i = 0; i < r; ++i)
      for (int k = 0; k < c; ++k)
        if (m[i][k] != 0)
          for (int j = 0; j < c; ++j) {
            x[i][j] = (x[i][j] + ((m[i][k] * o.m[k][j]) % mod) ) % mod;
          }
    memcpy(m, x, sizeof(m));
    return *this;
  }
};

void matrix_pow(matrix b, long long e, matrix &res) {
  memset(res.m, 0, sizeof res.m);
  for (int i = 0; i < b.r; ++i)
    res.m[i][i] = 1;

  if (e == 0) return;
  while (true) {
```

```
    if (e & 1) res *= b;
    if ((e >>= 1) == 0) break;
    b *= b;
  }
}
```

7 Misc

7.1 Template Java

```
import java.io.*;
import java.util.StringTokenizer;

public class Template {

    public static void main(String []args) throws IOException {
        Scanner in = new Scanner(System.in);
        OutputWriter out = new OutputWriter(System.out);
        Task solver = new Task();
        solver.solve(in, out);
        out.close();
    }

}

class Task{
    public void solve(Scanner in, OutputWriter out){

    }
}

class Scanner{
    public BufferedReader reader;
    public StringTokenizer st;

    public Scanner(InputStream stream){
        reader = new BufferedReader(new InputStreamReader(stream));
        st = null;
    }

    public String next(){
        while(st == null || !st.hasMoreTokens()){
            try{
                String line = reader.readLine();
                if(line == null) return null;
                st = new StringTokenizer(line);
            }catch (Exception e){
                throw (new RuntimeException());
            }
        }
        return st.nextToken();
    }

    public int nextInt(){
        return Integer.parseInt(next());
    }
    public long nextLong(){
        return Long.parseLong(next());
    }
    public double nextDouble(){
        return Double.parseDouble(next());
    }
}

class OutputWriter{
    BufferedWriter writer;

    public OutputWriter(OutputStream stream){
        writer = new BufferedWriter(new OutputStreamWriter(stream));
    }

    public void print(int i) throws IOException {
        writer.write(i);
    }

    public void print(String s) throws IOException {
        writer.write(s);
    }

    public void print(char []c) throws IOException {
        writer.write(c);
    }
    public void close() throws IOException {
        writer.close();
    }

}
```

7.2 fraction

```
struct frac{
  long long x, y;
  frac(long long a, long long b) {
    long long g = __gcd(a, b);
    x = a / g;
    y = b / g;
  }
  bool operator < (const frac &o) const {
    return (x * o.y < y * o.x);
  }
};
```

7.3 io

```
// taken from :
    https://github.com/lbv/pc-code/blob/master/solved/c-e/diablo/diablo.cpp
// this is very fast as well :
    https://github.com/lbv/pc-code/blob/master/code/input.cpp

typedef unsigned int u32;
#define BUF 524288
struct Reader {
  char buf[BUF]; char b; int bi, bz;
  Reader() { bi=bz=0; read(); }
  void read() {
    if (bi==bz) { bi=0; bz = fread(buf, 1, BUF, stdin); }
    b = bz ? buf[bi++] : 0; }
  void skip() { while (b > 0 && b <= 32) read(); }
  u32 next_u32() {
    u32 v = 0; for (skip(); b > 32; read()) v = v*10 + b-48; return v; }
  int next_int() {
    int v = 0; bool s = false;
    skip(); if (b == '-') { s = true; read(); }
    for (; 48<=b&&b<=57; read()) v = v*10 + b-48; return s ? -v : v; }
  char next_char() { skip(); char c = b; read(); return c; }
};
```

8 Number theory

8.1 convolution

```
typedef long long int LL;
typedef pair<LL, LL> PLL;

inline bool is_pow2(LL x) {
  return (x & (x-1)) == 0;
}

inline int ceil_log2(LL x) {
  int ans = 0;
  --x;
  while (x != 0) {
    x >>= 1;
    ans++;
  }
  return ans;
}

/* Returns the convolution of the two given vectors in time proportional
    to n*log(n).
 * The number of roots of unity to use nroots_unity must be set so that
     the product of the first
 * nroots_unity primes of the vector nth_roots_unity is greater than the
     maximum value of the
 * convolution. Never use sizes of vectors bigger than 2^24, if you need
     to change the values of
 * the nth roots of unity to appropriate primes for those sizes.
 */
vector<LL> convolve(const vector<LL> &a, const vector<LL> &b, int
    nroots_unity = 2) {
  int N = 1 << ceil_log2(a.size() + b.size());
  vector<LL> ans(N,0), fA(N), fB(N), fC(N);
  LL modulo = 1;
  for (int times = 0; times < nroots_unity; times++) {
    fill(fA.begin(), fA.end(), 0);
    fill(fB.begin(), fB.end(), 0);
    for (int i = 0; i < a.size(); i++) fA[i] = a[i];
    for (int i = 0; i < b.size(); i++) fB[i] = b[i];
    LL prime = nth_roots_unity[times].first;
    LL inv_modulo = mod_inv(modulo % prime, prime);
    LL normalize = mod_inv(N, prime);
```

```
    ntfft(fA, 1, nth_roots_unity[times]);
    ntfft(fB, 1, nth_roots_unity[times]);
    for (int i = 0; i < N; i++) fC[i] = (fA[i] * fB[i]) % prime;
    ntfft(fC, -1, nth_roots_unity[times]);
    for (int i = 0; i < N; i++) {
      LL curr = (fC[i] * normalize) % prime;
      LL k = (curr - (ans[i] % prime) + prime) % prime;
      k = (k * inv_modulo) % prime;
      ans[i] += modulo * k;
    }
    modulo *= prime;
  }
  return ans;
}
```

8.2 crt

```
/**
 * Chinese remainder theorem.
 * Find z such that z % x[i] = a[i] for all i.
 * */
long long crt(vector<long long> &a, vector<long long> &x) {
  long long z = 0;
  long long n = 1;
  for (int i = 0; i < x.size(); ++i)
    n *= x[i];

  for (int i = 0; i < a.size(); ++i) {
    long long tmp = (a[i] * (n / x[i])) % n;
    tmp = (tmp * mod_inv(n / x[i], x[i])) % n;
    z = (z + tmp) % n;
  }

  return (z + n) % n;
}
```

8.3 discrete logarithm

```
// Computes x which a ^ x = b mod n.
long long d_log(long long a, long long b, long long n) {
  long long m = ceil(sqrt(n));
  long long aj = 1;
  map<long long, long long> M;
  for (int i = 0; i < m; ++i) {
    if (!M.count(aj))
      M[aj] = i;
    aj = (aj * a) % n;
  }

  long long coef = mod_pow(a, n - 2, n);
  coef = mod_pow(coef, m, n);
  // coef = a ^ (-m)
  long long gamma = b;
  for (int i = 0; i < m; ++i) {
    if (M.count(gamma)) {
      return i * m + M[gamma];
    } else {
      gamma = (gamma * coef) % n;
    }
  }
  return -1;
}
```

8.4 ext euclidean

```
void ext_euclid(long long a, long long b, long long &x, long long &y,
    long long &g) {
  x = 0, y = 1, g = b;
  long long m, n, q, r;
  for (long long u = 1, v = 0; a != 0; g = a, a = r) {
    q = g / a, r = g % a;
    m = x - u * q, n = y - v * q;
    x = u, y = v, u = m, v = n;
  }
}
```

8.5 highest exponent factorial

```
int highest_exponent(int p, const int &n){
  int ans = 0;
```

```
  int t = p;
  while(t <= n){
    ans += n/t;
    t*=p;
  }
  return ans;
}
```

8.6 miller rabin

```
const int rounds = 20;

// checks whether a is a witness that n is not prime, 1 < a < n
bool witness(long long a, long long n) {
  // check as in Miller Rabin Primality Test described
  long long u = n - 1;
  int t = 0;
  while (u % 2 == 0) {
    t++;
    u >>= 1;
  }
  long long next = mod_pow(a, u, n);
  if (next == 1) return false;
  long long last;
  for (int i = 0; i < t; ++i) {
    last = next;
    next = mod_mul(last, last, n);
    if (next == 1) {
      return last != n - 1;
    }
  }
  return next != 1;
}


// Checks if a number is prime with prob 1 - 1 / (2 ^ it)
//  D(miller_rabin(999999999999997LL) == 1);
//  D(miller_rabin(9999999999971LL) == 1);
//  D(miller_rabin(7907) == 1);
bool miller_rabin(long long n, int it = rounds) {
  if (n <= 1) return false;
  if (n == 2) return true;
  if (n % 2 == 0) return false;
  for (int i = 0; i < it; ++i) {
    long long a = rand() % (n - 1) + 1;
    if (witness(a, n)) {
      return false;
    }
  }
  return true;
}
```

8.7 mod inv

```
long long mod_inv(long long n, long long m) {
  long long x, y, gcd;
  ext_euclid(n, m, x, y, gcd);
  if (gcd != 1)
    return 0;
  return (x + m) % m;
}
```

8.8 mod mul

```
// Computes (a * b) % mod
long long mod_mul(long long a, long long b, long long mod) {
  long long x = 0, y = a % mod;
  while (b > 0) {
    if (b & 1)
      x = (x + y) % mod;
    y = (y * 2) % mod;
    b /= 2;
  }
  return x % mod;
}
```

8.9 mod pow

```
// Computes ( a ^ exp ) % mod.
long long mod_pow(long long a, long long exp, long long mod) {
  long long ans = 1;
  while (exp > 0) {
```

```
    if (exp & 1)
      ans = mod_mul(ans, a, mod);
    a = mod_mul(a, a, mod);
    exp >>= 1;
  }
  return ans;
}
```

8.10 number theoretic transform

```
typedef long long int LL;
typedef pair<LL, LL> PLL;

/* The following vector of pairs contains pairs (prime, generator)
 * where the prime has an Nth root of unity for N being a power of two.
 * The generator is a number g s.t g^(p-1)=1 (mod p)
 * but is different from 1 for all smaller powers */
vector<PLL> nth_roots_unity {
  {1224736769,330732430},{1711276033,927759239},{167772161,167489322},
   {469762049,343261969},{754974721,643797295},{1107296257,883865065}};

PLL ext_euclid(LL a, LL b) {
  if (b == 0)
    return make_pair(1,0);
  pair<LL,LL> rc = ext_euclid(b, a % b);
  return make_pair(rc.second, rc.first - (a / b) * rc.second);
}

//returns -1 if there is no unique modular inverse
LL mod_inv(LL x, LL modulo) {
  PLL p = ext_euclid(x, modulo);
  if ( (p.first * x + p.second * modulo) != 1 )
    return -1;
  return (p.first+modulo) % modulo;
}


//Number theory fft. The size of a must be a power of 2
void ntfft(vector<LL> &a, int dir, const PLL &root_unity) {
  int n = a.size();
  LL prime = root_unity.first;
  LL basew = mod_pow(root_unity.second, (prime-1) / n, prime);
  if (dir < 0) basew = mod_inv(basew, prime);
  for (int m = n; m >= 2; m >>= 1) {
    int mh = m >> 1;
    LL w = 1;
    for (int i = 0; i < mh; i++) {
      for (int j = i; j < n; j += m) {
        int k = j + mh;
        LL x = (a[j] - a[k] + prime) % prime;
        a[j] = (a[j] + a[k]) % prime;
        a[k] = (w * x) % prime;
      }
      w = (w * basew) % prime;
    }
    basew = (basew * basew) % prime;
  }
  int i = 0;
  for (int j = 1; j < n - 1; j++) {
    for (int k = n >> 1; k > (i ^= k); k >>= 1);
    if (j < i) swap(a[i], a[j]);
  }
}
```

8.11 pollard rho factorize

```
long long pollard_rho(long long n) {
  long long x, y, i = 1, k = 2, d;
  x = y = rand() % n;
  while (1) {
    ++i;
    x = mod_mul(x, x, n);
    x += 2;
    if (x >= n) x -= n;
    if (x == y) return 1;
    d = __gcd(abs(x - y), n);
    if (d != 1) return d;
    if (i == k) {
      y = x;
      k *= 2;
    }
  }
  return 1;
}
```

```
// Returns a list with the prime divisors of n
vector<long long> factorize(long long n) {
  vector<long long> ans;
  if (n == 1)
    return ans;
  if (miller_rabin(n)) {
    ans.push_back(n);
  } else {
    long long d = 1;
    while (d == 1)
      d = pollard_rho(n);
    vector<long long> dd = factorize(d);
    ans = factorize(n / d);
    for (int i = 0; i < dd.size(); ++i)
      ans.push_back(dd[i]);
  }
  return ans;
}
```

8.12 totient sieve

```
  for (int i = 1; i < MN; i++)
    phi[i] = i;

  for (int i = 1; i < MN; i++)
    if (!sieve[i]) // is prime
      for (int j = i; j < MN; j += i)
        phi[j] -= phi[j] / i;
```

8.13 totient

```
long long totient(long long n) {
  if (n == 1) return 0;
  long long ans = n;
  for (int i = 0; primes[i] * primes[i] <= n; ++i) {
    if ((n % primes[i]) == 0) {
      while ((n % primes[i]) == 0) n /= primes[i];
      ans -= ans / primes[i];
    }
  }
  if (n > 1) {
    ans -= ans / n;
  }
  return ans;
}
```

9 Strings

9.1 minimal string rotation

```
// Lexicographically minimal string rotation
int lmsr() {
  string s;
  cin >> s;
  int n = s.size();
  s += s;
  vector<int> f(s.size(), -1);
  int k = 0;
  for (int j = 1; j < 2 * n; ++j) {
    int i = f[j - k - 1];
    while (i != -1 && s[j] != s[k + i + 1]) {
      if (s[j] < s[k + i + 1])
        k = j - i - 1;
      i = f[i];
    }
    if (i == -1 && s[j] != s[k + i + 1]) {
      if (s[j] < s[k + i + 1]) {
        k = j;
      }
      f[j - k] = -1;
    } else {
      f[j - k] = i + 1;
    }
  }
  return k;
}
```

9.2 suffix array

```
/**
 * O (n log^2 (n))
```

```
 * See http://web.stanford.edu/class/cs97si/suffix-array.pdf for reference
 * */

struct entry{
  int a, b, p;
  entry(){}
  entry(int x, int y, int z): a(x), b(y), p(z){}
  bool operator < (const entry &o) const {
    return (a == o.a) ? (b == o.b) ? ( p < o.p) : (b < o.b) : (a < o.a);
  }
};

struct SuffixArray{
  const int N;
  string s;
  vector<vector<int> > P;
  vector<entry> M;

  SuffixArray(const string &s) : N(s.length()) , s(s), P(1, vector<int>
      (N, 0)), M(N) {
    for (int i = 0; i < N; ++i)
      P[0][i] = (int) s[i];

    for (int skip = 1, level = 1; skip < N; skip *= 2, level++) {
      P.push_back(vector<int>(N, 0));
      for (int i = 0 ; i < N; ++i) {
        int next = ((i + skip) < N) ? P[level - 1][i + skip] : -10000;
        M[i] = entry(P[level - 1][i], next, i);
      }
      sort(M.begin(), M.end());
      for (int i = 0; i < N; ++i)
        P[level][M[i].p] = (i > 0 and M[i].a == M[i - 1].a and M[i].b ==
            M[i - 1].b) ? P[level][M[i - 1].p] : i;
    }
  }

  vector<int> getSuffixArray(){
    vector<int> &rank = P.back();
    vector<pair<int, int> > inv(rank.size());
    for (int i = 0; i < rank.size(); ++i)
      inv[i] = make_pair(rank[i], i);
    sort(inv.begin(), inv.end());
    vector<int> sa(rank.size());
    for (int i = 0; i < rank.size(); ++i)
      sa[i] = inv[i].second;
    return sa;
  }

  // returns the length of the longest common prefix of s[i...L-1] and
      s[j...L-1]
  int lcp(int i, int j) {
    int len = 0;
    if (i == j) return N - i;
    for (int k = P.size() - 1; k >= 0 && i < N && j < N; --k) {
      if (P[k][i] == P[k][j]) {
        i += 1 << k;
        j += 1 << k;
        len += 1 << k;
      }
    }
    return len;
  }
};
```

9.3 suffix automaton

```
/*
 * Suffix automaton:
 * This implementation was extended to maintain (online) the
 * number of different substrings. This is equivalent to compute
 * the number of paths from the initial state to all the other
 * states.
 *
 * The overall complexity is O(n)
 * can be tested here:
     https://www.urionlinejudge.com.br/judge/en/problems/view/1530
 * */

struct state {
  int len, link;
  long long num_paths;
  map<int, int> next;
};

const int MN = 200011;
state sa[MN << 1];
int sz, last;
long long tot_paths;
```

```cpp
void sa_init() {
  sz = 1;
  last = 0;
  sa[0].len = 0;
  sa[0].link = -1;
  sa[0].next.clear();
  sa[0].num_paths = 1;
  tot_paths = 0;
}

void sa_extend(int c) {
  int cur = sz++;
  sa[cur].len = sa[last].len + 1;
  sa[cur].next.clear();
  sa[cur].num_paths = 0;
  int p;
  for (p = last; p != -1 && !sa[p].next.count(c); p = sa[p].link) {
    sa[p].next[c] = cur;
    sa[cur].num_paths += sa[p].num_paths;
    tot_paths += sa[p].num_paths;
  }

  if (p == -1) {
    sa[cur].link = 0;
  } else {
    int q = sa[p].next[c];
    if (sa[p].len + 1 == sa[q].len) {
      sa[cur].link = q;
    } else {
      int clone = sz++;
      sa[clone].len = sa[p].len + 1;
      sa[clone].next = sa[q].next;
      sa[clone].num_paths = 0;
      sa[clone].link = sa[q].link;
      for (; p!= -1 && sa[p].next[c] == q; p = sa[p].link) {
        sa[p].next[c] = clone;
        sa[q].num_paths -= sa[p].num_paths;
        sa[clone].num_paths += sa[p].num_paths;
      }
      sa[q].link = sa[cur].link = clone;
    }
  }
  last = cur;
}
```

9.4 z algorithm

```cpp
using namespace std;
#include<bits/stdc++.h>

vector<int> compute_z(const string &s){
  int n = s.size();
  vector<int> z(n,0);
  int l,r;
  r = l = 0;
  for(int i = 1; i < n; ++i){
    if(i > r) {
      l = r = i;
      while(r < n and s[r - l] == s[r])r++;
      z[i] = r - l;r--;
    }else{
      int k = i-l;
      if(z[k] < r - i +1) z[i] = z[k];
      else {
        l = i;
        while(r < n and s[r - l] == s[r])r++;
        z[i] = r - l;r--;
      }
    }
  }
  return z;
}

int main(){

  //string line;cin>>line;
  string line = "alfalfa";
  vector<int> z = compute_z(line);

  for(int i = 0; i < z.size(); ++i ){
    if(i)cout<<" ";
    cout<<z[i];
  }
  cout<<endl;

  // must print "0 0 0 4 0 0 1"

  return 0;
}
```