# Team notebook

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# 1 Algorithms

# 1.1 Mo's algorithm on trees

```
void flat(vector<vector<edge>> &g, vector<int> &a,
    vector<int> &le, vector<int> &ri, vector<int> &cost,
    int node, int pi, int &ts, int w) {

    cost[node] = w;
    le[node] = ts;
    a[ts] = node;
    ts++;
    for (auto e : g[node]) {
        if (e.to == pi) continue;
        flat(g, a, le, ri, cost, e.to, node, ts, e.w);
    }
    ri[node] = ts;
    a[ts] = node;
    ts++;
}

/**
```

```
* Case when the cost is in the edges.
void compute_queries(vector<vector<edge>> &g) {
 // g is undirected
 int n = g.size();
 lca_tree.init(g, 0);
 vector\langle int \rangle a(2 * n), le(n), ri(n), cost(n);
 // a: nodes in the flatten array
 // le: left id of the given node
 // ri: right id of the given node
 // cost: cost of the edge from the node to the parent
 int ts = 0; // timestamp
 flat(g, a, le, ri, cost, 0, -1, ts, 0);
 int q; cin >> q;
 vector<query> queries(q);
 for (int i = 0; i < q; i++) {</pre>
   int u, v;
   cin >> u >> v;
   u--; v--;
   int lca = lca_tree.query(u, v);
   if (le[u] > le[v])
     swap(u, v);
   queries[i].id = i;
   queries[i].lca = lca;
   queries[i].u = u;
   queries[i].v = v;
   if (lca == u) {
     queries[i].a = le[u] + 1;
     queries[i].b = le[v];
   } else {
     queries[i].a = ri[u];
     queries[i].b = le[v];
 solve_mo(queries, a, le, cost); // this is the usal algorithm
```

# 1.2 Mo's algorithm

```
const int MN = 5 * 100000 + 100;
const int SN = 708;
struct query {
 int a, b, id;
 query() {}
 query(int x, int y, int i) : a(x), b(y), id(i) {}
 bool operator < (const query &o) const {</pre>
   return b < o.b;</pre>
 }
};
vector<query> s[SN];
int ans[MN];
struct DS {
 void clear() {}
 void insert(int x) {}
 void erase(int x) {}
 long long query() {}
};
DS data;
int main() {
 int n, q;
 while (cin >> n >> q) {
   for (int i = 0; i < SN; ++i)</pre>
     s[i].clear();
    vector<int> a(n);
   for (auto &i : a) cin >> i;
    for (int i = 0; i < q; ++i) {</pre>
     int b, e;
     cin >> b >> e;
     b--; e--;
     s[b / SN].emplace_back(b, e, i);
    for (int i = 0; i < SN; ++i) {</pre>
```

```
if (s[i].size())
       sort(s[i].begin(), s[i].end());
    }
   for (int b = 0; b < SN; ++b) {
     if (s[b].size() == 0) continue;
     int i = s[b][0].a;
     int j = s[b][0].a - 1;
     data.clear();
     for (int k = 0; k < (int)s[b].size(); ++k) {</pre>
       int L = s[b][k].a;
       int R = s[b][k].b;
       while (j < R) {
         j++;
         data.insert(a[j]);
       while (j > R) {
         data.erase(a[j]);
         j--;
       }
       while (i < L) {</pre>
         data.erase(a[i]);
         i++;
       }
       while (i > L) {
         i--;
         data.insert(a[i]);
       ans[s[b][k].id] = data.query();
   }
    for (int i = 0; i < q; ++i) {</pre>
     cout << ans[i] << endl;</pre>
 return 0;
};
```

# 1.3 sliding window

```
* Given an array ARR and an integer K, the problem boils down to
     computing for each index i: min(ARR[i], ARR[i-1], ..., ARR[i-K+1]).
* if mx == true, returns the maximun.
* http://people.cs.uct.ac.za/~ksmith/articles/sliding_window_minimum.html
vector<int> sliding_window_minmax(vector<int> & ARR, int K, bool mx) {
 deque< pair<int, int> > window;
 vector<int> ans;
 for (int i = 0; i < ARR.size(); i++) {</pre>
   if (mx) {
     while (!window.empty() && window.back().first <= ARR[i])</pre>
       window.pop_back();
     while (!window.empty() && window.back().first >= ARR[i])
       window.pop_back();
   }
   window.push_back(make_pair(ARR[i], i));
   while(window.front().second <= i - K)</pre>
     window.pop_front();
   ans.push_back(window.front().first);
 }
 return ans;
```

# 2 DP Optimizations

#### 2.1 convex hull trick

```
struct line {
  long long m, b;
  line (long long a, long long c) : m(a), b(c) {}
  long long eval(long long x) {
    return m * x + b;
  }
};
```

```
long double inter(line a, line b) {
 long double den = a.m - b.m;
 long double num = b.b - a.b;
 return num / den;
struct convex_hull_trick {
 vector<line> ch:
 int idx; // id of last "best" in query
  convex_hull_trick() {
   idx = 0;
 }
 void add(line cur) {
   // new line's slope is less than all the previous
   while (ch.size() > 1 &&
      (inter(cur, ch[ch.size() - 2]) >= inter(cur, ch[ch.size() - 1]))) {
       // f(x) is better in interval [inter(ch.back(), cur), inf)
       ch.pop_back();
   }
   ch.push_back(cur);
 long long query(long long x) {
   // current x is greater than all the previous x,
   // if that is not the case we can make binary search.
   idx = min<int>(idx, ch.size() - 1);
   while (idx + 1 < (int)ch.size() && ch[idx + 1].eval(x) <=
        ch[idx].eval(x))
     idx++:
   return ch[idx].eval(x);
};
```

# 3 Data structures

# 3.1 STL Treap

```
#include <ext/rope> //header with rope
```

```
using namespace std;
using namespace __gnu_cxx; //namespace with rope and some additional stuff
int main()
{
    ios_base::sync_with_stdio(false);
    rope <int> v; //use as usual STL container
    int n, m;
    cin >> n >> m;
    for(int i = 1; i <= n; ++i)</pre>
       v.push_back(i); //initialization
    int 1, r;
    for(int i = 0; i < m; ++i)</pre>
       cin >> 1 >> r;
       --1, --r;
       rope \langle int \rangle cur = v.substr(1, r - 1 + 1);
       v.erase(1, r - 1 + 1);
       v.insert(v.mutable_begin(), cur);
    for(rope <int>::iterator it = v.mutable_begin(); it !=
        v.mutable_end(); ++it)
       cout << *it << " ";
    return 0;
}
```

#### 3.2 STL order statistics tree II

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>

using namespace std;
using namespace __gnu_pbds;

typedef tree<int,null_type,less<int>,rb_tree_tag,
tree_order_statistics_node_update> order_set;

order_set X;

int get(int y) {
   int l=0,r=1e9+1;
   while(l<r) {
    int m=l+((r-1)>>1);
}
```

```
if(m-X.order_of_key(m+1)<y)</pre>
     1=m+1;
    else
      r=m;
 return 1;
}
main(){
  ios::sync_with_stdio(0);
  cin.tie(0);
 int n,m;
  cin>>n>>m;
  for(int i=0;i<m;i++) {</pre>
    char a;
    int b;
    cin>>a>>b;
    if(a=='L')
      cout<<get(b)<<endl;</pre>
    else
      X.insert(get(b));
 }
}
/***
Input
20 7
L 5
D 5
L 4
I. 5
D 5
L 4
L 5
Output
5
4
6
4
7
***/
```

#### 3.3 STL order statistics tree

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <bits/stdc++.h>
using namespace __gnu_pbds;
using namespace std;
typedef
tree<
 pair<int,int>,
 null_type,
 less<pair<int,int>>,
 rb_tree_tag,
 tree_order_statistics_node_update>
ordered_set;
main()
    ios::sync_with_stdio(0);
    cin.tie(0);
    int n;
    int sz=0;
    cin>>n;
    vector<int> ans(n,0);
    ordered_set t;
    int x,y;
    for(int i=0;i<n;i++)</pre>
       cin>>x>>y;
       ans[t.order_of_key(\{x,++sz\})]++;
       t.insert({x,sz});
    }
    for(int i=0;i<n;i++)</pre>
       cout<<ans[i]<<'\n';</pre>
}
/***
Input
1 1
5 1
```

```
7 1
3 3
5 5

Output
1
2
1
1
0
****/
```

# 3.4 binary index tree

```
struct binary_index_tree {
 int n;
 int t[2 * N];
 void add(int where, long long what){
   for (where++; where <= n; where += where & -where){</pre>
     t[where] += what;
  }
  void add(int from, int to, long long what) {
   add(from, what);
   add(to + 1, -what);
  long long query(int where){
   long long sum = t[0];
   for (where++; where > 0; where -= where & -where){
     sum += t[where];
   return sum;
 }
};
```

#### 3.5 hash table

/\*:

```
* Micro hash table, can be used as a set.

* Very efficient vs std::set

* */

const int MN = 1001;
struct ht {
  int _s[(MN + 10) >> 5];
  int len;
  void set(int id) {
    len++;
    _s[id >> 5] |= (1LL << (id & 31));
  }
  bool is_set(int id) {
    return _s[id >> 5] & (1LL << (id & 31));
  }
};</pre>
```

# 3.6 heavy light decomposition

```
// Heavy-Light Decomposition
struct TreeDecomposition {
 vector<int> g[MAXN], c[MAXN];
 int s[MAXN]; // subtree size
 int p[MAXN]; // parent id
 int r[MAXN]; // chain root id
 int t[MAXN]; // index used in segtree/bit/...
 int d[MAXN]; // depht
 int ts;
 void dfs(int v, int f) {
   p[v] = f;
   s[v] = 1;
   if (f != -1) d[v] = d[f] + 1;
   else d[v] = 0;
   for (int i = 0; i < g[v].size(); ++i) {</pre>
     int w = g[v][i];
     if (w != f) {
       dfs(w, v);
       s[v] += s[w];
     }
   }
 }
```

```
void hld(int v, int f, int k) {
 t[v] = ts++:
 c[k].push_back(v);
 r[v] = k;
  int x = 0, y = -1;
  for (int i = 0; i < g[v].size(); ++i) {</pre>
   int w = g[v][i];
   if (w != f) {
     if (s[w] > x) {
       x = s[w];
       y = w;
   }
  if (y != -1) {
   hld(y, v, k);
 for (int i = 0; i < g[v].size(); ++i) {</pre>
   int w = g[v][i];
   if (w != f && w != y) {
     hld(w, v, w);
void init(int n) {
 for (int i = 0; i < n; ++i) {</pre>
   g[i].clear();
}
void add(int a, int b) {
  g[a].push_back(b);
 g[b].push_back(a);
void build() {
  ts = 0;
  dfs(0, -1);
 hld(0, 0, 0);
```

# 3.7 persistent array

```
struct node {
 node *1, *r;
 int val;
 node (int x) : 1(NULL), r(NULL), val(x) {}
 node (): l(NULL), r(NULL), val(-1) {}
};
typedef node* pnode;
pnode update(pnode cur, int 1, int r, int at, int what) {
 pnode ans = new node();
 if (cur != NULL) {
   *ans = *cur;
 if (1 == r) {
   ans-> val = what;
   return ans;
 int m = (1 + r) >> 1;
 if (at <= m) ans-> 1 = update(ans-> 1, 1, m, at, what);
 else ans-> r = update(ans-> r, m + 1, r, at, what);
 return ans:
}
int get(pnode cur, int 1, int r, int at) {
 if (cur == NULL) return 0;
 if (1 == r) return cur-> val;
 int m = (1 + r) >> 1;
 if (at <= m) return get(cur-> 1, 1, m, at);
             return get(cur-> r, m + 1, r, at);
```

## 3.8 persistent seg tree

```
/**

* Important:

* When using lazy propagation remembert to create new

* versions for each push_down operation!!!

* */
```

```
struct node {
  node *1, *r;
  long long acc;
  int flip;
  node (int x) : 1(NULL), r(NULL), acc(x), flip(0) {}
 node () : 1(NULL), r(NULL), acc(0), flip(0) {}
};
typedef node* pnode;
pnode create(int 1, int r) {
  if (1 == r) return new node();
  pnode cur = new node();
 int m = (1 + r) >> 1;
  cur \rightarrow 1 = create(1, m);
  cur-> r = create(m + 1, r);
 return cur:
pnode copy_node(pnode cur) {
  pnode ans = new node();
  *ans = *cur;
  return ans;
void push_down(pnode cur, int 1, int r) {
  assert(cur);
 if (cur-> flip) {
    int len = r - l + 1;
    cur-> acc = len - cur-> acc;
    if (cur-> 1) {
     cur-> 1 = copy_node(cur-> 1);
     cur-> 1 -> flip ^= 1;
    if (cur-> r) {
     cur-> r = copy_node(cur-> r);
     cur-> r -> flip ^= 1;
    cur \rightarrow flip = 0;
 }
}
int get_val(pnode cur) {
```

```
assert(cur);
 assert((cur-> flip) == 0);
 if (cur) return cur-> acc:
 return 0;
}
pnode update(pnode cur, int 1, int r, int at, int what) {
 pnode ans = copy_node(cur);
 if (1 == r) {
   assert(1 == at);
   ans-> acc = what:
   ans-> flip = 0;
   return ans;
 int m = (1 + r) >> 1;
 push_down(ans, 1, r);
 if (at <= m) ans-> 1 = update(ans-> 1, 1, m, at, what);
 else ans-> r = update(ans-> r, m + 1, r, at, what);
 push_down(ans-> 1, 1, m);
 push_down(ans-> r, m + 1, r);
 ans-> acc = get_val(ans-> 1) + get_val(ans-> r);
 return ans;
}
pnode flip(pnode cur, int 1, int r, int a, int b) {
 pnode ans = new node();
 if (cur != NULL) {
   *ans = *cur;
 if (1 > b | | r < a)
   return ans:
 if (1 >= a && r <= b) {</pre>
   ans-> flip ^= 1;
   push_down(ans, 1, r);
   return ans;
 }
 int m = (1 + r) >> 1;
 ans-> 1 = flip(ans-> 1, 1, m, a, b);
 ans-> r = flip(ans-> r, m + 1, r, a, b);
 push_down(ans-> 1, 1, m);
 push_down(ans-> r, m + 1, r);
```

### 3.9 persistent trie

```
* Persistent version of trie:
 * could be used as a persistent BST of integers.
 * add: O(log2(n))
 * query: O(log2(n)), works equal than the non-persistent version.
 * */
const int MD = 31;
const int MAX_CHILD = 2;
struct node {
 node *child[MAX_CHILD];
};
typedef node* pnode;
pnode copy_node(pnode cur) {
 pnode ans = new node();
 if (cur)
   *ans = *cur;
 return ans:
```

```
pnode add(pnode cur, int val, int id = MD) {
  pnode ans = copy_node(cur);
  if (id == -1) return ans;
  int t = (val >> id) & 1;
  ans-> child[t] = add(ans-> child[t], val, id - 1);
  return ans;
}
```

# 3.10 segment tree

```
/**
 * Taken from: http://codeforces.com/blog/entry/18051
 * */
const int N = 1e5; // limit for array size
int n; // array size
int t[2 * N];
void build() { // build the tree
 for (int i = n - 1; i > 0; --i) t[i] = t[i << 1] + t[i << 1|1];
}
// Single modification, range query.
void modify(int p, int value) { // set value at position p
 for (t[p += n] = value; p > 1; p >>= 1) t[p>>1] = t[p] + t[p^1];
}
int query(int 1, int r) { // sum on interval [1, r)
 int res = 0;
 for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
   if (l&1) res += t[l++];
   if (r\&1) res += t[--r];
 }
 return res;
}
// Range modification, single query.
void modify(int 1, int r, int value) {
 for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
   if (1&1) t[1++] += value;
   if (r\&1) t[--r] += value:
 }
```

```
int query(int p) {
 int res = 0;
 for (p += n; p > 0; p >>= 1) res += t[p];
 return res;
* If at some point after modifications we need to inspect all the
 * elements in the array, we can push all the modifications to the
 * leaves using the following code. After that we can just traverse
 * elements starting with index n. This way we reduce the complexity
 * from O(n \log(n)) to O(n) similarly to using build instead of n
     modifications.
 * */
void push() {
 for (int i = 1; i < n; ++i) {
   t[i<<1] += t[i];
   t[i << 1|1] += t[i];
   t[i] = 0;
}
// Non commutative combiner functions.
void modify(int p, const S& value) {
 for (t[p += n] = value; p >>= 1;) t[p] = combine(t[p<<1], t[p<<1|1]);
S query(int 1, int r) {
 S resl, resr;
 for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
   if (l\&1) resl = combine(resl, t[l++]);
   if (r&1) resr = combine(t[--r], resr);
 return combine(resl, resr);
// To be continued ...
```

## 3.11 sparse table

```
// RMQ.
const int MN = 100000 + 10; // Max number of elements
const int ML = 18; // ceil(log2(MN));
struct st {
 int data[MN]:
 int M[MN][ML];
 int n;
 void read(int _n) {
   n = _n;
   for (int i = 0; i < n; ++i)
     cin >> data[i];
 }
 void build() {
   for (int i = 0; i < n; ++i)
     M[i][0] = data[i];
   for (int j = 1, p = 2, q = 1; p \le n; ++j, p \le 1, q \le 1)
     for (int i = 0; i + p - 1 < n; ++i)
       M[i][j] = max(M[i][j-1], M[i+q][j-1]);
 }
 int query(int b, int e) {
   int k = log2(e - b + 1);
   return max(M[b][k], M[e + 1 - (1<<k)][k]);</pre>
 }
};
```

# 3.12 splay tree

```
using namespace std;
#include<bits/stdc++.h>
#define D(x) cout<<x<<endl;

typedef int T;

struct node{
  node *left, *right, *parent;
  T key;
  node (T k) : key(k), left(0), right(0), parent(0) {}
};</pre>
```

```
struct splay_tree{
 node *root;
 void right_rot(node *x) {
   node *p = x->parent;
   if (x->parent = p->parent) {
     if (x->parent->left == p) x->parent->left = x;
     if (x->parent->right == p) x->parent->right = x;
   if (p->left = x->right) p->left->parent = p;
   x->right = p;
   p->parent = x;
 void left_rot(node *x) {
   node *p = x->parent;
   if (x->parent = p->parent) {
     if (x->parent->left == p) x->parent->left = x;
     if (x->parent->right == p) x->parent->right = x;
   if (p->right = x->left) p->right->parent = p;
   x\rightarrowleft = p;
   p->parent = x;
 void splay(node *x, node *fa = 0) {
   while( x->parent != fa and x->parent != 0) {
     node *p = x->parent;
     if (p->parent == fa)
       if (p->right == x)
        left_rot(x);
       else
        right_rot(x);
     else {
       node *gp = p->parent; //grand parent
       if (gp->left == p)
        if (p->left == x)
          right_rot(x), right_rot(x);
          left_rot(x),right_rot(x);
       else
         if (p->left == x)
```

```
right_rot(x), left_rot(x);
         else
           left_rot(x), left_rot(x);
     }
   }
   if (fa == 0) root = x;
 }
 void insert(T key) {
   node *cur = root;
   node *pcur = 0;
   while (cur) {
     pcur = cur;
     if (key > cur->key) cur = cur->right;
     else cur = cur->left;
   }
   cur = new node(key);
   cur->parent = pcur;
   if (!pcur) root = cur;
   else if (key > pcur->key ) pcur->right = cur;
   else pcur->left = cur;
   splay(cur);
 }
 node *find(T key) {
   node *cur = root;
   while (cur) {
     if (key > cur->key) cur = cur->right;
     else if(key < cur->key) cur = cur->left;
     else return cur;
   }
   return 0;
 }
  splay_tree(){ root = 0;};
};
```

#### 3.13 trie

```
const int MN = 26; // size of alphabet
const int MS = 100010; // Number of states.
struct trie{
```

```
struct node{
   int c;
   int a[MN];
 };
  node tree[MS];
  int nodes;
  void clear(){
   tree[nodes].c = 0;
   memset(tree[nodes].a, -1, sizeof tree[nodes].a);
   nodes++;
 }
  void init(){
   nodes = 0:
   clear();
  int add(const string &s, bool query = 0){
   int cur_node = 0;
   for(int i = 0; i < s.size(); ++i){</pre>
     int id = gid(s[i]);
     if(tree[cur_node].a[id] == -1){
       if(query) return 0;
       tree[cur_node].a[id] = nodes;
       clear():
     cur_node = tree[cur_node].a[id];
   if(!query) tree[cur_node].c++;
   return tree[cur_node].c;
};
```

# 4 Geometry

# 4.1 center 2 points + radious

```
vector<point> find_center(point a, point b, long double r) {
  point d = (a - b) * 0.5;
```

```
if (d.dot(d) > r * r) {
    return vector<point> ();
}
point e = b + d;
long double fac = sqrt(r * r - d.dot(d));
vector<point> ans;
point x = point(-d.y, d.x);
long double l = sqrt(x.dot(x));
x = x * (fac / l);
ans.push_back(e + x);
x = point(d.y, -d.x);
x = x * (fac / l);
ans.push_back(e + x);
return ans;
}
```

# 4.2 closest pair problem

```
struct point {
 double x, y;
 int id;
 point() {}
 point (double a, double b) : x(a), y(b) {}
}:
double dist(const point &o, const point &p) {
 double a = p.x - o.x, b = p.y - o.y;
 return sqrt(a * a + b * b);
}
double cp(vector<point> &p, vector<point> &x, vector<point> &y) {
 if (p.size() < 4) {</pre>
   double best = 1e100;
   for (int i = 0; i < p.size(); ++i)</pre>
     for (int j = i + 1; j < p.size(); ++j)</pre>
       best = min(best, dist(p[i], p[j]));
   return best;
 }
 int ls = (p.size() + 1) >> 1;
 double 1 = (p[ls - 1].x + p[ls].x) * 0.5;
 vector<point> xl(ls), xr(p.size() - ls);
 unordered_set<int> left;
```

```
for (int i = 0; i < ls; ++i) {</pre>
   xl[i] = x[i];
   left.insert(x[i].id);
  for (int i = ls; i < p.size(); ++i) {</pre>
   xr[i - ls] = x[i];
  vector<point> yl, yr;
  vector<point> pl, pr;
 yl.reserve(ls); yr.reserve(p.size() - ls);
  pl.reserve(ls); pr.reserve(p.size() - ls);
 for (int i = 0; i < p.size(); ++i) {</pre>
   if (left.count(y[i].id))
     yl.push_back(y[i]);
   else
     yr.push_back(y[i]);
   if (left.count(p[i].id))
     pl.push_back(p[i]);
   else
     pr.push_back(p[i]);
 }
  double dl = cp(pl, xl, yl);
  double dr = cp(pr, xr, yr);
  double d = min(dl, dr);
  vector<point> yp; yp.reserve(p.size());
 for (int i = 0; i < p.size(); ++i) {</pre>
   if (fabs(y[i].x - 1) < d)
     yp.push_back(y[i]);
 for (int i = 0; i < yp.size(); ++i) {</pre>
   for (int j = i + 1; j < yp.size() && j < i + 7; ++j) {
     d = min(d, dist(yp[i], yp[j]));
   }
 }
 return d;
double closest_pair(vector<point> &p) {
 vector<point> x(p.begin(), p.end());
  sort(x.begin(), x.end(), [](const point &a, const point &b) {
   return a.x < b.x;</pre>
 });
```

```
vector<point> y(p.begin(), p.end());
sort(y.begin(), y.end(), [](const point &a, const point &b) {
   return a.y < b.y;
});
return cp(p, x, y);
}</pre>
```

### 4.3 squares

```
typedef long double ld;
const ld eps = 1e-12;
int cmp(ld x, ld y = 0, ld tol = eps) {
   return ( x \le y + tol) ? (x + tol < y) ? -1 : 0 : 1;
}
struct point{
 ld x, y;
 point(ld a, ld b) : x(a), y(b) {}
 point() {}
};
struct square{
 ld x1, x2, y1, y2,
    a, b, c;
 point edges[4];
 square(ld _a, ld _b, ld _c) {
   a = _a, b = _b, c = _c;
   x1 = a - c * 0.5;
   x2 = a + c * 0.5;
   v1 = b - c * 0.5;
   v2 = b + c * 0.5;
   edges[0] = point(x1, y1);
   edges[1] = point(x2, y1);
   edges[2] = point(x2, y2);
   edges[3] = point(x1, y2);
 }
};
ld min_dist(point &a, point &b) {
 1d x = a.x - b.x,
    y = a.y - b.y;
```

```
return sqrt(x * x + y * y);
}
bool point_in_box(square s1, point p) {
 if (cmp(s1.x1, p.x) != 1 \&\& cmp(s1.x2, p.x) != -1 \&\&
     cmp(s1.y1, p.y) != 1 && cmp(s1.y2, p.y) != -1)
   return true;
 return false:
bool inside(square &s1, square &s2) {
 for (int i = 0; i < 4; ++i)
   if (point_in_box(s2, s1.edges[i]))
     return true;
 return false:
bool inside_vert(square &s1, square &s2) {
 if ((cmp(s1.y1, s2.y1) != -1 && cmp(s1.y1, s2.y2) != 1) ||
     (cmp(s1.y2, s2.y1) != -1 \&\& cmp(s1.y2, s2.y2) != 1))
   return true;
return false;
}
bool inside_hori(square &s1, square &s2) {
 if ((cmp(s1.x1, s2.x1) != -1 && cmp(s1.x1, s2.x2) != 1) ||
     (cmp(s1.x2, s2.x1) != -1 \&\& cmp(s1.x2, s2.x2) != 1))
   return true;
return false;
ld min_dist(square &s1, square &s2) {
 if (inside(s1, s2) || inside(s2, s1))
   return 0;
 ld ans = 1e100;
 for (int i = 0; i < 4; ++i)
   for (int j = 0; j < 4; ++j)
     ans = min(ans, min_dist(s1.edges[i], s2.edges[j]));
  if (inside_hori(s1, s2) || inside_hori(s2, s1)) {
   if (cmp(s1.y1, s2.y2) != -1)
     ans = min(ans, s1.y1 - s2.y2);
```

```
else
if (cmp(s2.y1, s1.y2) != -1)
   ans = min(ans, s2.y1 - s1.y2);
}

if (inside_vert(s1, s2) || inside_vert(s2, s1)) {
   if (cmp(s1.x1, s2.x2) != -1)
      ans = min(ans, s1.x1 - s2.x2);
   else
   if (cmp(s2.x1, s1.x2) != -1)
      ans = min(ans, s2.x1 - s1.x2);
}

return ans;
```

# 4.4 triangles

Let a, b, c be length of the three sides of a triangle.

$$p = (a + b + c) * 0.5$$

The inradius is defined by:

$$iR = \sqrt{\frac{(p-a)(p-b)(p-c)}{p}}$$

The radius of its circumcircle is given by the formula:

$$cR = \frac{abc}{\sqrt{(a+b+c)(a+b-c)(a+c-b)(b+c-a)}}$$

# 5 Graphs

# 5.1 bridges

```
struct edge{
  int to, id;
  edge(int a, int b) : to(a), id(b) {}
};

struct graph {
  vector<vector<edge> > g;
```

```
vector<int> vi, low, d, pi, is_b;
int ticks, edges;
graph(int n, int m) {
 g.assign(n, vector<edge>());
 is_b.assign(m, 0);
 vi.resize(n);
 low.resize(n);
 d.resize(n);
 pi.resize(n);
 edges = 0;
void add_edge(int u, int v) {
 g[u].push_back(edge(v, edges));
 g[v].push_back(edge(u, edges));
 edges++;
void dfs(int u) {
 vi[u] = true;
 d[u] = low[u] = ticks++;
 for (int i = 0; i < g[u].size(); ++i) {</pre>
   int v = g[u][i].to;
   if (v == pi[u]) continue;
   if (!vi[v]) {
     pi[v] = u;
     dfs(v);
     if (d[u] < low[v])</pre>
       is_b[g[u][i].id] = true;
     low[u] = min(low[u], low[v]);
   } else {
     low[u] = min(low[u], d[v]);
 }
// Multiple edges from a to b are not allowed.
// (they could be detected as a bridge).
// If you need to handle this, just count
// how many edges there are from a to b.
void comp_bridges() {
```

```
fill(pi.begin(), pi.end(), -1);
fill(vi.begin(), vi.end(), 0);
fill(low.begin(), low.end(), 0);
fill(d.begin(), d.end(), 0);
ticks = 0;
for (int i = 0; i < g.size(); ++i)
    if (!vi[i]) dfs(i);
}
};</pre>
```

# 5.2 dijkstra

```
struct edge {
 int to;
 long long w;
 edge () {}
  edge (int a, long long b) : to(a), w(b) {}
 bool operator < (const edge &o) const {</pre>
   return w > o.w;
 }
};
typedef vector<vector<edge>> graph;
const long long inf = 1000000LL * 10000000LL;
pair<vector<int>, vector<long long>> dijkstra(graph &g, int start) {
 int n = g.size();
 vector<long long> d(n, inf);
 vector<int> p(n, -1);
 d[start] = 0;
 priority_queue<edge> q;
 q.push(edge(start, 0));
  while (!q.empty()) {
   int node = q.top().to;
   long long dist = q.top().w;
   q.pop();
   if (dist > d[node]) continue;
   for (int i = 0; i < (int)g[node].size(); i++) {</pre>
     int to = g[node][i].to;
```

```
long long w_extra = g[node][i].w;

if (dist + w_extra < d[to]) {
    p[to] = node;
    d[to] = dist + w_extra;
    q.push(edge(to, d[to]));
    }
}

return {p, d};
}</pre>
```

#### 5.3 directed mst

```
const int inf = 1000000 + 10;
struct edge {
 int u, v, w;
 edge() {}
  edge(int a,int b,int c): u(a), v(b), w(c) {}
};
 * Computes the minimum spanning tree for a directed graph
 * - edges : Graph description in the form of list of edges.
 * each edge is: From node u to node v with cost w
 * - root : Id of the node to start the DMST.
 * - n : Number of nodes in the graph.
* */
int dmst(vector<edge> &edges, int root, int n) {
  int ans = 0;
 int cur nodes = n:
  while (true) {
   vector<int> lo(cur_nodes, inf), pi(cur_nodes, inf);
   for (int i = 0; i < edges.size(); ++i) {</pre>
     int u = edges[i].u, v = edges[i].v, w = edges[i].w;
     if (w < lo[v] and u != v) {
       lo[v] = w;
       pi[v] = u;
```

```
lo[root] = 0;
  for (int i = 0: i < lo.size(): ++i) {</pre>
   if (i == root) continue;
   if (lo[i] == inf) return -1;
  }
  int cur_id = 0;
  vector<int> id(cur_nodes, -1), mark(cur_nodes, -1);
  for (int i = 0; i < cur_nodes; ++i) {</pre>
   ans += lo[i];
   int u = i:
   while (u != root and id[u] < 0 and mark[u] != i) {</pre>
     mark[u] = i;
     u = pi[u];
   if (u != root and id[u] < 0) { // Cycle}
      for (int v = pi[u]; v != u; v = pi[v])
        id[v] = cur_id;
      id[u] = cur_id++;
   }
  }
  if (cur_id == 0)
   break;
  for (int i = 0; i < cur_nodes; ++i)</pre>
   if (id[i] < 0) id[i] = cur id++:
  for (int i = 0; i < edges.size(); ++i) {</pre>
   int u = edges[i].u, v = edges[i].v, w = edges[i].w;
   edges[i].u = id[u];
   edges[i].v = id[v];
   if (id[u] != id[v])
     edges[i].w -= lo[v];
  cur_nodes = cur_id;
 root = id[root];
}
return ans;
```

# 5.4 eulerian path

```
// Taken from https://github.com/lbv/pc-code/blob/master/code/graph.cpp
// Eulerian Trail
struct Euler {
 ELV adj; IV t;
  Euler(ELV Adj) : adj(Adj) {}
 void build(int u) {
   while(! adj[u].empty()) {
     int v = adj[u].front().v;
     adj[u].erase(adj[u].begin());
     build(v);
   t.push_back(u);
ጉ:
bool eulerian_trail(IV &trail) {
 Euler e(adj);
 int odd = 0, s = 0;
    for (int v = 0; v < n; v++) {
    int diff = abs(in[v] - out[v]);
    if (diff > 1) return false;
    if (diff == 1) {
    if (++odd > 2) return false;
    if (out[v] > in[v]) start = v;
    }
    }
    */
  e.build(s);
 reverse(e.t.begin(), e.t.end());
 trail = e.t;
 return true;
```

# 5.5 karp min mean cycle

```
/**
 * Finds the min mean cycle, if you need the max mean cycle
 * just add all the edges with negative cost and print
 * ans * -1
 *
 * test: uva, 11090 - Going in Cycle!!
```

```
* */
const int MN = 1000;
struct edge{
 int v;
 long long w;
 edge(){} edge(int v, int w) : v(v), w(w) {}
}:
long long d[MN][MN];
// This is a copy of g because increments the size
// pass as reference if this does not matter.
int karp(vector<vector<edge> > g) {
 int n = g.size();
 g.resize(n + 1); // this is important
 for (int i = 0; i < n; ++i)
   if (!g[i].empty())
     g[n].push_back(edge(i,0));
  ++n:
 for(int i = 0;i<n;++i)</pre>
   fill(d[i],d[i]+(n+1),INT_MAX);
 d[n - 1][0] = 0:
 for (int k = 1; k \le n; ++k) for (int u = 0; u \le n; ++u) {
   if (d[u][k - 1] == INT_MAX) continue;
   for (int i = g[u].size() - 1; i >= 0; --i)
     d[g[u][i].v][k] = min(d[g[u][i].v][k], d[u][k-1] + g[u][i].w);
 }
 bool flag = true;
 for (int i = 0; i < n && flag; ++i)</pre>
   if (d[i][n] != INT_MAX)
     flag = false;
 if (flag) {
   return true; // return true if there is no a cycle.
 }
 double ans = 1e15:
```

```
for (int u = 0; u + 1 < n; ++u) {
   if (d[u][n] == INT_MAX) continue;
   double W = -1e15;

   for (int k = 0; k < n; ++k)
      if (d[u][k] != INT_MAX)
      W = max(W, (double)(d[u][n] - d[u][k]) / (n - k));

   ans = min(ans, W);
}

// printf("%.2lf\n", ans);
   cout << fixed << setprecision(2) << ans << endl;
   return false;
}</pre>
```

# 5.6 konig's theorem

In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover

# 5.7 minimum path cover in DAG

Given a directed acyclic graph G = (V, E), we are to find the minimum number of vertex-disjoint paths to cover each vertex in V.

We can construct a bipartite graph  $G' = (Vout \cup Vin, E')$  from G, where :

```
Vout = \{v \in V : v \text{ has positive out} - degree\}
Vin = \{v \in V : v \text{ has positive in} - degree\}
E' = \{(u, v) \in Vout \times Vin : (u, v) \in E\}
```

Then it can be shown, via König's theorem, that G' has a matching of size m if and only if there exists n-m vertex-disjoint paths that cover each vertex in G, where n is the number of vertices in G and m is the maximum cardinality bipartite mathching in G'.

Therefore, the problem can be solved by finding the maximum cardinality matching in G' instead.

**NOTE:** If the paths are note necesarily disjoints, find the transitive closure and solve the problem for disjoint paths.

# 5.8 planar graph (euler)

Euler's formula states that if a finite, connected, planar graph is drawn in the plane without any edge intersections, and v is the number of vertices, e is the number of edges and f is the number of faces (regions bounded by edges, including the outer, infinitely large region), then:

$$f + v = e + 2$$

It can be extended to non connected planar graphs with c connected components:

$$f + v = e + c + 1$$

# 5.9 query with lca

```
struct lowest_ca {
 int T[MN], L[MN], W[MN];
 int P[MN][ML], MI[MN][ML], MA[MN][ML];
 void dfs(vector<vector<edge> > &g, int root, int pi = -1) {
   if (pi == -1) {
     L[root] = W[root] = 0;
     T[root] = -1;
   }
   for (int i = 0; i < (int)g[root].size(); ++i) {</pre>
     int to = g[root][i].v;
     if (to != pi) {
       T[to] = root;
       W[to] = g[root][i].w;
       L[to] = L[root] + 1;
       dfs(g, to, root);
     }
   }
 }
 void init(vector<vector<edge> > &g, int root) {
   // g is undirected
   dfs(g, root);
   int N = g.size(), i, j;
   for (i = 0: i < N: i++) {</pre>
     for (j = 0; 1 << j < N; j++) {
       P[i][i] = -1;
```

```
MI[i][j] = inf;
   }
 }
 for (i = 0; i < N; i++) {</pre>
   P[i][0] = T[i];
   MI[i][0] = W[i];
 }
 for (j = 1; 1 << j < N; j++)
   for (i = 0; i < N; i++)</pre>
     if (P[i][i - 1] != -1) {
       P[i][j] = P[P[i][j-1]][j-1];
       MI[i][j] = min(MI[i][j-1], MI[P[i][j-1]][j-1]);
     }
}
int query(int p, int q) {
 int tmp, log, i;
 int mmin = inf;
 if (L[p] < L[q])
   tmp = p, p = q, q = tmp;
 for (log = 1; 1 << log <= L[p]; log++);</pre>
 log--;
 for (i = log; i >= 0; i--)
   if (L[p] - (1 << i) >= L[q]) {
     mmin = min(mmin, MI[p][i]);
     p = P[p][i];
 if (p == q)
   // return p;
   return mmin;
 for (i = log; i >= 0; i--)
   if (P[p][i] != -1 && P[p][i] != P[q][i]) {
     mmin = min(mmin, min(MI[p][i], MI[q][i]));
     p = P[p][i], q = P[q][i];
 // return T[p];
 return min(mmin, min(MI[p][0], MI[q][0]));
```

```
}
};
```

# 5.10 tarjan scc

```
const int MN = 20002;
struct tarjan_scc {
 int scc[MN], low[MN], d[MN], stacked[MN];
 int ticks, current_scc;
 deque<int> s; // used as stack.
 tarjan_scc() {}
 void init () {
   memset(scc, -1, sizeof scc);
   memset(d, -1, sizeof d);
   memset(stacked, 0, sizeof stacked);
   s.clear();
   ticks = current_scc = 0;
 }
 void compute(vector<vector<int> > &g, int u) {
   d[u] = low[u] = ticks++;
   s.push_back(u);
   stacked[u] = true;
   for (int i = 0; i < g[u].size(); ++i) {</pre>
     int v = g[u][i];
     if (d[v] == -1)
       compute(g, v);
     if (stacked[v]) {
       low[u] = min(low[u], low[v]);
     }
   }
   if (d[u] == low[u]) { // root
     int v;
     do {
       v = s.back();s.pop_back();
       stacked[v] = false;
       scc[v] = current_scc;
     } while (u != v);
     current_scc++;
```

```
}
};
```

# 5.11 two sat (with kosaraju)

```
/**
* Given a set of clauses (a1 v a2)^(a2 v a3)....
* this algorithm find a solution to it set of clauses.
* test: http://lightoj.com/volume_showproblem.php?problem=1251
#include<bits/stdc++.h>
using namespace std;
#define MAX 100000
#define endl '\n'
vector<int> G[MAX];
vector<int> GT[MAX];
vector<int> Ftime;
vector<vector<int> > SCC;
bool visited[MAX];
int n;
void dfs1(int n){
 visited[n] = 1;
 for (int i = 0; i < G[n].size(); ++i) {</pre>
   int curr = G[n][i];
   if (visited[curr]) continue;
   dfs1(curr);
 Ftime.push_back(n);
void dfs2(int n, vector<int> &scc) {
 visited[n] = 1;
 scc.push_back(n);
 for (int i = 0;i < GT[n].size(); ++i) {</pre>
   int curr = GT[n][i];
```

```
if (visited[curr]) continue;
   dfs2(curr, scc);
 }
}
void kosaraju() {
 memset(visited, 0, sizeof visited);
 for (int i = 0; i < 2 * n; ++i) {
   if (!visited[i]) dfs1(i);
 }
 memset(visited, 0, sizeof visited);
 for (int i = Ftime.size() - 1; i >= 0; i--) {
   if (visited[Ftime[i]]) continue;
   vector<int> _scc;
   dfs2(Ftime[i],_scc);
   SCC.push_back(_scc);
}
/**
 * After having the SCC, we must traverse each scc, if in one SCC are -b
     y b, there is not a solution.
 * Otherwise we build a solution, making the first "node" that we find
     truth and its complement false.
 **/
bool two_sat(vector<int> &val) {
 kosaraju();
 for (int i = 0; i < SCC.size(); ++i) {</pre>
   vector<bool> tmpvisited(2 * n, false);
   for (int j = 0; j < SCC[i].size(); ++j) {</pre>
     if (tmpvisited[SCC[i][j] ^ 1]) return 0;
     if (val[SCC[i][j]] != -1) continue;
     else {
       val[SCC[i][j]] = 0;
       val[SCC[i][j] ^ 1] = 1;
     tmpvisited[SCC[i][j]] = 1;
 }
 return 1;
```

```
// Example of use
int main() {
 int m, u, v, nc = 0, t; cin >> t;
 // n = "nodes" number, m = clauses number
  while (t--) {
   cin >> m >> n:
   Ftime.clear();
   SCC.clear();
   for (int i = 0; i < 2 * n; ++i) {
     G[i].clear();
     GT[i].clear();
   // (a1 v a2) = (a1 -> a2) = (a2 -> a1)
   for (int i = 0; i < m; ++i) {</pre>
     cin >> u >> v;
     int t1 = abs(u) - 1;
     int t2 = abs(v) - 1;
     int p = t1 * 2 + ((u < 0)? 1 : 0);
     int q = t2 * 2 + ((v < 0)? 1 : 0);
     G[p ^ 1].push_back(q);
     G[q ^ 1].push_back(p);
     GT[p].push_back(q ^ 1);
     GT[q].push_back(p ^ 1);
   vector < int > val(2 * n, -1);
   cout << "Case " << ++nc <<": ";
   if (two sat(val)) {
     cout << "Yes" << endl;</pre>
     vector<int> sol;
     for (int i = 0; i < 2 * n; ++i)
       if (i % 2 == 0 and val[i] == 1)
         sol.push_back(i / 2 + 1);
     cout << sol.size() ;</pre>
     for (int i = 0; i < sol.size(); ++i) {</pre>
       cout << " " << sol[i]:</pre>
     cout << endl;</pre>
```

```
} else {
    cout << "No" << endl;
}
return 0;
}</pre>
```

### 6 Math

#### 6.1 Lucas theorem

For non-negative integers m and n and a prime p, the following congruence relation holds: :

$$\binom{m}{n} \equiv \prod_{i=0}^{k} \binom{m_i}{n_i} \pmod{p},$$

where:

$$m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0,$$

and:

$$n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$$

are the base p expansions of m and n respectively. This uses the convention that  $\binom{m}{n} = 0$  if  $m \le n$ .

#### 6.2 cumulative sum of divisors

```
/**
The function SOD(n) (sum of divisors) is defined
as the summation of all the actual divisors of
an integer number n. For example,

SOD(24) = 2+3+4+6+8+12 = 35.

The function CSOD(n) (cumulative SOD) of an integer n, is defined as
below:

csod(n) = \sum_{{i = 1}^{n}} sod(i)

It can be computed in O(sqrt(n)):
*/
long long csod(long long n) {
```

```
long long ans = 0;
for (long long i = 2; i * i <= n; ++i) {
   long long j = n / i;
   ans += (i + j) * (j - i + 1) / 2;
   ans += i * (j - i);
}
return ans;
}</pre>
```

#### 6.3 fft

```
* Fast Fourier Transform.
 * Useful to compute convolutions.
 * computes:
 * C(f \operatorname{star} g)[n] = \operatorname{sum}_m(f[m] * g[n - m])
 * for all n.
 * test: icpc live archive, 6886 - Golf Bot
 * */
using namespace std;
#include<bits/stdc++.h>
#define D(x) cout << #x " = " << (x) << endl
#define endl '\n'
const int MN = 262144 << 1;</pre>
int d[MN + 10], d2[MN + 10];
const double PI = acos(-1.0);
struct cpx {
 double real, image;
  cpx(double _real, double _image) {
   real = _real;
   image = _image;
  cpx(){}
};
cpx operator + (const cpx &c1, const cpx &c2) {
 return cpx(c1.real + c2.real, c1.image + c2.image);
```

```
}
cpx operator - (const cpx &c1, const cpx &c2) {
 return cpx(c1.real - c2.real, c1.image - c2.image);
cpx operator * (const cpx &c1, const cpx &c2) {
 return cpx(c1.real*c2.real - c1.image*c2.image, c1.real*c2.image +
      c1.image*c2.real);
}
int rev(int id, int len) {
 int ret = 0;
 for (int i = 0; (1 << i) < len; i++) {</pre>
   ret <<= 1;
   if (id & (1 << i)) ret |= 1;</pre>
 return ret;
}
cpx A[1 << 20];
void FFT(cpx *a, int len, int DFT) {
 for (int i = 0; i < len; i++)</pre>
   A[rev(i, len)] = a[i];
 for (int s = 1; (1 << s) <= len; s++) {
   int m = (1 << s):
   cpx wm = cpx(cos(DFT * 2 * PI / m), sin(DFT * 2 * PI / m));
   for(int k = 0; k < len; k += m) {</pre>
     cpx w = cpx(1, 0);
     for(int j = 0; j < (m >> 1); j++) {
       cpx t = w * A[k + j + (m >> 1)];
       cpx u = A[k + j];
       A[k + j] = u + t;
       A[k + j + (m >> 1)] = u - t;
       w = w * wm;
     }
   }
 }
 if (DFT == -1) for (int i = 0; i < len; i++) A[i].real /= len,
      A[i].image /= len;
 for (int i = 0; i < len; i++) a[i] = A[i];</pre>
 return;
}
```

```
cpx in[1 << 20];
void solve(int n) {
 memset(d, 0, sizeof d);
 for (int i = 0; i < n; ++i) {</pre>
   cin >> t;
   d[t] = true;
  int m;
  cin >> m;
  vector<int> q(m);
 for (int i = 0; i < m; ++i)</pre>
   cin >> q[i];
 for (int i = 0; i < MN; ++i) {</pre>
   if (d[i])
     in[i] = cpx(1, 0);
   else
     in[i] = cpx(0, 0);
 FFT(in, MN, 1);
 for (int i = 0; i < MN; ++i) {</pre>
   in[i] = in[i] * in[i];
 FFT(in, MN, -1);
 int ans = 0;
 for (int i = 0; i < q.size(); ++i) {</pre>
   if (in[q[i]].real > 0.5 || d[q[i]]) {
     ans++:
   }
  cout << ans << endl;</pre>
int main() {
  ios_base::sync_with_stdio(false);cin.tie(NULL);
 int n;
  while (cin >> n)
   solve(n):
 return 0;
```

### 6.4 fibonacci properties

Let A, B and n be integer numbers.

$$k = A - B \tag{1}$$

$$F_A F_B = F_{k+1} F_A^2 + F_k F_A F_{A-1} \tag{2}$$

$$\sum_{i=0}^{n} F_i^2 = F_{n+1} F_n \tag{3}$$

ev(n) = returns 1 if n is even.

$$\sum_{i=0}^{n} F_i F_{i+1} = F_{n+1}^2 - ev(n) \tag{4}$$

$$\sum_{i=0}^{n} F_i F_{i-1} = \sum_{i=0}^{n-1} F_i F_{i+1} \tag{5}$$

# 6.5 sigma function

the sigma function is defined as:

$$\sigma_x(n) = \sum_{d|n} d^x$$

when x=0 is called the divisor function, that counts the number of positive divisors of n.

Now, we are interested in find

$$\sum_{d|n} \sigma_0(d)$$

if n is written as prime factorization:

$$n = \prod_{i=1}^{k} P_i^{e_k}$$

we can demonstrate that:

$$\sum_{d|n} \sigma_0(d) = \prod_{i=1}^k g(e_k + 1)$$

where g(x) is the sum of the first x positive numbers:

$$g(x) = (x * (x+1))/2$$

# 7 Matrix

#### 7.1 matrix

```
const int MN = 111;
const int mod = 10000;
struct matrix {
  int r, c;
  int m[MN][MN];
  matrix (int _r, int _c) : r (_r), c (_c) {
   memset(m, 0, sizeof m);
  void print() {
   for (int i = 0; i < r; ++i) {</pre>
     for (int j = 0; j < c; ++j)
       cout << m[i][j] << " ";
     cout << endl:</pre>
   }
 }
  int x[MN][MN];
  matrix & operator *= (const matrix &o) {
   memset(x, 0, sizeof x);
   for (int i = 0; i < r; ++i)
     for (int k = 0; k < c; ++k)
       if (m[i][k] != 0)
         for (int j = 0; j < c; ++j) {
           x[i][j] = (x[i][j] + ((m[i][k] * o.m[k][j]) % mod)) % mod;
   memcpy(m, x, sizeof(m));
   return *this;
};
void matrix_pow(matrix b, long long e, matrix &res) {
  memset(res.m, 0, sizeof res.m);
```

```
for (int i = 0; i < b.r; ++i)
  res.m[i][i] = 1;

if (e == 0) return;
while (true) {
  if (e & 1) res *= b;
  if ((e >>= 1) == 0) break;
  b *= b;
}
```

# 8 Misc

### 8.1 Template Java

```
import java.io.*;
import java.util.StringTokenizer;
public class Template {
   public static void main(String []args) throws IOException {
       Scanner in = new Scanner(System.in);
       OutputWriter out = new OutputWriter(System.out);
       Task solver = new Task();
       solver.solve(in, out);
       out.close();
   }
}
class Task{
   public void solve(Scanner in, OutputWriter out){
class Scanner{
   public BufferedReader reader;
   public StringTokenizer st;
   public Scanner(InputStream stream){
       reader = new BufferedReader(new InputStreamReader(stream));
```

```
st = null:
   public String next(){
       while(st == null || !st.hasMoreTokens()){
          try{
              String line = reader.readLine();
              if(line == null) return null;
              st = new StringTokenizer(line);
          }catch (Exception e){
              throw (new RuntimeException());
          }
       }
       return st.nextToken();
   public int nextInt(){
       return Integer.parseInt(next());
   public long nextLong(){
       return Long.parseLong(next());
   public double nextDouble(){
       return Double.parseDouble(next());
}
class OutputWriter{
   BufferedWriter writer;
   public OutputWriter(OutputStream stream){
       writer = new BufferedWriter(new OutputStreamWriter(stream));
   public void print(int i) throws IOException {
       writer.write(i);
   public void print(String s) throws IOException {
       writer.write(s);
   public void print(char []c) throws IOException {
       writer.write(c);
```

```
public void close() throws IOException {
    writer.close();
}
```

#### 8.2 dates

```
// Time - Leap years
// A[i] has the accumulated number of days from months previous to i
const int A[13] = { 0, 0, 31, 59, 90, 120, 151, 181, 212, 243, 273, 304,
    334 };
// same as A, but for a leap year
const int B[13] = { 0, 0, 31, 60, 91, 121, 152, 182, 213, 244, 274, 305,
    335 };
// returns number of leap years up to, and including, y
int leap_vears(int y) { return y / 4 - y / 100 + y / 400; }
bool is_leap(int y) { return y % 400 == 0 || (y % 4 == 0 && y % 100 !=
    0); }
// number of days in blocks of years
const int p400 = 400*365 + leap_years(400);
const int p100 = 100*365 + leap_years(100);
const int p4 = 4*365 + 1;
const int p1 = 365;
int date_to_days(int d, int m, int y)
 return (y - 1) * 365 + leap_years(y - 1) + (is_leap(y) ? B[m] : A[m]) +
      d:
void days_to_date(int days, int &d, int &m, int &y)
 bool top100; // are we in the top 100 years of a 400 block?
 bool top4; // are we in the top 4 years of a 100 block?
 bool top1; // are we in the top year of a 4 block?
 y = 1;
 top100 = top4 = top1 = false;
 y += ((days-1) / p400) * 400;
 d = (days-1) \% p400 + 1;
```

```
if (d > p100*3) top100 = true, d -= 3*p100, y += 300;
else y += ((d-1) / p100) * 100, d = (d-1) % p100 + 1;

if (d > p4*24) top4 = true, d -= 24*p4, y += 24*4;
else y += ((d-1) / p4) * 4, d = (d-1) % p4 + 1;

if (d > p1*3) top1 = true, d -= p1*3, y += 3;
else y += (d-1) / p1, d = (d-1) % p1 + 1;

const int *ac = top1 && (!top4 || top100) ? B : A;
for (m = 1; m < 12; ++m) if (d <= ac[m + 1]) break;
d -= ac[m];
}</pre>
```

#### 8.3 fraction

```
struct frac{
  long long x, y;
  frac(long long a, long long b) {
    long long g = __gcd(a, b);
    x = a / g;
    y = b / g;
  }
  bool operator < (const frac &o) const {
    return (x * o.y < y * o.x);
  }
};</pre>
```

#### 8.4 io

```
// taken from :
    https://github.com/lbv/pc-code/blob/master/solved/c-e/diablo/diablo.cpp
// this is very fast as well :
    https://github.com/lbv/pc-code/blob/master/code/input.cpp

typedef unsigned int u32;
#define BUF 524288
struct Reader {
    char buf[BUF]; char b; int bi, bz;
    Reader() { bi=bz=0; read(); }
```

```
void read() {
   if (bi==bz) { bi=0; bz = fread(buf, 1, BUF, stdin); }
   b = bz ? buf[bi++] : 0; }
void skip() { while (b > 0 && b <= 32) read(); }
u32 next_u32() {
   u32 v = 0; for (skip(); b > 32; read()) v = v*10 + b-48; return v; }
int next_int() {
   int v = 0; bool s = false;
   skip(); if (b == '-') { s = true; read(); }
   for (; 48<=b&&b<=57; read()) v = v*10 + b-48; return s ? -v : v; }
char next_char() { skip(); char c = b; read(); return c; }
};</pre>
```

# 9 Number theory

#### 9.1 convolution

```
typedef long long int LL;
typedef pair<LL, LL> PLL;
inline bool is_pow2(LL x) {
 return (x & (x-1)) == 0;
}
inline int ceil_log2(LL x) {
 int ans = 0;
 --x;
 while (x != 0) {
   x >>= 1;
   ans++;
 }
 return ans;
/* Returns the convolution of the two given vectors in time proportional
    to n*log(n).
 * The number of roots of unity to use nroots_unity must be set so that
     the product of the first
 * nroots_unity primes of the vector nth_roots_unity is greater than the
     maximum value of the
 * convolution. Never use sizes of vectors bigger than 2^24, if you need
     to change the values of
```

```
* the nth roots of unity to appropriate primes for those sizes.
vector<LL> convolve(const vector<LL> &a, const vector<LL> &b, int
    nroots_unity = 2) {
 int N = 1 \ll ceil_log2(a.size() + b.size());
 vector<LL> ans(N,0), fA(N), fB(N), fC(N);
 LL modulo = 1;
 for (int times = 0; times < nroots_unity; times++) {</pre>
   fill(fA.begin(), fA.end(), 0);
   fill(fB.begin(), fB.end(), 0);
   for (int i = 0; i < a.size(); i++) fA[i] = a[i];
   for (int i = 0; i < b.size(); i++) fB[i] = b[i];</pre>
   LL prime = nth_roots_unity[times].first;
   LL inv_modulo = mod_inv(modulo % prime, prime);
   LL normalize = mod_inv(N, prime);
   ntfft(fA, 1, nth_roots_unity[times]);
   ntfft(fB, 1, nth_roots_unity[times]);
   for (int i = 0; i < N; i++) fC[i] = (fA[i] * fB[i]) % prime;</pre>
   ntfft(fC, -1, nth_roots_unity[times]);
   for (int i = 0; i < N; i++) {</pre>
     LL curr = (fC[i] * normalize) % prime;
     LL k = (curr - (ans[i] % prime) + prime) % prime;
     k = (k * inv_modulo) % prime;
     ans[i] += modulo * k;
   modulo *= prime;
 return ans;
```

#### 9.2 crt

```
/**
 * Chinese remainder theorem.
 * Find z such that z % x[i] = a[i] for all i.
 * */
long long crt(vector<long long> &a, vector<long long> &x) {
 long long z = 0;
 long long n = 1;
 for (int i = 0; i < x.size(); ++i)
    n *= x[i];

for (int i = 0; i < a.size(); ++i) {</pre>
```

```
long long tmp = (a[i] * (n / x[i])) % n;
tmp = (tmp * mod_inv(n / x[i], x[i])) % n;
z = (z + tmp) % n;
}
return (z + n) % n;
```

### 9.3 discrete logarithm

```
// Computes x which a \hat{x} = b \mod n.
long long d_log(long long a, long long b, long long n) {
 long long m = ceil(sqrt(n));
 long long aj = 1;
 map<long long, long long> M;
 for (int i = 0: i < m: ++i) {</pre>
   if (!M.count(aj))
     M[aj] = i;
   aj = (aj * a) % n;
 long long coef = mod_pow(a, n - 2, n);
 coef = mod_pow(coef, m, n);
 // coef = a ^{-} (-m)
 long long gamma = b;
 for (int i = 0; i < m; ++i) {</pre>
   if (M.count(gamma)) {
     return i * m + M[gamma];
   } else {
     gamma = (gamma * coef) % n;
   }
 }
 return -1;
```

### 9.4 ext euclidean

```
void ext_euclid(long long a, long long b, long long &x, long long &y,
    long long &g) {
```

```
x = 0, y = 1, g = b;
long long m, n, q, r;
for (long long u = 1, v = 0; a != 0; g = a, a = r) {
  q = g / a, r = g % a;
  m = x - u * q, n = y - v * q;
  x = u, y = v, u = m, v = n;
}
```

# 9.5 highest exponent factorial

```
int highest_exponent(int p, const int &n){
  int ans = 0;
  int t = p;
  while(t <= n){
    ans += n/t;
    t*=p;
  }
  return ans;
}</pre>
```

### 9.6 miller rabin

```
const int rounds = 20;
// checks whether a is a witness that n is not prime, 1 < a < n
bool witness(long long a, long long n) {
 // check as in Miller Rabin Primality Test described
 long long u = n - 1;
 int t = 0;
 while (u % 2 == 0) {
   t++:
   u >>= 1;
 long long next = mod_pow(a, u, n);
 if (next == 1) return false;
 long long last;
 for (int i = 0; i < t; ++i) {</pre>
   last = next:
   next = mod_mul(last, last, n);
   if (next == 1) {
```

```
return last != n - 1;
   }
 }
 return next != 1;
// Checks if a number is prime with prob 1 - 1 / (2 ^ it)
// D(miller_rabin(999999999999997LL) == 1);
// D(miller_rabin(999999999971LL) == 1);
// D(miller_rabin(7907) == 1);
bool miller_rabin(long long n, int it = rounds) {
 if (n <= 1) return false;</pre>
 if (n == 2) return true;
 if (n % 2 == 0) return false;
 for (int i = 0; i < it; ++i) {</pre>
   long long a = rand() \% (n - 1) + 1;
   if (witness(a, n)) {
     return false:
   }
 }
 return true;
```

#### 9.7 mod inv

```
long long mod_inv(long long n, long long m) {
  long long x, y, gcd;
  ext_euclid(n, m, x, y, gcd);
  if (gcd != 1)
    return 0;
  return (x + m) % m;
}
```

#### 9.8 mod mul

```
// Computes (a * b) % mod
long long mod_mul(long long a, long long b, long long mod) {
  long long x = 0, y = a % mod;
  while (b > 0) {
   if (b & 1)
```

```
x = (x + y) % mod;
y = (y * 2) % mod;
b /= 2;
}
return x % mod;
}
```

#### 9.9 mod pow

```
// Computes ( a ^ exp ) % mod.
long long mod_pow(long long a, long long exp, long long mod) {
  long long ans = 1;
  while (exp > 0) {
    if (exp & 1)
      ans = mod_mul(ans, a, mod);
    a = mod_mul(a, a, mod);
    exp >>= 1;
  }
  return ans;
}
```

#### 9.10 number theoretic transform

```
typedef long long int LL;
typedef pair<LL, LL> PLL;

/* The following vector of pairs contains pairs (prime, generator)
  * where the prime has an Nth root of unity for N being a power of two.
  * The generator is a number g s.t g^(p-1)=1 (mod p)
  * but is different from 1 for all smaller powers */
vector<PLL> nth_roots_unity {
  {1224736769,330732430},{1711276033,927759239},{167772161,167489322},
  {469762049,343261969},{754974721,643797295},{1107296257,883865065}};

PLL ext_euclid(LL a, LL b) {
  if (b == 0)
    return make_pair(1,0);
   pair<LL,LL> rc = ext_euclid(b, a % b);
   return make_pair(rc.second, rc.first - (a / b) * rc.second);
}
```

```
//returns -1 if there is no unique modular inverse
LL mod_inv(LL x, LL modulo) {
 PLL p = ext_euclid(x, modulo);
 if ( (p.first * x + p.second * modulo) != 1 )
   return -1;
 return (p.first+modulo) % modulo;
}
//Number theory fft. The size of a must be a power of 2
void ntfft(vector<LL> &a, int dir, const PLL &root_unity) {
 int n = a.size();
 LL prime = root_unity.first;
 LL basew = mod_pow(root_unity.second, (prime-1) / n, prime);
 if (dir < 0) basew = mod_inv(basew, prime);</pre>
 for (int m = n; m >= 2; m >>= 1) {
   int mh = m >> 1;
   LL w = 1:
   for (int i = 0: i < mh: i++) {</pre>
     for (int j = i; j < n; j += m) {
       int k = j + mh;
       LL x = (a[j] - a[k] + prime) % prime;
       a[j] = (a[j] + a[k]) \% prime;
       a[k] = (w * x) % prime;
     w = (w * basew) % prime;
   basew = (basew * basew) % prime;
 }
 int i = 0;
 for (int j = 1; j < n - 1; j++) {
   for (int k = n >> 1; k > (i ^= k); k >>= 1);
   if (j < i) swap(a[i], a[j]);</pre>
 }
}
```

# 9.11 pollard rho factorize

```
long long pollard_rho(long long n) {
  long long x, y, i = 1, k = 2, d;
  x = y = rand() % n;
  while (1) {
    ++i;
```

```
x = mod_mul(x, x, n);
   x += 2;
   if (x \ge n) x = n:
   if (x == y) return 1;
   d = \_gcd(abs(x - y), n);
   if (d != 1) return d;
   if (i == k) {
    y = x;
     k *= 2;
 return 1;
// Returns a list with the prime divisors of n
vector<long long> factorize(long long n) {
 vector<long long> ans;
 if (n == 1)
   return ans;
 if (miller_rabin(n)) {
   ans.push_back(n);
 } else {
   long long d = 1;
   while (d == 1)
     d = pollard_rho(n);
   vector<long long> dd = factorize(d);
   ans = factorize(n / d);
   for (int i = 0; i < dd.size(); ++i)</pre>
     ans.push_back(dd[i]);
 return ans;
```

# 9.12 primes

```
namespace primes {
  const int MP = 100001;
  bool sieve[MP];
  long long primes[MP];
  int num_p;
  void fill_sieve() {
   num_p = 0;
```

```
sieve[0] = sieve[1] = true;
  for (long long i = 2; i < MP; ++i) {
   if (!sieve[i]) {
     primes[num_p++] = i;
     for (long long j = i * i; j < MP; j += i)
       sieve[j] = true;
   }
 }
}
// Finds prime numbers between a and b, using basic primes up to sqrt(b)
vector<long long> seg_sieve(long long a, long long b) {
  long long ant = a;
  a = max(a, 3LL);
  vector<bool> pmap(b - a + 1);
  long long sqrt_b = sqrt(b);
  for (int i = 0; i < num_p; ++i) {</pre>
   long long p = primes[i];
   if (p > sqrt_b) break;
   long long j = (a + p - 1) / p;
   for (long long v = (j == 1) ? p + p : j * p; v <= b; v += p) {
     pmap[v - a] = true;
   }
  }
  vector<long long> ans;
  if (ant == 2) ans.push_back(2);
  int start = a % 2 ? 0 : 1;
  for (int i = start, I = b - a + 1; i < I; i += 2)</pre>
   if (pmap[i] == false)
     ans.push_back(a + i);
  return ans;
}
vector<pair<int, int>> factor(int n) {
  vector<pair<int, int>> ans;
  if (n == 0) return ans;
 for (int i = 0; primes[i] * primes[i] <= n; ++i) {</pre>
   if ((n % primes[i]) == 0) {
     int expo = 0;
     while ((n % primes[i]) == 0) {
       expo++;
       n /= primes[i];
     ans.emplace_back(primes[i], expo);
```

```
if (n > 1) {
    ans.emplace_back(n, 1);
}
return ans;
}
```

#### 9.13 totient sieve

```
for (int i = 1; i < MN; i++)
  phi[i] = i;

for (int i = 1; i < MN; i++)
  if (!sieve[i]) // is prime
  for (int j = i; j < MN; j += i)
    phi[j] -= phi[j] / i;</pre>
```

### 9.14 totient

```
long long totient(long long n) {
   if (n == 1) return 0;
   long long ans = n;
   for (int i = 0; primes[i] * primes[i] <= n; ++i) {
      if ((n % primes[i]) == 0) {
        while ((n % primes[i]) == 0) n /= primes[i];
        ans -= ans / primes[i];
    }
   if (n > 1) {
      ans -= ans / n;
   }
   return ans;
}
```

# 10 Strings

#### 10.1 Incremental Aho Corasick

```
class IncrementalAhoCorasic {
 static const int Alphabets = 26;
 static const int AlphabetBase = 'a';
 struct Node {
   Node *fail:
   Node *next[Alphabets];
   int sum;
   Node() : fail(NULL), next{}, sum(0) { }
 };
 struct String {
   string str;
   int sign;
 };
public:
 //totalLen = sum of (len + 1)
 void init(int totalLen) {
   nodes.resize(totalLen);
   nNodes = 0:
   strings.clear();
   roots.clear();
   sizes.clear();
   que.resize(totalLen);
 void insert(const string &str, int sign) {
   strings.push_back(String{ str, sign });
   roots.push_back(nodes.data() + nNodes);
   sizes.push_back(1);
   nNodes += (int)str.size() + 1;
   auto check = [&]() { return sizes.size() > 1 && sizes.end()[-1] ==
        sizes.end()[-2]; };
   if(!check())
     makePMA(strings.end() - 1, strings.end(), roots.back(), que);
   while(check()) {
     int m = sizes.back();
     roots.pop_back();
     sizes.pop_back();
     sizes.back() += m;
```

```
if(!check())
       makePMA(strings.end() - m * 2, strings.end(), roots.back(), que);
   }
 }
 int match(const string &str) const {
   int res = 0;
   for(const Node *t : roots)
     res += matchPMA(t, str);
   return res;
private:
  static void makePMA(vector<String>::const_iterator begin,
     vector<String>::const_iterator end, Node *nodes, vector<Node*>
     &que) {
   int nNodes = 0;
   Node *root = new(&nodes[nNodes ++]) Node();
   for(auto it = begin; it != end; ++ it) {
     Node *t = root;
     for(char c : it->str) {
       Node *&n = t->next[c - AlphabetBase];
       if(n == nullptr)
         n = new(&nodes[nNodes ++]) Node();
       t = n:
     }
     t->sum += it->sign;
   int qt = 0;
   for(Node *&n : root->next) {
     if(n != nullptr) {
       n->fail = root:
       que[qt ++] = n;
     } else {
       n = root;
     }
   for(int qh = 0; qh != qt; ++ qh) {
     Node *t = que[qh];
     int a = 0;
     for(Node *n : t->next) {
       if(n != nullptr) {
         que[qt ++] = n;
         Node *r = t- fail;
         while(r->next[a] == nullptr)
```

```
r = r - fail;
         n->fail = r->next[a];
         n->sum += r->next[a]->sum:
       ++ a;
   }
 }
  static int matchPMA(const Node *t, const string &str) {
   int res = 0:
   for(char c : str) {
     int a = c - AlphabetBase;
     while(t->next[a] == nullptr)
       t = t->fail;
     t = t-\text{next[a]}:
     res += t->sum;
   return res;
 }
 vector<Node> nodes;
 int nNodes:
 vector<String> strings;
 vector<Node*> roots:
 vector<int> sizes:
 vector<Node*> que;
};
int main() {
 int m:
 while(~scanf("%d", &m)) {
   IncrementalAhoCorasic iac;
   iac.init(600000);
   rep(i, m) {
     int ty;
     char s[300001];
     scanf("%d%s", &ty, s);
     if(ty == 1) {
       iac.insert(s, +1);
     } else if(ty == 2) {
       iac.insert(s, -1);
     } else if(ty == 3) {
       int ans = iac.match(s);
```

```
printf("%d\n", ans);
    fflush(stdout);
} else {
    abort();
}
}
return 0;
}
```

# 10.2 minimal string rotation

```
// Lexicographically minimal string rotation
int lmsr() {
 string s;
 cin >> s;
 int n = s.size();
 s += s;
 vector<int> f(s.size(), -1);
 int k = 0;
 for (int j = 1; j < 2 * n; ++j) {
   int i = f[j - k - 1];
   while (i != -1 && s[j] != s[k + i + 1]) {
     if (s[j] < s[k + i + 1])
      k = j - i - 1;
     i = f[i];
   if (i == -1 \&\& s[j] != s[k + i + 1]) {
     if (s[j] < s[k + i + 1]) {
       k = j;
     f[j - k] = -1;
   } else {
     f[j - k] = i + 1;
 }
 return k;
```

# 10.3 suffix array

```
/**
 * 0 (n log^2 (n))
 * See http://web.stanford.edu/class/cs97si/suffix-array.pdf for reference
 * */
struct entry{
  int a, b, p;
  entry(){}
  entry(int x, int y, int z): a(x), b(y), p(z){}
 bool operator < (const entry &o) const {</pre>
    return (a == o.a)? (b == o.b)? (p < o.p): (b < o.b): (a < o.a);
};
struct SuffixArray{
  const int N;
  string s;
  vector<vector<int> > P;
  vector<entry> M;
  SuffixArray(const string &s): N(s.length()), s(s), P(1, vector<int>
      (N, O)), M(N) {
    for (int i = 0; i < N; ++i)</pre>
     P[0][i] = (int) s[i];
    for (int skip = 1, level = 1; skip < N; skip *= 2, level++) {</pre>
     P.push_back(vector<int>(N, 0));
     for (int i = 0 ; i < N; ++i) {</pre>
       int next = ((i + skip) < N) ? P[level - 1][i + skip] : -10000;
       M[i] = entry(P[level - 1][i], next, i);
     sort(M.begin(), M.end());
     for (int i = 0; i < N; ++i)</pre>
       P[level][M[i].p] = (i > 0 \text{ and } M[i].a == M[i - 1].a \text{ and } M[i].b ==
            M[i - 1].b) ? P[level][M[i - 1].p] : i;
   }
  }
  vector<int> getSuffixArray(){
    vector<int> &rank = P.back();
    vector<pair<int, int> > inv(rank.size());
    for (int i = 0; i < rank.size(); ++i)</pre>
     inv[i] = make_pair(rank[i], i);
    sort(inv.begin(), inv.end());
    vector<int> sa(rank.size());
```

```
for (int i = 0; i < rank.size(); ++i)</pre>
     sa[i] = inv[i].second;
   return sa:
 // returns the length of the longest common prefix of s[i...L-1] and
      s[j...L-1]
  int lcp(int i, int j) {
   int len = 0;
   if (i == j) return N - i;
   for (int k = P.size() - 1; k >= 0 && i < N && j < N; --k) {
     if (P[k][i] == P[k][j]) {
       i += 1 << k;
       j += 1 << k;
       len += 1 << k;
   return len;
 }
};
```

#### 10.4 suffix automaton

```
/*
 * Suffix automaton:
 * This implementation was extended to maintain (online) the
 * number of different substrings. This is equivalent to compute
 * the number of paths from the initial state to all the other
 * states.
 * The overall complexity is O(n)
 * can be tested here:
     https://www.urionlinejudge.com.br/judge/en/problems/view/1530
 * */
struct state {
 int len, link;
 long long num_paths;
 map<int, int> next;
};
const int MN = 200011;
state sa[MN << 1];
```

```
int sz, last;
long long tot_paths;
void sa_init() {
 sz = 1;
 last = 0;
 sa[0].len = 0;
 sa[0].link = -1;
 sa[0].next.clear();
 sa[0].num_paths = 1;
 tot_paths = 0;
}
void sa_extend(int c) {
 int cur = sz++;
  sa[cur].len = sa[last].len + 1;
 sa[cur].next.clear();
 sa[cur].num_paths = 0;
 int p;
 for (p = last; p != -1 && !sa[p].next.count(c); p = sa[p].link) {
   sa[p].next[c] = cur;
   sa[cur].num_paths += sa[p].num_paths;
   tot_paths += sa[p].num_paths;
 }
 if (p == -1) {
   sa[cur].link = 0;
 } else {
   int q = sa[p].next[c];
   if (sa[p].len + 1 == sa[q].len) {
     sa[cur].link = q;
   } else {
     int clone = sz++;
     sa[clone].len = sa[p].len + 1;
     sa[clone].next = sa[q].next;
     sa[clone].num_paths = 0;
     sa[clone].link = sa[q].link;
     for (; p!= -1 && sa[p].next[c] == q; p = sa[p].link) {
       sa[p].next[c] = clone;
       sa[q].num_paths -= sa[p].num_paths;
       sa[clone].num_paths += sa[p].num_paths;
     sa[q].link = sa[cur].link = clone;
   }
 }
```

```
last = cur;
```

# 10.5 z algorithm

```
using namespace std;
#include<bits/stdc++.h>
vector<int> compute_z(const string &s){
 int n = s.size();
 vector<int> z(n,0);
 int 1,r;
 r = 1 = 0;
 for(int i = 1; i < n; ++i){</pre>
   if(i > r) {
     1 = r = i:
     while (r < n \text{ and } s[r - 1] == s[r])r++;
     z[i] = r - 1;r--;
   }else{
     int k = i-1;
     if(z[k] < r - i +1) z[i] = z[k];
     else {
       l = i;
       while (r < n \text{ and } s[r - 1] == s[r])r++;
       z[i] = r - 1;r--;
     }
 return z;
int main(){
 //string line;cin>>line;
  string line = "alfalfa";
  vector<int> z = compute_z(line);
 for(int i = 0; i < z.size(); ++i ){</pre>
   if(i)cout<<" ";
   cout<<z[i];
 }
  cout << end1;
```

```
// must print "0 0 0 4 0 0 1"

return 0;
}
```