

# 7. Math

# Fibonacci

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```
int fib(int n) {  
    if (n <= 1) return n;  
    return fib(n - 1) + fib(n - 2);  
}
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}
```

*Easy Right?*

Don't do this!  
**EVER!**

# Fibonacci

```
int fib(int n)
    if (n <= 1) return n;
    int f2 = 0, f1 = 1;
    for (int i = 2; i <= n; i++) {
        int prev_f1 = f1;
        f1 += f2; f2 = prev_f1;
    }
    return f1;
```



Use  
memoization to  
get linear time

# Fibonacci

$$F_0 = 0, F_1 = 1, F_n = F_{n-1} + F_{n-2}$$

$$(F_{n-2}, F_{n-1}) \times \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} = (F_{n-1}, F_n)$$

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Use fast exponentiation to make this  $O(\log n)$ !



# Fibonacci

$$F_0 = 0, F_1 = 1, F_n = F_{n-1} + F_{n-2}$$

This is also cool to know (see [here](#))

$$\text{Fib}(n) = ((1.6180339\dots)^n - (-0.6180339\dots)^n) / 2.236067977\dots$$

But it will give wrong results for  $n > 100$  (rounding errors)

# Primes

Useful in

1. cryptography
2. modulo arithmetic
3. as part of other algorithms (e.g. find divisors)

# Primes

To check if a number is prime:

```
boolean isPrime(int n):  
  
    for (int i = 2; i < n; i++) {  
        if (n % i == 0) return false;  
    }  
  
    return true
```

# Primes

To check if a number is prime:

```
boolean isPrime(int n):  
  
    for (int i = 2; i < n / 2; i++) {  
        if (n % i == 0) return false;  
    }  
  
    return true
```

# Primes

To check if a number is prime:

```
boolean isPrime(int n):  
    for (int i = 2; i <= sqrt(n); i++) {  
        if (n % i == 0) return false;  
    }  
    return true
```

# Primes

To check if a number is prime:

```
boolean isPrime(int n):
```

```
    if (n == 2) return true;
```

```
    for (int i = 3; i <= sqrt(n); i+=2) {
```

```
        if (n % i == 0) return false;
```

```
    return true;
```

# Primes

To check if a number is prime:

```
boolean isPrime(int n):  
  
    for (int i : smallerPrimes(n)) {  
        if (n % i == 0) return false;  
    }  
  
    return true;
```

# Primes

To find all primes smaller than n:

```
void sieve(int n):  
    int primes[] = new int[n];  
    for (int i = 2; i < n; i++)  
        if (primes[i] == 0)  
            mark all multiples
```



# Primes

To find all primes smaller than n:

```
void sieve(int n):
```

```
    char primes[] = new char[n];
```

```
    for (int i = 2; i < n; i++)
```

```
        if (primes[i] == 0)
```

```
            mark all multiples
```

# Primes

To find all primes smaller than n:

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void sieve(int n):
```

```
    char primes[] = new char[n];
```

```
    for (int i = 2; i < n; i++)
```

```
        if (primes[i] == 0)
```

```
            mark all multiples
```

Could use bit optimizations for 8x less memory (a bit tricky to implement).

# Primes

To find all primes smaller than  $n$ :

## 1. naive approach

- a.  $O(\sqrt{N} * N)$ ,  $O(\log n)$  memory
- b.  $O(\sqrt{N} / \log(N) * N)$ ,  $O(\log n)$  memory

## 2. sieve

- a.  $O(\log(\log(N)) * N)$ ,  $O(n)$  memory

# Gaussian Elimination


“Systems of linear equations arise in 75% of scientific computing problems”

S. Skienna

# Gaussian Elimination

1. Start with row  $i = 0$ , col  $j = 0$
2. Find first non-zero on col  $j = A[k][j]$
3. Swap  $k$  and  $i$
4. Divide row  $i$  by the  $A[i][j]$
5. Update all rows below row  $i$  and to the right of  $j$ 
  - a. subtract (row  $i$ ) \* (first nnz element of row) to update

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- 
- Row  
echelon  
form

# Gaussian Elimination

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5. Update all rows below row  $i$  and to the right of  $j$ 
  - a. subtract (row  $i$ ) \* (first nnz element of row) to update
6. Start with row  $i = 0$ 
  - a. find the first non-zero element ( $j$ )
  - b. compute the corresponding variable by
    - i.  $x[j] = A[i][j] - \text{sum}(k = j + 1 \dots m, x[k] * A[k][j])$

# Gaussian Elimination

1. Gaussian elimination is OK for small systems
2. But there are some pretty big systems
  - a. [The Web](#)
  - b. Large physics simulation



# Iterative Algorithms

Iterative solvers:

1. Guess a solution
2. Check if it's good
3. If not, guess another solution, based on
  - a. previous guesses
  - b. how far we are from the solution (residual)

# Conjugate Gradient

```
for i in xrange(10 ** 3):  
    Ap = a * p  
    alpha = rsold / (p.T * Ap)  
    x = x + multiply(alpha, p)  
    r = r - multiply(alpha, Ap)  
    rsnew = r.T * r  
    if rsnew < EPS ** 2:  
        return x  
    p = r + multiply(rsnew / rsold, p)  
    rsold = rsnew
```

# Sparse Algebra

1. Don't store the matrix as value, row, col
2. Store only nonzero values
3. Sparse systems arise in many situations
  - a. power grids
  - b. ocean modelling

# Sparse Algebra

Pretty important research topic

- trade-offs between:
  - storage
  - computation, communication overhead
- dedicated architectures
- various formats
  - Block/Bit Compressed Row/Column,

# Linear Algebra

It's usually best **NOT** to write your own:

1. [Intel MKL](#), [CUSP](#), [LAPACK](#), [ACML](#)
2. These (can) take into account:
  - a. various cache optimisations
  - b. sparsity
  - c. preconditioning

# Recap

# Overview

## 1. Complete Search

- a. First approach
- b. Bad complexity
- c. Can sometimes use pruning effectively

## 2. Divide and Conquer

- a. Useful Algo Design paradigm
- b. Search, Sort etc.

# Overview

## 3. Greedy

- a. Pick “locally” best option at each step
- b. Usually really fast
- c. May be wrong (unless greedy property holds)

## 4. Dynamic Programming

- a. Can speedup algorithms by effective caching



# Overview

## 5. Graphs

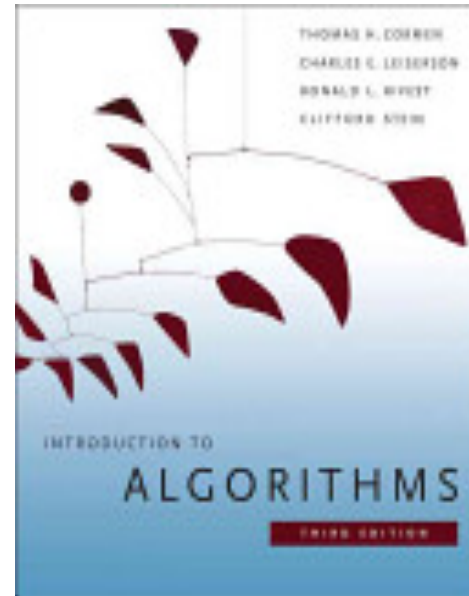
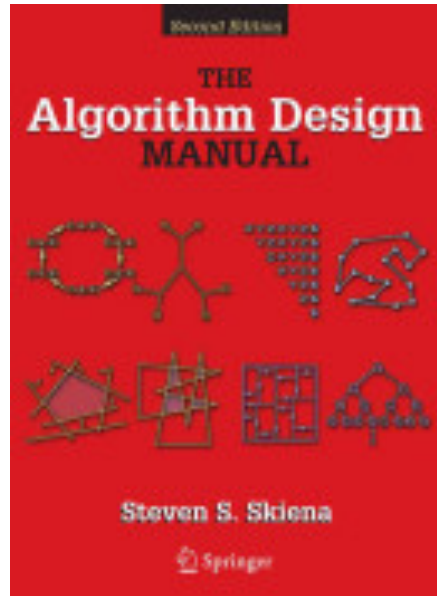
- a. Natural database for relations
- b. Used in many areas of CS
  - i. scheduling, logic problems, assignment etc.
- c. Classic algorithms
  - i. DFS, BFS
  - ii. Topological Sort
  - iii. Minimum Spanning Tree
  - iv. Dijkstra etc.

# Overview

## 6. Geometry and Maths

- a. Heavily used in scientific computing
- b. Very nice algorithms (e.g. Graham Scan)
- c. Often tricky to implement
  - i. floating point errors
  - ii. edge cases etc.

# Where To?



# Where To?

## Practice:

1. <https://projecteuler.net/problems>
2. <http://uhunt.felix-halim.net/>
3. [http://www.topcoder.com/tc?d1=tutorials&d2=alg\\_index&module=Static](http://www.topcoder.com/tc?d1=tutorials&d2=alg_index&module=Static)
4. <http://codeforces.com/>
5. <http://code.google.com/codejam/>

# Questions ?

Feedback?