# Team Notebook

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Contents		4	l Ge	Geometry 11 8 Misc		SC .	18	
			4.1	center 2 points $+$ radious $\dots$	11	8.1	Template Java	18
1	Algorithms	2	4.2	closest pair problem	11	8.2	dates	19
	1.1 Mo's algorithm on trees	2	4.3			8.3	fraction	
	1.2 Mo's algorithm	$\begin{bmatrix} 2\\2 \end{bmatrix}$ 5	6 Gr	raphs	11	8.4	io	19
			5.1	SCC kosaraju	1.0	) Nu	mber theory	<b>2</b> 0
<b>2</b>	DP Optimizations	3	5.2 5.3			9.1	convolution	20
	2.1 convex hull trick	3	5.4			9.2	diophantine equations	
	2.2 divide and conquer	3	5.5			9.3	discrete logarithm	21
	2.3 dp on trees	4	5.6		13	9.4	ext euclidean	21
			5.7		14	9.5	highest exponent factorial	21
3	Data structures	4	5.8	minimum path cover in DAG	14	9.6	miller rabin	21
	3.1 STL Treap	4	5.9	planar graph (euler)	14	9.7	mod integer	21
	3.2 STL order statistics tree II	5	5.10	0 query with lca	14	9.8	mod inv	
	3.3 STL order statistics tree	5		1 tarjan scc		9.9	$\mod \operatorname{mul} \ldots \ldots \ldots \ldots \ldots \ldots \ldots$	22
	3.4 binary index tree	5	5.13	2 two sat (with kosaraju)	15		$\mod pow  \dots \dots \dots \dots \dots \dots \dots \dots$	
		5 6	6 Ma	ath	16		number theoretic transform	
	3.6 heavy light decomposition	6	6.1	<del></del>			primes	
	3.7 persistent array	6	6.2			9.13	totient sieve	23
	3.8 persistent seg tree	6	6.3			9.14	totient	23
	3.9 persistent trie	7	6.4		17			•
		0	6.5			lo Stri		23
	3.10 segment tree	0	6.6				Incremental Aho Corasick	
	3.11 sparse table	9	6.7	sigma function	18		minimal string rotation	
	3.12 splay tree	9					suffix array	
	3.13 trie	10 7		atrix	18		suffix automaton	
	3.14 wavelet tree	10	7.1	matrix	18	10.5	z algorithm	25

## 1 Algorithms

### 1.1 Mo's algorithm on trees

```
problems:
   - https://codeforces.com/gym/101161 problem E
void flat(vector<vector<edge>> &g, vector<int> &a,
   vector<int> &le, vector<int> &ri, vector<int> &cost,
   int node. int pi. int &ts. int w) {
 cost[node] = w;
 le[node] = ts:
 a[ts] = node;
 ts++:
 for (auto e : g[node]) {
   if (e.to == pi) continue;
   flat(g, a, le, ri, cost, e.to, node, ts, e.w);
 ri[node] = ts;
 a[ts] = node:
 ts++;
* Case when the cost is in the edges.
void compute_queries(vector<vector<edge>> &g) {
 // g is undirected
 int n = g.size();
 lca_tree.init(g, 0);
 vector < int > a(2 * n), le(n), ri(n), cost(n):
 // a: nodes in the flatten array
 // le: left id of the given node
 // ri: right id of the given node
 // cost: cost of the edge from the node to the parent
 int ts = 0; // timestamp
 flat(g, a, le, ri, cost, 0, -1, ts, 0);
 int q; cin >> q;
 vector<query> queries(q);
 for (int i = 0: i < a: i++) {
   int u, v;
   cin >> u >> v;
   u--; v--;
   int lca = lca_tree.query(u, v);
```

```
if (le[u] > le[v])
    swap(u, v);
    queries[i].id = i;
    queries[i].lca = lca;
    queries[i].v = v;
    if (lca == u) {
        queries[i].a = le[u] + 1;
        queries[i].b = le[v];
    } else {
        queries[i].a = ri[u];
        queries[i].b = le[v];
}
}
solve_mo(queries, a, le, cost); // this is the usal algorithm
}
```

## 1.2 Mo's algorithm

```
const int MN = 5 * 100000 + 1:
const int SN = 708;
struct Querv {
  int a, b, id;
  Querv() {}
  Query(int x, int y, int i) : a(x), b(y), id(i) {}
  bool operator<(const Query &o) const {</pre>
   if (a / SN != o.a / SN) return a < o.a;
   return a / SN & 1 ? b < o.b : b > o.b:
 }
};
struct DS {
 DS() : {}
  void Insert(int x) {}
  void Erase(int x) {}
  long long Query() {}
Querv s[MN]:
int ans[MN];
DS active;
int main() {
```

```
int n:
 cin >> n:
 vector<int> a(n);
 for (auto &i : a) cin >> i:
 int a:
 cin >> q;
 for (int i = 0; i < q; ++i) {</pre>
   int b, e;
   cin >> b >> e;
   s[i] = Query(b, e, i);
 sort(s, s + q);
 int i = 0;
 int j = -1;
 for (int k = 0; k < (int)q; ++k) {
   int L = s[k].a;
   int R = s[k].b:
   while (j < R) active.Insert(a[++j]);</pre>
   while (j > R) active.Erase(a[j--]);
   while (i < L) active.Erase(a[i++]);</pre>
   while (i > L) active.Insert(a[--i]);
   ans[s[k].id] = active.Query();
 for (int i = 0; i < q; ++i) {</pre>
   cout << ans[i] << endl:</pre>
 return 0:
};
```

## 1.3 sliding window

```
/*
 * Given an array ARR and an integer K, the problem boils
    down to computing for each index i: min(ARR[i], ARR[i
    -1], ..., ARR[i-K+1]).
 * if mx == true, returns the maximun.
 * http://people.cs.uct.ac.za/~ksmith/articles/
    sliding_window_minimum.html
 * */
```

```
vector<int> sliding_window_minmax(vector<int> & ARR, int K,
    bool mx) {
 deque< pair<int, int> > window;
 vector<int> ans;
 for (int i = 0: i < ARR.size(): i++) {</pre>
   if (mx) {
     while (!window.empty() && window.back().first <= ARR[i</pre>
          1)
       window.pop_back();
     while (!window.emptv() && window.back().first >= ARR[i
          1)
       window.pop_back();
   window.push_back(make_pair(ARR[i], i));
   while(window.front().second <= i - K)</pre>
     window.pop_front();
   ans.push_back(window.front().first);
 return ans;
```

## 2 DP Optimizations

#### 2.1 convex hull trick

```
/**
 * Problems:
 * http://codeforces.com/problemset/problem/319/C
 * http://codeforces.com/contest/311/problem/B
 * https://csacademy.com/contest/archive/task/squared-ends
 * http://codeforces.com/contest/932/problem/F
 * */

struct line {
  long long m, b;
  line (long long a, long long c) : m(a), b(c) {}
  long long eval(long long x) {
    return m * x + b;
  }
};

long double inter(line a, line b) {
  long double den = a.m - b.m;
  long double num = b.b - a.b;
```

```
return num / den:
* min m_i * x_j + b_i, for all i.
      x i \le x \{i + 1\}
      m_i >= m_{i} + 1
struct ordered_cht {
 vector<line> ch;
  int idx: // id of last "best" in query
  ordered cht() {
   idx = 0;
  void insert_line(long long m, long long b) {
   line cur(m, b):
   // new line's slope is less than all the previous
   while (ch.size() > 1 &&
      (inter(cur, ch[ch.size() - 2]) >= inter(cur, ch[ch.
           size() - 1]))) {
       // f(x) is better in interval [inter(ch.back(), cur),
             inf)
       ch.pop_back();
   ch.push back(cur):
 long long eval(long long x) { // minimum
   // current x is greater than all the previous x,
   // if that is not the case we can make binary search.
   idx = min<int>(idx, ch.size() - 1);
   while (idx + 1 < (int)ch.size() && ch[idx + 1].eval(x) <=</pre>
         ch[idx].eval(x))
     idx++:
   return ch[idx].eval(x);
 }
};
* Dynammic convex hull trick
typedef long long int64;
typedef long double float128;
const int64 is_query = -(1LL<<62), inf = 1e18;</pre>
```

```
struct Line {
 int64 m, b;
 mutable function<const Line*()> succ:
 bool operator<(const Line& rhs) const {</pre>
   if (rhs.b != is_query) return m < rhs.m;</pre>
   const Line* s = succ();
   if (!s) return 0:
   int64 x = rhs.m:
   return b - s->b < (s->m - m) * x;
}:
struct HullDynamic : public multiset<Line> { // will
     maintain upper hull for maximum
  bool bad(iterator y) {
   auto z = next(v):
   if (v == begin()) {
     if (z == end()) return 0:
      return y->m == z->m && y->b <= z->b;
    auto x = prev(y);
    if (z == end()) return y \rightarrow m == x \rightarrow m \&\& y \rightarrow b <= x \rightarrow b;
   return (float128)(x\rightarrow b - y\rightarrow b)*(z\rightarrow m - y\rightarrow m) >= (float128)
         (y-b - z-b)*(y-m - x-m);
  void insert_line(int64 m, int64 b) {
   auto v = insert({ m, b });
   v->succ = [=] { return next(y) == end() ? 0 : &*next(y);
        }:
    if (bad(y)) { erase(y); return; }
    while (next(y) != end() && bad(next(y))) erase(next(y));
   while (y != begin() && bad(prev(y))) erase(prev(y));
 int64 eval(int64 x) {
   auto 1 = *lower_bound((Line) { x, is_query });
   return 1.m * x + 1.b:
};
```

#### 2.2 divide and conquer

```
/**
  * recurrence:
  * dp[k][i] = min dp[k-1][j] + c[i][j - 1], for all j > i;
  *
  * "comp" computes dp[k][i] for all i in O(n log n) (k is
      fixed)
```

```
* Problems:
* https://icpc.kattis.com/problems/branch
* http://codeforces.com/contest/321/problem/E
* */

void comp(int l, int r, int le, int re) {
   if (l > r) return;
   int mid = (l + r) >> 1;
   int best = max(mid + 1, le);
   dp[cur][mid] = dp[cur ^ 1][best] + cost(mid, best - 1);
   for (int i = best; i <= re; i++) {
     if (dp[cur][mid] > dp[cur ^ 1][i] + cost(mid, i - 1)) {
       best = i;
       dp[cur][mid] = dp[cur ^ 1][i] + cost(mid, i - 1);
   }
}

comp(l, mid - 1, le, best);
comp(mid + 1, r, best, re);
}
```

#### 2.3 dp on trees

```
/**
 * This trick is very useful when doing DP on trees,
     basically, you can save
 * the answer for each node as if it was the root of the
      tree. Partial results
 * are also stored in order to query subtrees (taking the
     root and exclude some
 * child).
 * problems:
 * - http://codeforces.com/gym/101161, problem I : Sky tax
 * - http://codeforces.com/contest/791/problem/D
struct edge {
 int to, p_id;
 edge (int a, int b) : to(a), p_id(b) {}
}:
struct state {
 bool seen:
 long long missing;
```

```
long long total;
 vector<long long> partial;
 state() { clear(): }
 void clear() {
   seen = false;
   missing = 0;
   total = 0;
   partial.clear();
void add_edge(int u, int v) {
 int id_u_v = g[u].size();
 int id_v_u = g[v].size();
 g[u].emplace_back(v, id_v_u); // id of the parent in the
      child's list (g[v][id] -> u)
 g[v].emplace_back(u, id_u_v); // id of the parent in the
      child's list (g[u][id] -> v)
int go(int node, int id_parent) {
 state &s = dp[node];
 if (!s.seen) {
   int ans = 1;
   s.partial.assign(g[node].size(), 0); // create the list
        of partial results.
   for (int i = 0; i < int(g[node].size()); i++) {</pre>
     int to = g[node][i].to;
     int pid = g[node][i].p_id;
     if (i != id_parent) {
      int tmp = go(to, pid);
      ans += tmp;
       s.partial[i] = tmp;
   s.missing = id_parent;
   s.total = ans:
   s.seen = true;
   return ans:
 } else {
   if (s.missing == id_parent) { // the same id_parent than
       before, so we can not complete the results yet
```

### 3 Data structures

### 3.1 STL Treap

```
#include <ext/rope> //header with rope
using namespace std;
using namespace __gnu_cxx; //namespace with rope and some
    additional stuff
int main()
   ios_base::sync_with_stdio(false);
   rope <int> v: //use as usual STL container
   int n, m;
   cin >> n >> m;
   for(int i = 1; i <= n; ++i)</pre>
       v.push_back(i); //initialization
   int 1. r:
   for(int i = 0: i < m: ++i)</pre>
       cin >> 1 >> r:
       --1, --r;
       rope \langle int \rangle cur = v.substr(1, r - 1 + 1);
       v.erase(1, r - 1 + 1);
       v.insert(v.mutable_begin(), cur);
   for(rope <int>::iterator it = v.mutable_begin(); it != v.
        mutable_end(); ++it)
       cout << *it << " ":
   return 0;
```

#### 3.2 STL order statistics tree II

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
typedef tree<int,null_type,less<int>,rb_tree_tag,
tree_order_statistics_node_update> order_set;
order_set X;
int get(int y) {
 int l=0.r=1e9+1:
 while(l<r) {</pre>
   int m=l+((r-l)>>1);
   if (m-X.order_of_key(m+1)<y)</pre>
     l=m+1;
   else
     r=m:
 }
 return 1;
main(){
 ios::sync_with_stdio(0);
 cin.tie(0);
 int n,m;
 cin>>n>>m;
 for(int i=0:i<m:i++) {</pre>
   char a;
   int b:
   cin>>a>>b;
   if(a=='L')
     cout<<get(b)<<endl;</pre>
   else
      X.insert(get(b));
}
/***
Input
20 7
L 5
D 5
T. 4
L 5
```

```
D 5
L 4
L 5
Output
5
4
6
4
7
***/
```

#### 3.3 STL order statistics tree

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <bits/stdc++.h>
using namespace __gnu_pbds;
using namespace std;
typedef
tree<
 pair<int,int>,
 null_type,
 less<pair<int,int>>,
 rb_tree_tag,
 tree_order_statistics_node_update>
ordered_set;
main()
   ios::sync_with_stdio(0);
   cin.tie(0):
   int n;
   int sz=0;
   cin>>n:
   vector<int> ans(n,0);
   ordered_set t;
   int x, y;
   for(int i=0;i<n;i++)</pre>
   {
       cin>>x>>y;
       ans[t.order_of_key({x,++sz})]++;
       t.insert({x,sz});
   for(int i=0;i<n;i++)</pre>
```

```
cout<<ans[i]<<'\n';
}
/***
Input
5
1 1
5 1
7 1
3 3
5 5

Output
1
2
1
0
****/</pre>
```

### 3.4 binary index tree

```
struct binary_index_tree {
   int n;
   int t[2 * N];

void add(int where, long long what) {
    for (where++; where <= n; where += where & -where) {
        t[where] += what;
    }
}

void add(int from, int to, long long what) {
   add(from, what);
   add(to + 1, -what);
}

long long query(int where) {
   long long sum = t[0];
   for (where++; where > 0; where -= where & -where) {
        sum += t[where];
    }
   return sum;
}
```

#### 3.5 dsu

```
struct Dsu {
 vector<int> p;
 Dsu(int n) {
   p.resize(n):
   for (int i = 0; i < n; i++) {</pre>
    p[i] = i:
 }
 int Find(int x) { return x == p[x] ? x : p[x] = Find(p[x])
      ; }
 int Join(int x, int y) {
   int px = Find(x), py = Find(y);
   if (px == py) return 0;
   p[px] = py;
   return 1;
 }
};
```

#### 3.6 heavy light decomposition

```
struct TreeDecomposition {
 vector<int> g[MAXN], c[MAXN];
 int s[MAXN]; // subtree size
 int p[MAXN]; // parent id
 int r[MAXN]; // chain root id
 int t[MAXN]; // index used in segtree/bit/...
 int d[MAXN]; // depht
 int ts:
 void dfs(int v. int f) {
   p[v] = f;
   s[v] = 1;
   if (f != -1) d[v] = d[f] + 1:
   else d[v] = 0;
   for (int i = 0; i < g[v].size(); ++i) {</pre>
     int w = g[v][i];
     if (w != f) {
      dfs(w, v);
       s[v] += s[w];
 void hld(int v, int f, int k) {
```

```
t[v] = ts++:
   c[k].push_back(v);
   r[v] = k:
   int x = 0, y = -1;
   for (int i = 0; i < g[v].size(); ++i) {</pre>
     int w = g[v][i];
     if (w != f) {
      if (s[w] > x) {
        x = s[w];
         v = w:
   if (y != -1) {
     hld(y, v, k);
   for (int i = 0; i < g[v].size(); ++i) {</pre>
    int w = g[v][i]:
     if (w != f && w != y) {
      hld(w, v, w);
   }
 }
 void init(int n) {
   for (int i = 0; i < n; ++i) {
     g[i].clear();
 }
 void add(int a, int b) {
   g[a].push_back(b);
   g[b].push_back(a);
 void build() {
   ts = 0;
   dfs(0, -1);
   hld(0, 0, 0):
};
```

### 3.7 persistent array

```
struct node {
  node *1, *r;
  int val;
```

```
node (int x) : 1(NULL), r(NULL), val(x) {}
 node (): 1(NULL), r(NULL), val(-1) {}
typedef node* pnode:
pnode update(pnode cur, int 1, int r, int at, int what) {
 pnode ans = new node();
 if (cur != NULL) {
   *ans = *cur:
 if (1 == r) {
   ans-> val = what:
   return ans;
 int m = (1 + r) >> 1;
 if (at <= m) ans-> 1 = update(ans-> 1, 1, m, at, what);
 else ans-> r = update(ans-> r, m + 1, r, at, what);
 return ans;
int get(pnode cur, int 1, int r, int at) {
 if (cur == NULL) return 0;
 if (1 == r) return cur-> val;
 int m = (1 + r) >> 1:
 if (at <= m) return get(cur-> 1, 1, m, at);
             return get(cur-> r, m + 1, r, at);
```

## 3.8 persistent seg tree

```
/**
 * Problems:
 * http://codeforces.com/contest/813/problem/E
 *
 * Important:
 * When using lazy propagation remembert to create new
 * versions for each push_down operation!!!
 * */

struct node {
   node *1, *r;
   long long acc;
   int flip;

   node (int x) : 1(NULL), r(NULL), acc(x), flip(0) {}
   node () : 1(NULL), r(NULL), acc(0), flip(0) {}
```

```
}:
typedef node* pnode;
pnode create(int 1, int r) {
 if (1 == r) return new node():
 pnode cur = new node();
 int m = (1 + r) >> 1;
 cur-> 1 = create(1, m);
 cur \rightarrow r = create(m + 1, r);
 return cur:
pnode copy_node(pnode cur) {
 pnode ans = new node();
 *ans = *cur;
 return ans:
}
void push_down(pnode cur, int 1, int r) {
 assert(cur);
 if (cur-> flip) {
   int len = r - l + 1;
   cur-> acc = len - cur-> acc;
   if (cur-> 1) {
     cur-> 1 = copy_node(cur-> 1);
     cur-> 1 -> flip ^= 1;
   if (cur-> r) {
     cur-> r = copy_node(cur-> r);
     cur-> r -> flip ^= 1;
   cur-> flip = 0;
}
int get_val(pnode cur) {
 assert(cur):
 assert((cur-> flip) == 0);
 if (cur) return cur-> acc:
 return 0:
}
pnode update(pnode cur, int 1, int r, int at, int what) {
 pnode ans = copy_node(cur);
 if (1 == r) {
   assert(1 == at);
   ans-> acc = what:
   ans-> flip = 0;
   return ans:
```

```
int m = (1 + r) >> 1:
 push_down(ans, 1, r);
 if (at <= m) ans-> 1 = update(ans-> 1, 1, m, at, what);
 else ans-> r = update(ans-> r, m + 1, r, at, what);
 push_down(ans-> 1, 1, m);
 push_down(ans-> r, m + 1, r);
 ans-> acc = get_val(ans-> 1) + get_val(ans-> r);
 return ans;
pnode flip(pnode cur, int 1, int r, int a, int b) {
 pnode ans = new node();
 if (cur != NULL) {
   *ans = *cur:
 if (1 > b | | r < a)
   return ans:
 if (1 >= a kk r <= b) {
   ans-> flip ^= 1;
   push_down(ans, 1, r);
   return ans:
 int m = (1 + r) >> 1:
 ans-> 1 = flip(ans-> 1, 1, m, a, b);
 ans-> r = flip(ans-> r, m + 1, r, a, b):
 push_down(ans-> 1, 1, m);
 push_down(ans-> r, m + 1, r);
 ans-> acc = get_val(ans-> 1) + get_val(ans-> r);
 return ans;
long long get_all(pnode cur, int 1, int r) {
 assert(cur):
 push_down(cur, 1, r);
 return cur-> acc:
void traverse(pnode cur, int 1, int r) {
 if (!cur) return:
 cout << 1 << " - " << r << " : " << (cur-> acc) << " " <<
      (cur-> flip) << endl:
 traverse(cur-> 1, 1, (1 + r) >> 1);
 traverse(cur-> 1, 1 + ((1 + r) >> 1), r):
```

#### 3.9 persistent trie

```
// both tries can be tested with the problem: http://
    codeforces.com/problemset/problem/916/D
// Persistent binary trie (BST for integers)
const int MD = 31:
struct node bin {
 node_bin *child[2];
 int val;
 node bin() : val(0) {
   child[0] = child[1] = NULL:
};
typedef node_bin* pnode_bin;
pnode_bin copy_node(pnode_bin cur) {
 pnode_bin ans = new node_bin();
 if (cur) *ans = *cur:
 return ans:
pnode_bin modify(pnode_bin cur, int key, int inc, int id =
    MD) {
 pnode_bin ans = copy_node(cur);
 ans->val += inc;
 if (id >= 0) {
   int to = (kev >> id) & 1:
   ans->child[to] = modify(ans->child[to], key, inc, id - 1)
 return ans;
int sum_smaller(pnode_bin cur, int key, int id = MD) {
 if (cur == NULL) return 0:
 if (id < 0) return 0: // strictly smaller
 // if (id == - 1) return cur->val; // smaller or equal
 int ans = 0;
 int to = (kev >> id) & 1:
 if (to) {
   if (cur->child[0]) ans += cur->child[0]->val;
   ans += sum smaller(cur->child[1], kev, id - 1):
 } else {
```

```
ans = sum_smaller(cur->child[0], key, id - 1);
 }
 return ans;
// Persistent trie for strings.
const int MAX_CHILD = 26;
struct node {
 node *child[MAX_CHILD];
 int val:
 node() : val(-1) {
   for (int i = 0; i < MAX_CHILD; i++) {</pre>
     child[i] = NULL;
 }
}:
typedef node* pnode:
pnode copy_node(pnode cur) {
 pnode ans = new node();
 if (cur) *ans = *cur;
 return ans:
pnode set_val(pnode cur, string &key, int val, int id = 0) {
 pnode ans = copv node(cur):
 if (id >= int(key.size())) {
   ans->val = val:
 } else {
   int t = kev[id] - 'a';
   ans->child[t] = set val(ans->child[t], kev, val, id + 1): }:
 }
 return ans;
pnode get(pnode cur, string &key, int id = 0) {
 if (id >= int(key.size()) || !cur)
   return cur:
 int t = kev[id] - 'a':
 return get(cur->child[t], key, id + 1);
```

### 3.10 segment tree

```
/**
 * Taken from: http://codeforces.com/blog/entry/18051
 * */
```

```
const int MN = 1e5; // limit for array size
struct seg_tree {
 int n; // array size
 int t[2 * MN]:
 seg_tree(int _n) : n(_n) {}
 void clear() {
   memset(t, 0, sizeof t):
 void build() { // build the tree
   for (int i = n - 1: i > 0: --i) t[i] = t[i << 1] + t[i]
        <<1|1];
 }
 // Single modification, range query.
 void modify(int p, int value) { // set value at position p
   for (t[p += n] = value; p > 1; p >>= 1) t[p>>1] = t[p] +
        t[p^1];
 }
 int query(int 1, int r) { // sum on interval [1, r)
   int res = 0:
   for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
    if (l&1) res += t[l++]:
    if (r&1) res += t[--r];
   return res;
// Range modification, single query.
void modify(int 1, int r, int value) {
 for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
   if (1&1) t[1++] += value;
   if (r&1) t[--r] += value:
 }
int query(int p) {
 int res = 0;
 for (p += n; p > 0; p >>= 1) res += t[p];
 return res;
/**
```

```
* If at some point after modifications we need to inspect
* elements in the array, we can push all the modifications
* leaves using the following code. After that we can just
* elements starting with index n. This way we reduce the
     complexity
* from O(n log(n)) to O(n) similarly to using build instead
      of n modifications.
void push() {
 for (int i = 1; i < n; ++i) {</pre>
   t[i<<1] += t[i]:
   t[i<<1|1] += t[i];
   t[i] = 0:
}
// Non commutative combiner functions.
void modify(int p, const S& value) {
 for (t[p += n] = value; p >>= 1; ) t[p] = combine(t[p << 1],
       t[p<<1|1]);
S querv(int 1, int r) {
 S resl, resr;
 for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
  if (l\&1) resl = combine(resl, t[l++]);
   if (r&1) resr = combine(t[--r], resr);
 return combine(resl, resr);
* segment tree for intervals
* */
const int MN = 100000 + 100:
struct seg_tree {
 int val[MN * 4 + 4];
 int pending [MN * 4 + 4];
 seg_tree() {
   memset(val, -1, sizeof val);
   memset(pending, -1, sizeof pending);
```

```
}
void propagate(int node, int b, int e) {
  if (pending[node] != -1) {
    val[node] = pending[node];
    if (b < e) {
     pending[node << 1] = pending[node];</pre>
     pending[node << 1 | 1] = pending[node];</pre>
   pending[node] = -1;
void set(int node, int b, int e, int from, int to, int v)
  if (b > to || e < from) return;</pre>
  if (b >= from && e <= to) {</pre>
   pending[node] = v:
   propagate(node, b, e);
    return;
  int mid = (b + e) >> 1;
  set(node << 1, b, mid, from, to, v);
  set(node << 1 | 1, mid + 1, e, from, to, v):
int query(int node, int b, int e, int pos) {
  propagate(node, b, e);
  if (b == e && b == pos) {
    return val[node]:
  int mid = (b + e) >> 1;
  if (pos <= mid)</pre>
   return query(node << 1, b, mid, pos);</pre>
  return query(node << 1 | 1, mid + 1, e, pos);</pre>
void set(int from, int to, int v) {
  return set(1, 0, MN - 1, from, to, v):
}
int query(int pos) {
  return query(1, 0, MN - 1, pos);
```

```
}
};
```

## 3.11 sparse table

```
// RMQ.
const int MN = 100000 + 10: // Max number of elements
const int ML = 18; // ceil(log2(MN));
struct st {
 int data[MN]:
 int M[MN][ML];
 int n:
 void init(const vector<int> &d) {
   n = d.size();
   for (int i = 0; i < n; ++i)
    data[i] = d[i]:
   build():
 }
 void build() {
   for (int i = 0; i < n; ++i)
     M[i][0] = data[i]:
   for (int j = 1, p = 2, q = 1; p \le n; ++j, p \le 1, q \le 1
     for (int i = 0; i + p - 1 < n; ++i)
      M[i][i] = max(M[i][i-1], M[i+q][i-1]):
 int query(int b, int e) {
   int k = log2(e - b + 1);
   return max(M[b][k], M[e + 1 - (1<<k)][k]);</pre>
 }
};
```

#### 3.12 splay tree

```
using namespace std;
#include<bits/stdc++.h>
#define D(x) cout<<x<<endl;

typedef int T;

struct node{
  node *left, *right, *parent;
  T key;</pre>
```

```
node (T k) : kev(k), left(0), right(0), parent(0) {}
struct splay_tree{
 node *root;
 void right_rot(node *x) {
   node *p = x->parent;
   if (x->parent = p->parent) {
     if (x->parent->left == p) x->parent->left = x:
     if (x->parent->right == p) x->parent->right = x;
   if (p->left = x->right) p->left->parent = p;
   x->right = p;
   p->parent = x;
 void left rot(node *x) {
   node *p = x->parent;
   if (x->parent = p->parent) {
     if (x->parent->left == p) x->parent->left = x;
     if (x->parent->right == p) x->parent->right = x;
   if (p->right = x->left) p->right->parent = p;
   x\rightarrowleft = p;
   p->parent = x;
 void splay(node *x, node *fa = 0) {
   while( x->parent != fa and x->parent != 0) {
     node *p = x->parent;
     if (p->parent == fa)
       if (p->right == x)
        left_rot(x);
       else
         right_rot(x);
     else {
       node *gp = p->parent; //grand parent
       if (gp - > left == p)
        if (p->left == x)
          right_rot(x), right_rot(x);
          left_rot(x),right_rot(x);
        if (p->left == x)
          right_rot(x), left_rot(x);
```

```
left rot(x). left rot(x):
   }
     (fa == 0) root = x:
void insert(T key) {
  node *cur = root;
  node *pcur = 0;
  while (cur) {
   pcur = cur:
   if (kev > cur->kev) cur = cur->right:
    else cur = cur->left;
  cur = new node(key);
  cur->parent = pcur;
  if (!pcur) root = cur;
  else if (key > pcur->key ) pcur->right = cur;
  else pcur->left = cur:
  splay(cur);
node *find(T key) {
  node *cur = root:
  while (cur) {
   if (key > cur->key) cur = cur->right;
   else if(key < cur->key) cur = cur->left;
   else return cur:
 return 0:
}
splay_tree(){ root = 0;};
```

#### 3.13 trie

```
const int MN = 26; // size of alphabet
const int MS = 100010; // Number of states.

struct trie{
    struct node{
    int c;
    int a[MN];
    };

node tree[MS];
    int nodes;
```

```
void clear(){
   tree[nodes].c = 0:
   memset(tree[nodes].a, -1, sizeof tree[nodes].a);
   nodes++:
 7
 void init(){
   nodes = 0:
   clear();
 int add(const string &s, bool query = 0){
   int cur_node = 0;
   for(int i = 0; i < s.size(); ++i){</pre>
     int id = gid(s[i]);
     if(tree[cur_node].a[id] == -1){
      if(query) return 0:
       tree[cur_node].a[id] = nodes;
       clear():
     cur_node = tree[cur_node].a[id];
   if(!query) tree[cur_node].c++;
   return tree[cur_nodel.c:
};
```

#### 3.14 wavelet tree

```
// this can be tested in the problem: http://www.spoj.com/
    problems/ILKQUERY/

struct wavelet {
    vector<int> values, ori;
    vector<int> map_left, map_right;
    int l, r, m;
    wavelet *left, *right;
    wavelet() : left(NULL), right(NULL) {}
    wavelet(int a, int b, int c) : l(a), r(b), m(c), left(NULL), right(NULL) {};
};

wavelet *init(vector<int> &data, vector<int> &ind, int lo,
    int hi) {
    if (lo > hi || (data.size() == 0)) return NULL;
    int mid = ((long long)(lo) + hi) / 2;
    if (lo + 1 == hi) mid = lo; // handle negative values
```

```
wavelet *node = new wavelet(lo. hi. mid):
 vector<int> data_1, data_r, ind_1, ind_r;
 int ls = 0, rs = 0:
 for (int i = 0; i < int(data.size()); i++) {</pre>
   int value = data[i]:
   if (value <= mid) {</pre>
     data_l.emplace_back(value);
     ind_l.emplace_back(ind[i]);
     ls++;
   } else {
     data r.emplace back(value):
     ind_r.emplace_back(ind[i]);
     rs++:
   node->map_left.emplace_back(ls);
   node->map_right.emplace_back(rs);
   node->values.emplace_back(value);
   node->ori.emplace back(ind[i]):
 if (lo < hi) {
   node->left = init(data_1, ind_1, lo, mid);
   node->right = init(data_r, ind_r, mid + 1, hi);
 return node;
int kth(wavelet *node, int to, int k) {
 // returns the kth element in the sorted version of (a[0].
       ..., a[to])
 if (node->1 == node->r) return node->m;
 int c = node->map left[to]:
 if (k < c)
   return kth(node->left, c - 1, k):
 return kth(node->right, node->map right[to] - 1, k - c):
int pos_kth_ocurrence(wavelet *node, int val, int k) {
 // returns the position on the original array of the kth
      ocurrence of the value "val"
 if (!node) return -1;
 if (node->1 == node->r) {
   if (int(node->ori.size()) <= k)</pre>
     return -1:
   return node->ori[k];
 if (val <= node->m)
```

```
return pos_kth_ocurrence(node->left, val, k);
return pos_kth_ocurrence(node->right, val, k);
```

## 4 Geometry

### 4.1 center 2 points + radious

```
vector<point> find_center(point a, point b, long double r) {
  point d = (a - b) * 0.5;
  if (d.dot(d) > r * r) {
    return vector<point> ();
  }
  point e = b + d;
  long double fac = sqrt(r * r - d.dot(d));
  vector<point> ans;
  point x = point(-d.y, d.x);
  long double 1 = sqrt(x.dot(x));
  x = x * (fac / 1);
  ans.push_back(e + x);
  x = point(d.y, -d.x);
  x = x * (fac / 1);
  ans.push_back(e + x);
  return ans;
}
```

### 4.2 closest pair problem

```
struct point {
   double x, y;
   int id;
   point() {}
   point (double a, double b) : x(a), y(b) {}
};

double dist(const point &o, const point &p) {
   double a = p.x - o.x, b = p.y - o.y;
   return sqrt(a * a + b * b);
}

double cp(vector<point> &p, vector<point> &x, vector<point> &y) {
   if (p.size() < 4) {
      double best = 1e100;
      for (int i = 0; i < p.size(); ++i)
            for (int j = i + 1; j < p.size(); ++j)</pre>
```

```
best = min(best, dist(p[i], p[j]));
   return best:
}
 int ls = (p.size() + 1) >> 1;
 double l = (p[ls - 1].x + p[ls].x) * 0.5;
 vector<point> xl(ls), xr(p.size() - ls);
 unordered_set<int> left;
 for (int i = 0: i < ls: ++i) {</pre>
   xl[i] = x[i];
   left.insert(x[i].id):
 for (int i = ls; i < p.size(); ++i) {</pre>
   xr[i - ls] = x[i]:
 vector<point> yl, yr;
 vector<point> pl, pr;
 vl.reserve(ls): vr.reserve(p.size() - ls):
 pl.reserve(ls); pr.reserve(p.size() - ls);
 for (int i = 0; i < p.size(); ++i) {</pre>
   if (left.count(y[i].id))
     vl.push_back(v[i]);
     yr.push_back(y[i]);
   if (left.count(p[i].id))
     pl.push back(p[i]):
     pr.push_back(p[i]);
 double dl = cp(pl, xl, yl);
 double dr = cp(pr, xr, vr);
 double d = min(dl, dr);
 vector<point> yp; yp.reserve(p.size());
 for (int i = 0; i < p.size(); ++i) {</pre>
   if (fabs(y[i].x - 1) < d)
     yp.push_back(y[i]);
 for (int i = 0: i < vp.size(): ++i) {</pre>
   for (int j = i + 1; j < vp.size() && j < i + 7; ++j) {
     d = min(d, dist(yp[i], yp[j]));
 return d:
double closest_pair(vector<point> &p) {
 vector<point> x(p.begin(), p.end());
```

```
sort(x.begin(), x.end(), [](const point &a, const point &b
        ) {
    return a.x < b.x;
});
    vector<point> y(p.begin(), p.end());
    sort(y.begin(), y.end(), [](const point &a, const point &b
        ) {
        return a.y < b.y;
});
    return cp(p, x, y);
}</pre>
```

#### 4.3 triangles

Let a, b, c be length of the three sides of a triangle.

$$p = (a+b+c)*0.5$$

The inradius is defined by:

$$iR = \sqrt{\frac{(p-a)(p-b)(p-c)}{p}}$$

The radius of its circumcircle is given by the formula:

$$cR = \frac{abc}{\sqrt{(a+b+c)(a+b-c)(a+c-b)(b+c-a)}}$$

## 5 Graphs

## 5.1 SCC kosaraju

```
struct SCC {
  vector<vector<int> > g, gr;
  vector<bool> used;
  vector<int> order, component;
  int total_components;

SCC(vector<vector<int> > &adj) {
    g = adj;
    int n = g.size();
    gr.resize(n);
    for (int i = 0; i < n; i++)
        for (auto to : g[i])</pre>
```

```
gr[to].push back(i):
   used.assign(n, false);
   for (int i = 0: i < n: i++)
     if (!used[i])
       GenTime(i):
   used.assign(n, false);
   component.assign(n, -1);
   total_components = 0;
   for (int i = n - 1; i >= 0; i--) {
     int v = order[i]:
     if (!used[v]) {
       vector<int> cur_component;
       Dfs(cur_component, v);
       for (auto node : cur_component)
         component[node] = total_components;
       total_components++;
 void GenTime(int node) {
   used[node] = true:
   for (auto to : g[node])
     if (!used[to])
       GenTime(to):
   order.push_back(node);
 void Dfs(vector<int> &cur, int node) {
   used[node] = true;
   cur.push back(node):
   for (auto to : gr[node])
     if (!used[to])
       Dfs(cur. to):
 }
 vector<vector<int>> CondensedGraph() {
   vector<vector<int>> ans(total_components);
   for (int i = 0: i < int(g.size()): i++) {</pre>
     for (int to : g[i]) {
       int u = component[i], v = component[to];
       if (u != v)
         ans[u].push_back(v);
     }
   }
   return ans:
 }
};
```

#### 5.2 board

## 5.3 bridges

```
struct Graph {
 vector<vector<Edge>> g;
 vector<int> vi, low, d, pi, is_b;
 int bridges_computed;
 int ticks, edges;
 Graph(int n, int m) {
   g.assign(n, vector<Edge>());
   is b.assign(m. 0):
   vi.resize(n):
   low.resize(n);
   d.resize(n):
   pi.resize(n);
   edges = 0;
   bridges_computed = 0;
 void AddEdge(int u, int v) {
   g[u].push_back(Edge(v, edges));
   g[v].push_back(Edge(u, edges));
   edges++;
 void Dfs(int u) {
   vi[u] = true:
   d[u] = low[u] = ticks++;
   for (int i = 0; i < (int)g[u].size(); ++i) {</pre>
```

```
int v = g[u][i].to:
     if (v == pi[u]) continue;
     if (!vi[v]) {
       pi[v] = u;
       Dfs(v);
       if (d[u] < low[v]) is b[g[u][i].id] = true;</pre>
       low[u] = min(low[u], low[v]);
     } else {
       low[u] = min(low[u], d[v]);
   }
 }
 // Multiple edges from a to b are not allowed.
 // (they could be detected as a bridge).
 // If you need to handle this, just count
 // how many edges there are from a to b.
 void CompBridges() {
   fill(pi.begin(), pi.end(), -1);
   fill(vi.begin(), vi.end(), 0);
   fill(low.begin(), low.end(), 0);
   fill(d.begin(), d.end(), 0);
   ticks = 0:
   for (int i = 0; i < (int)g.size(); ++i)</pre>
     if (!vi[i]) Dfs(i);
   bridges_computed = true;
 map<int, vector<Edge>> BridgesTree() {
   if (!bridges_computed) CompBridges();
   int n = g.size();
   Dsu dsu(g.size());
   for (int i = 0; i < n; i++)
     for (auto e : g[i])
       if (!is b[e.id]) dsu.Join(i, e.to):
   map<int, vector<Edge>> tree;
   for (int i = 0; i < n; i++)
    for (auto e : g[i])
       if (is b[e.id])
        tree[dsu.Find(i)].emplace_back(dsu.Find(e.to), e.id
             ):
   return tree;
};
```

#### 5.4 directed mst

```
const int inf = 1000000 + 10:
struct edge {
 int u, v, w;
 edge() {}
 edge(int a,int b,int c) : u(a), v(b), w(c) {}
/**
* Computes the minimum spanning tree for a directed graph
* - edges : Graph description in the form of list of edges.
     each edge is: From node u to node v with cost w
* - root : Id of the node to start the DMST.
* - n : Number of nodes in the graph.
* */
int dmst(vector<edge> &edges, int root, int n) {
 int ans = 0;
 int cur_nodes = n;
 while (true) {
   vector<int> lo(cur_nodes, inf), pi(cur_nodes, inf);
   for (int i = 0; i < edges.size(); ++i) {</pre>
     int u = edges[i].u, v = edges[i].v, w = edges[i].w;
     if (w < lo[v] and u != v) {</pre>
      lo[v] = w:
      pi[v] = u:
   lo[root] = 0:
   for (int i = 0: i < lo.size(): ++i) {</pre>
     if (i == root) continue;
     if (lo[i] == inf) return -1;
   int cur_id = 0;
   vector<int> id(cur_nodes, -1), mark(cur_nodes, -1);
   for (int i = 0; i < cur_nodes; ++i) {</pre>
     ans += lo[i]:
     int n = i:
     while (u != root and id[u] < 0 and mark[u] != i) {</pre>
      mark[u] = i:
      u = pi[u];
     if (u != root and id[u] < 0) { // Cycle</pre>
        for (int v = pi[u]; v != u; v = pi[v])
         id[v] = cur_id;
        id[u] = cur id++:
```

```
if (cur_id == 0)
    break;

for (int i = 0; i < cur_nodes; ++i)
    if (id[i] < 0) id[i] = cur_id++;

for (int i = 0; i < edges.size(); ++i) {
    int u = edges[i].u, v = edges[i].v, w = edges[i].w;
    edges[i].u = id[u];
    edges[i].v = id[v];
    if (id[u] != id[v])
        edges[i].w -= lo[v];
}
cur_nodes = cur_id;
root = id[root];
}

return ans;
}
</pre>
```

## 5.5 eulerian path

```
// Taken from https://github.com/lbv/pc-code/blob/master/
    code/graph.cpp
// Eulerian Trail
struct Euler {
 ELV adj; IV t;
 Euler(ELV Adj) : adj(Adj) {}
 void build(int u) {
   while(! adj[u].empty()) {
     int v = adi[u].front().v:
     adj[u].erase(adj[u].begin());
     build(v);
   t.push_back(u);
};
bool eulerian trail(IV &trail) {
 Euler e(adi):
 int odd = 0, s = 0;
    for (int v = 0: v < n: v++) {
    int diff = abs(in[v] - out[v]);
    if (diff > 1) return false;
    if (diff == 1) {
    if (++odd > 2) return false;
```

```
if (out[v] > in[v]) start = v;
}
    */
e.build(s);
reverse(e.t.begin(), e.t.end());
trail = e.t;
return true;
```

#### 5.6 karp min mean cycle

```
* Finds the min mean cycle, if you need the max mean cycle
* just add all the edges with negative cost and print
* ans * -1
* test: uva, 11090 - Going in Cycle!!
const int MN = 1000;
struct edge{
 int v:
 long long w;
 edge(){} edge(int v, int w) : v(v), w(w) {}
long long d[MN][MN];
// This is a copy of g because increments the size
// pass as reference if this does not matter.
int karp(vector<vector<edge> > g) {
int n = g.size();
 g.resize(n + 1); // this is important
 for (int i = 0: i < n: ++i)
   if (!g[i].empty())
     g[n].push_back(edge(i,0));
 for(int i = 0:i<n:++i)</pre>
   fill(d[i],d[i]+(n+1),INT_MAX);
 d[n - 1][0] = 0:
 for (int k = 1; k \le n; ++k) for (int u = 0; u \le n; ++u) {
   if (d[u][k - 1] == INT MAX) continue:
   for (int i = g[u].size() - 1; i >= 0; --i)
```

```
d[g[u][i].v][k] = min(d[g[u][i].v][k], d[u][k-1] + g[
        ul[i].w):
bool flag = true;
for (int i = 0; i < n && flag; ++i)</pre>
  if (d[i][n] != INT_MAX)
    flag = false;
if (flag) {
  return true: // return true if there is no a cycle.
double ans = 1e15;
for (int u = 0: u + 1 < n: ++u) {
  if (d[u][n] == INT_MAX) continue;
  double W = -1e15:
  for (int k = 0; k < n; ++k)
    if (d[u][k] != INT_MAX)
     W = max(W, (double)(d[u][n] - d[u][k]) / (n - k));
  ans = min(ans, W);
// printf("%.21f\n", ans);
cout << fixed << setprecision(2) << ans << endl;</pre>
return false;
```

### 5.7 konig's theorem

In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover

### 5.8 minimum path cover in DAG

Given a directed acyclic graph G = (V, E), we are to find the minimum number of vertex-disjoint paths to cover each vertex in V.  $\frac{5.10}{\text{struct}}$ 

We can construct a bipartite graph  $G' = (Vout \cup Vin, E')$  from G, where :

```
Vout = \{v \in V : v \text{ has positive out } - degree\} Vin = \{v \in V : v \text{ has positive } in - degree\} E' = \{(u, v) \in Vout \times Vin : (u, v) \in E\}
```

Then it can be shown, via König's theorem, that G' has a matching of size m if and only if there exists n-m vertex-disjoint paths that cover each vertex in G, where n is the number of vertices in G and m is the maximum cardinality bipartite mathching in G'.

Therefore, the problem can be solved by finding the maximum cardinality matching in G' instead.

**NOTE:** If the paths are note necessarily disjoints, find the transitive closure and solve the problem for disjoint paths.

### 5.9 planar graph (euler)

Euler's formula states that if a finite, connected, planar graph is drawn in the plane without any edge intersections, and v is the number of vertices, e is the number of edges and f is the number of faces (regions bounded by edges, including the outer, infinitely large region), then:

$$f + v = e + 2$$

It can be extended to non connected planar graphs with  $\boldsymbol{c}$  connected components:

$$f + v = e + c + 1$$

## 5.10 query with lca

```
struct lowest_ca {
  int T[MN], L[MN], W[MN];
  int P[MN][ML], MI[MN][ML], MA[MN][ML];
```

```
void dfs(vector<vector<edge> > &g. int root, int pi = -1)
  if (pi == -1) {
   L[root] = W[root] = 0:
   T[root] = -1;
  for (int i = 0; i < (int)g[root].size(); ++i) {</pre>
    int to = g[root][i].v;
    if (to != pi) {
     T[to] = root;
     W[to] = g[root][i].w;
     L[to] = L[root] + 1:
      dfs(g, to, root);
}
void init(vector<vector<edge> > &g, int root) {
  // g is undirected
  dfs(g, root);
  int N = g.size(), i, j;
  for (i = 0; i < N; i++) {</pre>
    for (j = 0; 1 << j < N; j++) {
     P[i][j] = -1;
     MI[i][j] = inf;
  for (i = 0: i < N: i++) {</pre>
   P[i][0] = T[i];
    MI[i][0] = W[i];
  for (j = 1; 1 << j < N; j++)
    for (i = 0; i < N; i++)
      if (P[i][i - 1] != -1) {
       P[i][j] = P[P[i][j - 1]][j - 1];
       MI[i][j] = min(MI[i][j-1], MI[P[i][j-1]][j-1]
            1]);
}
int query(int p, int q) {
  int tmp, log, i;
  int mmin = inf;
  if (L[p] < L[q])
    tmp = p, p = q, q = tmp;
```

```
for (log = 1; 1 << log <= L[p]; log++);</pre>
 log--;
 for (i = log; i >= 0; i--)
   if (L[p] - (1 << i) >= L[q]) {
     mmin = min(mmin, MI[p][i]):
     p = P[p][i];
   }
 if (p == q)
   // return p;
   return mmin:
 for (i = log; i >= 0; i--)
   if (P[p][i] != -1 && P[p][i] != P[q][i]) {
     mmin = min(mmin, min(MI[p][i], MI[q][i]));
     p = P[p][i], q = P[q][i];
 // return T[p]:
 return min(mmin, min(MI[p][0], MI[q][0]));
int get_child(int p, int q) { // p is ancestor of q
 if (p == q) return -1;
 int i, log;
 for (log = 1; 1 << log <= L[q]; log++) {}</pre>
 log--;
 for (i = log; i >= 0; i--)
   if (L[q] - (1 << i) > L[p]) {
     q = P[q][i]:
 assert(P[q][0] == p):
 return q;
int is_ancestor(int p, int q) {
 if (L[p] >= L[a])
   return false;
 int dist = L[q] - L[p];
 int cur = q;
 int step = 0;
 while (dist) {
   if (dist & 1)
     cur = P[cur][step]:
```

```
step++;
  dist >>= 1;
}

return cur == p;
}
};
```

#### 5.11 tarjan scc

```
const int MN = 20002;
struct tarjan_scc {
 int scc[MN], low[MN], d[MN], stacked[MN];
 int ticks. current scc:
 deque<int> s; // used as stack.
 tarjan_scc() {}
 void init () {
   memset(scc. -1, sizeof scc):
   memset(d, -1, sizeof d);
   memset(stacked, 0, sizeof stacked);
   s.clear():
   ticks = current_scc = 0;
 void compute(vector<vector<int> > &g, int u) {
   d[u] = low[u] = ticks++:
   s.push_back(u);
   stacked[u] = true:
   for (int i = 0; i < g[u].size(); ++i) {</pre>
    int v = g[u][i];
    if (d[v] == -1)
      compute(g, v);
     if (stacked[v]) {
      low[u] = min(low[u], low[v]):
   if (d[u] == low[u]) { // root
    int v:
     do {
      v = s.back();s.pop_back();
      stacked[v] = false:
      scc[v] = current_scc;
     } while (u != v);
     current_scc++;
```

## 5.12 two sat (with kosaraju)

```
* Given a set of clauses (a1 v a2)^(a2 v a3)....
* this algorithm find a solution to it set of clauses.
* test: http://lightoj.com/volume_showproblem.php?problem
**/
#include<bits/stdc++.h>
using namespace std;
#define MAX 100000
#define endl '\n'
vector<int> G[MAX]:
vector<int> GT[MAX];
vector<int> Ftime;
vector<vector<int> > SCC:
bool visited[MAX];
int n:
void dfs1(int n){
 visited[n] = 1;
 for (int i = 0; i < G[n].size(); ++i) {</pre>
   int curr = G[n][i];
   if (visited[curr]) continue;
   dfs1(curr):
 }
 Ftime.push_back(n);
void dfs2(int n, vector<int> &scc) {
 visited[n] = 1:
 scc.push_back(n);
 for (int i = 0;i < GT[n].size(); ++i) {</pre>
   int curr = GT[n][i];
   if (visited[curr]) continue;
   dfs2(curr. scc):
```

```
void kosaraju() {
 memset(visited, 0, sizeof visited);
 for (int i = 0; i < 2 * n : ++i) {
   if (!visited[i]) dfs1(i);
 memset(visited, 0, sizeof visited);
 for (int i = Ftime.size() - 1; i >= 0; i--) {
   if (visited[Ftime[i]]) continue;
   vector<int> scc:
   dfs2(Ftime[i]. scc):
   SCC.push_back(_scc);
}
/**
 * After having the SCC, we must traverse each scc, if in
      one SCC are -b v b, there is not a solution.
 * Otherwise we build a solution, making the first "node"
      that we find truth and its complement false.
 **/
bool two_sat(vector<int> &val) {
 kosaraju();
 for (int i = 0; i < SCC.size(); ++i) {</pre>
   vector<bool> tmpvisited(2 * n, false);
   for (int j = 0; j < SCC[i].size(); ++j) {</pre>
     if (tmpvisited[SCC[i][j] ^ 1]) return 0;
     if (val[SCC[i][j]] != -1) continue;
       val[SCC[i][j]] = 0;
       val[SCC[i][j] ^ 1] = 1;
     tmpvisited[SCC[i][i]] = 1:
 }
 return 1;
// Example of use
int main() {
 int m, u, v, nc = 0, t: cin >> t:
 // n = "nodes" number, m = clauses number
 while (t--) {
   cin >> m >> n;
```

```
Ftime.clear():
  SCC.clear():
  for (int i = 0; i < 2 * n; ++i) {
   G[i].clear():
   GT[i].clear();
  // (a1 v a2) = (a1 -> a2) = (a2 -> a1)
  for (int i = 0; i < m ; ++i) {</pre>
   cin >> u >> v;
   int t1 = abs(u) - 1:
   int t2 = abs(v) - 1:
   int p = t1 * 2 + ((u < 0)? 1 : 0);
    int q = t2 * 2 + ((v < 0)? 1 : 0);
    G[p ^ 1].push_back(q);
    G[q ^ 1].push_back(p);
    GT[p].push_back(q ^ 1);
    GT[q].push_back(p ^ 1);
  vector < int > val(2 * n, -1);
  cout << "Case " << ++nc <<": ";
  if (two_sat(val)) {
   cout << "Yes" << endl;</pre>
   vector<int> sol;
   for (int i = 0; i < 2 * n; ++i)
     if (i % 2 == 0 and val[i] == 1)
       sol.push_back(i / 2 + 1);
    cout << sol.size();</pre>
    for (int i = 0; i < sol.size(); ++i) {</pre>
     cout << " " << sol[i];
    cout << endl;</pre>
 } else {
    cout << "No" << endl:</pre>
return 0;
```

#### 6 Math

#### 6.1 Lucas theorem

For non-negative integers m and n and a prime p, the following congruence relation holds: :

$$\binom{m}{n} \equiv \prod_{i=0}^{k} \binom{m_i}{n_i} \pmod{p},$$

where:

$$m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0,$$

and:

$$n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$$

are the base p expansions of m and n respectively. This uses the convention that  $\binom{m}{n} = 0$  if  $m \le n$ .

## 6.2 counting

```
const int MN = 1e5 + 100;
long long fact[MN];

void fill_fact() {
  fact[0] = 1;
  for (int i = 1; i < MN; i++) {
    fact[i] = mult(fact[i - 1], i);
  }
}

long long perm_rep(vector<int> &frec) {
  int total = 0;
  long long den = 1;
  for (int i = 0; i < (int)frec.size(); i++) {
    den = mult(den, mod_inv(fact[frec[i]]));
    total += frec[i];
  }
  return mult(fact[total], den);
}</pre>
```

#### 6.3 cumulative sum of divisors

```
/**
The function SOD(n) (sum of divisors) is defined
as the summation of all the actual divisors of
an integer number n. For example,
 SOD(24) = 2+3+4+6+8+12 = 35.
The function CSOD(n) (cumulative SOD) of an integer n, is
    defined as below:
 csod(n) = \sum \{i = 1\}^{n} sod(i)
It can be computed in O(sqrt(n)):
*/
long long csod(long long n) {
 long long ans = 0;
 for (long long i = 2; i * i <= n; ++i) {
   long long j = n / i;
   ans += (i + j) * (j - i + 1) / 2;
   ans += i * (i - i):
 }
 return ans;
```

#### 6.4 fft

```
/**
 * Fast Fourier Transform.
 * Useful to compute convolutions.
 * computes:
 * C(f star g)[n] = sum_m(f[m] * g[n - m])
 * for all n.
 * test: icpc live archive, 6886 - Golf Bot
 * */

using namespace std;
#include<bits/stdc++.h>
#define D(x) cout << #x " = " << (x) << endl
#define endl '\n'

const int MN = 262144 << 1;
int d[MN + 10], d2[MN + 10];

const double PI = acos(-1.0);</pre>
```

```
struct cpx {
 double real, image;
 cpx(double _real, double _image) {
  real = real:
   image = _image;
 cpx(){}
cpx operator + (const cpx &c1, const cpx &c2) {
 return cpx(c1.real + c2.real, c1.image + c2.image);
cpx operator - (const cpx &c1, const cpx &c2) {
 return cpx(c1.real - c2.real, c1.image - c2.image);
cpx operator * (const cpx &c1, const cpx &c2) {
 return cpx(c1.real*c2.real - c1.image*c2.image, c1.real*c2
      .image + c1.image*c2.real);
int rev(int id, int len) {
 int ret = 0:
 for (int i = 0; (1 << i) < len; i++) {
   ret <<= 1:
   if (id & (1 << i)) ret |= 1;</pre>
 return ret;
cpx A[1 << 20];
void FFT(cpx *a, int len, int DFT) {
 for (int i = 0: i < len: i++)</pre>
   A[rev(i, len)] = a[i]:
 for (int s = 1; (1 << s) <= len; s++) {
   int m = (1 << s):
   cpx wm = cpx(cos(DFT * 2 * PI / m), sin(DFT * 2 * PI / m)
   for(int k = 0: k < len: k += m) {
     cpx w = cpx(1, 0);
     for(int j = 0; j < (m >> 1); j++) {
      cpx t = w * A[k + j + (m >> 1)];
      cpx u = A[k + j];
      A[k + j] = u + t;
      A[k + j + (m >> 1)] = u - t;
      w = w * wm:
   }
```

```
if (DFT == -1) for (int i = 0; i < len; i++) A[i].real /=
      len, A[i].image /= len;
 for (int i = 0; i < len; i++) a[i] = A[i];</pre>
 return;
cpx in[1 << 20];
void solve(int n) {
 memset(d, 0, sizeof d):
 for (int i = 0; i < n; ++i) {</pre>
   cin >> t:
   d[t] = true:
 int m:
 cin >> m;
 vector<int> a(m):
 for (int i = 0; i < m; ++i)</pre>
   cin >> q[i];
 for (int i = 0; i < MN; ++i) {</pre>
   if (d[i])
    in[i] = cpx(1, 0);
   else
     in[i] = cpx(0, 0);
 FFT(in, MN, 1):
 for (int i = 0; i < MN; ++i) {</pre>
  in[i] = in[i] * in[i];
 FFT(in, MN, -1);
 int ans = 0:
 for (int i = 0; i < q.size(); ++i) {</pre>
  if (in[q[i]].real > 0.5 || d[q[i]]) {
     ans++;
   }
 cout << ans << endl;</pre>
int main() {
ios_base::sync_with_stdio(false);cin.tie(NULL);
 int n;
 while (cin >> n)
  solve(n):
 return 0:
```

Sample Team Name (Sample University Name)

## 6.5 fibonacci properties

Let A, B and n be integer numbers.

$$k = A - B \tag{1}$$

$$F_A F_B = F_{k+1} F_A^2 + F_k F_A F_{A-1} \tag{2}$$

$$\sum_{i=0}^{n} F_i^2 = F_{n+1} F_n \tag{3}$$

ev(n) = returns 1 if n is even.

$$\sum_{i=0}^{n} F_i F_{i+1} = F_{n+1}^2 - ev(n) \tag{4}$$

$$\sum_{i=0}^{n} F_i F_{i-1} = \sum_{i=0}^{n-1} F_i F_{i+1}$$
 (5)

#### 6.6 polynomials

```
const double pi = acos(-1);
struct poly {
 deque<double> coef;
 double x_lo, x_hi;
 double evaluate(double x) {
   double ans = 0;
   for (auto it : coef)
     ans = (ans * x + it);
   return ans:
 }
 double volume(double x, double dx=1e-6) {
   dx = (x_hi - x_lo) / 1000000.0;
   double ans = 0;
   for (double ix = x lo: ix <= x: ix += dx) {
     double rad = evaluate(ix);
     ans += pi * rad * rad * dx;
   return ans;
```

} };

#### 6.7 sigma function

the sigma function is defined as:

$$\sigma_x(n) = \sum_{d|n} d^x$$

when x=0 is called the divisor function, that counts the number of positive divisors of n.

Now, we are interested in find

$$\sum_{d|n} \sigma_0(d)$$

if n is written as prime factorization:

$$n = \prod_{i=1}^{k} P_i^{e_k}$$

we can demonstrate that:

$$\sum_{d|n} \sigma_0(d) = \prod_{i=1}^k g(e_k + 1)$$

where g(x) is the sum of the first x positive numbers:

$$g(x) = (x * (x + 1))/2$$

#### 7 Matrix

### 7.1 matrix

```
const int MN = 111;
const int mod = 10000;
struct matrix {
  int r, c;
  int m[MN] [MN];
```

```
matrix (int _r, int _c) : r (_r), c (_c) {
   memset(m, 0, sizeof m);
 void print() {
   for (int i = 0: i < r: ++i) {
     for (int j = 0; j < c; ++j)
       cout << m[i][j] << " ";
     cout << endl:
   }
 }
 int x[MN][MN];
 matrix & operator *= (const matrix &o) {
   memset(x, 0, sizeof x);
   for (int i = 0; i < r; ++i)
     for (int k = 0; k < c; ++k)
       if (m[i][k] != 0)
        for (int j = 0; j < c; ++j) {
          x[i][j] = (x[i][j] + ((m[i][k] * o.m[k][j]) % mod
               ) ) % mod;
   memcpy(m, x, sizeof(m));
   return *this;
};
void matrix_pow(matrix b, long long e, matrix &res) {
 memset(res.m, 0, sizeof res.m);
 for (int i = 0; i < b.r; ++i)
   res.m[i][i] = 1;
 if (e == 0) return:
 while (true) {
   if (e & 1) res *= b;
   if ((e >>= 1) == 0) break:
```

### 8 Misc

#### 8.1 Template Java

```
import java.io.*;
import java.util.StringTokenizer;
public class Template {
```

```
public static void main(String []args) throws IOException
       Scanner in = new Scanner(System.in):
       OutputWriter out = new OutputWriter(System.out);
       Task solver = new Task():
       solver.solve(in, out);
       out.close();
class Task{
   public void solve(Scanner in, OutputWriter out){
}
class Scanner{
   public BufferedReader reader:
   public StringTokenizer st;
   public Scanner(InputStream stream){
       reader = new BufferedReader(new InputStreamReader(
       st = null;
   public String next(){
       while(st == null || !st.hasMoreTokens()){
           try{
              String line = reader.readLine();
              if(line == null) return null:
              st = new StringTokenizer(line);
          }catch (Exception e){
              throw (new RuntimeException()):
       return st.nextToken();
   public int nextInt(){
       return Integer.parseInt(next()):
   public long nextLong(){
       return Long.parseLong(next());
   public double nextDouble(){
       return Double.parseDouble(next());
```

#### 8.2 dates

```
11
// Time - Leap years
//
// A[i] has the accumulated number of days from months
    previous to i
const int A[13] = \{ 0, 0, 31, 59, 90, 120, 151, 181, 212, \dots \}
    243, 273, 304, 334 }:
// same as A, but for a leap year
const int B[13] = \{ 0, 0, 31, 60, 91, 121, 152, 182, 213, \dots \}
    244, 274, 305, 335 }:
// returns number of leap years up to, and including, y
int leap_years(int y) { return y / 4 - y / 100 + y / 400; }
bool is_leap(int y) { return y % 400 == 0 || (y % 4 == 0 &&
    v % 100 != 0); }
// number of days in blocks of years
const int p400 = 400*365 + leap_years(400);
const int p100 = 100*365 + leap_years(100);
const int p4 = 4*365 + 1;
const int p1 = 365;
```

```
int date to days(int d. int m. int v)
 return (y - 1) * 365 + leap_years(y - 1) + (is_leap(y) ? B
      [m] : A[m]) + d:
void days to date(int days, int &d, int &m, int &v)
 bool top100; // are we in the top 100 years of a 400 block
 bool top4; // are we in the top 4 years of a 100 block?
 bool top1: // are we in the top year of a 4 block?
 v = 1;
 top100 = top4 = top1 = false;
 y += ((days-1) / p400) * 400;
 d = (davs-1) \% p400 + 1:
 if (d > p100*3) top100 = true, d = 3*p100, v += 300;
 else v += ((d-1) / p100) * 100, d = (d-1) % p100 + 1;
 if (d > p4*24) top4 = true, d = 24*p4, y += 24*4;
 else y += ((d-1) / p4) * 4, d = (d-1) % p4 + 1;
 if (d > p1*3) top1 = true, d = p1*3, y += 3;
 else y += (d-1) / p1, d = (d-1) % p1 + 1;
 const int *ac = top1 && (!top4 || top100) ? B : A:
 for (m = 1; m < 12; ++m) if (d \le ac[m + 1]) break;
 d -= ac[m]:
```

#### 8.3 fraction

```
struct frac{
  long long x, y;
  frac(long long a, long long b) {
    long long g = __gcd(a, b);
    x = a / g;
    y = b / g;
  }
  bool operator < (const frac &o) const {
    return (x * o.y < y * o.x);
  }
};</pre>
```

#### 8.4 io

```
// taken from : https://github.com/lbv/pc-code/blob/master/
    solved/c-e/diablo/diablo.cpp
// this is very fast as well : https://github.com/lbv/pc-
    code/blob/master/code/input.cpp
typedef unsigned int u32;
#define BUF 524288
struct Reader {
 char buf[BUF]; char b; int bi, bz;
 Reader() { bi=bz=0; read(); }
 void read() {
   if (bi==bz) { bi=0; bz = fread(buf, 1, BUF, stdin); }
   b = bz ? buf[bi++] : 0; }
 void skip() { while (b > 0 && b <= 32) read(); }</pre>
 u32 next_u32() {
   u32 v = 0: for (skip(): b > 32: read()) v = v*10 + b-48:
        return v; }
 int next int() {
   int v = 0: bool s = false:
   skip(); if (b == '-') { s = true; read(); }
   for (: 48 \le b \& b \le 57: read()) v = v*10 + b-48: return s?
 char next_char() { skip(); char c = b; read(); return c; }
};
```

## 9 Number theory

#### 9.1 convolution

```
typedef long long int LL;
typedef pair<LL, LL> PLL;
inline bool is_pow2(LL x) {
  return (x & (x-1)) == 0;
}
inline int ceil_log2(LL x) {
  int ans = 0;
  --x;
  while (x != 0) {
    x >>= 1;
    ans++;
  }
  return ans;
}
```

```
/* Returns the convolution of the two given vectors in time
    proportional to n*log(n).
* The number of roots of unity to use nroots_unity must be
     set so that the product of the first
* nroots_unity primes of the vector nth_roots_unity is
     greater than the maximum value of the
* convolution. Never use sizes of vectors bigger than 2^24,
      if you need to change the values of
* the nth roots of unity to appropriate primes for those
     sizes
vector<LL> convolve(const vector<LL> &a. const vector<LL> &b
     , int nroots_unity = 2) {
 int N = 1 << ceil_log2(a.size() + b.size());</pre>
 vector<LL> ans(N,0), fA(N), fB(N), fC(N);
 LL modulo = 1;
 for (int times = 0: times < nroots unity: times++) {</pre>
   fill(fA.begin(), fA.end(), 0);
   fill(fB.begin(), fB.end(), 0):
   for (int i = 0: i < a.size(): i++) fA[i] = a[i]:</pre>
   for (int i = 0; i < b.size(); i++) fB[i] = b[i];</pre>
   LL prime = nth_roots_unity[times].first;
   LL inv_modulo = mod_inv(modulo % prime, prime);
   LL normalize = mod_inv(N, prime);
   ntfft(fA, 1, nth_roots_unity[times]);
   ntfft(fB, 1, nth_roots_unity[times]);
   for (int i = 0; i < N; i++) fC[i] = (fA[i] * fB[i]) %
        prime:
   ntfft(fC, -1, nth_roots_unity[times]);
   for (int i = 0: i < N: i++) {</pre>
     LL curr = (fC[i] * normalize) % prime;
     LL k = (curr - (ans[i] % prime) + prime) % prime;
     k = (k * inv_modulo) % prime;
     ans[i] += modulo * k;
   modulo *= prime:
 return ans;
```

## 9.2 diophantine equations

```
long long gcd(long long a, long long b, long long &x, long
    long &y) {
    if (a == 0) {
        x = 0;
        y = 1;
        return b;
    }
```

```
long long x1, v1:
 long long d = gcd(b \% a, a, x1, y1);
 x = y1 - (b / a) * x1;
 v = x1:
 return d;
bool find_any_solution(long long a, long long b, long long c
     , long long &x0,
   long long &y0, long long &g) {
 g = gcd(abs(a), abs(b), x0, y0);
 if (c % g) {
   return false;
 x0 *= c / g;
 y0 *= c / g;
 if (a < 0) x0 = -x0;
 if (b < 0) v0 = -v0:
 return true:
void shift_solution(long long &x, long long &y, long long a,
     long long b,
   long long cnt) {
 x += cnt * b;
 y -= cnt * a;
long long find_all_solutions(long long a, long long b, long
    long c,
   long long minx, long long maxx, long long miny,
   long long maxv) {
 long long x, y, g;
 if (!find_any_solution(a, b, c, x, y, g)) return 0;
 a /= g:
 b /= g;
 long long sign_a = a > 0 ? +1 : -1;
 long long sign_b = b > 0 ? +1 : -1;
 shift_solution(x, y, a, b, (minx - x) / b);
 if (x < minx) shift_solution(x, y, a, b, sign_b);</pre>
 if (x > maxx) return 0;
 long long lx1 = x;
 shift_solution(x, y, a, b, (maxx - x) / b);
 if (x > maxx) shift_solution(x, y, a, b, -sign_b);
 long long rx1 = x;
```

```
shift_solution(x, y, a, b, -(miny - y) / a);
if (y < miny) shift_solution(x, y, a, b, -sign_a);
if (y > maxy) return 0;
long long lx2 = x;

shift_solution(x, y, a, b, -(maxy - y) / a);
if (y > maxy) shift_solution(x, y, a, b, sign_a);
long long rx2 = x;

if (lx2 > rx2) swap(lx2, rx2);
long long lx = max(lx1, lx2);
long long rx = min(rx1, rx2);

if (lx > rx) return 0;
return (rx - lx) / abs(b) + 1;
```

#### 9.3 discrete logarithm

```
// Computes x which a \hat{x} = b \mod n.
long long d_log(long long a, long long b, long long n) {
 long long m = ceil(sqrt(n));
 long long aj = 1;
 map<long long, long long> M;
 for (int i = 0: i < m: ++i) {</pre>
   if (!M.count(aj))
     M[ai] = i;
   aj = (aj * a) % n;
 long long coef = mod_pow(a, n - 2, n);
 coef = mod_pow(coef, m, n);
 // coef = a ^ (-m)
 long long gamma = b;
 for (int i = 0; i < m; ++i) {</pre>
   if (M.count(gamma)) {
     return i * m + M[gamma];
   } else {
     gamma = (gamma * coef) % n;
 return -1;
```

## 9.4 ext euclidean

## 9.5 highest exponent factorial

```
int highest_exponent(int p, const int &n){
  int ans = 0;
  int t = p;
  while(t <= n){
    ans += n/t;
    t*=p;
  }
  return ans;
}</pre>
```

#### 9.6 miller rabin

```
const int rounds = 20:
// checks whether a is a witness that n is not prime, 1 < a
bool witness(long long a, long long n) {
 // check as in Miller Rabin Primality Test described
 long long u = n - 1;
 int t = 0;
 while (u % 2 == 0) {
   t++:
   u >>= 1:
 long long next = mod_pow(a, u, n);
 if (next == 1) return false:
 long long last;
 for (int i = 0; i < t; ++i) {</pre>
   last = next:
   next = mod_mul(last, last, n);
   if (next == 1) {
     return last != n - 1:
```

```
}
    return next != 1;
}

// Checks if a number is prime with prob 1 - 1 / (2 ^ it)
// D(miller_rabin(999999999999999997LL) == 1);
// D(miller_rabin(99999999999971LL) == 1);
// D(miller_rabin(7907) == 1);
bool miller_rabin(long long n, int it = rounds) {
    if (n <= 1) return false;
    if (n == 2) return true;
    if (n % 2 == 0) return false;
    for (int i = 0; i < it; ++i) {
        long long a = rand() % (n - 1) + 1;
        if (witness(a, n)) {
            return false;
        }
    }
    return true;
}</pre>
```

### 9.7 mod integer

```
template<class T, T mod>
struct mint_t {
   T val;
   mint_t() : val(0) {}
   mint_t(T v) : val(v % mod) {}

mint_t operator + (const mint_t& o) const {
   return (val + o.val) % mod;
   }
   mint_t operator - (const mint_t& o) const {
    return (val - o.val) % mod;
   }

mint_t operator * (const mint_t& o) const {
   return (val * o.val) % mod;
   }

mint_t operator * (const mint_t& o) const {
   return (val * o.val) % mod;
   }
};

typedef mint_t<long long, 998244353> mint;
```

#### 9.8 mod inv

```
long long mod_inv(long long n, long long m) {
  long long x, y, gcd;
```

```
ext_euclid(n, m, x, y, gcd);
if (gcd != 1)
  return 0;
return (x + m) % m;
```

#### 9.9 mod mul

```
// Computes (a * b) % mod
long long mod_mul(long long a, long long b, long long mod) {
  long long x = 0, y = a % mod;
  while (b > 0) {
    if (b & 1)
        x = (x + y) % mod;
        y = (y * 2) % mod;
        b /= 2;
    }
  return x % mod;
}
```

### $9.10 \mod pow$

```
// Computes ( a ^ exp ) % mod.
long long mod_pow(long long a, long long exp, long long mod)
    {
    long long ans = 1;
    while (exp > 0) {
        if (exp & 1)
            ans = mod_mul(ans, a, mod);
        a = mod_mul(a, a, mod);
        exp >>= 1;
    }
    return ans;
}
```

#### 9.11 number theoretic transform

```
* but is different from 1 for all smaller powers */
vector<PLL> nth roots unity {
 {1224736769,330732430},{1711276033,927759239},{167772161,167489322}
  \substack{\{469762049,343261969\},\{754974721,643797295\},\{1107296257,883889069\}\};} \mathbf{primes}
PLL ext_euclid(LL a, LL b) {
 if (b == 0)
   return make_pair(1,0);
 pair<LL.LL> rc = ext euclid(b, a % b):
 return make pair(rc.second, rc.first - (a / b) * rc.second
      );
//returns -1 if there is no unique modular inverse
LL mod inv(LL x, LL modulo) {
 PLL p = ext_euclid(x, modulo);
 if ( (p.first * x + p.second * modulo) != 1 )
  return -1:
 return (p.first+modulo) % modulo;
//Number theory fft. The size of a must be a power of 2
void ntfft(vector<LL> &a, int dir, const PLL &root_unity) {
 int n = a.size();
 LL prime = root unity.first:
 LL basew = mod_pow(root_unity.second, (prime-1) / n, prime
      ):
 if (dir < 0) basew = mod_inv(basew, prime);</pre>
 for (int m = n; m >= 2; m >>= 1) {
   int mh = m >> 1:
   LL w = 1;
   for (int i = 0; i < mh; i++) {</pre>
    for (int i = i: i < n: i += m) {
      int k = i + mh:
      LL x = (a[j] - a[k] + prime) % prime;
      a[j] = (a[j] + a[k]) \% prime;
      a[k] = (w * x) % prime;
       = (w * basew) % prime;
   basew = (basew * basew) % prime;
 for (int j = 1; j < n - 1; j++) {
   for (int k = n >> 1; k > (i ^= k); k >>= 1);
   if (j < i) swap(a[i], a[j]);</pre>
```

```
namespace primes {
 const int MP = 100001;
 bool sieve[MP];
 long long primes[MP];
 int num_p;
 void fill sieve() {
   num_p = 0;
   sieve[0] = sieve[1] = true;
   for (long long i = 2; i < MP; ++i) {</pre>
    if (!sieve[i]) {
       primes[num p++] = i:
      for (long long j = i * i; j < MP; j += i)</pre>
        sieve[j] = true;
  }
 }
 // Finds prime numbers between a and b, using basic primes
       up to sqrt(b)
 // a must be greater than 1.
 vector<long long> seg_sieve(long long a, long long b) {
   long long ant = a;
   a = max(a, 3LL);
   vector<bool> pmap(b - a + 1);
   long long sqrt_b = sqrt(b);
   for (int i = 0; i < num_p; ++i) {</pre>
    long long p = primes[i];
     if (p > sart b) break:
     long long j = (a + p - 1) / p;
     for (long long v = (j == 1) ? p + p : j * p; v <= b; v
          } (q =+
       pmap[v - a] = true;
   vector<long long> ans;
   if (ant == 2) ans.push_back(2);
   int start = a % 2 ? 0 : 1;
   for (int i = start, I = b - a + 1; i < I; i += 2)</pre>
    if (pmap[i] == false)
       ans.push_back(a + i);
   return ans:
 vector<pair<int, int>> factor(int n) {
   vector<pair<int, int>> ans;
```

```
if (n == 0) return ans;
for (int i = 0; primes[i] * primes[i] <= n; ++i) {
   if ((n % primes[i]) == 0) {
     int expo = 0;
     while ((n % primes[i]) == 0) {
       expo++;
       n /= primes[i];
     }
     ans.emplace_back(primes[i], expo);
   }
}
if (n > 1) {
   ans.emplace_back(n, 1);
}
   return ans;
}
```

#### 9.13 totient sieve

```
for (int i = 1; i < MN; i++)
  phi[i] = i;

for (int i = 1; i < MN; i++)
  if (!sieve[i]) // is prime
  for (int j = i; j < MN; j += i)
    phi[j] -= phi[j] / i;</pre>
```

#### 9.14 totient

```
long long totient(long long n) {
   if (n == 1) return 0;
   long long ans = n;
   for (int i = 0; primes[i] * primes[i] <= n; ++i) {
      if ((n % primes[i]) == 0) {
       while ((n % primes[i]) == 0) n /= primes[i];
      ans -= ans / primes[i];
   }
   }
   if (n > 1) {
      ans -= ans / n;
   }
   return ans;
}
```

## 10 Strings

#### 10.1 Incremental Aho Corasick

```
class IncrementalAhoCorasic {
 static const int Alphabets = 26;
 static const int AlphabetBase = 'a';
 struct Node {
  Node *fail:
   Node *next[Alphabets];
   int sum:
   Node() : fail(NULL), next{}, sum(0) { }
 struct String {
   string str:
   int sign;
 };
 //totalLen = sum of (len + 1)
 void init(int totalLen) {
   nodes.resize(totalLen);
   nNodes = 0:
   strings.clear():
   roots.clear();
   sizes.clear():
   que.resize(totalLen);
 void insert(const string &str, int sign) {
   strings.push_back(String{ str, sign });
   roots.push back(nodes.data() + nNodes):
   sizes.push_back(1);
   nNodes += (int)str.size() + 1:
   auto check = [&]() { return sizes.size() > 1 && sizes.end
        ()[-1] == sizes.end()[-2]; };
    makePMA(strings.end() - 1, strings.end(), roots.back(),
          que);
   while(check()) {
    int m = sizes.back():
    roots.pop_back();
     sizes.pop_back();
     sizes.back() += m;
    if(!check())
      makePMA(strings.end() - m * 2, strings.end(), roots.
           back(), que);
```

```
int match(const string &str) const {
   int res = 0:
   for(const Node *t : roots)
     res += matchPMA(t, str);
   return res:
private:
 static void makePMA(vector<String>::const_iterator begin,
      vector<String>::const_iterator end, Node *nodes,
      vector<Node*> &que) {
   int nNodes = 0;
   Node *root = new(&nodes[nNodes ++]) Node();
   for(auto it = begin; it != end; ++ it) {
     Node *t = root;
     for(char c : it->str) {
      Node *&n = t->next[c - AlphabetBase];
      if(n == nullptr)
        n = new(&nodes[nNodes ++]) Node();
     t->sum += it->sign;
   int qt = 0;
   for(Node *&n : root->next) {
     if(n != nullptr) {
      n->fail = root:
       que[qt ++] = n;
     } else {
       n = root;
   for(int qh = 0; qh != qt; ++ qh) {
     Node *t = que[qh];
     int a = 0:
     for(Node *n : t->next) {
      if(n != nullptr) {
        que[qt ++] = n;
        Node *r = t->fail:
        while(r->next[a] == nullptr)
          r = r->fail;
        n->fail = r->next[a]:
        n->sum += r->next[a]->sum;
       ++ a:
```

```
static int matchPMA(const Node *t, const string &str) {
   int res = 0:
   for(char c : str) {
     int a = c - AlphabetBase:
     while(t->next[a] == nullptr)
      t = t->fail:
     t = t-\text{next}[a];
     res += t->sum;
   return res;
 vector<Node> nodes:
 int nNodes:
 vector<String> strings;
 vector<Node*> roots:
 vector<int> sizes;
 vector<Node*> que:
int main() {
 int m;
 while("scanf("%d", &m)) {
   IncrementalAhoCorasic iac;
   iac.init(600000);
   rep(i, m) {
     int tv:
     char s[300001];
     scanf("%d%s", &tv, s):
     if(ty == 1) {
      iac.insert(s, +1);
     } else if(tv == 2) {
       iac.insert(s, -1);
     } else if(ty == 3) {
       int ans = iac.match(s):
       printf("%d\n", ans);
      fflush(stdout);
     } else {
       abort():
     }
 return 0;
```

### 10.2 minimal string rotation

```
// Lexicographically minimal string rotation
```

```
int lmsr() {
 string s:
 cin >> s:
 int n = s.size();
 vector<int> f(s.size(), -1);
 int k = 0:
 for (int j = 1; j < 2 * n; ++j) {
   int i = f[j - k - 1];
   while (i != -1 && s[j] != s[k + i + 1]) {
    if (s[i] < s[k + i + 1])
      k = i - i - 1:
    i = f[i];
   if (i == -1 \&\& s[j] != s[k + i + 1]) {
    if (s[i] < s[k + i + 1]) {
      k = j;
    f[i - k] = -1:
   } else {
    f[i - k] = i + 1;
 return k;
```

## 10.3 suffix array

```
* 0 (n log^2 (n))
* See http://web.stanford.edu/class/cs97si/suffix-array.pdf
      for reference
* */
struct entry{
 int a, b, p;
 entrv(){}
 entry(int x, int y, int z): a(x), b(y), p(z){}
 bool operator < (const entry &o) const {</pre>
   return (a == o.a) ? (b == o.b) ? (p < o.p) : (b < o.b) :
         (a < o.a):
 }
};
struct SuffixArray{
 const int N;
 string s;
 vector<vector<int> > P:
 vector<entry> M;
```

```
SuffixArray(const string &s) : N(s.length()), s(s), P(1,
      vector<int> (N, 0)), M(N) {
   for (int i = 0: i < N: ++i)
     P[0][i] = (int) s[i];
   for (int skip = 1, level = 1; skip < N; skip *= 2, level</pre>
        ++) {
     P.push_back(vector<int>(N, 0));
     for (int i = 0 ; i < N; ++i) {</pre>
       int next = ((i + skip) < N) ? P[level - 1][i + skip]</pre>
       M[i] = entry(P[level - 1][i], next, i);
     sort(M.begin(), M.end());
     for (int i = 0; i < N; ++i)</pre>
      P[level][M[i].p] = (i > 0 \text{ and } M[i].a == M[i - 1].a
            and M[i].b == M[i - 1].b) ? P[level][M[i - 1].p]
   }
 }
 vector<int> getSuffixArray(){
   vector<int> &rank = P.back();
   vector<pair<int, int> > inv(rank.size());
   for (int i = 0; i < rank.size(); ++i)</pre>
     inv[i] = make_pair(rank[i], i);
   sort(inv.begin(), inv.end());
   vector<int> sa(rank.size());
   for (int i = 0: i < rank.size(): ++i)</pre>
     sa[i] = inv[i].second;
   return sa;
 // returns the length of the longest common prefix of s[i
      ...L-1] and s[i...L-1]
 int lcp(int i, int j) {
   int len = 0:
   if (i == j) return N - i;
   for (int k = P.size() - 1; k >= 0 && i < N && j < N; --k)
     if (P[k][i] == P[k][i]) {
       i += 1 << k:
       j += 1 << k;
       len += 1 << k;
   }
   return len:
};
```

#### 10.4 suffix automaton

```
* Suffix automaton:
 * This implementation was extended to maintain (online) the
 * number of different substrings. This is equivalent to
 * the number of paths from the initial state to all the
 * states.
 * The overall complexity is O(n)
 * can be tested here: https://www.urionlinejudge.com.br/
      judge/en/problems/view/1530
struct state {
 int len, link;
 long long num paths:
 map<int, int> next;
};
const int MN = 200011;
state sa[MN << 1]:
int sz, last;
long long tot_paths;
void sa_init() {
 sz = 1:
 last = 0:
 sa[0].len = 0;
 sa[0].link = -1;
 sa[0].next.clear();
 sa[0].num_paths = 1;
 tot_paths = 0;
void sa_extend(int c) {
 int cur = sz++;
 sa[cur].len = sa[last].len + 1;
```

```
sa[cur].next.clear():
sa[cur].num_paths = 0;
for (p = last; p != -1 && !sa[p].next.count(c); p = sa[p].
    link) {
  sa[p].next[c] = cur:
  sa[cur].num_paths += sa[p].num_paths;
  tot_paths += sa[p].num_paths;
if (p == -1) {
 sa[cur].link = 0:
} else {
  int q = sa[p].next[c];
  if (sa[p].len + 1 == sa[q].len) {
   sa[cur].link = q;
  } else {
   int clone = sz++;
   sa[clone].len = sa[p].len + 1;
    sa[clone].next = sa[q].next;
    sa[clone].num_paths = 0;
    sa[clone].link = sa[q].link;
    for (; p!= -1 && sa[p].next[c] == q; p = sa[p].link) {
     sa[p].next[c] = clone;
     sa[q].num_paths -= sa[p].num_paths;
     sa[clone].num_paths += sa[p].num_paths;
    sa[q].link = sa[cur].link = clone:
}
last = cur;
```

### 10.5 z algorithm

```
using namespace std;
#include<bits/stdc++.h>
```

```
vector<int> compute_z(const string &s){
int n = s.size():
 vector<int> z(n,0);
 int l.r:
 r = 1 = 0;
 for(int i = 1: i < n: ++i){</pre>
  if(i > r) {
    1 = r = i;
     while (r < n \text{ and } s[r - 1] == s[r])r++;
     z[i] = r - 1;r--;
   }else{
     int k = i-l:
     if(z[k] < r - i +1) z[i] = z[k];
     else {
      1 = i:
      while(r < n and s[r - 1] == s[r])r++;
      z[i] = r - 1:r--:
   }
 }
 return z;
int main(){
 //string line;cin>>line;
 string line = "alfalfa";
 vector<int> z = compute z(line):
 for(int i = 0: i < z.size(): ++i ){</pre>
   if(i)cout<<" ";
   cout<<z[i];
 cout << end1;
 // must print "0 0 0 4 0 0 1"
 return 0;
```