## Simple hand written parsers

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Grammar categories
Bottom-up parsers
Top-down parsers
Parser generators
Recursive descent parser

#### Outline

- 1. Introduction to parsers
- 2. Abstract syntax tree
- 3. Lexers
- 4. Backtracking and lookahead with iterators
- 5. Example grammar and parser
- 6 Parser combinator

#### Grammars Grammar categories Bottom-up parsers

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#### Grammars

• grammar - a set of production rules in a formal language

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- grammar a set of production rules in a formal language
- production rule a rule specifying substitutions to be performed to generate a symbol sequence

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- grammar a set of production rules in a formal language
- production rule a rule specifying substitutions to be performed to generate a symbol sequence
- Backus–Naur Form (BNF)

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- grammar a set of production rules in a formal language
- production rule a rule specifying substitutions to be performed to generate a symbol sequence
- Backus–Naur Form (BNF)
- \(\langle digit \rangle ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9
- $\langle number \rangle ::= \langle digit \rangle \langle digit \rangle^*$

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- grammar a set of production rules in a formal language
- production rule a rule specifying substitutions to be performed to generate a symbol sequence
- Backus–Naur Form (BNF)
- \(\langle digit \rangle ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9
- $\langle number \rangle ::= \langle digit \rangle \langle digit \rangle^*$
- $\langle expression \rangle ::= \langle instantiation \rangle \mid \langle number \rangle$
- $\langle instantiation \rangle ::= \langle identifier \rangle$  (  $\langle expression \rangle$  ( ,  $\langle expression \rangle$  )\* )

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## Grammar categories

Noam Chomsky

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- Noam Chomsky
- Chomsky hierarchy

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- Noam Chomsky
- Chomsky hierarchy
  - type 0 unrestricted

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- Noam Chomsky
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  - type 0 unrestricted
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- Noam Chomsky
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  - type 2 context-free

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- Noam Chomsky
- Chomsky hierarchy
  - type 0 unrestricted
  - type 1 context-sensitive
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    - LR(k) Left to right, reverse Rightmost derivation
    - LL(k) Left to right, Leftmost derivation

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- Noam Chomsky
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    - LR(k) Left to right, reverse Rightmost derivation
    - LL(k) Left to right, Leftmost derivation
  - type 3 regular

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### Bottom-up parsers

• build production rules starting from tokens (terminal symbols)

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#### Bottom-up parsers

- build production rules starting from tokens (terminal symbols)
- LR(k) grammars

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## Top-down parsers

 $\bullet$  start parsing from production rules and go down to tokens

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## Top-down parsers

- start parsing from production rules and go down to tokens
- LL(k) grammars

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#### Parser generators

• generate parsing code straight from a formal grammar definition

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### Parser generators

- generate parsing code straight from a formal grammar definition
- typically produces table-based parsers

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#### Parser generators

- generate parsing code straight from a formal grammar definition
- typically produces table-based parsers
- the generated code is usually completely unreadable and undebuggable

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## Recursive descent parsers

- a kind of a top-down parser
- easy to write by hand

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## Recursive descent parsers

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- simple to implement for LL(k)

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## Recursive descent parsers

- a kind of a top-down parser
- easy to write by hand
- simple to implement for LL(k)
- can use backtracking for parsing more complex grammars, like LL(\*)

Basics Virtual-dispatch based AST: Static ASTs

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Basics Virtual-dispatch based ASTs Static ASTs

## **Basics**

AST - Abstract Syntax Tree

#### **Basics**

- AST Abstract Syntax Tree
- type erased ASTs

#### Basics

- AST Abstract Syntax Tree
- type erased ASTs
- statically typed ASTs

easy to extend dynamically

- easy to extend dynamically
- use virtual dispatch for visitation

- easy to extend dynamically
- use virtual dispatch for visitation
- or use dynamic casts for visitation

```
struct visitor {
    virtual bool visit(const assignment &) {}
};
struct expression {
    virtual ~expression() = default;
    virtual bool visit(visitor & v) = 0;
};
struct assignment : expression {
    virtual bool visit(visitor & v) override {
        return visitor.visit(*this);
```

```
struct expression {
    virtual ~expression() = default;
}:
struct assignment : expression {};
struct visitor {
    std::unordered_map<std::type_index, std::function<void (expression *)>> callbacks;
    bool visit(expression * expr) {
        return callbacks.at(typeid(expr))(expr);
    template<typename T, typename F>
    void add_visitor(F && f) {
        callbacks.emplace(typeid(T), [f = std::forward<F>(f)](expression * ptr) {
            f(dynamic_cast<T *>(ptr));
        });
};
```

### Static ASTs

• harder to extend, with lots of recompilation

#### Static ASTs

- harder to extend, with lots of recompilation
- type-checked for handling most (all?) cases

#### Static ASTs

- harder to extend, with lots of recompilation
- type-checked for handling most (all?) cases
- variant-based

## Static ASTs

```
using expression = variant<
    assignment,
    binary_expression.
    function call
void analyze_expression(const expression & expr) {
    fmap(expr, make_overload_set()
        [](const assignment &) { /* handle assignment */ },
        [](const binary_expression &) { /* handle binary expression */ },
        [](const function_call &) { /* handle function call */ }
    )):
```

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## The role of lexers

simplify parsers

### The role of lexers

- simplify parsers
- handle encodings

### The role of lexers

- simplify parsers
- handle encodings
- convert slow to compare types into easier to compare ones

## Lexer generators

generate lexers from regular expressions

# Lexer generators

- generate lexers from regular expressions
- flex
- lexertl
- ...a lot more

## Lexer generators

- generate lexers from regular expressions
- flex
- lexertl
- ...a lot more
- force a specific token format

Intro AST **Lexers** Iterators Example Combinators

The role of lexers Lexer generators Lexing by hand

# Lexing by hand

flexible

# Lexing by hand

- flexible
- basically require implementing a separate parser

# Lexing by hand

```
struct position {
    std size t offset;
    std : size_t line;
    std size_t column;
    std::string file;
};
struct range_type {
    position start;
    position end;
};
```

# Lexing by hand

```
enum class token_type {
    identifier, dot, comma,
    plus, minus, equals,
    open_paren, close_paren,
    string
};
struct token {
    token_type type;
    std::u32string string;
    range_type range;
};
```

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Example
Combinators

Parsing context Lookahead Backtracking

# Parsing context

## Parsing context

```
struct context {
    lexer::iterator begin; // or "current"?
    lexer::iterator end;
    // additional context data, like:
    std::vector<operator_context> operator_stack;
};
```

```
lexer::token expect(context & ctx, lexer::token_type type)
{
    if (ctx.begin == ctx.end || ctx.begin->type != type) {
        throw unexpected{};
    }
    return std::move(*ctx.begin++);
}
```

```
optional<token &> peek(context & ctx)
{
    if (ctx.begin != ctx.end) {
        return { *ctx.begin };
    }
    return none;
}
```

```
optional<token &> peek(context & ctx, lexer::token_type type)
{
    if (ctx.begin != ctx.end && ctx.begin->type == type) {
        return { *ctx.begin };
    }
    return none;
}
```

```
optional < token & > peek (context & ctx, std::size t k)
    auto begin = ctx.begin;
    while (k--) {
        if (begin == ctx.end) {
            return none;
        ++begin;
    }
       (begin != ctx.end) {
        return { *begin };
    }
    return none;
```

```
\langle expression \rangle ::= \overline{\langle various \ kinds \ of \ expressions \rangle \mid \langle tuple-expression \rangle}
\langle statement \rangle ::= \langle expression \rangle,;
\langle block \rangle ::= \{, (\langle statement \rangle \mid \langle block \rangle)^*, \}
\langle tuple-expression \rangle ::= \{, \langle expression \rangle, (,, \langle expression \rangle)^*, \}
```

```
expression parse_expression(context & ctx);
optional<statement> try_parse_statement(context & ctx) {
    try {
        auto ctx_temp = ctx;
        auto expr = parse_expression(ctx_tmp);
        expect(ctx_tmp, token_type::semicolon);
        ctx = ctx_temp;
        return { expr };
    }
    catch (unexpected &) {
        return none;
    }
}
```

```
expression pa
             rse expression(context & ctx):
optional < statement > try_parse_statement(context & ctx) {
    trv {
        auto ctx temp = ctx:
        auto expr = parse_expression(ctx_tmp);
        expect(ctx_tmp, token_type::semicolon);
        ctx = ctx_temp;
        return { expr }:
    catch (unexpected &) {
        return none:
                 (context & ctx) {
block
    block ret:
    expect(ctx, token_type::open_brace);
    while (!peek(ctx, token_type::close_brace) {
        if (auto stmt = try_parse_statement(ctx)) { block.add(stmt); }
        else { block.add(parse_block(ctx)); }
    expect(ctx, token_type::close_brace);
```

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AST Lexers Iterators Example Combinators

Use-case Grammar AST Parsers

### Use-case

build-system description language

### Use-case

- build-system description language
- simple
- extensible

#### Grammar

#### Grammar

```
\langle string \rangle ::= anything inside double quotes \langle identifier \rangle ::= alphanumeric string without quotes \langle id\text{-expression} \rangle ::= \langle identifier \rangle ( . \langle identifier \rangle )* \langle instantiation \rangle ::= \langle id\text{-expression} \rangle ( (\langle expression \rangle ( , \langle expression \rangle )? )
```

#### Grammar

```
 \langle simple-expression \rangle ::= \langle string \rangle \\ | \langle id\text{-}expression \rangle \\ | \langle instantiation \rangle \\ \langle expression \rangle ::= \langle simple-expression \rangle \ (\ (\ + \ | \ - \ ) \ \langle simple-expression \rangle \ ) + \\ \langle assignment \rangle ::= \langle id\text{-}expression \rangle = \langle expression \rangle \\ \langle program \rangle ::= \langle assignment \rangle^*
```

```
string_node {
    token value:
};
struct identifier {
    token value;
};
       id_expression {
    std::vector<identifier> identifiers:
};
```

```
expression;
struct instantiation {
    id_expression type_name;
    std::vector<expression> arguments;
};
      simple_expression = variant<</pre>
    string_node,
    id_expression,
    instantiation
```

```
enum class operation type { addition, subtraction };
struct operation {
    operation_type operation;
    simple_expression operand;
};
struct expression {
    simple_expression base;
    std::vector<operation> operations;
};
```

```
struct assignment {
    id_expression lhs;
    expression rhs;
};
```

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ombinators

## **Parsers**

#### **Parsers**

```
string parse_string(context & ctx) {
    return { expect(ctx, token_type::string) };
}
identifier parse_identifier(context & ctx) {
    return { expect(ctx, token_type::identifier) };
}
```

#### Parsers

```
string parse string(context & ctx) {
    return { expect(ctx, token_type::string) };
identifier parse_identifier(context & ctx) {
    return { expect(ctx, token_type::identifier) };
id expression parse id expression(context & ctx) {
    std::vector<identifier> identifiers:
    identifiers.push_back(parse_identifier(ctx));
    while (peek(ctx, token_type::dot)) {
        expect(ctx, token_type::dot);
        identifiers.push back(parse identifier(ctx));
    }
    return { std::move(identifiers) }:
```

```
simple_expression parse_simple_expression(rcontext & ctx) {
    auto peeked = peek(ctx);
    if (!peeked) {
        throw unexpected{};
    }
    switch (peeked->type) {
        case token_type::string:
            return { parse_string(ctx) }; // string
        case token_type::identifier: {
            throw unexpected{}:
```

```
auto id = parse_id_expression(ctx);
   (peek(ctx, token_type::open_paren)) {
    expect(ctx, token_type::open_paren);
    std::vector<expression> arguments;
       (peek(ctx) && peek(ctx)->type != token_type::close_paren) {
        arguments.push_back(parse_expression(ctx));
        while (peek(ctx) && peek(ctx)->type != token_type::close_paren) {
            expect(ctx. token_type::comma);
            arguments.push_back(parse_expression(ctx));
         close = expect(ctx, token_type::close_paren);
    return { instantiation{ std::move(id), std::move(arguments) } }; // instantiation
else {
    return { std : move(id) }: // id-expression
```

```
expression parse_expression(context & ctx) {
    auto expr = parse_simple_expression(ctx);
    auto peeked = peek(ctx);
       (peeked && (peeked->type == token_type::plus
          peeked->type == token_type::minus))
    {
            operations = parse_operations(ctx);
        return expression{ std::move(expr), std::move(operations) };
    return expression{ std::move(expr), {} };
```

```
std::vector<operation> parse_operations(context & ctx) {
    auto peeked = peek(ctx);
    std::vector<operation> operations;
    while (peeked && (peeked->type == token_type::plus
        | peeked->type == token_type::minus))
    {
        operation_type operation = expect(ctx, peeked->type).type == token_type::plus
            ? operation_type::addition
            : operation_type::removal;
        auto operand = parse_simple_expression(ctx);
        operations.push_back({ operation, std::move(operand) });
        peeked = peek(ctx):
    }
    return operations;
```

```
assignment parse_assignment(context & ctx) {
    auto id = parse_id_expression(ctx);
    expect(ctx, token_type::equals);
    auto value = parse_expression(ctx);
    return { std::move(id), std::move(value) }; // assignment
}
```

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Intro AST Lexers Iterators Example Combinators

General idea List Optional Alternative

### General idea

higher-level functions creating more complex parsers from simpler ones

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- higher-level functions creating more complex parsers from simpler ones
- Boost.Spirit.Qi operators

## General idea

- higher-level functions creating more complex parsers from simpler ones
- Boost.Spirit.Qi operators
- help reduce code duplication

## List

```
template<typename F>
auto list(F && f, token_type separator) {
    return [f = std::forward<F>(f), separator](auto && ctx) {
        auto ret = make_vector(f(ctx));
        while (peek(ctx, separator)) {
            expect(ctx, separator);
            ret.push_back(f(ctx));
        }
        return ret:
    };
```

# Optional

```
template<typename F>
auto optional(F && f) {
    return [f = std::forward<F>(f)](auto && ctx) {
        try {
            return make_optional(f(ctx));
        }
        catch (unexpected &) {
            return none:
```

# Alternative

```
template < typename F, typename G>
auto alternative(F && f, G && g) {
    return [f = std::forward<F>(f), g = std::forward<G>(g)](auto && ctx)
        -> variant < decltype(f(ctx)), decltype(g(ctx)) > {
        trv {
            return { f(ctx) }:
        }
        catch (unexpected &) {
            return { g(ctx) };
        }
```

#### Links

- https://github.com/griwes/despayre
- https://github.com/griwes/vapor