Let's break this down following your preferred explanation style:

## 1 Approach

We need to compress a given character array by replacing consecutive repeating characters with the character followed by its count. The algorithm should:

- **Iterate through the array** while keeping track of character groups.
- Count consecutive occurrences of the same character.
- Modify the input array in-place with the character and its count.
- Return the length of the modified array.

### **Key Constraints & Optimizations:**

- We use **constant extra space** (O(1)).
- The array is modified in-place, avoiding additional string storage.
- We use **two pointers**:
  - o i to traverse the array.
  - ansIndex to keep track of where to write the compressed characters.

# 2 Step-by-Step Explanation

```
class Solution {
public:
   int compress(vector<char>& chars) {
```

- We define a function that takes a vector<char>& chars as input.
- This function modifies chars in-place and returns its new length.

#### **Step 1: Initialize Pointers**

```
int i = 0;
int ansIndex = 0;
int n = chars.size();
```

- i → Points to the start of the current group.
- ansIndex → Tracks where to write the compressed characters.
- n → Stores the size of the input array.

### **Step 2: Traverse the Array**

```
while (i < n) {
```

• We loop through the array.

### **Step 3: Count Consecutive Characters**

```
int j = i + 1;
while (j < n && chars[i] == chars[j]) {
    j++;
}</pre>
```

- j starts at i+1 and moves forward as long as characters are the same.
- When j stops:
  - Either we've reached a **new character**.
  - Or we reached the **end of the array**.

### **Step 4: Store the Character**

```
chars[ansIndex++] = chars[i];
```

- We store the **current character** at ansIndex, then move ansIndex forward.
- **Doubt Explanation:** The chars[i] being stored here is the first occurrence of the repeated character.

### **Step 5: Store the Count (if > 1)**

```
int count = j - i;
```

• count stores how many times chars[i] appears.

```
if (count > 1) {
    string cnt = to_string(count);
    for (char ch : cnt) {
        chars[ansIndex++] = ch;
    }
}
```

- If count > 1, we convert it to a string and store each digit separately.
- Example:

```
"ccc" → count = 3, so store '3' in the array.
"aaaaa" → count = 5, so store '5'.
```

### **Step 6: Move to Next Character Group**

```
i = j;
```

• i jumps to j (the next new character).

### **Step 7: Return the New Length**

```
return ansIndex;
}
};
```

• The final compressed array is stored **up to** ansIndex, so we return it.

## 3 Dry Run with Input & Output

### **Example Input:**

```
chars = ["a","a","b","b","c","c"]
```

```
        Step
        i
        j
        count
        ansIndex
        chars (Modified)

        1 "aa"
        0
        2
        2
        1
        ["a", "2", "b", "b", "b", "c", "c", "c", "c"]

        2 "bb"
        2
        4
        2
        3
        ["a", "2", "b", "2", "b", "2", "c", "c", "3", "c"]

        3 "ccc"
        4
        7
        3
        5
        ["a", "2", "b", "2", "b", "2", "c", "c", "3", "c"]
```

### **Final Output:**

- New length: 6
- Compressed Array: ["a", "2", "b", "2", "c", "3"]

## 4 Topics & Patterns Used

### ✓ Two Pointers:

- i keeps track of the current group.
- ansIndex tracks where to write compressed data.

### **✓** String Manipulation:

- Convert int count to string.
- Store individual digits in the array.

### **✓** In-Place Modification:

We modify chars without using extra space.

## **Final Thoughts**

- This solution runs in **O(N) time**, iterating through the array **only once**.
- It doesn't use extra space, modifying chars in-place.
- Works for all cases, including when count is 10+ (e.g., "aaaaaaaaaaa" → "a11").

Let's go step by step through the **code flow**, explaining how the code executes **iteration by iteration**.

## **Given Input:**

```
chars = ["a", "a", "b", "b", "c", "c"]
```

### **Initial Variables Setup:**

```
int i = 0;
int ansIndex = 0;
int n = chars.size(); // n = 7
```

- i = 0 → Points to the first character 'a'
- ansIndex = 0 → Keeps track of where to write in chars
- $n = 7 \rightarrow Total number of characters in the input array$

## **Code Execution - Iteration by Iteration**

First Iteration (i = 0, processing "aa")

```
int j = i + 1; // j = 1
while (j < n && chars[i] == chars[j]) { j++; } // j moves to 2</pre>
```

- chars[0] = 'a'
- j moves while 'a' is found → Stops at j = 2

```
• Group: "aa"
```

```
• Count: j - i = 2 - 0 = 2
```

Store the character in chars[ansIndex] and move ansIndex:

```
chars[ansIndex++] = chars[i]; // chars[0] = 'a', ansIndex = 1
```

Store the count in chars[ansIndex] and move ansIndex:

```
string cnt = to_string(2); // cnt = "2"
chars[ansIndex++] = '2'; // chars[1] = '2', ansIndex = 2
```

**✓** Updated chars after this step:

```
["a", "2", "b", "b", "c", "c"]
```

Move i to j:

```
i = j; // i = 2
```

Second Iteration (i = 2, processing "bb")

```
int j = i + 1; // j = 3
while (j < n && chars[i] == chars[j]) { j++; } // j moves to 4</pre>
```

- chars[2] = 'b'
- j moves while 'b' is found → Stops at j = 4
- Group: "bb"
- Count: j i = 4 2 = 2

Store the character in chars[ansIndex] and move ansIndex:

```
chars[ansIndex++] = chars[i]; // chars[2] = 'b', ansIndex = 3
```

Store the count in chars[ansIndex] and move ansIndex:

```
string cnt = to_string(2); // cnt = "2"
chars[ansIndex++] = '2'; // chars[3] = '2', ansIndex = 4
```

### **☑** Updated chars after this step:

```
["a", "2", "b", "2", "c", "c"]
```

Move i to j:

```
i = j; // i = 4
```

## Third Iteration (i = 4, processing "ccc")

```
int j = i + 1; // j = 5 while (j < n \& chars[i] == chars[j]) { <math>j++; } // j moves to 7
```

- chars[4] = 'c'
- j moves while 'c' is found → Stops at j = 7
- Group: "ccc"
- Count: j i = 7 4 = 3

#### Store the character in chars[ansIndex] and move ansIndex:

```
chars[ansIndex++] = chars[i]; // chars[4] = 'c', ansIndex = 5
```

### Store the count in chars[ansIndex] and move ansIndex:

```
string cnt = to_string(3); // cnt = "3"
chars[ansIndex++] = '3'; // chars[5] = '3', ansIndex = 6
```

### **☑** Updated chars after this step:

```
["a", "2", "b", "2", "c", "3", "c"]
```

Move i to j:

```
i = j; // i = 7 (end of array)
```

### Loop Ends (i = 7, loop condition i < n fails)

The loop terminates, and the function returns:

```
return ansIndex; // returns 6
```

#### **Final Output:**

- Compressed Array: ["a", "2", "b", "2", "c", "3"]
- New Length: 6

## **Final Summary of Iterations**

Iteration	i	j	Group	Count	Updated chars
1	0 → 2	2	"aa"	2	["a", "2", "b", "b", "c", "c", "c"]
2	2 → 4	4	"bb"	2	["a", "2", "b", "2", "c", "c", "c"]
3	<b>4</b> → <b>7</b>	7	"ccc"	3	["a", "2", "b", "2", "c", "3", "c"]

- **☑** Final Compressed Length: 6
- ✓ Final Output Array: ["a", "2", "b", "2", "c", "3"]

# Key Takeaways

#### 1. Two Pointers Approach:

- o i tracks the current group of characters.
- o j expands the window to count occurrences.
- o ansIndex writes the compressed result.

### 2. In-Place Modification:

• We update chars without extra space (O(1) space complexity).

#### 3. O(N) Time Complexity:

• We process each character **only once**, making it an **efficient** algorithm.

# **\*\*** Do You Have Any Doubts?

Let me know if any step is unclear, and I'll clarify it further!  $\mathscr{Q}$ 

Yes, you're absolutely right! Here's the detailed breakdown:

Step-by-Step Execution of chars[ansIndex++] = 'a';

Let's assume:

```
int ansIndex = 0;
```

### 1 Execution of chars[ansIndex++] = 'a';

- First, chars[ansIndex] = 'a'; is executed.
- Since ansIndex = 0, this means:

```
chars[0] = 'a';
```

• Now, chars looks like:

```
["a", "a", "b", "b", "c", "c", "c"] // The first 'a' has been correctly placed
```

### 2 Then, ansIndex++ increments ansIndex

• ansIndex was 0, but after ansIndex++, it becomes 1.

Now, when we later assign:

```
chars[ansIndex++] = '2'; // Now ansIndex is 1
```

- chars[1] = '2'; (since ansIndex = 1)
- Now, chars looks like:

```
["a", "2", "b", "b", "c", "c"]
```

• Then, ansIndex++ moves to 2, ready for the next character.

### **S** Conclusion

### **Yes, you understood correctly!**

- 1. First, store 'a' at chars[0] (when ansIndex = 0).
- 2. Then, ansIndex increments to 1, so the next available position is chars[1].
- 3. Now, '2' gets stored at chars[1], as expected.

Great job thinking through this carefully! Let me know if you have any more questions. 😌