

# Linked List (One Way List)

Start

- Linear collection of data elements called Nodes
- Each node contain two parts:

Information (Info)

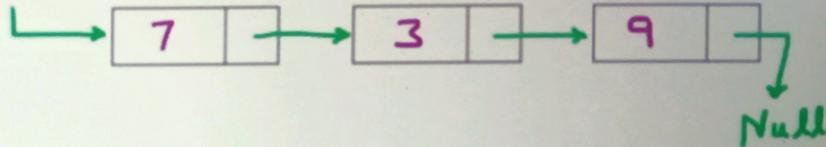
- Contain information of element

Next (Link) Pointer

- Contain address of next node

Null Pointer

- Signals end of list
- Denoted by X or NULL

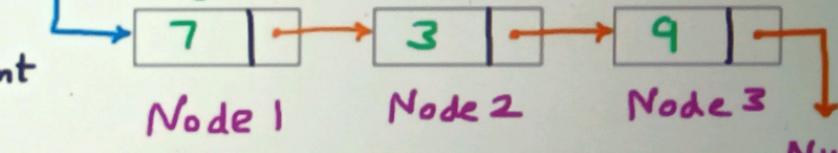


Start

Info Link

Info Link

Info Link



Node 1

Node 2

Node 3

Node 1

Node 2

Node 3

Start (Name) Pointer

- Contain address of first node

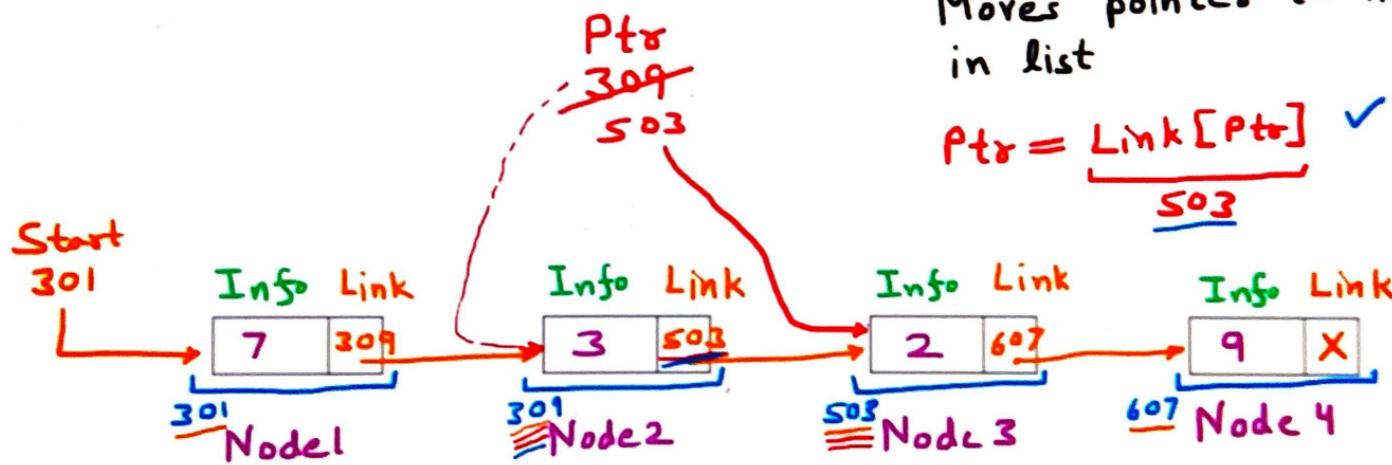
Null (Empty) List

Start

- List has no node
- Denoted by Null Pointer in Start

# Linked List (One Way List)

## Traversing Linked List

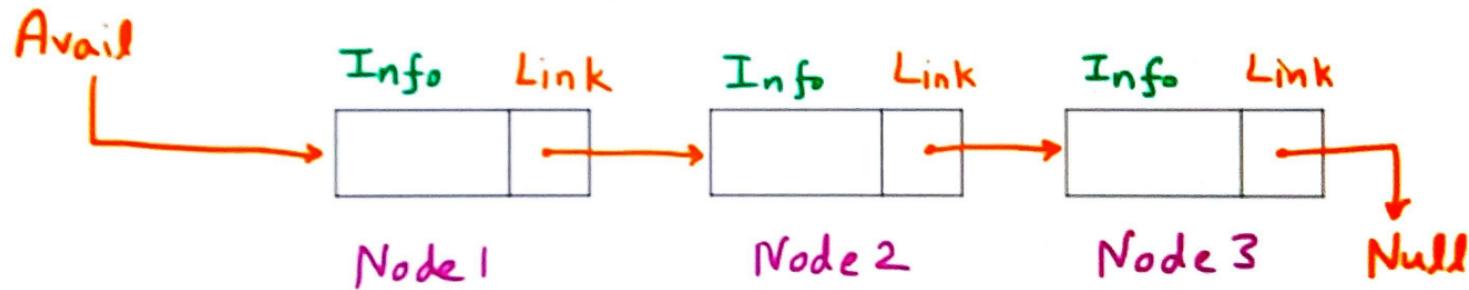


$\text{ptr} = \underline{\text{Link}}[\text{ptr}]$

Moves pointer to next node  
in list

# Available List

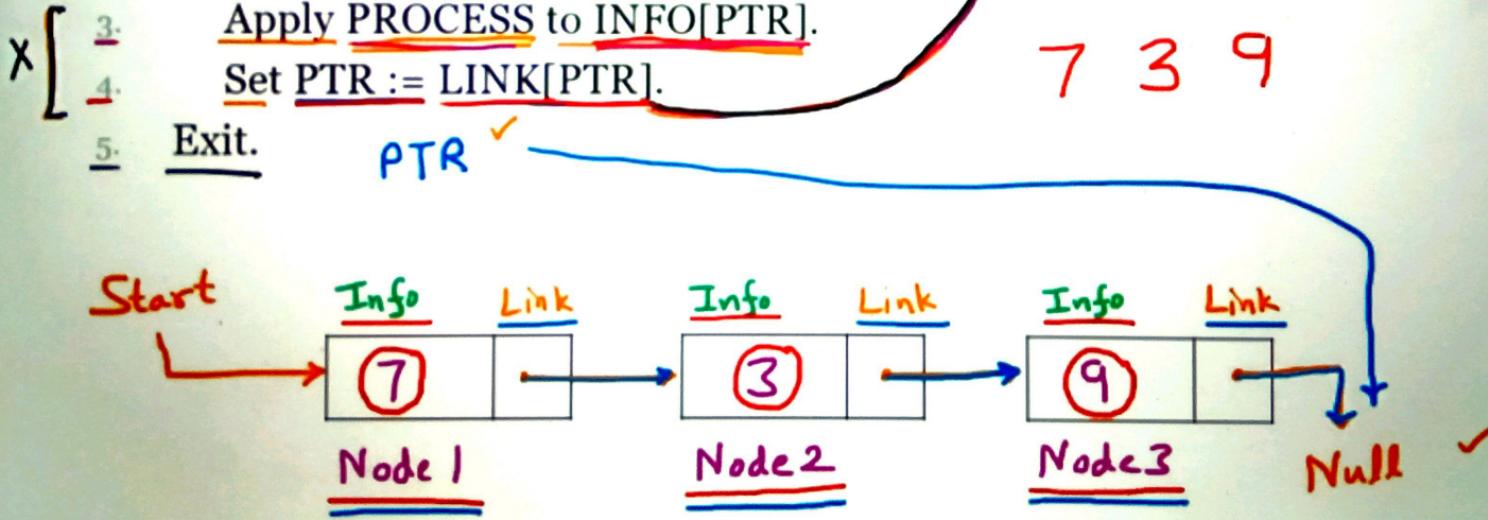
- Linked List which contain available memory spaces



## ALGORITHM: Traversing Linked List

Let LIST be a Linked List in Memory. This algorithm traverse LIST, applying an operation PROCESS to each element of LIST.  
The variable PTR points to node currently being processed.

1. Set PTR := START.
2. Repeat Steps 3 and 4 while PTR != NULL.
3. Apply PROCESS to INFO[PTR].
4. Set PTR := LINK[PTR].
5. Exit.



## ALGORITHM: SEARCH\_SORTED\_LL (INFO, LINK, START, ITEM, LOC)

LIST is the linked list in memory. The algorithm finds the location LOC of the node where ITEM first appears in LIST, or sets LOC=NULL.

1. Set PTR := START
2. Repeat Step 3 while PTR != NULL:
  - 3. If ITEM < INFO[PTR], then:
    - x Set PTR := LINK[PTR].
  - Else if ITEM = INFO[PTR], then:
    - ↳ Set LOC := PTR, and Exit.
  - Else:
    - Set LOC := NULL, and Exit.
4. Set LOC := NULL
5. Exit.

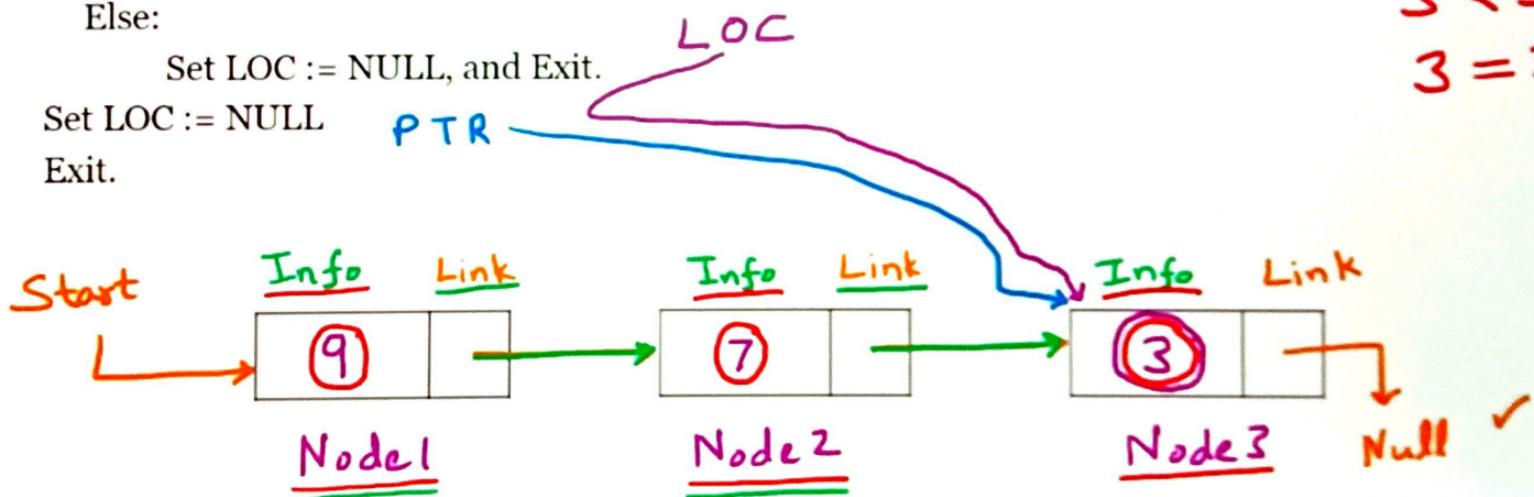
Item  
3

3 < 9

3 < 7

3 < 3 X

3 = 3 ✓



## ALGORITHM: SEARCH\_UNSORTED\_LL (INFO, LINK, START, ITEM, LOC)

LIST is the linked list in memory. The algorithm finds the location LOC of the node where ITEM first appears in LIST, or sets LOC=NULL.

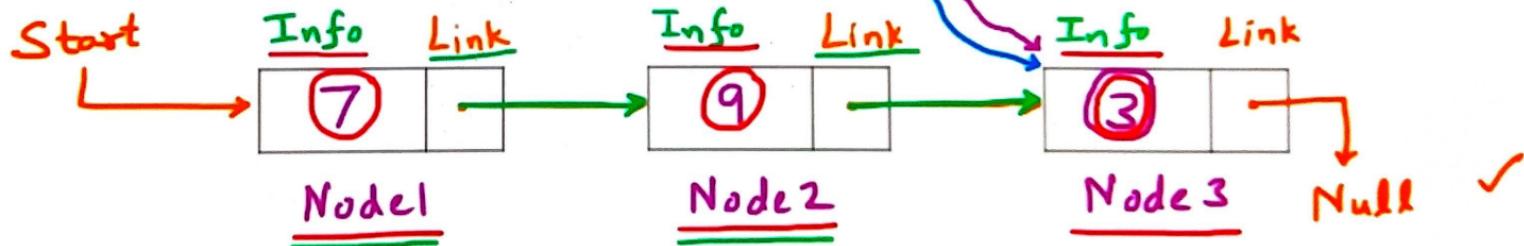
1. Set PTR := START.
2. Repeat Step 3 while PTR != NULL:
  3. If ITEM = INFO[PTR], then:  
~~xx~~ Set LOC := PTR, and Exit.
  - Else:  
Set PTR := LINK[PTR].
4. Set LOC := NULL.
5. Exit

Item  
③ ✓

$3=7$  X

$3=9$  X

$3=3$  ✓

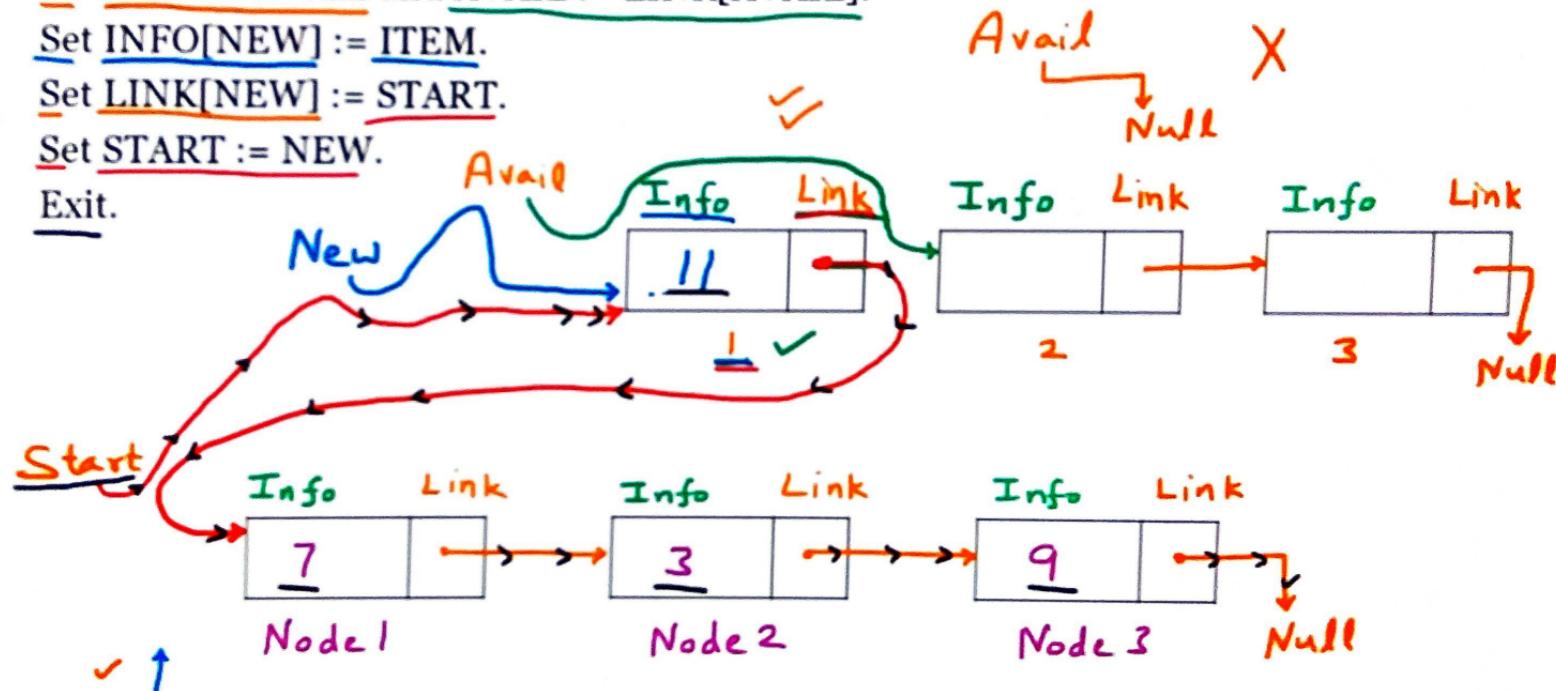


## ALGORITHM: INSFIRST (INFO, LINK, START, AVAIL, ITEM)

This algorithm inserts ITEM as the first node in the List.

1. If AVAIL = NULL, then: Write OVERFLOW and Exit.
2. Set NEW := AVAIL and AVAIL := LINK[AVAIL].
3. Set INFO[NEW] := ITEM.
4. Set LINK[NEW] := START.
5. Set START := NEW.
6. Exit.

Item  
11 ✓✓

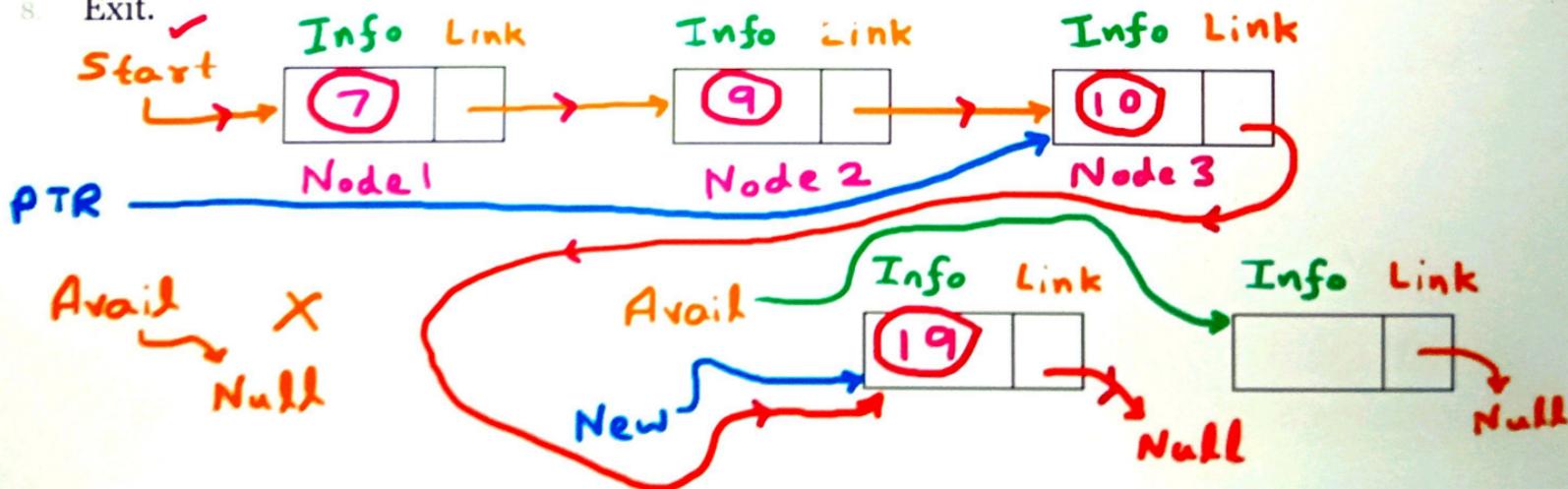


## ALGORITHM: INSERT END (INFO, LINK, START, AVAIL, ITEM)

This algorithm inserts ITEM as the END node in the List.

- 1 If AVAIL = NULL, then: Write OVERFLOW and Exit.
- 2 Set NEW := AVAIL and AVAIL := LINK[AVAIL].
- 3 Set INFO[NEW] := ITEM.
- 4 Set PTR := START.
- 5 Repeat Step 6 while LINK[PTR] != NULL
- 6 Set PTR := LINK[PTR].
- 7 Set LINK[PTR] := NEW and LINK[NEW] := NULL.
- 8 Exit.

Item  
19



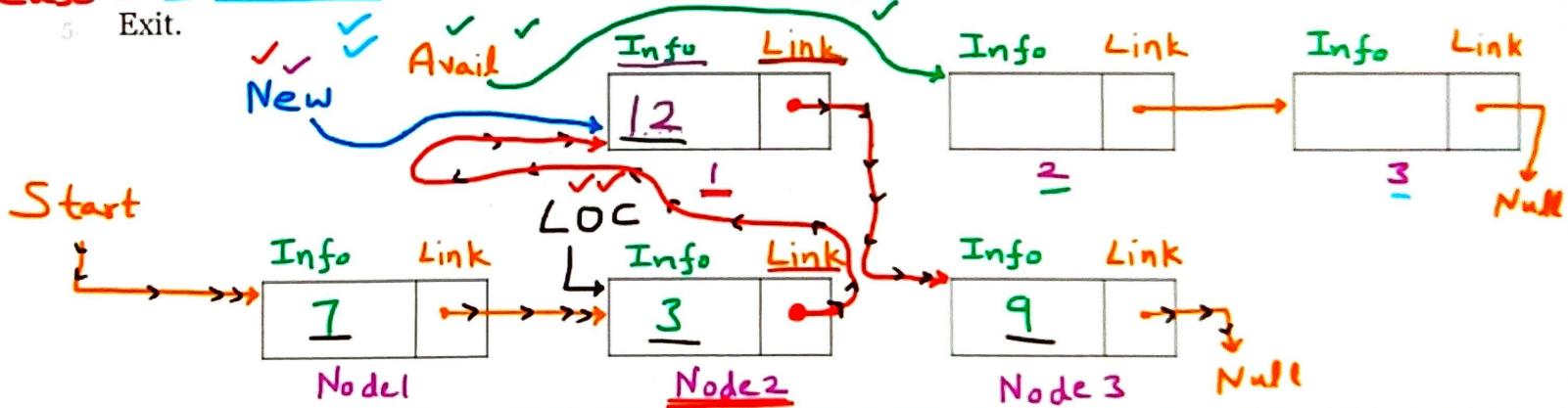
**ALGORITHM:** INSLOC (INFO, LINK, START, AVAIL, LOC, ITEM)

This algorithm inserts ITEM so that ITEM follows the node with location LOC or inserts ITEM as the first node when LOC = NULL.

1. If AVAIL = NULL, then: Write OVERFLOW and Exit.
  2. Set NEW := AVAIL and AVAIL := LINK[AVAIL].
  3. Set INFO[NEW] := ITEM.
  4. If LOC = NULL, then:

**Case 1** Set LINK[NEW] := START and START := NEW.  
Else:

**Case 2** Set LINK[NEW] := LINK[LOC] and LINK[LOC] := NEW.  
Exit.



## ALGORITHM: DEL (INFO, LINK, START, AVAIL, LOC, LOCP)

This algorithm Deletes the node N with location LOC. LOCP is the location of node which precedes N or, when N is the first node, LOCP = NULL.

- If LOCP = NULL, then:

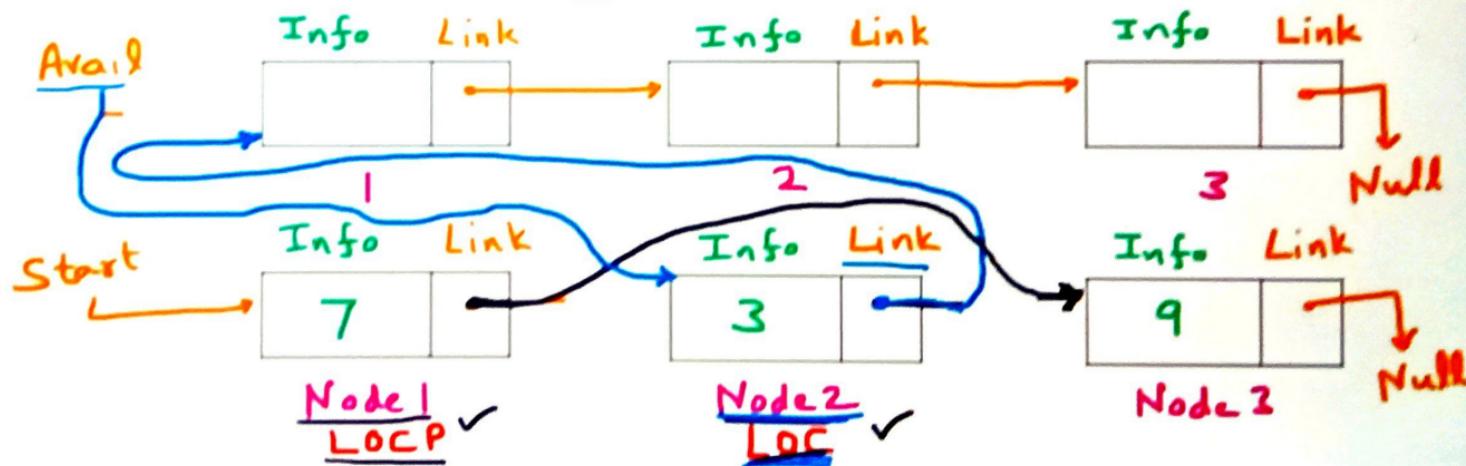
Case 1: Set START := LINK[START].

Else:

Case 2: Set LINK[LOCP] := LINK[LOC].

- Set LINK[LOC] := AVAIL and AVAIL := LOC.

Exit.



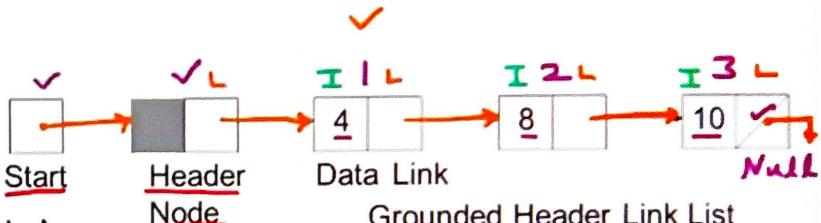
# Header Linked List

- Contain special node called Header Node at begining of list

## 2 Types of Header Linked List

### Grounded Header List

- Last node contain NULL pointer Start
- Term grounded indicate NULL pointer
- Start always point to Header Node
- LINK[START] = NULL means Grounded Header List is Empty



### Circular Header List

- Last node point back to Header Node
- LINK[START] = START means Circular Header List is Empty
- Header list always be circular, unless stated
- Header node acting as sentinel indicating end of list

