Rotate Function

Given an array of integers A and let *n* to be its length.

Assume B_k to be an array obtained by rotating the array A k positions clock-wise, we define a "rotation function" F on A as follow:

```
F(k) = 0 * B_k[0] + 1 * B_k[1] + ... + (n-1) * B_k[n-1].
```

Calculate the maximum value of F(0), F(1), ..., F(n-1).

Note:

n is guaranteed to be less than 10⁵.

Example:

```
A = [4, 3, 2, 6]
F(0) = (0 * 4) + (1 * 3) + (2 * 2) + (3 * 6) = 0 + 3 + 4 + 18 = 25
F(1) = (0 * 6) + (1 * 4) + (2 * 3) + (3 * 2) = 0 + 4 + 6 + 6 = 16
F(2) = (0 * 2) + (1 * 6) + (2 * 4) + (3 * 3) = 0 + 6 + 8 + 9 = 23
F(3) = (0 * 3) + (1 * 2) + (2 * 6) + (3 * 4) = 0 + 2 + 12 + 12 = 26
So the maximum value of F(0), F(1), F(2), F(3) is F(3) = 26.
```

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Solution 1

```
F(k) = 0 * Bk[0] + 1 * Bk[1] + ... + (n-1) * Bk[n-1]
F(k-1) = 0 * Bk-1[0] + 1 * Bk-1[1] + ... + (n-1) * Bk-1[n-1]
= 0 * Bk[1] + 1 * Bk[2] + ... + (n-2) * Bk[n-1] + (n-1) * Bk[0]
```

Then,

```
F(k) - F(k-1) = Bk[1] + Bk[2] + ... + Bk[n-1] + (1-n)Bk[0]
= (Bk[0] + ... + Bk[n-1]) - nBk[0]
= sum - nBk[0]
```

Thus,

```
F(k) = F(k-1) + sum - nBk[0]
```

What is Bk[o]?

```
k = 0; B[0] = A[0];
k = 1; B[0] = A[len-1];
k = 2; B[0] = A[len-2];
```

```
int allSum = 0;
int len = A.length;
int F = 0;
for (int i = 0; i < len; i++) {
    F += i * A[i];
    allSum += A[i];
}
int max = F;
for (int i = len - 1; i >= 1; i--) {
    F = F + allSum - len * A[i];
    max = Math.max(F, max);
}
return max;
```

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Solution 2

Consider we have 5 coins A,B,C,D,E

According to the problem statement

$$F(o) = (oA) + (1B) + (2C) + (3D) + (4E)$$

$$F(1) = (4A) + (oB) + (1C) + (2D) + (3E)$$

$$F(2) = (3A) + (4B) + (oC) + (1D) + (2*E)$$

This problem at a glance seem like a difficult problem. I am not very strong in mathematics, so this is how I visualize this problem

We can construct F(1) from F(0) by two step:

Step 1. taking away one count of each coin from F(0), this is done by subtracting "sum" from "iteration" in the code below after step 1 F(0) = (-1A) + (oB) + (1C) + (2D) + (3*E)

Step 2. Add n times the element which didn't contributed in F(o), which is A. This is done by adding "A[j-1]len" in the code below.

```
after step 2 F(0) = (4A) + (0B) + (1C) + (2D) + (3E)
```

At this point F(0) can be considered as F(1) and F(2) to F(4) can be constructed by repeating the above steps.

Hope this explanation helps, cheers!

```
public int maxRotateFunction(int[] A) {
        if(A.length == 0){
            return 0;
        int sum =0, iteration = 0, len = A.length;
        for(int i=0; i<len; i++){</pre>
            sum += A[i];
            iteration += (A[i] * i);
        }
        int max = iteration;
        for(int j=1; j<len; j++){
            // for next iteration lets remove one entry value of each entry and t
he prev 0 * k
            iteration = iteration - sum + A[j-1]*len;
            max = Math.max(max, iteration);
        }
        return max;
    }
```

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Solution 3

```
public class Solution {

public int maxRotateFunction(int[] A) {
   int n = A.length;
   int sum = 0;
   int candidate = 0;

for (int i = 0; i < n; i++) {
    sum += A[i];
    candidate += A[i] * i;
   }
   int best = candidate;

for (int i = n - 1; i > 0; i--) {
    candidate = candidate + sum - A[i] * n;
   best = Math.max(best, candidate);
   }
   return best;
}
```

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