## **Target Sum**

You are given a list of non-negative integers, a1, a2, ..., an, and a target, S. Now you have 2 symbols + and -. For each integer, you should choose one from + and - as its new symbol.

Find out how many ways to assign symbols to make sum of integers equal to target S.

## Example 1:

```
Input: nums is [1, 1, 1, 1, 1], S is 3.
Output: 5
Explanation:

-1+1+1+1+1 = 3
+1-1+1+1+1 = 3
+1+1-1+1+1 = 3
+1+1+1-1+1 = 3
+1+1+1-1+1 = 3
```

There are 5 ways to assign symbols to make the sum of nums be target 3.

#### Note:

- 1. The length of the given array is positive and will not exceed 20.
- 2. The sum of elements in the given array will not exceed 1000.
- 3. Your output answer is guaranteed to be fitted in a 32-bit integer.

### Solution 1

The recursive solution is very slow, because its runtime is exponential

The original problem statement is equivalent to:

Find a **subset** of **nums** that need to be positive, and the rest of them negative, such that the sum is equal to **target** 

Let P be the positive subset and N be the negative subset For example:

```
Given nums = \begin{bmatrix} 1, 2, 3, 4, 5 \end{bmatrix} and target = 3 then one possible solution is +1-2+3-4+5=3
```

Here positive subset is P = [1, 3, 5] and negative subset is N = [2, 4]

Then let's see how this can be converted to a subset sum problem:

```
sum(P) - sum(N) = target
sum(P) + sum(N) + sum(P) - sum(N) = target + sum(P) + sum(N)
2 * sum(P) = target + sum(nums)
```

So the original problem has been converted to a subset sum problem as follows: Find a **subset** P of nums such that sum(P) = (target + sum(nums)) / 2

Note that the above formula has proved that target + sum(nums) must be even We can use that fact to quickly identify inputs that do not have a solution (Thanks to @BrunoDeNadaiSarnaglia for the suggestion)

For detailed explanation on how to solve subset sum problem, you may refer to Partition Equal Subset Sum

Here is Java solution (15 ms)

Here is C++ solution (3 ms)

written by yuxiangmusic original link here

```
public class Solution {
    public int findTargetSumWays(int[] nums, int s) {
        int sum = 0;
        for(int i: nums) sum+=i;
        if(s>sum || s<-sum) return 0;</pre>
        int[] dp = new int[2*sum+1];
        dp[0+sum] = 1;
        for(int i = 0; i<nums.length; i++){</pre>
             int[] next = new int[2*sum+1];
             for(int k = 0; k<2*sum+1; k++){</pre>
                 if(dp[k]!=0){
                     next[k + nums[i]] += dp[k];
                     next[k - nums[i]] += dp[k];
                 }
            }
            dp = next;
        return dp[sum+s];
}
```

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# Solution 3

I'm quite surprised that simple DFS could pass the test since for DFS solution there are obvious a lot of overlap subproblems. So I used a map to record the intermediate result while we are walking along the recursion tree.

```
public class Solution {
    public int findTargetSumWays(int[] nums, int S) {
        if (nums == null || nums.length == 0){
            return 0;
        }
        return helper(nums, 0, 0, S, new HashMap<>());
    }
    private int helper(int[] nums, int index, int sum, int S, Map<String, Integer</pre>
> map){
        String encodeString = index + "->" + sum;
        if (map.containsKey(encodeString)){
            return map.get(encodeString);
        }
        if (index == nums.length){
            if (sum == S){
                return 1;
            }else {
                return 0;
            }
        }
        int curNum = nums[index];
        int add = helper(nums, index + 1, sum - curNum, S, map);
        int minus = helper(nums, index + 1, sum + curNum, S, map);
        map.put(encodeString, add + minus);
        return add + minus;
    }
}
```

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