Arranging Coins

You have a total of n coins that you want to form in a staircase shape, where every k-th row must have exactly k coins.

Given n, find the total number of **full** staircase rows that can be formed.

n is a non-negative integer and fits within the range of a 32-bit signed integer.

Example 1:

Example 2:

Solution 1

The idea is about quadratic equation, the formula to get the sum of arithmetic progression is

```
sum = (x + 1) * x / 2
```

so for this problem, if we know the sum, then we can know the x = (-1 + sqrt(8 * n + 1)) / 2

```
public class Solution {
    public int arrangeCoins(int n) {
        return (int)((-1 + Math.sqrt(1 + 8 * (long)n)) / 2);
    }
}
```

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[JAVA] Clean Code with Explanations and Running Time [2 Solutions]

Full Solutions and Explanations

Solution 1

```
public class Solution {
    public int arrangeCoins(int n) {
        int start = 0;
        int end = n;
        int mid = 0;
        while (start <= end){</pre>
            mid = (start + end) >>> 1;
            if ((0.5 * mid * mid + 0.5 * mid) <= n){
                 start = mid + 1;
            }else{
                end = mid - 1;
            }
        }
        return start - 1;
    }
}
```

Complexity Analysis

Uniform cost model is used as Cost Model and `n` is the input number. `b` in this case would be `2`.

Time Complexity:

- Best Case `O(log_b(n))`: With respect to the input, the algorithm will always depend on the value of input.
- Average Case `O(log_b(n))`: With respect to the input, the algorithm will always depend on the value of input.
- Worst Case `O(log_b(n))`: With respect to the input, the algorithm will always depend on the value of input.

Auxiliary Space:

• Worst Case `O(1)`: Additional variables are of constant size.

Algorithm

Approach: Binary Search

The problem is basically asking the maximum length of consecutive number that has the running sum lesser or equal to `n`. In other word, find `x` that satisfy the following condition:

```
1+2+3+4+5+6+7+...+x \le n

sum_{i=1}^x i \le n
```

Running sum can be simplified,

```
(x * (x + 1)) / 2 \le n
```

Binary search is used in this case to slowly narrow down the x that will satisfy the equation. Note that 0.5 * mid * mid + 0.5 * mid does not have overflow issue as the intermediate result is implicitly autoboxed to double data type.

Solution 2

```
public class Solution {
    public int arrangeCoins(int n) {
        return (int) ((Math.sqrt(1 + 8.0 * n) - 1) / 2);
    }
}
```

Complexity Analysis

Uniform cost model is used as Cost Model and `n` is the input number. `b` in this case would be `2`.

Time Complexity:

- Best Case `O(1)`: With respect to the input, the algorithm will always perform basic mathematical operation that run in constant time.
- Average Case `O(1)`: With respect to the input, the algorithm will always perform basic mathematical operation that run in constant time.
- Worst Case `O(1)`: With respect to the input, the algorithm will always perform basic mathematical operation that run in constant time.

Auxiliary Space:

• Worst Case `O(1)`: No extra space is used.

Algorithm

Approach: Mathematics

The problem is basically asking the maximum length of consecutive number that has the running sum lesser or equal to `n`. In other word, find `x` that satisfy the following condition:

```
1+2+3+4+5+6+7+...+x \le n
\sum_{i=1}^{n} x_i \le n
```

Running sum can be simplified,

```
(x * (x + 1)) / 2 \le n
```

Using quadratic formula, `x` is evaluated to be,

```
`x = 1 / 2 * (-sqrt(8 * n + 1)-1)` (Inapplicable) or `x = 1 / 2 * (sqrt(8 * n + 1)-1)` Negative root is ignored and positive root is used instead. Note that 8.0 * n is very important because it will cause Java to implicitly autoboxed the intermediate result into double data type. The code will not work if it is simply 8 * n. Alternatively, an explicit casting can be done 8 * (long) n).
```

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Solution 3

```
public int arrangeCoins(int n) {
    //convert int to long to prevent integer overflow
    long nLong = (long)n;

    long st = 0;
    long ed = nLong;

    long mid = 0;

    while (st <= ed) {
        mid = st + (ed - st) / 2;

        if (mid * (mid + 1) <= 2 * nLong) {
            st = mid + 1;
        }else {
            ed = mid - 1;
        }
    }

    return (int)(st - 1);
}</pre>
```

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From Leetcoder.