Concatenated Words

Given a list of words (**without duplicates**), please write a program that returns all concatenated words in the given list of words.

A concatenated word is defined as a string that is comprised entirely of at least two shorter words in the given array.

Example:

```
Input: ["cat","cats","catsdogcats","dog","dogcatsdog","hippopotamuses","rat","ratca
tdogcat"]

Output: ["catsdogcats","dogcatsdog","ratcatdogcat"]

Explanation: "catsdogcats" can be concatenated by "cats", "dog" and "cats";
  "dogcatsdog" can be concatenated by "dog", "cats" and "dog";
  "ratcatdogcat" can be concatenated by "rat", "cat", "dog" and "cat".
```

Note:

- 1. The number of elements of the given array will not exceed 10,000
- 2. The length sum of elements in the given array will not exceed 600,000.
- 3. All the input string will only include lower case letters.
- 4. The returned elements order does not matter.

Solution 1

Do you still remember how did you solve this problem? https://leetcode.com/problems/word-break/

If you do know one optimized solution for above question is using **DP**, this problem is just one more step further. We iterate through each **word** and see if it can be formed by using other **words**.

Of course it is also obvious that a word can only be formed by words shorter than it. So we can first sort the input by length of each word, and only try to form one word by using words in front of it.

```
public class Solution {
    public static List<String> findAllConcatenatedWordsInADict(String[] words) {
        List<String> result = new ArrayList<>();
        Set<String> preWords = new HashSet<>();
        Arrays.sort(words, new Comparator<String>() {
            public int compare (String s1, String s2) {
                return s1.length() - s2.length();
        });
        for (int i = 0; i < words.length; i++) {</pre>
            if (canForm(words[i], preWords)) {
                result.add(words[i]);
            preWords.add(words[i]);
        }
        return result;
    }
    private static boolean canForm(String word, Set<String> dict) {
        if (dict.isEmpty()) return false;
 boolean[] dp = new boolean[word.length() + 1];
 dp[0] = true;
 for (int i = 1; i <= word.length(); i++) {</pre>
     for (int j = 0; j < i; j++) {
  if (!dp[j]) continue;
  if (dict.contains(word.substring(j, i))) {
     dp[i] = true;
     break;
  }
 return dp[word.length()];
}
```

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Solution 2

Most of lines are adding words into Trie Tree

This solution is like putting two pointers to search through the tree. When find a word, put the other pointer back on root then continue searching.

But I'm not sure about the time complexity of my solution. Suppose word length is len and there are n words. Is the time complexity $O(len * n ^ 2)$?

```
public class Solution {
   class TrieNode {
        TrieNode[] children;
        String word;
        boolean isEnd;
        boolean combo; //if this word is a combination of simple words
        boolean added; //if this word is already added in result
        public TrieNode() {
            this.children = new TrieNode[26];
            this.word = new String();
            this.isEnd = false;
            this.combo = false;
            this.added = false;
        }
   private void addWord(String str) {
        TrieNode node = root;
        for (char ch : str.toCharArray()) {
            if (node.children[ch - 'a'] == null) {
                node.children[ch - 'a'] = new TrieNode();
            node = node.children[ch - 'a'];
        }
        node.isEnd = true;
       node.word = str;
   private TrieNode root;
   private List<String> result;
   public List<String> findAllConcatenatedWordsInADict(String[] words) {
        root = new TrieNode();
        for (String str : words) {
            if (str.length() == 0) {
                continue;
            addWord(str);
        result = new ArrayList<>();
        dfs(root, 0);
        return result;
   private void dfs(TrieNode node, int multi) {
     //multi counts how many single words combined in this word
        if(node.isEnd && !node.added && multi > 1) {
            node.combo = true;
            node.added = true;
            result.add(node.word);
        searchWord(node root multi):
```

```
private void searchWord(TrieNode node1, TrieNode node2, int multi) {
    if (node2.combo) {
        return;
    }
    if (node2.isEnd) {
        //take the pointer of node2 back to root
        dfs(node1, multi + 1);
    }
    for (int i = 0; i < 26; i++) {
        if (node1.children[i] != null && node2.children[i] != null) {
            searchWord(node1.children[i], node2.children[i], multi);
        }
    }
}</pre>
```

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Solution 3

Let's discuss whether a word should be included in our answer.

Consider the word as a topologically sorted directed graph, where each node is a letter, and an edge exists from i to j if word[i:j] is in our wordlist, [and there is no edge from i=o to j=len(word)-1]. We want to know if there is a path from beginning to end. If there is, then the word can be broken into concatenated parts from our wordlist. To answer this question, we DFS over this graph.

Code:

```
S = set(A)
ans = []
for word in A:
  if not word: continue
  stack = [0]
  seen = \{0\}
 M = len(word)
  while stack:
    node = stack.pop()
    if node == M:
      ans.append(word)
      break
    for j in xrange(M - node + 1):
      if (word[node:node+j] in S and
          node + j not in seen and
          (node > 0 or node + i != M):
        stack.append(node + j)
        seen.add(node + j)
return ans
```

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From Leetcoder.