

Contiguous Array

Given a binary array, find the maximum length of a contiguous subarray with equal number of 0 and 1.

Example 1:

Input: [0,1]

Output: 2

Explanation: [0, 1] is the longest contiguous subarray with equal number of 0 and 1.

Example 2:

Input: [0,1,0]

Output: 2

Explanation: [0, 1] (or [1, 0]) is a longest contiguous subarray with equal number of 0 and 1.

Note: The length of the given binary array will not exceed 50,000.

Solution 1

The idea is to change `0` in the original array to `-1`. Thus, if we find `SUM[i, j] == 0` then we know there are even number of `-1` and `1` between index `i` and `j`. Also put the `sum` to `index` mapping to a HashMap to make search faster.

```
public class Solution {
    public int findMaxLength(int[] nums) {
        for (int i = 0; i < nums.length; i++) {
            if (nums[i] == 0) nums[i] = -1;
        }

        Map<Integer, Integer> sumToIndex = new HashMap<>();
        sumToIndex.put(0, -1);
        int sum = 0, max = 0;

        for (int i = 0; i < nums.length; i++) {
            sum += nums[i];
            if (sumToIndex.containsKey(sum)) {
                max = Math.max(max, i - sumToIndex.get(sum));
            }
            else {
                sumToIndex.put(sum, i);
            }
        }

        return max;
    }
}
```

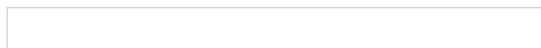
written by [shawngao](#) original link [here](#)

Solution 2

My idea is very similar to others, but let me try to explain it more visually. My thought was inspired by 121. Best Time to Buy and Sell Stock.

Let's have a variable `count` initially equals 0 and traverse through `nums`. Every time we meet a 0, we decrease `count` by 1, and increase `count` by 1 when we meet 1. It's pretty easy to conclude that we have a contiguous subarray with equal number of 0 and 1 when `count` equals 0.

What if we have a sequence `[0, 0, 0, 0, 1, 1]`? the maximum length is 4, the `count` starting from 0, will equal -1, -2, -3, -4, -3, -2, and won't go back to 0 again. But wait, the longest subarray with equal number of 0 and 1 started and ended when `count` equals -2. We can plot the changes of `count` on a graph, as shown below. Point (0,0) indicates the initial value of `count` is 0, so we count the sequence starting from index 1. The longest subarray is from index 2 to 6.



From above illustration, we can easily understand that two points with the same y-axis value indicates the sequence between these two points has equal number of 0 and 1.

Another example, sequence `[0, 0, 1, 0, 0, 0, 1, 1]`, as shown below,



There are 3 points have the same y-axis value -2. So subarray from index 2 to 4 has equal number of 0 and 1, and subarray from index 4 to 8 has equal number of 0 and 1. We can add them up to form the longest subarray from index 2 to 8, so the maximum length of the subarray is $8 - 2 = 6$.

Yet another example, sequence `[0, 1, 1, 0, 1, 1, 1, 0]`, as shown below. The longest subarray has the y-axis value of 0.



To find the maximum length, we need a dict to store the value of `count` (as the key) and its associated index (as the value). We only need to save a `count` value and its index at the first time, when the same `count` values appear again, we use the new index subtracting the old index to calculate the length of a subarray. A variable `max_length` is used to keep track of the current maximum length.

```
class Solution(object):
    def findMaxLength(self, nums):
        count = 0
        max_length=0
        table = {0: 0}
        for index, num in enumerate(nums, 1):
            if num == 0:
                count -= 1
            else:
                count += 1

            if count in table:
                max_length = max(max_length, index - table[count])
            else:
                table[count] = index

        return max_length
```

written by [NoAnyLove](#) original link [here](#)

Solution 3

Keeping track of the balance (number of ones minus number of zeros) and storing the first index where each balance occurred.

Python

Keeping the balance in units of 0.5 which makes the update expression short (not that `num * 2 - 1` or `1 if num else -1` would be terribly long):

```
def findMaxLength(self, nums):
    index = {0: -1}
    balance = maxlen = 0
    for i, num in enumerate(nums):
        balance += num - 0.5
        maxlen = max(maxlen, i - index.setdefault(balance, i))
    return maxlen
```

Just for fun as an ugly one-liner:

```
def findMaxLength(self, nums):
    return reduce(lambda f, b, m, (i, x): (f, b+x-.5, max(m, i-f.setdefault(b+x-.5, i))),
        enumerate(nums), ({0:-1}, 0, 0))[2]
```

Java

Using `putIfAbsent` so I only need one map function call per number.

```
public int findMaxLength(int[] nums) {
    Map<Integer, Integer> index = new HashMap<>();
    index.put(0, -1);
    int balance = 0, maxlen = 0;
    for (int i = 0; i < nums.length; i++) {
        balance += nums[i] * 2 - 1;
        Integer first = index.putIfAbsent(balance, i);
        if (first != null)
            maxlen = Math.max(maxlen, i - first);
    }
    return maxlen;
}
```

Could avoid using `Math.max` like this:

```
if (first != null && i - first > maxlen)
    maxlen = i - first;
```

written by [StefanPochmann](#) original link [here](#)

