COMP2402
Abstract Data Types and Algorithms

# Asymptotic Analysis and Other Math

Find the Smallest and Swap it with the First Position,
Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position

```
for i ∈ [0, length-1):
    nextsmallestindex ← i
    for j ∈ [i+1, length):
        if (item[j] < item[nextsmallest]):
            nextsmallest ← j
    swap item[i] with item[nextsmallest]</pre>
```

Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position

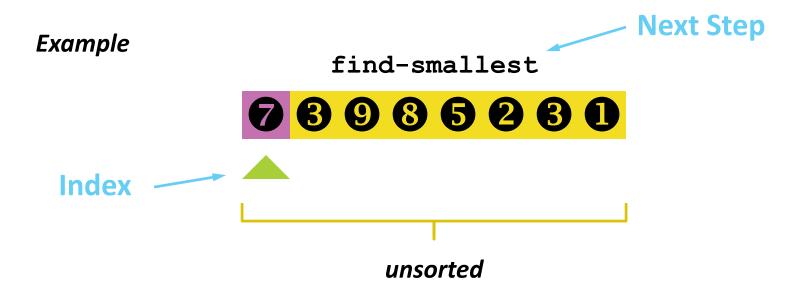


Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position

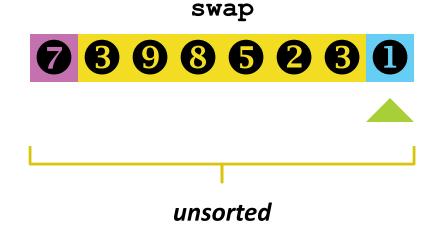


Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position



Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position

Example move-to-next

1 9 9 5 2 3 7



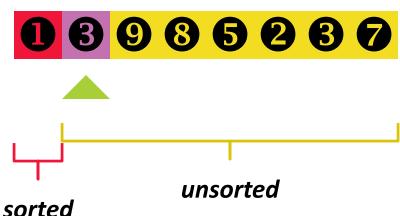
Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position

### Example find-smallest

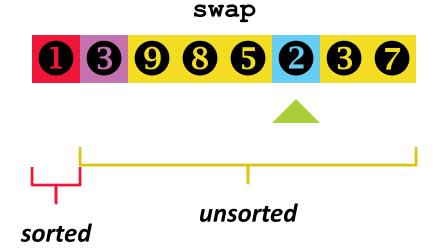


Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position



Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position

#### Example

move-to-next

1 2 9 8 6 8 6

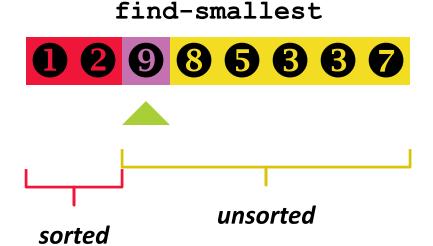
sorted

Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position

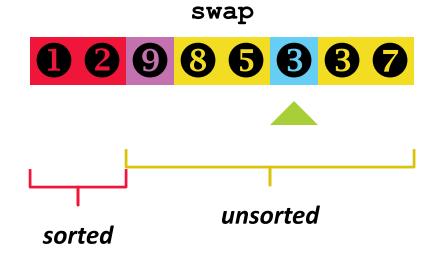


Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position

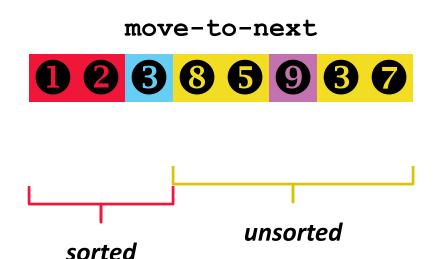


Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position



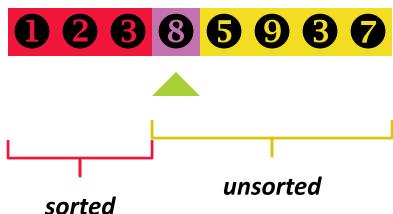
Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position

## Example find-smallest

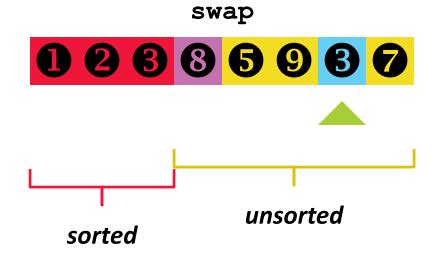


Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position



Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position

# move-to-next 1 2 3 3 5 9 8 7 unsorted

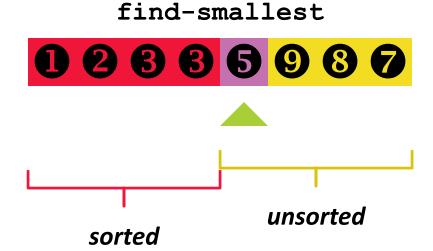
sorted

Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position

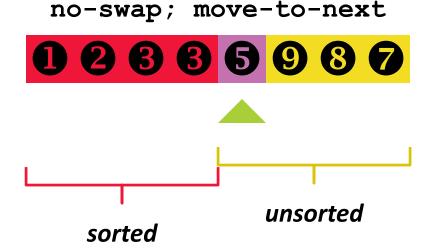


Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position

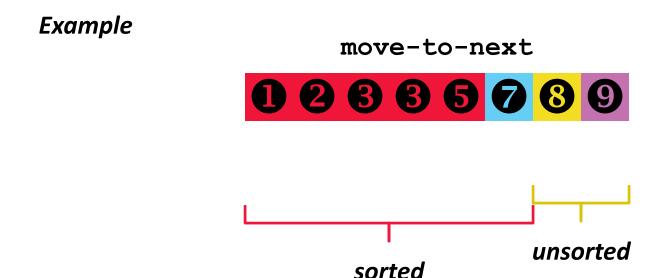


Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position



Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position

find-smallest

1 2 3 3 5 7 8 9

unsorted

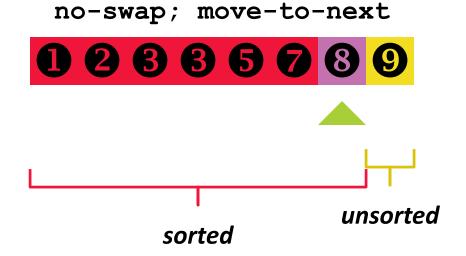
sorted

Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position

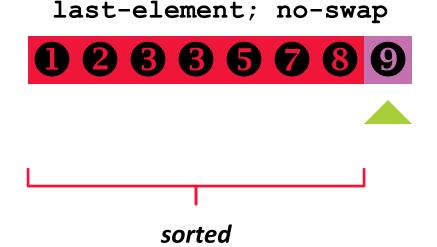


Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position



Find the Smallest and Swap it with the First Position,

Find the Next Smallest and Swap it with the Second Position,

...

Find the Second-to-Last Smallest and Swap it with the Second Last Position,
the Largest is Already in the Last Position

Example

finished.



Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

```
n ← length
do:
    flag ← false
    for i ∈ [1, n):
        if (item[i-1] > item[i]):
            swap item[i-1] with item[i]
            flag ← true
        n--
while (flag)
```

Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

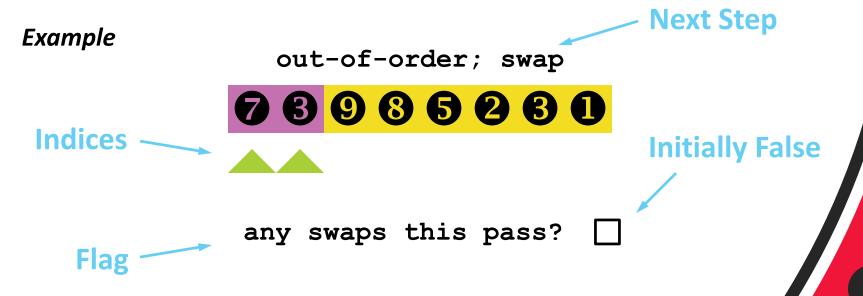
If Any Swaps were performed, Reset the Algorithm and Start Again



Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again



Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

move-to-next







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

no-swap; move-to-next







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

out-of-order; swap







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

move-to-next







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

out-of-order; swap







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

...

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

move-to-next







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

out-of-order; swap







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

move-to-next







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

out-of-order; swap







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

any-swaps? yes







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

reset-the-algorithm





Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example





Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

no-swap; move-to-next





Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

**Example** 

out-of-order; swap



Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

...

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

move-to-next







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

out-of-order; swap







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

...

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

move-to-next







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

out-of-order; swap







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

move-to-next







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

out-of-order; swap







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

•••

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

move-to-next







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

no-swap; move-to-next







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

any-swaps? yes







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

reset-the-algorithm





Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

no-swap; move-to-next





Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

out-of-order; swap





Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

move-to-next





Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

out-of-order; swap







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

...

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

move-to-next







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

out-of-order; swap







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

...

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

move-to-next







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

out-of-order; swap







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

move-to-next







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

no-swap; move-to-next







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

no-swap; move-to-next







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

any-swaps? yes







Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

• • •

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

reset-the-algorithm





any swaps this pass? [

**Fast Forward** 



Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

•••

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

no-swap; move-to-next





any swaps this pass?

**Fast Forward** 



Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

swap; move-to-next **8 2 5 3 0 7 8 9** 

**Fast Forward** 





Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

swap; move-to-next











Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

swap; move-to-next











Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

no-swap; move-to-next











Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

no-swap; move-to-next











Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

no-swap; move-to-next









Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

any-swaps? yes









Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

• • •

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

reset-the-algorithm





any swaps this pass? [

**Fast Forward** 



Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

 **Fast Forward** 





Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

•••

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

 **Fast Forward** 





Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

swap; move-to-next











Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

...

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

no-swap; move-to-next









Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

no-swap; move-to-next











Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

•••

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

no-swap; move-to-next



**Fast Forward** 





Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

**Example** 

no-swap; move-to-next











Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

any-swaps? yes



any swaps this pass?

**Fast Forward** 



Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

...

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

reset-the-algorithm





any swaps this pass? [

**Fast Forward** 



Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

**20335789** 

**Fast Forward** 





Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

Example

**1 2 3 3 5 7 8 9** 

**Fast Forward** 





Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

**Example** 



**Fast Forward** 



Compare the First Pair of Adjacent Elements and Swap If they are Out of Order Compare the Next Pair of Adjacent Elements and Swap If they are Out of Order

Compare the Last Pair of Adjacent Elements and Swap If they are Out of Order

If Any Swaps were performed, Reset the Algorithm and Start Again

finished.

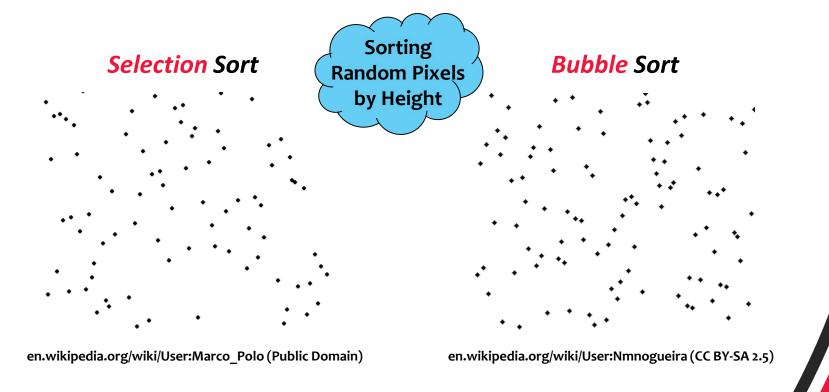
1 2 3 3 5 7 8 9

sorted

# Selection Sort vs. Bubble Sort

Selection Sort and Bubble Sort are Very Different Algorithms

(Finding Minimum Values vs. Exchanging Adjacent Elements)



**How do they Compare in terms of Efficiency?** 

a Fundamental Operation for Sorting Collections of Elements

is Accessing (i.e., Visiting) the Individual Elements

a Fundamental Operation for Sorting Collections of Elements

is Accessing (i.e., Visiting) the Individual Elements

**How Often does Selection Sort Visit an element?** 

**How Often does Bubble Sort Visit an element?** 

a Fundamental Operation for Sorting Collections of Elements

is Accessing (i.e., Visiting) the Individual Elements

**How Often does Selection Sort Visit an element?** 

**How Often does Bubble Sort Visit an element?** 

in Order to Produce a Generalizable Result

**Express the Solution as a Function of the Size of the Collection** 

we can Examine the Implementation of swap in the Source for java.util.Collections to see How Many get/set Calls are Made

```
2119: /**
2120:
       * Swaps the elements at the specified positions within the list. Equal
2121: * positions have no effect.
2122:
2123:
      * @param 1 the list to work on
2124: * @param i the first index to swap
2125: * @param j the second index
2126: * @throws UnsupportedOperationException if list.set is not supported
2127: * @throws IndexOutOfBoundsException if either i or j is < 0 or &qt;=
2128: *
                list.size()
2129:
       * @since 1.4
2130:
2131: public static void swap (List<?> 1, int i, int j)
2132: {
2133: List<Object> list = (List<Object>) 1;
2134:
        list.set(i, list.set(j, list.get(i)));
2135: }
```

## **Analyzing Selection Sort Performance**

#### **Every Outer Loop Requires**

#### **Swap and Finding Next Smallest**

```
(i.e., (3 + 2(n-1)) + (3 + 2(n-2)) + ...)
```

```
for i ∈ [0, length-1):
    nextsmallestindex ← i
    for j ∈ [i+1, length):
        if (item[j] < item[nextsmallest]):
            nextsmallest ← j
    swap item[i] with item[nextsmallest]</pre>
```

$$= \sum_{i=0}^{n-2} \sum_{i=0}^{n-2} (3 + 2(n-1-i))$$

$$= 3(n-1) \sum_{i=0}^{n-2} \sum_{i=0}^{n-2} 2(n-1-i)$$

$$= 3(n-1) + (2(n-1)(n-1)) + \sum_{i=0}^{n-2} \sum_{i=0}^{n-2} 2i$$

$$= 3(n-1) + (2(n-1)(n-1)) - 2 \sum_{i=0}^{n-2} i$$

$$= 3(n-1) + (2(n-1)(n-1)) - 2 ((n-2)(n-1)/2)$$

$$= n^2 + 2n - 3$$

# **Analyzing Bubble Sort Performance**

for an array of Length n:

```
n ← length
do:
    flag ← false
    for i ∈ [1, n):
        if (item[i-1] > item[i]):
            swap item[i-1] with item[i]
            flag ← true
        n--
while (flag)
```

What is the Maximum number of Do-While Loops? n

What is the Maximum number of Swaps on First Pass? n-1

What is the Maximum number of Swaps on Next Pass? n-2

(n.b., a Swap is still 3 Array Accesses)

# **Analyzing Bubble Sort Performance**

do:

n - length

flag ← false

for  $i \in [1, n)$ :

if (item[i-1] > item[i]):

flag ← true

swap item[i-1] with item[i]

#### **Every Outer Loop Requires**

#### that all Remaining Pairs are

#### **Compared and Swapped**

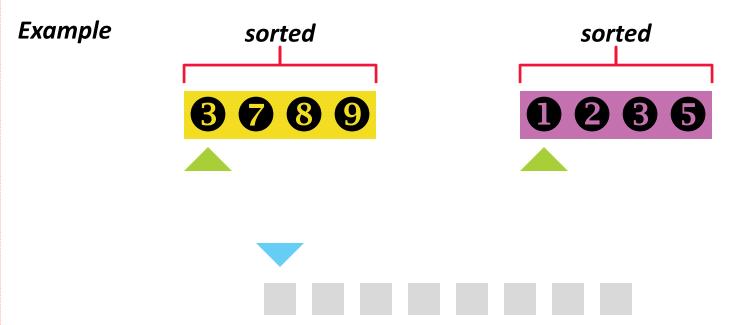
(i.e.,  $(n-1) + (n-2) + (n-3) + \ldots + 1$ )

```
while (flag)
= \sum_{i=1}^{n-1} \left( (2 + 3) (n-i) \right)
= 5 \left( \sum_{i=1}^{n-1} n - \sum_{i=1}^{n-1} i \right)
= 5 \left( (n-1) (n) - \sum_{i=1}^{n-1} i \right)
= 5 \left( (n-1) (n) - ((n-1) (n) / 2) \right)
= 5 \left( ((n-1) (n) / 2 \right)
```

$$= \frac{5}{2} (n^2 - n)$$

If you were Provided Two Sorted Lists (of lengths x and y)

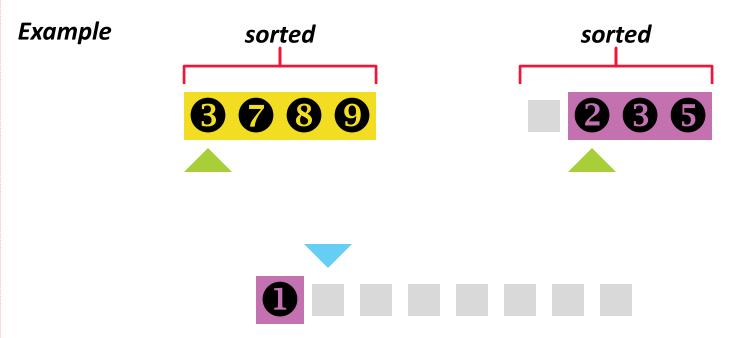
Could you Merge them into a Single Sorted List (of length x+y)?



Separate Indices into the New List and the Sorted Sublists

If you were Provided Two Sorted Lists (of lengths x and y)

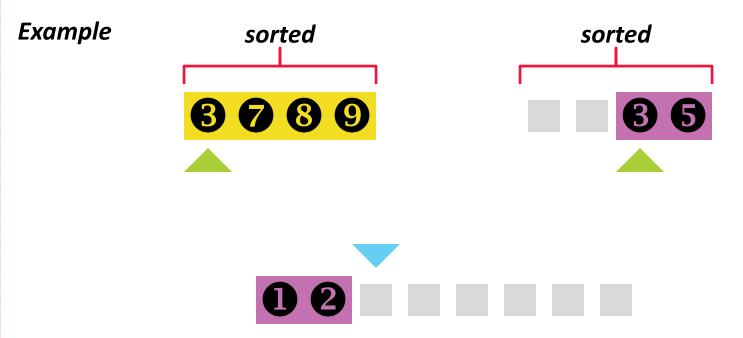
Could you Merge them into a Single Sorted List (of length x+y)?



Separate Indices into the New List and the Sorted Sublists

If you were Provided Two Sorted Lists (of lengths x and y)

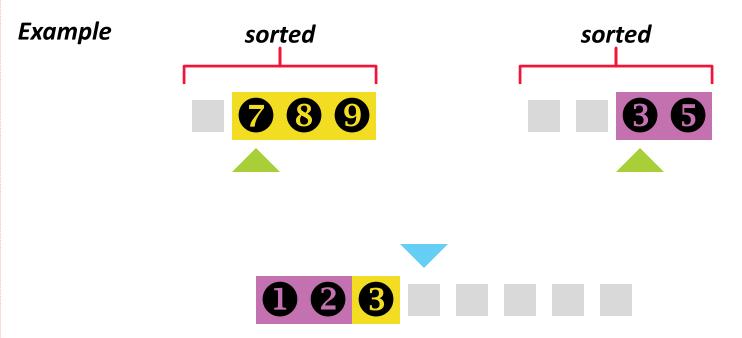
Could you Merge them into a Single Sorted List (of length x+y)?



Separate Indices into the New List and the Sorted Sublists

If you were Provided Two Sorted Lists (of lengths x and y)

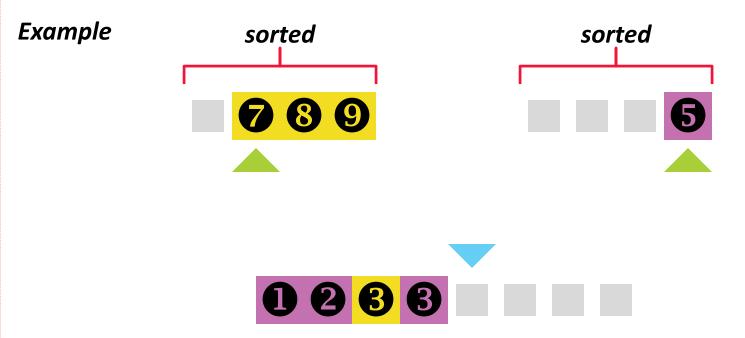
Could you Merge them into a Single Sorted List (of length x+y)?



Separate Indices into the New List and the Sorted Sublists

If you were Provided Two Sorted Lists (of lengths x and y)

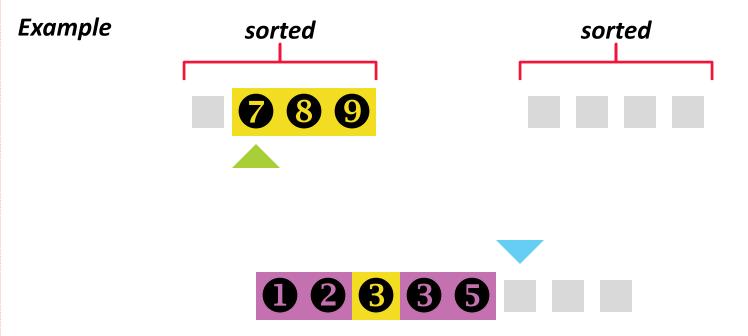
Could you Merge them into a Single Sorted List (of length x+y)?



Separate Indices into the New List and the Sorted Sublists

If you were Provided Two Sorted Lists (of lengths x and y)

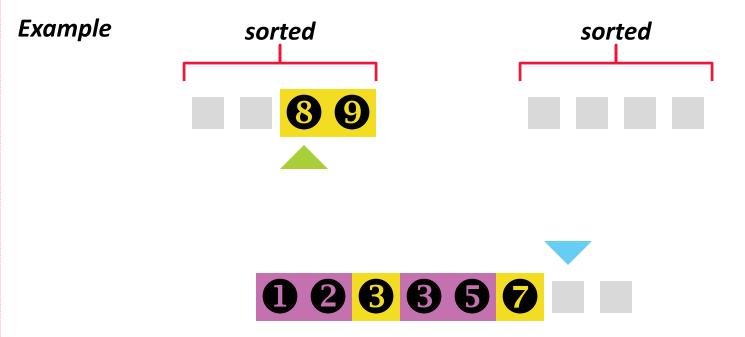
Could you Merge them into a Single Sorted List (of length x+y)?



Separate Indices into the New List and the Sorted Sublists

If you were Provided Two Sorted Lists (of lengths x and y)

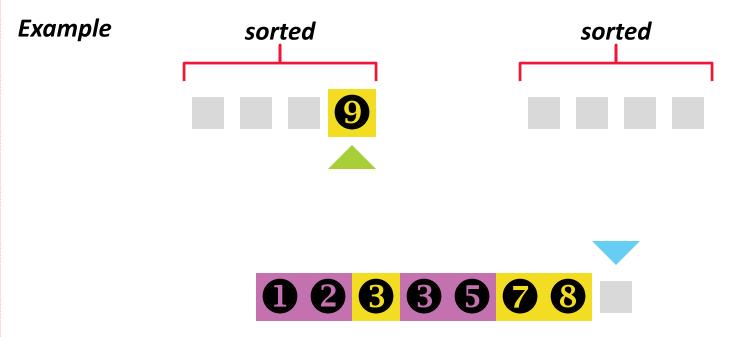
Could you Merge them into a Single Sorted List (of length x+y)?



Separate Indices into the New List and the Sorted Sublists

If you were Provided Two Sorted Lists (of lengths x and y)

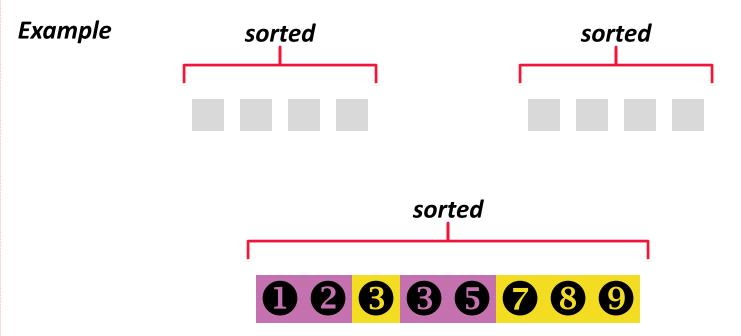
Could you Merge them into a Single Sorted List (of length x+y)?



Separate Indices into the New List and the Sorted Sublists

If you were Provided Two Sorted Lists (of lengths x and y)

Could you Merge them into a Single Sorted List (of length x+y)?



Separate Indices into the New List and the Sorted Sublists

# **Additional Leading Questions**

Two Unsorted Lists X and Y and the Length of X > the Length of Y

Which of the Two Lists is (in General) Easier to Sort?

How do you Sort a Singleton List (i.e., a List of Length 1)?

# **Additional Leading Questions**

Two Unsorted Lists X and Y and the Length of X > the Length of Y

Which of the Two Lists is (in General) Easier to Sort?

How do you Sort a Singleton List (i.e., a List of Length 1)?



**Procedure for Sorting a List:** 

Divide the Unsorted List (of Length L)

into Two Lists (of Length \[ \L/2 \] and \[ \L/2 \] respectively)

Sort the Sublists and then Merge them into a Single Sorted List

# **Additional Leading Questions**

Two Unsorted Lists X and Y and the Length of X > the Length of Y

Which of the Two Lists is (in General) Easier to Sort?

How do you Sort a Singleton List (i.e., a List of Length 1)?



Procedure for Sorting a List:

Divide the Unsorted List (of Length L)

into Two Lists (of Length \[ \L/2 \] and \[ \L/2 \] respectively)

Sort the Sublists and then Merge them into a Single Sorted List

Can the Procedure Itself be Used to Solve this Subproblem?

**Recursive Approach** 

If the Unsorted List is of Length 1, Return

Otherwise, Divide the list in Half (approximately) into Two Sublists

Recursively Call Merge Sort on Each Sublist and Merge the Return Values

```
sort(item):
    if (len(item) \leq 1):
        return item
    else
        right = sort(item[len(item)/2:len(item)])
        return merge(left, right)
                        merge(left, right):
                            j = 0, k = 0
                            for i ∈ [0, len(left)+len(right)):
                                if left[j] < right[k]:</pre>
                                    item[i] ← left[j] ; j++
                                else:
                                    item[i] ← right[k] ; k++
                            return item
```

**Example** 



```
sort( 73935230)
```

```
sort(item[0:len(item)/2])
sort(item[len(item)/2:len(item)])
merge(left, right)
```

**Example** 



```
sort( 73935230)
```

```
merge(sort( 7393), sort( 5231))
```

```
sort(item[0:len(item)/2])
sort(item[len(item)/2:len(item)])
merge(left, right)
```

**Example** 



```
sort(item[0:len(item)/2])
sort(item[len(item)/2:len(item)])
merge(left, right)
```



```
sort( 7 3 9 8 5 2 3 1 )

merge(sort( 7 3 9 3 ), sort( 5 2 3 1 ))

merge(sort( 7 3 ), sort( 9 3 ))

merge(sort( 7 ), sort( 3 ))
```

```
sort(item[0:len(item)/2])
sort(item[len(item)/2:len(item)])
merge(left, right)
```



```
sort( 7 3 9 3 5 2 3 1 )

merge( sort( 7 3 9 3 ), sort( 5 2 3 1 )))

merge( sort( 7 3 ), sort( 9 3 ))

merge( sort( 7 ), sort( 3 )))
```

```
sort(item[0:len(item)/2])
sort(item[len(item)/2:len(item)])
merge(left, right)
```



```
sort( 73935230)
merge(sort(7398),sort(5230))
   merge(sort(78),sort(98))
      merge( 7, 3)
      8 7
```

```
sort(item[0:len(item)/2])
sort(item[len(item)/2:len(item)])
merge(left, right)
```



```
sort( 73935230)
merge(sort(7398),sort(5230))
  merge(sort(78),sort(98))
     merge( 7, 3)
     8 7
```

```
sort(item[0:len(item)/2])
sort(item[len(item)/2:len(item)])
merge(left, right)
```



```
sort( 73935230)
merge(sort(7398),sort(5230))
  merge(sort(78),sort(98))
     merge( 7, 3 ), merge( 9, 8 )
     8 7
```

```
sort(item[0:len(item)/2])
sort(item[len(item)/2:len(item)])
merge(left, right)
```



```
sort( 73935230)
merge(sort(7398),sort(5230))
  merge(sort(78),sort(98))
     merge( 7, 3 ), merge( 9, 8 )
     8 9
```

```
sort(item[0:len(item)/2])
sort(item[len(item)/2:len(item)])
merge(left, right)
```



```
sort( 73935230)
merge(sort( 7393), sort( 5230))
merge( merge( sort( 73 )), sort( 98 ))
   merge( merge( 7, 8 ), merge( 9, 8 ))
   merge( 37, 39)
   8 8 8 9
```

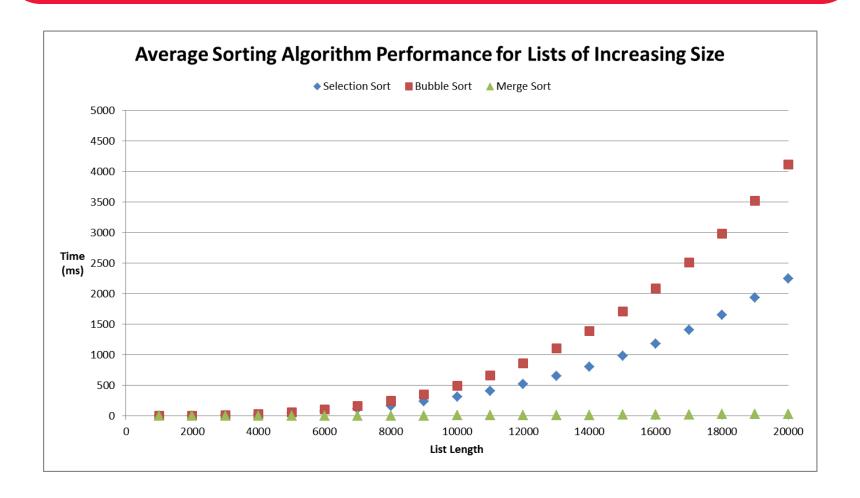
#### **Example**

0 2 3 3 5 7 8 9

```
78985280
```

```
sort(item[0:len(item)/2])
sort( 73985230)
                           sort(item[len(item)/2:len(item)])
                               merge(left, right)
merge(merge(sort(73)),sort(98)),merge(sort(52),sort(30)))
merge(merge(sort(5),sort(2)),merge(sort(3),sort(1))))
merge ( merge ( 3 7 , 3 9 ), merge ( 2 5 , 1 3 ) )
merge( 3789, 1235)
```

# **Measuring Real-Time Performance**



**Merge Sort Outclasses both Selection and Bubble Sort** 

(n.b., Not Just Faster; the Curve is Different!)

## **Analyzing Merge Sort Performance**

for an array of Length n (to simplify assume  $n = 2^m$  for  $m \in \mathbb{Z}^+$ )

How Many Accesses to Split an Array of Length n in two? n

How Many Accesses to Merge Two Arrays of Length <sup>n</sup>/<sub>2</sub>? n

If Merge Sorting an array of Length n Required T Accesses

How Many Accesses to Merge Sort array of Length  $^{n}/_{2}$ ?

```
sort(item):
                                      merge(left, right):
    if (len(item) \leq 1):
                                           j = 0, k = 0
         return item
                                           for i ∈ [0, len(left)+len(right)):
    else
                                                if left[j] < right[k]:</pre>
         left ← sort(item[...
                                                    item[i] ← left[j] ; j++
         right \( \text{sort(item[...} \)
                                               else:
         return merge(left, right)
                                                    item[i] ← right[k] ; k++
                                           return item
T(n) = T(^{n}/_{2}) + T(^{n}/_{2}) + 2n
          = 2 (T(^{n}/_{2})) + 2n
```

```
sort(item):
                                      merge(left, right):
    if (len(item) \leq 1):
                                           j = 0, k = 0
         return item
                                           for i ∈ [0, len(left)+len(right)):
    else
                                                if left[j] < right[k]:</pre>
         left ← sort(item[...
                                                    item[i] ← left[j] ; j++
         right \( \text{sort(item[...} \)
                                                else:
         return merge(left, right)
                                                    item[i] ← right[k] ; k++
                                           return item
T(n) = T(^{n}/_{2}) + T(^{n}/_{2}) + 2n
          = 2 (T(^{n}/_{2})) + 2n
```

$$T(^{n}/_{2}) = T(^{n}/_{4}) + T(^{n}/_{4}) + 2(^{n}/_{2})$$
  
=  $2(T(^{n}/_{4})) + 2(^{n}/_{2})$ 

```
sort(item):
                                      merge(left, right):
    if (len(item) \leq 1):
                                           j = 0, k = 0
         return item
                                           for i ∈ [0, len(left)+len(right)):
    else
                                               if left[j] < right[k]:</pre>
         left ← sort(item[...
                                                    item[i] ← left[j] ; j++
         right \( \text{sort(item[...} \)
                                               else:
         return merge(left, right)
                                                    item[i] ← right[k] ; k++
                                           return item
T(n) = T(^{n}/_{2}) + T(^{n}/_{2}) + 2n
          = 2 (T(^{n}/_{2})) + 2n
          = 2(2(T(^{n}/_{4})) + (^{n}/_{2})) + 2n
```

$$T(^{n}/_{2}) = T(^{n}/_{4}) + T(^{n}/_{4}) + 2(^{n}/_{2})$$
  
=  $2(T(^{n}/_{4})) + 2(^{n}/_{2})$ 

```
sort(item):
                                     merge(left, right):
    if (len(item) \leq 1):
                                          j = 0, k = 0
        return item
                                          for i ∈ [0, len(left)+len(right)):
    else
                                              if left[j] < right[k]:</pre>
        left ← sort(item[...
                                                   item[i] ← left[j] ; j++
        right \( \text{sort(item[...} \)
                                              else:
        return merge(left, right)
                                                   item[i] ← right[k] ; k++
                                          return item
T(n) = T(^{n}/_{2}) + T(^{n}/_{2}) + 2n
         = 2(T(^{n}/_{2})) + 2n
         = 2(2(T(^{n}/_{4})) + 2(^{n}/_{2})) + 2n
         = 2(2(T(^{n}/_{4}))) + 2n + 2n
```

```
sort(item):
                                     merge(left, right):
    if (len(item) \leq 1):
                                         j = 0, k = 0
        return item
                                         for i ∈ [0, len(left)+len(right)):
    else
                                              if left[j] < right[k]:</pre>
        left ← sort(item[...
                                                  item[i] ← left[j] ; j++
        right \( \text{sort(item[...} \)
                                              else:
        return merge(left, right)
                                                  item[i] ← right[k] ; k++
                                         return item
T(n) = T(^{n}/_{2}) + T(^{n}/_{2}) + 2n
         = 2 (T(^{n}/_{2})) + 2n
         = 2(2(T(^{n}/_{4})) + 2(^{n}/_{2})) + 2n
         = 2(2(T(n/4))) + 2n + 2n
```

$$T(^{n}/_{4}) = T(^{n}/_{8}) + T(^{n}/_{8}) + 2(^{n}/_{4})$$
  
=  $2(T(^{n}/_{8})) + 2(^{n}/_{4})$ 

let T (n) be the Number of Accesses to Sort array of Length n

```
sort(item):
                                     merge(left, right):
    if (len(item) \leq 1):
                                         j = 0, k = 0
        return item
                                         for i ∈ [0, len(left)+len(right)):
    else
                                              if left[j] < right[k]:</pre>
        left ← sort(item[...
                                                  item[i] ← left[j] ; j++
        right \( \text{sort(item[...} \)
                                              else:
        return merge(left, right)
                                                  item[i] ← right[k] ; k++
                                         return item
T(n) = T(^{n}/_{2}) + T(^{n}/_{2}) + 2n
         = 2(T(^{n}/_{2})) + 2n
         = 2(2(T(^{n}/_{4})) + 2(^{n}/_{2})) + 2n
         = 2(2(T(n/4))) + 2n + 2n
         = 2(2(2(T(^{n}/_{8}))) + 2(^{n}/_{4}))) + 2n + 2n
T(^{n}/_{4}) = T(^{n}/_{8}) + T(^{n}/_{8}) + 2(^{n}/_{4})
```

 $= 2 (T(^{n}/_{8})) + 2 (^{n}/_{4})$ 

```
sort(item):
                                    merge(left, right):
    if (len(item) \leq 1):
                                        j = 0, k = 0
        return item
                                        for i ∈ [0, len(left)+len(right)):
    else
                                            if left[j] < right[k]:</pre>
        left ← sort(item[...
                                                 item[i] ← left[j] ; j++
        right \( \text{sort(item[...} \)
                                            else:
        return merge(left, right)
                                                 item[i] ← right[k] ; k++
                                        return item
T(n) = T(^{n}/_{2}) + T(^{n}/_{2}) + 2n
         = 2(T(^{n}/_{2})) + 2n
         = 2(2(T(^{n}/_{4})) + 2(^{n}/_{2})) + 2n
         = 2(2(T(^{n}/_{4}))) + 2n + 2n
         = 2(2(2(T(^{n}/_{8}))) + 2(^{n}/_{4}))) + 2n + 2n
         = 2(2(2(T(^{n}/_{8})))) + 2n + 2n + 2n
```

```
sort(item):
                                   merge(left, right):
    if (len(item) \leq 1):
                                       j = 0, k = 0
        return item
                                       for i ∈ [0, len(left)+len(right)):
    else
                                           if left[j] < right[k]:</pre>
        left ← sort(item[...
                                               item[i] ← left[j] ; j++
        right ← sort(item[...
                                           else:
        return merge(left, right)
                                               item[i] ← right[k] ; k++
                                       return item
T(n) = T(^{n}/_{2}) + T(^{n}/_{2}) + 2n
         = 2 (T(n/2)) + 2n
         = 2(2(T(^{n}/_{4})) + 2(^{n}/_{2})) + 2n
         = 2(2(T(^{n}/_{4}))) + 2n + 2n
         = 2(2(2(T(^{n}/_{8}))) + 2(^{n}/_{4}))) + 2n + 2n
         = 2(2(2(T(^{n}/_{8})))) + 2n + 2n + 2n
     ...after k steps...
         = 2^{k} (T(n/_{2}^{k})) + 2nk
```

let T (n) be the Number of Accesses to Sort array of Length n

```
sort(item):
                                      merge(left, right):
    if (len(item) \leq 1):
                                          j = 0, k = 0
         return item
                                           for i ∈ [0, len(left)+len(right)):
    else
                                               if left[j] < right[k]:</pre>
         left ← sort(item[...
                                                    item[i] ← left[j] ; j++
         right ← sort(item[...
                                               else:
         return merge(left, right)
                                                   item[i] ← right[k] ; k++
                                           return item
T(n) = 2^{k} (T(^{n}/_{2^{k}})) + 2nk
```

...Recall the Assumption that  $n = 2^m$  for Integer m...

...How Many Steps (k) Until we Evaluate T (1)? m

```
= 2^{m}(T(^{n}/_{2^{m}})) + 2nm  n = 2^{m}

= n(T(1)) + 2nm  m = \log_{2} n

= n(T(1)) + 2n(\log_{2}(n))

...and this is actually 0!
```

#### **Complexity Analyses for the Sorting Algorithms can be Compared**

Selection Sort = 
$$n^2 + 2n - 3$$

Bubble Sort = 
$$\frac{5}{2}(n^2 - n)$$

Merge Sort = 
$$2(n)(\log_2 n)$$

at Very Large Values of n (i.e., very large lists)

**Higher-Order Terms Dominate (Grow Faster) than Lower Order** 

( at Very Large Values of n: n<sup>2</sup> >> n log n >> n >> log n )

**Remove the Lower Order Terms** 

#### **Complexity Analyses for the Sorting Algorithms can be Compared**

Selection Sort = 
$$n^2 + 2n - 3$$
  
Bubble Sort =  $\frac{5}{2}(n^2 - n)$   
Merge Sort =  $2(n)(\log_2 n)$ 

Constants and Constant Co-efficients can also be Removed Why?

#### **Complexity Analyses for the Sorting Algorithms can be Compared**

Selection Sort = 
$$n^2 + 2n - 3$$
  
Bubble Sort =  $\frac{5}{2}(n^2 - n)$   
Merge Sort =  $2(n)(\log_2 n)$ 

**Constants and Constant Co-efficients can also be Removed** 

Why?

3n<sup>2</sup> > 2n<sup>2</sup> but Difference is Relatively Insignificant

(more Concerned with Changing Order than constants)

## **Complexity Analysis and Big-Oh Notation**

#### **Algorithms are typically Compared using Big-Oh Notation**

## **Complexity Analysis and Big-Oh Notation**

**Algorithms are typically Compared using Big-Oh Notation** 

$$f(n) = O(g(n))$$
 "Function  $f(n)$  is Big-Oh  $g(n)$ "

there are Two Constants c and k such that

 $0 \le f(n) \le cg(n)$  for all  $n \ge k$ 

"Function f Does Not Grow Any Faster than Function g"

to Convert a Formula into it's Big-Oh Notation

Remove all Constants and Constant Co-efficients

and Remove Everything Except the Highest-Order Term

## **Search Algorithms Revisited**

given an Telephone Directory (i.e., Phonebook)

there are Several Approaches for Searching it (since it is Sorted)

#### **Naïve Approach**

starting from the first page of the phonebook, check for the name "Collier" on each page and, if you don't find it, move on to the next page and continue the search until you check the whole book

#### **Alternative Approach**

open the phonebook to the center and compare the name on that page with "Collier" and, if "Collier" isn't there, at least you know which half will contain it, so you start again with only half the book

**Describe Each Approach in Pseudocode** 

**Express the Time Complexity of each (using Big-Oh)** 

## **Linear Search Pseudocode**

```
linearsearch(list, key, low, high):
   for i ∈ [0, len(list)):
      if (list[i] == key):
        return i
   return "key not found"
```

# **Binary Search Pseudocode**

```
binarysearch(list, key, low, high):
   if (high < low):</pre>
       return "key not found"
   else
       center ← (low + high) / 2
       if (list[center] > key):
          return binarysearch(list, key, low, center-1)
       else if (list[center] < key):</pre>
          return binarysearch(list, key, center+1, high)
       else
          return center
```

given an Telephone Directory in which you must Search for a Key if you Use the Linear Search Algorithm:

What is the Best-Case Scenario?

(i.e., What would be the Easiest to Find?)

If the Directory has n Pages,

What is the Worst-Case Scenario?

given an Telephone Directory in which you must Search for a Key if you Use the Linear Search Algorithm:

What is the Best-Case Scenario?

(i.e., What would be the Easiest to Find?)

on the First Page

1 Page Searched

If the Directory has n Pages,

What is the Worst-Case Scenario?

given an Telephone Directory in which you must Search for a Key if you Use the Linear Search Algorithm:

What is the Best-Case Scenario?

(i.e., What would be the Easiest to Find?)

on the First Page

1 Page Searched

If the Directory has n Pages,

What is the Worst-Case Scenario?

on the Last Page

n Pages Searched

given an Telephone Directory in which you must Search for a Key if you Use the Linear Search Algorithm:

What is the Best-Case Scenario?

(i.e., What would be the Easiest to Find?)

on the First Page

1 Page Searched

If the Directory has n Pages,

What is the Worst-Case Scenario?

on the Last Page

n Pages Searched

What is the Average-Case Scenario?

given an Telephone Directory in which you must Search for a Key if you Use the Linear Search Algorithm:

What is the Best-Case Scenario?

(i.e., What would be the Easiest to Find?)

on the First Page

1 Page Searched

If the Directory has n Pages,

What is the Worst-Case Scenario?

on the Last Page

n Pages Searched

What is the Average-Case Scenario?

on Average you must Search Half (i.e., n/2)

given an Telephone Directory in which you must Search for a Key if you Use the Binary Search Algorithm:

What is the Best-Case Scenario?

(i.e., What would be the Easiest to Find?)

If the Directory has n Pages,

What is the Worst-Case Scenario?

given an Telephone Directory in which you must Search for a Key if you Use the Binary Search Algorithm:

What is the Best-Case Scenario?

on the First Page

(i.e., What would be the Easiest to Find?)

1 Page Searched

If the Directory has n Pages,

What is the Worst-Case Scenario?

given an Telephone Directory in which you must Search for a Key if you Use the Binary Search Algorithm:

What is the Best-Case Scenario?

(i.e., What would be the Easiest to Find?)

on the First Page

1 Page Searched

If the Directory has n Pages,

What is the Worst-Case Scenario?

?

How can we Determine the Worst-Case

**Time Complexity of Binary Search?** 

## **Analyzing Binary Search Performance**

let T (n) be the Number of Accesses to Search Array of Length n

```
binarysearch(list, key, low, high):
                                              Access Center and then
     center ← (low + high) / 2
                                           Search an Array of Length <sup>n</sup>/<sub>2</sub>
     if (list[center] > key):
         return binarysearch(list, key, low, center-1)
     else if (list[center] < key):</pre>
         return binarysearch(list, key, center+1, high)
     else
         return center
T(n) = 1 + T(^{n}/_{2})
T(n) = 1 + 1 + T(n/4) = 2 + T(n/4)
T(n) = 1 + 1 + 1 + ... + T(n/2k) = k + T(n/2k)
    ...Assume that n = 2^m for Integer m...
    ...How Many Steps (k) Until we Evaluate T (1)?
                                                                    m
```

 $T(n) = m + T(^{n}/_{2}^{m})$ 

 $T(n) = \log_2(n) + T(1)$ 

**Complexity Analyses for the Searching Algorithms** 

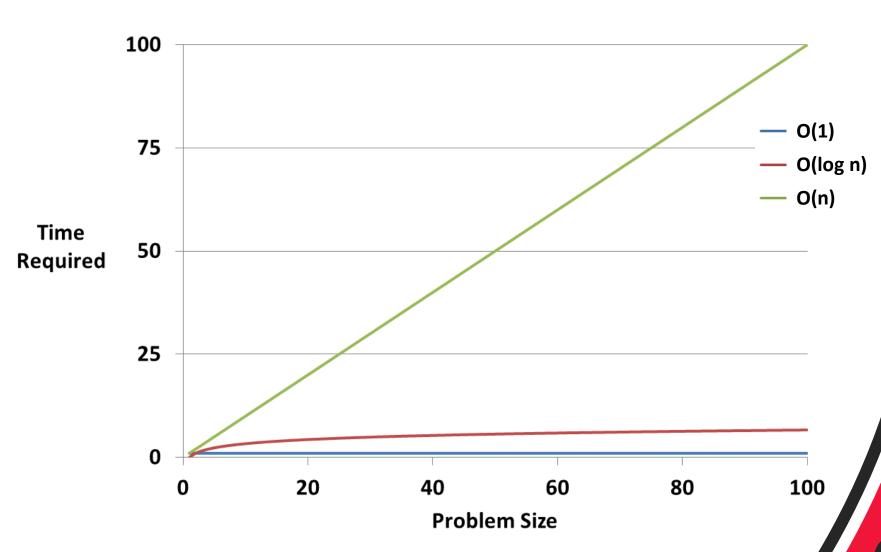
(assuming the Worst-Case Scenario)

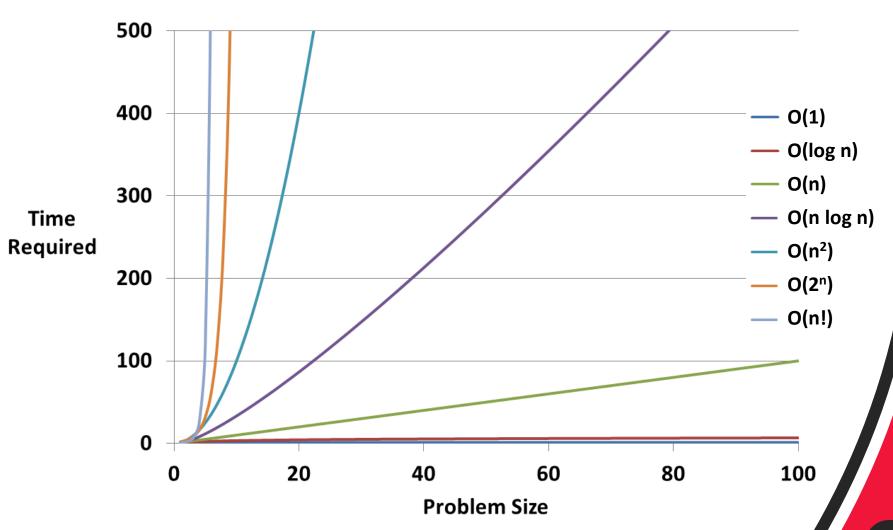
Binary Search = 
$$log_2$$
 n

If the Array to be Searched Doubles in Size

**How Much Worse is Linear Search? Binary Search?** 

n vs. 2n 
$$log_2(n)$$
 vs.  $log_2(2n)$   $1 + log_2(n)$ 





# **Time Complexity Terms**

Time Complexity	Term	Reaction
O(1)	"Constant"	
O(log n)	"Logarithmic"	
O(n)	"Linear"	
O(n log n)	"Quasilinear"	
O(n²)	"Polynomial"	
O(2 <sup>n</sup> )	"Exponential"	
O(n!)	"Factorial"	