PRINCIPLES OF COMPUTER NETWORKS
COMP 3203

Evangelos Kranakis

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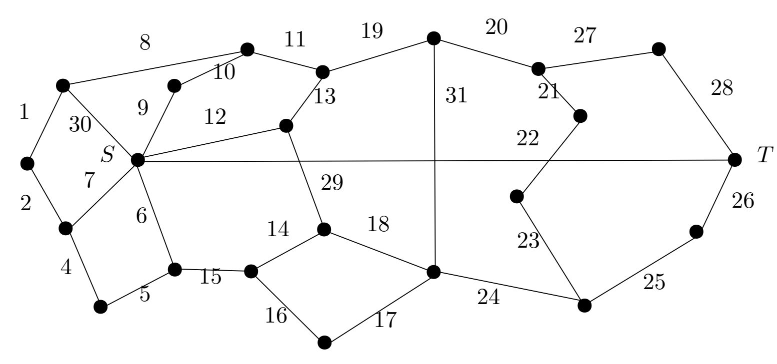
Solve the problems below. Your answers do not have to be long, but they should be complete, precise, concise and clear. Write the solutions on your own and acknowledge your sources in case you used "library" material. Look in the course web page on how to avoid plagiarism and for submission details (where/when/how). Exercises marked with ($\star$) are usually more challenging. **NB:** some of the exercises may require material that will be covered in forthcoming lectures while others marked with ($\star$) may involve additional material not covered in class. It is preferred that you type your work using your favourite software package but you submit only in pdf. Two excellent and free packages are LATEX (for typesetting mathematics) and Ipe (for drawing pictures).

Assignment C

1 [10 pts]

The links of a wireless network are labeled as depicted in the Figure below. Moreover there is a source node S and a destination node T.



Picture Revised

1. **[2 pts]** List the faces traversed when using face routing from node S to node T. For each face list the set of edges

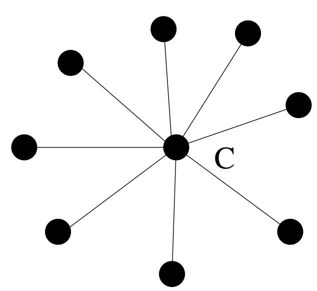
2. [**4 pts**] Apply the left-hand rule in face routing algorithm to give a path from node S to node T.

3. [**4 pts**] Apply the right-hand rule in face routing algorithm to give a path from node S to node T.

Note that if $e = (u, v)$ is the last edge you discover (which is crossing the line ST) during the traversal of a face then start the traversal of the next face from the vertex u: you must list this edge but do not have to repeat this edge in your traversal. In your answer you must list all the links being traversed including the last one crossed and the paths formed must use the corresponding face routing algorithm!

2 [10 pts]

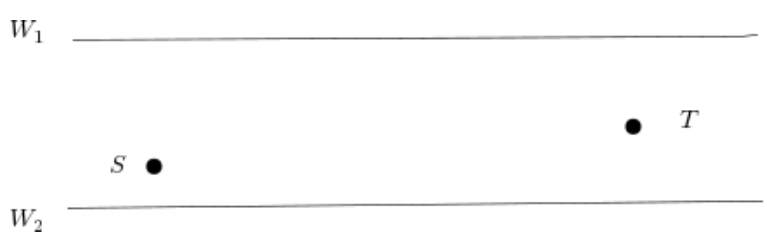
A wireless network consists of $k + 1$ omnidirectional sensors of identical range (radius) equal to one. They form a hub (star graph) with one sensor at the center C



and the remaining set S of k other sensors around it and within range of the center C. Assume that the sensors in S are outside the range of each other. Show that if all the sensors have the same range then $k \leq 5$.

3 [10 pts] ($\star$)

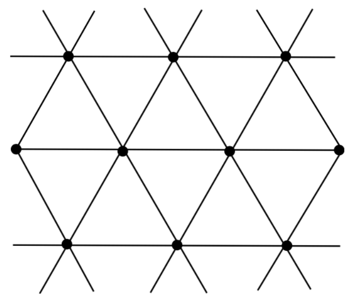
Consider two parallel mirrors, depicted as W_1, W_2, at distance $w = 5$ from each other. A source S and a target T are located between the mirrors at respective heights $s = 1$ and $t = 2$ from W_2.



Assume the horizontal distance between S and T is $d = 10$. At what angle, say ϕ, should a ray leave S so as to reach T after one reflection on W_1?

4 [10 pts]

We want to cover a large city by placing antennas at the vertices of a regular triangular network. Two neighboring antennas are at a distance a, the side length of the equilateral



triangles building up the network.

1. What point inside the triangle is the furthest away from all three corners and why?
2. What is the distance of this point from the corners of the triangle (as a function of the length a)?
3. If identical antennae are to be placed at the vertices of the triangles, what is the minimum radius that will ensure that every point in the plane is within the range of an antenna?

5 [10 pts]

In the questions below provide the formulas and explain your reasoning.

1. [**5 pts**] A wireless transmission system has bit-error rate p bits per sec $(0 < p < 1)$ and packet length of n bits. What is the probability that a packet has an error?
2. [**5 pts**] What is the maximum possible packet length so that the probability a packet has error is at most ϵ?

6 [10 pts]

There are n possible frequencies and $k \leq n$ synchronous wireless stations. Each station selects at random one of these frequencies in order to talk. Given k and a frequency f what is the probability that there is exactly one station that talks at frequency f while other stations talk at frequency different from this one?

7 [10 pts]

A node keeps unsent TCP packets at a buffer whose capacity is C (in number of packets it can hold). Initially the buffer has C_0 packets, where $C_0 > 0$. Because of the current flow patterns at this node, the amount of packets at the buffer after departure of the old and arrival of the new packets increases at the rate of r, where $r > 0$ (arrival and departure happens at one time unit).

1. **[5 pts]** After how many time units (expressed as a function of C) does the *buffer overflow*?

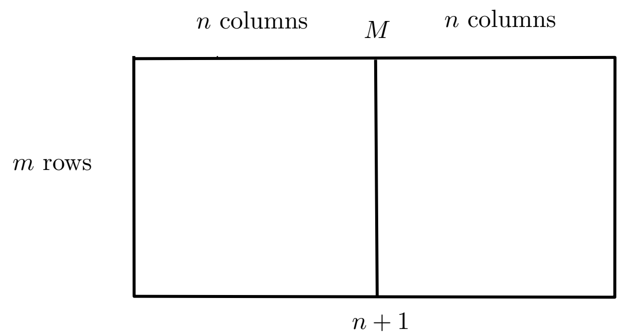
2. **[5 pts]** An early warning system sends a *potential buffer overflow* message when the buffer reaches $p\%$ of its capacity. After how many time units (expressed as a function of C, p) will a *potential buffer overflow* message be sent?

Make the correction!

8 [10 pts] ($\star$)

$m(2m+1)$

Consider an $m \times (2n+1)$ grid with m rows and $2n+1$ columns and a total of $m(2n+1)^2$ nodes. Label by (i, j) the node in the i-th row and j-th column of the grid. Also consider the middle column M of nodes $(1, n+1), \ldots, (m, n+1)$.



The nodes establish simultaneous communications sessions as follows. Each of the m nodes at leftmost column creates an arbitrary dedicated path traversing the grid to each node in the rightmost column.

1. **[2 pts]** Show that in total at least $\Omega(m^2)$ paths have to pass through nodes of the middle column M.

2. **[4 pts]** Show that there is node in the column M so that at least $\Omega(m)$ paths have to pass through it.[1]

3. **[4 pts]** Now assume that all but ℓ nodes between the nodes of the vertical column M are faulty and cannot route packets; so all packets have to be routed through

[1] The symbol $\Omega(f(n))$ means "at least $cf(n)$", where $c > 0$ is a constant independent of n.

non-faulty nodes in M; the senders know of these non-faulty nodes and can forward packets so as to avoid faulty nodes. Show that there is node in the column M so that at least $\Omega(m^2/\ell)$ paths have to pass through it.?

9 [10 pts]

Application of each new protocol adds a header of length h bits to a packet.

1. [**5 pts**] If a packet is undergoing applications of n protocols what percenteage of the resulting packet length is occupied by protocol headers. (Assume each protocol header has length h bits.)
2. [**5 pts**] After how many protocol applications is the length of the resulting packet at least triple the length of the original packet?