


Section 3.3 Linked Lists

1. Overview
2. Basic linked lists
3. Advanced linked lists
4. Insertion
5. Deletion

3.3.1 Overview

- ◆ Typical application processing in "real world":
 - read from a data source
 - file, database, user, etc.
 - store data in memory ←
 - iterate through data (maybe many times) and process it
 - store results to data sink
 - file, database, user, etc.

Overview (cont.)

- ◆ How we store data in memory is important! 
- ◆ We want
 - fastest possible access
 - least amount of memory
- ◆ Choice of data structure has major impact on performance

Overview (cont.)

- ◆ Option #1: array
 - advantages
 - elements are contiguous
 - faster access
 - disadvantages
 - once allocated, array cannot be resized
 - no growing, no shrinking
 - trade-offs
 - oversized array == waste of memory
 - undersized array == array overflow

Overview (cont.)

- ◆ Option #2: linked list
 - advantages
 - can be resized anytime
 - elements can be inserted, removed, shifted anywhere in the list
 - disadvantages
 - elements are not contiguous
 - slower access

3.3.2 Basic Linked Lists

- ◆ Singly linked list consists of:
 - a pointer to the first node in the list
 - head
 - a set of nodes, each consisting of:
 - data
 - pointer to next node

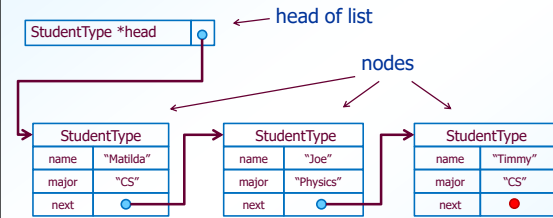
Basic Linked Lists (cont.)

- ◆ Doubly linked list consists of:
 - a pointer to the first node in the list
 - head
 - a pointer to the last node in the list
 - tail
 - a set of nodes, each consisting of:
 - data
 - pointer to next node
 - pointer to previous node

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Basic Linked Lists (cont.)



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Processing a Linked List

- ◆ Initialization
 - always initialize your pointers
 - use NULL or zero for empty pointers
 - check for NULL pointers in your code!
 - NULL is used as a sentinel
- ◆ Traversal
 - use an iteration pointer
- ◆ Do not lose the head of the list! 

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3.3.3 Advanced Linked Lists

- ◆ Problems with basic linked lists
 - element data mixed with list data
 - no encapsulation!
 - bad design
 - each element is hard-coded to point to a specific other
 - what if we need the same element in multiple lists?



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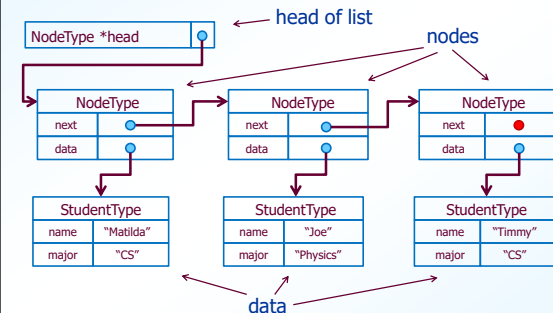
Advanced Linked Lists (cont.)

- ◆ Solution: separate the nodes from the data
- ◆ Why?
 - think "real world"
 - encapsulation
 - keep data-related stuff together, and list-related stuff together
 - compartmentalize what each element knows
 - ◆ should not know that it's in a linked list
 - ◆ should only have information related to itself
 - reuse
 - one element may be included in multiple linked lists

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Advanced Linked Lists (cont.)



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3.3.4 Linked List Insertion

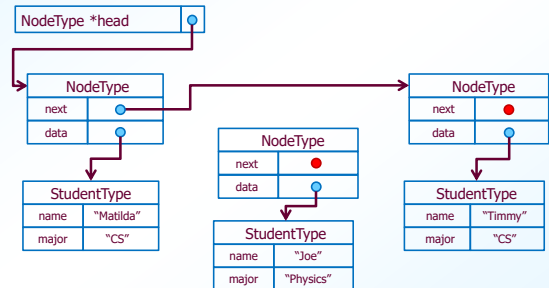
- ◆ We can insert an element anywhere in the list
 - shift pointer values
- ◆ Always consider four cases:
 - element is the first to be added
 - element is to be added in first position
 - element is to be added in middle of the list
 - element is to be added in last position

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Linked List Insertion (cont.)

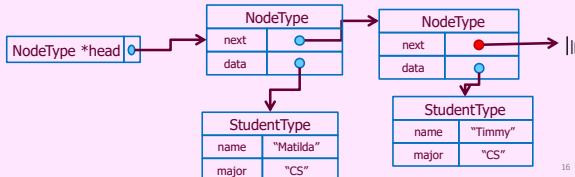
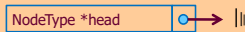
- ◆ Original list:



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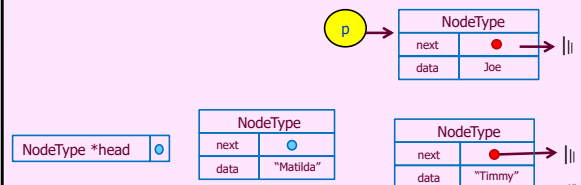
Case 1: insert as first element



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Case 1: insert as first element

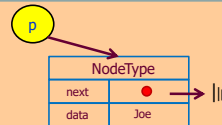
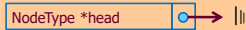
Step 1: Allocate memory for the node
Initialize it (data and next)



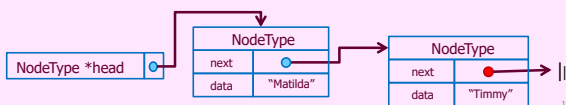
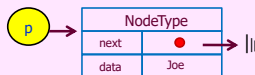
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Case 1: insert as first element

Step 1: Allocate memory for the node
Initialize it (data and next)



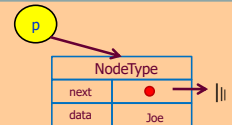
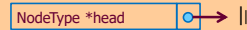
Step 1: Allocate memory for the node
Initialize it (data and next)



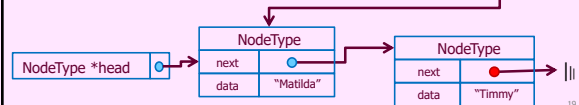
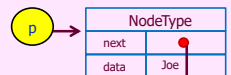
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Case 1: insert as first element

Step 2: p->next = head



Step 2: p->next = head



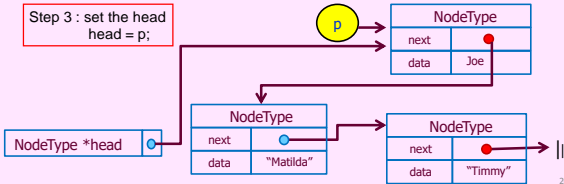
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Case 1: insert as first element

Step 3 : set the head
head = p;



Step 3 : set the head
head = p;



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Insert First

- Note:
 - Always maintain a handle to the allocated memory
 - Always maintain a handle to the linked list

```
int insertFirst(NodeType **head, DataType data)
```

```
{
```

```
    // allocate memory
```

```
    // set the data
```

```
    // make new node point to first node of list
```

```
    // update the head
```

```
    return(0);
```

```
}
```

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Insert First

- Note:
 - Always maintain a handle to the allocated memory
 - Always maintain a handle to the linked list

```
int insertFirst(NodeType **head, DataType data)
```

```
{
```

```
    NodeType *p = NULL;
```

```
    // allocate memory
    p = (NodeType *) malloc(sizeof(NodeType));
    if (p == NULL) return(1);
```

```
    // set the data
    p->data = data;
```

```
    // make new node point to first node of list
    p->next = *head;
```

```
    // update the head
    *head = p;
```

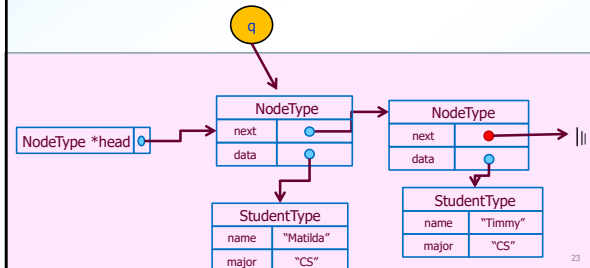
```
    return(0);
```

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Case 1: Case 2: Insert in the middle or end of LL

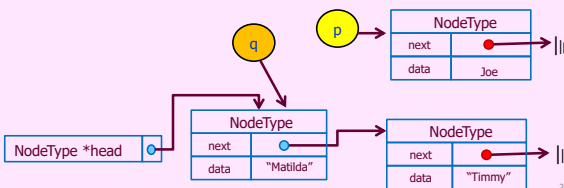
The operation is insert after an existing node, e.g., q
(Must have a handle to node q)



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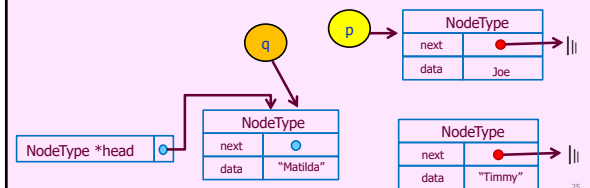
Case 2: insert in the middle or end of LL

Step 1: Allocate memory for the node
Initialize it (data and next)



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Case 2: insert in the middle or end of LL



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Insert After

- Note:
 - Always maintain a handle to the allocated memory
 - Always maintain a handle to the linked list

```
int insertAfter(NodeType *q, DataType data)
{
    // allocate memory

    // set the data

    // make new node point to node after q

    // make node of q point to new node

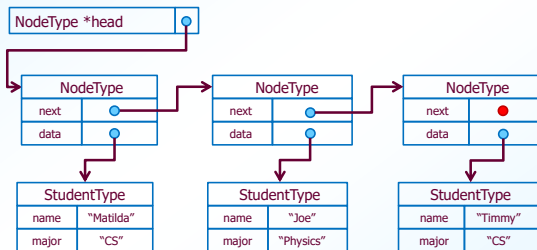
    return(0);
}
```

3.3.5 Linked List Deletion

- We can remove an element from anywhere in the list
 - shift pointer values
 - deallocate memory
 - node or data or both?
- Always consider five cases:
 - list is empty
 - element to be removed is the only element in the list
 - element is to be removed from the first position
 - element is to be removed from the middle of the list
 - element is to be removed from the last position

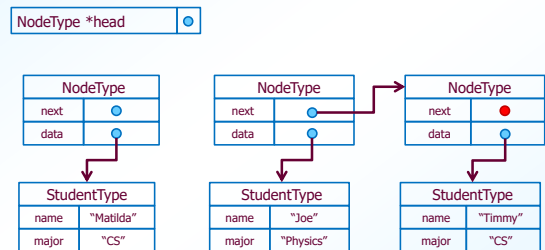
Linked List Deletion Delete First

- Original list:



Linked List Deletion Delete First

- Original list:



Linked List Deletion Delete First

- Note:
 - Always maintain a handle to the delete node
 - Always connect the remaining list to the head

```
int deleteFirst(NodeType **head, DataType *data)
{
    // keep a handle, p, to the node to be deleted

    // Update the head: set the head to the node after p

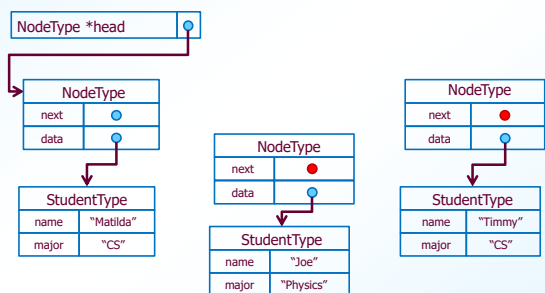
    // copy the data to the output

    // free the memory of p

    return(0);
}
```

Linked List Deletion after

- After deletion from the middle:



Linked List Deletion Delete After

- ◆ Note:
 - Always maintain a handle to the delete node
 - Always connect the remaining list to the head

```
int deleteFirst(NodeType *q, DataType *data)
{
    // keep a handle, p, to the node to be deleted

    // Update node q: set the q to point to the node after p

    // copy the data to the output

    // free the memory of p

    return(0);
}
```

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Linked List Cleanup

- ◆ Don't forget to explicitly deallocate your memory!
- ◆ Nodes
 - **always** deallocate the nodes when deallocating the list
- ◆ Data
 - only deallocate the data that will not be used again
 - **do not** deallocate data used elsewhere in the program



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Linked List Traversal

- ◆ Iterative
 - Start from head
 - Process the node
 - If list was not exhausted then move to next node
- ◆ Recursive
 - Check boundary condition
 - If boundary condition is not met then
 - Process node
 - Call yourself recursively with next node

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Doubly Linked List

- ◆ Node Characteristics
 - Next pointer
 - Previous pointer
 - Data
- ◆ List access characteristics
 - Head
 - Tail
- ◆ Processing
 - Can traverse the list in both directions!!
- ◆ Expense
 - Additional pointer

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Summary

- ◆ Linked list operation
 - Insertion
 - Deletion
 - Traversal
- ◆ When to use
 - When data is sparse
 - When data is dynamic (modified often by insertion and deletions)

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