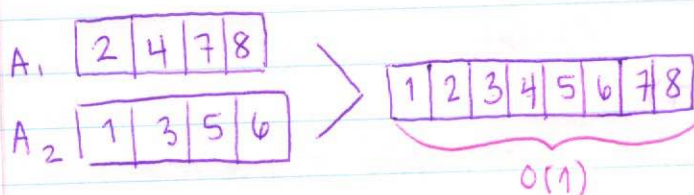


Nov 24
COMP 2402

Sorting Algorithms

→ MergeSort

merge ($A[0 \dots n/2]$, $A[n/2+1 \dots n+1]$)



MergeSort ($A[0 \dots n-1]$)

if ($n \leq 1$) return A

mergeSort ($A[0 \dots n/2]$)

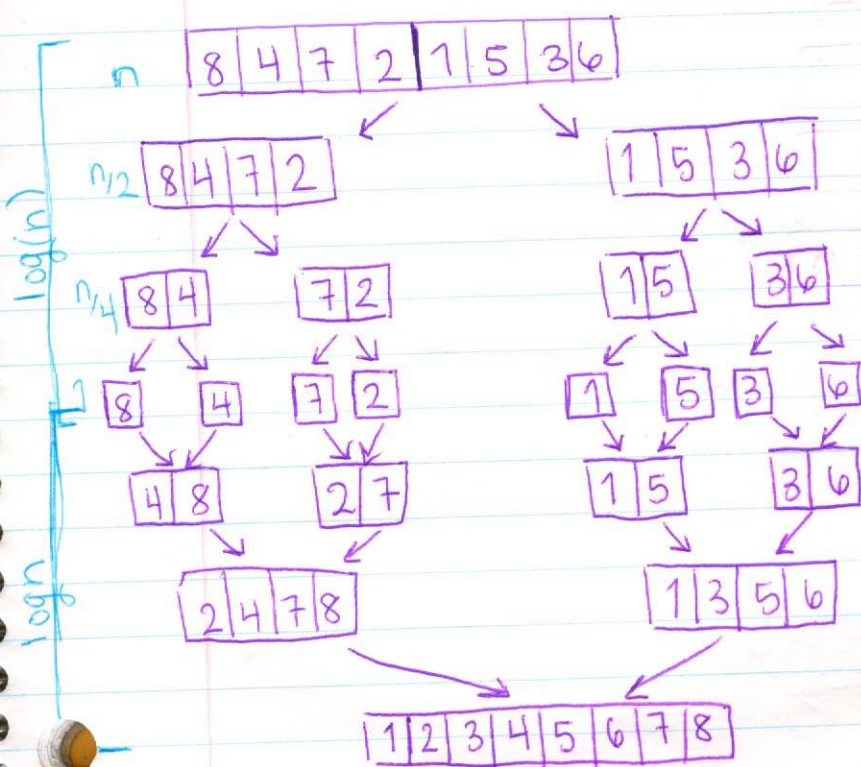
mergeSort ($A[n/2 \dots n]$)

merge ($A[0 \dots n/2]$, $A[n/2 \dots n]$) // $O(n)$

$T(n)$

$T(n/2)$

$T(n/2)$



$$T(n/2) = O(n/2) + 2T(n/4)$$

split

n sorted array

↳ 0 comparisons

$$T(n) = O(n) + 2T(n/2)$$

$$= O(n) + 2[O(n/2) + 2T(n/4)]$$

$$= O(n) + O(n) + 4T(n/4)$$

$$= O(n) + O(n) + 4[O(n/4) + 2T(n/8)]$$

$$= O(n) + O(n) + O(n) + 8T(n/8)$$

⋮

$$= K \cdot O(n) + 2^K T(n/2^K)$$

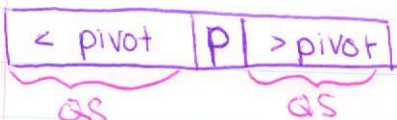
$$= O(n \log n) + O(n)$$

recurse until...

$$n/2^K = 1 \rightarrow \log(n) = K$$

Hilary

→ QuickSort



QuickSort($A[0 \dots n]$)

pivot = $A[0]$

$p = \text{partition}(A[0 \dots n], \text{pivot})$

QuickSort($A[0 \dots p]$)

QuickSort($A[p+1 \dots n]$)

partition($A[0 \dots n], \text{pivot}$)

p
5 4 9 7 3 1 6 8 2

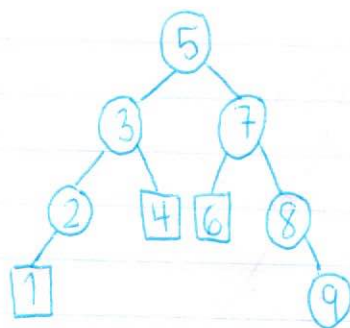
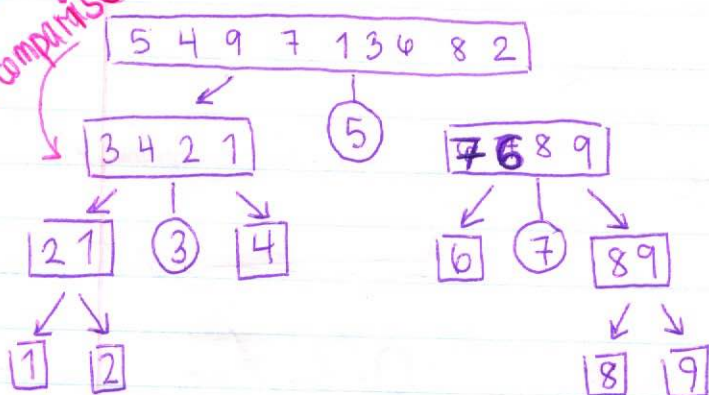
5 4 2 7 3 1 6 8 9

5 4 2 1 3 7 6 8 9

3 4 2 1 5 7 6 8 9

$A[0 \dots p] < A[p] < A[p+1 \dots n]$

n comparisons



$$\Pr \{i \text{ path } x\} = 1/(i-x+1)$$

$$E[\text{len}(\text{path}(x))] = H_{x+1} + H_{n-x} - 2$$

$$E[\# \text{ compare to pivot}] = H_{i+1} + H_{n-i}$$

$$\begin{aligned} E[\# \text{ of comparisons}] &= \sum_{i=0}^{n-1} H_{i+1} + H_{n-i} - 2 \\ &= H_1 + H_2 + \dots + H_n + H_1 + H_2 + \dots + H_n \\ &= 2 \sum_{i=1}^n H_i \\ &\leq 2 \sum_{i=1}^n H_n \\ &\leq 2n H_n \\ &\leq 2n (\ln(n) + 1) \\ &= 2n \ln(n) + O(n) \end{aligned}$$