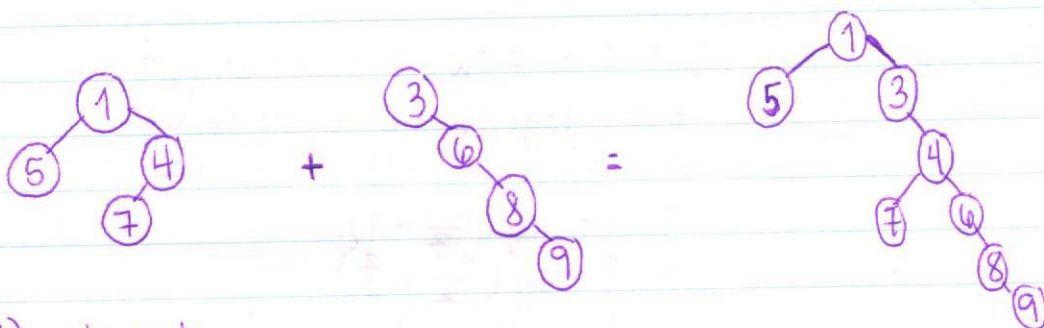


NOV 22
COMP 2402

Meldable Heaps

→ Randomized heap

$\text{meld}(h_1, h_2)$



if ($h_1 == \text{null}$) return h_2

if ($h_2 == \text{null}$) return h_1

if ($h_1.x > h_2.x$) ~~return~~ $\text{meld}(h_2, h_1)$

flip a coin

if (heads)

$h_1.\text{right} = \text{meld}(h_1.\text{right}, h_2)$

else

$h_1.\text{left} = \text{meld}(h_1.\text{left}, h_2)$

- Random walk : a path through a binary tree.

→ The expected length of a random walk in a binary tree of size n ($\mathbb{E}[W_n]$) is $\leq \log(n+1)$

Hilroy

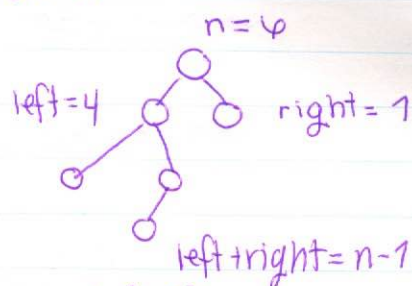
* will clean up later (correction)

Proof: $E[W_n] \leq \log(n+1)$, by induction on n

Base case: $n=1$

$$W_1 = 1$$

$$\log(1+1) = \log_2(2) = 1 \quad \checkmark$$



Assume for all $n' < n$, that $E[W_{n'}] \leq \log(n'+1)$

maximized
left=right = $\frac{n-1}{2}$

$$\begin{aligned} E[W_n] &= 1 + E[W_{\text{left}}] + E[W_{\text{right}}] \\ &\leq 1 + \log(\text{left}+1) + \log(\text{right}+1) \quad // \text{ by induction} \\ &\leq 1 + \frac{1}{2} \log\left(\frac{n-1}{2}+1\right) + \frac{1}{2} \log\left(\frac{n-1}{2}+1\right) * \\ &\leq 1 + \log\left(\frac{n-1}{2}+1\right) \\ &= 1 + \log\left(\frac{n-1}{2} + \frac{2}{2}\right) \\ &= 1 + \log\left(\frac{n+1}{2}\right) \\ &= 1 + \log(n+1) - \log_2(2) \\ &= 1 + \log(n+1) - 1 \\ &= \log(n+1) \end{aligned}$$

$$\begin{aligned} \text{meld}(h_1, h_2) &= E[W_{n_1}] + E[W_{n_2}] \\ &\leq \log(n_1+1) + \log(n_2+1) \\ &\leq 2 \log(n+1) \end{aligned}$$