

Nov 29
COMP 2402 1

Sorting Algorithms

MergeSort

- comparisons: $n \log(n)$ worst case
- in-place? no
- stable? yes

QuickSort

- comparisons: $1.38n \log(n) + O(n)$ expected
- in-place? yes
- stable? no

HeapSort

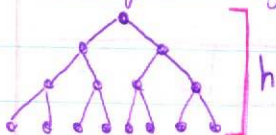
- comparisons: $2n \log(n) + O(n)$ worst case
- in-place? yes
- stable? no

A:

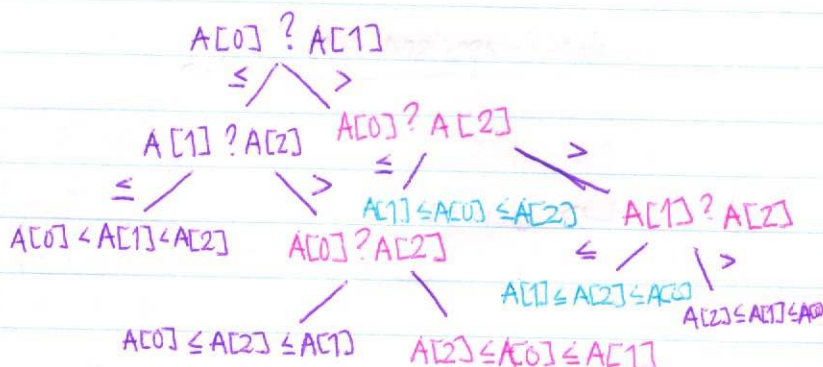
x	x	x
0	1	2

$n!$ permutations of A
 $n!$ leaves

↳ height $\geq \log(n!)$



m leaves $\leftrightarrow \log(m) \leq h$



$\log(n!) \text{ is } \Omega(n \log n)$

$$g = \Omega(f)$$

upper: $\log(n!) = \log(n \cdot (n-1) \cdot (n-2) \dots)$
 $= \log(n) \cdot \log(n-1) \cdot \log(n-2) \dots \log(1)$
 $\leq \log n + \log n + \log n + \dots \log(n)$
 $= n \log n$

lower bound: $\log(n!) = \log(1) + \log(2) + \dots + \log(n)$
 $\geq \log(\frac{n}{2} + 1) + \log(\frac{n}{2} + 1) + \dots \log(n)$ *cut out: $\log(1) \dots \log(\frac{n}{2} - 1)$*
 $\geq \log(\frac{n}{2}) + \log(\frac{n}{2}) + \dots \log(\frac{n}{2})$ *lowered everything to smallest term*
 $= \frac{n}{2} \cdot \log(\frac{n}{2})$

side bar:

$\log(n) \geq 2$ for $n \geq 4$

$\frac{1}{4} \log(n) \geq \frac{1}{2}$

$\frac{n}{4} \log(n) \geq \frac{n}{2}$

$\frac{n}{4} \log(n) - \frac{n}{2} \geq 0$

$\frac{n}{2} \log(n) - \frac{n}{2} \geq \frac{n}{4} \log n$

$= \frac{n}{2} \cdot \log(\frac{n}{2})$
 $= \frac{n}{2} (\log(n) - \log 2)$
 $= \frac{n}{2} (\log(n) - 1)$
 $= \frac{n}{2} \log(n) - \frac{n}{2}$
 $= \frac{1}{4} n \log n$

$\frac{1}{4} n \log n \leq \log(n!) \leq n \log n$

Non-Comparison-Based String

Counting Sort
 $\hookrightarrow O(n+K)$

a:

0	3	1	3	2	2	1	3
---	---	---	---	---	---	---	---

 $K=4$
0 1 2 3 4 5 6 7

c:

1	2	2	3
---	---	---	---

 // number of ^{times} each number appears
0 1 2 3

for (i=0; i<n; i++)
 $c[a[i]]++$

a:

0	1	1	2	2	3	3	3
---	---	---	---	---	---	---	---

Stable: c'

1	3	5	8
---	---	---	---

b:

0	1	1	2	2	3	3	3
---	---	---	---	---	---	---	---

0 1 2 3 4 5 6 7

for (i=1; i<K; i++)

$c'[i] += c[i-1]$

for (i=n-1; i>=0; i--)

$b[--c'[a[i]]] = a[i]$

$a = b$

RadixSort

668	0010	1001	1100
3018	1011	1100	1010
1409	0101	1000	0001
3233	1100	1010	0001
714	0010	1100	1010

countingsort

number sections
we have to
deal with

$$K = \frac{w}{d}$$

$$d=4 \rightarrow K = \frac{16}{4} = 4$$

CountingSort
0...K

Radix
0...2^K