

2110313 Operating Systems Multithreaded Programming

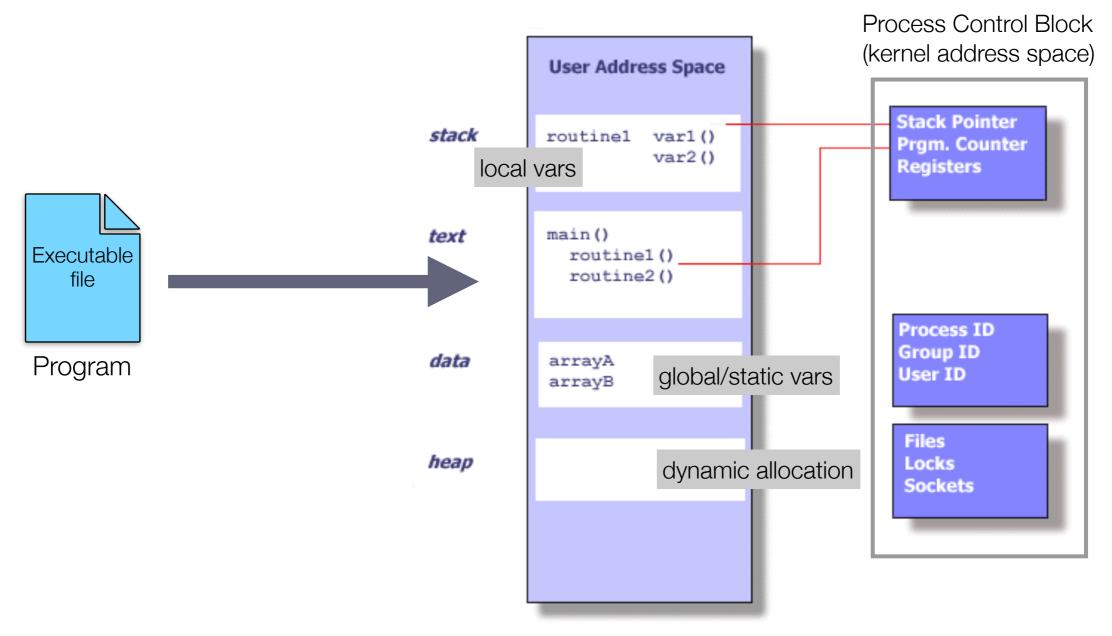
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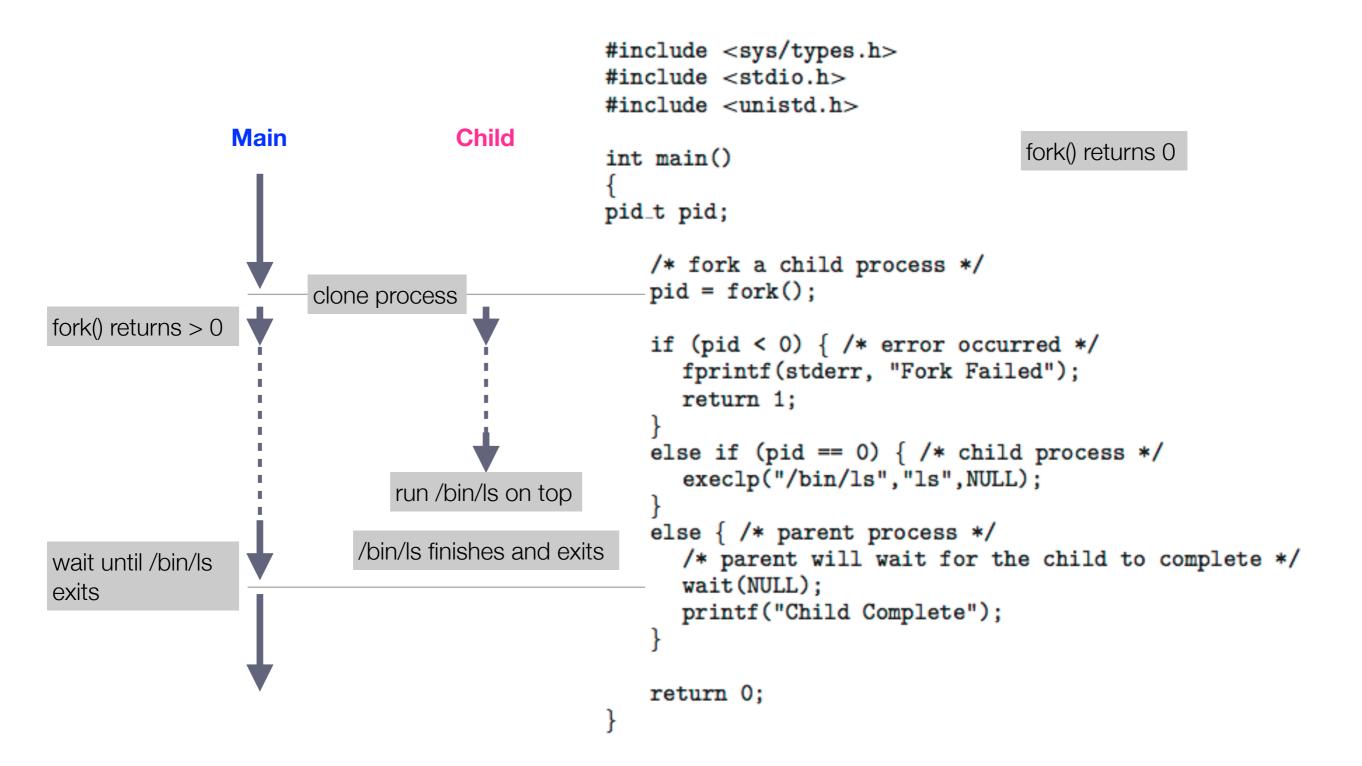


Process

Process = Program in Execution



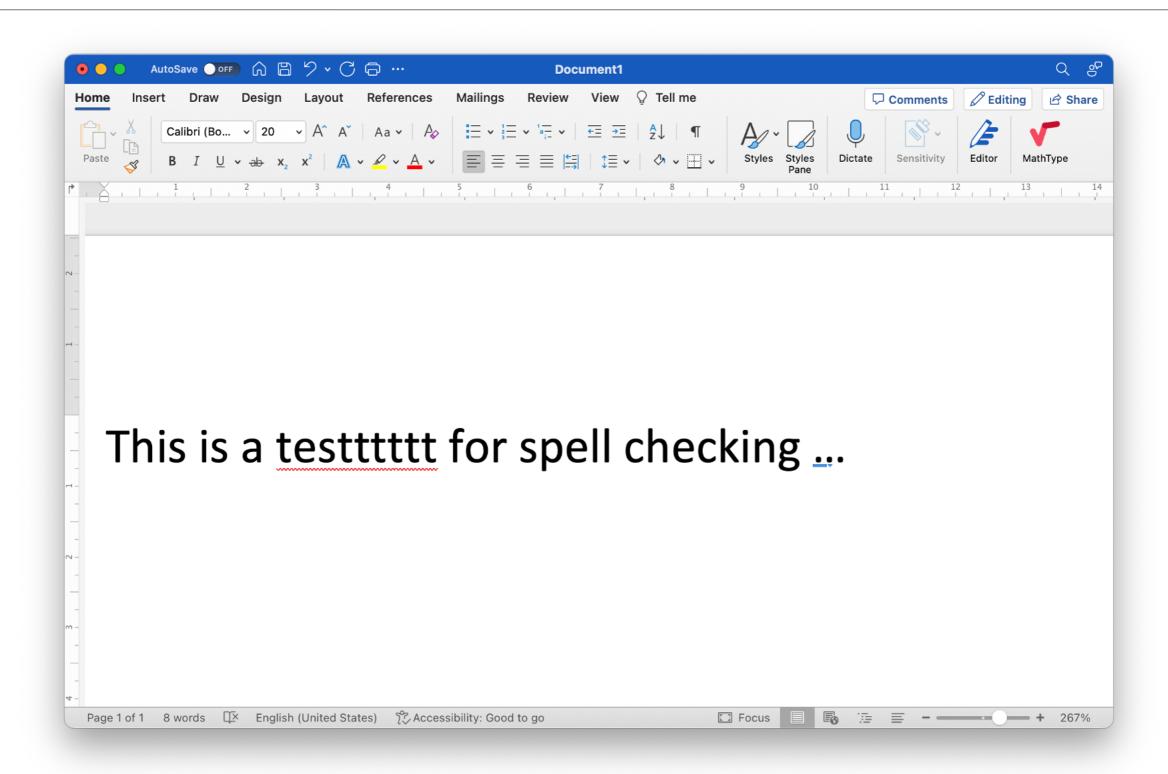
Source: https://computing.llnl.gov/tutorials/pthreads/



Introduction to Thread

- Most modern applications are multithreaded
- Threads run within application
- Multiple tasks with the application can be implemented by separate threads
 - Update display
 - Fetch data
 - Spell checking
 - Answer a network request

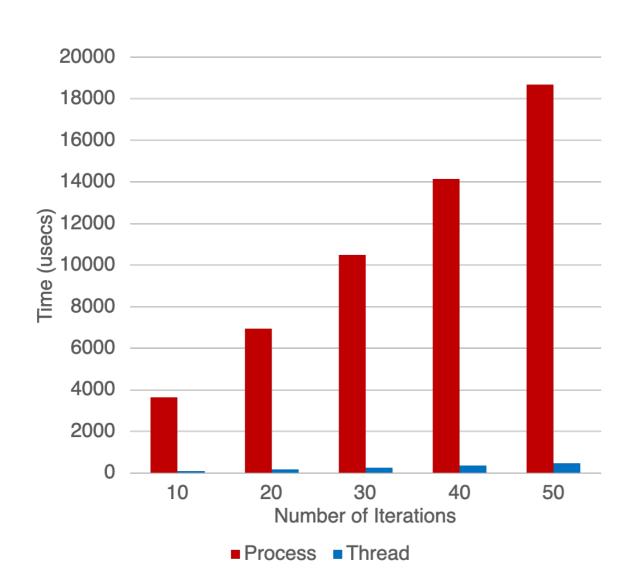
Microsoft Word - Spell Checking



Introduction to Thread

- Thread = Light-Weight Process
- Process creation is heavy-weight while thread creation is light-weight
- Can simplify code, increase efficiency
- Kernels are generally multithreaded

Process/Thread Creation Overheads



Iterations	Process	Thread
10	3633.80	87.80
20	6947.80	168.20
30	10501.00	269.80
40	14135.00	359.40
50	18693.00	470.60

Type as fast as you can, press q to quit The pressed key is

Count number of key pressed in a second

Keyhit Rating – No Thread

```
key hit count = 0
last second count = 0
t0 = current timestamp
While not stop:
      Check keyboard or timeout in 9,5 seconds
      If there is a key hit:
            increase key hit count by one
      t1 = current timestamp
      if t1 - t0 >= 1 second:
            key hit last second = key hit count - last second count
            rate = key hit last second / (t1 - t0)
            print rate
            last second count = key hit count
            t0 = t1
```

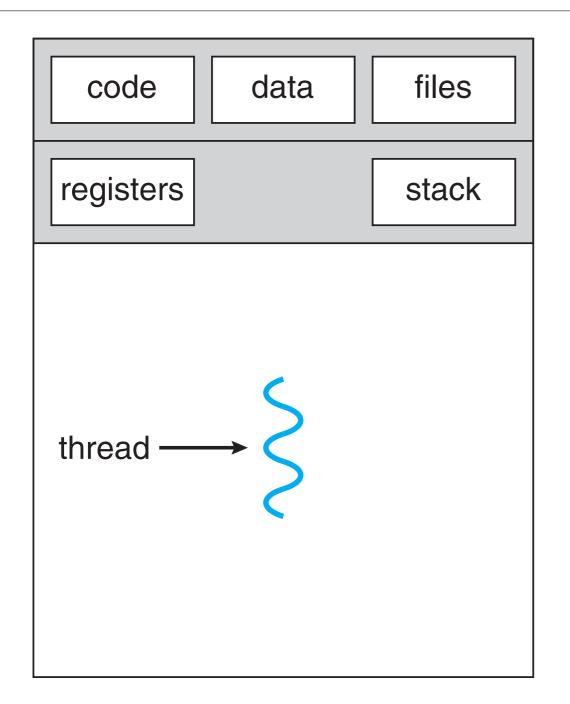
Keyhit Rating – Two Threads

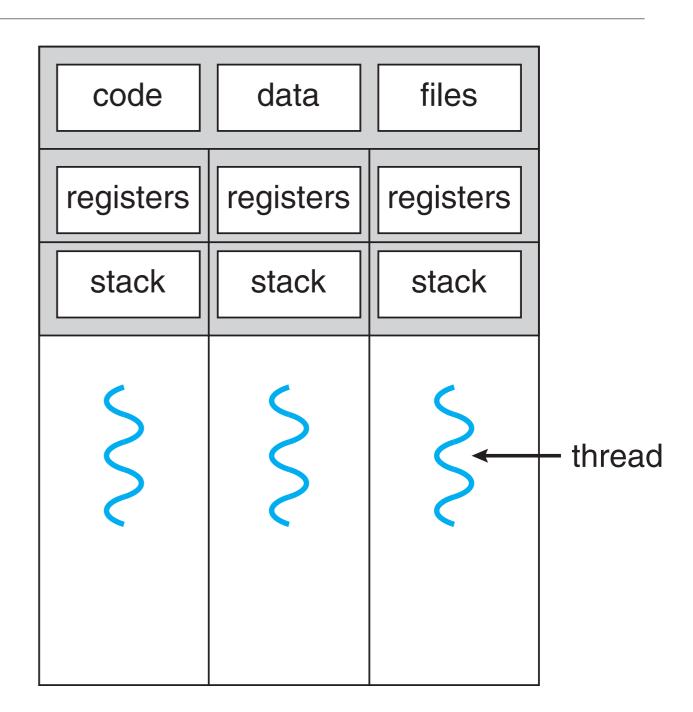
```
Key count thread:
      While not stop:
            Wait for keyhit
            increase key hit count by one
Rating thread:
      While not stop:
            sleep for 1 second
            key hit last second = key hit count - last second count
            rate = key hit last second / (t1 - t0)
            print rate
            last second count = key hit count
Main:
      key hit count = 0
      last second count = 0
      Start Rating thread
      Start Key count thread
```

Inside Thread

- Each thread has its program counter and stack (local variables)
- All threads share code, data (global / static variables), heap (dynamic allocation)
- Each thread executes at its own pace

Single-Threaded vs. Multithreaded Processes

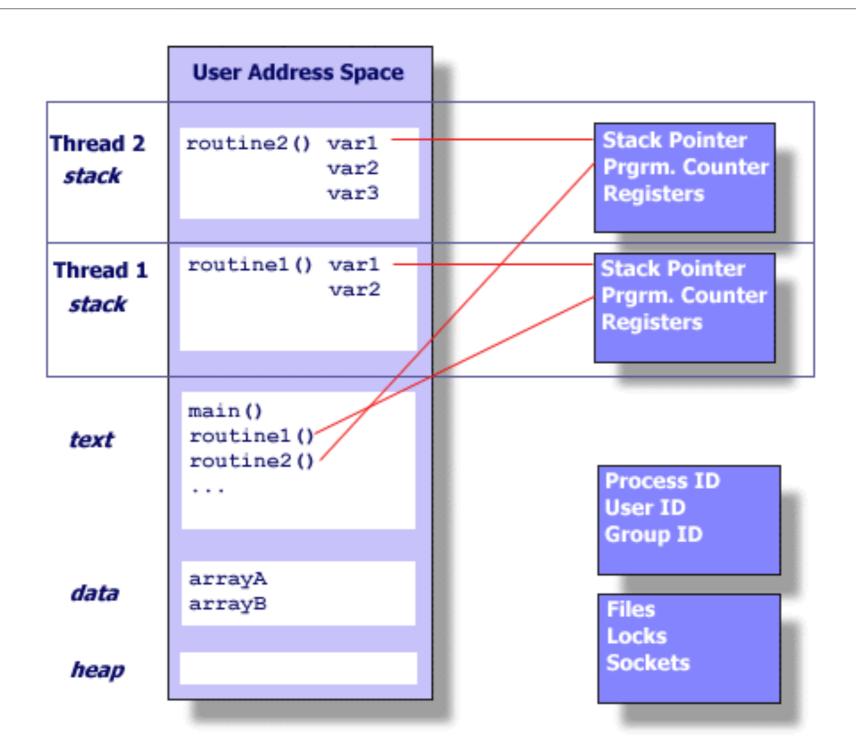




single-threaded process

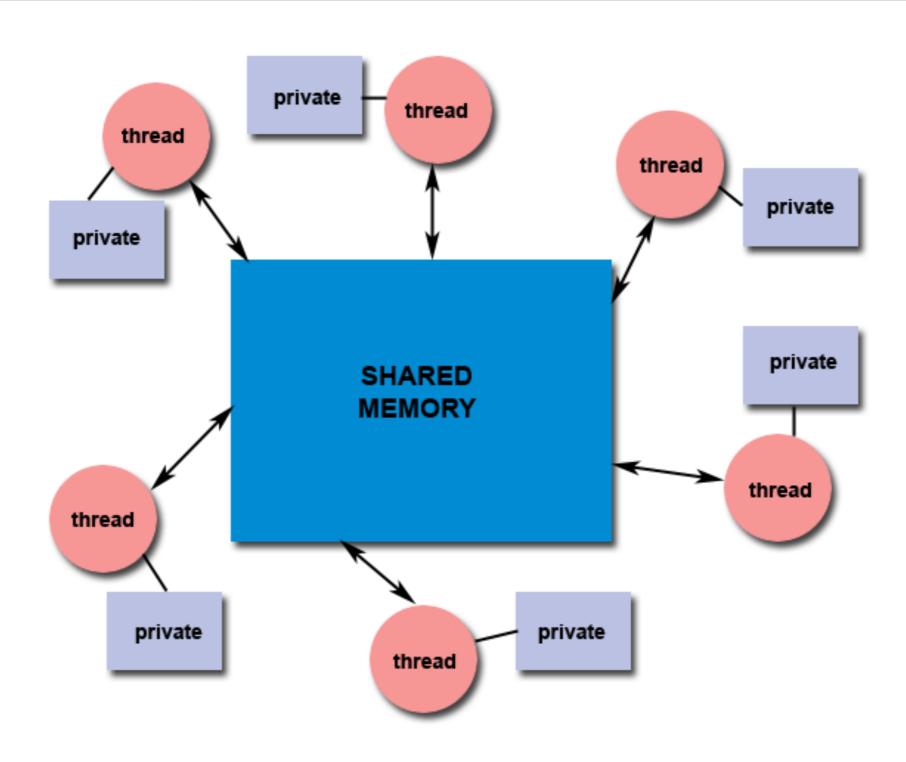
multithreaded process

Threads and Process Memory Spaces



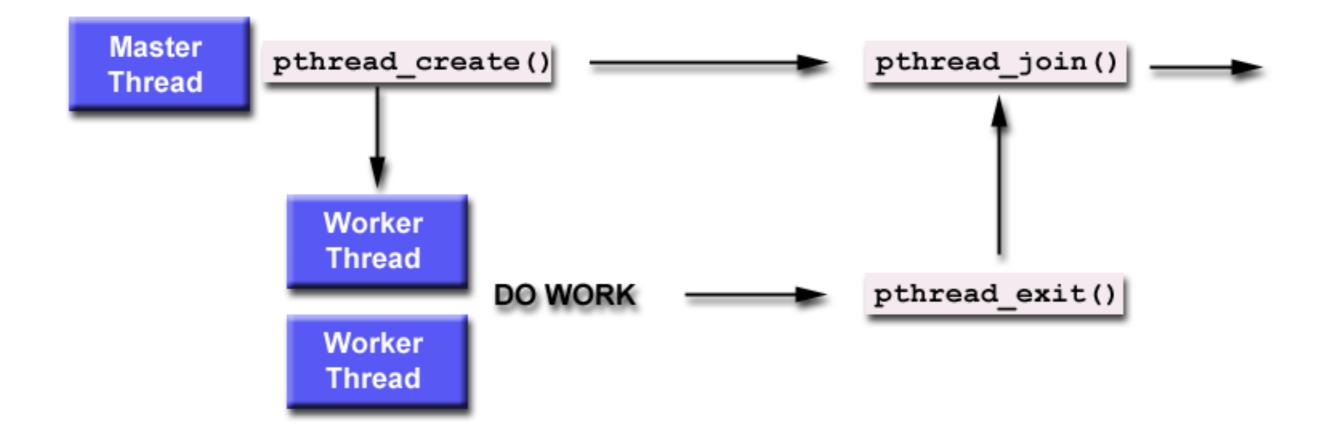
Source: https://computing.llnl.gov/tutorials/pthreads/

Shared Memory Model



Source: https://computing.llnl.gov/tutorials/pthreads/

Thread Creation / Termination

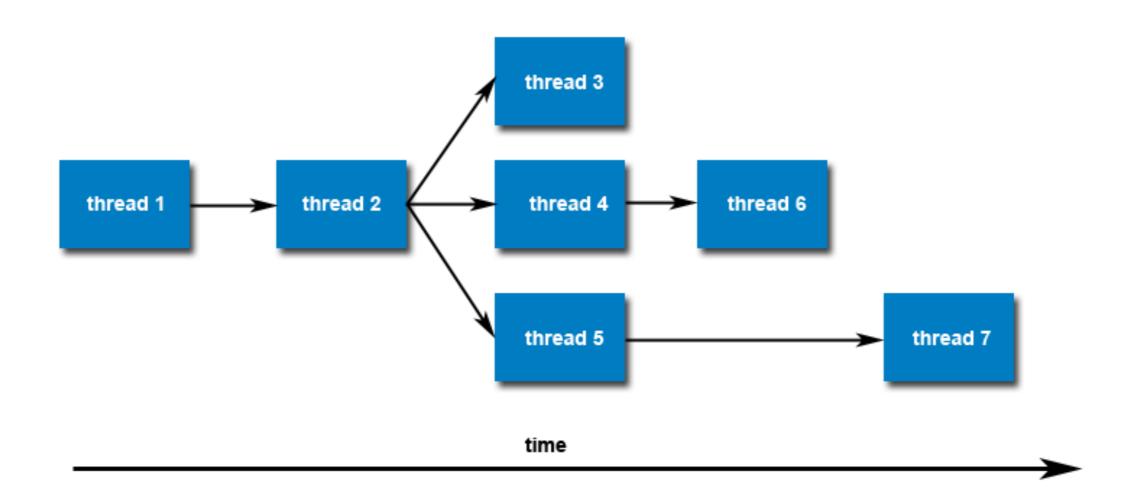


```
1 | #include <pthread.h>
   #include <stdio.h>
    #include <stdlib.h>
    #include <math.h>
    #define NUM_THREADS
                            4
 6
 7
    void *BusyWork(void *t)
 8
    {
 9
       int i;
       long tid;
10
       double result=0.0;
11
       tid = (long)t;
12
13
       printf("Thread %ld starting...\n",tid);
       for (i=0; i<1000000; i++)
14
15
          result = result + sin(i) * tan(i);
16
17
       printf("Thread %ld done. Result = %e\n",tid, result);
18
       pthread_exit((void*) t);
19
    }
20
```

```
int main (int argc, char *argv[])
22
23
    {
24
       pthread t thread[NUM THREADS];
       pthread attr t attr;
25
26
       int rc;
27
       long t;
28
       void *status;
29
30
       /* Initialize and set thread detached attribute */
31
       pthread attr init(&attr);
       pthread_attr_setdetachstate(&attr, PTHREAD_CREATE_JOINABLE);
32
33
34
       for(t=0; t<NUM THREADS; t++) {</pre>
35
          printf("Main: creating thread %ld\n", t);
          rc = pthread create(&thread[t], &attr, BusyWork, (void *)t);
36
37
          if (rc) {
38
             printf("ERROR; return code from pthread create() is %d\n", rc);
39
             exit(-1);
40
41
          }
42
       /* Free attribute and wait for the other threads */
43
44
       pthread attr destroy(&attr);
       for(t=0; t<NUM THREADS; t++) {</pre>
45
46
          rc = pthread join(thread[t], &status);
          if (rc) {
47
             printf("ERROR; return code from pthread join() is %d\n", rc);
48
49
             exit(-1);
50
          printf("Main: completed join with thread %ld having a status
51
                of %ld\n",t,(long)status);
52
53
          }
54
    printf("Main: program completed. Exiting.\n");
55
56
    pthread exit(NULL);
57
```

Once created, threads are peers, and may create other threads. There is no implied hierarchy or dependency between threads.

ญี่ไม่ชี่ คราเบ่นระกับ Thread

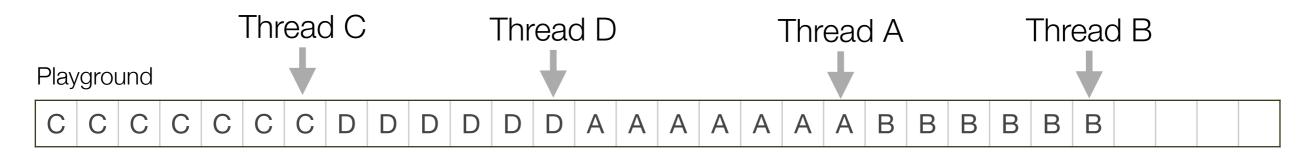


Source: https://computing.llnl.gov/tutorials/pthreads/

Thread Scheduling

- Order of starting execution of threads can be arbitrary
- CPU times are also allocated randomly similar to process scheduling
- Thus, we should write a program to NOT depend on the order of thread execution

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```
void *MyThread(void *arg)
20
    {
21
        struct timeval end;
22
23
        char my_call_sign = (char)*((char *)arg);
24
        sleep(1);
25
        for(int i=0 ; i < PLAYGROUND_SIZE ; i++) {</pre>
26
             for(int k=0; k < 10000; k++)
27
28
             playground[i] = my_call_sign;
29
30
        gettimeofday(&end, 0);
31
        double elapsed_thread = elapsed_time(&t_start, &end);
32
        printf("Thread %c ends at time %f\n", my_call_sign, elapsed_thread);
33
        pthread_exit(NULL);
34
35
```

```
(base) natawut@Natawuts-Mac-mini os-class % ./attack
main(): Created 4 threads.
Thread D ends at time 1.103172
Thread A ends at time 1.103241
Thread C ends at time 1.108047
Thread B ends at time 1.112472
[0] 218
[1] 3653
[2] 6294
[3] 75
```

(base) natawut@Natawuts-Mac-mini os-class %

Passing Arguments to Threads

```
long taskids[NUM_THREADS];

for(t=0; t<NUM_THREADS; t++)
{
   taskids[t] = t;
   printf("Creating thread %ld\n", t);
   rc = pthread_create(&threads[t], NULL, PrintHello, (void *) taskids[t]);
   ...
}</pre>
```

```
int rc;
long t;

for(t=0; t<NUM_THREADS; t++)
{
    printf("Creating thread %ld\n", t);
    rc = pthread_create(&threads[t], NULL, PrintHello, (void *) &t);
    ...
}</pre>
```

Source: https://computing.llnl.gov/tutorials/pthreads/

```
struct thread data{
   int thread id;
  int sum;
  char *message;
};
struct thread data thread data array[NUM THREADS];
void *PrintHello(void *threadarg)
   struct thread data *my data;
  my data = (struct thread data *) threadarg;
  taskid = my_data->thread id;
  sum = my data->sum;
  hello msg = my data->message;
   . . .
int main (int argc, char *argv[])
  thread_data_array[t].thread_id = t;
  thread_data_array[t].sum = sum;
  thread_data_array[t].message = messages[t];
  rc = pthread_create(&threads[t], NULL, PrintHello,
        (void *) &thread data array[t]);
                      pointer
```

Multicore Programming In thread 3111 845014

- Multicore or multiprocessor systems putting pressure on programmers, challenges include:
 - Dividing activities
 - Balance
 - Data splitting
 - Data dependency
 - Testing and debugging

Multicore Programming

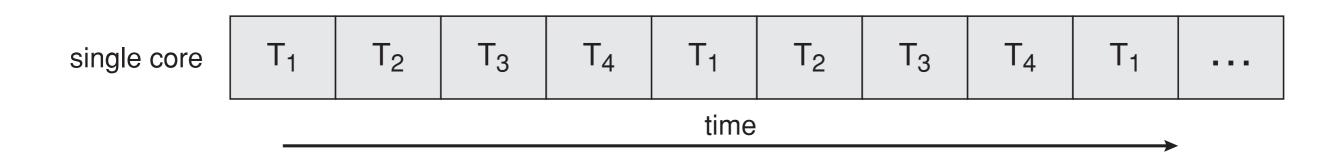
- Parallelism implies a system can perform more than one real multi processor core task simultaneously
- Concurrency supports more than one task making progress makully grann smultaneously
 - Single processor / core, scheduler providing concurrency
- Types of parallelism
- pes of parallelism

 Mulouz Mulatarinau (GPV)

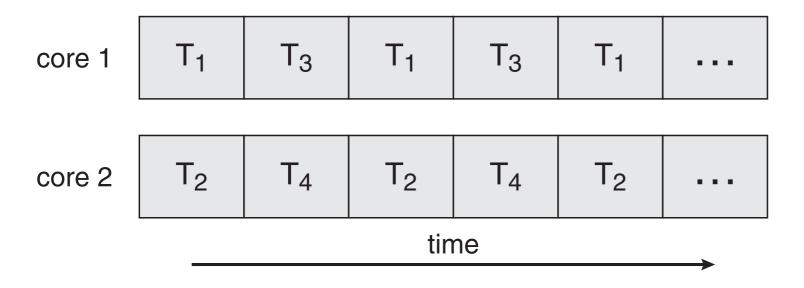
 Data parallelism distributes subsets of the same data across multiple cores, same operation on each
 - Task parallelism distributing threads across cores, each thread performing unique operation

Concurrency vs. Parallelism

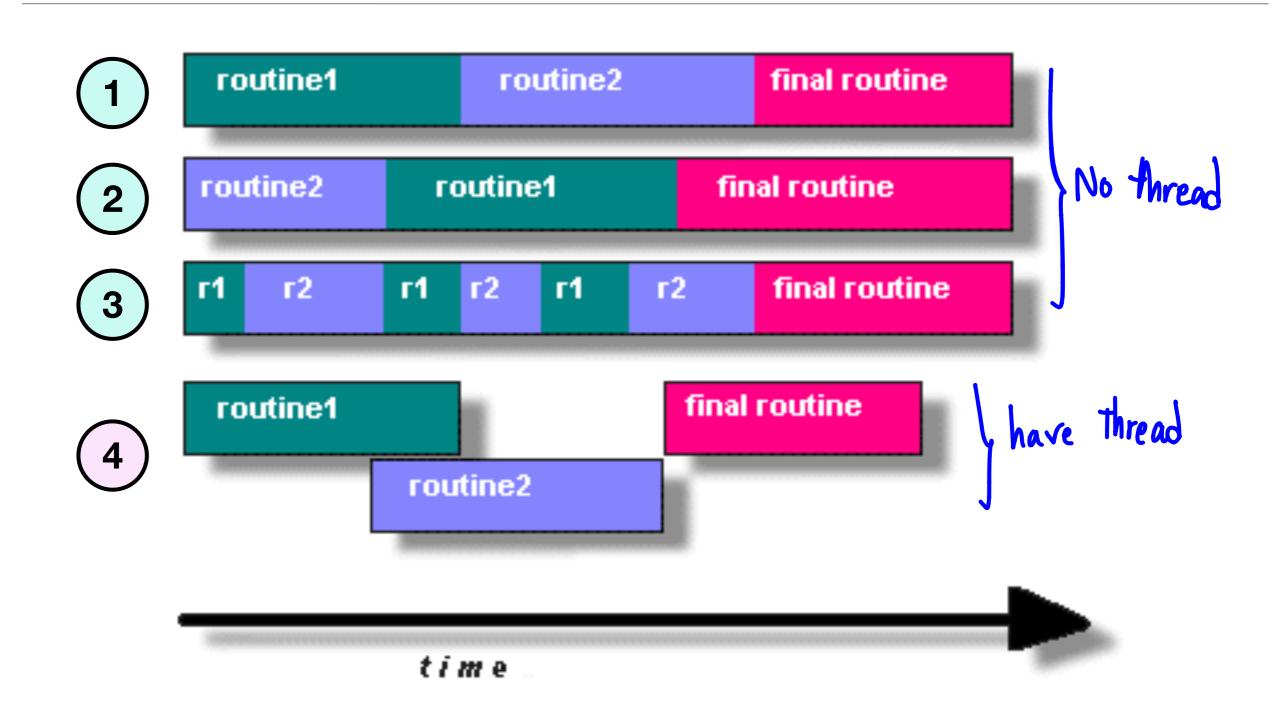
Concurrent execution on single-core system:



Parallelism on a multi-core system:

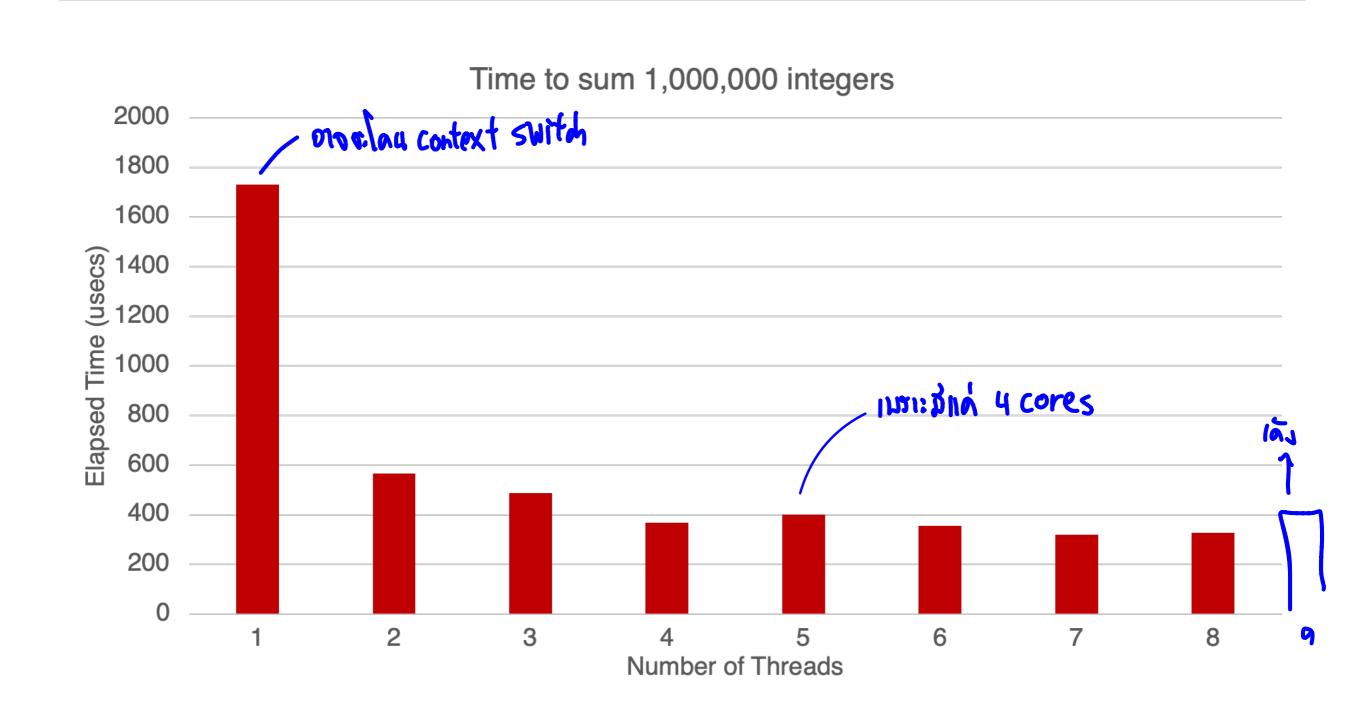


Execution Patterns of Interleave and Overlap



Source: https://computing.llnl.gov/tutorials/pthreads/

Multithreading Performance - Parallel Sum



Threading Issues

- User Thread vs. Kernel Thread
- Thread Libraries
- Thread Pools
- fork()/exec() and Thread
- Single and cancelation

User Thread vs. Kernel Thread

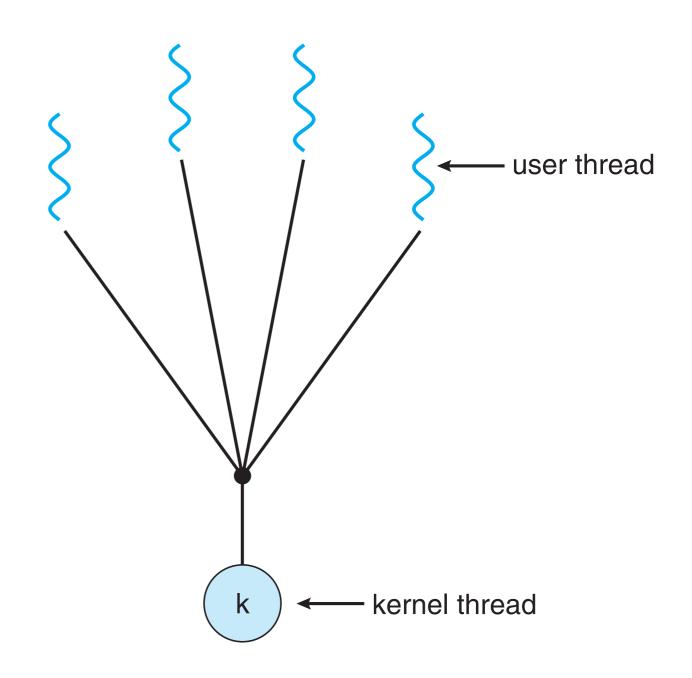
- User threads
 - management done by user-level threads library
- Example thread libraries:
 - POSIX Pthreads
 - Win32 threads
 - Java threads
 - Python threads

User Thread vs. Kernel Thread

- Kernel threads Supported by the Kernel
- Examples virtually all general purpose operating systems, including:
 - Windows
 - Linux
 - MacOS
- How are user threads mapped to kernel threads?

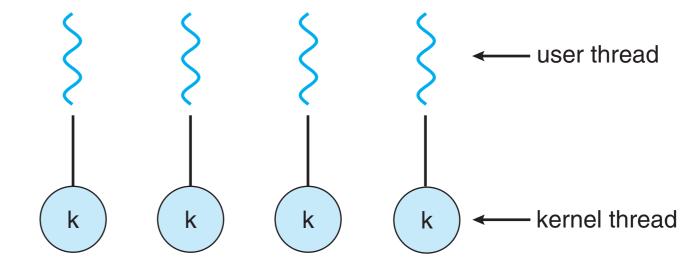
Many-to-One

- Many user-level threads
 mapped to single kernel thread
- One thread blocking causes all to block
- Multiple threads may not run in parallel on muticore system because only one may be in kernel at a time
- Examples:
 - Old OS (e.g. Solaris)
 - GNU Portable Threads



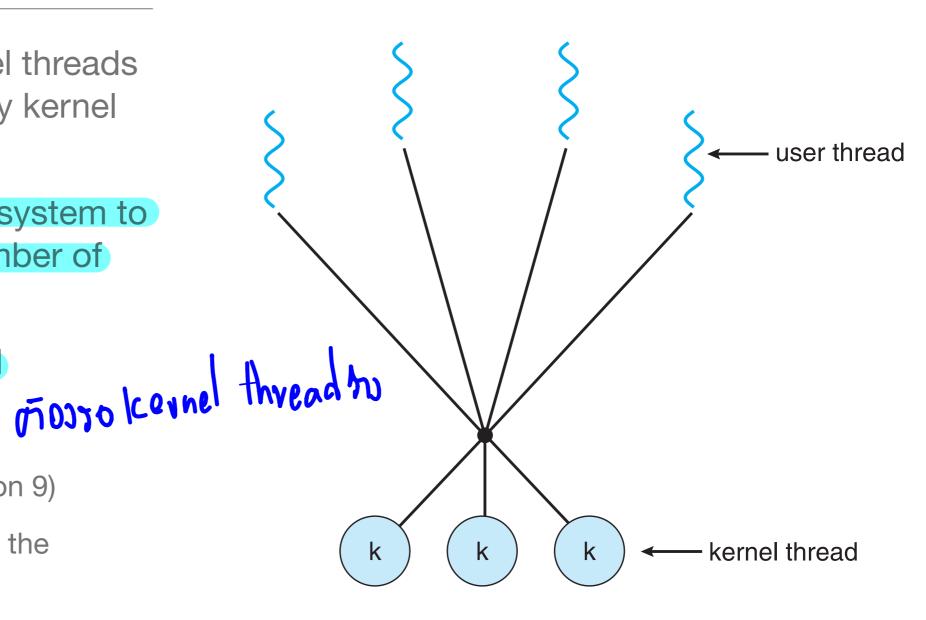
One-to-One

- Each user-level thread maps to kernel thread
- Creating a user-level thread creates a kernel thread
- More concurrency than manyto-one
- Number of threads per process sometimes restricted due to overhead
- Examples:
 - Windows NT/XP/2000
 - Linux



Many-to-Many

- Allows many user level threads to be mapped to many kernel threads
- Allows the operating system to create a sufficient number of kernel threads
- Similar to Thread Pool
- Examples:
 - Old OS (Solaris < version 9)
 - Windows NT/2000 with the ThreadFiber package
- Variation: Two-level model



Thread Pools

- Create a number of threads in a pool where they await work
- Advantages:
 - Usually slightly faster to service a request with an existing thread than create a new thread เรียกใช้ ใกล้า ผ่างใน
 - Allows the number of threads in the application(s) to be bound to
 the size of the pool Kregure in inquere
 - Separating task to be performed from mechanics of creating task allows different strategies for running task
 - i.e. Tasks could be scheduled to run periodically

fork()/exec() and Thread

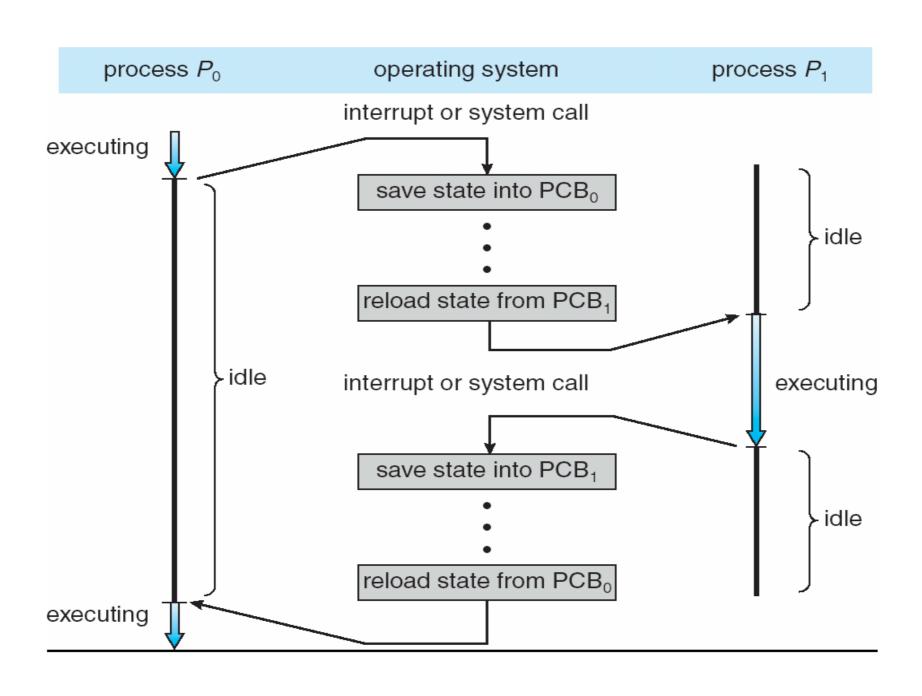
- Does fork() duplicate only the calling thread or all threads? No, wocess on Is clone thread
- Exec() usually works as normal replace the running process including all threads

Benefits of Using Threads

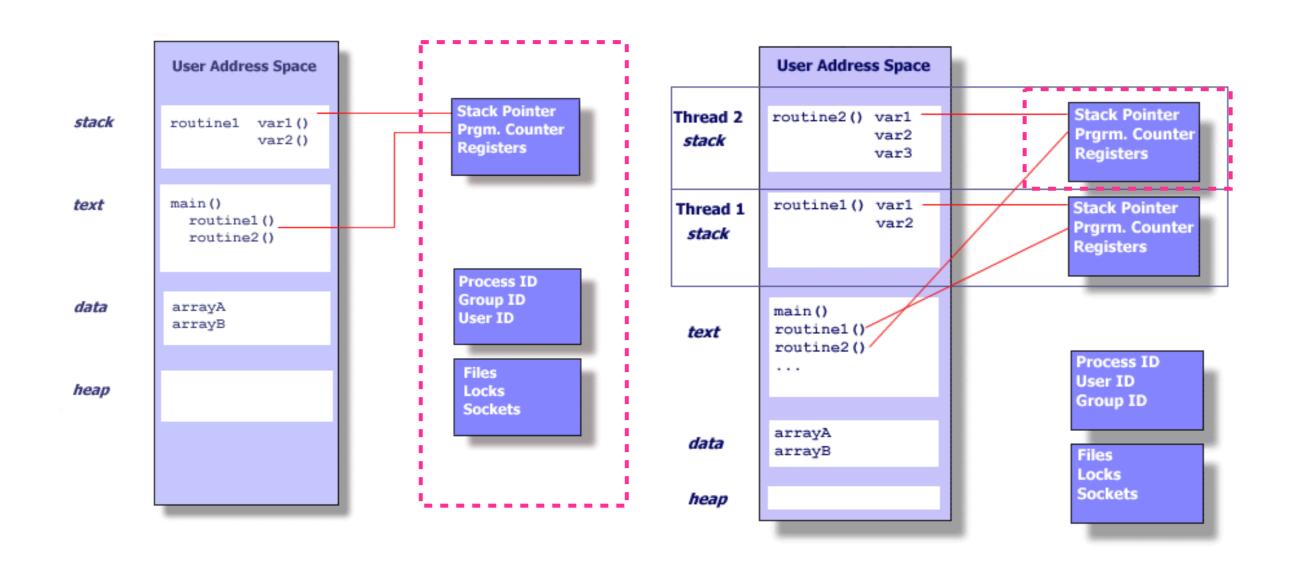
Light Weight

- Lesser overheads to create a new thread than a new process
- More efficient inter-thread vs. inter-process communication
- Lesser overheads for context switching (no need to deal with memory)

Context Switching Overheads



Context Switching vs. Thread Switching



Benefits of Using Threads

- Effective Usages of Multi-Core/Multi-Thread CPU with Parallel Programming
 - Utilize more CPU cores / threads
 - Overlapping of executions (CPU vs. CPU, CPU vs. I/O)
- Simplify Program
 - Tasks can be divided among threads
 - Each thread is responsible for its own job
 - Threads work on the same data using shared memory
 - Note that synchronization between threads is needed