

Description of Program:

Do the following sorts, Insertion, Heap, Quick, and Batchers.

Structure:

“.h” files *insertion*, *batcher*, *heap*, *quick* create the functions that will be defined in the “.c” files.

Stats. {h, c} create and define the functions, *cmp*, *move*, *swap*, and *reset*. And will keep track of the moves and comparisons used for later outputs.

Insertion Sort Function:



A list with 1 element is already sorted so it is the default.

Move the | right with now 2 elements.

If $e[1] > e[2]$, swap

Move | right 3 elements.

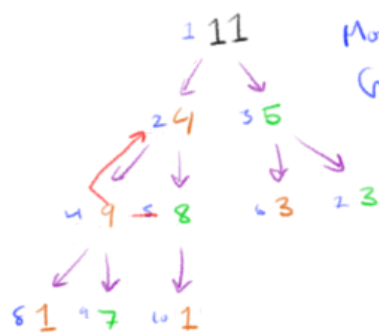
If $e[2] > e[3]$ swap.

If $e[1] > e[2]$ swap.

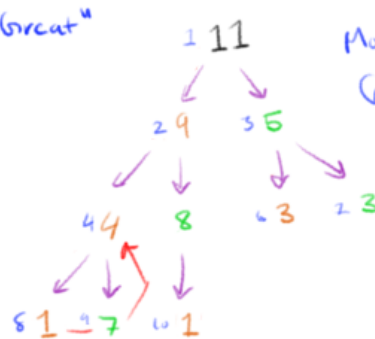
else nothing

...

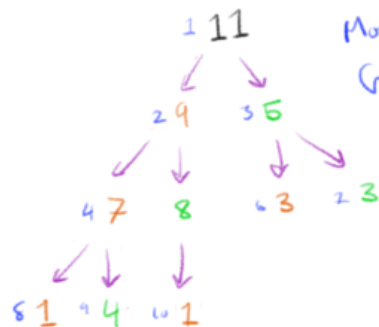
continue...



Mother 2 "Greatest"
Greatest: 9



Mother 4 "Greatest"
Greatest: 9



Mother 9 "Greatest"
Greatest: 18 "Default 9*2"

Mother > $\frac{\text{Last}}{2}$ Therefore Mother is sad and has no children



Heap Sort Function

Swap A[1] and A[Last]

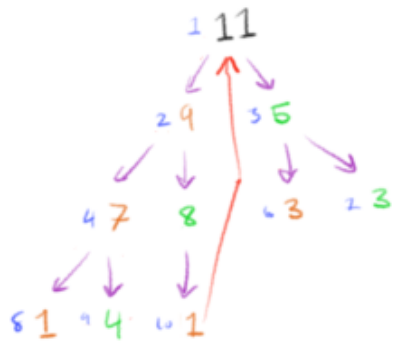
A[Last] is now greatest because A[1] if a Max Heap was greatest.

The greatest number is A[Last] now move down 1.

Last = Last - 1

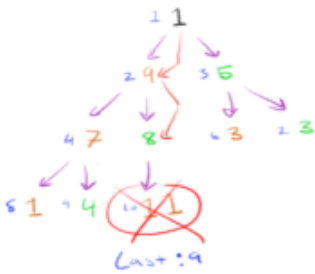
Fix heap. Make it back into a Max Heap so A[1] is the greatest of A[1 to Last]

Swap A[1] and A[Last]

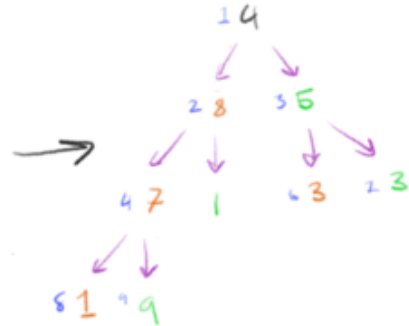
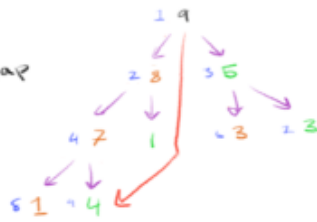


Parent is always greater than the child
just like in real life

[1 9 5 7 8 3 3 1 4 | 11] Last: 10 → 9



Fix_heap



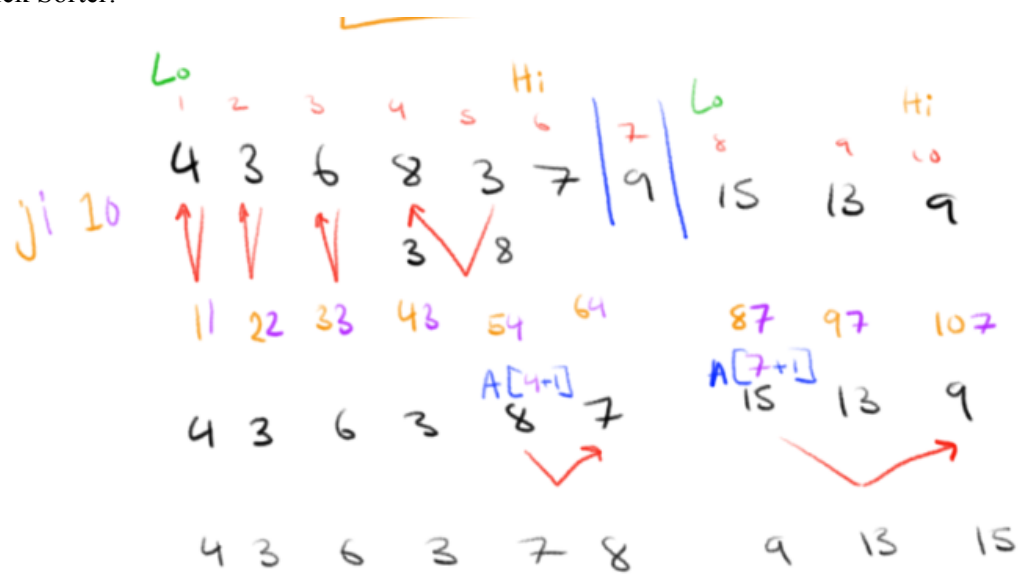
keep repeating till it is ordered

Quick Sort Function:

Partition Function:



Quick Sorter:



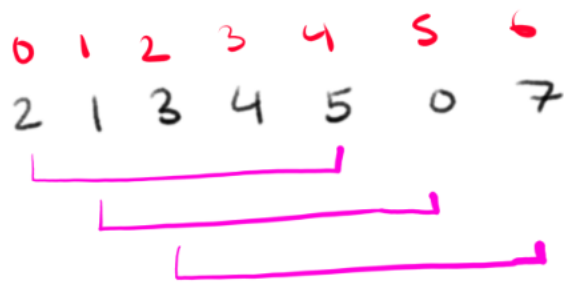
Partition function - Finds the location where to partition

- It determines the location by using a variable j to increment through array and variable i to increment under the condition that the value $A[j]$ is less than the maximum value, then it swaps with the $A[i]$.
- Variable i will eventually lag behind j as the array encounters values higher than $A[\text{last}]$.
- At the end swap the $A[\text{last}]$ with $A[i]$. Logically if $A[\text{last}]$ is greater than i elements in the list then it should be at position i in a sorted array

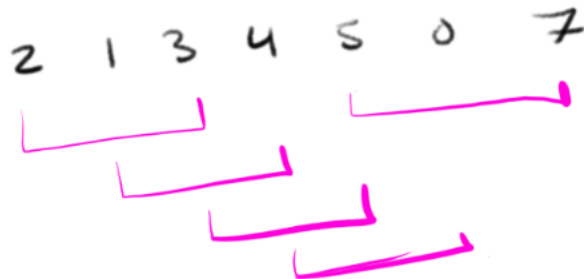
Quick Sorter - calls the Partition over and over from Partitions

- $(\text{Lo}, \text{Partition} - 1)$
- $(\text{Partition} + 1, \text{Hi})$
- Exclude Partition because it is already in its correct index. By calling it over and over until $\text{Lo} == \text{Hi}$ or the length of the section is 1, everything will become a partition and order itself.

Batcher's Odd-Even Merge Sort Function:



4 sort



2 sorting



1 sorting

0 1 2 3 4 5 6
2 1 3 4 5 0 7

—4 Sorting—

Compare(L[0], L[4]), return -1. Nothing.

Compare(L[1], L[5]), return 1. Swap.

0 1 2 3 4 5 6
2 0 3 4 5 1 7

Compare(L[2], L[6]), return -1. Nothing.

—2 Sorting—

Compare(L[0], L[2]), return -1. Nothing.

Compare(L[4], L[6]), return -1. Nothing.

Compare(L[1], L[3]), return -1. Nothing.

Compare(L[2], L[4]), return -1. Nothing.

Compare(L[3], L[5]), return 1. Swap.

0 1 2 3 4 5 6

2 0 3 1 5 4 7

—1 Sorting—

Compare(L[0], L[1]), return 1. Swap.

0 1 2 3 4 5 6

0 2 3 1 5 4 7

Compare(L[2], L[3]), return 1. Swap.

0 1 2 3 4 5 6

0 2 1 3 5 4 7

Compare(L[4], L[5]), return 1. Swap.

0 1 2 3 4 5 6

0 2 1 3 4 5 7

Compare(L[1], L[4]), return -1. Nothing.

Compare(L[3], L[6]), return -1. Nothing.

Compare(L[3], L[4]), return -1. Nothing.

Compare(L[5], L[6]), return -1. Nothing.

0 1 2 3 4 5 6

0 2 1 3 4 5 7