

Title:Implementation of Strassen's Matrix multiplication.

Theory/Description:

using Strassen's Matrix multiplication algorithm, the time consumption can be improved a little bit.

Strassen's Matrix multiplication can be performed only on **square matrices** where **n** is a **power of 2**. Order of both of the matrices are **$n \times n$** .

Divide **X**, **Y** and **Z** into four $(n/2) \times (n/2)$ matrices as represented below –

$Z = [IKJL]$ $Z = [IJKL]$ $X = [ACBD]$ $X = [ABCD]$ and $Y = [EGFH]$ $Y = [EFGH]$

Using Strassen's Algorithm compute the following –

$$\begin{aligned}M1 &:= (A+C) \times (E+F) \\M2 &:= (B+D) \times (G+H) \\M3 &:= (A-D) \times (E+H) \\M4 &:= A \times (F-H) \\M5 &:= (C+D) \times (E) \\M6 &:= (A+B) \times (H) \\M7 &:= D \times (G-E)\end{aligned}$$

Then,

$$\begin{aligned}I &:= M2 + M3 - M6 - M7 \\J &:= M4 + M6 \\K &:= M5 + M7 \\L &:= M1 - M3 - M4 - M5\end{aligned}$$

Analysis:

$$T(n) = \begin{cases} c & \text{.....} \\ f & n=1 \end{cases}$$

$7 \times T(n/2) + d \times n^2$ otherwise, where c and d are constants

Using this recurrence relation, we get $T(n) = O(n \log 7)$ $T(n) = O(n \log 7)$

Hence, the complexity of Strassen's matrix multiplication algorithm is $O(n \log 7)$ $O(n \log 7)$.

C code of two 2 by 2 matrix multiplication using Strassen's algorithm

```
*/  
#include<stdio.h>  
  
int main(){  
    int a[2][2], b[2][2], c[2][2], i, j;  
    int m1, m2, m3, m4 , m5, m6, m7;  
  
    printf("Enter the 4 elements of first matrix: ");  
    for(i = 0; i < 2; i++)  
        for(j = 0; j < 2; j++)  
            scanf("%d", &a[i][j]);  
  
    printf("Enter the 4 elements of second matrix: ");  
    for(i = 0; i < 2; i++)  
        for(j = 0; j < 2; j++)  
            scanf("%d", &b[i][j]);  
  
    printf("\nThe first matrix is\n");  
    for(i = 0; i < 2; i++){  
        printf("\n");  
        for(j = 0; j < 2; j++)  
            printf("%d\t", a[i][j]);  
    }  
  
    printf("\nThe second matrix is\n");  
    for(i = 0; i < 2; i++){  
        printf("\n");  
        for(j = 0; j < 2; j++)  
            printf("%d\t", b[i][j]);  
    }
```

m1= (a[0][0] + a[1][1]) * (b[0][0] + b[1][1]);

m2= (a[1][0] + a[1][1]) * b[0][0];

m3= a[0][0] * (b[0][1] - b[1][1]);

m4= a[1][1] * (b[1][0] - b[0][0]);

m5= (a[0][0] + a[0][1]) * b[1][1];

m6= (a[1][0] - a[0][0]) * (b[0][0]+b[0][1]);

m7= (a[0][1] - a[1][1]) * (b[1][0]+b[1][1]);

c[0][0] = m1 + m4- m5 + m7;

c[0][1] = m3 + m5;

c[1][0] = m2 + m4;

c[1][1] = m1 - m2 + m3 + m6;

printf("\nAfter multiplication using Strassen's algorithm \n");

for(i = 0; i < 2 ; i++){

printf("\n");

for(j = 0;j < 2; j++)

printf("%d\t", c[i][j]);

}

return 0;

}

VIVA-VOCE QUESTIONS:

1. Is there any optimum solution for Matrixmultiplication?

Ans: Yes. Divide and conquer method suggests Strassen's matrix multiplication method to be used. If we follow this method, the time complexity is $O(n*n*n*2.81)$ times rather $O(n*n*n*3)$ times.