
DAY-07

infix to prefix:

infix : $a*b/c*d+e*f-g*h$ --> string

prefix : $hg*fe*dcba*/*+-$ --> empty string

cur ele =

stack:

- infix expression, empty postfix expression and stack which is initially empty

step1: start scanning an infix expression from right to left

step2:

if(cur ele is an operand)

append it into the prefix expression

else//if cur ele is an operator

```
{
    while( !stack is not empty && priority(topmost ele) > priority(cur ele))
    {
        pop ele from the stack and append it into the prefix expression
    }
    push cur ele into the stack
}
```

step3: repeat step1 & step2 till the end of infix expression

step4: pop all remaining ele's from the stack one by one and append them into the prefix expression.

step5: reverse prefix expression will be the final answer

infix expression : $a*b/c*d+e*f-g*h$

postfix expression : $ab*c/d*ef*+gh*-$

prefix expression : $-+*/abcd*ef*gh$

* to convert prefix expression into its equivalent postfix:

- algorithm:

step1: start scanning prefix expression from right to left

step2:

if(cur ele is an operand)

push it onto the stack

{

 $\{$

- pop two ele's from the stack
- form a string by concatenating op1, op2 & opr(cur ele)
"op1+op2+opr"
- push result string onto stack

}

step3: repeat step1 & step2 till end of the prefix expression

prefix expression : - + * / * a b c d * e f * gh

postfix expression: $ab^*c/d^*ef^*+gh^*-$

$$ab^*c/d^*ef^*+gh^*-$$

+ postfix expression evaluation algorithm:

infix expression : 4 * 5 / 8 * 6 - 3 + 7

postfix expression :

$$\Rightarrow 4 * 5 / 8 * 6 - 3 + 7$$
$$\Rightarrow 45 * / 8 * 6 - 3 + 7$$
$$\Rightarrow 45 \cdot 8 / 6 - 3 + 7$$
$$\Rightarrow 45 \cdot 8 / 6 - 3 + 7$$
$$\Rightarrow 45 \cdot 8 / 6 \cdot 3 - + 7$$
$$\Rightarrow 45 \cdot 8 / 6 \cdot 3 - 7 +$$

postfix expression: 4 5 * 8 / 6 * 3 - 7 +

cur ele : 9 + 7

stack:

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step1: start scanning postfix expression from left to right

step2:

```
if( cur ele is an operand )
```

push it onto the stack

else//if cur ele is an operator

{

- pop two ele's from the stack
 - first popped ele will be the second operand
 - and second popped ele will be the first operand
- ```
op2 = peek(); pop();
op1 = peek(); pop();
```

- perform cur ele operation between op1 & op2  
res = op1 opr op2
- push result back onto the stack

}

**step3: repeat step2 & step3 till the end of postfix expression**

**step4: pop the final result from the stack and return it to the calling function**

- 
- + **"Queue"**: it is a collection/list of logically related similar type of elements in which element can be added from one end referred as "rear" end, whereas element can be deleted from another end referred as "front" end.
- in this list, element which was inserted first can be deleted first, so this list works in "first in first out" manner, and hence this list is also called as FIFO list/ "last in last out" i.e. LIFO list.
  - basically we can perform two operations onto queue data structure in  $O(1)$  time:
    1. enqueue: to insert/push/add an ele into the queue from rear end
    2. dequeue: to delete/remove/pop an ele from the queue which is at front end
- there are different types of queue:
1. "linear queue"
  2. "circular queue"
  3. "priority queue": it is a type of queue in which element can added randomly (i.e. without checking its priority), whereas element which is having highest priority can only be deleted first.
    - priority queue can be implemented by using a linked list, whereas it can implemented efficiently by using binary heap.
  4. "double ended queue (deque)": it is a type of queue in which elements can be added as well deleted from both the ends.
    - we can perform four basic operations on deque in  $O(1)$ :
      1. push\_front -> add\_first()
      2. push\_back -> add\_last()
      3. pop\_front -> delete\_first()
      4. pop\_back -> delete\_last()

- deque can be implemented by using doubly circular linked list

- + implementation of a linear queue by using an array : static queue
- + implementation of a linear queue by using a linked list : dynamic queue
- + implementation of a linear queue by using an array : static queue

```
struct queue
{
 int arr[5];
 int rear;
 int front;
}queue_t;
```

```
arr : int []
rear : int
front : int
```

```
is_queue_full : rear == SIZE-1
is_queue_empty : rear == -1 || front > rear
```

1. enqueue: to insert/push/add an ele into the queue from rear end:

```
step1: check is queue not full
step2: increment the value of rear by 1
step3: insert/push an ele into the queue from rear end
step4: if(front == -1)
 front = 0
```

2. dequeue: to delete/remove/pop an ele from the queue which is at front end:

```
step1: check is queue not empty
step2: increment the value of front by 1 (i.e. we are achieving deletion of an ele
from the queue).
```

+ "limitation of linear queue":

- vacant places cannot be reutilized in a linear queue

```
rear = 4 & front = 0
rear = 0 & front = 1
rear = 1 & front = 2
rear = 2 & front = 3
rear = 3 & front = 4
```

- if front is at next position of rear we can say queue is full:

```
queue_full: front == (rear + 1) % SIZE
```

```
for rear = 0, front = 1
front == (rear + 1) % SIZE
1 == (0+1) % 5
1 == 1 % 5
1 == 1
```

```
for rear = 1, front = 2
front == (rear + 1) % SIZE
2 == (1+1) % 5
```

```
2 == 2 % 5
2 == 2
```

```
for rear = 2, front = 3
front == (rear + 1) % SIZE
3 == (2+1) % 5
3 == 3 % 5
3 == 3
```

```
for rear = 3, front = 4
front == (rear + 1) % SIZE
4 == (3+1) % 5
4 == 4 % 5
4 == 4
```

```
for rear = 4, front = 0
front == (rear + 1) % SIZE
0 == (4+1) % 5
0 == 5 % 5
0 == 0
```

```
front++;
=> front = front + 1
```

```
for incrementing value of front by 1
front = (front+1)%SIZE
```

```
if front = 0 => front = (front+1)%SIZE => (0+1)%5 => 1%5 => 1
if front = 1 => front = (front+1)%SIZE => (1+1)%5 => 2%5 => 2
if front = 2 => front = (front+1)%SIZE => (2+1)%5 => 3%5 => 3
if front = 3 => front = (front+1)%SIZE => (3+1)%5 => 4%5 => 4
if front = 4 => front = (front+1)%SIZE => (4+1)%5 => 5%5 => 0
```