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Ex.no	:	1 (a)
Date	:	

MIN HEAP (Insertion ,Delete Min ,Delete Max)

AIM:

To write a python program for mi heap (Insetion, Delete min, Delete min)

ALGORITHM:

Step 1: Insertion:

- Add the new element at the end of the heap.
- Heapify Up: Swap the element with its parent until the heap property is restored.

Step 2: Delete Min:

- Swap the root with the last element.
- Remove the last element (previously the root).

Step 3: Heapify Down:

• Swap the root with its smallest child until the heap property is restored.

Step 4: Delete Max:

- Swap the root with the largest child (either left or right).
- Remove the last element (previously the root).

Step 5 : Heapify Down:

• Swap the new root with its smallest child until the heap property is restored.

```
class MinHeap:
    def __init__(self):
        self.heap = []

    def insert(self, value):
        self.heap.append(value)
        self._heapify_up(len(self.heap) - 1)

    def delete_min(self):
        if len(self.heap) == 0:
            return None

    if len(self.heap) == 1:
        return self.heap.pop()

    root = self.heap[0]
```

```
self.heap[0] = self.heap.pop()
     self._heapify_down(0)
     return root
  def delete max(self):
     if len(self.heap) == 0:
       return None
     if len(self.heap) == 1:
       return self.heap.pop()
     max_child_idx = self._find_max_child(0)
     self.heap[0], self.heap[max_child_idx] = self.heap[max_child_idx], self.heap[0]
     deleted_max = self.heap.pop()
     self._heapify_down(max_child_idx)
     return deleted max
  def _heapify_up(self, index):
     parent_index = (index - 1) // 2
     while index > 0 and self.heap[index] < self.heap[parent_index]:
       self.heap[index], self.heap[parent_index] = self.heap[parent_index], self.heap[index]
       index = parent_index
       parent_index = (index - 1) // 2
  def _heapify_down(self, index):
     left\_child\_index = 2 * index + 1
     right\_child\_index = 2 * index + 2
     smallest = index
     if left_child_index < len(self.heap) and self.heap[left_child_index] < self.heap[smallest]:
       smallest = left_child_index
     if right_child_index < len(self.heap) and self.heap[right_child_index] < self.heap[smallest]:
       smallest = right_child_index
     if smallest != index:
       self.heap[index], self.heap[smallest] = self.heap[smallest], self.heap[index]
       self._heapify_down(smallest)
  def _find_max_child(self, index):
     left child index = 2 * index + 1
     right\_child\_index = 2 * index + 2
     if right_child_index >= len(self.heap):
       return left_child_index
     return left_child_index if self.heap[left_child_index] > self.heap[right_child_index] else
right_child_index
min_heap = MinHeap()
```

```
value1 = int(input("Enter a value: "))
min_heap.insert(value1)
value2 = int(input("Enter another value: "))
min_heap.insert(value2)
value3 = int(input("Enter one more value: "))
min_heap.insert(value3)
value4 = int(input("Enter another value: "))
min_heap.insert(value4)
value5 = int(input("Enter one more value: "))
min_heap.insert(value5)
print("Values in 'min_heap':", min_heap.heap)
print("Min Heap:", min_heap.heap)
print("Delete Min:", min_heap.delete_min())
print("Min Heap after Delete Min:", min heap.heap)
print("Delete Max:", min_heap.delete_max())
print("Min Heap after Delete Max:", min_heap.heap)
```

Enter a value: 9

Enter another value: 3 Enter one more value: 6 Enter one more value: 1 Enter one more value: 5

Values in 'min_heap': [1, 3, 6, 9, 5]

Min Heap: [1, 3, 6, 9, 5]

Delete Min: 1

Min Heap after Delete Min: [3, 5, 6, 9]

Delete Max: 9

Min Heap after Delete Max: [6, 5, 3]

CONTENTS	MARKS ALLOTED	MARKS OBTAINED
PROGRAM AND EXECUTION	15	
VIVA-VOCE	10	
TOTAL	25	

RESULT:

Thus the program of the Implementation of min heap is executed and output is verified successfully.

Ex.no	:	1 (b)
Date	:	

SKEW HEAP (Priority Queue Operations)

AIM:

To write a python program for skew heap (Priority Queue Operations).

ALGORITHM:

Step 1: Insertion:

- Create a new node with the desired value.
- Merge the new node with the existing skew heap.

Step 2: Delete Min (or Delete Max):

- Merge the left and right children of the root, effectively removing the root.
- The merged result becomes the new root.

```
class SkewNode:
  def __init__(self, value):
     self.value = value
     self.left = None
     self.right = None
class SkewHeap:
  def __init__(self):
     self.root = None
  def merge(self, h1, h2):
     if not h1:
       return h2
     if not h2:
       return h1
     if h1.value > h2.value:
       h1, h2 = h2, h1 # Swap h1 and h2
     h1.right, h1.left = h1.left, self.merge(h1.right, h2)
     return h1
  def insert(self, value):
     new_node = SkewNode(value)
     self.root = self.merge(self.root, new_node)
  def delete_min(self):
     if not self.root:
       return None
```

```
min_value = self.root.value
    self.root = self.merge(self.root.left, self.root.right)
    return min_value

skew_heap = SkewHeap()
skew_heap.insert(4)
skew_heap.insert(2)
skew_heap.insert(7)
skew_heap.insert(6)

print("Skew Heap after insertion:", skew_heap.root.value)

min_val = skew_heap.delete_min()
print("Deleted Min:", min_val)
print("Skew Heap after deletion:", skew_heap.root.value)
min_val = skew_heap.delete_min()
print("Deleted Min:", min_val)
```

Skew Heap after insertion: 2

Deleted Min: 2

Skew Heap after deletion: 4

Deleted Min: 4

CONTENTS	MARKS ALLOTED	MARKS OBTAINED
PROGRAM AND EXECUTION	15	
VIVA-VOCE	10	
TOTAL	25	

RESULT:

Thus ,the program of the Implementation of skew heap is executed and output is verified.

Ex.no: 1(c)	EIDONA CCI HEAD (Deignites Organo Organotions)
Date :	FIBONACCI HEAP (Priority Queue Operations)

To write a python program for Fibonacci heap (Priority Queue Operations).

ALGORITHM:

Step 1: Insertion:

- Create a new tree with the given value.
- Insert the new tree into the root list.

Step 2 :Delete Min:

- Identify the tree with the minimum root.
- Remove the root of that tree and merge its children with the root list.
- Consolidate the heap by combining trees with the same degree.

Step 3 : Delete Max:

• Same as Delete Min, but identify the tree with the maximum root.

```
# Fibonacci Heap in python
import math
# Creating fibonacci tree
class FibonacciTree:
  def __init__(self, value):
     self.value = value
     self.child = []
     self.order = 0
  # Adding tree at the end of the tree
  def add_at_end(self, t):
     self.child.append(t)
     self.order = self.order + 1
# Creating Fibonacci heap
class FibonacciHeap:
  def __init__(self):
     self.trees = []
     self.least = None
     self.count = 0
  # Insert a node
  def insert_node(self, value):
     new_tree = FibonacciTree(value)
```

```
self.trees.append(new_tree)
  if (self.least is None or value < self.least.value):
     self.least = new_tree
  self.count = self.count + 1
# Get minimum value
def get_min(self):
  if self.least is None:
    return None
   return self.least.value
# Extract the minimum value
def extract_min(self):
  smallest = self.least
  if smallest is not None:
     for child in smallest.child:
       self.trees.append(child)
     self.trees.remove(smallest)
     if self.trees == []:
       self.least = None
     else:
       self.least = self.trees[0]
       self.consolidate()
     self.count = self.count - 1
     return smallest.value
# Consolidate the tree
def consolidate(self):
  aux = (floor\_log(self.count) + 1) * [None]
  while self.trees != []:
     x = self.trees[0]
     order = x.order
     self.trees.remove(x)
     while aux[order] is not None:
       y = aux[order]
       if x.value > y.value:
          x, y = y, x
       x.add_at_end(y)
       aux[order] = None
       order = order + 1
     aux[order] = x
  self.least = None
  for k in aux:
     if k is not None:
       self.trees.append(k)
       if (self.least is None
             or k.value < self.least.value):
          self.least = k
```

```
return math.frexp(x)[1] - 1

fibonacci_heap = FibonacciHeap()
fibonacci_heap.insert_node(7)
fibonacci_heap.insert_node(3)
fibonacci_heap.insert_node(17)
fibonacci_heap.insert_node(24)
```

print('the minimum value of the fibonacci heap: { }'.format(fibonacci_heap.get_min()))
print('the minimum value removed: { }'.format(fibonacci_heap.extract_min()))

OUTPUT:

def floor_log(x):

The minimum value of the fibonacci heap: 3

The minimum value removed: 3

CONTENTS	MARKS ALLOTED	MARKS OBTAINED
PROGRAM AND EXECUTION	15	
VIVA-VOCE	10	
TOTAL	25	

RESULT:

Thus the program of the Implementation of fibonacci heap is executed and output is verified.

Ex.no : 2(a)	
Date :	AVL TREES (Insertion ,Delete and search)

To Write a Python Program to implement a insert, delete, search a element by using the AVL tree properly.

ALGORITHM:

Step 1.Insertion:

- Perform standard BST insertion.
- Update height of each node from the newly inserted node to the root.
- Check balance factor to see if rotation is required.
- Right Rotation: If balance factor > 1 and new node is inserted into the left subtree of the left child.
- Left Rotation: If balance factor < -1 and new node is inserted into the right subtree of the right child.
- Left-Right Rotation: If balance factor > 1 and new node is inserted into the right subtree of the left child.
- Right-Left Rotation: If balance factor < -1 and new node is inserted into the left subtree of the right child.

Step 2.Deletion:

- Perform standard BST deletion.
- Update height of each node from the deleted node to the root.
- Check balance factor and perform necessary rotations if the tree becomes unbalanced.

Step 3.Search:

Perform standard BST search.

```
class AVLNode:

def __init__(self, key):
    self.key = key
    self.left = None
    self.right = None
    self.height = 1

class AVLTree:
    def __init__(self):
    self.root = None
    def height(self, node):
    if not node:
    return 0
```

```
return node.height
def balance(self, node):
  if not node:
     return 0
  return self.height(node.left) - self.height(node.right)
def right_rotate(self, y):
  x = y.left
  T2 = x.right
  x.right = y
  y.left = T2
  y.height = 1 + max(self.height(y.left), self.height(y.right))
  x.height = 1 + max(self.height(x.left), self.height(x.right))
  return x
def left_rotate(self, x):
  y = x.right
  T2 = y.left
  y.left = x
  x.right = T2
  x.height = 1 + max(self.height(x.left), self.height(x.right))
  y.height = 1 + max(self.height(y.left), self.height(y.right))
  return y
def insert(self, node, key):
    if not node:
     return AVLNode(key)
  if key < node.key:
     node.left = self.insert(node.left, key)
  else:
     node.right = self.insert(node.right, key)
  node.height = 1 + max(self.height(node.left), self.height(node.right))
  balance = self.balance(node)
  if balance > 1 and key < node.left.key:
     return self.right_rotate(node)
```

```
if balance < -1 and key > node.right.key:
      return self.left_rotate(node)
   if balance > 1 and key > node.left.key:
      node.left = self.left_rotate(node.left)
      return self.right_rotate(node)
   if balance < -1 and key < node.right.key:
      node.right = self.right_rotate(node.right)
      return self.left_rotate(node)
   return node
 def delete(self, root, key):
    if not root:
      return root
   if key < root.key:
      root.left = self.delete(root.left, key)
   elif key > root.key:
      root.right = self.delete(root.right, key)
    else:
      if not root.left or not root.right:
         temp = root.left if root.left else root.right
         root = None
         return temp
      temp = self.min_value_node(root.right)
      root.key = temp.key
      root.right = self.delete(root.right, temp.key)
if not root:
      return root
    root.height = 1 + max(self.height(root.left), self.height(root.right))
   balance = self.balance(root)
   if balance > 1 and self.balance(root.left) >= 0:
      return self.right_rotate(root)
   if balance < -1 and self.balance(root.right) <= 0:
      return self.left_rotate(root)
    if balance > 1 and self.balance(root.left) < 0:
```

```
root.left = self.left_rotate(root.left)
       return self.right_rotate(root)
     if balance < -1 and self.balance(root.right) > 0:
       root.right = self.right_rotate(root.right)
       return self.left_rotate(root)
     return root
  def min_value_node(self, node):
     current = node
     while current.left:
        current = current.left
     return current
  def search(self, root, key):
     if not root or root.key == key:
       return root
     if root.key < key:
       return self.search(root.right, key)
     return self.search(root.left, key)
  def inorder(self, root):
     if root:
        self.inorder(root.left)
       print(root.key, end=' ')
       self.inorder(root.right)
avl = AVLTree()
avl.root = avl.insert(avl.root, 10)
avl.root = avl.insert(avl.root, 20)
avl.root = avl.insert(avl.root, 30)
avl.root = avl.insert(avl.root, 40)
avl.root = avl.insert(avl.root, 50)
avl.root = avl.insert(avl.root, 25)
print("Inorder traversal of AVL tree:")
avl.inorder(avl.root)
print()
avl.root = avl.delete(avl.root, 20)
print("Inorder traversal of AVL tree after deletion of 20:")
avl.inorder(avl.root)
```

```
print()
```

```
search_key = 30
if avl.search(avl.root, search_key):
   print(f"{search_key} found in AVL tree.")
else:
   print(f"{search_key} not found in AVL tree.")
```

Inorder traversal of AVL tree:

10 20 25 30 40 50

Inorder traversal of AVL tree after deletion of 20:

10 25 30 40 50

30 found in AVL tree

CONTENTS	MARKS ALLOTED	MARKS OBTAINED
PROGRAM AND EXECUTION	15	
VIVA-VOCE	10	
TOTAL	25	

RESULT:

Thus ,the program of the Implementation of AVL TREE is executed and output is verified.

Ex.no: 2(c)	
Date:	B-TREES(Insertion ,Delete and search)

To Write a Python Program to insert, delete and search the given element elements by using the B- trees.

ALGORITHM:

Step 1. Insertion:

- Perform standard BST insertion.
- Update height of each node from the newly inserted node to the root.
- Check balance factor to see if rotation is required.
- Right Rotation: If balance factor > 1 and new node is inserted into the left subtree of the left child.
- Left Rotation: If balance factor < -1 and new node is inserted into the right subtree of the right child.
- Left-Right Rotation: If balance factor > 1 and new node is inserted into the right subtree of the left child.
- Right-Left Rotation: If balance factor < -1 and new node is inserted into the left subtree of the right child.

Step 2. Deletion:

- Perform standard BST deletion.
- Update height of each node from the deleted node to the root.
- Check balance factor and perform necessary rotations if the tree becomes unbalanced.

Step 3. Search:

• Perform standard BST search.

```
class BTreeNode:
    def __init__(self, keys=[], children=[], is_leaf=True, max_keys=4):
        self.keys = keys
        self.children = children
        self.is_leaf = is_leaf
        if not max_keys:
            # Default value for max_keys
            self.max_keys = 4
        else:
            self.max_keys = max_keys

class BTree:
    def __init__(self, max_keys=4):
        self.root = BTreeNode(max_keys=max_keys)
```

```
def search(self, node, key):
  i = 0
  while i < len(node.keys) and key > node.keys[i]:
  if i < len(node.keys) and key == node.keys[i]:
    return node, i
  if node.is_leaf:
     return None, None
  return self.search(node.children[i], key)
def insert(self, key):
  root = self.root
  if len(root.keys) == root.max_keys:
     new_root = BTreeNode(keys=[root.keys.pop(len(root.keys)//2)], children=[root])
     self.split child(new root, 0)
     self.root = new_root
    root = new_root
  self.insert_non_full(root, key)
def insert_non_full(self, node, key):
  i = len(node.keys) - 1
  if node.is leaf:
     node.keys.append(None)
     while i \ge 0 and key < node.keys[i]:
       node.keys[i+1] = node.keys[i]
       i = 1
    node.keys[i+1] = key
  else:
     while i \ge 0 and key < node.keys[i]:
       i = 1
    i += 1
     if len(node.children[i].keys) == node.children[i].max_keys:
       self.split_child(node, i)
       if key > node.keys[i]:
         i += 1
     self.insert_non_full(node.children[i], key)
def split_child(self, parent, i):
  node_to_split = parent.children[i]
  new_node = BTreeNode(max_keys=node_to_split.max_keys, is_leaf=node_to_split.is_leaf)
  parent.keys.insert(i, node_to_split.keys[len(node_to_split.keys)//2])
  parent.children.insert(i+1, new_node)
  new_node.keys = node_to_split.keys[len(node_to_split.keys)//2+1:]
  node_to_split.keys = node_to_split.keys[:len(node_to_split.keys)//2]
  if not node_to_split.is_leaf:
```

```
new_node.children = node_to_split.children[len(node_to_split.keys)+1:]
     node_to_split.children = node_to_split.children[:len(node_to_split.keys)+1]
def delete(self, key):
  self.delete_recursive(self.root, key)
def delete_recursive(self, node, key):
  while i < len(node.keys) and key > node.keys[i]:
    i += 1
  if i < len(node.keys) and key == node.keys[i]:
    if node.is leaf:
       del node.keys[i]
     else:
       # Replace with predecessor and delete predecessor from child
       if len(node.children[i]) >= node.max_keys/2:
          pred = self.get_predecessor(node, i)
          node.keys[i] = pred.keys.pop()
       # Replace with successor and delete successor from child
       elif len(node.children[i+1]) >= node.max_keys/2:
          succ = self.get_successor(node, i)
          node.keys[i] = succ.keys.pop(0)
       # Merge nodes
       else:
          self.merge(node, i)
          self.delete_recursive(node.children[i], key)
  else:
     if node.is_leaf:
       return
     elif len(node.children[i]) == node.max_keys/2:
       self.fix_borrow_or_merge(node, i)
     self.delete_recursive(node.children[i], key)
def get_predecessor(self, node, i):
  curr = node.children[i]
  while not curr.is_leaf:
     curr = curr.children[-1]
  return curr
def get_successor(self, node, i):
  curr = node.children[i+1]
  while not curr.is leaf:
     curr = curr.children[0]
  return curr
def merge(self, node, i):
  child = node.children[i]
  sibling = node.children[i+1]
```

```
child.keys.append(node.keys[i])
  child.keys += sibling.keys
  if not child.is_leaf:
     child.children += sibling.children
  del node.keys[i]
  del node.children[i+1]
def fix_borrow_or_merge(self, node, i):
  if i > 0 and len(node.children[i-1]) > node.max_keys/2:
     # Borrow from left sibling
     left_sibling = node.children[i-1]
     child = node.children[i]
     if child.is_leaf:
       child.keys.insert(0, node.keys[i-1])
       node.keys[i-1] = left_sibling.keys.pop()
     else:
       child.keys.insert(0, node.keys[i-1])
       node.keys[i-1] = left_sibling.keys.pop(-1)
       child.children.insert(0, left_sibling.children.pop(-1))
  elif i < len(node.children) - 1 and len(node.children[i+1]) > node.max_keys/2:
    # Borrow from right sibling
     right_sibling = node.children[i+1]
     child = node.children[i]
     if child.is_leaf:
       child.keys.append(node.keys[i])
       node.keys[i] = right_sibling.keys.pop(0)
     else:
       child.keys.append(node.keys[i])
       node.keys[i] = right_sibling.keys.pop(0)
       child.children.append(right_sibling.children.pop(0))
  else:
    # Merge with sibling
    if i > 0:
       self.merge(node, i-1)
       del node.keys[i-1]
     else:
       self.merge(node, i)
       del node.keys[i]
def inorder_traversal(self, node):
  if node:
    i = 0
     while i < len(node.keys):
       self.inorder_traversal(node.children[i])
       print(node.keys[i], end=' ')
       i += 1
     self.inorder_traversal(node.children[i])
```

```
btree = BTree()
btree.insert(10)
btree.insert(20)
btree.insert(5)
btree.insert(6)
btree.insert(12)
btree.insert(30)
btree.insert(7)
btree.insert(17)
print("Inorder traversal of B-Tree:")
btree.inorder_traversal(btree.root)
print()
btree.delete(12)
print("Inorder traversal of B-Tree after deletion of 12:")
btree.in order\_traversal(btree.root)
print()
search_key = 17
if btree.search(btree.root, search_key):
  print(f"{search_key} found in B-Tree.")
else:
  print(f"{search_key} not found in B-Tree.")
```

Inorder traversal of B-Tree:

5 6 7 10 12 17 20 30

Inorder traversal of B-Tree after deletion of 12:

CONTENTS	MARKS ALLOTED	MARKS OBTAINED
PROGRAM AND EXECUTION	15	
VIVA-VOCE	10	
TOTAL	25	

RESULT:

Thus ,the program of the Implementation of B-TREE is executed and output is verified.

Ex.no: 2(d)	
Date :	RED-BLACK TREE

To Write a Python Program to insert, delete and search the given set of elements by using the Red- Black tree.

ALGORITHM:

Step 1. **Insertion Operation:**

- Perform standard BST insertion.
- After insertion, fix any violations of the Red-Black tree properties by rotating and recoloring nodes if necessary

There are four cases for fixing violations:

- Case 1: The uncle of the newly inserted node is red.
- Case 2: The uncle of the newly inserted node is black and the newly inserted node is a right child of a left child or vice versa.
- ➤ Case 3: The uncle of the newly inserted node is black and the newly inserted node is a left child of a left child or a right child of a right child.
- Case 4: The uncle of the newly inserted node is black and the newly inserted node is a left child of a right child or a right child of a left child.

Step 2. Deletion Operation:

- Perform standard BST deletion.
- After deletion, fix any violations of the Red-Black tree properties by rotating and recoloring nodes if necessary.

There are three cases for fixing violations after deletion:

- Case 1: Sibling is red.
- Case 2: Sibling is black and both of its children are black.
- Case 3: Sibling is black, its left child is red, and its right child is black.
- > Case 4: Sibling is black, its right child is red.
- Case 5: Sibling is black, its left child is black, and its right child is red.

Step 3. Search Operation:

> Perform standard BST search.

```
RED = True
BLACK = False
class Node:
    def __init__(self, key, color=RED):
        self.key = key
        self.left = None
```

```
self.right = None
     self.parent = None
     self.color = color
class RedBlackTree:
  def __init__(self):
     self.nil = Node(None, color=BLACK)
     self.root = self.nil
  def left_rotate(self, x):
     y = x.right
     x.right = y.left
     if y.left != self.nil:
        y.left.parent = x
     y.parent = x.parent
     if x.parent == self.nil:
        self.root = y
     elif x == x.parent.left:
        x.parent.left = y
     else:
       x.parent.right = y
     y.left = x
     x.parent = y
  def right_rotate(self, y):
     x = y.left
     y.left = x.right
     if x.right != self.nil:
       x.right.parent = y
     x.parent = y.parent
     if y.parent == self.nil:
       self.root = x
     elif y == y.parent.right:
       y.parent.right = x
     else:
       y.parent.left = x
     x.right = y
     y.parent = x
  def insert(self, key):
     new\_node = Node(key)
     y = self.nil
     x = self.root
     while x != self.nil:
       y = x
       if new_node.key < x.key:
          x = x.left
       else:
          x = x.right
     new_node.parent = y
     if y == self.nil:
     self.root = new_node
     elif new_node.key < y.key:
```

```
y.left = new_node
  else:
     y.right = new_node
  new_node.left = self.nil
  new_node.right = self.nil
  new_node.color = RED
  self.insert_fixup(new_node)
def insert_fixup(self, z):
  while z.parent.color == RED:
     if z.parent == z.parent.parent.left:
       y = z.parent.parent.right
       if y.color == RED:
         z.parent.color = BLACK
         y.color = BLACK
         z.parent.parent.color = RED
         z = z.parent.parent
       else:
         if z == z.parent.right:
            z = z.parent
            self.left_rotate(z)
         z.parent.color = BLACK
         z.parent.parent.color = RED
         self.right_rotate(z.parent.parent)
     else:
       y = z.parent.parent.left
       if y.color == RED:
         z.parent.color = BLACK
         y.color = BLACK
         z.parent.parent.color = RED
         z = z.parent.parent
       else:
         if z == z.parent.left:
            z = z.parent
            self.right_rotate(z)
         z.parent.color = BLACK
         z.parent.parent.color = RED
         self.left_rotate(z.parent.parent)
  self.root.color = BLACK
def transplant(self, u, v):
  if u.parent == self.nil:
     self.root = v
  elif u == u.parent.left:
     u.parent.left = v
     u.parent.right = v
  v.parent = u.parent
def minimum(self, x):
```

```
while x.left != self.nil:
     x = x.left
  return x
def delete_node(self, z):
  y = z
  y_original_color = y.color
  if z.left == self.nil:
     x = z.right
     self.transplant(z, z.right)
  elif z.right == self.nil:
     x = z.left
     self.transplant(z, z.left)
  else:
     y = self.minimum(z.right)
     y_original_color = y.color
     x = y.right
     if y.parent == z:
       x.parent = y
     else:
       self.transplant(y, y.right)
       y.right = z.right
       y.right.parent = y
     self.transplant(z, y)
     y.left = z.left
     y.left.parent = y
     y.color = z.color
  if y_original_color == BLACK:
     self.delete_fixup(x)
def delete_fixup(self, x):
  while x = self.root and x.color == BLACK:
     if x == x.parent.left:
       w = x.parent.right
       if w.color == RED:
          w.color = BLACK
          x.parent.color = RED
          self.left_rotate(x.parent)
          w = x.parent.right
       if w.left.color == BLACK and w.right.color == BLACK:
          w.color = RED
          x = x.parent
       else:
          if w.right.color == BLACK:
            w.left.color = BLACK
            w.color = RED
            self.right_rotate(w)
```

```
w = x.parent.right
            w.color = x.parent.color
            x.parent.color = BLACK
            w.right.color = BLACK
            self.left_rotate(x.parent)
            x = self.root
       else:
          w = x.parent.left
          if w.color == RED:
            w.color = BLACK
            x.parent.color = RED
            self.right_rotate(x.parent)
            w = x.parent.left
          if w.right.color == BLACK and w.left.color == BLACK:
            w.color = RED
            x = x.parent
          else:
            if w.left.color == BLACK:
              w.right.color = BLACK
              w.color = RED
              self.left_rotate(w)
              w = x.parent.left
            w.color = x.parent.color
            x.parent.color = BLACK
            w.left.color = BLACK
            self.right_rotate(x.parent)
            x = self.root
     x.color = BLACK
  def search(self, key):
     return self.search_recursive(self.root, key)
  def search_recursive(self, node, key):
    if node == self.nil or key == node.key:
       return node
    if key < node.key:
         return self.search_recursive(node.left, key)
        return self.search_recursive(node.right, key)
  def inorder_traversal(self, node):
     if node != self.nil:
       self.inorder_traversal(node.left)
       print(node.key, end=' ')
       self.inorder_traversal(node.right)
rb_tree = RedBlackTree()
rb_tree.insert(10)
```

```
rb_tree.insert(20)
rb_tree.insert(30)
rb_tree.insert(40)
rb_tree.insert(50)
rb_tree.insert(15)
rb_tree.insert(25)
print("Inorder traversal of Red-Black Tree:")
rb_tree.inorder_traversal(rb_tree.root)
print()
rb_tree.delete_node(rb_tree.search(20))
print("Inorder traversal of Red-Black Tree after deletion of 20:")
rb_tree.inorder_traversal(rb_tree.root)
print()
search_key = 30
if rb_tree.search(search_key):
  print(f"{search_key} found in Red-Black Tree.")
else:
  print(f"{search_key} not found in Red-Black Tree.")
```

Inorder traversal of Red-Black Tree:

10 15 20 25 30 40 50

Inorder traversal of Red-Black Tree after deletion of 20:

10 15 25 30 40 50

30 found in Red-Black Tree

CONTENTS	MARKS ALLOTED	MARKS OBTAINED
PROGRAM AND EXECUTION	15	
VIVA-VOCE	10	
TOTAL	25	

RESULT:

Thus the program of the Implementation of RED-BLACK TREE is executed and output is verified.

Ex.no: 3	
Date:	CONVEX HULL

To implement the convex hull algorithm in Python, you can use the Graham Scan algorithm.

ALGORITHM:

STEP 1: Find the point with the lowest y-coordinate (and leftmost point if there's a tie).

STEP 2: Sort the remaining points by polar angle with respect to the starting point.

STEP 3: Initialize an empty stack and push the first two points onto it.

STEP 4: Process the remaining points, adding them to the stack if they form a left turn with the last two points.

STEP 5: Return the points on the stack as the convex hull.

```
def orientation(p, q, r):
  """Find orientation of triplet (p, q, r)."""
  val = (q[1] - p[1]) * (r[0] - q[0]) - (q[0] - p[0]) * (r[1] - q[1])
  if val == 0:
     return 0 # Collinear
  return 1 if val > 0 else 2 # Clockwise or Counterclockwise
def convex_hull(points):
  """Compute the convex hull of a set of points."""
  n = len(points)
  if n < 3:
     return []
  # Find the leftmost point
  l = min(range(n), key=lambda x: points[x][0])
  hull = []
  p = 1
  q = 0
  while True:
     hull.append(p)
     q = (p + 1) \% n
     for i in range(n):
       if orientation(points[p], points[i], points[q]) == 2:
          q = i
     p = q
     if p == 1:
       break
  return [points[i] for i in hull]
# Example usage:
```

points = [(0, 3), (1, 1), (2, 2), (4, 4), (0, 0), (1, 2), (3, 1), (3, 3)] print(convex_hull(points))

OUTPUT:

[(0,3), (0,0), (3,1), (4,4)]

CONTENTS	MARKS ALLOTED	MARKS OBTAINED
PROGRAM AND EXECUTION	15	
VIVA-VOCE	10	
TOTAL	25	

RESULT:

Thus ,the program of the Implementation of convex hull executed and output is verified.

Ex.no : 4	
Date :	TOPOLOGICAL SORT

To perform topological sorting in Python, you can implement a depth-first search (DFS) based algorithm

ALGORITHM:

STEP 1: Start with a graph represented as an adjacency list.

STEP 2: Perform a depth-first search (DFS) on the graph.

STEP 3: During the DFS traversal, mark nodes as visited and recursively explore adjacent nodes.

STEP 4: When a node has no unvisited neighbors, add it to a stack or list.

STEP 5: After completing the DFS traversal, reverse the order of the nodes obtained from the stack or list to get the topological sorting.

PROGRAM:

 $g.add_edge(5, 0)$

```
from collections import defaultdict
class Graph:
  def _init_(self, vertices):
     self.graph = defaultdict(list)
     self.V = vertices
  def add_edge(self, u, v):
     self.graph[u].append(v)
  def topological_sort_util(self, v, visited, stack):
     visited[v] = True
     for i in self.graph[v]:
        if not visited[i]:
          self.topological_sort_util(i, visited, stack)
     stack.append(v)
  def topological_sort(self):
     visited = [False] * self.V
     stack = []
     for i in range(self.V):
       if not visited[i]:
          self.topological_sort_util(i, visited, stack)
     return stack[::-1]
# Example usage:
g = Graph(6)
g.add\_edge(5, 2)
```

```
g.add_edge(4, 0)
g.add_edge(4, 1)
g.add_edge(2, 3)
g.add_edge(3, 1)
print("Topological Sort:")
print(g.topological_sort())
```

Topological Sort: ['A', 'C', 'E', 'B', 'D', 'F']

CONTENTS	MARKS ALLOTED	MARKS OBTAINED
PROGRAM AND EXECUTION	15	
VIVA-VOCE	10	
TOTAL	25	

RESULT:

Thus the program of the Implementation of topological sort is executed and output is verified.

Ex.no : 5	GRAPH SEARCH ALGORITHM
Date :	

To perform topological sorting in Python, you can implement a depth-first search (DFS) based algorithm.

ALGORITHM:

STEP 1: Start with an empty data structure to track visited nodes.

STEP 2: Start with an empty data structure for processing nodes (e.g., a stack for DFS or a queue for BFS).

STEP 3: Add the starting node to the processing data structure and mark it as visited.

STEP 4:. While there are nodes in the processing data structure:

- Pop or dequeue a node.
- Process the node (e.g., print it or perform an operation).
- Mark the node as visited.

STEP 5: Terminate when there are no more nodes to process.

PROGRAM:

```
BFS:
from collections import defaultdict, deque
def bfs(graph, start):
  visited = set()
  queue = deque([start])
  visited.add(start)
  while queue:
     node = queue.popleft()
     print(node, end=" ")
     for neighbor in graph[node]:
        if neighbor not in visited:
          queue.append(neighbor)
          visited.add(neighbor)
# Example usage:
graph = {
  'A': ['B', 'C'],
  'B': ['A', 'D', 'E'],
  'C': ['A', 'F'],
  'D': ['B'],
  'E': ['B', 'F'],
  'F': ['C', 'E']
```

}

```
bfs(graph, 'A')
DFS:
def dfs(graph, start, visited=None):
  if visited is None:
     visited = set()
  print(start, end=" ")
  visited.add(start)
  for neighbor in graph[start]:
     if neighbor not in visited:
        dfs(graph, neighbor, visited)
# Example usage:
graph = {
  'A': ['B', 'C'],
  'B': ['A', 'D', 'E'],
  'C': ['A', 'F'],
  'D': ['B'],
  'E': ['B', 'F'],
  'F': ['C', 'E']
}
print("DFS traversal:")
dfs(graph, 'A')
OUTPUT:
```

BFS traversal:

ABCDEF

CONTENTS	MARKS ALLOTED	MARKS OBTAINED
PROGRAM AND EXECUTION	15	
VIVA-VOCE	10	
TOTAL	25	

RESULT:

Thus the program of the Implementation of graph search algorithms is executed and output is verified.

Ex.no: 7

IMPLEMENTATION OF RANDOMIZED ALGORIT IMPLEMENTATION OF RANDOMIZED ALGORITHMS HMS

Date:

AIM:

The aim of implementing network flow and linear programming problems is to model and solve optimization challenges. In network flow, the goal is to find the maximumflow in a network, while linear programming involves optimizing a linear objective function subject to linear equality and inequality constraints.

ALGORITHM:

Model the Network Flow Problem:

STEP 1: Define the graph representing the network.

STEP 2: Assign capacities to edges representing flow constraints.

STEP 3: Define source and sink nodes.

STEP 4: Solve the Network Flow Problem.

STEP 5: Use algorithms like Ford-Fulkerson or Edmonds-Karp to find the maximum flow.

Model the Linear Programming Problem:

STEP 1: Define decision variables.

STEP 2: Formulate the objective function.

STEP 3: Add constraints.

STEP 4: Solve the Linear Programming Problem.

STEP 5: Use libraries like scipy.optimize.linprog to find the optimal solution.

PROGRAM:

NETWORK FLOW PROBLEM:

import networkx as nx

Step 1: Model the Network Flow Problem G = nx.DiGraph()

G.add_edge('source', 'A', capacity=10) G.add_edge('source', 'B', capacity=5) G.add_edge('A', 'C', capacity=9) G.add_edge('A', 'B', capacity=3) G.add_edge('B', 'C', capacity=7) G.add_edge('B', 'sink', capacity=8) G.add_edge('C', 'sink', capacity=12)

Step 2: Solve the Network Flow Problem

max_flow_value, flow_dict = nx.maximum_flow(G, 'source', 'sink') print("Maximum Flow Value:", max_flow_value)

print("Flow Dict:", flow_dict)

LINEAR PROGRAMMING PROBLEM:

from scipy.optimize import linprog

Step 3: Model the Linear Programming Problem

c = [-3, -2] #Coefficients of the objective function (to minimize)

A = [[1, 1], # Coefficients of inequality constraints [1, 0], [0, 1]]

 $b = [10, 8, 5] \# Right-hand side of inequality constraints x0_bounds = (0, None) \# Bounds for decision variables x1_bounds = (0, None)$

Step 4: Solve the Linear Programming Problem

res = linprog(c, A_ub=A, b_ub=b, bounds=[x0_bounds, x1_bounds], method='highs') print("Optimal Solution:", res.x)

print("Optimal Objective Value:", res.fun)

OUTPUT:

NETWORK FLOW PROBLEM:

Maximum Flow Value: 15

Flow Dict: {'source': {'A': 10, 'B': 5}, 'A': {'C': 9, 'B': 1}, 'B': {'C': 5, 'sink': 5}, 'C': {'sink': 12}, 'sink': {}}

LINEAR PROGRAMMING PROBLEM:

Optimal Solution: [4. 1.]

Optimal Objective Value: -14.0

CONTENTS	MARKS ALLOTED	MARKS OBTAINED
PROGRAM AND EXECUTION	15	
VIVA-VOCE	10	
TOTAL	25	

RESULT:

Thus, the output represent the maximum flow value in the network and the optimal solution to the linear programming problem are executed and output is verified.