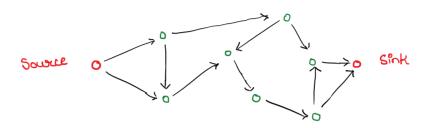
Flow Networks

Suppose that a Network consists of a Source, a Sink, and these two nodes are connected via switches. The bandwidth of each connection isn't the same. At the same time, each of the switch does cut-through switching without any queue capacity.



This system can be modeled using a directed graph with two special nodes, the source and the sink.

Let the graph be G(V,E), source be s and sink be t.

- By definition, no edge terminates at s and no edge begins at t.

Each edge has a capacity c(e), which is the maximum possible capacity of that edge. We define fle), as the capacity of the edge in use currently.

* <u>Floω</u>

Following the example, each bridge has a data stream flowing through it. As no data can be accumulated, Inward flow = Outward flow

(Similar to Kirchoff's law from electricity)

- The outwards flow for a node is represented by + (v), and is given by>

$$\phi_{\rightarrow}(h) = \sum_{n \in \mathcal{N}} \phi_{(n,h)}$$

The value of inwards flow is similarly defined, and is represented as flow.

- Value of the flow 10

The total amount of "data" flowing through the network, It is given by IfI, and is calculated as shown.

$$| \downarrow \downarrow | = \sum_{\substack{v \in V \\ (s,v) \in E}} \downarrow (s,v) = \downarrow \rightarrow (s) = \sum_{\substack{v \in V \\ (v,t) \in E}} \downarrow (v,t) = \downarrow \leftarrow (t)$$